

Tutorial 1: Propositional Logics

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(1) if ((i+j ≤ p+q) && (i ≤ p) &&  
    ((j > q) || (List1[i] ≤ List2[j])))  
(2)   List3[k] = List1[i]  
(3)   i = i+1  
(4) else  
(5)   List3[k] = List2[j]  
(6)   j = j+1  
(7) k = k+1
```

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(1) if (((i+j ≤ p+q) && (i ≤ p) &&  
    (j > q)) || ((i+j ≤ p+q) && (i ≤  
    p) && (List1[i] ≤ List2[j])))  
(2)   List3[k] = List1[i]  
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(4) else  
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(6)   j = j+1  
(7) k = k+1
```

Consider the above two pieces of code. They are taken from two different versions of *Mergesort*. Do they do the same thing? (Note that $\&\&$ = “and”, $\|\|$ = “or”.)

Observation: The two pieces of code are the same except for line 1.

(1) $((i+j \leq p+q) \ \&\& \ (i \leq p) \ \&\& \ ((j > q) \ || \ (List1[i] \leq List2[j])))$

(1') $((((i+j \leq p+q) \ \&\& \ (i \leq p) \ \&\& \ (j > q)) \ || \ ((i+j \leq p+q) \ \&\& \ (i \leq p) \ \&\& \ (List1[i] \leq List2[j]))))$

Are they equivalent?

Let's rewrite using

$s \sim (i+j \leq p+q) \quad t \sim (i \leq p) \quad u \sim (j > q) \quad v \sim (List[i] \leq List2[j])$

(1) $s \text{ and } t \text{ and } (u \text{ or } v)$

(1') $(s \text{ and } t \text{ and } u) \text{ or } (s \text{ and } t \text{ and } v)$

Now set $w \sim (s \text{ and } t)$

(1) $w \text{ and } (u \text{ or } v) \xleftrightarrow{\text{equal?}} (1') (w \text{ and } u) \text{ or } (w \text{ and } v)$

We just transformed our code into **compound propositions** and **next** want to develop a theory of how to determine whether two such propositions are equal (equivalent).

Question 1

Are $s \oplus t$ and $(s \wedge \neg t) \vee (\neg s \wedge t)$ logically equivalent? If yes, prove it by truth table. Otherwise, give a counterexample to show that they are not logically equivalent.

Solution:

s	t	$s \oplus t$	$s \wedge \neg t$	$\neg s \wedge t$	$(s \wedge \neg t) \vee (\neg s \wedge t)$
T	T	F	F	F	F
T	F	T	T	F	T
F	T	T	F	T	T
F	F	F	F	F	F

Question 2

Are $w \wedge (s \oplus t)$ and $(w \vee s) \oplus (w \vee t)$ logically equivalent? If yes, prove it by truth table. Otherwise, give a counterexample to show that they are not logically equivalent.

Solution:

No, they are not logically equivalent. Set $w = F$, $s = F$ and $t = T$. $w \wedge (s \oplus t) = F$ but $(w \vee s) \oplus (w \vee t) = T$.

Question 3

For each of the following pairs of propositions, either prove they are logically equivalent, or give a counterexample.

- (a) $p \rightarrow q$ and $q \rightarrow p$
- (b) $p \rightarrow q$ and $\neg q \rightarrow \neg p$
- (c) $(p \rightarrow q) \rightarrow q$ and $p \rightarrow q$
- (d) $(p \rightarrow p) \rightarrow q$ and q
- (e) $q \rightarrow (p \wedge \neg p)$ and $\neg q$
- (f) $q \rightarrow (p \wedge \neg p)$ and $\neg p$

Question 3

Solution:

- (a) Not equivalent. Set $p = T, q = F$, then $p \rightarrow q = F$ but $q \rightarrow p = T$.
- (b) Equivalent. $p \rightarrow q \equiv \neg p \vee q \equiv (\neg\neg q) \vee (\neg p) \equiv \neg q \rightarrow \neg p$.
Or truth table.

p	q	$p \rightarrow q$	$\neg q$	$\neg p$	$\neg q \rightarrow \neg p$
T	T	T	F	F	T
T	F	F	T	F	F
F	T	T	F	T	T
F	F	T	T	T	T

- (c) Not equivalent. Set $p = F, q = F$, then $(p \rightarrow q) \rightarrow q = F$ but $p \rightarrow q = T$

Question 3

(d) Equivalent. $(p \rightarrow p) \rightarrow q \equiv T \rightarrow q \equiv q$. Or truth table.

p	q	$p \rightarrow p$	$(p \rightarrow p) \rightarrow q$	q
T	T	T	T	T
T	F	T	F	F
F	T	T	T	T
F	F	T	F	F

(e) Equivalent. $q \rightarrow (p \wedge \neg p) \equiv q \rightarrow F \equiv \neg q$. Or truth table.

p	q	$p \wedge \neg p$	$q \rightarrow (p \wedge \neg p)$	$\neg q$
T	T	F	F	F
T	F	F	T	T
F	T	F	F	F
F	F	F	T	T

(f) Not equivalent. Set $p = T, q = F$, then $q \rightarrow (p \wedge \neg p) = T$ but $\neg p = F$

Question 4

For each of the following pair of propositions, either prove that the two propositions are logically equivalent without using truth table, or give a counterexample to show that the propositions are not logically equivalent.

(a) (i) $p \rightarrow (q \rightarrow r)$
(ii) $(p \wedge q) \rightarrow r$

(b) (i) $(\neg(p \rightarrow \neg q)) \rightarrow (\neg(p \wedge r) \vee s)$
(ii) $(\neg r \rightarrow p) \rightarrow ((p \wedge q) \rightarrow s)$

Question 4

Solution:

(a) Consider

$$\begin{aligned} p \rightarrow (q \rightarrow r) &= \neg p \vee (q \rightarrow r) \\ &= \neg p \vee (\neg q \vee r) \\ &= \neg p \vee \neg q \vee r \end{aligned}$$

Consider

$$\begin{aligned} (p \wedge q) \rightarrow r &= \neg(p \wedge q) \vee r \\ &= (\neg p \vee \neg q) \vee r \\ &= \neg p \vee \neg q \vee r \end{aligned}$$

Thus, $p \rightarrow (q \rightarrow r)$ is equivalent to $(p \wedge q) \rightarrow r$.

(b) Set $p = T, q = T, r = F, s = F$,

$(\neg(p \rightarrow \neg q)) \rightarrow (\neg(p \wedge r) \vee s)$ becomes true and

$(\neg r \rightarrow p) \rightarrow ((p \wedge q) \rightarrow s)$ becomes false. Thus they are not equivalent.