

## Dijkstra's Algorithm on Example Graph

**Graph Description:** Nodes:  $A, B, C, D, E, F$  Edges:

- $A \rightarrow B(1)$
- $A \rightarrow C(4)$
- $B \rightarrow D(2)$
- $B \rightarrow E(7)$
- $C \rightarrow D(5)$
- $D \rightarrow F(6)$
- $F \rightarrow E(3)$

**Initialization:**

$$d[A] = 0$$

$$d[B] = \infty$$

$$d[C] = \infty$$

$$d[D] = \infty$$

$$d[E] = \infty$$

$$d[F] = \infty$$

$$Q = \{A, B, C, D, E, F\}$$

### Iterations of the given psuedo code:

the algorithm given doesn't save the shortest path. After the some modification to the algorithm, the shortest path would be

$A \rightarrow B \rightarrow D \rightarrow F$ . The  $Q$  and  $d$  are updated as follows:

1.  $Q = \{B, C, D, E, F\}$  and  $d = \{0, 1, 4, \infty, \infty, \infty\}$
2.  $Q = \{C, D, E, F\}$  and  $d = \{0, 1, 4, 3, 8, \infty\}$
3.  $Q = \{D, E, F\}$  and  $d = \{0, 1, 4, 3, 8, 9\}$
4.  $Q = \{E, F\}$  and  $d = \{0, 1, 4, 3, 8, 9\}$
5.  $Q = \{F\}$  and  $d = \{0, 1, 4, 3, 8, 9\}$
6.  $Q = \{\}$  and  $d = \{0, 1, 4, 3, 8, 9\}$  and the loop will break since the  $Q$  is empty.  $d$  is the shortest distances between the start to all other nodes.

**Shortest Paths and Costs:**

$$A \rightarrow B : 1$$

$$A \rightarrow C : 4$$

$$A \rightarrow D : 3$$

$$A \rightarrow E : 8$$

$$A \rightarrow F : 9$$

#### 4. A\* Algorithm on Example Graph

**Heuristic Values:**

$$h(A) = 10, h(B) = 8, h(C) = 6, h(D) = 4, h(E) = 2, h(F) = 0$$

**Initialization:**

$$d[A] = 0$$

$$d[B] = \infty$$

$$d[C] = \infty$$

$$d[D] = \infty$$

$$d[E] = \infty$$

$$d[F] = \infty$$

$$Q = \{A, B, C, D, E, F\}$$

**Iterations of the given psuedo code for A\* Algorithm:**

1.  $Q = \{B, C, D, E, F\}$  and  $d = \{0, 1, 4, \infty, \infty, \infty\}$
2.  $Q = \{C, D, E, F\}$  and  $d = \{0, 1, 4, 3, 8, \infty\}$
3.  $Q = \{C, E, F\}$  and  $d = \{0, 1, 4, 3, 8, 9\}$
4.  $Q = \{C, E\}$  and  $d = \{0, 1, 4, 3, 8, 9\}$  and the loop will break since the target node has been reached. Evrery item we popped in the loop is the shortest path.

**Shortest Path:**

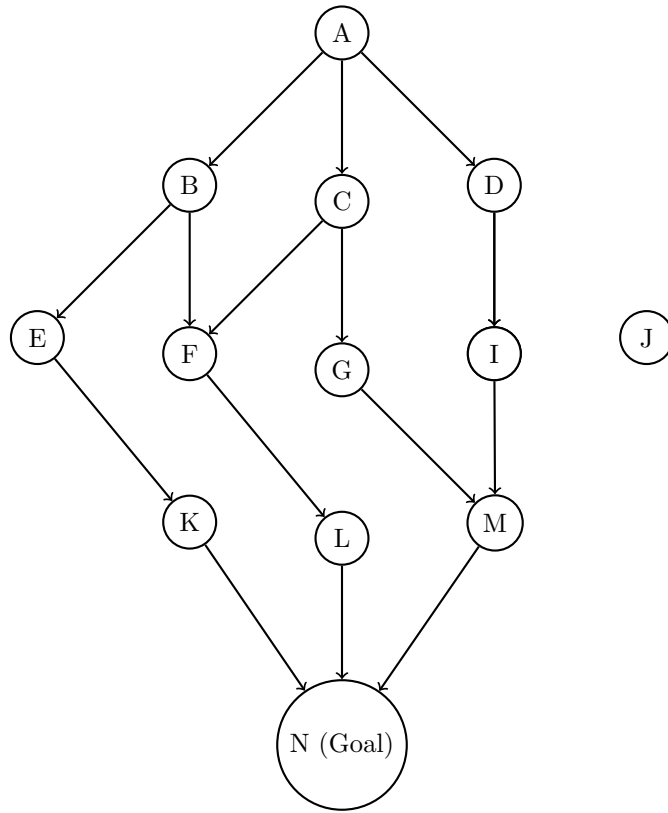
$$A \rightarrow B \rightarrow D \rightarrow F$$

$$\text{Cost: } 9$$

#### 5. Graphs Where A\* Isn't Helpful

**Graph 1: Highly Connected Graphs.**

A\* may explore multiple alternative paths due to the abundance of connections, especially when edge weights are nearly equal. This increases the number of nodes explored unnecessarily.



### Misleading Heuristic:

In highly connected graphs, even a reasonable heuristic can mislead A\* into exploring unnecessary nodes. Below is a dense graph where A\* fails to prioritize the shortest path efficiently.

A heuristic misaligned with the graph's geometry can guide A\* in the wrong direction, causing it to explore irrelevant paths.

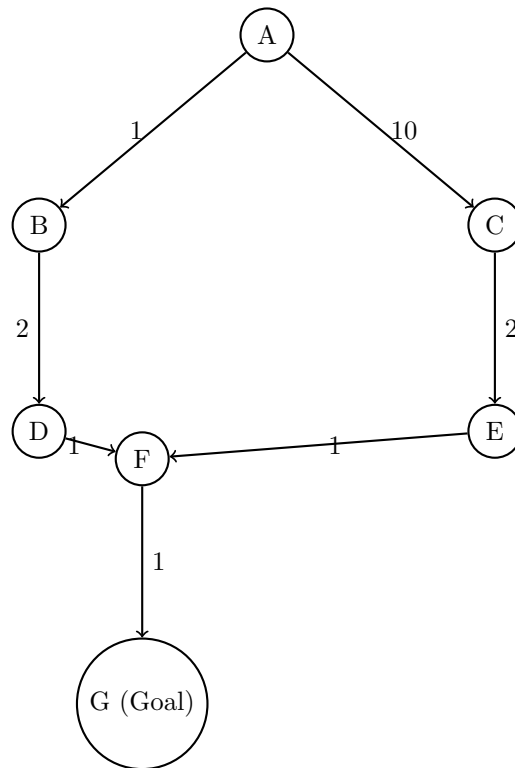
### Heuristic Values

The heuristic for the nodes is:

$$h(A) = 10, h(B) = 8, h(C) = 1, h(D) = 8, h(E) = 1, h(F) = 5, h(G) = 0$$

### Problem

- The heuristic overestimates costs in  $B$  and  $D$ , guiding A\* toward the path  $C \rightarrow E \rightarrow F \rightarrow G$ , even though the shortest path is  $A \rightarrow B \rightarrow D \rightarrow F \rightarrow G$ .
- The algorithm explores  $C$  and  $E$ , wasting time on an unnecessary detour.



**To improve efficiency:**

- Use bidirectional A\* to reduce the number of explored nodes. Start one search from the start and another from the goal, meeting in the middle.
- Ensure the heuristic is admissible (never overestimates the cost) and consistent (satisfies the triangle inequality). In cases like grids, use Manhattan or Euclidean distances to reflect actual geometry.