University of Groningen Balloon Addicts

1 SETUP,

```
1
```

```
1 Setup
                                                   3.2.2 0-1 Knapsack . . . . 11
  1.1 header.h . . . . . . . . . . . .
                                                   3.2.3
                                                          Coin change . . . . . 11
  1.2 Bash for c++ compile with
                                               header.h . . . . . . . . . . . .
                                                          Tree diameter . . . . 12
  1.3 Bash for run tests c++ . .
                                                   3.3.2 Tree Node Count . . 12
  1.4 Bash for run tests python .
                                               3.4 Num. Th. / Comb. . . . . . 12
       1.4.1 Aux. helper C++ .
                                                   3.4.1 Basic stuff . . . . . . 12
       1.4.2 Aux. helper python
                                                   3.4.2 Mod. exponentiation 13
2 Python
                                                          GCD . . . . . . . . . . . 13
  2.1 Graphs . . . . . . . . . . . . .
                                                          Sieve of Eratosthenes 13
       2.1.1 BFS . . . . . . . . .
                                                   3.4.5 Fibonacci % prime . 13
       2.1.2 Diikstra . . . . . .
                                                   3.4.6 nCk % prime . . . . 13
  2.2 Num. Th. / Comb. . . . . .
                                                   3.4.7 Chin. rem. th. . . . 13
       2.2.1 nCk % prime . . . .
                                               2.2.2 Sieve of E. . . . . . .
                                                   3.5.1 Z alg. . . . . . . . . . 13
  2.3 Strings . . . . . . . . . . . . .
                                                   3.5.2 KMP . . . . . . . . 14
       2.3.1 LCS . . . . . . . . .
                                                   3.5.3 Aho-Corasick . . . . 14
       2.3.2 KMP . . . . . . . .
                                                   3.5.4 Long. palin. subs . . 15
       2.3.3 Edit distance . . . .
                                                   Geometry . . . . . . . . . . . . 15
  2.4 Other Algorithms . . . . . .
                                                   3.6.1 essentials.cpp . . . . 15
       2.4.1 Rotate matrix . . . .
                                                   3.6.2 Two segs. itersec. . . 15
  2.5 Other Data Structures . . .
                                                          Convex Hull . . . . . 15
       2.5.1 Segment Tree . . . .
                                               3.7 Other Algorithms . . . . . . 16
       2.5.2 Trie . . . . . . . . . .
                                                   3.7.1 2-sat . . . . . . . . . . . 16
3 C++
                                                   3.7.2 Matrix Solve . . . . 16
  3.1 Graphs . . . . . . . . . . . . .
                                                   3.7.3 Matrix Exp. . . . . . 17
       3.1.1 BFS . . . . . . . . .
                                                   3.7.4 Finite field . . . . . 17
       3.1.2 DFS . . . . . . . . .
                                                   3.7.5 Complex field . . . . 17
       3.1.3 Dijkstra . . . . . . .
                                                   3.7.6 FFT . . . . . . . . . . . 17
       3.1.4 Floyd-Warshall . . .
                                                          Polyn. inv. div. . . . 18
       3.1.5 Kruskal . . . . . . .
                                                   3.7.8
                                                         Linear recurs. . . . . 18
       3.1.6 Hungarian algorithm
                                                   3.7.9
                                                          Convolution . . . . . 19
              Suc. shortest path.
       3.1.7
                                               3.8 Other Data Structures . . . 19
       3.1.8 Bipartite check . . .
                                                          Disjoint set . . . . .
       3.1.9 Find cycle directed.
                                                         Fenwick tree . . . .
       3.1.10 Find cycle directed.
                                                          Fenwick2d tree . . .
       3.1.11 Tarjan's SCC . . . .
                                                          Trie . . . . . . . . . .
       3.1.12 SCC edges . . . . .
                                                   3.8.5 Treap . . . . . . . .
       3.1.13 Find Bridges . . . .
                                           4 Other Mathematics
       3.1.14 Artic. points . . . .
                                                  Helpful functions . . . . . .
       3.1.15 Topological sort . . 10
                                                   4.1.1 Euler's Totient Fuc-
       3.1.16 Bellmann-Ford . . . 10
                                                          ntion . . . . . . . . .
       3.1.17 Ford-Fulkerson . . . 10
                                                   4.1.2 Pascal's trinagle . .
  3.2 Dynamic Programming . . 11
                                               4.2 Theorems and definitions .
       3.2.1 Longest Incr. Subseq. 11
                                              4.3 Geometry Formulas . . . .
```

1 Setup

1.1 header.h

```
1 #pragma once // Delete this when copying this file
2 #include <bits/stdc++.h>
3 using namespace std;
5 #define ll long long
6 #define ull unsigned ll
7 #define ld long double
8 #define pl pair<11, 11>
9 #define pi pair <int, int>
                              // use pl where possible/necessary
10 #define vl vector<ll>
11 #define vi vector<int> // change to vl where possible/necessary
12 #define vb vector <bool>
13 #define vvi vector <vi>
14 #define vvl vector <vl>
15 #define vpl vector <pl>
16 #define vpi vector <pi>
17 #define vld vector <ld>
18 #define vvpi vector < vpi>
19 #define in_fast(el, cont) (cont.find(el) != cont.end())
20 #define in(el, cont) (find(cont.begin(), cont.end(), el) != cont.end())
22 constexpr int INF = 200000011;
23 constexpr ll LLINF = 900000000000000010LL;
25 template <typename T, template <typename ELEM, typename ALLOC = std::
      allocator < ELEM > > class Container >
26 std::ostream& operator<<(std::ostream& o, const Container<T>& container) {
    typename Container <T>::const_iterator beg = container.begin();
    if (beg != container.end()) {
      o << *beg++;
      while (beg != container.end()) {
        o << " " << *beg++;
31
32
    }
33
    return o;
35 }
37 // int main() {
      ios::sync_with_stdio(false); // do not use cout + printf
      cin.tie(NULL);
      cout << fixed << setprecision(12);</pre>
41 // return 0;
42 // }
```

1.2 Bash for c++ compile with header.h

```
1 #!/bin/bash
2 if [ $# -ne 1 ]; then echo "Usage: $0 <input_file>"; exit 1; fi
3 f="$1";d=code/;o=a.out
4 [ -f $d/$f ] || { echo "Input file not found: $f"; exit 1; }
5 g++ -I$d $d/$f -o $o && echo "Compilation successful. Executable '$o' created." || echo "Compilation failed."
```

1.3 Bash for run tests c++

```
g++ $1/$1.cpp -o $1/$1.out
2 for file in $1/*.in; do diff <($1/$1.out < "$file") "${file%.in}.ans"; done
```

1.4 Bash for run tests python

```
_{1} for file in 1/*.in; do diff <(python3 1/$1.py < "$file") "${file%.in}.ans "; done
```

1.4.1 Aux. helper C++

```
1 #include "header.h"
3 int main() {
      // Read in a line including white space
      string line;
      getline(cin, line);
      // When doing the above read numbers as follows:
      getline(cin, line);
      stringstream ss(line);
      ss >> n:
12
      // Count the number of 1s in binary represnatation of a number
13
      ull number:
14
      __builtin_popcountll(number);
15
16 }
```

1.4.2 Aux. helper python

```
1 # Read until EOF
2 while True:
3    try:
4    pattern = input()
5    except EOFError:
6    break
```

2 Python

2.1 Graphs

2.1.1 BFS

```
1 from collections import deque
2 def bfs(g, roots, n):
      q = deque(roots)
      explored = set(roots)
      distances = [float("inf")]*n
      distances[0][0] = 0
      while len(a) != 0:
          node = q.popleft()
          if node in explored: continue
10
          explored.add(node)
11
12
          for neigh in g[node]:
              if neigh not in explored:
13
                   q.append(neigh)
14
                   distances[neigh] = distances[node] + 1
      return distances
```

2.1.2 Dijkstra

```
1 from heapq import *
2 def dijkstra(n, root, g): # g = {node: (cost, neigh)}
    dist = [float("inf")]*n
    dist[root] = 0
    prev = [-1]*n
    pq = [(0, root)]
    heapify(pq)
    visited = set([])
    while len(pq) != 0:
      _, node = heappop(pq)
      if node in visited: continue
14
      visited.add(node)
15
16
      # In case of disconnected graphs
17
      if node not in g:
18
        continue
20
      for cost, neigh in g[node]:
21
        alt = dist[node] + cost
        if alt < dist[neigh]:</pre>
23
          dist[neigh] = alt
          prev[neigh] = node
          heappush(pq, (alt, neigh))
26
    return dist
```

2.2 Num. Th. / Comb.

2.2.1 nCk % prime

```
1 # Note: p must be prime and k < p
2 def fermat_binom(n, k, p):
      if k > n:
          return 0
      # calculate numerator
      n_{11}m = 1
      for i in range(n-k+1, n+1):
          num *= i % p
      num %= p
      # calculate denominator
      denom = 1
11
      for i in range(1,k+1):
12
          denom *= i % p
13
      denom %= p
14
      # numerator * denominator^(p-2) (mod p)
15
      return (num * pow(denom, p-2, p)) % p
```

2.2.2 Sieve of E. O(n) so actually faster than C++ version, but more memory

```
_{1} MAX SIZE = 10**8+1
2 isprime = [True] * MAX_SIZE
_3 prime = []
4 SPF = [None] * (MAX SIZE)
6 def manipulated_seive(N): # Up to N (not included)
    isprime[0] = isprime[1] = False
    for i in range(2, N):
      if isprime[i] == True:
        prime.append(i)
        SPF[i] = i
11
      i = 0
12
      while (j < len(prime) and
        i * prime[j] < N and</pre>
          prime[j] <= SPF[i]):</pre>
        isprime[i * prime[j]] = False
        SPF[i * prime[j]] = prime[j]
        i += 1
```

2.3 Strings

2.3.1 LCS

```
def longestCommonSubsequence(text1, text2): # 0(m*n) time, 0(m) space
    n = len(text1)
    m = len(text2)

# Initializing two lists of size m
    prev = [0] * (m + 1)
    cur = [0] * (m + 1)
```

```
for idx1 in range(1, n + 1):
9
10
           for idx2 in range(1, m + 1):
               # If characters are matching
1.1
               if text1[idx1 - 1] == text2[idx2 - 1]:
12
                   cur[idx2] = 1 + prev[idx2 - 1]
13
               else:
14
                   # If characters are not matching
15
                   cur[idx2] = max(cur[idx2 - 1], prev[idx2])
16
17
           prev = cur.copy()
18
19
      return cur[m]
```

2.3.2 KMP

```
1 class KMP:
      def partial(self, pattern):
           """ Calculate partial match table: String -> [Int]"""
          for i in range(1, len(pattern)):
              j = ret[i - 1]
              while j > 0 and pattern[j] != pattern[i]: j = ret[j - 1]
              ret.append(j + 1 if pattern[j] == pattern[i] else j)
          return ret
10
      def search(self, T, P):
11
          """KMP search main algorithm: String -> String -> [Int]
12
          Return all the matching position of pattern string P in T"""
13
          partial, ret, j = self.partial(P), [], 0
14
          for i in range(len(T)):
15
              while j > 0 and T[i] != P[j]: j = partial[j - 1]
              if T[i] == P[j]: j += 1
17
              if i == len(P):
18
                  ret.append(i - (j - 1))
19
                  j = partial[j - 1]
20
          return ret
```

2.3.3 Edit distance

```
def editDistance(str1, str2):
    # Get the lengths of the input strings
    m = len(str1)
    n = len(str2)

# Initialize a list to store the current row
    curr = [0] * (n + 1)

# Initialize the first row with values from 0 to n
    for j in range(n + 1):
        curr[j] = j

# Initialize a variable to store the previous value
```

```
previous = 0
    # Loop through the rows of the dynamic programming matrix
    for i in range (1. m + 1):
      # Store the current value at the beginning of the row
      previous = curr[0]
      curr[0] = i
21
      # Loop through the columns of the dynamic programming matrix
22
      for j in range (1, n + 1):
        # Store the current value in a temporary variable
24
        temp = curr[i]
        # Check if the characters at the current positions in str1 and str2
            are the same
        if str1[i - 1] == str2[j - 1]:
          curr[j] = previous
          # Update the current cell with the minimum of the three adjacent
              cells
          curr[j] = 1 + min(previous, curr[j - 1], curr[j])
32
33
        # Update the previous variable with the temporary value
        previous = temp
35
    # The value in the last cell represents the minimum number of operations
37
    return curr[n]
```

2.4 Other Algorithms

2.4.1 Rotate matrix

```
1 def rotate_matrix(m):
2    return [[m[j][i] for j in range(len(m))] for i in range(len(m[0])
-1,-1,-1)]
```

2.5 Other Data Structures

2.5.1 Segment Tree

```
13 def updateTreeNode(p, value, n): # function to update a tree node
      # set value at position p
      tree[p + n] = value
      p = p + n
16
17
      i = p # move upward and update parents
18
10
      while i > 1:
          tree[i >> 1] = tree[i] + tree[i ^ 1]
20
          i >>= 1
23 def query(1, r, n): # function to get sum on interval [1, r)
      # loop to find the sum in the range
      r += n
      while 1 < r:
          if 1 & 1:
              res += tree[1]
              1 += 1
          if r & 1:
              r -= 1
33
              res += tree[r]
          1 >>= 1
          r >>= 1
      return res
```

2.5.2 Trie

```
1 class TrieNode:
      def init (self):
          self.children = [None]*26
           self.isEndOfWord = False
6 class Trie:
      def __init__(self):
          self.root = self.getNode()
10
      def getNode(self):
          return TrieNode()
11
12
      def _charToIndex(self,ch):
13
           return ord(ch)-ord('a')
14
15
16
      def insert(self.kev):
17
          pCrawl = self.root
18
           length = len(key)
19
          for level in range(length):
20
               index = self._charToIndex(key[level])
21
               if not pCrawl.children[index]:
22
                   pCrawl.children[index] = self.getNode()
23
               pCrawl = pCrawl.children[index]
24
           pCrawl.isEndOfWord = True
26
```

```
def search(self, key):

pCrawl = self.root

length = len(key)

for level in range(length):

index = self._charToIndex(key[level])

if not pCrawl.children[index]:

return False

pCrawl = pCrawl.children[index]

return pCrawl.isEndOfWord
```

3 C++

3.1 Graphs

3.1.1 BFS

```
1 #include "header.h"
2 #define graph unordered_map<11, unordered_set<11>>
3 vi bfs(int n, graph& g, vi& roots) {
       vi parents(n+1, -1); // nodes are 1..n
       unordered_set <int> visited;
      queue < int > q;
      for (auto x: roots) {
          q.emplace(x);
          visited.insert(x);
       while (not q.empty()) {
11
           int node = q.front();
12
13
          q.pop();
14
          for (auto neigh: g[node]) {
15
               if (not in(neigh, visited)) {
16
                   parents[neigh] = node;
                   q.emplace(neigh);
                   visited.insert(neigh);
               }
20
          }
21
22
      return parents;
23
24 }
     reconstruct_path(vi parents, int start, int goal) {
25 Vi
26
       int curr = goal;
27
       while (curr != start) {
28
           path.push_back(curr);
29
           if (parents[curr] == -1) return vi(); // No path, empty vi
30
           curr = parents[curr];
32
      path.push back(start):
33
      reverse(path.begin(), path.end());
34
      return path;
```

3.1.2 DFS Cycle detection / removal

```
1 #include "header.h"
2 void removeCyc(ll node, unordered_map<11, vector<pair<11, 11>>>& neighs,
      vector < bool > & visited.
3 vector < bool > & recStack, vector < 11 > & ans) {
      if (!visited[node]) {
           visited[node] = true:
           recStack[node] = true;
           auto it = neighs.find(node);
           if (it != neighs.end()) {
               for (auto util: it->second) {
                   11 nnode = util.first:
10
                   if (recStack[nnode]) {
11
                        ans.push_back(util.second);
12
                   } else if (!visited[nnode]) {
13
                       removeCyc(nnode, neighs, visited, recStack, ans);
14
               }
16
           }
17
18
19
      recStack[node] = false;
20 }
```

3.1.3 Dijkstra

```
1 #include "header.h"
2 vector<int> dijkstra(int n, int root, map<int, vector<pair<int, int>>>& g) {
    unordered_set <int> visited;
    vector < int > dist(n. INF):
      priority_queue < pair < int , int >> pq;
      dist[root] = 0;
      pq.push({0, root});
      while (!pq.empty()) {
           int node = pq.top().second;
           int d = -pq.top().first;
10
           pq.pop();
11
12
13
           if (in(node, visited)) continue;
           visited.insert(node);
14
15
           for (auto e : g[node]) {
16
               int neigh = e.first;
17
               int cost = e.second;
18
               if (dist[neigh] > dist[node] + cost) {
                   dist[neigh] = dist[node] + cost;
                   pq.push({-dist[neigh], neigh});
21
           }
23
      }
24
      return dist:
25
26 }
```

3.1.4 Floyd-Warshall

3.1.5 Kruskal Minimum spanning tree of undirected weighted graph

```
1 #include "header.h"
2 #include "disjoint_set.h"
3 // O(E log E)
4 pair<set<pair<11, 11>>, 11> kruskal(vector<tuple<11, 11, 11>>& edges, 11 n)
      set <pair <11. 11>> ans:
      11 cost = 0;
      sort(edges.begin(), edges.end());
      DisjointSet < 11 > fs(n);
10
      ll dist, i, j;
11
12
      for (auto edge: edges) {
           dist = get<0>(edge);
13
14
          i = get < 1 > (edge);
          j = get < 2 > (edge);
           if (fs.find_set(i) != fs.find_set(j)) {
17
               fs.union_sets(i, j);
18
               ans.insert({i, j});
19
               cost += dist;
          }
21
22
      return pair < set < pair < 11, 11>>, 11> {ans, cost};
23
```

3.1.6 Hungarian algorithm

```
8 * max(|C|)
   * @param C a matrix of dimensions JxW such that C[j][w] = cost to assign j-
     job to w-th worker (possibly negative)
   * Creturn a vector of length J, with the j-th entry equaling the minimum
  * to assign the first (j+1) jobs to distinct workers
15 template <class T> vector<T> hungarian(const vector<vector<T>> &C) {
      const int J = (int)size(C). W = (int)size(C[0]):
      assert(J <= W):
17
18
      // job[w] = job assigned to w-th worker, or -1 if no job assigned
      // note: a W-th worker was added for convenience
      vector < int > job(W + 1, -1);
20
      vector<T> ys(J), yt(W + 1); // potentials
^{21}
      // -yt[W] will equal the sum of all deltas
22
      vector <T> answers;
23
      const T inf = numeric_limits <T>::max();
24
      for (int j_cur = 0; j_cur < J; ++j_cur) { // assign j_cur-th job</pre>
25
          int w_cur = W;
26
          iob[w_cur] = j_cur;
27
          // min reduced cost over edges from Z to worker w
28
          vector <T> min_to(W + 1, inf);
29
          vector<int> prv(W + 1, -1); // previous worker on alternating path
30
          vector < bool > in_Z(W + 1);  // whether worker is in Z
31
          while (job[w_cur] != -1) { // runs at most j_cur + 1 times
32
              in Z[w cur] = true:
33
               const int j = job[w_cur];
34
              T delta = inf;
35
              int w_next;
36
              for (int w = 0; w < W; ++w) {
37
                   if (!in Z[w]) {
                       if (ckmin(min_to[w], C[j][w] - ys[j] - yt[w]))
                           prv[w] = w cur:
40
                       if (ckmin(delta, min_to[w])) w_next = w;
41
                   }
42
43
44
               // delta will always be non-negative,
               // except possibly during the first time this loop runs
45
              // if any entries of C[j_cur] are negative
              for (int w = 0; w \le W; ++w) {
47
                   if (in_Z[w]) ys[job[w]] += delta, yt[w] -= delta;
48
                   else min to[w] -= delta:
49
50
              w_cur = w_next;
51
          }
52
          // update assignments along alternating path
53
          for (int w; w_cur != W; w_cur = w) job[w_cur] = job[w = prv[w_cur]];
          answers.push_back(-yt[W]);
      return answers:
```

3.1.7 Suc. shortest path Calculates max flow, min cost

```
1 #include "header.h"
2 // map<node, map<node, pair<cost, capacity>>>
3 #define graph unordered_map<int, unordered_map<int, pair<ld, int>>>
5 const ld inftv = 1e601: // Change if necessary
6 ld fill(int n, vld& potential) { // Finds max flow, min cost
    priority_queue < pair < ld, int >> pq;
    vector <bool> visited(n+2, false);
    vi parent(n+2, 0);
    vld dist(n+2, infty);
    dist[0] = 0.1:
    pg.emplace(make_pair(0.1, 0));
    while (not pq.empty()) {
      int node = pq.top().second;
14
      pq.pop();
15
      if (visited[node]) continue;
16
      visited[node] = true;
17
      for (auto& x : g[node]) {
18
        int neigh = x.first:
19
        int capacity = x.second.second;
        ld cost = x.second.first;
21
        if (capacity and not visited[neigh]) {
22
          ld d = dist[node] + cost + potential[node] - potential[neigh];
23
          if (d + 1e-10l < dist[neigh]) {</pre>
             dist[neigh] = d;
25
            pq.emplace(make_pair(-d, neigh));
26
             parent[neigh] = node;
27
    }}}
28
29
    for (int i = 0; i < n+2; i++) {</pre>
      potential[i] = min(infty, potential[i] + dist[i]);
31
32
    if (not parent[n+1]) return infty;
    ld ans = 0.1:
    for (int x = n+1: x: x = parent[x]) {
      ans += g[parent[x]][x].first;
      g[parent[x]][x].second--;
      g[x][parent[x]].second++;
39
    return ans;
40
41 }
```

3.1.8 Bipartite check

```
if (side[st] == -1) {
10
               q.push(st);
11
12
               side[st] = 0:
               while (!a.emptv()) {
13
                   int v = q.front();
14
                   q.pop();
15
                   for (int u : adj[v]) {
16
                        if (side[u] == -1) {
17
                            side[u] = side[v] ^ 1:
18
                            q.push(u);
                       } else {
20
                            is_bipartite &= side[u] != side[v];
21
22
                       }
23 }}}}
```

3.1.9 Find cycle directed

```
1 #include "header.h"
2 int n;
3 \text{ const int } mxN = 2e5+5;
4 vvi adi(mxN):
5 vector < char > color;
6 vi parent:
7 int cycle_start, cycle_end;
8 bool dfs(int v) {
       color[v] = 1:
       for (int u : adj[v]) {
           if (color[u] == 0) {
11
               parent[u] = v;
12
13
               if (dfs(u)) return true;
           } else if (color[u] == 1) {
14
               cycle_end = v;
               cycle_start = u;
16
17
               return true:
           }
18
      }
       color[v] = 2;
       return false;
21
22 }
23 void find_cycle() {
       color.assign(n, 0);
       parent.assign(n, -1);
       cvcle_start = -1;
26
       for (int v = 0; v < n; v++) {
27
           if (color[v] == 0 && dfs(v))break:
28
29
      if (cycle_start == -1) {
30
           cout << "Acyclic" << endl;</pre>
31
      } else {
32
           vector < int > cycle;
33
           cycle.push_back(cycle_start);
34
           for (int v = cycle_end; v != cycle_start; v = parent[v])
35
36
               cycle.push_back(v);
           cycle.push_back(cycle_start);
37
```

```
38          reverse(cycle.begin(), cycle.end());
39
40          cout << "Cycle_Found:__";
41          for (int v : cycle) cout << v << "__";
42          cout << endl;
43          }
44 }</pre>
```

3.1.10 Find cycle directed

```
1 #include "header.h"
2 int n;
3 const int mxN = 2e5 + 5:
4 vvi adj(mxN);
5 vector < bool > visited:
6 vi parent:
7 int cycle_start, cycle_end;
s bool dfs(int v, int par) { // passing vertex and its parent vertex
      visited[v] = true;
      for (int u : adj[v]) {
           if(u == par) continue; // skipping edge to parent vertex
11
           if (visited[u]) {
12
               cycle_end = v;
               cycle_start = u;
               return true;
15
16
           parent[u] = v;
17
           if (dfs(u, parent[u]))
18
               return true;
19
20
      return false;
21
22 }
23 void find_cycle() {
       visited.assign(n. false):
       parent.assign(n, -1);
25
       cvcle_start = -1;
26
       for (int v = 0; v < n; v++) {
27
           if (!visited[v] && dfs(v, parent[v])) break;
28
29
      if (cycle_start == -1) {
30
           cout << "Acvclic" << endl;</pre>
31
      } else {
32
           vector<int> cycle;
33
           cycle.push_back(cycle_start);
34
           for (int v = cycle_end; v != cycle_start; v = parent[v])
               cycle.push_back(v);
           cvcle.push_back(cycle_start);
37
           cout << "Cvcle..Found:..":
38
           for (int v : cycle) cout << v << "";</pre>
           cout << endl:</pre>
41
42 }
```

3.1.11 Tarjan's SCC

```
1 #include "header.h"
3 struct Tarian {
    vvi &edges;
    int V, counter = 0, C = 0;
    vi n. 1:
    vector < bool > vs;
    stack<int> st:
    Tarjan(vvi &e): edges(e), V(e.size()), n(V, -1), l(V, -1), vs(V, false)
        {}
    void visit(int u. vi &com) {
      l[u] = n[u] = counter++;
      st.push(u);
      vs[u] = true:
      for (auto &&v : edges[u]) {
14
        if (n[v] == -1) visit(v. com);
        if (vs[v]) 1[u] = min(1[u], 1[v]);
17
18
      if (1[u] == n[u]) {
        while (true) {
19
          int v = st.top();
20
          st.pop();
21
          vs[v] = false:
22
          com[v] = C: // <== ACT HERE
23
          if (u == v) break;
        C++;
26
27
28
    int find_sccs(vi &com) { // component indices will be stored in 'com'
      com.assign(V, -1);
30
      C = 0:
31
      for (int u = 0; u < V; ++u)
        if (n[u] == -1) visit(u, com);
      return C:
34
35
    // scc is a map of the original vertices of the graph to the vertices
    // of the SCC graph, scc_graph is its adjacency list.
    // SCC indices and edges are stored in 'scc' and 'scc_graph'.
    void scc collapse(vi &scc. vvi &scc graph) {
      find sccs(scc):
40
      scc_graph.assign(C, vi());
41
      set < pi > rec; // recorded edges
      for (int u = 0; u < V; ++u) {
43
        assert(scc[u] != -1):
44
        for (int v : edges[u]) {
45
          if (scc[v] == scc[u] ||
            rec.find({scc[u], scc[v]}) != rec.end()) continue;
47
          scc_graph[scc[u]].push_back(scc[v]);
18
          rec.insert({scc[u], scc[v]});
50
      }
51
    // Function to find sources and sinks in the SCC graph
```

```
// The number of edges needed to be added is max(sources.size(), sinks.())
void findSourcesAndSinks(const vvi &scc_graph, vi &sources, vi &sinks) {
    vi in_degree(C, 0), out_degree(C, 0);
    for (int u = 0; u < C; u++) {
        for (auto v : scc_graph[u]) {
            in_degree[v]++;
            out_degree[u]++;
        }
        for (int i = 0; i < C; ++i) {
        if (in_degree[i] == 0) sources.push_back(i);
        if (out_degree[i] == 0) sinks.push_back(i);
    }
}</pre>
```

3.1.12 SCC edges Prints out the missing edges to make the input digraph strongly connected

```
1 #include "header.h"
2 const int N=1e5+10;
3 int n.a[N].cnt[N].vis[N]:
4 vector < int > hd,tl;
5 int dfs(int x){
      vis[x]=1:
      if(!vis[a[x]])return vis[x]=dfs(a[x]);
      return vis[x]=x:
9 }
10 int main(){
       scanf("%d".&n):
      for(int i=1:i<=n:i++){</pre>
12
           scanf("%d",&a[i]);
           cnt[a[i]]++;
      int k=0:
16
      for(int i=1;i<=n;i++){</pre>
17
           if(!cnt[i]){
18
               k++:
               hd.push_back(i);
               tl.push_back(dfs(i));
21
           }
22
      }
23
      int tk=k:
24
       for(int i=1;i<=n;i++){</pre>
25
           if(!vis[i]){
26
               k++:
               hd.push_back(i);
28
                tl.push back(dfs(i)):
29
           }
30
31
      if(k==1&&!tk)k=0:
32
33
      printf("%d\n",k);
      for(int i=0;i<k;i++)printf("%d<sub>||</sub>%d\n",tl[i],hd[(i+1)%k]);
34
      return 0:
```

```
36 }
```

3.1.13 Find Bridges

```
1 #include "header.h"
2 int n; // number of nodes
3 vvi adj; // adjacency list of graph
4 vector < bool > visited:
5 vi tin, low;
6 int timer:
7 void dfs(int v, int p = -1) {
      visited[v] = true;
      tin[v] = low[v] = timer++:
      for (int to : adj[v]) {
           if (to == p) continue;
11
           if (visited[to]) {
12
               low[v] = min(low[v], tin[to]);
          } else {
14
               dfs(to, v);
15
               low[v] = min(low[v], low[to]);
16
               if (low[to] > tin[v])
                   IS_BRIDGE(v, to);
           }
19
      }
20
21 }
22 void find_bridges() {
      timer = 0;
      visited.assign(n, false);
24
      tin.assign(n, -1);
      low.assign(n, -1);
26
      for (int i = 0; i < n; ++i) {</pre>
           if (!visited[i]) dfs(i):
28
      }
29
30 }
```

3.1.14 Artic. points (i.e. cut off points)

```
1 #include "header.h"
2 int n; // number of nodes
3 vvi adj; // adjacency list of graph
4 vector < bool > visited;
5 vi tin, low;
6 int timer:
7 void dfs(int v, int p = -1) {
      visited[v] = true;
      tin[v] = low[v] = timer++:
      int children=0:
      for (int to : adj[v]) {
11
          if (to == p) continue;
12
          if (visited[to]) {
13
               low[v] = min(low[v], tin[to]):
14
          } else {
15
               dfs(to, v);
```

```
low[v] = min(low[v], low[to]);
               if (low[to] >= tin[v] && p!=-1) IS_CUTPOINT(v);
               ++children:
21
      if(p == -1 && children > 1)
22
          IS CUTPOINT(v):
23
24 }
25 void find cutpoints() {
      timer = 0;
      visited.assign(n, false);
27
      tin.assign(n, -1);
      low.assign(n, -1);
      for (int i = 0; i < n; ++i) {
           if (!visited[i]) dfs (i);
32
33 }
```

3.1.15 Topological sort

```
1 #include "header.h"
2 int n: // number of vertices
3 vvi adi: // adiacency list of graph
4 vector < bool > visited:
5 vi ans:
6 void dfs(int v) {
      visited[v] = true;
      for (int u : adj[v]) {
          if (!visited[u]) dfs(u);
      ans.push_back(v);
11
12 }
13 void topological_sort() {
      visited.assign(n. false):
      ans.clear():
15
      for (int i = 0; i < n; ++i) {
16
          if (!visited[i]) dfs(i);
17
18
      reverse(ans.begin(), ans.end());
19
20 }
```

3.1.16 Bellmann-Ford Same as Dijkstra but allows neg. edges

```
1 #include "header.h"
2 // Switch vi and vvpi to vl and vvpl if necessary
3 void bellmann_ford_extended(vvpi &e, int source, vi &dist, vb &cyc) {
4    dist.assign(e.size(), INF);
5    cyc.assign(e.size(), false); // true when u is in a <0 cycle
6    dist[source] = 0;
7    for (int iter = 0; iter < e.size() - 1; ++iter){
8        bool relax = false;
9        for (int u = 0; u < e.size(); ++u)
10        if (dist[u] == INF) continue;</pre>
```

```
else for (auto &e : e[u])
           if(dist[u]+e.second < dist[e.first])</pre>
13
             dist[e.first] = dist[u]+e.second. relax = true:
      if(!relax) break:
14
15
    bool ch = true;
16
    while (ch) {
                         // keep going untill no more changes
      ch = false:
                         // set dist to -INF when in cycle
      for (int u = 0: u < e.size(): ++u)</pre>
        if (dist[u] == INF) continue;
        else for (auto &e : e[u])
21
          if (dist[e.first] > dist[u] + e.second
             && !cvc[e.first]) {
23
             dist[e.first] = -INF;
             ch = true; //return true for cycle detection only
             cyc[e.first] = true;
26
27
    }
29 }
```

3.1.17 Ford-Fulkerson Basic Max. flow

```
1 #include "header.h"
2 #define V 6 // Num. of vertices in given graph
4 /* Returns true if there is a path from source 's' to sink
5 't' in residual graph. Also fills parent[] to store the
6 path */
7 bool bfs(int rGraph[V][V], int s, int t, int parent[]) {
    bool visited[V]:
    memset(visited, 0, sizeof(visited));
    queue < int > q;
    q.push(s);
    visited[s] = true:
    parent[s] = -1;
    // Standard BFS Loop
    while (!q.empty()) {
16
      int u = q.front();
17
      q.pop();
18
19
      for (int v = 0; v < V; v++) {
20
        if (visited[v] == false && rGraph[u][v] > 0) {
21
          if (v == t) {
22
             parent[v] = u:
            return true;
24
25
          q.push(v);
26
          parent[v] = u;
          visited[v] = true:
20
      }
30
31
    return false;
```

```
Returns the maximum flow from s to t in the given graph
36 int fordFulkerson(int graph[V][V], int s, int t) {
    int u. v:
    int rGraph[V]
        [V]:
    for (u = 0; u < V; u++)
     for (v = 0; v < V; v++)
        rGraph[u][v] = graph[u][v];
    int parent[V]; // This array is filled by BFS and to
          // store path
45
    int max_flow = 0; // There is no flow initially
    while (bfs(rGraph, s, t, parent)) {
      int path_flow = INT_MAX;
      for (v = t; v != s; v = parent[v]) {
        u = parent[v];
        path_flow = min(path_flow, rGraph[u][v]);
53
      for (v = t: v != s: v = parent[v]) {
54
        u = parent[v];
        rGraph[u][v] -= path_flow;
        rGraph[v][u] += path_flow;
57
58
      max_flow += path_flow;
    return max_flow;
61
62 }
```

3.2 Dynamic Programming

3.2.1 Longest Incr. Subseq.

```
1 #include "header.h"
2 template < class T>
3 vector <T> index_path_lis(vector <T>& nums) {
    int n = nums.size():
    vector <T> sub:
      vector <int> subIndex:
    vector <T> path(n, -1);
    for (int i = 0; i < n; ++i) {</pre>
        if (sub.empty() || sub[sub.size() - 1] < nums[i]) {</pre>
      path[i] = sub.empty() ? -1 : subIndex[sub.size() - 1];
      sub.push_back(nums[i]);
11
      subIndex.push back(i):
      int idx = lower_bound(sub.begin(), sub.end(), nums[i]) - sub.begin();
      path[i] = idx == 0 ? -1 : subIndex[idx - 1]:
      sub[idx] = nums[i];
      subIndex[idx] = i;
```

```
vector <T> ans:
    int t = subIndex[subIndex.size() - 1];
    while (t != -1) {
        ans.push_back(t);
        t = path[t];
24
    }
25
    reverse(ans.begin(), ans.end());
    return ans;
29 // Length only
30 template < class T>
31 int length_lis(vector<T> &a) {
    set <T> st;
    typename set<T>::iterator it;
    for (int i = 0; i < a.size(); ++i) {</pre>
      it = st.lower_bound(a[i]);
      if (it != st.end()) st.erase(it):
      st.insert(a[i]);
38
    return st.size();
40 }
```

3.2.2 0-1 Knapsack

```
1 #include "header.h"
2 // given a number of coins, calculate all possible distinct sums
3 int main() {
    int n;
    vi coins(n); // all possible coins to use
    int sum = 0:  // sum of the coins
    vi dp(sum + 1, 0);
                          // dp[x] = 1 if sum x can be made
    dp[0] = 1:
                               // sum 0 can be made
    for (int c = 0; c < n; ++c)
                                     // first iteration: sums with first
     for (int x = sum: x \ge 0: --x)
                                      // coin. next first 2 coins etc
        if (dp[x]) dp[x + coins[c]] = 1; // if sum x valid, x+c valid
11
```

3.2.3 Coin change Number of coins required to achieve a given value

```
#include "header.h"

// Returns total distinct ways to make sum using n coins of

// different denominations

int count(vi& coins, int n, int sum) {

// 2d dp array where n is the number of coin

// denominations and sum is the target sum

vector<vector<int>> dp(n + 1, vector<int>(sum + 1, 0));

dp[0][0] = 1;

for (int i = 1; i <= n; i++) {

for (int j = 0; j <= sum; j++) {

// without using the current coin,

dp[i][j] += dp[i - 1][j];</pre>
```

3.3 Trees

3.3.1 Tree diameter

```
1 #include "header.h"
2 const int mxN = 2e5 + 5;
3 int n, d[mxN]; // distance array
4 vi adj[mxN]; // tree adjacency list
5 void dfs(int s. int e) {
    d[s] = 1 + d[e]:
                         // recursively calculate the distance from the
        starting node to each node
    for (auto u : adj[s]) { // for each adjacent node
      if (u != e) dfs(u, s); // don't move backwards in the tree
10 }
11 int main() {
    // read input, create adi list
    dfs(0, -1):
                                  // first dfs call to find farthest node from
         arbitrary node
    dfs(distance(d, max_element(d, d + n)), -1); // second dfs call to find
        farthest node from that one
    cout << *max_element(d, d + n) - 1 << '\n'; // distance from second node</pre>
        to farthest is the diameter
16 }
```

3.3.2 Tree Node Count

```
#include "header.h"
// calculate amount of nodes in each node's subtree
const int mxN = 2e5 + 5;
int n, cnt[mxN];
vi adj[mxN];
void dfs(int s = 0, int e = -1) {
cnt[s] = 1; // count leaves as one
for (int u : adj[s]) {
dfs(u, s);
cnt[s] += cnt[u]; // add up nodes of the subtrees
}
```

3.4 Num. Th. / Comb.

3.4.1 Basic stuff

```
1 #include "header.h"
2 ll gcd(ll a, ll b) { while (b) { a %= b; swap(a, b); } return a; }
3 11 1cm(11 a, 11 b) { return (a / gcd(a, b)) * b; }
4 ll mod(ll a, ll b) { return ((a % b) + b) % b; }
5 // \text{ Finds } x, y \text{ s.t. ax + by = d = gcd(a, b)}.
6 void extended_euclid(ll a, ll b, ll &x, ll &y, ll &d) {
    11 yy = x = 1;
    while (b) {
      11 \ a = a / b:
      11 t = b; b = a \% b; a = t;
      t = xx; xx = x - q * xx; x = t;
      t = yy; yy = y - q * yy; y = t;
14 }
15
16 }
17 //  solves ab = 1 (mod n), -1 on failure
18 ll mod_inverse(ll a, ll n) {
    11 x. v. d:
    extended_euclid(a, n, x, y, d);
    return (d > 1 ? -1 : mod(x, n));
22 }
23 // All modular inverses of [1..n] mod P in O(n) time.
24 vi inverses(ll n. ll P) {
   vi I(n+1, 1LL):
    for (11 i = 2; i <= n; ++i)</pre>
      I[i] = mod(-(P/i) * I[P\%i], P):
    return I;
29 }
30 // (a*b)%m
31 ll mulmod(ll a, ll b, ll m){
   11 x = 0. v=a\%m:
    while(b>0){
      if(b\&1) x = (x+y)\%m;
      y = (2*y)\%m, b /= 2;
36
    return x % m;
39 // Finds b^e % m in O(lg n) time, ensure that b < m to avoid overflow!
40 ll powmod(ll b, ll e, ll m) {
    11 p = e < 2 ? 1 : powmod((b*b)\%m, e/2, m);
    return e&1 ? p*b%m : p;
43 }
44 // Solve ax + by = c, returns false on failure.
45 bool linear_diophantine(ll a, ll b, ll c, ll &x, ll &y) {
  11 d = gcd(a, b);
   if (c % d) {
      return false:
   } else {
      x = c / d * mod_inverse(a / d, b / d);
      y = (c - a * x) / b;
      return true;
53 }
54 }
```

9 }

3.4.2 Mod. exponentiation Or use pow() in python

```
#include "header.h"
2 ll mod_pow(ll base, ll exp, ll mod) {
3    if (mod == 1) return 0;
4       if (exp == 0) return 1;
5       if (exp == 1) return base;
6
7    ll res = 1;
8    base %= mod;
9    while (exp) {
10       if (exp % 2 == 1) res = (res * base) % mod;
11       exp >>= 1;
12       base = (base * base) % mod;
13    }
14
15    return res % mod;
16 }
```

3.4.3 GCD Or math.gcd in python, std::gcd in C++

```
#include "header.h"
2 ll gcd(ll a, ll b) {
3    if (a == 0) return b;
4    return gcd(b % a, a);
5 }
```

3.4.4 Sieve of Eratosthenes

```
#include "header.h"
vl primes;
void getprimes(ll n) { // Up to n (not included)

vector<bool> p(n, true);

p[0] = false;

p[1] = false;

for(ll i = 0; i < n; i++) {

if(p[i]) {

primes.push_back(i);

for(ll j = i*2; j < n; j+=i) p[j] = false;
}
}}</pre>
```

3.4.5 Fibonacci % prime

```
1 #include "header.h"
2 const 11 MOD = 1000000007;
3 unordered_map<11, 11> Fib;
4 l1 fib(l1 n) {
5     if (n < 2) return 1;
6     if (Fib.find(n) != Fib.end()) return Fib[n];
7     Fib[n] = (fib((n + 1) / 2) * fib(n / 2) + fib((n - 1) / 2) * fib((n - 2) / 2)) % MOD;
8     return Fib[n];</pre>
```

3.4.6 nCk % prime

```
1 #include "header.h"
2 ll binom(ll n, ll k) {
      ll ans = 1;
      for (11 i = 1: i <= min(k,n-k): ++i) ans = ans*(n+1-i)/i:
7 ll mod_nCk(ll n, ll k, ll p ){
      ll ans = 1;
      while(n){
          11 np = n\%p, kp = k\%p;
          if(kp > np) return 0;
11
          ans *= binom(np,kp);
12
          n /= p; k /= p;
13
14
15
      return ans;
16 }
```

3.4.7 Chin. rem. th.

```
1 #include "header.h"
2 #include "elementary.cpp"
_3 // Solves x = a1 mod m1, x = a2 mod m2, x is unique modulo lcm(m1, m2).
4 // Returns {0, -1} on failure, {x, lcm(m1, m2)} otherwise.
5 pair<11, 11> crt(11 a1, 11 m1, 11 a2, 11 m2) {
6 ll s, t, d;
    extended_euclid(m1, m2, s, t, d);
    if (a1 % d != a2 % d) return {0, -1};
    return {mod(s*a2 %m2 * m1 + t*a1 %m1 * m2, m1 * m2) / d, m1 / d * m2};
10 }
11
_{12} // Solves x = ai mod mi. x is unique modulo lcm mi.
13 // Returns {0, -1} on failure, {x, lcm mi} otherwise.
14 pair <11, 11> crt(vector <11> &a, vector <11> &m) {
pair<11, 11> res = {a[0], m[0]};
    for (ull i = 1; i < a.size(); ++i) {</pre>
      res = crt(res.first, res.second, mod(a[i], m[i]), m[i]);
      if (res.second == -1) break;
19
   }
    return res:
20
```

3.5 Strings

3.5.1 Z alg. KMP alternative

```
#include "../header.h"
2 void Z_algorithm(const string &s, vi &Z) {
```

```
Z.assign(s.length(), -1);
int L = 0, R = 0, n = s.length();
for (int i = 1; i < n; ++i) {
   if (i > R) {
      L = R = i;
      while (R < n && s[R - L] == s[R]) R++;
      Z[i] = R - L; R--;
   } else if (Z[i - L] >= R - i + 1) {
      L = i;
      while (R < n && s[R - L] == s[R]) R++;
      Z[i] = R - L; R--;
   } else Z[i] = Z[i - L];
}</pre>
```

3.5.2 KMP

```
1 #include "header.h"
void compute_prefix_function(string &w, vi &prefix) {
    prefix.assign(w.length(), 0);
    int k = prefix[0] = -1:
    for(int i = 1: i < w.length(): ++i) {</pre>
      while (k >= 0 \&\& w[k + 1] != w[i]) k = prefix[k];
      if(w[k + 1] == w[i]) k++;
      prefix[i] = k:
10
11 }
12 void knuth_morris_pratt(string &s, string &w) {
    int a = -1:
    vi prefix;
    compute prefix function(w. prefix):
    for(int i = 0; i < s.length(); ++i) {</pre>
16
      while (q >= 0 \&\& w[q + 1] != s[i]) q = prefix[q];
17
      if(w[q + 1] == s[i]) q++;
18
      if(q + 1 == w.length()) {
       // Match at position (i - w.length() + 1)
        q = prefix[q];
21
23
    }
24 }
```

3.5.3 Aho-Corasick Also can be used as Knuth-Morris-Pratt algorithm

```
#include "header.h"

map < char, int > cti;

map < char, int > cti;

map template < int ALPHABET_SIZE, int (*mp)(char) >

map template < int child[ALPHABET_SIZE], failure = 0, match_par = -1;

map template < int child[ALPHABET_SIZE], failure = 0, match_par = -1;

map template < int child[ALPHABET_SIZE], failure = 0, match_par = -1;

map template < int child[ALPHABET_SIZE], failure = 0, match_par = -1;

map template < int child[ALPHABET_SIZE], failure = 0, match_par = -1;

map template < int child[ALPHABET_SIZE], failure = 0, match_par = -1;

map template < int child[ALPHABET_SIZE], failure = 0, match_par = -1;

map template < int child[ALPHABET_SIZE], failure = 0, match_par = -1;

map template < int child[ALPHABET_SIZE], failure = 0, match_par = -1;

map template < int child[ALPHABET_SIZE], failure = 0, match_par = -1;

map template < int child[ALPHABET_SIZE], failure = 0, match_par = -1;

map template < int child[ALPHABET_SIZE], failure = 0, match_par = -1;

map template < int child[ALPHABET_SIZE], failure < int child[ALPHABET_SIZE], fai
```

```
Node() { for (int i = 0; i < ALPHABET_SIZE; ++i) child[i] = -1; }
    };
11
    vector < Node > a:
12
    vector<string> &words;
    AC_FSM(vector<string> &words) : words(words) {
      a.push_back(Node());
       construct_automaton();
16
17
    void construct automaton() {
18
      for (int w = 0, n = 0; w < words.size(); ++w, n = 0) {
19
        for (int i = 0: i < words[w].size(): ++i) {</pre>
20
           if (a[n].child[mp(words[w][i])] == -1) {
21
             a[n].child[mp(words[w][i])] = a.size();
22
             a.push_back(Node());
23
            = a[n].child[mp(words[w][i])];
25
26
        a[n].match.push_back(w);
27
28
      queue < int > q:
      for (int k = 0; k < ALPHABET_SIZE; ++k) {</pre>
30
        if (a[0].child[k] == -1) a[0].child[k] = 0:
31
        else if (a[0].child[k] > 0) {
32
           a[a[0].child[k]].failure = 0;
33
           q.push(a[0].child[k]);
34
35
      }
36
      while (!a.emptv()) {
37
        int r = q.front(); q.pop();
38
        for (int k = 0, arck; k < ALPHABET_SIZE; ++k) {</pre>
39
           if ((arck = a[r].child[k]) != -1) {
40
             q.push(arck);
41
             int v = a[r].failure:
42
             while (a[v].child[k] == -1) v = a[v].failure;
43
             a[arck].failure = a[v].child[k]:
44
             a[arck].match_par = a[v].child[k];
             while (a[arck].match_par != -1
46
                 && a[a[arck].match par].match.emptv())
47
               a[arck].match_par = a[a[arck].match_par].match_par;
48
49
50
        }
      }
51
52
    void aho corasick(string &sentence, vvi &matches){
53
      matches.assign(words.size(), vi());
54
      int state = 0, ss = 0:
55
      for (int i = 0; i < sentence.length(); ++i, ss = state) {</pre>
56
        while (a[ss].child[mp(sentence[i])] == -1)
57
           ss = a[ss].failure:
        state = a[state].child[mp(sentence[i])]
59
             = a[ss].child[mp(sentence[i])]:
60
        for (ss = state; ss != -1; ss = a[ss].match_par)
61
           for (int w : a[ss].match)
62
             matches[w].push_back(i + 1 - words[w].length());
63
      }
64
```

```
66 };
67 int char_to_int(char c) {
68  return cti[c];
69 }
```

3.5.4 Long. palin. subs Manacher - O(n)

```
1 #include "header.h"
void manacher(string &s. vi &pal) {
    int n = s.length(), i = 1, 1, r;
    pal.assign(2 * n + 1, 0);
    while (i < 2 * n + 1) {
      if ((i&1) && pal[i] == 0) pal[i] = 1;
      1 = i / 2 - pal[i] / 2; r = (i-1) / 2 + pal[i] / 2;
      while (1 - 1 >= 0 \&\& r + 1 < n \&\& s[1 - 1] == s[r + 1])
        --1, ++r, pal[i] += 2;
10
11
      for (1 = i - 1, r = i + 1; 1 >= 0 && r < 2 * n + 1; --1, ++r)
12
        if (1 <= i - pal[i]) break:</pre>
13
        if (1 / 2 - pal[1] / 2 > i / 2 - pal[i] / 2)
14
15
          pal[r] = pal[1];
        else { if (1 >= 0)
            pal[r] = min(pal[1], i + pal[i] - r);
          break:
      i = r;
22 } }
```

3.6 Geometry

3.6.1 essentials.cpp

```
1 #include "../header.h"
2 using C = ld; // could be long long or long double
3 constexpr C EPS = 1e-10; // change to 0 for C=11
4 struct P { // may also be used as a 2D vector
    P(C x = 0, C y = 0) : x(x), y(y) {}
    P operator + (const P &p) const { return {x + p.x, y + p.y}; }
   P operator - (const P &p) const { return {x - p.x, y - p.y}; }
   P operator* (C c) const { return {x * c, y * c}; }
   P operator/ (C c) const { return {x / c, y / c}; }
   C operator* (const P &p) const { return x*p.x + y*p.y; }
12 C operator (const P &p) const { return x*p.y - p.x*y; }
   P perp() const { return P{y, -x}; }
C lensq() const { return x*x + y*y; }
15 ld len() const { return sqrt((ld)lensq()); }
   static ld dist(const P &p1, const P &p2) {
     return (p1-p2).len(); }
    bool operator == (const P &r) const {
```

```
return ((*this)-r).lensq() <= EPS*EPS; }</pre>
20 };
21 C det(P p1, P p2) { return p1^p2; }
22 C det(P p1, P p2, P o) { return det(p1-o, p2-o); }
23 C det(const vector <P> &ps) {
    C sum = 0; P prev = ps.back();
    for(auto &p : ps) sum += det(p, prev), prev = p;
    return sum:
28 // Careful with division by two and C=11
29 C area(P p1, P p2, P p3) { return abs(det(p1, p2, p3))/C(2); }
30 C area(const vector<P> &poly) { return abs(det(poly))/C(2); }
31 int sign(C c) { return (c > C(0)) - (c < C(0)); }
32 int ccw(P p1, P p2, P o) { return sign(det(p1, p2, o)); }
_{34} // Only well defined for C = ld.
35 P unit(const P &p) { return p / p.len(); }
_{36} P rotate(P p, 1d a) { return P{p.x*cos(a)-p.y*sin(a), p.x*sin(a)+p.y*cos(a)}
      }; }
```

3.6.2 Two segs. itersec.

```
1 #include "header.h"
2 #include "essentials.cpp"
3 bool intersect(P a1, P a2, P b1, P b2) {
4    if (max(a1.x, a2.x) < min(b1.x, b2.x)) return false;
5    if (max(b1.x, b2.x) < min(a1.x, a2.x)) return false;
6    if (max(a1.y, a2.y) < min(b1.y, b2.y)) return false;
7    if (max(b1.y, b2.y) < min(a1.y, a2.y)) return false;
8    bool l1 = ccw(a2, b1, a1) * ccw(a2, b2, a1) <= 0;
9    bool l2 = ccw(b2, a1, b1) * ccw(b2, a2, b1) <= 0;
10    return l1 && l2;
11 }</pre>
```

3.6.3 Convex Hull

```
1 #include "header.h"
2 #include "essentials.cpp"
3 struct ConvexHull { // O(n lg n) monotone chain.
    vector<size_t> h, c; // Indices of the hull are in 'h', ccw.
    const vector <P> &p;
    ConvexHull(const vector<P> &_p) : n(_p.size()), c(n), p(_p) {
      std::iota(c.begin(), c.end(), 0):
      std::sort(c.begin(), c.end(), [this](size_t l, size_t r) -> bool {
          return p[1].x != p[r].x ? p[1].x < p[r].x : p[1].y < p[r].y; });
      c.erase(std::unique(c.begin(), c.end(), [this](size_t l, size_t r) {
10
          return p[1] == p[r]; }), c.end());
      for (size_t s = 1, r = 0; r < 2; ++r, s = h.size()) {</pre>
11
        for (size_t i : c) {
12
          while (h.size() > s \&\& ccw(p[h.end()[-2]], p[h.end()[-1]], p[i]) \le
13
            h.pop_back();
14
```

```
h.push_back(i);
        reverse(c.begin(), c.end());
17
18
      if (h.size() > 1) h.pop_back();
19
20
    size_t size() const { return h.size(); }
^{21}
    template <class T, void U(const P &, const P &, const P &, T &)>
22
    void rotating calipers(T &ans) {
      if (size() <= 2)</pre>
        U(p[h[0]], p[h.back()], p[h.back()], ans);
25
26
        for (size_t i = 0, j = 1, s = size(); i < 2 * s; ++i) {</pre>
27
           while (\det(p[h[(i + 1) \% s]) - p[h[i \% s]), p[h[(j + 1) \% s]] - p[h[
              i]]) >= 0)
             j = (j + 1) \% s;
          U(p[h[i % s]], p[h[(i + 1) % s]], p[h[j]], ans);
32
33 }:
     Example: furthest pair of points. Now set ans = OLL and call
     ConvexHull(pts).rotating calipers < 11. update > (ans):
36 void update(const P &p1, const P &p2, const P &o, 11 &ans) {
    ans = \max(ans, (11)\max((p1 - o).lensq(), (p2 - o).lensq()));
38 }
```

3.7 Other Algorithms

3.7.1 2-sat

```
1 #include "../header.h"
2 #include "../Graphs/tarjan.cpp"
3 struct TwoSAT {
    int n:
    vvi imp; // implication graph
    Tarjan tj;
    TwoSAT(int _n) : n(_n), imp(2 * _n, vi()), tj(imp) { }
    // Only copy the needed functions:
    void add_implies(int c1, bool v1, int c2, bool v2) {
11
      int u = 2 * c1 + (v1 ? 1 : 0),
12
        v = 2 * c2 + (v2 ? 1 : 0);
      imp[u].push_back(v); // u => v
      imp[v^1].push_back(u^1); // -v => -u
15
16
    void add equivalence(int c1. bool v1. int c2. bool v2) {
17
      add_implies(c1, v1, c2, v2);
      add_implies(c2, v2, c1, v1);
19
20
    void add_or(int c1, bool v1, int c2, bool v2) {
21
      add_implies(c1, !v1, c2, v2);
22
    void add_and(int c1, bool v1, int c2, bool v2) {
```

```
add_true(c1, v1); add_true(c2, v2);
26
27
    void add_xor(int c1, bool v1, int c2, bool v2) {
      add or(c1, v1, c2, v2):
      add_or(c1, !v1, c2, !v2);
29
    }
30
    void add_true(int c1, bool v1) {
      add_implies(c1, !v1, c1, v1);
32
33
    // on true: a contains an assignment.
    // on false: no assignment exists.
    bool solve(vb &a) {
      vi com:
      tj.find_sccs(com);
      for (int i = 0; i < n; ++i)</pre>
        if (com[2 * i] == com[2 * i + 1])
41
          return false;
42
43
      vvi bvcom(com.size()):
      for (int i = 0; i < 2 * n; ++i)
45
        bvcom[com[i]].push back(i):
46
47
      a.assign(n, false);
48
      vb vis(n. false):
49
50
      for(auto &&component : bycom){
        for (int u : component) {
51
          if (vis[u / 2]) continue:
          vis[u / 2] = true;
          a[u / 2] = (u \% 2 == 1):
54
55
      }
      return true:
59 }:
```

3.7.2 Matrix Solve

```
1 #include "header.h"
2 #define REP(i, n) for(auto i = decltype(n)(0); i < (n); i++)</pre>
3 using T = double;
4 constexpr T EPS = 1e-8;
5 template < int R, int C>
6 using M = array < array < T, C > , R >; // matrix
7 template < int R. int C>
s T ReducedRowEchelonForm(M<R,C> &m, int rows) { // return the determinant
    int r = 0; T det = 1;
                                      // MODIFIES the input
    for(int c = 0: c < rows && r < rows: <math>c++) {
      int p = r;
      for(int i=r+1; i<rows; i++) if(abs(m[i][c]) > abs(m[p][c])) p=i;
      if(abs(m[p][c]) < EPS){ det = 0; continue; }</pre>
      swap(m[p], m[r]); det = -det;
14
      T s = 1.0 / m[r][c], t; det *= m[r][c];
      REP(j,C) m[r][j] *= s;  // make leading term in row 1
16
```

```
REP(i,rows) if (i!=r){ t = m[i][c]; REP(j,C) m[i][j] -= t*m[r][j]; }
      ++r;
    }
    return det:
21 }
22 bool error, inconst; // error => multiple or inconsistent
23 template <int R, int C> // Mx = a; M:R*R, v:R*C => x:R*C
24 M<R,C> solve(const M<R,R> &m, const M<R,C> &a, int rows){
    M < R.R+C > a:
    REP(r,rows){
      REP(c,rows) q[r][c] = m[r][c];
27
      REP(c,C) q[r][R+c] = a[r][c];
    ReducedRowEchelonForm <R,R+C>(q,rows);
30
    M<R,C> sol; error = false, inconst = false;
    REP(c,C) for(auto j = rows-1; j >= 0; --j){
      T t=0: bool allzero=true:
33
      for (auto k = j+1; k < rows; ++k)
        t += q[j][k]*sol[k][c], allzero &= abs(q[j][k]) < EPS;
35
      if(abs(q[i][i]) < EPS)
        error = true, inconst |= allzero && abs(q[i][R+c]) > EPS;
37
38
      else sol[i][c] = (q[i][R+c] - t) / q[i][i]; // usually q[i][i]=1
    return sol;
40
41 }
```

3.7.3 Matrix Exp.

```
1 #include "header.h"
2 #define ITERATE_MATRIX(w) for (int r = 0; r < (w); ++r) \</pre>
               for (int c = 0: c < (w): ++c)
4 template <class T, int N>
5 struct M {
    array <array <T,N>,N> m;
    M() \{ ITERATE_MATRIX(N) m[r][c] = 0; \}
    static M id() {
      M I; for (int i = 0; i < N; ++i) I.m[i][i] = 1; return I;
    M operator*(const M &rhs) const {
11
12
      ITERATE MATRIX(N) for (int i = 0: i < N: ++i)
13
          out.m[r][c] += m[r][i] * rhs.m[i][c];
14
      return out;
15
   }
16
    M raise(ll n) const {
17
      if(n == 0) return id();
      if(n == 1) return *this:
      auto r = (*this**this).raise(n / 2);
      return (n%2 ? *this*r : r):
  }
22
23 };
```

3.7.4 Finite field For FFT

```
1 #include "header.h"
#include "../Number, Theory/elementary.cpp"
3 template<11 p,11 w> // prime, primitive root
4 struct Field { using T = Field; ll x; Field(ll x=0) : x{x} {}}
    T operator+(T r) const { return \{(x+r.x)\%p\}; }
    T operator-(T r) const { return {(x-r.x+p)%p}; }
    T operator*(T r) const { return {(x*r.x)%p}; }
    T operator/(T r) const { return (*this)*r.inv(); }
    T inv() const { return {mod_inverse(x,p)}; }
    static T root(11 k) { assert( (p-1)\%k==0 ); //(p-1)\%k==0?
      auto r = powmod(w,(p-1)/abs(k),p); // k-th root of unity
      return k>=0 ? T{r} : T{r}.inv();
    bool zero() const { return x == OLL: }
16 using F1 = Field < 1004535809,3 >;
17 using F2 = Field<1107296257,10>; // 1<<30 + 1<<25 + 1
18 using F3 = Field < 2281701377,3 >; // 1 < < 31 + 1 < < 27 + 1
```

3.7.5 Complex field For FFR

```
1 #include "header.h"
2 const double m_pi = M_PIf64x;
3 struct Complex { using T = Complex; double u,v;
    Complex(double u=0, double v=0) : u\{u\}, v\{v\} {}
    T operator+(T r) const { return {u+r.u. v+r.v}: }
    T operator - (T r) const { return {u-r.u, v-r.v}; }
    T operator*(T r) const { return {u*r.u - v*r.v, u*r.v + v*r.u}; }
    T operator/(T r) const {
      auto norm = r.u*r.u+r.v*r.v;
      return \{(u*r.u + v*r.v)/norm. (v*r.u - u*r.v)/norm\}:
    T operator*(double r) const { return T{u*r, v*r}; }
    T operator/(double r) const { return T{u/r, v/r}; }
    T inv() const { return T{1,0}/ *this; }
    T conj() const { return T{u, -v}; }
    static T root(ll k){ return {cos(2*m_pi/k), sin(2*m_pi/k)}; }
    bool zero() const { return max(abs(u), abs(v)) < 1e-6; }</pre>
17
18 };
```

3.7.6 FFT

```
1 #include "header.h"
2 #include "complex_field.cpp"
3 #include "fin_field.cpp"
4 void brinc(int &x, int k) {
5    int i = k - 1, s = 1 << i;
6    x ^= s;
7    if ((x & s) != s) {
8       --i; s >>= 1;
9    while (i >= 0 && ((x & s) == s))
10    x = x &~ s, --i, s >>= 1;
```

```
if (i >= 0) x |= s:
12
13 }
using T = Complex; // using T=F1,F2,F3
15 vector <T> roots;
16 void root_cache(int N) {
    if (N == (int)roots.size()) return;
    roots.assign(N, T{0});
    for (int i = 0: i < N: ++i)
      roots[i] = ((i\&-i) == i)
        ? T\{\cos(2.0*m_pi*i/N), \sin(2.0*m_pi*i/N)\}
21
        : roots[i&-i] * roots[i-(i&-i)]:
22
23 }
24 void fft(vector<T> &A, int p, bool inv = false) {
    int N = 1 << p;
    for(int i = 0, r = 0; i < N; ++i, brinc(r, p))
      if (i < r) swap(A[i], A[r]);</pre>
      Uncomment to precompute roots (for T=Complex). Slower but more precise.
     root_cache(N);
            sh=p-1
                         , --sh
    for (int m = 2; m <= N; m <<= 1) {
      T w. w m = T::root(inv ? -m : m):
      for (int k = 0; k < N; k += m) {
        w = T\{1\};
        for (int j = 0; j < m/2; ++ j) {
36 //
            T w = (!inv ? roots[j << sh] : roots[j << sh].conj());
          T t = w * A[k + j + m/2];
          A[k + i + m/2] = A[k + i] - t
          A[k + j] = A[k + j] + t;
          w = w * w m:
41
    if(inv){ T inverse = T(N).inv(); for(auto &x : A) x = x*inverse; }
44
45 }
     convolution leaves A and B in frequency domain state
     C may be equal to A or B for in-place convolution
48 void convolution(vector<T> &A, vector<T> &B, vector<T> &C){
    int s = A.size() + B.size() - 1;
    int q = 32 - \_builtin\_clz(s-1), N=1 << q; // fails if s=1
    A.resize(N,{}); B.resize(N,{}); C.resize(N,{});
    fft(A, q, false); fft(B, q, false);
    for (int i = 0; i < N; ++i) C[i] = A[i] * B[i];
    fft(C, q, true); C.resize(s);
55 }
56 void square_inplace(vector<T> &A) {
    int s = 2*A.size()-1, q = 32 - __builtin_clz(s-1), N=1<<q;
    A.resize(N,{}); fft(A, q, false);
    for (auto &x : A) x = x*x;
    fft(A, q, true); A.resize(s);
60
61 }
```

this code computes a_n in $O(k \log k \log n)$ time.

```
1 #include "header.h'
2 #include "poly.cpp'
```

```
1 #include "header.h"
2 #include "fft.cpp"
3 vector<T> &rev(vector<T> &A) { reverse(A.begin(), A.end()); return A; }
4 void copy into (const vector <T > &A. vector <T > &B. size t n) {
    std::copy(A.begin(), A.begin()+min({n, A.size(), B.size()}), B.begin());
6 }
     Multiplicative inverse of A modulo x^n. Requires A[0] != 0!!
9 vector<T> inverse(const vector<T> &A. int n) {
    vector <T> Ai{A[0].inv()};
    for (int k = 0; (1<<k) < n; ++k) {
      vector < T > As(4 << k, T(0)), Ais(4 << k, T(0));
      copy_into(A, As, 2<<k); copy_into(Ai, Ais, Ai.size());</pre>
      fft(As, k+2, false); fft(Ais, k+2, false);
      for (int i = 0; i < (4 << k); ++i) As[i] = As[i]*Ais[i]*Ais[i];
      fft(As, k+2, true); Ai.resize(2<<k, {});
16
      for (int i = 0; i < (2 < k); ++i) Ai[i] = T(2) * Ai[i] - As[i];
17
18
19
    Ai.resize(n);
    return Ai:
20
21 }
     Polynomial division. Returns \{Q, R\} such that A = QB+R, deg R < deg B.
     Requires that the leading term of B is nonzero.
  pair < vector < T > , vector < T > > divmod(const vector < T > & A , const vector < T > & B) {
    size_t n = A.size()-1, m = B.size()-1;
    if (n < m) return {vector <T>(1, T(0)), A};
    vector <T > X(A), Y(B), Q, R:
    convolution(rev(X), Y = inverse(rev(Y), n-m+1), Q);
    Q.resize(n-m+1): rev(Q):
30
31
    X.resize(Q.size()), copy_into(Q, X, Q.size());
    Y.resize(B.size()), copy_into(B, Y, B.size());
    convolution(X, Y, X);
    R.resize(m), copy_into(A, R, m);
    for (size_t i = 0; i < m; ++i) R[i] = R[i] - X[i];</pre>
    while (R.size() > 1 && R.back().zero()) R.pop_back();
    return {Q, R};
41 vector <T > mod(const vector <T > &A, const vector <T > &B) {
    return divmod(A, B).second;
43 }
```

$$a_n = \sum_{i=0}^{k-1} c_i a_{n-i-1}$$

```
3.7.7 Polyn. inv. div.
```

```
3 // x^k \mod f
4 vector<T> xmod(const vector<T> f, ll k) {
    vector <T> r{T(1)}:
    for (int b = 62: b >= 0: --b) {
      if (r.size() > 1)
        square_inplace(r), r = mod(r, f);
      if ((k>>b)&1) {
        r.insert(r.begin(), T(0));
        if (r.size() == f.size()) {
11
          T c = r.back() / f.back();
12
          for (size_t i = 0; i < f.size(); ++i)</pre>
            r[i] = r[i] - c * f[i];
          r.pop_back();
    return r;
19
     Given A[0,k) and C[0,k), computes the n-th term of:
     A[n] = \sum_{i=1}^{n} A[n-i-1]
    nth_term(const vector<T> &A, const vector<T> &C, 11 n) {
    int k = (int)A.size():
    if (n < k) return A[n];</pre>
    vector <T> f(k+1, T{1});
    for (int i = 0; i < k; ++i)
    f[i] = T\{-1\} * C[k-i-1];
    f = xmod(f, n):
    T r = T{0}:
    for (int i = 0; i < k; ++i)
      r = r + f[i] * A[i];
    return r;
36 }
```

3.7.9 Convolution Precise up to 9e15

```
1 #include "header.h"
2 #include "fft.cpp"
3 void convolution_mod(const vi &A, const vi &B, 11 MOD, vi &C) {
    int s = A.size() + B.size() - 1; ll m15 = (1LL << 15) - 1LL;
    int q = 32 - \_builtin\_clz(s-1), N=1 << q; // fails if s=1
    vector \langle T \rangle Ac(N), Bc(N), R1(N), R2(N);
    for (size_t i = 0; i < A.size(); ++i) Ac[i] = T{A[i]&m15, A[i]>>15};
    for (size_t i = 0; i < B.size(); ++i) Bc[i] = T{B[i]&m15, B[i]>>15};
    fft(Ac, q, false); fft(Bc, q, false);
    for (int i = 0, j = 0; i < N; ++i, j = (N-1)&(N-i)) {
      T as = (Ac[i] + Ac[j].conj()) / 2;
11
      T = (Ac[i] - Ac[i].coni()) / T{0, 2};
12
      T bs = (Bc[i] + Bc[j].conj()) / 2;
      T bl = (Bc[i] - Bc[j].conj()) / T{0, 2};
      R1[i] = as*bs + al*bl*T{0,1}, R2[i] = as*bl + al*bs;
15
    fft(R1, q, true); fft(R2, q, true);
```

3.8 Other Data Structures

3.8.1 Disjoint set (i.e. union-find)

```
1 template <typename T>
2 class DisjointSet {
       typedef T * iterator;
      T *parent, n, *rank;
       public:
           // O(n), assumes nodes are [0, n)
           DisjointSet(T n) {
                this->parent = new T[n]:
                this -> n = n;
                this->rank = new T[n]:
11
                for (T i = 0; i < n; i++) {</pre>
12
                    parent[i] = i;
13
                    rank[i] = 0;
14
1.5
           }
16
17
           // O(\log n)
18
           T find_set(T x) {
                if (x == parent[x]) return x;
20
                return parent[x] = find_set(parent[x]);
21
           }
22
23
           // O(log n)
24
           void union_sets(T x, T y) {
25
               x = this \rightarrow find set(x):
26
27
               v = this->find_set(v);
               if (x == y) return;
30
                if (rank[x] < rank[y]) {</pre>
                   Tz = x;
                    x = v;
                    y = z;
35
                parent[y] = x;
                if (rank[x] == rank[y]) rank[x]++;
40 };
```

3.8.2 Fenwick tree (i.e. BIT) eff. update + prefix sum calc.

```
1 #include "header.h"
2 #define maxn 200010
3 int t,n,m,tree[maxn],p[maxn];
5 void update(int k, int z) {
      while (k <= maxn) {
          tree[k] += z;
          k += k & (-k);
10 }
11
12 int sum(int k) {
      int ans = 0;
      while(k) {
15
          ans += tree[k];
          k = k & (-k):
16
18
      return ans;
19 }
```

3.8.3 Fenwick2d tree

```
1 #include "header.h"
2 template <class T>
3 struct FenwickTree2D {
    vector < vector <T> > tree;
    FenwickTree2D(int n): n(n) { tree.assign(n + 1, vector<T>(n + 1, 0)): }
    T query(int x1, int y1, int x2, int y2) {
      return query(x2,y2)+query(x1-1,y1-1)-query(x2,y1-1)-query(x1-1,y2);
    T query(int x, int y) {
11
      for (int i = x; i > 0; i -= (i & (-i)))
        for (int j = y; j > 0; j -= (j & (-j)))
          s += tree[i][j];
14
      return s;
15
16
    void update(int x, int y, T v) {
      for (int i = x; i \le n; i += (i & (-i)))
18
        for (int j = y; j <= n; j += (j & (-j)))
19
          tree[i][j] += v;
   }
21
22 };
```

3.8.4 Trie

```
#include "header.h"
const int ALPHABET_SIZE = 26;
inline int mp(char c) { return c - 'a'; }
##include "header.h"
const int ALPHABET_SIZE = 26;
for include "header.h"
const include
```

```
5 struct Node {
    Node* ch[ALPHABET_SIZE];
    bool isleaf = false;
    Node() {
       for(int i = 0; i < ALPHABET_SIZE; ++i) ch[i] = nullptr;</pre>
    }
11
    void insert(string &s, int i = 0) {
12
      if (i == s.length()) isleaf = true;
13
      else {
14
         int v = mp(s[i]);
15
         if (ch[v] == nullptr)
16
           ch[v] = new Node();
17
         ch[v] \rightarrow insert(s, i + 1);
18
19
    }
20
21
    bool contains(string &s, int i = 0) {
22
      if (i == s.length()) return isleaf;
      else {
         int v = mp(s[i]);
25
26
         if (ch[v] == nullptr) return false;
         else return ch[v]->contains(s, i + 1);
      }
28
    }
29
30
    void cleanup() {
31
      for (int i = 0: i < ALPHABET SIZE: ++i)</pre>
         if (ch[i] != nullptr) {
33
           ch[i]->cleanup();
34
           delete ch[i];
35
    }
37
38 };
```

3.8.5 Treap A binary tree whose nodes contain two values, a key and a priority, such that the key keeps the BST property

```
#include "header.h"

struct Node {
    ll v;
    int sz, pr;
    Node *l = nullptr, *r = nullptr;
    Node(ll val) : v(val), sz(1) { pr = rand(); }

    int size(Node *p) { return p ? p->sz : 0; }
    void update(Node* p) {
        if (!p) return;
        p->sz = 1 + size(p->l) + size(p->r);
        // Pull data from children here

    if (!p) return;
    // Push data to children here
```

```
17 }
18 void merge(Node *&t, Node *1, Node *r) {
     propagate(1), propagate(r);
    if (!1) t = r:
     else if (!r) t = 1;
    else if (1->pr > r->pr)
         merge(1->r, 1->r, r), t = 1;
    else merge(r\rightarrow 1, 1, r\rightarrow 1), t = r;
    update(t):
25
26 }
27 void spliti(Node *t, Node *&1, Node *&r, int index) {
     propagate(t);
    if (!t) { l = r = nullptr; return; }
    int id = size(t->1);
     if (index <= id) // id \in [index, \infty), so move it right</pre>
       spliti(t->1, 1, t->1, index), r = t;
    else
33
       spliti(t\rightarrow r, t\rightarrow r, r, index - id), l = t;
34
36 }
37 void splitv(Node *t, Node *&1, Node *&r, 11 val) {
     propagate(t):
    if (!t) { 1 = r = nullptr; return; }
    if (val \leftarrow t->v) // t->v \in [val, \infty), so move it right
       splitv(t->1, 1, t->1, val), r = t;
41
42
       splitv(t->r, t->r, r, val), l = t;
    update(t);
45 }
46 void clean(Node *p) {
    if (p) { clean(p->1), clean(p->r); delete p; }
48 }
```

4 Other Mathematics

4.1 Helpful functions

4.1.1 Euler's Totient Fucntion $n = p_1^{k_1-1} \cdot (p_1-1) \cdot \ldots \cdot p_r^{k_r-1} \cdot (p_r-1)$, where $p_1^{k_1} \cdot \ldots \cdot p_r^{k_r}$ is the prime factorization of n.

```
13     }
14     if (n > 1) ans *= n-1;
15     return ans;
16 }
17 vi phis(int n) { // All \Phi(i) up to n}
18     vi phi(n + 1, OLL);
19     iota(phi.begin(), phi.end(), OLL);
20     for (ll i = 2LL; i <= n; ++i)
21     if (phi[i] == i)
22         for (ll j = i; j <= n; j += i)
23         phi[j] -= phi[j] / i;
24     return phi;
25 }</pre>
```

Formulas $\Phi(n)$ counts all numbers in $1, \ldots, n-1$ coprime to n. $a^{\varphi(n)} \equiv 1 \mod n$, a and n are coprimes. $\forall e > \log_2 m : n^e \mod m = n^{\Phi(m) + e \mod \Phi(m)} \mod m$. $\gcd(m, n) = 1 \Rightarrow \Phi(m \cdot n) = \Phi(m) \cdot \Phi(n)$.

4.1.2 Pascal's trinagle $\binom{n}{k}$ is k-th element in the n-th row, indexing both from 0

4.2 Theorems and definitions

Fermat's little theorem

 $a^p \equiv a \mod p$

Subfactorial

$$!n = n! \sum_{i=0}^{n} \frac{(-1)^{i}}{i!}$$

$$!(0) = 1, !n = n \cdot !(n-1) + (-1)^{n}$$

Binomials and other partitionings

$$\binom{n}{k} = \binom{n-1}{k} + \binom{n-1}{k-1} = \prod_{i=1}^{k} \frac{n-i+1}{i}$$

This last product may be computed incrementally since any product of k' consecutive values is divisible by k'!.

Basic identities: The hockeystick identity:

$$\sum_{k=r}^{n} \binom{k}{r} = \binom{n+1}{r+1}$$

or

$$\sum_{k \le n} \binom{r+k}{k} = \binom{r+n+1}{n}$$

Also

$$\sum_{k=0}^{n} \binom{k}{m} = \binom{n+1}{m+1}$$
$$\sum_{i=0}^{n} \binom{n}{i} = 2^{n}$$

For $n, m \ge 0$ and p prime: write n, m in base p, i.e. $n = n_k p^k + \cdots + n_1 p + n_0$ and $m = m_k p^k + \cdots + m_1 p + m_0$. Then by Lucas theorem we have $\binom{n}{m} \equiv \prod_{i=0}^k \binom{n_i}{m_i} \mod p$, with the convention that $n_i < m_i \implies \binom{n_i}{m_i} = 0$.

Fibonacci (See also number theory section)

$$\sum_{0 \le k \le n} {n-k \choose k} = F_{n+1}$$

$$F_n = \frac{1}{\sqrt{5}} \left(\frac{1+\sqrt{5}}{2}\right)^n - \frac{1}{\sqrt{5}} \left(\frac{1-\sqrt{5}}{2}\right)^n$$

$$\sum_{i=1}^{n} F_i = F_{n+2} - 1, \ \sum_{i=1}^{n} F_i^2 = F_n F_{n+1}$$

$$\gcd(F_m, F_n) = F_{\gcd(m,n)}$$

$$\gcd(F_n, F_{n+1}) = \gcd(F_n, F_{n+2}) = 1$$

Bit stuff $a + b = a \oplus b + 2(a \& b) = a|b + a \& b$.

kth bit is set in x iff $x \mod 2^{k-1} \ge 2^k$, or iff $x \mod 2^{k-1} - x \mod 2^k \ne 0$ (i.e. $= 2^k$) It comes handy when you need to look at the bits of the numbers which are pair sums or subset sums etc.

 $n \mod 2^i = n\&(2^i - 1).$

$$\forall k: \ 1 \oplus 2 \oplus \ldots \oplus (4k-1) = 0$$

Stirling's numbers First kind: $S_1(n,k)$ count permutations on n items with k cycles. $S_1(n,k) = S_1(n-1,k-1) + (n-1)S_1(n-1,k)$ with $S_1(0,0) = 1$. Note $\sum_{k=0}^{n} S_1(n,k)x^k = x(x+1)\dots(x+n-1)$.

Second kind: $S_2(n,k)$ count partitions of n distinct elements into exactly k non-empty groups. $S_2(n,k) = S_2(n-1,k-1) + kS_2(n-1,k)$ with $S_2(n,1) = S_2(n,n) = 1$ and $S_2(n,k) = \frac{1}{k!} \sum_{i=0}^{k} (-1)^{k-i} {k \choose i} i^n$

4.3 Geometry Formulas

$$[ABC] = rs = \frac{1}{2}ab\sin\gamma = \frac{abc}{4R} = \sqrt{s(s-a)(s-b)(s-c)} = \frac{1}{2} |(B-A, C-A)^T|$$

$$s = \frac{a+b+c}{2} \qquad 2R = \frac{a}{\sin \alpha}$$
 cosine rule:
$$c^2 = a^2 + b^2 - 2ab\cos \gamma$$
 Euler:
$$1 + CC = V - E + F$$
 Pick:
$$\text{Area} = \text{interior points} + \frac{\text{boundary points}}{2} - 1$$

$$p \cdot q = |p||q|\cos(\theta) \qquad |p \times q| = |p||q|\sin(\theta)$$

Given a non-self-intersecting closed polygon on n vertices, given as (x_i, y_i) , its centroid (C_x, C_y) is given as:

$$C_x = \frac{1}{6A} \sum_{i=0}^{n-1} (x_i + x_{i+1})(x_i y_{i+1} - x_{i+1} y_i), \quad C_y = \frac{1}{6A} \sum_{i=0}^{n-1} (y_i + y_{i+1})(x_i y_{i+1} - x_{i+1} y_i)$$

$$A = \frac{1}{2} \sum_{i=0}^{n-1} (x_i y_{i+1} - x_{i+1} y_i) = \text{polygon area}$$

Inclusion-Exclusion

For appropriate f compute $\sum_{S\subseteq T} (-1)^{|T\setminus S|} f(S)$, or if only the size of S matters, $\sum_{s=0}^{n} (-1)^{n-s} {n \choose s} f(s)$. In some contexts we might use Stirling numbers, not binomial coefficients!

Some useful applications:

Graph coloring Let I(S) count the number of independent sets contained in $S \subseteq V$ $(I(\emptyset) = 1, I(S) = I(S \setminus v) + I(S \setminus N(v)))$. Let $c_k = \sum_{S \subseteq V} (-1)^{|V \setminus S|} I(S)$. Then V is k-colorable iff v > 0. Thus we can compute the chromatic number of a graph in $O^*(2^n)$ time.

Burnside's lemma

Given a group G acting on a set X, the number of elements in X up to symmetry is

$$\frac{1}{|G|} \sum_{g \in G} |X^g|$$

with X^g the elements of X invariant under g. For example, if f(n) counts "configurations" of some sort of length n, and we want to count them up to rotational symmetry using $G = \mathbb{Z}/n\mathbb{Z}$, then

$$g(n) = \frac{1}{n} \sum_{k=0}^{n-1} f(\gcd(n,k)) = \frac{1}{n} \sum_{k||n} f(k)\phi(n/k)$$

I.e. for coloring with c colors we have $f(k) = k^c$.

Relatedly, in Pólya's enumeration theorem we imagine X as a set of n beads with G permuting the beads (e.g. a necklace, with G all rotations and reflections of the n-cycle, i.e. the dihedral group D_n). Suppose further that we had Y colors, then the number of G-invariant colorings Y^X/G is counted by

$$\frac{1}{|G|} \sum_{g \in G} |Y|^{c(g)}$$

with c(g) counting the number of cycles of g when viewed as a permutation of X. We can generalize this to a weighted version: if the color i can occur exactly r_i times, then this is counted by the coefficient of $t_n^{r_1} cdots t_n^{r_n}$ in the polynomial

$$Z(t_1, \dots, t_n) = \frac{1}{|G|} \sum_{g \in G} \prod_{m \ge 1} (t_1^m + \dots + t_n^m)^{c_m(g)}$$

where $c_m(g)$ counts the number of length m cycles in g acting as a permutation on X. Note we get the original formula by setting all $t_i = 1$. Here Z is the cycle index. Note: you can cleverly deal with even/odd sizes by setting some t_i to -1.