file

```
16 #define vpi vector<pi>
17 #define vld vector<ld>
18 #define vvpi vector < vpi>
19 #define in fast(el. cont) (cont.find(el) != cont.
      end())
20 #define in(el, cont) (find(cont.begin(), cont.end
      (), el) != cont.end())
21 #define all(x) x.begin(), x.end()
22 #define rall(x) x.rbegin(), x.rend()
_{24} constexpr int INF = 2000000010;
25 constexpr ll LLINF = 900000000000000010LL;
27 // int main() {
28 // ios::sync_with_stdio(false); // do not use
      cout + printf
29 // cin.tie(NULL);
30 // cout << fixed << setprecision(12);
31 // return 0;
32 // }
```

1.2 Bash for c++ compile with header.h

```
1 #!/bin/bash
2 if [ $# -ne 1 ]; then echo "Usage: $0 <input_file</pre>
     >"; exit 1;fi
3 f="$1":d=code/:o=a.out
4 [ -f $d/$f ] || { echo "Input file not found: $f
      "; exit 1; }
5 g++ -I$d $d/$f -o $o && echo "Compilation
      successful. Executable '$0' created." || echo
       "Compilation failed."
```

1.3 Bash for run tests c++

```
1 g++ $1/$1.cpp -o $1/$1.out
2 for file in $1/*.in; do diff <($1/$1.out < "$file</pre>
      ") "${file%.in}.ans"; done
```

1.4 Bash for run tests python

```
1 for file in $1/*.in; do diff <(python3 $1/$1.py <
      "$file") "${file%.in}.ans"; done
```

1.4.1 Aux. helper C++

```
1 #include "header.h"
3 int main() {
      // Read in a line including white space
      string line:
      getline(cin, line);
      // When doing the above read numbers as
          follows:
      int n:
      getline(cin, line);
      stringstream ss(line);
      ss >> n;
12
      // Count the number of 1s in binary
          represnatation of a number
      ull number:
       __builtin_popcountll(number);
18 // __int128
19 using 111 = __int128;
20 ostream& operator << ( ostream& o, __int128 n ) {</pre>
    auto t = n < 0? -n : n; char b[128], *d = end(b)
    do *--d = '0'+t\%10, t /= 10; while (t);
    if(n<0) *--d = '-':
    o.rdbuf()->sputn(d,end(b)-d);
    return o;
26 }
```

1.4.2 Aux. helper python

```
1 from functools import lru_cache
3 # Read until EOF
4 while True:
          pattern = input()
      except EOFError:
          break
10 @lru cache(maxsize=None)
11 def smth_memoi(i, j, s):
      # Example in-built cache
      return "sol"
15 # Fast T
16 import io, os
17 def fast_io():
      finput = io.BvtesIO(os.read(0.
          os.fstat(0).st_size)).readline
      s = finput().decode()
      return s
```

```
23 # Fast O
24 import sys
25 def fast out():
      sys.stdout.write(str(n)+"\n")
```

2 Pvthon

2.1 Graphs

2.1.1 BFS

```
1 from collections import deque
2 def bfs(g, roots, n):
     q = deque(roots)
     explored = set()
      distances = [0 if v in roots else float('inf'
         ) for v in range(n)]
     while len(q) != 0:
          node = q.popleft()
          if node in explored: continue
          explored.add(node)
          for neigh in g[node]:
              if neigh not in explored:
                  q.append(neigh)
                  distances[neigh] = distances[node
      return distances
```

2.1.2 Dijkstra

```
1 from heapq import *
2 def dijkstra(n, root, g): # g = {node: (cost,
     neigh)}
3 dist = [float("inf")]*n
   dist[root] = 0
   prev = [-1]*n
   pq = [(0, root)]
   heapifv(pg)
   visited = set([])
   while len(pq) != 0:
     _, node = heappop(pq)
     if node in visited: continue
      visited.add(node)
      # In case of disconnected graphs
```

37

39

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```
if node not in g:
continue

for cost, neigh in g[node]:
    alt = dist[node] + cost
    if alt < dist[neigh]:
        dist[neigh] = alt
        prev[neigh] = node
        heappush(pq, (alt, neigh))
    return dist</pre>
```

2.1.3 Topological Sort

```
1 #Python program to print topological sorting of a
2 from collections import defaultdict
4 #Class to represent a graph
5 class Graph:
      def __init__(self,vertices):
          self.graph = defaultdict(list) #
              dictionary containing adjacency List
          self.V = vertices #No. of vertices
8
      # function to add an edge to graph
10
      def addEdge(self,u,v):
11
          self.graph[u].append(v)
12
13
      # A recursive function used by
14
          topologicalSort
      def topologicalSortUtil(self,v,visited,stack)
15
16
          # Mark the current node as visited.
17
          visited[v] = True
18
19
          # Recur for all the vertices adjacent to
              this vertex
          for i in self.graph[v]:
21
              if visited[i] == False:
22
                  self.topologicalSortUtil(i,
23
                      visited, stack)
          # Push current vertex to stack which
25
              stores result
          stack.insert(0,v)
26
27
      # The function to do Topological Sort. It
28
          uses recursive
      # topologicalSortUtil()
29
      def topologicalSort(self):
30
          # Mark all the vertices as not visited
31
          visited = [False]*self.V
32
          stack =[]
```

```
# Call the recursive helper function to
        store Topological
    # Sort starting from all vertices one by
    for i in range(self.V):
        if visited[i] == False:
            self.topologicalSortUtil(i,
                visited.stack)
    # Print contents of stack
    return stack
def isCyclicUtil(self, v, visited, recStack):
    # Mark current node as visited and
    # adds to recursion stack
    visited[v] = True
    recStack[v] = True
    # Recur for all neighbours
    # if any neighbour is visited and in
    # recStack then graph is cyclic
    for neighbour in self.graph[v]:
        if visited[neighbour] == False:
            if self.isCyclicUtil(neighbour,
                visited. recStack) == True:
                return True
        elif recStack[neighbour] == True:
            return True
    # The node needs to be popped from
    # recursion stack before function ends
    recStack[v] = False
    return False
# Returns true if graph is cyclic else false
def isCvclic(self):
    visited = [False] * (self.V + 1)
    recStack = [False] * (self.V + 1)
    for node in range(self.V):
        if visited[node] == False:
            if self.isCyclicUtil(node,
                visited. recStack) == True:
                return True
    return False
```

2.1.4 Kruskal (UnionFind)

```
1 class UnionFind:
2    def __init__(self, n):
3         self.parent = [-1]*n
4
5    def find(self, x):
```

```
if self.parent[x] < 0:</pre>
               return x
           self.parent[x] = self.find(self.parent[x
           return self.parent[x]
10
      def connect(self, a, b):
           ra = self.find(a)
12
           rb = self.find(b)
           if ra == rb:
               return False
           if self.parent[ra] > self.parent[rb]:
               self.parent[rb] += self.parent[ra]
               self.parent[ra] = rb
           else:
               self.parent[ra] += self.parent[rb]
               self.parent[rb] = ra
           return True
24 # Full MST is len(spanning==n-1)
25 def kruskal(n, edges):
      uf = UnionFind(n)
      spanning = []
      edges.sort(key = lambda d: -d[2])
      while edges and len(spanning) < n-1:
          u, v, w = edges.pop()
          if not uf.connect(u, v):
31
               continue
           spanning.append((u, v, w))
33
      return spanning
36 # Example
_{37} \text{ edges} = [(1, 2, 10), (2, 3, 20)]
```

2.2 Num. Th. / Comb.

2.2.1 nCk % prime

```
# Note: p must be prime and k  n:
        return 0
    # calculate numerator
    num = 1
    for i in range(n-k+1, n+1):
        num *= i % p
    num %= p
    # calculate denominator
    denom = 1
    for i in range(1,k+1):
        denom *= i % p
    denom %= p
    # numerator * denominator^(p-2) (mod p)
    return (num * pow(denom, p-2, p)) % p
```

2.2.2 Sieve of E. O(n) so actually faster than C++ version, but more memory

```
_{1} MAX STZE = 10**8+1
2 isprime = [True] * MAX SIZE
3 prime = []
4 SPF = [None] * (MAX SIZE)
6 def manipulated_seive(N): # Up to N (not
      included)
    isprime[0] = isprime[1] = False
    for i in range(2. N):
      if isprime[i] == True:
        prime.append(i)
        SPF[i] = i
11
      while (j < len(prime) and
       i * prime[j] < N and</pre>
          prime[i] <= SPF[i]):</pre>
        isprime[i * prime[j]] = False
        SPF[i * prime[j]] = prime[j]
        j += 1
```

2.3 Strings

2.3.1 LCS

```
1 def longestCommonSubsequence(text1, text2): # 0(
      m*n) time, O(m) space
      n = len(text1)
      m = len(text2)
      # Initializing two lists of size m
      prev = [0] * (m + 1)
      cur = \lceil 0 \rceil * (m + 1)
      for idx1 in range(1, n + 1):
          for idx2 in range(1, m + 1):
              # If characters are matching
11
              if text1[idx1 - 1] == text2[idx2 -
                   cur[idx2] = 1 + prev[idx2 - 1]
                   # If characters are not matching
15
                   cur[idx2] = max(cur[idx2 - 1],
                       prev[idx2])
          prev = cur.copy()
18
19
      return cur[m]
```

2.3.2 KMP

```
1 class KMP:
      def partial(self, pattern):
2
          """ Calculate partial match table: String
               -> [Int]"""
          ret = [0]
          for i in range(1, len(pattern)):
              j = ret[i - 1]
              while j > 0 and pattern[j] != pattern
                  [i]: j = ret[j - 1]
              ret.append(j + 1 if pattern[j] ==
                  pattern[i] else j)
          return ret
11
      def search(self, T, P):
          """KMP search main algorithm: String ->
12
              String -> [Int]
          Return all the matching position of
13
              pattern string P in T"""
          partial, ret, j = self.partial(P), [], 0
          for i in range(len(T)):
              while j > 0 and T[i] != P[j]: j =
16
                  partial[j - 1]
              if T[i] == P[i]: i += 1
17
              if j == len(P):
                  ret.append(i - (j - 1))
                  j = partial[j - 1]
          return ret
```

2.3.3 Edit distance

```
def editDistance(str1, str2):
   # Get the lengths of the input strings
   m = len(str1)
   n = len(str2)
   # Initialize a list to store the current row
   curr = \lceil 0 \rceil * (n + 1)
   # Initialize the first row with values from 0
   for j in range(n + 1):
     curr[i] = i
   # Initialize a variable to store the previous
       value
   previous = 0
   # Loop through the rows of the dynamic
       programming matrix
   for i in range (1, m + 1):
     # Store the current value at the beginning of
           the row
     previous = curr[0]
     curr[0] = i
```

```
# Loop through the columns of the dynamic
      programming matrix
  for i in range(1, n + 1):
    # Store the current value in a temporary
    temp = curr[i]
    # Check if the characters at the current
        positions in str1 and str2 are the same
    if str1[i - 1] == str2[i - 1]:
      curr[j] = previous
      # Update the current cell with the
          minimum of the three adjacent cells
      curr[j] = 1 + min(previous, curr[j - 1],
          curr[i])
    # Update the previous variable with the
       temporary value
    previous = temp
# The value in the last cell represents the
    minimum number of operations
return curr[n]
```

2.4 Other Algorithms

2.4.1 Rotate matrix

```
1 def rotate_matrix(m):
2    return [[m[j][i] for j in range(len(m))] for
        i in range(len(m[0])-1,-1,-1)]
```

2.5 Geometry

2.5.1 Convex Hull

```
def vec(a,b):
    return (b[0]-a[0],b[1]-a[1])
def det(a,b):
    return a[0]*b[1] - b[0]*a[1]

def convexhull(P):
    if (len(P) == 1):
        return [(p[0][0], p[0][1])]

h = sorted(P)
lower = []
i = 0
while i < len(h):
    if len(lower) > 1:
```

```
a = vec(lower[-2], lower[-1])
              b = vec(lower[-1], h[i])
              if det(a,b) <= 0 and len(lower) > 1:
                   lower.pop()
                   continue
19
          lower.append(h[i])
20
          i += 1
21
22
      upper = []
23
      i = 0
24
      while i < len(h):
25
          if len(upper) > 1:
              a = vec(upper[-2], upper[-1])
27
              b = vec(upper[-1], h[i])
              if det(a,b) >= 0:
                   upper.pop()
                   continue
31
          upper.append(h[i])
32
          i += 1
33
34
      reversedupper = list(reversed(upper[1:-1:]))
35
      reversedupper.extend(lower)
36
      return reversedupper
```

2.5.2 Geometry

```
2 def vec(a,b):
      return (b[0]-a[0],b[1]-a[1])
5 def det(a,b):
      return a[0]*b[1] - b[0]*a[1]
      lower = []
      i = 0
      while i < len(h):
          if len(lower) > 1:
              a = vec(lower[-2], lower[-1])
              b = vec(lower[-1], h[i])
              if det(a,b) <= 0 and len(lower) > 1:
14
                  lower.pop()
15
                  continue
          lower.append(h[i])
17
          i += 1
18
19
      # find upper hull
20
      # det <= 0 -> replace
      upper = []
22
      i = 0
23
      while i < len(h):
^{24}
          if len(upper) > 1:
25
              a = vec(upper[-2], upper[-1])
              b = vec(upper[-1], h[i])
27
              if det(a,b) >= 0:
```

2.6 Other Data Structures

2.6.1 Segment Tree

```
_{1} N = 100000 # limit for array size
2 tree = [0] * (2 * N) # Max size of tree
4 def build(arr. n): # function to build the tree
      # insert leaf nodes in tree
      for i in range(n):
          tree[n + i] = arr[i]
      # build the tree by calculating parents
      for i in range(n - 1, 0, -1):
          tree[i] = tree[i << 1] + tree[i << 1 | 1]
11
13 def updateTreeNode(p, value, n): # function to
      update a tree node
      # set value at position p
      tree[p + n] = value
16
      i = p # move upward and update parents
19
          tree[i >> 1] = tree[i] + tree[i ^ 1]
          i >>= 1
23 def query(1, r, n): # function to get sum on
      interval [1, r)
      res = 0
      # loop to find the sum in the range
      1 += n
      r += n
      while 1 < r:
         if 1 & 1:
29
              res += tree[1]
              1 += 1
31
          if r & 1:
              r -= 1
34
              res += tree[r]
          1 >>= 1
35
36
          r >>= 1
      return res
```

2.6.2 Trie

```
class TrieNode:
def __init__(self):
```

```
self.children = [None] *26
           self.isEndOfWord = False
6 class Trie:
      def __init__(self):
          self.root = self.getNode()
      def getNode(self):
10
          return TrieNode()
12
13
      def charToIndex(self.ch):
          return ord(ch)-ord('a')
15
      def insert(self,key):
          pCrawl = self.root
          length = len(key)
          for level in range(length):
               index = self._charToIndex(key[level])
21
               if not pCrawl.children[index]:
                  pCrawl.children[index] = self.
                       getNode()
               pCrawl = pCrawl.children[index]
          pCrawl.isEndOfWord = True
25
26
27
      def search(self, key):
          pCrawl = self.root
28
          length = len(kev)
          for level in range(length):
30
              index = self._charToIndex(key[level])
              if not pCrawl.children[index]:
                   return False
               pCrawl = pCrawl.children[index]
          return pCrawl.isEndOfWord
```

3 C++

3.1 Graphs

3.1.1 BFS

```
while (not q.empty()) {
          int node = q.front();
13
          q.pop();
14
          for (auto neigh: g[node]) {
15
              if (not in(neigh, visited)) {
                   parents[neigh] = node;
                  q.emplace(neigh);
                   visited.insert(neigh):
              }
          }
21
22
23
      return parents;
24 }
25 vi reconstruct_path(vi parents, int start, int
      goal) {
      vi path;
      int curr = goal;
      while (curr != start) {
          path.push_back(curr);
          if (parents[curr] == -1) return vi(); //
              No path, empty vi
          curr = parents[curr];
32
      path.push_back(start);
33
      reverse(path.begin(), path.end());
34
      return path;
35
36 }
```

3.1.2 DFS Cycle detection / removal

```
1 #include "header.h"
void removeCyc(ll node, unordered_map<ll, vector<</pre>
      pair<ll, ll>>>& neighs, vector<bool>& visited
3 vector < bool > & recStack, vector < 11 > & ans) {
      if (!visited[node]) {
          visited[node] = true:
          recStack[node] = true;
          auto it = neighs.find(node);
          if (it != neighs.end()) {
              for (auto util: it->second) {
                  11 nnode = util.first:
                   if (recStack[nnode]) {
                       ans.push_back(util.second);
                  } else if (!visited[nnode]) {
                       removeCyc(nnode, neighs,
14
                           visited, recStack, ans);
                  }
              }
17
18
      recStack[node] = false;
19
```

3.1.3 Dijkstra

```
1 #include "header.h"
2 vector < int > dijkstra(int n, int root, map < int,</pre>
       vector<pair<int. int>>>& g) {
    unordered_set <int> visited;
    vector < int > dist(n, INF);
      priority_queue < pair < int , int >> pq;
      dist[root] = 0;
      pq.push({0, root});
       while (!pq.emptv()) {
           int node = pq.top().second;
           int d = -pq.top().first;
10
           pq.pop();
11
12
           if (in(node, visited)) continue;
           visited.insert(node);
14
           for (auto e : g[node]) {
               int neigh = e.first;
17
               int cost = e.second:
               if (dist[neigh] > dist[node] + cost)
19
                   dist[neigh] = dist[node] + cost;
                   pq.push({-dist[neigh], neigh});
21
           }
24
      return dist:
26 }
```

3.1.4 Floyd-Warshall

3.1.5 Kruskal Minimum spanning tree of undirected weighted graph

```
#include "header.h"
#include "disjoint_set.h"
// O(E log E)
```

```
4 pair < set < pair < 11, 11 >> , 11 > kruskal (vector < tuple</pre>
       <11, 11, 11>>& edges, 11 n) {
       set <pair <11, 11>> ans;
       11 cost = 0:
       sort(edges.begin(), edges.end());
       DisjointSet < 11 > fs(n);
       ll dist. i. i:
       for (auto edge: edges) {
           dist = get<0>(edge);
           i = get<1>(edge);
           j = get < 2 > (edge);
15
           if (fs.find_set(i) != fs.find_set(j)) {
               fs.union_sets(i, j);
               ans.insert({i, j});
                cost += dist;
           }
21
       return pair < set < pair < 11, 11>>, 11> {ans, cost
24 }
```

3.1.6 Hungarian algorithm

```
1 #include "header.h"
3 template <class T> bool ckmin(T &a, const T &b) {
       return b < a ? a = b, 1 : 0; }
4 /**
* Given J jobs and W workers (J <= W), computes
       the minimum cost to assign each
* prefix of jobs to distinct workers.
7 * @tparam T a type large enough to represent
       integers on the order of J *
9 * @param C a matrix of dimensions JxW such that
       C[j][w] = cost to assign j-th
* job to w-th worker (possibly negative)
12 * Creturn a vector of length J, with the j-th
       entry equaling the minimum cost
* to assign the first (j+1) jobs to distinct
15 template <class T> vector<T> hungarian(const
      vector < vector < T >> &C) {
      const int J = (int)size(C), W = (int)size(C
          [0]);
      assert(J <= W);</pre>
      // job[w] = job assigned to w-th worker, or
          -1 if no job assigned
      // note: a W-th worker was added for
          convenience
```

```
vector < int > job(W + 1, -1);
      vector < T > vs(J), vt(W + 1); // potentials
21
      // -vt[W] will equal the sum of all deltas
22
      vector <T> answers:
23
      const T inf = numeric_limits <T>::max();
^{24}
      for (int j_cur = 0; j_cur < J; ++j_cur) { //</pre>
25
           assign j_cur-th job
          int w_cur = W;
          iob[w_cur] = j_cur;
27
          // min reduced cost over edges from Z to
              worker w
          vector <T> min_to(W + 1, inf);
          vector<int> prv(W + 1, -1); // previous
              worker on alternating path
          vector < bool > in_Z(W + 1);  // whether
              worker is in Z
          while (job[w_cur] != -1) {    // runs at
              most j_cur + 1 times
              in_Z[w_cur] = true;
              const int j = job[w_cur];
              T delta = inf;
35
              int w next:
36
              for (int w = 0; w < W; ++w) {
                  if (!in_Z[w]) {
                      if (ckmin(min_to[w], C[j][w]
                           - vs[i] - vt[w]))
                           prv[w] = w_cur;
                       if (ckmin(delta. min to[w]))
                           w next = w:
                  }
              }
              // delta will always be non-negative,
              // except possibly during the first
                  time this loop runs
              // if any entries of C[i cur] are
                  negative
              for (int w = 0; w \le W; ++w) {
                  if (in_Z[w]) ys[job[w]] += delta,
                       vt[w] -= delta;
                   else min_to[w] -= delta;
              }
              w_cur = w_next;
51
52
          // update assignments along alternating
53
          for (int w: w cur != W: w cur = w) job[
              w_cur] = job[w = prv[w_cur]];
          answers.push_back(-yt[W]);
55
      return answers;
57
```

```
\textbf{3.1.7} \quad \textbf{Suc. shortest path} \quad \text{Calculates max flow, min} \\ \text{cost}
```

```
1 #include "header.h"
2 // map<node, map<node, pair<cost, capacity>>>
3 #define graph unordered map<int, unordered map<
      int, pair<ld, int>>>
4 graph g;
5 const ld infty = 1e60l; // Change if necessary
6 ld fill(int n, vld& potential) { // Finds max
      flow, min cost
    priority_queue < pair < ld, int >> pq;
    vector < bool > visited(n+2, false);
    vi parent(n+2, 0);
    vld dist(n+2, infty);
    dist[0] = 0.1;
    pg.emplace(make pair(0.1, 0)):
    while (not pq.empty()) {
      int node = pq.top().second;
      pq.pop();
15
      if (visited[node]) continue;
      visited[node] = true:
      for (auto& x : g[node]) {
        int neigh = x.first;
        int capacity = x.second.second;
20
        ld cost = x.second.first;
21
        if (capacity and not visited[neigh]) {
          ld d = dist[node] + cost + potential[node
23
              ] - potential[neigh];
          if (d + 1e-10l < dist[neigh]) {</pre>
24
            dist[neigh] = d;
             pq.emplace(make_pair(-d, neigh));
27
            parent[neigh] = node;
    }}}
28
    for (int i = 0; i < n+2; i++) {</pre>
      potential[i] = min(infty, potential[i] + dist
    if (not parent[n+1]) return infty;
    ld ans = 0.1:
    for (int x = n+1; x; x=parent[x]) {
      ans += g[parent[x]][x].first;
      g[parent[x]][x].second--:
      g[x][parent[x]].second++;
   }
39
    return ans;
41 }
```

3.1.8 Bipartite check

```
#include "header.h"
int main() {
   int n;
   vvi adj(n);
```

```
vi side(n. -1): // will have 0's for one
          side 1's for other side
      bool is bipartite = true: // becomes false
          if not bipartite
      queue < int > q;
      for (int st = 0; st < n; ++st) {</pre>
           if (side[st] == -1) {
              q.push(st);
11
               side[st] = 0:
               while (!q.empty()) {
                   int v = q.front();
                   q.pop();
                   for (int u : adj[v]) {
16
                       if (side[u] == -1) {
                           side[u] = side[v] ^ 1;
                           q.push(u);
19
                       } else {
20
                           is_bipartite &= side[u]
                               != side[v];
                       }
23 }}}}
```

3.1.9 Find cycle directed

```
1 #include "header.h"
2 int n;
3 \text{ const int } mxN = 2e5+5:
4 vvi adj(mxN);
5 vector < char > color;
6 vi parent;
7 int cycle_start, cycle_end;
8 bool dfs(int v) {
      color[v] = 1;
      for (int u : adj[v]) {
           if (color[u] == 0) {
11
               parent[u] = v:
               if (dfs(u)) return true;
          } else if (color[u] == 1) {
               cycle_end = v;
               cycle_start = u;
16
               return true;
17
           }
19
      color[v] = 2;
      return false:
23 void find_cycle() {
      color.assign(n, 0);
24
      parent.assign(n, -1);
      cvcle_start = -1;
      for (int v = 0: v < n: v++) {
           if (color[v] == 0 && dfs(v))break;
29
      if (cvcle start == -1) {
```

```
cout << "Acyclic" << endl;</pre>
      } else {
           vector<int> cvcle:
33
           cvcle.push back(cvcle start):
34
           for (int v = cycle_end; v != cycle_start;
35
                v = parent[v])
               cycle.push_back(v);
           cycle.push_back(cycle_start);
           reverse(cycle.begin(), cycle.end());
38
           cout << "Cycle_Found: ";
40
           for (int v : cvcle) cout << v << "":</pre>
41
           cout << endl;</pre>
42
43
44 }
```

3.1.10 Find cycle undirected

```
1 #include "header.h"
2 int n;
3 const int mxN = 2e5 + 5:
4 vvi adj(mxN);
5 vector < bool > visited;
6 vi parent;
7 int cycle_start, cycle_end;
8 bool dfs(int v, int par) { // passing vertex and
      its parent vertex
      visited[v] = true;
      for (int u : adj[v]) {
          if(u == par) continue; // skipping edge
11
              to parent vertex
          if (visited[u]) {
              cvcle_end = v;
               cvcle_start = u;
              return true;
15
16
          parent[u] = v;
          if (dfs(u, parent[u]))
              return true;
19
20
      return false;
21
23 void find_cycle() {
      visited.assign(n, false);
      parent.assign(n, -1);
25
      cycle_start = -1;
      for (int v = 0; v < n; v++) {
          if (!visited[v] && dfs(v, parent[v]))
28
29
      if (cvcle start == -1) {
30
          cout << "Acyclic" << endl;</pre>
32
          vector < int > cycle;
```

3.1.11 Tarjan's SCC

```
1 #include "header.h"
3 struct Tarian {
    vvi &edges;
    int V. counter = 0, C = 0;
    vi n. 1:
    vector < bool > vs;
    stack<int> st:
    Tarjan(vvi &e) : edges(e), V(e.size()), n(V,
        -1), l(V, -1), vs(V, false) {}
    void visit(int u, vi &com) {
      l[u] = n[u] = counter++;
      st.push(u);
      vs[u] = true:
      for (auto &&v : edges[u]) {
        if (n[v] == -1) visit(v, com);
        if (vs[v]) 1[u] = min(1[u], 1[v]);
16
17
      if (1[u] == n[u]) {
        while (true) {
19
          int v = st.top();
          st.pop();
21
          vs[v] = false;
          com[v] = C: // <== ACT HERE
          if (u == v) break:
24
        C++:
26
      }
27
    int find_sccs(vi &com) { // component indices
        will be stored in 'com'
      com.assign(V. -1):
      for (int u = 0; u < V; ++u)
        if (n[u] == -1) visit(u, com);
35
    // scc is a map of the original vertices of the
         graph to the vertices
    // of the SCC graph, scc_graph is its adjacency
         list.
```

```
// SCC indices and edges are stored in 'scc'
        and 'scc_graph'.
    void scc_collapse(vi &scc, vvi &scc_graph) {
      find sccs(scc):
      scc_graph.assign(C, vi());
      set < pi > rec; // recorded edges
      for (int u = 0; u < V; ++u) {
        assert(scc[u] != -1);
        for (int v : edges[u]) {
          if (scc[v] == scc[u] ||
            rec.find({scc[u], scc[v]}) != rec.end()
                ) continue:
           scc_graph[scc[u]].push_back(scc[v]);
          rec.insert({scc[u], scc[v]});
49
      }
51
    }
52
    // Function to find sources and sinks in the
        SCC graph
    // The number of edges needed to be added is
        max(sources.size(), sinks.())
    void findSourcesAndSinks(const vvi &scc graph.
        vi &sources, vi &sinks) {
      vi in_degree(C, 0), out_degree(C, 0);
      for (int u = 0; u < C; u++) {</pre>
        for (auto v : scc_graph[u]) {
          in_degree[v]++;
          out degree[u]++:
        }
62
      for (int i = 0; i < C; ++i) {</pre>
        if (in_degree[i] == 0) sources.push_back(i)
        if (out_degree[i] == 0) sinks.push_back(i);
   }
68 };
```

3.1.12 SCC edges Prints out the missing edges to make the input digraph strongly connected

```
#include "header.h"
const int N=1e5+10;
int n,a[N],cnt[N],vis[N];
vector<int> hd,tl;
int dfs(int x){
    vis[x]=1;
    if(!vis[a[x]])return vis[x]=dfs(a[x]);
    return vis[x]=x;
}
int main(){
scanf("%d",&n);
for(int i=1;i<=n;i++){
scanf("%d",&a[i]);</pre>
```

```
cnt[a[i]]++;
16
       int k=0:
       for(int i=1;i<=n;i++){</pre>
           if(!cnt[i]){
18
                k++:
19
                hd.push_back(i);
                tl.push_back(dfs(i));
21
           }
22
23
       int tk=k:
24
       for(int i=1:i<=n:i++){</pre>
           if(!vis[i]){
26
                k++:
                hd.push_back(i);
                tl.push_back(dfs(i));
29
           }
30
31
       if(k==1&&!tk)k=0;
32
       printf("%d\n".k):
       for (int i=0; i < k; i++) printf ("%d<sub>11</sub>%d\n", tl[i], hd
           [(i+1)%k]);
       return 0;
36 }
```

3.1.13 Find Bridges

```
1 #include "header.h"
2 int n; // number of nodes
3 vvi adj; // adjacency list of graph
4 vector < bool > visited;
5 vi tin. low:
6 int timer;
7 void dfs(int v, int p = -1) {
      visited[v] = true:
      tin[v] = low[v] = timer++;
      for (int to : adj[v]) {
          if (to == p) continue:
          if (visited[to]) {
              low[v] = min(low[v], tin[to]):
13
          } else {
14
              dfs(to, v);
              low[v] = min(low[v], low[to]);
              if (low[to] > tin[v])
                  IS BRIDGE(v. to):
          }
21 }
22 void find_bridges() {
      timer = 0;
      visited.assign(n, false);
      tin.assign(n, -1);
      low.assign(n, -1);
      for (int i = 0: i < n: ++i) {
```

3.1.14 Articulation points (i.e. cut off points)

```
1 #include "header.h"
2 int n: // number of nodes
3 vvi adj; // adjacency list of graph
4 vector < bool > visited;
5 vi tin, low;
6 int timer;
7 \text{ void dfs(int v. int p = -1)}
      visited[v] = true:
      tin[v] = low[v] = timer++;
      int children=0:
      for (int to : adj[v]) {
          if (to == p) continue;
          if (visited[to]) {
               low[v] = min(low[v], tin[to]);
14
15
          } else {
               dfs(to, v):
16
               low[v] = min(low[v], low[to]);
               if (low[to] >= tin[v] && p!=-1)
                   IS_CUTPOINT(v);
               ++children:
19
          }
20
21
      if(p == -1 && children > 1)
          IS CUTPOINT(v):
23
24 }
25 void find cutpoints() {
       timer = 0:
      visited.assign(n, false):
      tin.assign(n, -1);
      low.assign(n, -1);
      for (int i = 0; i < n; ++i) {</pre>
          if (!visited[i]) dfs (i);
31
33 }
```

3.1.15 Topological sort

```
#include "header.h"
int n; // number of vertices
vvi adj; // adjacency list of graph
vector<bool> visited;
vi ans;
void dfs(int v) {
visited[v] = true;
for (int u : adj[v]) {
if (!visited[u]) dfs(u);
```

```
11     ans.push_back(v);
12 }
13 void topological_sort() {
14     visited.assign(n, false);
15     ans.clear();
16     for (int i = 0; i < n; ++i) {
17         if (!visited[i]) dfs(i);
18     }
19     reverse(ans.begin(), ans.end());
20 }</pre>
```

3.1.16 Bellmann-Ford Same as Dijkstra but allows neg. edges

```
1 #include "header.h"
2 // Switch vi and vvpi to vl and vvpl if necessary
3 void bellmann_ford_extended(vvpi &e, int source,
      vi &dist, vb &cvc) {
    dist.assign(e.size(). INF):
    cyc.assign(e.size(), false); // true when u is
        in a <0 cvcle
    dist[source] = 0:
    for (int iter = 0; iter < e.size() - 1; ++iter)</pre>
      bool relax = false;
      for (int u = 0; u < e.size(); ++u)
       if (dist[u] == INF) continue:
        else for (auto &e : e[u])
          if(dist[u]+e.second < dist[e.first])</pre>
            dist[e.first] = dist[u]+e.second. relax
      if(!relax) break:
15
    bool ch = true;
    while (ch) {
                         // keep going untill no
        more changes
      ch = false:
                         // set dist to -INF when in
           cvcle
      for (int u = 0; u < e.size(); ++u)</pre>
        if (dist[u] == INF) continue:
        else for (auto &e : e[u])
          if (dist[e.first] > dist[u] + e.second
            && !cvc[e.first]) {
            dist[e.first] = -INF;
            ch = true; //return true for cycle
                detection only
            cyc[e.first] = true;
27
   }
28
29 }
```

3.1.17 Ford-Fulkerson Basic Max. flow

```
1 #include "header.h"
2 #define V 6 // Num. of vertices in given graph
4 /* Returns true if there is a path from source 's
      ' to sink
5 't' in residual graph. Also fills parent[] to
      store the
6 path */
7 bool bfs(int rGraph[V][V], int s, int t, int
      parent[]) {
   bool visited[V];
    memset(visited, 0, sizeof(visited));
    queue < int > q;
    a.push(s):
    visited[s] = true;
    parent[s] = -1;
    // Standard BFS Loop
    while (!a.emptv()) {
      int u = q.front();
      q.pop();
19
      for (int v = 0; v < V; v++) {
20
        if (visited[v] == false && rGraph[u][v] >
            0) {
          if (v == t) {
            parent[v] = u;
23
            return true;
          q.push(v);
          parent[v] = u;
          visited[v] = true:
      }
    return false;
33 }
34
35 // Returns the maximum flow from s to t in the
      given graph
36 int fordFulkerson(int graph[V][V], int s, int t)
    int u. v:
    int rGraph[V]
    for (u = 0: u < V: u++)
     for (v = 0; v < V; v++)
        rGraph[u][v] = graph[u][v];
    int parent[V]; // This array is filled by BFS
        and to
          // store path
    int max_flow = 0; // There is no flow initially
    while (bfs(rGraph, s, t, parent)) {
      int path_flow = INT_MAX;
```

```
for (v = t; v != s; v = parent[v]) {
    u = parent[v];
    path_flow = min(path_flow, rGraph[u][v]);
}

for (v = t; v != s; v = parent[v]) {
    u = parent[v];
    rGraph[u][v] -= path_flow;
    rGraph[v][u] += path_flow;
}

max_flow += path_flow;
}

return max_flow;
}
```

3.1.18 Dinic max flow $O(V^2E)$, O(Ef)

```
2 using F = 11; using W = 11; // types for flow and
       weight/cost
3 struct Sf
                             // neighbour
      const int v:
      const int r:
                       // index of the reverse edge
                       // current flow
                      // capacity
      const F cap;
      const W cost;
                      // unit cost
      S(int v. int ri. F c. W cost = 0):
          v(v), r(ri), f(0), cap(c), cost(cost) {}
      inline F res() const { return cap - f; }
12 }:
13 struct FlowGraph : vector < vector < S >> {
      FlowGraph(size t n) : vector < vector < S >> (n) {}
      void add_edge(int u, int v, F c, W cost = 0){
           auto &t = *this:
          t[u].emplace_back(v, t[v].size(), c, cost
          t[v].emplace_back(u, t[u].size()-1, c, -
              cost):
18
      void add arc(int u. int v. F c. W cost = 0){
          auto &t = *this:
          t[u].emplace_back(v, t[v].size(), c, cost
          t[v].emplace_back(u, t[u].size()-1, 0, -
              cost):
      void clear() { for (auto &E : *this) for (
          auto &e : E) e.f = OLL: }
24 };
25 struct Dinic{
      FlowGraph & edges: int V.s.t:
      vi 1; vector < vector < S > :: iterator > its; //
          levels and iterators
      Dinic(FlowGraph &edges, int s, int t) :
```

```
edges(edges), V(edges.size()), s(s), t(t)
               , 1(V,-1), its(V) {}
      ll augment(int u. F c) { // we reuse the same
           iterators
           if (u == t) return c; ll r = OLL;
           for(auto &i = its[u]; i != edges[u].end()
              : i++){
               auto &e = *i;
              if (e.res() && l[u] < l[e.v]) {
                   auto d = augment(e.v, min(c, e.
                       res())):
                  if (d > 0) { e.f += d; edges[e.v
                       ][e.r].f -= d; c -= d:
                       r += d: if (!c) break: }
           return r;
      }
      11 run() {
          11 \text{ flow} = 0. \text{ f}:
           while(true) {
               fill(1.begin(), 1.end(),-1); 1[s]=0;
                   // recalculate the lavers
               queue < int > q; q.push(s);
               while(!q.empty()){
                   auto u = q.front(); q.pop(); its[
                       u] = edges[u].begin();
                   for(auto &&e : edges[u]) if(e.res
                       () && l[e.v]<0)
                       l[e.v] = l[u]+1, q.push(e.v);
               }
               if (1[t] < 0) return flow;</pre>
               while ((f = augment(s, INF)) > 0)
                   flow += f;
54 }:
```

3.1.19 Edmonds-Karp Max flow $O(VE^2)$

```
fill(all(par), -1);
      par[source] = 0;
      int ptr = 1;
      a[0] = source:
17
      rep(i,0,ptr) {
       int x = q[i];
        for (auto e : graph[x]) {
          if (par[e.first] == -1 \&\& e.second > 0) {
            par[e.first] = x;
            q[ptr++] = e.first;
            if (e.first == sink) goto out;
        }
28
      return flow;
29 out:
      T inc = numeric_limits <T>::max();
      for (int y = sink; y != source; y = par[y])
        inc = min(inc, graph[par[y]][y]);
33
34
      flow += inc:
      for (int y = sink; y != source; y = par[y]) {
        int p = par[v];
        if ((graph[p][y] -= inc) <= 0) graph[p].</pre>
            erase(v);
        graph[y][p] += inc;
    }
41 }
```

3.2 Dynamic Programming

3.2.1 Longest Incr. Subseq.

```
1 #include "header.h"
2 template < class T>
3 vector<T> index_path_lis(vector<T>& nums) {
    int n = nums.size();
    vector <T> sub:
      vector < int > subIndex;
    vector <T> path(n, -1);
    for (int i = 0; i < n; ++i) {
        if (sub.empty() || sub[sub.size() - 1] <</pre>
            nums[i]) {
      path[i] = sub.empty() ? -1 : subIndex[sub.
          size() - 1];
      sub.push_back(nums[i]);
      subIndex.push_back(i);
12
       } else {
13
      int idx = lower_bound(sub.begin(), sub.end(),
           nums[i]) - sub.begin();
      path[i] = idx == 0 ? -1 : subIndex[idx - 1];
15
      sub[idx] = nums[i]:
```

```
subIndex[idx] = i;
        }
19
    vector <T> ans:
    int t = subIndex[subIndex.size() - 1];
    while (t != -1) {
        ans.push_back(t);
        t = path[t];
    reverse(ans.begin(), ans.end());
    return ans:
28 }
29 // Length only
30 template < class T>
31 int length_lis(vector<T> &a) {
    set <T> st;
    typename set<T>::iterator it;
    for (int i = 0; i < a.size(); ++i) {</pre>
      it = st.lower_bound(a[i]);
      if (it != st.end()) st.erase(it);
      st.insert(a[i]);
   }
    return st.size();
```

3.2.2 0-1 Knapsack

```
1 #include "header.h"
2 // given a number of coins, calculate all
      possible distinct sums
3 int main() {
   vi coins(n); // all possible coins to use
   int sum = 0;
                    // sum of the coins
   vi dp(sum + 1, 0);
                              // dp[x] = 1 if sum
        x can be made
   dp[0] = 1;
                               // sum 0 can be
       made
  for (int c = 0; c < n; ++c)
                                       // first
       iteration: sums with first
     for (int x = sum: x \ge 0: --x)
                                         // coin.
         next first 2 coins etc
       if (dp[x]) dp[x + coins[c]] = 1; // if sum
            x valid, x+c valid
```

3.2.3 Coin change Number of coins required to achieve a given value

```
1 #include "header.h"
2 // Returns total distinct ways to make sum using
    n coins of
3 // different denominations
```

3.3 Trees

3.3.1 Tree diameter

```
1 #include "header.h"
2 \text{ const int } mxN = 2e5 + 5;
3 int n, d[mxN]; // distance array
4 vi adj[mxN]; // tree adjacency list
5 void dfs(int s, int e) {
6 d[s] = 1 + d[e]:
                      // recursively calculate
        the distance from the starting node to each
for (auto u : adj[s]) { // for each adjacent
      if (u != e) dfs(u, s); // don't move
          backwards in the tree
9 }
11 int main() {
12 // read input, create adj list
    dfs(0, -1);
                                  // first dfs call
         to find farthest node from arbitrary node
    dfs(distance(d, max_element(d, d + n)), -1);
        // second dfs call to find farthest node
        from that one
   cout << *max element(d, d + n) - 1 << '\n'; //
         distance from second node to farthest is
        the diameter
```

3.3.2 Tree Node Count

3.4 Numerical

3.4.1 Template (for this section)

```
#include <bits/stdc++.h>
using namespace std;
#define rep(i, a, b) for(int i = a; i < (b); ++i)
#define all(x) begin(x), end(x)
#define sz(x) (int)(x).size()
typedef long long ll;
typedef pair<int, int> pii;
typedef vector<int> vi;
```

3.4.2 Polynomial

```
1 #include "template.cpp"
3 struct Poly {
    vector <double > a:
    double operator()(double x) const {
      double val = 0;
      for (int i = sz(a); i--;) (val *= x) += a[i];
      return val;
   }
    void diff() {
      rep(i,1,sz(a)) a[i-1] = i*a[i];
      a.pop_back();
    void divroot(double x0) {
      double b = a.back(), c: a.back() = 0:
15
      for(int i=sz(a)-1; i--;) c = a[i], a[i] = a[i
          +1]*x0+b, b=c;
      a.pop_back();
18 }
19 };
```

3.4.3 Poly Roots

```
1 /**
2 * Description: Finds the real roots to a
       polynomial.
3 * Usage: polyRoots({{2,-3,1}},-1e9,1e9) // solve
        x^2-3x+2 = 0
4 * Time: O(n^2 \log(1/\epsilon))
6 #include "Polynomial.h"
7 #include "template.cpp"
9 vector < double > polyRoots(Poly p, double xmin,
      double xmax) {
   if (sz(p.a) == 2) { return {-p.a[0]/p.a[1]}; }
    vector < double > ret;
    Poly der = p;
    der.diff();
    auto dr = polyRoots(der, xmin, xmax);
    dr.push_back(xmin-1);
    dr.push back(xmax+1);
    sort(all(dr));
    rep(i,0,sz(dr)-1) {
      double 1 = dr[i], h = dr[i+1];
      bool sign = p(1) > 0;
      if (sign ^(p(h) > 0)) {
        rep(it,0,60) { // while (h - 1 > 1e-8)
          double m = (1 + h) / 2, f = p(m);
23
          if ((f <= 0) ^ sign) l = m;</pre>
          else h = m;
26
        ret.push_back((1 + h) / 2);
28
    }
29
31 }
```

3.4.4 Golden Section Search

3.4.5 Hill Climbing

```
1 /**
2 * Description: Poor man's optimization for
      unimodal functions.
3 */
4 #include "template.cpp"
6 typedef array < double, 2> P;
8 template < class F > pair < double, P > hillClimb(P
      start, F f) {
   pair < double , P > cur(f(start), start);
   for (double jmp = 1e9; jmp > 1e-20; jmp /= 2) {
     rep(j,0,100) rep(dx,-1,2) rep(dy,-1,2) {
       P p = cur.second;
        p[0] += dx*jmp;
       p[1] += dv*imp;
        cur = min(cur, make_pair(f(p), p));
   }
   return cur;
```

3.4.6 Integration

3.4.7 Integration Adaptive

```
1 /**
2 * Description: Fast integration using an
       adaptive Simpson's rule.
    double sphereVolume = quad(-1, 1, [](double x)
    return quad(-1, 1, [\&](double y) {
    return quad(-1, 1, [\&](double z) {
    return x*x + y*y + z*z < 1; });});});
  * Status: mostly untested
10 #include "template.cpp"
12 typedef double d;
13 #define S(a,b) (f(a) + 4*f((a+b) / 2) + f(b)) * (
      b-a) / 6
15 template <class F>
16 d rec(F& f, d a, d b, d eps, d S) {
  dc = (a + b) / 2:
    d S1 = S(a, c), S2 = S(c, b), T = S1 + S2;
   if (abs(T - S) \le 15 * eps | | b - a < 1e-10)
     return T + (T - S) / 15;
    return rec(f, a, c, eps / 2, S1) + rec(f, c, b,
         eps / 2, S2):
22 }
23 template < class F>
24 d quad(d a, d b, F f, d eps = 1e-8) {
    return rec(f, a, b, eps, S(a, b));
26 }
```

3.5 Num. Th. / Comb.

3.5.1 Basic stuff

```
1 #include "header.h"
2 ll gcd(ll a, ll b) { while (b) { a %= b; swap(a, b); } return a; }
```

```
3 11 1cm(11 a, 11 b) { return (a / gcd(a, b)) * b;
4 ll mod(ll a, ll b) { return ((a % b) + b) % b; }
5 // Finds x, v s.t. ax + bv = d = gcd(a, b).
6 void extended_euclid(ll a, ll b, ll &x, ll &y, ll
       &d) {
    11 xx = y = 0;
    11 yy = x = 1;
    while (b) {
      11 t = b; b = a % b; a = t;
      t = xx; xx = x - q * xx; x = t;
      t = yy; yy = y - q * yy; y = t;
13
16 }
17 //  solves ab = 1 (mod n), -1 on failure
18 ll mod_inverse(ll a, ll n) {
   11 x, y, d;
    extended_euclid(a, n, x, y, d);
    return (d > 1 ? -1 : mod(x, n));
22 }
23 // All modular inverses of [1..n] mod P in O(n)
24 vi inverses(ll n, ll P) {
    vi I(n+1, 1LL);
   for (11 i = 2; i <= n; ++i)
      I[i] = mod(-(P/i) * I[P\%i], P):
    return I;
30 // (a*b)%m
31 ll mulmod(ll a, ll b, ll m){
   11 x = 0, y=a\%m;
   while(b>0){
      if(b\&1) x = (x+v)\%m:
      y = (2*y)\%m, b /= 2;
    }
    return x % m:
39 // Finds b^e % m in O(lg n) time, ensure that b <
       m to avoid overflow!
40 ll powmod(ll b, ll e, ll m) {
   11 p = e<2 ? 1 : powmod((b*b)\%m,e/2,m);
    return e&1 ? p*b%m : p;
43 }
44 // Solve ax + by = c, returns false on failure.
45 bool linear_diophantine(ll a, ll b, ll c, ll &x,
      11 &v) {
    11 d = gcd(a, b);
    if (c % d) {
      return false:
      x = c / d * mod_inverse(a / d, b / d);
      y = (c - a * x) / b;
      return true;
```

```
54 }
56 // Description: Tonelli-Shanks algorithm for
      modular square roots. Finds x s.t. x^2 = a
       \propto p$ ($-x$ gives the other solution). 0
      (\log^2 p) worst case, O(\log p) for most $p$
57 ll sqrtmod(ll a, ll p) {
    a \% = p: if (a < 0) a += p:
    if (a == 0) return 0;
    assert(powmod(a, (p-1)/2, p) == 1); // else no
        solution
    if (p \% 4 == 3) return powmod(a, (p+1)/4, p);
    // a^{(n+3)/8} or 2^{(n+3)/8} * 2^{(n-1)/4} works if
        p % 8 == 5
    11 s = p - 1, n = 2;
    int r = 0. m:
    while (s \% 2 == 0)
      ++r, s /= 2;
    /// find a non-square mod p
    while (powmod(n, (p - 1) / 2, p) != p - 1) ++n;
    11 x = powmod(a, (s + 1) / 2, p);
    ll b = powmod(a, s, p), g = powmod(n, s, p);
    for (;; r = m) {
      11 t = b:
      for (m = 0; m < r \&\& t != 1; ++m)
       t = t * t % p;
      if (m == 0) return x:
      11 \text{ gs} = powmod(g, 1LL << (r - m - 1), p);}
      g = gs * gs % p;
      x = x * gs % p;
      b = b * g % p;
81 }
```

3.5.2 Mod. exponentiation Or use pow() in python

```
#include "header.h"
2 ll mod_pow(ll base, ll exp, ll mod) {
3    if (mod == 1) return 0;
4    if (exp == 0) return 1;
5    if (exp == 1) return base;
6
7    ll res = 1;
8    base %= mod;
9    while (exp) {
10        if (exp % 2 == 1) res = (res * base) % mod;
11        exp >>= 1;
12        base = (base * base) % mod;
13    }
14
15    return res % mod;
16 }
```

3.5.3 GCD Or math.gcd in python, std::gcd in C++

```
#include "header.h"
2 ll gcd(ll a, ll b) {
3    if (a == 0) return b;
4    return gcd(b % a, a);
5 }
```

3.5.4 Sieve of Eratosthenes

3.5.5 Fibonacci % prime

3.5.6 nCk % prime

```
14 }
15 return ans;
16 }
```

3.5.7 Chin. rem. th.

```
1 #include "header.h"
2 #include "elementary.cpp"
_3 // Solves x = a1 mod m1, x = a2 mod m2, x is
      unique modulo lcm(m1, m2).
4 // Returns {0, -1} on failure, {x, lcm(m1, m2)}
      otherwise.
5 pair<11, 11> crt(11 a1, 11 m1, 11 a2, 11 m2) {
6 ll s, t, d;
    extended_euclid(m1, m2, s, t, d);
    if (a1 % d != a2 % d) return {0, -1}:
    return {mod(s*a2 %m2 * m1 + t*a1 %m1 * m2, m1 *
         m2) / d, m1 / d * m2};
10 }
12 // Solves x = ai mod mi. x is unique modulo lcm
13 // Returns {0, -1} on failure, {x, lcm mi}
      otherwise.
14 pair <11, 11> crt(vector <11> &a, vector <11> &m) {
    pair<11, 11> res = {a[0], m[0]};
    for (ull i = 1; i < a.size(); ++i) {</pre>
      res = crt(res.first, res.second, mod(a[i], m[
          il), m[il):
      if (res.second == -1) break;
   return res:
```

3.5.8 Derangements Permutations of a set such that none of the elements appear in their original position:

$$!n = (n-1)(!(n-1)+!(n-2)) = \left[\frac{n!}{e}\right]$$

$$!n = 1 - e^{-1}, \ n \to \infty$$
(2)

3.6 Strings

3.6.1 Z alg. KMP alternative

```
#include "../header.h"
void Z_algorithm(const string &s, vi &Z) {
   Z.assign(s.length(), -1);
   int L = 0, R = 0, n = s.length();
   for (int i = 1; i < n; ++i) {</pre>
```

```
6    if (i > R) {
7         L = R = i;
8         while (R < n && s[R - L] == s[R]) R++;
9         Z[i] = R - L; R--;
10    } else if (Z[i - L] >= R - i + 1) {
11         L = i;
12         while (R < n && s[R - L] == s[R]) R++;
13         Z[i] = R - L; R--;
14     } else Z[i] = Z[i - L];
15    }
16 }</pre>
```

3.6.2 KMP

```
1 #include "header.h"
void compute_prefix_function(string &w, vi &
      prefix) {
   prefix.assign(w.length(), 0);
    int k = prefix[0] = -1;
    for(int i = 1; i < w.length(); ++i) {</pre>
      while (k >= 0 \&\& w[k + 1] != w[i]) k = prefix[
      if(w[k + 1] == w[i]) k++;
      prefix[i] = k;
11 }
12 void knuth_morris_pratt(string &s, string &w) {
    int a = -1:
    vi prefix;
    compute_prefix_function(w, prefix);
    for(int i = 0; i < s.length(); ++i) {</pre>
      while (q >= 0 \&\& w[q + 1] != s[i]) q = prefix[
      if(w[q + 1] == s[i]) q++;
      if(q + 1 == w.length()) {
        // Match at position (i - w.length() + 1)
        q = prefix[q];
23 }
```

3.6.3 Aho-Corasick Also can be used as Knuth-Morris-Pratt algorithm

```
#include "header.h"

map < char, int > cti;
int cti_size;
template < int ALPHABET_SIZE, int (*mp)(char) >
struct AC_FSM {
struct Node {
```

matches.assign(words.size(), vi());

```
int child[ALPHABET_SIZE], failure = 0,
                                                                int state = 0. ss = 0:
                                                                for (int i = 0; i < sentence.length(); ++i,</pre>
          match_par = -1;
      vi match:
                                                                    ss = state) {
                                                                  while (a[ss].child[mp(sentence[i])] == -1)
      Node() { for (int i = 0: i < ALPHABET SIZE:
          ++i) child[i] = -1; }
                                                                    ss = a[ss].failure;
                                                                  state = a[state].child[mp(sentence[i])]
    };
    vector < Node > a;
                                                                       = a[ss].child[mp(sentence[i])];
                                                                  for (ss = state; ss != -1; ss = a[ss].
    vector < string > & words;
    AC FSM(vector<string> &words) : words(words) {
                                                                      match par)
      a.push_back(Node());
                                                                    for (int w : a[ss].match)
      construct automaton():
                                                                       matches[w].push back(i + 1 - words[w].
                                                          63
16
                                                                           length());
    void construct_automaton() {
18
      for (int w = 0, n = 0; w < words.size(); ++w,</pre>
                                                             }
                                                          65
           n = 0) {
        for (int i = 0; i < words[w].size(); ++i) {</pre>
                                                         67 int char_to_int(char c) {
           if (a[n].child[mp(words[w][i])] == -1) {
                                                             return cti[c]:
21
             a[n].child[mp(words[w][i])] = a.size();
22
             a.push_back(Node());
                                                         70 int main() {
                                                              11 n:
          n = a[n].child[mp(words[w][i])];
                                                              string line;
25
                                                              while(getline(cin, line)) {
26
        a[n].match.push_back(w);
                                                                stringstream ss(line);
                                                                ss >> n;
28
      queue < int > q;
29
      for (int k = 0; k < ALPHABET_SIZE; ++k) {</pre>
                                                         77
                                                                vector < string > patterns(n);
30
        if (a[0].child[k] == -1) a[0].child[k] = 0;
                                                                for (auto& p: patterns) getline(cin, p);
         else if (a[0].child[k] > 0) {
          a[a[0].child[k]].failure = 0;
                                                                string text;
          q.push(a[0].child[k]);
                                                                getline(cin, text);
34
35
                                                                cti = {}, cti_size = 0;
36
      while (!q.empty()) {
                                                                for (auto c: text) {
        int r = q.front(); q.pop();
                                                                  if (not in(c, cti)) {
        for (int k = 0, arck; k < ALPHABET_SIZE; ++</pre>
                                                                    cti[c] = cti size++:
                                                         86
           if ((arck = a[r].child[k]) != -1) {
                                                                }
                                                          88
                                                                for (auto& p: patterns) {
             q.push(arck);
             int v = a[r].failure;
                                                                  for (auto c: p) {
42
             while (a[v].child[k] == -1) v = a[v].
                                                                    if (not in(c, cti)) {
                                                         91
                 failure:
                                                                      cti[c] = cti_size++;
             a[arck].failure = a[v].child[k];
                                                                    }
             a[arck].match_par = a[v].child[k];
                                                                  }
             while (a[arck].match_par != -1
                                                                }
                 && a[a[arck].match_par].match.empty
                                                                vvi matches:
               a[arck].match_par = a[a[arck].
                                                                AC_FSM <128+1, char_to_int > ac_fms(patterns);
                   match_par].match_par;
                                                                ac_fms.aho_corasick(text, matches);
          }
                                                                for (auto& x: matches) cout << x << endl;</pre>
                                                         101
                                                         102
51
52
                                                         103 }
    void aho_corasick(string &sentence, vvi &
         matches) {
```

```
3.6.4 Long. palin. subs Manacher - O(n)
```

```
1 #include "header.h"
void manacher(string &s, vi &pal) {
    int n = s.length(), i = 1, 1, r;
    pal.assign(2 * n + 1, 0);
    while (i < 2 * n + 1)  {
      if ((i&1) && pal[i] == 0) pal[i] = 1;
      l = i / 2 - pal[i] / 2; r = (i-1) / 2 + pal[i]
          1 / 2:
      while (1 - 1 >= 0 \&\& r + 1 < n \&\& s[1 - 1] ==
           s[r + 1]
        --1, ++r, pal[i] += 2;
11
      for (1 = i - 1, r = i + 1; 1 >= 0 && r < 2 *
12
          n + 1; --1, ++r) {
        if (1 <= i - pal[i]) break;</pre>
        if (1 / 2 - pal[1] / 2 > i / 2 - pal[i] /
            2)
          pal[r] = pal[1];
        else { if (1 \ge 0)
            pal[r] = min(pal[1], i + pal[i] - r);
19
      }
      i = r:
22 } }
```

3.7 Geometry

3.7.1 essentials.cpp

```
1 #include "../header.h"
2 using C = ld; // could be long long or long
      double
3 constexpr C EPS = 1e-10; // change to 0 for C=11
4 struct P { // may also be used as a 2D vector
P(C x = 0, C y = 0) : x(x), y(y) {}
7 P operator+ (const P &p) const { return {x + p.
        x, y + p.y; }
   P operator - (const P &p) const { return {x - p.
        x, y - p.y; }
   P operator* (C c) const { return {x * c, y * c
   P operator/ (C c) const { return {x / c, y / c
    C operator* (const P &p) const { return x*p.x +
    C operator (const P &p) const { return x*p.y -
         p.x*y; }
P perp() const { return P{y, -x}; }
C lensq() const { return x*x + y*y; }
15 ld len() const { return sqrt((ld)lensq()); }
```

size t n:

```
static ld dist(const P &p1, const P &p2) {
      return (p1-p2).len(); }
    bool operator == (const P &r) const {
      return ((*this)-r).lensq() <= EPS*EPS: }</pre>
21 C det(P p1, P p2) { return p1^p2; }
22 C det(P p1, P p2, P o) { return det(p1-o, p2-o);
    det(const vector <P> &ps) {
    C sum = 0; P prev = ps.back();
    for(auto &p : ps) sum += det(p, prev), prev = p
    return sum;
28 // Careful with division by two and C=11
29 C area(P p1, P p2, P p3) { return abs(det(p1, p2,
       p3))/C(2); }
30 C area(const vector <P> &poly) { return abs(det(
      poly))/C(2); }
31 int sign(C c) { return (c > C(0)) - (c < C(0)); }
32 int ccw(P p1, P p2, P o) { return sign(det(p1, p2
      . o)): }
_{34} // Only well defined for C = 1d.
35 P unit(const P &p) { return p / p.len(); }
36 P rotate(P p, ld a) { return P{p.x*cos(a)-p.y*sin
      (a), p.x*sin(a)+p.y*cos(a)}; }
```

3.7.2 Two segs. itersec.

```
#include "header.h"
#include "essentials.cpp"
bool intersect(P a1, P a2, P b1, P b2) {

if (max(a1.x, a2.x) < min(b1.x, b2.x)) return
    false;

if (max(b1.x, b2.x) < min(a1.x, a2.x)) return
    false;

if (max(a1.y, a2.y) < min(b1.y, b2.y)) return
    false;

if (max(b1.y, b2.y) < min(a1.y, a2.y)) return
    false;

bool 11 = ccw(a2, b1, a1) * ccw(a2, b2, a1) <=
    0;

bool 12 = ccw(b2, a1, b1) * ccw(b2, a2, b1) <=
    0;

return 11 && 12;

11 }</pre>
```

3.7.3 Convex Hull

```
1 #include "header.h"
2 #include "essentials.cpp"
3 struct ConvexHull { // O(n lg n) monotone chain.
```

```
vector < size_t > h, c; // Indices of the hull
        are in 'h', ccw.
    const vector <P> &p:
    ConvexHull(const vector <P> &_p) : n(_p.size()),
         c(n), p(_p) {
      std::iota(c.begin(), c.end(), 0);
      std::sort(c.begin(), c.end(), [this](size_t 1
          , size_t r) -> bool { return p[1].x != p[
          r].x ? p[1].x < p[r].x : p[1].y < p[r].y;
      c.erase(std::unique(c.begin(), c.end(), [this
10
          ](size_t l, size_t r) { return p[l] == p[
          r]; }), c.end());
      for (size_t s = 1, r = 0; r < 2; ++r, s = h.
          size()) {
        for (size t i : c) {
12
          while (h.size() > s && ccw(p[h.end()
13
              [-2], p[h.end()[-1]], p[i]) <= 0)
            h.pop_back();
          h.push_back(i);
15
16
        reverse(c.begin(), c.end());
      if (h.size() > 1) h.pop_back();
20
    size_t size() const { return h.size(); }
21
    template <class T. void U(const P &. const P &.
         const P &. T &)>
    void rotating_calipers(T &ans) {
      if (size() <= 2)
        U(p[h[0]], p[h.back()], p[h.back()], ans);
        for (size_t i = 0, j = 1, s = size(); i < 2</pre>
27
             * s: ++i) {
          while (det(p[h[(i + 1) % s]] - p[h[i % s
              ]], p[h[(j + 1) \% s]] - p[h[j]]) >=
29
            i = (i + 1) \% s;
          U(p[h[i % s]], p[h[(i + 1) % s]], p[h[j
30
              ]], ans);
        }
31
   }
32
34 // Example: furthest pair of points. Now set ans
      = OLL and call
35 // ConvexHull(pts).rotating_calipers<11, update>(
36 void update(const P &p1, const P &p2, const P &o,
       11 &ans) {
    ans = max(ans, (11)max((p1 - o).lensq(), (p2 -
        o).lensq())):
39 int main() {
   ios::sync_with_stdio(false); // do not use
```

```
cout + printf
    cin.tie(NULL);
    int n:
    cin >> n;
    while (n) {
      vector <P> ps;
           int x, y;
       for (int i = 0; i < n; i++) {</pre>
               cin >> x >> y;
               ps.push_back({x, y});
52
           ConvexHull ch(ps);
           cout << ch.h.size() << endl;</pre>
           for(auto& p: ch.h) {
               cout << ps[p].x << "" << ps[p].y <<
           }
       cin >> n;
    }
    return 0;
```

3.8 Other Algorithms

3.8.1 2-sat

```
1 #include "../header.h"
2 #include "../Graphs/tarjan.cpp"
3 struct TwoSAT {
    vvi imp; // implication graph
    Tarjan tj;
    TwoSAT(int _n) : n(_n), imp(2 * _n, vi()), tj(
        imp) { }
    // Only copy the needed functions:
    void add_implies(int c1, bool v1, int c2, bool
      int u = 2 * c1 + (v1 ? 1 : 0),
        v = 2 * c2 + (v2 ? 1 : 0);
      imp[u].push_back(v); // u => v
      imp[v^1].push_back(u^1); // -v => -u
    void add_equivalence(int c1, bool v1, int c2,
        bool v2) {
      add_implies(c1, v1, c2, v2);
      add_implies(c2, v2, c1, v1);
20
    void add_or(int c1, bool v1, int c2, bool v2) {
      add_implies(c1, !v1, c2, v2);
```

```
void add_and(int c1, bool v1, int c2, bool v2)
      add true(c1, v1): add true(c2, v2):
    void add_xor(int c1, bool v1, int c2, bool v2)
      add_or(c1, v1, c2, v2);
      add or(c1, !v1, c2, !v2):
    void add true(int c1. bool v1) {
      add implies(c1, !v1, c1, v1):
33
    // on true: a contains an assignment.
    // on false: no assignment exists.
    bool solve(vb &a) {
      vi com:
      tj.find_sccs(com);
      for (int i = 0: i < n: ++i)
        if (com[2 * i] == com[2 * i + 1])
41
          return false:
      vvi bycom(com.size());
44
      for (int i = 0; i < 2 * n; ++i)
        bycom[com[i]].push_back(i);
46
      a.assign(n. false):
      vb vis(n, false);
49
      for(auto &&component : bycom){
       for (int u : component) {
          if (vis[u / 2]) continue;
          vis[u / 2] = true:
          a[u / 2] = (u \% 2 == 1);
      return true;
59 };
```

3.8.2 Matrix Solve

```
for(int c = 0; c < rows && r < rows; c++) {</pre>
      int p = r;
11
      for(int i=r+1: i<rows: i++) if(abs(m[i][c]) >
           abs(m[p][c])) p=i:
      if(abs(m[p][c]) < EPS){ det = 0; continue; }</pre>
      swap(m[p], m[r]); det = -det;
14
      T s = 1.0 / m[r][c], t; det *= m[r][c];
      REP(j,C) m[r][j] *= s; // make leading
16
           term in row 1
      REP(i,rows) if (i!=r){ t = m[i][c]; REP(j,C)
          m[i][j] -= t*m[r][j]; }
18
      ++r;
    return det:
21 }
22 bool error, inconst; // error => multiple or
      inconsistent
23 template <int R, int C> // Mx = a; M:R*R, v:R*C =>
24 M<R,C> solve(const M<R,R> &m, const M<R,C> &a,
      int rows){
    M < R.R+C > a:
    REP(r,rows){
      REP(c,rows) q[r][c] = m[r][c];
      REP(c,C) q[r][R+c] = a[r][c];
    ReducedRowEchelonForm <R,R+C>(q,rows);
    M<R.C> sol: error = false. inconst = false:
    REP(c,C) for(auto j = rows-1; j >= 0; --j){
     T t=0: bool allzero=true:
      for (auto k = j+1; k < rows; ++k)
        t += q[j][k]*sol[k][c], allzero &= abs(q[j])
            l[k]) < EPS:
      if(abs(q[i][i]) < EPS)</pre>
        error = true, inconst |= allzero && abs(q[j
            l[R+c]) > EPS:
      else sol[j][c] = (q[j][R+c] - t) / q[j][j];
          // usually q[j][j]=1
39
    return sol;
```

3.8.3 Matrix Exp.

3.8.4 Finite field For FFT

```
1 #include "header.h"
2 #include "../Number Theory/elementary.cpp"
3 template<11 p,11 w> // prime, primitive root
4 struct Field { using T = Field; ll x; Field(ll x
      =0): x\{x\}\{
5 T operator+(T r) const { return {(x+r.x)%p}; }
    T operator - (T r) const { return {(x-r.x+p)%p};
    T operator*(T r) const { return {(x*r.x)%p}; }
    T operator/(T r) const { return (*this)*r.inv()
    T inv() const { return {mod_inverse(x,p)}; }
    static T root(ll k) { assert( (p-1)%k==0 );
        // (p-1)%k == 0?
      auto r = powmod(w,(p-1)/abs(k),p);
                                               // k-
          th root of unity
      return k>=0 ? T{r} : T{r}.inv():
    bool zero() const { return x == OLL; }
16 using F1 = Field<1004535809,3 >;
17 using F2 = Field<1107296257,10>; // 1<<30 + 1<<25
18 using F3 = Field < 2281701377,3 >; // 1 < < 31 + 1 < < 27
```

3.8.5 Complex field For FFR

```
#include "header.h"
const double m_pi = M_PIf64x;
struct Complex { using T = Complex; double u,v;
Complex(double u=0, double v=0) : u{u}, v{v} {};
T operator+(T r) const { return {u+r.u, v+r.v};
}
T operator-(T r) const { return {u-r.u, v-r.v};
}
```

```
T operator*(T r) const { return {u*r.u - v*r.v,
         u*r.v + v*r.u}; }
   T operator/(T r) const {
      auto norm = r.u*r.u+r.v*r.v;
      return {(u*r.u + v*r.v)/norm, (v*r.u - u*r.v)
11
   T operator*(double r) const { return T{u*r, v*r
   T operator/(double r) const { return T{u/r, v/r
       }: }
   T inv() const { return T{1,0}/ *this; }
   T conj() const { return T{u, -v}; }
   static T root(ll k){ return {cos(2*m_pi/k), sin
        (2*m_pi/k); }
bool zero() const { return max(abs(u), abs(v))
        < 1e-6: }
18 };
```

3.8.6 FFT

1 #include "header.h"

```
2 #include "complex_field.cpp"
3 #include "fin_field.cpp"
4 void brinc(int &x, int k) {
5 int i = k - 1, s = 1 << i;</pre>
   if ((x & s) != s) {
      --i; s >>= 1;
      while (i >= 0 && ((x & s) == s))
       x = x &^{\sim} s, --i, s >>= 1;
      if (i >= 0) x |= s:
using T = Complex; // using T=F1,F2,F3
15 vector <T> roots;
16 void root_cache(int N) {
    if (N == (int)roots.size()) return:
    roots.assign(N, T{0});
    for (int i = 0: i < N: ++i)
      roots[i] = ((i\&-i) == i)
        ? T\{\cos(2.0*m_pi*i/N), \sin(2.0*m_pi*i/N)\}
        : roots[i&-i] * roots[i-(i&-i)]:
24 void fft(vector<T> &A. int p. bool inv = false) {
   int N = 1 << p:
   for(int i = 0, r = 0; i < N; ++i, brinc(r, p))
      if (i < r) swap(A[i], A[r]);</pre>
28 // Uncomment to precompute roots (for T=Complex)
      . Slower but more precise.
29 // root cache(N):
            , sh=p-1
  for (int m = 2; m <= N; m <<= 1) {
      T w, w m = T::root(inv ? -m : m):
```

```
for (int k = 0: k < N: k += m) {
        w = T\{1\};
        for (int i = 0: i < m/2: ++i) {
35
            T w = (!inv ? roots[j << sh] : roots[j <<
      sh].conj());
          T t = w * A[k + j + m/2];
          A[k + j + m/2] = A[k + j] - t;
          A[k + j] = A[k + j] + t;
          w = w * w m:
42
      }
43
    if(inv){ T inverse = T(N).inv(); for(auto &x :
        A) x = x*inverse:
45 }
_{
m 46} // convolution leaves A and B in frequency domain
47 // C may be equal to A or B for in-place
      convolution
48 void convolution(vector <T> &A, vector <T> &B,
      vector <T> &C) {
    int s = A.size() + B.size() - 1:
    int q = 32 - __builtin_clz(s-1), N=1<<q; //</pre>
        fails if s=1
    A.resize(N,{}); B.resize(N,{}); C.resize(N,{});
    fft(A, q, false); fft(B, q, false);
    for (int i = 0; i < N; ++i) C[i] = A[i] * B[i];</pre>
    fft(C, q, true): C.resize(s):
56 void square_inplace(vector <T > &A) {
    int s = 2*A.size()-1, q = 32 - _builtin_clz(s
        -1), N=1<<q;
    A.resize(N,{}); fft(A, q, false);
    for (auto &x : A) x = x*x;
    fft(A, q, true); A.resize(s);
```

3.8.7 Polyn. inv. div.

```
vector<T> As(4 << k, T(0)), Ais(4 << k, T(0));
      copy_into(A, As, 2<<k); copy_into(Ai, Ais, Ai
           .size()):
      fft(As, k+2, false); fft(Ais, k+2, false);
      for (int i = 0; i < (4<<k); ++i) As[i] = As[i
          ] * A is [i] * A is [i];
      fft(As, k+2, true); Ai.resize(2<<k, {});</pre>
      for (int i = 0; i < (2 << k); ++i) Ai[i] = T(2)
           * Ai[i] - As[i]:
    Ai.resize(n):
    return Ai:
22 // Polynomial division. Returns {Q, R} such that
      A = QB+R, deg R < deg B.
_{23} // Requires that the leading term of B is nonzero
24 pair < vector < T >, vector < T >> divmod (const vector < T >
       &A, const vector <T> &B) {
    size_t n = A.size()-1, m = B.size()-1;
    if (n < m) return {vector <T>(1, T(0)), A};
    vector < T > X(A), Y(B), Q, R;
    convolution(rev(X), Y = inverse(rev(Y), n-m+1),
    Q.resize(n-m+1); rev(Q);
    X.resize(Q.size()), copy_into(Q, X, Q.size());
    Y.resize(B.size()), copy_into(B, Y, B.size());
    convolution(X, Y, X):
    R.resize(m), copy_into(A, R, m);
    for (size_t i = 0; i < m; ++i) R[i] = R[i] - X[</pre>
    while (R.size() > 1 && R.back().zero()) R.
        pop_back();
    return {Q, R};
41 vector <T > mod(const vector <T > &A, const vector <T >
       &B) {
    return divmod(A, B).second;
43 }
```

3.8.8 Linear recurs. Given a linear recurrence of the form

$$a_n = \sum_{i=0}^{k-1} c_i a_{n-i-1}$$

this code computes a_n in $O(k \log k \log n)$ time.

```
1 #include "header.h"
2 #include "poly.cpp"
3 // x^k mod f
```

```
4 vector<T> xmod(const vector<T> f, ll k) {
     vector \langle T \rangle r\{T(1)\};
     for (int b = 62; b >= 0; --b) {
      if (r.size() > 1)
         square_inplace(r), r = mod(r, f);
      if ((k>>b)&1) {
        r.insert(r.begin(), T(0));
        if (r.size() == f.size()) {
11
          T c = r.back() / f.back():
           for (size_t i = 0; i < f.size(); ++i)</pre>
             r[i] = r[i] - c * f[i]:
           r.pop_back();
16
     return r;
20 }
_{21} // Given A[0,k) and C[0, k), computes the n-th
      term of:
22 // A[n] = \sum i C[i] * A[n-i-1]
23 T nth_term(const vector<T> &A, const vector<T> &C
      . 11 n) {
    int k = (int)A.size();
    if (n < k) return A[n];</pre>
    vector <T> f(k+1, T{1});
    for (int i = 0; i < k; ++i)</pre>
     f[i] = T\{-1\} * C[k-i-1]:
    f = xmod(f, n);
    T r = T\{0\};
    for (int i = 0; i < k; ++i)</pre>
      r = r + f[i] * A[i]:
    return r;
35
36 }
```

3.8.9 Convolution Precise up to 9e15

```
1 #include "header.h"
2 #include "fft.cpp"
3 void convolution_mod(const vi &A, const vi &B, 11
       MOD, vi &C) {
   int s = A.size() + B.size() - 1: 11 m15 = (1LL
        <<15) -1LL;
   int q = 32 - __builtin_clz(s-1), N=1<<q; //</pre>
        fails if s=1
   vector < T > Ac(N), Bc(N), R1(N), R2(N);
    for (size_t i = 0; i < A.size(); ++i) Ac[i] = T</pre>
        \{A[i]\&m15, A[i]>>15\};
  for (size_t i = 0; i < B.size(); ++i) Bc[i] = T</pre>
        {B[i]&m15, B[i]>>15};
   fft(Ac, q, false); fft(Bc, q, false);
  for (int i = 0, j = 0; i < N; ++i, j = (N-1)&(N
        -i)) {
```

```
T as = (Ac[i] + Ac[j].conj()) / 2;
      T = (Ac[i] - Ac[i].coni()) / T{0, 2};
      T bs = (Bc[i] + Bc[j].conj()) / 2;
      T bl = (Bc[i] - Bc[i].coni()) / T{0. 2}:
      R1[i] = as*bs + al*bl*T{0,1}, R2[i] = as*bl +
            al*bs:
16
    fft(R1, q, true); fft(R2, q, true);
17
    11 p15 = (1LL << 15) \% MOD . p30 = (1LL << 30) \% MOD : C.
        resize(s);
    for (int i = 0; i < s; ++i) {</pre>
      11 1 = 11round(R1[i].u), m = 11round(R2[i].u)
          , h = llround(R1[i].v);
      C[i] = (1 + m*p15 + h*p30) \% MOD;
22
23 }
```

3.8.10 Partitions of n Finds all possible partitions of a number

```
1 #include "header.h"
void printArray(int p[], int n) {
3 for (int i = 0: i < n: i++)
      cout << p[i] << "";
    cout << endl;</pre>
6 }
8 void printAllUniqueParts(int n) {
    int p[n]; // An array to store a partition
   int k = 0; // Index of last element in a
    p[k] = n: // Initialize first partition as
        number itself
    // This loop first prints current partition
        then generates next
    // partition. The loop stops when the current
        partition has all 1s
    while (true) {
      printArrav(p, k + 1):
      // Find the rightmost non-one value in p[].
          Also, update the
      // rem_val so that we know how much value can
10
           be accommodated
      int rem val = 0:
      while (k >= 0 \&\& p[k] == 1) {
        rem_val += p[k];
22
      }
25
      // if k < 0, all the values are 1 so there
          are no more partitions
      if (k < 0) return:
```

```
// Decrease the p[k] found above and adjust
          the rem val
      p[k]--:
      rem_val++;
      // If rem_val is more, then the sorted order
          is violated. Divide
      // rem val in different values of size p[k]
          and copy these values at
      // different positions after p[k]
      while (rem_val > p[k]) {
        p[k + 1] = p[k];
        rem_val = rem_val - p[k];
      }
40
41
      // Copy rem_val to next position and
          increment position
      p[k + 1] = rem val:
      k++:
44
  }
45
46 }
```

3.8.11 Ternary search

```
1 /**
2 * Description:
3 * Find the smallest i in $[a,b]$ that maximizes
       f(i), assuming that f(a) < dots < f(i) 
       ge \dots \ge f(b)$.
4 * To reverse which of the sides allows non-
       strict inequalities, change the < marked
       with (A) to <=. and reverse the loop at (B).
* To minimize $f$, change it to >, also at (B).
   int ind = ternSearch(0,n-1,[\k](int i){return a
        [i];});
8 * Time: O(\log(b-a))
10 #include "../Numerical/template.cpp"
12 template < class F>
int ternSearch(int a, int b, F f) {
14 assert(a <= b):
    while (b - a >= 5) {
     int mid = (a + b) / 2:
      if (f(mid) < f(mid+1)) a = mid; // (A)</pre>
      else b = mid+1;
    rep(i,a+1,b+1) if (f(a) < f(i)) a = i: // (B)
    return a;
21
22 }
```

3.9 Other Data Structures

3.9.1 Disjoint set (i.e. union-find)

```
1 template <typename T>
2 class DisjointSet {
      typedef T * iterator;
      T *parent, n, *rank;
      public:
          // O(n), assumes nodes are [0, n)
          DisjointSet(T n) {
              this->parent = new T[n];
              this -> n = n;
              this->rank = new T[n]:
              for (T i = 0; i < n; i++) {</pre>
                   parent[i] = i;
                   rank[i] = 0;
              }
          }
17
18
          // O(log n)
          T find_set(T x) {
              if (x == parent[x]) return x;
              return parent[x] = find_set(parent[x
          }
          // O(log n)
24
          void union_sets(T x, T y) {
25
              x = this->find_set(x);
              v = this->find_set(v);
27
              if (x == y) return;
              if (rank[x] < rank[y]) {</pre>
                  Tz = x;
                  x = y;
                  y = z;
              parent[v] = x;
              if (rank[x] == rank[y]) rank[x]++;
40 };
```

3.9.2 Fenwick tree (i.e. BIT) eff. update + prefix sum calc.

```
#include "header.h"
#define maxn 200010
int t,n,m,tree[maxn],p[maxn];

void update(int k, int z) {
```

3.9.3 Fenwick2d tree

```
1 #include "header.h"
2 template <class T>
3 struct FenwickTree2D {
    vector < vector <T> > tree;
   int n:
   FenwickTree2D(int n) : n(n) { tree.assign(n +
        1, vector <T>(n + 1, 0)); }
   T query(int x1, int y1, int x2, int y2) {
      return query (x2, y2) + query (x1-1, y1-1) - query (x2
          .v1-1)-querv(x1-1,v2):
    T query(int x, int y) {
     T s = 0:
      for (int i = x; i > 0; i = (i & (-i)))
        for (int j = y; j > 0; j -= (j & (-j)))
          s += tree[i][i]:
      return s;
15
   }
16
    void update(int x, int y, T v) {
      for (int i = x; i <= n; i += (i & (-i)))
        for (int j = y; j <= n; j += (j & (-j)))
          tree[i][j] += v;
21
  }
```

3.9.4 Trie

```
#include "header.h"
const int ALPHABET_SIZE = 26;
inline int mp(char c) { return c - 'a'; }

struct Node {
Node* ch[ALPHABET_SIZE];
bool isleaf = false;
Node() {
```

```
for(int i = 0; i < ALPHABET_SIZE; ++i) ch[i]</pre>
           = nullptr;
10
11
    void insert(string &s, int i = 0) {
      if (i == s.length()) isleaf = true;
      else {
        int v = mp(s[i]);
        if (ch[v] == nullptr)
           ch[v] = new Node();
         ch[v] \rightarrow insert(s, i + 1);
19
20
    }
21
    bool contains(string &s, int i = 0) {
      if (i == s.length()) return isleaf;
23
      else {
        int v = mp(s[i]);
        if (ch[v] == nullptr) return false;
         else return ch[v]->contains(s, i + 1);
28
    }
29
    void cleanup() {
      for (int i = 0; i < ALPHABET_SIZE; ++i)</pre>
        if (ch[i] != nullptr) {
          ch[i]->cleanup();
          delete ch[i]:
        }
   }
38 };
```

3.9.5 Treap A binary tree whose nodes contain two values, a key and a priority, such that the key keeps the BST property

```
#include "header.h"

struct Node {

11 v;

4   int sz, pr;

5   Node *1 = nullptr, *r = nullptr;

6   Node(11 val) : v(val), sz(1) { pr = rand(); }

7 };

8   int size(Node *p) { return p ? p->sz : 0; }

9   void update(Node* p) {

10     if (!p) return;

11   p->sz = 1 + size(p->1) + size(p->r);

12     // Pull data from children here

13 }

14   void propagate(Node *p) {

15     if (!p) return;

16     // Push data to children here

17 }

18   void merge(Node *&t, Node *1, Node *r) {
```

```
propagate(1), propagate(r);
    if (!1) t = r;
    else if (!r) t = 1;
     else if (1->pr > r->pr)
         merge(1->r, 1->r, r), t = 1;
    else merge(r->1, 1, r->1), t = r;
     update(t);
26 }
27 void spliti(Node *t. Node *&l. Node *&r. int
       index) {
    propagate(t);
    if (!t) { l = r = nullptr; return; }
    int id = size(t->1);
    if (index <= id) // id \in [index, \infty), so</pre>
         move it right
       spliti(t\rightarrow 1, 1, t\rightarrow 1, index), r = t;
       spliti(t\rightarrow r, t\rightarrow r, r, index - id), l = t;
    update(t);
37 void splitv(Node *t, Node *&l, Node *&r, 11 val)
     propagate(t);
    if (!t) { l = r = nullptr; return; }
    if (val \langle = t - \rangle v) // t - \rangle v \in [val, \infty), so
         move it right
       splitv(t\rightarrow 1, 1, t\rightarrow 1, val), r = t;
       splitv(t\rightarrow r, t\rightarrow r, r, val), l = t;
    update(t):
44
46 void clean(Node *p) {
    if (p) { clean(p->1), clean(p->r); delete p; }
48 }
```

3.9.6 Segment tree

```
#include "../header.h"
template <class T, const T&(*op)(const T&, const T&)>
struct SegmentTree {
   int n; vector<T> tree; T id;
   SegmentTree(int _n, T _id) : n(_n), tree(2 * n, _id), id(_id) { }
   void update(int i, T val) {
   for (tree[i+n] = val, i = (i+n)/2; i > 0; i _/= 2)
        tree[i] = op(tree[2*i], tree[2*i+1]);
   }
   T query(int 1, int r) {
        T lhs = T(id), rhs = T(id);
        for (1 += n, r += n; 1 < r; 1 >>= 1, r >>= 1)
        {
        if ( 1&1 ) lhs = op(lhs, tree[1++]);
    }
}
```

3.9.7 Lazy segment tree Uptimizes range updates

```
1 #include "../header.h"
2 using T=int; using U=int; using I=int;
      exclusive right bounds
3 T t_id; U u_id;
4 T op(T a, T b){ return a+b; }
5 void join(U &a, U b){ a+=b: }
6 void apply(T &t, U u, int x){ t+=x*u; }
7 T convert(const I &i){ return i: }
8 struct LazySegmentTree {
   struct Node { int 1, r, 1c, rc; T t; U u;
      Node(int 1, int r, T t=t_id):1(1),r(r),1c(-1)
         ,rc(-1),t(t),u(u_id)
   int N; vector < Node > tree; vector < I > &init;
    LazySegmentTree(vector < I > &init) : N(init.size
        ()), init(init){
      tree.reserve(2*N-1); tree.push_back({0,N});
         build(0, 0, N);
15
   void build(int i, int l, int r) { auto &n =
       tree[i];
      if (r > 1+1) \{ int m = (1+r)/2;
       ,r});
       build(n.lc,1,m);
                            build(n.rc,m,r);
        n.t = op(tree[n.lc].t, tree[n.rc].t);
     } else n.t = convert(init[1]);
    void push(Node &n, U u) { apply(n.t, u, n.r-n.l)
        ; join(n.u,u); }
    void push(Node &n){push(tree[n.lc],n.u);push(
       tree[n.rc],n.u);n.u=u_id;}
   T query(int 1, int r, int i = 0) { auto &n =
       tree[i]:
      if(r <= n.1 || n.r <= 1) return t_id;</pre>
      if(1 <= n.1 && n.r <= r) return n.t:
      return push(n), op(query(1,r,n.lc),query(1,r,
         n.rc));
30
   void update(int 1, int r, U u, int i = 0) {
       auto &n = tree[i];
      if(r <= n.1 || n.r <= 1) return:
      if(1 <= n.1 && n.r <= r) return push(n,u);</pre>
      push(n); update(1,r,u,n.lc); update(1,r,u,n.
34
```

```
35     n.t = op(tree[n.lc].t, tree[n.rc].t);
36    }
37 };
```

3.9.8 Suffix tree

```
1 #include "../header.h"
2 using T = char;
3 using M = map<T,int>; // or array<T,</pre>
      ALPHABET_SIZE >
4 using V = string;
                        // could be vector<T> as
      well
5 using It = V::const_iterator;
6 struct Node{
7 It b, e; M edges; int link; // end is
        exclusive
    Node(It b, It e) : b(b), e(e), link(-1) {}
   int size() const { return e-b; }
11 struct SuffixTree{
    const V &s; vector < Node > t;
    int root,n,len,remainder,llink; It edge;
    SuffixTree(const V &s) : s(s) { build(); }
    int add_node(It b, It e){ return t.push_back({b
        ,e}), t.size()-1; }
    int add_node(It b){ return add_node(b,s.end());
    void link(int node){ if(llink) t[llink].link =
        node; llink = node; }
    void build(){
      len = remainder = 0; edge = s.begin();
      n = root = add_node(s.begin(), s.begin());
      for(auto i = s.begin(): i != s.end(): ++i){
        ++remainder; llink = 0;
        while(remainder){
          if(len == 0) edge = i;
          if(t[n].edges[*edge] == 0){
              new leaf
            t[n].edges[*edge] = add node(i): link(n
                );
          } else {
            auto x = t[n].edges[*edge]; // neXt
                node [with edge]
            if(len >= t[x].size()){}
                                        // walk to
                next node
              len -= t[x].size(); edge += t[x].size
                  (); n = x;
              continue;
            if(*(t[x].b + len) == *i){
               along edge
              ++len; link(n); break;
                           // split edge
            auto split = add_node(t[x].b, t[x].b+
                len):
```

```
t[n].edges[*edge] = split;
t[x].b += len;
t[split].edges[*i] = add_node(i);
t[split].edges[*t[x].b] = x;
link(split);
}
--remainder;
if(n == root && len > 0)
--len, edge = i - remainder + 1;
else n = t[n].link > 0 ? t[n].link : root
;
}
```

3.9.9 UnionFind

```
1 #include "header.h"
2 struct UnionFind {
    std::vector<int> par, rank, size;
    UnionFind(int n) : par(n), rank(n, 0), size(n,
        1), c(n) {
      for(int i = 0; i < n; ++i) par[i] = i;</pre>
    int find(int i) { return (par[i] == i ? i : (
        par[i] = find(par[i])); }
    bool same(int i, int j) { return find(i) ==
        find(j); }
    int get_size(int i) { return size[find(i)]; }
    int count() { return c; }
    int merge(int i, int j) {
      if((i = find(i)) == (j = find(j))) return -1;
      if(rank[i] > rank[j]) swap(i, j);
15
      par[i] = j;
      size[j] += size[i];
      if(rank[i] == rank[j]) rank[j]++;
      return j;
20
21 };
```

4 Other Mathematics

4.1 Helpful functions

4.1.1 Euler's Totient Fucntion $n = p_1^{k_1-1} \cdot (p_1-1) \cdot \dots \cdot p_r^{k_r-1} \cdot (p_r-1)$, where $p_1^{k_1} \cdot \dots \cdot p_r^{k_r}$ is the prime factorization of n.

```
1 # include "header.h"
2 11 phi(11 n) { // \Phi(n)
      ll ans = 1:
      for (11 i = 2; i*i <= n; i++) {</pre>
           if (n % i == 0) {
               ans *= i-1:
               n /= i;
               while (n % i == 0) {
                   ans *= i;
                   n /= i;
           }
      if (n > 1) ans *= n-1:
       return ans;
17 vi phis(int n) { // All \Phi(i) up to n
    vi phi(n + 1, OLL);
    iota(phi.begin(), phi.end(), OLL);
    for (11 i = 2LL; i <= n; ++i)</pre>
      if (phi[i] == i)
        for (ll j = i; j <= n; j += i)</pre>
           phi[j] -= phi[j] / i;
    return phi;
25 }
```

Formulas $\Phi(n)$ counts all numbers in $1, \ldots, n-1$ coprime to n. $a^{\varphi(n)} \equiv 1 \mod n$, a and n are coprimes. $\forall e > \log_2 m : n^e \mod m = n^{\Phi(m)+e \mod \Phi(m)} \mod m$. $\gcd(m,n) = 1 \Rightarrow \Phi(m \cdot n) = \Phi(m) \cdot \Phi(n)$.

4.1.2 Pascal's trinagle $\binom{n}{k}$ is k-th element in the n-th row, indexing both from 0

4.2 Theorems and definitions

Fermat's little theorem

$$a^p \equiv a \mod p$$

Subfactorial

$$!n = n! \sum_{i=0}^{n} \frac{(-1)^{i}}{i!}$$

$$!(0) = 1, !n = n \cdot !(n-1) + (-1)^n$$

Binomials and other partitionings

$$\binom{n}{k} = \binom{n-1}{k} + \binom{n-1}{k-1} = \prod_{i=1}^{k} \frac{n-i+1}{i}$$

This last product may be computed incrementally since any product of k' consecutive values is divisible by k'!. Basic identities: The hockeystick identity:

$$\sum_{k=r}^{n} \binom{k}{r} = \binom{n+1}{r+1}$$

or

$$\sum_{k \le n} \binom{r+k}{k} = \binom{r+n+1}{n}$$

Also

$$\sum_{k=0}^{n} \binom{k}{m} = \binom{n+1}{m+1}$$

$$\sum_{i=0}^{n} \binom{n}{i} = 2^n$$

For $n, m \geq 0$ and p prime: write n, m in base p, i.e. $n = n_k p^k + \dots + n_1 p + n_0$ and $m = m_k p^k + \dots + m_1 p + m_0$. Then by Lucas theorem we have $\binom{n}{m} \equiv \prod_{i=0}^k \binom{n_i}{m_i} \mod p$, with the convention that $n_i < m_i \implies \binom{n_i}{m_i} = 0$.

Fibonacci (See also number theory section)

$$\sum_{0 \le k \le n} \binom{n-k}{k} = F_{n+1}$$

$$F_n = \frac{1}{\sqrt{5}} \left(\frac{1 + \sqrt{5}}{2} \right)^n - \frac{1}{\sqrt{5}} \left(\frac{1 - \sqrt{5}}{2} \right)^n$$

$$\sum_{i=1}^n F_i = F_{n+2} - 1, \sum_{i=1}^n F_i^2 = F_n F_{n+1}$$

$$\gcd(F_m, F_n) = F_{\gcd(m,n)}$$

$$\gcd(F_n, F_{n+1}) = \gcd(F_n, F_{n+2}) = 1$$

Bit stuff $a+b=a\oplus b+2(a\&b)=a|b+a\&b$. kth bit is set in x iff $x \mod 2^{k-1} \geq 2^k$, or iff $x \mod 2^{k-1}-x \mod 2^k \neq 0$ (i.e. $=2^k$) It comes handy when you need to look at the bits of the numbers which are pair sums or subset sums etc.

$$n \mod 2^i = n\&(2^i - 1).$$

$$\forall k: \ 1 \oplus 2 \oplus \ldots \oplus (4k-1) = 0$$

Stirling's numbers First kind: $S_1(n,k)$ count permutations on n items with k cycles. $S_1(n,k) = S_1(n-1,k-1) + (n-1)S_1(n-1,k)$ with $S_1(0,0) = 1$. Note:

$$\sum_{k=0}^{n} S_1(n,k)x^k = x(x+1)\dots(x+n-1)$$

$$\sum_{k=0}^{n} S_1(n,k) = n!$$

Second kind: $S_2(n, k)$ count partitions of n distinct elements into exactly k non-empty groups.

$$S_2(n,k) = S_2(n-1,k-1) + kS_2(n-1,k)$$

$$S_2(n,1) = S_2(n,n) = 1$$

$$S_2(n,k) = \frac{1}{k!} \sum_{i=1}^{k} (-1)^{k-i} {k \choose i} i^n$$

4.3 Geometry Formulas

$$[ABC] = rs = \frac{1}{2}ab\sin\gamma$$

$$= \frac{abc}{4R} = \sqrt{s(s-a)(s-b)(s-c)} = \frac{1}{2} \left| (B-A, C-A)^T \right|$$

$$s = \frac{a+b+c}{2} \qquad 2R = \frac{a}{\sin \alpha}$$
 cosine rule:
$$c^2 = a^2 + b^2 - 2ab\cos \gamma$$
 Euler:
$$1 + CC = V - E + F$$
 Pick:
$$\operatorname{Area} = \operatorname{itr} \operatorname{pts} + \frac{\operatorname{bdry} \operatorname{pts}}{2} - 1$$

$$p \cdot q = |p||q|\cos(\theta) \qquad |p \times q| = |p||q|\sin(\theta)$$

Given a non-self-intersecting closed polygon on n vertices, given as (x_i, y_i) , its centroid (C_x, C_y) is given as:

$$C_x = \frac{1}{6A} \sum_{i=0}^{n-1} (x_i + x_{i+1})(x_i y_{i+1} - x_{i+1} y_i),$$

$$C_y = \frac{1}{6A} \sum_{i=0}^{n-1} (y_i + y_{i+1})(x_i y_{i+1} - x_{i+1} y_i)$$

$$A = \frac{1}{2} \sum_{i=0}^{n-1} (x_i y_{i+1} - x_{i+1} y_i) = \text{polygon area}$$

Inclusion-Exclusion For appropriate f compute $\sum_{S\subseteq T} (-1)^{|T\setminus S|} f(S)$, or if only the size of S matters, $\sum_{s=0}^{n} (-1)^{n-s} \binom{n}{s} f(s)$. In some contexts we might use Stirling numbers, not binomial coefficients!

Some useful applications:

Graph coloring Let I(S) count the number of independent sets contained in $S \subseteq V$ ($I(\emptyset) = 1$, $I(S) = I(S \setminus v) + I(S \setminus N(v))$). Let $c_k = \sum_{S \subseteq V} (-1)^{|V \setminus S|} I(S)$. Then V is k-colorable iff v > 0. Thus we can compute the chromatic number of a graph in $O^*(2^n)$ time.

Burnside's lemma Given a group G acting on a set X, the number of elements in X up to symmetry is

$$\frac{1}{|G|} \sum_{g \in G} |X^g|$$

with X^g the elements of X invariant under g. For example, if f(n) counts "configurations" of some sort of length n, and we want to count them up to rotational symmetry using $G = \mathbb{Z}/n\mathbb{Z}$, then

$$g(n) = \frac{1}{n} \sum_{k=0}^{n-1} f(\gcd(n,k)) = \frac{1}{n} \sum_{k||n} f(k)\phi(n/k)$$

I.e. for coloring with c colors we have $f(k) = k^c$.

Relatedly, in Pólya's enumeration theorem we imagine X as a set of n beads with G permuting the beads (e.g. a necklace, with G all rotations and reflections of the n-cycle, i.e. the dihedral group D_n). Suppose further that we had Y colors, then the number of G-invariant colorings Y^X/G is counted by

$$\frac{1}{|G|} \sum_{g \in G} |Y|^{c(g)}$$

with c(g) counting the number of cycles of g when viewed as a permutation of X. We can generalize this to a weighted version: if the color i can occur exactly r_i times, then this is counted by the coefficient of $t_1^{r_1} ldots t_n^{r_n}$ in the polynomial

$$Z(t_1, \dots, t_n) = \frac{1}{|G|} \sum_{g \in G} \prod_{m \ge 1} (t_1^m + \dots + t_n^m)^{c_m(g)}$$

where $c_m(g)$ counts the number of length m cycles in g acting as a permutation on X. Note we get the original formula by setting all $t_i = 1$. Here Z is the cycle index. Note: you can cleverly deal with even/odd sizes by setting some t_i to -1.

Lucas Theorem If p is prime, then:

$$\frac{p^a}{k} \equiv 0 \pmod{p}$$

Thus for non-negative integers $m = m_k p^k + \ldots + m_1 p + m_0$ and $n = n_k p^k + \ldots + n_1 p + n_0$:

$$\frac{m}{n} = \prod_{i=0}^k \frac{m_i}{n_i} \mod p$$

Note: The fraction's mean integer division.

Catalan Numbers - Number of correct bracket sequence consisting of n opening and n closing brackets.

The number of ways to completely parenthesize n+1 factors.

The number of triangulations of a convex polygon with n+2 sides (i.e. the number of partitions of polygon into disjoint triangles by using the diagonals).

The number of ways to connect the 2n points on a circle to form n disjoint i.e. non-intersecting chords.

$$C_n = \frac{1}{n+1} \binom{2n}{n}$$

$$C_0 = 1, \ C_1 = 1, \ C_n = \sum_{k=0}^{n-1} C_k C_{n-1-k}$$

Narayana numbers The number of expressions containing n pairs of parentheses, which are correctly matched and which contain k distinct nestings.

$$N(n,k) = \frac{1}{n} \frac{n}{k} \frac{n}{k-1}$$