

Martynas Kemežys IF-8/1

(C8230)

1. Palyginti

0.	$f(n) = a, g(n) = \frac{100}{\ln n}, a = \text{const}$
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$$\lim_{n \rightarrow \infty} \frac{a}{\frac{100}{\ln n}} = \frac{a}{100} \lim_{n \rightarrow \infty} \ln n = \frac{a}{100} \cdot \infty = \infty$$

$$\text{Ita: } a = \Omega\left(\frac{100}{\ln n}\right)$$

2. Suprastinti

3.	$O(n \ln^2 n + \sqrt[6]{n^4})$
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$$\lim_{n \rightarrow \infty} \frac{n \ln^2 n}{\sqrt[6]{n^4}} = \lim_{n \rightarrow \infty} \frac{n \cdot \ln^2 n}{n^{\frac{4}{6}}} = \lim_{n \rightarrow \infty} \frac{n \cdot \ln^2 n}{n^{\frac{2}{3}}} =$$

$$= \lim_{n \rightarrow \infty} n^{\frac{1}{3}} \ln^2 n = \lim_{n \rightarrow \infty} n^{\frac{1}{3}} \cdot \lim_{n \rightarrow \infty} \ln^2 n =$$

=

$$\lim_{n \rightarrow \infty} n^{\frac{1}{3}} = \infty$$

$$= \infty \cdot \infty = \infty$$

$$\lim_{n \rightarrow \infty} \ln^2 n = \infty$$

$$\text{Ita: } O(n \ln^2 n + \sqrt[6]{n^4}) = O(n \ln^2 n)$$

3. Ivertinti sumą

2.	$\sum_{i=1}^m \ln^3 i$
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$$\begin{aligned}
 \int_0^m \ln^3 x dx &= x \ln^3 x - \int x d \ln^3 x = x \ln^3 x - 3 \int \frac{x \cdot \ln^2 x}{x} dx = \\
 &= x \ln^3 x - 3 \int \ln^2 x dx = x \ln^3 x - 3x \ln^2 x + 3 \int x d \ln^2 x = \\
 &= x \ln^3 x - 3x \ln^2 x + 6 \int \ln x dx = x \ln^3 x - 3x \ln^2 x + 6x \ln x - \\
 &\quad \left(6 \int \frac{1}{x} dx \right) \Big|_0^m =
 \end{aligned}$$

Atg.

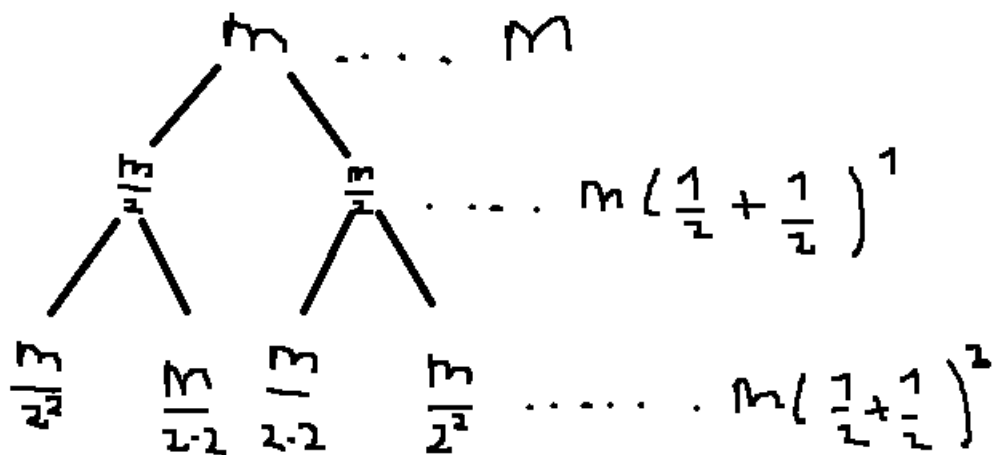
$$\begin{aligned}
 x \ln^3 x - 3x \ln^2 x + 6x \ln x - 6 \int \frac{1}{x} dx &\leq \sum_{i=0}^{m+1} \ln^3 i \leq \\
 &\leq x \ln^3 x - 3x \ln^2 x + 6x \ln x - 6 \int \frac{1}{x} dx
 \end{aligned}$$

4. Surasti lygties $T(n) = aT(n/b) + f(n)$ sprendinį, taikant medžio metodą:

0.	$T(n) = T(\sqrt{n}) + T(\sqrt{n-1}) + \log_6 n$
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Keičiame: $n = 6^m$

$$S(m) = S\left(\frac{m}{2}\right) + S\left(\frac{m}{2}\right) + m$$



$$S(m) = m \sum_{i=0}^h \left(\frac{1}{2} + \frac{1}{2}\right)^i$$

$$h = \log_2 m$$

$$S(m) < 2S\left(\frac{m}{2}\right) + m$$

$$S(m) \geq 2S\left(\frac{m}{2}\right) + m$$

5. Surasti lygties $T(n) = aT(n/b) + f(n)$ sprendinį, taikant pagrindinę teoremą:

3.	$T(n) = 3T\left(\frac{n}{3}\right) + \frac{\ln n}{n}$
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$$a = 3$$

$$b = 3$$

$$f(n) = \frac{\ln n}{n}$$

$$\lim_{n \rightarrow \infty} \frac{\ln n}{n^{\log_3 3 + \epsilon}} = \lim_{n \rightarrow \infty} \frac{\ln n}{n^{1+\epsilon}} =$$

$$= n^{1+\epsilon} = \infty, \text{ kai } 0 < \epsilon < 1$$

$$3 \ln \frac{n}{3} < \ln n$$

$$\ln n < c \ln n$$



$$1 < c < 1$$

$$\therefore T(n) = \Theta(n^{\log_3 3})$$

6. Įvertinti sudėtingumą:

2.	<pre> public void AA(int[] C, int m, int n) { int s = 0; int p = BB(n - m + 1); if (p < n - m) { AA(C, m, n - p); AA(C, m, n - p); for (int i = m; i <= n; i++) { C(i) = C(i) + 1; } AA(C, n - p + 1, n); AA(C, n - p + 1, n); } } public int BB(int k) { int s = 0; for (int i = 1; i <= k; i++) { for (int j = 1; j <= k; j++) { s = s + 1; } } return Sqrt(s) / 2; } </pre>
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2.	<pre> public void AA(int[] C, int m, int n) { int s = 0; int p = BB(n - m + 1); if (p < n - m) { AA(C, m, n - p); AA(C, m, n - p); for (int i = m; i <= n; i++) { C(i) = C(i) + 1; } AA(C, n - p + 1, n); AA(C, n - p + 1, n); } } public int BB(int k) { int s = 0; for (int i = 1; i <= k; i++) { for (int j = 1; j <= k; j++) { s = s + 1; } } return Sqrt(s) / 2; } </pre>	kaina	kiekis
	int s = 0;	c	1
	int p = BB(n - m + 1);	c	1
	if (p < n - m) {	c	1
	AA(C, m, n - p);	AA(m, n - p)	1
	AA(C, m, n - p);	AA(m, n - p)	1
	for (int i = m; i <= n; i++)	C	n - m + 2
	{		
	C(i) = C(i) + 1;	c	n - m + 1
	}		
	AA(C, n - p + 1, n);	AA(n - p + 1, n)	1
	AA(C, n - p + 1, n);	AA(m, n - p)	1
	}		
	}		
	public int BB(int k)		
	{		
	int s = 0;	c	1
	for (int i = 1; i <= k; i++)	c	k + 1
	{		
	for (int j = 1; j <= k; j++)	c	k
	{		
	s = s + 1;	c	k - 1
	}		
	}		
	return Sqrt(s) / 2;	c	1
	}		

$$\begin{aligned}
T(m, n) &= 3C + 3TAA(m, n-p) + \\
&+ C(n-m+1) + C(n-m+2) + \\
&+ TAA(n-p+1, n) = 3TAA(m, n-p) + \\
&+ TAA(n-p+1, n) + 2C(n-m+1.5) + 3C
\end{aligned}$$

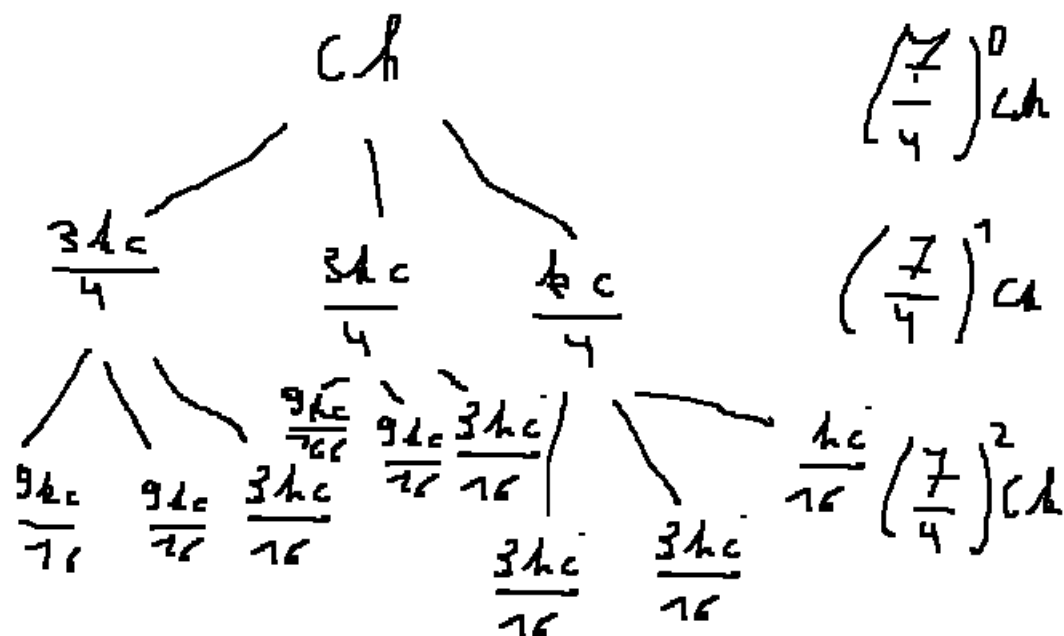
$$T(m, n) = T(k)$$

$$k = n - m + 1$$

$$p = \frac{n-m+1}{4} = \frac{k}{4}$$

$$\begin{aligned}
T(k) &= 3T(n-m-p+1) + T(p-1+1) + \\
&+ \underbrace{2C(n-m+1.5) + 3C}_{Ck}
\end{aligned}$$

$$T(k) = 3T\left(\frac{3k}{4}\right) + T\left(\frac{k}{4}\right) + Ck$$



$$\log_4 b \leq h \leq \log_{\frac{4}{3}} b$$

$$T(h) = ch \sum_{i=0}^h \left(\frac{7}{4}\right)^i = \frac{\left(\frac{7}{4}\right)^{h+1} - 1}{\frac{7}{4} - 1} = \frac{4 \left(\frac{7}{4}\right)^{h+1} - 4}{3}$$

$$4 \left(\frac{7}{4}\right)^{\log_4 b + 1} \leq T(h) \leq \frac{4 \left(\frac{7}{4}\right)^{\log_{\frac{4}{3}} b + 1} - 4}{3}$$

$$T(h) = \Theta(c \log^4 b)$$

$$T_{AA}(m, n) = \Theta(c \log^4(n-m))$$