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# 1 Setup

## 1.1 header.h

```
3 using namespace std;
5 #define 11 long long
6 #define ull unsigned ll
7 #define ld long double
8 #define pl pair<ll, ll>
9 #define pi pair <int, int>
10 #define vl vector<ll>
11 #define vi vector<int>
12 #define vvi vector <vi>
13 #define vvl vector <vl>
14 #define vpl vector <pl>
15 #define vpi vector <pi>
16 #define vld vector <ld>
17 #define in_fast(el, cont) (cont.find(el) != cont.end())
18 #define in(el, cont) (find(cont.begin(), cont.end(), el) != cont.end())
20 constexpr int INF = 200000010;
21 constexpr 11 LLINF = 900000000000000010LL;
23 template <typename T, template <typename ELEM, typename ALLOC = std::
      allocator < ELEM > > class Container >
_{24} std::ostream& operator<<(std::ostream& o, const Container<T>& container) {
   typename Container <T>::const_iterator beg = container.begin();
   if (beg != container.end()) {
    o << *beg++;
      while (beg != container.end()) {
        o << " " << *beg++;
29
30
    }
    return o;
33 }
35 // int main() {
36 // ios::sync_with_stdio(false); // do not use cout + printf
37 // cin.tie(NULL);
38 // cout << fixed << setprecision(12);
39 // return 0;
40 // }
```

# 1.2 Bash for c++ compile with header.h

```
#!/bin/bash
if [ $# -ne 1 ]; then echo "Usage: $0 <input_file>"; exit 1; fi
if [ $f="$1";d=code/;o=a.out
if [ -f $d/$f ] || { echo "Input file not found: $f"; exit 1; }
if g++ -I$d $d/$f -o $0 && echo "Compilation successful. Executable '$o' created." || echo "Compilation failed."
```

# 1.3 Bash for run tests c++

<sup>1 #</sup>pragma once // Delete this when copying this file 2 #include <bits/stdc++.h>

```
_2 for file in 1/*.in; do diff ((1/1).out < "file") "file%.in%.ans"; done
```

# 1.4 Bash for run tests python

```
_1 for file in $1/*.in; do diff <(python3 $1/$1.py < "$file") "${file%.in}.ans "; done
```

#### 1.4.1 Auxiliary helper C++

```
1 #include "header.h"
3 int main() {
      // Read in a line including white space
      string line;
      getline(cin, line);
      // When doing the above read numbers as follows:
      getline(cin, line);
      stringstream ss(line);
      ss >> n;
11
12
      // Count the number of 1s in binary represnatation of a number
14
      ull number;
      __builtin_popcountll(number);
15
16 }
```

#### 1.4.2 Auxiliary helper python

```
# Read until EOF

while True:

try:

pattern = input()

except EOFError:

break
```

# 2 Python

# 2.1 Graphs

#### 2.1.1 BFS

```
from collections import deque
def bfs(g, roots, n):
    q = deque(roots)
    explored = set(roots)
    distances = [float("inf")]*n
    distances[0][0] = 0

while len(q) != 0:
    node = q.popleft()
```

## 2.1.2 Dijkstra

```
1 from heapq import *
2 def dijkstra(n, root, g): # g = {node: (cost, neigh)}
    dist = [float("inf")]*n
    dist[root] = 0
    prev = [-1]*n
    pq = [(0, root)]
    heapify(pq)
    visited = set([])
    while len(pq) != 0:
      _, node = heappop(pq)
13
      if node in visited: continue
14
      visited.add(node)
16
      # In case of disconnected graphs
17
      if node not in g:
18
        continue
19
20
      for cost, neigh in g[node]:
21
        alt = dist[node] + cost
22
        if alt < dist[neigh]:</pre>
23
          dist[neigh] = alt
24
          prev[neigh] = node
          heappush(pq, (alt, neigh))
26
    return dist
```

# 2.2 Dynamic Programming

- 2.3 Trees
- 2.4 Number Theory / Combinatorics

## 2.4.1 nCk % prime

```
# Note: p must be prime and k  n:
    return 0

# calculate numerator
num = 1
for i in range(n-k+1, n+1):
    num *= i % p
num %= p
# calculate denominator
```

# **2.4.2 Sieve of Eratosthenes** O(n) so actually faster than C++ version, but more memory

```
_{1} MAX SIZE = 10**8+1
2 isprime = [True] * MAX_SIZE
3 prime = []
4 SPF = [None] * (MAX_SIZE)
6 def manipulated_seive(N): # Up to N (not included)
    isprime[0] = isprime[1] = False
    for i in range(2, N):
      if isprime[i] == True:
        prime.append(i)
        SPF[i] = i
      while (j < len(prime) and
14
        i * prime[j] < N and
          prime[j] <= SPF[i]):</pre>
15
        isprime[i * prime[j]] = False
        SPF[i * prime[j]] = prime[j]
        i += 1
```

# 2.5 Strings

#### 2.5.1 LCS

```
1 def longestCommonSubsequence(text1, text2): # 0(m*n) time, 0(m) space
      n = len(text1)
      m = len(text2)
      # Initializing two lists of size m
      prev = [0] * (m + 1)
      cur = \lceil 0 \rceil * (m + 1)
      for idx1 in range(1, n + 1):
          for idx2 in range(1, m + 1):
              # If characters are matching
11
              if text1[idx1 - 1] == text2[idx2 - 1]:
12
                   cur[idx2] = 1 + prev[idx2 - 1]
               else:
                   # If characters are not matching
                   cur[idx2] = max(cur[idx2 - 1], prev[idx2])
          prev = cur.copy()
18
19
      return cur[m]
```

#### 2.5.2 KMP

```
1 class KMP:
      def partial(self, pattern):
          """ Calculate partial match table: String -> [Int] """
          for i in range(1, len(pattern)):
               i = ret[i - 1]
               while j > 0 and pattern[j] != pattern[i]: j = ret[j - 1]
               ret.append(j + 1 if pattern[j] == pattern[i] else j)
          return ret
      def search(self, T, P):
11
          """KMP search main algorithm: String -> String -> [Int]
12
          Return all the matching position of pattern string P in T"""
13
          partial, ret, j = self.partial(P), [], 0
14
          for i in range(len(T)):
15
               while j > 0 and T[i] != P[j]: j = partial[j - 1]
16
               if T[i] == P[j]: j += 1
17
18
              if i == len(P):
                   ret.append(i - (j - 1))
19
                   j = partial[j - 1]
20
          return ret
```

# 2.6 Geometry

# 2.7 Other Algorithms

#### 2.7.1 Rotate matrix

## 2.8 Other Data Structures

#### 2.8.1 Segment Tree

```
_{1} N = 100000 # limit for array size
2 tree = [0] * (2 * N) # Max size of tree
4 def build(arr, n): # function to build the tree
      # insert leaf nodes in tree
      for i in range(n):
          tree[n + i] = arr[i]
      # build the tree by calculating parents
10
      for i in range(n - 1, 0, -1):
          tree[i] = tree[i << 1] + tree[i << 1 | 1]</pre>
11
13 def updateTreeNode(p, value, n): # function to update a tree node
      # set value at position p
14
      tree[p + n] = value
15
16
      p = p + n
17
      i = p # move upward and update parents
```

```
while i > 1:
          tree[i >> 1] = tree[i] + tree[i ^ 1]
          i >>= 1
23 def query(1, r, n): # function to get sum on interval [1, r)
24
      # loop to find the sum in the range
25
      1 += n
      r += n
27
      while 1 < r:
28
          if 1 & 1:
29
              res += tree[1]
              1 += 1
31
          if r & 1:
32
              r -= 1
              res += tree[r]
          1 >>= 1
          r >>= 1
36
      return res
```

# 3 C++

# 3.1 Graphs

## 3.1.1 BFS

```
1 #include "header.h"
2 #define graph unordered_map<11, unordered_set<11>>
3 vi bfs(int n, graph& g, vi& roots) {
      vi parents(n+1, -1); // nodes are 1..n
      unordered_set <int> visited;
      queue < int > q;
      for (auto x: roots) {
          q.emplace(x);
          visited.insert(x);
10
      while (not q.empty()) {
11
          int node = q.front();
12
          q.pop();
13
14
          for (auto neigh: g[node]) {
               if (not in(neigh, visited)) {
                   parents[neigh] = node;
17
                   q.emplace(neigh);
18
                   visited.insert(neigh);
              }
          }
21
22
23
      return parents;
25 vi reconstruct_path(vi parents, int start, int goal) {
      vi path;
26
27
      int curr = goal;
      while (curr != start) {
          path.push_back(curr);
          if (parents[curr] == -1) return vi(); // No path, empty vi
```

# **3.1.2 DFS** Cycle detection / removal

```
1 #include "header.h"
2 void removeCyc(ll node, unordered_map<ll, vector<pair<ll, 11>>>& neighs,
      vector < bool > & visited.
3 vector < bool > & recStack, vector < 11 > & ans) {
      if (!visited[node]) {
          visited[node] = true;
          recStack[node] = true;
          auto it = neighs.find(node);
          if (it != neighs.end()) {
               for (auto util: it->second) {
                   11 nnode = util.first:
                   if (recStack[nnode]) {
                       ans.push_back(util.second);
12
                   } else if (!visited[nnode]) {
13
                       removeCyc(nnode, neighs, visited, recStack, ans);
14
                   }
               }
16
          }
17
18
      recStack[node] = false;
19
20 }
```

## 3.1.3 Dijkstra

```
1 #include "header.h"
2 vector<int> dijkstra(int n, int root, map<int, vector<pair<int, int>>>& g) {
    unordered_set <int> visited;
    vector < int > dist(n. INF):
      priority_queue < pair < int , int >> pq;
      dist[root] = 0;
      pq.push({0, root});
      while (!pq.empty()) {
           int node = pq.top().second;
           int d = -pq.top().first;
10
           pq.pop();
11
12
           if (in(node, visited)) continue;
13
           visited.insert(node);
14
15
           for (auto e : g[node]) {
16
               int neigh = e.first;
17
               int cost = e.second;
18
               if (dist[neigh] > dist[node] + cost) {
19
                   dist[neigh] = dist[node] + cost;
20
                   pq.push({-dist[neigh], neigh});
21
           }
```

```
24 }
25 return dist;
26 }
```

## 3.1.4 Floyd-Warshall

#### 3.1.5 Kruskal Minimum spanning tree of undirected weighted graph

```
1 #include "header.h"
2 #include "disjoint_set.h"
3 // O(E log E)
4 pair < set < pair < 11, 11 >> , 11 > kruskal (vector < tuple < 11, 11, 11 >> & edges, 11 n)
      set <pair <11. 11>> ans:
      11 cost = 0:
      sort(edges.begin(), edges.end());
      DisjointSet <11> fs(n);
10
      ll dist, i, j;
11
      for (auto edge: edges) {
12
           dist = get<0>(edge);
           i = get<1>(edge);
14
           j = get < 2 > (edge);
15
16
           if (fs.find_set(i) != fs.find_set(j)) {
17
               fs.union_sets(i, j);
               ans.insert({i, j});
20
               cost += dist;
          }
22
      return pair<set<pair<11, 11>>, 11> {ans, cost};
23
24 }
```

## 3.1.6 Hungarian algorithm

```
* prefix of jobs to distinct workers.
   * Otparam T a type large enough to represent integers on the order of J \ast
   * Oparam C a matrix of dimensions JxW such that C[j][w] = cost to assign j-
     job to w-th worker (possibly negative)
12 * Creturn a vector of length J. with the i-th entry equaling the minimum
  * to assign the first (j+1) jobs to distinct workers
15 template <class T> vector<T> hungarian(const vector<vector<T>> &C) {
      const int J = (int)size(C), W = (int)size(C[0]);
      assert(J <= W);
17
      // job[w] = job assigned to w-th worker, or -1 if no job assigned
      // note: a W-th worker was added for convenience
      vector < int > job(W + 1, -1);
20
      vector<T> ys(J), yt(W + 1); // potentials
21
22
      // -vt[W] will equal the sum of all deltas
      vector <T> answers;
23
24
      const T inf = numeric_limits<T>::max();
      for (int j_cur = 0; j_cur < J; ++j_cur) { // assign j_cur-th job</pre>
25
          int w_cur = W;
26
          iob[w cur] = i cur:
27
          // min reduced cost over edges from Z to worker w
28
          vector<T> min_to(W + 1, inf);
          vector<int> prv(W + 1, -1); // previous worker on alternating path
30
          vector < bool > in Z (W + 1);  // whether worker is in Z
31
          while (job[w_cur] != -1) { // runs at most j_cur + 1 times
               in_Z[w_cur] = true;
33
               const int j = job[w_cur];
34
              T delta = inf:
35
               int w_next;
36
               for (int w = 0; w < W; ++w) {
37
                   if (!in Z[w]) {
38
                       if (ckmin(min_to[w], C[j][w] - ys[j] - yt[w]))
39
                           prv[w] = w cur:
40
                       if (ckmin(delta, min_to[w])) w_next = w;
41
                  }
              }
43
               // delta will always be non-negative,
44
               // except possibly during the first time this loop runs
               // if any entries of C[j_cur] are negative
46
               for (int w = 0; w \le W; ++w) {
                   if (in_Z[w]) ys[job[w]] += delta, yt[w] -= delta;
48
                   else min_to[w] -= delta;
49
50
               w_cur = w_next;
51
52
          // update assignments along alternating path
          for (int w; w_cur != W; w_cur = w) job[w_cur] = job[w = prv[w_cur]];
54
          answers.push_back(-yt[W]);
55
      }
56
57
      return answers;
58 }
```

3.1.7 Successive shortest path Calculates max flow, min cost

```
1 #include "header.h"
2 // map<node, map<node, pair<cost, capacity>>>
3 #define graph unordered_map<int, unordered_map<int, pair<ld, int>>>
5 const ld infty = 1e60l; // Change if necessary
6 ld fill(int n, vld& potential) { // Finds max flow, min cost
    priority_queue < pair < ld, int >> pq;
    vector < bool > visited(n+2, false);
    vi parent(n+2, 0);
    vld dist(n+2, infty);
    dist[0] = 0.1;
11
    pq.emplace(make_pair(0.1, 0));
    while (not pq.empty()) {
      int node = pq.top().second;
      pq.pop();
15
      if (visited[node]) continue;
16
      visited[node] = true;
17
      for (auto& x : g[node]) {
        int neigh = x.first;
19
        int capacity = x.second.second;
20
        ld cost = x.second.first:
21
        if (capacity and not visited[neigh]) {
22
          ld d = dist[node] + cost + potential[node] - potential[neigh];
23
          if (d + 1e-101 < dist[neigh]) {</pre>
24
             dist[neigh] = d;
25
            pq.emplace(make_pair(-d, neigh));
26
            parent[neigh] = node;
    }}}
28
29
    for (int i = 0; i < n+2; i++) {
30
      potential[i] = min(infty, potential[i] + dist[i]);
31
    if (not parent[n+1]) return infty;
33
    ld ans = 0.1;
    for (int x = n+1; x; x=parent[x]) {
      ans += g[parent[x]][x].first;
      g[parent[x]][x].second--;
      g[x][parent[x]].second++;
    }
    return ans;
40
41 }
```

#### 3.1.8 Bipartite check

```
int v = q.front();
14
15
                   a.pop():
                   for (int u : adj[v]) {
16
                        if (side[u] == -1) {
17
                            side[u] = side[v] ^ 1:
18
                            q.push(u);
19
                        } else {
20
                            is_bipartite &= side[u] != side[v];
21
                        }
22
23 }}}}
```

## 3.1.9 Find cycle directed

```
1 #include "header.h"
2 int n:
3 \text{ const int } mxN = 2e5+5;
4 vvi adj(mxN);
5 vector < char > color:
6 vi parent;
7 int cycle_start, cycle_end;
8 bool dfs(int v) {
       color[v] = 1;
       for (int u : adj[v]) {
           if (color[u] == 0) {
11
                parent[u] = v:
12
13
                if (dfs(u)) return true;
           } else if (color[u] == 1) {
14
                cycle_end = v;
1.5
                cvcle_start = u;
16
                return true;
17
           }
18
19
       color[v] = 2:
20
       return false:
21
22 }
23 void find_cycle() {
       color.assign(n, 0);
24
25
       parent.assign(n, -1);
       cycle_start = -1;
       for (int v = 0; v < n; v++) {
27
           if (color[v] == 0 && dfs(v))break:
28
      }
29
       if (cycle_start == -1) {
30
           cout << "Acyclic" << endl;</pre>
31
      } else {
32
           vector<int> cycle;
33
34
           cycle.push_back(cycle_start);
           for (int v = cycle_end; v != cycle_start; v = parent[v])
35
                cycle.push_back(v);
36
           cycle.push_back(cycle_start);
37
           reverse(cycle.begin(), cycle.end());
38
           cout << "Cycle_Found:_";</pre>
40
           for (int v : cycle) cout << v << "";
41
           cout << endl;</pre>
42
43
44 }
```

#### 3.1.10 Find cycle directed

```
1 #include "header.h"
2 int n:
3 \text{ const int } mxN = 2e5 + 5;
4 vvi adj(mxN);
5 vector < bool > visited;
6 vi parent;
7 int cycle_start, cycle_end;
8 bool dfs(int v, int par) { // passing vertex and its parent vertex
       visited[v] = true;
       for (int u : adj[v]) {
           if(u == par) continue; // skipping edge to parent vertex
           if (visited[u]) {
12
               cycle_end = v;
13
               cvcle_start = u;
14
               return true;
16
           parent[u] = v;
17
           if (dfs(u, parent[u]))
18
               return true;
19
20
       return false;
21
22 }
23 void find_cycle() {
       visited.assign(n, false);
^{24}
       parent.assign(n, -1);
25
       cvcle start = -1:
       for (int v = 0; v < n; v++) {
27
           if (!visited[v] && dfs(v, parent[v])) break;
28
29
       if (cycle_start == -1) {
30
           cout << "Acvclic" << endl:</pre>
31
       } else {
32
           vector < int > cycle;
33
           cycle.push_back(cycle_start);
34
           for (int v = cycle_end; v != cycle_start; v = parent[v])
35
                cycle.push_back(v);
           cycle.push_back(cycle_start);
           cout << "Cycle_Found:";</pre>
38
           for (int v : cycle) cout << v << "";</pre>
39
           cout << endl;</pre>
41
42 }
```

# 3.1.11 Tarjan's SCC

```
l[u] = n[u] = counter++;
11
      st.push(u):
      vs[u] = true:
      for (auto &&v : edges[u]) {
14
        if (n[v] == -1) visit(v, com);
16
        if (vs[v]) 1[u] = min(1[u], 1[v]);
17
      if (1[u] == n[u]) {
        while (true) {
19
          int v = st.top();
20
          st.pop():
21
          vs[v] = false;
22
          com[v] = C: // <== ACT HERE
          if (u == v) break;
24
25
        C++:
      }
27
28
    int find_sccs(vi &com) { // component indices will be stored in 'com'
29
      com.assign(V, -1);
30
      C = 0:
31
      for (int u = 0; u < V; ++u)
32
        if (n[u] == -1) visit(u, com);
34
    }
35
    // scc is a map of the original vertices of the graph to the vertices
    // of the SCC graph, scc_graph is its adjacency list.
    // SCC indices and edges are stored in 'scc' and 'scc_graph'.
    void scc_collapse(vi &scc, vvi &scc_graph) {
      find_sccs(scc);
40
      scc_graph.assign(C, vi());
41
      set < pi > rec; // recorded edges
42
      for (int u = 0; u < V; ++u) {
        assert(scc[u] != -1);
44
        for (int v : edges[u]) {
45
          if (scc[v] == scc[u] ||
46
            rec.find({scc[u], scc[v]}) != rec.end()) continue;
47
          scc_graph[scc[u]].push_back(scc[v]);
48
          rec.insert({scc[u], scc[v]});
50
      }
51
52
    // Function to find sources and sinks in the SCC graph
    // The number of edges needed to be added is max(sources.size(), sinks.())
    void findSourcesAndSinks(const vvi &scc_graph, vi &sources, vi &sinks) {
      vi in_degree(C, 0), out_degree(C, 0);
      for (int u = 0; u < C; u++) {
57
        for (auto v : scc_graph[u]) {
58
          in_degree[v]++;
          out degree[u]++:
        }
      }
      for (int i = 0: i < C: ++i) {</pre>
63
        if (in_degree[i] == 0) sources.push_back(i);
        if (out_degree[i] == 0) sinks.push_back(i);
66
   }
68 };
```

# 3.1.12 SCC edges Prints out the missing edges to make the input digraph strongly connected

```
1 #include "header.h"
2 const int N=1e5+10:
3 int n,a[N],cnt[N],vis[N];
4 vector < int > hd,tl;
5 int dfs(int x){
       vis[x]=1;
       if(!vis[a[x]])return vis[x]=dfs(a[x]);
       return vis[x]=x;
9 }
10 int main(){
       scanf("%d",&n);
11
       for(int i=1:i<=n:i++){</pre>
           scanf("%d",&a[i]);
13
           cnt[a[i]]++;
14
       int k=0;
16
       for(int i=1;i<=n;i++){</pre>
17
           if(!cnt[i]){
18
                k++;
19
                hd.push_back(i);
20
                tl.push_back(dfs(i));
21
           }
22
23
       int tk=k;
24
       for(int i=1:i<=n:i++){</pre>
25
           if(!vis[i]){
26
               k++;
27
               hd.push_back(i);
                tl.push_back(dfs(i));
29
           }
30
31
       if (k==1&&!tk)k=0;
32
       printf("%d\n",k);
33
       for(int i=0;i<k;i++)printf("%du%d\n",tl[i],hd[(i+1)%k]);</pre>
34
       return 0;
35
```

## 3.1.13 Find Bridges

```
1 #include "header.h"
1 int n; // number of nodes
3 vvi adj; // adjacency list of graph
4 vector < bool > visited:
5 vi tin, low;
6 int timer;
7 \text{ void } dfs(int v, int p = -1) {
      visited[v] = true;
      tin[v] = low[v] = timer++:
      for (int to : adj[v]) {
10
          if (to == p) continue;
11
          if (visited[to]) {
12
               low[v] = min(low[v], tin[to]);
13
          } else {
14
               dfs(to, v);
               low[v] = min(low[v], low[to]);
```

```
if (low[to] > tin[v])
18
                   IS_BRIDGE(v, to);
           }
      }
20
21 }
22 void find_bridges() {
      timer = 0;
      visited.assign(n, false);
      tin.assign(n, -1);
25
      low.assign(n, -1);
26
      for (int i = 0: i < n: ++i) {
27
           if (!visited[i]) dfs(i);
28
29
30 }
```

#### **3.1.14** Find articulation points (i.e. cut off points)

```
1 #include "header.h"
2 int n: // number of nodes
3 vvi adj; // adjacency list of graph
4 vector < bool > visited;
5 vi tin, low;
6 int timer;
7 \text{ void } dfs(int v, int p = -1) {
       visited[v] = true:
       tin[v] = low[v] = timer++;
      int children=0;
      for (int to : adj[v]) {
11
           if (to == p) continue;
12
           if (visited[to]) {
13
               low[v] = min(low[v], tin[to]);
14
           } else {
15
16
               dfs(to, v);
               low[v] = min(low[v], low[to]);
               if (low[to] >= tin[v] && p!=-1) IS_CUTPOINT(v);
               ++children:
19
           }
20
21
      if (p == -1 && children > 1)
22
           IS_CUTPOINT(v);
23
24 }
25 void find_cutpoints() {
       timer = 0:
      visited.assign(n, false);
      tin.assign(n, -1);
28
      low.assign(n, -1);
      for (int i = 0; i < n; ++i) {</pre>
30
           if (!visited[i]) dfs (i);
31
32
33
```

# 3.2 Dynamic Programming

## 3.2.1 Longest Increasing Subsequence

```
1 #include "header.h"
```

```
2 template < class T>
3 vector<T> index path lis(vector<T>& nums) {
    int n = nums.size();
    vector <T> sub;
      vector < int > subIndex:
    vector <T> path(n, -1);
    for (int i = 0; i < n; ++i) {</pre>
        if (sub.emptv() || sub[sub.size() - 1] < nums[i]) {</pre>
      path[i] = sub.empty() ? -1 : subIndex[sub.size() - 1];
      sub.push_back(nums[i]);
      subIndex.push_back(i);
       } else {
      int idx = lower_bound(sub.begin(), sub.end(), nums[i]) - sub.begin();
      path[i] = idx == 0 ? -1 : subIndex[idx - 1];
15
      sub[idx] = nums[i]:
      subIndex[idx] = i;
19
    vector <T> ans:
    int t = subIndex[subIndex.size() - 1];
21
    while (t != -1) {
         ans.push_back(t);
         t = path[t];
    reverse(ans.begin(), ans.end());
    return ans:
28 }
29 // Length only
30 template < class T>
31 int length_lis(vector<T> &a) {
    set <T> st:
   typename set<T>::iterator it;
    for (int i = 0; i < a.size(); ++i) {</pre>
     it = st.lower_bound(a[i]);
    if (it != st.end()) st.erase(it):
      st.insert(a[i]);
    return st.size();
```

#### 3.2.2 0-1 Knapsack

```
1 #include "header.h"
2 // given a number of coins, calculate all possible distinct sums
3 int main() {
4 int n:
   vi coins(n); // all possible coins to use
  int sum = 0;  // sum of the coins
                         // dp[x] = 1 if sum x can be made
  vi dp(sum + 1, 0):
   dp[0] = 1;
                             // sum 0 can be made
  for (int c = 0; c < n; ++c)
                                  // first iteration: sums with first
    for (int x = sum: x >= 0: --x)
                                   // coin. next first 2 coins etc
       if (dp[x]) dp[x + coins[c]] = 1; // if sum x valid, x+c valid
11
```

#### 3.3 Trees

#### 3.3.1 Tree diameter

```
1 #include "header.h"
2 \text{ const int } mxN = 2e5 + 5;
3 int n, d[mxN]; // distance array
4 vi adj[mxN]; // tree adjacency list
5 void dfs(int s, int e) {
                         // recursively calculate the distance from the
   d[s] = 1 + d[e]:
        starting node to each node
   for (auto u : adi[s]) { // for each adiacent node
      if (u != e) dfs(u, s): // don't move backwards in the tree
10 }
11 int main() {
   // read input, create adj list
   dfs(0, -1);
                                   // first dfs call to find farthest node from
         arbitrary node
   dfs(distance(d, max_element(d, d + n)), -1); // second dfs call to find
        farthest node from that one
    cout << *max_element(d, d + n) - 1 << '\n'; // distance from second node</pre>
        to farthest is the diameter
16 }
```

#### 3.3.2 Tree Node Count

```
#include "header.h"

// calculate amount of nodes in each node's subtree

const int mxN = 2e5 + 5;

int n, cnt[mxN];

vi adj[mxN];

void dfs(int s = 0, int e = -1) {

cnt[s] = 1; // count leaves as one

for (int u : adj[s]) {

dfs(u, s);

cnt[s] += cnt[u]; // add up nodes of the subtrees
}
```

# 3.4 Number Theory / Combinatorics

#### **3.4.1** Modular exponentiation Or use pow() in python

```
#include "header.h"
2 ll mod_pow(ll base, ll exp, ll mod) {
3   if (mod == 1) return 0;
4   if (exp == 0) return 1;
5   if (exp == 1) return base;
6
7   ll res = 1;
8   base %= mod;
9   while (exp) {
10    if (exp % 2 == 1) res = (res * base) % mod;
11   exp >>= 1;
12   base = (base * base) % mod;
```

```
return res % mod;
15
16 }
```

## **3.4.2** GCD Or math.gcd in python, std::gcd in C++

```
1 #include "header.h"
2 ll gcd(ll a. ll b) {
3 if (a == 0) return b;
   return gcd(b % a, a);
5 }
```

#### 3.4.3 Sieve of Eratosthenes

```
1 #include "header.h"
2 vl primes;
3 void getprimes(ll n) { // Up to n (not included)
      vector < bool > p(n, true);
      p[0] = false;
      p[1] = false;
      for(11 i = 0; i < n; i++) {</pre>
           if(p[i]) {
               primes.push_back(i);
               for(ll j = i*2; j < n; j+=i) p[j] = false;</pre>
10
11 }}}
```

#### 3.4.4 Fibonacci % prime

```
1 #include "header.h"
2 const 11 MOD = 1000000007;
3 unordered_map<11, 11> Fib;
4 ll fib(ll n) {
      if (n < 2) return 1;
      if (Fib.find(n) != Fib.end()) return Fib[n];
      Fib[n] = (fib((n + 1) / 2) * fib(n / 2) + fib((n - 1) / 2) * fib((n - 2))
          / 2)) % MOD;
      return Fib[n];
9 }
```

## 3.4.5 nCk % prime

```
1 #include "header.h"
2 ll binom(ll n, ll k) {
      11 \text{ ans} = 1:
      for (ll i = 1; i <= min(k,n-k); ++i) ans = ans*(n+1-i)/i;
      return ans;
6 }
7 ll mod_nCk(ll n, ll k, ll p ){
      ll ans = 1:
      while(n){
           11 np = n\%p, kp = k\%p;
10
          if(kp > np) return 0;
           ans *= binom(np,kp);
12
```

```
n /= p; k /= p;
13
14
15
       return ans;
16 }
```

#### Strings 3.5

**3.5.1** Aho-Corasick algorithm Also can be used as Knuth-Morris-Pratt algorithm

```
1 #include "header.h"
2
3 map < char, int > cti;
4 int cti_size;
5 template <int ALPHABET_SIZE, int (*mp)(char)>
6 struct AC_FSM {
    struct Node {
      int child[ALPHABET_SIZE], failure = 0, match_par = -1;
      Node() { for (int i = 0; i < ALPHABET_SIZE; ++i) child[i] = -1; }
    }:
11
12
    vector < Node > a;
    vector<string> &words;
    AC_FSM(vector<string> &words) : words(words) {
      a.push_back(Node());
16
      construct_automaton();
    }
17
    void construct_automaton() {
      for (int w = 0, n = 0; w < words.size(); ++w, n = 0) {</pre>
19
        for (int i = 0; i < words[w].size(); ++i) {</pre>
20
           if (a[n].child[mp(words[w][i])] == -1) {
21
             a[n].child[mp(words[w][i])] = a.size();
22
             a.push_back(Node());
24
           n = a[n].child[mp(words[w][i])];
25
26
        a[n].match.push_back(w);
27
28
      queue < int > q:
29
      for (int k = 0: k < ALPHABET SIZE: ++k) {
30
        if (a[0].child[k] == -1) a[0].child[k] = 0;
31
        else if (a[0].child[k] > 0) {
32
33
           a[a[0].child[k]].failure = 0;
           q.push(a[0].child[k]);
34
35
36
37
      while (!q.empty()) {
        int r = q.front(); q.pop();
38
        for (int k = 0, arck; k < ALPHABET_SIZE; ++k) {</pre>
39
           if ((arck = a[r].child[k]) != -1) {
40
             q.push(arck);
41
             int v = a[r].failure;
42
             while (a[v].child[k] == -1) v = a[v].failure:
43
             a[arck].failure = a[v].child[k];
44
             a[arck].match_par = a[v].child[k];
45
             while (a[arck].match_par != -1
46
                 && a[a[arck].match_par].match.empty())
47
               a[arck].match_par = a[a[arck].match_par].match_par;
           }
```

```
}
51
52
    void aho_corasick(string &sentence, vvi &matches){
53
      matches.assign(words.size(), vi());
54
      int state = 0, ss = 0;
55
      for (int i = 0; i < sentence.length(); ++i, ss = state) {</pre>
56
        while (a[ss].child[mp(sentence[i])] == -1)
           ss = a[ss].failure;
        state = a[state].child[mp(sentence[i])]
59
             = a[ss].child[mp(sentence[i])];
         for (ss = state; ss != -1; ss = a[ss].match_par)
61
          for (int w : a[ss].match)
             matches[w].push_back(i + 1 - words[w].length());
63
64
66 };
67 int char_to_int(char c) {
    return cti[c]:
69 }
```

#### 3.5.2 KMP

```
1 #include "header.h"
void compute_prefix_function(string &w, vi &prefix) {
    prefix.assign(w.length(), 0);
    int k = prefix[0] = -1;
    for(int i = 1; i < w.length(); ++i) {</pre>
      while (k >= 0 \&\& w[k + 1] != w[i]) k = prefix[k];
      if(w[k + 1] == w[i]) k++;
      prefix[i] = k;
10
11 }
12 void knuth_morris_pratt(string &s, string &w) {
    int q = -1;
    vi prefix;
14
    compute_prefix_function(w, prefix);
    for(int i = 0; i < s.length(); ++i) {</pre>
      while (q >= 0 \&\& w[q + 1] != s[i]) q = prefix[q];
17
      if(w[q + 1] == s[i]) q++;
18
      if (q + 1 == w.length()) {
19
        // Match at position (i - w.length() + 1)
21
         q = prefix[q];
22
    }
23
24 }
```

# 3.6 Geometry

# 3.6.1 essentials.cpp

```
#include "../header.h"

using C = ld; // could be long long or long double

constexpr C EPS = 1e-10; // change to 0 for C=11

struct P { // may also be used as a 2D vector
```

```
P(C x = 0, C y = 0) : x(x), y(y) {}
    P operator+ (const P &p) const { return {x + p.x, y + p.y}; }
    P operator - (const P &p) const { return {x - p.x, y - p.y}; }
    P operator* (C c) const { return {x * c, y * c}; }
    P operator/ (C c) const { return {x / c, y / c}; }
    C operator* (const P &p) const { return x*p.x + y*p.y; }
    C operator^ (const P &p) const { return x*p.y - p.x*y; }
    P perp() const { return P{y, -x}; }
    C lensq() const { return x*x + y*y; }
    ld len() const { return sqrt((ld)lensq()); }
    static ld dist(const P &p1, const P &p2) {
      return (p1-p2).len(); }
    bool operator == (const P &r) const {
19
      return ((*this)-r).lensq() <= EPS*EPS; }</pre>
    det(P p1, P p2) { return p1^p2; }
22 C det(P p1, P p2, P o) { return det(p1-o, p2-o); }
    det(const vector <P> &ps) {
    C sum = 0; P prev = ps.back();
    for(auto &p : ps) sum += det(p, prev), prev = p;
    return sum;
27 }
28 // Careful with division by two and C=11
    area(P p1, P p2, P p3) { return abs(det(p1, p2, p3))/C(2); }
    area(const vector <P> &poly) { return abs(det(poly))/C(2); }
31 int sign(C c) { return (c > C(0)) - (c < C(0)); }
32 int ccw(P p1, P p2, P o) { return sign(det(p1, p2, o)); }
_{34} // Only well defined for C = 1d.
35 P unit(const P &p) { return p / p.len(); }
_{36} P rotate(P p, 1d a) { return P{p.x*cos(a)-p.y*sin(a), p.x*sin(a)+p.y*cos(a)}
      }; }
```

#### 3.6.2 Convex Hull

```
1 #include "header.h"
2 #include "essentials.cpp"
3 struct ConvexHull { // O(n lg n) monotone chain.
    size_t n;
    vector<size_t> h, c; // Indices of the hull are in 'h', ccw.
    const vector <P> &p;
    ConvexHull(const vector < P > & _ p) : n(_p.size()), c(n), p(_p) {
      std::iota(c.begin(), c.end(), 0);
      std::sort(c.begin(), c.end(), [this](size_t 1, size_t r) -> bool {
          return p[1].x != p[r].x ? p[1].x < p[r].x : p[1].y < p[r].y; });
      c.erase(std::unique(c.begin(), c.end(), [this](size_t l, size_t r) {
10
          return p[1] == p[r]; }), c.end());
1.1
      for (size_t s = 1, r = 0; r < 2; ++r, s = h.size()) {
12
        for (size_t i : c) {
          while (h.size() > s \&\& ccw(p[h.end()[-2]], p[h.end()[-1]], p[i]) \le
13
            h.pop_back();
14
15
          h.push_back(i);
16
        reverse(c.begin(), c.end());
17
18
      if (h.size() > 1) h.pop_back();
19
```

```
size t size() const { return h.size(): }
21
    template <class T, void U(const P &, const P &, const P &, T &)>
    void rotating_calipers(T &ans) {
      if (size() <= 2)
        U(p[h[0]], p[h.back()], p[h.back()], ans);
25
26
        for (size t i = 0, i = 1, s = size(); i < 2 * s; ++i) {
          while (\det(p[h[(i + 1) \% s]) - p[h[i \% s]), p[h[(j + 1) \% s]] - p[h[
              j]]) >= 0)
            j = (j + 1) \% s;
          U(p[h[i % s]], p[h[(i + 1) % s]], p[h[j]], ans);
32
33 };
     Example: furthest pair of points. Now set ans = OLL and call
     ConvexHull(pts).rotating_calipers<11, update>(ans);
36 void update (const P &p1, const P &p2, const P &o, 11 &ans) {
    ans = \max(\text{ans.}(11)\max((p1 - o).lensq(), (p2 - o).lensq())):
38 }
```

# 3.7 Other Algorithms

## 3.8 Other Data Structures

# **3.8.1** Disjoint set (i.e. union-find)

```
1 template <typename T>
2 class DisjointSet {
      typedef T * iterator;
      T *parent, n, *rank;
      public:
          // O(n), assumes nodes are [0, n)
          DisjointSet(T n) {
              this->parent = new T[n];
              this -> n = n:
              this->rank = new T[n];
              for (T i = 0; i < n; i++) {</pre>
                   parent[i] = i;
                   rank[i] = 0:
              }
          }
          // O(log n)
          T find set(T x) {
19
               if (x == parent[x]) return x;
20
               return parent[x] = find_set(parent[x]);
          }
22
23
24
          // O(\log n)
          void union_sets(T x, T y) {
25
              x = this->find_set(x);
              y = this->find_set(y);
              if (x == y) return;
```

```
if (rank[x] < rank[y]) {
        T z = x;
        x = y;
        y = z;
}

parent[y] = x;
if (rank[x] == rank[y]) rank[x]++;
}

// T z = x;
// T
```

**3.8.2** Fenwick tree (i.e. BIT) eff. update + prefix sum calc.

```
1 #include "header.h"
2 #define maxn 200010
3 int t,n,m,tree[maxn],p[maxn];
5 void update(int k, int z) {
      while (k <= maxn) {</pre>
           tree[k] += z;
           k += k & (-k);
           // cout << "k: " << k << endl;
10
11 }
13 int sum(int k) {
      int ans = 0:
      while(k) {
           ans += tree[k];
           k = k & (-k):
18
      return ans;
19
```

# 4 Other Mathematics

# 4.1 Helpful functions

**4.1.1 Euler's Totient Fucntion**  $n = p_1^{k_1-1} \cdot (p_1-1) \cdot \ldots \cdot p_r^{k_r-1} \cdot (p_r-1)$ , where  $p_1^{k_1} \cdot \ldots \cdot p_r^{k_r}$  is the prime factorization of n.

```
if (n > 1) ans *= n-1;
return ans;

for it phis(int n) { // All \Phi(i) up to n}

vi phi(n + 1, OLL);

iota(phi.begin(), phi.end(), OLL);

for (ll i = 2LL; i <= n; ++i)

if (phi[i] == i)

for (ll j = i; j <= n; j += i)

phi[j] -= phi[j] / i;

return phi;
</pre>
```

Formulas  $\Phi(n)$  counts all numbers in  $1, \ldots, n-1$  coprime to n.  $a^{\varphi(n)} \equiv 1 \mod n$ , a and n are coprimes.  $\forall e > \log_2 m : n^e \mod m = n^{\Phi(m) + e \mod \Phi(m)} \mod m$ .  $\gcd(m, n) = 1 \Rightarrow \Phi(m \cdot n) = \Phi(m) \cdot \Phi(n)$ .

## 4.2 Theorems and definitions

Fermat's little theorem  $a^p \equiv a \mod p$ 

**Subfactorial** 
$$!n = n! \sum_{i=0}^{n} \frac{(-1)^i}{i!}, !(0) = 1, !n = n \cdot !(n-1) + (-1)^n$$

**Least common multiple**  $lcm(a, b) = a \cdot b/gcd(a, b)$ 

Binomials and other partitionings We have  $\binom{n}{k} = \binom{n-1}{k} + \binom{n-1}{k-1} = \prod_{i=1}^k \frac{n-i+1}{i}$ . This last product may be computed incrementally since any product of k' consecutive values is divisible by k'!. Basic identities: The hockeystick identity:  $\sum_{k=r}^{n} \binom{k}{r} = \binom{n+1}{r+1}$  or  $\sum_{k \leq n} \binom{r+k}{k} = \binom{r+n+1}{n}$ . Also  $\sum_{k=0}^{n} \binom{k}{m} = \binom{n+1}{m+1}$ .

For  $n, m \ge 0$  and p prime. Write n, m in base p, i.e.  $n = n_k p^k + \cdots + n_1 p + n_0$  and  $m = m_k p^k + \cdots + m_1 p + m_0$ . Then by Lucas theorem we have  $\binom{n}{m} \equiv \prod_{i=0}^k \binom{n_i}{m_i} \mod p$ , with the convention that  $n_i < m_i \implies \binom{n_i}{m_i} = 0$ .

Fibonacci (See also number theory section)

$$\sum_{0 \le k \le n} {n-k \choose k} = F_{n+1}, F_n = \frac{1}{\sqrt{5}} \left(\frac{1+\sqrt{5}}{2}\right)^n - \frac{1}{\sqrt{5}} \left(\frac{1-\sqrt{5}}{2}\right)^n,$$

$$\sum_{i=1}^n F_i = F_{n+2} - 1, \sum_{i=1}^n F_i^2 = F_n F_{n+1},$$

$$\gcd(F_m, F_n) = F_{\gcd(m,n)}, \gcd(F_n, F_{n+1}) = \gcd(F_n, F_{n+2}) = 1$$

Bit stuff  $a + b = a \oplus b + 2(a \& b) = a|b + a \& b$ .

kth bit is set in x iff  $x \mod 2^{k-1} \ge 2^k$ , or iff  $x \mod 2^{k-1} - x \mod 2^k \ne 0$  (i.e.  $= 2^k$ ) It comes handy when you need to look at the bits of the numbers which are pair sums or subset sums etc.

```
n \mod 2^i = n\&(2^i - 1).
\forall k: \ 1 \oplus 2 \oplus \ldots \oplus (4k - 1) = 0
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