

Dynamic Programming

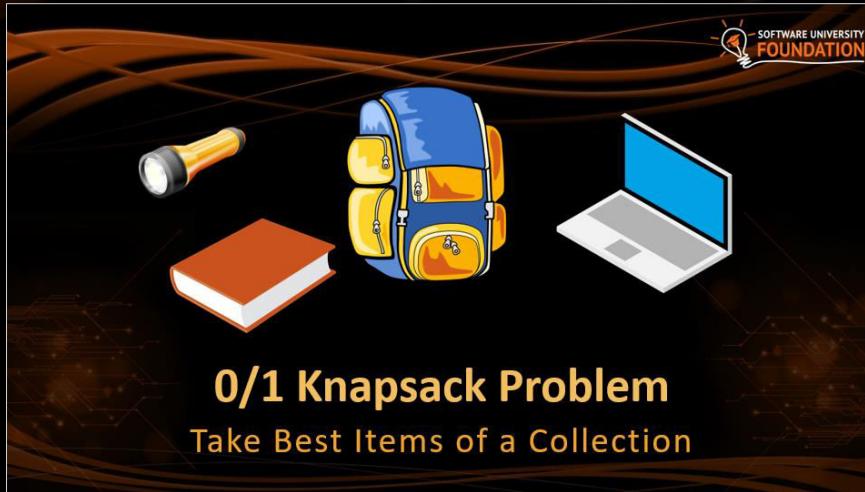
Solving Optimization Problems
with Dynamic Programming Part 2



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Technical Trainers
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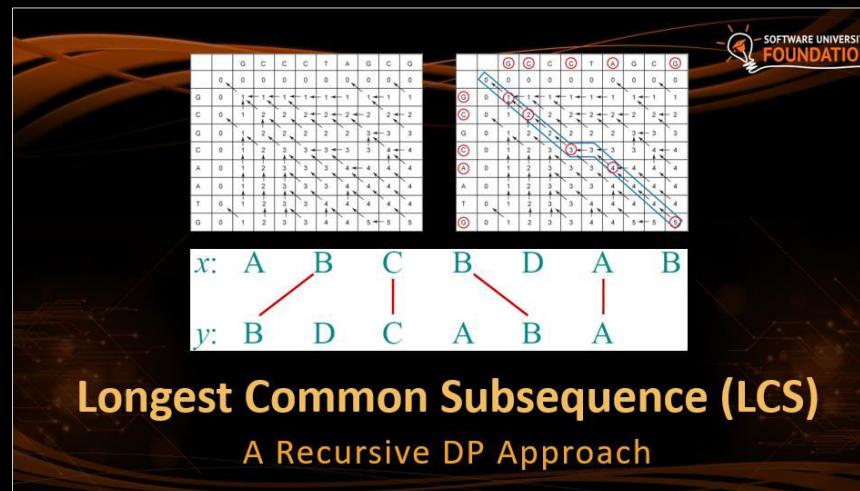
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Take Best Items of a Collection



Subset Sum
Sum with Limited Coins



Longest Common Subsequence (LCS)
A Recursive DP Approach

Sequence x : A B C B D A B

Sequence y : B D C A B A

G	C	C	T	A	G	C	G
G	0	1	2	3	4	5	6
C	0	1	2	3	4	5	6
G	0	1	2	3	4	5	6
A	0	1	2	3	4	5	6
A	0	1	2	3	4	5	6
T	0	1	2	3	4	5	6
G	0	1	2	3	4	5	6

G	C	C	T	A	G	C	G
G	0	1	2	3	4	5	6
C	0	1	2	3	4	5	6
G	0	1	2	3	4	5	6
A	0	1	2	3	4	5	6
A	0	1	2	3	4	5	6
T	0	1	2	3	4	5	6
G	0	1	2	3	4	5	6

Have a Question?



sli.do

#DsAlgo



0/1 Knapsack Problem

Take Best Items of a Collection

The "Knapsack" Problem

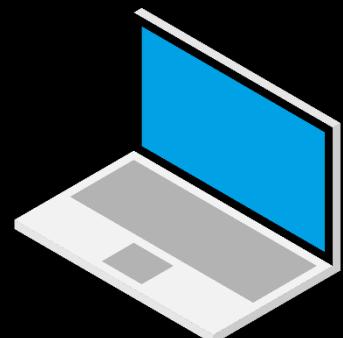
- We have **n** items
 - Each item has price **c_i** and weight **w_i**
- We have a knapsack of given **capacity** (maximum weight)
 - We want to take maximal priced items within the knapsack capacity
- Take a subset of the items, so that:
 - The total price is maximal
 - The total weight \leq capacity

The "Knapsack" Problem – Example



Value: 3

Weight: 2



Value: 1

Weight: 2



Value: 3

Weight: 1



Max Capacity: 4

The "Knapsack" Problem – Example



Value: 3

Weight: 2



Value: 1

Weight: 2

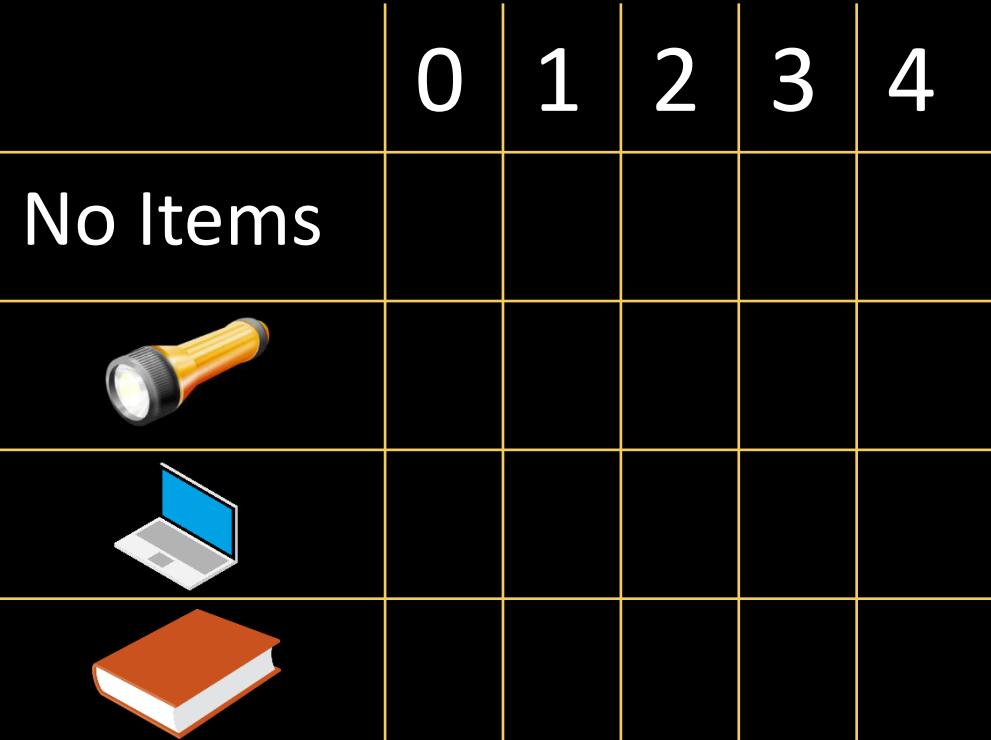


Value: 3

Weight: 1



Max Capacity: 4



The "Knapsack" Problem – Example



Value: 3

Weight: 2



Value: 1

Weight: 2



Value: 3

Weight: 1



Max Capacity: 4

No Items	0	1	2	3	4
	0	0	0	0	0
					
					
					

The "Knapsack" Problem – Example



Value: 3

Weight: 2



Value: 1

Weight: 2



Value: 3

Weight: 1



Max Capacity: 4

No Items	0	1	2	3	4
	0	0	0	0	0
	0	0	0	0	0
	0	0	0	0	0

The "Knapsack" Problem – Example



Value: 3

Weight: 2



Value: 1

Weight: 2



Value: 3

Weight: 1



Max Capacity: 4

Excluding current item value:
`matrix[index - 1][capacity]`

Including current item value:
`value + matrix[index - 1][capacity - weight]`

Choosing if we are going to take the item
`max(excluding, including)`

	0	1	2	3	4
No Items	0	0	0	0	0
Flashlight	0	3	0	0	0
Laptop	0	0	0	0	0
Book	0	0	0	0	0

The "Knapsack" Problem – Example



Value: 3

Weight: 2



Value: 1

Weight: 2



Value: 3

Weight: 1



Max Capacity: 4

Excluding current item value:
excluding = 0

Including current item value:
including = 0

Weight > Capacity

Choosing if we are going to take the item
max(excluding, including)

	0	1	2	3	4
No Items	0	0	0	0	0
Flashlight	0	0	0	0	0
Laptop	0	0	0	0	0
Book	0	0	0	0	0

The "Knapsack" Problem – Example



Value: 3

Weight: 2



Value: 1

Weight: 2



Value: 3

Weight: 1



Max Capacity: 4

Excluding current item value:
`matrix[index - 1][capacity]`

Including current item value:
`value + matrix[index - 1][capacity - weight]`

Choosing if we are going to take the item
`max(excluding, including)`

	0	1	2	3	4
No Items	0	0	0	0	0
Flashlight	0	0	X	0	0
Laptop	0	0	0	0	0
Book	0	0	0	0	0

The "Knapsack" Problem – Example



Value: 3

Weight: 2



Value: 1

Weight: 2



Value: 3

Weight: 1



Max Capacity: 4

Excluding current item value:
`matrix[index - 1][capacity]`

Including current item value:
`value + matrix[index - 1][capacity - weight]`

Choosing if we are going to take the item
`max(excluding, including)`

	0	1	2	3	4
No Items	0	0	0	0	0
Flashlight	0	0	3	3	
Laptop	0				
Book	0				

The "Knapsack" Problem – Example



Value: 3

Weight: 2



Value: 1

Weight: 2



Value: 3

Weight: 1



Max Capacity: 4

Excluding current item value:
`matrix[index - 1][capacity]`

Including current item value:
`value + matrix[index - 1][capacity - weight]`

Choosing if we are going to take the item
`max(excluding, including)`

	0	1	2	3	4
No Items	0	0	0	0	0
Flashlight	0	0	3	3	
Laptop	0				
Book	0				

The "Knapsack" Problem – Example



Value: 3

Weight: 2



Value: 1

Weight: 2



Value: 3

Weight: 1



Max Capacity: 4

Excluding current item value:
`matrix[index - 1][capacity]`

Including current item value:
`value + matrix[index - 1][capacity - weight]`

Choosing if we are going to take the item
`max(excluding, including)`

	0	1	2	3	4
No Items	0	0	0	0	0
Flashlight	0	0	3	3	3
Laptop	0	0	3	3	3
Book	0	0	3	3	3

The "Knapsack" Problem – Example



Value: 3

Weight: 2



Value: 1

Weight: 2



Value: 3

Weight: 1



Max Capacity: 4

Excluding current item value:
`matrix[index - 1][capacity]`

Including current item value:
`value + matrix[index - 1][capacity - weight]`

Choosing if we are going to take the item
`max(excluding, including)`

	0	1	2	3	4
No Items	0	0	0	0	0
Flashlight	0	0	3	3	3
Laptop	0	0	0	0	0
Book	0	0	0	0	0

The table illustrates the state of the knapsack problem. The columns represent the maximum capacity (0 to 4), and the rows represent the items (No Items, Flashlight, Laptop, Book). The values in the cells indicate the maximum value that can be achieved for each combination of item and capacity. A green circle highlights the value 3 for the first row (No Items) at capacity 3. A red circle highlights the value 3 for the first row (No Items) at capacity 4. A red 'X' marks the cell for the first row (No Items) at capacity 1. A blue circle highlights the value 0 for the second row (Flashlight) at capacity 2. A green checkmark marks the cell for the third row (Laptop) at capacity 3. A green checkmark marks the cell for the fourth row (Book) at capacity 4.

The "Knapsack" Problem – Example



Value: 3

Weight: 2



Value: 1

Weight: 2



Value: 3

Weight: 1



Max Capacity: 4

Excluding current item value:
`matrix[index - 1][capacity]`

Including current item value:
`value + matrix[index - 1][capacity - weight]`

Choosing if we are going to take the item
`max(excluding, including)`

	0	1	2	3	4
No Items	0	0	0	0	0
Flashlight	0	0	3	3	3
Laptop	0	0	3	3	3
Book	0	0	3	3	3

The "Knapsack" Problem – Example



Value: 3

Weight: 2



Value: 1

Weight: 2



Value: 3

Weight: 1



Max Capacity: 4

Excluding current item value:
`matrix[index - 1][capacity]`

Including current item value:
`value + matrix[index - 1][capacity - weight]`

Choosing if we are going to take the item
`max(excluding, including)`

	0	1	2	3	4
No Items	0	0	0	0	0
Flashlight	0	0	3	3	3
Laptop	0	0	3	3	
Book	0				

The "Knapsack" Problem – Example



Value: 3

Weight: 2



Value: 1

Weight: 2



Value: 3

Weight: 1



Max Capacity: 4

Excluding current item value:
`matrix[index - 1][capacity]`

Including current item value:
`value + matrix[index - 1][capacity - weight]`

Choosing if we are going to take the item
`max(excluding, including)`

	0	1	2	3	4
No Items	0	0	0	0	0
Flashlight	0	0	3	3	3
Laptop	0	0	3	3	4
Book	0	0	3	3	4

The "Knapsack" Problem – Example



Value: 3

Weight: 2



Value: 1

Weight: 2



Value: 3

Weight: 1



Max Capacity: 4

Excluding current item value:
`matrix[index - 1][capacity]`

Including current item value:
`value + matrix[index - 1][capacity - weight]`

Choosing if we are going to take the item
`max(excluding, including)`

	0	1	2	3	4
No Items	0	0	0	0	0
Flashlight	0	0	3	3	3
Laptop	0	0	3	3	4
Book	0	3	3	4	4

The "Knapsack" Problem – Example



Value: 3

Weight: 2



Value: 1

Weight: 2



Value: 3

Weight: 1



Max Capacity: 4

Excluding current item value:
`matrix[index - 1][capacity]`

Including current item value:
`value + matrix[index - 1][capacity - weight]`

Choosing if we are going to take the item
`max(excluding, including)`

	0	1	2	3	4
No Items	0	0	0	0	0
Flashlight	0	0	3	3	3
Laptop	0	0	3	3	4
Book	0	3	3	0	0

The "Knapsack" Problem – Example



Value: 3

Weight: 2



Value: 1

Weight: 2



Value: 3

Weight: 1



Max Capacity: 4

Excluding current item value:
`matrix[index - 1][capacity]`

Including current item value:
`value + matrix[index - 1][capacity - weight]`

Choosing if we are going to take the item
`max(excluding, including)`

	0	1	2	3	4
No Items	0	0	0	0	0
Flashlight	0	0	3	3	3
Laptop	0	0	3	3	4
Book	0	3	3	6	6

The "Knapsack" Problem – Example



Value: 3

Weight: 2



Value: 1

Weight: 2



Value: 3

Weight: 1

Maximum Value: 6



Max Capacity: 4

	0	1	2	3	4
No Items	0	0	0	0	0
	0	0	3	3	3
	0	0	3	3	4
	0	3	3	6	6

The "Knapsack" Problem – Example



Weight: 2



Weight: 2



Weight: 1

Maximum Value: 6



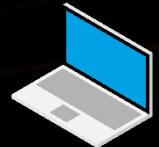
Max Capacity: 4

	0	1	2	3	4
No Items					
	✗	✓	✓	✓	
	✗	✗	✗	✓	
	✓	✗	✓	✓	✓

The "Knapsack" Problem – Example



Weight: 2



Weight: 2



Weight: 1

Maximum Value: 6



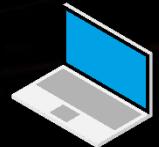
Max Capacity: 4

	0	1	2	3	4
No Items					
	✗	✓	✓	✓	✓
	✗	✗	✗	✓	
	✓	✗	✓	✓	○

The "Knapsack" Problem – Example



Weight: 2



Weight: 2



Weight: 1

Maximum Value: 6



Max Capacity: 4

	0	1	2	3	4
No Items					
	✗	✓	✓	✓	✓
	✗	✗	✗	✗	✓
	✓	✗	✓	✓	✓

The "Knapsack" Problem – Example



Weight: 2



Weight: 2



Weight: 1

Maximum Value: 6



Max Capacity: 4

0	1	2	3	4
No Items				
	✗	✓	✓	✓
✗		✗	✗	✓
✓		✓	✗	✓

The "Knapsack" Problem – Example



Weight: 2



Weight: 2



Weight: 1

Maximum Value: 6



Max Capacity: 4

	0	1	2	3	4
No Items					
✓	✗	✓			✓
✗	✗	✗			✓
✓	✓	✗	✓		

The "Knapsack" Problem – Example



Max Capacity: 4

Maximum Value: 6

With Items:



Total Capacity: 3

The "Knapsack" Problem – Variants

- Many variants are available for the "Knapsack" problem
 - E.g. take each item at most once / use each item multiple times
 - Integer knapsack problem
 - Items have integer weights
 - Efficient dynamic programming solution exists $O(n*c)$ where
 - **n** – number of items
 - **c** – the knapsack capacity
 - Non-integer knapsack problem → no efficient solution
 - Bin packing problem – m knapsacks of capacities (c_1, \dots, c_m)

0/1 Knapsack – Iterative Implementation

```
public static List<Item> FillKnapsack(Item[] items, int capacity)
{
    int[,] maxValues = new int[items.Length + 1, capacity + 1];
    bool[,] itemIncluded = new bool[items.Length + 1, capacity + 1];

    // Calculate maxValues

    // Return the items in the knapsack
}
```

Knapsack – Iterative Implementation (2)

```
for (int i = 0; i < items.Length; i++)
{
    for (int currCapacity = 1; currCapacity <= capacity; currCapacity++)
    {
        if (items[i].Weight > currCapacity) continue;
        int valueIncluded = items[i].Value
            + maxValues[i, currCapacity - items[i].Weight];
        if (valueIncluded > maxValues[i, currCapacity])
        {
            maxValues[i + 1, currCapacity] = valueIncluded;
            itemIncluded[i + 1, currCapacity] = true;
        }
        else
            maxValues[i + 1, currCapacity] = maxValues[i, currCapacity];
    }
}
```

Knapsack – Iterative Implementation (3)

```
// return the taken items in the knapsack
List<Item> takenItems = new List<Item>();
for (int i = items.Length; i > 0; i--)
{
    if (!itemIncluded[i, capacity]) continue;

    Item item = items[i - 1];
    takenItems.Add(item);

    capacity -= item.Weight;
}

return takenItems;
```



The 0/1 Knapsack Problem

Live Coding (Lab)



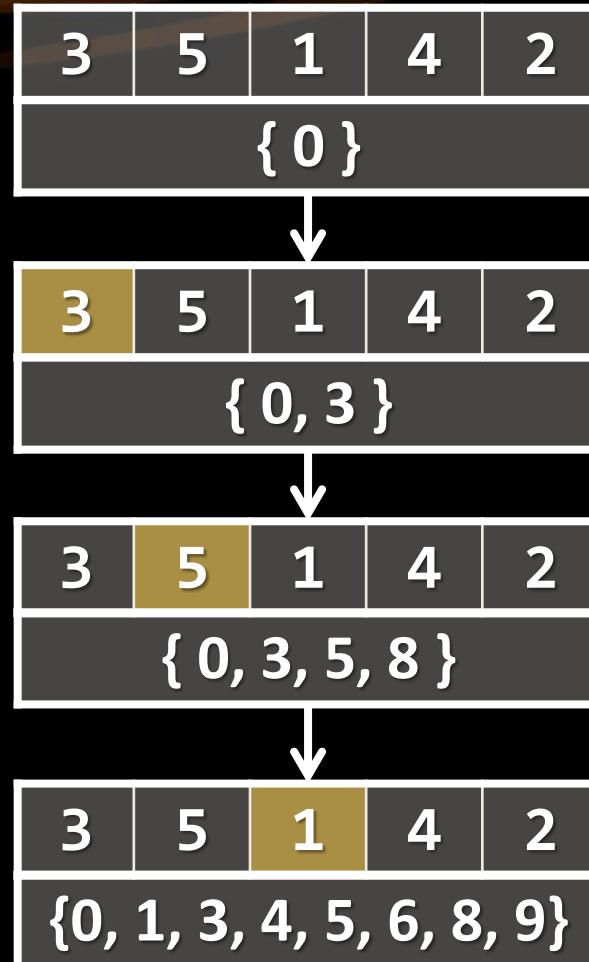
Subset Sum

Sum with Limited Coins

Subset Sum Problem and Its Variations

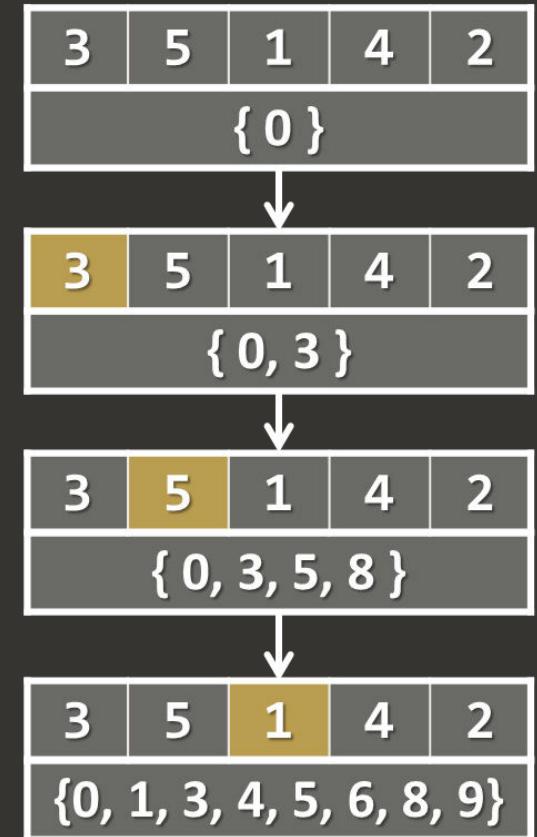
- Subset sum problem (zero subset sum problem)
 - Given a set of integers, find a non-empty subset whose sum 0
 - E.g. {8, 3, -50, 1, -2, -1, 15, -2} → {3, 1, -2, -2}
 - Given a set of integers and an integer **S**,
find a subset whose sum is **S**
 - E.g. {8, 3, 2, 1, 12, 1}, **S=16** → {3, 1, 12}
- Given a set of integers, find all possible sums
- Can you split given sum into set of coins?

Subset Sum Problem (No Repeats)

- Solving the subset sum problem:
 - **nums** = { 3, 5, 1, 4, 2 }, **targetSum** = 6
 - Start with **possibleSums** = { 0 }
 - Step 1: obtain all possible sums ending at { 3 }
 - **possibleSums** = { 0 } \cup { 0+3 } = { 0, 3 }
 - Step 2: obtain all possible sums ending at { 5 }
 - **possibleSums** = { 0, 3 } \cup { 0+5, 3+5 } = { 0, 3, 5, 8 }
 - Step 3: obtain all possible sums ending at { 1 }
 - **possibleSums** = { 0, 3, 5, 8 } \cup { 0+1, 3+1, 5+1, 8+1 } = { 0, 1, 3, 4, 5, 6, 8, 9 }
- 

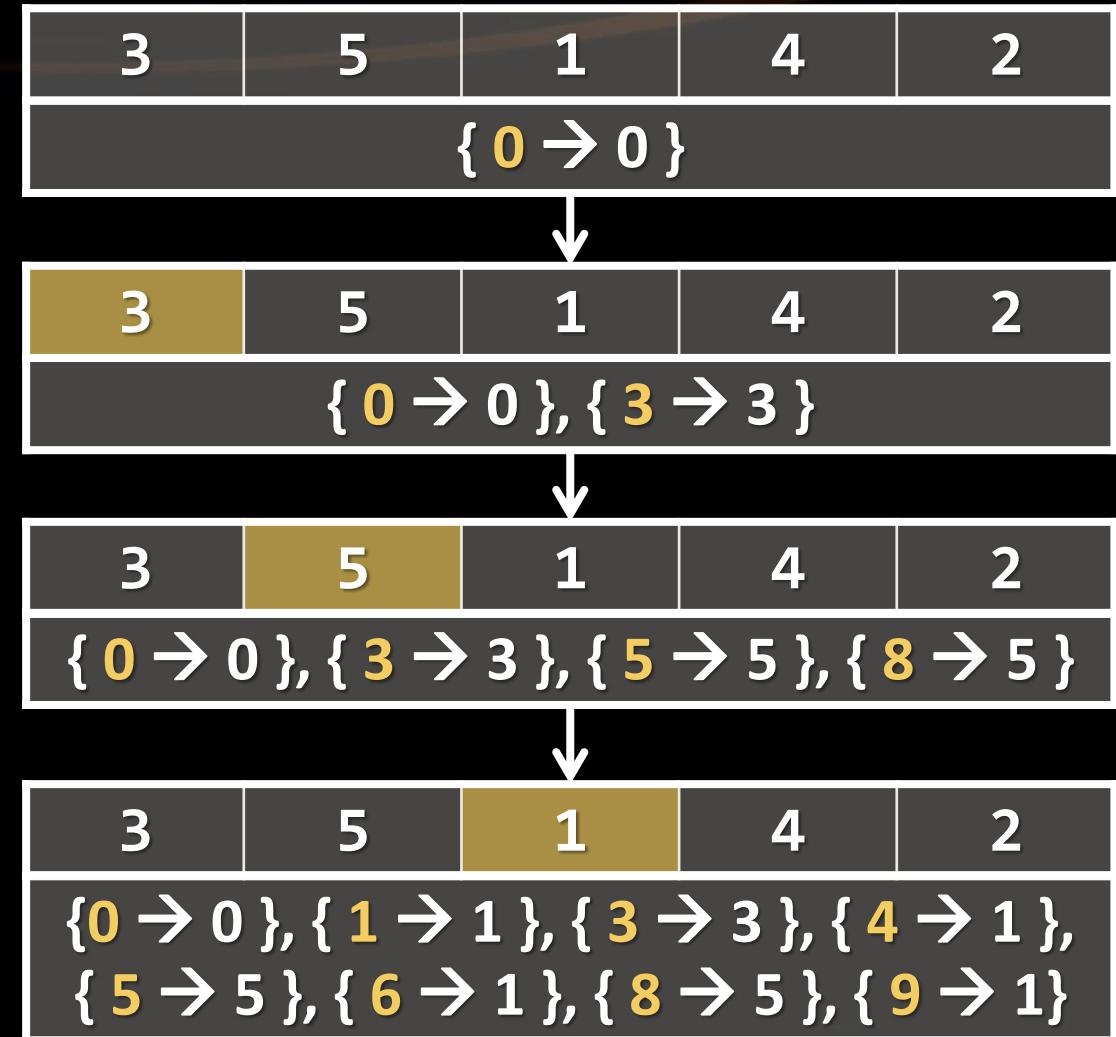
Subset Sum Problem (No Repeats)

```
static ISet<int> CalcPossibleSumsSet(int[] nums, int targetSum)
{
    var possibleSums = new HashSet<int>() { 0 };
    for (int i = 0; i < nums.Length; i++)
    {
        var newSums = new HashSet<int>();
        foreach (var sum in possibleSums)
        {
            int newSum = sum + nums[i];
            newSums.Add(newSum);
        }
        possibleSums.UnionWith(newSums);
    }
    return possibleSums;
}
```



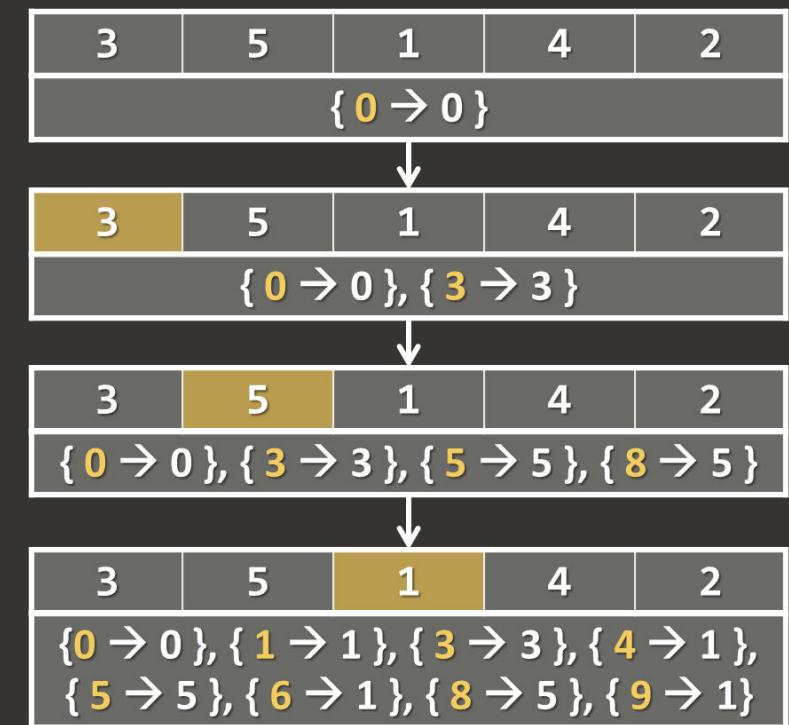
Subset Sum: How to Recover the Subset?

- Keep for each obtained sum in **possibleSums** how it is obtained
- Use a **dictionary** instead of set:
 - **possibleSums[s] → num**
 - The sum **s** is obtained by adding **num** to some previously obtained subset sum
 - **s - num** gives us the previous sum



Subset Sum (No Repeats + Subset Recovery)

```
static IDictionary<int, int> CalcPossibleSums(  
    int[] nums, int targetSum) {  
    var possibleSums = new Dictionary<int, int>();  
    possibleSums.Add(0, 0);  
    for (int i = 0; i < nums.Length; i++) {  
        var newSums = new Dictionary<int, int>();  
        foreach (var sum in possibleSums.Keys) {  
            int newSum = sum + nums[i];  
            if (!possibleSums.ContainsKey(newSum))  
                newSums.Add(newSum, nums[i]);  
        }  
        foreach (var sum in newSums)  
            possibleSums.Add(sum.Key, sum.Value);  
    }  
    return possibleSums;  
}
```



Subset Sum (No Repeats): Subset Recovery

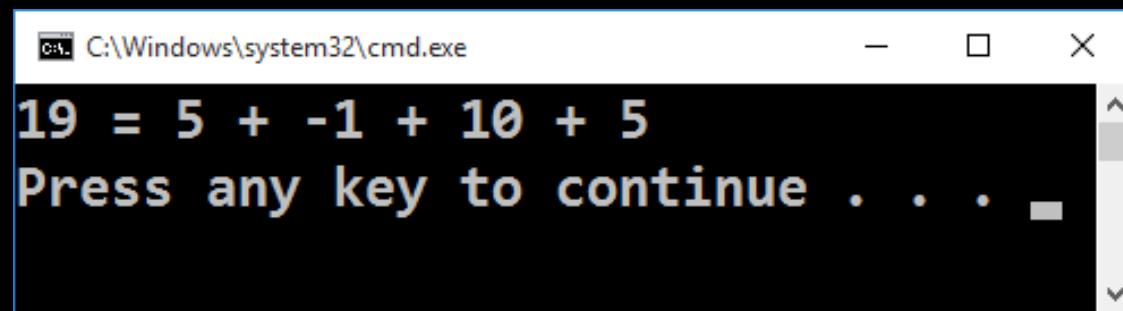
```
static List<int> FindSubset(int[] nums, int targetSum,  
IDictionary<int, int> possibleSums)  
{  
    var subset = new List<int>();  
    while (targetSum > 0)  
    {  
        var lastNum = possibleSums[targetSum];  
        subset.Add(lastNum);  
        targetSum -= lastNum;  
    }  
    subset.Reverse();  
    return subset;  
}
```

{ 9 → 1 } → { 8 → 5 } → { 3 → 3 } → 0

```
int[] nums = { 3, 5, -1, 10, 5, 7 };
int targetSum = 19;

var possibleSums = CalcPossibleSums(nums, targetSum);

var subset = FindSubset(nums, targetSum, possibleSums);
Console.WriteLine(targetSum + " = ");
Console.WriteLine(String.Join(" + ", subset));
```



The screenshot shows a Windows Command Prompt window titled 'C:\Windows\system32\cmd.exe'. The window contains the following text:
19 = 5 + -1 + 10 + 5
Press any key to continue . . .

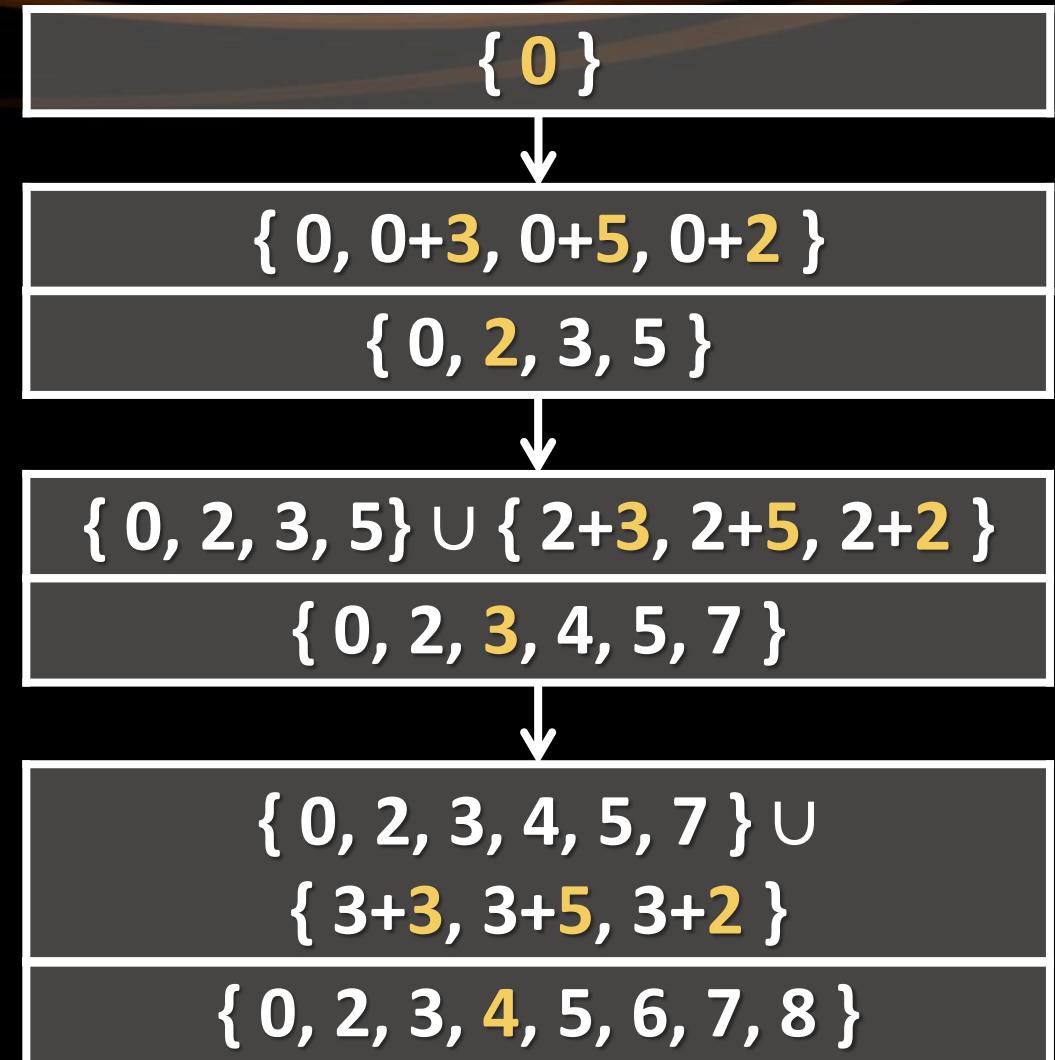
Subset Sum (with No Repeats)

Live Demo

Subset Sum Problem (With Repeats)

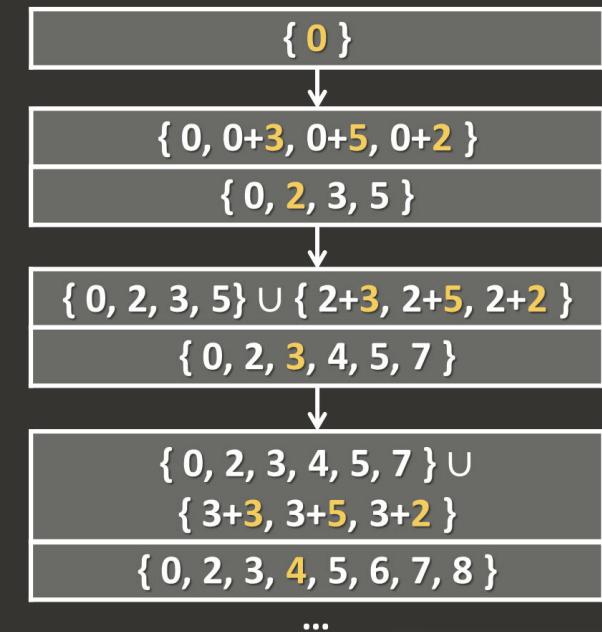
- Solving the subset sum problem (BFS-based algorithm):
 - nums** = { 3, 5, 2 }, **targetSum** = 6

```
possibleSums = { 0 };
for i = 0 ... targetSum
{
    Append to possibleSums[i]
    all numbers from nums[]
}
```



Subset Sum (With Repeats)

```
static bool[] CalcPossibleSums(int[] nums, int targetSum)
{
    var possible = new bool[targetSum + 1];
    possible[0] = true;
    for (int sum = 0; sum < possible.Length; sum++)
        if (possible[sum])
            for (int i = 0; i < nums.Length; i++)
            {
                int newSum = sum + nums[i];
                if (newSum <= targetSum)
                    possible[newSum] = true;
            }
    return possible;
}
```



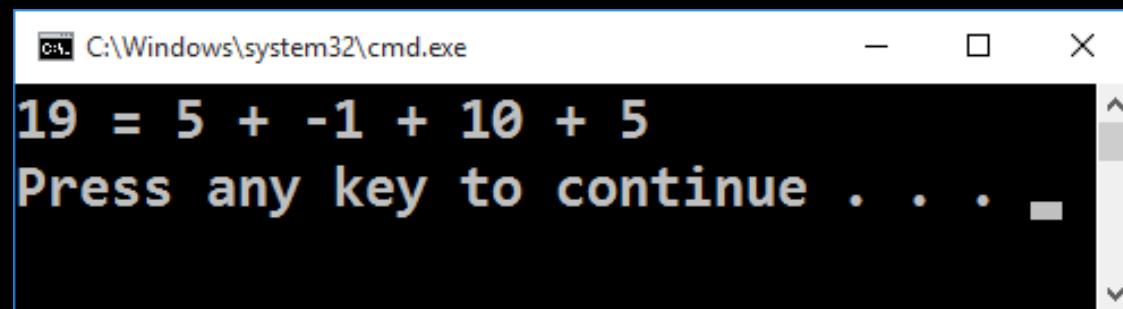
Subset Sum (With Repeats): Recovery

```
private static IEnumerable<int> FindSubset(  
    int[] nums, int targetSum, bool[] possibleSums)  
{  
    var subset = new List<int>();  
    while (targetSum > 0)  
        for (int i = 0; i < nums.Length; i++)  
        {  
            int newSum = targetSum - nums[i];  
            if (newSum >= 0 && possibleSums[newSum])  
            {  
                targetSum = newSum;  
                subset.Add(nums[i]);  
            }  
        }  
    return subset;  
}
```

```
int[] nums = { 3, 5, -1, 10, 5, 7 };
int targetSum = 19;

var possibleSums = CalcPossibleSums(nums, targetSum);

var subset = FindSubset(nums, targetSum, possibleSums);
Console.WriteLine(targetSum + " = ");
Console.WriteLine(String.Join(" + ", subset));
```

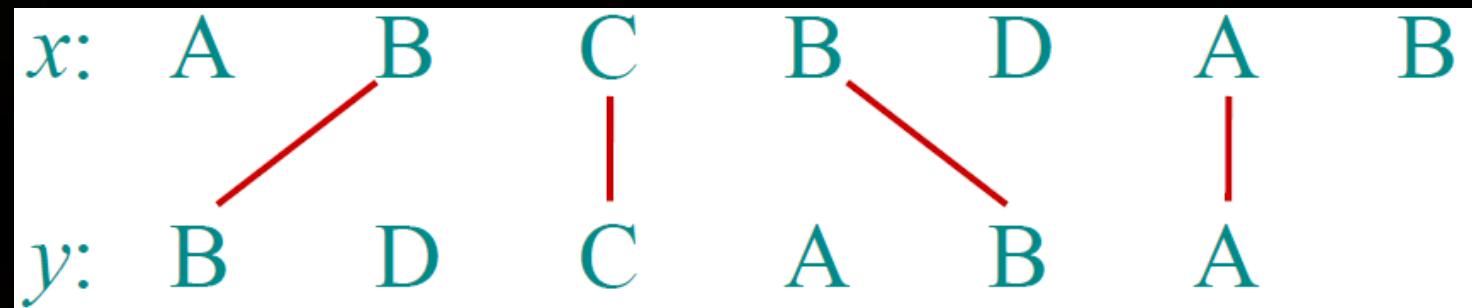
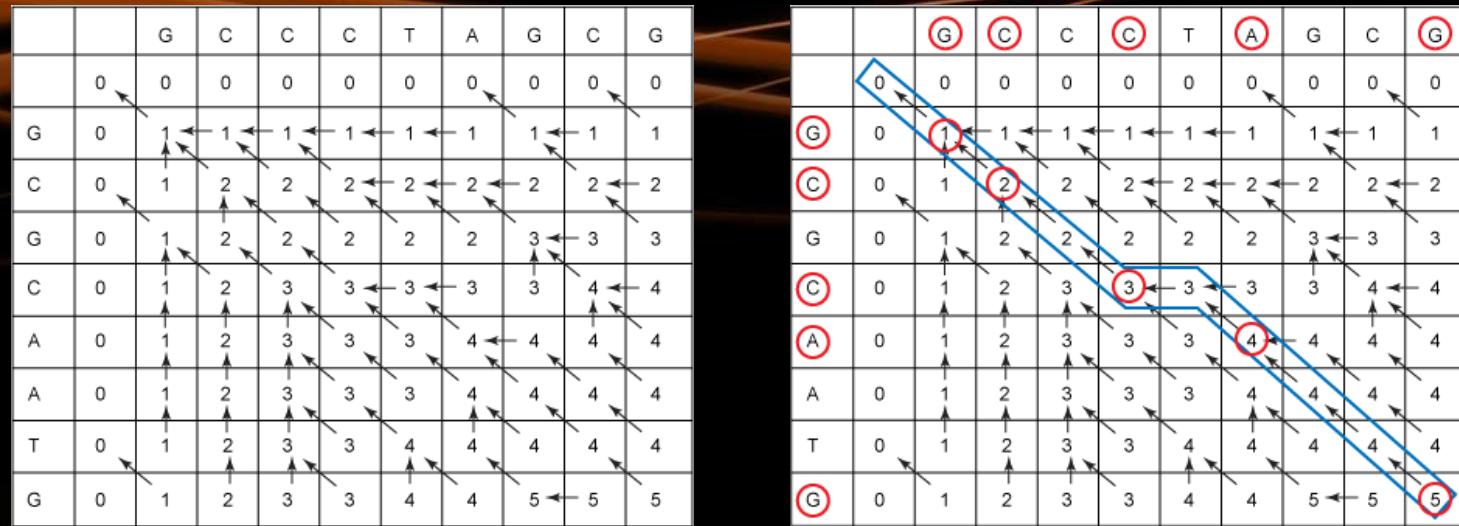


C:\Windows\system32\cmd.exe

```
19 = 5 + -1 + 10 + 5
Press any key to continue . . .
```

Subset Sum (With Repeats)

Live Demo

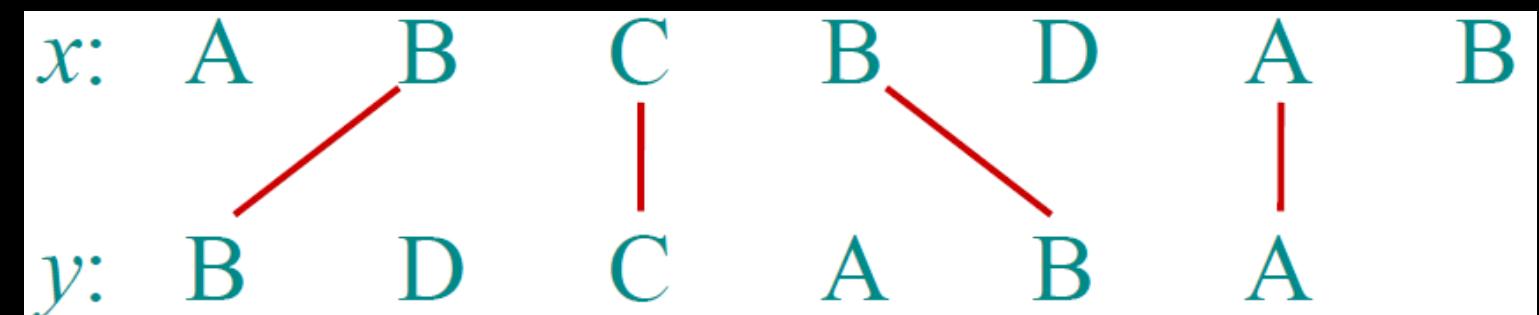


Longest Common Subsequence (LCS)

A Recursive DP Approach

Longest Common Subsequence (LCS)

- Longest common subsequence (LCS) problem:
 - Given two sequences $x[1 \dots m]$ and $y[1 \dots n]$
 - Find a longest common subsequence (LCS) to them both
- Example:
 - $x = "A\textcolor{blue}{B}C\textcolor{blue}{B}D\textcolor{blue}{A}B"$
 - $y = "\textcolor{blue}{B}DC\textcolor{blue}{A}B\textcolor{blue}{A}"$
 - $LCS = "\textcolor{blue}{B}C\textcolor{blue}{B}A"$



LCS – Recursive Approach

- $S_1 = \text{GCCCTAGCG}$, $S_2 = \text{GCGCAATG}$
 - Let $C_1 =$ the right-most character of S_1 ($C_1 = G$)
 - Let $C_2 =$ the right-most character of S_2 ($C_2 = G$)
 - Let $S_1' = S_1$ with C_1 "chopped-off" ($S_1' = \text{GCCCTAGC}$)
 - Let $S_2' = S_2$ with C_2 "chopped-off" ($S_2' = \text{GCGCAAT}$)
- There are three recursive sub-problems:
 - $L_1 = \text{LCS}(S_1', S_2)$
 - $L_2 = \text{LCS}(S_1, S_2')$
 - $L_3 = \text{LCS}(S_1', S_2')$

LCS – Recursive Formula

- Let $\text{lcs}[x, y]$ be the longest common subsequence of $S_1[0 \dots x]$ and $S_2[0 \dots y]$
- LCS has the following recursive properties:

```
lcs[-1, y] = 0
```

```
lcs[x, -1] = 0
```

```
lcs[x, y] = max(
```

```
    lcs[x-1, y],
```

```
    lcs[x, y-1],
```

```
    lcs[x-1, y-1]+1 when S1[x] == S2[y]
```

```
)
```

G	C	C	C	T	A	G	C	G
0	0	0	0	0	0	0	0	0
G	0	1	1	1	1	1	1	1
C	0	1	2	2	2	2	2	2
G	0	1	2	2	2	2	3	3
C	0	1	2	3	3	3	3	4
A	0	1	2	3	3	4	4	4
A	0	1	2	3	3	4	4	4
T	0	1	2	3	3	4	4	4
G	0	1	2	3	3	4	4	5

Calculating the LCS Table

```
static int CalcLCS(int x, int y) {  
    if (x < 0 || y < 0)  
        return 0;  
    if (lcs[x, y] == NOT_CALCULATED) {  
        int lcsFirstMinusOne = CalcLCS(x - 1, y);  
        int lcsSecondMinusOne = CalcLCS(x, y - 1);  
        lcs[x, y] = Math.Max(  
            lcsFirstMinusOne, lcsSecondMinusOne);  
        if (firstStr[x] == secondStr[y]) {  
            int lcsBothMinusOne = 1 + CalcLCS(x - 1, y - 1);  
            lcs[x, y] = Math.Max(lcs[x, y], lcsBothMinusOne);  
        }  
    }  
    return lcs[x, y];  
}
```

G	C	C	C	T	A	G	C	G
0	0	0	0	0	0	0	0	0
G	0	1	1	1	1	1	1	1
C	0	1	2	2	2	2	2	2
G	0	1	2	2	2	2	3	3
C	0	1	2	3	3	3	3	4
A	0	1	2	3	3	3	4	4
A	0	1	2	3	3	3	4	4
T	0	1	2	3	3	4	4	4
G	0	1	2	3	3	4	4	5

Reconstructing the LCS Sequence

```
static string PrintLCS(int x, int y) {  
    var lcsLetters = new List<char>();  
    while (x >= 0 && y >= 0)  
        if ((firstStr[x] == secondStr[y]) &&  
            (CalcLCS(x - 1, y - 1) + 1 == lcs[x, y])) {  
            lcsLetters.Add(firstStr[x]);  
            x--;  
            y--;  
        }  
        else if (CalcLCS(x - 1, y) == lcs[x, y])  
            x--;  
        else  
            y--;  
    lcsLetters.Reverse();  
    return string.Join("", lcsLetters);  
}
```

	G	C	C	C	T	A	G	C	G
0	0	0	0	0	0	0	0	0	0
G	0	1	1	1	1	1	1	1	1
C	0	1	2	2	2	2	2	2	2
G	0	1	2	2	2	2	2	3	3
C	0	1	2	3	3	3	3	3	4
A	0	1	2	3	3	3	4	4	4
A	0	1	2	3	3	3	4	4	4
T	0	1	2	3	3	4	4	4	4
G	0	1	2	3	3	4	4	5	5

	G	C	C	C	T	A	G	C	G
0	0	0	0	0	0	0	0	0	0
G	0	1	1	1	1	1	1	1	1
C	0	1	2	2	2	2	2	2	2
G	0	1	2	2	2	2	2	3	3
C	0	1	2	3	3	3	3	4	4
A	0	1	2	3	3	3	4	4	4
A	0	1	2	3	3	3	4	4	4
T	0	1	2	3	3	4	4	4	4
G	0	1	2	3	3	4	4	5	5

Longest Common Subsequence (LCS)

Live Coding (Lab)

Summary

- DP → Solve a problem by **solving overlapping subproblems**
- Memoization → Save subproblem **solutions** for later use
- Optimal Substructure
 - Subproblems should have **optimal solutions**
 - Combine optimal solutions for subproblems
 - Get optimal solution for original problem
- Top down approach – **Recursive**
- Bottom up approach – **Iterative**



Dynamic Programming



Questions?



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