#### Lecture 02b:

#### What is Entropy?



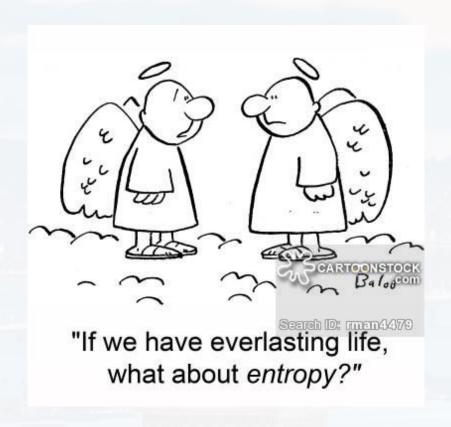
Markus Hohle
University California, Berkeley

Machine Learning Algorithms

MSSE 277B, 3 Units

Spring 2025

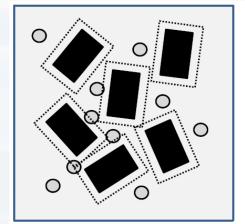




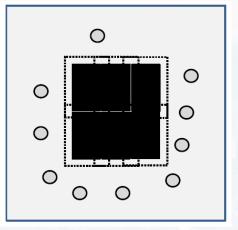




What do you think: in which image is entropy higher?



low entropy



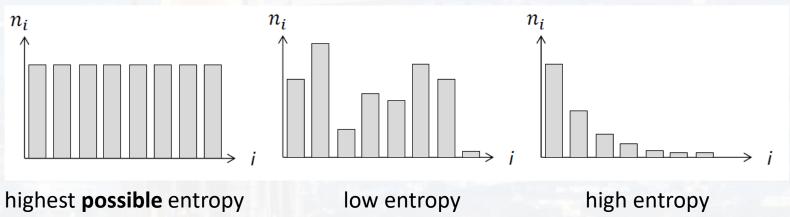
high entropy

#### first of all:

Entropy is **not** a measure of disorder!

: states

 $n_i$ : number of particles in state i





Often people explain entropy with an ordered vs messy office...

first of all:

Entropy is **not** a measure of disorder!



...and then say, that entropy (disorder) grows with time (in closed systems).

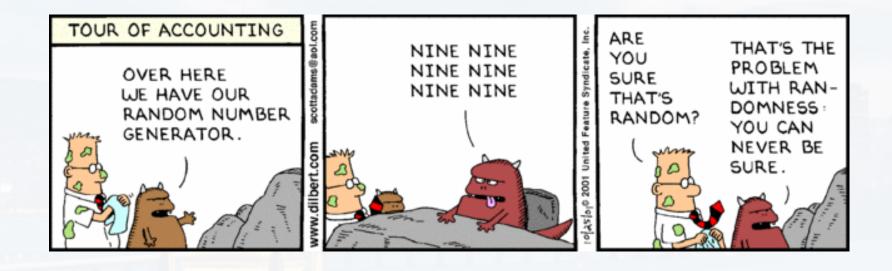
- But how is it possible, that an office can do that, just by itself?
- What if my office just looks messy,
   but I can still pull any file you are asking me for?

order/disorder is not a physical quantity!

Those examples have nothing to do with entropy conceptionally!



actually, the idea of entropy is more like that:



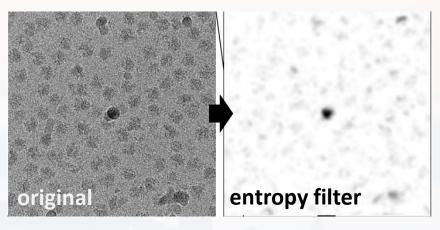


**Entropy:** 

data analysis:

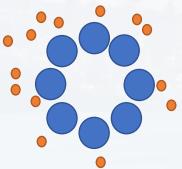
- image processing
- noise reduction
- feature detection

#### **Cryo-EM image of ribosomes**



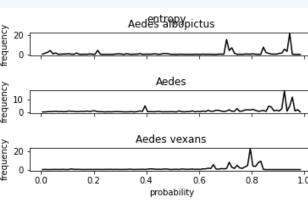
biophysics:

- molecular driving forces
- formation of macromolecules
- "ordering forces"



AI:

- optimization
- cross entropy





**Entropy:** 

statistics/information theory:

- maximum entropy, given constrains

Distribution name	Probability density / mass function	Maximum Entropy constraint	Support
Uniform (discrete)	$f(k) = \frac{1}{b-a+1}$	None	$\{a,a+1,\ldots,b-1,b\}$
Uniform (continuous)	$f(x) = \frac{1}{b-a}$	None	[a,b]
Bernoulli	$f(k) = p^k (1-p)^{1-k}$	$\mathbb{E}[\ K\ ]=p$	{0,1}
Geometric	$f(k)=(1-p)^{k-1}\;p$	$\mathbb{E}[K]=rac{1}{p}$	$\mathbb{N} \smallsetminus \{0\} = \{1,2,3,\dots\}$
Exponential	$f(x) = \lambda \exp(-\lambda x)$	$\mathbb{E}[X]=rac{1}{\lambda}$	$[0,\infty)$
Laplace	$f(x) = rac{1}{2b} \expigg(-rac{ x-\mu }{b}igg)$	$\mathbb{E}[\  X-\mu \ ]=b$	$(-\infty,\infty)$
Asymmetric Laplace	$f(x) = rac{\lambda \; \expig(-\left(x-m ight) \lambda \; s \; \kappa^sig)}{\left(\kappa + rac{1}{\kappa} ight)}$ where $s \equiv \mathrm{sgn}(x-m)$	$\mathbb{E}[\;(X-m)\;s\;\kappa^s\;]=rac{1}{\lambda}$	$(-\infty,\infty)$
Pareto	$f(x)=rac{lpha\ x_m^lpha}{x^{lpha+1}}$	$\mathbb{E}[\; \ln X ] = rac{1}{lpha} + \ln(x_m)$	$[x_m,\infty)$
Normal	$f(x) = rac{1}{\sqrt{2\pi\sigma^2}} \exp\Biggl(-rac{(x-\mu)^2}{2\sigma^2}\Biggr)$	$egin{aligned} \mathbb{E}[\ X\ ] &= \mu\ , \ \mathbb{E}[\ X^2\ ] = \sigma^2 + \mu^2 \end{aligned}$	$(-\infty,\infty)$
	( - )		1000
Gamma	$f(x) = rac{x^{k-1} \exp\left(-rac{x}{ heta} ight)}{ heta^k \; \Gamma(k)}$	$egin{aligned} \mathbb{E}[\;X\;] &= k\; heta\;, \ \mathbb{E}[\;\ln X\;] &= \psi(k) + \ln  heta \end{aligned}$	$[0,\infty)$



What is entropy, really?



N: number of dice

 $n_i$ : number of dice exposing a certain number i

(= having a certain state i)

I: number of states a die can have

What is the probability *P* to observe the *system* in a certain state?

What is the probability  $p_i$  to observe a die in a certain state?

$$\Omega = \frac{N!}{n_1! \, n_2! \dots n_I!}$$

assumption: all *i* are equally likely:

$$P = 1/\Omega$$





$$\Omega = \frac{N!}{n_1! \, n_2! \dots n_I!}$$

N: number of dice

 $n_i$ : number of dice exposing a certain number i

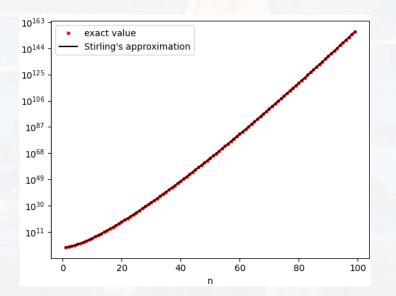
(= having a certain state i)

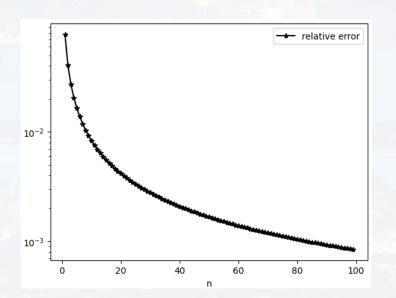
I: number of states a die can have

is large, even for small systems!

Stirling's approximation

$$n! \approx \sqrt{2\pi n} \left(\frac{n}{e}\right)^n$$









$$\Omega = \frac{N!}{n_1! \, n_2! \dots n_I!}$$

N: number of dice

number of dice exposing a certain number i  $n_i$ :

(= having a certain state i)

I: number of states a die can have

is large, even for small systems!

#### for large $n_i$ :

$$\Omega = \frac{N!}{n_1! \, n_2! \, ... \, n_I!} \approx \frac{N^N}{n_1^{n_1} n_2^{n_2} \, ... \, n_I^{n_I}}$$
$$\approx \frac{1}{p_1^{n_1} p_2^{n_2} \, ... \, p_I^{n_I}}$$

$$p_i pprox rac{n_i}{N}$$

$$ln\Omega = -\sum_{i}^{I} n_{i} \, lnp_{i} \qquad \frac{\ln\Omega}{N} = -\sum_{i}^{I} p_{i} \, lnp_{i} \qquad S = -\sum_{i}^{I} p_{i} \, lnp_{i}$$

$$S = -\sum_{i}^{I} p_{i} \ln p_{i}$$

entropy per particle





$$S = -\sum_{i}^{I} p_{i} \, ln p_{i}$$

N: number of dice

 $n_i$ : number of dice exposing a certain number i

(= having a certain state i)

I: number of states a die can have

 $p_i$ :  $n_i/N$ 

assumption: all *i* are equally likely:  $P = 1/\Omega$ 

subsets of  $\Omega$ : - sum M of all numbers on the dice

- dice can only be distinguished by their state

N = 2: min(M) = 1 + 1 = 2

max(M) = 6 + 6 = 12

most likely M = 7 (or  $2 \times mean(I)$ ), because there are **six** possibilities to obtain it:

$$1 + 6$$
;  $1 + 6$ ;  $2 + 5$ ;  $5 + 2$ ;  $3 + 4$ ;  $4 + 3$ 

N: min(M) = N

max(M) = I M

most likely  $M = N \times mean(I)$ 





$$S = -\sum_{i}^{I} p_{i} \, ln p_{i}$$

N: number of dice

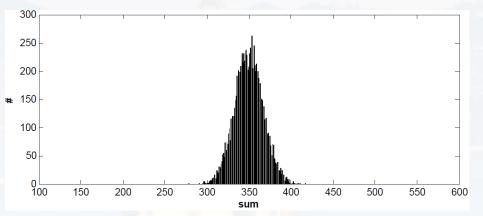
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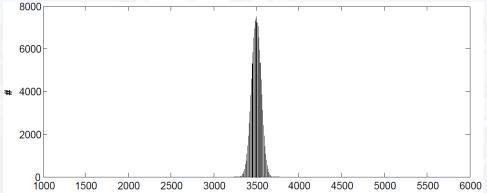
 $p_i$ :  $n_i/N$ 

assumption: all *i* are equally likely:



N = 100

- some subsets of  $\Omega$ , hence some states of the system are way more likely than other states
- becomes more extreme for large N



N = 1000





$$S = -\sum_{i}^{I} p_{i} \, ln p_{i}$$

N: number of dice

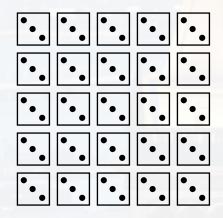
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I: number of states a die can have

 $p_i$ :  $n_i/N$ 

#### we can also see this as dynamical process:



t = 0 : all dice have the same state





$$S = -\sum_{i}^{I} p_{i} \, ln p_{i}$$

N: number of dice

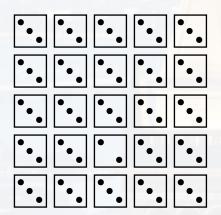
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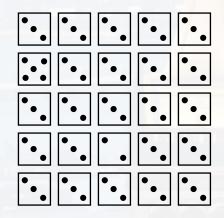
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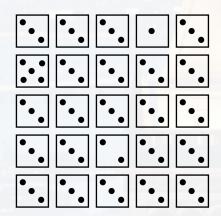
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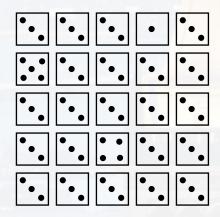
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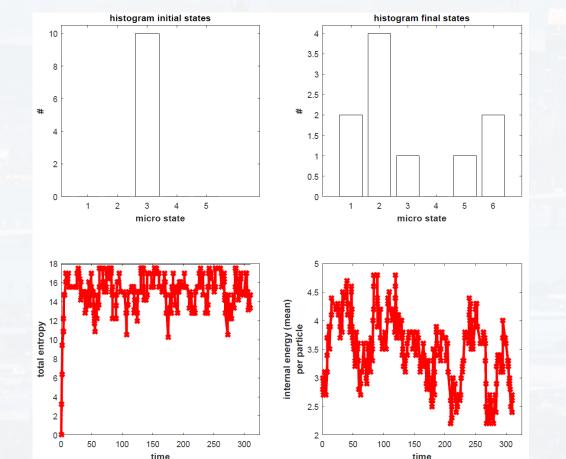
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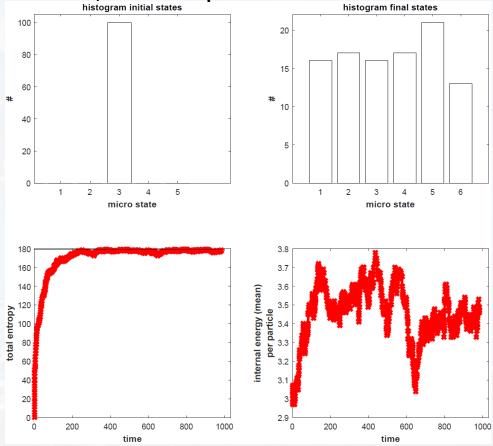
I: number of states a die can have

 $p_i$ :  $n_i/N$ 

10 dice 300 timesteps













$$S = -\sum_{i}^{I} p_{i} \, ln p_{i}$$

N: number of dice

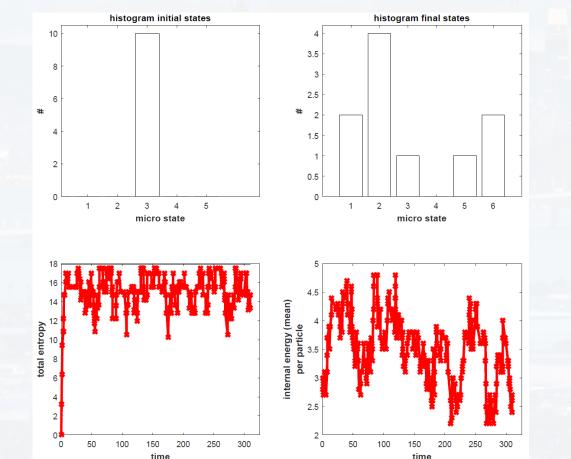
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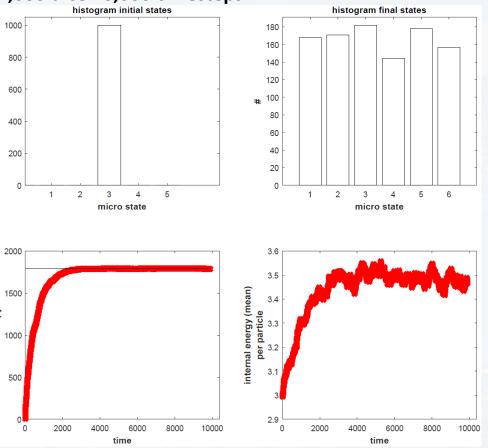
I: number of states a die can have

 $p_i$ :  $n_i/N$ 

#### 10 dice 300 timesteps

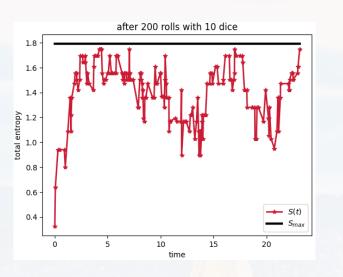


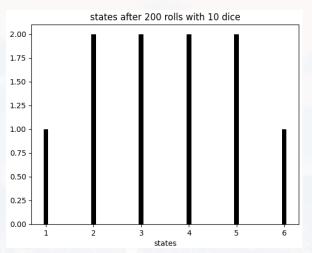
#### 1,000 dice 10,000 timesteps



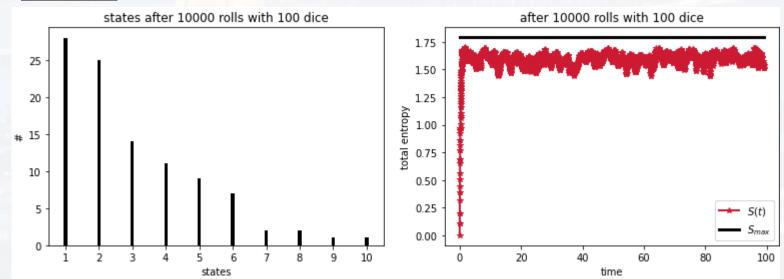


#### check out random\_machine.ipynb





#### constrain: M is conserved, I is free, but >0





#### conclusions:

- entropy increases with time (in a closed system) because it is the **most likely** state
- the larger the system, the more deterministic it looks
- for small systems: entropy can fluctuate in both ways and does not increase!
   (Stirling's approximation)
- large systems: (thermodynamic) arrow of time (question: what if we are at  $S_{max}$  already)
- small systems: symmetry in time!
- even if i are not equally likely (constrains, some states are more accessible)
   → different weights, but same principle
- uniform distribution has highest entropy (do the math : ) )



#### Entropy is a mathematically precise measure of information!

- entropy high → information low
- entropy low → information high

	I		
S = -	$-\sum_{i}$	$p_i$	$lnp_i$

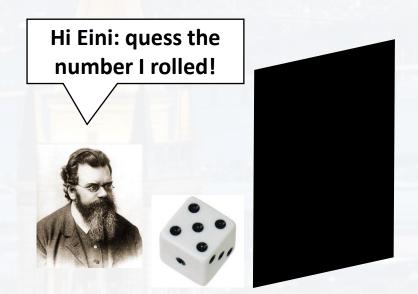


fair die: events are

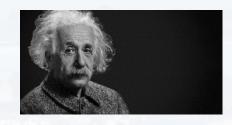
 $p_n$  are

1, 2, 3, 4, 5, 6

 $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$ ,  $p_5$ ,  $p_6$  = 1/6 for all  $p_n$ 



I have no idea, so all numbers are equally likely if it's a fair die.





#### Entropy is a mathematically precise measure of information!

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$$S = -\sum_{i}^{I} p_{i} \, ln p_{i}$$



fair die: events are

 $p_n$  are

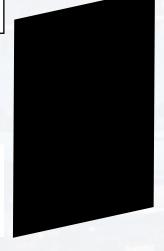
1, 2, 3, 4, 5, 6

 $p_1, p_2, p_3, p_4, p_5, p_6 = 1/6$  for all  $p_n$ 









$$S = -\frac{1}{6}ln\left(\frac{1}{6}\right) * 6 = 1.79 \dots$$





#### Entropy is a mathematically precise measure of information!

- entropy high → information low
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I		
$S = -\sum_{i} f_{i}$	$p_i$	$lnp_i$

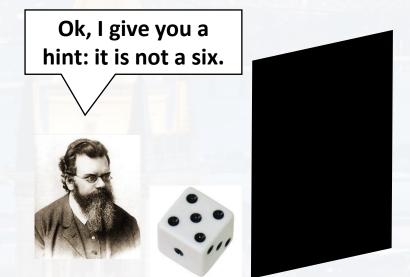


fair die: events are

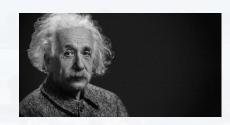
 $p_n$  are

1, 2, 3, 4, 5, 6

 $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$ ,  $p_5$ ,  $p_6$  = 1/6 for all  $p_n$ 



Alright, so then  $p_6=0$  and all the other  $p_5=1/5$ .





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	I		
S = -	$-\sum_{i}$	$p_i$	$lnp_i$

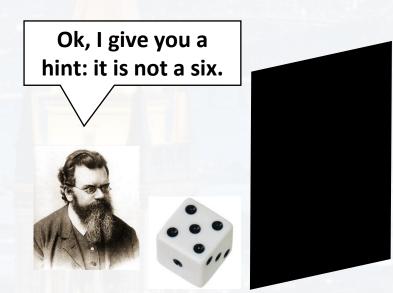


fair die: events are

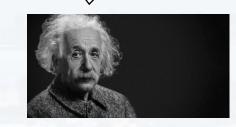
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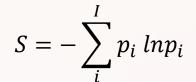
Hence, 
$$S = -\frac{1}{5} ln(\frac{1}{5}) * 5 = 1.61 ...$$





#### Entropy is a mathematically precise measure of information!

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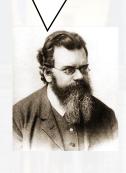
fair die: events are

 $p_n$  are

1, 2, 3, 4, 5, 6

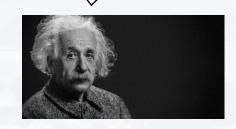
 $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$ ,  $p_5$ ,  $p_6$  = 1/6 for all  $p_n$ 

Come on, don't be so nerdy. It is an odd number.





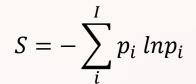
This helps a lot:  $p_2=p_4=p_6=0$  and thus,  $p_1=p_3=p_5=1/3$ 





#### Entropy is a mathematically precise measure of information!

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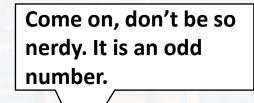


fair die: events are

 $p_n$  are

1, 2, 3, 4, 5, 6

 $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$ ,  $p_5$ ,  $p_6$  = 1/6 for all  $p_n$ 

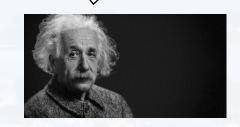








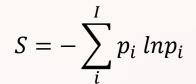
Hence, 
$$S = -\frac{1}{3} ln \left(\frac{1}{3}\right) * 3 = 1.10 ...$$





#### Entropy is a mathematically precise measure of information!

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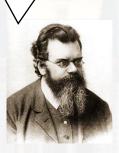
fair die: events are

 $p_n$  are

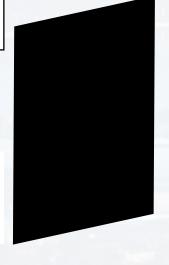
1, 2, 3, 4, 5, 6

 $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$ ,  $p_5$ ,  $p_6$  = 1/6 for all  $p_n$ 

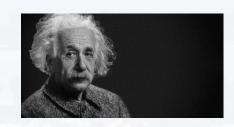
I give up. It's a five. Next time I am gonna play with Schroedinger.







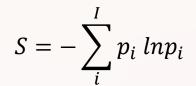
Fantastic! So all  $oldsymbol{p}_n=\mathbf{0}$  except for n = 5





#### Entropy is a mathematically precise measure of information!

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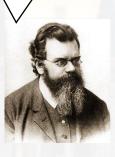
**fair** die: events are

 $p_n$  are

1, 2, 3, 4, 5, 6

 $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$ ,  $p_5$ ,  $p_6$  = 1/6 for all  $p_n$ 

I give up. It's a five. Next time I am gonna play with Schroedinger.







Hence,  $S = -0 \ ln(0) * 5 - 1 * ln(1) = 0$ 

But don't mention your cat!



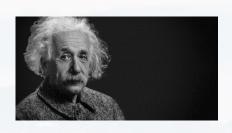
#### Entropy is a mathematically precise measure of information!

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I	
$S = -\sum_{i}$	$p_i lnp_i$







**Information:** 

none (any number)

not a six

an odd number

the actual number

**Entropy:** 

S = 1.79

S = 1.61

S = 1.10

S = 0.00



#### M. Hohle:

### Thank you for your attention!

