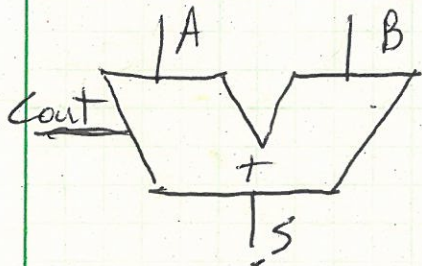


Overview:

- adders
 - half adder
 - full adder
- carry propagate
- subtractors
- Comparators
 - equality
 - magnitude

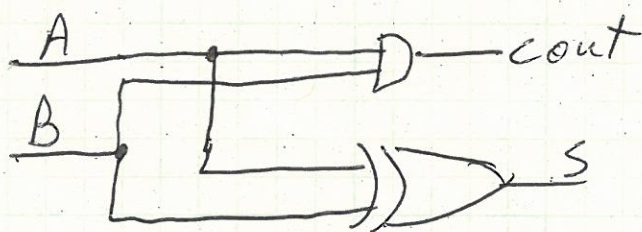
Half Adder (1-Bit)

A	B	Cout	S
0	0	0	0
0	1	0	1
1	0	0	1
1	1	1	0

Eqns:

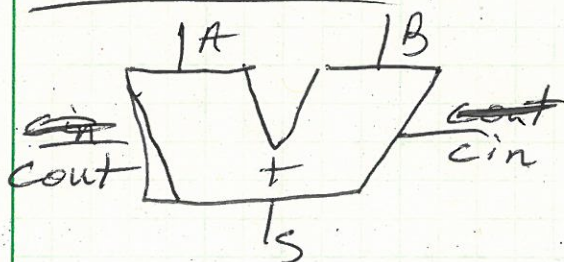
$$\Rightarrow \begin{aligned} \text{Cout} &= AB \\ S &= A \oplus B \end{aligned}$$

Any Problems with this?

Draw schematic:

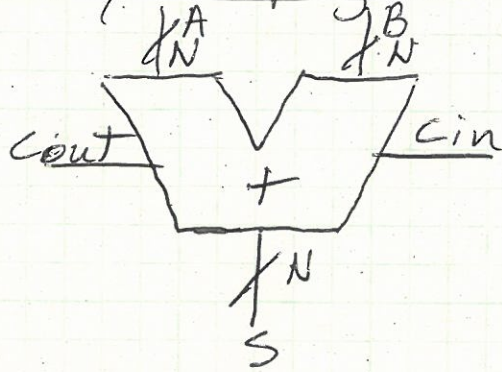
Limited in function.
Doesn't contain circuitry
to take/make use of
cin/cout.

So we use a Full Adder

Full Adder: (1-Bit)

Cin	A	B	Cout	S
0	0	0	0	0
0	0	1	0	1
0	1	0	0	1
0	1	1	1	0
1	0	0	0	1
1	0	1	1	0
1	1	0	1	0
1	1	1	1	1

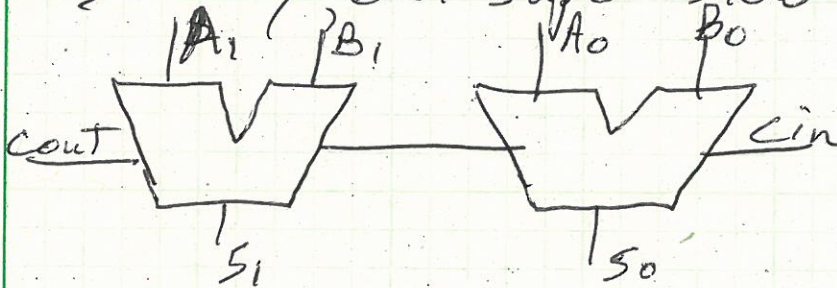
Eqns: $S = \text{Cin} \oplus A \oplus B$ $\text{Cout} = \text{Cin}B + AB + \text{Cin}A$

Carry PropagateTypes:

- ripple carry
- carry Lookahead
- pre fix

Ripple Carry:

- N - Full adders chained
 - delay grows linearly
- ⇒ easy but super slow

Carry Lookahead

- Divides adder into blocks to determine carry out ~~asap~~ ASAP
 - $N > 16$ - we start to see timing gains
 - faster than Ripple carry
- ⇒ delay still grows linearly

- Generate - if carry out produced independently of cin
 $G = AB$ (means we will generate a carry)
- Propagate - cout produced whenever there is cin
 $P = A + B$ (means we will propagate (ie pass) a carry)

Examples:

$$\begin{array}{r} 1010 \\ + 0101 \\ \hline \end{array} \quad \leftarrow C_{in} = 1$$

Count? **yes!** $P = A + B$

$$\begin{array}{r} 1000 \\ + 1000 \\ \hline \end{array}$$

Count? **yes!** $G = AB$

$$\begin{array}{r} 1011 \\ 0101 \\ \hline \end{array}$$

Count? **yes!** why?

because of multiplication, we will have a carry in the i th most Bit

$$C_i = A_i B_i + (A_i + B_i) C_{i-1} = G_i + P_i C_{i-1}$$

$$\begin{array}{r} 3210 \\ 1011 \\ 0101 \\ \hline \end{array}$$

$$G_0 = 1$$

$$P_{1,2,3} = 1$$

Since we generate ~~any~~ carry in Bit 0 & then propagate through $1 \rightarrow 3$ we have a Count?

$$\begin{array}{r} 1010 \\ + 0110 \\ \hline \end{array}$$

Count? **yes!**

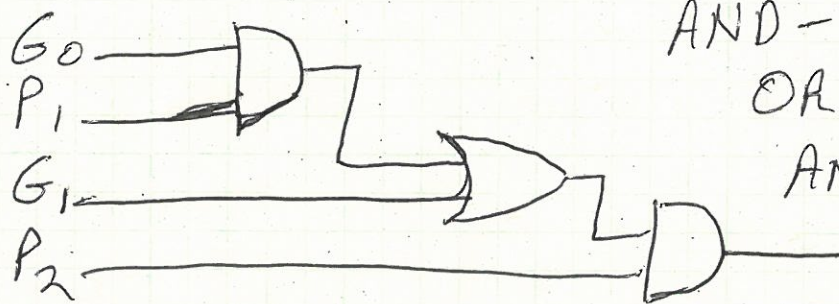
$$G_1 = 1$$

$$P_{2,3} = 1$$

$$\begin{array}{r} 1011 \\ 0111 \\ \hline \end{array} \quad \begin{array}{c} \leftarrow P P P G \\ \quad \quad \quad \leftarrow G \end{array}$$

$$\begin{array}{r} 1001 \\ + 0010 \\ \hline \end{array} \quad \begin{array}{c} \text{add} \\ \leftarrow C_{in} = 1 \\ \downarrow \end{array}$$

Count? **No!** **No**



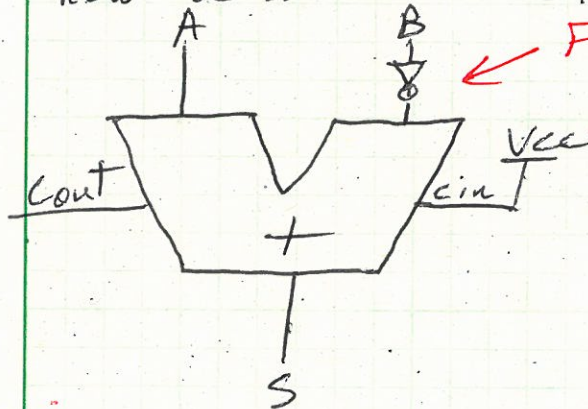
AND - next column propagate
 OR - Next col Generate
 AND - Next col Prop.

prefix:

- Look ahead by 2^x columns
- makes adder faster
- delays grow logarithmically

Subtractor:

how do we subtract? using 2's complement.

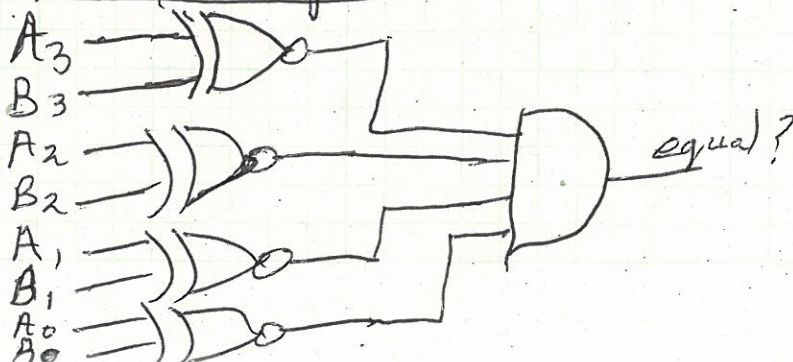


We can still use the same adder (Full adder) we used before

Comparator:

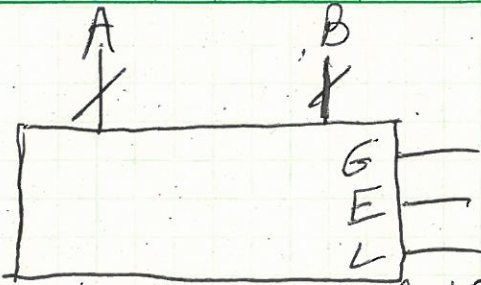
- equality
- magnitude

Equality Comparator:



XNOR:

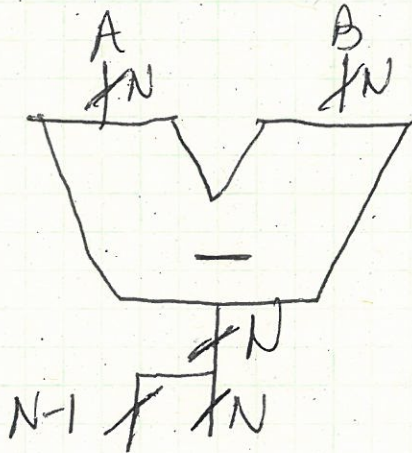
A	B	Y
0	0	1
0	1	0
1	0	0
1	1	1



g = greater
 e = equal
 l = less

Why is this useful?
 How does it work

Boolean if statements in code



Implement this function using things we learned today? No gates just Logic Functions

If $a < 4$
 $z = y + 3$
 else

