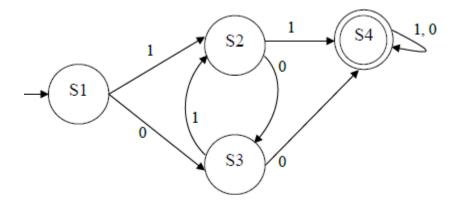
1.

A finite state machine (FSM) can be used to define a language: a string is allowed in a language if it is accepted by the FSM that represents the rules of the language.

Figure 1 shows the state transition diagram for an FSM.

Figure 1



An FSM can be represented as a state transition diagram or as a state transition table. The table below is an incomplete state transition table for **Figure 1**.

(a) Complete the table.

Original state	Input	New state
S3		
S3		

(1)

Any language that can be defined using an FSM can also be defined using a regular expression.
The FSM in Figure 1 defines the language that allows all strings containing at least, either two consecutive 1s or two consecutive 0s.
The strings 0110, 00 and 01011 are all accepted by the FSM and so are valid strings in the language.
The strings 1010 and 01 are not accepted by the FSM and so are not valid strings in the language.
Write a regular expression that is equivalent to the FSM shown in Figure 1 .

(b)

(3)

(c) Backus-Naur Form (BNF) can be used to define the rules of a language.

Figure 2 shows an attempt to write a set of BNF production rules to define a language of full names.

Figure 2 Note: underscores (_) have been used to denote spaces. Note: rule numbers have been included but are not part of the BNF rules. Rule number 1 <fullname> ::= <title>_<name>_<endtitle> | <name> <title>_<name> | <name>_<endtitle> <title> ::= MRS | MS | MISS | MR | DR | SIR 3 <endtitle> ::= ESQUIRE | OBE | CBE <name> ::= <word> | <name>_<word> 5 <word> ::= <char><word> 6 <char> ::= A | B | C | D | E | F | G | H | I | J | K | L | M | N | O | P | Q | R | S | T | U | V | W | X | Y | Z

BNF can be used to define languages that are not possible to define using regular expressions. The language defined in **Figure 2** could not have been defined using regular expressions.

Complete the table below by writing either a **Y** for **Yes** or **N** for **No** in each row.

Rule number (given in Figure 2)	Could be defined using a regular expression
1	
2	
3	
4	
5	
6	

(1)

(d) There is an error in rule 5 in **Figure 2** which means that no names are defined by the language.

Explain what is wrong with the production rule and rewrite the production rule so that the
language does define some names – the names 'BEN D JONES', 'JO GOLOMBEK' and
ALULIM' should all be defined.

(2)

(Total 7 marks)

2. The Backus-Naur Form (BNF) production rules below define the syntax of a number of programming language constructs.

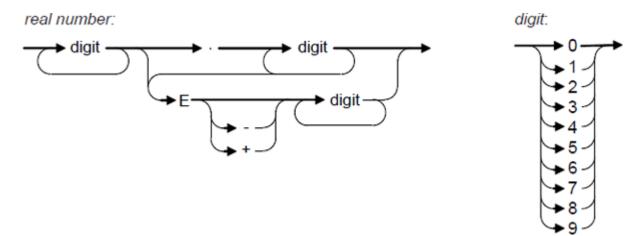
A <letter> is any alphabetic character from a to z or A to z.

	<variable></variable>	Valid? (✓ any number of rows)	
	а		
	money_paid		
	taxrate2		
	2ndPlayerName		
The product	ion rule for an <integer></integer>	is recursive.	
Explain why	recursion has been used	in this production rule.	
Here is an e	xample of a valid <forloo< th=""><th></th><th></th></forloo<>		
	FOR CC	ount = 1 TO 10	
The BNF pro	FOR co	ount = 1 TO 10	her or not a <forloop> is</forloop>
The BNF pro syntactically However, it i correct <for< td=""><td>FOR co</td><td>ount = 1 TO 10 De used to check whet ming language statem</td><th>ent that is a syntactically</th></for<>	FOR co	ount = 1 TO 10 De used to check whet ming language statem	ent that is a syntactically
The BNF prosyntactically However, it is correct <for compiled.="" describe="" on<="" td=""><td>FOR conduction rules above can be correct. is possible that a programme loop may still produce a</td><td>ount = 1 TO 10 De used to check whet ming language statem in error when the programmer.</td><th>ent that is a syntactically</th></for>	FOR conduction rules above can be correct. is possible that a programme loop may still produce a	ount = 1 TO 10 De used to check whet ming language statem in error when the programmer.	ent that is a syntactically
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The BNF prosyntactically However, it is correct <for compiled.="" describe="" on<="" td=""><td>poduction rules above can be correct. is possible that a program at 100p> may still produce at the example of why a syntantic conditions.</td><td>ount = 1 TO 10 De used to check whet ming language statem in error when the programmer.</td><th>ent that is a syntactically ram that it is part of is</th></for>	poduction rules above can be correct. is possible that a program at 100p> may still produce at the example of why a syntantic conditions.	ount = 1 TO 10 De used to check whet ming language statem in error when the programmer.	ent that is a syntactically ram that it is part of is

The table below contains a list of variable names. Place a tick in each row if the stated

(a)

In a particular programming language, the correct syntax for a real number is defined by the syntax diagrams in the diagram below.



(a) Write **Yes** or **No** in the spaces in the empty column of the table below to identify whether or not the numbers listed in the table are valid real numbers which conform to the correct syntax for this language.

Real number	Valid? (Yes / No)
203.412	
-12.87	
12.43E-12	

(3)

(b) In the same language:

A *digit* is defined as any single numeric symbol from this list: 0, 1, 2, 3, 4, 5, 6, 7, 8, 9. A *whole number* is defined as a sequence of one or more *digits*.

An integer is defined as a whole number or a + or a - symbol followed by a whole number.

Write Backus-Naur Form (BNF) production rules for digit, whole number and in		

(3)

(Total 6 marks)