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44 1 1 1	possible/necessary	<pre>typename ALLOC = std::allocator<elem> > class</elem></pre>
1.1 header.h	<pre>12 #define vb vector < bool ></pre>	Container >
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1 #pragma once // Delete this when copying this	16 #define vpi vector <pi>16 #define vpi vector <pi>20</pi></pi>	container.begin();
file	17 #define vld vector <ld></ld>	if (beg != container.end()) {
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5 #define ll long long	end())	31
6 #define ull unsigned ll	<pre>20 #define in(el, cont) (find(cont.begin(), cont.end</pre>	32 } 33 }
7 #define ld long double	21	34 return o;
8 #define pl pair<11, 11>	22 constexpr int INF = 200000010;	35 }
<pre>9 #define pi pair<int, int=""> // use pl where</int,></pre>	23 constexpr ll LLINF = 90000000000000010LL;	36

1.2 Bash for c++ compile with header.h

```
#!/bin/bash

if [ $# -ne 1 ]; then echo "Usage: $0 <input_file >"; exit 1; fi

f="$1"; d=code/; o=a.out

[-f $d/$f] || { echo "Input file not found: $f "; exit 1; }

g++ -I$d $d/$f -o $o && echo "Compilation successful. Executable '$o' created." || echo "Compilation failed."
```

1.3 Bash for run tests c++

```
1 g++ $1/$1.cpp -o $1/$1.out
2 for file in $1/*.in; do diff <($1/$1.out < "$file
") "${file%.in}.ans"; done</pre>
```

1.4 Bash for run tests python

```
_{\rm 1} for file in $1/*.in; do diff <(python3 $1/$1.py < "$file") "${file%.in}.ans"; done
```

1.4.1 Aux. helper C++

```
#include "header.h"

int main() {

// Read in a line including white space
string line;
getline(cin, line);
// When doing the above read numbers as
follows:
int n;
getline(cin, line);
```

```
stringstream ss(line);
ss >> n;

// Count the number of 1s in binary
represnatation of a number
ull number;
__builtin_popcountll(number);
}
```

1.4.2 Aux. helper python

```
# Read until EOF
while True:
    try:
    pattern = input()
    except EOFError:
    break
```

2 Python

2.1 Graphs

2.1.1 BFS

```
1 from collections import deque
2 def bfs(g, roots, n):
      q = deque(roots)
      explored = set(roots)
      distances = [float("inf")]*n
      distances[0][0] = 0
      while len(q) != 0:
          node = q.popleft()
          if node in explored: continue
11
          explored.add(node)
          for neigh in g[node]:
              if neigh not in explored:
                  q.append(neigh)
                   distances[neigh] = distances[node
      return distances
```

2.1.2 Dijkstra

```
from heapq import *
2 def dijkstra(n, root, g): # g = {node: (cost, neigh)}
3  dist = [float("inf")]*n
4  dist[root] = 0
5  prev = [-1]*n
```

```
pq = [(0, root)]
heapify(pq)
visited = set([])
while len(pq) != 0:
  _, node = heappop(pq)
  if node in visited: continue
  visited.add(node)
  # In case of disconnected graphs
  if node not in g:
    continue
  for cost, neigh in g[node]:
    alt = dist[node] + cost
    if alt < dist[neigh]:</pre>
      dist[neigh] = alt
      prev[neigh] = node
      heappush(pq, (alt, neigh))
return dist
```

2.2 Num. Th. / Comb.

2.2.1 nCk % prime

```
1  # Note: p must be prime and k  n:
4         return 0
5     # calculate numerator
6     num = 1
7     for i in range(n-k+1, n+1):
8         num *= i % p
9     num %= p
10     # calculate denominator
11     denom = 1
12     for i in range(1,k+1):
13         denom *= i % p
14     denom %= p
15     # numerator * denominator^(p-2) (mod p)
16     return (num * pow(denom, p-2, p)) % p
```

2.2.2 Sieve of E. O(n) so actually faster than C++ version, but more memory

```
_{1} MAX STZE = 10**8+1
2 isprime = [True] * MAX SIZE
3 prime = []
4 SPF = [None] * (MAX SIZE)
6 def manipulated_seive(N): # Up to N (not
      included)
    isprime[0] = isprime[1] = False
    for i in range(2. N):
      if isprime[i] == True:
        prime.append(i)
        SPF[i] = i
11
      while (j < len(prime) and
       i * prime[j] < N and</pre>
          prime[i] <= SPF[i]):</pre>
        isprime[i * prime[j]] = False
        SPF[i * prime[j]] = prime[j]
        j += 1
```

2.3 Strings

2.3.1 LCS

```
1 def longestCommonSubsequence(text1, text2): # 0(
      m*n) time, O(m) space
      n = len(text1)
      m = len(text2)
      # Initializing two lists of size m
      prev = [0] * (m + 1)
      cur = [0] * (m + 1)
      for idx1 in range(1, n + 1):
          for idx2 in range(1, m + 1):
              # If characters are matching
11
              if text1[idx1 - 1] == text2[idx2 -
                  cur[idx2] = 1 + prev[idx2 - 1]
                  # If characters are not matching
15
                  cur[idx2] = max(cur[idx2 - 1],
                      prev[idx2])
          prev = cur.copy()
18
19
      return cur[m]
```

2.3.2 KMP

```
1 class KMP:
      def partial(self, pattern):
2
           """ Calculate partial match table: String
3
               -> [Int]"""
          ret = [0]
          for i in range(1, len(pattern)):
              j = ret[i - 1]
              while j > 0 and pattern[j] != pattern
                  [i]: j = ret[j-1]
              ret.append(j + 1 if pattern[j] ==
                  pattern[i] else i)
          return ret
11
      def search(self, T, P):
          """KMP search main algorithm: String ->
12
              String -> [Int]
          Return all the matching position of
13
              pattern string P in T"""
          partial, ret, j = self.partial(P), [], 0
          for i in range(len(T)):
15
              while j > 0 and T[i] != P[j]: j =
16
                  partial[j - 1]
              if T[i] == P[i]: i += 1
17
              if j == len(P):
                  ret.append(i - (j - 1))
                  j = partial[j - 1]
          return ret
```

2.3.3 Edit distance

```
def editDistance(str1, str2):
   # Get the lengths of the input strings
   m = len(str1)
   n = len(str2)
   # Initialize a list to store the current row
   curr = [0] * (n + 1)
   # Initialize the first row with values from 0
   for j in range(n + 1):
     curr[i] = i
   # Initialize a variable to store the previous
       value
   previous = 0
   # Loop through the rows of the dynamic
       programming matrix
   for i in range (1, m + 1):
     # Store the current value at the beginning of
           the row
     previous = curr[0]
     curr[0] = i
```

```
# Loop through the columns of the dynamic
      programming matrix
  for i in range(1, n + 1):
    # Store the current value in a temporary
    temp = curr[i]
    # Check if the characters at the current
        positions in str1 and str2 are the same
    if str1[i - 1] == str2[i - 1]:
      curr[j] = previous
      # Update the current cell with the
          minimum of the three adjacent cells
      curr[j] = 1 + min(previous, curr[j - 1],
          curr[i])
    # Update the previous variable with the
       temporary value
    previous = temp
# The value in the last cell represents the
    minimum number of operations
return curr[n]
```

2.4 Other Algorithms

2.4.1 Rotate matrix

```
1 def rotate_matrix(m):
2    return [[m[j][i] for j in range(len(m))] for
        i in range(len(m[0])-1,-1,-1)]
```

2.5 Other Data Structures

2.5.1 Segment Tree

```
# set value at position p
      tree[p + n] = value
      p = p + n
16
17
      i = p # move upward and update parents
18
      while i > 1:
19
          tree[i >> 1] = tree[i] + tree[i ^ 1]
          i >>= 1
21
23 def query(1, r, n): # function to get sum on
      interval [1, r)
      res = 0
      # loop to find the sum in the range
      r += n
      while l < r:
28
          if 1 & 1:
29
              res += tree[1]
              1 += 1
          if r & 1:
              r -= 1
33
              res += tree[r]
          1 >>= 1
          r >>= 1
36
      return res
```

2.5.2 Trie

```
1 class TrieNode:
      def init (self):
          self.children = [None] *26
          self.isEndOfWord = False
6 class Trie:
      def __init__(self):
          self.root = self.getNode()
      def getNode(self):
          return TrieNode()
11
12
      def _charToIndex(self,ch):
13
          return ord(ch)-ord('a')
14
15
16
      def insert(self.kev):
17
          pCrawl = self.root
18
          length = len(key)
19
          for level in range(length):
20
              index = self._charToIndex(key[level])
21
               if not pCrawl.children[index]:
22
                   pCrawl.children[index] = self.
23
                       getNode()
               pCrawl = pCrawl.children[index]
          pCrawl.isEndOfWord = True
```

```
26
      def search(self, key):
27
          pCrawl = self.root
28
          length = len(kev)
29
          for level in range(length):
               index = self._charToIndex(key[level])
31
32
               if not pCrawl.children[index]:
                   return False
33
               pCrawl = pCrawl.children[index]
          return pCrawl.isEndOfWord
```

3 C++

3.1 Graphs

1 #include "header.h"

3.1.1 BFS

```
2 #define graph unordered_map<11, unordered_set<11</pre>
3 vi bfs(int n, graph& g, vi& roots) {
      vi parents(n+1, -1); // nodes are 1..n
      unordered_set <int> visited;
       queue < int > q;
      for (auto x: roots) {
          q.emplace(x);
           visited.insert(x);
      7
10
       while (not q.empty()) {
11
           int node = q.front():
12
          q.pop();
13
14
           for (auto neigh: g[node]) {
15
               if (not in(neigh, visited)) {
16
                   parents[neigh] = node;
17
                   g.emplace(neigh):
18
                   visited.insert(neigh);
               }
20
          }
23
      return parents;
24 }
25 vi reconstruct path(vi parents, int start, int
      goal) {
      vi path;
      int curr = goal:
       while (curr != start) {
           path.push_back(curr);
29
          if (parents[curr] == -1) return vi(); //
30
               No path, empty vi
           curr = parents[curr]:
31
```

```
path.push_back(start);
reverse(path.begin(), path.end());
return path;
}
```

3.1.2 DFS Cycle detection / removal

```
1 #include "header.h"
void removeCvc(ll node, unordered map<ll, vector<</pre>
      pair<11, 11>>>& neighs, vector<bool>& visited
3 vector < bool > & recStack. vector < 11 > & ans) {
      if (!visited[node]) {
          visited[node] = true:
           recStack[node] = true;
           auto it = neighs.find(node);
           if (it != neighs.end()) {
               for (auto util: it->second) {
                   11 nnode = util.first:
                   if (recStack[nnode]) {
                       ans.push_back(util.second);
                   } else if (!visited[nnode]) {
                       removeCyc(nnode, neighs,
                           visited, recStack, ans);
                   }
               }
           }
17
      recStack[node] = false;
20 }
```

3.1.3 Dijkstra

```
1 #include "header.h"
2 vector<int> dijkstra(int n, int root, map<int,</pre>
      vector<pair<int. int>>>& g) {
    unordered_set <int> visited;
    vector < int > dist(n. INF):
      priority_queue < pair < int , int >> pq;
      dist[root] = 0;
      pq.push({0, root});
      while (!pq.empty()) {
          int node = pq.top().second;
           int d = -pq.top().first;
           pq.pop();
11
12
           if (in(node, visited)) continue;
           visited.insert(node);
           for (auto e : g[node]) {
               int neigh = e.first;
17
               int cost = e.second:
```

3.1.4 Floyd-Warshall

3.1.5 Kruskal Minimum spanning tree of undirected weighted graph

```
1 #include "header.h"
2 #include "disjoint_set.h"
3 // O(E log E)
4 pair < set < pair < 11, 11 >> , 11 > kruskal (vector < tuple</pre>
      <11, 11, 11>>& edges, 11 n) {
      set <pair <11, 11>> ans;
      11 cost = 0:
      sort(edges.begin(), edges.end());
      DisjointSet < 11 > fs(n);
11
      ll dist, i, j;
      for (auto edge: edges) {
12
13
          dist = get<0>(edge);
          i = get<1>(edge);
14
          i = get < 2 > (edge);
16
17
           if (fs.find_set(i) != fs.find_set(j)) {
               fs.union_sets(i, j);
18
               ans.insert({i, j});
               cost += dist;
          }
```

```
23     return pair<set<pair<11, 11>>, 11> {ans, cost
          };
24 }
```

3.1.6 Hungarian algorithm

```
1 #include "header.h"
3 template <class T> bool ckmin(T &a, const T &b) {
       return b < a ? a = b, 1 : 0; }
5 * Given J jobs and W workers (J <= W), computes</pre>
       the minimum cost to assign each
6 * prefix of jobs to distinct workers.
7 * @tparam T a type large enough to represent
       integers on the order of J *
  * max(|C|)
  * Oparam C a matrix of dimensions JxW such that
       C[i][w] = cost to assign i-th
* job to w-th worker (possibly negative)
* @return a vector of length J, with the j-th
       entry equaling the minimum cost
* to assign the first (j+1) jobs to distinct
       workers
15 template <class T> vector<T> hungarian(const
      vector < vector < T >> &C) {
      const int J = (int)size(C), W = (int)size(C
          [01]:
      assert(J <= W);</pre>
      // job[w] = job assigned to w-th worker, or
          -1 if no job assigned
      // note: a W-th worker was added for
          convenience
      vector < int > job(W + 1, -1);
      vector<T> ys(J), yt(W + 1); // potentials
      // -yt[W] will equal the sum of all deltas
22
      vector <T> answers;
      const T inf = numeric limits <T>::max():
      for (int j_cur = 0; j_cur < J; ++j_cur) { //</pre>
           assign j_cur-th job
          int w_cur = W;
26
          job[w_cur] = i_cur;
27
          // min reduced cost over edges from Z to
              worker w
          vector <T> min_to(W + 1, inf);
          vector<int> prv(W + 1, -1); // previous
30
              worker on alternating path
          vector < bool > in_Z(W + 1);  // whether
31
              worker is in Z
          while (job[w_cur] != -1) {    // runs at
32
              most j_cur + 1 times
              in Z[w cur] = true:
```

```
const int j = job[w_cur];
              T delta = inf;
36
              int w next:
              for (int w = 0: w < W: ++w) {
                  if (!in_Z[w]) {
                      if (ckmin(min_to[w], C[j][w]
                          - ys[j] - yt[w]))
                          prv[w] = w_cur;
                      if (ckmin(delta, min to[w]))
                          w_next = w;
                  }
              }
              // delta will always be non-negative,
              // except possibly during the first
                  time this loop runs
              // if any entries of C[j_cur] are
                  negative
              for (int w = 0; w \le W; ++w) {
                  if (in_Z[w]) ys[job[w]] += delta,
                       yt[w] -= delta;
                  else min_to[w] -= delta;
              w_cur = w_next;
          // update assignments along alternating
          for (int w; w_cur != W; w_cur = w) job[
              w curl = iob[w = prv[w curl]:
          answers.push_back(-yt[W]);
      return answers;
```

3.1.7 Suc. shortest path Calculates max flow, min cost

```
1 #include "header.h"
2 // map<node, map<node, pair<cost, capacity>>>
3 #define graph unordered_map<int, unordered_map<</pre>
     int. pair<ld. int>>>
5 const ld infty = 1e60l; // Change if necessary
6 ld fill(int n, vld& potential) { // Finds max
     flow, min cost
   priority_queue < pair < ld, int >> pq;
   vector < bool > visited(n+2, false);
   vi parent(n+2, 0);
   vld dist(n+2, infty);
   dist[0] = 0.1;
   pq.emplace(make_pair(0.1, 0));
   while (not pq.empty()) {
     int node = pq.top().second;
     pq.pop();
     if (visited[node]) continue:
```

```
visited[node] = true:
      for (auto& x : g[node]) {
        int neigh = x.first:
19
        int capacity = x.second.second;
        ld cost = x.second.first;
21
        if (capacity and not visited[neigh]) {
22
          ld d = dist[node] + cost + potential[node
              ] - potential[neigh];
          if (d + 1e-101 < dist[neigh]) {</pre>
            dist[neigh] = d;
            pq.emplace(make_pair(-d, neigh));
26
            parent[neigh] = node:
    }}}
28
    for (int i = 0; i < n+2; i++) {
      potential[i] = min(infty, potential[i] + dist
    if (not parent[n+1]) return infty;
    1d ans = 0.1:
    for (int x = n+1; x; x=parent[x]) {
      ans += g[parent[x]][x].first;
      g[parent[x]][x].second--;
      g[x][parent[x]].second++;
    return ans;
41 }
```

3.1.8 Bipartite check

```
1 #include "header.h"
2 int main() {
     int n;
     vvi adj(n);
     vi side(n, -1);
                       // will have 0's for one
         side 1's for other side
     bool is bipartite = true: // becomes false
         if not bipartite
     queue < int > q:
     for (int st = 0; st < n; ++st) {</pre>
          if (side[st] == -1) {
             q.push(st);
              side[st] = 0;
              while (!a.emptv()) {
                  int v = q.front();
                  q.pop();
                  for (int u : adj[v]) {
                      if (side[u] == -1) {
                          side[u] = side[v] ^ 1;
                          q.push(u);
                      } else {
                          is_bipartite &= side[u]
                              != side[v]:
```

```
22 }
23 }}}}}
```

3.1.9 Find cycle directed

```
1 #include "header.h"
2 int n:
3 \text{ const int } mxN = 2e5+5;
4 vvi adj(mxN);
5 vector < char > color:
6 vi parent;
7 int cycle_start, cycle_end;
8 bool dfs(int v) {
      color[v] = 1;
      for (int u : adj[v]) {
           if (color[u] == 0) {
11
               parent[u] = v;
12
               if (dfs(u)) return true;
13
           } else if (color[u] == 1) {
               cycle_end = v;
15
               cycle_start = u;
               return true;
17
           }
       color[v] = 2;
       return false:
23 void find_cycle() {
       color.assign(n, 0);
      parent.assign(n, -1);
      cvcle start = -1:
      for (int v = 0; v < n; v++) {
           if (color[v] == 0 && dfs(v))break;
28
20
      if (cycle_start == -1) {
30
           cout << "Acyclic" << endl;</pre>
31
      } else {
32
           vector<int> cycle;
33
           cycle.push_back(cycle_start);
34
           for (int v = cycle_end; v != cycle_start;
35
                v = parent[v])
               cycle.push_back(v);
           cycle.push_back(cycle_start);
37
           reverse(cycle.begin(), cycle.end());
           cout << "Cycle__Found:__";</pre>
           for (int v : cvcle) cout << v << "":</pre>
41
           cout << endl:
42
43
44 }
```

3.1.10 Find cycle directed

```
1 #include "header.h"
2 int n;
3 const int mxN = 2e5 + 5:
4 vvi adj(mxN);
5 vector < bool > visited;
6 vi parent:
7 int cycle_start, cycle_end;
8 bool dfs(int v, int par) { // passing vertex and
      its parent vertex
      visited[v] = true:
      for (int u : adj[v]) {
          if(u == par) continue; // skipping edge
               to parent vertex
          if (visited[u]) {
               cvcle_end = v;
               cycle_start = u;
               return true:
           parent[u] = v:
          if (dfs(u, parent[u]))
               return true;
19
20
      return false;
22 }
23 void find cvcle() {
      visited.assign(n, false);
      parent.assign(n. -1):
      cvcle_start = -1;
      for (int v = 0; v < n; v++) {
           if (!visited[v] && dfs(v, parent[v]))
              break:
      if (cycle_start == -1) {
           cout << "Acvclic" << endl;</pre>
32
           vector<int> cycle;
           cycle.push_back(cycle_start);
           for (int v = cycle_end; v != cycle_start;
               v = parent[v])
               cycle.push_back(v);
           cvcle.push back(cvcle start):
           cout << "Cycle_Found:_";</pre>
           for (int v : cycle) cout << v << "";</pre>
39
           cout << endl:
      }
41
42 }
```

3.1.11 Tarjan's SCC

```
1 #include "header.h"
2
3 struct Tarjan {
4  vvi &edges;
5  int V, counter = 0, C = 0;
```

```
vi n. 1:
    vector < bool > vs;
    stack<int> st:
    Tarjan(vvi &e) : edges(e), V(e.size()), n(V.
        -1), l(V, -1), vs(V, false) {}
    void visit(int u, vi &com) {
      1[u] = n[u] = counter++:
      st.push(u);
      vs[u] = true:
      for (auto &&v : edges[u]) {
        if (n[v] == -1) visit(v, com);
15
        if (vs[v]) 1[u] = min(1[u], 1[v]);
17
      if (1[u] == n[u]) {
        while (true) {
          int v = st.top();
          st.pop();
21
          vs[v] = false:
          com[v] = C: // <== ACT HERE
          if (u == v) break:
25
        C++:
26
28
    int find_sccs(vi &com) { // component indices
        will be stored in 'com'
      com.assign(V, -1);
      for (int u = 0; u < V; ++u)
        if (n[u] == -1) visit(u, com);
    }
35
    // scc is a map of the original vertices of the
         graph to the vertices
    // of the SCC graph, scc graph is its adjacency
    // SCC indices and edges are stored in 'scc'
        and 'scc graph'.
    void scc_collapse(vi &scc, vvi &scc_graph) {
      find_sccs(scc);
      scc_graph.assign(C, vi());
      set < pi > rec; // recorded edges
42
      for (int u = 0: u < V: ++u) {
        assert(scc[u] != -1):
        for (int v : edges[u]) {
          if (scc[v] == scc[u] ||
            rec.find({scc[u], scc[v]}) != rec.end()
47
                ) continue;
          scc_graph[scc[u]].push_back(scc[v]);
          rec.insert({scc[u], scc[v]});
      }
51
    // Function to find sources and sinks in the
        SCC graph
```

```
// The number of edges needed to be added is
        max(sources.size(), sinks.())
    void findSourcesAndSinks(const vvi &scc graph.
        vi &sources. vi &sinks) {
      vi in_degree(C, 0), out_degree(C, 0);
      for (int u = 0; u < C; u++) {</pre>
        for (auto v : scc_graph[u]) {
          in_degree[v]++;
          out degree[u]++:
      for (int i = 0: i < C: ++i) {</pre>
63
        if (in_degree[i] == 0) sources.push_back(i)
64
        if (out_degree[i] == 0) sinks.push_back(i);
66
   }
68 };
```

3.1.12 SCC edges Prints out the missing edges to make the input digraph strongly connected

```
1 #include "header.h"
2 const int N=1e5+10:
3 int n,a[N],cnt[N],vis[N];
4 vector <int> hd,tl;
5 int dfs(int x){
      vis[x]=1:
      if(!vis[a[x]])return vis[x]=dfs(a[x]);
      return vis[x]=x:
9 }
10 int main(){
      scanf("%d",&n);
      for(int i=1;i<=n;i++){</pre>
           scanf("%d",&a[i]);
13
           cnt[a[i]]++;
14
      int k=0:
      for(int i=1;i<=n;i++){</pre>
           if(!cnt[i]){
               k++:
               hd.push_back(i);
               tl.push_back(dfs(i));
21
           }
      int tk=k:
      for(int i=1;i<=n;i++){</pre>
           if(!vis[i]){
26
27
               k++;
               hd.push_back(i);
               tl.push_back(dfs(i));
           }
31
      if(k==1&&!tk)k=0:
```

3.1.13 Find Bridges

```
1 #include "header.h"
2 int n; // number of nodes
3 vvi adj; // adjacency list of graph
4 vector < bool > visited;
5 vi tin. low:
6 int timer:
7 void dfs(int v, int p = -1) {
      visited[v] = true:
      tin[v] = low[v] = timer++;
      for (int to : adj[v]) {
          if (to == p) continue;
          if (visited[to]) {
              low[v] = min(low[v], tin[to]);
          } else {
              dfs(to, v);
              low[v] = min(low[v], low[to]);
              if (low[to] > tin[v])
                  IS_BRIDGE(v, to);
          }
      }
21 }
22 void find_bridges() {
      timer = 0;
      visited.assign(n, false);
      tin.assign(n, -1);
      low.assign(n, -1);
      for (int i = 0; i < n; ++i) {</pre>
          if (!visited[i]) dfs(i);
29
30 }
```

3.1.14 Artic. points (i.e. cut off points)

```
#include "header.h"
int n; // number of nodes
vvi adj; // adjacency list of graph
vector<bool> visited;
vi tin, low;
int timer;
void dfs(int v, int p = -1) {
    visited[v] = true;
    tin[v] = low[v] = timer++;
    int children=0;
    for (int to : adj[v]) {
        if (to == p) continue;
    }
}
```

```
if (visited[to]) {
               low[v] = min(low[v], tin[to]);
          } else {
               dfs(to, v):
               low[v] = min(low[v], low[to]);
17
               if (low[to] >= tin[v] && p!=-1)
                   IS_CUTPOINT(v);
               ++children;
          }
20
       if(p == -1 && children > 1)
22
           IS CUTPOINT (v):
23
24 }
25 void find_cutpoints() {
      timer = 0;
      visited.assign(n, false);
27
      tin.assign(n, -1);
      low.assign(n, -1);
      for (int i = 0; i < n; ++i) {</pre>
          if (!visited[i]) dfs (i):
31
32
33 }
```

3.1.15 Topological sort

```
1 #include "header.h"
2 int n: // number of vertices
3 vvi adj; // adjacency list of graph
4 vector < bool > visited;
5 vi ans:
6 void dfs(int v) {
      visited[v] = true;
      for (int u : adj[v]) {
          if (!visited[u]) dfs(u);
      ans.push_back(v);
11
13 void topological sort() {
      visited.assign(n, false);
      ans.clear();
      for (int i = 0: i < n: ++i) {</pre>
16
          if (!visited[i]) dfs(i);
18
      reverse(ans.begin(), ans.end()):
19
20 }
```

3.1.16 Bellmann-Ford Same as Dijkstra but allows neg. edges

```
dist.assign(e.size(), INF);
    cyc.assign(e.size(), false); // true when u is
        in a <0 cvcle
    dist[source] = 0:
    for (int iter = 0; iter < e.size() - 1; ++iter)</pre>
        {
      bool relax = false;
      for (int u = 0; u < e.size(); ++u)</pre>
        if (dist[u] == INF) continue:
        else for (auto &e : e[u])
          if(dist[u]+e.second < dist[e.first])</pre>
             dist[e.first] = dist[u]+e.second. relax
      if(!relax) break;
15
    bool ch = true;
16
    while (ch) {
                         // keep going untill no
        more changes
      ch = false:
                         // set dist to -INF when in
            cvcle
      for (int u = 0; u < e.size(); ++u)</pre>
        if (dist[u] == INF) continue:
        else for (auto &e : e[u])
          if (dist[e.first] > dist[u] + e.second
            && !cyc[e.first]) {
            dist[e.first] = -INF;
            ch = true: //return true for cvcle
                 detection only
            cyc[e.first] = true;
28
29 }
```

3.1.17 Ford-Fulkerson Basic Max. flow

```
1 #include "header.h"
2 #define V 6 // Num. of vertices in given graph
4 /* Returns true if there is a path from source 's
     , to sink
5 't' in residual graph. Also fills parent[] to
      store the
6 path */
7 bool bfs(int rGraph[V][V], int s, int t, int
     parent[]) {
   bool visited[V];
   memset(visited, 0, sizeof(visited));
   queue < int > q;
   q.push(s);
   visited[s] = true;
   parent[s] = -1:
   // Standard BFS Loop
   while (!a.emptv()) {
```

```
int u = q.front();
      q.pop();
      for (int v = 0: v < V: v++) {
        if (visited[v] == false && rGraph[u][v] >
          if (v == t) {
            parent[v] = u;
            return true:
          a.push(v):
          parent[v] = u:
          visited[v] = true;
    }
    return false:
35 // Returns the maximum flow from s to t in the
      given graph
36 int fordFulkerson(int graph[V][V], int s, int t)
    int u, v;
    int rGraph[V]
    for (u = 0: u < V: u++)
      for (v = 0; v < V; v++)
        rGraph[u][v] = graph[u][v];
    int parent[V]; // This array is filled by BFS
        and to
          // store path
    int max_flow = 0; // There is no flow initially
    while (bfs(rGraph, s, t, parent)) {
      int path_flow = INT_MAX;
      for (v = t; v != s; v = parent[v]) {
        u = parent[v]:
        path_flow = min(path_flow, rGraph[u][v]);
52
      for (v = t; v != s; v = parent[v]) {
        u = parent[v]:
        rGraph[u][v] -= path_flow;
        rGraph[v][u] += path_flow;
      max_flow += path_flow;
    }
    return max flow:
62 }
```

3.2 Dynamic Programming

3.2.1 Longest Incr. Subseq.

```
1 #include "header.h"
2 template < class T>
3 vector <T> index path lis(vector <T>& nums) {
    int n = nums.size();
    vector <T> sub;
      vector < int > subIndex:
    vector <T> path(n, -1);
    for (int i = 0; i < n; ++i) {</pre>
        if (sub.empty() || sub[sub.size() - 1] <</pre>
            nums[i]) {
      path[i] = sub.empty() ? -1 : subIndex[sub.
          size() - 1];
      sub.push_back(nums[i]);
      subIndex.push back(i):
       } else {
13
      int idx = lower_bound(sub.begin(), sub.end(),
           nums[i]) - sub.begin();
      path[i] = idx == 0 ? -1 : subIndex[idx - 1];
      sub[idx] = nums[i]:
      subIndex[idx] = i;
    vector <T> ans;
    int t = subIndex[subIndex.size() - 1];
    while (t != -1) {
        ans.push_back(t);
        t = path[t]:
    reverse(ans.begin(), ans.end());
    return ans:
29 // Length only
30 template < class T>
31 int length_lis(vector <T> &a) {
    set <T> st:
    typename set<T>::iterator it;
    for (int i = 0; i < a.size(); ++i) {</pre>
      it = st.lower bound(a[i]):
      if (it != st.end()) st.erase(it);
      st.insert(a[i]):
    return st.size();
40 }
```

3.2.2 0-1 Knapsack

```
1 #include "header.h"
2 // given a number of coins, calculate all
        possible distinct sums
3 int main() {
4    int n;
5    vi coins(n); // all possible coins to use
6    int sum = 0; // sum of the coins
```

3.2.3 Coin change Number of coins required to achieve a given value

```
1 #include "header.h"
_{2} // Returns total distinct ways to make sum using
      n coins of
3 // different denominations
4 int count(vi& coins, int n, int sum) {
      // 2d dp array where n is the number of coin
      // denominations and sum is the target sum
      vector < vector < int > dp(n + 1, vector < int > (
          sum + 1, 0));
      dp[0][0] = 1;
      for (int i = 1: i <= n: i++) {</pre>
          for (int j = 0; j <= sum; j++) {</pre>
10
11
               // without using the current coin,
               dp[i][i] += dp[i - 1][i];
13
               // using the current coin
               if ((j - coins[i - 1]) >= 0)
                   dp[i][j] += dp[i][j - coins[i -
                       111:
          }
      return dp[n][sum];
20
```

3.3 Trees

3.3.1 Tree diameter

```
#include "header.h"
const int mxN = 2e5 + 5;
int n, d[mxN]; // distance array
vi adj[mxN]; // tree adjacency list
void dfs(int s, int e) {
d[s] = 1 + d[e]; // recursively calculate
the distance from the starting node to each
node
```

```
for (auto u : adj[s]) { // for each adjacent
      if (u != e) dfs(u, s): // don't move
         backwards in the tree
9 }
10 }
11 int main() {
12 // read input, create adj list
    dfs(0, -1):
                                 // first dfs call
         to find farthest node from arbitrary node
    dfs(distance(d, max_element(d, d + n)), -1);
        // second dfs call to find farthest node
        from that one
15 cout << *max_element(d, d + n) - 1 << '\n'; //
        distance from second node to farthest is
        the diameter
16 }
```

3.3.2 Tree Node Count

3.4 Num. Th. / Comb.

3.4.1 Basic stuff

```
1 #include "header.h"
2 ll gcd(ll a, ll b) { while (b) { a %= b; swap(a, b); } return a; }
3 ll lcm(ll a, ll b) { return (a / gcd(a, b)) * b; }
4 ll mod(ll a, ll b) { return ((a % b) + b) % b; }
5 // Finds x, y s.t. ax + by = d = gcd(a, b).
6 void extended_euclid(ll a, ll b, ll &x, ll &y, ll &d) {
7 ll xx = y = 0;
8 ll yy = x = 1;
9 while (b) {
10 ll q = a / b;
11 t = b; b = a % b; a = t;
```

```
t = xx; xx = x - q * xx; x = t;
      t = yy; yy = y - q * yy; y = t;
16 }
     solves ab = 1 \pmod{n}, -1 on failure
    mod_inverse(ll a, ll n) {
    11 x, y, d;
    extended_euclid(a, n, x, y, d);
    return (d > 1 ? -1 : mod(x, n));
23 // All modular inverses of [1..n] mod P in O(n)
24 vi inverses(ll n, ll P) {
    vi I(n+1, 1LL);
    for (11 i = 2; i <= n; ++i)
     I[i] = mod(-(P/i) * I[P\%i], P);
29 }
30 // (a*b)%m
31 ll mulmod(ll a, ll b, ll m){
  11 x = 0. v = a\%m:
    while(b>0){
      if(b\&1) x = (x+y)\%m;
      y = (2*y)\%m, b /= 2;
   return x % m;
39 // Finds b^e % m in O(lg n) time, ensure that b <
       m to avoid overflow!
40 ll powmod(ll b, ll e, ll m) {
    11 p = e < 2 ? 1 : powmod((b*b)%m, e/2, m);
   return e&1 ? p*b%m : p;
44 // Solve ax + by = c. returns false on failure.
45 bool linear_diophantine(ll a, ll b, ll c, ll &x,
      11 &v) {
   11 d = gcd(a, b);
   if (c % d) {
     return false;
      x = c / d * mod_inverse(a / d, b / d);
      y = (c - a * x) / b;
      return true:
54 }
```

3.4.2 Mod. exponentiation Or use pow() in python

```
1 #include "header.h"
2 ll mod_pow(ll base, ll exp, ll mod) {
3    if (mod == 1) return 0;
4    if (exp == 0) return 1;
5    if (exp == 1) return base;
```

```
11 res = 1;
8 base %= mod;
9 while (exp) {
10    if (exp % 2 == 1) res = (res * base) % mod;
11    exp >>= 1;
12    base = (base * base) % mod;
13  }
14
15    return res % mod;
16 }
```

3.4.3 GCD Or math.gcd in python, std::gcd in C++

```
#include "header.h"
2 ll gcd(ll a, ll b) {
3   if (a == 0) return b;
4   return gcd(b % a, a);
5 }
```

3.4.4 Sieve of Eratosthenes

3.4.5 Fibonacci % prime

3.4.6 nCk % prime

```
1 #include "header.h"
2 ll binom(ll n, ll k) {
      ll ans = 1:
      for (ll i = 1: i \leq min(k,n-k): ++i) ans = ans
           *(n+1-i)/i;
      return ans:
6 }
7 ll mod_nCk(ll n, ll k, ll p ){
      11 ans = 1:
      while(n){
          ll np = n\%p, kp = k\%p;
          if(kp > np) return 0:
          ans *= binom(np,kp);
          n /= p; k /= p;
14
      return ans;
```

3.4.7 Chin. rem. th.

```
1 #include "header.h"
2 #include "elementary.cpp"
_3 // Solves x = a1 mod m1, x = a2 mod m2, x is
      unique modulo lcm(m1, m2).
4 // Returns {0, -1} on failure, {x, lcm(m1, m2)}
      otherwise.
5 pair<11, 11> crt(11 a1, 11 m1, 11 a2, 11 m2) {
6 ll s, t, d;
    extended euclid(m1, m2, s, t, d):
    if (a1 % d != a2 % d) return {0, -1};
    return {mod(s*a2 %m2 * m1 + t*a1 %m1 * m2, m1 *
         m2) / d. m1 / d * m2:
12 // Solves x = ai mod mi. x is unique modulo lcm
13 // Returns {0, -1} on failure, {x, lcm mi}
      otherwise.
14 pair<ll, ll> crt(vector<ll> &a, vector<ll> &m) {
pair<11. 11> res = {a[0], m[0]}:
   for (ull i = 1; i < a.size(); ++i) {</pre>
      res = crt(res.first. res.second. mod(a[i]. m[
          i]), m[i]);
      if (res.second == -1) break;
20
    return res;
```

3.5 Strings

3.5.1 Z alg. KMP alternative

```
#include "../header.h"
void Z_algorithm(const string &s, vi &Z) {
    Z.assign(s.length(), -1);
    int L = 0, R = 0, n = s.length();
    for (int i = 1; i < n; ++i) {
        if (i > R) {
            L = R = i;
            while (R < n && s[R - L] == s[R]) R++;
            Z[i] = R - L; R--;
        } else if (Z[i - L] >= R - i + 1) {
            L = i;
            while (R < n && s[R - L] == s[R]) R++;
            Z[i] = R - L; R--;
        } else if (Z[i - L] >= R - i + 1) {
            L = i;
            while (R < n && s[R - L] == s[R]) R++;
            Z[i] = R - L; R--;
        } else Z[i] = Z[i - L];
    }
}</pre>
```

3.5.2 KMP

```
1 #include "header.h"
void compute_prefix_function(string &w, vi &
      prefix) {
g prefix.assign(w.length(), 0);
    int k = prefix[0] = -1;
    for(int i = 1; i < w.length(); ++i) {</pre>
      while (k >= 0 \&\& w[k + 1] != w[i]) k = prefix[
      if(w[k + 1] == w[i]) k++;
      prefix[i] = k:
10
12 void knuth_morris_pratt(string &s, string &w) {
    int q = -1;
    vi prefix:
    compute_prefix_function(w, prefix);
    for(int i = 0: i < s.length(): ++i) {</pre>
      while (q \ge 0 \&\& w[q + 1] != s[i]) q = prefix[
      if(w[q + 1] == s[i]) q++;
      if(q + 1 == w.length()) {
19
       // Match at position (i - w.length() + 1)
        a = prefix[a]:
   }
23
```

3.5.3 Aho-Corasick Also can be used as Knuth-Morris-Pratt algorithm

```
1 #include "header.h"
```

```
3 map < char, int > cti;
4 int cti_size;
5 template <int ALPHABET_SIZE, int (*mp)(char)>
6 struct AC FSM {
    struct Node {
      int child[ALPHABET_SIZE], failure = 0,
          match_par = -1;
      Node() { for (int i = 0: i < ALPHABET SIZE:
          ++i) child[i] = -1; }
    }:
1.1
    vector < Node > a:
    vector < string > & words;
    AC_FSM(vector<string> &words) : words(words) {
      a.push_back(Node());
      construct_automaton();
16
17
    void construct_automaton() {
      for (int w = 0, n = 0; w < words.size(); ++w,</pre>
           n = 0) {
        for (int i = 0; i < words[w].size(); ++i) {</pre>
          if (a[n].child[mp(words[w][i])] == -1) {
            a[n].child[mp(words[w][i])] = a.size();
            a.push_back(Node());
          n = a[n].child[mp(words[w][i])];
        a[n].match.push_back(w);
28
      queue < int > q:
      for (int k = 0; k < ALPHABET_SIZE; ++k) {</pre>
        if (a[0].child[k] == -1) a[0].child[k] = 0;
        else if (a[0].child[k] > 0) {
          a[a[0].child[k]].failure = 0;
          g.push(a[0].child[k]);
      }
36
      while (!q.empty()) {
        int r = q.front(); q.pop();
        for (int k = 0, arck; k < ALPHABET_SIZE; ++</pre>
          if ((arck = a[r].child[k]) != -1) {
            q.push(arck);
            int v = a[r].failure:
42
            while (a[v].child[k] == -1) v = a[v].
                failure:
44
            a[arck].failure = a[v].child[k];
            a[arck].match_par = a[v].child[k];
             while (a[arck].match_par != -1
                && a[a[arck].match_par].match.empty
              a[arck].match par = a[a[arck].
                  match_par].match_par;
          }
        }
```

```
52
     void aho corasick(string &sentence, vvi &
         matches){
       matches.assign(words.size(), vi());
       int state = 0, ss = 0;
       for (int i = 0; i < sentence.length(); ++i,</pre>
           ss = state) {
         while (a[ss].child[mp(sentence[i])] == -1)
           ss = a[ss].failure;
         state = a[state].child[mp(sentence[i])]
             = a[ss].child[mp(sentence[i])];
         for (ss = state; ss != -1; ss = a[ss].
             match par)
           for (int w : a[ss].match)
             matches[w].push_back(i + 1 - words[w].
                 length());
    }
67 int char_to_int(char c) {
    return cti[c]:
70 int main() {
     11 n:
     while(getline(cin, line)) {
       stringstream ss(line):
       ss >> n:
       vector < string > patterns(n);
       for (auto& p: patterns) getline(cin, p);
       string text;
       getline(cin. text):
       cti = {}, cti_size = 0;
       for (auto c: text) {
        if (not in(c, cti)) {
           cti[c] = cti_size++;
88
       for (auto& p: patterns) {
       for (auto c: p) {
          if (not in(c, cti)) {
             cti[c] = cti size++:
       }
       vvi matches:
       AC FSM <128+1. char to int > ac fms(patterns):
       ac_fms.aho_corasick(text, matches);
       for (auto& x: matches) cout << x << endl;</pre>
101
    }
```

```
102 103 }
```

3.5.4 Long. palin. subs Manacher - O(n)

```
1 #include "header.h"
void manacher(string &s, vi &pal) {
    int n = s.length(), i = 1, 1, r;
    pal.assign(2 * n + 1, 0);
    while (i < 2 * n + 1) {
      if ((i&1) && pal[i] == 0) pal[i] = 1;
      l = i / 2 - pal[i] / 2; r = (i-1) / 2 + pal[i
         1 / 2:
      while (1 - 1 >= 0 \&\& r + 1 < n \&\& s[1 - 1] ==
           s[r + 1]
        --1, ++r, pal[i] += 2;
11
      for (1 = i - 1, r = i + 1; 1 >= 0 && r < 2 *
          n + 1; --1, ++r) {
        if (1 <= i - pal[i]) break:</pre>
        if (1 / 2 - pal[1] / 2 > i / 2 - pal[i] /
          pal[r] = pal[1];
        else { if (1 \ge 0)
            pal[r] = min(pal[l], i + pal[i] - r);
          break:
      i = r;
```

3.6 Geometry

3.6.1 essentials.cpp

```
C operator (const P &p) const { return x*p.y -
         p.x*v; }
    P perp() const { return P{y, -x}; }
    C lensq() const { return x*x + v*v: }
    ld len() const { return sqrt((ld)lensq()); }
    static ld dist(const P &p1, const P &p2) {
      return (p1-p2).len(); }
    bool operator == (const P &r) const {
     return ((*this)-r).lensq() <= EPS*EPS: }
21 C det(P p1, P p2) { return p1^p2; }
22 C det(P p1, P p2, P o) { return det(p1-o, p2-o);
23 C det(const vector <P> &ps) {
for(auto &p : ps) sum += det(p, prev), prev = p
    return sum;
28 // Careful with division by two and C=11
29 C area(P p1, P p2, P p3) { return abs(det(p1, p2,
       p3))/C(2): 
30 C area(const vector<P> &poly) { return abs(det(
      poly))/C(2); }
31 int sign(C c) { return (c > C(0)) - (c < C(0)); }
32 int ccw(P p1, P p2, P o) { return sign(det(p1, p2
      , o)); }
_{34} // Only well defined for C = ld.
35 P unit(const P &p) { return p / p.len(); }
36 P rotate(P p, ld a) { return P{p.x*cos(a)-p.y*sin
      (a), p.x*sin(a)+p.y*cos(a)}; }
```

3.6.2 Two segs. itersec.

3.6.3 Convex Hull

```
1 #include "header.h"
2 #include "essentials.cpp"
3 struct ConvexHull { // O(n lg n) monotone chain.
    vector < size t > h. c: // Indices of the hull
        are in 'h'. ccw.
    const vector <P> &p;
    ConvexHull(const vector <P> &_p) : n(_p.size()),
         c(n), p(_p) {
      std::iota(c.begin(), c.end(), 0);
      std::sort(c.begin(), c.end(), [this](size_t 1
          , size_t r) -> bool { return p[1].x != p[
          r].x ? p[1].x < p[r].x : p[1].y < p[r].y;
      c.erase(std::unique(c.begin(), c.end(), [this
          ](size_t l, size_t r) { return p[l] == p[
          r]; }), c.end());
      for (size_t s = 1, r = 0; r < 2; ++r, s = h.
          size()) {
        for (size_t i : c) {
          while (h.size() > s && ccw(p[h.end()
              [-2]], p[h.end()[-1]], p[i]) <= 0)
            h.pop_back();
          h.push_back(i);
        reverse(c.begin(), c.end());
      if (h.size() > 1) h.pop_back();
19
    size_t size() const { return h.size(); }
    template <class T, void U(const P &, const P &,
         const P &. T &)>
    void rotating_calipers(T &ans) {
      if (size() <= 2)
        U(p[h[0]], p[h.back()], p[h.back()], ans);
25
        for (size_t i = 0, j = 1, s = size(); i < 2</pre>
             * s: ++i) {
          while (\det(p[h[(i + 1) \% s]) - p[h[i \% s
              ]], p[h[(j + 1) \% s]] - p[h[j]]) >=
              0)
            j = (j + 1) \% s;
          U(p[h[i % s]], p[h[(i + 1) % s]], p[h[j
              ]], ans);
        }
  }
32
34 // Example: furthest pair of points. Now set ans
      = OLL and call
35 // ConvexHull(pts).rotating_calipers<11, update>(
36 void update(const P &p1, const P &p2, const P &o,
       11 &ans) {
    ans = max(ans, (11)max((p1 - o).lensq(), (p2 -
```

```
o).lensq()));
38 }
39 int main() {
    ios::svnc with stdio(false): // do not use
        cout + printf
    cin.tie(NULL);
    int n;
    cin >> n:
    while (n) {
      vector <P> ps;
          int x, y;
      for (int i = 0; i < n; i++) {</pre>
              cin >> x >> y;
               ps.push_back({x, y});
51
52
          ConvexHull ch(ps);
53
          cout << ch.h.size() << endl;</pre>
          for(auto& p: ch.h) {
               cout << ps[p].x << "" << ps[p].y <<
                   endl:
      cin >> n;
    return 0;
```

3.7 Other Algorithms

3.7.1 2-sat

```
1 #include "../header.h"
2 #include "../Graphs/tarjan.cpp"
3 struct TwoSAT {
   int n:
   vvi imp; // implication graph
   Tarjan tj;
   TwoSAT(int _n) : n(_n), imp(2 * _n, vi()), tj(
       imp) { }
   // Only copy the needed functions:
   void add_implies(int c1, bool v1, int c2, bool
     int u = 2 * c1 + (v1 ? 1 : 0),
      v = 2 * c2 + (v2 ? 1 : 0):
     imp[u].push_back(v); // u => v
     imp[v^1].push_back(u^1); // -v => -u
   void add_equivalence(int c1, bool v1, int c2,
       bool v2) {
      add_implies(c1, v1, c2, v2);
```

```
add_implies(c2, v2, c1, v1);
    void add_or(int c1, bool v1, int c2, bool v2) {
21
      add implies(c1. !v1. c2. v2):
23
    void add_and(int c1, bool v1, int c2, bool v2)
      add_true(c1, v1); add_true(c2, v2);
26
    void add_xor(int c1, bool v1, int c2, bool v2)
      add_or(c1, v1, c2, v2);
      add_or(c1, !v1, c2, !v2);
   }
30
    void add_true(int c1, bool v1) {
      add_implies(c1, !v1, c1, v1);
    // on true: a contains an assignment.
    // on false: no assignment exists.
    bool solve(vb &a) {
      vi com:
      tj.find_sccs(com);
      for (int i = 0; i < n; ++i)</pre>
        if (com[2 * i] == com[2 * i + 1])
          return false;
42
43
      vvi bvcom(com.size()):
      for (int i = 0; i < 2 * n; ++i)
45
        bycom[com[i]].push_back(i);
      a.assign(n, false);
      vb vis(n, false);
      for(auto &&component : bycom){
       for (int u : component) {
          if (vis[u / 2]) continue;
          vis[u / 2] = true;
          a[u / 2] = (u \% 2 == 1):
        }
      return true;
```

3.7.2 Matrix Solve

```
8 T ReducedRowEchelonForm(M<R.C> &m. int rows) {
      // return the determinant
9 int r = 0: T det = 1:
                                      // MODIFIES
        the input
    for(int c = 0; c < rows && r < rows; c++) {</pre>
      for(int i=r+1; i<rows; i++) if(abs(m[i][c]) >
           abs(m[p][c])) p=i;
      if(abs(m[p][c]) < EPS){ det = 0; continue; }</pre>
      swap(m[p], m[r]);   det = -det;
      T s = 1.0 / m[r][c], t; det *= m[r][c];
      term in row 1
      REP(i,rows) if (i!=r){ t = m[i][c]; REP(j,C)
          m[i][j] -= t*m[r][j]; }
      ++r:
    }
    return det;
22 bool error, inconst; // error => multiple or
23 template <int R.int C> // Mx = a: M:R*R, v:R*C =>
24 M<R,C> solve(const M<R,R> &m, const M<R,C> &a,
      int rows){
    M < R, R + C > q;
    REP(r.rows){
      REP(c,rows) q[r][c] = m[r][c];
      REP(c,C) q[r][R+c] = a[r][c];
   }
29
    ReducedRowEchelonForm <R,R+C>(q,rows);
    M<R,C> sol; error = false, inconst = false;
    REP(c,C) for(auto j = rows-1; j >= 0; --j){
      T t=0; bool allzero=true;
      for(auto k = j+1: k < rows: ++k)
        t += q[j][k]*sol[k][c], allzero &= abs(q[j
            ][k]) < EPS;
      if(abs(q[j][j]) < EPS)</pre>
        error = true, inconst |= allzero && abs(q[i
            ][R+c]) > EPS;
      else sol[j][c] = (q[j][R+c] - t) / q[j][j];
          // usually q[i][i]=1
    return sol:
41 }
```

3.7.3 Matrix Exp.

```
array <array <T,N>,N> m;
    M() \{ ITERATE_MATRIX(N) m[r][c] = 0; \}
    static M id() {
      M I: for (int i = 0: i < N: ++i) I.m[i][i] =
          1; return I;
    M operator*(const M &rhs) const {
12
      ITERATE MATRIX(N) for (int i = 0: i < N: ++i)
          out.m[r][c] += m[r][i] * rhs.m[i][c];
      return out:
15
16
17
    M raise(ll n) const {
      if(n == 0) return id():
      if(n == 1) return *this;
      auto r = (*this**this).raise(n / 2);
      return (n%2 ? *this*r : r):
22
23 };
```

3.7.4 Finite field For FFT

```
1 #include "header.h"
2 #include "../Number_Theory/elementary.cpp"
3 template <11 p,11 w> // prime, primitive root
4 struct Field { using T = Field; ll x; Field(ll x
      =0): x\{x\}\{
   T operator+(T r) const { return {(x+r.x)%p}; }
    T operator - (T r) const { return {(x-r.x+p)%p};
    T operator*(T r) const { return {(x*r.x)%p}; }
    T operator/(T r) const { return (*this)*r.inv()
    T inv() const { return {mod_inverse(x,p)}; }
    static T root(ll k) { assert( (p-1)%k==0 );
        // (p-1) \% k == 0?
      auto r = powmod(w,(p-1)/abs(k),p);
                                               // k-
          th root of unity
      return k>=0 ? T{r} : T{r}.inv();
    bool zero() const { return x == OLL; }
15 };
16 using F1 = Field<1004535809.3 >:
using F2 = Field<1107296257,10>; // 1<<30 + 1<<25
18 using F3 = Field < 2281701377,3 >; // 1 < < 31 + 1 < < 27
       + 1
```

3.7.5 Complex field For FFR

```
1 #include "header.h"
2 const double m_pi = M_PIf64x;
3 struct Complex { using T = Complex; double u,v;
```

```
Complex (double u=0, double v=0) : u{u}, v{v} {}}
    T operator+(T r) const { return {u+r.u, v+r.v};
    T operator-(T r) const { return {u-r.u, v-r.v};
    T operator*(T r) const { return {u*r.u - v*r.v,
         u*r.v + v*r.u: }
    T operator/(T r) const {
      auto norm = r.u*r.u+r.v*r.v:
      return {(u*r.u + v*r.v)/norm, (v*r.u - u*r.v)
          /norm}:
11
    T operator*(double r) const { return T{u*r, v*r
   T operator/(double r) const { return T{u/r, v/r
        }: }
   T inv() const { return T{1,0}/ *this; }
    T conj() const { return T{u, -v}; }
    static T root(ll k){ return {cos(2*m_pi/k), sin
        (2*m_pi/k); }
    bool zero() const { return max(abs(u), abs(v))
        < 1e-6: }
18 };
```

3.7.6 FFT

```
1 #include "header.h"
2 #include "complex field.cpp"
3 #include "fin_field.cpp"
4 void brinc(int &x, int k) {
5 int i = k - 1, s = 1 << i;</pre>
  x ^= s;
  if ((x & s) != s) {
      --i; s >>= 1;
      while (i >= 0 && ((x & s) == s))
        x = x &^{\sim} s, --i, s >>= 1;
      if (i >= 0) x |= s;
11
12 }
using T = Complex; // using T=F1,F2,F3
15 vector <T> roots:
16 void root_cache(int N) {
    if (N == (int)roots.size()) return;
    roots.assign(N, T{0});
    for (int i = 0; i < N; ++i)</pre>
      roots[i] = ((i\&-i) == i)
20
        ? T\{\cos(2.0*m_pi*i/N), \sin(2.0*m_pi*i/N)\}
        : roots[i&-i] * roots[i-(i&-i)];
24 void fft(vector<T> &A, int p, bool inv = false) {
   for(int i = 0, r = 0; i < N; ++i, brinc(r, p))
      if (i < r) swap(A[i], A[r]);</pre>
28 // Uncomment to precompute roots (for T=Complex)
      . Slower but more precise.
```

```
29 // root cache(N):
30 //
                       , --sh
            , sh=p-1
   for (int m = 2: m <= N: m <<= 1) {
      T w. w m = T::root(inv ? -m : m):
      for (int k = 0; k < N; k += m) {
        w = T\{1\};
        for (int j = 0; j < m/2; ++j) {
            T w = (!inv ? roots[j << sh] : roots[j <<
36 //
      shl.coni()):
          T t = w * A[k + j + m/2];
          A[k + j + m/2] = A[k + j] - t;
          A[k + j] = A[k + j] + t;
          w = w * w_m;
      }
    }
    if(inv){ T inverse = T(N).inv(); for(auto &x :
        A) x = x*inverse;
_{
m 46} // convolution leaves A and B in frequency domain
47 // C may be equal to A or B for in-place
      convolution
48 void convolution(vector<T> &A, vector<T> &B,
      vector<T> &C){
    int s = A.size() + B.size() - 1;
    int q = 32 - __builtin_clz(s-1), N=1<<q; //</pre>
        fails if s=1
    A.resize(N,{}); B.resize(N,{}); C.resize(N,{});
    fft(A, q, false); fft(B, q, false);
    for (int i = 0; i < N; ++i) C[i] = A[i] * B[i];
    fft(C, q, true); C.resize(s);
56 void square_inplace(vector<T> &A) {
   int s = 2*A.size()-1, q = 32 - __builtin_clz(s
        -1), N=1<<a:
    A.resize(N,{}); fft(A, q, false);
    for (auto &x : A) x = x*x;
    fft(A, q, true); A.resize(s);
```

3.7.7 Polyn. inv. div.

```
9 vector<T> inverse(const vector<T> &A, int n) {
     vector <T> Ai{A[0].inv()};
     for (int k = 0; (1<<k) < n; ++k) {
       vector \langle T \rangle As (4 << k, T(0)). Ais (4 << k, T(0)):
       copy_into(A, As, 2<<k); copy_into(Ai, Ais, Ai
13
           .size());
       fft(As, k+2, false); fft(Ais, k+2, false);
       for (int i = 0; i < (4<<k); ++i) As[i] = As[i
15
           1*Ais[i]*Ais[i]:
       fft(As, k+2, true); Ai.resize(2<<k, {});</pre>
16
       for (int i = 0; i < (2 << k); ++i) Ai[i] = T(2)
17
            * Ai[i] - As[i]:
    Ai.resize(n);
    return Ai;
21 }
22 // Polynomial division. Returns {Q, R} such that
       A = QB+R, deg R < deg B.
23 // Requires that the leading term of B is nonzero
24 pair < vector < T > , vector < T >> divmod (const vector < T >
       &A. const vector <T> &B) {
     size_t n = A.size()-1, m = B.size()-1;
    if (n < m) return {vector < T > (1, T(0)), A};
    vector\langle T \rangle X(A), Y(B), Q, R;
    convolution(rev(X), Y = inverse(rev(Y), n-m+1),
    Q.resize(n-m+1); rev(Q);
31
    X.resize(Q.size()), copy_into(Q, X, Q.size());
    Y.resize(B.size()), copy_into(B, Y, B.size());
    convolution(X, Y, X);
    R.resize(m), copy_into(A, R, m);
    for (size_t i = 0; i < m; ++i) R[i] = R[i] - X[</pre>
        i];
    while (R.size() > 1 && R.back().zero()) R.
         pop_back();
    return {Q, R};
39
41 vector <T > mod(const vector <T > &A, const vector <T >
    return divmod(A. B).second:
43 }
```

3.7.8 Linear recurs. Given a linear recurrence of the form

$$a_n = \sum_{i=0}^{k-1} c_i a_{n-i-1}$$

this code computes a_n in $O(k \log k \log n)$ time.

```
1 #include "header.h"
2 #include "poly.cpp"
3 // x^k \mod f
4 vector <T> xmod(const vector <T> f, ll k) {
  vectorT> r\{T(1)\};
    for (int b = 62; b >= 0; --b) {
      if (r.size() > 1)
        square_inplace(r), r = mod(r, f);
      if ((k>>b)&1) {
        r.insert(r.begin(), T(0));
        if (r.size() == f.size()) {
11
         T c = r.back() / f.back();
12
13
          for (size_t i = 0; i < f.size(); ++i)</pre>
            r[i] = r[i] - c * f[i];
          r.pop_back();
      }
    return r;
_{21} // Given A[0,k) and C[0, k), computes the n-th
      term of:
22 // A[n] = \sum_{i=1}^{n} C[i] * A[n-i-1]
23 T nth_term(const vector <T> &A, const vector <T> &C
      , 11 n) {
    int k = (int)A.size();
    if (n < k) return A[n]:
    vector <T> f(k+1, T{1});
    for (int i = 0: i < k: ++i)
     f[i] = T\{-1\} * C[k-i-1];
    f = xmod(f, n);
    T r = T\{0\};
    for (int i = 0: i < k: ++i)
      r = r + f[i] * A[i];
    return r;
```

3.7.9 Convolution Precise up to 9e15

```
fft(Ac, q, false); fft(Bc, q, false);
    for (int i = 0, j = 0; i < N; ++i, j = (N-1)&(N
      T as = (Ac[i] + Ac[j].conj()) / 2;
      T = (Ac[i] - Ac[i].coni()) / T{0, 2};
      T bs = (Bc[i] + Bc[j].conj()) / 2;
      T bl = (Bc[i] - Bc[j].conj()) / T{0, 2};
      R1[i] = as*bs + al*bl*T{0,1}, R2[i] = as*bl +
17
    fft(R1, q, true); fft(R2, q, true);
    11 p15 = (1LL << 15) \% MOD, p30 = (1LL << 30) \% MOD; C.
        resize(s);
    for (int i = 0; i < s; ++i) {
      11 1 = llround(R1[i].u), m = llround(R2[i].u)
          , h = llround(R1[i].v);
      C[i] = (1 + m*p15 + h*p30) \% MOD;
23 }
```

3.7.10 Partitions of n Finds all possible partitions of a number

```
1 #include "header.h"
void printArray(int p[], int n) {
   for (int i = 0; i < n; i++)</pre>
      cout << p[i] << "";
   cout << endl;</pre>
8 void printAllUniqueParts(int n) {
   int p[n]; // An array to store a partition
   int k = 0; // Index of last element in a
   p[k] = n; // Initialize first partition as
        number itself
   // This loop first prints current partition
       then generates next
   // partition. The loop stops when the current
       partition has all 1s
   while (true) {
     printArray(p, k + 1);
     // Find the rightmost non-one value in p[].
          Also, update the
     // rem_val so that we know how much value can
          be accommodated
      int rem_val = 0;
      while (k >= 0 \&\& p[k] == 1) {
       rem_val += p[k];
```

```
// if k < 0, all the values are 1 so there
          are no more partitions
      if (k < 0) return:
28
      // Decrease the p[k] found above and adjust
29
          the rem val
      p[k]--:
      rem_val++;
31
32
      // If rem_val is more, then the sorted order
          is violated. Divide
      // rem_val in different values of size p[k]
          and copy these values at
      // different positions after p[k]
      while (rem_val > p[k]) {
        p[k + 1] = p[k];
        rem_val = rem_val - p[k];
      }
40
41
      // Copy rem_val to next position and
42
          increment position
      p[k + 1] = rem_val;
      k++:
45
46 }
```

3.8 Other Data Structures

3.8.1 Disjoint set (i.e. union-find)

```
1 template <typename T>
2 class DisjointSet {
      typedef T * iterator;
      T *parent, n, *rank;
      public:
          // O(n), assumes nodes are [0, n)
          DisjointSet(T n) {
              this->parent = new T[n];
              this -> n = n:
              this->rank = new T[n];
              for (T i = 0: i < n: i++) {
                  parent[i] = i;
                  rank[i] = 0:
              }
          }
17
18
          // O(log n)
          T find_set(T x) {
19
              if (x == parent[x]) return x;
              return parent[x] = find_set(parent[x
                  ]);
          }
```

3.8.2 Fenwick tree (i.e. BIT) eff. update + prefix sum calc.

```
1 #include "header.h"
2 #define maxn 200010
3 int t,n,m,tree[maxn],p[maxn];
5 void update(int k. int z) {
      while (k <= maxn) {
          tree[k] += z:
          k += k & (-k):
10 }
12 int sum(int k) {
      int ans = 0:
      while(k) {
          ans += tree[k]:
          k = k & (-k);
16
17
      return ans:
```

3.8.3 Fenwick2d tree

```
#include "header.h"
template <class T>
struct FenwickTree2D {
  vector< vector<T> > tree;
  int n;
FenwickTree2D(int n) : n(n) { tree.assign(n +
        1, vector<T>(n + 1, 0)); }
T query(int x1, int y1, int x2, int y2) {
```

```
return query(x2,y2)+query(x1-1,y1-1)-query(x2
          ,v1-1)-query(x1-1,v2);
   T querv(int x, int v) {
10
      T s = 0:
      for (int i = x; i > 0; i -= (i & (-i)))
        for (int j = y; j > 0; j -= (j & (-j)))
          s += tree[i][i]:
      return s:
16
17
    void update(int x, int v, T v) {
      for (int i = x; i \le n; i += (i & (-i)))
        for (int j = v; j \le n; j += (j & (-j)))
          tree[i][i] += v:
21 }
22 };
```

3.8.4 Trie

```
1 #include "header.h"
2 const int ALPHABET SIZE = 26:
3 inline int mp(char c) { return c - 'a'; }
5 struct Node {
    Node* ch[ALPHABET_SIZE];
    bool isleaf = false;
    Node() {
      for(int i = 0; i < ALPHABET_SIZE; ++i) ch[i]</pre>
          = nullptr:
   }
10
    void insert(string &s. int i = 0) {
      if (i == s.length()) isleaf = true;
      else {
       int v = mp(s[i]);
        if (ch[v] == nullptr)
          ch[v] = new Node();
        ch[v] \rightarrow insert(s, i + 1):
18
19
    }
20
    bool contains(string &s, int i = 0) {
      if (i == s.length()) return isleaf;
      else {
        int v = mp(s[i]);
        if (ch[v] == nullptr) return false;
        else return ch[v]->contains(s, i + 1);
28
    }
29
    void cleanup() {
      for (int i = 0; i < ALPHABET_SIZE; ++i)</pre>
        if (ch[i] != nullptr) {
          ch[i]->cleanup():
```

3.8.5 Treap A binary tree whose nodes contain two values, a key and a priority, such that the key keeps the BST property

```
1 #include "header.h"
2 struct Node {
    11 v;
    int sz, pr;
    Node *l = nullptr, *r = nullptr;
    Node(ll val) : v(val), sz(1) { pr = rand(); }
8 int size(Node *p) { return p ? p->sz : 0; }
9 void update(Node* p) {
     if (!p) return;
     p->sz = 1 + size(p->1) + size(p->r);
     // Pull data from children here
13 }
14 void propagate(Node *p) {
     if (!p) return;
     // Push data to children here
18 void merge(Node *&t, Node *1, Node *r) {
    propagate(1), propagate(r);
    if (!1)
                 t = r;
    else if (!r) t = 1;
    else if (1->pr > r->pr)
         merge(1->r, 1->r, r), t = 1;
    else merge(r\rightarrow 1, 1, r\rightarrow 1), t = r;
26 }
27 void spliti(Node *t, Node *&l, Node *&r, int
     propagate(t);
    if (!t) { l = r = nullptr; return; }
    int id = size(t->1):
    if (index <= id) // id \in [index, \infty), so</pre>
         move it right
       spliti(t\rightarrow 1, 1, t\rightarrow 1, index), r = t;
       spliti(t->r, t->r, r, index - id), l = t;
    update(t);
35
36 }
37 void splitv(Node *t, Node *&1, Node *&r, 11 val)
    propagate(t);
    if (!t) { l = r = nullptr; return; }
    if (val \leftarrow t->v) // t->v \in [val, \infty), so
         move it right
       splitv(t\rightarrow 1, 1, t\rightarrow 1, val), r = t;
```

```
42    else
43         splitv(t->r, t->r, r, val), l = t;
44         update(t);
45    }
46    void clean(Node *p) {
47         if (p) { clean(p->l), clean(p->r); delete p; }
48    }
```

4 Other Mathematics

4.1 Helpful functions

4.1.1 Euler's Totient Fucntion $n = p_1^{k_1-1} \cdot (p_1-1) \cdot \dots \cdot p_r^{k_r-1} \cdot (p_r-1)$, where $p_1^{k_1} \cdot \dots \cdot p_r^{k_r}$ is the prime factorization of n.

```
1 # include "header.h"
2 ll phi(ll n) { // \Phi(n)
      11 \text{ ans} = 1;
      for (11 i = 2; i*i <= n; i++) {
          if (n % i == 0) {
               ans *= i-1;
               n /= i;
               while (n \% i == 0) {
                   ans *= i:
                   n /= i;
          }
13
       if (n > 1) ans *= n-1:
       return ans;
16 }
17 vi phis(int n) { // All \Phi(i) up to n
    vi phi(n + 1, OLL);
    iota(phi.begin(), phi.end(), OLL);
    for (11 i = 2LL; i <= n; ++i)
      if (phi[i] == i)
        for (ll j = i; j <= n; j += i)
          phi[j] -= phi[j] / i;
23
    return phi;
25 }
```

Formulas $\Phi(n)$ counts all numbers in $1, \ldots, n-1$ coprime to n. $a^{\varphi(n)} \equiv 1 \mod n$, a and n are coprimes. $\forall e > \log_2 m : n^e \mod m = n^{\Phi(m)+e \mod \Phi(m)} \mod m$. $\gcd(m,n) = 1 \Rightarrow \Phi(m \cdot n) = \Phi(m) \cdot \Phi(n)$.

4.1.2 Pascal's trinagle $\binom{n}{k}$ is k-th element in the n-th row, indexing both from 0

4.2 Theorems and definitions

Fermat's little theorem

$$a^p \equiv a \mod p$$

Subfactorial

$$!n = n! \sum_{i=0}^{n} \frac{(-1)^i}{i!}$$

$$!(0) = 1, !n = n \cdot !(n-1) + (-1)^n$$

Binomials and other partitionings

$$\binom{n}{k} = \binom{n-1}{k} + \binom{n-1}{k-1} = \prod_{i=1}^{k} \frac{n-i+1}{i}$$

This last product may be computed incrementally since any product of k' consecutive values is divisible by k'!. Basic identities: The hockeystick identity:

$$\sum_{k=r}^{n} \binom{k}{r} = \binom{n+1}{r+1}$$

or

$$\sum_{k \le n} \binom{r+k}{k} = \binom{r+n+1}{n}$$

Also

$$\sum_{k=0}^{n} \binom{k}{m} = \binom{n+1}{m+1}$$

$$\sum_{i=0}^{n} \binom{n}{i} = 2^n$$

For $n, m \geq 0$ and p prime: write n, m in base p, i.e. $n = n_k p^k + \dots + n_1 p + n_0$ and $m = m_k p^k + \dots + m_1 p + m_0$. Then by Lucas theorem we have $\binom{n}{m} \equiv \prod_{i=0}^k \binom{n_i}{m_i} \mod p$, with the convention that $n_i < m_i \implies \binom{n_i}{m_i} = 0$.

Fibonacci (See also number theory section)

$$\sum_{0 \le k \le n} \binom{n-k}{k} = F_{n+1}$$

$$F_n = \frac{1}{\sqrt{5}} \left(\frac{1+\sqrt{5}}{2} \right)^n - \frac{1}{\sqrt{5}} \left(\frac{1-\sqrt{5}}{2} \right)^n$$

$$\sum_{i=1}^{n} F_i = F_{n+2} - 1, \ \sum_{i=1}^{n} F_i^2 = F_n F_{n+1}$$

$$\gcd(F_m, F_n) = F_{\gcd(m,n)}$$

$$\gcd(F_n,F_{n+1})=\gcd(F_n,F_{n+2})=1$$

Bit stuff $a+b=a\oplus b+2(a\&b)=a|b+a\&b$. kth bit is set in x iff $x \mod 2^{k-1} \ge 2^k$, or iff $x \mod 2^{k-1}-x \mod 2^k \ne 0$ (i.e. $=2^k$) It comes handy when you need to look at the bits of the numbers which are pair sums or subset sums etc.

$$n \mod 2^i = n\&(2^i - 1).$$

$$\forall k: \ 1 \oplus 2 \oplus \ldots \oplus (4k-1) = 0$$

Stirling's numbers First kind: $S_1(n,k)$ count permutations on n items with k cycles. $S_1(n,k) = S_1(n-1,k-1) + (n-1)S_1(n-1,k)$ with $S_1(0,0) = 1$. Note:

$$\sum_{k=0}^{n} S_1(n,k)x^k = x(x+1)\dots(x+n-1)$$

$$\sum_{k=0}^{n} S_1(n,k) = n!$$

Second kind: $S_2(n, k)$ count partitions of n distinct elements into exactly k non-empty groups.

$$S_2(n,k) = S_2(n-1,k-1) + kS_2(n-1,k)$$

$$S_2(n,1) = S_2(n,n) = 1$$

$$S_2(n,k) = \frac{1}{k!} \sum_{i=0}^{k} (-1)^{k-i} {k \choose i} i^n$$

4.3 Geometry Formulas

$$[ABC] = rs = \frac{1}{2}ab\sin\gamma = \frac{abc}{4R} = \sqrt{s(s-a)(s-b)(s-c)} = \frac{1}{2}\left|(B - A_{\rm g} \operatorname{Gaph}A_{\rm h})^T O^*(2^n)\right| \text{ time.}$$

$$s = \frac{a+b+c}{2} \qquad \qquad 2R = \frac{a}{\sin \alpha}$$
 cosine rule:
$$c^2 = a^2 + b^2 - 2ab\cos \gamma$$
 Euler:
$$1 + CC = V - E + F$$

Pick: Area = interior points +
$$\frac{\text{bour}}{\text{constant}}$$

$$p \cdot q = |p||q|\cos(\theta)$$
 $|p \times q| = |p||q|\sin(\theta)$

Given a non-self-intersecting closed polygon on n vertices, given as (x_i, y_i) , its centroid (C_x, C_y) is given as:

$$C_x = \frac{1}{6A} \sum_{i=0}^{n-1} (x_i + x_{i+1})(x_i y_{i+1} - x_{i+1} y_i), \quad C_y = \frac{1}{6A} \sum_{i=0}^{n-1} (y_i y_{i+1} - y_i)$$

$$A = \frac{1}{2} \sum_{i=0}^{n-1} (x_i y_{i+1} - x_{i+1} y_i) = \text{polygon area}$$

Inclusion-Exclusion For appropriate f compute $\sum_{S\subseteq T} (-1)^{|T\setminus S|} f(S)$, or if only the size of S matters, $\sum_{s=0}^{n} (-1)^{n-s} \binom{n}{s} f(s)$. In some contexts we might use Stirling numbers, not binomial coefficients!

Some useful applications:

Graph coloring Let I(S) count the number of independent sets contained in $S \subseteq V$ ($I(\emptyset) = 1$, $I(S) = I(S \setminus v) + I(S \setminus N(v))$). Let $c_k = \sum_{S \subseteq V} (-1)^{|V \setminus S|} I(S)$. Then V is k-colorable iff v > 0. Thus we can compute the chromatic number of a $A_{S} A_{D} A_{D}^{T} O^{*}(2^{n})$ time.

Burnside's lemma Given a group G acting on a set X, the number of elements in X up to symmetry is

$$\frac{1}{|G|} \sum_{g \in G} |X^g|$$

with X^g the elements of X invariant under g. For example, if f(n) counts "configurations" of some sort of length n, and we want to count them up to rotational symmetry using $G = \mathbb{Z}/n\mathbb{Z}$, then

$$g(n) = \frac{1}{n} \sum_{k=0}^{n-1} f(\gcd(n,k)) = \frac{1}{n} \sum_{k||n} f(k)\phi(n/k)$$

I.e. for coloring with c colors we have $f(k) = k^c$.

Relatedly, in Pólya's enumeration theorem we imagine X as a set of n beads with G permuting the beads (e.g. a necklace, with G all rotations and reflections of the n-cycle, i.e. the dihedral group D_n). Suppose further that we had Y colors, then the number of G-invariant colorings Y^X/G is counted by

$$\frac{1}{|G|} \sum_{g \in G} |Y|^{c(g)}$$

with c(g) counting the number of cycles of g when viewed as a permutation of X. We can generalize this to a weighted version: if the color i can occur exactly r_i times, then this is counted by the coefficient of $t_1^{r_1} \dots t_n^{r_n}$ in the polynomial

$$Z(t_1, \dots, t_n) = \frac{1}{|G|} \sum_{g \in G} \prod_{m \ge 1} (t_1^m + \dots + t_n^m)^{c_m(g)}$$

where $c_m(g)$ counts the number of length m cycles in g acting as a permutation on X. Note we get the original formula by setting all $t_i = 1$. Here Z is the cycle index. Note: you can cleverly deal with even/odd sizes by setting some t_i to -1.

Lucas Theorem If p is prime, then:

$$\frac{p^a}{k} \equiv 0 \pmod{p}$$

Thus for non-negative integers $m = m_k p^k + \ldots + m_1 p + m_0$ and $n = n_k p^k + \ldots + n_1 p + n_0$:

$$\frac{m}{n} = \prod_{i=0}^k \frac{m_i}{n_i} \mod p$$

Note: The fraction's mean integer division.

Catalan Numbers - Number of correct bracket sequence consisting of n opening and n closing brackets.

The number of ways to completely parenthesize n+1 factors.

The number of triangulations of a convex polygon with n+2 sides (i.e. the number of partitions of polygon into disjoint triangles by using the diagonals).

The number of ways to connect the 2n points on a circle to form n disjoint i.e. non-intersecting chords.

$$C_n = \frac{1}{n+1} \binom{2n}{n}$$

$$C_0 = 1, \ C_1 = 1, \ C_n = \sum_{k=0}^{n-1} C_k C_{n-1-k}$$

Narayana numbers The number of expressions containing n pairs of parentheses, which are correctly matched and which contain k distinct nestings.

$$N(n,k) = \frac{1}{n} \frac{n}{k} \frac{n}{k-1}$$