19/20 Gone matakes in Q3.

proofs in Q1 are nice, but in a few places can be made shorter/easier.

(ge my solutions)

Interactive Theorem Proving Assignment 1
Martin Skilleter (u6679334)
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```
namespace hidden
inductive Natural
| zero : Natural
| succ : Natural -> Natural
open Natural
instance : has_zero Natural := .
{ zero := zero}
                                            (an you explain what is betweense

if you down

(success b = succe (add ab)

instead?
instance : has_one Natural :=
{ one := succ zero}/
def add : Natural -> Natural -> Natural
a zero := a
| a (succ b) := succ (add a b).
instance : has add Natural :=
{ add := add } //
def times : Natural -> Natural -> Natural
| a zero := zero
| a (succ b) := add (times a b) a
instance : has mul Natural :=
{ mul := times}
def pow : Natural -> Natural -> Natural
a zero := (succ zero)
| a (succ b) := times (pow a b) a
instance : has_pow Natural Natural :=
{ pow := pow}
```

```
-lid remove the whole "show" line here:
The calc is already completely explicit about
What's leing proved.
theorem add_associativity (a b c : Natural) : (a + b) + c = a + (b + c) :=
Natural\rec_on c
(show (a + b) + 0 = a + (b + 0), from calc
     (a + b) + 0 = a + b : rfl
              ... = a + (b + 0) : rfl
)
(assume c, assume ih : (a + b) + c = a + (b + c),
show (a + b) + (c + 1) = a + (b + (c + 1)), from calc
     (a + b) + (c + 1) = ((a + b) + c) + 1 : rfl
                   \dots = (a + (b + c)) + 1 : by rw ih
                    ... = a + ((b + c) + 1) : rfl
                    ... = a + (b + (c + 1)) : rfl
lemma zero_commutativity (a : Natural) : a + 0 = 0 + a :=
Natural.rec_on a
(show zero + 0 = 0 + zero, from rfl
)
(assume a, assume ih : a + 0 = 0 + a,
show (a + 1) + 0 = 0 + (a + 1), from calc
     (a + 1) + 0 = a + 1 : rfl
             \dots = (a + 0) + 1 : rfl
             ... = (0 + a) + 1 : by rw ih
              \dots = 0 + (a + 1) : rfl
)
lemma one_commutativity (a : Natural) : a + 1 = 1 + a :=
Natural.rec_on a
(show (0 + 1 : Natural) = 1 + 0, by rw zero_commutativity)
(assume a, assume ih : a + 1 = 1 + a,
show (a + 1) + 1 = 1 + (a + 1), from calc
     (a + 1) + 1 = (1 + a) + 1: by rw ih
             \dots = 1 + (a + 1) : rfl
)
```

```
theorem add commutativity (a b : Natural) : a + b = b + a :=
Natural.rec_on b
(show (a + 0 : Natural) = 0 + a, by rw zero_commutativity)
(assume b, assume ih : a + b = b + a,
show a + (b + 1) = (b + 1) + a, from calc
     a + (b + 1) = (a + b) + 1: by rw add associativity
             ... = (b + a) + 1 : by rw ih
             ... = b + (a + 1) : by rw add_associativity
             \dots = b + (1 + a) : by rw one_commutativity
             \dots = (b + 1) + a : by rw add_associativity
)
-- Question 1a
theorem left_distributivity (a b c : Natural) : a * (b + c) = a * b + a * c :=
Natural.rec_on c
(show a * (b + 0) = a * b + a * 0, from calc
     a * (b + 0) = a * b : rfl
             ... = a * b + 0 : rfl
             ... = a * b + a * 0 : rfl
(assume c, assume ih : a * (b + c) = a * b + a * c,
show a * (b + (c + 1)) = a * b + a * (c + 1), from calc
     a * (b + (c + 1)) = a * ((b + c) + 1) : by rw add_associativity
                   ... = a * (b + c) + a : rfl
                   ... = (a * b + a * c) + a : by rw ih
                   \dots = a * b + (a * c + a) : by rw add_associativity
                   \dots = a * b + a * (c + 1) : rfl
)
lemma times_zero (a : Natural) : 0 * a = 0 :=
Natural.rec_on a
(show (0 * 0 : Natural) = 0, by refl)
(assume a, assume ih : 0 * a = 0,
show 0 * (a + 1) = 0, from calc
     0 * (a + 1) = 0 * a + 0 * 1 : by rw left_distributivity
             \dots = 0 + 0 * 1 : by rw ih
             ... = 0 + 0 : rfl
             \dots = 0 : rfl
)
```

```
theorem right_distributivity (a b c : Natural) : (a + b) * c = a * c + b * c :=
Natural.rec_on c
(show (a + b) * 0 = a * 0 + b * 0, by refl)
(assume c, assume ih : (a + b) * c = a * c + b * c,
show (a + b) * (c + 1) = a * (c + 1) + b * (c + 1), from calc
     (a + b) * (c + 1) = ((a + b) * c) + (a + b) : rfl
                    ... = (a * c + b * c) + (a + b) : by rw ih
                    ... = ((a * c + b * c) + a) + b: by rw \leftarrow add_associativity
                    \dots = (a * c + (b * c + a)) + b : by rw \leftarrow add_associativity
                    \dots = (a * c + (a + b * c)) + b : by rw add_commutativity (b)
* c)
                    \dots = ((a * c + a) + b * c) + b : by rw \leftarrow add_associativity
                    \dots = (a * c + a) + (b * c + b) : by rw \leftarrow add_associativity
                    \dots = a * (c + 1) + b * (c + 1) : by refl
)
lemma right_times_one (a : Natural) : a * 1 = a :=
(show a * 1 = a, from calc
     a * 1 = a * 0 + a : rfl
       ... = 0 + a : rfl
       ... = a + 0 : by rw add_commutativity
       \dots = a : rfl
)
lemma left_times_one (a : Natural) : 1 * a = a :=
Natural.rec_on a
(show (1 * 0 : Natural) = 0, by refl)
(assume a, assume ih : 1 * a = a,
show 1 * (a + 1) = a + 1, from calc
     1 * (a + 1) = 1 * a + 1 : rfl
            ... = a + 1 : by rw ih
)
```

```
-- Question 1b
theorem times_commutativity (a b : Natural) : a * b = b * a :=
Natural.rec_on b
(show a * 0 = 0 * a, from calc
     a * 0 = 0 : rfl
       ... = 0 * a : by rw times_zero
)
(assume b, assume ih : a * b = b * a,
show a * (b + 1) = (b + 1) * a, from calc
     a * (b + 1) = a * b + a : rfl
             \dots = b * a + a : by rw ih
             ... = a + b * a : by rw add_commutativity
             ... = 1 * a + b * a : by rw left_times_one
             ... = (1 + b) * a : by rw right_distributivity
             \dots = (b + 1) * a : by rw add_commutativity
)
-- Question 1c
theorem times_associativity (a b c : Natural) : (a * b) * c = a * (b * c) :=
Natural.rec_on c
(show (a * b) * 0 = a * (b * 0), from calc
     (a * b) * 0 = 0 : rfl
             ... = a * 0 : rfl
             \dots = a * (b * 0) : rfl
)
(assume c, assume ih : (a * b) * c = a * (b * c),
show (a * b) * (c + 1) = a * (b * (c + 1)), from calc
     (a * b) * (c + 1) = (a * b) * c + a * b : rfl
                   ... = a * (b * c) + a * b : by rw ih
                   \dots = a * ((b * c) + b) : by rw left_distributivity
                   ... = a * (b * (c + 1)) : rfl
)
                                                             great.
end hidden
```

```
import data.finset
import data.real.basic
import data.real.cau seg completion
import data.nat.gcd
import analysis.exponential
noncomputable theory
open nat
open finset
open int
open real
open complex
example : decidable_linear_ordered_comm_group R := by apply_instance
-- n is the length of the arithmetic sequence we want, a is the starting number
-- and k is the common difference (k \neq 0 or else we can choose a to be prime
∡nd we are done)
theorem green tao : Prop :=
\forall (n : \mathbb{N}), \exists (a k : \mathbb{N}), \forall (i < n), (k \neq 0) \land prime (a + i * k)
                                                       Not wrong, but awkward to put this under the quantities for i. You could also use: k: N+
-- The nth partial sum of a series \Sigma f
def P_n (f : \mathbb{N} \to \mathbb{R}) (n : \mathbb{N}) : \mathbb{R} :=
finset.sum (finset.Ico 1 (n+1)) f
-- The series with f inside the sum def converges (f:\mathbb{N}\to\mathbb{R}) (L:\mathbb{R}): Prop := \forall (\varepsilon>0), \exists (N:\mathbb{N}), \forall (n:\mathbb{N}), n \geq N \rightarrow \underline{\text{complex.abs}} (P_n \neq n - L) < \varepsilon
-- The series with f inside the sum
-- Apery's Theorem says that the sum from n=1 to ∞ of 1/n^3 has a limit and have to
                                                                                     untraduce the

R) ≠ L complex number

here, and it will

interfere with proof.
that the limit is irrational
theorem apery's_theorem : Prop :=
\exists (L : \mathbb{R}), converges (\lambda n, 1/(n^3)) L \wedge \forall (p q : \mathbb{Z}), (p/q : \mathbb{R}) \neq L
def find_prime_factors (n: N) : list N :=
((list.range (n+1)).filter prime).filter ( | n)
def rad (n : \mathbb{N}) : \mathbb{N} :=
(find_prime_factors n).prod
```

```
perhaps dance à cote?
                                                                  - This warrants more explanation, You need to know
                                                                         XITE 15 montone in E

    We use formulation 2 of the abc conjecture from

          https://en.wikipedia.org/wiki/Abc_conjecture
          -- Note that because 1/n \to 0 as n \to \infty, this is equivalent to the \varepsilon
          formulation
          theorem abc_conjecture : Prop :=
           \forall (n: \mathbb{N}), \exists (K: \mathbb{R}), \forall (a b c: \mathbb{N}),
          n > 0 \bigwedge \forall (z : \mathbb{N}), (z \mid a \land z \mid b \land z \mid c \rightarrow z = 1) \land a + b = c \rightarrow (c : \mathbb{R})
          \checkmark K*(rad a*b*c : \mathbb{R})^(1+1/n)
          -- We cite the paper "An/Elementary Problem Equivalent to the Riemann
                                                                                    wong place.
          Hypothesis"
          http://www.math.lsa.umich.edu/plagarijas/doc/elementaryrh.pdf See my commit.
          -- by Jeffrey Lagarias, which can be found at
          -- We begin by defining the harmonic series
                                                         N, SU 1/n=0.
          def H_n (n : \mathbb{N}) : \mathbb{R} := P_n (\lambdam, 1/m) n
          def sum_of_divisors (n : N) : R :=
          theorem riemann_hypothesis : Prop :=
          \forall (n : \mathbb{N}), sum_of_divisors n \leq H_n n + exp (H_n n)*log (H_n n) \land
                     sum_of_divisors n = H_n n + exp (H_n n)*log (H_n n) \leftrightarrow n = 1
interestingly, this really should be a det, not a theorem.
       Sadly, Mis one is filse too,

(1) There are some essential parentheses missing!
         2) What about n=0?
                             (see my commit)
```

```
import data.list
namespace hidden
variable {\alpha : Type}
def len \{\alpha : \mathsf{Type}\} : \mathsf{list} \ \alpha \to \mathbb{N}
| [] := 0
| (hd :: tl) := 1 + len tl
def concat {\alpha : Type} : list (list \alpha) → list \alpha
| [] := []
| (hd :: tl) := hd ++ (concat tl)
def nonempty \{\alpha : \mathsf{Type}\} : \mathsf{list} \ \alpha \to \mathsf{Prop}
| [] := false
| (_ :: _) := true
-- If the 2nd of two lists being concatenated is non-empty then their
concatenation is non-empty
lemma nonempty_tail (L M : list \alpha) (h : nonempty M) : nonempty (L ++ M) :=
begin
   cases L,
   {simp,
   exact h},
   {simp}
end
```

```
-- If the 1st of two lists being concatenated is non-empty then their
concatenation is non-empty
-- This is the lemma that is actually used in the proof of our theorem, but the
other lemma
-- is needed for this proof
lemma nonempty_head (L M : list \alpha) (h : nonempty L) : nonempty (L ++ M) :=
begin
   cases M,
   {simp,
   exact h},
                                         I no need to start a rested
begin ... end block
   {have h₂ : nonempty (M_hd :: M_tl), begin
       dsimp [nonempty],
       trivial,
   end,
   apply nonempty_tail,
   exact h<sub>2</sub>,
   }
end
-- Given a non-empty list of lists, all of the elements of which are also
non-empty, we show that the full concatenation is non-empty
theorem nonempty_concat_of_nonempty_is_nonempty
(L : list (list \alpha)) (h : nonempty L) (w : \forall (m \in L), nonempty m) : nonempty
(concat L) :=
begin
 cases L,
 {cases h},
 {dsimp [concat],
 apply nonempty_head,
  apply w,
 simp,
                                 great.
}
end
end hidden
```

```
import data.finset
import data.nat.choose
import tactic.squeeze

open nat
open finset

-- We apply the add_pow theorem, which is actually the binomial expansion under a pseudonym
theorem binomial_expansion_for_2 (n : N) :
finset.sum (finset.range (n+1)) (\lambda (m : N), choose n m) = 2^n := begin
    have h := add_pow 1 1 n,
    simp only [nat.one_pow, mul_one, one_mul, nat.cast_id, nat.pow_eq_pow] at h,
    simp only [h, eq_self_iff_true],
end
```