

FAI LAB 7

First-Order logic

Paolo Morettin

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- **Objects:** *Paolo, Roberto, 3*
- **Relations:** *SurvivedPhD(·), Paper(·), Wrote(·, ·), Coauthors(·, ·)*
- **Functions:** *CV(·), PublicationDate(·), ErdosNumber(·), (· + ·), ...*

SurvivedPhD(Paolo) \wedge Coauthors(Paolo, Roberto)
|CV(Roberto)| > |CV(Paolo)|

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Variables and quantifiers we can compactly describe knowledge over many/infinite objects:

$$\begin{aligned} \forall x . \forall y . [Coauthor(x, y) \leftrightarrow Coauthor(y, x)] \\ \forall x . [HasPhD(x) \rightarrow \exists y . (Paper(y) \wedge AuthorOf(x, y))] \\ \forall x . \forall y . [Coauthors(x, y) \rightarrow \exists z . (AuthorOf(x, z) \wedge AuthorOf(y, z))] \end{aligned}$$

Example: Erdős numbers

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Exercises: NL to FOL

Write the following natural language sentence in FOL using the following predicates/functions: *IsFriend(·, ·)*, *IsEnemy(·, ·)*, *Hates(·, ·)*

- *Karen has a friend who has a friend who hates Karen.*
- *Hate might not be mutual.*
- *The enemy of an enemy is a friend.*
- *People can't have friends if they hate themselves.*

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$$\exists x . \exists y . [IsFriend(x, Karen) \wedge IsFriend(y, x) \wedge Hates(y, Karen)]$$

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- *The enemy of an enemy is a friend.*

$$\forall x . \forall y . \forall z . [(IsEnemy(x, z) \wedge IsEnemy(y, x)) \rightarrow IsFriend(y, z)]$$

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$$\exists x . \exists y . [\text{IsFriend}(x, \text{Karen}) \wedge \text{IsFriend}(y, x) \wedge \text{Hates}(y, \text{Karen})]$$

- *Hate might not be mutual.*

$$\exists x . \exists y . [\text{Hates}(x, y) \wedge \neg \text{Hates}(y, x)]$$

- *The enemy of an enemy is a friend.*

$$\forall x . \forall y . \forall z . [(\text{IsEnemy}(x, z) \wedge \text{IsEnemy}(y, x)) \rightarrow \text{IsFriend}(y, z)]$$

- *People can't have friends if they hate themselves.*

$$\forall x . [\text{Hates}(x, x) \rightarrow \forall y . \neg \text{IsFriend}(y, x)]$$

Exercises: NL to FOL

Write the following natural language sentence in FOL using the following predicates/functions: *Movie(·)*, *FeaturedIn(·, ·)*, *Loves(·, ·)*, *Year(·)*

- *Alicia loves every movie featuring Julia.*
- *Alicia and Julia are featured in a movie together.*
- *Alicia doesn't love every movie filmed in the 80s.*
- *Julia never featured in a movie before 1987.*

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$$\exists x . [\text{Movie}(x) \wedge (\text{Year}(x) \in [1980, 1989]) \wedge \neg \text{Loves}(\text{Alicia}, x)]$$

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- *Julia never featured in a movie before 1987.*

$$\neg \exists x . [\text{Movie}(x) \wedge (\text{Year}(x) < 1987) \wedge \text{FeaturedIn}(\text{Julia}, x)]$$

Exercises: FOL semantics

- The formula $\exists x . (x + x = x)$ is satisfiable.
- The formula $\exists x . (x + x = x)$ is valid.
- The formula $\neg \exists x . (x + x = x)$ is satisfiable.
- The formula $\neg \exists x . (x + x = x)$ is valid.
- The formula $\exists x . (x < x)$ is satisfiable.
- The formula $\exists x . [P(x) \wedge \neg P(x)]$ is satisfiable.
- The formula $\forall x . [P(x) \vee \neg P(x)]$ is satisfiable.
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- The formula $\exists x . (x + x = x)$ is satisfiable. **True**
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FOL semantics: a warning

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$$\exists x . P(x, x)$$