

APPENDIX B

INSTRUCTION LIST

ALPHABETIC BY MNEMONIC

WITH OP CODES, EXECUTION CYCLES

AND MEMORY REQUIREMENTS

The following notation applies to this summary:

A	Accumulator
X, Y	Index Registers
M	Memory
P	Processor Status Register
S	Stack Pointer
✓	Change
—	No Change
+	Add
∧	Logical AND
−	Subtract
⊕	Logical Exclusive Or
↑	Transfer from Stack
↓	Transfer to Stack
→	Transfer to
←	Transfer to
V	Logical OR
PC	Program Counter
PCH	Program Counter High
PCL	Program Counter Low
OPER	OPERAND
#	IMMEDIATE ADDRESSING MODE

Note: At the top of each table is located in parentheses a reference number (Ref: XX) which directs the user to that Section in the MCS6500 Microcomputer Family Programming Manual in which the instruction is defined and discussed.

ADC*Add memory to accumulator with carry***ADC**Operation: $A + M + C \rightarrow A, C$

N Z C I D V

✓ ✓ ✓ — — ✓

(Ref: 2.2.1)

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles
Immediate	ADC # Oper	69	2	2
Zero Page	ADC Oper	65	2	3
Zero Page, X	ADC Oper, X	75	2	4
Absolute	ADC Oper	6D	3	4
Absolute, X	ADC Oper, X	7D	3	4*
Absolute, Y	ADC Oper, Y	79	3	4*
(Indirect, X)	ADC (Oper, X)	61	2	6
(Indirect), Y	ADC (Oper), Y	71	2	5*

* Add 1 if page boundary is crossed.

BIT

BIT Test bits in memory with accumulator

BIT

Operation: $A \wedge M$, $M_7 \rightarrow N$, $M_6 \rightarrow V$

Bit 6 and 7 are transferred to the status register. N Z C I D V

If the result of $A \wedge M$ is zero then $Z = 1$, otherwise $M_7 \checkmark - - - M_6$

$Z = \emptyset$

(Ref: 4.2.1.1)

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles
Zero Page	BIT Oper	24	2	3
Absolute	BIT Oper	2C	3	4

SBC**SBC** *Subtract memory from accumulator with borrow***SBC**Operation: $A \leftarrow M - \overline{C} \rightarrow A$

N Z C I D V

Note: \overline{C} = Borrow

(Ref: 2.2.2)

✓ ✓ ✓ — — ✓

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles
Immediate	SBC #Oper	E9	2	2
Zero Page	SBC Oper	E5	2	3
Zero Page, X	SBC Oper, X	F5	2	4
Absolute	SBC Oper	ED	3	4
Absolute, X	SBC Oper, X	FD	3	4*
Absolute, Y	SBC Oper, Y	F9	3	4*
(Indirect, X)	SBC (Oper, X)	E1	2	6
(Indirect), Y	SBC (Oper), Y	F1	2	5*

* Add 1 when page boundary is crossed.

APPENDIX C

INSTRUCTION ADDRESSING

MODES AND

RELATED EXECUTION TIMES

INSTRUCTION ADDRESSING MODES AND RELATED EXECUTION TIMES (in clock cycles)

	Accumulator	Immediate	Zero Page	Zero Page, X	Zero Page, Y	Absolute	Absolute, X	Absolute, Y	Implied	Relative	(Indirect, X)	(Indirect, Y)	Absolute Indirect
✓ ADC	.	2	3	4	.	4	4*	4*	.	.	6	5*	.
✓ AND	.	2	3	4	.	4	4*	4*	.	.	6	5*	.
✓ ASL	2	.	5	6	.	6	7	.	.	2**	.	.	.
BCC	2**	.	.	.
BCS	2**	.	.	.
BEQ	2**	.	.	.
BIT	.	.	3	.	.	4
✓ BMI	2**	.	.	.
✓ BNE	2**	.	.	.
✓ BPL	2**	.	.	.
✓ BRK
BVC	2**	.	.	.
BVS	2**	.	.	.
✓ CLC	2	.	.	.
✓ CLD	2	.	.	.
✓ CLI	2	.	.	.
✓ CLV	2	.	.	.
✓ CMP	.	2	3	4	.	4	4*	4*	2	.	6	5*	.
✓ CPX	.	2	3	.	.	4
✓ CPY	.	2	3	.	.	4
✓ DEC	.	.	5	6	.	6	7
✓ DEX
✓ DEY
✓ EOR	.	2	3	4	.	4	4*	4*	2	.	6	5	.
✓ INC	.	.	5	6	.	6	7
✓ INX
✓ INY
✓ JMP	3	5
JSR
LDA
LDX
LDY
LSR	2	.	5	6	.	6	7
NOP	.	2	3	4	.	4	4*	4*	2	.	6	5*	.
ORA
PHA
PHP
PLA
PLP
ROL	2	.	5	6	.	6	7
ROR	2	.	5	6	.	6	7
RTI
RTS
SBC	.	2	3	4	.	4	4*	4*	2	.	6	5*	.
SEC
SED
SEI
STA	.	.	3	4	.	4	4	5	2	.	6	6	.
STX*	.	.	3	4	.	4	4
STY**	.	.	3	4	.	4	4
TAX
TAY
TSX
TXA
TXS
TYA

* Add one cycle if indexing across page boundary

** Add one cycle if branch is taken, Add one additional if branching operation crosses page boundary

APPENDIX E

SUMMARY OF ADDRESSING MODES

This appendix is to serve the user in providing a reference for the MCS650X addressing modes. Each mode of address is shown with a symbolic illustration of the bus status at each cycle during the instruction fetch and execution. The example number as found in the text is provided for reference purposes.

E.1 IMPLIED ADDRESSING

Example 5.3: Illustration of implied addressing

<u>Clock Cycle</u>	<u>Address Bus</u>	<u>Program Counter</u>	<u>Data Bus</u>	<u>Comments</u>
1	PC	PC + 1	OP CODE	Fetch OP CODE
2	PC + 1	PC + 1	New OP CODE	Ignore New OP CODE; Decode Old OP CODE
3	PC + 1	PC + 2	New OP CODE	Fetch New OP CODE; Execute Old OP CODE

E.2 IMMEDIATE ADDRESSING

Example 5.4: Illustration of immediate addressing

<u>Clock Cycle</u>	<u>Address Bus</u>	<u>Program Counter</u>	<u>Data Bus</u>	<u>Comments</u>
1	PC	PC + 1	OP CODE	Fetch OP CODE
2	PC + 1	PC + 2	Data	Fetch Data, Decode OP CODE
3	PC + 2	PC + 3	New OP CODE	Fetch New OP CODE, Execute Old OP CODE

E.3 ABSOLUTE ADDRESSING

Example 5.5: Illustration of absolute addressing

<u>Clock Cycle</u>	<u>Address Bus</u>	<u>Program Counter</u>	<u>Data Bus</u>	<u>Comments</u>
1	PC	PC + 1	OP CODE	Fetch OP CODE
2	PC + 1	PC + 2	ADL	Fetch ADL, Decode OP CODE
3	PC + 2	PC + 3	ADH	Fetch ADH, Retail ADL
4	ADH, ADL	PC + 3	Data	Fetch Data
5	PC + 3	PC + 4	New OP CODE	Fetch New OP CODE, Execute Old OP CODE

E.4 ZERO PAGE ADDRESSING

Example 5.6: Illustration of zero page addressing

<u>Clock Cycle</u>	<u>Address Bus</u>	<u>Program Counter</u>	<u>Data Bus</u>	<u>Comments</u>
1	PC	PC + 1	OP CODE	Fetch OP CODE
2	PC + 1	PC + 2	ADL	Fetch ADL, De- code OP CODE
3	00, ADL	PC + 2	Data	Fetch Data
4	PC + 2	PC + 3	New OP CODE	Fetch New OP CODE, Exe- cute Old OP CODE

E.5 RELATIVE ADDRESSING – (Branch Positive, no crossing of page boundaries)

Example 5.8: Illustration of relative addressing--branch positive taken, no crossing of page boundaries

<u>Cycle</u>	<u>Address Bus</u>	<u>Data Bus</u>	<u>External Operation</u>	<u>Internal Operation</u>
1	0100	OP CODE	Fetch OP CODE	Finish Previous Oper- ation, Increment Pro- gram Counter to 101
2	0101	+50	Fetch Offset	Interpret Instruction, Increment Program Counter to 102
3	0102	Next OP CODE	Fetch Next OP CODE	Check Flags, Add Rela- tive to PCL, Increment Program Counter to 103
4	0152	Next OP CODE	Fetch Next OP CODE	Transfer Results to PCL, Increment Program Counter to 153

E.6 ABSOLUTE INDEXED ADDRESSING – (with page crossing)

Step 5 is deleted and the data in step 4 is valid when no page crossing occurs.

Example 6.7: Absolute Indexed; With Page Crossing

<u>Cycle</u>	<u>Address Bus</u>	<u>Data Bus</u>	<u>External Operation</u>	<u>Internal Operation</u>
1	0100	OP CODE	Fetch OP CODE	Finish Previous Operation Increment PC to 101
2	0101	BAL	Fetch BAL	Interpret Instruction Increment PC to 102
3	0102	BAH	Fetch BAH	Add BAL + Index Increment PC to 103
4	BAH, BAL +X	Data (Ignore)	Fetch Data (Data is ignored)	Add BAH + Carry
5	BAH+1, BAL+X	Data	Fetch Data	
6	0103	Next OP CODE	Fetch Next OP CODE	Finish Operation

E.7 ZERO PAGE INDEXED ADDRESSING

Example 6.8: Illustration of Zero Page Indexing

<u>Cycle</u>	<u>Address Bus</u>	<u>Data Bus</u>	<u>External Operation</u>	<u>Internal Operation</u>
1	0100	OP CODE	Fetch OP CODE	Finish Previous Operation
2	0101	BAL	Fetch Base Address Low (BAL)	Interpret Instruct- ion
3	00,BAL	Data (Dis- carded	Fetch Discarded Data	Add: BAL + X
4	00,BAL +X	Data	Fetch Data	
5	0102	Next OP CODE	Fetch Next OP CODE	Finish Operation

E.8 INDEXED INDIRECT ADDRESSING

Example 6.10: Illustration of Indexed Indirect Addressing

<u>Cycle</u>	<u>Address Bus</u>	<u>Data Bus</u>	<u>External Operation</u>	<u>Internal Operation</u>
1	0100	OP CODE	Fetch OP CODE	Finish Previous Operation
2	0101	BAL	Fetch BAL	Interpret Instruction
3	00,BAL	DATA (Discarded)	Fetch Discarded DATA	Add BAL + X
4	00,BAL + X	ADL	Fetch ADL	Add 1 to BAL + X
5	00,BAL + X + 1		Fetch ADH	Hold ADL
6	ADH,ADL	DATA	Fetch DATA	
7	0102	Next OP	Fetch Next OP CODE	Finish Operation

E.9 INDIRECT INDEXED ADDRESSING (with page crossing)

Step 6 is deleted and the data in step 5 is valid when no page crossing occurs.

Example 6.12: Indirect Indexed Addressing (With Page Crossing)

<u>Cycle</u>	<u>Address Bus</u>	<u>Data Bus</u>	<u>External Operation</u>	<u>Internal Operation</u>
1	0100	OP CODE	Load OP CODE	Finish Previous Operation
2	0101	IAL	Fetch IAL	Interpret Instruction
3	00, IAL	BAL	Fetch BAL	Add 1 to IAL
4	00, IAL + 1	BAH	Fetch BAH	Add BAL to Y
5	BAH, BAL + Y	DATA (Discarded)	Fetch DATA (Discarded)	Add 1 to BAH
6	BAH + 1 BAL + Y	DATA	Fetch Data	
7	0102	Next OP CODE	Fetch Next OP CODE	Finish This Operation

APPENDIX H

REVIEW OF BINARY

AND

BINARY CODED DECIMAL

ARITHMETIC

The microprocessor automatically takes this into account and corrects for the fact that

<u>Decimal</u>		<u>BCD</u>		<u>Hex</u>
79	=	01111001		79 = 01111001
+12	=	00010010		12 = 00010010
91	=	10010001		88 = 10001011

The only difference between Hex and BCD representation is that the microprocessor automatically adjusts for the fact that BCD does not allow for Hex values A - F during add and subtract operations.

The offset which follows a branch instruction is in signed two's complement form which means that

$$\begin{aligned}
 \$+50 &= +80 = 01010000 \\
 \text{and } \$-50 &= -80 = 10110000 \\
 \text{Proof} &= 00000000
 \end{aligned}$$

The sign for this operation is in bit 7 where an 0 equals positive and a 1 equals negative.

This bit is correct for the two's complement representation but also flags the microprocessor whether to carry or borrow from the address high byte.

The following 4 examples represent the combinations of offsets which might occur (all notations are in hexadecimal):

Example H.4.1: Forward reference, no page crossing

0105	INE
0106	+55
0107	Next CP CODE

To calculate next instruction if the branch is taken

Offset	+55	01010101
Address Low		
for next		
OP CODE	07	00000111
	5C	01011100

with no carry, giving 015C as the result.

Example H.4.2: Backward reference, no page crossing

015A	BNE
015B	-55
015C	Next OP CODE

To calculate if branch is taken,

Offset	-55 = AB = 10101011
+ Address Low for	
Next OP CODE	$\frac{+5C}{07} = \frac{5C}{07} = \frac{01011100}{00000111}$

The carry is expected because of the negative offset and is ignored, thus giving 0107 as the result.

Example H.4.3: Backward reference if page boundary crossed

0105	BNE
0106	-55
0107	Next OP CODE

To calculate if branch is taken, first calculate a low byte

Offset	-55 = AB = 10101011
Address Low for	
Next OP CODE	$\frac{07}{B2} = \frac{07}{B2} = \frac{00000111}{10110010}$

There is no carry from a negative offset; therefore, a carry must be made:

-1 =	-1 = FF = 11111111
+ Address High	$= \frac{01}{00} = \frac{01}{00} = \frac{00000001}{00000000}$

This gives 00 B2 as a result.

Example H.4.4: Forward reference across page boundary

00B0	BNE
00B1	+55
00B2	Next OP CODE

To calculate next instruction if branch is taken,

Offset	55 = 01010101
Address Low for Next	
OP CODE	B2 = 10110010
	07 = 00000111

with carry on positive number.

	+1	1 = 00000001
Address High	00 = 00000000	
	1 = 00000001	

which gives 0107.