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_							
1 Setup		3 using names	3 using namespace std;		possible/necessary 12 #define vb vector <bool></bool>		
		5 #define ll	long long		ne vvi vector <vi></vi>		
1.1	header.h	6 #define ull	6 #define ull unsigned ll 7 #define ld long double 8 #define pl pair <ll, ll=""></ll,>		14 #define vvl vector <vl></vl>		
					<pre>15 #define vpl vector<pl> 16 #define vpi vector<pi></pi></pl></pre>		
			pair (if, if) pair (int, int) // use pl where		ne vpr vector vpr> ne vld vector <ld></ld>		
1 #pr	agma once // Delete this when copying this	possible	possible/necessary		18 #define vvpi vector < vpi >		
ш.	file		10 #define vl vector <ll> 11 #define vi vector<int> // change to vl where</int></ll>		ne in_fast(el, cont) (cont.find(el) != con nd())	ıt.	
2 #in	clude <bits stdc++.h=""></bits>	11 #dCIIIE VI	// change to vi where	eı	· · · · · · · · · · · · · · · · · · ·		

```
20 #define in(el, cont) (find(cont.begin(), cont.end
      (), el) != cont.end())
22 constexpr int INF = 200000010;
23 constexpr 11 LLINF = 900000000000000010LL;
25 template <typename T, template <typename ELEM,
      typename ALLOC = std::allocator < ELEM > > class
       Container >
26 std::ostream& operator << (std::ostream& o, const
      Container < T > & container ) {
    typename Container <T>::const_iterator beg =
        container.begin();
    if (beg != container.end()) {
      o << *beg++;
      while (beg != container.end()) {
        o << " " << *beg++;
    return o:
35 }
37 // int main() {
38 // ios::sync_with_stdio(false); // do not use
      cout + printf
      cin.tie(NULL);
40 // cout << fixed << setprecision(12);
41 // return 0:
42 // }
```

# 1.2 Bash for c++ compile with header.h

# 1.3 Bash for run tests c++

```
1 g++ $1/$1.cpp -o $1/$1.out
2 for file in $1/*.in; do diff <($1/$1.out < "$file
") "${file%.in}.ans"; done</pre>
```

# 1.4 Bash for run tests python

```
for file in $1/*.in; do diff <(python3 $1/$1.py < "$file") "${file%.in}.ans"; done
```

# 1.4.1 Aux. helper C++

```
1 #include "header.h"
3 int main() {
      // Read in a line including white space
      string line;
      getline(cin, line);
      // When doing the above read numbers as
          follows:
      int n:
      getline(cin, line);
      stringstream ss(line):
      // Count the number of 1s in binary
          represnatation of a number
      ull number:
      __builtin_popcountll(number);
16 }
17
18 // int128
19 using 111 = __int128;
20 ostream& operator << ( ostream& o, __int128 n ) {</pre>
    auto t = n < 0 ? -n : n; char b[128], *d = end(b)
    do *--d = '0'+t\%10, t /= 10; while (t);
    if(n<0) *--d = '-';
    o.rdbuf()->sputn(d,end(b)-d);
    return o:
```

# 1.4.2 Aux. helper python

```
from functools import lru_cache

# Read until EOF

while True:
    try:
    pattern = input()
    except EOFError:
    break

# Clru_cache(maxsize=None)
def smth_memoi(i, j, s):
    # Example in-built cache
```

# 2 Python

# 2.1 Graphs

#### 2.1.1 BFS

```
1 from collections import deque
2 def bfs(g, roots, n):
     q = deque(roots)
      explored = set()
      distances = [0 if v in roots else float('inf'
         ) for v in range(n)]
      while len(q) != 0:
          node = q.popleft()
          if node in explored: continue
          explored.add(node)
          for neigh in g[node]:
             if neigh not in explored:
                  q.append(neigh)
                  distances[neigh] = distances[node
                      1 + 1
      return distances
```

#### 2.1.2 Dijkstra

```
1 from heapq import *
2 def dijkstra(n, root, g): # g = {node: (cost, neigh)}
3  dist = [float("inf")]*n
4  dist[root] = 0
5  prev = [-1]*n
6
7  pq = [(0, root)]
8  heapify(pq)
```

27

28

30

31

32

33

36

42

44

45

47

48

53

55

67

70

73

```
visited = set([])
    while len(pq) != 0:
      _, node = heappop(pq)
13
      if node in visited: continue
14
15
      visited.add(node)
16
      # In case of disconnected graphs
17
      if node not in g:
         continue
19
20
      for cost, neigh in g[node]:
21
         alt = dist[node] + cost
22
         if alt < dist[neigh]:</pre>
23
           dist[neigh] = alt
^{24}
           prev[neigh] = node
25
          heappush(pq, (alt, neigh))
    return dist
```

# 2.1.3 Topological Sort

```
1 #Python program to print topological sorting of a
2 from collections import defaultdict
4 #Class to represent a graph
5 class Graph:
      def __init__(self, vertices):
          self.graph = defaultdict(list) #
              dictionary containing adjacency List
          self.V = vertices #No. of vertices
      # function to add an edge to graph
10
      def addEdge(self,u,v):
11
          self.graph[u].append(v)
12
13
      # A recursive function used by
          topologicalSort
      def topologicalSortUtil(self,v,visited,stack)
15
16
          # Mark the current node as visited.
17
          visited[v] = True
18
19
          # Recur for all the vertices adjacent to
              this vertex
          for i in self.graph[v]:
              if visited[i] == False:
22
                  self.topologicalSortUtil(i,
23
                      visited.stack)
          # Push current vertex to stack which
25
              stores result
```

```
stack.insert(0,v)
# The function to do Topological Sort. It
   uses recursive
# topologicalSortUtil()
def topologicalSort(self):
   # Mark all the vertices as not visited
```

stack =[] # Call the recursive helper function to store Topological # Sort starting from all vertices one by for i in range(self.V): if visited[i] == False: self.topologicalSortUtil(i,

visited, stack)

# Print contents of stack return stack

visited = [False]\*self.V

def isCyclicUtil(self, v, visited, recStack): # Mark current node as visited and

# adds to recursion stack visited[v] = True recStack[v] = True

# Recur for all neighbours # if any neighbour is visited and in # recStack then graph is cyclic for neighbour in self.graph[v]: if visited[neighbour] == False: if self.isCvclicUtil(neighbour. visited. recStack) == True:

> return True elif recStack[neighbour] == True: return True

# The node needs to be popped from # recursion stack before function ends recStack[v] = False return False

# Returns true if graph is cyclic else false def isCyclic(self):

visited = [False] \* (self.V + 1) recStack = [False] \* (self.V + 1) for node in range(self.V): if visited[node] == False: if self.isCvclicUtil(node.

> visited, recStack) == True: return True

return False

#### 2.1.4 Kruskal (UnionFind)

```
class UnionFind:
      def __init__(self, n):
           self.parent = [-1]*n
      def find(self, x):
           if self.parent[x] < 0:</pre>
               return x
           self.parent[x] = self.find(self.parent[x
           return self.parent[x]
      def connect(self. a. b):
          ra = self.find(a)
           rb = self.find(b)
          if ra == rb:
               return False
           if self.parent[ra] > self.parent[rb]:
               self.parent[rb] += self.parent[ra]
17
               self.parent[ra] = rb
               self.parent[ra] += self.parent[rb]
               self.parent[rb] = ra
           return True
24 # Full MST is len(spanning==n-1)
25 def kruskal(n. edges):
      uf = UnionFind(n)
      spanning = []
      edges.sort(key = lambda d: -d[2])
      while edges and len(spanning) < n-1:
           u, v, w = edges.pop()
           if not uf.connect(u, v):
               continue
32
           spanning.append((u, v, w))
      return spanning
36 # Example
37 \text{ edges} = [(1, 2, 10), (2, 3, 20)]
```

# Num. Th. / Comb.

#### 2.2.1 nCk % prime

```
1 # Note: p must be prime and k < p</pre>
2 def fermat_binom(n, k, p):
      if k > n:
          return 0
      # calculate numerator
      n 11 m = 1
      for i in range(n-k+1, n+1):
          num *= i % p
      num %= p
```

```
# calculate denominator
denom = 1
for i in range(1,k+1):
    denom *= i % p
denom %= p
# numerator * denominator^(p-2) (mod p)
return (num * pow(denom, p-2, p)) % p
```

# **2.2.2** Sieve of E. O(n) so actually faster than C++ version, but more memory

```
1 MAX SIZE = 10**8+1
2 isprime = [True] * MAX_SIZE
3 \text{ prime} = []
4 SPF = [None] * (MAX SIZE)
6 def manipulated_seive(N): # Up to N (not
      included)
    isprime[0] = isprime[1] = False
    for i in range(2, N):
      if isprime[i] == True:
        prime.append(i)
        SPF[i] = i
      while (j < len(prime) and
13
       i * prime[j] < N and</pre>
          prime[j] <= SPF[i]):</pre>
15
        isprime[i * prime[j]] = False
        SPF[i * prime[j]] = prime[j]
        i += 1
```

# 2.3 Strings

#### 2.3.1 LCS

```
1 def longestCommonSubsequence(text1, text2): # 0(
      m*n) time, O(m) space
      n = len(text1)
      m = len(text2)
      # Initializing two lists of size m
      prev = [0] * (m + 1)
      cur = [0] * (m + 1)
      for idx1 in range(1, n + 1):
          for idx2 in range(1, m + 1):
              # If characters are matching
11
              if text1[idx1 - 1] == text2[idx2 -
12
                  cur[idx2] = 1 + prev[idx2 - 1]
              else:
                  # If characters are not matching
```

#### 2.3.2 KMP

```
1 class KMP:
      def partial(self, pattern):
          """ Calculate partial match table: String
               -> [Int]"""
          ret = [0]
          for i in range(1, len(pattern)):
              j = ret[i - 1]
              while j > 0 and pattern[j] != pattern
                  [i]: j = ret[j - 1]
              ret.append(j + 1 if pattern[j] ==
                  pattern[i] else i)
          return ret
10
      def search(self, T, P):
11
          """KMP search main algorithm: String ->
12
              String -> [Int]
          Return all the matching position of
13
              pattern string P in T"""
          partial, ret, j = self.partial(P), [], 0
          for i in range(len(T)):
              while j > 0 and T[i] != P[j]: j =
16
                   partial[i - 1]
              if T[i] == P[i]: i += 1
17
              if i == len(P):
                  ret.append(i - (j - 1))
                  j = partial[j - 1]
20
          return ret
```

#### 2.3.3 Edit distance

```
# Initialize a variable to store the previous
    value
previous = 0
# Loop through the rows of the dynamic
    programming matrix
for i in range(1, m + 1):
  # Store the current value at the beginning of
  previous = curr[0]
  curr[0] = i
  # Loop through the columns of the dynamic
      programming matrix
  for j in range (1, n + 1):
    # Store the current value in a temporary
        variable
    temp = curr[j]
    # Check if the characters at the current
        positions in str1 and str2 are the same
    if str1[i - 1] == str2[i - 1]:
      curr[j] = previous
    else:
      # Update the current cell with the
          minimum of the three adjacent cells
      curr[j] = 1 + min(previous, curr[j - 1],
          curr[i])
    # Update the previous variable with the
        temporary value
    previous = temp
# The value in the last cell represents the
    minimum number of operations
return curr[n]
```

# 2.4 Other Algorithms

#### 2.4.1 Rotate matrix

# 2.5 Geometry

#### 2.5.1 Convex Hull

```
def vec(a,b):
    return (b[0]-a[0],b[1]-a[1])
def det(a,b):
```

```
return a[0]*b[1] - b[0]*a[1]
6 def convexhull(P):
      if (len(P) == 1):
          return [(p[0][0], p[0][1])]
      h = sorted(P)
      lower = []
11
      i = 0
      while i < len(h):
          if len(lower) > 1:
14
              a = vec(lower[-2], lower[-1])
              b = vec(lower[-1], h[i])
16
              if det(a,b) <= 0 and len(lower) > 1:
                  lower.pop()
                   continue
19
          lower.append(h[i])
20
          i += 1
21
22
      upper = []
23
      i = 0
^{24}
      while i < len(h):
25
          if len(upper) > 1:
              a = vec(upper[-2], upper[-1])
27
              b = vec(upper[-1], h[i])
              if det(a,b) >= 0:
29
                   upper.pop()
                   continue
          upper.append(h[i])
32
          i += 1
33
34
      reversedupper = list(reversed(upper[1:-1:]))
35
      reversedupper.extend(lower)
      return reversedupper
```

#### 2.5.2 Geometry

```
2 def vec(a,b):
      return (b[0]-a[0],b[1]-a[1])
5 def det(a,b):
      return a[0]*b[1] - b[0]*a[1]
      lower = []
      i = 0
      while i < len(h):
          if len(lower) > 1:
11
              a = vec(lower[-2], lower[-1])
12
              b = vec(lower[-1], h[i])
13
              if det(a,b) <= 0 and len(lower) > 1:
                  lower.pop()
                  continue
16
          lower.append(h[i])
```

```
i += 1
20
      # find upper hull
      # det <= 0 -> replace
      upper = []
      i = 0
23
      while i < len(h):
          if len(upper) > 1:
25
              a = vec(upper[-2], upper[-1])
26
              b = vec(upper[-1], h[i])
              if det(a,b) >= 0:
                   upper.pop()
                   continue
          upper.append(h[i])
31
          i += 1
```

#### 2.6 Other Data Structures

#### 2.6.1 Segment Tree

```
_{1} N = 100000 # limit for array size
2 tree = [0] * (2 * N) # Max size of tree
4 def build(arr. n): # function to build the tree
      # insert leaf nodes in tree
      for i in range(n):
          tree[n + i] = arr[i]
      # build the tree by calculating parents
      for i in range(n - 1, 0, -1):
          tree[i] = tree[i << 1] + tree[i << 1 | 1]
11
13 def updateTreeNode(p, value, n): # function to
      update a tree node
      # set value at position p
      tree[p + n] = value
      p = p + n
      i = p # move upward and update parents
      while i > 1:
          tree[i >> 1] = tree[i] + tree[i ^ 1]
          i >>= 1
23 def query(1, r, n): # function to get sum on
      interval [1, r)
      # loop to find the sum in the range
      r += n
      while 1 < r:
29
          if 1 & 1:
              res += tree[1]
              1 += 1
31
          if r & 1:
```

#### 2.6.2 Trie

```
class TrieNode:
      def init (self):
          self.children = [None] *26
          self.isEndOfWord = False
6 class Trie:
      def __init__(self):
          self.root = self.getNode()
      def getNode(self):
          return TrieNode()
11
12
      def charToIndex(self.ch):
13
          return ord(ch)-ord('a')
14
16
      def insert(self,key):
          pCrawl = self.root
          length = len(kev)
19
          for level in range(length):
              index = self._charToIndex(key[level])
21
              if not pCrawl.children[index]:
                  pCrawl.children[index] = self.
                      getNode()
              pCrawl = pCrawl.children[index]
          pCrawl.isEndOfWord = True
26
      def search(self, key):
27
          pCrawl = self.root
28
          length = len(key)
          for level in range(length):
              index = self. charToIndex(kev[level])
              if not pCrawl.children[index]:
                  return False
              pCrawl = pCrawl.children[index]
          return pCrawl.isEndOfWord
```

# 3 C++

# 3.1 Graphs

#### 3.1.1 BFS

```
1 #include "header.h"
2 #define graph unordered_map<11, unordered_set<11</pre>
3 vi bfs(int n, graph& g, vi& roots) {
      vi parents(n+1, -1); // nodes are 1..n
      unordered_set < int > visited;
      queue < int > q;
      for (auto x: roots) {
          q.emplace(x);
          visited.insert(x);
10
      while (not q.empty()) {
11
          int node = q.front();
          q.pop();
13
14
          for (auto neigh: g[node]) {
               if (not in(neigh, visited)) {
16
                   parents[neigh] = node;
17
                   q.emplace(neigh);
                   visited.insert(neigh);
              }
          7
21
22
      return parents;
23
24 }
25 vi reconstruct_path(vi parents, int start, int
      vi path;
26
      int curr = goal;
      while (curr != start) {
          path.push_back(curr);
29
          if (parents[curr] == -1) return vi(); //
               No path, empty vi
           curr = parents[curr];
32
33
      path.push_back(start);
      reverse(path.begin(), path.end());
34
      return path;
35
36 }
```

# **3.1.2 DFS** Cycle detection / removal

# 3.1.3 Dijkstra

```
1 #include "header.h"
2 vector < int > dijkstra(int n, int root, map < int,</pre>
      vector<pair<int, int>>>& g) {
    unordered_set <int> visited;
    vector<int> dist(n. INF):
      priority_queue < pair < int , int >> pq;
      dist[root] = 0:
      pq.push({0, root});
      while (!pq.empty()) {
           int node = pq.top().second;
          int d = -pq.top().first;
10
          pq.pop();
11
           if (in(node, visited)) continue;
13
           visited.insert(node):
15
           for (auto e : g[node]) {
16
               int neigh = e.first:
               int cost = e.second;
               if (dist[neigh] > dist[node] + cost)
                   dist[neigh] = dist[node] + cost;
                   pq.push({-dist[neigh], neigh});
21
22
      return dist;
25
26 }
```

#### 3.1.4 Floyd-Warshall

# ${\bf 3.1.5}$ Kruskal Minimum spanning tree of undirected weighted graph

```
1 #include "header.h"
2 #include "disjoint_set.h"
3 // O(E log E)
4 pair < set < pair < 11, 11 >> , 11 > kruskal (vector < tuple</pre>
       <11, 11, 11>>& edges, 11 n) {
       set < pair < 11 . 11 >> ans:
       11 cost = 0;
       sort(edges.begin(), edges.end());
       DisjointSet <11> fs(n):
10
       ll dist. i. i:
       for (auto edge: edges) {
           dist = get<0>(edge);
14
           i = get <1>(edge);
           j = get < 2 > (edge);
15
           if (fs.find_set(i) != fs.find_set(j)) {
               fs.union_sets(i, j);
                ans.insert({i, j});
19
                cost += dist;
20
21
22
       return pair < set < pair < 11, 11>>, 11> {ans, cost
24 }
```

#### 3.1.6 Hungarian algorithm

```
12 * Oreturn a vector of length J, with the j-th
       entry equaling the minimum cost
* to assign the first (j+1) jobs to distinct
15 template <class T> vector<T> hungarian(const
      vector < vector < T >> &C) {
      const int J = (int)size(C), W = (int)size(C
          [0]);
      assert(J <= W);
17
      // job[w] = job assigned to w-th worker, or
18
          -1 if no job assigned
      // note: a W-th worker was added for
19
          convenience
      vector < int > job(W + 1, -1);
      vector<T> ys(J), yt(W + 1); // potentials
^{21}
      // -yt[W] will equal the sum of all deltas
22
      vector <T> answers;
23
      const T inf = numeric_limits <T>::max();
24
      for (int j_cur = 0; j_cur < J; ++j_cur) { //</pre>
           assign j_cur-th job
26
          int w cur = W:
          job[w_cur] = j_cur;
          // min reduced cost over edges from Z to
              worker w
          vector <T> min_to(W + 1, inf);
          vector<int> prv(W + 1, -1); // previous
              worker on alternating path
          vector < bool > in_Z(W + 1);  // whether
              worker is in Z
          while (job[w_cur] != -1) { // runs at
              most j_cur + 1 times
              in_Z[w_cur] = true;
              const int j = job[w_cur];
              T delta = inf:
35
              int w next:
              for (int w = 0; w < W; ++w) {
                  if (!in Z[w]) {
                      if (ckmin(min_to[w], C[i][w]
                          - ys[j] - yt[w]))
                          prv[w] = w_cur;
                      if (ckmin(delta, min_to[w]))
                          w next = w:
                  }
              }
              // delta will always be non-negative,
              // except possibly during the first
                  time this loop runs
              // if any entries of C[j_cur] are
                  negative
              for (int w = 0; w \le W; ++w) {
                  if (in Z[w]) vs[iob[w]] += delta.
                       vt[w] -= delta;
                  else min_to[w] -= delta;
              }
```

# 3.1.7 Suc. shortest path Calculates max flow, min cost

1 #include "header.h"

```
2 // map<node, map<node, pair<cost, capacity>>>
3 #define graph unordered_map<int, unordered_map<</pre>
      int, pair<ld, int>>>
5 const ld infty = 1e601; // Change if necessary
6 ld fill(int n, vld& potential) { // Finds max
      flow, min cost
    priority_queue < pair < ld, int >> pq;
    vector < bool > visited(n+2, false);
    vi parent(n+2, 0);
    vld dist(n+2, infty);
    dist[0] = 0.1:
    pq.emplace(make_pair(0.1, 0));
    while (not pq.empty()) {
      int node = pq.top().second;
14
      pq.pop();
      if (visited[node]) continue:
      visited[node] = true;
      for (auto& x : g[node]) {
        int neigh = x.first;
19
        int capacity = x.second.second;
        ld cost = x.second.first:
        if (capacity and not visited[neigh]) {
22
          ld d = dist[node] + cost + potential[node
              1 - potential[neigh]:
          if (d + 1e-101 < dist[neigh]) {</pre>
             dist[neigh] = d;
             pq.emplace(make_pair(-d, neigh));
             parent[neigh] = node;
28
    }}}
    for (int i = 0; i < n+2; i++) {</pre>
      potential[i] = min(infty, potential[i] + dist
    if (not parent[n+1]) return infty;
    1d ans = 0.1;
    for (int x = n+1; x; x=parent[x]) {
      ans += g[parent[x]][x].first;
```

```
37          g[parent[x]][x].second--;
38          g[x][parent[x]].second++;
39          }
40          return ans;
41     }
```

#### 3.1.8 Bipartite check

```
1 #include "header.h"
2 int main() {
      int n:
      vvi adj(n);
      vi side(n. -1):
                        // will have 0's for one
          side 1's for other side
      bool is bipartite = true: // becomes false
          if not bipartite
      queue < int > q;
      for (int st = 0: st < n: ++st) {</pre>
          if (side[st] == -1) {
10
              q.push(st);
               side[st] = 0:
12
               while (!q.empty()) {
                   int v = q.front();
                   q.pop();
15
                   for (int u : adj[v]) {
                       if (side[u] == -1) {
                           side[u] = side[v] ^ 1;
                           q.push(u);
20
                       } else {
                           is_bipartite &= side[u]
                               != side[v]:
                       }
23 }}}}
```

#### 3.1.9 Find cycle directed

```
1 #include "header.h"
2 int n:
3 \text{ const int } mxN = 2e5+5;
4 vvi adj(mxN);
5 vector < char > color;
6 vi parent;
7 int cvcle start. cvcle end:
8 bool dfs(int v) {
       color[v] = 1;
       for (int u : adj[v]) {
           if (color[u] == 0) {
               parent[u] = v;
12
               if (dfs(u)) return true:
          } else if (color[u] == 1) {
               cycle_end = v;
15
               cycle_start = u;
```

```
return true:
          }
      color[v] = 2:
      return false;
22 }
23 void find_cycle() {
      color.assign(n, 0);
      parent.assign(n. -1):
      cycle_start = -1;
      for (int v = 0; v < n; v++) {
27
          if (color[v] == 0 && dfs(v))break:
29
      if (cycle_start == -1) {
30
          cout << "Acvclic" << endl;</pre>
32
          vector < int > cycle;
33
           cycle.push_back(cycle_start);
34
          for (int v = cycle_end; v != cycle_start;
               v = parent[v])
               cycle.push_back(v);
          cvcle.push back(cvcle start):
          reverse(cycle.begin(), cycle.end());
           cout << "Cycle_Found:_";
          for (int v : cycle) cout << v << "";</pre>
41
          cout << endl:
44 }
```

#### 3.1.10 Find cycle directed

```
1 #include "header.h"
2 int n:
3 \text{ const int } mxN = 2e5 + 5;
4 vvi adj(mxN);
5 vector < bool > visited:
6 vi parent:
7 int cycle_start, cycle_end;
8 bool dfs(int v, int par) { // passing vertex and
      its parent vertex
      visited[v] = true;
      for (int u : adi[v]) {
          if(u == par) continue; // skipping edge
              to parent vertex
          if (visited[u]) {
              cvcle_end = v;
              cycle_start = u;
              return true;
16
          parent[u] = v;
17
          if (dfs(u, parent[u]))
18
              return true:
```

```
return false:
22 }
23 void find_cycle() {
       visited.assign(n, false):
       parent.assign(n, -1);
      cycle_start = -1;
      for (int v = 0; v < n; v++) {
          if (!visited[v] && dfs(v, parent[v]))
28
      if (cvcle start == -1) {
30
           cout << "Acvclic" << endl:
           vector<int> cycle;
33
           cycle.push_back(cycle_start);
           for (int v = cycle_end; v != cycle_start;
               v = parent[v])
               cvcle.push_back(v);
           cycle.push_back(cycle_start);
37
           cout << "Cycle_Found: ";
           for (int v : cycle) cout << v << "";</pre>
           cout << endl:</pre>
      }
42 }
```

#### 3.1.11 Tarjan's SCC

```
1 #include "header.h"
3 struct Tarjan {
    vvi &edges;
    int V. counter = 0. C = 0:
    vi n, 1;
    vector < bool > vs;
    stack<int> st:
    Tarjan(vvi &e) : edges(e), V(e.size()), n(V,
        -1), l(V, -1), vs(V, false) {}
    void visit(int u. vi &com) {
     1[u] = n[u] = counter++;
      st.push(u):
      vs[u] = true;
      for (auto &&v : edges[u]) {
        if (n[v] == -1) visit(v, com);
15
        if (vs[v]) 1[u] = min(1[u], 1[v]);
16
17
      if (1[u] == n[u]) {
        while (true) {
          int v = st.top();
20
          st.pop();
          vs[v] = false;
          com[v] = C: // <== ACT HERE
          if (u == v) break;
        }
        C++:
```

```
}
28
    int find_sccs(vi &com) { // component indices
        will be stored in 'com'
      com.assign(V, -1);
      C = 0:
      for (int u = 0; u < V; ++u)
        if (n[u] == -1) visit(u, com);
      return C:
    // scc is a map of the original vertices of the
         graph to the vertices
    // of the SCC graph, scc_graph is its adjacency
    // SCC indices and edges are stored in 'scc'
        and 'scc_graph'.
    void scc_collapse(vi &scc, vvi &scc_graph) {
      find sccs(scc):
      scc_graph.assign(C, vi());
      set <pi>rec; // recorded edges
      for (int u = 0; u < V; ++u) {
       assert(scc[u] != -1):
       for (int v : edges[u]) {
          if (scc[v] == scc[u] ||
           rec.find({scc[u], scc[v]}) != rec.end()
                ) continue;
          scc_graph[scc[u]].push_back(scc[v]);
          rec.insert({scc[u]. scc[v]}):
      }
51
    // Function to find sources and sinks in the
        SCC graph
    // The number of edges needed to be added is
        max(sources.size(), sinks.())
    void findSourcesAndSinks(const vvi &scc_graph,
        vi &sources, vi &sinks) {
      vi in_degree(C, 0), out_degree(C, 0);
      for (int u = 0; u < C; u++) {
        for (auto v : scc_graph[u]) {
          in_degree[v]++;
          out_degree[u]++;
      for (int i = 0; i < C; ++i) {</pre>
        if (in_degree[i] == 0) sources.push_back(i)
        if (out_degree[i] == 0) sinks.push_back(i);
  }
68 }:
```

**3.1.12** SCC edges Prints out the missing edges to make the input digraph strongly connected

```
1 #include "header.h"
2 const int N=1e5+10;
3 int n,a[N],cnt[N],vis[N];
4 vector<int> hd,tl;
5 int dfs(int x){
      vis[x]=1:
      if(!vis[a[x]])return vis[x]=dfs(a[x]);
       return vis[x]=x:
9 }
10 int main(){
       scanf("%d",&n);
      for(int i=1;i<=n;i++){</pre>
           scanf("%d",&a[i]);
           cnt[a[i]]++:
14
15
      int k=0;
      for(int i=1:i<=n:i++){</pre>
17
           if(!cnt[i]){
               k++:
19
               hd.push_back(i);
               tl.push_back(dfs(i));
21
           }
22
      }
23
      int tk=k:
      for(int i=1:i<=n:i++){</pre>
25
           if(!vis[i]){
               k++:
27
               hd.push_back(i);
28
               tl.push_back(dfs(i));
          }
30
31
      if(k==1&&!tk)k=0:
32
      printf("%d\n",k);
33
      for (int i=0; i < k; i++) printf ("%du%d\n", tl[i], hd
           [(i+1)%k]);
35
      return 0;
36 }
```

#### 3.1.13 Find Bridges

```
1 #include "header.h"
2 int n; // number of nodes
3 vvi adj; // adjacency list of graph
4 vector < bool > visited;
5 vi tin. low:
6 int timer;
7 void dfs(int v, int p = -1) {
      visited[v] = true:
      tin[v] = low[v] = timer++;
      for (int to : adj[v]) {
          if (to == p) continue;
          if (visited[to]) {
12
              low[v] = min(low[v], tin[to]);
13
          } else {
```

```
dfs(to, v):
               low[v] = min(low[v], low[to]);
               if (low[to] > tin[v])
17
                   IS BRIDGE(v. to):
          }
      7
21 }
22 void find_bridges() {
       timer = 0:
       visited.assign(n, false);
      tin.assign(n, -1);
      low.assign(n, -1);
      for (int i = 0; i < n; ++i) {</pre>
          if (!visited[i]) dfs(i);
28
29
30 }
```

# **3.1.14** Artic. points (i.e. cut off points)

```
1 #include "header.h"
2 int n: // number of nodes
3 vvi adj; // adjacency list of graph
4 vector < bool > visited;
5 vi tin. low:
6 int timer;
7 void dfs(int v, int p = -1) {
      visited[v] = true:
      tin[v] = low[v] = timer++;
      int children=0:
      for (int to : adj[v]) {
          if (to == p) continue;
          if (visited[to]) {
13
               low[v] = min(low[v], tin[to]);
          } else {
               dfs(to, v);
               low[v] = min(low[v], low[to]);
               if (low[to] >= tin[v] && p!=-1)
                   IS_CUTPOINT(v);
               ++children;
19
          }
20
21
      if(p == -1 && children > 1)
          IS CUTPOINT(v):
24 }
25 void find cutpoints() {
       timer = 0:
       visited.assign(n, false);
27
      tin.assign(n, -1);
28
       low.assign(n, -1);
      for (int i = 0; i < n; ++i) {
          if (!visited[i]) dfs (i):
31
32
33 }
```

#### 3.1.15 Topological sort

```
1 #include "header.h"
2 int n; // number of vertices
3 vvi adj; // adjacency list of graph
4 vector < bool > visited;
5 vi ans:
6 void dfs(int v) {
      visited[v] = true;
      for (int u : adj[v]) {
          if (!visited[u]) dfs(u);
       ans.push_back(v);
12 }
13 void topological sort() {
       visited.assign(n, false);
15
       ans.clear();
       for (int i = 0: i < n: ++i) {</pre>
           if (!visited[i]) dfs(i);
17
18
      reverse(ans.begin(), ans.end());
20 }
```

# **3.1.16 Bellmann-Ford** Same as Dijkstra but allows neg. edges

```
1 #include "header.h"
2 // Switch vi and vvpi to vl and vvpl if necessary
3 void bellmann_ford_extended(vvpi &e, int source,
      vi &dist. vb &cvc) {
    dist.assign(e.size(), INF);
    cyc.assign(e.size(), false); // true when u is
        in a <0 cycle
    dist[source] = 0;
    for (int iter = 0: iter < e.size() - 1: ++iter)</pre>
      bool relax = false:
      for (int u = 0: u < e.size(): ++u)
        if (dist[u] == INF) continue;
        else for (auto &e : e[u])
          if(dist[u]+e.second < dist[e.first])</pre>
            dist[e.first] = dist[u]+e.second, relax
                 = true:
      if(!relax) break;
15
    bool ch = true;
    while (ch) {
                         // keep going untill no
        more changes
      ch = false;
                         // set dist to -INF when in
      for (int u = 0; u < e.size(); ++u)</pre>
        if (dist[u] == INF) continue;
        else for (auto &e : e[u])
          if (dist[e.first] > dist[u] + e.second
```

### 3.1.17 Ford-Fulkerson Basic Max. flow

```
1 #include "header.h"
2 #define V 6 // Num. of vertices in given graph
4 /* Returns true if there is a path from source 's
5 't' in residual graph. Also fills parent[] to
      store the
6 path */
7 bool bfs(int rGraph[V][V], int s, int t, int
      parent[]) {
  bool visited[V]:
    memset(visited, 0, sizeof(visited));
    queue < int > q;
   q.push(s);
    visited[s] = true;
    parent[s] = -1;
    // Standard BFS Loop
    while (!q.empty()) {
      int u = q.front();
17
      a.pop():
18
      for (int v = 0; v < V; v++) {
        if (visited[v] == false && rGraph[u][v] >
21
            0) {
          if (v == t) {
            parent[v] = u:
            return true;
25
          q.push(v);
          parent[v] = u;
          visited[v] = true;
    return false;
33 }
35 // Returns the maximum flow from s to t in the
      given graph
36 int fordFulkerson(int graph[V][V], int s, int t)
    int u, v;
```

```
int rGraph[V]
        [V];
    for (u = 0: u < V: u++)
     for (v = 0: v < V: v++)
        rGraph[u][v] = graph[u][v];
43
    int parent[V]; // This array is filled by BFS
        and to
          // store path
    int max_flow = 0; // There is no flow initially
    while (bfs(rGraph, s, t, parent)) {
      int path flow = INT MAX:
      for (v = t; v != s; v = parent[v]) {
        u = parent[v]:
        path_flow = min(path_flow, rGraph[u][v]);
52
      for (v = t; v != s; v = parent[v]) {
        u = parent[v]:
55
        rGraph[u][v] -= path_flow;
        rGraph[v][u] += path_flow;
57
      max_flow += path_flow;
    return max flow:
62 }
```

# 3.1.18 Dinic max flow $O(V^2E)$ , O(Ef)

```
2 using F = 11; using W = 11; // types for flow and
       weight/cost
3 struct Sf
      const int v;
                            // neighbour
      const int r;
                      // index of the reverse edge
      F f:
                      // current flow
                      // capacity
      const F cap;
      const W cost; // unit cost
      S(int v. int ri. F c. W cost = 0):
          v(v), r(ri), f(0), cap(c), cost(cost) {}
      inline F res() const { return cap - f; }
11
12 };
13 struct FlowGraph : vector < vector < S >> {
      FlowGraph(size t n) : vector < vector < S >> (n) {}
      void add_edge(int u, int v, F c, W cost = 0){
           auto &t = *this:
          t[u].emplace_back(v, t[v].size(), c, cost
          t[v].emplace_back(u, t[u].size()-1, c, -
17
              cost);
      void add arc(int u, int v, F c, W cost = 0){
19
          auto &t = *this:
          t[u].emplace_back(v, t[v].size(), c, cost
              ):
```

```
t[v].emplace_back(u, t[u].size()-1, 0, -
               cost);
22
      void clear() { for (auto &E : *this) for (
23
          auto &e : E) e.f = OLL; }
24 };
25 struct Dinic{
      FlowGraph & edges; int V,s,t;
      vi 1: vector < vector < S > :: iterator > its: //
          levels and iterators
      Dinic(FlowGraph &edges, int s, int t) :
           edges(edges), V(edges.size()), s(s), t(t)
              , 1(V,-1), its(V) {}
      ll augment(int u. F c) { // we reuse the same
           iterators
          if (u == t) return c; ll r = OLL;
           for(auto &i = its[u]; i != edges[u].end()
               : i++){
               auto &e = *i:
               if (e.res() && l[u] < l[e.v]) {
                   auto d = augment(e.v, min(c, e.
                       res())):
                  if (d > 0) { e.f += d; edges[e.v
                      ][e.r].f -= d; c -= d;
                       r += d: if (!c) break: }
          } }
          return r:
      }
      11 run() {
41
          11 \text{ flow} = 0. \text{ f}:
           while(true) {
               fill(1.begin(), 1.end(),-1); 1[s]=0;
                   // recalculate the lavers
               queue < int > q; q.push(s);
               while(!a.emptv()){
                   auto u = q.front(); q.pop(); its[
                       u] = edges[u].begin();
                   for(auto &&e : edges[u]) if(e.res
                       () && 1[e.v]<0)
                       l[e.v] = l[u]+1, q.push(e.v);
               if (1[t] < 0) return flow;</pre>
               while ((f = augment(s. INF)) > 0)
                   flow += f:
          }
54 }:
```

# 3.2 Dynamic Programming

#### 3.2.1 Longest Incr. Subseq.

```
#include "header.h"
template < class T>
vector < T > index path lis(vector < T > & nums) f
```

```
int n = nums.size();
    vector <T> sub;
      vector < int > subIndex:
    vector <T > path(n. -1):
    for (int i = 0; i < n; ++i) {</pre>
        if (sub.empty() || sub[sub.size() - 1] <</pre>
            nums[i]) {
      path[i] = sub.empty() ? -1 : subIndex[sub.
          size() - 1]:
      sub.push_back(nums[i]);
      subIndex.push back(i):
       } else {
      int idx = lower_bound(sub.begin(), sub.end(),
           nums[i]) - sub.begin();
      path[i] = idx == 0 ? -1 : subIndex[idx - 1];
      sub[idx] = nums[i];
      subIndex[idx] = i:
    vector <T> ans:
    int t = subIndex[subIndex.size() - 1];
    while (t != -1) {
        ans.push_back(t);
        t = path[t];
    reverse(ans.begin(), ans.end());
    return ans:
29 // Length only
30 template < class T>
31 int length_lis(vector<T> &a) {
    set <T> st:
    tvpename set<T>::iterator it:
    for (int i = 0; i < a.size(); ++i) {</pre>
    it = st.lower bound(a[i]):
      if (it != st.end()) st.erase(it);
      st.insert(a[i]);
   return st.size();
```

#### **3.2.2 0-1** Knapsack

# **3.2.3 Coin change** Number of coins required to achieve a given value

```
1 #include "header.h"
2 // Returns total distinct ways to make sum using
      n coins of
3 // different denominations
4 int count(vi& coins, int n, int sum) {
      // 2d dp array where n is the number of coin
      // denominations and sum is the target sum
      vector < vector < int > > dp(n + 1, vector < int > (
         sum + 1, 0));
      dp[0][0] = 1:
      for (int i = 1; i <= n; i++) {</pre>
         for (int i = 0: i <= sum: i++) {
              // without using the current coin,
              dp[i][j] += dp[i - 1][j];
              // using the current coin
              if ((j - coins[i - 1]) >= 0)
                   dp[i][j] += dp[i][j - coins[i -
                      1]]:
          }
      return dp[n][sum];
21 }
```

# 3.3 Trees

#### 3.3.1 Tree diameter

```
#include "header.h"
const int mxN = 2e5 + 5;
int n, d[mxN]; // distance array
vi adj[mxN]; // tree adjacency list
void dfs(int s, int e) {
d[s] = 1 + d[e]; // recursively calculate the distance from the starting node to each node
for (auto u : adj[s]) { // for each adjacent node

if (u != e) dfs(u, s); // don't move backwards in the tree
```

#### 3.3.2 Tree Node Count

#### 3.4 Numerical

#### 3.4.1 Template (for this section)

```
#include <bits/stdc++.h>
using namespace std;
#define rep(i, a, b) for(int i = a; i < (b); ++i)
#define all(x) begin(x), end(x)
#define sz(x) (int)(x).size()
typedef long long ll;
typedef pair<int, int> pii;
typedef vector<int> vi;
```

#### 3.4.2 Polynomial

```
#include "template.cpp"

struct Poly {
    vector<double> a;
    double operator()(double x) const {
```

```
double val = 0;
for (int i = sz(a); i--;) (val *= x) += a[i];
return val;

void diff() {
   rep(i,1,sz(a)) a[i-1] = i*a[i];
   a.pop_back();
}

void divroot(double x0) {
   double b = a.back(), c; a.back() = 0;
   for(int i=sz(a)-1; i--;) c = a[i], a[i] = a[i +1]*x0+b, b=c;
   a.pop_back();
}

a.pop_back();
}
};
```

# 3.4.3 Poly Roots

```
2 * Description: Finds the real roots to a
       polynomial.
3 * Usage: polyRoots({{2,-3,1}},-1e9,1e9) // solve
        x^2-3x+2 = 0
4 * Time: O(n^2 \log(1/\epsilon))
6 #include "Polynomial.h"
7 #include "template.cpp"
9 vector < double > polyRoots(Poly p, double xmin,
      double xmax) {
    if (sz(p.a) == 2) { return {-p.a[0]/p.a[1]}; }
    vector < double > ret;
12 Polv der = p:
    auto dr = polyRoots(der, xmin, xmax);
    dr.push_back(xmin-1);
    dr.push_back(xmax+1);
    sort(all(dr)):
17
    rep(i,0,sz(dr)-1) {
      double 1 = dr[i], h = dr[i+1];
      bool sign = p(1) > 0;
      if (sign ^(p(h) > 0)) {
21
        rep(it,0,60) { // while (h - 1 > 1e-8)
          double m = (1 + h) / 2, f = p(m);
          if ((f <= 0) ^ sign) l = m;</pre>
          else h = m;
25
        ret.push_back((1 + h) / 2);
27
29
    return ret:
```

#### 3.4.4 Golden Section Search

```
1 /**
2 * Description: Finds the argument minimizing the
        function $f$ in the interval $[a,b]$
3 * assuming $f$ is unimodal on the interval, i.e.
        has only one local minimum and no local
* maximum. The maximum error in the result is
       $eps$. Works equally well for maximization
5 * with a small change in the code. See
       TernarySearch.h in the Various chapter for a
6 * discrete version.
7 * Usage:
    double func(double x) { return 4+x+.3*x*x; }
    double xmin = gss(-1000.1000.func):
* Time: O(\log((b-a) / \epsilon))
12 #include "template.cpp"
14 /// It is important for r to be precise,
      otherwise we don't necessarily maintain the
      inequality a < x1 < x2 < b.
15 double gss(double a, double b, double (*f)(double
      )) {
    double r = (sqrt(5)-1)/2, eps = 1e-7;
    double x1 = b - r*(b-a), x2 = a + r*(b-a);
    double f1 = f(x1), f2 = f(x2);
    while (b-a > eps)
     if (f1 < f2) { //change to > to find maximum
        b = x2; x2 = x1; f2 = f1;
        x1 = b - r*(b-a): f1 = f(x1):
      } else {
        a = x1; x1 = x2; f1 = f2;
        x2 = a + r*(b-a); f2 = f(x2);
     }
27
    return a;
```

### 3.4.5 Hill Climbing

```
p[1] += dy*jmp;
cur = min(cur, make_pair(f(p), p));
}
return cur;
}
```

#### 3.4.6 Integration

```
2 * Description: Simple integration of a function
       over an interval using
3 * Simpson's rule. The error should be
       proportional to $h^4$, although in
4 * practice you will want to verify that the
       result is stable to desired
5 * precision when epsilon changes.
6 */
7 #include "template.cpp"
9 template < class F>
10 double quad(double a, double b, F f, const int n
      = 1000) {
   double h = (b - a) / 2 / n, v = f(a) + f(b);
  rep(i,1,n*2)
     v += f(a + i*h) * (i&1 ? 4 : 2);
  return v * h / 3:
```

#### 3.4.7 Integration Adaptive

```
2 * Description: Fast integration using an
       adaptive Simpson's rule.
    double sphereVolume = quad(-1, 1, [](double x)
    return quad(-1, 1, [\&](double y) {
    return quad(-1, 1, [\&](double z) {
    return x*x + y*y + z*z < 1; });});});
   * Status: mostly untested
10 #include "template.cpp"
12 typedef double d;
13 #define S(a,b) (f(a) + 4*f((a+b) / 2) + f(b)) * (
      b-a) / 6
15 template <class F>
16 d rec(F& f, d a, d b, d eps, d S) {
d c = (a + b) / 2;
d S1 = S(a, c), S2 = S(c, b), T = S1 + S2;
if (abs(T - S) \le 15 * eps | | b - a < 1e-10)
```

# 3.5 Num. Th. / Comb.

#### 3.5.1 Basic stuff

```
1 #include "header.h"
2 11 gcd(11 a, 11 b) { while (b) { a %= b; swap(a,
      b): } return a: }
3 11 1cm(11 a, 11 b) { return (a / gcd(a, b)) * b;
4 ll mod(ll a, ll b) { return ((a % b) + b) % b: }
_5 // Finds x, y s.t. ax + by = d = gcd(a, b).
6 void extended_euclid(ll a, ll b, ll &x, ll &y, ll
   11 xx = y = 0;
   11 yy = x = 1;
   while (b) {
     ll q = a / b;
    ll t = b; b = a % b; a = t;
      t = xx; xx = x - q * xx; x = t;
      t = yy; yy = y - q * yy; y = t;
    solves ab = 1 \pmod{n}, -1 on failure
18 ll mod_inverse(ll a, ll n) {
    11 x, y, d;
    extended_euclid(a, n, x, y, d);
    return (d > 1 ? -1 : mod(x, n));
_{23} // All modular inverses of [1..n] mod P in O(n)
24 vi inverses(ll n, ll P) {
    vi I(n+1, 1LL);
   for (ll i = 2: i <= n: ++i)
      I[i] = mod(-(P/i) * I[P\%i], P);
   return I:
29 }
30 // (a*b)%m
31 ll mulmod(ll a, ll b, ll m){
   11 x = 0, y=a%m;
    while(b>0){
     if(b\&1) x = (x+v)\%m:
      y = (2*y)%m, b /= 2;
   return x % m:
```

```
39 // Finds b^e % m in O(lg n) time, ensure that b <
       m to avoid overflow!
40 ll powmod(ll b. ll e. ll m) {
   11 p = e<2 ? 1 : powmod((b*b)\%m,e/2,m);
  return e&1 ? p*b%m : p;
43 }
44 // Solve ax + by = c, returns false on failure.
45 bool linear_diophantine(ll a, ll b, ll c, ll &x,
      11 &v) {
   11 d = gcd(a, b);
    if (c % d) {
     return false;
   } else {
      x = c / d * mod_inverse(a / d, b / d);
      y = (c - a * x) / b;
     return true:
54 }
```

# **3.5.2** Mod. exponentiation Or use pow() in python

```
#include "header.h"
2 ll mod_pow(ll base, ll exp, ll mod) {
3    if (mod == 1) return 0;
4     if (exp == 0) return 1;
5    if (exp == 1) return base;
6
7    ll res = 1;
8    base %= mod;
9    while (exp) {
10        if (exp % 2 == 1) res = (res * base) % mod;
11        exp >>= 1;
12        base = (base * base) % mod;
13    }
14
15    return res % mod;
16 }
```

#### **3.5.3** GCD Or math.gcd in python, std::gcd in C++

```
#include "header.h"
2 ll gcd(ll a, ll b) {
3    if (a == 0) return b;
4    return gcd(b % a, a);
5 }
```

#### 3.5.4 Sieve of Eratosthenes

```
1 #include "header.h"
2 vl primes;
```

#### 3.5.5 Fibonacci % prime

#### 3.5.6 nCk % prime

#### 3.5.7 Chin. rem. th.

```
#include "header.h"

#include "elementary.cpp"

3 // Solves x = a1 mod m1, x = a2 mod m2, x is
        unique modulo lcm(m1, m2).

4 // Returns {0, -1} on failure, {x, lcm(m1, m2)}
        otherwise.
```

```
5 pair<11, 11> crt(11 a1, 11 m1, 11 a2, 11 m2) {
6 ll s, t, d;
    extended_euclid(m1, m2, s, t, d);
   if (a1 % d != a2 % d) return {0, -1}:
    return {mod(s*a2 %m2 * m1 + t*a1 %m1 * m2, m1 *
         m2) / d, m1 / d * m2};
10 }
12 // Solves x = ai mod mi. x is unique modulo lcm
13 // Returns {0, -1} on failure, {x, lcm mi}
      otherwise.
14 pair<11, 11> crt(vector<11> &a, vector<11> &m) {
    pair<11, 11> res = \{a[0], m[0]\};
    for (ull i = 1; i < a.size(); ++i) {</pre>
      res = crt(res.first, res.second, mod(a[i], m[
          il), m[il):
      if (res.second == -1) break;
20
    return res;
21 }
```

# 3.6 Strings

#### **3.6.1 Z** alg. KMP alternative

```
#include "../header.h"
void Z_algorithm(const string &s, vi &Z) {
    Z.assign(s.length(), -1);
    int L = 0, R = 0, n = s.length();
    for (int i = 1; i < n; ++i) {
        if (i > R) {
            L = R = i;
            while (R < n && s[R - L] == s[R]) R++;
            Z[i] = R - L; R--;
        } else if (Z[i - L] >= R - i + 1) {
            L = i;
            while (R < n && s[R - L] == s[R]) R++;
            Z[i] = R - L; R--;
        } else Z[i] = Z[i - L];
} else Z[i] = Z[i - L];
}</pre>
```

#### 3.6.2 KMP

```
while(k >= 0 && w[k + 1] != w[i]) k = prefix[
          k];
      if(w[k + 1] == w[i]) k++:
      prefix[i] = k:
   }
11 }
void knuth_morris_pratt(string &s, string &w) {
    int q = -1;
    vi prefix:
    compute_prefix_function(w, prefix);
    for(int i = 0; i < s.length(); ++i) {</pre>
      while (q >= 0 \&\& w[q + 1] != s[i]) q = prefix[
      if(w[q + 1] == s[i]) q++;
      if(q + 1 == w.length()) {
        // Match at position (i - w.length() + 1)
        q = prefix[q];
   }
23
24 }
```

# **3.6.3 Aho-Corasick** Also can be used as Knuth-Morris-Pratt algorithm

```
1 #include "header.h"
3 map < char, int > cti;
4 int cti_size;
5 template <int ALPHABET_SIZE, int (*mp)(char)>
6 struct AC FSM {
    struct Node {
      int child[ALPHABET_SIZE], failure = 0,
          match_par = -1;
      Node() { for (int i = 0; i < ALPHABET_SIZE;</pre>
           ++i) child[i] = -1; }
    }:
    vector < Node > a:
    vector<string> &words;
    AC FSM(vector < string > & words) : words(words) {
      a.push_back(Node());
      construct_automaton();
17
    void construct_automaton() {
      for (int w = 0, n = 0; w < words.size(); ++w.
        for (int i = 0; i < words[w].size(); ++i) {</pre>
          if (a[n].child[mp(words[w][i])] == -1) {
21
            a[n].child[mp(words[w][i])] = a.size();
22
             a.push_back(Node());
           n = a[n].child[mp(words[w][i])];
26
        a[n].match.push_back(w);
```

```
queue < int > q;
      for (int k = 0; k < ALPHABET_SIZE; ++k) {</pre>
        if (a[0].child[k] == -1) a[0].child[k] = 0:
        else if (a[0].child[k] > 0) {
          a[a[0].child[k]].failure = 0;
          q.push(a[0].child[k]);
35
      }
      while (!q.empty()) {
        int r = q.front(); q.pop();
        for (int k = 0, arck; k < ALPHABET_SIZE; ++</pre>
          if ((arck = a[r].child[k]) != -1) {
            q.push(arck);
            int v = a[r].failure;
            while (a[v].child[k] == -1) v = a[v].
                 failure:
            a[arck].failure = a[v].child[k];
            a[arck].match_par = a[v].child[k];
            while (a[arck].match_par != -1
                 && a[a[arck].match_par].match.empty
               a[arck].match_par = a[a[arck].
                   match_par].match_par;
      }
    }
52
    void aho_corasick(string &sentence, vvi &
        matches){
      matches.assign(words.size(), vi());
      int state = 0. ss = 0:
      for (int i = 0; i < sentence.length(); ++i,</pre>
          ss = state) {
        while (a[ss].child[mp(sentence[i])] == -1)
          ss = a[ss].failure;
        state = a[state].child[mp(sentence[i])]
            = a[ss].child[mp(sentence[i])];
        for (ss = state; ss != -1; ss = a[ss].
            match_par)
          for (int w : a[ss].match)
            matches[w].push_back(i + 1 - words[w].
                length()):
   }
67 int char_to_int(char c) {
    return cti[c]:
69 }
70 int main() {
    string line;
    while(getline(cin, line)) {
      stringstream ss(line);
```

```
ss >> n:
       vector < string > patterns(n);
77
       for (auto& p: patterns) getline(cin, p);
 79
       string text;
       getline(cin, text);
       cti = {}. cti size = 0:
       for (auto c: text) {
         if (not in(c, cti)) {
           cti[c] = cti size++:
         }
       for (auto& p: patterns) {
         for (auto c: p) {
           if (not in(c, cti)) {
             cti[c] = cti_size++;
           }
         }
       }
95
       vvi matches;
       AC_FSM <128+1, char_to_int > ac_fms(patterns);
       ac fms.aho corasick(text. matches):
       for (auto& x: matches) cout << x << endl;</pre>
100
102
103 }
```

# **3.6.4** Long. palin. subs Manacher - O(n)

```
1 #include "header.h"
void manacher(string &s, vi &pal) {
    int n = s.length(), i = 1, 1, r;
    pal.assign(2 * n + 1, 0);
    while (i < 2 * n + 1) {</pre>
      if ((i&1) && pal[i] == 0) pal[i] = 1;
      l = i / 2 - pal[i] / 2; r = (i-1) / 2 + pal[i]
         1 / 2:
      while (1 - 1 >= 0 \&\& r + 1 < n \&\& s[1 - 1] ==
           s[r + 1]
        --1, ++r, pal[i] += 2;
11
      for (1 = i - 1, r = i + 1; 1 >= 0 && r < 2 *
          n + 1; --1, ++r) {
        if (1 <= i - pal[i]) break:</pre>
        if (1 / 2 - pal[1] / 2 > i / 2 - pal[i] /
          pal[r] = pal[1]:
        else { if (1 \ge 0)
            pal[r] = min(pal[l], i + pal[i] - r);
17
          break:
```

```
19 }
20 }
21 i = r;
22 } }
```

# 3.7 Geometry

# 3.7.1 essentials.cpp #include "../header.h"

```
2 using C = ld; // could be long long or long
3 constexpr C EPS = 1e-10; // change to 0 for C=11
4 struct P { // may also be used as a 2D vector
5 C x, y;
6 P(C \times = 0, C \times = 0) : x(x), y(y) {}
7 P operator+ (const P &p) const { return {x + p.
        x, y + p.y; }
   P operator - (const P &p) const { return {x - p.
        x, y - p.y; }
   P operator* (C c) const { return {x * c, y * c
   P operator/ (C c) const { return {x / c, y / c
    C operator* (const P &p) const { return x*p.x +
    C operator (const P &p) const { return x*p.v -
         p.x*v; }
    P perp() const { return P{y, -x}; }
    C lensq() const { return x*x + y*y; }
    ld len() const { return sqrt((ld)lensq()); }
    static ld dist(const P &p1, const P &p2) {
      return (p1-p2).len(); }
    bool operator == (const P &r) const {
      return ((*this)-r).lensq() <= EPS*EPS; }</pre>
21 C det(P p1, P p2) { return p1^p2; }
22 C det(P p1, P p2, P o) { return det(p1-o, p2-o):
23 C det(const vector <P> &ps) {
24    C sum = 0;    P prev = ps.back();
    for(auto &p : ps) sum += det(p, prev), prev = p
    return sum;
28 // Careful with division by two and C=11
29 C area(P p1, P p2, P p3) { return abs(det(p1, p2,
       p3))/C(2): }
30 C area(const vector <P> &poly) { return abs(det(
      poly))/C(2); }
31 int sign(C c) { return (c > C(0)) - (c < C(0)): }
32 int ccw(P p1, P p2, P o) { return sign(det(p1, p2
      . o)); }
```

#### 3.7.2 Two segs. itersec.

```
#include "header.h"

#include "essentials.cpp"

bool intersect(P a1, P a2, P b1, P b2) {

if (max(a1.x, a2.x) < min(b1.x, b2.x)) return
    false;

if (max(b1.x, b2.x) < min(a1.x, a2.x)) return
    false;

if (max(a1.y, a2.y) < min(b1.y, b2.y)) return
    false;

if (max(b1.y, b2.y) < min(a1.y, a2.y)) return
    false;

bool 11 = ccw(a2, b1, a1) * ccw(a2, b2, a1) <=
    0;

bool 12 = ccw(b2, a1, b1) * ccw(b2, a2, b1) <=
    0;

return 11 && 12;

11 }</pre>
```

#### 3.7.3 Convex Hull

```
1 #include "header.h"
2 #include "essentials.cpp"
3 struct ConvexHull { // O(n lg n) monotone chain.
    vector < size_t > h, c; // Indices of the hull
        are in 'h', ccw.
  const vector <P> &p;
   ConvexHull(const vector <P> &_p) : n(_p.size()),
         c(n), p(_p) {
      std::iota(c.begin(), c.end(), 0);
      std::sort(c.begin(), c.end(), [this](size_t 1
          , size_t r) -> bool { return p[1].x != p[
          r].x ? p[1].x < p[r].x : p[1].y < p[r].y;
      c.erase(std::unique(c.begin(), c.end(), [this
          ](size_t 1, size_t r) { return p[1] == p[
          rl: }). c.end()):
      for (size_t s = 1, r = 0; r < 2; ++r, s = h.
          size()) {
        for (size t i : c) {
          while (h.size() > s && ccw(p[h.end()
              [-2], p[h.end()[-1]], p[i]) <= 0)
            h.pop_back();
          h.push_back(i);
16
        reverse(c.begin(), c.end());
```

```
if (h.size() > 1) h.pop_back();
    size t size() const { return h.size(); }
    template <class T, void U(const P &, const P &,
         const P &, T &)>
    void rotating_calipers(T &ans) {
      if (size() <= 2)
        U(p[h[0]], p[h.back()], p[h.back()], ans);
        for (size t i = 0, i = 1, s = size(): i < 2
             * s: ++i) {
          while (det(p[h[(i + 1) % s]] - p[h[i % s
              ]], p[h[(j + 1) \% s]] - p[h[j]]) >=
            j = (j + 1) \% s;
          U(p[h[i % s]], p[h[(i + 1) % s]], p[h[j
              ]], ans);
33 };
     Example: furthest pair of points. Now set ans
      = OLL and call
35 // ConvexHull(pts).rotating_calipers<11, update>(
36 void update(const P &p1, const P &p2, const P &o,
       ll &ans) {
    ans = max(ans, (11)max((p1 - o).lensq(), (p2 -
        o).lensq()));
38 }
39 int main() {
    ios::sync_with_stdio(false); // do not use
        cout + printf
    cin.tie(NULL);
    int n:
    cin >> n;
    while (n) {
      vector <P> ps;
          int x, y;
47
      for (int i = 0; i < n; i++) {</pre>
              cin >> x >> y;
              ps.push_back({x, y});
          }
          ConvexHull ch(ps):
          cout << ch.h.size() << endl;</pre>
54
          for(auto& p: ch.h) {
55
              cout << ps[p].x << "" << ps[p].y <<
                  endl;
      cin >> n:
59
    return 0;
```

# 3.8 Other Algorithms

#### 3.8.1 2-sat

62 }

```
1 #include "../header.h"
#include "../Graphs/tarjan.cpp"
3 struct TwoSAT {
4 int n:
    vvi imp; // implication graph
    Tarjan tj;
    TwoSAT(int _n) : n(_n), imp(2 * _n, vi()), tj(
        imp) { }
   // Only copy the needed functions:
    void add_implies(int c1, bool v1, int c2, bool
      int u = 2 * c1 + (v1 ? 1 : 0),
       v = 2 * c2 + (v2 ? 1 : 0):
      imp[u].push_back(v); // u => v
      imp[v^1].push_back(u^1); // -v => -u
15
    void add_equivalence(int c1, bool v1, int c2,
        bool v2) {
      add implies(c1, v1, c2, v2):
      add_implies(c2, v2, c1, v1);
    void add_or(int c1, bool v1, int c2, bool v2) {
      add_implies(c1, !v1, c2, v2);
22
23
    void add_and(int c1, bool v1, int c2, bool v2)
      add_true(c1, v1); add_true(c2, v2);
    void add_xor(int c1, bool v1, int c2, bool v2)
      add_or(c1, v1, c2, v2);
      add or(c1, !v1, c2, !v2);
    void add_true(int c1, bool v1) {
      add_implies(c1, !v1, c1, v1);
    // on true: a contains an assignment.
    // on false: no assignment exists.
    bool solve(vb &a) {
      vi com;
      tj.find_sccs(com);
      for (int i = 0; i < n; ++i)
       if (com[2 * i] == com[2 * i + 1])
         return false;
42
```

```
vvi bycom(com.size());
for (int i = 0; i < 2 * n; ++i)
bycom[com[i]].push_back(i);

a.assign(n, false);
vb vis(n, false);
for (auto &&component : bycom){
    for (int u : component) {
        if (vis[u / 2]) continue;
        vis[u / 2] = true;
        a[u / 2] = (u % 2 == 1);
}

return true;
}
</pre>
```

#### 3.8.2 Matrix Solve

```
1 #include "header.h"
2 #define REP(i, n) for(auto i = decltype(n)(0): i
      < (n): i++)
3 using T = double;
4 constexpr T EPS = 1e-8;
5 template < int R, int C>
6 using M = array < array < T, C > , R >; // matrix
7 template < int R. int C>
8 T ReducedRowEchelonForm(M<R,C> &m, int rows) {
     // return the determinant
9 int r = 0; T det = 1;
                                     // MODIFIES
        the input
   for(int c = 0: c < rows && r < rows: c++) {
      int p = r;
      for(int i=r+1; i<rows; i++) if(abs(m[i][c]) >
           abs(m[p][c])) p=i;
      if(abs(m[p][c]) < EPS){ det = 0; continue; }</pre>
      swap(m[p], m[r]); det = -det;
      T s = 1.0 / m[r][c], t; det *= m[r][c];
      term in row 1
      REP(i,rows) if (i!=r){ t = m[i][c]; REP(j,C)
          m[i][j] -= t*m[r][j]; }
      ++r:
19
    return det:
21 }
22 bool error, inconst; // error => multiple or
      inconsistent
23 template <int R, int C> // Mx = a; M:R*R, v:R*C =>
24 M<R.C> solve(const M<R.R> &m. const M<R.C> &a.
     int rows){
25 M<R.R+C> q:
26 REP(r.rows){
```

```
REP(c,rows) q[r][c] = m[r][c];
      REP(c,C) q[r][R+c] = a[r][c];
    ReducedRowEchelonForm <R.R+C>(a.rows):
    M<R,C> sol; error = false, inconst = false;
    REP(c,C) for(auto j = rows-1; j >= 0; --j){
      T t=0; bool allzero=true;
      for (auto k = j+1; k < rows; ++k)
       t += q[j][k]*sol[k][c], allzero &= abs(q[i
           ][k]) < EPS;
      if(abs(q[i][i]) < EPS)
        error = true, inconst |= allzero && abs(q[j
           ][R+c]) > EPS;
      else sol[i][c] = (q[i][R+c] - t) / q[i][i];
         // usually q[i][i]=1
   return sol;
41 }
```

# 3.8.3 Matrix Exp.

```
1 #include "header.h"
2 #define ITERATE_MATRIX(w) for (int r = 0; r < (w)</pre>
                 for (int c = 0: c < (w): ++c)
4 template <class T, int N>
5 struct M {
    array <array <T, N>, N> m;
    M() \{ ITERATE_MATRIX(N) m[r][c] = 0; \}
    static M id() {
      M I; for (int i = 0; i < N; ++i) I.m[i][i] =
          1; return I;
    M operator*(const M &rhs) const {
      M out:
12
      ITERATE MATRIX(N) for (int i = 0: i < N: ++i)</pre>
          out.m[r][c] += m[r][i] * rhs.m[i][c];
14
      return out:
    M raise(ll n) const {
      if(n == 0) return id();
      if(n == 1) return *this;
      auto r = (*this**this).raise(n / 2);
      return (n%2 ? *this*r : r):
22 }
23 };
```

#### 3.8.4 Finite field For FFT

```
1 #include "header.h"
2 #include "../Number_Theory/elementary.cpp"
3 template<1l p,ll w> // prime, primitive root
```

```
4 struct Field { using T = Field; ll x; Field(ll x
      =0) : x\{x\} \{\}
   T operator+(T r) const { return {(x+r.x)%p}; }
   T operator - (T r) const { return \{(x-r,x+p)\%p\}:
    T operator*(T r) const { return {(x*r.x)%p}; }
    T operator/(T r) const { return (*this)*r.inv()
    T inv() const { return {mod inverse(x,p)}: }
    static T root(ll k) { assert( (p-1)%k==0 );
        // (p-1)%k == 0?
      auto r = powmod(w,(p-1)/abs(k),p);
                                                // k-
          th root of unity
      return k>=0 ? T{r} : T{r}.inv();
13
  bool zero() const { return x == OLL; }
15 }:
16 using F1 = Field < 1004535809,3 >;
17 using F2 = Field<1107296257,10>; // 1<<30 + 1<<25</pre>
18 using F3 = Field < 2281701377,3 >; // 1 < < 31 + 1 < < 27
```

# 3.8.5 Complex field For FFR

```
1 #include "header.h"
2 const double m_pi = M_PIf64x;
3 struct Complex { using T = Complex; double u, v;
4 Complex(double u=0, double v=0) : u{u}, v{v} {}}
   T operator+(T r) const { return {u+r.u, v+r.v};
   T operator-(T r) const { return {u-r.u, v-r.v};
   T operator*(T r) const { return {u*r.u - v*r.v,
         u*r.v + v*r.u}; }
   T operator/(T r) const {
      auto norm = r.u*r.u+r.v*r.v;
      return {(u*r.u + v*r.v)/norm. (v*r.u - u*r.v)
          /norm}:
    T operator*(double r) const { return T{u*r, v*r
  T operator/(double r) const { return T{u/r, v/r
        }: }
14  T inv() const { return T{1,0}/ *this; }
   T conj() const { return T{u, -v}; }
    static T root(ll k){ return {cos(2*m_pi/k), sin
        (2*m_pi/k); }
   bool zero() const { return max(abs(u), abs(v))
        < 1e-6; }
```

#### 3.8.6 FFT

```
1 #include "header.h"
2 #include "complex_field.cpp"
3 #include "fin field.cpp"
4 void brinc(int &x, int k) {
int i = k - 1, s = 1 << i;
7 if ((x & s) != s) {
      --i: s >>= 1:
      while (i >= 0 && ((x & s) == s))
      x = x &^{\sim} s, --i, s >>= 1;
      if (i >= 0) x |= s:
12 }
13 }
14 using T = Complex: // using T=F1.F2.F3
15 vector <T> roots;
16 void root_cache(int N) {
    if (N == (int)roots.size()) return;
    roots.assign(N, T{0});
    for (int i = 0: i < N: ++i)
      roots[i] = ((i\&-i) == i)
        ? T\{\cos(2.0*m_pi*i/N), \sin(2.0*m_pi*i/N)\}
        : roots[i&-i] * roots[i-(i&-i)];
24 void fft(vector<T> &A, int p, bool inv = false) {
_{25} int N = 1<<p:
  for(int i = 0, r = 0; i < N; ++i, brinc(r, p))
      if (i < r) swap(A[i], A[r]):
28 // Uncomment to precompute roots (for T=Complex)
      . Slower but more precise.
29 // root cache(N):
          , sh=p-1
                       , --sh
31 for (int m = 2: m <= N: m <<= 1) {
      T w, w_m = T::root(inv ? -m : m);
      for (int k = 0; k < N; k += m) {
        w = T\{1\}:
       for (int j = 0; j < m/2; ++ j) {
            T w = (!inv ? roots[j << sh] : roots[j <<
      shl.coni()):
          T t = w * A[k + j + m/2];
          A[k + j + m/2] = A[k + j] - t;
          A[k + j] = A[k + j] + t;
          w = w * w_m;
    if(inv){ T inverse = T(N).inv(); for(auto &x :
        A) x = x*inverse;
46 // convolution leaves A and B in frequency domain
47 // C may be equal to A or B for in-place
      convolution
48 void convolution(vector<T> &A, vector<T> &B,
      vector <T> &C) {
   int s = A.size() + B.size() - 1;
```

## 3.8.7 Polyn. inv. div.

```
1 #include "header.h"
2 #include "fft.cpp"
3 vector <T> &rev(vector <T> &A) { reverse(A.begin(),
       A.end()); return A; }
4 void copy_into(const vector <T> &A, vector <T> &B,
      size t n) {
    std::copy(A.begin(), A.begin()+min({n, A.size()
         , B.size()}), B.begin());
6 }
8 // Multiplicative inverse of A modulo x^n.
      Requires A[0] != 0!!
9 vector<T> inverse(const vector<T> &A, int n) {
     vector <T> Ai{A[0].inv()};
    for (int k = 0: (1<<k) < n: ++k) {
      vector <T> As (4 << k, T(0)), Ais (4 << k, T(0));
      copy_into(A, As, 2<<k); copy_into(Ai, Ais, Ai</pre>
          .size()):
      fft(As, k+2, false); fft(Ais, k+2, false);
14
      for (int i = 0; i < (4<<k); ++i) As[i] = As[i</pre>
          1*Ais[i]*Ais[i]:
      fft(As, k+2, true); Ai.resize(2<<k, {});</pre>
      for (int i = 0: i < (2 << k): ++i) Ai[i] = T(2)
17
            * Ai[i] - As[i];
    Ai.resize(n);
    return Ai;
21 }
22 // Polynomial division. Returns {Q, R} such that
      A = QB+R, deg R < deg B.
23 // Requires that the leading term of B is nonzero
24 pair < vector < T > , vector < T >> divmod(const vector < T >
       &A. const vector <T> &B) {
    size_t n = A.size()-1, m = B.size()-1;
    if (n < m) return {vector < T > (1, T(0)), A};
```

**3.8.8 Linear recurs.** Given a linear recurrence of the form

$$a_n = \sum_{i=0}^{k-1} c_i a_{n-i-1}$$

this code computes  $a_n$  in  $O(k \log k \log n)$  time.

```
1 #include "header.h"
2 #include "poly.cpp"
3 // x^k \mod f
4 vector<T> xmod(const vector<T> f, ll k) {
5 vector <T> r{T(1)};
    for (int b = 62; b >= 0; --b) {
      if (r.size() > 1)
        square_inplace(r), r = mod(r, f);
      if ((k>>b)&1) {
        r.insert(r.begin(), T(0));
        if (r.size() == f.size()) {
11
         T c = r.back() / f.back();
          for (size_t i = 0; i < f.size(); ++i)</pre>
            r[i] = r[i] - c * f[i]:
          r.pop_back();
16
      }
19
    return r:
_{21} // Given A[0,k) and C[0, k), computes the n-th
      term of:
_{22} // A[n] = \sum_i C[i] * A[n-i-1]
23 T nth_term(const vector <T > &A, const vector <T > &C
      . 11 n) {
```

```
int k = (int)A.size();
if (n < k) return A[n];

vector<T> f(k+1, T{1});
for (int i = 0; i < k; ++i)
    f[i] = T{-1} * C[k-i-1];

f = xmod(f, n);

T r = T{0};

for (int i = 0; i < k; ++i)
    r = r + f[i] * A[i];

return r;
}</pre>
```

## **3.8.9 Convolution** Precise up to 9e15

```
1 #include "header.h"
2 #include "fft.cpp"
3 void convolution mod(const vi &A. const vi &B. 11
       MOD, vi &C) {
4 int s = A.size() + B.size() - 1; ll m15 = (1LL
        <<15) -1LL:
   int q = 32 - \_builtin_clz(s-1), N=1 << q; //
         fails if s=1
    vector\langle T \rangle Ac(N), Bc(N), R1(N), R2(N);
    for (size_t i = 0; i < A.size(); ++i) Ac[i] = T</pre>
        {A[i]&m15, A[i]>>15};
   for (size_t i = 0; i < B.size(); ++i) Bc[i] = T</pre>
         {B[i]&m15, B[i]>>15};
    fft(Ac, q, false); fft(Bc, q, false);
    for (int i = 0, j = 0; i < N; ++i, j = (N-1)&(N
      T as = (Ac[i] + Ac[j].conj()) / 2;
      T = (Ac[i] - Ac[j].conj()) / T{0, 2};
      T bs = (Bc[i] + Bc[j].conj()) / 2;
      T b1 = (Bc[i] - Bc[i].coni()) / T{0, 2}:
      R1[i] = as*bs + al*bl*T{0,1}, R2[i] = as*bl +
           al*bs:
    }
16
    fft(R1, q, true); fft(R2, q, true);
    11 p15 = (1LL << 15) \% MOD, p30 = (1LL << 30) \% MOD; C.
        resize(s);
    for (int i = 0; i < s; ++i) {</pre>
      11 1 = 1 \text{lround}(R1[i].u), m = 1 \text{lround}(R2[i].u)
           , h = llround(R1[i].v);
      C[i] = (1 + m*p15 + h*p30) \% MOD;
22
```

**3.8.10** Partitions of n Finds all possible partitions of a number

```
1 #include "header.h"
```

```
void printArray(int p[], int n) {
    for (int i = 0; i < n; i++)</pre>
      cout << p[i] << "";
    cout << endl:
6 }
8 void printAllUniqueParts(int n) {
    int p[n]; // An array to store a partition
    int k = 0: // Index of last element in a
        partition
    p[k] = n: // Initialize first partition as
        number itself
    // This loop first prints current partition
        then generates next
    // partition. The loop stops when the current
        partition has all 1s
    while (true) {
      printArray(p, k + 1);
17
      // Find the rightmost non-one value in p[].
18
          Also, update the
      // rem_val so that we know how much value can
           be accommodated
      int rem val = 0:
      while (k >= 0 \&\& p[k] == 1) {
21
        rem_val += p[k];
        k--:
      }
^{24}
25
      // if k < 0, all the values are 1 so there
          are no more partitions
      if (k < 0) return:
27
28
      // Decrease the p[k] found above and adjust
          the rem val
      p[k]--;
30
      rem val++:
31
32
      // If rem_val is more, then the sorted order
          is violated. Divide
      // rem_val in different values of size p[k]
34
          and copy these values at
      // different positions after p[k]
      while (rem_val > p[k]) {
        p[k + 1] = p[k]:
        rem_val = rem_val - p[k];
39
        k++;
      7
41
      // Copy rem_val to next position and
          increment position
      p[k + 1] = rem_val;
      k++:
45
```

# 3.9 Other Data Structures

46 }

# **3.9.1** Disjoint set (i.e. union-find)

```
1 template <typename T>
2 class DisjointSet {
      typedef T * iterator;
      T *parent, n, *rank;
      public:
          // O(n), assumes nodes are [0, n)
          DisjointSet(T n) {
              this->parent = new T[n];
              this -> n = n;
9
              this->rank = new T[n]:
              for (T i = 0: i < n: i++) {
                  parent[i] = i:
13
                  rank[i] = 0:
14
          }
16
18
          // O(log n)
          T find_set(T x) {
              if (x == parent[x]) return x:
              return parent[x] = find_set(parent[x
21
                  ]);
          }
          // O(log n)
          void union_sets(T x, T y) {
25
              x = this->find_set(x);
26
              y = this->find_set(y);
              if (x == y) return;
              if (rank[x] < rank[y]) {</pre>
                  Tz = x;
                  x = y;
33
                  y = z;
              parent[v] = x;
              if (rank[x] == rank[y]) rank[x]++;
40 };
```

**3.9.2 Fenwick tree** (i.e. BIT) eff. update + prefix sum calc.

```
1 #include "header.h"
```

```
2 #define maxn 200010
3 int t,n,m,tree[maxn],p[maxn];

5 void update(int k, int z) {
6     while (k <= maxn) {
7          tree[k] += z;
8          k += k & (-k);
9     }
10 }

11     int sum(int k) {
13         int ans = 0;
14         while(k) {
15             ans += tree[k];
16             k -= k & (-k);
17     }
18     return ans;
19 }</pre>
```

#### 3.9.3 Fenwick2d tree

```
1 #include "header.h"
2 template <class T>
3 struct FenwickTree2D {
vector < vector <T> > tree;
    FenwickTree2D(int n) : n(n) { tree.assign(n +
        1, vector (T > (n + 1, 0));}
7  T query(int x1, int y1, int x2, int y2) {
      return query (x2, y2) + query (x1-1, y1-1) - query (x2
          ,v1-1)-query(x1-1,v2);
   T query(int x, int y) {
      T s = 0:
      for (int i = x: i > 0: i -= (i & (-i)))
      for (int j = v; j > 0; j = (j & (-j)))
          s += tree[i][i]:
15
      return s;
   }
16
   void update(int x, int y, T v) {
      for (int i = x; i <= n; i += (i & (-i)))
       for (int j = y; j <= n; j += (j & (-j)))
          tree[i][i] += v:
20
21 }
22 }:
```

#### 3.9.4 Trie

```
#include "header.h"
const int ALPHABET_SIZE = 26;
inline int mp(char c) { return c - 'a'; }

struct Node {
```

```
Node* ch[ALPHABET_SIZE];
    bool isleaf = false;
      for(int i = 0; i < ALPHABET_SIZE; ++i) ch[i]</pre>
          = nullptr;
11
    void insert(string &s, int i = 0) {
      if (i == s.length()) isleaf = true;
       else {
       int v = mp(s[i]);
        if (ch[v] == nullptr)
          ch[v] = new Node();
        ch[v] \rightarrow insert(s, i + 1);
    }
20
21
    bool contains(string &s, int i = 0) {
22
      if (i == s.length()) return isleaf;
      else {
        int v = mp(s[i]);
25
        if (ch[v] == nullptr) return false:
        else return ch[v]->contains(s, i + 1);
    }
29
    void cleanup() {
      for (int i = 0: i < ALPHABET SIZE: ++i)</pre>
        if (ch[i] != nullptr) {
          ch[i]->cleanup();
          delete ch[i];
        }
   }
38 };
```

**3.9.5 Treap** A binary tree whose nodes contain two values, a key and a priority, such that the key keeps the BST property

```
16 // Push data to children here
17 }
18 void merge(Node *&t, Node *1, Node *r) {
    propagate(1). propagate(r):
    if (!1)
               t = r:
    else if (!r) t = 1;
    else if (1->pr > r->pr)
         merge(1->r, 1->r, r), t = 1;
    else merge(r->1, 1, r->1), t = r:
    update(t);
26 }
27 void spliti(Node *t, Node *&1, Node *&r, int
       index) {
    propagate(t);
    if (!t) { l = r = nullptr; return; }
    int id = size(t->1);
    if (index <= id) // id \in [index, \infty), so</pre>
       spliti(t\rightarrow 1, 1, t\rightarrow 1, index), r = t;
       spliti(t->r, t->r, r, index - id), l = t;
    update(t);
37 void splitv(Node *t, Node *&1, Node *&r, 11 val)
    propagate(t);
    if (!t) { l = r = nullptr; return; }
    if (val \langle = t - \rangle v) // t - \rangle v \in [val, \inftv), so
         move it right
       splitv(t\rightarrow 1, 1, t\rightarrow 1, val), r = t;
       splitv(t->r, t->r, r, val), l = t;
   update(t):
46 void clean(Node *p) {
   if (p) { clean(p->1), clean(p->r); delete p; }
48 }
```

#### 3.9.6 Segment tree

#### 3.9.7 Lazy segment tree Uptimizes range updates

```
1 #include "../header.h"
2 using T=int; using U=int; using I=int;
     exclusive right bounds
3 T t id: U u id:
4 T op(T a, T b){ return a+b; }
5 void join(U &a, U b){ a+=b; }
6 void apply(T &t, U u, int x){ t+=x*u; }
7 T convert(const I &i) { return i; }
8 struct LazySegmentTree {
   struct Node { int 1, r, 1c, rc; T t; U u;
     Node(int 1, int r, T t=t_id):1(1),r(r),1c(-1)
         ,rc(-1),t(t),u(u_id){}
   };
   int N; vector < Node > tree; vector < I > & init;
   LazySegmentTree(vector <I > &init) : N(init.size
       ()), init(init){
     tree.reserve(2*N-1); tree.push_back({0,N});
         build(0, 0, N);
   void build(int i, int l, int r) { auto &n =
       tree[i];
     if (r > 1+1) \{ int m = (1+r)/2;
       .r}):
       build(n.lc.l.m): build(n.rc.m.r):
       n.t = op(tree[n.lc].t, tree[n.rc].t);
     } else n.t = convert(init[1]):
   void push(Node &n, U u){ apply(n.t, u, n.r-n.l)
       ; join(n.u,u); }
   void push(Node &n){push(tree[n.lc],n.u);push(
       tree[n.rc].n.u):n.u=u id:}
  T query(int 1, int r, int i = 0) { auto &n =
       tree[i]:
     if(r <= n.1 || n.r <= 1) return t id:
     if(1 <= n.1 && n.r <= r) return n.t;</pre>
     return push(n), op(query(1,r,n.lc),query(1,r,
         n.rc)):
  void update(int 1, int r, U u, int i = 0) {
       auto &n = tree[i]:
```

#### 3.9.8 Suffix tree

```
1 #include "../header.h"
2 using T = char;
3 using M = map<T,int>; // or array<T,</pre>
      ALPHABET SIZE >
4 using V = string:
                       // could be vector<T> as
      well
5 using It = V::const_iterator;
6 struct Node{
   It b, e; M edges; int link; // end is
    Node(It b, It e) : b(b), e(e), link(-1) {}
    int size() const { return e-b; }
10 };
11 struct SuffixTree{
    const V &s; vector < Node > t;
   int root,n,len,remainder,llink; It edge;
    SuffixTree(const V &s) : s(s) { build(); }
    int add node(It b. It e) { return t.push back({b
        ,e}), t.size()-1; }
    int add_node(It b){ return add_node(b,s.end());
    void link(int node){ if(llink) t[llink].link =
        node: llink = node: }
    void build(){
      len = remainder = 0; edge = s.begin();
      n = root = add_node(s.begin(), s.begin());
      for(auto i = s.begin(); i != s.end(); ++i){
21
        ++remainder; llink = 0;
        while (remainder) {
          if(len == 0) edge = i;
          if(t[n].edges[*edge] == 0){
              new leaf
            t[n].edges[*edge] = add_node(i); link(n
               );
          } else {
            auto x = t[n].edges[*edge]; // neXt
                node [with edge]
            if(len >= t[x].size()){ // walk to
                next node
              len -= t[x].size(); edge += t[x].size
                 (); n = x;
              continue:
            if(*(t[x].b + len) == *i){ // walk}
                along edge
```

```
++len; link(n); break;
                 // split edge
            auto split = add_node(t[x].b, t[x].b+
36
            t[n].edges[*edge] = split;
            t[x].b += len;
            t[split].edges[*i] = add_node(i);
            t[split].edges[*t[x].b] = x;
40
            link(split):
41
          --remainder:
          if(n == root && len > 0)
            --len, edge = i - remainder + 1;
          else n = t[n].link > 0 ? t[n].link : root
   }
50 };
```

#### 3.9.9 UnionFind

```
1 #include "header.h"
2 struct UnionFind {
    std::vector<int> par, rank, size;
    UnionFind(int n) : par(n), rank(n, 0), size(n,
        1), c(n) {
     for(int i = 0; i < n; ++i) par[i] = i;</pre>
   int find(int i) { return (par[i] == i ? i : (
        par[i] = find(par[i]))); }
    bool same(int i, int j) { return find(i) ==
        find(j); }
    int get_size(int i) { return size[find(i)]; }
    int count() { return c; }
    int merge(int i, int j) {
      if((i = find(i)) == (i = find(i))) return -1;
13
      if(rank[i] > rank[j]) swap(i, j);
      par[i] = j;
      size[j] += size[i];
      if(rank[i] == rank[j]) rank[j]++;
     return j;
20 }
21 };
```

# 4 Other Mathematics

# 4.1 Helpful functions

**4.1.1** Euler's Totient Function  $n = p_1^{k_1-1} \cdot (p_1-1) \cdot \dots \cdot p_r^{k_r-1} \cdot (p_r-1)$ , where  $p_1^{k_1} \cdot \dots \cdot p_r^{k_r}$  is the prime factorization of n.

```
1 # include "header.h"
2 ll phi(ll n) { // \Phi(n)
      ll ans = 1;
      for (11 i = 2: i*i <= n: i++) {
       if (n % i == 0) {
              ans *= i-1:
             n /= i:
              while (n % i == 0) {
                  ans *= i:
                 n /= i;
         }
      if (n > 1) ans *= n-1;
      return ans:
17 vi phis(int n) { // All \Phi(i) up to n
    vi phi(n + 1, OLL);
    iota(phi.begin(), phi.end(), OLL);
    for (11 i = 2LL: i <= n: ++i)
     if (phi[i] == i)
      for (11 j = i; j <= n; j += i)
          phi[j] -= phi[j] / i;
24 return phi;
```

Formulas  $\Phi(n)$  counts all numbers in  $1, \ldots, n-1$  coprime to n.  $a^{\varphi(n)} \equiv 1 \mod n$ , a and n are coprimes.  $\forall e > \log_2 m : n^e \mod m = n^{\Phi(m)+e \mod \Phi(m)} \mod m$ .  $\gcd(m,n) = 1 \Rightarrow \Phi(m \cdot n) = \Phi(m) \cdot \Phi(n)$ .

**4.1.2** Pascal's trinagle  $\binom{n}{k}$  is k-th element in the n-th row, indexing both from 0

```
#include "header.h"
void printPascal(int n) {
for (int line = 1; line <= n; line++) {
    int C = 1; // used to represent C(line, i
    )

for (int i = 1; i <= line; i++) {
    // The first value in a line is
    always 1</pre>
```

# 4.2 Theorems and definitions

#### Fermat's little theorem

$$a^p \equiv a \mod p$$

**Subfactorial** 

$$!n = n! \sum_{i=0}^{n} \frac{(-1)^{i}}{i!}$$

$$!(0) = 1, !n = n \cdot !(n-1) + (-1)^n$$

#### Binomials and other partitionings

$$\binom{n}{k} = \binom{n-1}{k} + \binom{n-1}{k-1} = \prod_{i=1}^{k} \frac{n-i+1}{i}$$

This last product may be computed incrementally since any product of k' consecutive values is divisible by k'!. Basic identities: The hockeystick identity:

$$\sum_{k=r}^{n} \binom{k}{r} = \binom{n+1}{r+1}$$

or

$$\sum_{k \le n} \binom{r+k}{k} = \binom{r+n+1}{n}$$

Also

$$\sum_{k=0}^{n} \binom{k}{m} = \binom{n+1}{m+1}$$

$$\sum_{i=0}^{n} \binom{n}{i} = 2^n$$

For  $n, m \geq 0$  and p prime: write n, m in base p, i.e.  $n = n_k p^k + \dots + n_1 p + n_0$  and  $m = m_k p^k + \dots + m_1 p + m_0$ . Then by Lucas theorem we have  $\binom{n}{m} \equiv \prod_{i=0}^k \binom{n_i}{m_i} \mod p$ , with the convention that  $n_i < m_i \implies \binom{n_i}{m_i} = 0$ .

Fibonacci (See also number theory section)

$$\sum_{0 \le k \le n} \binom{n-k}{k} = F_{n+1}$$

$$F_n = \frac{1}{\sqrt{5}} \left( \frac{1 + \sqrt{5}}{2} \right)^n - \frac{1}{\sqrt{5}} \left( \frac{1 - \sqrt{5}}{2} \right)^n$$

$$\sum_{i=1}^n F_i = F_{n+2} - 1, \sum_{i=1}^n F_i^2 = F_n F_{n+1}$$

$$\gcd(F_m, F_n) = F_{\gcd(m,n)}$$

$$\gcd(F_n, F_{n+1}) = \gcd(F_n, F_{n+2}) = 1$$

Bit stuff  $a+b=a\oplus b+2(a\&b)=a|b+a\&b$ . kth bit is set in x iff  $x \mod 2^{k-1} \geq 2^k$ , or iff  $x \mod 2^{k-1}-x \mod 2^k \neq 0$  (i.e.  $=2^k$ ) It comes handy when you need to look at the bits of the numbers which are pair sums or subset sums etc.

$$n \mod 2^i = n\&(2^i - 1).$$

$$\forall k: \ 1 \oplus 2 \oplus \ldots \oplus (4k-1) = 0$$

Stirling's numbers First kind:  $S_1(n,k)$  count permutations on n items with k cycles.  $S_1(n,k) = S_1(n-1,k-1) + (n-1)S_1(n-1,k)$  with  $S_1(0,0) = 1$ . Note:

$$\sum_{k=0}^{n} S_1(n,k)x^k = x(x+1)\dots(x+n-1)$$

$$\sum_{k=0}^{n} S_1(n,k) = n!$$

**Second kind:**  $S_2(n, k)$  count partitions of n distinct elements into exactly k non-empty groups.

$$S_2(n,k) = S_2(n-1,k-1) + kS_2(n-1,k)$$

$$S_2(n,1) = S_2(n,n) = 1$$

$$S_2(n,k) = \frac{1}{k!} \sum_{i=1}^{k} (-1)^{k-i} \binom{k}{i} i^n$$

# 4.3 Geometry Formulas

$$[ABC] = rs = \frac{1}{2}ab\sin\gamma$$

$$= \frac{abc}{4R} = \sqrt{s(s-a)(s-b)(s-c)} = \frac{1}{2} \left| (B-A, C-A)^T \right|$$

$$s = \frac{a+b+c}{2} \qquad 2R = \frac{a}{\sin \alpha}$$
 cosine rule: 
$$c^2 = a^2 + b^2 - 2ab\cos \gamma$$
 Euler: 
$$1 + CC = V - E + F$$
 Pick: 
$$\operatorname{Area} = \operatorname{itr} \operatorname{pts} + \frac{\operatorname{bdry} \operatorname{pts}}{2} - 1$$
 
$$p \cdot q = |p||q|\cos(\theta) \qquad |p \times q| = |p||q|\sin(\theta)$$

Given a non-self-intersecting closed polygon on n vertices, given as  $(x_i, y_i)$ , its centroid  $(C_x, C_y)$  is given as:

$$C_x = \frac{1}{6A} \sum_{i=0}^{n-1} (x_i + x_{i+1})(x_i y_{i+1} - x_{i+1} y_i),$$

$$C_y = \frac{1}{6A} \sum_{i=0}^{n-1} (y_i + y_{i+1})(x_i y_{i+1} - x_{i+1} y_i)$$

$$A = \frac{1}{2} \sum_{i=0}^{n-1} (x_i y_{i+1} - x_{i+1} y_i) = \text{polygon area}$$

**Inclusion-Exclusion** For appropriate f compute  $\sum_{S\subseteq T} (-1)^{|T\setminus S|} f(S)$ , or if only the size of S matters,  $\sum_{s=0}^{n} (-1)^{n-s} \binom{n}{s} f(s)$ . In some contexts we might use Stirling numbers, not binomial coefficients!

Some useful applications:

**Graph coloring** Let I(S) count the number of independent sets contained in  $S \subseteq V$  ( $I(\emptyset) = 1$ ,  $I(S) = I(S \setminus v) + I(S \setminus N(v))$ ). Let  $c_k = \sum_{S \subseteq V} (-1)^{|V \setminus S|} I(S)$ . Then V is k-colorable iff v > 0. Thus we can compute the chromatic number of a graph in  $O^*(2^n)$  time.

**Burnside's lemma** Given a group G acting on a set X, the number of elements in X up to symmetry is

$$\frac{1}{|G|} \sum_{g \in G} |X^g|$$

with  $X^g$  the elements of X invariant under g. For example, if f(n) counts "configurations" of some sort of length n, and we want to count them up to rotational symmetry using  $G = \mathbb{Z}/n\mathbb{Z}$ , then

$$g(n) = \frac{1}{n} \sum_{k=0}^{n-1} f(\gcd(n,k)) = \frac{1}{n} \sum_{k||n} f(k)\phi(n/k)$$

I.e. for coloring with c colors we have  $f(k) = k^c$ .

Relatedly, in Pólya's enumeration theorem we imagine X as a set of n beads with G permuting the beads (e.g. a necklace, with G all rotations and reflections of the n-cycle, i.e. the dihedral group  $D_n$ ). Suppose further that we had Y colors, then the number of G-invariant colorings  $Y^X/G$  is counted by

$$\frac{1}{|G|} \sum_{g \in G} |Y|^{c(g)}$$

with c(g) counting the number of cycles of g when viewed as a permutation of X. We can generalize this to a weighted version: if the color i can occur exactly  $r_i$  times, then this is counted by the coefficient of  $t_1^{r_1} \dots t_n^{r_n}$  in the polynomial

$$Z(t_1, \dots, t_n) = \frac{1}{|G|} \sum_{g \in G} \prod_{m \ge 1} (t_1^m + \dots + t_n^m)^{c_m(g)}$$

where  $c_m(g)$  counts the number of length m cycles in g acting as a permutation on X. Note we get the original formula by setting all  $t_i = 1$ . Here Z is the cycle index. Note: you can cleverly deal with even/odd sizes by setting some  $t_i$  to -1.

**Lucas Theorem** If p is prime, then:

$$\frac{p^a}{k} \equiv 0 \pmod{p}$$

Thus for non-negative integers  $m = m_k p^k + \ldots + m_1 p + m_0$ and  $n = n_k p^k + \ldots + n_1 p + n_0$ :

$$\frac{m}{n} = \prod_{i=0}^k \frac{m_i}{n_i} \mod p$$

Note: The fraction's mean integer division.

Catalan Numbers - Number of correct bracket sequence consisting of n opening and n closing brackets.

The number of ways to completely parenthesize n+1 factors.

The number of triangulations of a convex polygon with n+2 sides (i.e. the number of partitions of polygon into disjoint triangles by using the diagonals).

The number of ways to connect the 2n points on a circle to form n disjoint i.e. non-intersecting chords.

$$C_n = \frac{1}{n+1} \binom{2n}{n}$$

$$C_0 = 1, \ C_1 = 1, \ C_n = \sum_{k=0}^{n-1} C_k C_{n-1-k}$$

Narayana numbers The number of expressions containing n pairs of parentheses, which are correctly matched and which contain k distinct nestings.

$$N(n,k) = \frac{1}{n} \frac{n}{k} \frac{n}{k-1}$$