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Mork Shnozahi
CPKS 350
                           1. Design an algorithm that finds such
Homework 6
                               a y for any given a and B. Also, analyze
                               its complexity.
                               Table delnition.
                               Let DP[i][i][k] be the longth of the longest
                                Subsequence of strings a[1...1] and B[1...i], k can represent
                                        on additional Statale
                                 Base case:
                                       For i = 0 or j=0
                                        DP[i][i][k]=o Farall k
                                   Recurrence relation
                                    Far 170 and 120
                                    if alia == BCJ3:
                      012 ... 1
                                            1fa[i] != 19, or k !=1:
                    0000 ... 0
                                                DP[i][j][o] = DP[i-1][j-1][o]+1
                                               PP [ i] [ i] [ i] = PP [ i - 1] [ j-1] [ 1] -1] ( if a ( i) = = 6
                                            else
                                                                                and occi-1] == 'a)
                                                  DP[:][:][:][:]=mox(0P[:-1][:-1][:],0P[:-1][:]]
                                      else
                                         PPCIJEJJEOJ=MAX (PP[i-1][JJCOJ, PP[i][J-1][O])
                                         0PC1JEJJ(1)=mox (0PC1-1) CJJE1J, PP[1] CJ-1] CJ), if
                                    Staley From DP(NJ[M][W], abbi patta is
                                                                                       ali) or PLi]
                                     New Completed
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* The Toble has during form 1×19 x2 which is for each entry DP[i] CiJ(k) since a constant # or operations is performed the time complexity is O(nm).

!= 'b')

2. Pesign on algorithm Hat comprk the number

1 with D = Max

a \in L_1, B \in L_2

· The a high lovel way at explainly how to conceptualize on algorithm.

psendocade and be

NI + Construct NFA(L1) NZ + Construct NFA(L2)

PECENSharetproduct Automotion (N1 ,N2)

If hos cycle(P)
rotern "D is infinite!

Déo

for each Blake pair (31,52) in p:

If Is Accepting (\$1,N3) and Is Accepting (\$2,N2):

Length & Find Largest Politic (pls3,52))

D & Mar (D, length)

return D.

constructly on NFA form a regular expression has a complexity of O(r) for a regular expression of Size'r. The automotohor P can be QUAM) when I and more values NI and NZ,

The has cycle deck is O(v +E) for the franches V and edged E. Lossly computing he largest path

can be O(v+E) also.

3. Is there my locality sensitive hosts
Scheme for strings?

Both MN-Hasting and Sim Hosh confe Considered locally sensitive tosting Scheres because they tend to ensure that similar items bosh to the some bucket more often than dissimilar illens, one on, fes thee are locally sensitive tosting LSH Schere for strys.