

**Einführung in Wissensbasierte Systeme WS 2018/19, 3.0 VU,  
184.737**

**Exercise Sheet 1 – Logic, Part 1**

**Exercise 1.1:** Check using semantics whether the following statements hold (do *not* use truth tables). If a statement does not hold, provide a counterexample.

1.  $\neg(p \rightarrow q) \models p \rightarrow \neg q$

2.  $p \rightarrow q \models (p \rightarrow \neg q) \rightarrow \neg p$

3.  $p \rightarrow q \models q \rightarrow \neg p$

**Exercise 1.2:** Let  $\phi$  be a tautology,  $\psi$  a contradiction, and  $\chi$  a contingency (i.e.,  $\chi$  is satisfiable as well as falsifiable). Which of the following formulas are (i) tautological, (ii) contradictory, (iii) contingent, or (iv) logically equivalent to  $\chi$ ? Justify your answers.

1.  $\phi \wedge \chi$ .

2.  $\phi \vee \chi$ .

3.  $\psi \wedge \chi$ .

4.  $\psi \vee \chi$ .

5.  $\phi \vee \psi$ .

6.  $\chi \rightarrow \psi$ .

**Exercise 1.5:** Prove or refute whether the following formulas are tautologies. Establish your claim by purely semantic means (i.e., without applying a calculus like **TC1**).

(a)  $\exists x (P(x) \rightarrow P(f(x)))$

(b)  $\forall x \exists y R(x, y) \rightarrow \exists y \forall x R(x, y)$

*Hint:* Truth tables do not work.

**Exercise 1.6:** Translate the following arguments into entailment problems of first-order logic and check the validity of each statement. If the argument is valid, provide a proof in TC1. Otherwise, prove that there exists no closed tableau for the (negation of the) corresponding formula in TC1.

- (a) Influenza is caused by a virus. Pseudomonas is a bacterium and not a virus, therefore it doesn't cause influenza.
- (b) Every chicken eats corn. Berta eats grass. Therefore Berta is not a chicken.

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## Exercise Sheet 2 – Nonmonotonic Reasoning

**Exercise 2.1:** Consider the following theory:

$$T = \{\forall x(\neg P(x) \rightarrow Q(x)), \neg P(b), R(a), \forall x(R(x) \rightarrow P(x))\}.$$

- (a) Determine  $\text{CWA}(T)$  as well as  $\text{CWA}^Q(T)$ .
- (b) Prove or refute whether  $\text{CWA}(T)$  is consistent. What about  $\text{CWA}^Q(T)$ ?

**Exercise 2.3:** Consider the open default theory  $T = (W, \Delta)$ , where

$$W = \{\exists x(P(x) \vee Q(x)), \forall x(P(x) \vee R(x)), \forall x(R(x) \rightarrow Q(x))\}$$

$$\Delta = \left\{ \frac{P(x) : \neg Q(x)}{\neg Q(x)}, \frac{Q(x) : \neg P(x)}{\neg P(x)}, \frac{\top : R(x)}{R(x)} \right\}.$$

Compute the closure of  $T$  and determine the possible candidates for being an extension. Then compute the classical reduct  $\Delta_E$  for each candidate  $E$ , and determine all extensions of  $T$ .

**Exercise 2.5:** The property *semi-monotonicity* is a weaker form of monotonicity and is defined as follows:

Let  $T = (W, \Delta)$  and  $T' = (W, \Delta')$  be default theories with  $\Delta \subseteq \Delta'$ , and let  $E$  be an extension of  $T$ . Then  $T'$  has an extension  $E'$  such that  $E \subseteq E'$ .

Prove or disprove whether semi-monotonicity holds in default logic.

**Exercise 2.6:** You are given the propositional default theory  $T = (W, \Delta)$ , where

$$W = \{a\},$$

$$\Delta = \left\{ \frac{a : b}{b}, \frac{b : c, \neg d}{c}, \frac{b : d, \neg c}{d}, \frac{c : e}{\neg e} \right\}.$$

Now consider the following (erroneous) reasoning trying to find an extension  $E$  of  $T$ :

1. Starting with the certain knowledge  $a \in W$ , we can apply the default  $a : b/b$  to derive  $b$ .
2. Having derived  $b$ , we can now apply the second default  $b : c, \neg d/c$  to derive  $c$ .
3. Due to  $c$ , the third default is blocked, but we can apply the fourth default  $c : e/\neg e$  and derive  $\neg e$ .
4. No more defaults are applicable, so we arrive at  $E = Cn(\{a, b, c, \neg e\})$ .

However,  $E$  is *not* an extension of  $T$ .

Your task is to analyse the above reasoning and explain in detail why this approach does not work. Furthermore, find the actual extensions of  $T$ .



**Exercise 2.7:** Consider the following information:

1. People who live in a city usually rent a flat.
2. People who live in a city and rent a flat are usually not rich.
3. Rich people usually do not rent a flat.
4. Emily lives in a city.
5. Emily is rich.

Formalise the given information in terms of an open default theory and compute all of its extensions. Use

- $city(x)$  for “ $x$  lives in a city”,
- $flat(x)$  for “ $x$  rents a flat”,
- $rich(x)$  for “ $x$  is rich”, and
- the constant symbol  $e$  for “Emily”.

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**Exercise Sheet 3 – Answer-Set Programming and Probabilistic Reasoning**

**Exercise 3.1:** Let  $P$  be a Horn logic program and  $M_1, M_2 \subseteq HB(P)$  classical models of  $P$ . Prove that  $M_1 \cap M_2$  is also a classical model of  $P$ . What about  $M_1 \cup M_2$ ?

**Exercise 3.2:** Consider the following disjunctive logic program  $P$ :

$$P = \left\{ \begin{array}{l} a \vee b \leftarrow \text{not } c. \\ c \leftarrow b, \text{not } d. \\ d \leftarrow \text{not } c. \\ a \leftarrow d. \end{array} \right\}.$$

- (i) Determine all answer sets of  $P$ . For each proposed answer set  $S$ , argue formally that it is indeed an answer set (using the Gelfond-Lifschitz reduct  $P^S$ ).
- (ii) Is it possible to add new rules  $Q$  to  $P$  such that  $S' = \{b, c\}$  is an answer set of  $P \cup Q$ ? If yes, give the new rules  $Q$ , if not, argue why.
- (iii) The same as in (??), but for  $S' = \{a, b\}$ .

**Exercise 3.3:** For a program  $P$ , we denote by  $AS(P)$  the set of all answer sets of  $P$ . Let  $P, Q$  be programs. We say that  $P, Q$  are

(i) *equivalent* if  $AS(P) = AS(Q)$  and

(ii) *strongly equivalent* if  $AS(P \cup R) = AS(Q \cup R)$  for every program  $R$ .

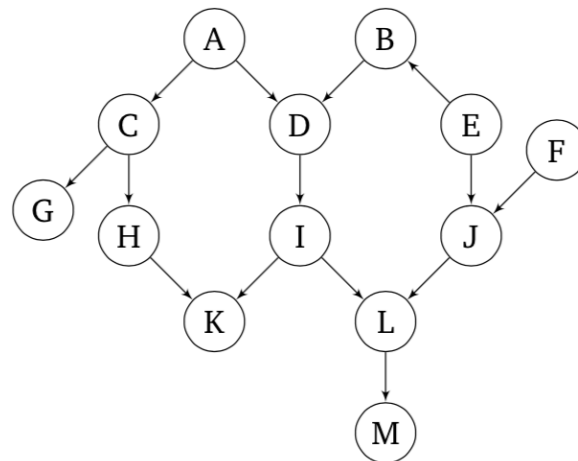
Prove or refute whether

(a) (i) implies (ii),

(b) (ii) implies (i).

**Exercise 3.4:** Let  $S_1, S_2$  be answer sets of a program  $P$ . Show that  $S_1 \subseteq S_2$  implies  $S_1 = S_2$ .

**Exercise 3.6:** Consider the following graph of a Bayesian network:



Answer the following questions and give an explanation of your answer.

- (i) Is  $B$  conditionally independent of  $G$ ?
- (ii) Is  $B$  conditionally independent of  $G$  given evidence  $M$ ?
- (iii) Which evidences are needed such that  $G$  is conditionally independent of  $M$ ?
- (iv) Which evidences are needed such that  $A$  is conditionally independent of  $L$ ?

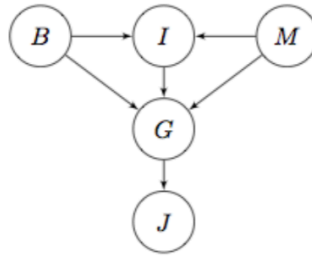
**Exercise 3.7:** Given the following Bayesian network with the Boolean variables

- $B = \text{BrokeElectionLaw}$ ,
- $I = \text{Indicted}$ ,
- $M = \text{PoliticallyMotivatedProsecutor}$ ,
- $G = \text{FoundGuilty}$ , and
- $J = \text{Jailed}$ ,

where the probabilities for  $B$  and  $M$  are  $P(B) = 0.9$  and  $P(M) = 0.1$ , respectively. The conditional probabilities for  $P(G|B, I, M)$ ,  $P(J|G)$ , and  $P(I|B, M)$  are as follows:

$B$	$M$	$P(I B, M)$
0	0	0.1
0	1	0.5
1	0	0.5
1	1	0.9

$G$	$P(J G)$
1	0.9
0	0.0



$B$	$I$	$M$	$P(G B, I, M)$
0	0	0	0.0
0	0	1	0.0
0	1	0	0.1
0	1	1	0.2
1	0	0	0.0
1	0	1	0.0
1	1	0	0.8
1	1	1	0.9

- (i) Which, if any, of the following relations hold in view of the *network structure alone* (ignoring the conditional probability tables for now)?
- $P(B, I, M) = P(B)P(I)P(M)$ .
  - $P(J|G) = P(J|G, I)$ .
  - $P(M|G, B, I) = P(M|G, B, I, J)$ .
- (ii) Calculate  $P(b, i, \neg m, g, j)$ .
- (iii) Calculate the probability that someone goes to jail given that he or she broke the law, has been indicted, and faces a politically motivated prosecutor.