

# Dynamic Binary Instrumentation

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# Principle of operation

- Library overloading
- Modify assembly code to insert instrumentation
- Offline: rewrite the executable file
- Online: self modifying code
  - Paused: i.e. ptrace()
  - Non-stop (kernel)
    - Code update while executing!
    - tracepoints (static\_key)
    - FTrace

# Library overloading

- Compile a library with the symbols to overload
- Preload the library with LD\_PRELOAD
- The overloaded function will be called instead of the original library function. You may call the real library function.
- Does not work with static executables

# Exemple: overloading malloc()

Function pointer for the real malloc

```
void *(*malloc_real)(size_t size);
```

```
void *malloc(size_t size)  
{
```

```
    void *ptr;  
    if (malloc_real == NULL) {  
        malloc_real = dlsym(RTLD_NEXT, "malloc");  
    }  
    ptr = malloc_real(size);  
    const char *msg = "malloc() called";  
    write(1, msg, strlen(msg));  
    return ptr;  
}
```

Redefine malloc()

Load the real malloc if not already loaded

Custom instrumentation. Beware: you must not call malloc() here, otherwise an infinite recursion will occur! Call malloc\_real() instead.

# Dyninst

- API to inspect and modify assembly code
- Online with ptrace() or offline mode
- Can inject generated code
- Web site: <https://github.com/dyninst/dyninst>
- Example: dyninst-demo

# Linux tracepoint

- When not enabled: nop
- Enabling changes the code to a jump
- Based on gcc asm goto
  - Requires a constant key address
  - False branch code is not optimized out
- Meta-info about instrumented sites in section `__jump_table`
- See example 31-asm-goto

# Linux FTrace

- Compilation with -pg (man gcc)
- All function prologue contains call to mcount
  - The space is 5 bytes
- Call to mcount is patched to nop
- The nop is patched to jump to instrumentation

# Runtime patching on x86

- Processor must see valid instructions at all time!
- Patch the first instruction to 0xCC
  - Call to Interrupt Procedure
  - Interrupt handler jumps to the next instruction
- Synchronize cores
- Change the last 4 bytes
- Synchronize cores
- Change the first byte
- Synchronize cores
- New code is effective!



# Useful resources

- <http://ref.x86asm.net/coder64.html>
- <http://www.felixcloutier.com/x86/>
- Tools: objdump, nm, readelf
- See linux/arch/x86/kernel/jump\_label.c