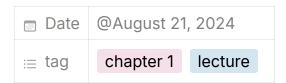
Computer Networks and the Internet



▼ 1.1 What is the Internet?

- The Internet is a public computer network used to discuss computer networks and protocols.
- It can be described as:
 - 1. Nuts-and-bolts: Basic hardware and software components.
 - 2. **Networking infrastructure**: Provides services to distributed applications.

1.1.1 A Nuts-and-Bolts Description

- The Internet is a network that connects millions of devices worldwide.
- End systems (hosts) include:
 - Traditional devices (PCs, servers).
 - Non-traditional devices (laptops, smartphones, TVs, sensors, etc.).
 - Estimated 850 million end systems in 2011, with about 2 billion users globally.

Communication Links and Packet Switches:

- End systems are connected by communication links (e.g., coaxial cable, copper wire, optical fiber).
- Data is transmitted in bits/second.

- Packet switches (routers and link-layer switches) forward data in packets.
- Routers: Used in the core network.
- Link-layer switches: Used in access networks.
- Packets follow a route or path through the network.

Packet-Switched Networks:

- Similar to transportation networks.
- Data is segmented like cargo in trucks, which travel through a network of highways and roads to their destination.

Internet Service Providers (ISPs):

- End systems access the Internet via ISPs (residential, corporate, university ISPs).
- ISPs offer various access types (broadband, DSL, dial-up).
- ISPs are interconnected, with lower-tier ISPs connecting through national/international upper-tier ISPs (e.g., AT&T, Sprint).
- Each ISP network is independently managed and runs the IP protocol.

Protocols and Standards

- The Internet runs on protocols, with TCP and IP being the most important.
 - IP protocol: Defines packet formats between routers and end systems.
 - Collectively called **TCP/IP** protocols.

Standards:

- Protocols must be standardized for interoperability.
- Internet standards are developed by the Internet Engineering Task
 Force (IETF).
- Standards documents are known as Requests for Comments (RFCs).
- Over 6,000 RFCs define protocols such as TCP, IP, HTTP, and SMTP.

• IEEE Standards:

 IEEE 802 LAN/MAN Standards Committee defines network components like Ethernet and WiFi.

1.1.2 A Services Description

Internet Applications:

- Examples: Electronic mail, Web surfing, social networks, instant messaging, VoIP, video streaming, distributed games, P2P file sharing, television over the Internet, remote login.
- Distributed applications involve multiple end systems exchanging data.

Applications vs. Packet Switches:

- Applications run on end systems, not on packet switches.
- Packet switches facilitate data exchange but do not handle application data directly.

• Developing Internet Applications:

- Requires programming end systems (e.g., Java, C, Python).
- Programs on different end systems need to communicate data.
- End systems use an Application Programming Interface (API) to instruct the Internet to deliver data to destination programs.

Internet API:

- A set of rules for data delivery.
- Analogy: Like postal service rules for sending a letter (envelope, address, stamp).

Internet Services:

- Provides multiple services to applications.
- Developers must choose from these services for their applications.
- Details on Internet services will be covered in Chapter 2.

• Further Learning:

 The book will explain packet switching, TCP/IP, routers, communication links, distributed applications, and Internet attachment of devices like toasters and weather sensors.

1.1.3 What Is a Protocol?

Human Analogy:

- **Example 1**: Asking someone for the time:
 - Greeting is essential (e.g., "Hi").
 - Response indicates willingness to communicate.
 - Lack of response or inappropriate responses may end the interaction.
- Example 2: Asking a question in class:
 - Raising a hand to indicate a question.
 - Teacher acknowledges and responds.

Network Protocols:

 Similar to human protocols but involve hardware or software components.

• Examples:

- Flow of bits between network interface cards.
- Congestion control in end systems.
- Path determination by routers.

Key Elements of a Protocol:

- Defines format and order of messages exchanged.
- Specifies actions taken on message transmission and receipt.

• Internet Protocols:

- Examples include Web server communication:
 - Sending a connection request.

- Receiving a connection reply.
- Sending a GET request for a Web page.
- Receiving the Web page.

• Importance:

- Understanding protocols is crucial for mastering computer networking.
- Protocols can range from simple to complex.

▼ 1.2 The Network Edge

End Systems

- **Definition**: Computers, smartphones, and other devices connected to the Internet.
- Types:
 - Desktop computers (PCs, Macs, Linux boxes)
 - Servers (Web, e-mail)
 - Mobile devices (laptops, smartphones, tablets)
- Hosts: End systems running application programs (e.g., Web browsers, e-mail clients).

1.2.1 Access Networks

- **Definition**: Network connecting an end system to the first router (edge router).
- Types:
 - Home Access:
 - DSL:
 - Uses existing telephone lines.
 - Transmission rates: 12-24 Mbps downstream, 1.8-2.5 Mbps upstream.
 - Cable:

- Uses cable TV infrastructure.
- Transmission rates: Up to 42.8 Mbps downstream, 30.7 Mbps upstream.
- Shared broadcast medium affecting speeds.

FTTH (Fiber to the Home):

- Direct optical fiber path from CO to home.
- · Speed potential: Gigabits per second.

Other Technologies:

- Satellite: Internet speeds over 1 Mbps.
- Dial-Up: Up to 56 kbps, slow compared to DSL and cable.

Access in the Enterprise and Home

• Ethernet:

- Common in corporate and university settings.
- Speed: Typically 100 Mbps to 10 Gbps.

WiFi:

- Wireless LAN technology.
- Speed: Up to 54 Mbps.
- Common in homes and public places.

Wide-Area Wireless Access

• 3G:

Provides speeds in excess of 1 Mbps.

• 4G/LTE:

- Higher-speed access, potentially over 10 Mbps.
- LTE speeds: Tens of Mbps reported.

1.2.2 Physical Media

Definition and Role

- Physical media refer to the physical medium through which data travels from one end system to another.
- The bit travels through various transmitter-receiver pairs, each using different physical media.

Categories of Physical Media

- Guided Media: Waves are guided along a solid medium.
 - Examples: Fiber-optic cable, twisted-pair copper wire, coaxial cable.
- **Unguided Media**: Waves propagate in the atmosphere or space.
 - Examples: Wireless LAN, digital satellite channels.

Cost Considerations

- Material cost is often minor compared to installation labor.
- Installation of multiple media types in buildings can be cost-effective for future use.

Types of Physical Media

• Twisted-Pair Copper Wire

- Commonly used in telephone networks and LANs.
- Consists of two insulated copper wires twisted together to reduce interference.
- Data rates: 10 Mbps to 10 Gbps (modern technology).
- DSL technology allows Internet access at tens of Mbps.

Coaxial Cable

- Contains two concentric copper conductors.
- Used in cable television systems and cable Internet access.

 Provides high data transmission rates and can be used as a shared medium.

Fiber Optics

- Uses thin, flexible fibers to conduct light pulses.
- Supports high bit rates (tens to hundreds of gigabits per second).
- Immune to electromagnetic interference, low signal attenuation.
- Common in long-haul and backbone networks.
- High cost of optical devices limits short-haul applications.

Terrestrial Radio Channels

- Carry signals via electromagnetic spectrum.
- No physical wire required; can penetrate walls and provide mobile connectivity.
- Classified into short-distance, local-area, and wide-area channels.
- Characteristics affected by path loss, shadow fading, multipath fading, and interference.

Satellite Radio Channels

- Use communication satellites to link ground stations.
- Two types:
 - Geostationary Satellites: Fixed above one spot on Earth; higher signal propagation delay.
 - Low-Earth Orbit (LEO) Satellites: Rotate around Earth; require multiple satellites for coverage.
- Geostationary satellites used for high-speed links in remote areas; LEO satellites may become more relevant for future Internet access.

▼ 1.3 The Network Core

 The network core consists of packet switches and links connecting end systems.

1.3.1 Packet Switching

Function

- End systems exchange messages that are broken into packets.
- Packets travel through communication links and packet switches (routers and link-layer switches).
- Transmission rate: Packet transmission time = L/R seconds, where L is packet size and R is link rate.

L/RL / R

LL

RR

Store-and-Forward Transmission

- Packets must be fully received before forwarding begins.
- Total delay for one packet: 2L/R seconds.
- For multiple packets, delay accumulates based on transmission and processing at each router.
- $\circ \ d_{EndToEnd} = NL/R$

Queuing Delays and Packet Loss

- Packets wait in output buffers if links are busy.
- Buffer overflow leads to packet loss.
- Queuing delays vary with network congestion.

Forwarding Tables and Routing Protocols

Forwarding Process

- Routers use forwarding tables to direct packets based on destination IP addresses.
- The router's forwarding table maps addresses to outbound links.

Routing Protocols

- Automated protocols set and update forwarding tables.
- They determine paths and configure routers to optimize packet routing.

Practical Example

Use Traceroute to observe end-to-end routing paths in the Internet.

1.3.2 Circuit Switching

Overview:

- Definition: Resources (buffers, link transmission rate) are reserved for the duration of the communication session.
- Analogy: Similar to reserving a table at a restaurant versus waiting for a table.

Traditional Use:

- **Example**: Telephone networks.
- Connection: Establishes a dedicated end-to-end connection with reserved transmission rate.

Mechanism:

- **Circuit**: Maintains connection state and reserves a constant transmission rate.
- Illustration: Figure shows circuits on links with dedicated bandwidth.

Multiplexing:

- **Frequency-Division Multiplexing (FDM)**: Divides frequency spectrum among connections.
- Time-Division Multiplexing (TDM): Divides time into frames and assigns time slots to connections.

• Efficiency Issues:

- Resource Wastage: Idle periods result in underutilization of reserved resources.
- Complexity: Requires complex signaling software.

Numerical Example:

 File Transmission: 640,000 bits file with 1.536 Mbps bit rate and 24 TDM slots takes 10.5 seconds (including circuit establishment).

Packet Switching vs. Circuit Switching

Packet Switching:

- Overview: Sends packets without reserving resources; uses resources on demand.
- Efficiency: More efficient as it allows for higher user numbers and better sharing of transmission capacity.

• Comparison:

- Real-Time Services: Circuit switching is preferred for services requiring predictable delays (e.g., voice calls).
- Capacity Sharing: Packet switching allows for dynamic allocation of bandwidth based on demand.

• Examples:

- Circuit Switching: Limited to number of simultaneous users due to reserved bandwidth.
- Packet Switching: Can handle more users with minimal delay, even under high demand.

• Trends:

 Current Use: Packet switching is increasingly dominant, with many circuit-switched networks migrating to packet switching for efficiency.

1.3.3 A Network of Networks

Overview

• Access ISPs: Provide connectivity to end systems (e.g., PCs, smartphones) via wired or wireless technologies (DSL, cable, FTTH, Wi-Fi, cellular).

- Access ISP Examples: Universities, companies, telcos, and cable companies.
- Objective: Connect billions of end systems through interconnecting access ISPs.

Network Structures

1. Network Structure 1

- Design: Access ISPs connect to a single global transit ISP.
- Cost: High for global transit ISP; charges access ISPs based on traffic.
- Role: Global transit ISP acts as a provider to access ISPs.

2. Network Structure 2

- Design: Multiple global transit ISPs interconnect with access ISPs.
- Benefit: Access ISPs can choose from multiple providers.
- Requirement: Global transit ISPs must interconnect to enable communication between access ISPs connected to different providers.

3. Network Structure 3.

- **Design**: Includes regional ISPs connecting to tier-1 ISPs.
- Components:
 - Tier-1 ISPs: Major ISPs with extensive coverage (e.g., Level 3, AT&T, Sprint).
 - Regional ISPs: Connect to tier-1 ISPs and other regional ISPs.
 - Hierarchy: Access ISPs connect to regional ISPs, which connect to tier-1 ISPs.

4. Network Structure 4

• Components:

- Points of Presence (PoPs): Routers where customer ISPs connect.
- **Multi-Homing**: ISPs connect to multiple providers for redundancy.

- Peering: ISPs connect directly to avoid intermediary costs; typically settlement-free.
- Internet Exchange Points (IXPs): Facilitate peering among multiple ISPs.

5. Network Structure 5

- **Design**: Adds content provider networks to the existing hierarchy.
- **Example**: Google's private network connects its data centers globally and peers with lower-tier ISPs.
- Advantage: Reduces payments to upper-tier ISPs and improves service control.

Key Points

- Internet Complexity: Consists of tier-1 ISPs, regional ISPs, and access ISPs.
- Tier-1 ISPs: Large, globally connected ISPs.
- **Content Providers**: Establish private networks to reduce costs and enhance control over content delivery.
- **Customer-Provider Relationships**: Access ISPs and content providers pay higher-tier ISPs for connectivity.

▼ 1.4 Delay, Loss, and Throughput in Packet Switched Networks

1.4.1 Overview of Delay in Packet-Switched Networks

Overview

- The Internet provides services to distributed applications on end systems.
- Ideal network performance aims for instantaneous data transfer without loss.

Real networks face constraints: limited throughput, delays, and packet loss.

Types of Delay

1. Processing Delay

- Time to examine the packet header and check for errors.
- Typically in microseconds or less.

2. Queuing Delay

- Time spent waiting in the queue before transmission.
- Depends on the number of packets ahead in the queue.
- Ranges from microseconds to milliseconds.

3. Transmission Delay

- Time to push all bits of a packet into the link.
- Calculated as L/R, where L is packet length and R is transmission rate.
- Typically microseconds to milliseconds.

4. Propagation Delay

- Time for a bit to travel from sender to receiver.
- Calculated as d/s, where d is distance and s is propagation speed.
- Ranges from milliseconds in wide-area networks.
 - $2.10^8 meters/secto 3.10^8 meters/sec$

Comparison of Delays

- Transmission Delay: Time to push the packet into the link; independent of distance.
- **Propagation Delay**: Time for bits to propagate; dependent on distance.

Practical Considerations

• Total Nodal Delay: dnodal=dproc+dqueue+dtrans+dprop

$$d_{
m nodal} = d_{
m proc} + d_{
m queue} + d_{
m trans} + d_{
m prop}$$

- dproc: Often negligible but impacts maximum throughput.
- dqueue: Varies with traffic load.
- dtrans: Significant for low-speed links.
- dprop: Can dominate in satellite links.

1.4.2 Queuing Delay and Packet Loss

Queuing Delay

• **Definition**: Time a packet spends waiting in the queue before transmission.

Characteristics:

- Varies from packet to packet.
- Analyzed using statistical measures: average delay, variance, and probability of exceeding a threshold.

• Factors Influencing Queuing Delay:

- \circ **Traffic Intensity**: Ratio La/R , where La is the packet arrival rate and R is the transmission rate.
 - If La/R>1: Queue length increases indefinitely, causing delay to approach infinity.
 - If La/R <= 1: Queuing delay is impacted by traffic patterns (periodic or bursty).

Traffic Patterns:

- Periodic Arrivals: Predictable queuing delays.
- Random Arrivals: More complex; average delay depends on traffic intensity.

Packet Loss

• **Definition**: Occurs when a packet arrives at a full queue and is dropped due to finite queue capacity.

Impact:

- Packet loss becomes more frequent as traffic intensity exceeds 1.
- From an end-system perspective, a lost packet never reaches its destination.
- Performance is assessed based on both delay and packet loss probability.

• Mitigation:

Lost packets may be retransmitted to ensure complete data transfer.

1.4.3 End-to-End Delay

End-to-End Delay

- **Definition**: Total delay from the source host to the destination host, considering multiple routers.
- Formula:

$$d_{end-to-end}=N imes(d_{proc}+d_{trans}+d_{prop})$$
 where $d_{trans}=rac{L}{R}$, L is the packet size, and N is the number of routers.

Traceroute

 Purpose: Measures end-to-end delay by sending special packets through routers.

• Process:

- Sends multiple packets, each marked with a unique number.
- Each router responds with a message containing its name and address.
- The source calculates the round-trip delay for each packet.

Output:

Displays router numbers, names, addresses, and round-trip delays.

- Example output includes delays for each router along the route.
- **Usage**: Available at <u>Traceroute.org</u> and graphical tools like PingPlotter.

Additional Delays

- End System Delays:
 - Protocol-based delays (e.g., medium access in shared networks).
 - Media packetization delay (e.g., VoIP applications) where time is spent filling packets with data.

1.4.4 Throughput in Computer Networks

Definition

• **Throughput:** The rate at which data is successfully transferred from Host A to Host B, measured in bits per second (bps).

Examples and Analysis

- Two-Link Network:
 - Throughput: Determined by the bottleneck link.
 - \circ Formula: Throughput $= \min(R_s, R_c)$
 - \circ Transfer Time: $\mathrm{Time} = \frac{F}{\min(R_s,R_c)}$, where F is the file size, Rs is the server link rate, and Rc is the client link rate.
- Network with Multiple Links:
 - Throughput: Determined by the bottleneck link in the path.
 - **Formula**: Throughput = $\min(R_1, R_2, ..., R_N)$, where Ri are the rates of the links.
- Core Network with High-Speed Links:

- Throughput: Limited by the access network links.
- Formula: Throughput $= \min(R_s, R_c)$, assuming core links have high capacity.

• Shared Core Link with Multiple Downloads:

- Throughput: Can be limited by the shared core link.
- **Example**: If the common link rate is R and is shared equally among N downloads, throughput per download is $\frac{R}{N}$

Key Points

- **Throughput** is affected by the transmission rates of all links along the data path.
- **Bottleneck Link**: The link with the lowest transmission rate along the path typically determines the overall throughput.
- **Traffic Impact**: Throughput can be reduced if a shared link is heavily used by multiple data flows.

▼ 1.5 Protocol Layers and Their Service Models

1.5.1 Layered Architecture

- **Analogy**: Complex systems, like the airline system, can be described in layers (e.g., ticketing, baggage, gate, takeoff/landing).
- Layered Functionality:
 - Each layer performs specific actions and uses services from the layer below.
 - Layers provide modularity and ease of updates without affecting other components.

Protocol Layering

• Network Protocols:

- Organized in layers for structure.
- Each layer offers specific services and uses services from the layer below.

Implementation:

- Layers can be implemented in software, hardware, or a combination.
- Distributed across end systems, packet switches, and network components.

Advantages:

- Modularity and ease of updating.
- Structured approach to discussing system components.

Drawbacks:

- Possible duplication of functionality.
- Potential need for inter-layer information, violating separation goals.

Internet Protocol Stack

Layers:

- 1. **Application Layer:** Includes protocols like HTTP, SMTP, FTP, and DNS. Handles network applications and messages.
- 2. **Transport Layer**: Provides TCP (connection-oriented, reliable delivery) and UDP (connectionless, no-frills service).
- 3. **Network Layer:** Handles datagrams with IP protocol and routing protocols. Often referred to as the IP layer.
- 4. **Link Layer**: Moves frames between nodes, with protocols like Ethernet, WiFi, and DOCSIS.
- 5. **Physical Layer**: Moves bits across the link, with protocols varying by transmission medium (e.g., twisted-pair, fiber).

The OSI Model

Seven Layers:

- 1. Application
- 2. Presentation
- 3. Session
- 4. Transport
- 5. Network
- 6. Data Link
- 7. Physical

• Comparison:

- The Internet model roughly matches the OSI model's functionality except for Presentation and Session layers.
- **Presentation Layer:** Data compression, encryption, and description.
- Session Layer: Data exchange synchronization, checkpointing, and recovery.

Internet Approach:

 Functions of Presentation and Session layers are handled at the application level if needed.

1.5.2 Encapsulation

Overview:

 Path: Data moves through protocol stacks of sending and receiving end systems, as well as intervening link-layer switches and routers.

• Layers Implemented:

- Link-layer Switches: Layers 1 and 2.
- Routers: Layers 1 through 3.
- Hosts: Implement all five layers.

• Encapsulation Process:

At Sending Host:

- Application Layer: Message (M) is generated.
- Transport Layer: Adds transport-layer header information (H_t), creating a transport-layer segment.
- Network Layer: Adds network-layer header information (H_n), forming a network-layer datagram.
- Link Layer: Adds link-layer header information, creating a link-layer frame.

Analogy:

- o Interoffice Memo:
 - Application-Layer Message: Memo.
 - Transport-Layer Segment: Memo in an interoffice envelope with header information.
 - Datagram: Interoffice envelope in a postal envelope with postal addresses.
 - De-encapsulation: Receiving end extracts the memo from the postal envelope and interoffice envelope.

• Complexity:

 Large Messages: May be divided into multiple segments and datagrams, requiring reassembly at the receiving end.

▼ 1.6 Networks Under Attack

- Importance of Network Security:
 - The Internet is critical for institutions and individuals.
 - Network security involves defending against attacks and designing secure architectures.

• Prevalent Security Issues:

- Malware:
 - Types:

- **Viruses**: Require user interaction to spread (e.g., via email attachments).
- Worms: Spread without user interaction, exploiting vulnerabilities in network applications.
- Effects: Deleting files, installing spyware, creating botnets for attacks.

Denial-of-Service (DoS) Attacks:

Types:

- Vulnerability Attacks: Exploit flaws in applications or systems to crash services.
- Bandwidth Flooding: Overwhelm the target's access link with excessive traffic.
- **Connection Flooding:** Saturate the target with numerous bogus connections.
- Distributed DoS (DDoS): Uses multiple sources to overwhelm a target, making it harder to detect and defend against.

Packet Sniffing:

- Vulnerability: Wireless and wired environments can be vulnerable to packet sniffers that capture sensitive data.
- Defense: Cryptography and secure protocols to protect data.

IP Spoofing:

- Vulnerability: Crafting packets with false source addresses to deceive receivers.
- Defense: End-point authentication to verify message origins.

Historical Context:

- Original Internet Design: Based on mutual trust and transparency, lacking built-in security.
- Current Challenges: Addressing security issues in a context where mutual trust is not always present.

• Future Focus:

 Develop defenses against various attacks including malware, sniffing, spoofing, and DDoS.

▼ 1.7 History of Computer Networking and the Internet

1.7.1 The Development of Packet Switching: 1961–1972

• Origins:

- In the early 1960s, the telephone network used circuit switching for constant-rate voice transmission.
- Increasing use of computers and timesharing led to the need for a more efficient method for bursty traffic.

• Key Developments:

- Packet Switching Invention:
 - **Leonard Kleinrock** (MIT): Demonstrated the effectiveness of packet switching using queuing theory (1961-1964).
 - Paul Baran (Rand Institute): Investigated packet switching for secure military communications (1964).
 - Donald Davies and Roger Scantlebury (National Physical Laboratory, England): Developed packet switching concepts concurrently.

ARPAnet:

- J.C.R. Licklider and Lawrence Roberts: Led ARPA's computer science program and planned the ARPAnet.
- 1969: ARPAnet's initial installation with four nodes: UCLA, Stanford Research Institute, UC Santa Barbara, and University of Utah.
- **1972**: ARPAnet grew to 15 nodes and demonstrated its first host-to-host protocol, NCP.

• Ray Tomlinson: Developed the first e-mail program in 1972.

1.7.2 Proprietary Networks and Internetworking: 1972–1980

- Expansion of Networks:
 - Early 1970s: Emergence of additional packet-switching networks:
 - ALOHANet: Microwave network linking Hawaiian universities.
 - DARPA's Packet-Satellite and Packet-Radio Networks: Innovations in satellite and radio packet-switching.
 - **Telenet**: Commercial packet-switching network based on ARPAnet technology.
 - Cyclades: French packet-switching network developed by Louis Pouzin.
 - Tymnet and GE Information Services: Time-sharing networks.
 - IBM's SNA (1969–1974): Parallel development to ARPAnet.
- Internetworking Developments:
 - Vinton Cerf and Robert Kahn: Pioneered interconnecting networks, leading to the concept of "internetting."
 - Early TCP and IP:
 - **TCP**: Initially combined reliable data delivery with forwarding functions.
 - Separation of IP and TCP: Led to the development of UDP for nonflow-controlled transport.
 - By the end of the 1970s: TCP, UDP, and IP were conceptually established.
- Notable Contributions:
 - ALOHA Protocol: First multiple-access protocol developed by Norman Abramson for ALOHAnet.

Ethernet Protocol:

- Developed by Metcalfe and Boggs: Based on ALOHA protocol principles for wire-based shared networks.
- Purpose: Designed to connect multiple PCs, printers, and shared disks.

1.7.3 A Proliferation of Networks: 1980–1990

- Growth of Hosts:
 - End of 1970s: Approximately 200 hosts on ARPAnet.
 - End of 1980s: 100,000 hosts on the public Internet.
- Key Network Developments:
 - BITNET: Provided e-mail and file transfers among Northeast universities.
 - CSNET: Linked university researchers without ARPAnet access.
 - NSFNET (1986): Linked NSF-sponsored supercomputing centers;
 backbone speed increased from 56 kbps to 1.5 Mbps by decade's end.

ARPAnet and Internet Architecture:

- January 1, 1983: TCP/IP officially deployed as the standard protocol for ARPAnet (replacing NCP).
- Late 1980s: Extensions to TCP for host-based congestion control;
 development of DNS for mapping Internet names to IP addresses.

• International Development:

- Minitel Project (Early 1980s):
 - France: Public packet-switched network based on X.25; free terminals distributed to French households.
 - Success: Offered over 20,000 services by mid-1990s, including home banking and research databases.

1.7.4 The Internet Explosion: The 1990s

• End of ARPAnet:

- ARPAnet ceased to exist.
- NSFNET lifted commercial use restrictions in 1991; decommissioned in 1995.

Emergence of the World Wide Web:

- Invention: Tim Berners-Lee developed the Web at CERN (1989-1991).
- Components: HTML, HTTP, Web server, and browser.
- **Early Growth**: About 200 Web servers by late 1993; Web browsers with GUI interfaces developed, including Mosaic (later Netscape).

Browser Wars:

- Netscape: Gained popularity in the mid-1990s.
- Microsoft: Entered the browser market in 1996, eventually dominating the browser market.

Internet Innovations and Growth:

- Major Applications:
 - E-mail (with attachments and Web-accessible)
 - Web browsing and Internet commerce
 - Instant messaging with contact lists
 - Peer-to-peer file sharing (e.g., Napster)
- **Impact**: Hundreds of applications by the end of the 1990s.

Financial Market Trends:

- Dot-Com Bubble: Many startups went public before becoming profitable; market collapse in 2000-2001.
- Winners: Microsoft, Cisco, Yahoo, eBay, Google, and Amazon emerged successful.

1.7.5 The New Millennium

Broadband Internet Access:

- **Deployment**: Cable modems, DSL, fiber to the home.
- Applications: User-generated video (YouTube), on-demand streaming (Netflix), multi-person video conferencing (Skype).

High-Speed Wireless Networks:

- WiFi and Cellular: High-speed public WiFi (54 Mbps+), medium-speed
 3G/4G.
- Impact: Ubiquitous connectivity, rise of hand-held devices (iPhones, Androids, iPads).

Online Social Networks:

- Platforms: Facebook, Twitter.
- Features: Massive people networks, APIs for networked applications and distributed games.

• Private Networks by Online Service Providers:

- Examples: Google, Microsoft.
- Function: Connect global data centers, bypass the Internet by peering directly with lower-tier ISPs.

Cloud Computing:

- Providers: Amazon EC2, Google App Engine, Microsoft Azure.
- Usage: Internet applications, email, Web hosting.
- Benefits: Scalable computing/storage, high-performance private networks.