

## Project2: Optimizing the Performance of a Pipelined Processor

CS 359, COMPUTER ARCHITECTURE, Yanyan Shen, Spring 2017

### 1 Preknowledge

I have read the CSAPP book and concluded several basic points which are essential to solving the problems in Part A, B, and C.

- There are 8 registers in the Y86 system

*%eax, %ecx, %edx, %ebx*

*%esi, %edi, %esp, %ebp*

Each of these registers stores a word. Among them register *%esp* is used as a stack pointer by the push, pop, call, and return instructions.

- There are three single-bit condition codes, ZF, SF, and OF, storing information about the effect of the most recent arithmetic or logical instruction.
- The Y86 instruction set is largely a subset of the IA32 instruction set but it still have some differences. The picture below shows the Y86 instruction set.

Byte	0	1	2	3	4	5
halt	0	0				
nop	1	0				
rrmovl <b>rA</b> , <b>rB</b>	2	0	<b>rA</b>	<b>rB</b>		
irmovl <b>V</b> , <b>rB</b>	3	0	F	<b>rB</b>	<b>V</b>	
rmmovl <b>rA</b> , <b>D(rB)</b>	4	0	<b>rA</b>	<b>rB</b>	<b>D</b>	
mrmovl <b>D(rB)</b> , <b>rA</b>	5	0	<b>rA</b>	<b>rB</b>	<b>D</b>	
OPl <b>rA</b> , <b>rB</b>	6	<b>fn</b>	<b>rA</b>	<b>rB</b>		
jXX <b>Dest</b>	7	<b>fn</b>	<b>Dest</b>			
cmovXX <b>rA</b> , <b>rB</b>	2	<b>fn</b>	<b>rA</b>	<b>rB</b>		
call <b>Dest</b>	8	0	<b>Dest</b>			
ret	9	0				
pushl <b>rA</b>	A	0	<b>rA</b>	F		
popl <b>rA</b>	B	0	<b>rA</b>	F		

- In the Y86 system, the stack starts at a certain address and grows toward lower addresses, which prevents space conflict.

## 2 Part A

### 2.1 Description

- The program **sum.y**s was used to sum linked list elements iteratively. We are supposed to add the sum of a list from head to tail using the Y86 coding rules.
- The program **rsum.y**s is similar to the **sum.y**s, except it sums linked list elements recursively. We are supposed to add the sum of a list from head to tail using the Y86 coding rules.
- The program **copy.y**s is used for two purposes. First it copies a block of words from one part of memory to another area of memory. Second it computes the checksum (Xor) of all the words copied.

### 2.2 Solution

sum.y

```
1      irmovl Stack,%esp
2      rrmovl %esp,%ebp
3      irmovl ele1 , %edx
4      pushl %edx
5      call sum_list
6      halt
7
8      .align 4
9      ele1:
10     .long 0x00a
11     .long ele2
12     ele2:
13     .long 0x0b0
14     .long ele3
15     ele3:
16     .long 0xc00
17     .long 0
18
19     sum_list:
20     pushl %ebp
21     xorl %eax,%eax
22     rrmovl %esp,%ebp
23     mrmovl 8(%ebp),%edx
24     andl %edx,%edx
25     je L4
26     L5:
27     mrmovl (%edx),%esi
28     addl %esi,%eax
29     mrmovl 4(%edx),%edx
30     andl %edx,%edx
31     jne L5
32     L4:
33     rrmovl %ebp,%esp
34     popl %ebp
```

```

35         ret
36
37         .pos 0x100
38         Stack:

```

#### rsum.ys

```

1         .pos 0
2         init:
3         irmovl Stack, %esp
4         irmovl Stack, %ebp
5         jmp Main
6
7         .align 4
8         ele1:
9         .long 0x00a
10        .long ele2
11        ele2:
12        .long 0x0b0
13        .long ele3
14        ele3:
15        .long 0xc00
16        .long 0
17
18        Main:
19        irmovl ele1, %edx
20        pushl %edx
21        call rsum_list
22        halt
23
24        rsum_list:
25        pushl %ebp
26        rrmovl %esp,%ebp
27        mrmovl 0x8(%ebp),%edx
28        xorl %eax,%eax
29        pushl %ebx
30        andl %edx,%edx
31        je End
32        irmovl $0xc,%esi
33        subl %esi,%esp
34        mrmovl (%edx),%ebx
35        mrmovl 0x4(%edx),%ecx
36        pushl %ecx
37        call rsum_list
38        addl %ebx,%eax
39
40        End:
41        mrmovl 0xffffffffc(%ebp),%ebx
42        rrmovl %ebp, %esp
43        popl %ebp
44        ret

```

```

45
46         .pos      0x100
47         Stack:

```

# copy.ys

```

1         .pos 0
2         init:
3         irmovl Stack, %esp
4         irmovl Stack, %ebp
5         jmp Main
6
7         .align 4
8         src:
9         .long 0x00a
10        .long 0x0b0
11        .long 0xc00
12        # Destination block
13        dest:
14        .long 0x111
15        .long 0x222
16        .long 0x333
17
18        Main:
19        irmovl $3,%eax
20        pushl %eax
21        irmovl dest,%edx
22        pushl %edx
23        irmovl src,%ecx
24        pushl %ecx
25        call Copy
26        halt
27
28        Copy:
29        pushl %ebp
30        rrmovl %esp,%ebp
31        mrmovl 8(%ebp),%ecx
32        mrmovl 12(%ebp),%ebx
33        mrmovl 16(%ebp),%edx
34        irmovl $0,%eax
35        andl %edx,%edx
36        je End
37
38        Loop:
39        mrmovl (%ecx),%esi
40        rmmovl %esi, (%ebx)
41        xorl %esi,%eax
42        irmovl $4,%edi
43        addl %edi,%ecx
44        addl %edi,%ebx
45        irmovl $-1,%edi

```

```

46         addl %edi,%edx
47         jne Loop
48
49         End:
50         popl %ebp
51         ret
52
53         .pos 0x100
54         Stack:

```

## 2.3 Analysis

### 2.3.1 sum.js

- For the program **sum.js**, the line 1 ~ 19 and line 37 ~ 38 are the preparatory work and do not need further explanation. For the main part, which is the *sum.js*, first we store the *%ebp* part to stack, then we set *%eax* to zero as the initial sum value, *%edx* as the initial list position, then if the list has reached its end, we jump to state L5, which restore the value. If not, we goto the LOOP L5 part.
- In the L5 part, first  $\%eax + \%esi$  to add to the sum, then  $(\%edx) + 4$  to goto the list's next elements. After doing that, again we judge whether the list has reached its end and the condition is exactly the same.
- Based the outcome of this program, *%eax* stores the overall sum value, which is *0xcba*, *%ebp* get popped and get its initial value, which is the *0x100*, *%esi* stores the last elements of the list, which is *%0xc00*, which proves that our program is correct.

### 2.3.2 rsum.js

- For the program **rsum.js**, the line 1 ~ 17 and line 46 ~ 47 are the preparatory work and do not need further explanation. For the main part, first we let *%edx* stores the *el*, which is the beginning, then we save *%edx* and call *rsum\_list*.
- For the *rsum\_list* part(line24 ~ 44), which is the recursive version of the *sum\_list*, first we use *%edx* to get starting address, *%eax* to get the initial result which is 0, then we save *%ebx* and compare whether *%edx* is zero or not. If so, goto *end*, else, goto the next part, which is the recursive part.
- For the recursive part, the main idea is to push  $\%edx \rightarrow next$  as the first parameter and then call *rsum\_list* by recursion.
- Based the outcome of this program, *%ebp* pop out and its initial value is *0x100*, *%eax* stores the overall sum value *0xcba*, which proves that our program is correct.

### 2.3.3 copy.js

- For the program **copy.js**, the line 1 ~ 17 and line 53 ~ 54 are the preparatory work and do not need further explanation. For the main part(line 18 ~ 26), first we change *%eax* to 3, *%edx* to the *dest*, *%ecx* to the *src* and store the value to the stack, then we call Copy function and halt.

- For the Copy function(line 28 ~ 47), first we let `%ebp` store `%esp`, `%ecx` store `src`, `%ebx` store `dest`, `%edx` store the length, `%eax` stores the initial value of the result which is 0, then we compare whether the copy function has reached its end, if so, goto the End part, else, goto the Loop part.
- For the Loop part, first we get the address of the `src`, then we begin the copy process, while we add the `src` and `dest` by 4 to move to the next address, in the meantime we let `len - 1` to serve as flag to decide when to stop.
- Based the outcome of this program, the address 0x0020 value is changed from 0x111 to 0xa, the address 0x0024 value is changed from 0x222 to 0xb0, the address 0x0028 value is changed from 0x333 to 0xc00, which proves that our program is correct.

## 2.4 Outcome

```

marshallee@ubuntu: ~/Documents/computer_architecture/project2/Project2/project2-h
marshallee@ubuntu:~/Documents/computer_architecture/project2/Project2/project2-h
andout/sim/sim/misc$ ./yis sum.yo
Stopped in 30 steps at PC = 0x15. Status 'HLT', CC Z=1 S=0 O=0
Changes to registers:
%eax: 0x00000000      0x00000cba
%esp: 0x00000000      0x00000fc
%ebp: 0x00000000      0x00000100
%esi: 0x00000000      0x00000c00
Changes to memory:
0x00f4: 0x00000000      0x00000100
0x00f8: 0x00000000      0x00000015
0x00fc: 0x00000000      0x00000018

```

```

marshallee@ubuntu: ~/Documents/computer_architecture/project2/Project2/project2-h
marshallee@ubuntu:~/Documents/computer_architecture/project2/Project2/project2-h
andout/sim/sim/misc$ ./yis rsum.yo
Stopped in 72 steps at PC = 0x39. Status 'HLT', CC Z=0 S=0 O=0
Changes to registers:
%eax: 0x00000000      0x00000cba
%esp: 0x00000000      0x00000fc
%ebp: 0x00000000      0x00000100
%esi: 0x00000000      0x0000000c
Changes to memory:
0x009c: 0x00000000      0x00000c00
0x00a0: 0x00000000      0x000000bc
0x00a4: 0x00000000      0x0000006a
0x00b8: 0x00000000      0x000000b0
0x00bc: 0x00000000      0x000000d8
0x00c0: 0x00000000      0x0000006a
0x00c4: 0x00000000      0x00000024
0x00d4: 0x00000000      0x0000000a
0x00d8: 0x00000000      0x000000f4
0x00dc: 0x00000000      0x0000006a
0x00e0: 0x00000000      0x0000001c
0x00f4: 0x00000000      0x00000100
0x00f8: 0x00000000      0x00000039
0x00fc: 0x00000000      0x00000014

```

## 2.5 Review

# 3 Part B

## 3.1 Description

In this part, we are asked to add a new instruction *iaddl* to the SEQ processor, this instruction is meant to add a constant to a register.

```
marshallee@ubuntu: ~/Documents/computer_architecture/project2/Project2/project2-h
marshallee@ubuntu:~/Documents/computer_architecture/project2/Project2/project2-h
andout/sim/sim/misc$ ./yis copy.yo
Stopped in 48 steps at PC = 0x49. Status 'HLT', CC Z=1 S=0 O=0
Changes to registers:
%eax: 0x00000000      0x00000cba
%ecx: 0x00000000      0x00000020
%ebx: 0x00000000      0x0000002c
%esp: 0x00000000      0x000000f4
%ebp: 0x00000000      0x00000100
%esi: 0x00000000      0x00000c00
%edi: 0x00000000      0xffffffff

Changes to memory:
0x0020: 0x00000111      0x0000000a
0x0024: 0x00000222      0x000000b0
0x0028: 0x00000333      0x00000c00
0x00ec: 0x00000000      0x00000100
0x00f0: 0x00000000      0x00000049
0x00f4: 0x00000000      0x00000014
0x00f8: 0x00000000      0x00000020
0x00fc: 0x00000000      0x00000003
```

### 3.2 Solution

### 3.3 Analysis

It can be implemented by first using *irmovl* instruction to let the register contains the constant number, then we can use *addl* instruction to add the constant number to the destination register.

### 3.4 Outcome

### 3.5 Review



