typename Container <T>::const\_iterator beg =

18 #define vvpi vector < vpi>

6 #define ull unsigned ll

```
container.begin();
28  if (beg != container.end()) {
29    o << *beg++;
30    while (beg != container.end()) {
31         o << " " << *beg++;
32    }
33    }
34    return o;
35 }
36
37 // int main() {
38    // ios::sync_with_stdio(false);    // do not use cout + printf
39    // cin.tie(NULL);
40    // cout << fixed << setprecision(12);
41    // return 0;
42    // }</pre>
```

## 1.2 Bash for c++ compile with header.h

## 1.3 Bash for run tests c++

```
1 g++ $1/$1.cpp -o $1/$1.out
2 for file in $1/*.in; do diff <($1/$1.out < "$file
") "${file%.in}.ans"; done</pre>
```

# 1.4 Bash for run tests python

```
_{1} for file in 1/*.in; do diff <(python3 1/$1.py < "file") "${file%.in}.ans"; done
```

### 1.4.1 Aux. helper C++

```
1 #include "header.h"
3 int main() {
      // Read in a line including white space
      string line:
      getline(cin, line);
      // When doing the above read numbers as
          follows:
      int n;
      getline(cin. line):
      stringstream ss(line);
      ss >> n:
12
      // Count the number of 1s in binary
13
          represnatation of a number
      ull number;
      __builtin_popcountll(number);
16 }
```

#### 1.4.2 Aux. helper python

```
from functools import lru_cache

# Read until EOF
while True:
try:
pattern = input()
except EOFError:
break

Olru_cache(maxsize=None)
def smth_memoi(i, j, s):
# Example in-built cache
return "sol"
```

# 2 Python

# 2.1 Graphs

### 2.1.1 BFS

```
if node in explored: continue
explored.add(node)
for neigh in g[node]:
    if neigh not in explored:
        q.append(neigh)
distances[neigh] = distances[node] + 1
return distances
```

## 2.1.2 Dijkstra

```
1 from heapq import *
2 def dijkstra(n, root, g): # g = {node: (cost,
    dist = [float("inf")]*n
    dist[root] = 0
    prev = [-1]*n
    pq = [(0, root)]
    heapify(pq)
    visited = set([])
    while len(pq) != 0:
      _, node = heappop(pq)
13
      if node in visited: continue
      visited.add(node)
      # In case of disconnected graphs
18
      if node not in g:
        continue
19
      for cost, neigh in g[node]:
        alt = dist[node] + cost
        if alt < dist[neigh]:</pre>
          dist[neigh] = alt
          prev[neigh] = node
          heappush(pq, (alt, neigh))
    return dist
```

## 2.1.3 Topological Sort

57

58

62

63

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66

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70

71

73

```
# function to add an edge to graph
      def addEdge(self,u,v):
11
          self.graph[u].append(v)
12
13
      # A recursive function used by
14
          topologicalSort
      def topologicalSortUtil(self,v,visited,stack)
15
16
          # Mark the current node as visited.
17
          visited[v] = True
18
19
          # Recur for all the vertices adjacent to
20
              this vertex
          for i in self.graph[v]:
              if visited[i] == False:
22
                   self.topologicalSortUtil(i,
23
                      visited, stack)
          # Push current vertex to stack which
25
              stores result
          stack.insert(0.v)
26
      # The function to do Topological Sort. It
28
          uses recursive
      # topologicalSortUtil()
29
      def topologicalSort(self):
30
          # Mark all the vertices as not visited
31
          visited = [False]*self.V
32
          stack =[]
33
34
          # Call the recursive helper function to
              store Topological
          # Sort starting from all vertices one by
          for i in range(self.V):
              if visited[i] == False:
                   self.topologicalSortUtil(i,
                       visited, stack)
40
          # Print contents of stack
41
          return stack
42
43
      def isCyclicUtil(self, v, visited, recStack):
44
45
          # Mark current node as visited and
46
47
          # adds to recursion stack
          visited[v] = True
          recStack[v] = True
          # Recur for all neighbours
51
          # if any neighbour is visited and in
52
          # recStack then graph is cyclic
53
```

for neighbour in self.graph[v]:

if visited[neighbour] == False:

54

55

```
if self.isCyclicUtil(neighbour,
                visited, recStack) == True:
                return True
       elif recStack[neighbour] == True:
            return True
   # The node needs to be popped from
   # recursion stack before function ends
   recStack[v] = False
   return False
# Returns true if graph is cyclic else false
def isCyclic(self):
   visited = [False] * (self.V + 1)
   recStack = [False] * (self.V + 1)
   for node in range(self.V):
       if visited[node] == False:
            if self.isCyclicUtil(node,
                visited, recStack) == True:
                return True
   return False
```

## 2.1.4 Kruskal (UnionFind)

```
class UnionFind:
      def __init__(self, n):
          self.parent = [-1]*n
      def find(self, x):
          if self.parent[x] < 0:</pre>
              return x
          self.parent[x] = self.find(self.parent[x
              ])
          return self.parent[x]
10
      def connect(self, a, b):
11
          ra = self.find(a)
          rb = self.find(b)
13
          if ra == rb:
14
               return False
15
          if self.parent[ra] > self.parent[rb]:
               self.parent[rb] += self.parent[ra]
               self.parent[ra] = rb
18
19
               self.parent[ra] += self.parent[rb]
               self.parent[rb] = ra
21
          return True
24 # Full MST is len(spanning==n-1)
25 def kruskal(n, edges):
      uf = UnionFind(n)
      spanning = []
      edges.sort(key = lambda d: -d[2])
      while edges and len(spanning) < n-1:
```

## 2.2 Num. Th. / Comb.

## 2.2.1 nCk % prime

```
# Note: p must be prime and k  n:
        return 0

# calculate numerator

num = 1

for i in range(n-k+1, n+1):
        num *= i % p

num %= p

# calculate denominator

denom = 1

for i in range(1,k+1):
        denom *= i % p

denom %= p

# numerator * denominator^(p-2) (mod p)

return (num * pow(denom, p-2, p)) % p
```

# **2.2.2** Sieve of E. O(n) so actually faster than C++ version, but more memory

```
1 MAX SIZE = 10**8+1
2 isprime = [True] * MAX_SIZE
3 \text{ prime} = []
4 SPF = [None] * (MAX SIZE)
6 def manipulated_seive(N): # Up to N (not
      included)
    isprime[0] = isprime[1] = False
    for i in range(2, N):
      if isprime[i] == True:
         prime.append(i)
         SPF[i] = i
11
      i = 0
12
      while (j < len(prime) and
        i * prime[j] < N and
14
           prime[j] <= SPF[i]):</pre>
15
         isprime[i * prime[j]] = False
         SPF[i * prime[j]] = prime[j]
17
         j += 1
```

# 2.3 Strings

#### 2.3.1 LCS

```
1 def longestCommonSubsequence(text1, text2): # 0(
      m*n) time. O(m) space
      n = len(text1)
      m = len(text2)
      # Initializing two lists of size m
      prev = [0] * (m + 1)
      cur = \lceil 0 \rceil * (m + 1)
      for idx1 in range(1, n + 1):
          for idx2 in range(1, m + 1):
10
              # If characters are matching
11
              if text1[idx1 - 1] == text2[idx2 -
                   cur[idx2] = 1 + prev[idx2 - 1]
               else:
                   # If characters are not matching
                   cur[idx2] = max(cur[idx2 - 1],
                       prev[idx2])
17
          prev = cur.copy()
19
      return cur[m]
```

#### 2.3.2 KMP

```
1 class KMP:
      def partial(self, pattern):
          """ Calculate partial match table: String
               -> [Int]"""
          ret = [0]
          for i in range(1, len(pattern)):
              j = ret[i - 1]
              while j > 0 and pattern[j] != pattern
                  [i]: j = ret[j - 1]
              ret.append(j + 1 if pattern[j] ==
                  pattern[i] else j)
          return ret
10
      def search(self, T, P):
11
          """KMP search main algorithm: String ->
12
              String -> [Int]
          Return all the matching position of
              pattern string P in T"""
          partial, ret, j = self.partial(P), [], 0
14
          for i in range(len(T)):
15
              while j > 0 and T[i] != P[j]: j =
                  partial[j - 1]
              if T[i] == P[j]: j += 1
              if i == len(P):
```

```
ret.append(i - (j - 1))
j = partial[j - 1]
return ret
```

#### 2.3.3 Edit distance

def editDistance(str1, str2):

```
# Get the lengths of the input strings
    m = len(str1)
    n = len(str2)
    # Initialize a list to store the current row
    curr = [0] * (n + 1)
    # Initialize the first row with values from 0
    for j in range(n + 1):
      curr[i] = i
    # Initialize a variable to store the previous
        value
    previous = 0
    # Loop through the rows of the dynamic
        programming matrix
    for i in range (1, m + 1):
      # Store the current value at the beginning of
           the row
      previous = curr[0]
      curr[0] = i
      # Loop through the columns of the dynamic
          programming matrix
      for j in range(1, n + 1):
        # Store the current value in a temporary
24
            variable
        temp = curr[i]
        # Check if the characters at the current
            positions in str1 and str2 are the same
        if str1[i - 1] == str2[j - 1]:
          curr[j] = previous
30
        else:
          # Update the current cell with the
              minimum of the three adjacent cells
          curr[j] = 1 + min(previous, curr[j - 1],
              curr[i])
        # Update the previous variable with the
            temporary value
        previous = temp
    # The value in the last cell represents the
        minimum number of operations
```

38 return curr[n]

## 2.4 Other Algorithms

#### 2.4.1 Rotate matrix

```
1 def rotate_matrix(m):
2    return [[m[j][i] for j in range(len(m))] for
                    i in range(len(m[0])-1,-1,-1)]
```

## 2.5 Geometry

#### 2.5.1 Convex Hull

```
1 def vec(a,b):
      return (b[0]-a[0],b[1]-a[1])
3 def det(a,b):
      return a[0]*b[1] - b[0]*a[1]
6 def convexhull(P):
      if (len(P) == 1):
          return [(p[0][0], p[0][1])]
      h = sorted(P)
      lower = []
      i = 0
      while i < len(h):
          if len(lower) > 1:
              a = vec(lower[-2], lower[-1])
              b = vec(lower[-1], h[i])
              if det(a,b) <= 0 and len(lower) > 1:
                  lower.pop()
                   continue
          lower.append(h[i])
21
          i += 1
      upper = []
23
      i = 0
      while i < len(h):
          if len(upper) > 1:
26
              a = vec(upper[-2], upper[-1])
              b = vec(upper[-1], h[i])
              if det(a,b) >= 0:
                   upper.pop()
                   continue
          upper.append(h[i])
32
          i += 1
      reversedupper = list(reversed(upper[1:-1:]))
      reversedupper.extend(lower)
      return reversedupper
```

## 2.5.2 Geometry

```
2 def vec(a,b):
      return (b[0]-a[0],b[1]-a[1])
5 def det(a,b):
      return a[0]*b[1] - b[0]*a[1]
      lower = []
      i = 0
      while i < len(h):
          if len(lower) > 1:
              a = vec(lower[-2], lower[-1])
12
              b = vec(lower[-1], h[i])
              if det(a,b) \le 0 and len(lower) > 1:
                  lower.pop()
15
                   continue
16
          lower.append(h[i])
17
          i += 1
19
      # find upper hull
20
      # det <= 0 -> replace
21
      upper = []
22
      i = 0
23
      while i < len(h):
          if len(upper) > 1:
25
              a = vec(upper[-2], upper[-1])
26
              b = vec(upper[-1], h[i])
              if det(a,b) >= 0:
                  upper.pop()
                  continue
          upper.append(h[i])
          i += 1
```

## 2.6 Other Data Structures

#### 2.6.1 Segment Tree

```
N = 100000 # limit for array size
tree = [0] * (2 * N) # Max size of tree

def build(arr, n): # function to build the tree
    # insert leaf nodes in tree
    for i in range(n):
        tree[n + i] = arr[i]

# build the tree by calculating parents
for i in range(n - 1, 0, -1):
        tree[i] = tree[i << 1] + tree[i << 1 | 1]

def updateTreeNode(p, value, n): # function to update a tree node
# set value at position p</pre>
```

```
tree[p + n] = value
16
      p = p + n
17
      i = p # move upward and update parents
18
          tree[i >> 1] = tree[i] + tree[i ^ 1]
20
21
          i >>= 1
23 def querv(1, r, n): # function to get sum on
      interval [1, r)
      # loop to find the sum in the range
      r += n
      while 1 < r:
          if 1 & 1:
              res += tree[1]
              1 += 1
          if r & 1:
              r -= 1
              res += tree[r]
34
          1 >>= 1
          r >>= 1
      return res
```

#### 2.6.2 Trie

```
1 class TrieNode:
      def __init__(self):
          self.children = [None] *26
          self.isEndOfWord = False
6 class Trie:
      def __init__(self):
          self.root = self.getNode()
      def getNode(self):
10
          return TrieNode()
11
12
      def charToIndex(self.ch):
13
          return ord(ch)-ord('a')
14
15
16
      def insert(self,key):
17
          pCrawl = self.root
18
          length = len(key)
19
          for level in range(length):
               index = self._charToIndex(key[level])
21
               if not pCrawl.children[index]:
                   pCrawl.children[index] = self.
                       getNode()
               pCrawl = pCrawl.children[index]
          pCrawl.isEndOfWord = True
25
```

```
def search(self, key):

pCrawl = self.root

length = len(key)

for level in range(length):
 index = self._charToIndex(key[level])

if not pCrawl.children[index]:

return False

pCrawl = pCrawl.children[index]

return pCrawl.isEndOfWord
```

## 3 C++

## 3.1 Graphs

#### 3.1.1 BFS

```
1 #include "header.h"
2 #define graph unordered_map<11, unordered_set<11</pre>
3 vi bfs(int n, graph& g, vi& roots) {
      vi parents(n+1, -1): // nodes are 1..n
      unordered_set < int > visited;
      queue < int > q;
      for (auto x: roots) {
          g.emplace(x):
           visited.insert(x);
9
10
      while (not q.empty()) {
11
          int node = q.front();
12
           q.pop();
13
           for (auto neigh: g[node]) {
               if (not in(neigh, visited)) {
                   parents[neigh] = node;
                   q.emplace(neigh);
                   visited.insert(neigh):
              }
          }
21
      return parents;
25 vi reconstruct_path(vi parents, int start, int
      goal) {
      vi path;
      int curr = goal;
      while (curr != start) {
           path.push_back(curr);
           if (parents[curr] == -1) return vi(); //
               No path, empty vi
           curr = parents[curr];
31
32
      path.push_back(start);
```

```
reverse(path.begin(), path.end());
return path;

36 }
```

## **3.1.2 DFS** Cycle detection / removal

```
1 #include "header.h"
2 void removeCvc(ll node, unordered map<ll, vector<
      pair<11, 11>>>& neighs, vector<bool>& visited
3 vector < bool > & recStack, vector < 11 > & ans) {
      if (!visited[node]) {
          visited[node] = true:
          recStack[node] = true:
          auto it = neighs.find(node);
          if (it != neighs.end()) {
              for (auto util: it->second) {
                  11 nnode = util.first;
                  if (recStack[nnode]) {
                       ans.push_back(util.second);
                  } else if (!visited[nnode]) {
                      removeCyc(nnode, neighs,
                           visited, recStack, ans);
                  }
              }
          }
17
18
      recStack[node] = false;
19
```

#### 3.1.3 Dijkstra

```
1 #include "header.h"
2 vector<int> dijkstra(int n, int root, map<int,</pre>
      vector<pair<int, int>>>& g) {
    unordered_set <int> visited;
    vector < int > dist(n. INF):
      priority_queue < pair < int , int >> pq;
      dist[root] = 0:
      pq.push({0, root});
      while (!pq.empty()) {
          int node = pq.top().second;
          int d = -pq.top().first;
11
          pq.pop();
          if (in(node, visited)) continue;
          visited.insert(node):
14
15
          for (auto e : g[node]) {
16
              int neigh = e.first;
17
              int cost = e.second;
               if (dist[neigh] > dist[node] + cost)
```

### 3.1.4 Floyd-Warshall

# **3.1.5** Kruskal Minimum spanning tree of undirected weighted graph

```
1 #include "header.h"
2 #include "disjoint_set.h"
3 // O(E log E)
4 pair < set < pair < 11, 11 >> , 11 > kruskal (vector < tuple</pre>
       <11. 11. 11>>& edges. 11 n) {
       set <pair <11, 11>> ans;
       11 cost = 0:
       sort(edges.begin(), edges.end());
       DisjointSet < 11 > fs(n);
10
       ll dist, i, j;
11
12
       for (auto edge: edges) {
           dist = get<0>(edge);
13
           i = get < 1 > (edge):
14
           i = get < 2 > (edge);
15
16
           if (fs.find_set(i) != fs.find_set(j)) {
               fs.union_sets(i, j);
18
               ans.insert({i, j});
                cost += dist:
           }
21
22
       return pair < set < pair < 11, 11>>, 11> {ans, cost
23
           };
24 }
```

#### 3.1.6 Hungarian algorithm

```
1 #include "header.h"
3 template <class T> bool ckmin(T &a. const T &b) {
       return b < a ? a = b, 1 : 0; }
4 /**
5 * Given J jobs and W workers (J <= W), computes</pre>
       the minimum cost to assign each
6 * prefix of jobs to distinct workers.
7 * @tparam T a type large enough to represent
       integers on the order of J *
9 * @param C a matrix of dimensions JxW such that
       C[j][w] = cost to assign j-th
* job to w-th worker (possibly negative)
12 * @return a vector of length J, with the j-th
       entry equaling the minimum cost
* to assign the first (j+1) jobs to distinct
       workers
14 */
15 template <class T> vector <T> hungarian(const
      vector < vector < T >> &C) {
      const int J = (int)size(C), W = (int)size(C
          [0]):
      assert(J <= W);</pre>
      // job[w] = job assigned to w-th worker, or
          -1 if no job assigned
      // note: a W-th worker was added for
          convenience
      vector < int > job(W + 1, -1);
20
      vector<T> ys(J), yt(W + 1); // potentials
      // -yt[W] will equal the sum of all deltas
      vector <T> answers;
      const T inf = numeric_limits <T>::max();
      for (int j_cur = 0; j_cur < J; ++j_cur) { //</pre>
           assign j_cur-th job
          int w_cur = W;
          job[w_cur] = j_cur;
27
          // min reduced cost over edges from Z to
              worker w
          vector <T> min_to(W + 1, inf);
          vector<int> prv(W + 1, -1); // previous
              worker on alternating path
          vector < bool > in_Z(W + 1);  // whether
              worker is in Z
          while (job[w_cur] != -1) {  // runs at
              most j_cur + 1 times
              in_Z[w_cur] = true;
33
              const int j = job[w_cur];
              T delta = inf;
              int w next:
              for (int w = 0; w < W; ++w) {
                  if (!in_Z[w]) {
38
                      if (ckmin(min_to[w], C[j][w]
```

```
- ys[j] - yt[w]))
                          prv[w] = w_cur;
                      if (ckmin(delta. min to[w]))
                          w next = w:
                  }
              }
              // delta will always be non-negative,
              // except possibly during the first
                  time this loop runs
              // if any entries of C[j_cur] are
                  negative
              for (int w = 0: w \le W: ++w) {
                  if (in_Z[w]) ys[job[w]] += delta,
                       yt[w] -= delta;
                  else min_to[w] -= delta;
              }
              w_cur = w_next;
51
52
          // update assignments along alternating
          for (int w; w_cur != W; w_cur = w) job[
              w curl = iob[w = prv[w curll:
          answers.push_back(-yt[W]);
56
      return answers:
57
58 }
```

# **3.1.7** Suc. shortest path Calculates max flow, min cost

```
1 #include "header.h"
2 // map<node, map<node, pair<cost, capacity>>>
3 #define graph unordered_map<int, unordered_map<</pre>
      int, pair<ld, int>>>
4 graph g;
5 const ld infty = 1e60l; // Change if necessary
6 ld fill(int n, vld& potential) { // Finds max
      flow, min cost
    priority_queue < pair < ld, int >> pq;
    vector <bool > visited(n+2, false);
    vi parent(n+2, 0);
    vld dist(n+2, infty);
    dist[0] = 0.1:
    pq.emplace(make_pair(0.1, 0));
    while (not pq.empty()) {
      int node = pq.top().second;
14
      pq.pop();
      if (visited[node]) continue;
      visited[node] = true;
17
      for (auto& x : g[node]) {
18
        int neigh = x.first;
19
        int capacity = x.second.second;
        ld cost = x.second.first;
21
        if (capacity and not visited[neigh]) {
```

```
ld d = dist[node] + cost + potential[node
              ] - potential[neigh];
          if (d + 1e-10l < dist[neigh]) {</pre>
             dist[neigh] = d:
             pq.emplace(make_pair(-d, neigh));
             parent[neigh] = node;
    1111
    for (int i = 0: i < n+2: i++) {</pre>
      potential[i] = min(infty, potential[i] + dist
    if (not parent[n+1]) return infty;
    ld ans = 0.1:
    for (int x = n+1; x; x=parent[x]) {
      ans += g[parent[x]][x].first;
      g[parent[x]][x].second--;
      g[x][parent[x]].second++;
    return ans;
41 }
```

### 3.1.8 Bipartite check

```
1 #include "header.h"
2 int main() {
      int n;
      vvi adi(n):
      vi side(n. -1):
                          // will have 0's for one
          side 1's for other side
      bool is_bipartite = true; // becomes false
          if not bipartite
      aueue < int > a:
      for (int st = 0; st < n; ++st) {</pre>
9
          if (side[st] == -1) {
10
               q.push(st);
               side[st] = 0;
12
               while (!q.empty()) {
13
                   int v = q.front();
14
                   q.pop();
                   for (int u : adj[v]) {
                       if (side[u] == -1) {
17
                           side[u] = side[v] ^ 1;
18
                           q.push(u);
                       } else {
                           is bipartite &= side[u]
                               != side[v];
                       }
23 }}}}
```

## 3.1.9 Find cycle directed

```
1 #include "header.h"
2 int n;
3 const int mxN = 2e5+5:
4 vvi adj(mxN);
5 vector < char > color;
6 vi parent:
7 int cycle_start, cycle_end;
8 bool dfs(int v) {
      color[v] = 1:
      for (int u : adj[v]) {
           if (color[u] == 0) {
               parent[u] = v;
               if (dfs(u)) return true;
          } else if (color[u] == 1) {
14
               cvcle_end = v;
               cycle_start = u;
               return true;
18
19
      color[v] = 2;
      return false;
22 }
23 void find_cycle() {
      color.assign(n, 0);
      parent.assign(n, -1):
      cycle_start = -1;
      for (int v = 0: v < n: v++) {
27
           if (color[v] == 0 && dfs(v))break;
30
      if (cycle_start == -1) {
           cout << "Acyclic" << endl;</pre>
32
           vector<int> cycle;
33
           cycle.push_back(cycle_start);
           for (int v = cycle_end; v != cycle_start;
35
                v = parent[v])
               cycle.push_back(v);
           cycle.push_back(cycle_start);
37
           reverse(cycle.begin(), cycle.end());
39
           cout << "Cvcle..Found:..":
           for (int v : cycle) cout << v << "";</pre>
           cout << endl:
42
44 }
```

## 3.1.10 Find cycle directed

```
#include "header.h"
int n;
const int mxN = 2e5 + 5;
vvi adj(mxN);
vector<bool> visited;
vi parent;
```

```
7 int cycle_start, cycle_end;
8 bool dfs(int v, int par) { // passing vertex and
      its parent vertex
      visited[v] = true:
      for (int u : adj[v]) {
10
          if(u == par) continue; // skipping edge
11
              to parent vertex
          if (visited[u]) {
              cvcle end = v:
               cycle_start = u;
               return true:
          parent[u] = v;
17
          if (dfs(u, parent[u]))
              return true;
19
20
      return false;
21
23 void find_cycle() {
      visited.assign(n, false);
      parent.assign(n, -1);
25
      cvcle start = -1:
      for (int v = 0; v < n; v++) {</pre>
          if (!visited[v] && dfs(v, parent[v]))
              break:
      if (cycle_start == -1) {
           cout << "Acvclic" << endl:</pre>
      } else {
32
          vector < int > cycle;
33
           cycle.push_back(cycle_start);
34
          for (int v = cycle_end; v != cycle_start;
               v = parent[v])
               cycle.push_back(v);
           cycle.push_back(cycle_start);
37
          cout << "Cycle_Found:_";
          for (int v : cycle) cout << v << "";</pre>
           cout << endl:
41
```

## 3.1.11 Tarjan's SCC

```
#include "header.h"

struct Tarjan {
    vvi &edges;
    int V, counter = 0, C = 0;
    vi n, 1;
    vector<bool> vs;
    stack<int> st;
    Tarjan(vvi &e) : edges(e), V(e.size()), n(V, -1), 1(V, -1), vs(V, false) {}
    void visit(int u, vi &com) {
```

```
l[u] = n[u] = counter++:
      st.push(u);
      vs[u] = true:
13
      for (auto &&v : edges[u]) {
        if (n[v] == -1) visit(v, com);
        if (vs[v]) 1[u] = min(1[u], 1[v]);
      if (1[u] == n[u]) {
18
        while (true) {
          int v = st.top();
21
          st.pop();
          vs[v] = false:
          com[v] = C; // <== ACT HERE
          if (u == v) break:
        C++;
26
      }
    int find_sccs(vi &com) { // component indices
        will be stored in 'com'
      com.assign(V, -1);
      C = 0:
      for (int u = 0; u < V; ++u)</pre>
        if (n[u] == -1) visit(u, com);
      return C;
    // scc is a map of the original vertices of the
         graph to the vertices
    // of the SCC graph, scc_graph is its adjacency
         list.
    // SCC indices and edges are stored in 'scc'
        and 'scc_graph'.
    void scc_collapse(vi &scc, vvi &scc_graph) {
      find sccs(scc):
      scc graph.assign(C, vi()):
      set <pi>rec; // recorded edges
      for (int u = 0; u < V; ++u) {
        assert(scc[u] != -1):
        for (int v : edges[u]) {
          if (scc[v] == scc[u] ||
            rec.find({scc[u], scc[v]}) != rec.end()
                ) continue:
          scc_graph[scc[u]].push_back(scc[v]);
          rec.insert({scc[u]. scc[v]}):
        }
51
52
    // Function to find sources and sinks in the
    // The number of edges needed to be added is
        max(sources.size(), sinks.())
    void findSourcesAndSinks(const vvi &scc graph.
        vi &sources, vi &sinks) {
      vi in_degree(C, 0), out_degree(C, 0);
      for (int u = 0; u < C; u++) {
```

# **3.1.12** SCC edges Prints out the missing edges to make the input digraph strongly connected

```
1 #include "header.h"
2 const int N=1e5+10:
3 int n,a[N],cnt[N],vis[N];
4 vector<int> hd.tl:
5 int dfs(int x){
      vis[x]=1;
       if(!vis[a[x]])return vis[x]=dfs(a[x]);
      return vis[x]=x;
10 int main(){
       scanf("%d",&n);
      for(int i=1;i<=n;i++){</pre>
           scanf("%d",&a[i]);
           cnt[a[i]]++;
1.4
      int k=0;
       for(int i=1:i<=n:i++){</pre>
           if(!cnt[i]){
               k++;
               hd.push_back(i);
               tl.push_back(dfs(i));
      }
      int tk=k:
      for(int i=1;i<=n;i++){</pre>
           if(!vis[i]){
               hd.push back(i):
               tl.push_back(dfs(i));
      if(k==1&&!tk)k=0;
       printf("%d\n",k):
      for(int i=0;i<k;i++)printf("%du%d\n",tl[i],hd</pre>
           [(i+1)%k]);
      return 0;
35
```

### 3.1.13 Find Bridges

```
1 #include "header.h"
2 int n: // number of nodes
3 vvi adj; // adjacency list of graph
4 vector < bool > visited;
5 vi tin. low:
6 int timer;
7 void dfs(int v, int p = -1) {
      visited[v] = true;
      tin[v] = low[v] = timer++;
      for (int to : adj[v]) {
          if (to == p) continue;
          if (visited[to]) {
              low[v] = min(low[v], tin[to]):
          } else {
              dfs(to, v);
15
              low[v] = min(low[v], low[to]);
              if (low[to] > tin[v])
                   IS BRIDGE(v. to):
          }
21 }
22 void find_bridges() {
      timer = 0;
      visited.assign(n, false);
      tin.assign(n, -1);
      low.assign(n, -1);
      for (int i = 0; i < n; ++i) {</pre>
27
          if (!visited[i]) dfs(i);
29
30 }
```

## **3.1.14** Artic. points (i.e. cut off points)

```
1 #include "header.h"
2 int n; // number of nodes
3 vvi adj; // adjacency list of graph
4 vector < bool > visited:
5 vi tin, low;
6 int timer:
7 \text{ void } dfs(int v, int p = -1) {
      visited[v] = true;
      tin[v] = low[v] = timer++:
      int children=0;
      for (int to : adi[v]) {
          if (to == p) continue;
          if (visited[to]) {
              low[v] = min(low[v], tin[to]);
          } else {
              dfs(to, v);
16
              low[v] = min(low[v], low[to]);
              if (low[to] >= tin[v] && p!=-1)
                   IS_CUTPOINT(v);
              ++children:
```

#### 3.1.15 Topological sort

```
1 #include "header.h"
2 int n: // number of vertices
3 vvi adi: // adiacency list of graph
4 vector <bool> visited;
5 vi ans:
6 void dfs(int v) {
      visited[v] = true;
      for (int u : adj[v]) {
          if (!visited[u]) dfs(u);
10
       ans.push_back(v);
11
12 }
13 void topological_sort() {
      visited.assign(n, false);
       ans.clear():
      for (int i = 0; i < n; ++i) {</pre>
           if (!visited[i]) dfs(i);
17
19
      reverse(ans.begin(), ans.end());
20 }
```

# **3.1.16 Bellmann-Ford** Same as Dijkstra but allows neg. edges

```
if (dist[u] == INF) continue;
        else for (auto &e : e[u])
          if(dist[u]+e.second < dist[e.first])</pre>
12
            dist[e.first] = dist[u]+e.second. relax
                 = true:
      if(!relax) break;
    bool ch = true;
    while (ch) {
                         // keep going untill no
        more changes
      ch = false:
                         // set dist to -INF when in
           cvcle
      for (int u = 0; u < e.size(); ++u)</pre>
        if (dist[u] == INF) continue;
        else for (auto &e : e[u])
          if (dist[e.first] > dist[u] + e.second
            && !cyc[e.first]) {
            dist[e.first] = -INF;
            ch = true; //return true for cycle
                detection only
            cvc[e.first] = true;
   }
29 }
```

#### **3.1.17 Ford-Fulkerson** Basic Max. flow

```
1 #include "header.h"
2 #define V 6 // Num. of vertices in given graph
4 /* Returns true if there is a path from source 's
5 't' in residual graph. Also fills parent[] to
      store the
6 path */
7 bool bfs(int rGraph[V][V], int s, int t, int
      parent[]) {
8 bool visited[V]:
   memset(visited, 0, sizeof(visited));
   queue < int > q;
   q.push(s);
   visited[s] = true;
    parent[s] = -1;
   // Standard BFS Loop
    while (!q.empty()) {
     int u = q.front();
     q.pop();
      for (int v = 0; v < V; v++) {
       if (visited[v] == false && rGraph[u][v] >
          if (v == t) {
            parent[v] = u:
```

```
return true:
          q.push(v);
          parent[v] = u:
          visited[v] = true;
    return false:
35 // Returns the maximum flow from s to t in the
36 int fordFulkerson(int graph[V][V], int s, int t)
    int u, v;
    int rGraph[V]
    for (u = 0; u < V; u++)</pre>
     for (v = 0: v < V: v++)
        rGraph[u][v] = graph[u][v];
42
43
    int parent[V]; // This array is filled by BFS
        and to
          // store path
    int max_flow = 0; // There is no flow initially
    while (bfs(rGraph, s, t, parent)) {
      int path flow = INT MAX:
      for (v = t; v != s; v = parent[v]) {
        u = parent[v]:
        path_flow = min(path_flow, rGraph[u][v]);
51
52
      for (v = t; v != s; v = parent[v]) {
54
        u = parent[v]:
        rGraph[u][v] -= path_flow;
        rGraph[v][u] += path_flow;
57
      max_flow += path_flow;
    return max_flow;
62 }
```

## **3.1.18 Dinic max flow** $O(V^2E)$ , O(Ef)

```
S(int v, int ri, Fc, W cost = 0):
          v(v), r(ri), f(0), cap(c), cost(cost) {}
      inline F res() const { return cap - f; }
11
12 }:
13 struct FlowGraph : vector < vector < S >> {
      FlowGraph(size_t n) : vector < vector < S >> (n) {}
      void add_edge(int u, int v, F c, W cost = 0){
           auto &t = *this:
          t[u].emplace back(v, t[v].size(), c, cost
16
          t[v].emplace_back(u, t[u].size()-1, c, -
              cost):
      void add_arc(int u, int v, F c, W cost = 0){
          auto &t = *this;
          t[u].emplace_back(v, t[v].size(), c, cost
          t[v].emplace_back(u, t[u].size()-1, 0, -
              cost):
22
      void clear() { for (auto &E : *this) for (
23
          auto &e : E) e.f = OLL: }
24 };
25 struct Dinic{
      FlowGraph & edges; int V,s,t;
      vi l; vector < vector < S > :: iterator > its; //
          levels and iterators
      Dinic(FlowGraph &edges, int s, int t) :
          edges(edges), V(edges.size()), s(s), t(t)
29
              , 1(V,-1), its(V) {}
      ll augment(int u, F c) { // we reuse the same
           iterators
          if (u == t) return c: ll r = OLL:
31
          for(auto &i = its[u]; i != edges[u].end()
32
              : i++){
              auto &e = *i:
              if (e.res() && 1[u] < 1[e.v]) {</pre>
                   auto d = augment(e.v, min(c, e.
                   if (d > 0) { e.f += d; edges[e.v
36
                       ][e.r].f -= d; c -= d;
                       r += d; if (!c) break; }
          return r:
40
      11 run() {
42
          11 \text{ flow} = 0, f;
          while(true) {
               fill(1.begin(), 1.end(),-1); 1[s]=0;
                   // recalculate the layers
              queue < int > q; q.push(s);
               while(!a.emptv()){
46
                   auto u = q.front(); q.pop(); its[
                       u] = edges[u].begin();
                   for(auto &&e : edges[u]) if(e.res
```

# 3.2 Dynamic Programming

## 3.2.1 Longest Incr. Subseq.

```
1 #include "header.h"
2 template < class T>
3 vector<T> index path lis(vector<T>& nums) {
    int n = nums.size();
    vector <T> sub:
      vector < int > subIndex:
    vector <T> path(n, -1);
    for (int i = 0; i < n; ++i) {</pre>
        if (sub.empty() || sub[sub.size() - 1] <</pre>
            nums[i]) {
      path[i] = sub.empty() ? -1 : subIndex[sub.
           size() - 1];
      sub.push_back(nums[i]);
      subIndex.push back(i):
       } else {
      int idx = lower_bound(sub.begin(), sub.end(),
           nums[i]) - sub.begin():
      path[i] = idx == 0 ? -1 : subIndex[idx - 1];
      sub[idx] = nums[i]:
      subIndex[idx] = i;
17
    }
19
    vector <T> ans;
    int t = subIndex[subIndex.size() - 1];
    while (t != -1) {
        ans.push_back(t);
        t = path[t]:
24
    reverse(ans.begin(), ans.end());
    return ans:
29 // Length only
30 template < class T>
31 int length_lis(vector <T> &a) {
    set <T> st:
    typename set<T>::iterator it;
    for (int i = 0; i < a.size(); ++i) {</pre>
      it = st.lower bound(a[i]):
      if (it != st.end()) st.erase(it);
      st.insert(a[i]);
   }
```

```
39 return st.size();
40 }
```

## **3.2.2 0-1** Knapsack

```
1 #include "header.h"
2 // given a number of coins, calculate all
      possible distinct sums
3 int main() {
   vi coins(n); // all possible coins to use
   int sum = 0:
                    // sum of the coins
    vi dp(sum + 1, 0);
                              // dp[x] = 1 if sum
        x can be made
    dp[0] = 1;
                                // sum 0 can be
       made
   for (int c = 0; c < n; ++c)
       iteration: sums with first
      for (int x = sum; x \ge 0; --x)
                                          // coin,
         next first 2 coins etc
       if (dp[x]) dp[x + coins[c]] = 1; // if sum
            x valid, x+c valid
12 }
```

# **3.2.3 Coin change** Number of coins required to achieve a given value

```
1 #include "header.h"
2 // Returns total distinct ways to make sum using
      n coins of
3 // different denominations
4 int count(vi& coins, int n, int sum) {
      // 2d dp array where n is the number of coin
      // denominations and sum is the target sum
      vector < vector < int > > dp(n + 1, vector < int > (
          sum + 1, 0));
      dp[0][0] = 1:
      for (int i = 1: i <= n: i++) {</pre>
          for (int j = 0; j <= sum; j++) {</pre>
              // without using the current coin,
              dp[i][i] += dp[i - 1][i];
              // using the current coin
              if ((i - coins[i - 1]) >= 0)
                   dp[i][j] += dp[i][j - coins[i -
                      1]];
19
      return dp[n][sum];
```

#### 3.3 Trees

#### 3.3.1 Tree diameter

```
1 #include "header.h"
2 \text{ const int } mxN = 2e5 + 5;
3 int n, d[mxN]; // distance array
4 vi adj[mxN]; // tree adjacency list
5 void dfs(int s, int e) {
6 d[s] = 1 + d[e];
                      // recursively calculate
        the distance from the starting node to each
         node
for (auto u : adj[s]) { // for each adjacent
      if (u != e) dfs(u, s); // don't move
          backwards in the tree
11 int main() {
12 // read input, create adj list
    dfs(0, -1):
                                 // first dfs call
         to find farthest node from arbitrary node
    dfs(distance(d, max_element(d, d + n)), -1);
        // second dfs call to find farthest node
        from that one
   cout << *max element(d, d + n) - 1 << '\n': //
         distance from second node to farthest is
        the diameter
16 }
```

#### 3.3.2 Tree Node Count

## 3.4 Num. Th. / Comb.

#### 3.4.1 Basic stuff

```
1 #include "header.h"
```

```
2 ll gcd(ll a, ll b) { while (b) { a %= b; swap(a,
      b); } return a; }
3 11 1cm(11 a, 11 b) { return (a / gcd(a, b)) * b;
4 ll mod(ll a, ll b) { return ((a % b) + b) % b; }
_5 // Finds x, y s.t. ax + by = d = gcd(a, b).
6 void extended_euclid(ll a, ll b, ll &x, ll &y, ll
       &d) {
   11 xx = v = 0:
   11 yy = x = 1;
    while (b) {
      11 a = a / b:
      11 t = b; b = a % b; a = t;
      t = xx; xx = x - q * xx; x = t;
      t = yy; yy = y - q * yy; y = t;
17 // solves ab = 1 (mod n), -1 on failure
18 ll mod inverse(ll a. ll n) {
    11 x, y, d;
    extended_euclid(a, n, x, y, d);
    return (d > 1 ? -1 : mod(x, n));
23 // All modular inverses of [1..n] mod P in O(n)
24 vi inverses(ll n, ll P) {
25 vi I(n+1, 1LL):
  for (ll i = 2; i <= n; ++i)
      I[i] = mod(-(P/i) * I[P\%i], P);
   return I;
29 }
30 // (a*b)%m
31 ll mulmod(ll a, ll b, ll m){
32 11 x = 0, v=a\%m:
  while(b>0){
     if(b\&1) x = (x+y)\%m;
      y = (2*y)%m, b /= 2;
   return x % m;
_{39} // Finds b^e % m in O(lg n) time, ensure that b <
       m to avoid overflow!
40 ll powmod(ll b. ll e. ll m) {
11 p = e < 2 ? 1 : powmod((b*b)\%m, e/2, m);
   return e&1 ? p*b%m : p;
43 }
44 // Solve ax + by = c, returns false on failure.
45 bool linear_diophantine(ll a, ll b, ll c, ll &x,
      11 &v) {
11 d = gcd(a, b);
47 if (c % d) {
   return false;
  } else {
      x = c / d * mod_inverse(a / d, b / d);
```

```
51    y = (c - a * x) / b;
52    return true;
53    }
54 }
```

## **3.4.2** Mod. exponentiation Or use pow() in python

```
#include "header.h"
2 ll mod_pow(ll base, ll exp, ll mod) {
3   if (mod == 1) return 0;
4    if (exp == 0) return 1;
5    if (exp == 1) return base;
6
7   ll res = 1;
8   base %= mod;
9   while (exp) {
10    if (exp % 2 == 1) res = (res * base) % mod;
11   exp >>= 1;
12   base = (base * base) % mod;
13   }
14
15   return res % mod;
16 }
```

## **3.4.3** GCD Or math.gcd in python, std::gcd in C++

```
#include "header.h"
2 ll gcd(ll a, ll b) {
3    if (a == 0) return b;
4    return gcd(b % a, a);
5 }
```

#### 3.4.4 Sieve of Eratosthenes

### 3.4.5 Fibonacci % prime

## 3.4.6 nCk % prime

```
1 #include "header.h"
2 ll binom(ll n, ll k) {
      ll ans = 1:
      for(ll i = 1; i \le min(k,n-k); ++i) ans = ans
          *(n+1-i)/i:
      return ans;
7 ll mod_nCk(ll n, ll k, ll p ){
      ll ans = 1:
      while(n){
          11 np = n\%p, kp = k\%p;
          if(kp > np) return 0;
          ans *= binom(np,kp);
12
          n /= p; k /= p;
      return ans:
15
16 }
```

### 3.4.7 Chin. rem. th.

# 3.5 Strings

## **3.5.1 Z** alg. KMP alternative

```
#include "../header.h"

void Z_algorithm(const string &s, vi &Z) {
    Z.assign(s.length(), -1);
    int L = 0, R = 0, n = s.length();
    for (int i = 1; i < n; ++i) {
        if (i > R) {
            L = R = i;
            while (R < n && s[R - L] == s[R]) R++;
            Z[i] = R - L; R--;
        } else if (Z[i - L] >= R - i + 1) {
            L = i;
            while (R < n && s[R - L] == s[R]) R++;
            Z[i] = R - L; R--;
            Yelse if (Z[i - L] >= R - i + 1) {
            L = i;
            while (R < n && s[R - L] == s[R]) R++;
            Z[i] = R - L; R--;
            Yelse Z[i] = Z[i - L];
        }
}</pre>
```

#### 3.5.2 KMP

```
1 #include "header.h"
void compute_prefix_function(string &w, vi &
     prefix) {
prefix.assign(w.length(), 0);
   int k = prefix[0] = -1;
   for(int i = 1; i < w.length(); ++i) {</pre>
      if(w[k + 1] == w[i]) k++;
     prefix[i] = k:
  }
12 void knuth_morris_pratt(string &s, string &w) {
   vi prefix;
   compute_prefix_function(w, prefix);
   for(int i = 0; i < s.length(); ++i) {</pre>
     while (q \ge 0 \&\& w[q + 1] != s[i]) q = prefix[
         q];
```

# **3.5.3 Aho-Corasick** Also can be used as Knuth-Morris-Pratt algorithm

```
1 #include "header.h"
3 map<char, int> cti;
4 int cti_size;
5 template <int ALPHABET SIZE. int (*mp)(char)>
6 struct AC_FSM {
    struct Node {
      int child[ALPHABET_SIZE], failure = 0,
          match_par = -1;
      vi match:
      Node() { for (int i = 0; i < ALPHABET_SIZE;</pre>
          ++i) child[i] = -1: }
    vector < Node > a;
    vector < string > & words;
    AC FSM(vector<string> &words) : words(words) {
      a.push_back(Node());
      construct_automaton();
    }
17
    void construct_automaton() {
      for (int w = 0, n = 0; w < words.size(); ++w.
           n = 0) {
        for (int i = 0; i < words[w].size(); ++i) {</pre>
          if (a[n].child[mp(words[w][i])] == -1) {
21
             a[n].child[mp(words[w][i])] = a.size();
22
             a.push_back(Node());
          n = a[n].child[mp(words[w][i])];
26
        a[n].match.push_back(w);
27
28
29
      aueue < int > a:
      for (int k = 0; k < ALPHABET_SIZE; ++k) {</pre>
30
        if (a[0].child[k] == -1) a[0].child[k] = 0:
31
        else if (a[0].child[k] > 0) {
          a[a[0].child[k]].failure = 0;
          q.push(a[0].child[k]);
34
35
36
      while (!q.empty()) {
37
        int r = q.front(); q.pop();
        for (int k = 0, arck; k < ALPHABET_SIZE; ++</pre>
            k) {
```

```
if ((arck = a[r].child[k]) != -1) {
            q.push(arck);
42
            int v = a[r].failure;
            while (a[v].child[k] == -1) v = a[v].
43
                failure:
            a[arck].failure = a[v].child[k];
            a[arck].match_par = a[v].child[k];
            while (a[arck].match_par != -1
46
                 && a[a[arck].match par].match.emptv
              a[arck].match_par = a[a[arck].
                  match_par].match_par;
          }
        }
      }
    }
52
    void aho_corasick(string &sentence, vvi &
        matches){
      matches.assign(words.size(), vi());
      int state = 0. ss = 0:
      for (int i = 0; i < sentence.length(); ++i,</pre>
          ss = state) {
        while (a[ss].child[mp(sentence[i])] == -1)
          ss = a[ss].failure;
        state = a[state].child[mp(sentence[i])]
            = a[ss].child[mp(sentence[i])];
        for (ss = state; ss != -1; ss = a[ss].
            match par)
          for (int w : a[ss].match)
            matches[w].push back(i + 1 - words[w].
                length());
67 int char to int(char c) {
  return cti[c]:
69 }
70 int main() {
    string line;
    while(getline(cin, line)) {
      stringstream ss(line);
74
      ss >> n:
      vector < string > patterns(n);
      for (auto& p: patterns) getline(cin, p);
      string text;
      getline(cin, text);
      cti = {}. cti size = 0:
      for (auto c: text) {
        if (not in(c, cti)) {
          cti[c] = cti_size++;
```

## **3.5.4** Long. palin. subs Manacher - O(n)

```
1 #include "header.h"
void manacher(string &s. vi &pal) {
   int n = s.length(), i = 1, 1, r;
   pal.assign(2 * n + 1, 0);
   while (i < 2 * n + 1) {
     if ((i&1) && pal[i] == 0) pal[i] = 1;
     l = i / 2 - pal[i] / 2; r = (i-1) / 2 + pal[i]
         1 / 2:
     while (1 - 1 >= 0 \&\& r + 1 < n \&\& s[1 - 1] ==
          s[r + 1])
       --1, ++r, pal[i] += 2;
     for (1 = i - 1, r = i + 1; 1 >= 0 && r < 2 *
         n + 1: --1. ++r) {
       if (1 <= i - pal[i]) break;</pre>
       if (1 / 2 - pal[1] / 2 > i / 2 - pal[i] /
          pal[r] = pal[1];
        else { if (1 \ge 0)
            pal[r] = min(pal[1], i + pal[i] - r);
     i = r;
```

# 3.6 Geometry

## 3.6.1 essentials.cpp

```
1 #include "../header.h"
2 using C = ld; // could be long long or long
double
```

```
3 constexpr C EPS = 1e-10; // change to 0 for C=11
4 struct P { // may also be used as a 2D vector
    P(C x = 0, C v = 0) : x(x), v(v) {}
    P operator+ (const P &p) const { return {x + p.
        x, y + p.y; }
    P operator - (const P &p) const { return {x - p.
        x, y - p.y; }
    P operator* (C c) const { return {x * c, y * c}
    P operator/ (C c) const { return {x / c, y / c
    C operator* (const P &p) const { return x*p.x +
         v*p.v: }
    C operator^ (const P &p) const { return x*p.y -
         p.x*v; }
    P perp() const { return P{y, -x}; }
    C lensq() const { return x*x + y*y; }
    ld len() const { return sqrt((ld)lensq()); }
    static ld dist(const P &p1, const P &p2) {
      return (p1-p2).len(); }
    bool operator == (const P &r) const {
      return ((*this)-r).lensq() <= EPS*EPS; }</pre>
20 };
21 C det(P p1, P p2) { return p1^p2; }
22 C det(P p1, P p2, P o) { return det(p1-o, p2-o);
23 C det(const vector <P> &ps) {
    C sum = 0; P prev = ps.back();
    for(auto &p : ps) sum += det(p, prev), prev = p
    return sum;
_{28} // Careful with division by two and C=11
29 C area(P p1, P p2, P p3) { return abs(det(p1, p2,
       p3))/C(2); }
30 C area(const vector <P> &poly) { return abs(det(
      poly))/C(2); }
31 int sign(C c) { return (c > C(0)) - (c < C(0)); }
32 int ccw(P p1, P p2, P o) { return sign(det(p1, p2
_{34} // Only well defined for C = ld.
35 P unit(const P &p) { return p / p.len(); }
36 P rotate(P p, ld a) { return P{p.x*cos(a)-p.y*sin
      (a), p.x*sin(a)+p.v*cos(a)}; }
```

#### 3.6.2 Two segs. itersec.

```
1 #include "header.h"
2 #include "essentials.cpp"
3 bool intersect(P a1, P a2, P b1, P b2) {
4   if (max(a1.x, a2.x) < min(b1.x, b2.x)) return
      false;</pre>
```

#### 3.6.3 Convex Hull

```
1 #include "header.h"
2 #include "essentials.cpp"
3 struct ConvexHull { // O(n lg n) monotone chain.
    vector < size_t > h, c; // Indices of the hull
        are in 'h', ccw.
    const vector <P> &p;
    ConvexHull(const vector < P > & _ p) : n(_p.size()),
         c(n), p(_p) {
      std::iota(c.begin(), c.end(), 0);
      std::sort(c.begin(), c.end(), [this](size_t 1
          , size_t r) -> bool { return p[1].x != p[
          r].x ? p[1].x < p[r].x : p[1].y < p[r].y;
      c.erase(std::unique(c.begin(), c.end(), [this
          ](size_t 1, size_t r) { return p[1] == p[
          rl: }). c.end()):
      for (size_t s = 1, r = 0; r < 2; ++r, s = h.
          size()) {
        for (size_t i : c) {
12
          while (h.size() > s && ccw(p[h.end()
13
              [-2]], p[h.end()[-1]], p[i]) <= 0)
            h.pop_back();
          h.push_back(i);
16
        reverse(c.begin(), c.end());
17
19
      if (h.size() > 1) h.pop_back();
20
    size_t size() const { return h.size(); }
    template <class T, void U(const P &, const P &,
         const P &, T &)>
    void rotating_calipers(T &ans) {
      if (size() <= 2)
        U(p[h[0]], p[h.back()], p[h.back()], ans);
25
        for (size_t i = 0, j = 1, s = size(); i < 2</pre>
             * s: ++i) {
          while (det(p[h[(i + 1) % s]] - p[h[i % s
```

```
]], p[h[(j + 1) \% s]] - p[h[j]]) >=
              0)
            j = (j + 1) \% s;
          U(p[h[i % s]], p[h[(i + 1) % s]], p[h[j
              ]], ans);
        }
   }
33 }:
34 // Example: furthest pair of points. Now set ans
      = OLL and call
35 // ConvexHull(pts).rotating_calipers<11, update>(
36 void update(const P &p1, const P &p2, const P &o,
       ll &ans) {
   ans = max(ans, (11)max((p1 - o).lensq(), (p2 -
        o).lensq()));
38 }
39 int main() {
    ios::sync_with_stdio(false); // do not use
        cout + printf
    cin.tie(NULL);
    int n;
    cin >> n;
    while (n) {
      vector <P> ps;
          int x, y;
      for (int i = 0: i < n: i++) {</pre>
               cin >> x >> y;
               ps.push_back({x, y});
          ConvexHull ch(ps):
          cout << ch.h.size() << endl;</pre>
          for(auto& p: ch.h) {
               cout << ps[p].x << "" << ps[p].y <<
                   endl;
      cin >> n;
    }
    return 0;
```

## 3.7 Other Algorithms

#### 3.7.1 2-sat

```
#include "../header.h"

#include "../Graphs/tarjan.cpp"

struct TwoSAT {

int n;

vvi imp; // implication graph

Tarjan ti:
```

```
TwoSAT(int _n) : n(_n), imp(2 * _n, vi()), tj(
        imp) { }
    // Only copy the needed functions:
    void add_implies(int c1, bool v1, int c2, bool
      int u = 2 * c1 + (v1 ? 1 : 0),
       v = 2 * c2 + (v2 ? 1 : 0):
      imp[u].push_back(v);  // u => v
      imp[v^1].push_back(u^1); // -v => -u
16
    void add_equivalence(int c1, bool v1, int c2,
17
        bool v2) {
      add_implies(c1, v1, c2, v2);
      add_implies(c2, v2, c1, v1);
19
20
    void add_or(int c1, bool v1, int c2, bool v2) {
      add_implies(c1, !v1, c2, v2);
23
    void add_and(int c1, bool v1, int c2, bool v2)
      add_true(c1, v1); add_true(c2, v2);
    void add_xor(int c1, bool v1, int c2, bool v2)
      add_or(c1, v1, c2, v2);
      add or(c1, !v1, c2, !v2);
    void add_true(int c1, bool v1) {
      add_implies(c1, !v1, c1, v1);
    }
    // on true: a contains an assignment.
    // on false: no assignment exists.
    bool solve(vb &a) {
      vi com;
38
      ti.find sccs(com):
      for (int i = 0; i < n; ++i)</pre>
        if (com[2 * i] == com[2 * i + 1])
          return false;
43
      vvi bvcom(com.size()):
44
      for (int i = 0: i < 2 * n: ++i)
        bycom[com[i]].push_back(i);
47
      a.assign(n, false);
48
      vb vis(n, false);
49
      for(auto &&component : bycom){
        for (int u : component) {
51
          if (vis[u / 2]) continue;
          vis[u / 2] = true:
53
          a[u / 2] = (u \% 2 == 1);
        }
```

```
57 return true;
58 }
59 };
```

#### 3.7.2 Matrix Solve

1 #include "header.h"

```
2 #define REP(i, n) for(auto i = decltype(n)(0); i
      <(n); i++)
3 using T = double;
4 constexpr T EPS = 1e-8;
5 template < int R, int C>
6 using M = array<array<T,C>,R>; // matrix
7 template < int R. int C>
8 T ReducedRowEchelonForm(M<R,C> &m, int rows) {
      // return the determinant
    int r = 0; T det = 1;
                                       // MODIFIES
        the input
    for(int c = 0: c < rows && r < rows: c++) {
      int p = r;
      for(int i=r+1: i<rows: i++) if(abs(m[i][c]) >
           abs(m[p][c])) p=i;
      if(abs(m[p][c]) < EPS){ det = 0; continue; }</pre>
      swap(m[p], m[r]); det = -det;
      T s = 1.0 / m[r][c], t; det *= m[r][c];
      REP(j,C) m[r][j] *= s;  // make leading
           term in row 1
      REP(i,rows) if (i!=r){ t = m[i][c]; REP(j,C)
          m[i][j] -= t*m[r][j]; }
      ++r:
18
    return det:
21 }
22 bool error, inconst; // error => multiple or
      inconsistent
23 template <int R, int C> // Mx = a; M:R*R, v:R*C =>
24 M<R,C> solve(const M<R,R> &m, const M<R,C> &a,
      int rows){
   M < R.R+C > a:
    REP(r,rows){
      REP(c,rows) q[r][c] = m[r][c];
      REP(c,C) q[r][R+c] = a[r][c];
28
29
    ReducedRowEchelonForm <R.R+C>(a.rows):
    M<R,C> sol; error = false, inconst = false;
    REP(c,C) for(auto j = rows-1; j \ge 0; --j){
     T t=0; bool allzero=true;
      for (auto k = j+1; k < rows; ++k)
        t += q[j][k]*sol[k][c], allzero &= abs(q[j])
            1[k]) < EPS:</pre>
      if(abs(q[j][j]) < EPS)
        error = true, inconst |= allzero && abs(q[j
37
            l(R+cl) > EPS:
```

## 3.7.3 Matrix Exp.

```
1 #include "header.h"
2 #define ITERATE_MATRIX(w) for (int r = 0; r < (w)</pre>
      : ++r) \
                for (int c = 0; c < (w); ++c)
4 template <class T. int N>
5 struct M {
    array <array <T,N>,N> m;
    M() { ITERATE MATRIX(N) m[r][c] = 0; }
    static M id() {
      M I; for (int i = 0; i < N; ++i) I.m[i][i] =
          1: return I:
   M operator*(const M &rhs) const {
      ITERATE_MATRIX(N) for (int i = 0; i < N; ++i)</pre>
          out.m[r][c] += m[r][i] * rhs.m[i][c]:
      return out;
15
   M raise(ll n) const {
      if(n == 0) return id();
      if(n == 1) return *this;
      auto r = (*this**this).raise(n / 2);
      return (n%2 ? *this*r : r);
22 }
23 };
```

#### 3.7.4 Finite field For FFT

```
1 #include "header.h"
2 #include "../Number LTheory/elementary.cpp"
3 template<11 p,11 w> // prime, primitive root
4 struct Field { using T = Field; ll x; Field(ll x
     =0) : x\{x\} \{\}
5 T operator+(T r) const { return {(x+r.x)%p}; }
   T operator - (T r) const { return \{(x-r.x+p)\%p\};
   T operator*(T r) const { return {(x*r.x)%p}; }
   T operator/(T r) const { return (*this)*r.inv()
   T inv() const { return {mod_inverse(x,p)}; }
   static T root(ll k) { assert( (p-1)%k==0 );
       // (p-1)%k == 0?
      auto r = powmod(w,(p-1)/abs(k),p);
                                               // k-
          th root of unity
     return k>=0 ? T{r} : T{r}.inv():
```

```
13  }
14  bool zero() const { return x == 0LL; }
15 };
16  using F1 = Field<1004535809,3 >;
17  using F2 = Field<1107296257,10>; // 1<<30 + 1<<25 + 1
18  using F3 = Field<2281701377,3 >; // 1<<31 + 1<<27 + 1</pre>
```

## 3.7.5 Complex field For FFR

```
1 #include "header.h"
2 const double m_pi = M_PIf64x;
3 struct Complex { using T = Complex: double u.v:
    Complex (double u=0, double v=0) : u\{u\}, v\{v\} {}
5 T operator+(T r) const { return {u+r.u, v+r.v};
   T operator-(T r) const { return {u-r.u, v-r.v};
    T operator*(T r) const { return {u*r.u - v*r.v,
         u*r.v + v*r.u}: }
    T operator/(T r) const {
      auto norm = r.u*r.u+r.v*r.v;
      return {(u*r.u + v*r.v)/norm, (v*r.u - u*r.v)
          /norm}:
11
    T operator*(double r) const { return T{u*r, v*r}
    T operator/(double r) const { return T{u/r, v/r
    T inv() const { return T{1,0}/ *this; }
    T conj() const { return T{u, -v}; }
    static T root(ll k){ return {cos(2*m_pi/k), sin
        (2*m pi/k)}: }
    bool zero() const { return max(abs(u), abs(v))
        < 1e-6; }
18 };
```

#### 3.7.6 FFT

```
#include "header.h"

#include "complex_field.cpp"

#include "fin_field.cpp"

void brinc(int &x, int k) {

int i = k - 1, s = 1 << i;

x ^= s;

if ((x & s) != s) {

--i; s >>= 1;

while (i >= 0 && ((x & s) == s))

x = x & s, --i, s >>= 1;

if (i >= 0) x |= s;

}

13 }
```

```
using T = Complex; // using T=F1,F2,F3
15 vector <T> roots;
16 void root cache(int N) {
  if (N == (int)roots.size()) return:
    roots.assign(N, T{0});
    for (int i = 0; i < N; ++i)</pre>
      roots[i] = ((i\&-i) == i)
        ? T\{\cos(2.0*m_pi*i/N), \sin(2.0*m_pi*i/N)\}
        : roots[i&-i] * roots[i-(i&-i)]:
24 void fft(vector<T> &A. int p. bool inv = false) {
25 int N = 1<<p:
   for(int i = 0, r = 0; i < N; ++i, brinc(r, p))
    if (i < r) swap(A[i], A[r]);</pre>
28 // Uncomment to precompute roots (for T=Complex)
      . Slower but more precise.
29 // root_cache(N);
           , sh=p-1
31 for (int m = 2; m <= N; m <<= 1) {
      T w. w m = T::root(inv ? -m : m):
      for (int k = 0; k < N; k += m) {
        w = T\{1\}:
        for (int j = 0; j < m/2; ++j) {
            T w = (!inv ? roots[j << sh] : roots[j <<
      shl.coni()):
         T t = w * A[k + j + m/2];
37
          A[k + j + m/2] = A[k + j] - t;
          A[k + j] = A[k + j] + t;
          w = w * w_m;
41
      }
42
43
    if(inv){ T inverse = T(N).inv(); for(auto &x :
        A) x = x*inverse;
45 }
_{
m 46} // convolution leaves A and B in frequency domain
47 // C may be equal to A or B for in-place
      convolution
48 void convolution(vector <T> &A, vector <T> &B,
      vector <T> &C) {
    int s = A.size() + B.size() - 1;
    int q = 32 - __builtin_clz(s-1), N=1<<q; //</pre>
        fails if s=1
    A.resize(N,{}); B.resize(N,{}); C.resize(N,{});
    fft(A, q, false); fft(B, q, false);
   for (int i = 0; i < N; ++i) C[i] = A[i] * B[i];</pre>
   fft(C, q, true); C.resize(s);
56 void square_inplace(vector <T> &A) {
   int s = 2*A.size()-1, q = 32 - _builtin_clz(s)
        -1). N=1<<a:
    A.resize(N,{}); fft(A, q, false);
    for (auto &x : A) x = x*x;
    fft(A, q, true); A.resize(s);
```

#### 3.7.7 Polyn. inv. div.

61 }

```
1 #include "header.h"
2 #include "fft.cpp"
3 vector <T> &rev(vector <T> &A) { reverse(A.begin(),
       A.end()); return A; }
4 void copy_into(const vector <T> &A, vector <T> &B,
      size_t n) {
std::copy(A.begin(), A.begin()+min({n, A.size()
        , B.size()}), B.begin());
6 }
8 // Multiplicative inverse of A modulo x^n.
      Requires A[0] != 0!!
9 vector <T> inverse(const vector <T> &A. int n) {
    vector<T> Ai{A[0].inv()};
    for (int k = 0; (1<<k) < n; ++k) {
      vector <T> As (4 << k, T(0)), Ais (4 << k, T(0));
      copy_into(A, As, 2<<k); copy_into(Ai, Ais, Ai
           .size()):
      fft(As, k+2, false); fft(Ais, k+2, false);
      for (int i = 0; i < (4<<k); ++i) As[i] = As[i
          ] * A is [i] * A is [i];
      fft(As, k+2, true); Ai.resize(2<<k, {});
      for (int i = 0; i < (2 << k); ++i) Ai[i] = T(2)
           * Ai[i] - As[i]:
    Ai.resize(n);
    return Ai:
22 // Polynomial division. Returns {Q, R} such that
      A = QB+R, deg R < deg B.
23 // Requires that the leading term of B is nonzero
24 pair < vector < T >, vector < T >> divmod(const vector < T >
       &A, const vector <T> &B) {
    size t n = A.size()-1, m = B.size()-1:
    if (n < m) return {vector < T > (1, T(0)), A};
    vector\langle T \rangle X(A), Y(B), Q, R;
    convolution(rev(X), Y = inverse(rev(Y), n-m+1),
    Q.resize(n-m+1); rev(Q);
    X.resize(Q.size()), copy_into(Q, X, Q.size());
    Y.resize(B.size()), copy_into(B, Y, B.size());
    convolution(X, Y, X);
    R.resize(m), copy_into(A, R, m);
    for (size_t i = 0; i < m; ++i) R[i] = R[i] - X[</pre>
    while (R.size() > 1 && R.back().zero()) R.
        pop_back();
```

**3.7.8** Linear recurs. Given a linear recurrence of the form

$$a_n = \sum_{i=0}^{k-1} c_i a_{n-i-1}$$

this code computes  $a_n$  in  $O(k \log k \log n)$  time.

1 #include "header.h"

```
2 #include "poly.cpp"
3 // x^k \mod f
4 vector <T> xmod(const vector <T> f, ll k) {
    vector <T> r{T(1)}:
    for (int b = 62; b >= 0; --b) {
      if (r.size() > 1)
         square_inplace(r), r = mod(r, f);
      if ((k>>b)&1) {
        r.insert(r.begin(), T(0));
        if (r.size() == f.size()) {
11
          T c = r.back() / f.back():
          for (size_t i = 0; i < f.size(); ++i)</pre>
            r[i] = r[i] - c * f[i];
          r.pop_back();
    return r;
19
20 }
_{21} // Given A[0,k) and C[0, k), computes the n-th
      term of:
22 // A[n] = \sum i C[i] * A[n-i-1]
23 T nth_term(const vector<T> &A, const vector<T> &C
      , ll n) {
    int k = (int)A.size();
    if (n < k) return A[n];</pre>
    vector \langle T \rangle f(k+1, T{1}):
    for (int i = 0; i < k; ++i)
     f[i] = T\{-1\} * C[k-i-1];
    f = xmod(f, n);
31
    T r = T{0}:
    for (int i = 0; i < k; ++i)</pre>
     r = r + f[i] * A[i]:
35
    return r:
```

#### **3.7.9 Convolution** Precise up to 9e15

```
1 #include "header.h"
2 #include "fft.cpp"
3 void convolution_mod(const vi &A, const vi &B, 11
   int s = A.size() + B.size() - 1; ll m15 = (1LL
        <<15) -1LL;
   int q = 32 - __builtin_clz(s-1), N=1<<q; //</pre>
        fails if s=1
    vector\langle T \rangle Ac(N), Bc(N), R1(N), R2(N);
   for (size_t i = 0; i < A.size(); ++i) Ac[i] = T</pre>
         {A[i]\&m15, A[i]>>15};
    for (size_t i = 0; i < B.size(); ++i) Bc[i] = T</pre>
        {B[i]&m15, B[i]>>15};
    fft(Ac, q, false); fft(Bc, q, false);
    for (int i = 0, j = 0; i < N; ++i, j = (N-1)&(N
      T as = (Ac[i] + Ac[j].conj()) / 2;
      T al = (Ac[i] - Ac[j].conj()) / T{0, 2};
      T bs = (Bc[i] + Bc[j].conj()) / 2;
      T bl = (Bc[i] - Bc[j].conj()) / T{0, 2};
      R1[i] = as*bs + al*bl*T{0,1}, R2[i] = as*bl +
    fft(R1, q, true); fft(R2, q, true);
    11 p15 = (1LL << 15) \% MOD, p30 = (1LL << 30) \% MOD; C.
        resize(s):
    for (int i = 0; i < s; ++i) {</pre>
      11 1 = 1 \text{lround}(R1[i].u), m = 1 \text{lround}(R2[i].u)
          , h = llround(R1[i].v);
      C[i] = (1 + m*p15 + h*p30) \% MOD;
   }
22
23 }
```

# **3.7.10** Partitions of n Finds all possible partitions of a number

```
#include "header.h"
void printArray(int p[], int n) {
for (int i = 0; i < n; i++)
cout << p[i] << "u";
cout << endl;
}

void printAllUniqueParts(int n) {
int p[n]; // An array to store a partition
int k = 0; // Index of last element in a partition
p[k] = n; // Initialize first partition as number itself

// This loop first prints current partition then generates next</pre>
```

```
// partition. The loop stops when the current
        partition has all 1s
    while (true) {
      printArrav(p, k + 1);
      // Find the rightmost non-one value in p[].
          Also, update the
      // rem_val so that we know how much value can
           be accommodated
      int rem_val = 0;
      while (k >= 0 \&\& p[k] == 1) {
        rem val += p[k]:
      }
      // if k < 0, all the values are 1 so there
          are no more partitions
      if (k < 0) return;</pre>
      // Decrease the p[k] found above and adjust
          the rem_val
      p[k]--:
      rem_val++;
      // If rem_val is more, then the sorted order
          is violated. Divide
      // rem_val in different values of size p[k]
          and copy these values at
      // different positions after p[k]
      while (rem_val > p[k]) {
       p[k + 1] = p[k];
       rem_val = rem_val - p[k];
      }
40
      // Copy rem_val to next position and
          increment position
      p[k + 1] = rem val:
```

## 3.8 Other Data Structures

#### **3.8.1** Disjoint set (i.e. union-find)

```
this -> n = n:
              this->rank = new T[n];
              for (T i = 0: i < n: i++) {
                  parent[i] = i;
                  rank[i] = 0;
              }
          }
          // O(log n)
          T find set(T x) {
19
              if (x == parent[x]) return x;
              return parent[x] = find_set(parent[x
                  1):
          }
23
          // O(\log n)
24
          void union_sets(T x, T y) {
25
              x = this->find_set(x);
              y = this->find_set(y);
              if (x == v) return;
              if (rank[x] < rank[y]) {</pre>
                  Tz = x:
                  x = y;
                  y = z;
              }
              parent[y] = x;
              if (rank[x] == rank[y]) rank[x]++;
40 };
```

# **3.8.2 Fenwick tree** (i.e. BIT) eff. update + prefix sum calc.

```
#include "header.h"
#define maxn 200010
int t,n,m,tree[maxn],p[maxn];

void update(int k, int z) {
    while (k <= maxn) {
        tree[k] += z;
        k += k & (-k);
    }
}

int sum(int k) {
    int ans = 0;
    while(k) {
    ans += tree[k];
    k -= k & (-k);
}</pre>
```

## 3.8.3 Fenwick2d tree

```
1 #include "header.h"
2 template <class T>
3 struct FenwickTree2D {
    vector < vector <T> > tree;
    FenwickTree2D(int n) : n(n) { tree.assign(n +
        1, vectorT>(n + 1, 0); }
   T query(int x1, int y1, int x2, int y2) {
      return query (x2, y2) + query (x1-1, y1-1) - query (x2
          ,y1-1)-query(x1-1,y2);
   T query(int x, int y) {
     T s = 0:
      for (int i = x: i > 0: i -= (i & (-i)))
        for (int j = v; j > 0; j = (j & (-j)))
14
          s += tree[i][i]:
      return s;
15
   }
    void update(int x, int y, T v) {
      for (int i = x; i <= n; i += (i & (-i)))
        for (int j = y; j <= n; j += (j & (-j)))
          tree[i][i] += v:
  }
21
22 };
```

#### 3.8.4 Trie

```
1 #include "header.h"
2 const int ALPHABET SIZE = 26:
3 inline int mp(char c) { return c - 'a'; }
5 struct Node {
    Node* ch[ALPHABET_SIZE];
    bool isleaf = false:
    Node() {
      for(int i = 0; i < ALPHABET_SIZE; ++i) ch[i]</pre>
          = nullptr:
10
11
    void insert(string &s, int i = 0) {
      if (i == s.length()) isleaf = true;
      else {
        int v = mp(s[i]);
        if (ch[v] == nullptr)
16
          ch[v] = new Node():
        ch[v]->insert(s, i + 1);
      }
19
   }
```

```
bool contains(string &s, int i = 0) {
      if (i == s.length()) return isleaf:
23
      else {
        int v = mp(s[i]);
        if (ch[v] == nullptr) return false;
        else return ch[v]->contains(s, i + 1);
28
    }
29
    void cleanup() {
31
      for (int i = 0: i < ALPHABET SIZE: ++i)</pre>
        if (ch[i] != nullptr) {
           ch[i]->cleanup();
           delete ch[i];
   }
38 };
```

**3.8.5** Treap A binary tree whose nodes contain two values, a key and a priority, such that the key keeps the BST property

```
1 #include "header.h"
2 struct Node {
  11 v;
   int sz. pr:
    Node *1 = nullptr, *r = nullptr;
    Node(ll val) : v(val), sz(1) { pr = rand(); }
7 };
8 int size(Node *p) { return p ? p->sz : 0; }
9 void update(Node* p) {
    if (!p) return;
    p->sz = 1 + size(p->1) + size(p->r);
    // Pull data from children here
14 void propagate(Node *p) {
  if (!p) return:
   // Push data to children here
17 }
18 void merge(Node *&t, Node *1, Node *r) {
    propagate(1), propagate(r);
   if (!1) t = r:
    else if (!r) t = 1;
    else if (1->pr > r->pr)
        merge(1->r, 1->r, r), t = 1;
    else merge(r->1, 1, r->1), t = r;
    update(t):
25
27 void spliti(Node *t, Node *&l, Node *&r, int
      index) {
    propagate(t);
  if (!t) { l = r = nullptr; return; }
  int id = size(t->1):
```

```
if (index <= id) // id \in [index, \infty), so</pre>
          move it right
       spliti(t\rightarrow 1, 1, t\rightarrow 1, index), r = t:
       spliti(t\rightarrow r, t\rightarrow r, r, index - id), l = t;
     update(t);
36 }
37 void splitv(Node *t, Node *&l, Node *&r, 11 val)
     propagate(t);
    if (!t) { l = r = nullptr: return: }
     if (val \langle = t - \rangle v) // t - \rangle v \in [val, \infty), so
         move it right
       splitv(t\rightarrow 1, 1, t\rightarrow 1, val), r = t;
       splitv(t->r, t->r, r, val), l = t;
    update(t);
46 void clean(Node *p) {
    if (p) { clean(p->1), clean(p->r); delete p; }
48 }
```

## 3.8.6 Segment tree

```
1 #include "../header.h"
2 template <class T, const T&(*op)(const T&, const</pre>
      T&)>
3 struct SegmentTree {
int n; vector<T> tree; T id;
    SegmentTree(int _n, T _id) : n(_n), tree(2 * n,
         _id), id(_id) { }
    void update(int i, T val) {
      for (tree[i+n] = val, i = (i+n)/2; i > 0; i
        tree[i] = op(tree[2*i], tree[2*i+1]);
    T query(int 1, int r) {
11
      T lhs = T(id), rhs = T(id);
      for (1 += n, r += n; 1 < r; 1 >>= 1, r >>= 1)
        if ( 1&1 ) lhs = op(lhs, tree[1++]);
        if (!(r\&1)) rhs = op(tree[r--], rhs);
14
15
      return op(1 == r ? op(lhs, tree[1]) : lhs,
          rhs):
17 }
18 };
```

## 3.8.7 Lazy segment tree Uptimizes range updates

```
3 T t id: U u id:
4 T op(T a, T b) { return a+b; }
5 void join(U &a, U b){ a+=b; }
6 void apply(T &t. U u. int x) { t+=x*u: }
7 T convert(const I &i){ return i; }
8 struct LazySegmentTree {
    struct Node { int 1, r, 1c, rc; T t; U u;
      Node(int 1, int r, T t=t_id):1(1),r(r),1c(-1)
          .rc(-1).t(t).u(u id){}
    };
12
   int N; vector < Node > tree; vector < I > &init;
    LazySegmentTree(vector < I > &init) : N(init.size
        ()), init(init){
      tree.reserve(2*N-1); tree.push_back({0,N});
          build(0, 0, N);
    void build(int i, int l, int r) { auto &n =
        tree[i]:
      if (r > 1+1) \{ int m = (1+r)/2;
        n.lc = tree.size(): n.rc = n.lc+1:
        10
            .r}):
        build(n.lc,1,m);
                              build(n.rc,m,r);
        n.t = op(tree[n.lc].t, tree[n.rc].t);
      } else n.t = convert(init[1]);
22
23
    void push(Node &n, U u) { apply(n.t, u, n.r-n.l)
        : ioin(n.u.u): }
    void push(Node &n){push(tree[n.lc],n.u);push(
        tree[n.rc].n.u):n.u=u id:}
   T query(int 1, int r, int i = 0) { auto &n =
        tree[i];
      if(r <= n.1 || n.r <= 1) return t id:
      if(1 <= n.1 && n.r <= r) return n.t;</pre>
      return push(n), op(querv(l.r.n.lc),querv(l.r.
          n.rc)):
   }
30
    void update(int 1, int r, U u, int i = 0) {
        auto &n = tree[i];
      if(r <= n.1 || n.r <= 1) return;</pre>
      if(1 <= n.1 && n.r <= r) return push(n,u);</pre>
      push(n); update(1,r,u,n.lc); update(1,r,u,n.
34
      n.t = op(tree[n.lc].t, tree[n.rc].t):
36
37 }:
```

#### 3.8.8 Suffix tree

```
5 using It = V::const_iterator;
6 struct Node{
    It b, e; M edges; int link; // end is
        exclusive
    Node(It b, It e) : b(b), e(e), link(-1) {}
    int size() const { return e-b; }
10 }:
11 struct SuffixTree{
    const V &s: vector < Node > t:
    int root,n,len,remainder,llink; It edge;
    SuffixTree(const V &s) : s(s) { build(); }
    int add_node(It b, It e){ return t.push_back({b
        ,e}), t.size()-1; }
    int add_node(It b){ return add_node(b,s.end());
    void link(int node){ if(llink) t[llink].link =
        node; llink = node; }
    void build(){
      len = remainder = 0; edge = s.begin();
      n = root = add_node(s.begin(), s.begin());
      for(auto i = s.begin(); i != s.end(); ++i){
        ++remainder: llink = 0:
        while(remainder){
          if(len == 0) edge = i;
          if(t[n].edges[*edge] == 0){
              new leaf
            t[n].edges[*edge] = add_node(i); link(n
                ):
          } else {
            auto x = t[n].edges[*edge];
                node [with edge]
            if(len >= t[x].size()){
                                        // walk to
                next node
              len -= t[x].size(); edge += t[x].size
                  (): n = x:
              continue:
            }
            if(*(t[x].b + len) == *i){
                along edge
              ++len; link(n); break;
                  // split edge
            auto split = add_node(t[x].b, t[x].b+
                len):
            t[n].edges[*edge] = split;
            t[x].b += len;
            t[split].edges[*i] = add_node(i);
            t[split].edges[*t[x].b] = x;
            link(split);
          --remainder;
          if(n == root && len > 0)
            --len, edge = i - remainder + 1:
          else n = t[n].link > 0 ? t[n].link : root
        }
```

```
48 }
49 }
50 };
```

## 4 Other Mathematics

# 4.1 Helpful functions

**4.1.1** Euler's Totient Fucntion  $n = p_1^{k_1-1} \cdot (p_1-1) \cdot \dots \cdot p_r^{k_r-1} \cdot (p_r-1)$ , where  $p_1^{k_1} \cdot \dots \cdot p_r^{k_r}$  is the prime factorization of n.

```
1 # include "header.h"
2 11 phi(11 n) { // \Phi(n)
      ll ans = 1;
      for (11 i = 2; i*i <= n; i++) {</pre>
          if (n % i == 0) {
              ans *= i-1;
              n /= i;
               while (n % i == 0) {
                   ans *= i:
                   n /= i;
          }
13
      if (n > 1) ans *= n-1;
14
15
      return ans:
16 }
     phis(int n) { // All \Phi(i) up to n
    vi phi(n + 1, OLL);
    iota(phi.begin(), phi.end(), OLL);
    for (11 i = 2LL; i <= n; ++i)</pre>
      if (phi[i] == i)
        for (11 j = i; j <= n; j += i)
          phi[j] -= phi[j] / i;
    return phi;
24
25 }
```

Formulas  $\Phi(n)$  counts all numbers in  $1, \ldots, n-1$  coprime to n.  $a^{\varphi(n)} \equiv 1 \mod n$ , a and n are coprimes.  $\forall e > \log_2 m : n^e \mod m = n^{\Phi(m)+e \mod \Phi(m)} \mod m$ .  $\gcd(m,n) = 1 \Rightarrow \Phi(m \cdot n) = \Phi(m) \cdot \Phi(n)$ .

**4.1.2** Pascal's trinagle  $\binom{n}{k}$  is k-th element in the n-th row, indexing both from 0

## 4.2 Theorems and definitions

#### Fermat's little theorem

$$a^p \equiv a \mod p$$

Subfactorial

$$!n = n! \sum_{i=0}^{n} \frac{(-1)^{i}}{i!}$$

$$!(0) = 1, !n = n \cdot !(n-1) + (-1)^n$$

## Binomials and other partitionings

$$\binom{n}{k} = \binom{n-1}{k} + \binom{n-1}{k-1} = \prod_{i=1}^{k} \frac{n-i+1}{i}$$

This last product may be computed incrementally since any product of k' consecutive values is divisible by k'!. Basic identities: The hockeystick identity:

$$\sum_{k=r}^{n} \binom{k}{r} = \binom{n+1}{r+1}$$

or

$$\sum_{k \le n} \binom{r+k}{k} = \binom{r+n+1}{n}$$

Also

$$\sum_{k=0}^{n} \binom{k}{m} = \binom{n+1}{m+1}$$

$$\sum_{i=0}^{n} \binom{n}{i} = 2^n$$

For  $n, m \geq 0$  and p prime: write n, m in base p, i.e.  $n = n_k p^k + \cdots + n_1 p + n_0$  and  $m = m_k p^k + \cdots + m_1 p + m_0$ . Then by Lucas theorem we have  $\binom{n}{m} \equiv \prod_{i=0}^k \binom{n_i}{m_i} \mod p$ , with the convention that  $n_i < m_i \implies \binom{n_i}{m_i} = 0$ .

Fibonacci (See also number theory section)

$$\sum_{0 \le k \le n} \binom{n-k}{k} = F_{n+1}$$

$$F_n = \frac{1}{\sqrt{5}} \left( \frac{1 + \sqrt{5}}{2} \right)^n - \frac{1}{\sqrt{5}} \left( \frac{1 - \sqrt{5}}{2} \right)^n$$

$$\sum_{i=1}^n F_i = F_{n+2} - 1, \ \sum_{i=1}^n F_i^2 = F_n F_{n+1}$$

$$\gcd(F_m, F_n) = F_{\gcd(m,n)}$$

$$\gcd(F_n, F_{n+1}) = \gcd(F_n, F_{n+2}) = 1$$

Bit stuff  $a+b=a\oplus b+2(a\&b)=a|b+a\&b$ . kth bit is set in x iff  $x \mod 2^{k-1} \geq 2^k$ , or iff  $x \mod 2^{k-1}-x \mod 2^k \neq 0$  (i.e.  $=2^k$ ) It comes handy when you need to look at the bits of the numbers which are pair sums or subset sums etc.

$$n \mod 2^i = n\&(2^i - 1).$$

$$\forall k: 1 \oplus 2 \oplus \ldots \oplus (4k-1) = 0$$

Stirling's numbers First kind:  $S_1(n,k)$  count permutations on n items with k cycles.  $S_1(n,k) = S_1(n-1,k-1) + (n-1)S_1(n-1,k)$  with  $S_1(0,0) = 1$ . Note:

$$\sum_{k=0}^{n} S_1(n,k)x^k = x(x+1)\dots(x+n-1)$$

$$\sum_{k=0}^{n} S_1(n,k) = n!$$

**Second kind:**  $S_2(n, k)$  count partitions of n distinct elements into exactly k non-empty groups.

$$S_2(n,k) = S_2(n-1,k-1) + kS_2(n-1,k)$$
$$S_2(n,1) = S_2(n,n) = 1$$

$$S_2(n,k) = \frac{1}{k!} \sum_{i=1}^{k} (-1)^{k-i} \binom{k}{i} i^n$$

# 4.3 Geometry Formulas

$$[ABC] = rs = \frac{1}{2}ab\sin\gamma$$

$$= \frac{abc}{4R} = \sqrt{s(s-a)(s-b)(s-c)} = \frac{1}{2} \left| (B-A, C-A)^T \right|$$

$$s = \frac{a+b+c}{2} \qquad 2R = \frac{a}{\sin \alpha}$$
 cosine rule: 
$$c^2 = a^2 + b^2 - 2ab\cos \gamma$$
 Euler: 
$$1 + CC = V - E + F$$
 Pick: 
$$\operatorname{Area} = \operatorname{itr} \operatorname{pts} + \frac{\operatorname{bdry} \operatorname{pts}}{2} - 1$$
 
$$p \cdot q = |p||q|\cos(\theta) \qquad |p \times q| = |p||q|\sin(\theta)$$

Given a non-self-intersecting closed polygon on n vertices, given as  $(x_i, y_i)$ , its centroid  $(C_x, C_y)$  is given as:

$$C_x = \frac{1}{6A} \sum_{i=0}^{n-1} (x_i + x_{i+1})(x_i y_{i+1} - x_{i+1} y_i),$$

$$C_y = \frac{1}{6A} \sum_{i=0}^{n-1} (y_i + y_{i+1})(x_i y_{i+1} - x_{i+1} y_i)$$

$$A = \frac{1}{2} \sum_{i=0}^{n-1} (x_i y_{i+1} - x_{i+1} y_i) = \text{polygon area}$$

**Inclusion-Exclusion** For appropriate f compute  $\sum_{S\subseteq T} (-1)^{|T\setminus S|} f(S)$ , or if only the size of S matters,  $\sum_{s=0}^{n} (-1)^{n-s} {n \choose s} f(s)$ . In some contexts we might use Stirling numbers, not binomial coefficients!

Some useful applications:

**Graph coloring** Let I(S) count the number of independent sets contained in  $S \subseteq V$  ( $I(\emptyset) = 1$ ,  $I(S) = I(S \setminus v) + I(S \setminus N(v))$ ). Let  $c_k = \sum_{S \subseteq V} (-1)^{|V \setminus S|} I(S)$ . Then V is k-colorable iff v > 0. Thus we can compute the chromatic number of a graph in  $O^*(2^n)$  time.

**Burnside's lemma** Given a group G acting on a set X, the number of elements in X up to symmetry is

$$\frac{1}{|G|} \sum_{g \in G} |X^g|$$

with  $X^g$  the elements of X invariant under g. For example, if f(n) counts "configurations" of some sort of length n, and we want to count them up to rotational symmetry using  $G = \mathbb{Z}/n\mathbb{Z}$ , then

$$g(n) = \frac{1}{n} \sum_{k=0}^{n-1} f(\gcd(n,k)) = \frac{1}{n} \sum_{k||n} f(k)\phi(n/k)$$

I.e. for coloring with c colors we have  $f(k) = k^c$ .

Relatedly, in Pólya's enumeration theorem we imagine X as a set of n beads with G permuting the beads (e.g. a necklace, with G all rotations and reflections of the n-cycle, i.e. the dihedral group  $D_n$ ). Suppose further that we had Y colors, then the number of G-invariant colorings  $Y^X/G$  is counted by

$$\frac{1}{|G|} \sum_{g \in G} |Y|^{c(g)}$$

with c(g) counting the number of cycles of g when viewed as a permutation of X. We can generalize this to a weighted version: if the color i can occur exactly  $r_i$  times, then this is counted by the coefficient of  $t_1^{r_1} \dots t_n^{r_n}$  in the polynomial

$$Z(t_1, \dots, t_n) = \frac{1}{|G|} \sum_{g \in G} \prod_{m \ge 1} (t_1^m + \dots + t_n^m)^{c_m(g)}$$

where  $c_m(g)$  counts the number of length m cycles in g acting as a permutation on X. Note we get the original formula by setting all  $t_i = 1$ . Here Z is the cycle index. Note: you can cleverly deal with even/odd sizes by setting some  $t_i$  to -1.

**Lucas Theorem** If p is prime, then:

$$\frac{p^a}{k} \equiv 0 \pmod{p}$$

Thus for non-negative integers  $m = m_k p^k + \ldots + m_1 p + m_0$ and  $n = n_k p^k + \ldots + n_1 p + n_0$ :

$$\frac{m}{n} = \prod_{i=0}^k \frac{m_i}{n_i} \mod p$$

Note: The fraction's mean integer division.

Catalan Numbers - Number of correct bracket sequence consisting of n opening and n closing brackets.

The number of ways to completely parenthesize n+1 factors.

The number of triangulations of a convex polygon with n+2 sides (i.e. the number of partitions of polygon into disjoint triangles by using the diagonals).

The number of ways to connect the 2n points on a circle to form n disjoint i.e. non-intersecting chords.

$$C_n = \frac{1}{n+1} \binom{2n}{n}$$

$$C_0 = 1, \ C_1 = 1, \ C_n = \sum_{k=0}^{n-1} C_k C_{n-1-k}$$

Narayana numbers The number of expressions containing n pairs of parentheses, which are correctly matched and which contain k distinct nestings.

$$N(n,k) = \frac{1}{n} \frac{n}{k} \frac{n}{k-1}$$