6 - Clustering problems

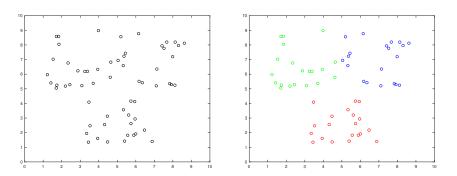
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Clustering problem - definition

Given a set S of patterns and an integer number k, find a partition of S in k subsets S_1, \ldots, S_k (clusters) that are homogeneous and well separated.



Clustering problem is central in unsupervisioned machine learning.

Clustering problem - optimization model

Assume that patterns are vectors $p_1, \ldots, p_\ell \in \mathbb{R}^n$.

Consider a distance $d: \mathbb{R}^n \times \mathbb{R}^n \to \mathbb{R}$ between vectors in \mathbb{R}^n . In the following $d(x,y) = \|x-y\|_2^2$ or $d(x,y) = \|x-y\|_1$.

For each cluster S_j introduce a centroid $x_j \in \mathbb{R}^n$ (unknown). Define clusters so that each pattern is associated to the closest centroid.

We aim to find k centroids in order to minimize the sum of the distances between each pattern and the closest centroid:

$$\begin{cases} \min \sum_{i=1}^{\ell} \min_{j=1,\ldots,k} d(p_i, x_j) \\ x_j \in \mathbb{R}^n & \forall j = 1, \ldots, k \end{cases}$$

Consider the distance $d(x, y) = ||x - y||_2^2$.

The optimization problem to solve is

$$\begin{cases} \min \sum_{i=1}^{\ell} \min_{j=1,\dots,k} \|p_i - x_j\|_2^2 \\ x_j \in \mathbb{R}^n \quad \forall j = 1,\dots,k \end{cases}$$

If k = 1 (one cluster), then it is a convex quadratic programming problem without constraints:

$$\begin{cases} \min \sum_{i=1}^{\ell} \|p_i - x\|_2^2 = \sum_{i=1}^{\ell} (x - p_i)^{\mathsf{T}} (x - p_i) \\ x \in \mathbb{R}^n \end{cases}$$
 (1)

The global optimum is the stationary point:

$$2\ell x - 2\sum_{i=1}^{\ell} p_i = 0 \iff x = \frac{\sum_{i=1}^{\ell} p_i}{\ell}$$
 (mean or baricenter)

If k > 1 (at least two clusters), then the problem is nonconvex and nonsmooth:

$$\begin{cases}
\min_{x} \sum_{i=1}^{\ell} \min_{j=1,\dots,k} \|p_i - x_j\|_2^2 \\
x_j \in \mathbb{R}^n \quad \forall j = 1,\dots,k
\end{cases} \tag{2}$$

Theorem

Problem (2) is equivalent to the following nonconvex smooth problem:

$$\begin{cases}
\min_{x,\alpha} \sum_{i=1}^{\ell} \sum_{j=1}^{k} \alpha_{ij} \| p_i - x_j \|_2^2 \\
\sum_{j=1}^{k} \alpha_{ij} = 1 \quad \forall i = 1, \dots, \ell \\
\alpha_{ij} \ge 0 \quad \forall i = 1, \dots, \ell, j = 1, \dots, k \\
x_j \in \mathbb{R}^n \quad \forall j = 1, \dots, k.
\end{cases}$$
(3)

Proof. Notice that
$$\min_{j=1,\ldots,k} \{a_j\} = \min \left\{ \sum_{j=1}^k \alpha_j a_j : \sum_{j=1}^k \alpha_j = 1, \ \alpha \geq 0 \right\}.$$

The k-means algorithm is based on the following properties of problem (3):

▶ If x_j are fixed, then (3) is decomposable into ℓ very simple LP problems: for any $i = 1, ..., \ell$, the optimal solution is

$$\alpha_{ij}^* = \begin{cases} 1 & \text{if } j \text{ is the first index s.t. } \|p_i - x_j\|_2 = \min_{h=1,\dots,k} \|p_i - x_h\|_2 \\ & (x_j \text{ is the first closest centroid to } p_i), \\ 0 & \text{otherwise.} \end{cases}$$

▶ If α_{ij} are fixed, then (3) is decomposable into k very simple convex QP problems similar to (1): for any j = 1, ..., k, the optimal solution is

$$x_j^* = rac{\sum\limits_{i=1}^{\ell} lpha_{ij} p_i}{\sum\limits_{i=1}^{\ell} lpha_{ij}}$$
 (mean of patterns).

The k-means algorithm consists in an alternating minimization of

$$f(x,\alpha) = \sum_{i=1}^{\ell} \sum_{j=1}^{k} \alpha_{ij} \|p_i - x_j\|_2^2$$
 with respect to the two blocks of variables x and α .

0. (Inizialization) Set t=0, choose centroids $x_1^0,\ldots,x_k^0\in\mathbb{R}^n$ and assign patterns to clusters: for any $i=1,\ldots,\ell$

$$\alpha_{ij}^0 = \begin{cases} 1 & \text{if } j \text{ is the first index s.t. } \| p_i - x_j^0 \|_2 = \min_{h=1,\dots,k} \| p_i - x_h^0 \|_2 \\ 0 & \text{otherwise.} \end{cases}$$

1. (Update centroids) For each j = 1, ..., k compute the mean

$$x_j^{t+1} = \left(\sum_{i=1}^{\ell} \alpha_{ij}^t p_i\right) / \left(\sum_{i=1}^{\ell} \alpha_{ij}^t\right).$$

2. (Update clusters) For any $i = 1, ..., \ell$ compute

$$\alpha_{ij}^{t+1} = \begin{cases} 1 & \text{if } j \text{ is the first index s.t. } \|p_i - x_j^{t+1}\|_2 = \min_{h=1,\dots,k} \|p_i - x_h^{t+1}\|_2 \\ 0 & \text{otherwise.} \end{cases}$$

3. (Stopping criterion) If $f(x^{t+1}, \alpha^{t+1}) = f(x^t, \alpha^t)$ then STOP else t = t + 1, go to Step 1.

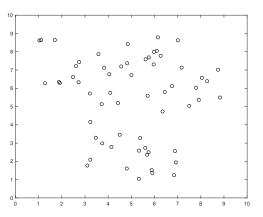
Theorem

The k-means algorithm stops after a finite number of iterations at a solution (x^*, α^*) of the KKT system of problem (3) such that

$$\begin{split} f(x^*,\alpha^*) &\leq f(x^*,\alpha), & \forall \ \alpha \geq 0 \ \text{ s.t. } \sum_{j=1}^k \alpha_{ij} = 1 \quad \forall \ i=1,\ldots,\ell, \\ f(x^*,\alpha^*) &\leq f(x,\alpha^*), & \forall \ x \in \mathbb{R}^{kn}. \end{split}$$

Remark. The k-means algorithm does not guarantee to find a global optimum.

Exercise 6.1. Consider the k-means algorithm, with k = 3, for the set of patterns given in the file 6-1.txt.



- a) Run the algorithm starting from centroids $x_1 = (5,7)$, $x_2 = (6,3)$, $x_3 = (4,3)$.
- **b)** Run the algorithm starting from centroids $x_1 = (5,7)$, $x_2 = (6,3)$, $x_3 = (4,4)$.
- c) Is it possible to improve the solutions obtained in a) and b)?

Consider now the distance $d(x, y) = ||x - y||_1$. The optimization problem to solve is

$$\begin{cases} \min \sum_{i=1}^{\ell} \min_{j=1,\dots,k} \|p_i - x_j\|_1 \\ x_j \in \mathbb{R}^n & \forall j = 1,\dots,k \end{cases}$$

If k=1 (one cluster), then it is a convex problem decomposable into n convex problems of one variable:

$$\begin{cases}
\min \sum_{i=1}^{\ell} \|p_i - x\|_1 = \sum_{i=1}^{\ell} \sum_{h=1}^{n} |x_h - (p_i)_h| = \sum_{h=1}^{n} \underbrace{\sum_{i=1}^{\ell} |x_h - (p_i)_h|}_{f_h(x_h)} \\
x \in \mathbb{R}^n
\end{cases} (4)$$

Given ℓ real numbers $a_1 < a_2 < \cdots < a_{\ell}$, what is the optimal solution of

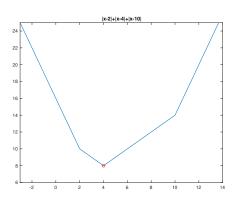
$$\begin{cases} \min \sum_{i=1}^{\ell} |x - a_i| = f(x) \\ x \in \mathbb{R} \end{cases}$$
?

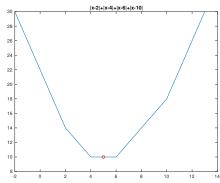
The objective function is convex and piecewise linear:

$$f(x) = \begin{cases} -\ell x + \sum_{i=1}^{\ell} a_i & \text{if } x < a_1 \\ (2 - \ell)x + \sum_{i=2}^{\ell} a_i - a_1 & \text{if } x \in [a_1, a_2] \\ \dots & \dots \\ (2r - \ell)x + \sum_{i=r+1}^{\ell} a_i - \sum_{i=1}^{r} a_1 & \text{if } x \in [a_r, a_{r+1}] \\ \dots & \dots \\ (\ell - 2)x + a_{\ell} - \sum_{i=1}^{\ell-1} a_i & \text{if } x \in [a_{\ell-1}, a_{\ell}] \\ \ell x - \sum_{i=1}^{\ell} a_i & \text{if } x > a_{\ell} \end{cases}$$

The global optimum is
$$\operatorname{median}(a_1,\ldots,a_\ell) = \begin{cases} a_{(\ell+1)/2} & \text{if } \ell \text{ is odd,} \\ \\ \frac{a_{\ell/2} + a_{1+\ell/2}}{2} & \text{if } \ell \text{ is even.} \end{cases}$$

The global optimum is
$$\operatorname{median}(a_1,\ldots,a_\ell) = \begin{cases} a_{(\ell+1)/2} & \text{if } \ell \text{ is odd,} \\ \frac{a_{\ell/2} + a_{1+\ell/2}}{2} & \text{if } \ell \text{ is even.} \end{cases}$$





If k > 1 (at least two clusters), then the problem is nonconvex and nonsmooth:

$$\begin{cases}
\min_{x} \sum_{i=1}^{\ell} \min_{j=1,\dots,k} \|p_i - x_j\|_1 \\
x_j \in \mathbb{R}^n \quad \forall j = 1,\dots,k
\end{cases}$$
(5)

Theorem

Problem (5) is equivalent to the following problem:

$$\begin{cases}
\min_{x,\alpha} \sum_{i=1}^{\ell} \sum_{j=1}^{k} \alpha_{ij} \| p_i - x_j \|_1 \\
\sum_{j=1}^{k} \alpha_{ij} = 1 \quad \forall i = 1, \dots, \ell \\
\alpha_{ij} \ge 0 \quad \forall i = 1, \dots, \ell, j = 1, \dots, k \\
x_j \in \mathbb{R}^n \quad \forall j = 1, \dots, k.
\end{cases}$$
(6)

Theorem

Problem (6) is, in turn, equivalent to the nonconvex bilinear problem:

$$\begin{cases} \min_{x,\alpha,u} \sum_{i=1}^{\ell} \sum_{j=1}^{k} \sum_{h=1}^{n} \alpha_{ij} u_{ijh} \\ u_{ijh} \geq (p_i)_h - (x_j)_h & \forall i = 1, \dots, \ell, \ j = 1, \dots, k, \ h = 1, \dots, n \\ u_{ijh} \geq (x_j)_h - (p_i)_h & \forall i = 1, \dots, \ell, \ j = 1, \dots, k, \ h = 1, \dots, n \\ \sum_{j=1}^{k} \alpha_{ij} = 1 & \forall i = 1, \dots, \ell \\ \alpha_{ij} \geq 0 & \forall i = 1, \dots, \ell, \ j = 1, \dots, k \\ x_j \in \mathbb{R}^n & \forall j = 1, \dots, k. \end{cases}$$
(7)

The k-median algorithm is based on the following properties of problem (6):

▶ If x_j are fixed, then (6) is decomposable into ℓ very simple LP problems: for any $i = 1, ..., \ell$, the optimal solution is

$$\alpha_{ij}^* = \begin{cases} 1 & \text{if } j \text{ is the first index s.t. } \|p_i - x_j\|_1 = \min_{h=1,\dots,k} \|p_i - x_h\|_1 \\ & (x_j \text{ is the first closest centroid to } p_i), \\ 0 & \text{otherwise.} \end{cases}$$

▶ If $\alpha_{ij} \in \{0,1\}$ are fixed, then (6) is decomposable into k very simple convex problems similar to (4): for any $j=1,\ldots,k$, the optimal solution is

$$x_j^* = \text{median}(p_i : \alpha_{ij} = 1).$$

The k-median algorithm consists in an alternating minimization of

$$f(x, \alpha) = \sum_{i=1}^{\ell} \sum_{j=1}^{k} \alpha_{ij} \|p_i - x_j\|_1$$
 with respect to the two blocks of variables x and α .

0. (Inizialization) Set t=0, choose centroids $x_1^0,\ldots,x_k^0\in\mathbb{R}^n$ and assign patterns to clusters: for any $i=1,\ldots,\ell$

$$\alpha_{ij}^0 = \begin{cases} 1 & \text{if } j \text{ is the first index s.t. } \|p_i - x_j^0\|_1 = \min_{h=1,\dots,k} \|p_i - x_h^0\|_1 \\ 0 & \text{otherwise.} \end{cases}$$

1. (Update centroids) For each j = 1, ..., k compute

$$x_i^{t+1} = \text{median}(p_i : \alpha_{ii}^t = 1).$$

2. (Update clusters) For any $i = 1, ..., \ell$ compute

$$\alpha_{ij}^{t+1} = \begin{cases} 1 & \text{if } j \text{ is the first index s.t. } \| p_i - x_j^{t+1} \|_1 = \min_{h=1,...,k} \| p_i - x_h^{t+1} \|_1 \\ 0 & \text{otherwise.} \end{cases}$$

3. (Stopping criterion) If $f(x^{t+1}, \alpha^{t+1}) = f(x^t, \alpha^t)$ then STOP else t = t+1, go to Step 1.

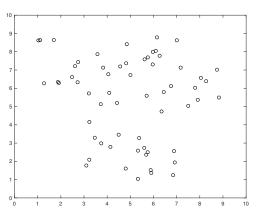
Theorem

The k-median algorithm stops after a finite number of iterations at a stationary point (x^*, α^*) of problem (6) such that

$$f(x^*, \alpha^*) \le f(x^*, \alpha),$$
 $\forall \alpha \ge 0 \text{ s.t. } \sum_{j=1}^k \alpha_{ij} = 1 \quad \forall i = 1, \dots, \ell,$
 $f(x^*, \alpha^*) \le f(x, \alpha^*),$ $\forall x \in \mathbb{R}^{kn}.$

Remark. The k-means algorithm does not guarantee to find a global optimum.

Exercise 6.2. Consider the k-median algorithm, with k = 3, for the set of patterns given in the file 6-1.txt.



- a) Run the algorithm starting from centroids $x_1 = (5,7)$, $x_2 = (6,3)$, $x_3 = (4,3)$.
- **b)** Run the algorithm starting from centroids $x_1 = (5,7)$, $x_2 = (6,3)$, $x_3 = (4,4)$.
- c) Is it possible to improve the solutions obtained in a) and b)?