



UNIVERSIDAD  
POLITÉCNICA  
DE MADRID

# Don't Let Me Down! Offloading Robot VFs Up to the Cloud

June 20, 2023

K. Gillani (UC3M)  
**J. Martín-Pérez** (UPM)  
M. Groshev (UC3M)  
A. de la Oliva (UC3M)  
R. Gazda (ID)

- 1 Introduction
  - Don't Let Me Down!
- 2 Problem Statement
  - Graph-based formulation
  - Classic Decisions  $a(n)$ ,  $a(n_1, n_2)$
  - Wireless Decision  $a(r, R)$
- 3 Don't Let Me Down!
  - the  $\tau$  Metric
  - the Algorithm
- 4 Experimental evaluation
  - Ware Housing Scenario
  - Industrial scenario
  - Industrial scenario II
- 5 Conclusions & Future Work

Networked robotics interact with servers for, e.g.:

- remote driving; or
- assembly robots.

A remote driving Service  $S$  has VFs:

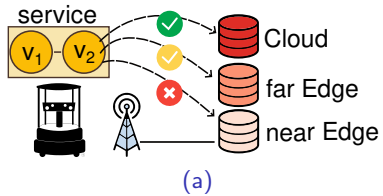
- the robot drivers  $v_1$ ; and
- the remote controller  $v_2$ .

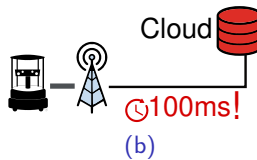
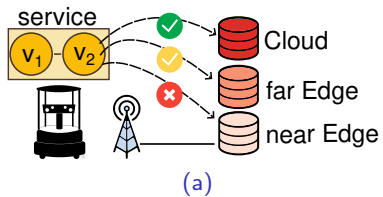
The literature tends to offload – e.g. the remote control  $v_2$  – to Edge, not to Cloud.

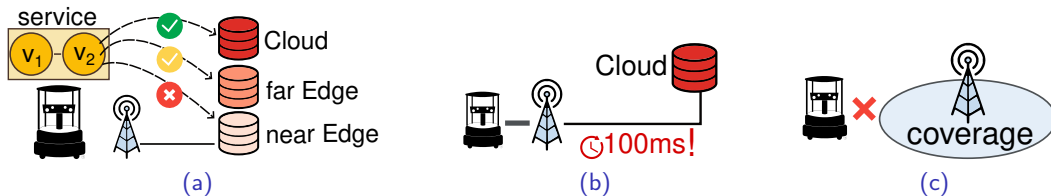
The literature tends to offload – e.g. the remote control  $v_2$  – to Edge, not to Cloud.

*But channel impairments are still there!*

Communication between VFs ( $v_1, v_2$ ) is harnessed.







**Figure 1:** Don't Let Me Down! fosters offloading the robot service VFs up to the cloud (a), yet preventing the cloud large latencies (b) and running out of coverage (c).



Hardware graph  $G$  with:

- robots  $r \in V(G)$ ;
- radio PoA  $R \in V(G)$ ;
- switches/routers  $w \in V(G)$ ; and
- server nodes  $n \in V(G)$

links are edges:

- wireless links  $(r, R) \in E(G)$ ; and
- wired links  $(n_1, n_2) \in E(G)$ .

Robotic service graph  $S$ :

- VFs  $v \in V(S)$ ; and
- VLs  $(v_1, v_2) \in V(S)$ .

The service  $S$  asks for

- a maximum delay  $D(S)$ ;
- computational resources  $C(v)$ ,  $\forall v \in S$ ; and
- bandwidth resources  $\lambda(v_1, v_2)$ ,  $\forall (v_1, v_2) \in E(S)$ .

The decisions to take are

- assigning VFs to servers  $a(n) \subset S$ ; and
- assigning VLs to links  $a(n_1, n_2) \subset E(S)$ .

to minimize the deployment cost

$$\min_{a(n)} \sum_{n \in V(G)} \kappa_n |a(n)| \quad (1)$$

with  $\kappa_n \in \mathbb{R}^+$  the server  $n$  monetary cost, typically satisfying

$$\kappa_{\text{Cloud}} < \kappa_{\text{Edge}} \quad (2)$$

Decide adequate PoA association/handovers  $a(r, R) \in E(S)$  as the robot moves.

$$\sum_{(v_1, v_2) \in a(r_i, R_i)} \lambda(v_1, v_2) \leq T(r_i, R_i), \quad \forall (v_1, v_2) \in a(r_i, R_i) \quad (3)$$

so wireless capacity is enough

$$T(r_i, R_i) = (1 - \delta(r_i, R_i)) \lambda(r_i, R_i) \log_2 \left( 1 + \frac{\sigma_{R_i}(r_i)}{N} \right) \quad (4)$$

with  $\delta(r_i, R_i)$  the PER and  $\sigma_{R_i}(r_i)$  the signal strength.

A shortest-path based algorithm whose core idea is

*Don't Let the VFs go Down to the Edge!*

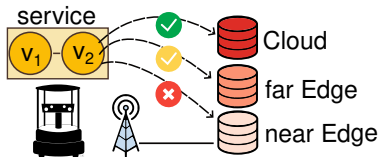


Figure 2: Don't Let Me Down! sends VFs Up to the Cloud.

Don't Let Me Down! is based on the  $\tau$  metric

$$\tau : \mathcal{P}(V(G)) \longrightarrow \mathbb{R}^+ \quad (5)$$

$$\mathcal{P}_0 \longmapsto \kappa_{\mathcal{P}_0[-1]} + \sum_{(n_1, n_2) \in \mathcal{P}_0} \left( \frac{1}{\lambda(n_1, n_2)} + d(n_1, n_2) \right) \quad (6)$$

that fosters the use of

- cheap servers as the Cloud;
- non-congested links; and
- fast links.

Don't Let Me Down! is summarized in three steps:

- 1 *prune PoAs* without enough capacity  $R \notin \{\hat{R}_i\}$

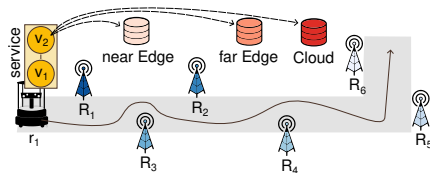
$$\{\hat{R}_i\} = \{R_i : \lambda(v_1, v_2) \leq T(r_i, R_i)\}_i \quad (7)$$

- 2 find server to *offload VFs* using

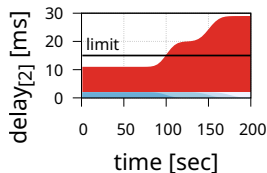
$$\text{Dijkstra}(\text{source} = r, \text{weight} = \tau) \quad (8)$$

- 3 *select PoA* with best free/delay ratio

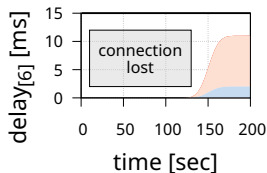
$$\arg \min_{\hat{R}_i} \left\{ \frac{1}{\lambda(r_i, \hat{R}_i)} + d(r_i, \hat{R}_i) \right\} \quad (9)$$



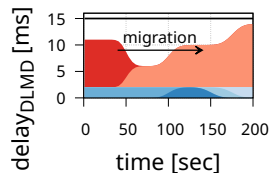
(a)



(b)



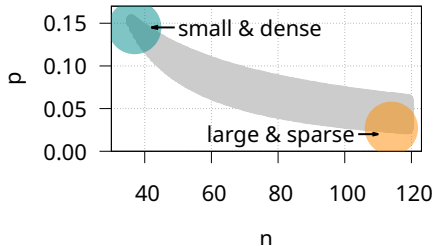
(c)



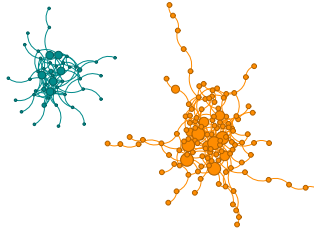
(d)

Figure 3: Delay experienced by the robot using [1], [2] and Don't Let Me Down (DLMD).





(a)



(b)



(c)

Figure 4: Erdős-Rényi setups (a) network graphs (b) with PoAs of an industrial area (c).

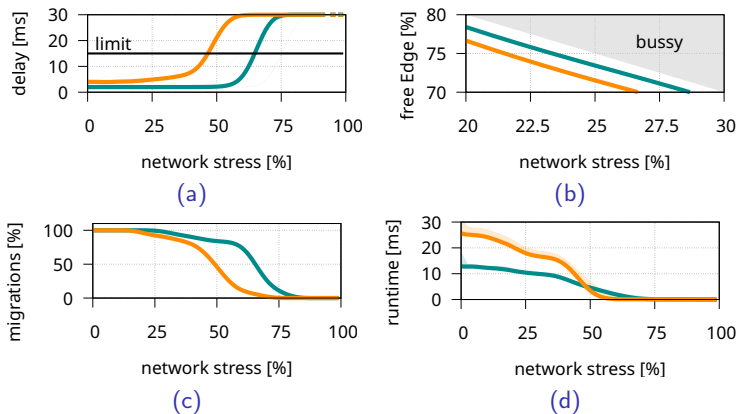




Figure 5: DLMD stress test in Fig. 4 small (green) and large (orange) networks.

## Don't Let Me Down!

- takes  $\sim 10$  [ms];
- reaches optimality in small scenarios (backup);
- offloads & handovers to meet  $< 15$  [ms]; and
- minimizes Edge consumption.

## Still to

- study w/ weighted  $\frac{1}{\lambda} + d$  metric for  $\tau$ ;
- include SINR+propagation models; and
- study probabilistic mobility models.

-  C. Delgado, L. Zanzi, X. Li, and X. Costa-Pérez, “OROS: Orchestrating ROS-driven Collaborative Connected Robots in Mission-Critical Operations,” *IEEE*, 2022.
-  B. Németh, N. Molner, J. Martín-Pérez, C. J. Bernardos, A. de la Oliva, and B. Sonkoly, “Delay and Reliability-Constrained VNF Placement on Mobile and Volatile 5G Infrastructure,” *IEEE Transactions on Mobile Computing*, vol. 21, no. 9, 2022.

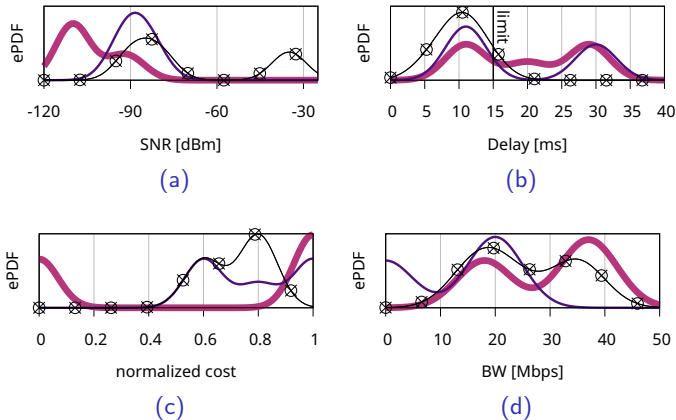


Figure 6: Metrics' ePDFs using DLMD (circle), optimal (cross), [1] (thickest), and [2] (thick)