Interactive Theorem Proving Assignment 4
Martin Skilleter (u6679334)
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## import tactic.basic

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/-
It's time to learn about universes!
This assignment is hopefully fairly quick, but will walk you through the
universe polymorphism issues involved in defining categories in Lean.
 -- Here's a simple attempt at defining a category.
 -- The parameter `C` describes the objects of the category.
class category (C : Type) :=
 (hom : C \rightarrow C \rightarrow Type)
 (id: \Pi X: C, hom XX)
 (comp : \Pi {X Y Z : C}, hom X Y \rightarrow hom Y Z \rightarrow hom X Z)
 (comp\_id : \Pi \{X Y : C\} (f : hom X Y), comp f (id Y) = f)
 (id\_comp : \Pi \{X Y : C\} (f : hom X Y), comp (id X) f = f)
 (assoc : \Pi {W X Y Z : C} (f : hom W X) (g : hom X Y) (h : hom Y Z), comp (comp
 f g) h = comp f (comp g h))
 -- However, this definition is no good: we can't define the category of types:
 instance category_of_types_broken : category Type :=
 { hom := \lambda X Y, X \rightarrow Y }
 -- To fix this, you're going to need to modify the definition above,
 -- to fix the universe levels.
 -- Question 0: Explain why the definition above wasn't useful, perhaps
 -- mentioning Russell's paradox.
/- Answer: Because of the hierarchy of the universes, Type cannot be of type
Type.
Instead, it must be of Type 1 (otherwise, we would have an analogue of Russell's
Paradox:
the set of all sets being an element of itself). In a dependent type theory
this means that nothing can be of Type itself, which is implemented by setting
Type to be of
Type 1, Type 1 of Type 2 etc. This is not reflected in the definition,
as the universe level of the category above is fixed. It should instead be
polymorphic, so
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that elements belonging to higher universe levels can be made into categories.

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One possible solution to this would be to change Type to Type 1 (at the very
least, this would
allow the category of types to be defined. However, we would then have an
identical problem to
the one above when we tried to create a category for a type in a universe level
higher
than 1).
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 -- Here's one attempt:
class category_1 (C : Type 1) :=
 (hom : C \rightarrow C \rightarrow Type)
 (id : \Pi X : C, hom X X)
 (comp : \Pi {X Y Z : C}, hom X Y \rightarrow hom Y Z \rightarrow hom X Z)
 (comp\_id : \Pi \{X Y : C\} (f : hom X Y), comp f (id Y) = f)
 (id\_comp : \Pi \{X Y : C\} (f : hom X Y), comp (id X) f = f)
 (assoc : \Pi {W X Y Z : C} (f : hom W X) (g : hom X Y) (h : hom Y Z), comp (comp
 f g) h = comp f (comp g h))
 -- Question 1: fill in the remaining fields of this definition.
 instance category_of_types : category_1 Type :=
 { hom := \lambda X Y, X \rightarrow Y,
  id := \lambda X, \lambda X, X,
  comp := \lambda X Y Z, \lambda f g, \lambda x, g (f x),
 comp_id := by {intros X Y f, simp},
 id_comp := by {intros X Y f, simp},
 assoc := by {intros W X Y Z f g h, simp}
 }
 -- However, this variant has its own problems. A standard solution to Russell's
 -- paradox about "the set of all sets" (resolved in Lean's dependent type
 -- by the typing judgement `Type : Type 1`) is to consider all subsets of some
 -- fixed 'big' set, rather than "all sets".
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-- Question 2: Decide whether `C` below should be `category` or `category_1`,
-- and fill in the remaining fields.
instance category_of_subsets (X : Type) : category (set X) :=
\{ \text{ hom } := \lambda P Q, \{ x // P x \} \rightarrow \{ x // Q x \}, \}
 id := \lambda P, \lambda x, x,
 comp := \lambda P Q R, \lambda p q, \lambda x, q (p x),
 comp_id := by {intros P Q f, simp},
id_comp := by {intros P Q f, simp},
assoc := by {intros P Q R f g h, simp}
}
-- We've now got a conundrum: for some reasonable examples, we want the
-- objects and morphisms to live in the same universe, while for other
-- examples we want the objects to live one universe level higher than
-- the morphisms.
-- In ZFC based category theory, these two variants of the notion of a category
-- are called "small categories" and "large categories".
-- Rather than duplicate everything we want to prove about categories
-- (or worse: we'll need to worry about four different sorts of functors,
-- as the source and target categories could individually be small or large!)
-- in the mathlib category theory library we've made a "universe polymorphic"
-- definition, which is essentially this one:
universes v u
class pcategory (C : Type u) :=
(hom : C \rightarrow C \rightarrow Type \ v)
(id : \Pi X : C, hom X X)
(comp : \Pi {X Y Z : C}, hom X Y \rightarrow hom Y Z \rightarrow hom X Z)
(comp\_id : \Pi \{X Y : C\} (f : hom X Y), comp f (id Y) = f)
(id\_comp : \Pi \{X Y : C\} (f : hom X Y), comp (id X) f = f)
(assoc : \Pi {W X Y Z : C} (f : hom W X) (g : hom X Y) (h : hom Y Z), comp (comp
f g) h = comp f (comp g h))
-- Question 3:
-- Work out the appropriate values of the universe parameters `p`, `q`, `r`,
-- below, and then verify that the same fields you used above can be used to
-- complete the following two definitions:
-- (Hint: substitute fixed large values, like p=5, q=10, and read the error
messages...)
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instance category_of_types' : pcategory.{0 1} Type :=
{ hom := \lambda X Y, X \rightarrow Y,
 comp := \lambda X Y Z, \lambda f g, \lambda x, g (f x),
 comp_id := by {intros X Y f, simp},
 id_comp := by {intros X Y f, simp},
 assoc := by {intros W X Y Z f g h, simp}
}
instance category_of_subsets' (X : Type) : pcategory.{0 0} (set X) :=
{ hom := \lambda P Q, { x // P x } \rightarrow { x // Q x },
 id := \lambda P, \lambda x, x,
 comp := \lambda P Q R, \lambda p q, \lambda x, q (p x),
 comp_id := by {intros P Q f, simp},
 id_comp := by {intros P Q f, simp},
 assoc := by {intros P Q R f g h, simp}
}
-- Question 4:
-- Complete the following definitions:
universes V<sub>1</sub> U<sub>1</sub> V<sub>2</sub> U<sub>2</sub>
structure Functor (C : Type u_1) [pcategory.{v_1 u_1} C] (D : Type u_2)
[pcategory.\{v_2 u_2\} D] :=
(obj : C \rightarrow D)
(map : \Pi {X Y : C}, pcategory.hom X Y \rightarrow pcategory.hom (obj X) (obj Y))
(id_map : □ X : C, map (pcategory.id X) = pcategory.id (obj X))
(map\_comp : \Pi \times Y \times Z : C, \Pi f : pcategory.hom \times Y, \Pi g : pcategory.hom Y Z,
map (pcategory.comp f g) = pcategory.comp (map f) (map g))
-- Hint: there are two missing fields here, giving the axioms for functors.
def List : Functor Type Type :=
{ obj := \lambda \alpha, list \alpha,
 map := \lambda \alpha \beta : Type, \lambda f : pcategory.hom \alpha \beta, \lambda L, list.map f L,
 id_map := by {intros X, dsimp [pcategory.id], funext, induction x, simp, dsimp
[list.map], rw [x_ih]},
 map_comp := by {intros X Y Z f g, simp, funext, dsimp [pcategory.comp],
induction x,
                   simp, dsimp [list.map], rw [x_ih],}}
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-- Question 5:
-- Complete the following definitions:
structure NaturalTransformation
 {C : Type u_1} [pcategory.{v_1 u_1} C] {D : Type u_2} [pcategory.{v_2 u_2} D]
 (F G : Functor C D) :=
(app : \Pi X : C, pcategory.hom (F.obj X) (G.obj X))
(naturality : ∏ X Y : C, ∏ f : pcategory.hom X Y, pcategory.comp (F.map f)
(app Y) = pcategory.comp (app X) (G.map f))
namespace NaturalTransformation
def id
 {C : Type u_1} [pcategory.{v_1 u_1} C] {D : Type u_2} [pcategory.{v_2 u_2} D]
 (F : Functor C D) : NaturalTransformation F F :=
{app := \lambda X : C, pcategory.id (F.obj X),
naturality := by {intros X Y f, rw [pcategory.comp_id, pcategory.id_comp]}}
def comp
 {C : Type u_1} [pcategory.{v_1 u_1} C] {D : Type u_2} [pcategory.{v_2 u_2} D]
 {F G H : Functor C D}
 (\alpha: NaturalTransformation F G) (\beta: NaturalTransformation G H):
NaturalTransformation F H :=
{app := \lambda X : C, pcategory.comp (\alpha.app X) (\beta.app X),
naturality := by {intros X Y f, rw [←pcategory.assoc, \alpha.naturality,
pcategory.assoc, pcategory.assoc],
apply congr_arg, rw [\beta.naturality]}}
@[extensionality] lemma Natural_ext
{C : Type u_1} [pcategory.{v_1 u_1} C] {D : Type u_2} [pcategory.{v_2 u_2} D]
{F G : Functor C D} {N M : NaturalTransformation F G} :
(\Pi X : C, N.app X = M.app X) \rightarrow N = M :=
begin
 intros h,
 induction N,
 induction M,
 simp at h,
 simp [h],
 funext,
 exact h x,
end
```

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@[simp] lemma id_app
{C : Type u_1} [pcategory.{v_1 u_1} C] {D : Type u_2} [pcategory.{v_2 u_2} D] {F :
Functor C D} :
(id F).app = \lambda X : C, pcategory.id (F.obj X) := rfl
-- Hint: you may find `conv` helpful.
end NaturalTransformation
-- Here's the crux of the whole assignment: understand what the correct values
-- of `p` and `q` are in this definition, then complete the definition.
-- Hint: you may like to prove an extensionality lemma for natural
transformations,
-- and appropriate simp lemmas.
instance {C : Type u_1} [pcategory.{v_1 u_1} C] {D : Type u_2} [pcategory.{v_2 u_2}
D] :
 pcategory.\{(\max u_1 v_2) (\max u_1 u_2 v_1 v_2)\} (Functor C D) :=
{ hom := \lambda F G, NaturalTransformation F G,
 id := NaturalTransformation.id,
 comp := \lambda F G H, NaturalTransformation.comp,
 comp_id := by {intros F G N, ext X, dsimp [NaturalTransformation.comp], rw
[pcategory.comp_id]},
 id_comp := by {intros F G N, ext X, dsimp [NaturalTransformation.comp], rw
[pcategory.id comp]},
 assoc := by {intros F G H I N M O, ext X, dsimp [NaturalTransformation.comp],
rw [pcategory.assoc]}
}
constant C : Type u<sub>1</sub>
constant D : Type u₂
variables [C_cat : pcategory.{v1 u1} C] [D_cat : pcategory.{v2 u2} D]
variables F G : @Functor C C_cat D D_cat
constant N : NaturalTransformation F G
set_option pp.universes true
#check @Functor C C_cat D D_cat -- Type (max u<sub>1</sub> u<sub>2</sub> v<sub>1</sub> v<sub>2</sub>)
#check @NaturalTransformation C C_cat D D_cat F G -- Type (max u<sub>1</sub> v<sub>2</sub>)
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- -- Question 6:
- -- What universe does `NaturalTransformation F G` live in? Why?
- -- Hint: a really good answer will use the word 'impredicativity' somewhere
- -- along the way.

/- Answer: NaturalTransformation F G lives in the universe 'max  $(u_1 \ v_2)$ '. We begin by noting that Functor C D is of type Type max  $(u_1 \ u_2 \ v_1 \ v_2)$ . This is because the field obj has type Type (max  $u_1 \ u_2$ ) (because C : Type  $u_1$  and D : Type  $u_2$  and so) and the field map has type Type (max  $v_1 \ v_2$ ) (because the morphisms in the category of C live in the universe  $v_1$ , and the morphisms in the category of D live in the universe  $v_2$ ).

The other fields are of type Prop = Sort 0 : Type 0, because they are functions from some type

into Prop. Impredicativity means that regardless of the universe level of the domain, the type of this will always be Type 0. We want this feature in our universe levels in Lean because we can interpret a function  $f:A\to B$  for some

types A: Type u and B: Prop by the statement "given some element of A, I can give you an

element of B". In the case where B : Prop, an element of B is a proof of B and so  $f:A\to Prop$ 

represents the statement "if A has an element then B", so f is itself a proposition (and hence

of type Prop), even if this is a lower universe level than A.

All of these facts together mean that Functor C D is of type

Type max  $(max u_1 u_2) (max v_1 v_2) = Type max (u_1 u_2 v_1 v_2)$ .

[Note: I realised after I had written this that the universe level of Functor does not affect

the universe level of NaturalTransformation. However, the explanation about impredicativity above

is still used below when we discuss why the field naturality of NaturalTransformation does not

affect its universe level, so I left it in.]

We can now apply the same reasoning as above to determine the universe level of NaturalTransformation. Because naturality is a Prop, it will not affect the universe level of

NaturalTransformation (once again by impredicativity), so it suffices to determine the universe

level of app. pcategory.hom  $(F.obj\ X)\ (G.obj\ X)$  is a morphism in the category D and so is of type

Type  $v_2$ . Because X: C, app is a function from something of type  $u_1$  to something of type

Type  $v_2$  and so its type is Type (max  $u_1$   $v_2$ ) (meaning it lives in the universe level max  $(u_1$   $v_2$ )).

Then NaturalTransformation has the same universe level as app and so lives in the universe level

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\max (u_1 v_2).
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