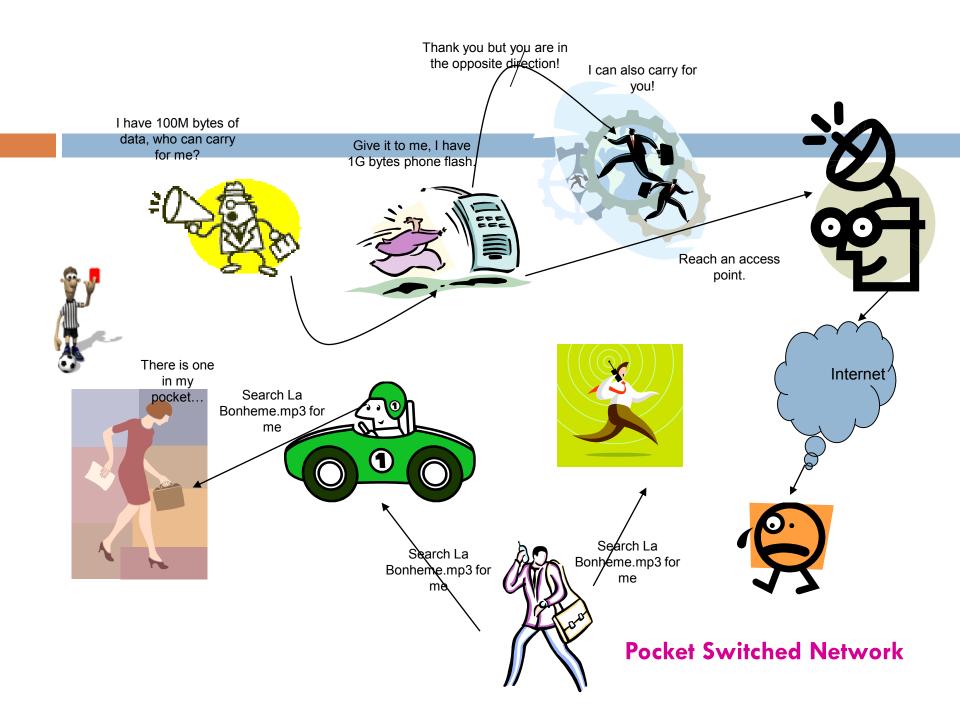
LECTURE 14: MOBILITY AND OPPORTUNISTIC NETWORKING

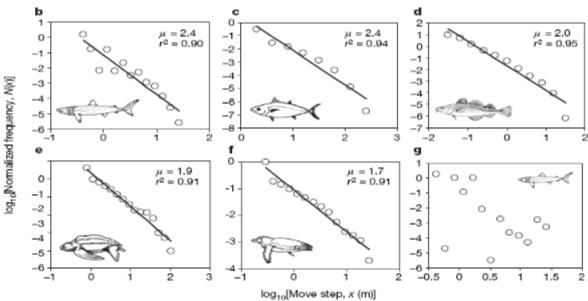


People Are the Network

Mobility Pattern

Pattern: most of the studies on animal mobility pattern including experimental data and theoretic analysis found that their mobility pattern follow the Levy flight:

$$\Pr(d) \propto d^{-2}$$



-0g₁₀ N(x) $r^2 = 0.92$ 40 35 -1.5 Frequency 30 20 -0.5 $Log_{10} x$ 15 10 11 12 13 14 zooplankton (g m⁻³)

45

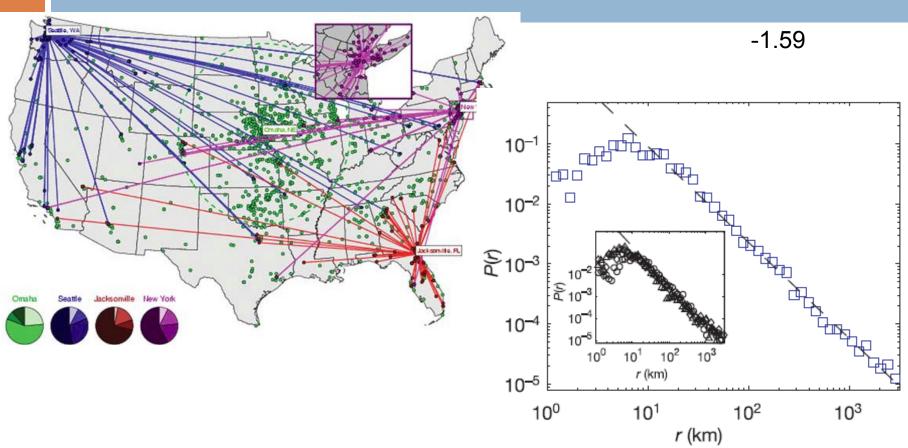
0.5

 $\mu = 1.8$

Levy flight search patterns of wandering albatrosses, Nature 381, (1996) Revisiting Le´vy flight search patterns of wandering albatrosses, bumblebees and deer, NATURE| Vol 449|25 October 2007

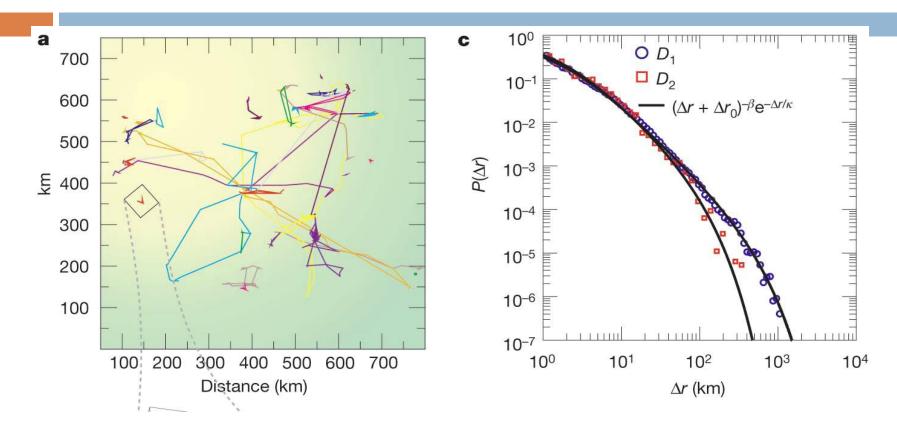
Scaling laws of marine predator search behaviour, Nature (2008)

Human Mobility



D. Brockmann, L. Hufnagel and T. Geisel, The scaling laws of human travel, Nature, 439, 462-465, (2006).

Human Mobility



Understanding individual human mobility patterns Marta C. Gonza´lez, Ce´sar A. Hidalgo & Albert-La´szlo´ Baraba´si, NATURE| Vol 453|5 June 2008

EXPERIMENTAL SETUP

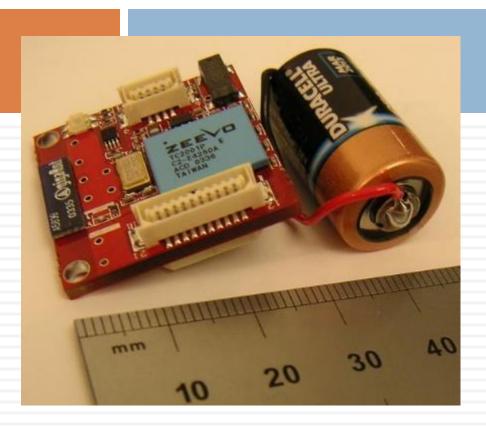
iMotes

ARM processor Bluetooth radio 64k flash memory

Bluetooth Inquiries

5 seconds every 2 minutes Log {MAC address, start time, end time} tuple of each contact

Experimental devices



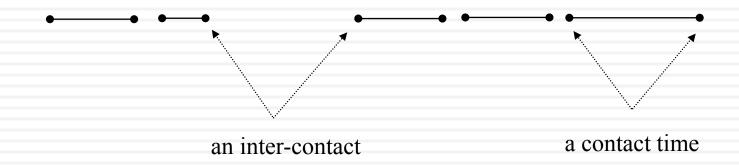


Experimental data set	Cambridge04	Infocom05	Hong Kong	Cambridge05	Infocom06
Device	iMote	iMote	iMote	iMote	iMote
Network type	Bluetooth	Bluetooth	Bluetooth	Bluetooth	Bluetooth
Duration (days)	3	3	5	11	3
Granularity (seconds)	120	120	120	600	120
Number of Experimental Devices	12	41	37	54	98
Number of internal contacts	4,229	22,459	560	10,873	191,336
Average # Contacts/pair/day	10	4.6	0.084	0.345	6.7
Number of External Devices	148	264	868	11,357	14,036
Number of external contacts	2,441	1,173	2,507	30,714	63,244

User Population	Toronto	UCSD	Dartmouth	MIT BT	MIT GSM
Device	PDA	PDA	Laptop/PDA	Cell Phone	Cell Phone
Network type	Bluetooth	WiFi	WiFi	BT	GSM
Duration (days)	16	77	114	246	246
Granularity (seconds)	120	120	300	300	10
Devices participating	23	273	6648	100	100
Number of internal contacts	2,802	195,364	4,058,284	54,667	572,190
Average # Contacts/pair/day	0.35	0.034	0.00080	0.022	0.23
Recorded external devices	N/A	N/A	N/A	N/A	N/A
Number of external contacts	N/A	N/A	N/A	N/A	N/A

PAIR WISE RELATIONSHIP

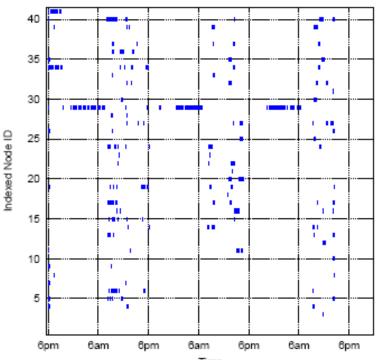
Duration of the experiment



Contacts seen by an iMote

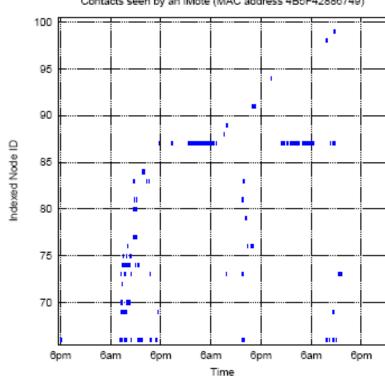
iMotes

Contacts seen by an iMote (MAC address 4B5F42886749)



External Devices

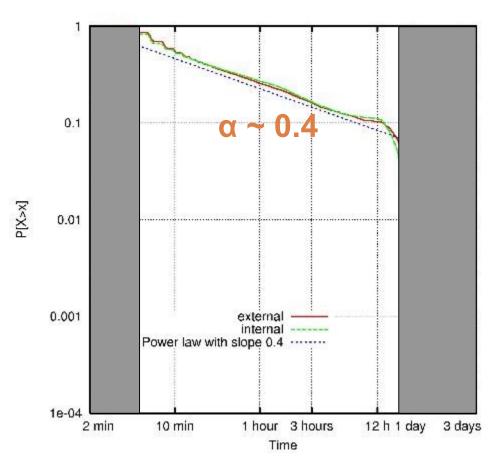




- x-axis shows the time of the day and the y-axis shows the node ID
- Contacts are sparse
- iMote sees node 28 every night and they stay together the whole night ©

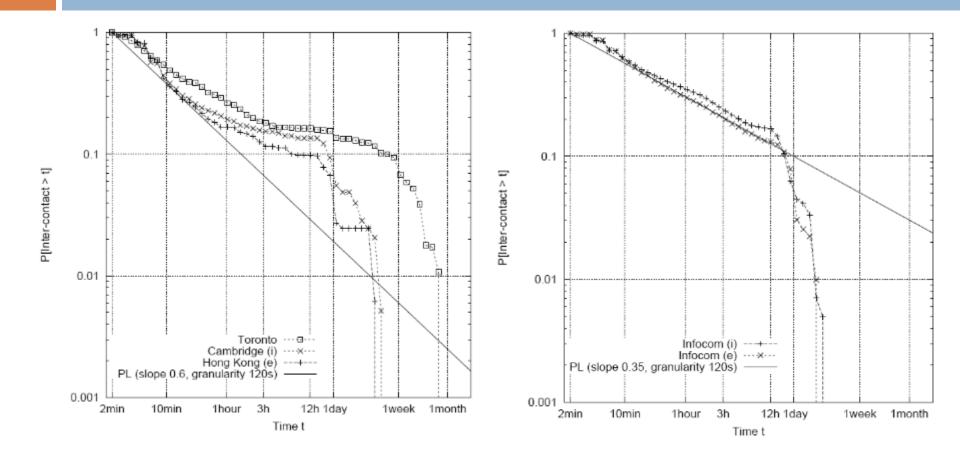
Heavy Tailed Distributions

Inter-contact time distribution



- It compares with a Power law on range [5mn;1day]
 slope α ~ 0.4.
- Internal and external contacts have almost the same law.

OTHER EXAMPLES



Implication on Opportunistic Forwarding:

 \Box For $\alpha > 2$

Any stateless algorithm achieves a finite expected delay.

□ For $\alpha \ge \frac{m+1}{m}$ and $\# \{ \text{ nodes } \} \ge 2m$

There exist a forwarding algorithm with m copies and a finite expected delay.

 \Box For α < 1

No stateless algorithm (even flooding) achieve a bounded delay (Orey's theorem).

Social Structures Vs Network Structures

□ Community structures

- Social communities, i.e. affiliations
- Topological cohesive groups or modules

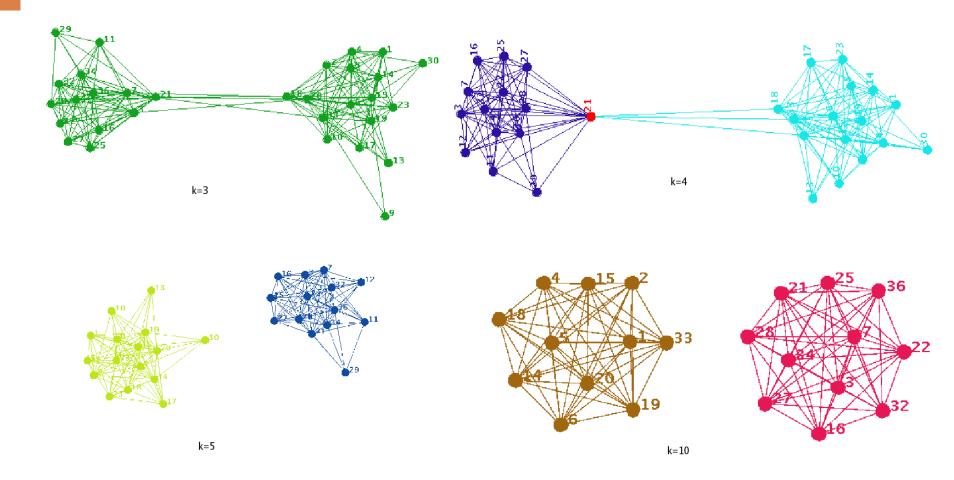
□ Centralities

- Social hubs, celebrities and postman
- Betweenness, closeness, inference power centrality

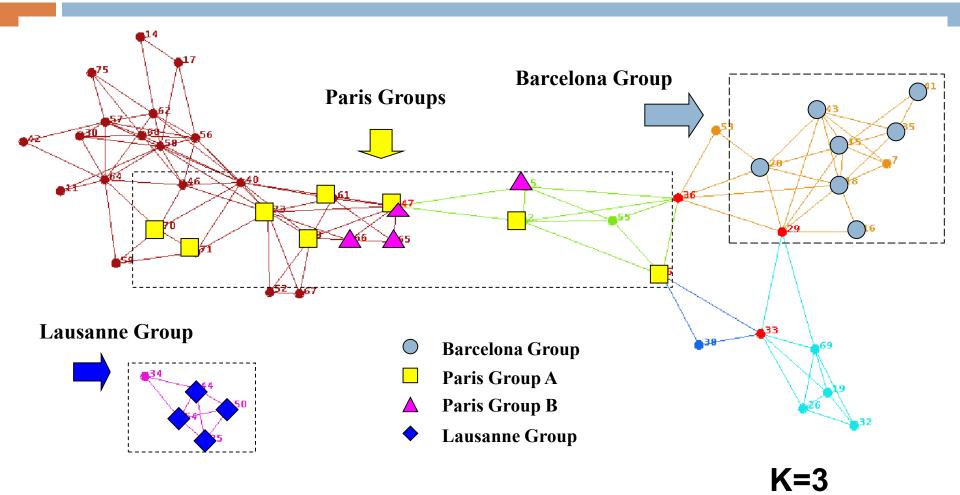
K-clique Community Definition

- Union of k-cliques reachable through a series of adjacent k-cliques [Palla et al]
- □ Adjacent k-cliques share k-1 nodes
- Members in a community reachable through well-connected well subsets
- Examples
 - 2-clique (connected components)
 - 3-clique (overlapping triangles)
 - Overlapping feature

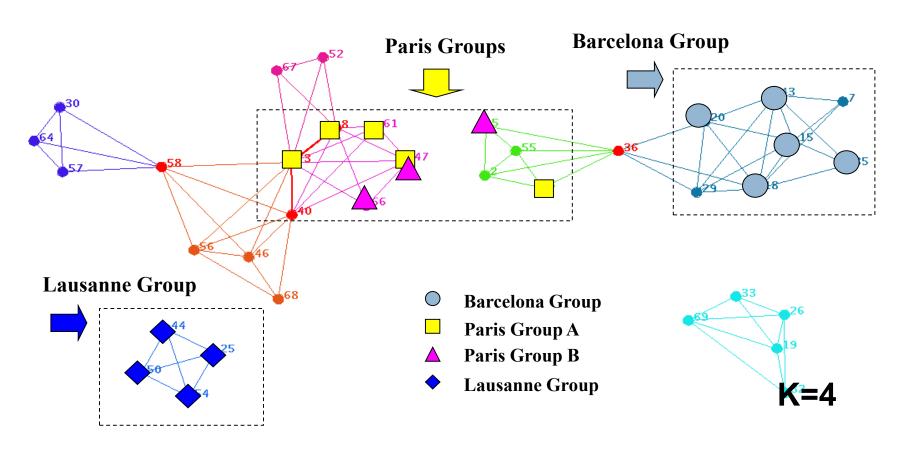
K-clique Communities in Cambridge Dataset



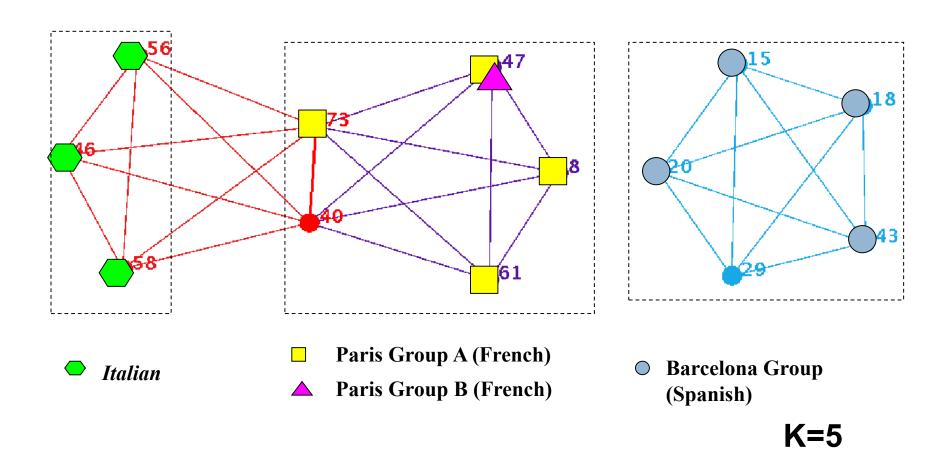
K-clique Communities in Infocom06 Dataset



K-clique Communities in Infocom06 Dataset



K-clique Communities in Infocom06 Dataset



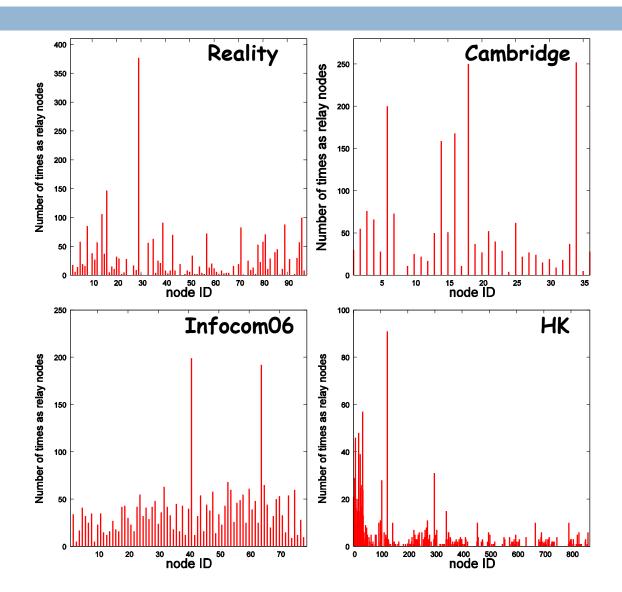
Other Community Detection Methods

- Betweenness [Newman04]
- Modularity [Newman06]
- Information theory[Rosvall06]
- Statistical mechanics[Reichardt]
- Weighted Network Analysis[Newman05]
- Survey Papers[Danon05][Newman04]

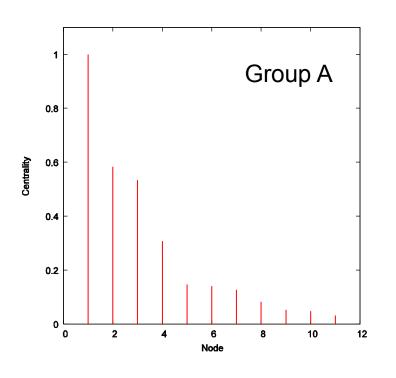
Centrality in Temporal Network

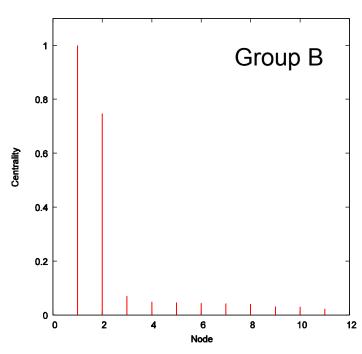
- □ Large number of unlimited flooding
- □ Uniform sourced and temporal traffic distribution
- □ Number of times on shortest delay deliveries
- □ Analogue to Freeman centrality [Freeman]

Homogenous Centrality

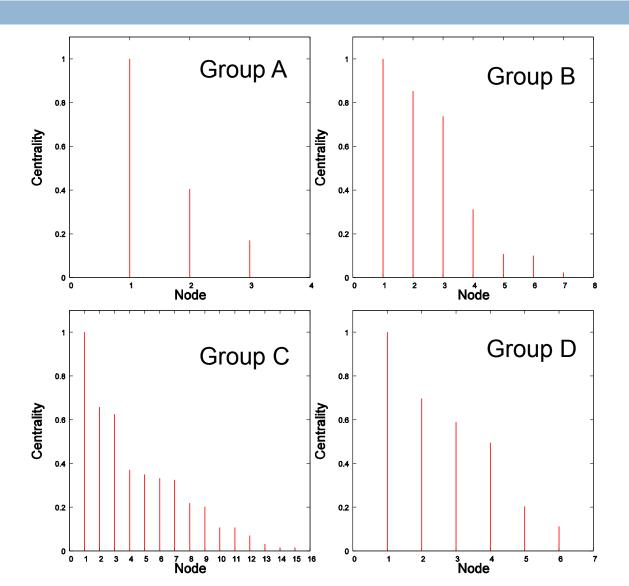


Within Group Centrality Cambridge Dataset

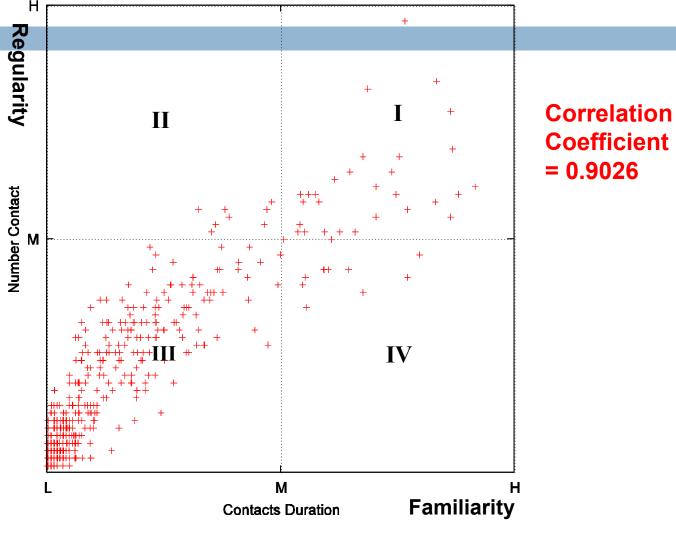




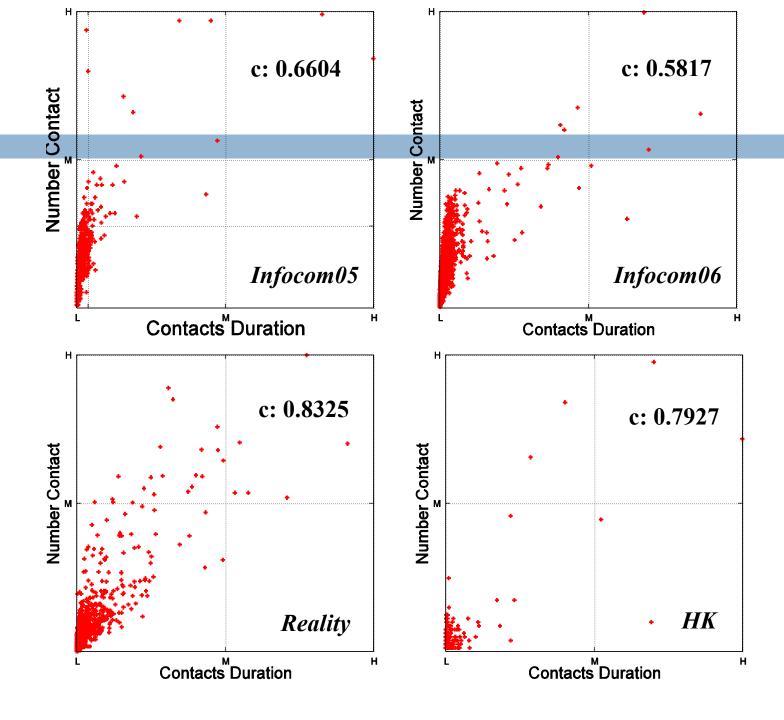
Within Group Centrality Reality Dataset



Regularity and Familiarity



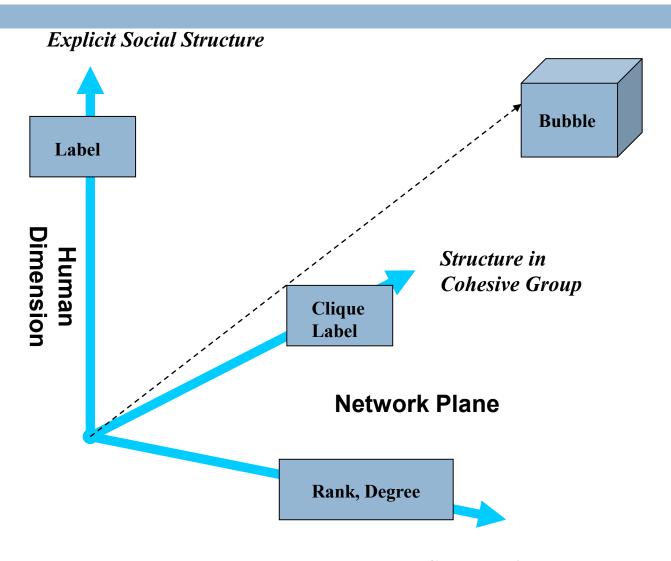
I: Community II. Familiar Strangers III. Strangers IV. Friends



Interaction and Forwarding

- Third generation human interaction model
- Categories of human contact patterns
- Clique and community
- Popularity/Centrality
 - Dual natures of mobile network
- Social network
- Physical network
 - Benchmark Forwarding strategies
- Flooding, Wait, and Multiple-copy-multiple-hop (MCP), PROPHET

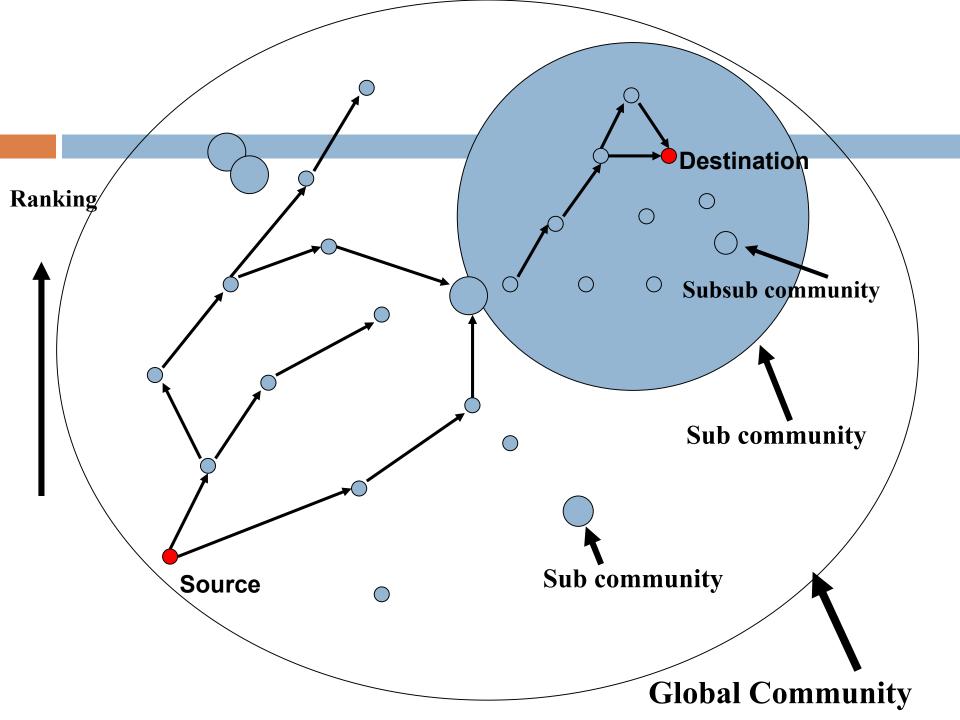
Design Space



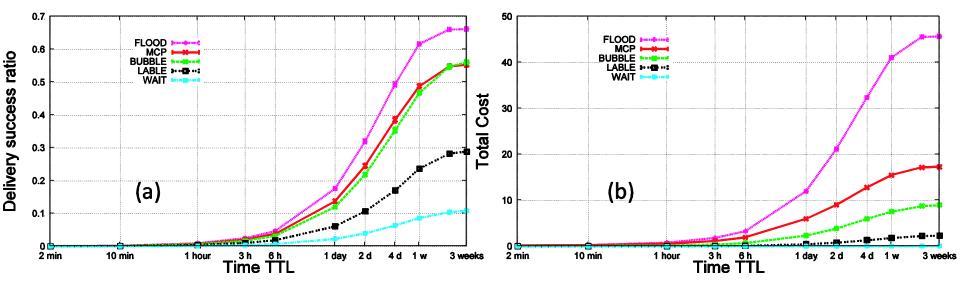
Structure in Degree

Centrality meets Community

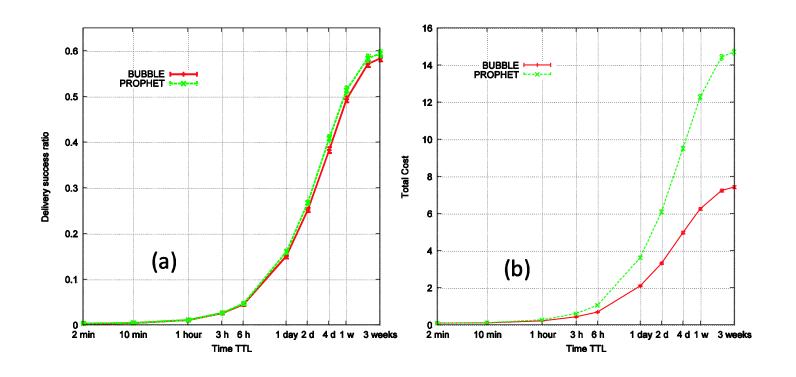
- □ Population divided into communities
- Node has a global and local ranking
- Global popular node like a postman, or politician in a city
- □ Local popular node
- **BUBBLE**



Centrality meets Community



Centrality meets Community

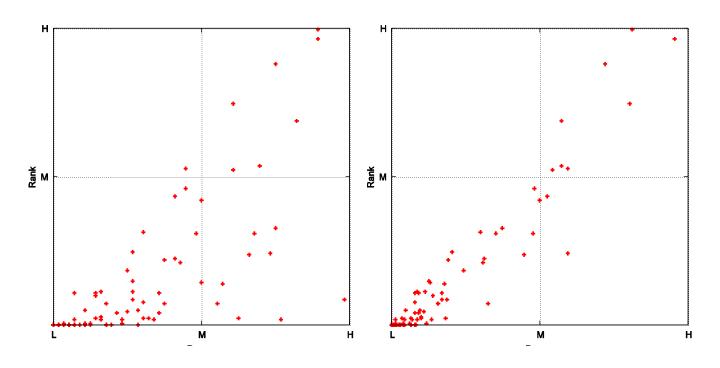


Making Centrality Practical

How can each node know its own centrality in decentralised way?

How well does past centrality predict the future?

Approximating Centrality



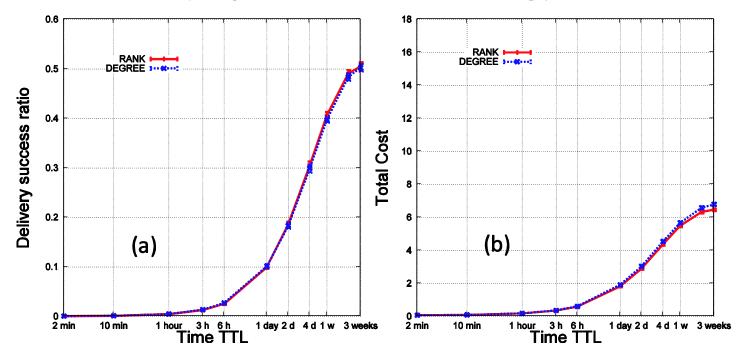
- □ Total degree, per-6-hour degree
- Correlation coefficients, 0.7401 and 0.9511

Approximating Centrality

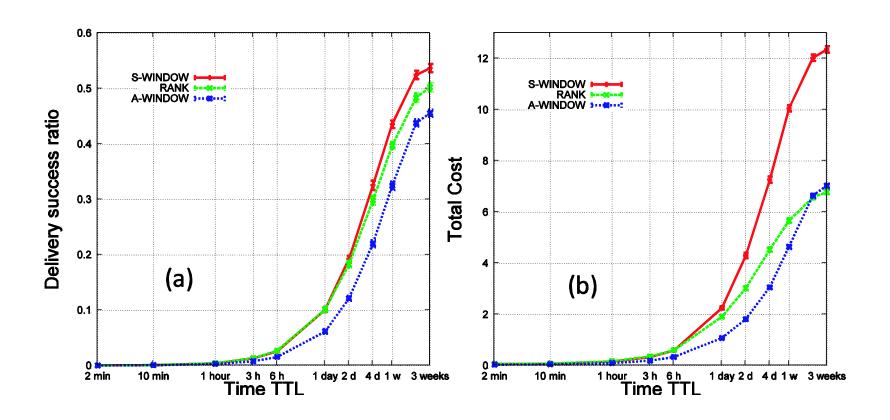
DEGREE

S-Window

A-Window (Exponential Smoothing)



Approximating Centrality

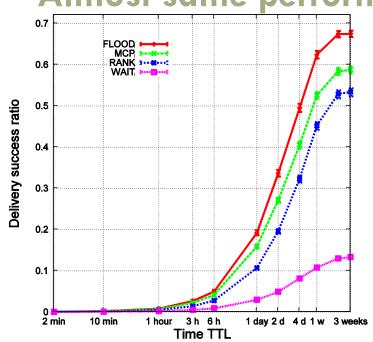


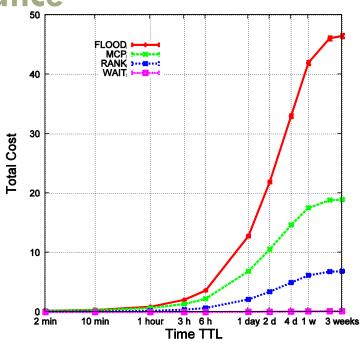
Predictability of Human Mobility

Three sessions of Reality dataset

Two sessions using the ranking calculated from the first session





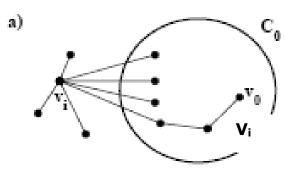


Distributed Community Detection

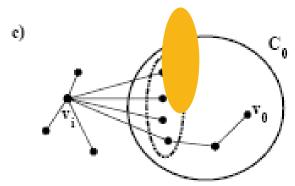
```
SIMPLE, K-CLIQUE, MODULARITY
Terminology: Familiar Set (F), Local Community (C)
Update and exchange local information during
encounter
Build up Familiar Set and Local Community

\[
\sumset CommunityAccept(), MergeCommunities()
\]
```

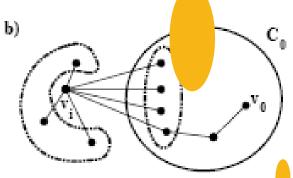
SIMPLE



We want to see whether v₁ should be added to C₀



Then we count the number of vertices in both v_1 's familiar set and also in the local community of $V_0 = \binom{n}{2}$



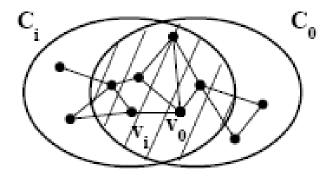
So we first count the number of vertices in v_i 's familiar set = $\left(\bigcirc + \bigcirc \right)$

d)
$$V_i$$
An ve admit $v_i C_0$ if λ

$$() > (C + ()) \times \lambda$$

CommunityAccept (v_i)

SIMPLE



We only consider merging the two communities $C_0 \& C_i$ if the fraction of them in common 1/1/1 > V

MergeCommunities (C_o , C_i)

Results and Evaluations

Data Set	SIMPLE	K-CLIQUE	MODULARITY
Reality	0.79/0.76	0.87	0.82
UCSD	0.47/0.56	0.55	0.40
Cambridge	0.85/0.85	0.85	0.87
Complexity	O(n)	O(n ²)	$O(n^4)/O(n^2k^2)$

Newman weighted analysis Palla et al, k-Clique

$$\sigma_{Jaccard} = \frac{|\Gamma_i \cap \Gamma_j|}{|\Gamma_i \cup \Gamma_j|}$$

Distributed BUBBLE RAP (DiBuBB)

