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## 1 Setup

**1.0.1 Tips Test session:** Check `__int128`, GNU builtins, and end of line whitespace requirements.

**C++ var. limits:** `int`  $-2^{31}$ ,  $2^{31} - 1$

`ll`  $-2^{63}$ ,  $2^{63} - 1$

`ull`  $0$ ,  $2^{64} - 1$

`__int128`  $-2^{127}$ ,  $2^{127} - 1$

`ld`  $-1.7e308$ ,  $1.7e308$ , 18 digits precision

**1.0.2 Xmodmap setup** remove `Lock = Caps_Lock`

`keysym Escape = Caps_Lock`

`keysym Caps_Lock = Escape`

add `Lock = Caps_Lock`

### 1.0.3 header.h

```
1 #pragma once
2 #include <bits/stdc++.h>
3 using namespace std;
4
5 #define ll long long
6 #define ull unsigned ll
7 #define ld long double
8 #define pl pair<ll, ll>
9 #define pi pair<int, int>
10 #define vl vector<ll>
11 #define vi vector<int>
12 #define vb vector<bool>
13 #define vvi vector<vi>
14 #define vvl vector<vl>
15 #define vpl vector<pl>
16 #define vpi vector<pi>
17 #define vld vector<ld>
18 #define vvp vector<vp>
19 #define in(e1, cont) (cont.find(e1) != cont.end())
20 // sets/maps
21 #define all(x) x.begin(), x.end()
22 #define rep(i, a, b) for(int i = a; i < (b); ++i)
23 constexpr int INF = INT_MAX;
24 constexpr ll LLINF = LONG_LONG_MAX;
25
26 // int main() {
27 //   ios::sync_with_stdio(false); // do not use
28 //   cout << printf
29 //   cin.tie(NULL);
30 //   cout << fixed << setprecision(12);
31 //   return 0;
32 // }
```

### 1.0.4 Aux. helper C++

```
1 #include "header.h"
2 int main() {
3     // Read in a line including white space
4     string line;
5     getline(cin, line);
6     // When doing the above read numbers as
7     // follows:
8     int n;
9     getline(cin, line);
10    stringstream ss(line);
11    ss >> n;
12
13    // Count the number of 1s in binary
14    // representation of a number
15    ull number;
16    __builtin_popcountll(number);
17 }
18
19 // __int128
20 using lll = __int128;
21 ostream& operator<<(ostream& o, __int128 n) {
22     auto t = n < 0 ? -n : n; char b[128], *d = end(b);
23     ;
24     do *--d = '0' + t % 10, t /= 10; while (t);
25     if(n < 0) *--d = '-';
26     o.rdbuf() -> sputn(d, end(b) - d);
27     return o;
28 }
```

### 1.0.5 Aux. helper python

```
1 from functools import lru_cache
2
3 # Read until EOF
4 while True:
5     try:
6         pattern = input()
7     except EOFError:
8         break
9
10 @lru_cache(maxsize=None)
11 def smth_memoi(i, j, s):
12     # Example in-built cache
13     return "sol"
14
15 # Fast I
16 import io, os
17 def fast_io():
18     finput = io.BytesIO(os.read(0,
19                               os.fstat(0).st_size)).readline
20     s = finput().decode()
21     return s
22
```

```
23 # Fast O
24 import sys
25 def fast_out():
26     n = 5
27     sys.stdout.write(str(n)+"\n")
```

## 2 Python

### 2.1 Graphs

#### 2.1.1 BFS

```
1 from collections import deque
2 def bfs(g, roots, n):
3     q = deque(roots)
4     explored = set()
5     distances = [0 if v in roots else float('inf')
6                  for v in range(n)]
7     while len(q) != 0:
8         node = q.popleft()
9         if node in explored: continue
10        explored.add(node)
11        for neigh in g[node]:
12            if neigh not in explored:
13                q.append(neigh)
14                distances[neigh] = float('inf')
15            distances[neigh] = distances[
16                node] + 1
17    return distances
```

#### 2.1.2 Dijkstra

```
1 from heapq import *
2 def dijkstra(n, root, g): # g = {node: (cost,
3                             # neigh)}
4     dist = [float("inf")] * n
5     dist[root] = 0
6     prev = [-1] * n
7
8     pq = [(0, root)]
9     heapify(pq)
10    visited = set([])
11
12    while len(pq) != 0:
13        _, node = heappop(pq)
14
15        if node in visited: continue
16        visited.add(node)
17
18    # In case of disconnected graphs
```

---

```

18 if node not in g:
19     continue
20
21 for cost, neigh in g[node]:
22     alt = dist[node] + cost
23     if alt < dist[neigh]:
24         dist[neigh] = alt
25         prev[neigh] = node
26         heappush(pq, (alt, neigh))
27 return dist

```

---

### 2.1.3 Topological Sort topological sorting of a DAG

---

```

1 from collections import defaultdict
2 class Graph:
3     def __init__(self, vertices):
4         self.graph = defaultdict(list) #adjacency
5             List
6         self.V = vertices #No. V
7
8     def addEdge(self, u, v):
9         self.graph[u].append(v)
10
11     def topologicalSortUtil(self, v, visited, stack):
12         :
13         visited[v] = True
14         # Recur for all the vertices adjacent to
15         this vertex
16         for i in self.graph[v]:
17             if visited[i] == False:
18                 self.topologicalSortUtil(i,
19                     visited, stack)
20         stack.insert(0, v)
21
22     def topologicalSort(self):
23         visited = [False]*self.V
24         stack = []
25         for i in range(self.V):
26             if visited[i] == False:
27                 self.topologicalSortUtil(i,
28                     visited, stack)
29         return stack
30
31     def isCyclicUtil(self, v, visited, recStack):
32         visited[v] = True
33         recStack[v] = True
34         for neighbour in self.graph[v]:
35             if visited[neighbour] == False:
36                 if self.isCyclicUtil(neighbour,
37                     visited, recStack) == True:
38                     return True
39             elif recStack[neighbour] == True:
40                 return True
41         recStack[v] = False

```

---



---

```

36 return False
37
38 def isCyclic(self):
39     visited = [False] * (self.V + 1)
40     recStack = [False] * (self.V + 1)
41     for node in range(self.V):
42         if visited[node] == False:
43             if self.isCyclicUtil(node,
44                 visited, recStack) == True:
45                 return True
46     return False

```

---

### 2.1.4 Kruskal (UnionFind) Min. span. tree

---

```

1 class UnionFind:
2     def __init__(self, n):
3         self.parent = [-1]*n
4
5     def find(self, x):
6         if self.parent[x] < 0:
7             return x
8         self.parent[x] = self.find(self.parent[x])
9         return self.parent[x]
10
11     def connect(self, a, b):
12         ra = self.find(a)
13         rb = self.find(b)
14         if ra == rb:
15             return False
16         if self.parent[ra] > self.parent[rb]:
17             self.parent[rb] += self.parent[ra]
18             self.parent[ra] = rb
19         else:
20             self.parent[ra] += self.parent[rb]
21             self.parent[rb] = ra
22         return True
23
24 # Full MST is len(spanning)==n-1
25 def kruskal(n, edges):
26     uf = UnionFind(n)
27     spanning = []
28     # Sort edges by asc. weight (check+-)
29     edges.sort(key = lambda d: -d[2])
30     while edges and len(spanning) < n-1:
31         u, v, w = edges.pop()
32         if not uf.connect(u, v):
33             continue
34         spanning.append((u, v, w))
35     return spanning

```

---

### 2.1.5 Prim Min. span. tree - good for dense graphs

---

```

1 from heapq import heappush, heappop, heapify
2 def prim(G, n):
3     s = next(iter(G.keys()))
4     V = set([s])
5     M = []
6     c = 0
7
8     E = [(w,s,v) for v,w in G[s].items()]
9     heapify(E)
10
11     while E and len(M) < n-1:
12         w,u,v = heappop(E)
13         if v in V: continue
14         V.add(v)
15         M.append((u,v))
16         c += w
17         u = v
18         [heappush(E, (w,u,v)) for v,w in G[u].items()
19             if v not in V]
20
21     if len(M) == n-1:
22         return M, c
23     else:
24         return None, None

```

---

## 2.2 Num. Th. / Comb.

### 2.2.1 nCk % prime p must be prime and k < p

---

```

1 def fermat_binom(n, k, p):
2     if k > n:
3         return 0
4     num = 1
5     for i in range(n-k+1, n+1):
6         num *= i % p
7     num %= p
8     denom = 1
9     for i in range(1, k+1):
10         denom *= i % p
11     denom %= p
12     # numerator * denominator^(p-2) (mod p)
13     return (num * pow(denom, p-2, p)) % p

```

---

### 2.2.2 Sieve of E. $O(n)$ so actually faster than C++ version, but more memory

---

```

1 MAX_SIZE = 10**8+1
2 isprime = [True] * MAX_SIZE
3 prime = []
4 SPF = [None] * (MAX_SIZE)
5 def manipulated_seive(N): # Up to N (not
6     included)

```

---

---

```

6 isprime[0] = isprime[1] = False
7 for i in range(2, N):
8     if isprime[i] == True:
9         prime.append(i)
10        SPF[i] = i
11        j = 0
12        while (j < len(prime) and
13              i * prime[j] < N and
14              prime[j] <= SPF[i]):
15            isprime[i * prime[j]] = False
16            SPF[i * prime[j]] = prime[j]
17            j += 1

```

---

### 2.2.3 Modular Inverse of a mod b

---

```

1 def modinv(a, b):
2     if b == 1: return 1
3     b0, x0, x1 = b, 0, 1
4     while a > 1:
5         q, a, b = a//b, b, a%b
6         x0, x1 = x1 - q * x0, x0
7     if x1 < 0: x1 += b0
8     return x1

```

---

**2.2.4 Chinese rem.** an  $x$  such that  $\forall y, m: yx = 1 \bmod m$  requires all  $m, m'$  to be  $\geq 1$  and coprime

---

```

1 def chinese_remainder(ys, ms):
2     N, x = 1, 0
3     for m in ms: N *= m
4     for y, m in zip(ys, ms):
5         n = N // m
6         x += n * y * modinv(n, m)
7     return x % N

```

---

### 2.2.5 Bezout

---

```

1 def bezout_id(a, b):
2     r, x, s, y, t, z = b, a, 0, 1, 1, 0
3     while r:
4         q = x // r
5         x, r = r, x % r
6         y, s = s, y - q * s
7         z, t = t, z - q * t
8     return y % (b // x), z % (-a // x)

```

---

### 2.2.6 Gen. chinese rem.

---

```

1 def general_chinese_remainder(a, b, m, n):
2     g = gcd(m, n)
3
4     if a == b and m == n:
5         return a, m
6     if (a % g) != (b % g):
7         return None, None
8
9     u, v = bezout_id(m, n)
10    x = (a*v*n + b*u*m) // g
11    return int(x) % lcm(m, n), int(lcm(m, n))

```

---

## 2.3 Strings

### 2.3.1 Longest common substr. (Consecutive)

$O(mn)$  time,  $O(m)$  space

---

```

1 from functools import lru_cache
2 @lru_cache
3 def lcs(s1, s2):
4     if len(s1) == 0 or len(s2) == 0:
5         return 0
6     return max(
7         lcs(s1[:-1], s2), lcs(s1, s2[:-1]),
8         (s1[-1] == s2[-1]) + lcs(s1[:-1], s2[:-1])
9     )

```

---

### 2.3.2 Longest common subseq. (Non-consecutive)

---

```

1 def longestCommonSubsequence(text1, text2):
2     n = len(text1)
3     m = len(text2)
4     prev = [0] * (m + 1)
5     cur = [0] * (m + 1)
6     for idx1 in range(1, n + 1):
7         for idx2 in range(1, m + 1):
8             # matching
9             if text1[idx1 - 1] == text2[idx2 - 1]:
10                cur[idx2] = 1 + prev[idx2 - 1]
11            else:
12                # not matching
13                cur[idx2] = max(cur[idx2 - 1], prev[idx2])
14        prev = cur.copy()
15    return cur[m]

```

---

### 2.3.3 KMP Return all matching pos. of P in T

---

```

1 class KMP:
2     def partial(self, pattern):
3         """_Calc._partial_match_table:_String_>_
4            [Int]"""
5         ret = [0]
6         for i in range(1, len(pattern)):
7             j = ret[i - 1]
8             while j > 0 and pattern[j] != pattern[i]:
9                 j = ret[j - 1]
10            ret.append(j + 1 if pattern[j] == pattern[i] else j)
11        return ret
12
13    def search(self, T, P):
14        """_KMPString_>_String_>_[Int]"""
15        partial, ret, j = self.partial(P), [], 0
16        for i in range(len(T)):
17            while j > 0 and T[i] != P[j]:
18                j = partial[j - 1]
19            if T[i] == P[j]:
20                j += 1
21            if j == len(P):
22                ret.append(i - (j - 1))
23                j = partial[j - 1]
24        return ret

```

---

**2.3.4 Longest common pref.** with the suffix array built we can do, e.g., longest common prefix of  $x, y$  with suffixarray where  $x, y$  are suffixes of the string used  $O(\log n)$

---

```

1 def lcp(x, y, P):
2     res = 0
3     if x == y:
4         return n - x
5     for k in range(len(P) - 1, -1, -1):
6         if x >= n or y >= n:
7             break
8         if P[k][x] == P[k][y]:
9             x += 1 << k
10            y += 1 << k
11            res += 1 << k
12    return res

```

---

### 2.3.5 Edit distance

---

```

1 def editDistance(str1, str2):
2     m = len(str1)
3     n = len(str2)
4     curr = [0] * (n + 1)
5     for j in range(n + 1):
6         curr[j] = j

```

---

```

7 previous = 0
8 # dp rows
9 for i in range(1, m + 1):
10     previous = curr[0]
11     curr[0] = i
12
13 # dp cols
14 for j in range(1, n + 1):
15     temp = curr[j]
16     if str1[i - 1] == str2[j - 1]:
17         curr[j] = previous
18     else:
19         curr[j] = 1 + min(previous, curr[j - 1],
20                             curr[j])
21     previous = temp
22 return curr[n]

```

**2.3.6 Bitstring** Slower than a set for many elements, but hashable. Also see Hashing

```

1 def add_element(bit_string, index):
2     return bit_string | (1 << index)
3 def remove_element(bit_string, index):
4     return bit_string & ~(1 << index)
5 def contains_element(bit_string, index):
6     return (bit_string & (1 << index)) != 0

```

## 2.4 Geometry

### 2.4.1 Convex Hull

```

1 def vec(a,b):
2     return (b[0]-a[0],b[1]-a[1])
3 def det(a,b):
4     return a[0]*b[1] - b[0]*a[1]
5 def convexhull(P):
6     if (len(P) == 1):
7         return [(p[0][0], p[0][1])]
8
9     h = sorted(P)
10    lower = []
11    i = 0
12    while i < len(h):
13        if len(lower) > 1:
14            a = vec(lower[-2], lower[-1])
15            b = vec(lower[-1], h[i])
16            if det(a,b) <= 0 and len(lower) > 1:
17                lower.pop()
18                continue
19            lower.append(h[i])
20            i += 1
21

```

```

22 upper = []
23 i = 0
24 while i < len(h):
25     if len(upper) > 1:
26         a = vec(upper[-2], upper[-1])
27         b = vec(upper[-1], h[i])
28         if det(a,b) >= 0:
29             upper.pop()
30             continue
31         upper.append(h[i])
32         i += 1
33
34 reversedupper = list(reversed(upper[1:-1]))
35 reversedupper.extend(lower)
36 return reversedupper

```

### 2.4.2 Geometry

```

1
2 def vec(a,b):
3     return (b[0]-a[0],b[1]-a[1])
4
5 def det(a,b):
6     return a[0]*b[1] - b[0]*a[1]
7
8 lower = []
9 i = 0
10 while i < len(h):
11     if len(lower) > 1:
12         a = vec(lower[-2], lower[-1])
13         b = vec(lower[-1], h[i])
14         if det(a,b) <= 0 and len(lower) > 1:
15             lower.pop()
16             continue
17         lower.append(h[i])
18         i += 1
19
20 # find upper hull
21 # det <= 0 -> replace
22 upper = []
23 i = 0
24 while i < len(h):
25     if len(upper) > 1:
26         a = vec(upper[-2], upper[-1])
27         b = vec(upper[-1], h[i])
28         if det(a,b) >= 0:
29             upper.pop()
30             continue
31         upper.append(h[i])
32         i += 1

```

## 3 C++

## 3.1 Graphs

### 3.1.1 BFS

```

1 #include "header.h"
2 #define graph unordered_map<ll, unordered_set<ll
>>
3 vi bfs(int n, graph& g, vi& roots) {
4     vi parents(n+1, -1); // nodes are 1..n
5     unordered_set<int> visited;
6     queue<int> q;
7     for (auto x: roots) {
8         q.emplace(x);
9         visited.insert(x);
10    }
11    while (not q.empty()) {
12        int node = q.front();
13        q.pop();
14
15        for (auto neigh: g[node]) {
16            if (not in(neigh, visited)) {
17                parents[neigh] = node;
18                q.emplace(neigh);
19                visited.insert(neigh);
20            }
21        }
22    }
23    return parents;
24 }
25 vi reconstruct_path(vi parents, int start, int
goal) {
26     vi path;
27     int curr = goal;
28     while (curr != start) {
29         path.push_back(curr);
30         if (parents[curr] == -1) return vi(); //
No path, empty vi
31         curr = parents[curr];
32     }
33     path.push_back(start);
34     reverse(path.begin(), path.end());
35     return path;
36 }

```

### 3.1.2 DFS Cycle detection / removal

```

1 #include "header.h"
2 void removeCyc(ll node, unordered_map<ll, vector<
pair<ll, ll>>>& neighs, vector<bool>& visited
,
3 vector<bool>& recStack, vector<ll>& ans) {
4     if (!visited[node]) {
5         visited[node] = true;
6         recStack[node] = true;

```

```

7     auto it = neighs.find(node);
8     if (it != neighs.end()) {
9         for (auto util: it->second) {
10             ll nnode = util.first;
11             if (recStack[nnode]) {
12                 ans.push_back(util.second);
13             } else if (!visited[nnode]) {
14                 removeCyc(nnode, neighs,
15                           visited, recStack, ans);
16             }
17         }
18     }
19     recStack[node] = false;
20 }

```

### 3.1.3 Dijkstra

```

1 #include "header.h"
2 vector<int> dijkstra(int n, int root, map<int,
3   vector<pair<int, int>>& g) {
4     unordered_set<int> visited;
5     vector<int> dist(n, INF);
6     priority_queue<pair<int, int>> pq;
7     dist[root] = 0;
8     pq.push({0, root});
9     while (!pq.empty()) {
10         int node = pq.top().second;
11         int d = -pq.top().first;
12         pq.pop();
13
14         if (in(node, visited)) continue;
15         visited.insert(node);
16
17         for (auto e : g[node]) {
18             int neigh = e.first;
19             int cost = e.second;
20             if (dist[neigh] > dist[node] + cost) {
21                 dist[neigh] = dist[node] + cost;
22                 pq.push({-dist[neigh], neigh});
23             }
24         }
25     }
26     return dist;

```

### 3.1.4 Floyd-Warshall

```

1 #include "header.h"
2 // g[i][j] = inf if not path from i to j
3 // if g[i][i] < 0, i is contained in a negative
4 // cycle

```

```

4 void warshall(vvl& g) {
5     for (int k=0; k<g.size(); ++k) {
6         for (int i=0; i<g.size(); ++i) {
7             for (int j=0; j<g.size(); ++j) {
8                 if (g[i][k] < LLONG_MAX and g[k][j] <
9                     LLONG_MAX and g[i][k] + g[k][j] < g[i][j]) {
10                     g[i][j] = g[i][k] + g[k][j];
11                 }
12             }
13         }
14     }
15 }

```

### 3.1.5 Kruskal Minimum spanning tree of undirected weighted graph. $O(E \log E)$

```

1 #include "header.h"
2 #include "disjoint_set.h"
3 pair<set<pair<ll, ll>>, ll> kruskal(vector<tuple<
4   ll, ll, ll>>& edges, ll n) {
5     set<pair<ll, ll>> ans;
6     ll cost = 0;
7
8     sort(edges.begin(), edges.end());
9     DisjointSet<ll> fs(n);
10
11     ll dist, i, j;
12     for (auto edge: edges) {
13         dist = get<0>(edge);
14         i = get<1>(edge);
15         j = get<2>(edge);
16
17         if (fs.find_set(i) != fs.find_set(j)) {
18             fs.union_sets(i, j);
19             ans.insert({i, j});
20             cost += dist;
21         }
22     }
23     return pair<set<pair<ll, ll>>, ll> {ans, cost};

```

### 3.1.6 Hungarian algorithm Given J jobs and W workers ( $J \leq W$ ), computes the minimum cost to assign each prefix of jobs to distinct workers.

```

1 #include "header.h"
2 template <class T> bool ckmin(T &a, const T &b) {
3     return b < a ? a = b, 1 : 0; }
4 /**
5  * @tparam T: type large enough to represent
6  * integers of  $0(J * \max(|C|))$ 
7  * @param C: JxW matrix such that  $C[j][w]$  = cost
8  * to assign j-th
9  * job to w-th worker (possibly negative)

```

```

7  * @return a vector (length J), with the j-th
8  * entry = min. cost
9  * to assign the first (j+1) jobs to distinct
10 * workers
11 */
12 template <class T> vector<T> hungarian(const
13   vector<vector<T>>& C) {
14     const int J = (int)C.size(), W = (int)C[0].size();
15     assert(J <= W);
16     // a W-th worker added for convenience
17     vector<int> job(W + 1, -1);
18     vector<T> ys(J), yt(W + 1); // potentials
19     vector<T> answers;
20     const T inf = numeric_limits<T>::max();
21     for (int j_cur = 0; j_cur < J; ++j_cur) {
22         int w_cur = W;
23         job[w_cur] = j_cur;
24         vector<T> min_to(W + 1, inf);
25         vector<int> prv(W + 1, -1);
26         vector<bool> in_Z(W + 1);
27         while (job[w_cur] != -1) { // runs at
28             most j_cur + 1 times
29             in_Z[w_cur] = true;
30             const int j = job[w_cur];
31             T delta = inf;
32             int w_next;
33             for (int w = 0; w < W; ++w) {
34                 if (!in_Z[w]) {
35                     if (ckmin(min_to[w], C[j][w] -
36                             ys[j] + yt[w])) {
37                         prv[w] = w_cur;
38                         if (ckmin(delta, min_to[w]))
39                             w_next = w;
40                     }
41                 }
42             }
43             for (int w = 0; w <= W; ++w) {
44                 if (in_Z[w]) ys[job[w]] += delta,
45                     yt[w] -= delta;
46                 else min_to[w] -= delta;
47             }
48             w_cur = w_next;
49         }
50         job[w_cur] = job[w = prv[w_cur]];
51         answers.push_back(-yt[W]);
52     }
53     return answers;
54 }

```

### 3.1.7 Suc. shortest path Calculates max flow, min cost

```

1 #include "header.h"

```

```

2 // map<node, map<node, pair<cost, capacity>>>
3 #define graph unordered_map<int, unordered_map<
  int, pair<ld, int>>>
4 graph g;
5 const ld infity = 1e60l; // Change if necessary
6 ld fill(int n, vld& potential) { // Finds max
  flow, min cost
7   priority_queue<pair<ld, int>> pq;
8   vector<bool> visited(n+2, false);
9   vi parent(n+2, 0);
10  vld dist(n+2, infity);
11  dist[0] = 0.1;
12  pq.emplace(make_pair(0.1, 0));
13  while (not pq.empty()) {
14    int node = pq.top().second;
15    pq.pop();
16    if (visited[node]) continue;
17    visited[node] = true;
18    for (auto& x : g[node]) {
19      int neigh = x.first;
20      int capacity = x.second.second;
21      ld cost = x.second.first;
22      if (capacity and not visited[neigh]) {
23        ld d = dist[node] + cost + potential[node]
          - potential[neigh];
24        if (d + 1e-10l < dist[neigh]) {
25          dist[neigh] = d;
26          pq.emplace(make_pair(-d, neigh));
27          parent[neigh] = node;
28        }
29      }
30    }
31    potential[i] = min(infity, potential[i] + dist
      [i]);
32  }
33  if (not parent[n+1]) return infity;
34  ld ans = 0.1;
35  for (int x = n+1; x; x=parent[x]) {
36    ans += g[parent[x]][x].first;
37    g[parent[x]][x].second--;
38    g[x][parent[x]].second++;
39  }
40  return ans;
41 }

```

### 3.1.8 Bipartite check

```

1 #include "header.h"
2 int main() {
3   int n;
4   vvi adj(n);
5
6   vi side(n, -1); // will have 0's for one
  side 1's for other side

```

```

7   bool is_bipartite = true; // becomes false
  if not bipartite
8   queue<int> q;
9   for (int st = 0; st < n; ++st) {
10    if (side[st] == -1) {
11      q.push(st);
12      side[st] = 0;
13      while (!q.empty()) {
14        int v = q.front();
15        q.pop();
16        for (int u : adj[v]) {
17          if (side[u] == -1) {
18            side[u] = side[v] ^ 1;
19            q.push(u);
20          } else {
21            is_bipartite &= side[u]
              != side[v];
22          }
23        }
24      }
25    }
26  }

```

**3.1.9 Bipartite matching (Hopcroft-Karp)** Fast bipartite matching algorithm. Graph  $g$  should be a list of neighbors of the left partition, and  $btoa$  should be a vector full of -1's of the same size as the right partition. Returns the size of the matching.  $btoa[i]$  will be the match for vertex  $i$  on the right side, or -1 if it's not matched. Time:  $O(\sqrt{VE})$

```

1 // Usage: vi btoa(m, -1); hopcroftKarp(g, btoa);
2
3 bool dfs(int a, int L, vector<vi>& g, vi& btoa,
  vi& A, vi& B) {
4   if (A[a] != L) return 0;
5   A[a] = -1;
6   for (int b : g[a]) if (B[b] == L + 1) {
7     B[b] = 0;
8     if (btoa[b] == -1 || dfs(btoa[b], L + 1, g,
        btoa, A, B))
9       return btoa[b] = a, 1;
10  }
11  return 0;
12
13 int hopcroftKarp(vector<vi>& g, vi& btoa) {
14   int res = 0;
15   vi A(g.size()), B(btoa.size()), cur, next;
16   for (;;) {
17     fill(all(A), 0); fill(all(B), 0);
18     /// Find the starting nodes for BFS (i.e.
        layer 0).
19     cur.clear();
20     for (int a : btoa) if (a != -1) A[a] = -1;
21     rep(a, 0, sz(g)) if (A[a] == 0) cur.push_back(a);
22     /// Find all layers using bfs.

```

```

22   for (int lay = 1; lay++) {
23     bool islast = 0;
24     next.clear();
25     for (int a : cur) for (int b : g[a]) {
26       if (btoa[b] == -1) {
27         B[b] = lay; islast = 1;
28       } else if (btoa[b] != a && !B[b]) {
29         B[b] = lay;
30         next.push_back(btoa[b]);
31       }
32     }
33     if (islast) break;
34     if (next.empty()) return res;
35     for (int a : next) A[a] = lay;
36     cur.swap(next);
37   }
38   /// Use DFS to scan for augmenting paths.
39   rep(a, 0, sz(g))
40     res += dfs(a, 0, g, btoa, A, B);

```

### 3.1.10 Find cycle directed

```

1 #include "header.h"
2 int n;
3 const int mxN = 2e5+5;
4 vvi adj(mxN);
5 vector<char> color;
6 vi parent;
7 int cycle_start, cycle_end;
8 bool dfs(int v) {
9   color[v] = 1;
10  for (int u : adj[v]) {
11    if (color[u] == 0) {
12      parent[u] = v;
13      if (dfs(u)) return true;
14    } else if (color[u] == 1) {
15      cycle_end = u;
16      cycle_start = v;
17      return true;
18    }
19  }
20  color[v] = 2;
21  return false;
22 }
23 void find_cycle() {
24   color.assign(n, 0);
25   parent.assign(n, -1);
26   cycle_start = -1;
27   for (int v = 0; v < n; v++) {
28     if (color[v] == 0 && dfs(v)) break;
29   }
30   if (cycle_start == -1) {
31     cout << "Acyclic" << endl;
32   } else {

```



```

33     vector<int> cycle;
34     cycle.push_back(cycle_start);
35     for (int v = cycle_end; v != cycle_start;
          v = parent[v])
36         cycle.push_back(v);
37     cycle.push_back(cycle_start);
38     reverse(cycle.begin(), cycle.end());
39
40     cout << "Cycle Found: ";
41     for (int v : cycle) cout << v << " ";
42     cout << endl;
43 }
44 }

```

### 3.1.11 Find cycle undirected

```

1 #include "header.h"
2 int n;
3 const int mxN = 2e5 + 5;
4 vvi adj(mxN);
5 vector<bool> visited;
6 vi parent;
7 int cycle_start, cycle_end;
8 bool dfs(int v, int par) { // passing vertex and
    its parent vertex
9     visited[v] = true;
10    for (int u : adj[v]) {
11        if (u == par) continue; // skipping edge
            to parent vertex
12        if (visited[u]) {
13            cycle_end = v;
14            cycle_start = u;
15            return true;
16        }
17        parent[u] = v;
18        if (dfs(u, parent[u]))
19            return true;
20    }
21    return false;
22 }
23 void find_cycle() {
24     visited.assign(n, false);
25     parent.assign(n, -1);
26     cycle_start = -1;
27     for (int v = 0; v < n; v++) {
28         if (!visited[v] && dfs(v, parent[v]))
29             break;
30     }
31     if (cycle_start == -1) {
32         cout << "Acyclic" << endl;
33     } else {
34         vector<int> cycle;
35         cycle.push_back(cycle_start);
36         for (int v = cycle_end; v != cycle_start;
            v = parent[v])

```

```

36         cycle.push_back(v);
37         cycle.push_back(cycle_start);
38         cout << "Cycle Found: ";
39         for (int v : cycle) cout << v << " ";
40         cout << endl;
41     }
42 }

```

### 3.1.12 Tarjan's SCC

```

1 #include "header.h"
2 struct Tarjan {
3     vvi &edges;
4     int V, counter = 0, C = 0;
5     vi n, l;
6     vector<bool> vs;
7     stack<int> st;
8     Tarjan(vvi &e) : edges(e), V(e.size()), n(V,
        -1), l(V, -1), vs(V, false) {}
9     void visit(int u, vi &com) {
10         l[u] = n[u] = counter++;
11         st.push(u);
12         vs[u] = true;
13         for (auto &&v : edges[u]) {
14             if (n[v] == -1) visit(v, com);
15             if (vs[v]) l[u] = min(l[u], l[v]);
16         }
17         if (l[u] == n[u]) {
18             while (true) {
19                 int v = st.top();
20                 st.pop();
21                 vs[v] = false;
22                 com[v] = C; //<== ACT HERE
23                 if (u == v) break;
24             }
25             C++;
26         }
27     }
28     int find_sccs(vi &com) { // component indices
        will be stored in 'com'
29         com.assign(V, -1);
30         C = 0;
31         for (int u = 0; u < V; ++u)
32             if (n[u] == -1) visit(u, com);
33         return C;
34     }
35     // scc is a map of the original vertices of the
        graph to the vertices of the SCC graph,
        scc_graph is its adjacency list. SCC
        indices and edges are stored in 'scc' and '
        scc_graph'.
36 void scc_collapse(vi &scc, vvi &scc_graph) {
37     find_sccs(scc);
38     scc_graph.assign(C, vi());

```

```

39     set<pi> rec; // recorded edges
40     for (int u = 0; u < V; ++u) {
41         assert(scc[u] != -1);
42         for (int v : edges[u]) {
43             if (scc[v] == scc[u] ||
44                 rec.find({scc[u], scc[v]}) != rec.end()
45                     ) continue;
46             scc_graph[scc[u]].push_back(scc[v]);
47             rec.insert({scc[u], scc[v]});
48         }
49     }
50     // The number of edges needed to be added is
        max(sources.size(), sinks.size())
51 void findSourcesAndSinks(const vvi &scc_graph,
        vi &sources, vi &sinks) {
52     vi in_degree(C, 0), out_degree(C, 0);
53     for (int u = 0; u < C; u++) {
54         for (auto v : scc_graph[u]) {
55             in_degree[v]++;
56             out_degree[u]++;
57         }
58     }
59     for (int i = 0; i < C; ++i) {
60         if (in_degree[i] == 0) sources.push_back(i);
61         if (out_degree[i] == 0) sinks.push_back(i);
62     }
63 }
64 };

```

### 3.1.13 SCC edges Prints out the missing edges to make the input digraph strongly connected

```

1 #include "header.h"
2 const int N=1e5+10;
3 int n,a[N],cnt[N],vis[N];
4 vector<int> hd,tl;
5 int dfs(int x){
6     vis[x]=1;
7     if(!vis[a[x]])return vis[x]=dfs(a[x]);
8     return vis[x]=x;
9 }
10 int main(){
11     scanf("%d",&n);
12     for(int i=1;i<=n;i++){
13         scanf("%d",&a[i]);
14         cnt[a[i]]++;
15     }
16     int k=0;
17     for(int i=1;i<=n;i++){
18         if(!cnt[i]){
19             k++;
20             hd.push_back(i);

```



```

21         tl.push_back(dfs(i));
22     }
23 }
24 int tk=k;
25 for(int i=1;i<=n;i++){
26     if(!vis[i]){
27         k++;
28         hd.push_back(i);
29         tl.push_back(dfs(i));
30     }
31 }
32 if(k==1&&!tk)k=0;
33 printf("%d\n",k);
34 for(int i=0;i<k;i++)printf("%d_\n",tl[i],hd
35     [(i+1)%k]);
36 return 0;
37 }

```

### 3.1.14 Topological sort

```

1 #include "header.h"
2 int n; // number of vertices
3 vvi adj; // adjacency list of graph
4 vector<bool> visited;
5 vi ans;
6 void dfs(int v) {
7     visited[v] = true;
8     for (int u : adj[v]) {
9         if (!visited[u]) dfs(u);
10    }
11    ans.push_back(v);
12 }
13 void topological_sort() {
14     visited.assign(n, false);
15     ans.clear();
16     for (int i = 0; i < n; ++i) {
17         if (!visited[i]) dfs(i);
18     }
19     reverse(ans.begin(), ans.end());
20 }

```

### 3.1.15 Bellmann-Ford Same as Dijkstra but allows neg. edges

```

1 #include "header.h"
2 // Switch vi and vvpj to vl and vvpl if necessary
3 void bellmann_ford_extended(vvpj &e, int source,
4     int goal, vi &dist, vb &cyc) {
5     dist.assign(e.size(), INF);
6     cyc.assign(e.size(), false); // true when u
7     // is in a <0 cycle
8     dist[source] = 0;
9 }

```

```

8 // Perform n-1 relaxations
9 for (int iter = 0; iter < e.size() - 1; ++
10     iter) {
11     bool relax = false;
12     for (int u = 0; u < e.size(); ++u) {
13         if (dist[u] == INF) continue;
14         for (auto &edge : e[u]) {
15             int v = edge.first, w = edge.
16                 second;
17             if (dist[u] + w < dist[v]) {
18                 dist[v] = dist[u] + w;
19                 relax = true;
20             }
21         }
22     }
23     if (!relax) break;
24 }
25 // Step to detect any reachable negative
26 // cycles
27 for (int u = 0; u < e.size(); ++u) {
28     if (dist[u] == INF) continue;
29     for (auto &edge : e[u]) {
30         int v = edge.first, w = edge.second;
31         if (dist[u] + w < dist[v]) {
32             // If we can still relax, mark
33             // the node in the negative
34             // cycle
35             dist[v] = -INF;
36             cyc[v] = true;
37         }
38     }
39 }
40 // Propagate neg. cycle detection to all
41 // reachable nodes (if necessary)
42 bool change = true;
43 while (change) {
44     change = false;
45     for (int u = 0; u < e.size(); ++u) {
46         if (!cyc[u]) continue;
47         for (auto &edge : e[u]) {
48             int v = edge.first;
49             if (!cyc[v]) {
50                 cyc[v] = true;
51                 dist[v] = -INF;
52                 change = true;
53             }
54         }
55     }
56 }
57 }

```

### 3.1.16 Dinic max flow $O(V^2E)$ , $O(Ef)$

```

1 #include "header.h"

```

```

2 using F = ll; using W = ll; // types for flow and
3 // weight/cost
4 struct S{
5     const int v; // neighbour
6     const int r; // index of the reverse edge
7     F f; // current flow
8     const F cap; // capacity
9     const W cost; // unit cost
10    S(int v, int ri, F c, W cost = 0) :
11        v(v), r(ri), f(0), cap(c), cost(cost) {}
12    inline F res() const { return cap - f; }
13 };
14 struct FlowGraph : vector<vector<S>> {
15     FlowGraph(size_t n) : vector<vector<S>>(n) {}
16     void add_edge(int u, int v, F c, W cost = 0){
17         auto &t = *this;
18         t[u].emplace_back(v, t[v].size(), c, cost
19             );
20         t[v].emplace_back(u, t[u].size()-1, c, -
21             cost);
22     }
23     void add_arc(int u, int v, F c, W cost = 0){
24         auto &t = *this;
25         t[u].emplace_back(v, t[v].size(), c, cost
26             );
27         t[v].emplace_back(u, t[u].size()-1, 0, -
28             cost);
29     }
30     void clear() { for (auto &E : *this) for (
31         auto &e : E) e.f = 0LL; }
32 };
33 struct Dinic{
34     FlowGraph &edges; int V,s,t;
35     vi l; vector<vector<S>::iterator> its; //
36     // levels and iterators
37     Dinic(FlowGraph &edges, int s, int t) :
38         edges(edges), V(edges.size()), s(s), t(t)
39         , l(V,-1), its(V) {}
40     ll augment(int u, F c) { // we reuse the same
41         // iterators
42         if (u == t) return c; ll r = 0LL;
43         for(auto &i = its[u]; i != edges[u].end()
44             ; i++){
45             auto &e = *i;
46             if (e.res() && l[u] < l[e.v]) {
47                 auto d = augment(e.v, min(c, e.
48                     res()));
49                 if (d > 0) { e.f += d; edges[e.v
50                     ][e.r].f -= d; c -= d;
51                     r += d; if (!c) break; }
52             }
53         }
54         return r;
55     }
56     ll run() {
57         ll flow = 0, f;
58     }
59 }

```

```

43 while(true) {
44     fill(l.begin(), l.end(), -1); l[s]=0;
45     queue<int> q; q.push(s);
46     while(!q.empty()){
47         auto u = q.front(); q.pop(); its[
48             u] = edges[u].begin();
49         for(auto &&e : edges[u]) if(e.res
50             () && l[e.v]<0)
51             l[e.v] = l[u]+1, q.push(e.v);
52     }
53     if (l[t] < 0) return flow;
54     while ((f = augment(s, INF)) > 0)
55         flow += f;
56 }
57 };

```

**3.1.17 Edmonds-Karp** (Max) flow algorithm with time  $O(VE^2)$ . To get edge flow values, compare capacities before and after, and take the positive values only.

```

1 #include "header.h"
2 template<class T> T edmondsKarp(vector<
3     unordered_map<int, T>>&
4     graph, int source, int sink) {
5     assert(source != sink);
6     T flow = 0;
7     vi par(sz(graph)), q = par;
8     for (;;) {
9         fill(all(par), -1);
10        par[source] = 0;
11        int ptr = 1;
12        q[0] = source;
13
14        rep(i,0,ptr) {
15            int x = q[i];
16            for (auto e : graph[x]) {
17                if (par[e.first] == -1 && e.second > 0) {
18                    par[e.first] = x;
19                    q[ptr++] = e.first;
20                    if (e.first == sink) goto out;
21                }
22            }
23        }
24        return flow;
25    out:
26        T inc = numeric_limits<T>::max();
27        for (int y = sink; y != source; y = par[y])
28            inc = min(inc, graph[par[y]][y]);
29
30        flow += inc;
31        for (int y = sink; y != source; y = par[y]) {
32            int p = par[y];

```

```

33         if ((graph[p][y] -= inc) <= 0) graph[p].
34             erase(y);
35         graph[y][p] += inc;
36     }
37 }

```

## 3.2 Dynamic Programming

### 3.2.1 Longest Incr. Subseq.

```

1 #include "header.h"
2 template<class T>
3 vector<T> index_path_lis(vector<T>& nums) {
4     int n = nums.size();
5     vector<T> sub;
6     vector<int> subIndex;
7     vector<T> path(n, -1);
8     for (int i = 0; i < n; ++i) {
9         if (sub.empty() || sub[sub.size() - 1] <
10            nums[i]) {
11             path[i] = sub.empty() ? -1 : subIndex[sub.
12                 size() - 1];
13             sub.push_back(nums[i]);
14             subIndex.push_back(i);
15             } else {
16                 int idx = lower_bound(sub.begin(), sub.end(),
17                     nums[i]) - sub.begin();
18                 path[i] = idx == 0 ? -1 : subIndex[idx - 1];
19                 sub[idx] = nums[i];
20                 subIndex[idx] = i;
21             }
22     }
23     vector<T> ans;
24     int t = subIndex[subIndex.size() - 1];
25     while (t != -1) {
26         ans.push_back(t);
27         t = path[t];
28     }
29     reverse(ans.begin(), ans.end());
30     return ans;
31 }
32 // Length only
33 template<class T>
34 int length_lis(vector<T> &a) {
35     set<T> st;
36     typename set<T>::iterator it;
37     for (int i = 0; i < a.size(); ++i) {
38         it = st.lower_bound(a[i]);
39         if (it != st.end()) st.erase(it);
40         st.insert(a[i]);
41     }
42     return st.size();
43 }

```

**3.2.2 0-1 Knapsack** Given a number of coins, calculate all possible distinct sums

```

1 #include "header.h"
2 int main() {
3     int n;
4     vi coins(n); // possible coins to use
5     int sum = 0; // their sum of the coins
6     vi dp(sum + 1, 0); // dp[x] = 1 if sum x can be
7         made
8     dp[0] = 1;
9     for (int c = 0; c < n; ++c)
10         for (int x = sum; x >= 0; --x)
11             if (dp[x]) dp[x + coins[c]] = 1;
12 }

```

**3.2.3 Coin change** Total distinct ways to make sum using  $n$  coins of different vals

```

1 #include "header.h"
2 int count(vi& coins, int n, int sum) {
3     vvi dp(n + 1, vi(sum + 1, 0));
4     dp[0][0] = 1;
5     for (int i = 1; i <= n; i++) {
6         for (int j = 0; j <= sum; j++) {
7             // without using the current coin,
8             dp[i][j] += dp[i - 1][j];
9             // using the current coin
10            if ((j - coins[i - 1]) >= 0)
11                dp[i][j] += dp[i][j - coins[i -
12                    1]];
13        }
14    }
15    return dp[n][sum];
16 }

```

**3.2.4 Longest common subseq.** Optimization for each unique element appearing  $k$ -times

```

1 #include "../header.h"
2 #include "../DataStructures/fenwick_tree.cpp"
3 int lcs(int k, vector<int>& A, vector<int>& B) {
4     int lenA = A.size();
5     int lenB = B.size();
6
7     // Determine the number of distinct elements
8     // from max element in A and B
9     int n = max(*max_element(A.begin(), A.end()),
10        *max_element(B.begin(), B.end())) + 1;
11
12     vector<vector<int>> C(n);
13     for (int j = 0; j < lenB; ++j) {

```

```

12     C[B[j]].push_back(j);
13 }
14
15 int ans = 0;
16 FenwickTree<int> fenwick(lenB + 1);
17 for (int i = 0; i < lenA; ++i) {
18     int a = A[i];
19     for (int j = C[a].size() - 1; j >= 0; --j)
20     {
21         int pos = C[a][j];
22         int x = fenwick.query(pos) + 1;
23         fenwick.update(pos + 1, x); //
24         Convert to 1-based index
25         ans = max(ans, x);
26     }
27 }
28 return ans;
29 }

```

## 3.3 Numerical

### 3.3.1 Template (for this section)

```

1 #include <bits/stdc++.h>
2 using namespace std;
3 #define rep(i, a, b) for(int i = a; i < (b); ++i)
4 #define all(x) begin(x), end(x)
5 #define sz(x) (int)(x).size()
6 typedef long long ll;
7 typedef pair<int, int> pii;
8 typedef vector<int> vi;

```

### 3.3.2 Polynomial

```

1 #include "template.cpp"
2 struct Poly {
3     vector<double> a;
4     double operator()(double x) const {
5         double val = 0;
6         for (int i = sz(a); i--;) (val += x) += a[i];
7         return val;
8     }
9     void diff() {
10         rep(i, 1, sz(a)) a[i-1] = i*a[i];
11         a.pop_back();
12     }
13     void divroot(double x0) {
14         double b = a.back(), c; a.back() = 0;
15         for(int i=sz(a)-1; i--;) c = a[i], a[i] = a[i]
16             +1*x0+b, b=c;
17         a.pop_back();
18     }
19 };

```

**3.3.3 Poly Roots** Finds the real roots to a polynomial.  $O(n^2 \log(1/\epsilon))$

```

1 // Usage: polyRoots({{2,-3,1}},-1e9,1e9) = solve
2 // x^2-3x+2 = 0
3 #include "Polynomial.h"
4 #include "template.cpp"
5 vector<double> polyRoots(Poly p, double xmin,
6     double xmax) {
7     if (sz(p.a) == 2) { return {-p.a[0]/p.a[1]}; }
8     vector<double> ret;
9     Poly der = p;
10    der.diff();
11    auto dr = polyRoots(der, xmin, xmax);
12    dr.push_back(xmin-1);
13    dr.push_back(xmax+1);
14    sort(all(dr));
15    rep(i,0,sz(dr)-1) {
16        double l = dr[i], h = dr[i+1];
17        bool sign = p(l) > 0;
18        if (sign ^ (p(h) > 0)) {
19            rep(it,0,60) { // while (h - l > 1e-8)
20                double m = (l + h) / 2, f = p(m);
21                if ((f <= 0) ^ sign) l = m;
22                else h = m;
23            }
24            ret.push_back((l + h) / 2);
25        }
26    }
27    return ret;
28 }

```

**3.3.4 Golden Section Search** Finds the argument minimizing the function  $f$  in the interval  $[a, b]$  assuming  $f$  is unimodal on the interval, i.e. has only one local minimum and no local maximum. The maximum error in the result is  $\epsilon$ . Works equally well for maximization with a small change in the code. See TernarySearch.h in the Various chapter for a discrete version.  $O(\log((b-a)/\epsilon))$

```

1 /** Usage:
2     double func(double x) { return 4+x+.3*x*x; }
3     double xmin = gss(-1000,1000,func); */
4 #include "template.cpp"
5 // It is important for r to be precise, otherwise
6 // we don't necessarily maintain the inequality
7 // a < x1 < x2 < b.
8 double gss(double a, double b, double (*f)(double)) {
9     double r = (sqrt(5)-1)/2, eps = 1e-7;
10    double x1 = b - r*(b-a), x2 = a + r*(b-a);
11    double f1 = f(x1), f2 = f(x2);
12    while (b-a > eps)

```

```

11    if (f1 < f2) { //change to > to find maximum
12        b = x2; x2 = x1; f2 = f1;
13        x1 = b - r*(b-a); f1 = f(x1);
14    } else {
15        a = x1; x1 = x2; f1 = f2;
16        x2 = a + r*(b-a); f2 = f(x2);
17    }
18    return a;
19 }

```

**3.3.5 Hill Climbing** Poor man's optimization for unimodal functions.

```

1 #include "template.cpp"
2 typedef array<double, 2> P;
3 template<class F> pair<double, P> hillClimb(P
4     start, F f) {
5     pair<double, P> cur(f(start), start);
6     for (double jmp = 1e9; jmp > 1e-20; jmp /= 2) {
7         rep(j,0,100) rep(dx,-1,2) rep(dy,-1,2) {
8             P p = cur.second;
9             p[0] += dx*jmp;
10            p[1] += dy*jmp;
11            cur = min(cur, make_pair(f(p), p));
12        }
13    }
14    return cur;
15 }

```

**3.3.6 Integration** Simple integration of a function over an interval using Simpson's rule. The error should be proportional to  $h^4$ , although in practice you will want to verify that the result is stable to desired precision when epsilon changes.

```

1 #include "template.cpp"
2 template<class F>
3 double quad(double a, double b, F f, const int n
4     = 1000) {
5     double h = (b - a) / 2 / n, v = f(a) + f(b);
6     rep(i,1,n*2)
7         v += f(a + i*h) * (i&1 ? 4 : 2);
8     return v * h / 3;
9 }

```

### 3.3.7 Integration Adaptive Fast integration using an adaptive Simpson's rule.

---

```

1 /** Usage:
2 double sphereVolume = quad(-1, 1, [](double x) {
3 return quad(-1, 1, [&](double y) {
4 return quad(-1, 1, [&](double z) {
5 return x*x + y*y + z*z < 1; });});}); */
6 #include "template.cpp"
7 typedef double d;
8 #define S(a,b) (f(a) + 4*f((a+b) / 2) + f(b)) * (
9 b-a) / 6
9 template <class F>
10 d rec(F& f, d a, d b, d eps, d S) {
11     d c = (a + b) / 2;
12     d S1 = S(a, c), S2 = S(c, b), T = S1 + S2;
13     if (abs(T - S) <= 15 * eps || b - a < 1e-10)
14         return T + (T - S) / 15;
15     return rec(f, a, c, eps / 2, S1) + rec(f, c, b,
16         eps / 2, S2);
17 }
18 template<class F>
19 d quad(d a, d b, F f, d eps = 1e-8) {
20     return rec(f, a, b, eps, S(a, b));
21 }

```

---

## 3.4 Num. Th. / Comb.

### 3.4.1 Basic stuff

---

```

1 #include "header.h"
2 ll gcd(ll a, ll b) { while (b) { a %= b; swap(a,
3     b); } return a; }
4 ll lcm(ll a, ll b) { return (a / gcd(a, b)) * b; }
5 ll mod(ll a, ll b) { return ((a % b) + b) % b; }
6 // Finds x, y s.t. ax + by = d = gcd(a, b).
7 void extended_euclid(ll a, ll b, ll &x, ll &y, ll
8     &d) {
9     ll xx = y = 0;
10    ll yy = x = 1;
11    while (b) {
12        ll q = a / b;
13        ll t = b; b = a % b; a = t;
14        t = xx; xx = x - q * xx; x = t;
15        t = yy; yy = y - q * yy; y = t;
16    }
17    d = a;
18 }
19 // solves ab = 1 (mod n), -1 on failure
20 ll mod_inverse(ll a, ll n) {
21     ll x, y, d;
22     extended_euclid(a, n, x, y, d);
23     return (d > 1 ? -1 : mod(x, n));
24 }

```

---

```

22 }
23 // All modular inverses of [1..n] mod P in O(n)
24 // time.
25 vi inverses(ll n, ll P) {
26     vi I(n+1, 1LL);
27     for (ll i = 2; i <= n; ++i)
28         I[i] = mod(-(P/i) * I[P%i], P);
29     return I;
30 }
31 // (a*b)%m
32 ll mulmod(ll a, ll b, ll m){
33     ll x = 0, y=a%m;
34     while(b>0){
35         if(b&1) x = (x+y)%m;
36         y = (2*y)%m, b /= 2;
37     }
38     return x % m;
39 }
40 // Finds b^e % m in O(lg n) time, ensure that b <
41 // m to avoid overflow!
42 ll powmod(ll b, ll e, ll m) {
43     ll p = e<2 ? 1 : powmod((b*b)%m,e/2,m);
44     return e&1 ? p*b%m : p;
45 }
46 // Solve ax + by = c, returns false on failure.
47 bool linear_diophantine(ll a, ll b, ll c, ll &x,
48     ll &y) {
49     ll d = gcd(a, b);
50     if (c % d) {
51         return false;
52     } else {
53         x = c / d * mod_inverse(a / d, b / d);
54         y = (c - a * x) / b;
55         return true;
56     }
57 }
58 // Description: Tonelli-Shanks algorithm for
59 // modular square roots. Finds $x$ s.t. $x^2 = a
60 // \pmod p$ ($-x$ gives the other solution). 0
61 // (\log^2 p) worst case, O(\log p) for most $p$
62 ll sqrtmod(ll a, ll p) {
63     a %= p; if (a < 0) a += p;
64     if (a == 0) return 0;
65     assert(powmod(a, (p-1)/2, p) == 1); // else no
66     // solution
67     if (p % 4 == 3) return powmod(a, (p+1)/4, p);
68     // a^(n+3)/8 or 2^(n+3)/8 * 2^(n-1)/4 works if
69     // p % 8 == 5
70     ll s = p - 1, n = 2;
71     int r = 0, m;
72     while (s % 2 == 0)
73         ++r, s /= 2;
74     /// find a non-square mod p
75     while (powmod(n, (p - 1) / 2, p) != p - 1) ++n;

```

---

```

69     ll x = powmod(a, (s + 1) / 2, p);
70     ll b = powmod(a, s, p), g = powmod(n, s, p);
71     for (; r = m) {
72         ll t = b;
73         for (m = 0; m < r && t != 1; ++m)
74             t = t * t % p;
75         if (m == 0) return x;
76         ll gs = powmod(g, 1LL << (r - m - 1), p);
77         g = gs * gs % p;
78         x = x * gs % p;
79         b = b * g % p;
80     }
81 }

```

---

### 3.4.2 Mod. exponentiation Or use pow() in python

---

```

1 #include "header.h"
2 ll mod_pow(ll base, ll exp, ll mod) {
3     if (mod == 1) return 0;
4     if (exp == 0) return 1;
5     if (exp == 1) return base;
6
7     ll res = 1;
8     base %= mod;
9     while (exp) {
10        if (exp % 2 == 1) res = (res * base) % mod;
11        exp >>= 1;
12        base = (base * base) % mod;
13    }
14
15    return res % mod;
16 }

```

---

### 3.4.3 GCD Or math.gcd in python, std::gcd in C++

---

```

1 #include "header.h"
2 ll gcd(ll a, ll b) {
3     if (a == 0) return b;
4     return gcd(b % a, a);
5 }

```

---

### 3.4.4 Sieve of Eratosthenes

---

```

1 #include "header.h"
2 vl primes;
3 void getprimes(ll n) { // Up to n (not included)
4     vector<bool> p(n, true);
5     p[0] = false;
6     p[1] = false;
7     for(ll i = 0; i < n; i++) {
8         if(p[i]) {

```

---

```

9     primes.push_back(i);
10    for(11 j = i*2; j < n; j+=i) p[j] =
        false;
11 }}

```

### 3.4.5 Fibonacci % prime Starting 1,1,2,3,...

```

1 #include "header.h"
2 const 11 MOD = 1000000007;
3 unordered_map<11, 11> Fib;
4 11 fib(11 n) {
5     if (n < 2) return 1;
6     if (Fib.find(n) != Fib.end()) return Fib[n];
7     Fib[n] = (fib((n + 1) / 2) * fib(n / 2) + fib
        ((n - 1) / 2) * fib((n - 2) / 2)) % MOD;
8     return Fib[n];
9 }

```

### 3.4.6 nCk % prime

```

1 #include "header.h"
2 11 binom(11 n, 11 k) {
3     11 ans = 1;
4     for(11 i = 1; i <= min(k,n-k); ++i) ans = ans
        *(n+1-i)/i;
5     return ans;
6 }
7 11 mod_nCk(11 n, 11 k, 11 p ){
8     11 ans = 1;
9     while(n){
10         11 np = n%p, kp = k%p;
11         if(kp > np) return 0;
12         ans *= binom(np,kp);
13         n /= p; k /= p;
14     }
15     return ans;
16 }

```

## 3.5 Strings

### 3.5.1 Z alg. KMP alternative (same complexities)

```

1 #include "../header.h"
2 void Z_algorithm(const string &s, vi &Z) {
3     Z.assign(s.length(), -1);
4     int L = 0, R = 0, n = s.length();
5     for (int i = 1; i < n; ++i) {
6         if (i > R) {
7             L = R = i;
8             while (R < n && s[R - L] == s[R]) R++;
9             Z[i] = R - L; R--;

```

```

10         } else if (Z[i - L] >= R - i + 1) {
11             L = i;
12             while (R < n && s[R - L] == s[R]) R++;
13             Z[i] = R - L; R--;
14         } else Z[i] = Z[i - L];
15     }
16 }

```

### 3.5.2 KMP

```

1 #include "header.h"
2 void compute_prefix_function(string &w, vi &
    prefix) {
3     prefix.assign(w.length(), 0);
4     int k = prefix[0] = -1;
5
6     for(int i = 1; i < w.length(); ++i) {
7         while(k >= 0 && w[k + 1] != w[i]) k = prefix[
            k];
8         if(w[k + 1] == w[i]) k++;
9         prefix[i] = k;
10    }
11 }
12 vi knuth_morris_pratt(string &s, string &w) {
13     int q = -1;
14     vi prefix, positions;
15     compute_prefix_function(w, prefix);
16     for(int i = 0; i < s.length(); ++i) {
17         while(q >= 0 && w[q + 1] != s[i]) q = prefix[
            q];
18         if(w[q + 1] == s[i]) q++;
19         if(q + 1 == w.length()) {
20             // Match at position (i - w.length() + 1)
21             positions.push_back(i - w.length() +
                1);
22             q = prefix[q];
23         }
24     }
25     return positions;
26 }

```

### 3.5.3 Aho-Corasick Also can be used as Knuth-Morris-Pratt algorithm

```

1 #include "header.h"
2 map<char, int> cti;
3 int cti_size;
4 template <int ALPHABET_SIZE, int (*mp)(char)>
5 struct AC_FSM {
6     struct Node {
7         int child[ALPHABET_SIZE], failure = 0,
            match_par = -1;
8         vi match;

```

```

9         Node() { for (int i = 0; i < ALPHABET_SIZE;
            ++i) child[i] = -1; }
10    };
11    vector<Node> a;
12    vector<string> &words;
13    AC_FSM(vector<string> &words) : words(words) {
14        a.push_back(Node());
15        construct_automaton();
16    }
17    void construct_automaton() {
18        for (int w = 0, n = 0; w < words.size(); ++w,
            n = 0) {
19            for (int i = 0; i < words[w].size(); ++i) {
20                if (a[n].child[mp(words[w][i])] == -1) {
21                    a[n].child[mp(words[w][i])] = a.size();
22                    a.push_back(Node());
23                }
24                n = a[n].child[mp(words[w][i])];
25            }
26            a[n].match.push_back(w);
27        }
28        queue<int> q;
29        for (int k = 0; k < ALPHABET_SIZE; ++k) {
30            if (a[0].child[k] == -1) a[0].child[k] = 0;
31            else if (a[0].child[k] > 0) {
32                a[a[0].child[k]].failure = 0;
33                q.push(a[0].child[k]);
34            }
35        }
36        while (!q.empty()) {
37            int r = q.front(); q.pop();
38            for (int k = 0, arck; k < ALPHABET_SIZE; ++
                k) {
39                if ((arck = a[r].child[k]) != -1) {
40                    q.push(arck);
41                    int v = a[r].failure;
42                    while (a[v].child[k] == -1) v = a[v].
                        failure;
43                    a[arck].failure = a[v].child[k];
44                    a[arck].match_par = a[v].child[k];
45                    while (a[arck].match_par != -1
                        && a[a[arck].match_par].match.empty
                            ())
46                        a[arck].match_par = a[a[arck].
                            match_par].match_par;
47                }
48            }
49        }
50    }
51 }
52 void aho_corasick(string &sentence, vvi &
    matches){
53     matches.assign(words.size(), vi());
54     int state = 0, ss = 0;
55     for (int i = 0; i < sentence.length(); ++i,
        ss = state) {

```

```

56 while (a[ss].child[mp(sentence[i])] == -1)
57     ss = a[ss].failure;
58 state = a[state].child[mp(sentence[i])]
59     = a[ss].child[mp(sentence[i])];
60 for (ss = state; ss != -1; ss = a[ss].
    match_par)
61     for (int w : a[ss].match)
62         matches[w].push_back(i + 1 - words[w].
            length());
63 }
64 }
65 };
66 int char_to_int(char c) {
67     return cti[c];
68 }
69 int main() {
70     ll n;
71     string line;
72     while(getline(cin, line)) {
73         stringstream ss(line);
74         ss >> n;
75
76         vector<string> patterns(n);
77         for (auto& p: patterns) getline(cin, p);
78
79         string text;
80         getline(cin, text);
81
82         cti = {}, cti_size = 0;
83         for (auto c: text) {
84             if (not in(c, cti)) {
85                 cti[c] = cti_size++;
86             }
87         }
88         for (auto& p: patterns) {
89             for (auto c: p) {
90                 if (not in(c, cti)) {
91                     cti[c] = cti_size++;
92                 }
93             }
94         }
95
96         vvi matches;
97         AC_FSM <128+1, char_to_int> ac_fms(patterns);
98         ac_fms.aho_corasick(text, matches);
99         for (auto& x: matches) cout << x << endl;
100     }
101 }
102 }

```

### 3.5.4 Long. palin. subs Manacher - $O(n)$

```

1 #include "header.h"
2 void manacher(string &s, vi &pal) {

```

```

3 int n = s.length(), i = 1, l, r;
4 pal.assign(2 * n + 1, 0);
5 while (i < 2 * n + 1) {
6     if ((i&1) && pal[i] == 0) pal[i] = 1;
7     l = i / 2 - pal[i] / 2; r = (i-1) / 2 + pal[i]
        ] / 2;
8
9     while (l - 1 >= 0 && r + 1 < n && s[l - 1] ==
        s[r + 1])
10         --l, ++r, pal[i] += 2;
11
12     for (l = i - 1, r = i + 1; l >= 0 && r < 2 *
        n + 1; --l, ++r) {
13         if (l <= i - pal[i]) break;
14         if (l / 2 - pal[l] / 2 > i / 2 - pal[i] /
            2)
15             pal[r] = pal[l];
16         else { if (l >= 0)
17             pal[r] = min(pal[l], i + pal[i] - r);
18             break;
19         }
20     }
21     i = r;
22 } }

```

## 3.6 Geometry

### 3.6.1 essentials.cpp

```

1 #include "../header.h"
2 using C = ld; // could be ll or ld
3 constexpr C EPS = 1e-10; // change to 0 for C=ll
4 struct P { // may also be used as a 2D vector
5     C x, y;
6     P(C x = 0, C y = 0) : x(x), y(y) {}
7     P operator+ (const P &p) const { return {x + p.
        x, y + p.y}; }
8     P operator- (const P &p) const { return {x - p.
        x, y - p.y}; }
9     P operator* (C c) const { return {x * c, y * c
        }; }
10    P operator/ (C c) const { return {x / c, y / c
        }; }
11    C operator* (const P &p) const { return x*p.x +
        y*p.y; }
12    C operator^ (const P &p) const { return x*p.y -
        p.x*y; }
13    P perp() const { return P{y, -x}; }
14    C lensq() const { return x*x + y*y; }
15    ld len() const { return sqrt((ld)lensq()); }
16    static ld dist(const P &p1, const P &p2) {
17        return (p1-p2).len(); }
18    bool operator==(const P &r) const {
19        return ((*this)-r).lensq() <= EPS*EPS; }

```

```

20 };
21 C det(P p1, P p2) { return p1^p2; }
22 C det(P p1, P p2, P o) { return det(p1-o, p2-o);
    }
23 C det(const vector<P> &ps) {
24     C sum = 0; P prev = ps.back();
25     for(auto &p : ps) sum += det(p, prev), prev = p
        ;
26     return sum;
27 }
28 // Careful with division by two and C=ll
29 C area(P p1, P p2, P p3) { return abs(det(p1, p2,
    p3))/C(2); }
30 C area(const vector<P> &poly) { return abs(det(
    poly))/C(2); }
31 int sign(C c){ return (c > C(0)) - (c < C(0)); }
32 int ccw(P p1, P p2, P o) { return sign(det(p1, p2
    , o)); }
33
34 // Only well defined for C = ld.
35 P unit(const P &p) { return p / p.len(); }
36 P rotate(P p, ld a) { return P{p.x*cos(a)-p.y*sin
    (a), p.x*sin(a)+p.y*cos(a)}; }

```

### 3.6.2 Two segs. itersec.

```

1 #include "header.h"
2 #include "essentials.cpp"
3 bool intersect(P a1, P a2, P b1, P b2) {
4     if (max(a1.x, a2.x) < min(b1.x, b2.x)) return
        false;
5     if (max(b1.x, b2.x) < min(a1.x, a2.x)) return
        false;
6     if (max(a1.y, a2.y) < min(b1.y, b2.y)) return
        false;
7     if (max(b1.y, b2.y) < min(a1.y, a2.y)) return
        false;
8     bool l1 = ccw(a2, b1, a1) * ccw(a2, b2, a1) <=
        0;
9     bool l2 = ccw(b2, a1, b1) * ccw(b2, a2, b1) <=
        0;
10    return l1 && l2;
11 }

```

### 3.6.3 Convex Hull

```

1 #include "header.h"
2 #include "essentials.cpp"
3 struct ConvexHull { // O(n lg n) monotone chain.
4     size_t n;
5     vector<size_t> h, c; // Indices of the hull
        are in 'h', ccw.
6     const vector<P> &p;

```



```

7  ConvexHull(const vector<P> &p) : n(p.size()),
   c(n), p(p) {
8      std::iota(c.begin(), c.end(), 0);
9      std::sort(c.begin(), c.end(), [this](size_t l
, size_t r) -> bool { return p[l].x != p[
r].x ? p[l].x < p[r].x : p[l].y < p[r].y;
});
10     c.erase(std::unique(c.begin(), c.end(), [this
](size_t l, size_t r) { return p[l] == p[
r]; }), c.end());
11     for (size_t s = 1, r = 0; r < 2; ++r, s = h.
size()) {
12         for (size_t i : c) {
13             while (h.size() > s && ccw(p[h.end()
[-2]], p[h.end()[-1]], p[i]) <= 0)
14                 h.pop_back();
15             h.push_back(i);
16         }
17         reverse(c.begin(), c.end());
18     }
19     if (h.size() > 1) h.pop_back();
20 }
21 size_t size() const { return h.size(); }
22 template <class T, void U(const P &, const P &,
const P &, T &>
23 void rotating_calipers(T &ans) {
24     if (size() <= 2)
25         U(p[h[0]], p[h.back()], p[h.back()], ans);
26     else
27         for (size_t i = 0, j = 1, s = size(); i < 2
* s; ++i) {
28             while (det(p[h[(i + 1) % s]] - p[h[i % s
]], p[h[(j + 1) % s]] - p[h[j]]) >=
0)
29                 j = (j + 1) % s;
30             U(p[h[i % s]], p[h[(i + 1) % s]], p[h[j
]], ans);
31         }
32 }
33 };
34 // Example: furthest pair of points. Now set ans
= OLL and call
35 // ConvexHull(pts).rotating_calipers<ll, update>(
ans);
36 void update(const P &p1, const P &p2, const P &o,
ll &ans) {
37     ans = max(ans, (ll)max((p1 - o).lensq(), (p2 -
o).lensq()));
38 }
39 int main() {
40     ios::sync_with_stdio(false); // do not use
cout + printf
41     cin.tie(NULL);
42
43     int n;

```

```

44     cin >> n;
45     while (n) {
46         vector<P> ps;
47         int x, y;
48         for (int i = 0; i < n; i++) {
49             cin >> x >> y;
50             ps.push_back({x, y});
51         }
52
53         ConvexHull ch(ps);
54         cout << ch.h.size() << endl;
55         for(auto& p: ch.h) {
56             cout << ps[p].x << " " << ps[p].y <<
endl;
57         }
58         cin >> n;
59     }
60
61     return 0;
62 }

```

## 3.7 Other Algorithms

### 3.7.1 2-sat

```

1  #include "../header.h"
2  #include "../Graphs/tarjan.cpp"
3  struct TwoSAT {
4      int n;
5      vvi imp; // implication graph
6      Tarjan tj;
7
8      TwoSAT(int _n) : n(_n), imp(2 * _n, vi()), tj(
imp) { }
9
10     // Only copy the needed functions:
11     void add_implies(int c1, bool v1, int c2, bool
v2) {
12         int u = 2 * c1 + (v1 ? 1 : 0),
13             v = 2 * c2 + (v2 ? 1 : 0);
14         imp[u].push_back(v); // u => v
15         imp[v^1].push_back(u^1); // -v => -u
16     }
17     void add_equivalence(int c1, bool v1, int c2,
bool v2) {
18         add_implies(c1, v1, c2, v2);
19         add_implies(c2, v2, c1, v1);
20     }
21     void add_or(int c1, bool v1, int c2, bool v2) {
22         add_implies(c1, !v1, c2, v2);
23     }
24     void add_and(int c1, bool v1, int c2, bool v2)
{
25         add_true(c1, v1); add_true(c2, v2);

```

```

26     }
27     void add_xor(int c1, bool v1, int c2, bool v2)
{
28         add_or(c1, v1, c2, v2);
29         add_or(c1, !v1, c2, !v2);
30     }
31     void add_true(int c1, bool v1) {
32         add_implies(c1, !v1, c1, v1);
33     }
34
35     // on true: a contains an assignment.
36     // on false: no assignment exists.
37     bool solve(vb &a) {
38         vi com;
39         tj.find_sccs(com);
40         for (int i = 0; i < n; ++i)
41             if (com[2 * i] == com[2 * i + 1])
42                 return false;
43
44         vvi bycom(com.size());
45         for (int i = 0; i < 2 * n; ++i)
46             bycom[com[i]].push_back(i);
47
48         a.assign(n, false);
49         vb vis(n, false);
50         for(auto &&component : bycom){
51             for (int u : component) {
52                 if (vis[u / 2]) continue;
53                 vis[u / 2] = true;
54                 a[u / 2] = (u % 2 == 1);
55             }
56         }
57         return true;
58     }
59 };

```

### 3.7.2 Finite field For FFT

```

1  #include "header.h"
2  #include "../NumberTheory/elementary.cpp"
3  template<ll p, ll w> // prime, primitive root
4  struct Field { using T = Field; ll x; Field(ll x
=0) : x{x} {}
5      T operator+(T r) const { return {(x+r.x)%p}; }
6      T operator-(T r) const { return {(x-r.x+p)%p};
}
7      T operator*(T r) const { return {(x*r.x)%p}; }
8      T operator/(T r) const { return (*this)*r.inv()
; }
9      T inv() const { return {mod_inverse(x,p)}; }
10     static T root(ll k) { assert((p-1)%k==0);
// (p-1)%k == 0?
11     auto r = powmod(w, (p-1)/abs(k), p); // k-
th root of unity

```



```

12     return k>=0 ? T{r} : T{r}.inv();
13 }
14 bool zero() const { return x == 0LL; }
15 };
16 using F1 = Field<1004535809,3 >;
17 using F2 = Field<1107296257,10>; // 1<<30 + 1<<25
18     + 1
19 using F3 = Field<2281701377,3 >; // 1<<31 + 1<<27
20     + 1

```

---

### 3.7.3 Complex field For FFR

---

```

1 #include "header.h"
2 const double m_pi = M_PI*64x;
3 struct Complex { using T = Complex; double u,v;
4     Complex(double u=0, double v=0) : u{u}, v{v} {}
5     T operator+(T r) const { return {u+r.u, v+r.v};
6     }
7     T operator-(T r) const { return {u-r.u, v-r.v};
8     }
9     T operator*(T r) const { return {u*r.u - v*r.v,
10         u*r.v + v*r.u}; }
11     T operator/(T r) const {
12         auto norm = r.u*r.u + r.v*r.v;
13         return {(u*r.u + v*r.v)/norm, (v*r.u - u*r.v)
14             /norm};
15     }
16     T operator*(double r) const { return T{u*r, v*r};
17     };
18     T operator/(double r) const { return T{u/r, v/r};
19     };
20     T inv() const { return T{1,0}/ *this; }
21     T conj() const { return T{u, -v}; }
22     static T root(11 k){ return {cos(2*m_pi/k), sin
23         (2*m_pi/k)}; }
24     bool zero() const { return max(abs(u), abs(v))
25         < 1e-6; }
26 };

```

---

### 3.7.4 FFT

```

1 #include "header.h"
2 #include "complex_field.cpp"
3 #include "fin_field.cpp"
4 void brinc(int &x, int k) {
5     int i = k - 1, s = 1 << i;
6     x ^= s;
7     if ((x & s) != s) {
8         --i; s >>= 1;
9         while (i >= 0 && ((x & s) == s))
10             x = x &~ s, --i, s >>= 1;
11         if (i >= 0) x |= s;
12     }

```

```

13 }
14 using T = Complex; // using T=F1,F2,F3
15 vector<T> roots;
16 void root_cache(int N) {
17     if (N == (int)roots.size()) return;
18     roots.assign(N, T{0});
19     for (int i = 0; i < N; ++i)
20         roots[i] = ((i&-i) == i)
21             ? T{cos(2.0*m_pi*i/N), sin(2.0*m_pi*i/N)}
22             : roots[i&-i] * roots[i-(i&-i)];
23 }
24 void fft(vector<T> &A, int p, bool inv = false) {
25     int N = 1<<p;
26     for(int i = 0, r = 0; i < N; ++i, brinc(r, p))
27         if (i < r) swap(A[i], A[r]);
28     // Uncomment to precompute roots (for T=Complex)
29     // . Slower but more precise.
30     // root_cache(N);
31     // , sh=p-1, --sh
32     for (int m = 2; m <= N; m <= 1) {
33         T w, w_m = T::root(inv ? -m : m);
34         for (int k = 0; k < N; k += m) {
35             w = T{1};
36             for (int j = 0; j < m/2; ++j) {
37                 // T w = (!inv ? roots[j<<sh] : roots[j<<
38                 // sh].conj());
39                 T t = w * A[k + j + m/2];
40                 A[k + j + m/2] = A[k + j] - t;
41                 A[k + j] = A[k + j] + t;
42                 w = w * w_m;
43             }
44         }
45     }
46     if(inv){ T inverse = T(N).inv(); for(auto &x :
47         A) x = x*inverse; }
48 // convolution leaves A and B in frequency domain
49 // state
50 // C may be equal to A or B for in-place
51 // convolution
52 void convolution(vector<T> &A, vector<T> &B,
53     vector<T> &C){
54     int s = A.size() + B.size() - 1;
55     int q = 32 - __builtin_clz(s-1), N=1<<q; //
56     fails if s=1
57     A.resize(N,{}); B.resize(N,{}); C.resize(N,{});
58     fft(A, q, false); fft(B, q, false);
59     for (int i = 0; i < N; ++i) C[i] = A[i] * B[i];
60     fft(C, q, true); C.resize(s);
61 }
62 void square_inplace(vector<T> &A) {
63     int s = 2*A.size()-1, q = 32 - __builtin_clz(s
64         -1), N=1<<q;
65     A.resize(N,{}); fft(A, q, false);
66     for(auto &x : A) x = x*x;

```

```

60     fft(A, q, true); A.resize(s);
61 }

```

### 3.7.5 Polyn. inv. div.

```

1 #include "header.h"
2 #include "fft.cpp"
3 vector<T> &rev(vector<T> &A) { reverse(A.begin(),
4     A.end()); return A; }
5 void copy_into(const vector<T> &A, vector<T> &B,
6     size_t n) {
7     std::copy(A.begin(), A.begin()+min({n, A.size()
8         , B.size()}), B.begin());
9 }
10 // Multiplicative inverse of A modulo x^n.
11 // Requires A[0] != 0!!
12 vector<T> inverse(const vector<T> &A, int n) {
13     vector<T> Ai{A[0].inv()};
14     for (int k = 0; (1<<k) < n; ++k) {
15         vector<T> As(4<<k, T(0)), Ais(4<<k, T(0));
16         copy_into(A, As, 2<<k); copy_into(Ai, Ais, Ai
17             .size());
18         fft(As, k+2, false); fft(Ais, k+2, false);
19         for (int i = 0; i < (4<<k); ++i) As[i] = As[i]
20             *Ais[i]*Ais[i];
21         fft(As, k+2, true); Ai.resize(2<<k, {});
22         for (int i = 0; i < (2<<k); ++i) Ai[i] = T(2)
23             * Ai[i] - As[i];
24     }
25     Ai.resize(n);
26     return Ai;
27 }
28 // Polynomial division. Returns {Q, R} such that
29 // A = QB+R, deg R < deg B.
30 // Requires that the leading term of B is nonzero
31 pair<vector<T>, vector<T>> divmod(const vector<T>
32     &A, const vector<T> &B) {
33     size_t n = A.size()-1, m = B.size()-1;
34     if (n < m) return {vector<T>(1, T(0)), A};
35     vector<T> X(A), Y(B), Q, R;
36     convolution(rev(X), Y = inverse(rev(Y), n-m+1),
37         Q);
38     Q.resize(n-m+1); rev(Q);
39     X.resize(Q.size()), copy_into(Q, X, Q.size());
40     Y.resize(B.size()), copy_into(B, Y, B.size());
41     convolution(X, Y, X);
42     R.resize(m), copy_into(A, R, m);
43     for (size_t i = 0; i < m; ++i) R[i] = R[i] - X[i]
44         *i;
45     while (R.size() > 1 && R.back().zero()) R.
46         pop_back();

```

```

38 return {Q, R};
39 }
40 vector<T> mod(const vector<T> &A, const vector<T>
    &B) {
41     return divmod(A, B).second;
42 }

```

**3.7.6 Linear recurs.** Given a linear recurrence of the form

$$a_n = \sum_{i=0}^{k-1} c_i a_{n-i-1}$$

this code computes  $a_n$  in  $O(k \log k \log n)$  time.

```

1 #include "header.h"
2 #include "poly.cpp"
3 // x^k mod f
4 vector<T> xmod(const vector<T> f, ll k) {
5     vector<T> r{T(1)};
6     for (int b = 62; b >= 0; --b) {
7         if (r.size() > 1)
8             square_inplace(r), r = mod(r, f);
9         if ((k >> b) & 1) {
10             r.insert(r.begin(), T(0));
11             if (r.size() == f.size()) {
12                 T c = r.back() / f.back();
13                 for (size_t i = 0; i < f.size(); ++i)
14                     r[i] = r[i] - c * f[i];
15                 r.pop_back();
16             }
17         }
18     }
19     return r;
20 }
21 // Given A[0,k) and C[0, k), computes the n-th
    term of:
22 // A[n] = \sum_i C[i] * A[n-i-1]
23 T nth_term(const vector<T> &A, const vector<T> &C
    , ll n) {
24     int k = (int)A.size();
25     if (n < k) return A[n];
26
27     vector<T> f(k+1, T{1});
28     for (int i = 0; i < k; ++i)
29         f[i] = T{-1} * C[k-i-1];
30     f = xmod(f, n);
31
32     T r = T{0};
33     for (int i = 0; i < k; ++i)
34         r = r + f[i] * A[i];
35     return r;
36 }

```

### 3.7.7 Convolution Precise up to 9e15

```

1 #include "header.h"
2 #include "fft.cpp"
3 void convolution_mod(const vi &A, const vi &B, ll
    MOD, vi &C) {
4     int s = A.size() + B.size() - 1; ll m15 = (1LL
        <<15)-1LL;
5     int q = 32 - __builtin_clz(s-1), N=1<<q; //
        fails if s=1
6     vector<T> Ac(N), Bc(N), R1(N), R2(N);
7     for (size_t i = 0; i < A.size(); ++i) Ac[i] = T
        {A[i]&m15, A[i]>>15};
8     for (size_t i = 0; i < B.size(); ++i) Bc[i] = T
        {B[i]&m15, B[i]>>15};
9     fft(Ac, q, false); fft(Bc, q, false);
10    for (int i = 0, j = 0; i < N; ++i, j = (N-1)&(N
        -i)) {
11        T as = (Ac[i] + Ac[j].conj()) / 2;
12        T al = (Ac[i] - Ac[j].conj()) / T{0, 2};
13        T bs = (Bc[i] + Bc[j].conj()) / 2;
14        T bl = (Bc[i] - Bc[j].conj()) / T{0, 2};
15        R1[i] = as*bs + al*bl*T{0,1}, R2[i] = as*bl +
            al*bs;
16    }
17    fft(R1, q, true); fft(R2, q, true);
18    ll p15 = (1LL<<15)%MOD, p30 = (1LL<<30)%MOD; C.
        resize(s);
19    for (int i = 0; i < s; ++i) {
20        ll l = llround(R1[i].u), m = llround(R2[i].u)
            , h = llround(R1[i].v);
21        C[i] = (l + m*p15 + h*p30) % MOD;
22    }
23 }

```

### 3.7.8 Partitions of $n$ Finds all possible partitions of a number

```

1 #include "header.h"
2 void printArray(int p[], int n) {
3     for (int i = 0; i < n; i++)
4         cout << p[i] << " ";
5     cout << endl;
6 }
7 void printAllUniqueParts(int n) {
8     int p[n]; // array to store a partition
9     int k = 0; // idx of last element in a
        partition
10    p[k] = n;
11
12    // The loop stops when the current partition
        has all 1s
13    while (true) {
14        printArray(p, k + 1);

```

```

15    int rem_val = 0;
16    while (k >= 0 && p[k] == 1) {
17        rem_val += p[k];
18        k--;
19    }
20    // no more partitions
21    if (k < 0) return;
22
23    p[k]--;
24    rem_val++;
25
26    // sorted order is violated (fix)
27    while (rem_val > p[k]) {
28        p[k + 1] = p[k];
29        rem_val = rem_val - p[k];
30        k++;
31    }
32
33    p[k + 1] = rem_val;
34    k++;
35 }
36 }

```

**3.7.9 Ternary search** Find the smallest  $i$  in  $[a, b]$  that maximizes  $f(i)$ , assuming that  $f(a) < \dots < f(i) \geq \dots \geq f(b)$ . To reverse which of the sides allows non-strict inequalities, change the  $<$  marked with (A) to  $\leq$ , and reverse the loop at (B). To minimize  $f$ , change it to  $>$ , also at (B).  $O(\log(b-a))$

```

1 // Usage: int ind = ternSearch(0,n-1,[\&](int i){
    return a[i];});
2 #include "../Numerical/template.cpp"
3 template<class F>
4 int ternSearch(int a, int b, F f) {
5     assert(a <= b);
6     while (b - a >= 5) {
7         int mid = (a + b) / 2;
8         if (f(mid) < f(mid+1)) a = mid; // (A)
9         else b = mid+1;
10    }
11    rep(i,a+1,b+1) if (f(a) < f(i)) a = i; // (B)
12    return a;
13 }

```

**3.7.10 Hashing** Also see Primes in Other Mathematics. For a proper rolling hash over a string, fix the modulus, and draw the base  $b$  uniformly at random from  $\{0, 1, \dots, p-1\}$ . Note that when comparing rolling hashes of strings of different lengths, it is useful to hash

the empty character to 0, and hash all actual characters to nonzero values. Some primes:

$$10^3 + \{-9, -3, 9, 13\},$$

$$10^6 + \{-17, 3, 33\}, 10^9 + \{7, 9, 21, 33, 87\}$$

---

```

1 #include "../header.h"
2 template <class T>
3 size_t hash_combine(size_t seed, const T &v) {
4     return seed ^ (std::hash<T>()(v) + 0x9e3779b9 +
5         (seed << 6) + (seed >> 2));
6 }
7 namespace std {
8     template <typename U, typename V>
9     struct hash<pair<U, V>> {
10         size_t operator()(const pair<U, V> &k) const {
11             return hash_combine(hash_combine(0, k.first),
12                 k.second);
13         }
14     };
15 }

```

---

## 3.8 Other Data Structures

### 3.8.1 Disjoint set (i.e. union-find)

---

```

1 template <typename T>
2 class DisjointSet {
3     typedef T * iterator;
4     T *parent, n, *rank;
5     public:
6         // O(n), assumes nodes are [0, n)
7         DisjointSet(T n) {
8             this->parent = new T[n];
9             this->n = n;
10            this->rank = new T[n];
11            for (T i = 0; i < n; i++) {
12                parent[i] = i;
13                rank[i] = 0;
14            }
15        }
16
17        // O(log n)
18        T find_set(T x) {
19            if (x == parent[x]) return x;
20            return parent[x] = find_set(parent[x]);
21        }
22
23        // O(log n)
24        void union_sets(T x, T y) {
25            x = this->find_set(x);

```

```

26            y = this->find_set(y);
27
28            if (x == y) return;
29            if (rank[x] < rank[y]) {
30                T z = x;
31                x = y;
32                y = z;
33            }
34            parent[y] = x;
35            if (rank[x] == rank[y]) rank[x]++;
36        }
37 };

```

---

**3.8.2 Fenwick tree** (i.e. BIT) eff. update + prefix sum calc. Can be generalized to arbitrary dimensions by duplicating loops.

---

```

1 // #include "header.h"
2 template < class T >
3 struct FenwickTree { // use 1 based indices !!!
4     int n ; vector<T> tree ;
5     FenwickTree ( int n ) : n ( n ) { tree .
6         assign ( n + 1 , 0 ) ; }
7     T query ( int l , int r ) { return query ( r
8         ) - query ( l - 1 ) ; }
9     T query ( int r ) {
10         T s = 0 ;
11         for ( ; r > 0 ; r -= ( r & ( - r ) ) ) s +=
12             tree [ r ] ;
13         return s ;
14     }
15     void update ( int i , T v ) {
16         for ( ; i <= n ; i += ( i & ( - i ) ) )
17             tree [ i ] += v ;
18     }
19 };

```

---

### 3.8.3 Trie

---

```

1 #include "header.h"
2 const int ALPHABET_SIZE = 26;
3 inline int mp(char c) { return c - 'a'; }
4 struct Node {
5     Node* ch[ALPHABET_SIZE];
6     bool isleaf = false;
7     Node() {
8         for(int i = 0; i < ALPHABET_SIZE; ++i) ch[i]
9             = nullptr;
10     }
11     void insert(string &s, int i = 0) {
12         if (i == s.length()) isleaf = true;

```

```

13     else {
14         int v = mp(s[i]);
15         if (ch[v] == nullptr)
16             ch[v] = new Node();
17         ch[v]->insert(s, i + 1);
18     }
19 }
20
21 bool contains(string &s, int i = 0) {
22     if (i == s.length()) return isleaf;
23     else {
24         int v = mp(s[i]);
25         if (ch[v] == nullptr) return false;
26         else return ch[v]->contains(s, i + 1);
27     }
28 }
29
30 void cleanup() {
31     for (int i = 0; i < ALPHABET_SIZE; ++i)
32         if (ch[i] != nullptr) {
33             ch[i]->cleanup();
34             delete ch[i];
35         }
36 }
37 };

```

---

**3.8.4 Treap** A binary tree whose nodes contain two values, a key and a priority, such that the key keeps the BST property

---

```

1 #include "header.h"
2 struct Node {
3     ll v;
4     int sz, pr;
5     Node *l = nullptr, *r = nullptr;
6     Node(ll val) : v(val), sz(1) { pr = rand(); }
7 };
8 int size(Node *p) { return p ? p->sz : 0; }
9 void update(Node* p) {
10     if (!p) return;
11     p->sz = 1 + size(p->l) + size(p->r);
12     // Pull data from children here
13 }
14 void propagate(Node *p) {
15     if (!p) return;
16     // Push data to children here
17 }
18 void merge(Node *&t, Node *l, Node *r) {
19     propagate(l), propagate(r);
20     if (!l) t = r;
21     else if (!r) t = l;
22     else if (l->pr > r->pr)
23         merge(l->r, l->r, r), t = l;
24     else merge(r->l, l, r->l), t = r;

```

```

25 update(t);
26 }
27 void spliti(Node *t, Node *&l, Node *&r, int
    index) {
28     propagate(t);
29     if (!t) { l = r = nullptr; return; }
30     int id = size(t->l);
31     if (index <= id) // id \in [index, \infty), so
        move it right
32     spliti(t->l, l, t->l, index), r = t;
33     else
34     spliti(t->r, t->r, r, index - id), l = t;
35     update(t);
36 }
37 void splitv(Node *t, Node *&l, Node *&r, ll val)
    {
38     propagate(t);
39     if (!t) { l = r = nullptr; return; }
40     if (val <= t->v) // t->v \in [val, \infty), so
        move it right
41     splitv(t->l, l, t->l, val), r = t;
42     else
43     splitv(t->r, t->r, r, val), l = t;
44     update(t);
45 }
46 void clean(Node *p) {
47     if (p) { clean(p->l), clean(p->r); delete p; }
48 }

```

### 3.8.5 Segment tree

```

1 #include "../header.h"
2 // example: SegmentTree<int, min> st(n, INT_MAX);
3 const int& addOp(const int& a, const int& b) {
4     static int result;
5     result = a + b;
6     return result;
7 }
8 template <class T, const T&(*op)(const T&, const
    T&>
9 struct SegmentTree {
10     int n; vector<T> tree; T id;
11     SegmentTree(int _n, T _id) : n(_n), tree(2 * n,
        _id), id(_id) { }
12     void update(int i, T val) {
13         for (tree[i+n] = val, i = (i+n)/2; i > 0; i
            /= 2)
14             tree[i] = op(tree[2*i], tree[2*i+1]);
15     }
16     T query(int l, int r) {
17         T lhs = T(id), rhs = T(id);
18         for (l += n, r += n; l < r; l >>= 1, r >>= 1)
            {
19             if (l & 1) lhs = op(lhs, tree[l++]);

```

```

20         if (!r & 1) rhs = op(tree[r--], rhs);
21     }
22     return op(l == r ? op(lhs, tree[l]) : lhs,
        rhs);
23 }
24 };

```

### 3.8.6 Lazy segment tree Optimizes range updates

```

1 #include "../header.h"
2 using T=int; using U=int; using I=int; //
    exclusive right bounds
3 T t_id; U u_id;
4 T op(T a, T b){ return a+b; }
5 void join(U &a, U b){ a+=b; }
6 void apply(T &t, U u, int x){ t+=x*u; }
7 T convert(const I &i){ return i; }
8 struct LazySegmentTree {
9     struct Node { int l, r, lc, rc; T t; U u;
10         Node(int l, int r, T t=t_id):l(l),r(r),lc(-1),rc(-1),t(t),u(u_id){}
11     };
12     int N; vector<Node> tree; vector<I> &init;
13     LazySegmentTree(vector<I> &init) : N(init.size()
        ), init(init){
14         tree.reserve(2*N-1); tree.push_back({0,N});
15         build(0, 0, N);
16     }
17     void build(int i, int l, int r) { auto &n =
        tree[i];
18         if (r > l+1) { int m = (l+r)/2;
19             n.lc = tree.size(); n.rc = n.lc+1;
20             tree.push_back({l,m}); tree.push_back({m,
                r});
21             build(n.lc,l,m); build(n.rc,m,r);
22             n.t = op(tree[n.lc].t, tree[n.rc].t);
23             } else n.t = convert(init[l]);
24     }
25     void push(Node &n, U u){ apply(n.t, u, n.r-n.l)
        ; join(n.u,u); }
26     void push(Node &n){push(tree[n.lc],n.u);push(
        tree[n.rc],n.u);n.u=u_id;}
27     T query(int l, int r, int i = 0) { auto &n =
        tree[i];
28         if(r <= n.l || n.r <= l) return t_id;
29         if(l <= n.l && n.r <= r) return n.t;
30         return push(n, op(query(l,r,n.lc),query(l,r,
            n.rc)));
31     }
32     void update(int l, int r, U u, int i = 0) {
33         auto &n = tree[i];
34         if(r <= n.l || n.r <= l) return;
35         if(l <= n.l && n.r <= r) return push(n,u);
36         push(n); update(l,r,u,n.lc);update(l,r,u,n.
            rc);

```

```

35     n.t = op(tree[n.lc].t, tree[n.rc].t);
36 }
37 };

```

**3.8.7 Dynamic segment tree** Sparse, i.e., large values, i.e., not stored as an array

```

1 #include "../header.h"
2 using T=ll; using U=ll; // exclusive
    right bounds
3 T t_id; U u_id;
4 T op(T a, T b){ return a+b; }
5 void join(U &a, U b){ a+=b; }
6 void apply(T &t, U u, int x){ t+=x*u; }
7 T part(T t, int r, int p){ return t/r*p; }
8 struct DynamicSegmentTree {
9     struct Node { int l, r, lc, rc; T t; U u;
10         Node(int l, int r):l(l),r(r),lc(-1),rc(-1),t(
            t_id),u(u_id){}
11     };
12     vector<Node> tree;
13     DynamicSegmentTree(int N) { tree.push_back({0,N
        }); }
14     void push(Node &n, U u){ apply(n.t, u, n.r-n.l)
        ; join(n.u,u); }
15     void push(Node &n){push(tree[n.lc],n.u);push(
        tree[n.rc],n.u);n.u=u_id;}
16     T query(int l, int r, int i = 0) { auto &n =
        tree[i];
17         if(r <= n.l || n.r <= l) return t_id;
18         if(l <= n.l && n.r <= r) return n.t;
19         if(n.lc < 0) return part(n.t, n.r-n.l, min(n.
            r,r)-max(n.l,l));
20         return push(n, op(query(l,r,n.lc),query(l,r,
            n.rc)));
21     }
22     void update(int l, int r, U u, int i = 0) {
23         auto &n = tree[i];
24         if(r <= n.l || n.r <= l) return;
25         if(l <= n.l && n.r <= r) return push(n,u);
26         if(n.lc < 0) { int m = (n.l + n.r) / 2;
27             n.lc = tree.size(); n.rc = n.lc+1;
28             tree.push_back({tree[i].l, m}); tree.
                push_back({m, tree[i].r});
29         }
30         push(tree[i]); update(l,r,u,tree[i].lc);
31         update(l,r,u,tree[i].rc);
32         tree[i].t = op(tree[tree[i].lc].t, tree[tree[
            i].rc].t);
33     }
34 };

```

### 3.8.8 Suffix array

```

1 #include "../header.h"
2 struct SuffixArray {
3     string s;
4     int n;
5     vvi P;
6     SuffixArray(string &s) : s(s), n(s.length())
7     { construct(); }
8     void construct() {
9         P.push_back(vi(n, 0));
10        compress();
11        vi occ(n + 1, 0), s1(n, 0), s2(n, 0);
12        for (int k = 1, cnt = 1; cnt / 2 < n; ++k,
13            cnt *= 2) {
14            P.push_back(vi(n, 0));
15            fill(occ.begin(), occ.end(), 0);
16            for (int i = 0; i < n; ++i)
17                occ[i+cnt<n ? P[k-1][i+cnt]+1 : 0]++;
18            partial_sum(occ.begin(), occ.end(), occ.
19                begin());
20            for (int i = n - 1; i >= 0; --i)
21                s1[--occ[i+cnt<n ? P[k-1][i+cnt]+1 : 0]]
22                    = i;
23            fill(occ.begin(), occ.end(), 0);
24            for (int i = 0; i < n; ++i)
25                occ[P[k-1][s1[i]]]++;
26            partial_sum(occ.begin(), occ.end(), occ.
27                begin());
28            for (int i = n - 1; i >= 0; --i)
29                s2[--occ[P[k-1][s1[i]]]] = s1[i];
30            for (int i = 1; i < n; ++i) {
31                P[k][s2[i]] = same(s2[i], s2[i - 1], k,
32                    cnt)
33                    ? P[k][s2[i - 1]] : i;
34            }
35        }
36    }
37    bool same(int i, int j, int k, int l) {
38        return P[k - 1][i] == P[k - 1][j]
39            && (i + l < n ? P[k - 1][i + l] : -1)
40            == (j + l < n ? P[k - 1][j + l] : -1);
41    }
42    void compress() {
43        vi cnt(256, 0);
44        for (int i = 0; i < n; ++i) cnt[s[i]]++;
45        for (int i = 0, mp = 0; i < 256; ++i)
46            if (cnt[i] > 0) cnt[i] = mp++;
47        for (int i = 0; i < n; ++i) P[0][i] = cnt[s[i]
48            ];
49    }
50    const vi &get_array() { return P.back(); }
51    int lcp(int x, int y) {
52        int ret = 0;
53        if (x == y) return n - x;
54        for (int k = P.size() - 1; k >= 0 && x < n &&
55            y < n; --k)

```

```

48        if (P[k][x] == P[k][y]) {
49            x += 1 << k;
50            y += 1 << k;
51            ret += 1 << k;
52        }
53        return ret;
54    }
55 };

```

### 3.8.9 Suffix tree

```

1 #include "../header.h"
2 using T = char;
3 using M = map<T,int>; // or array<T,ALPHABET_SIZE>
4 using V = string; // could be vector<T> as well
5 using It = V::const_iterator;
6 struct Node{
7     It b, e; M edges; int link; // end is exclusive
8     Node(It b, It e) : b(b), e(e), link(-1) {}
9     int size() const { return e-b; }
10 };
11 struct SuffixTree{
12     const V &s; vector<Node> t;
13     int root,n,len,remainder,llink; It edge;
14     SuffixTree(const V &s) : s(s) { build(); }
15     int add_node(It b, It e){ return t.push_back({b
16         ,e}), t.size()-1; }
17     int add_node(It b){ return add_node(b,s.end()); }
18     void link(int node){ if(llink) t[llink].link =
19         node; llink = node; }
20     void build(){
21         len = remainder = 0; edge = s.begin();
22         n = root = add_node(s.begin(), s.begin());
23         for(auto i = s.begin(); i != s.end(); ++i){
24             ++remainder; llink = 0;
25             while(remainder){
26                 if(len == 0) edge = i;
27                 if(t[n].edges[*edge] == 0){
28                     t[n].edges[*edge] = add_node(i); link(n
29                         );
30                 } else {
31                     auto x = t[n].edges[*edge];
32                     if(len >= t[x].size()){
33                         len -= t[x].size(); edge += t[x].size
34                             (); n = x;
35                     }
36                     continue;
37                 }
38                 if(*(t[x].b + len) == *i){
39                     ++len; link(n); break;
40                 }
41                 auto split = add_node(t[x].b, t[x].b+
42                     len);

```

```

37         t[n].edges[*edge] = split;
38         t[x].b += len;
39         t[split].edges[*i] = add_node(i);
40         t[split].edges[*t[x].b] = x;
41         link(split);
42     }
43     --remainder;
44     if(n == root && len > 0)
45         --len, edge = i - remainder + 1;
46     else n = t[n].link > 0? t[n].link: root;
47 }
48 }
49 }
50 };

```

### 3.8.10 Suffix automaton

```

1 #include "../header.h"
2 using T = char; using M = map<T,int>; using V =
3     string;
4 struct Node { // s: start, len: length, link:
5     suffix link, e: edges
6     int s, len, link; M e; bool term; // term
7     : terminal node?
8     Node(int s, int len, int link=-1):s(s), len(len
9         ), link(link), term(0) {}
10 };
11 struct SuffixAutomaton{
12     const V &s; vector<Node> t; int l; // string;
13     tree; last added state
14     SuffixAutomaton(const V &s) : s(s) { build(); }
15     void build(){
16         l = t.size(); t.push_back({0,-1}); //
17         root node
18         for(auto c : s){
19             int p=l, x=t.size(); t.push_back({0,t[l].
20                 len + 1}); // new node
21             while(p>0 && t[p].e[c] == 0) t[p].e[c] = x
22                 , p = t[p].link;
23             if(p<0) t[x].link = 0; // at
24             root
25             else {
26                 int q = t[p].e[c]; // the c-
27                 child of q
28                 if(t[q].len == t[p].len + 1)t[x].link =
29                     q;
30                 else {
31                     // cloning of q
32                     int cl = t.size(); t.push_back(t[q]);
33                     t[cl].len = t[p].len + 1;
34                     t[cl].s = t[q].s + t[q].len - t[p].len
35                         - 1;
36                     t[x].link = t[q].link = cl;
37                     while(p >= 0 && t[p].e.count(c) > 0 &&
38                         t[p].e[c] == q)

```

```

25         t[p].e[c] = cl, p = t[p].link; //
           relink suffix
26     }
27 }
28 l = x; // update last
29 }
30 while(l>=0) t[l].term = true, l = t[l].link;
31 }
32 };

```

### 3.8.11 UnionFind

```

1 #include "header.h"
2 struct UnionFind {
3     std::vector<int> par, rank, size;
4     int c;
5     UnionFind(int n) : par(n), rank(n, 0), size(n,
6         1), c(n) {
7         for(int i = 0; i < n; ++i) par[i] = i;
8     }
9     int find(int i) { return (par[i] == i ? i : (
10         par[i] = find(par[i]))); }
11     bool same(int i, int j) { return find(i) ==
12         find(j); }
13     int get_size(int i) { return size[find(i)]; }
14     int count() { return c; }
15     int merge(int i, int j) {
16         if((i = find(i)) == (j = find(j))) return -1;
17         --c;
18         if(rank[i] > rank[j]) swap(i, j);
19         par[i] = j;
20         size[j] += size[i];
21         if(rank[i] == rank[j]) rank[j]++;
22         return j;
23 };

```

**3.8.12 Indexed set** Similar to set, but allows accessing elements by index using `find_by_order()` in  $O(\log n)$

```

1 #include "../header.h"
2 #include <ext/pb_ds/assoc_container.hpp>
3 using namespace __gnu_pbds;
4 using namespace std;
5 typedef tree<int, null_type, less<int>, rb_tree_tag,
6     tree_order_statistics_node_update>
7     indexed_set;

```

**3.8.13 Order Statistics Tree** A set (not multiset!) with support for finding the  $n$ 'th element, and find-

ing the index of an element. To get a map, change `null_type.O(log N)`

```

1 #include <bits/extc++.h> // !!!!
2 using namespace __gnu_pbds;
3 using namespace std;
4
5 template<class T>
6 using Tree = tree<T, null_type, less<T>,
7     rb_tree_tag,
8     tree_order_statistics_node_update>;
9
10 void example() {
11     Tree<int> t, t2; t.insert(8);
12     auto it = t.insert(10).first;
13     assert(it == t.lower_bound(9));
14     assert(t.order_of_key(10) == 1);
15     assert(t.order_of_key(11) == 2);
16     assert(*t.find_by_order(0) == 8);
17     t.join(t2); // assuming T < T2 or T > T2, merge
18                 t2 into t
19 }

```

**3.8.14 Range minimum queries** Answers range minimum queries in constant time after  $O(V \log V)$  preproc.

```

1 template<class T>
2 struct RMQ {
3     vector<vector<T>> jmp;
4     RMQ(const vector<T>& V) : jmp(1, V) {
5         for (int pw = 1, k = 1; pw * 2 <= sz(V); pw
6             *= 2, ++k) {
7             jmp.emplace_back(sz(V) - pw * 2 + 1);
8             rep(j, 0, sz(jmp[k]))
9                 jmp[k][j] = min(jmp[k-1][j], jmp[k-1][j+pw]);
10         }
11     }
12     T query(int a, int b) { // returns min(V[a],
13         ..., V[b-1])
14         assert(a < b); // or return inf if a == b
15         int dep = 31 - __builtin_clz(b-a);
16         return min(jmp[dep][a], jmp[dep][b-(1<<dep)]);
17     }
18 };

```

### 3.8.15 Pareto Front

```

1 #include "../header.h"
2 struct pareto_front {
3     map<ll, ll> m;
4     void insert(ll a, ll b) {
5         auto it = m.lower_bound(a);

```

```

6         if (it != m.end() && it->second >= b)
7             return;
8         while (!m.empty() && (it = m.upper_bound(a))
9             != m.begin())
10             if ((--it)->first <= a && it->second <= b)
11                 m.erase(it); else break;
12         m[a] = b;
13     }
14     // max { b | (a, b) \in m, a >= u }, or -LLINF
15     otherwise.
16     ll max_tail(ll u) {
17         auto it = m.lower_bound(u);
18         return (it != m.end() ? it->second : -LLINF);
19     }
20 };

```

## 4 Other Mathematics

### 4.1 Helpful functions

**4.1.1 Euler's Totient Function**  $n = p_1^{k_1-1} \cdot (p_1 - 1) \cdot \dots \cdot p_r^{k_r-1} \cdot (p_r - 1)$ , where  $p_1^{k_1} \cdot \dots \cdot p_r^{k_r}$  is the prime factorization of  $n$ .

```

1 # include "header.h"
2 ll phi(ll n) { // \Phi(n)
3     ll ans = 1;
4     for (ll i = 2; i*i <= n; i++) {
5         if (n % i == 0) {
6             ans *= i-1;
7             n /= i;
8             while (n % i == 0) {
9                 ans *= i;
10                n /= i;
11            }
12        }
13    }
14    if (n > 1) ans *= n-1;
15    return ans;
16 }
17 vi phis(int n) { // All \Phi(i) up to n
18     vi phi(n+1, 0LL);
19     iota(phi.begin(), phi.end(), 0LL);
20     for (ll i = 2LL; i <= n; ++i)
21         if (phi[i] == i)
22             for (ll j = i; j <= n; j += i)
23                 phi[j] -= phi[j] / i;
24     return phi;
25 }

```



**Formulas**  $\Phi(n)$  counts all numbers in  $1, \dots, n-1$  coprime to  $n$ .

$a^{\varphi(n)} \equiv 1 \pmod n$ ,  $a$  and  $n$  are coprimes.

$\forall e > \log_2 m : n^e \pmod m = n^{\Phi(m)+e} \pmod{\Phi(m)} \pmod m$ .

$\gcd(m, n) = 1 \Rightarrow \Phi(m \cdot n) = \Phi(m) \cdot \Phi(n)$ .

**4.1.2 Pascal's trinagle**  $\binom{n}{k}$  is  $k$ -th element in the  $n$ -th row, indexing both from 0

```
1 #include "header.h"
2 void printPascal(int n) {
3     for (int line = 1; line <= n; line++) {
4         int C = 1; // used to represent C(line, i)
5         for (int i = 1; i <= line; i++) {
6             cout << C << " ";
7             C = C * (line - i) / i;
8         }
9         cout << "\n";
10    }
11 }
```

## 4.2 Theorems and definitions

**Subfactorial (Derangements)** Permutations of a set such that none of the elements appear in their original position:

$$!n = n! \sum_{i=0}^n \frac{(-1)^i}{i!}$$

$$!(0) = 1, !n = n \cdot !(n-1) + (-1)^n$$

$$!n = (n-1)(!(n-1) + !(n-2)) = \left\lfloor \frac{n!}{e} \right\rfloor \quad (1)$$

$$!n = 1 - e^{-1}, n \rightarrow \infty \quad (2)$$

### Binomials and other partitionings

$$\binom{n}{k} = \binom{n-1}{k} + \binom{n-1}{k-1} = \prod_{i=1}^k \frac{n-i+1}{i}$$

This last product may be computed incrementally since any product of  $k'$  consecutive values is divisible by  $k'!$ .

Basic identities: The hockeystick identity:

$$\sum_{k=r}^n \binom{k}{r} = \binom{n+1}{r+1}$$

or

$$\sum_{k \leq n} \binom{r+k}{k} = \binom{r+n+1}{n}$$

Also

$$\sum_{k=0}^n \binom{k}{m} = \binom{n+1}{m+1}$$

$$\sum_{i=0}^n \binom{n}{i} = 2^n$$

For  $n, m \geq 0$  and  $p$  prime: write  $n, m$  in base  $p$ , i.e.  $n = n_k p^k + \dots + n_1 p + n_0$  and  $m = m_k p^k + \dots + m_1 p + m_0$ . Then by Lucas theorem we have  $\binom{n}{m} \equiv \prod_{i=0}^k \binom{n_i}{m_i} \pmod p$ , with the convention that  $n_i < m_i \Rightarrow \binom{n_i}{m_i} = 0$ .

**Fibonacci** (See also number theory section)

$$\sum_{0 \leq k \leq n} \binom{n-k}{k} = F_{n+1}$$

$$F_n = \frac{1}{\sqrt{5}} \left( \frac{1+\sqrt{5}}{2} \right)^n - \frac{1}{\sqrt{5}} \left( \frac{1-\sqrt{5}}{2} \right)^n$$

$$\sum_{i=1}^n F_i = F_{n+2} - 1, \sum_{i=1}^n F_i^2 = F_n F_{n+1}$$

$$\gcd(F_m, F_n) = F_{\gcd(m, n)}$$

$$\gcd(F_n, F_{n+1}) = \gcd(F_n, F_{n+2}) = 1$$

**Bit stuff**  $a + b = a \oplus b + 2(a \& b) = a|b + a \& b$ .

$k$ th bit is set in  $x$  iff  $x \bmod 2^{k-1} \geq 2^k$ , or iff  $x \bmod 2^{k-1} - x \bmod 2^k \neq 0$  (i.e.  $= 2^k$ ) It comes handy when you need to look at the bits of the numbers which are pair sums or subset sums etc.

$$n \bmod 2^i = n \& (2^i - 1).$$

$$\forall k : 1 \oplus 2 \oplus \dots \oplus (4k-1) = 0$$

## 4.3 Geometry Formulas

$$\text{Euler:} \quad 1 + CC = V - E + F$$

$$\text{Pick:} \quad \text{Area} = \text{itr pts} + \frac{\text{bdry pts}}{2} - 1$$

Given a non-self-intersecting closed polygon on  $n$  vertices, given as  $(x_i, y_i)$ , its centroid  $(C_x, C_y)$  is given as:

$$C_x = \frac{1}{6A} \sum_{i=0}^{n-1} (x_i + x_{i+1})(x_i y_{i+1} - x_{i+1} y_i),$$

$$C_y = \frac{1}{6A} \sum_{i=0}^{n-1} (y_i + y_{i+1})(x_i y_{i+1} - x_{i+1} y_i)$$

$$A = \frac{1}{2} \sum_{i=0}^{n-1} (x_i y_{i+1} - x_{i+1} y_i) = \text{polygon area}$$

**Inclusion-Exclusion** For appropriate  $f$  compute  $\sum_{S \subseteq T} (-1)^{|T \setminus S|} f(S)$ , or if only the size of  $S$  matters,  $\sum_{s=0}^n (-1)^{n-s} \binom{n}{s} f(s)$ . In some contexts we might use Stirling numbers, not binomial coefficients!

Some useful applications:

**Graph coloring** Let  $I(S)$  count the number of independent sets contained in  $S \subseteq V$  ( $I(\emptyset) = 1$ ,  $I(S) = I(S \setminus v) + I(S \setminus N(v))$ ). Let  $c_k = \sum_{S \subseteq V} (-1)^{|V \setminus S|} I(S)$ . Then  $V$  is  $k$ -colorable iff  $v > 0$ . Thus we can compute the chromatic number of a graph in  $O^*(2^n)$  time.

**Burnside's lemma** Given a group  $G$  acting on a set  $X$ , the number of elements in  $X$  up to symmetry is

$$\frac{1}{|G|} \sum_{g \in G} |X^g|$$

with  $X^g$  the elements of  $X$  invariant under  $g$ . For example, if  $f(n)$  counts "configurations" of some sort of length  $n$ , and we want to count them up to rotational symmetry using  $G = \mathbb{Z}/n\mathbb{Z}$ , then



$$g(n) = \frac{1}{n} \sum_{k=0}^{n-1} f(\gcd(n, k)) = \frac{1}{n} \sum_{k|n} f(k) \phi(n/k)$$

I.e. for coloring with  $c$  colors we have  $f(k) = k^c$ .

Relatedly, in Pólya's enumeration theorem we imagine  $X$  as a set of  $n$  beads with  $G$  permuting the beads (e.g. a necklace, with  $G$  all rotations and reflections of the  $n$ -cycle, i.e. the dihedral group  $D_n$ ). Suppose further that we had  $Y$  colors, then the number of  $G$ -invariant colorings  $Y^X/G$  is counted by

$$\frac{1}{|G|} \sum_{g \in G} |Y|^{c(g)}$$

with  $c(g)$  counting the number of cycles of  $g$  when viewed as a permutation of  $X$ . We can generalize this to a weighted version: if the color  $i$  can occur exactly  $r_i$  times, then this is counted by the coefficient of  $t_1^{r_1} \dots t_n^{r_n}$  in the polynomial

$$Z(t_1, \dots, t_n) = \frac{1}{|G|} \sum_{g \in G} \prod_{m \geq 1} (t_1^m + \dots + t_n^m)^{c_m(g)}$$

where  $c_m(g)$  counts the number of length  $m$  cycles in  $g$  acting as a permutation on  $X$ . Note we get the original formula by setting all  $t_i = 1$ . Here  $Z$  is the cycle index. Note: you can cleverly deal with even/odd sizes by setting some  $t_i$  to  $-1$ .

#### 4.4 Recurrences

If  $a_n = c_1 a_{n-1} + \dots + c_k a_{n-k}$ , and  $r_1, \dots, r_k$  are distinct roots of  $x^k - c_1 x^{k-1} - \dots - c_k$ , there are  $d_1, \dots, d_k$  s.t.

$$a_n = d_1 r_1^n + \dots + d_k r_k^n.$$

Non-distinct roots  $r$  become polynomial factors, e.g.  $a_n = (d_1 n + d_2) r^n$ .

#### 4.5 Sums

$$1^3 + 2^3 + 3^3 + \dots + n^3 = \frac{n^2(n+1)^2}{4}$$

$$1^4 + 2^4 + 3^4 + \dots + n^4 = \frac{n(n+1)(2n+1)(3n^2+3n-1)}{30}$$

#### 4.6 Series

$$e^x = 1 + x + \frac{x^2}{2!} + \frac{x^3}{3!} + \dots, (-\infty < x < \infty)$$

$$\ln(1+x) = x - \frac{x^2}{2} + \frac{x^3}{3} - \frac{x^4}{4} + \dots, (-1 < x \leq 1)$$

$$\sqrt{1+x} = 1 + \frac{x}{2} - \frac{x^2}{8} + \frac{2x^3}{32} - \frac{5x^4}{128} + \dots, (-1 \leq x \leq 1)$$

$$\sin x = x - \frac{x^3}{3!} + \frac{x^5}{5!} - \frac{x^7}{7!} + \dots, (-\infty < x < \infty)$$

$$\cos x = 1 - \frac{x^2}{2!} + \frac{x^4}{4!} - \frac{x^6}{6!} + \dots, (-\infty < x < \infty)$$

#### 4.7 Quadrilaterals

With side lengths  $a, b, c, d$ , diagonals  $e, f$ , diagonals angle  $\theta$ , area  $A$  and magic flux  $F = b^2 + d^2 - a^2 - c^2$ :

$$4A = 2ef \cdot \sin \theta = F \tan \theta = \sqrt{4e^2 f^2 - F^2}$$

For cyclic quadrilaterals the sum of opposite angles is  $180^\circ$ ,  $ef = ac + bd$ , and  $A = \sqrt{(p-a)(p-b)(p-c)(p-d)}$ .

#### 4.8 Triangles

Side lengths:  $a, b, c$

Semiperimeter:  $p = \frac{a+b+c}{2}$

Area:

$$[ABC] = rp = \frac{1}{2} ab \sin \gamma$$

$$= \frac{abc}{4R} = \sqrt{p(p-a)(p-b)(p-c)} = \frac{1}{2} |(B-A, C-A)^T|$$

Circumradius:  $R = \frac{abc}{4A}$ , Inradius:  $r = \frac{A}{p}$

Length of median (divides triangle into two equal-area triangles):  $m_a = \frac{1}{2} \sqrt{2b^2 + 2c^2 - a^2}$

Length of bisector (divides angles in two):  $s_a =$

$$\sqrt{bc \left[ 1 - \left( \frac{a}{b+c} \right)^2 \right]}$$

Law of tangents:  $\frac{a+b}{a-b} = \frac{\tan \frac{\alpha+\beta}{2}}{\tan \frac{\alpha-\beta}{2}}$

#### 4.9 Trigonometry

$$\tan(v+w) = \frac{\tan v + \tan w}{1 - \tan v \tan w}$$

$$\sin v + \sin w = 2 \sin \frac{v+w}{2} \cos \frac{v-w}{2}$$

$$\cos v + \cos w = 2 \cos \frac{v+w}{2} \cos \frac{v-w}{2}$$

$$(V+W) \tan(v-w)/2 = (V-W) \tan(v+w)/2$$

where  $V, W$  are lengths of sides opposite angles  $v, w$ .

$$a \cos x + b \sin x = r \cos(x - \phi)$$

$$a \sin x + b \cos x = r \sin(x + \phi)$$

where  $r = \sqrt{a^2 + b^2}$ ,  $\phi = \text{atan2}(b, a)$ .

#### 4.10 Combinatorics

Combinations and Permutations

$$P(n, r) = \frac{n!}{(n-r)!}$$

$$C(n, r) = \binom{n}{r} = \frac{n!}{r!(n-r)!}$$

$$C(n, r) = C(n, n-r)$$

#### 4.11 Cycles

Let  $g_S(n)$  be the number of  $n$ -permutations whose cycle lengths all belong to the set  $S$ . Then

$$\sum_{n=0}^{\infty} g_S(n) \frac{x^n}{n!} = \exp \left( \sum_{n \in S} \frac{x^n}{n} \right)$$

#### 4.12 Labeled unrooted trees

# on  $n$  vertices:  $n^{n-2}$

# on  $k$  existing trees of size  $n_i$ :  $n_1 n_2 \dots n_k n^{k-2}$

# with degrees  $d_i$ :  $(n-2)! / ((d_1-1)! \dots (d_n-1)!)$

### 4.13 Partition function

Number of ways of writing  $n$  as a sum of positive integers, disregarding the order of the summands.

$$p(0) = 1, \quad p(n) = \sum_{k \in \mathbb{Z} \setminus \{0\}} (-1)^{k+1} p(n - k(3k-1)/2)$$

$$p(n) \sim 0.145/n \cdot \exp(2.56\sqrt{n})$$

$n$	0	1	2	3	4	5	6	7	8	9	20	50	100
$p(n)$	1	1	2	3	5	7	11	15	22	30	627	$\sim 2e5$	$\sim 2e8$

### 4.14 Numbers

**Bernoulli numbers** EGF of Bernoulli numbers is  $B(t) = \frac{t}{e^t - 1}$  (FFT-able).  $B[0, \dots] = [1, -\frac{1}{2}, \frac{1}{6}, 0, -\frac{1}{30}, 0, \frac{1}{42}, \dots]$   
Sums of powers:

$$\sum_{i=1}^n i^m = \frac{1}{m+1} \sum_{k=0}^m \binom{m+1}{k} B_k \cdot (n+1)^{m+1-k}$$

Euler-Maclaurin formula for infinite sums:

$$\begin{aligned} \sum_{i=m}^{\infty} f(i) &= \int_m^{\infty} f(x) dx - \sum_{k=1}^{\infty} \frac{B_k}{k!} f^{(k-1)}(m) \\ &\approx \int_m^{\infty} f(x) dx + \frac{f(m)}{2} - \frac{f'(m)}{12} + \frac{f'''(m)}{720} + O(f^{(5)}(m)) \end{aligned}$$

**Stirling's numbers First kind:**  $S_1(n, k)$  count permutations on  $n$  items with  $k$  cycles.  $S_1(n, k) = S_1(n-1, k-1) + (n-1)S_1(n-1, k)$  with  $S_1(0, 0) = 1$ . Note:

$$\sum_{k=0}^n S_1(n, k) x^k = x(x+1) \dots (x+n-1)$$

$$\sum_{k=0}^n S_1(n, k) = n!$$

$S_1(8, k) = 8, 0, 5040, 13068, 13132, 6769, 1960, 322, 28, 1$   
 $S_1(n, 2) = 0, 0, 1, 3, 11, 50, 274, 1764, 13068, 109584, \dots$

**Second kind:**  $S_2(n, k)$  count partitions of  $n$  distinct elements into exactly  $k$  non-empty groups.

$$S_2(n, k) = S_2(n-1, k-1) + kS_2(n-1, k)$$

$$S_2(n, 1) = S_2(n, n) = 1$$

$$S_2(n, k) = \frac{1}{k!} \sum_{i=0}^k (-1)^{k-i} \binom{k}{i} i^n$$

**Narayana numbers** The number of expressions containing  $n$  pairs of parentheses, which are correctly matched and which contain  $k$  distinct nestings.

$$N(n, k) = \frac{1}{n} \frac{n}{k} \frac{n}{k-1}$$

**Eulerian numbers** Number of permutations  $\pi \in S_n$  in which exactly  $k$  elements are greater than the previous element.  $k$   $j$ :s s.t.  $\pi(j) > \pi(j+1)$ ,  $k+1$   $j$ :s s.t.  $\pi(j) \geq j$ ,  $k$   $j$ :s s.t.  $\pi(j) > j$ .

$$E(n, k) = (n-k)E(n-1, k-1) + (k+1)E(n-1, k)$$

$$E(n, 0) = E(n, n-1) = 1$$

$$E(n, k) = \sum_{j=0}^k (-1)^j \binom{n+1}{j} (k+1-j)^n$$

**Bell numbers** Total number of partitions of  $n$  distinct elements.  $B(n) = 1, 1, 2, 5, 15, 52, 203, 877, 4140, 21147, \dots$   
For  $p$  prime,

$$B(p^m + n) \equiv mB(n) + B(n+1) \pmod{p}$$

**Catalan numbers** Number of correct bracket sequence consisting of  $n$  opening and  $n$  closing brackets.

The number of ways to completely parenthesize  $n+1$  factors.

The number of triangulations of a convex polygon with  $n+2$  sides (i.e. the number of partitions of polygon into disjoint triangles by using the diagonals).

The number of ways to connect the  $2n$  points on a circle to form  $n$  disjoint i.e. non-intersecting chords.

$$C_n = \frac{1}{n+1} \binom{2n}{n} = \binom{2n}{n} - \binom{2n}{n+1} = \frac{(2n)!}{(n+1)!n!}$$

$$C_0 = 1, \quad C_{n+1} = \frac{2(2n+1)}{n+2} C_n, \quad C_{n+1} = \sum C_i C_{n-i}$$

$$C_n = 1, 1, 2, 5, 14, 42, 132, 429, 1430, 4862, 16796, 58786, \dots$$

- sub-diagonal monotone paths in an  $n \times n$  grid.
- strings with  $n$  pairs of parenthesis, correctly nested.
- binary trees with  $n+1$  leaves (0 or 2 children).
- ordered trees with  $n+1$  vertices.
- ways a convex polygon with  $n+2$  sides can be cut into triangles by connecting vertices with straight lines.
- permutations of  $[n]$  with no 3-term increasing subseq.

### 4.15 Probability

Stochastic variables

$$P(X=r) = C(n, r) \cdot p^r \cdot (1-p)^{n-r}$$

**Bayes' Theorem**  $P(B|A) = \frac{P(A|B)P(B)}{P(A)}$

$$P(B|A) = \frac{P(A|B)P(B)}{P(A|B)P(B) + P(A|\bar{B})P(\bar{B})}$$

$$P(B_k|A) = \frac{P(A|B_k)P(B_k)}{P(A|B_1)P(B_1) + \dots + P(A|B_n)P(B_n)}$$

**Expectation** Let  $X$  be a discrete random variable with probability  $p_X(x)$  of assuming the value  $x$ . It will then have an expected value (mean)  $\mu = \mathbb{E}(X) = \sum_x x p_X(x)$  and variance  $\sigma^2 = V(X) = \mathbb{E}(X^2) - (\mathbb{E}(X))^2 = \sum_x (x - \mathbb{E}(X))^2 p_X(x)$  where  $\sigma$  is the standard deviation. If  $X$  is instead continuous it will have a probability density function  $f_X(x)$  and the sums above will instead be integrals with  $p_X(x)$  replaced by  $f_X(x)$ .

Expectation is linear:

$$\mathbb{E}(aX + bY) = a\mathbb{E}(X) + b\mathbb{E}(Y)$$

For independent  $X$  and  $Y$ ,

$$V(aX + bY) = a^2 V(X) + b^2 V(Y).$$

### 4.16 Number Theory

**Bezout's Theorem**

$$a, b \in \mathbb{Z}^+ \implies \exists s, t \in \mathbb{Z} : \gcd(a, b) = sa + tb$$

**Bézout's identity** For  $a \neq 0$ ,  $b \neq 0$ , then  $d = \gcd(a, b)$  is the smallest positive integer for which there are integer solutions to

$$ax + by = d$$

If  $(x, y)$  is one solution, then all solutions are given by

$$\left(x + \frac{kb}{\gcd(a, b)}, y - \frac{ka}{\gcd(a, b)}\right), \quad k \in \mathbb{Z}$$

### Partial Coprime Divisor Property

$$(\gcd(a, b) = 1) \wedge (a \mid bc) \implies (a \mid c)$$

### Coprime Modulus Equivalence Property

$$(\gcd(c, m) = 1) \wedge (ac \equiv bc \pmod{m}) \implies (a \equiv b \pmod{m})$$

### Fermat's Little Theorem

$$(\text{prime}(p)) \wedge (p \nmid a) \implies (a^{p-1} \equiv 1 \pmod{p})$$

$$(\text{prime}(p)) \implies (a^p \equiv a \pmod{p})$$

### Euler's Theorem

$$a^{\phi(m)-1} \equiv a^{-1} \pmod{m}, \text{ if } \gcd(a, m) = 1$$

$$a^{-1} \equiv a^{m-2} \pmod{m}, \text{ if } m \text{ is prime}$$

**Pythagorean Triples** The Pythagorean triples are uniquely generated by

$$a = k \cdot (m^2 - n^2), \quad b = k \cdot (2mn), \quad c = k \cdot (m^2 + n^2),$$

with  $m > n > 0$ ,  $k > 0$ ,  $m \perp n$ , and either  $m$  or  $n$  even.

**Primes**  $p = 962592769$  is such that  $2^{21} \mid p - 1$ , which may be useful. For hashing use 970592641 (31-bit number), 31443539979727 (45-bit), 3006703054056749 (52-bit). There are 78498 primes less than 1 000 000.

Primitive roots exist modulo any prime power  $p^a$ , except for  $p = 2, a > 2$ , and there are  $\phi(\phi(p^a))$  many. For  $p = 2, a > 2$ , the group  $\mathbb{Z}_{2^a}^\times$  is instead isomorphic to  $\mathbb{Z}_2 \times \mathbb{Z}_{2^{a-2}}$ .

**Estimates**  $\sum_{d \mid n} d = O(n \log \log n)$ .

The number of divisors of  $n$  is at most around 100 for  $n < 5e4$ , 500 for  $n < 1e7$ , 2000 for  $n < 1e10$ , 200 000 for  $n < 1e19$ .

### Mobius Function

$$\mu(n) = \begin{cases} 0 & n \text{ is not square free} \\ 1 & n \text{ has even number of prime factors} \\ -1 & n \text{ has odd number of prime factors} \end{cases}$$

Mobius Inversion:

$$g(n) = \sum_{d \mid n} f(d) \Leftrightarrow f(n) = \sum_{d \mid n} \mu(d) g(n/d)$$

Other useful formulas/forms:

$$\sum_{d \mid n} \mu(d) = [n = 1] \text{ (very useful)}$$

$$g(n) = \sum_{n \mid d} f(d) \Leftrightarrow f(n) = \sum_{n \mid d} \mu(d/n) g(d)$$

$$g(n) = \sum_{1 \leq m \leq n} f\left(\left\lfloor \frac{n}{m} \right\rfloor\right) \Leftrightarrow f(n) = \sum_{1 \leq m \leq n} \mu(m) g\left(\left\lfloor \frac{n}{m} \right\rfloor\right)$$

## 4.17 Discrete distributions

**Binomial distribution** The number of successes in  $n$  independent yes/no experiments, each which yields success with probability  $p$  is  $\text{Bin}(n, p)$ ,  $n = 1, 2, \dots$ ,  $0 \leq p \leq 1$ .

$$p(k) = \binom{n}{k} p^k (1-p)^{n-k}$$

$$\mu = np, \sigma^2 = np(1-p)$$

$\text{Bin}(n, p)$  is approximately  $\text{Po}(np)$  for small  $p$ .

**First success distribution** The number of trials needed to get the first success in independent yes/no experiments, each which yields success with probability  $p$  is  $\text{Fs}(p)$ ,  $0 \leq p \leq 1$ .

$$p(k) = p(1-p)^{k-1}, \quad k = 1, 2, \dots$$

$$\mu = \frac{1}{p}, \sigma^2 = \frac{1-p}{p^2}$$

**Poisson distribution** The number of events occurring in a fixed period of time  $t$  if these events occur with a known average rate  $\kappa$  and independently of the time since the last event is  $\text{Po}(\lambda)$ ,  $\lambda = t\kappa$ .

$$p(k) = e^{-\lambda} \frac{\lambda^k}{k!}, \quad k = 0, 1, 2, \dots$$

$$\mu = \lambda, \sigma^2 = \lambda$$

## 4.18 Continuous distributions

**Uniform distribution** If the probability density function is constant between  $a$  and  $b$  and 0 elsewhere it is  $\text{U}(a, b)$ ,  $a < b$ .

$$f(x) = \begin{cases} \frac{1}{b-a} & a < x < b \\ 0 & \text{otherwise} \end{cases}$$

$$\mu = \frac{a+b}{2}, \sigma^2 = \frac{(b-a)^2}{12}$$

**Exponential distribution** The time between events in a Poisson process is  $\text{Exp}(\lambda)$ ,  $\lambda > 0$ .

$$f(x) = \begin{cases} \lambda e^{-\lambda x} & x \geq 0 \\ 0 & x < 0 \end{cases}$$

$$\mu = \frac{1}{\lambda}, \sigma^2 = \frac{1}{\lambda^2}$$

**Normal distribution** Most real random values with mean  $\mu$  and variance  $\sigma^2$  are well described by  $\mathcal{N}(\mu, \sigma^2)$ ,  $\sigma > 0$ .

$$f(x) = \frac{1}{\sqrt{2\pi\sigma^2}} e^{-\frac{(x-\mu)^2}{2\sigma^2}}$$

If  $X_1 \sim \mathcal{N}(\mu_1, \sigma_1^2)$  and  $X_2 \sim \mathcal{N}(\mu_2, \sigma_2^2)$  then

$$aX_1 + bX_2 + c \sim \mathcal{N}(\mu_1 + \mu_2 + c, a^2\sigma_1^2 + b^2\sigma_2^2)$$