OS Interfaces and Abstractions

API

- Source level interface
 - Includes files, data types, constants, macros, routines, and parameters
- Basis of software portability
 - Recompile program for the desired architecture
 - Resulting binary runs on that architecture and OS
- An API compliant program will compile and run on any compliant system
 - APIs are primarily for programmers
 - Should work on all types of machines

ABI

- Binding interface for specifying:
 - Dynamic loadable libraries (DLLs)
 - Data formats, calling sequences, linkage conventions
- Binding interface of an API to a hardware
- Basis for binary compatibility
 - Usually one ABI for an OS
 - One binary servers all customers for that hardware (all x86, linux, MacOS)
- Installing a new version of the OS should allow your current apps to still work
- ABI compliant program will run on any compliant system

Libraries and Interfaces

- Normal libraries are accessed through an API
 - Source-level definitions of how to access the library
 - Readily portable between different machines
- Dynamically loadable libraries also called through an API
 - But the dynamic loading mechanism is ABI specific
 - Issues of word length, stack format, linkages

Interfaces and Interoperability

- Strong, stable interfaces are key to allowing programs to operate together
- Also important to OS evolution, since we don't want to break existing programs
- Interface (API and ABI) between the OS and apps shouldn't change

Interface Interoperability

- Stability
 - Programs use sys calls, libraries, and external files
 - API requirements are frozen at compile time
 - Execution platform must support those interfaces
 - Future upgrades must support old interfaces

Compliance

- Comply with the interface, don't try to find loopholes to interact with hardware
- Complete interoperability testing is impossible (cannot test all apps on all platforms)
- New apps are continuously added, which will comply with the interface

Side Effects

- Side effect: occurs when an action has an unexpected outcome
- Effects not specified by interfaces and aren't intended
- Focus on the interface, don't try to exploit side effects

Abstractions

- Simplifies the complex implementations for interacting with hardware
- Hides error handling, performance optimization
- Eliminates behaviors not important to the programmer
- 1. Memory abstractions
- 2. Processor abstractions
- 3. Communications abstractions

Memory Abstractions

- Variables, Files, Database records, messages

Complications

- Persistent memory: memory which persists after a computer is shutdown
- Transient Memory:
- Coherence and atomicity
- Latency (reading or writing to flash drive is very slow, writes takes different amounts of time)
- Same abstractions can be implemented in differently in different devices (read only, solid state, etc)

Source of complications

- OS doesn't have abstract devices
- Physical devices specific to that machine
- We must abstract the hardware regardless of what the implementation looks like

Memory Example

- File, can read or write to the file
- Coherence: If we write the file, we expect out next read to reflect the results of the write
- Atomicity: We expect the entre read/write to occur, entire chunk of 500 bytes should be written instead of just a portion
- If there are several read/writes to the files, we expect them to occur in some order

Implementation flash drives

- Flash drive
- Write once semantics:
 - Rewriting requires an erase cycle
 - Erases a whole block
 - Erases are slow
- Atomicity occurs at the work level
- Blocks can only be erased so many times
- OS needs to abstract the above behaviors

Abstraction considerations

- Flash drive storage is split into blocks
- Different structures for the file system since its difficult to overwrite data words in place
- Garbage collection to deal with blocks filled with inactive data
- Maintaining a pool of empty blocks (we can quickly write to these blocks)
- Wear leveling in use of blocks (make sure blocks are evenly spaced)
- Something that provides atomicity to multiple writes

Interpreter Abstractions

- An interpreter is code that performs commands
- Repertoire: a set of things the interpreter can do (can the CPU do floating point operations?)
- Environment reference: Describes the current state on which the next instruction should be performed
- Interrupts: Out of order instruction execution or canceling execution

Process Abstraction

- Every program is abstracted to a process
- Repertoire : Source code defines what the code can do
- Environment: Stack, heap, register contents
- Interrupts: Ctrl C

Process Abstractions

- Simple with one process, but difficult with multiple
- OS has limited physical memory to hold the environment information
- There is usually one set of registers
- No other interpreters should interfere with a process's resources

Abstraction Implementation

- Schedulers to share the CPU among various processes
- Memory management hardware and software
 - To multiplex memory use among the processes
 - Gives the impression of full use of memory (even though memory is actually shared by different processes)
- Access control mechanisms for other memory abstractions
 - So other processes can't read or alter my files
 - OS needs to isolate processes

Communications Abstractions

- Communication link allows one interpreter to talk to another on the same or different machines
- Memory and cables or networks and interprocess communication mechanisms

Communication properties

- Highly variable performance (network latency can be slow or packets lost)
- Often async, resulting in issues with syncing both parties
- Receiver may only perform the operation since the send occurred
- Additional complications when working with a remote machine

Same Machine Process Communication

- Easy if both processes are on the same machine
 - Done by reading and writing to RAM
 - Or transfer control of memory containing the memory from the sender to receiver

Different Machine Communication

- Need to optimize costs of copying, resulting in difficult memory management
- Inclusion of complex network protocols in the OS itself
- Difficulties with message loss and retransmission
- Security concerns

Generalizing Abstractions

- Different things are made to appear the same
- Applications can all deal with a single class, or a common unifying model

- Portable Document Format (PDF) for printed output
- SCSI/SATA for disks

Federation Frameworks

- A generalized abstraction
- A structure that allows many similar but different things to be treated uniformly
- By creating one interface that works with all implementations (at a high level, users can use flash drives and hard drives in the same way)
- Least common denominator which works for multiple models:
 - Common model has optional features
 - Ex. printers: some printers can print double sided, others can't

Abstractions and Layering

- Common to create increasingly complex services by layering abstractions
 - Generic file system layers on a particular file system
- Layering allows good modularity
 - Easy to build multiple services on a lower layer
 - Multiple file systems on one disk

Downside of Layering

- More abstractions results in performance penalties
- Often expensive to move form one layer to another
 - SInce frequently requires changing data structures or representations
 - Abstractions require additional instructions
- Another downside is that lower layer may limit what the upper layer can do
 - Any abstract network link may hide causes of packet losses

Other OS Abstractions

- Different abstractions provide different types of abstractions
- OS must do work to provide an abstraction (higher level = more work)
- Programmers and users have to choose the right abstraction to work with

Summary

- Stable interfaces needed for proper performance of an OS
 - APIs for program development and ABIs for user experience
- Abstractions make OS systems easier for both programmers and users
- The most important OS abstractions involve memory, interpreters, and communication