

# FAI LAB 7

## First-Order logic

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- **Objects:** *Paolo, Roberto, 3*
- **Relations:** *SurvivedPhD( $\cdot$ ), Paper( $\cdot$ ), Wrote( $\cdot, \cdot$ ), Coauthors( $\cdot, \cdot$ )*
- **Functions:** *CV( $\cdot$ ), PublicationDate( $\cdot$ ), ErdosNumber( $\cdot$ ), ( $\cdot + \cdot$ ), ...*

*SurvivedPhD(Paolo)  $\wedge$  Coauthors(Paolo, Roberto)*  
*|CV(Roberto)| > |CV(Paolo)|*

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**Variables and quantifiers** we can compactly describe knowledge over many/infinite objects:

$\forall x. \forall y. [Coauthor(x, y) \leftrightarrow Coauthor(y, x)]$

$\forall x. [HasPhD(x) \rightarrow \exists y. (Paper(y) \wedge AuthorOf(x, y))]$

$\forall x. \forall y. [Coauthors(x, y) \rightarrow \exists z. (AuthorOf(x, z) \wedge AuthorOf(y, z))]$

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- *Paul Erdős' coauthors' EN is one.*

$$\forall x . [Coauthors(Erdos, x) \wedge (EN(x) = 1)]$$

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# Exercises: NL to FOL

Write the following natural language sentence in FOL using the following predicates/functions: *IsFriend*( $\cdot, \cdot$ ), *IsEnemy*( $\cdot, \cdot$ ), *Hates*( $\cdot, \cdot$ )

- *Karen has a friend who has a friend who hates Karen.*
- *Hate might not be mutual.*
- *The enemy of an enemy is a friend.*
- *People can't have friends if they hate themselves.*

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- *Karen has a friend who has a friend who hates Karen.*

$$\exists x . \exists y . [IsFriend(x, Karen) \wedge IsFriend(y, x) \wedge Hates(y, Karen)]$$

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$$\forall x . \forall y . \forall z . [(IsEnemy(x, z) \wedge IsEnemy(y, x)) \rightarrow IsFriend(y, z)]$$

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- *The enemy of an enemy is a friend.*

$$\forall x . \forall y . \forall z . [(IsEnemy(x, z) \wedge IsEnemy(y, x)) \rightarrow IsFriend(y, z)]$$

- *People can't have friends if they hate themselves.*

$$\forall x . [Hates(x, x) \rightarrow \forall y . \neg IsFriend(y, x)]$$

# Exercises: NL to FOL

Write the following natural language sentence in FOL using the following predicates/functions: *Movie(·)*, *FeaturedIn(·, ·)*, *Loves(·, ·)*, *Year(·)*

- *Alicia loves every movie featuring Julia.*
- *Alicia and Julia are featured in a movie together.*
- *Alicia doesn't love every movie filmed in the 80s.*
- *Julia never featured in a movie before 1987.*

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$$\forall x . [(Movie(x) \wedge FeaturedIn(Julia, x)) \rightarrow Loves(Alicia, x)]$$

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$$\exists x . [Movie(x) \wedge FeaturedIn(Alicia, x) \wedge FeaturedIn(Julia, x)]$$

- *Alicia doesn't love every movie filmed in the 80s.*

$$\exists x . [Movie(x) \wedge (Year(x) \in [1980, 1989]) \wedge \neg Loves(Alicia, x)]$$

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- *Julia never featured in a movie before 1987.*

$$\neg \exists x . [Movie(x) \wedge (Year(x) < 1987) \wedge FeaturedIn(Julia, x)]$$

# Exercises: FOL semantics

- The formula  $\exists x . (x + x = x)$  is satisfiable.
- The formula  $\exists x . (x + x = x)$  is valid.
- The formula  $\neg \exists x . (x + x = x)$  is satisfiable.
- The formula  $\neg \exists x . (x + x = x)$  is valid.
- The formula  $\exists x . (x < x)$  is satisfiable.
- The formula  $\exists x . [P(x) \wedge \neg P(x)]$  is satisfiable.
- The formula  $\forall x . [P(x) \vee \neg P(x)]$  is satisfiable.
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- The formula  $\exists x . (x + x = x)$  is satisfiable. **True**
- The formula  $\exists x . (x + x = x)$  is valid.
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$$\exists x . (x < x)$$



$$\exists x . (x < x)$$

$$\exists x . (x \therefore x)$$

$$\exists x . (x \triangle x)$$

$$\exists x . (x < x)$$

$$\exists x . (x \therefore x)$$

$$\exists x . (x \triangleleft x)$$

$$\exists x . P(x, x)$$