

Weber Problem

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Introduction to the Weber Problem

Weber Problem: Find the "minimum" point (x^*, y^*) which minimizes the sum of weighted distances from itself to n fixed points with co-ordinate (a_i, b_i) , where each point i has a weight w_i

Formulation of Weber Problem

$$\min_{x,y} \{ W(x,y) = \sum_{i=1}^n w_i d_i(x,y) \}$$

$$d_i(x,y) = \sqrt{(x - a_i)^2 + (y - b_i)^2}$$

Solving the Weber Problem Geometrically

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- The Varignon Frame.
- This works because it is essentially the dual of the linear programming formulation.
- The vector (u_i, v_i) is the negative of the force vector exerted on the knot by weight w_i .

Dual of the Weber problem

$$\max_{(U,V)} D(U, V) = - \sum_{i=1}^n (a_i u_i + b_i v_i)$$

s.t.

$$\sum_{i=1}^n u_i = 0$$

$$\sum_{i=1}^n v_i = 0$$

$$\sqrt{u_i^2 + v_i^2} \leq w_i$$

(1)

- Take derivative of the objective function and set to 0
- Start with a candidate solution
- Determine next iteration
- repeat until convergence
- Note: can have problems if a candidate location is one of the fixed point. (Because the algorithm will only approach the fixed point.)

$$\frac{dW(x, y)}{dx} = \sum_{i=1}^n \frac{w_i(x - a_i)}{d_i(x, y)} = 0$$

$$dW(x, y) \frac{dy = \sum_{i=1}^n \frac{w_i(y - b_i)}{d_i(x, y)} = 0$$

$$x^{(k+1)}, y^{(k+1)} = \left(\frac{\sum_{i=1}^n \frac{w_i a_i}{d_i(x^{(k)}, y^{(k)})}}{\sum_{i=1}^n \frac{w_i}{d_i(x^{(k)}, y^{(k)})}}, \frac{\sum_{i=1}^n \frac{w_i b_i}{d_i(x^{(k)}, y^{(k)})}}{\sum_{i=1}^n \frac{w_i}{d_i(x^{(k)}, y^{(k)})}} \right) \quad (2)$$

- Minimizing sum of squared euclidean distances
- This is the centroid

$$\min_{x,y} C(x,y) = \sum_{i=1}^n w_i d_i^2(x,y)$$

Manhattan distance - only allow travel in a grid e.g. Roads in a grid layout

$$\min_{x,y} R(x,y) = \sum_{i=1}^n w_i d_i^R(x,y)$$

$$d_i^R(x,y) = |x - a_i| + |y - b_i|$$

- p -norm are generalizations of Euclidean distances

$$l_{pi} = \sqrt[p]{|x - a_i|^p + |y - b_i|^p} \quad (3)$$

- Where m new facilities will be placed and one new facility y
- Weights between facility j and demand point i is w_{ij}
- Weights between facility j and demand point s is v_{js}

$$\min_{(x,y)_{j=1,\dots(),m}} \sum_{i=1}^n \sum_{j=1}^m w_{ij} \sqrt{(x_j - a_i)^2 + (y_j - b_i)^2} +$$

$$\sum_{j=1}^{m-1} \sum_{s=j+1}^m v_{js} \sqrt{(x_j - a_s)^2 + (y_j - b_s)^2}$$

- Forbidden regions - Cannot locate in F
- $\min_{x,y} F(x,y) = \sum_{i=1}^n w_i d_i(x,y)$
- Barriers - Path cannot cross B
- $\min_{x,y} B(x,y) = \sum_{i=1}^n w_i d_i^B(x,y)$

- Locate a line l which is as close as possible to the set of points
- $\min_{l \in \mathcal{L}} L(l) = \sum_{i=1}^n w_i \gamma_i(l)$
- $\gamma_i(l)$ is the closest distance to point i to line l