# Basic idea of quantum computing and the operation of a specific quantum computing algorithm Quantum Engineering II

3rd of June 2021

### Quantum Parallelism

Quantum Parallelism: Parallel computation on superposition states

$$\frac{|0\rangle + |1\rangle}{\sqrt{2}} = \hat{H}|0\rangle. \tag{1}$$

Interference can make calculation that depends on all results.



#### Deutsch-Jozsa algorithm

Determine is the function of  $2^n$  entries even or odd?

$$f(x) = \frac{(110010)}{(000000)}$$

Classical: Worst case  $\frac{2^n}{2} + 1$  evaluations



# Deutsch-Jozsa algorithm

Begin with a query and anwer registers

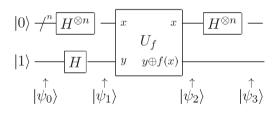
$$|\psi_0\rangle = |0\rangle^{\otimes n}|1\rangle$$

Put into superposition with Hadamard

$$|\psi_1\rangle = \sum_{x \in \{0,1\}^n} \frac{|x\rangle_n}{\sqrt{2^n}} |-\rangle$$

Calculate function  $|-\rangle \rightarrow (-1)^{f(x)}|-\rangle$ 

$$|\psi_2\rangle = \sum_{\mathsf{x}\in\{0,1\}^n} (-1)^{f(\mathsf{x})} \frac{|\mathsf{x}\rangle_n}{\sqrt{2^n}} |-\rangle$$



## Deutsch-Jozsa algorithm

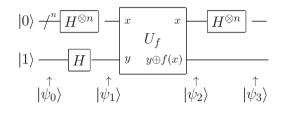
Use Hardamard on all again

$$H^{\otimes n}|x\rangle_n = \sum_z (-1)^{\mathbf{x}\cdot\mathbf{z}}|z\rangle_n/\sqrt{2^n}$$

$$|\psi_3\rangle = \sum_{z} \left( \sum_{x} \frac{(-1)^{\mathbf{x} \cdot \mathbf{z} + f(x)}}{2^n} \right) |z\rangle_n |-\rangle$$

$$c_z = \sum_{x} \frac{(-1)^{\mathbf{x} \cdot \mathbf{z} + f(x)}}{2^n}$$

$$c_0 = \sum_{x} \frac{(-1)^{f(x)}}{2^n} egin{cases} \pm 1 & \text{(constant)} \\ 0 & \text{(balanced)} \end{cases}$$



If projection onto  $|0\rangle_n$  then (constant) if not then (balanced)



#### Pros and Cons

Pros

Only one iteration is needed compared  $\frac{2^n}{2} + 1$ 

Cons

Not important problem

Classically you can be lucky after 2 iterations