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1 Setup

1.1 header.h

```
1 #pragma once
2 #include <bits/stdc++.h>
3 using namespace std;
4
5 #define ll long long
6 #define ull unsigned ll
7 #define ld long double
8 #define pl pair<ll, ll>
9 #define pi pair<int, int>
10 #define vl vector<ll>
```

```
11 #define vi vector<int>
12 #define vvi vector <vi>
13 #define vvl vector <vl>
14 #define vpl vector <pl>
15 #define vpi vector <pi>
16 #define vld vector<ld>
17 #define in_fast(el, cont) (cont.find(el) != cont.end())
18 #define in(el, cont) (find(cont.begin(), cont.end(), el) != cont.end())
20 constexpr int INF = 200000010;
21 constexpr 11 LLINF = 900000000000000010LL;
23 template <typename T, template <typename ELEM, typename ALLOC = std::
      allocator < ELEM > > class Container >
24 std::ostream& operator << (std::ostream& o, const Container <T>& container) {
    typename Container <T>::const_iterator beg = container.begin();
    if (beg != container.end()) {
      o << *beg++;
      while (beg != container.end()) {
        o << " " << *beg++;
    }
31
    return o;
33 }
35 // int main() {
36 // ios::sync_with_stdio(false); // do not use cout + printf
37 // cin.tie(NULL);
38 // cout << fixed << setprecision(12);
39 // return 0;
40 // }
```

1.2 Bash for c++ compile with header.h

1.3 Bash for run tests c++

```
g++ $1/$1.cpp -o $1/$1.out
2 for file in $1/*.in; do diff <($1/$1.out < "$file") "${file%.in}.ans"; done
```

1.4 Bash for run tests python

```
1 for file in $1/*.in; do diff <(python3 $1/$1.py < "$file") "${file%.in}.ans
     ": done
```

1.4.1 Auxiliary helper C++

```
1 #include "header.h"
3 int main() {
      // Read in a line including white space
      string line;
      getline(cin, line);
      // When doing the above read numbers as follows:
      getline(cin, line);
      stringstream ss(line);
      ss >> n:
12
      // Count the number of 1s in binary representation of a number
13
      ull number;
14
      __builtin_popcountll(number);
15
16 }
```

1.4.2 Auxiliary helper python

```
1 # Read until EOF
2 while True:
     try:
          pattern = input()
     except EOFError:
          break
```

2 Python

2.1 Graphs

2.1.1 BFS

```
1 from collections import deque
2 def bfs(g, roots, n):
      q = deque(roots)
      explored = set(roots)
      distances = [float("inf")]*n
      distances[0][0] = 0
      while len(q) != 0:
          node = q.popleft()
          if node in explored: continue
          explored.add(node)
11
          for neigh in g[node]:
12
              if neigh not in explored:
                  q.append(neigh)
                  distances[neigh] = distances[node] + 1
      return distances
```

2.1.2 Dijkstra

12 13

16

17

20

21

```
1 from heapq import *
2 def dijkstra(n, root, g): # g = {node: (cost, neigh)}
    dist = [float("inf")]*n
    dist[root] = 0
    prev = [-1]*n
    pq = [(0, root)]
    heapify(pq)
    visited = set([])
    while len(pq) != 0:
      _, node = heappop(pq)
      if node in visited: continue
14
      visited.add(node)
      # In case of disconnected graphs
18
      if node not in g:
        continue
19
      for cost, neigh in g[node]:
        alt = dist[node] + cost
        if alt < dist[neigh]:</pre>
23
          dist[neigh] = alt
24
          prev[neigh] = node
          heappush(pq, (alt, neigh))
    return dist
```

Dynamic Programming

- 2.3 Trees
- 2.4 Number Theory / Combinatorics
- 2.4.1 nCk % prime

```
1 # Note: p must be prime and k < p</pre>
2 def fermat_binom(n, k, p):
      if k > n:
          return O
      # calculate numerator
      num = 1
      for i in range(n-k+1, n+1):
          num *= i % p
      num %= p
      # calculate denominator
11
      for i in range(1.k+1):
          denom *= i % p
14
      denom %= p
      # numerator * denominator^(p-2) (mod p)
      return (num * pow(denom, p-2, p)) % p
```

2.4.2 Sieve of Eratosthenes O(n) so actually faster than C++ version, but more memory

```
_{1} MAX_SIZE = 10**8+1
2 isprime = [True] * MAX SIZE
3 prime = []
4 SPF = [None] * (MAX SIZE)
6 def manipulated_seive(N): # Up to N (not included)
    isprime[0] = isprime[1] = False
    for i in range(2, N):
      if isprime[i] == True:
        prime.append(i)
        SPF[i] = i
      j = 0
12
      while (j < len(prime) and
13
        i * prime[j] < N and
14
          prime[j] <= SPF[i]):</pre>
        isprime[i * prime[j]] = False
16
        SPF[i * prime[j]] = prime[j]
17
```

2.5 Strings

2.5.1 LCS

```
1 def longestCommonSubsequence(text1, text2): # 0(m*n) time, 0(m) space
      n = len(text1)
      m = len(text2)
      # Initializing two lists of size m
      prev = [0] * (m + 1)
      cur = \lceil 0 \rceil * (m + 1)
      for idx1 in range(1, n + 1):
          for idx2 in range(1, m + 1):
10
              # If characters are matching
11
              if text1[idx1 - 1] == text2[idx2 - 1]:
                   cur[idx2] = 1 + prev[idx2 - 1]
               else:
                   # If characters are not matching
                   cur[idx2] = max(cur[idx2 - 1], prev[idx2])
          prev = cur.copy()
18
19
      return cur[m]
```

2.5.2 KMP

```
class KMP:
def partial(self, pattern):
    """ Calculate partial match table: String -> [Int]"""
    ret = [0]
    for i in range(1, len(pattern)):
        j = ret[i - 1]
    while j > 0 and pattern[j] != pattern[i]: j = ret[j - 1]
```

```
ret.append(j + 1 if pattern[j] == pattern[i] else j)
9
          return ret
10
      def search(self, T, P):
11
          """KMP search main algorithm: String -> String -> [Int]
12
          Return all the matching position of pattern string P in T"""
13
          partial, ret, j = self.partial(P), [], 0
14
          for i in range(len(T)):
               while j > 0 and T[i] != P[j]: j = partial[j - 1]
16
              if T[i] == P[j]: j += 1
17
              if i == len(P):
18
                   ret.append(i - (j - 1))
                   j = partial[j - 1]
20
          return ret
```

2.6 Geometry

2.7 Other Algorithms

2.7.1 Rotate matrix

2.8 Other Data Structures

2.8.1 Segment Tree

```
_{1} N = 100000 # limit for array size
2 tree = [0] * (2 * N) # Max size of tree
4 def build(arr, n): # function to build the tree
      # insert leaf nodes in tree
      for i in range(n):
          tree[n + i] = arr[i]
      # build the tree by calculating parents
      for i in range(n - 1, 0, -1):
10
          tree[i] = tree[i << 1] + tree[i << 1 | 1]</pre>
11
12
13 def updateTreeNode(p, value, n): # function to update a tree node
      # set value at position p
14
15
      tree[p + n] = value
      p = p + n
16
17
      i = p # move upward and update parents
18
19
      while i > 1:
          tree[i >> 1] = tree[i] + tree[i ^ 1]
20
          i >>= 1
21
23 def query(1, r, n): # function to get sum on interval [1, r)
24
      # loop to find the sum in the range
25
      1 += n
      r += n
```

```
while l < r:
29
           if 1 & 1:
               res += tree[1]
               1 += 1
31
           if r & 1:
32
               r -= 1
33
                res += tree[r]
34
           1 >>= 1
           r >>= 1
36
       return res
37
```

3 C++

3.1 Graphs

3.1.1 BFS

```
1 #include "header.h"
2 #define graph unordered_map<11, unordered_set<11>>
3 vi bfs(int n, graph& g, vi& roots) {
      vi parents(n+1, -1); // nodes are 1..n
      unordered_set <int> visited;
      queue < int > q;
      for (auto x: roots) {
          q.emplace(x);
           visited.insert(x);
10
      while (not q.empty()) {
11
           int node = q.front();
12
          q.pop();
13
14
           for (auto neigh: g[node]) {
15
               if (not in(neigh, visited)) {
16
                   parents[neigh] = node;
17
                   q.emplace(neigh);
18
                   visited.insert(neigh);
19
              }
          }
21
22
23
      return parents;
24 }
     reconstruct_path(vi parents, int start, int goal) {
      vi path:
26
      int curr = goal;
27
      while (curr != start) {
           path.push_back(curr);
29
           if (parents[curr] == -1) return vi(); // No path, empty vi
30
           curr = parents[curr];
31
32
      path.push_back(start);
      reverse(path.begin(), path.end());
34
      return path;
35
36 }
```

3.1.2 DFS Cycle detection / removal

```
1 #include "header.h"
2 void removeCyc(ll node, unordered_map<ll, vector<pair<ll, 11>>>& neighs,
      vector < bool > & visited.
3 vector < bool > & recStack, vector < 11 > & ans) {
      if (!visited[node]) {
           visited[node] = true:
           recStack[node] = true;
           auto it = neighs.find(node);
           if (it != neighs.end()) {
               for (auto util: it->second) {
                   11 nnode = util.first:
10
                   if (recStack[nnode]) {
11
                        ans.push_back(util.second);
12
                   } else if (!visited[nnode]) {
13
                        removeCyc(nnode, neighs, visited, recStack, ans);
14
15
               }
16
           }
17
18
      recStack[node] = false;
19
20 }
```

3.1.3 Diikstra

```
1 #include "header.h"
2 vector < int > dijkstra(int n, int root, map < int, vector < pair < int, int >>> & g) {
    unordered_set <int> visited;
    vector<int> dist(n, INF);
       priority_queue < pair < int , int >> pq;
       dist[root] = 0;
       pq.push({0, root});
       while (!pq.empty()) {
           int node = pq.top().second;
           int d = -pq.top().first;
10
           pq.pop();
11
12
           if (in(node, visited)) continue;
13
           visited.insert(node);
14
15
           for (auto e : g[node]) {
16
               int neigh = e.first;
17
               int cost = e.second;
18
               if (dist[neigh] > dist[node] + cost) {
19
                    dist[neigh] = dist[node] + cost;
20
                   pq.push({-dist[neigh], neigh});
21
22
           }
23
24
       return dist;
25
26 }
```

3.1.4 Floyd-Warshall

```
1 #include "header.h"
2 // g[i][j] = infty if not path from i to j
```

3.1.5 Kruskal Minimum spanning tree of undirected weighted graph

```
1 #include "header.h"
2 #include "disjoint set.h"
3 // O(E log E)
4 pair < set < pair < 11, 11 >> , 11 > kruskal (vector < tuple < 11, 11, 11 >> & edges, 11 n)
       set <pair <11, 11>> ans;
      11 cost = 0:
       sort(edges.begin(), edges.end());
      DisjointSet < ll > fs(n);
10
      ll dist, i, j;
11
      for (auto edge: edges) {
           dist = get<0>(edge);
13
           i = get<1>(edge);
14
           j = get < 2 > (edge);
15
16
           if (fs.find_set(i) != fs.find_set(j)) {
17
               fs.union_sets(i, j);
18
               ans.insert({i, j});
19
20
               cost += dist:
           }
21
22
      return pair<set<pair<11, 11>>, 11> {ans, cost};
23
24 }
```

3.1.6 Hungarian algorithm

```
* to assign the first (j+1) jobs to distinct workers
15 template <class T> vector<T> hungarian(const vector<vector<T>> &C) {
      const int J = (int)size(C), W = (int)size(C[0]);
      assert(J <= W):
17
      // job[w] = job assigned to w-th worker, or -1 if no job assigned
18
      // note: a W-th worker was added for convenience
      vector < int > iob(W + 1, -1):
20
      vector<T> ys(J), yt(W + 1); // potentials
21
      // -yt[W] will equal the sum of all deltas
22
      vector <T> answers:
23
      const T inf = numeric_limits<T>::max();
24
      for (int j_cur = 0; j_cur < J; ++j_cur) { // assign j_cur-th job</pre>
25
          int w_cur = W;
26
27
          job[w_cur] = i_cur;
          // min reduced cost over edges from Z to worker w
28
          vector <T> min_to(W + 1, inf);
          vector<int> prv(W + 1, -1); // previous worker on alternating path
30
          vector < bool > in Z(W + 1): // whether worker is in Z
31
          while (job[w_cur] != -1) { // runs at most j_cur + 1 times
32
33
              in_Z[w_cur] = true;
              const int j = job[w_cur];
34
              T delta = inf;
25
              int w next:
36
              for (int w = 0; w < W; ++w) {
                   if (!in Z[w]) {
                       if (ckmin(min_to[w], C[j][w] - ys[j] - yt[w]))
39
                           prv[w] = w_cur;
40
                       if (ckmin(delta, min_to[w])) w_next = w;
                   }
42
43
               // delta will always be non-negative,
44
               // except possibly during the first time this loop runs
45
              // if any entries of C[j_cur] are negative
              for (int w = 0; w \le W; ++w) {
                   if (in_Z[w]) ys[job[w]] += delta, yt[w] -= delta;
                   else min to[w] -= delta:
49
               w_cur = w_next;
          }
52
          // update assignments along alternating path
53
          for (int w; w_cur != W; w_cur = w) job[w_cur] = job[w = prv[w_cur]];
          answers.push_back(-yt[W]);
55
56
57
      return answers;
```

3.2 Dynamic Programming

3.2.1 Longest Increasing Subsequence

```
#include "header.h"
template < class T >
vector < T > index_path_lis(vector < T > & nums) {
   int n = nums.size();
   vector < T > sub;
   vector < int > subIndex;
   vector < T > path(n, -1);
```

```
for (int i = 0; i < n; ++i) {</pre>
        if (sub.emptv() || sub[sub.size() - 1] < nums[i]) {</pre>
      path[i] = sub.empty() ? -1 : subIndex[sub.size() - 1];
      sub.push_back(nums[i]);
      subIndex.push back(i):
       } else {
13
      int idx = lower_bound(sub.begin(), sub.end(), nums[i]) - sub.begin();
14
      path[i] = idx == 0 ? -1 : subIndex[idx - 1]:
      sub[idx] = nums[i];
16
      subIndex[idx] = i:
17
        }
    vector <T> ans:
    int t = subIndex[subIndex.size() - 1];
    while (t != -1) {
        ans.push_back(t);
        t = path[t];
25
    reverse(ans.begin(), ans.end());
    return ans;
27
     Length only
30 template < class T>
31 int length lis(vector <T > &a) {
    set <T> st;
    typename set<T>::iterator it:
    for (int i = 0; i < a.size(); ++i) {</pre>
      it = st.lower_bound(a[i]);
      if (it != st.end()) st.erase(it);
      st.insert(a[i]);
    return st.size();
39
40 }
```

3.3 Trees

3.4 Number Theory / Combinatorics

3.4.1 Modular exponentiation Or use pow() in python

3.4.2 GCD Or math.gcd in python, std::gcd in C++

```
#include "header.h"
2 ll gcd(ll a, ll b) {
3    if (a == 0) return b;
4    return gcd(b % a, a);
5 }
```

3.4.3 Sieve of Eratosthenes

```
#include "header.h"

void getprimes(ll n) { // Up to n (not included)

vector<bool> p(n, true);

p[0] = false;

p[1] = false;

for(ll i = 0; i < n; i++) {

if(p[i]) {

primes.push_back(i);

for(ll j = i*2; j < n; j+=i) p[j] = false;

}
}</pre>
```

3.4.4 Fibonacci % prime

```
#include "header.h"
const ll MOD = 1000000007;
unordered_map<ll, ll> Fib;
ll fib(ll n) {
    if (n < 2) return 1;
    if (Fib.find(n) != Fib.end()) return Fib[n];
    Fib[n] = (fib((n + 1) / 2) * fib(n / 2) + fib((n - 1) / 2) * fib((n - 2) / 2)) % MOD;
    return Fib[n];
}</pre>
```

3.4.5 nCk % prime

```
1 #include "header.h"
2 ll binom(ll n. ll k) {
      for (ll i = 1; i <= min(k,n-k); ++i) ans = ans*(n+1-i)/i;
      return ans:
7 ll mod nCk(ll n, ll k, ll p ){
      ll ans = 1:
      while(n){
          11 np = n\%p, kp = k\%p;
10
          if(kp > np) return 0;
11
          ans *= binom(np,kp);
12
          n /= p; k /= p;
13
14
15
      return ans;
16 }
```

3.5.1 Aho-Corasick algorithm Also can be used as Knuth-Morris-Pratt algorithm

```
1 #include "header.h"
3 map < char, int > cti;
4 int cti_size;
5 template <int ALPHABET_SIZE, int (*mp)(char)>
6 struct AC_FSM {
     struct Node {
      int child[ALPHABET_SIZE], failure = 0, match_par = -1;
      Node() { for (int i = 0; i < ALPHABET_SIZE; ++i) child[i] = -1; }
    };
11
    vector < Node > a;
12
    vector<string> &words;
13
    AC_FSM(vector<string> &words) : words(words) {
14
      a.push_back(Node());
15
       construct_automaton();
16
17
     void construct_automaton() {
18
      for (int w = 0, n = 0; w < words.size(); ++w, n = 0) {</pre>
19
         for (int i = 0; i < words[w].size(); ++i) {</pre>
20
           if (a[n].child[mp(words[w][i])] == -1) {
21
             a[n].child[mp(words[w][i])] = a.size();
22
             a.push_back(Node());
23
24
           n = a[n].child[mp(words[w][i])];
25
26
27
         a[n].match.push_back(w);
       queue < int > q;
29
       for (int k = 0; k < ALPHABET_SIZE; ++k) {</pre>
30
         if (a[0].child[k] == -1) a[0].child[k] = 0;
31
         else if (a[0].child[k] > 0) {
32
           a[a[0].child[k]].failure = 0;
           q.push(a[0].child[k]);
34
35
36
       while (!q.empty()) {
37
         int r = q.front(); q.pop();
38
         for (int k = 0, arck; k < ALPHABET_SIZE; ++k) {</pre>
39
           if ((arck = a[r].child[k]) != -1) {
40
             q.push(arck);
             int v = a[r].failure;
             while (a[v].child[k] == -1) v = a[v].failure:
             a[arck].failure = a[v].child[k];
44
             a[arck].match_par = a[v].child[k];
45
             while (a[arck].match_par != -1
                 && a[a[arck].match_par].match.empty())
47
               a[arck].match_par = a[a[arck].match_par].match_par;
48
        }
50
51
52
    void aho_corasick(string &sentence, vvi &matches){
53
      matches.assign(words.size(), vi());
      int state = 0, ss = 0;
```

```
for (int i = 0; i < sentence.length(); ++i, ss = state) {</pre>
56
57
        while (a[ss].child[mp(sentence[i])] == -1)
           ss = a[ss].failure;
58
        state = a[state].child[mp(sentence[i])]
59
             = a[ss].child[mp(sentence[i])];
60
        for (ss = state; ss != -1; ss = a[ss].match_par)
61
           for (int w : a[ss].match)
62
             matches[w].push_back(i + 1 - words[w].length());
64
65
66 }:
67 int char_to_int(char c) {
    return cti[c]:
69 }
```

3.5.2 KMP

```
1 #include "header.h"
void compute_prefix_function(string &w, vi &prefix) {
    prefix.assign(w.length(), 0);
    int k = prefix[0] = -1;
    for(int i = 1; i < w.length(); ++i) {</pre>
      while (k >= 0 \&\& w[k + 1] != w[i]) k = prefix[k];
      if(w[k + 1] == w[i]) k++:
      prefix[i] = k;
10
11 }
12 void knuth_morris_pratt(string &s, string &w) {
    int q = -1;
    vi prefix;
    compute_prefix_function(w, prefix);
    for(int i = 0; i < s.length(); ++i) {</pre>
      while (q \ge 0 \&\& w[q + 1] != s[i]) q = prefix[q];
      if(w[q + 1] == s[i]) q++;
      if(q + 1 == w.length()) {
19
        // Match at position (i - w.length() + 1)
20
        q = prefix[q];
21
22
23
```

3.6 Geometry

3.6.1 essentials.cpp

```
C operator* (const P &p) const { return x*p.x + y*p.y; }
    C operator^ (const P &p) const { return x*p.y - p.x*y; }
    P perp() const { return P{y, -x}; }
    C lensq() const { return x*x + y*y; }
    ld len() const { return sqrt((ld)lensq()); }
    static ld dist(const P &p1, const P &p2) {
      return (p1-p2).len(); }
17
    bool operator == (const P &r) const {
      return ((*this)-r).lensq() <= EPS*EPS; }</pre>
19
20 };
21 C det(P p1, P p2) { return p1^p2; }
22 C det(P p1, P p2, P o) { return det(p1-o, p2-o); }
    det(const vector <P> &ps) {
    C sum = 0; P prev = ps.back();
    for(auto &p : ps) sum += det(p, prev), prev = p;
    return sum:
27 }
28 // Careful with division by two and C=11
29 C area(P p1, P p2, P p3) { return abs(det(p1, p2, p3))/C(2); }
30 C area(const vector <P> &poly) { return abs(det(poly))/C(2); }
31 int sign(C c){ return (c > C(0)) - (c < C(0)); }</pre>
32 int ccw(P p1, P p2, P o) { return sign(det(p1, p2, o)); }
_{34} // Only well defined for C = 1d.
35 P unit(const P &p) { return p / p.len(); }
_{36} P rotate(P p, ld a) { return P{p.x*cos(a)-p.y*sin(a), p.x*sin(a)+p.y*cos(a)}
      }: }
```

3.6.2 Convex Hull

```
1 #include "header.h"
2 #include "essentials.cpp"
3 struct ConvexHull { // O(n lg n) monotone chain.
    vector < size_t > h, c; // Indices of the hull are in 'h', ccw.
    const vector <P> &p:
    ConvexHull(const vector <P> &_p) : n(_p.size()), c(n), p(_p) {
      std::iota(c.begin(), c.end(), 0);
      std::sort(c.begin(), c.end(), [this](size_t 1, size_t r) -> bool {
          return p[1].x != p[r].x ? p[1].x < p[r].x : p[1].y < p[r].y; });
      c.erase(std::unique(c.begin(), c.end(), [this](size_t l, size_t r) {
10
          return p[1] == p[r]; }), c.end());
      for (size_t s = 1, r = 0; r < 2; ++r, s = h.size()) {</pre>
11
        for (size_t i : c) {
12
          while (h.size() > s && ccw(p[h.end()[-2]], p[h.end()[-1]], p[i]) <=</pre>
13
            h.pop_back();
          h.push_back(i);
16
17
         reverse(c.begin(), c.end());
18
      if (h.size() > 1) h.pop_back();
19
20
21
    size t size() const { return h.size(): }
    template <class T, void U(const P &, const P &, const P &, T &)>
    void rotating_calipers(T &ans) {
      if (size() <= 2)
24
        U(p[h[0]], p[h.back()], p[h.back()], ans);
25
```

```
26
        for (size t i = 0, i = 1, s = size(); i < 2 * s; ++i) {
          while (\det(p[h[(i + 1) \% s]) - p[h[i \% s]), p[h[(j + 1) \% s]] - p[h[
            j = (j + 1) \% s;
29
          U(p[h[i % s]], p[h[(i + 1) % s]], p[h[j]], ans);
30
31
    }
32
33 };
34 // Example: furthest pair of points. Now set ans = OLL and call
35 // ConvexHull(pts).rotating_calipers<11, update>(ans);
36 void update(const P &p1, const P &p2, const P &o, 11 &ans) {
    ans = \max(ans, (11)\max((p1 - o).lensq(), (p2 - o).lensq()));
38 }
```

3.7 Other Algorithms

3.8 Other Data Structures

3.8.1 Disjoint set (i.e. union-find)

```
1 template <typename T>
2 class DisjointSet {
       typedef T * iterator;
      T *parent, n, *rank;
       public:
           // O(n), assumes nodes are [0, n)
           DisjointSet(T n) {
               this->parent = new T[n];
               this -> n = n:
               this->rank = new T[n];
10
11
               for (T i = 0; i < n; i++) {</pre>
12
                    parent[i] = i;
13
                    rank[i] = 0;
14
15
           }
16
17
           // O(log n)
           T find_set(T x) {
19
               if (x == parent[x]) return x;
20
               return parent[x] = find_set(parent[x]);
21
           }
22
           // O(log n)
24
25
           void union_sets(T x, T y) {
               x = this->find_set(x);
26
               v = this->find_set(v);
27
28
               if (x == y) return;
29
30
               if (rank[x] < rank[y]) {</pre>
                   Tz = x;
                    x = y;
33
34
                    y = z;
               }
35
```

3.8.2 Fenwick tree (i.e. BIT) eff. update + prefix sum calc.

```
1 #include "header.h"
2 #define maxn 200010
3 int t,n,m,tree[maxn],p[maxn];
5 void update(int k, int z) {
      while (k <= maxn) {
          tree[k] += z:
          k += k & (-k);
          // cout << "k: " << k << endl;
11 }
13 int sum(int k) {
      int ans = 0;
      while(k) {
          ans += tree[k];
          k = k & (-k);
17
      return ans;
```

4 Other Mathematics

4.1 Helpful functions

4.1.1 Euler's Totient Fucntion $n = p_1^{k_1-1} \cdot (p_1-1) \cdot \ldots \cdot p_r^{k_r-1} \cdot (p_r-1)$, where $p_1^{k_1} \cdot \ldots \cdot p_r^{k_r}$ is the prime factorization of n.

```
1 # include "header.h"
2 11 phi(11 n) { // \Phi(n)
      11 \text{ ans} = 1;
      for (ll i = 2; i*i <= n; i++) {</pre>
           if (n % i == 0) {
               ans *= i-1;
               n /= i;
               while (n \% i == 0) {
                    ans *= i;
                   n /= i;
11
12
      if (n > 1) ans *= n-1;
14
      return ans;
15
16 }
     phis(int n) { // All \Phi(i) up to n
    vi phi(n + 1, OLL);
    iota(phi.begin(), phi.end(), OLL);
```

```
for (11 i = 2LL; i <= n; ++i)
if (phi[i] == i)
for (11 j = i; j <= n; j += i)
phi[j] -= phi[j] / i;
return phi;
}</pre>
```

Formulas $\Phi(n)$ counts all numbers in $1, \ldots, n-1$ coprime to n. $a^{\varphi(n)} \equiv 1 \mod n$, a and n are coprimes. $\forall e > \log_2 m : n^e \mod m = n^{\Phi(m) + e \mod \Phi(m)} \mod m$. $\gcd(m, n) = 1 \Rightarrow \Phi(m \cdot n) = \Phi(m) \cdot \Phi(n)$.

4.2 Theorems and definitions

Fermat's little theorem $a^p \equiv a \mod p$

Subfactorial
$$!n = n! \sum_{i=0}^{n} \frac{(-1)^{i}}{i!}, !(0) = 1, !n = n \cdot !(n-1) + (-1)^{n}$$

Least common multiple $lcm(a, b) = a \cdot b/gcd(a, b)$

Binomials and other partitionings We have $\binom{n}{k} = \binom{n-1}{k} + \binom{n-1}{k-1} = \prod_{i=1}^k \frac{n-i+1}{i}$. This last product may be computed incrementally since any product of k' consecutive values is divisible k'!. Basic identities: The hockeystick identity: $\sum_{k=r}^{n} \binom{k}{r} = \binom{n+1}{r+1}$ or $\sum_{k \leq n} \binom{r+k}{k} = \binom{r+n+1}{n}$. Also $\sum_{k=0}^{n} \binom{k}{m} = \binom{n+1}{m+1}$.

For $n, m \ge 0$ and p prime. Write n, m in base p, i.e. $n = n_k p^k + \dots + n_1 p + n_0$ and $m = m_k p^k + \dots + m_1 p + m_0$. Then by Lucas theorem we have $\binom{n}{m} \equiv \prod_{i=0}^k \binom{n_i}{m_i} \mod p$, with the convention that $n_i < m_i \implies \binom{n_i}{m_i} = 0$.

Fibonacci (See also number theory section)

$$\sum_{0 \le k \le n} {n-k \choose k} = F_{n+1}, F_n = \frac{1}{\sqrt{5}} \left(\frac{1+\sqrt{5}}{2}\right)^n - \frac{1}{\sqrt{5}} \left(\frac{1-\sqrt{5}}{2}\right)^n,$$

$$\sum_{i=1}^n F_i = F_{n+2} - 1, \sum_{i=1}^n F_i^2 = F_n F_{n+1},$$

$$\gcd(F_m, F_n) = F_{\gcd(m,n)}, \gcd(F_n, F_{n+1}) = \gcd(F_n, F_{n+2}) = 1$$

Bit stuff $a + b = a \oplus b + 2(a \& b) = a|b + a \& b$.

kth bit is set in x iff $x \mod 2^{k-1} \ge 2^k$, or iff $x \mod 2^{k-1} - x \mod 2^k \ne 0$ (i.e. $= 2^k$) It comes handy when you need to look at the bits of the numbers which are pair sums or subset sums etc.

```
n \mod 2^i = n\&(2^i - 1).
 \forall k: 1 \oplus 2 \oplus \ldots \oplus (4k - 1) = 0
```