

Planning, Learning and Decision Making

Lecture 11. POMDPs (conc.)

Representing $J^{(k)}$

- The cost-to-go at each iteration of VI is always PWLC
- Can always be written in the form

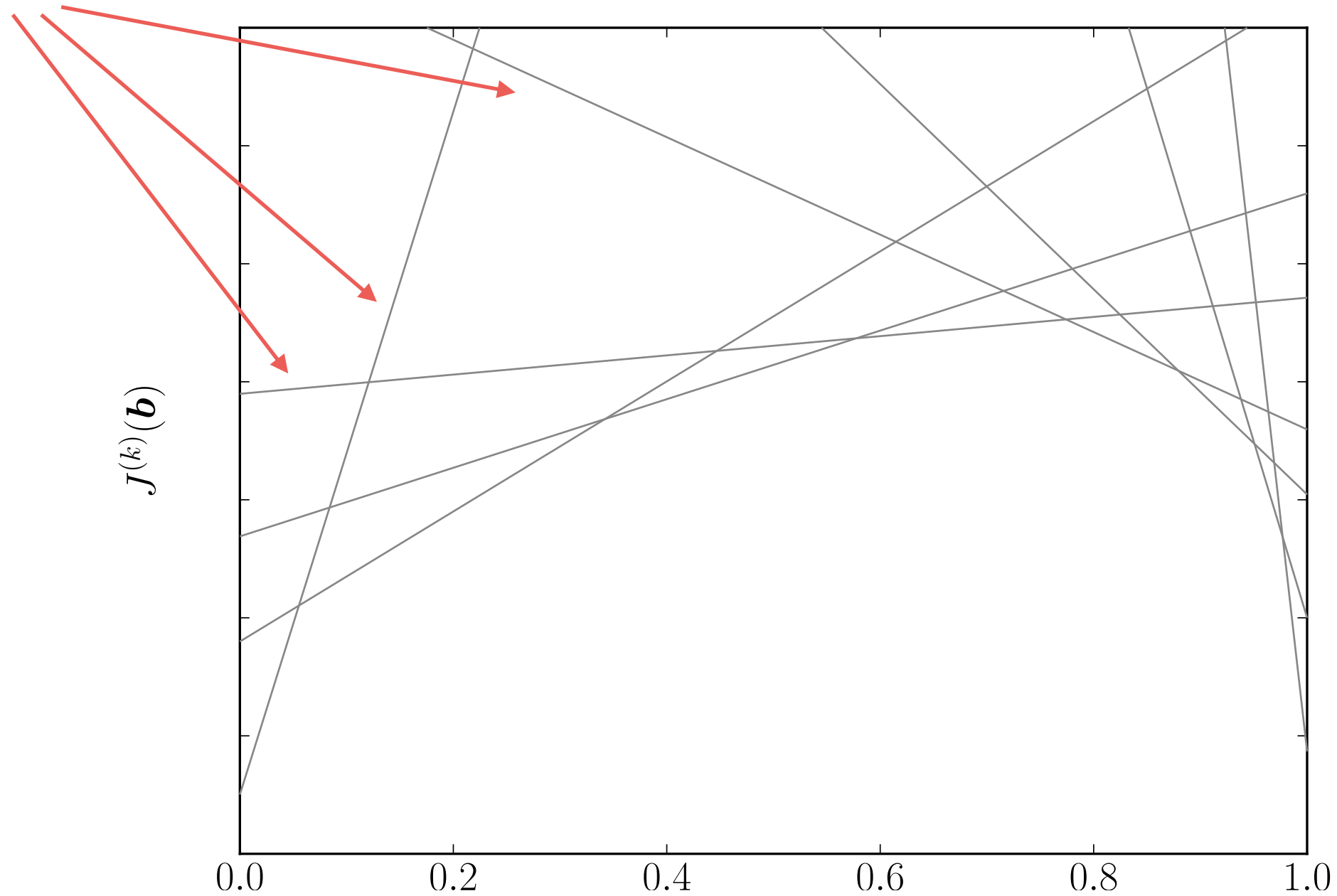
$$\begin{aligned} J^{(k)}(\mathbf{b}) &= \min_{\alpha \in \Gamma} \mathbf{b} \cdot \alpha \\ &= \min_{\alpha \in \Gamma} \sum_{x \in \mathcal{X}} b(x) \alpha(x) \end{aligned}$$

Set of vectors
used in the
representation



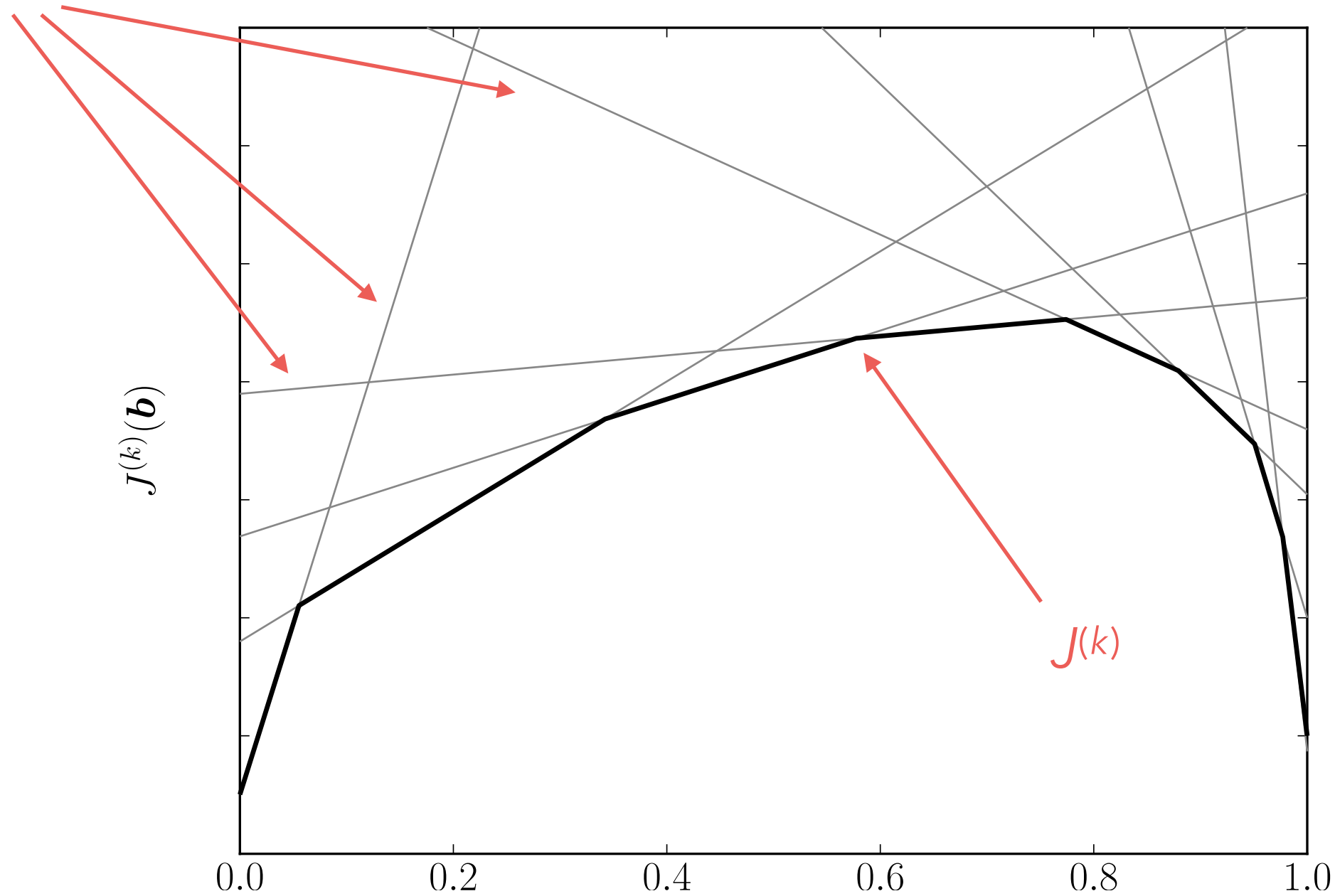
Representing $J^{(k)}$

α -vectors



Representing $J^{(k)}$

α -vectors



Value iteration

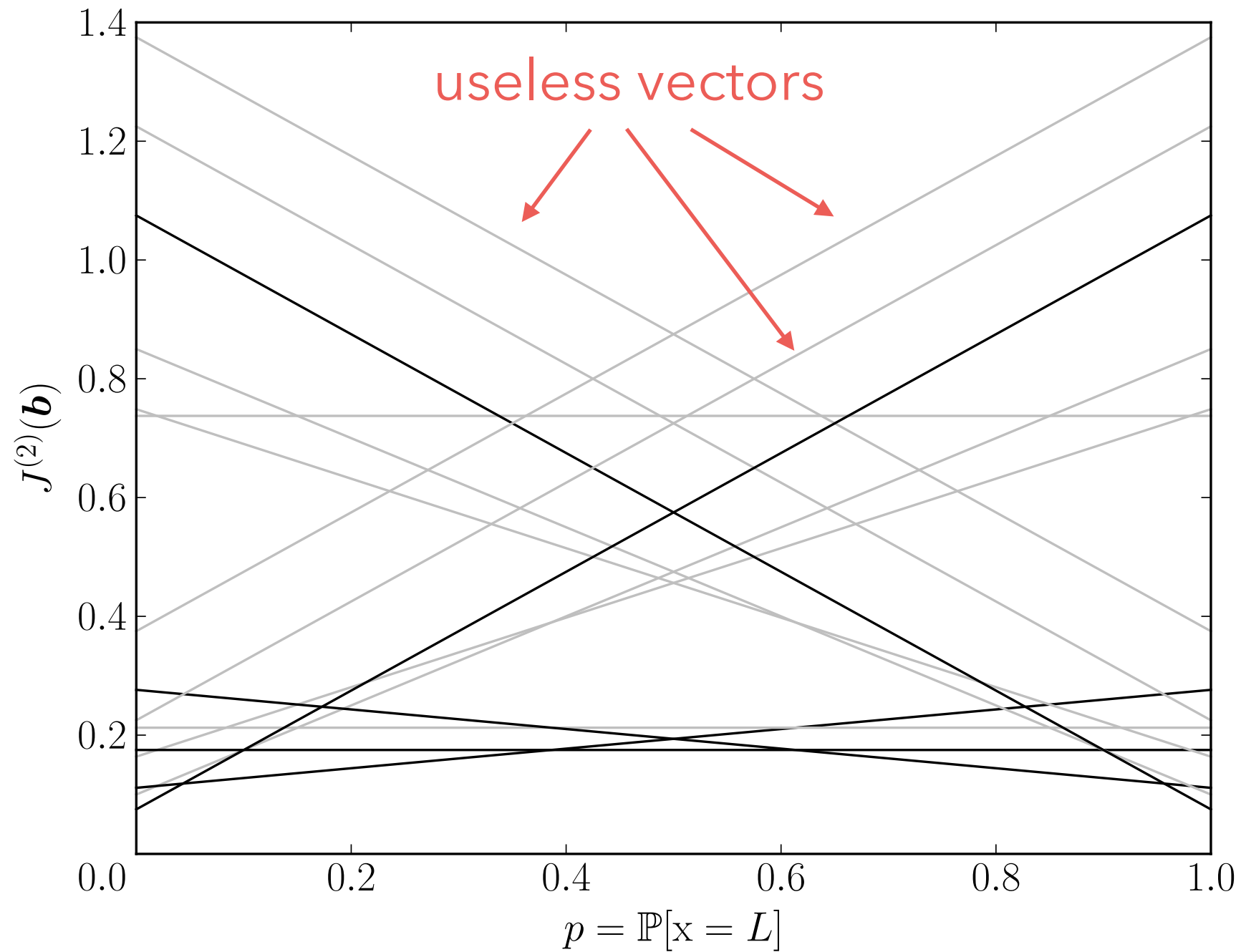
- Compute, at each iteration $k + 1$, the set $\Gamma^{(k+1)}$ from $\Gamma^{(k)}$
 - For each $\alpha \in \Gamma^{(k)}$, compute

$$\alpha_{a,z}^{(k)} = \frac{1}{|\mathcal{Z}|} C_{:,a} + \gamma P_a \text{diag}(\mathbf{O}_{z,a}) \alpha$$

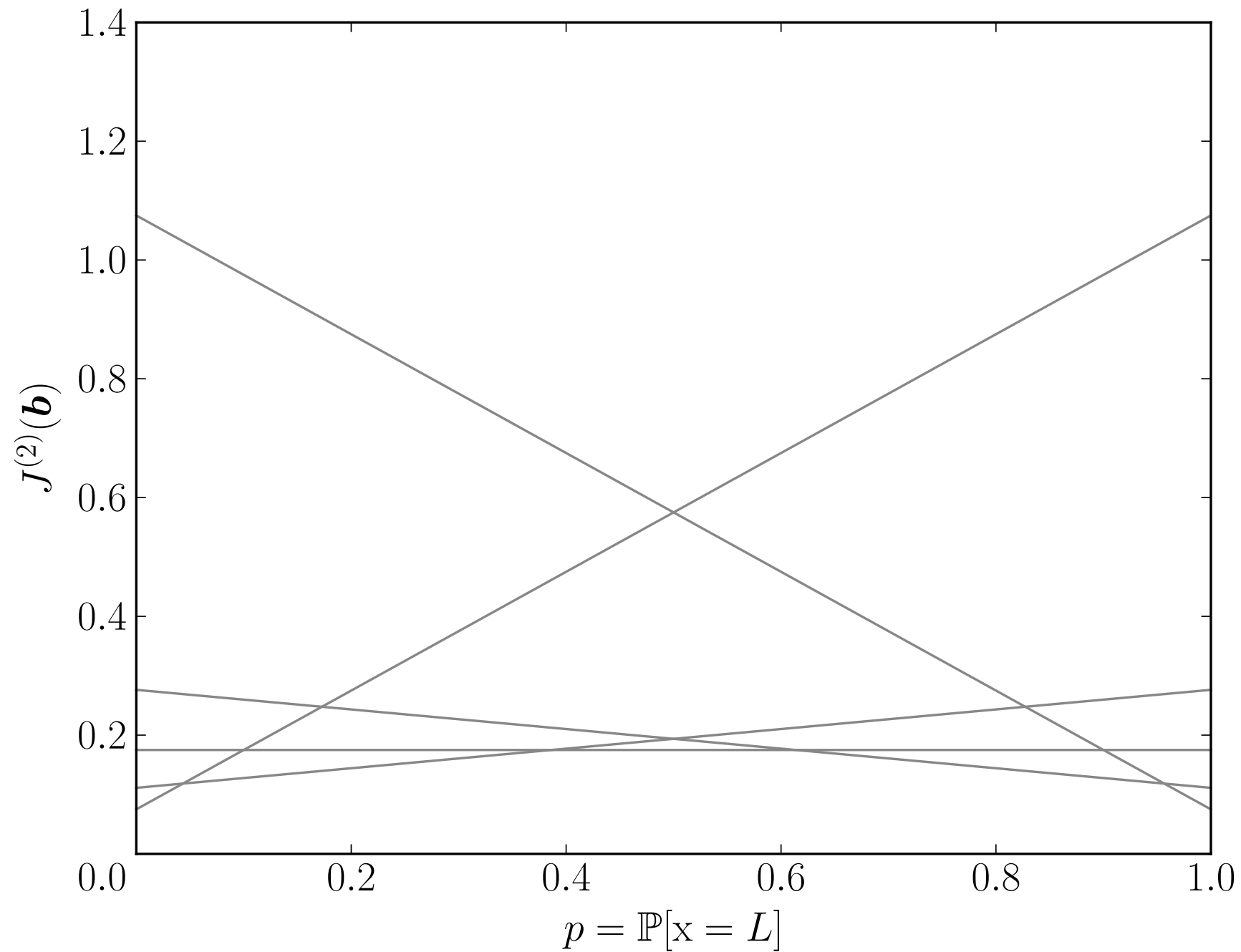
- Compute **all possible combinations** of $\alpha_{a,z}^{(k)}$, for each z
- For each combination, let

$$\alpha_a^{(k)} = \sum_{z \in \mathcal{Z}} \alpha_{z,a}^{(k)}$$

Value iteration



Value iteration



Value iteration

- Two approaches to build $\Gamma^{(k+1)}$ from $\Gamma^{(k)}$:
- **Region based methods:** Start with empty $\Gamma^{(k+1)}$ and only add vectors that are necessary



A vector is necessary if
it represents J in a non
empty belief region
(witness region)

Value iteration

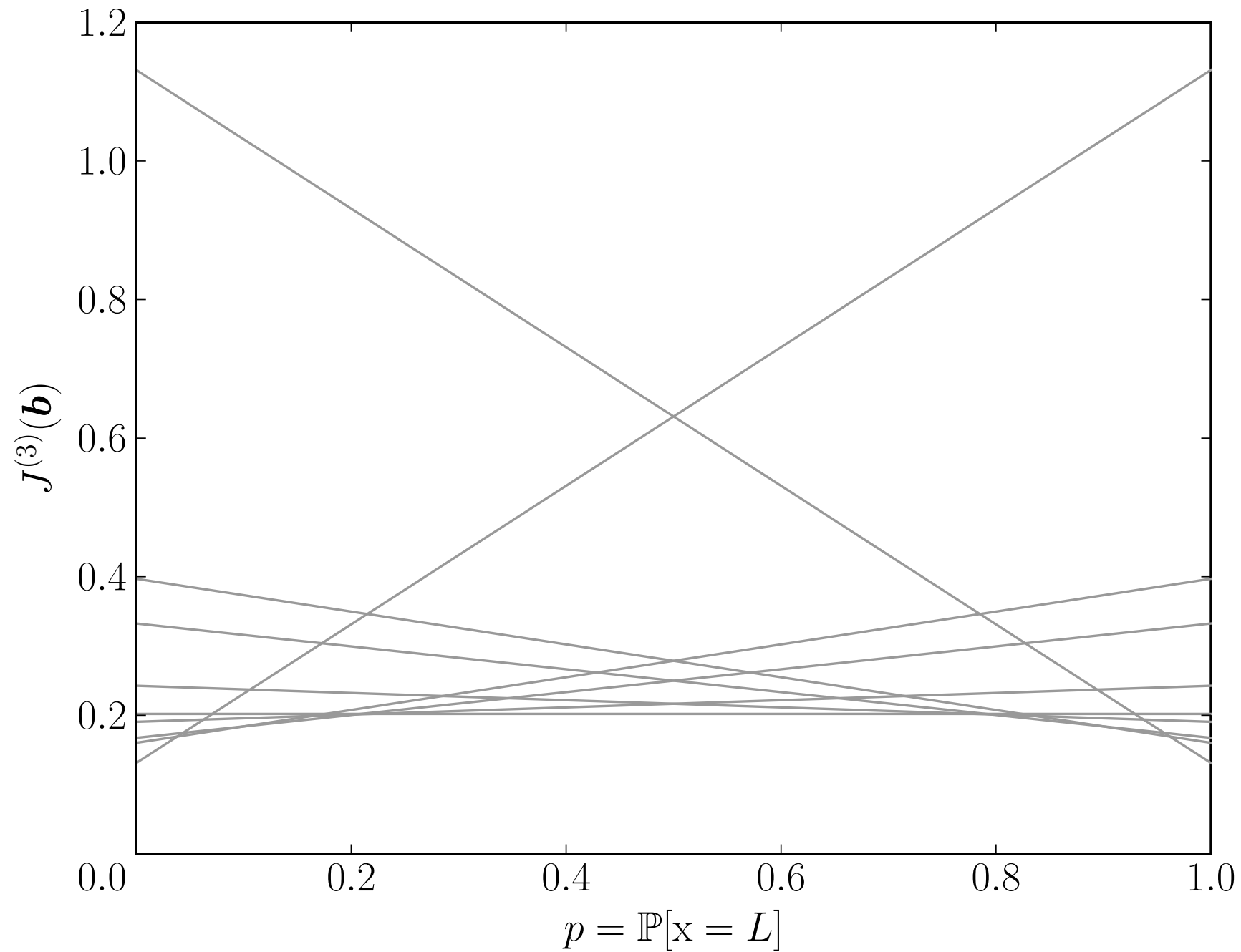
- Two approaches to build $\Gamma^{(k+1)}$ from $\Gamma^{(k)}$:
 - **Region based methods:** Start with empty $\Gamma^{(k+1)}$ and only add vectors that are necessary

Example: Witness algorithm

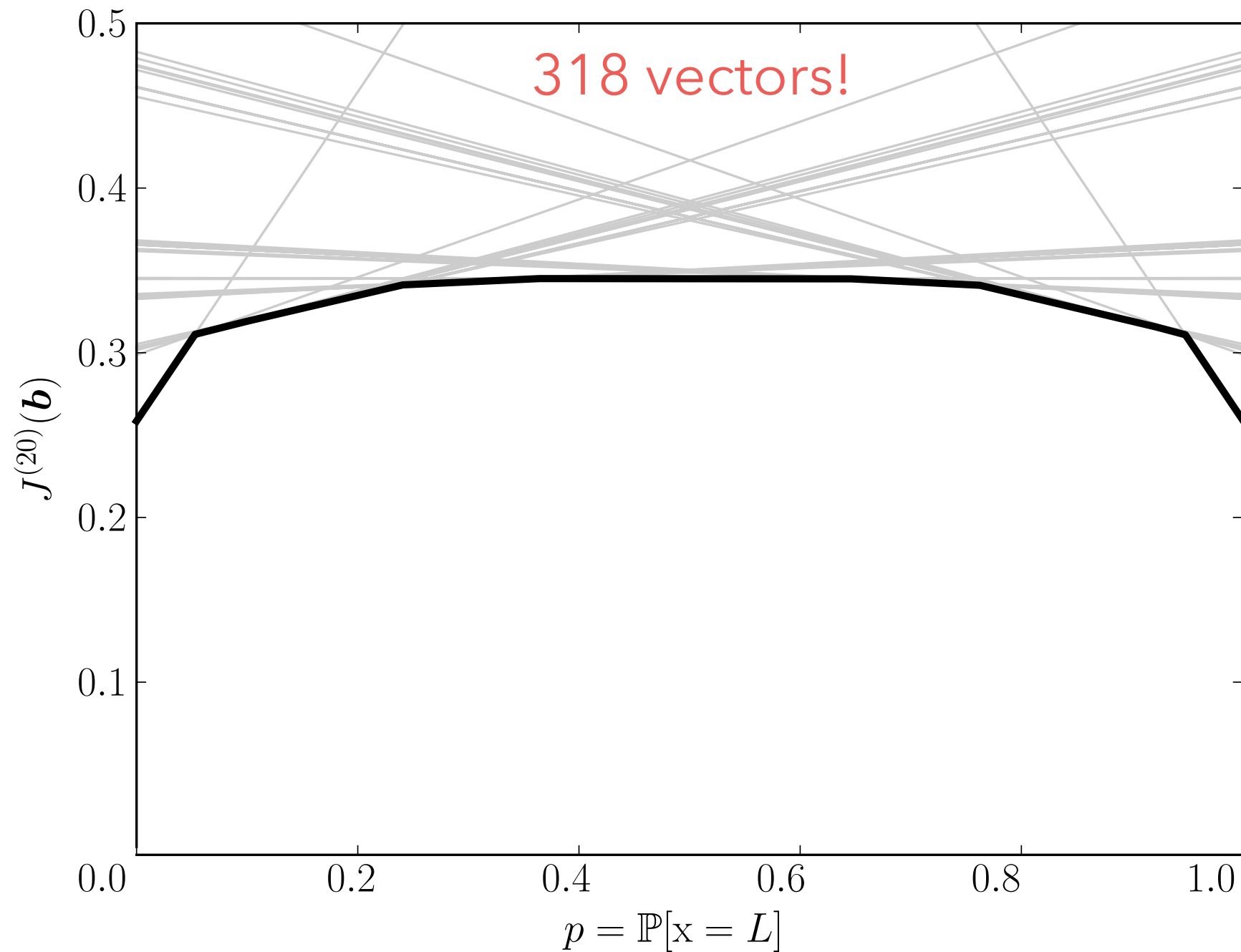
- **Pruning-based methods:** Start with complete $\Gamma^{(k+1)}$ and remove vectors that are unnecessary

Example: Incremental pruning

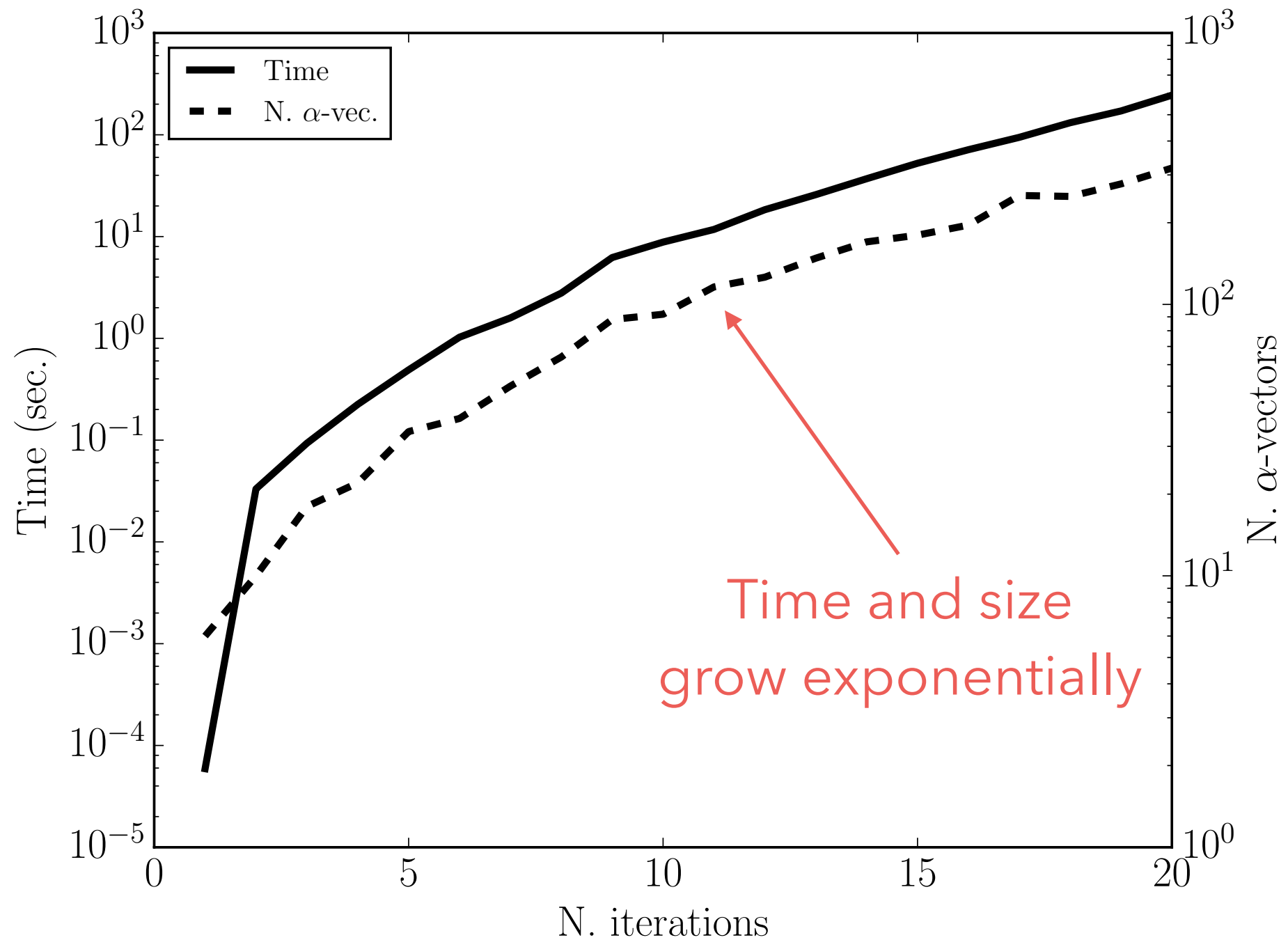
Value iteration



Value iteration



Computation time



Non-exact solutions

Idea n. 1 - Use the MDP

- MLS heuristic:

$$\pi_{\text{MLS}}(\mathbf{b}) = \pi_{\text{MDP}}(\operatorname{argmax}_{x \in \mathcal{X}} \mathbf{b}(x))$$

- AV heuristic:

$$\pi_{\text{AV}}(\mathbf{b}) = \operatorname{argmax}_{a \in \mathcal{A}} \sum_{x \in \mathcal{X}} \mathbf{b}(x) \mathbb{I}(a = \pi_{\text{MDP}}(x))$$

- Q-MDP heuristic:

$$\pi_{\text{Q-MDP}}(\mathbf{b}) = \operatorname{argmin}_{a \in \mathcal{A}} \sum_{x \in \mathcal{X}} \mathbf{b}(x) Q_{\text{MDP}}(x, a)$$

- FIB heuristic:

$$\pi_{\text{FIB}}(\mathbf{b}) = \operatorname{argmax}_{a \in \mathcal{A}} \sum_{x \in \mathcal{X}} \mathbf{b}(x) Q_{\text{FIB}}(x, a)$$

Idea n. 1 - Use the MDP

- MLS heuristic:

$$\pi_{\text{MLS}}(\mathbf{b}) = \pi_{\text{MDP}}(\operatorname{argmax}_{x \in \mathcal{X}} \mathbf{b}(x))$$

- AV heuristic:

$$\pi_{\text{AV}}(\mathbf{b}) = \operatorname{argmax}_{a \in \mathcal{A}} \sum_{x \in \mathcal{X}} \mathbf{b}(x) \pi_{\text{MDP}}(x, a)$$

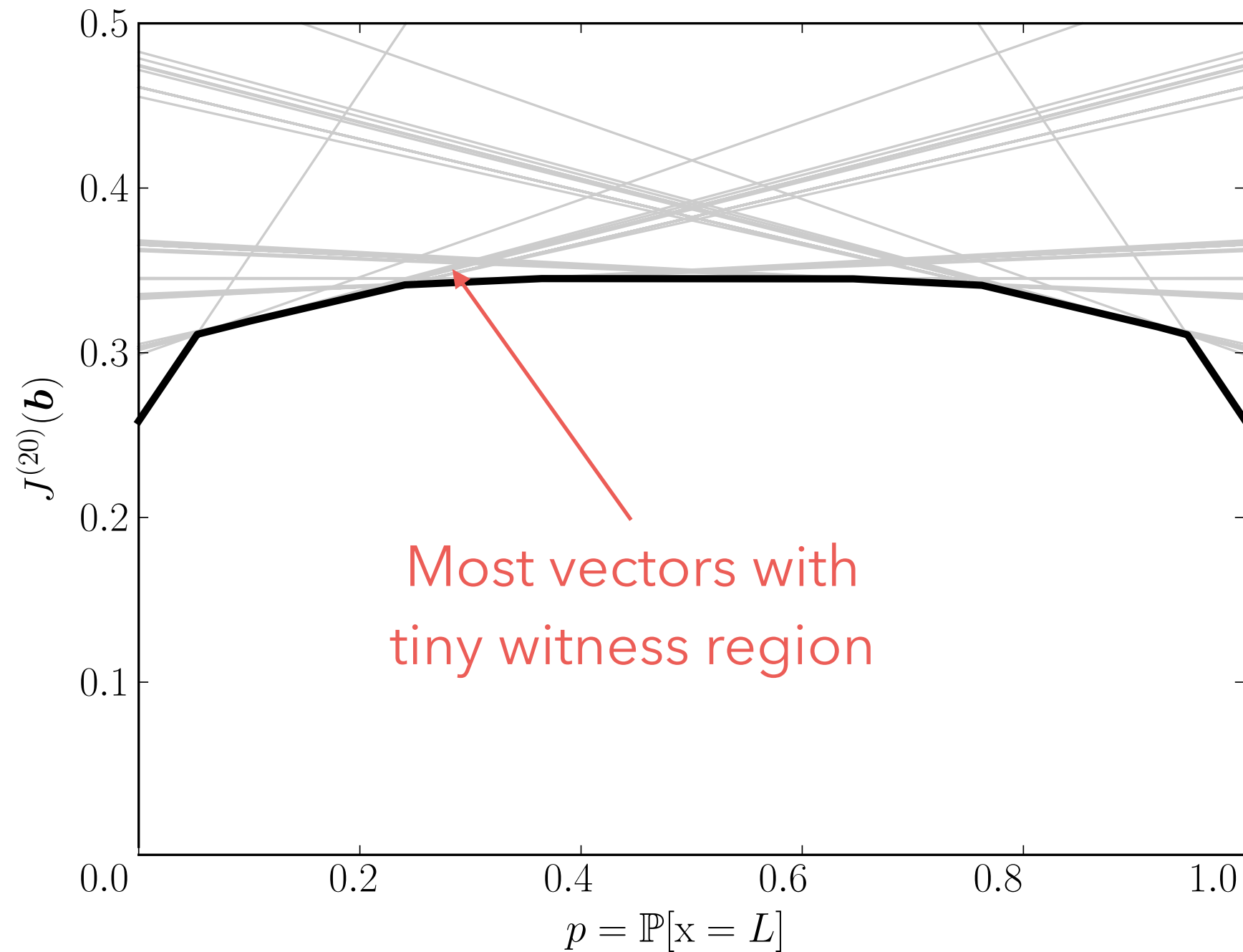
- Q-MDP heuristic:

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- FIB heuristic:

$$\pi_{\text{FIB}}(\mathbf{b}) = \operatorname{argmax}_{a \in \mathcal{A}} \sum_{x \in \mathcal{X}} \mathbf{b}(x) Q_{\text{FIB}}(x, a)$$

Idea n. 2

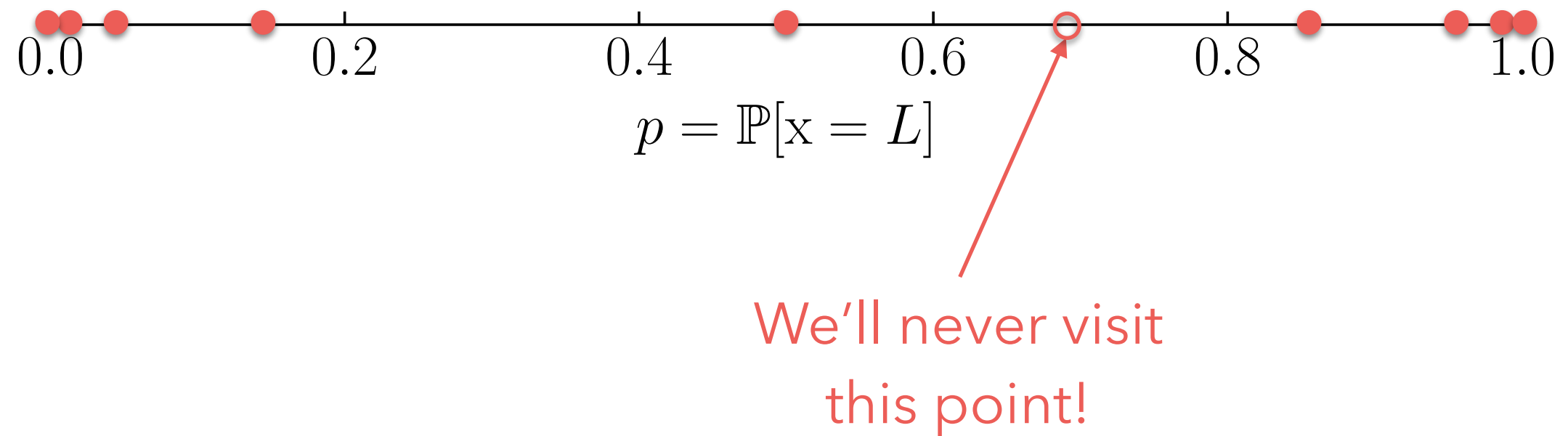


Idea n. 2

- Most α -vectors play little role in representing J
- What if we only compute the vectors that “truly matter”?

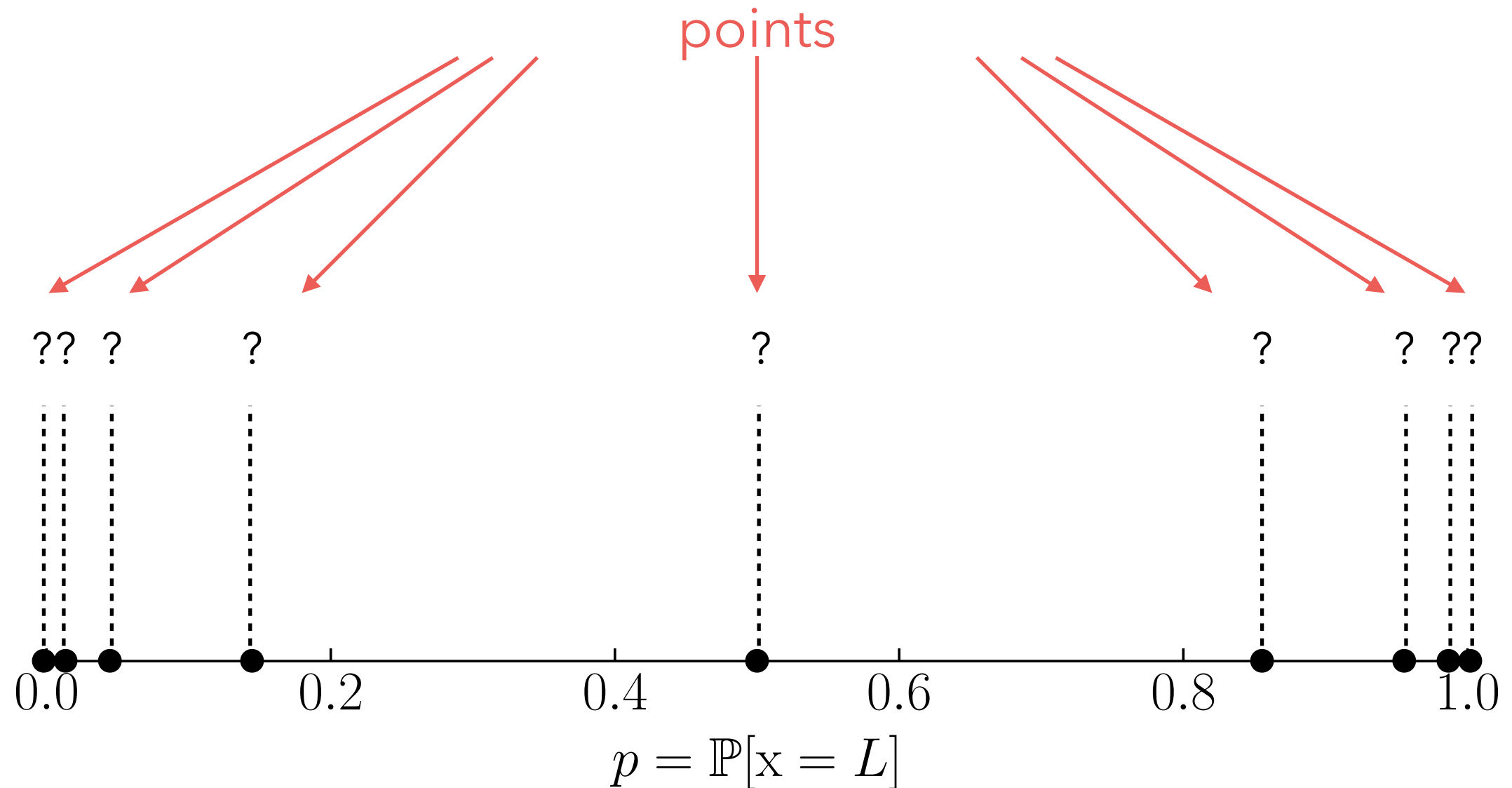
Which ones?

Idea n. 2



Idea n. 2

We only care about
the vectors in these
points



Point-based methods

- Select a finite set $\mathcal{B}_{\text{sample}}$ of beliefs to perform update
- For each belief, compute the corresponding α -vector

$$\alpha(\mathbf{b}) = \min_{a \in \mathcal{A}} \left[\mathbf{C}_{:,a} + \gamma \sum_{z \in \mathcal{Z}} P_a \text{diag}(\mathbf{O}_{z,a}) \min_{\alpha \in \Gamma} \alpha \cdot \mathbf{b}'_{za} \right]$$

Updated
belief



Point-based methods

- Select a finite set $\mathcal{B}_{\text{sample}}$ of beliefs to perform update
 - For each belief, compute the corresponding α -vector

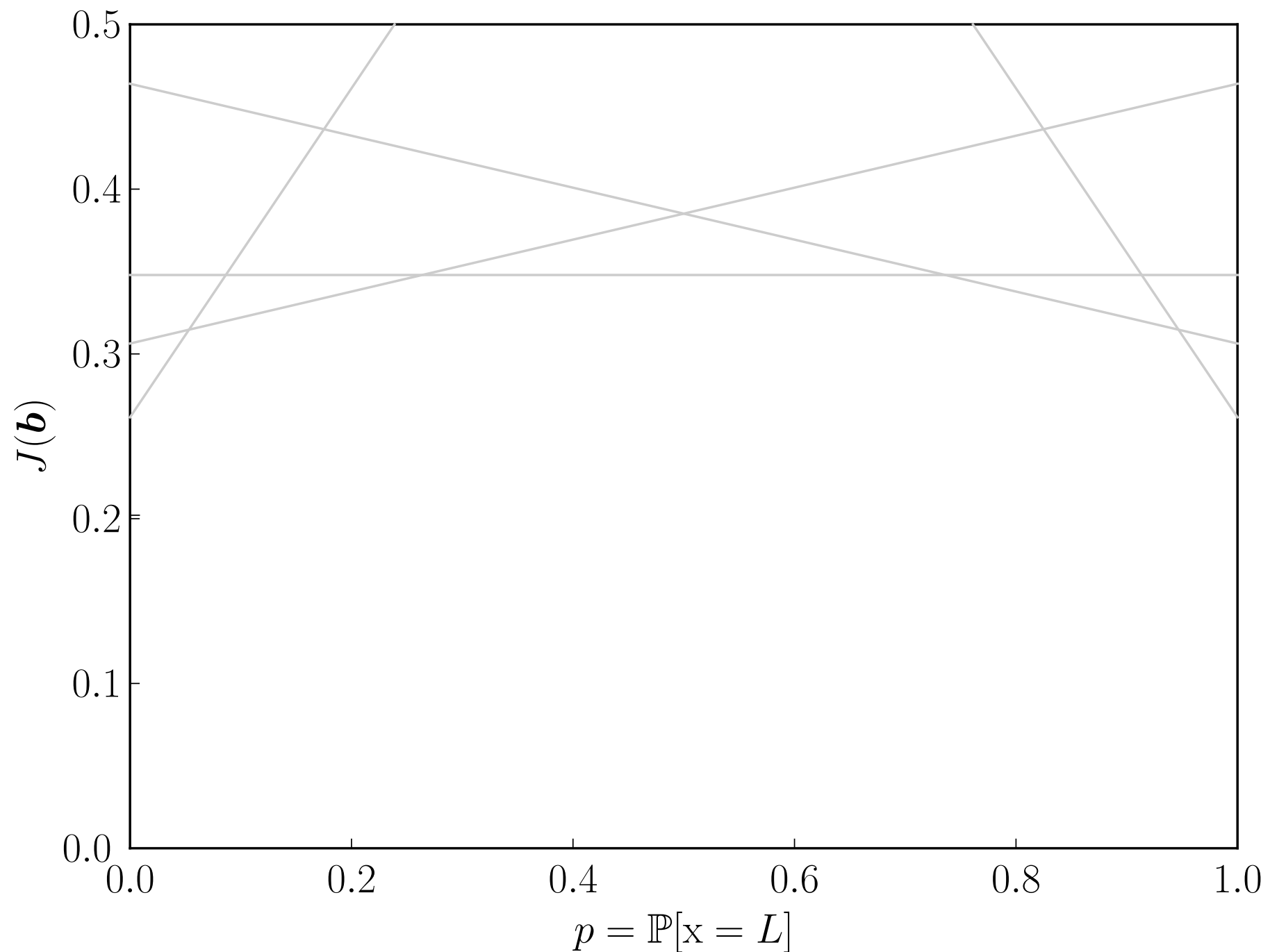
$$\alpha(\mathbf{b}) = \min_{a \in \mathcal{A}} \left[\mathbf{C}_{:,a} + \gamma \sum_{z \in \mathcal{Z}} P_a \text{diag}(\mathbf{O}_{z,a}) \min_{\alpha \in \Gamma} \alpha \cdot \mathbf{b}'_{za} \right]$$

- If necessary, rebuild $\mathcal{B}_{\text{sample}}$

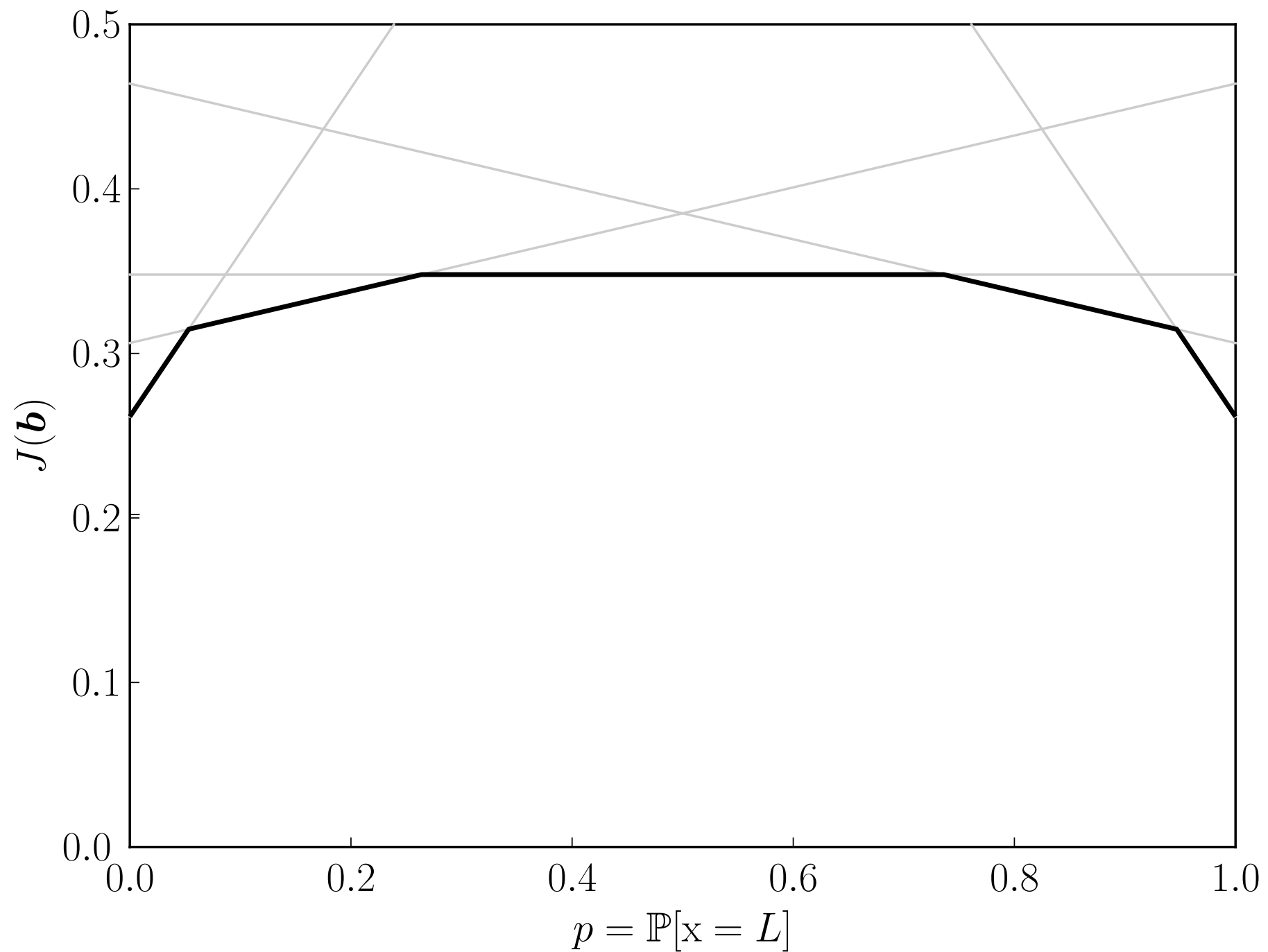
Point-based methods

- Many point-based methods:
 - PBVI (Pineau et al., 2003)
 - Perseus (Spaan & Vlassis, 2005)
 - HSVI (Smith & Simmons, 2005)
 - FSVI (Shani et al., 2007)
 - SARSOP (Kurniawati et al., 2008)
 - GapMin (Poupart et al., 2011)
- Much code available

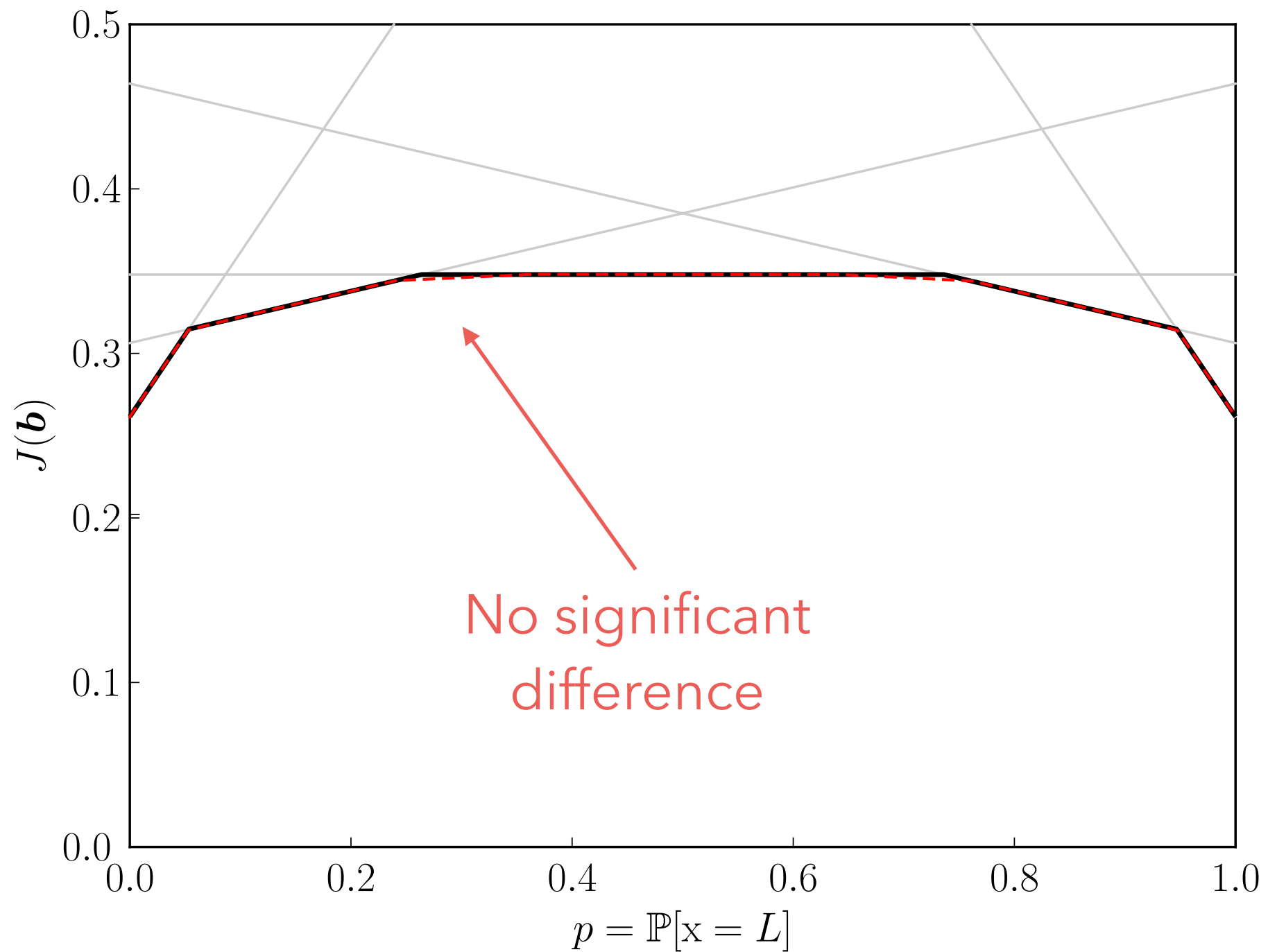
Point-based methods



Point-based methods



Point-based methods



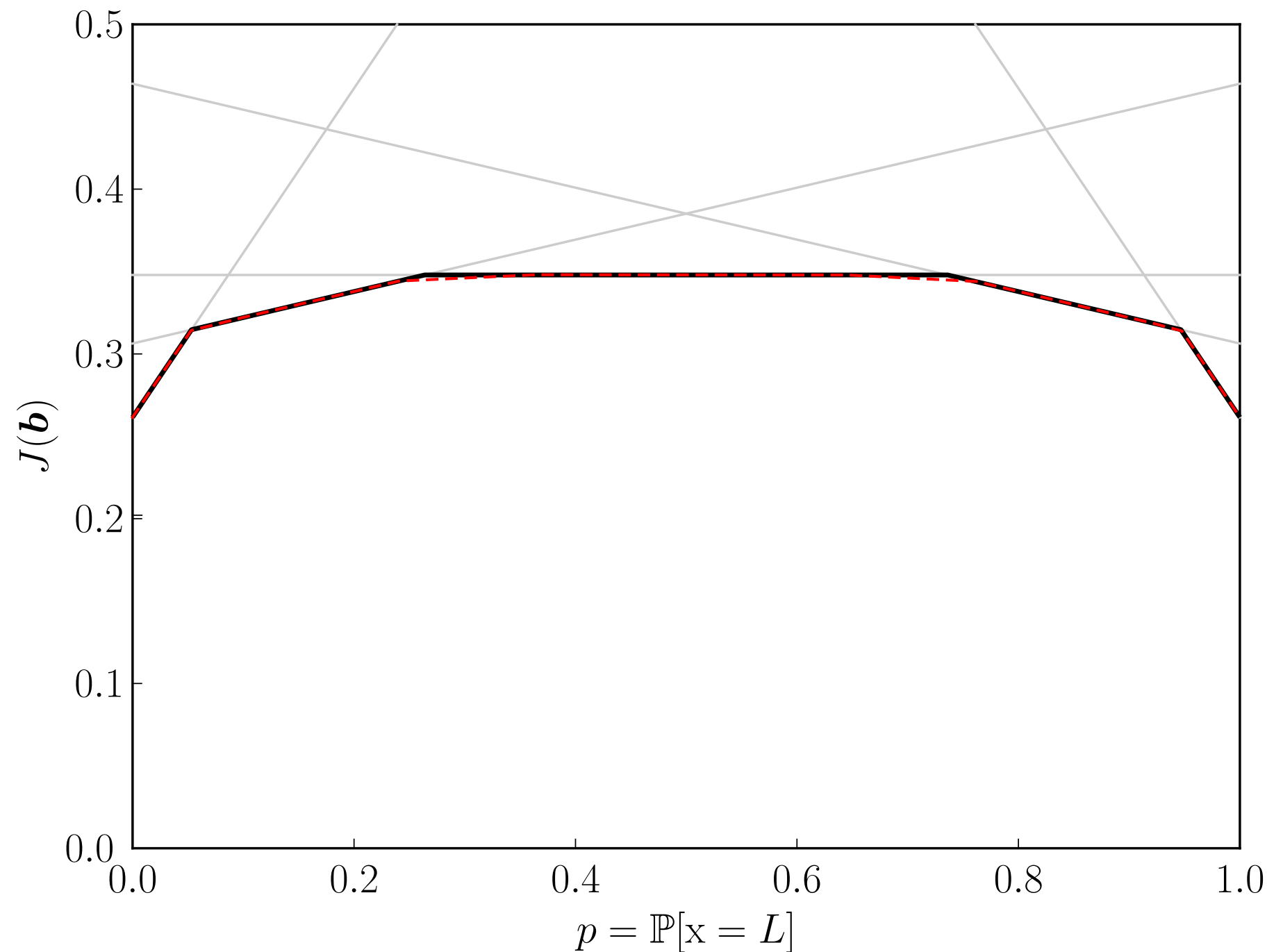
Point-based methods

VI:

- 318 vectors
- ~4 minutes

PERSEUS:

- 5 vectors
- 226 ms



What about policy iteration?

Policy iteration?

- Value iteration for POMDPs
 - How do we represent a cost-to-go function?
 - At each iteration of VI, the cost-to-go is PWLC
- Policy iteration for POMDPs
 - How do we represent a policy?

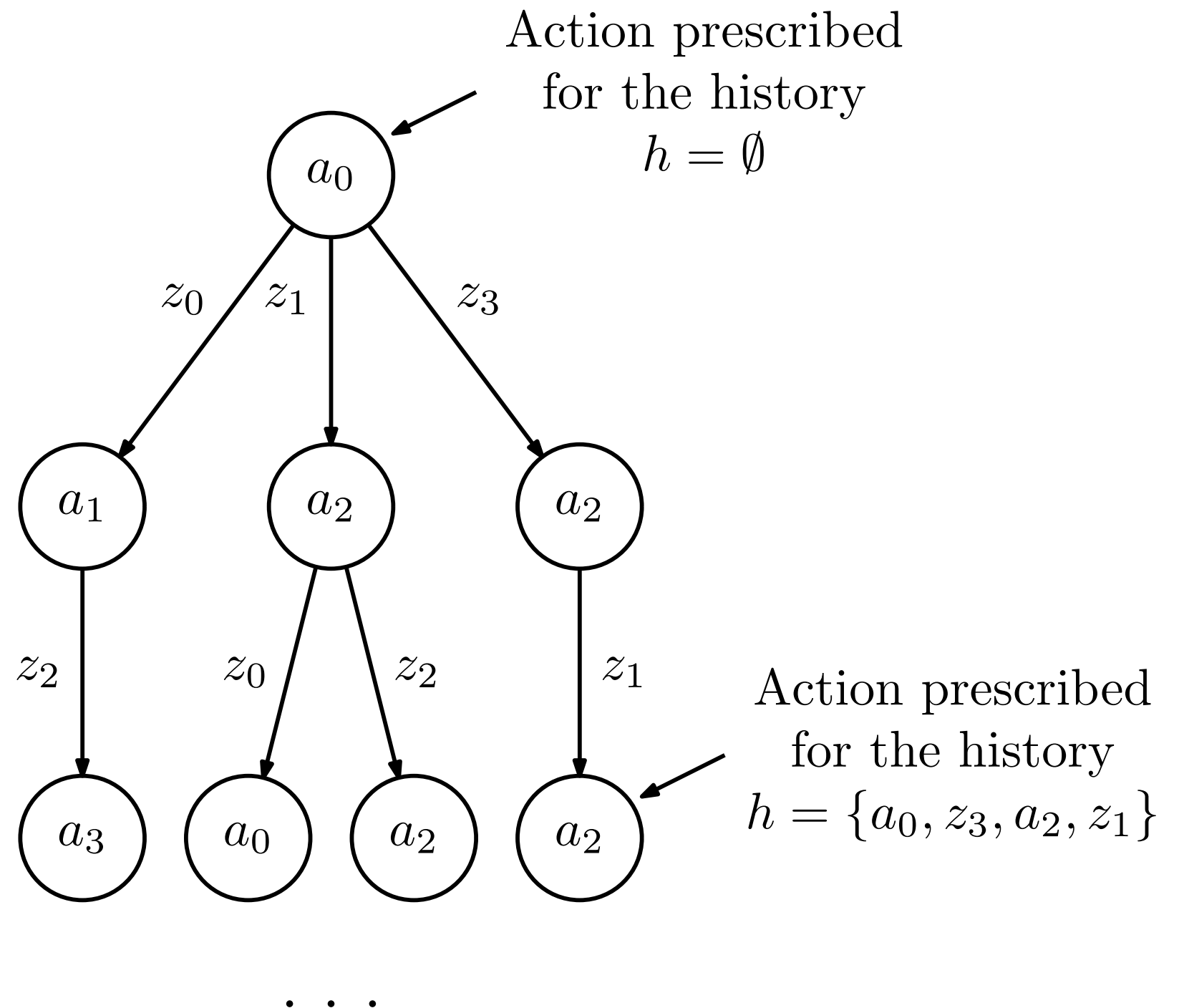
Policy graphs

- How can we represent a POMDP policy?
 - Compute it in runtime from J
 - Alternatively, we can consider policies as mapping **histories** to **actions**

Policy graphs

- We can represent the possible histories in a **policy tree**
 - Each node contains the **action** for that history
 - Branches correspond to **observations** from the node

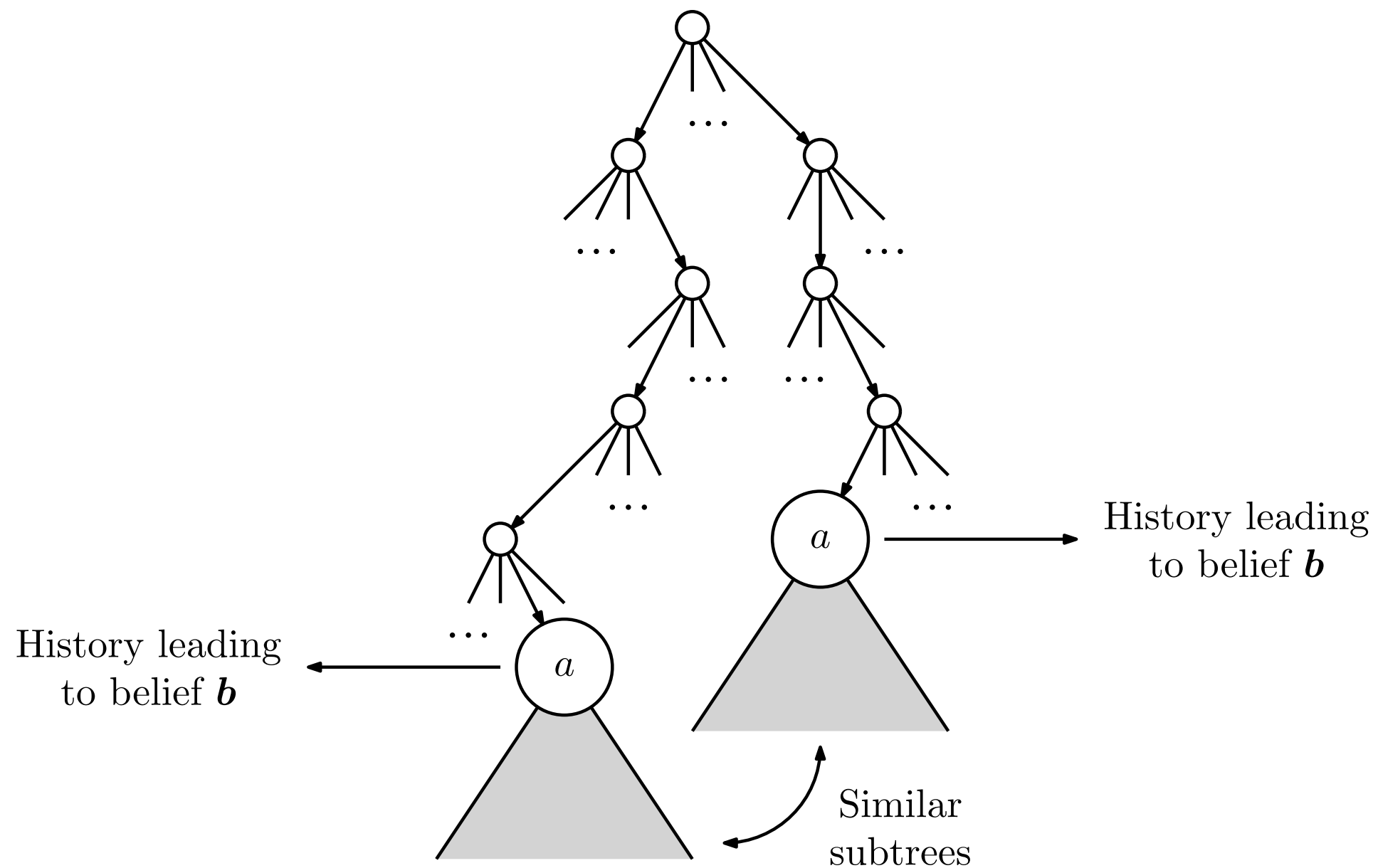
Policy graphs



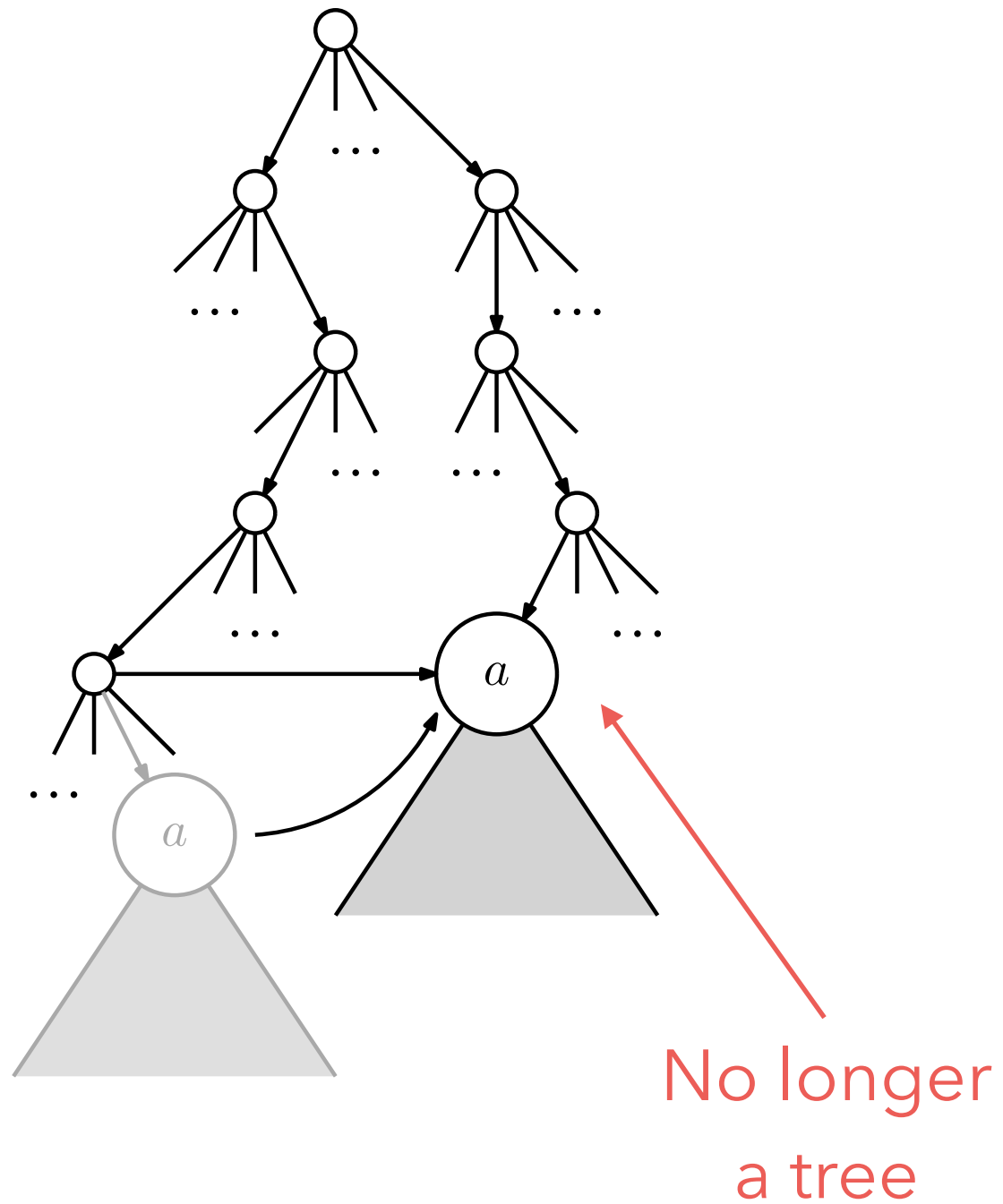
Policy graphs

- There is a lot of redundancy in policy trees
 - Histories leading to the same belief will have equivalent subtrees

Policy graphs



Policy graphs

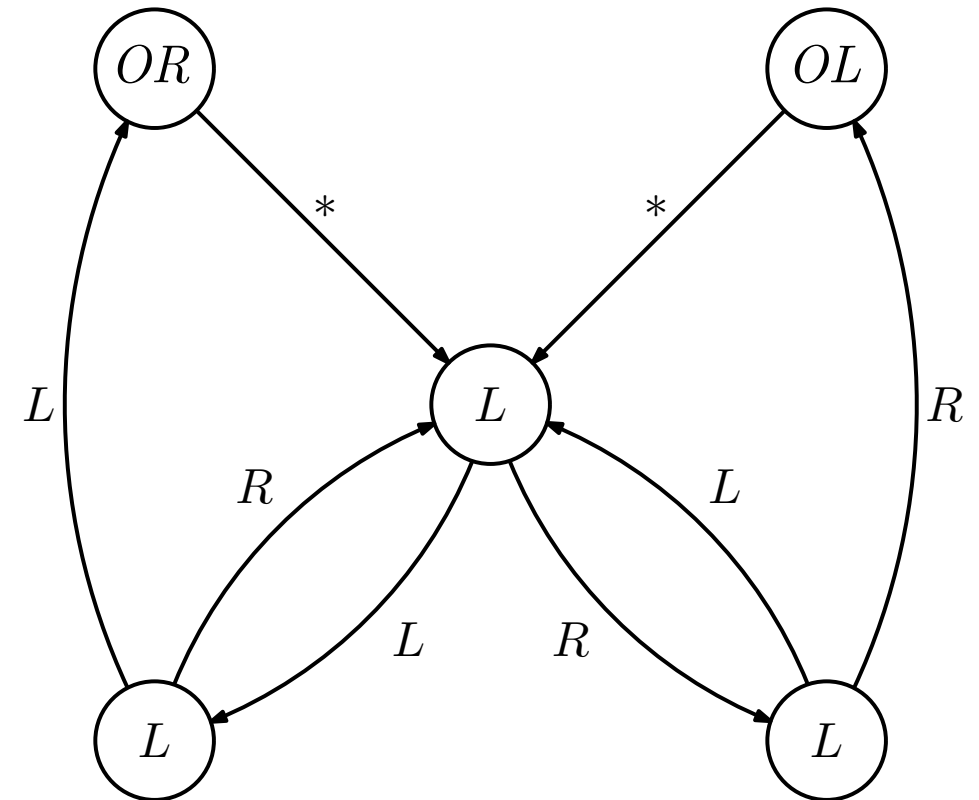
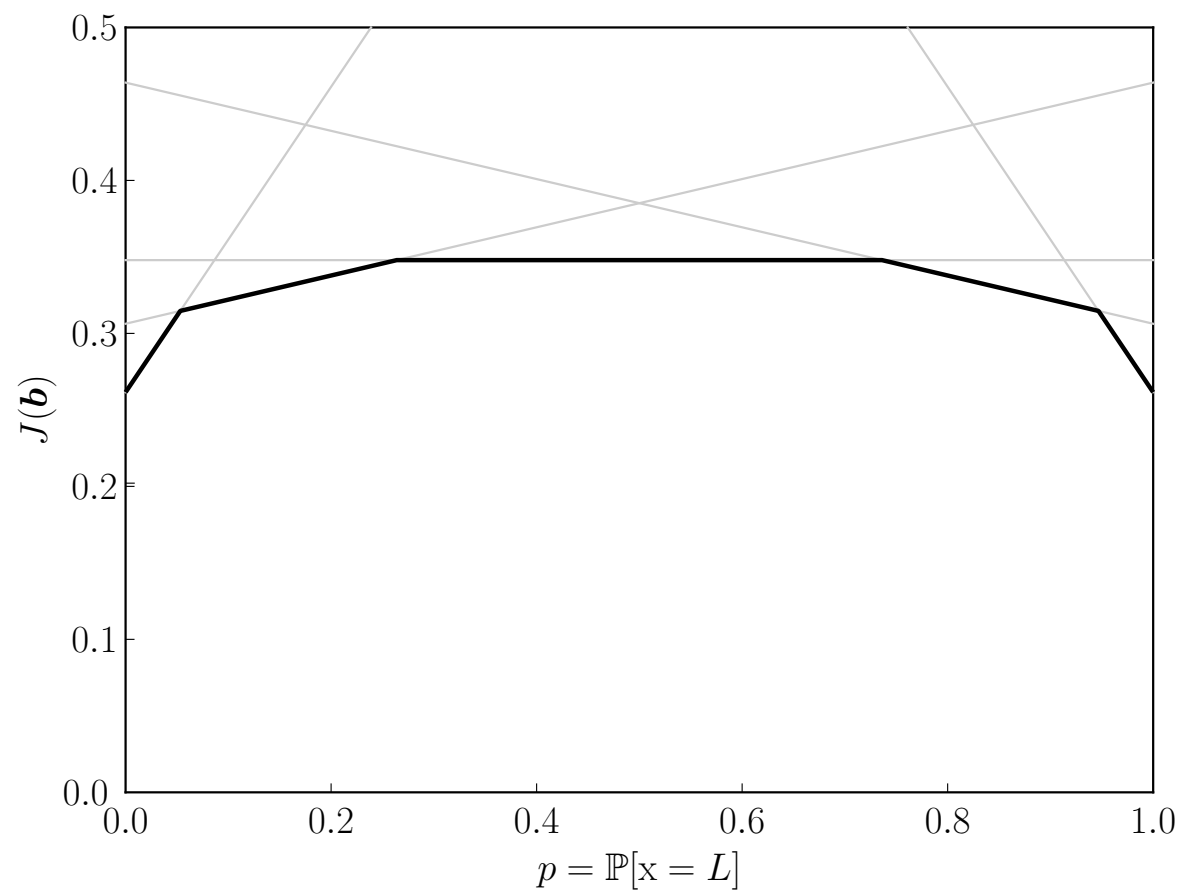


Policy graphs

- Policy graphs provide convenient representations for POMDP policies
- Close relation between policy graphs and α -vector
 - Each node in the graph corresponds to an α -vector and vice-versa

Policy graphs

- Example:



Key points about POMDPs

- Very general model for decision making under uncertainty
- Very hard to solve
- Beliefs provide a **summary** of the history
- POMDP \longleftrightarrow Belief MDP
 - We can use VI \rightarrow Cost-to-go is PWLC (finite representation)
 - We can use PI \rightarrow Policy graphs w/ finite number of nodes
- Approximate methods:
 - MDP heuristics
 - Point-based methods

Comments on complexity

How hard is an MDP?

- MDPs can be solved by a linear program
 - LP is known to be **polynomial-time (P)**
 - MDPs are solvable in **polynomial-time (P)**

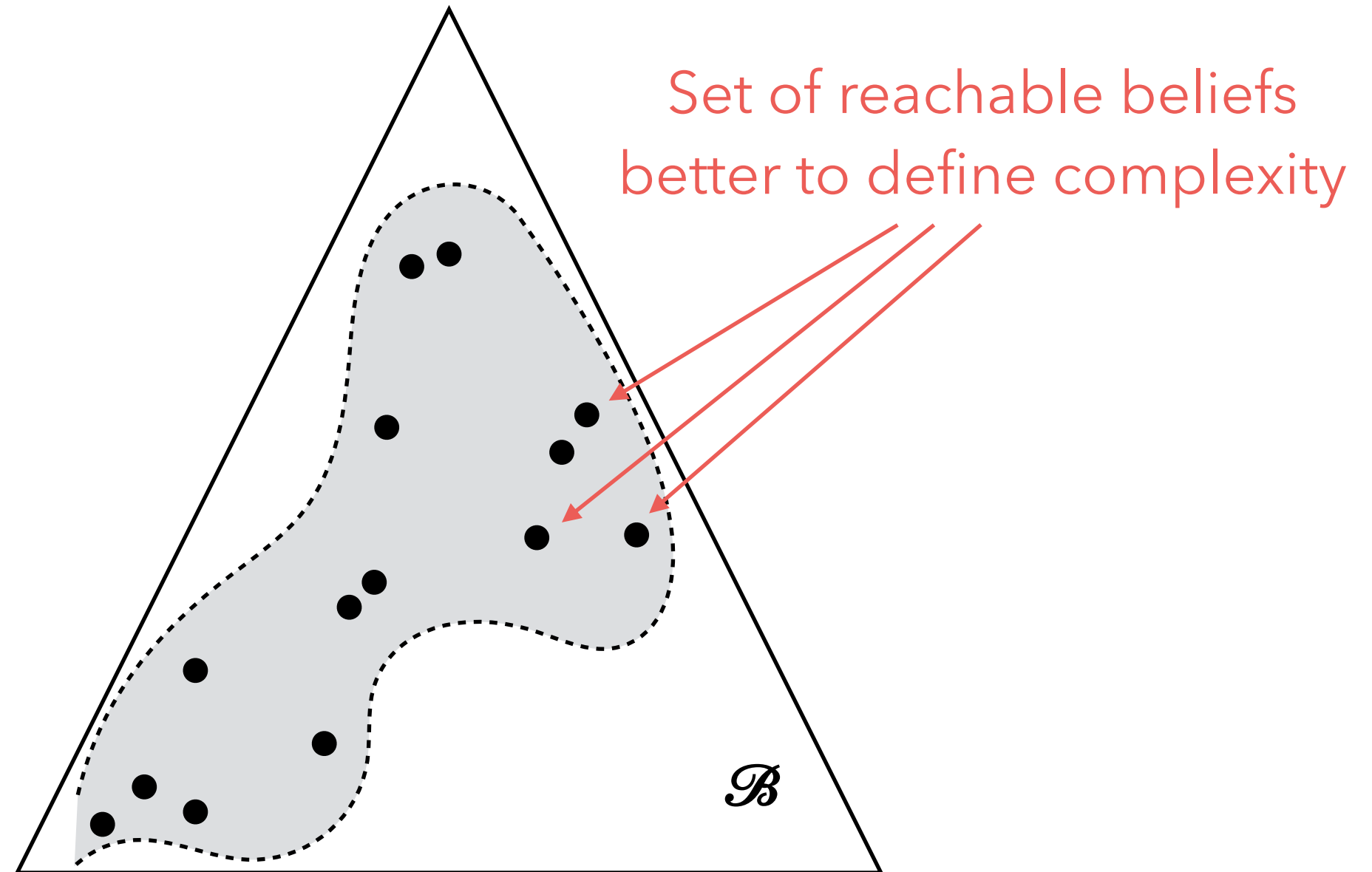
How hard is a POMDP?

- Infinite horizon POMDPs are **undecidable**
- Finite-horizon POMDPs are **PSPACE-complete**
 - ... little hope for exact solution methods
- POMDPs are **non-approximable**
 - ... in the worst case, you can't even guarantee a good approximation!

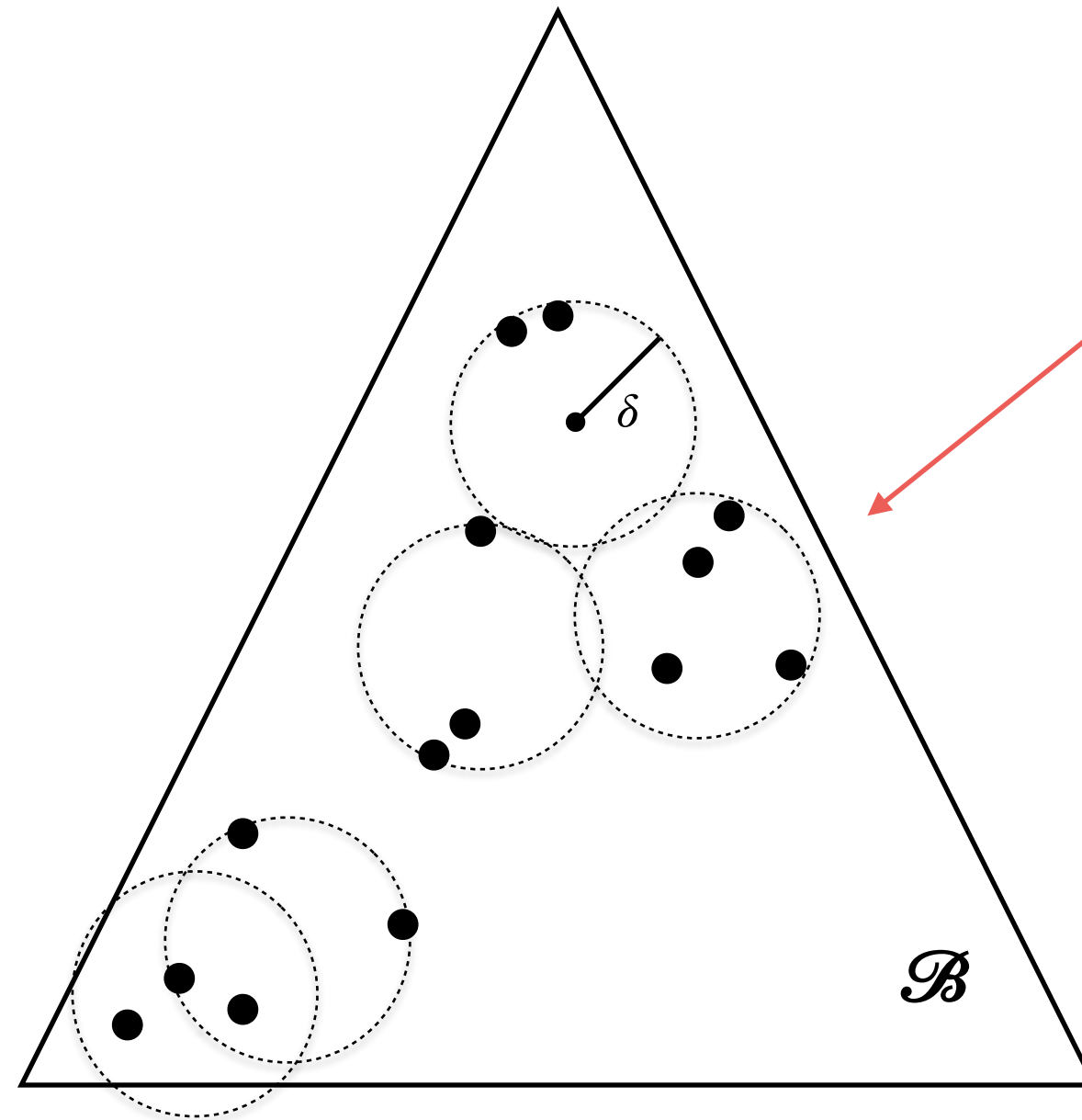
How hard is a POMDP?

- However, point-based methods work quite well!

How hard is a POMDP?

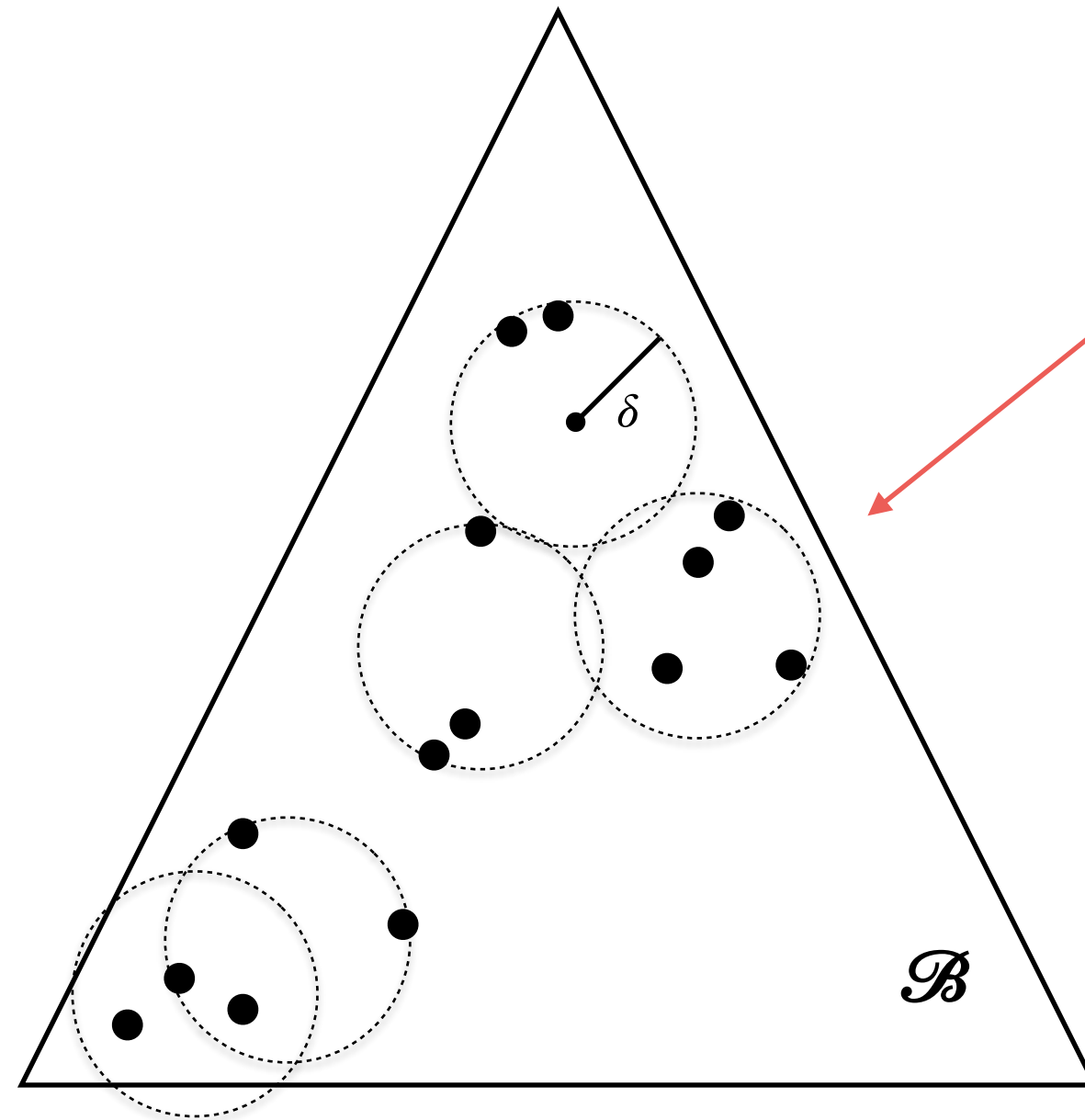


How hard is a POMDP?



Need 5 balls or
radius δ to "cover"
reachable points

How hard is a POMDP?



Covering number
 $C(\delta) = 5$

How hard is a POMDP?

- Complexity of POMDP planning better captured by the covering number of the reachable belief space
- Some point-based methods are built on such argument (they sample beliefs to cover reachable space)
 - Ex: SARSOP (Kurniawati et al., 2008)