Exercises for 3/12/13

- 1. Give one-pass algorithms, and specify when they can be used, for:
 - (a) Set union
 - (b) Set intersection
 - (c) Bag intersection
 - (d) Bag union
 - (e) Outerjoin
- 2. For each of the operations below, write an iterator for a one-pass algorithm. Is the operation *blocking*, i.e., must the whole relation be read before the first output tuple is produced?
 - (a) Projection
 - (b) δ
 - (c) γ
 - (d) Set union
 - (e) Bag intersection
- 3. Nested-loop join
 - (a) Write iterator methods for the block-based version of nested-loop ioin
 - (b) B(R) = B(S) = 10.000, M = 1001. Calculate the I/O cost of a nested-loop join
 - (c) For the same relations, what value of M is needed to compute $R\bowtie S$ with nested-loop join and no more than (a) 100.000 (b) 25.000 and (c) 15.000 disk I/O s?
- 4. Describe 2-pass sort-based algorithms for the following. If the relation(s) have 10.000 blocks each, how much memory is needed?
 - (a) γ
 - (b) Set intersection
 - (c) Set union

- 5. Suppose that there is an index on R.a. How can the index be used to improve the execution of the following operations? When is it more efficient than sort- or hash-based algorithms?
 - (a) $R \cup_S S$
 - (b) $R \cap_S S$
 - (c) $\delta(R)$
- 6. Let B(R) = 10.000 and T(R) = 500.000 with an index on R.a and V(R,a) = k. Give the I/O cost of $\sigma_{a=0}(R)$ as a function of k in the case that (a) we use the index and (b) we don't use it
- 7. For each of the following schedules,
 - Construct the precedence graph
 - Is the schedule conflict-serializable?
 - If so, what are all the equivalent serial schedules?
 - (a) $r_1(A)$; $r_2(A)$; $r_3(B)$; $w_1(A)$; $r_2(C)$; $r_2(B)$; $w_2(B)$; $w_1(C)$
 - (b) $r_1(A)$; $w_1(B)$; $r_2(B)$: $w_2(C)$; $r_3(C)$; $w_3(A)$
 - (c) $w_3(A)$; $r_1(A)$; $w_1(B)$; $r_2(B)$: $w_2(C)$; $r_3(C)$;
 - (d) $r_1(A); r_2(A); w_1(B); w_2(B); r_1(B); r_2(B); w_2(C); w_1(D)$
 - (e) $r_1(A)$; $r_2(A)$; $r_1(B)$; $r_2(B)$; $r_3(A)$; $r_4(B)$; $w_1(A)$; $w_2(B)$