

Design and Analysis of Algorithms

Matrix Multiplication

- one condition to multiply two matrices (the #col of first matrix is equal to #row of second matrix)

$$\begin{array}{cc} 1 & 2 \\ 3 & 4 \\ 5 & 6 \end{array} \begin{array}{c} * \\ \\ \end{array} \begin{array}{ccc} 1 & 2 & 3 \\ 4 & 5 & 6 \end{array} = \begin{array}{ccc} 9 & 12 & 15 \\ 19 & 26 & 33 \\ 29 & 40 & 51 \end{array}$$

$3 \times 2 \qquad 2 \times 3 \qquad 3 \times 3$

- time = $3 \times 3 \times 2$ two mul operation for 3×3 elements
- to produce 9 we do two mul operation $1 \times 1 + 2 \times 4$ the summation don't take many time so ignore it.
- two mul operation for one value and we need to produce 3×3 values so we need $2 \times 3 \times 3$ operation mul

Matrix Chain Multiplication

- if we need to multiply three matrices $A_1(3 \times 100) * A_2(100 \times 5) * A_3(5 \times 5)$
- we can $(A_1 * A_2) * A_3$ multiply $A_1 * A_2$ and the result mul A_3
 - $A_1 * A_2 = (3 \times 5)$ time = $3 * 5 * 100 = 1500$
 - $(3 \times 5) * A_3(5 \times 5) = (3 \times 5)$ time = $3 * 5 * 5 = 75$
 - Total time = $1500 + 75 = 1575$
- or we can $A_1 * (A_2 * A_3)$ multiply $A_2 * A_3$ and the A_1 mul result
 - $A_2 * A_3 = (100 \times 5)$ time = $100 * 5 * 5 = 2500$
 - $A_1(3 \times 100) * (100 \times 5) = (3 \times 5)$ time = $3 * 5 * 100 = 1500$
 - Total time = $2500 + 1500 = 4000$
- we cannot mul $A_1 * A_3$ first because i don't know if the condition is true or not
- so i have two way to mul three matrices (matrix chain multiplication is choose the way that give minimum cost). Here the first way is minimum cost.

Matrix Chain Multiplication

- if we have five matrices and i need to mul with minimum cost
 - $A_1 \quad * \quad A_2 \quad * \quad A_3 \quad * \quad A_4 \quad * \quad A_5$
 - $4*10 \quad \quad \quad 10*3 \quad \quad \quad 3*12 \quad \quad \quad 12*20 \quad \quad \quad 20*7$
- the way is said to divide it $(A_1*A_2) * (A_3*A_4*A_5)$ or $(A_1*A_2*A_3)*(A_4*A_5)$
- i will work on $(A_1*A_2) * (A_3*A_4*A_5)$
 - A_1*A_2 will produce matrix $4*3$ res1
 - $A_3*A_4*A_5$ will produce matrix $3*7$ res2
 - so the time to mul res1*res2 is $4*7 * 3$ $4*7$ element each with 3 mul operation
 - so the total time to mul the five matrices is
 - $M[1,5] = \begin{matrix} M[1,2] & + & M[3,5] & + & 4*3*7 \\ \text{time to mul } A_1*A_2 & & \text{time to mul } A_3*A_4*A_5 & & \text{time to mul res1*res2} \\ \text{res } 4*3 & & \text{res } 3*7 & & \end{matrix}$
- $M[i,j] = M[i,k] + M[k+1,j] + (P_i * P_k * P_j)$ $M[i,i] = 0$ no mul found only one matrix

Matrix Chain Multiplication

- now let know how the algorithm work(DP work by store the result in tubule).
- we want to start with $i=j$ then $i<j$ starting with a spread of 1 working our way up
- $M[i,j] = \{0$ if $i=j$ not need to mul the metrix with it self
$$\min_{i \leq k < j} \{M[i,j] = M[i,k] + M[k+1,j] + (P_i * P_{k+1} * P_{j+1})\}$$
$$\}$$
- i is the rows and j is the columns (this table will contain at the end with minimum cost)

i/j	1	2	3	4	5
1					
2	X				
3	X	X			
4	X	X	X		
5	X	X	X	x	

A2 can't multiply with A1 so Ai cannot mul Aj if $i > j$ so put x in the table

Matrix Chain Multiplication

- بملي الجدول قطر قطر
- step 1: fill the table for $i=j$
- $M[i,j] = \{0$ if $i=j$ not need to mul the metrix with it self
 $\min_{i \leq k < j} \{M[i,j] = M[i,k] + M[k+1,j] + (P_i * P_{k+1} * P_{j+1})\}$
 $\}$

i/j	1	2	3	4	5
1	0				
2	X	0			
3	X	X	0		
4	X	X	X	0	
5	X	X	X	X	0

Matrix Chain Multiplication

- step 2: fill the table for
- $A_1 \quad * \quad A_2 \quad * \quad A_3 \quad * \quad A_4 \quad * \quad A_5$
 $4*10 \quad 10*3 \quad 3*12 \quad 12*20 \quad 20*7$
 $p_1 \ p_2 \quad p_2 \ p_3 \quad p_3 \ p_4 \quad p_4 \ p_5 \quad p_5 \ p_6$

- $M[i,j] = \{0$ if $i=j$ not need to mul the metrix with it self
 $\min_{i \leq k < j} \{M[i,j] = M[i,k] + M[k+1,j] + (P_i * P_{k+1} * P_{j+1})\}$
 $\}$

- $i=1 \quad j=2$
- $M[1,2] = \min(1 \leq k < 2) \{ M[1,1] + M[1+1,2] + (P_1 * P_2 * P_3) \}$
- $= 0 + 0 + 4 * 10 * 3 = 120$

i/j	1	2	3	4	5
1	0	120			
2	X	0			
3	X	X	0		
4	X	X	X	0	
5	X	X	X	X	0

Matrix Chain Multiplication

- step 2: fill the table for

$$\begin{array}{ccccccccc}
 A_1 & * & A_2 & * & A_3 & * & A_4 & * & A_5 \\
 4*10 & & 10*3 & & 3*12 & & 12*20 & & 20*7 \\
 p_1 \ p_2 & & p_2 \ p_3 & & p_3 \ p_4 & & p_4 \ p_5 & & p_5 \ p_6
 \end{array}$$

- $M[i,j] = \{0$ if $i=j$ not need to mul the metrix with it self
 $\min_{i \leq k < j} \{M[i,j] = M[i,k] + M[k+1,j] + (P_i * P_{k+1} * P_{j+1})\}$
 $\}$

- $i=2 \quad j=3$

- $M[2,3] = \min(2 \leq k < 3) \{ M[2,2] + M[2+1,3] + (P_2 * P_3 * P_4) \}$

- $= 0 + 0 + 10 * 3 * 12 = 360$

i/j	1	2	3	4	5
1	0	120			
2	X	0	360		
3	X	X	0		
4	X	X	X	0	
5	X	X	X	X	0

Matrix Chain Multiplication

- step 2: fill the table for

$$\begin{array}{ccccccccc}
 A_1 & * & A_2 & * & A_3 & * & A_4 & * & A_5 \\
 4*10 & & 10*3 & & 3*12 & & 12*20 & & 20*7 \\
 p_1 \ p_2 & & p_2 \ p_3 & & p_3 \ p_4 & & p_4 \ p_5 & & p_5 \ p_6
 \end{array}$$

- $M[i,j] = \{0$ if $i=j$ not need to mul the metrix with it self
 $\min_{i \leq k < j} \{M[i,j] = M[i,k] + M[k+1,j] + (P_i * P_{k+1} * P_{j+1})\}$
 $\}$

- $i=3 \quad j=4$

- $M[3,4] = \min(3 \leq k < 4) \{ M[3,3] + M[3+1,4] + (P_3 * P_4 * P_5) \}$

- $= 0 + 0 + 3 * 12 * 20 = 720$

i/j	1	2	3	4	5
1	0	120			
2	X	0	360		
3	X	X	0	720	
4	X	X	X	0	
5	X	X	X	X	0

Matrix Chain Multiplication

- step 2: fill the table for

- $A_1 \quad * \quad A_2 \quad * \quad A_3 \quad * \quad A_4 \quad * \quad A_5$
 $4*10 \quad 10*3 \quad 3*12 \quad 12*20 \quad 20*7$
 $p_1 \ p_2 \quad p_2 \ p_3 \quad p_3 \ p_4 \quad p_4 \ p_5 \quad p_5 \ p_6$

- $M[i,j] = \{0 \quad \text{if } i=j \quad \text{not need to mul the matrix with it self}$
 $\min_{i \leq k < j} \{M[i,j] = M[i,k] + M[k+1,j] + (P_i * P_{k+1} * P_{j+1})\}$
 $\}$

- $i=4 \quad j=5$

- $M[4,5] = \min(4 \leq k < 5) \{ M[4,4] + M[4+1,5] + (P_4 * P_5 * P_6) \}$
 $= 0 + 0 + 12 * 20 * 7 = 1680$

i/j	1	2	3	4	5
1	0	120			
2	X	0	360		
3	X	X	0	720	
4	X	X	X	0	1680
5	X	X	X	X	0

Matrix Chain Multiplication

- step 3: fill the table for

A1	*	A2	*	A3	*	A4	*	A5
4*10		10*3		3*12		12*20		20*7
p1 p2		p2 p3		p3 p4		p4 p5		p5 p6

- $M[i,j] = \begin{cases} 0 & \text{if } i=j \\ \min_{i \leq k < j} \{M[i,k] + M[k+1,j] + (P_i * P_{k+1} * P_{j+1})\} & \text{not need to mul the matrix with it self} \end{cases}$

- $i=1$ $j=3$

- $M[1,3] = \min(1 \leq k < 3)$ have two solution

- $A1 * (A2 * A3)$ when $k=1$

$$\{M[1,1] + M[1+1,3] + (P_1 * P_2 * P_4)\} = 0 + 360 + 4 * 10 * 12 = 840$$

- or $(A1 * A2) * A3$ when $k=2$

$$\{M[1,2] + M[2+1,3] + (P_1 * P_3 * P_4)\} = 120 + 0 + 4 * 3 * 12 = 264$$

- so $264 < 840$

i/j	1	2	3	4	5
1	0	120	264		
2	X	0	360		
3	X	X	0	720	
4	X	X	X	0	1680
5	X	X	X	X	0

Matrix Chain Multiplication

- step 3: fill the table for

A1	*	A2	*	A3	*	A4	*	A5
4*10		10*3		3*12		12*20		20*7
p1 p2		p2 p3		p3 p4		p4 p5		p5 p6

- $M[i,j] = \begin{cases} 0 & \text{if } i=j \\ \min_{i \leq k < j} \{M[i,k] + M[k+1,j] + (P_i * P_{k+1} * P_{j+1})\} & \text{not need to mul the matrix with it self} \end{cases}$

- $i=2$ $j=4$

- $M[2,4] = \min(2 \leq k < 4)$ have two solution

- $A2 * (A3 * A4)$ when $k=2$

$$\{ M[2,2] + M[2+1,4] + (P_2 * P_3 * P_5) \} = 0 + 720 + 10 * 3 * 20 = 1320$$

- or $(A2 * A3) * A4$ when $k=3$

$$\{ M[2,3] + M[3+1,4] + (P_2 * P_4 * P_5) \} = 360 + 0 + 10 * 12 * 20 = 2760$$

- so $1320 < 2760$

i/j	1	2	3	4	5
1	0	120	264		
2	X	0	360	1320	
3	X	X	0	720	
4	X	X	X	0	1680
5	X	X	X	X	0

Matrix Chain Multiplication

- step 3: fill the table for

A1	*	A2	*	A3	*	A4	*	A5
4*10		10*3		3*12		12*20		20*7
p1 p2		p2 p3		p3 p4		p4 p5		p5 p6

- $M[i,j] = \begin{cases} 0 & \text{if } i=j \\ \min_{i \leq k < j} \{M[i,k] + M[k+1,j] + (P_i * P_{k+1} * P_{j+1})\} & \text{not need to mul the metrix with it self} \end{cases}$

- i=3 j=5

- $M[3,5] = \min(3 \leq k < 5)$ have two solution

- $A3 * (A4 * A5)$ when $k=3$

- $\{ M[3,3] + M[3+1,5] + (P_2 * P_3 * P_5) \} = 0 + 1680 + 3 * 12 * 7 = 1932$

- or $(A3 * A4) * A5$ when $k=4$

- $\{ M[3,4] + M[4+1,5] + (P_3 * P_5 * P_6) \} = 720 + 0 + 3 * 20 * 7 = 1140$

- so $1140 < 1932$

i/j	1	2	3	4	5
1	0	120	264		
2	X	0	360	1320	
3	X	X	0	720	1140
4	X	X	X	0	1680
5	X	X	X	X	0

Matrix Chain Multiplication

- step 4: fill the table for
- | | | | | | | | | |
|-------|---|-------|---|-------|---|-------|---|-------|
| A1 | * | A2 | * | A3 | * | A4 | * | A5 |
| 4*10 | | 10*3 | | 3*12 | | 12*20 | | 20*7 |
| p1 p2 | | p2 p3 | | p3 p4 | | p4 p5 | | p5 p6 |

- $M[i,j] = \begin{cases} 0 & \text{if } i=j \text{ not need to mul the metrix with it self} \\ \min_{i \leq k < j} \{M[i,k] + M[k+1,j] + (P_i * P_{k+1} * P_{j+1})\} & \end{cases}$

- $i=1 \quad j=4$

- $M[1,4] = \min(1 \leq k < 4)$ have three solution

- $A1*(A2*A3*A4)$ when $k=1$
- $\{ M[1,1] + M[1+1,4] + (P1 * P2 * P5) \} = 0+1320+4*10*20 = 2120$
- or $(A1*A2)*(A3*A4)$ when $k=2$
- $\{ M[1,2] + M[2+1,4] + (P1 * P3 * P5) \} = 120+720+4*3*20 = 1080$
- or $(A1*A2*A3)*A4$ when $k=3$
- $\{ M[1,3] + M[3+1,4] + (P1 * P4 * P5) \} = 264+0+4*12*20 = 1224$
- so $1080 < 1224 < 2120$

i/j	1	2	3	4	5
1	0	120	264	1080	
2	X	0	360	1320	
3	X	X	0	720	1140
4	X	X	X	0	1680
5	X	X	X	X	0

Matrix Chain Multiplication

- step 4: fill the table for
- | | | | | | | | | |
|-------|---|-------|---|-------|---|-------|---|-------|
| A1 | * | A2 | * | A3 | * | A4 | * | A5 |
| 4*10 | | 10*3 | | 3*12 | | 12*20 | | 20*7 |
| p1 p2 | | p2 p3 | | p3 p4 | | p4 p5 | | p5 p6 |

- $M[i,j] = \begin{cases} 0 & \text{if } i=j \\ \min_{i \leq k < j} \{M[i,k] + M[k+1,j] + (P_i * P_{k+1} * P_{j+1})\} & \text{not need to mul the metrix with it self} \end{cases}$

- $i=2$ $j=5$
- $M[2,5] = \min(2 \leq k < 5)$ have three solution

- $A2*(A3*A4*A5)$ when $k=2$
- $\{ M[2,2] + M[2+1,5] + (P2 * P3 * P6) \}$ $= 0+1140+10*3*7 = 1350$
- or $(A2*A3)*(A4*A5)$ when $k=3$
- $\{ M[2,3] + M[3+1,5] + (P2 * P4 * P6) \}$ $= 360+1680+10*12*7 = 2880$
- or $(A2*A3*A4)*A5$ when $k=4$
- $\{ M[2,4] + M[4+1,5] + (P2 * P5 * P6) \}$ $= 1320+0+10*20*7 = 2720$
- so $1350 < 2720 < 2880$

i/j	1	2	3	4	5
1	0	120	264	1080	
2	X	0	360	1320	1350
3	X	X	0	720	1140
4	X	X	X	0	1680
5	X	X	X	X	0

Matrix Chain Multiplication

- step 5: fill the table for

A1	*	A2	*	A3	*	A4	*	A5
4*10		10*3		3*12		12*20		20*7
p1 p2		p2 p3		p3 p4		p4 p5		p5 p6

- $M[i,j] = \begin{cases} 0 & \text{if } i=j \\ \min_{i \leq k < j} \{M[i,k] + M[k+1,j] + (P_i * P_{k+1} * P_{j+1})\} & \text{not need to mul the metrix with it self} \end{cases}$

- $i=1$ $j=5$

- $M[1,5] = \min(1 \leq k < 5)$ have four solution

- $A1*(A2*A3*A4*A5)$ when $k=1$

- $\{M[1,1] + M[1+1,5] + (P1 * P2 * P6)\} = 0+1350+4*10*7 = 1630$

- or $(A1*A2)*(A3*A4*A5)$ when $k=2$

- $\{M[1,2] + M[2+1,5] + (P1 * P3 * P6)\} = 120+1140+4*3*7 = 1344$

- or $(A1*A2*A3)*(A4*A5)$ when $k=3$

- $\{M[1,3] + M[3+1,5] + (P1 * P4 * P6)\} = 264+1680+4*12*7 = 2280$

- or $(A1*A2*A3*A4)*A5$ when $k=4$

- $\{M[1,4] + M[4+1,5] + (P1 * P5 * P6)\} = 1080+0+4*20*7 = 1640$

- so $1344 < 1630 < 1640 < 2280$

i/j	1	2	3	4	5
1	0	120	264	1080	1344
2	X	0	360	1320	1350
3	X	X	0	720	1140
4	X	X	X	0	1680
5	X	X	X	X	0

Matrix Chain Multiplication

- $\text{result} = 1344 = (A1 * A2) * (A3 * A4 * A5)$
- $(A3 * A4 * A5) = 1140 = (A3 * A4) * A5$
- $\text{final result} = (A1 * A2) * ((A3 * A4) * A5)$

i/j	1	2	3	4	5
1	0	120	264	1080	1344
2	X	0	360	1320	1350
3	X	X	0	720	1140
4	X	X	X	0	1680
5	X	X	X	X	0