

The background features a complex network of thin grey lines connecting various points, forming a web-like structure. Scattered throughout are numerous triangles of different sizes and orientations, some with solid grey outlines and others with dashed or dotted lines. The overall aesthetic is technical and modern.

Introduction to Assembly and Reverse Engineering

By Sam Goodwin

What is Assembly Language?

- **Very** low level programming language
- Pretty much macros for writing machine code
 - **Machine Code** - Computer instructions that directly control the CPU
- Assembly code is run through an **assembler** that turns it into machine code that can be executed by your computer.
- Old and pretty sucky, but much easier than writing/reading machine code ourselves.
- Represents the lowest level operations of your computer

```
000000000000000010c0 <register_tm_clones>:
10c0: 48 8d 3d 49 2f 00 00    lea    0x2f49(%rip),%rdi    # 4010 <__TMC_END__>
10c7: 48 8d 35 42 2f 00 00    lea    0x2f42(%rip),%rsi    # 4010 <__TMC_END__>
10ce: 48 29 fe                sub    %rdi,%rsi
10d1: 48 8b f0                mov    %rsi,%rax
10d4: 48 c1 ee 3f            shr    $0x3f,%rsi
10d8: 48 c1 f8 03            sar    $0x3,%rax
10dc: 48 01 c6                add    %rax,%rsi
10df: 48 d1 fe                sar    %rsi
10e2: 74 14                  je     10f8 <register_tm_clones+0x38>
10e4: 48 08 05 2f 00 00      mov    0x2f05(%rip),%rax    # 3ff0 <__ITM_registerTMCcloneTable>
10eb: 48 85 c0                test   %rax,%rax
10ee: 74 08                  je     10f8 <register_tm_clones+0x38>
10f0: ff e0                  jmpq   *%rax
10f2: 66 0f 1f 44 00 00      nopw   0x0(%rax,%rax,1)
10f8: c3                     retq
10f9: 0f 1f 80 00 00 00 00    nopl   0x0(%rax)

00000000000000001100 <_do_global_ctors_aux>:
1100: f3 0f 1e fa            endbr64
1104: 80 3d 05 2f 00 00      cmpb   $0x0,0x2f05(%rip)    # 4010 <__TMC_END__>
110b: 75 2b                  jne     1138 <_do_global_ctors_aux+0x38>
110d: 55                     push    %rbp
110e: 48 83 d3 e2 2e 00 00    cmpq   $0x0,0x2ee2(%rip)    # 3ff8 <_cxa_finalize@GLIBC_2.2.5>
1115: 00
1116: 48 89 e5                mov     %rsp,%rbp
1119: 74 0c                  je      1127 <_do_global_ctors_aux+0x27>
111b: 48 8b 3d e6 2e 00 00    mov     0x2ee6(%rip),%rdi    # 4008 <_dso_handle>
1122: e8 10 ff ff ff         callq   1048 <_cxa_finalize@plt>
1127: e8 64 ff ff ff         callq   1090 <register_tm_clones>
112c: c6 05 dd 2e 00 00 01    movb    $0x1,0x2edd(%rip)    # 4010 <__TMC_END__>
1133: 5d                     pop     %rbp
1134: c3                     retq
1135: 0f 1f 00                nopl    (%rax)
1138: c3                     retq
1139: 0f 1f 80 00 00 00 00    nopl    0x0(%rax)

00000000000000001140 <frame_dummy>:
1140: f3 0f 1e fa            endbr64
1144: e9 77 ff ff ff         jmpq    10c0 <register_tm_clones>

00000000000000001149 <main>:
1149: f3 0f 1e fa            endbr64
114d: 55                     push    %rbp
114e: 48 89 e5                mov     %rsp,%rbp
1151: 48 8d 3d ac 0e 00 00    lea     0xeac(%rip),%rdi    # 2004 <_IO_stdin_used+0x4>
1158: e8 f1 fe ff ff         callq   1058 <puts@plt>
115d: b8 00 00 00 00         mov     $0x0,%eax
1162: 5d                     pop     %rbp
1163: c3                     retq
1164: 66 2e 0f 1f 84 00 00    nopw    %cs:0x0(%rax,%rax,1)
116b: 00 00 00
116e: 66 90                  xchgb   %ax,%ax

00000000000000001170 <_libc_csu_init>:
1170: f3 0f 1e fa            endbr64
1174: 41 57                  push    %r15
```

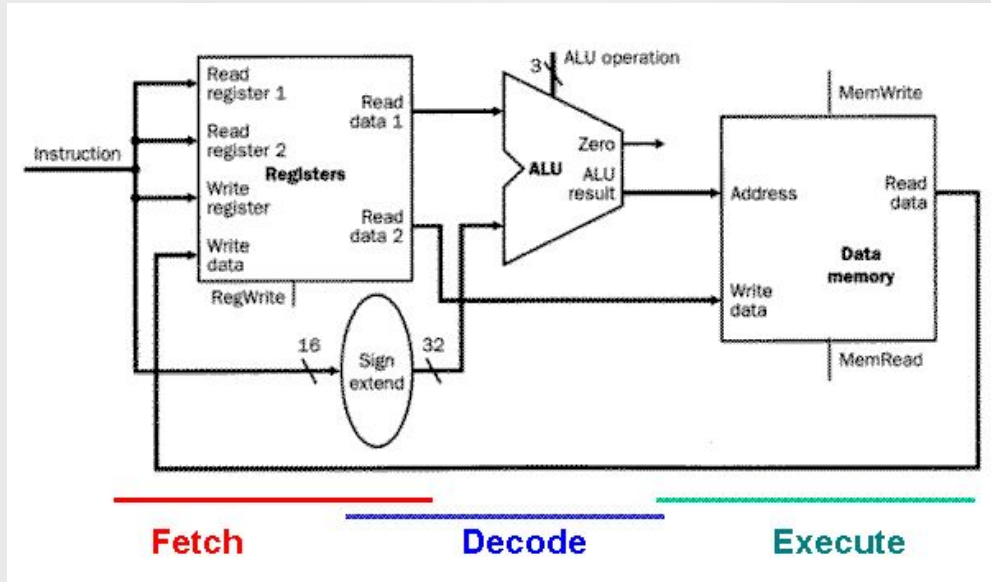
mov rax, rbx

Assembler

0x48 0x89 0xd8

What is Assembly Language? Cont.

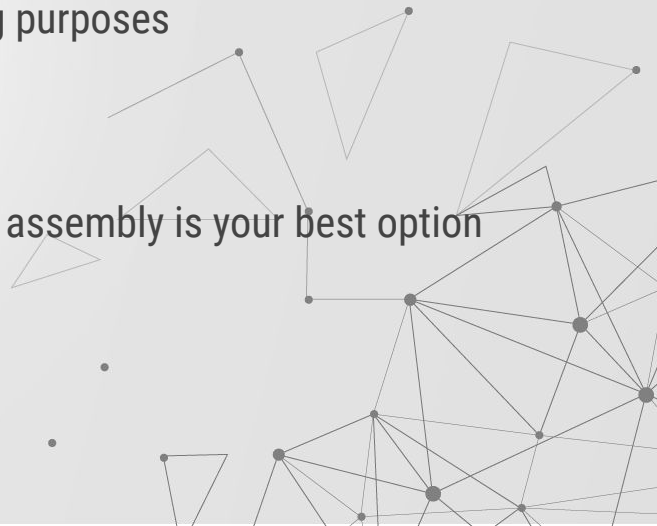
- Assembly language changes based on your **processor architecture**
 - x86 instructions are different from ARM, MIPS, RISC-V, ...
- This is why you may have heard some computers “speak different languages”



MIPS load/store instruction datapath

Why Should I Care About Assembly?

- Who wants to take the time to learn a frustrating caveman language?
 - Mastery of assembly means mastery of the low-level computer stack, very impressive!
 - Still used in some low-memory embedded systems environments
 - Useful for understanding a program's behaviour for debugging purposes
 - Pwn and RE categories are in almost all CTFs
 - If you want to reverse engineer a pre-compiled binary, reading assembly is your best option
- without getting into some **advanced software**

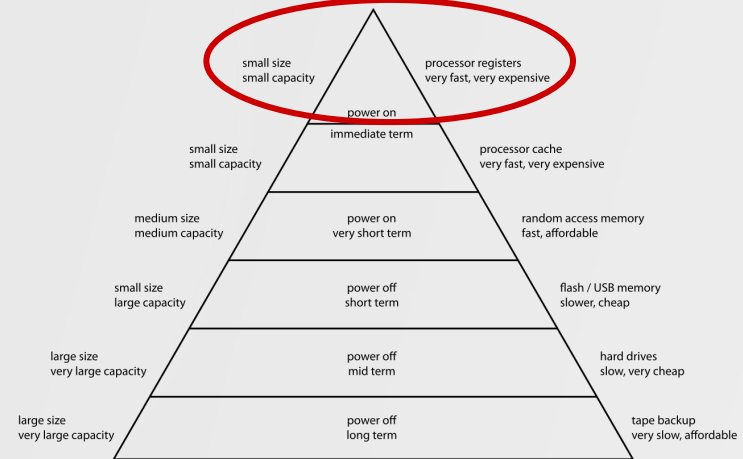


CPU Registers

- The fastest memory your computer has. Baked right on top of the CPU.
- Very small, only enough room for the most important information
 - x64 has 16 64-bit registers (128 bytes)
 - Typically holds the current instruction and some stack information. Can also store local variables, function arguments, and more. These also vary by architecture.

```
$rax : 0x0000555555555149 → <main+0> endbr64
$rbx : 0x0000555555555170 → <__libc_csu_init+0> endbr64
$rcx : 0x0000555555555170 → <__libc_csu_init+0> endbr64
$rdx : 0x00007fffffffe248 → 0x00007fffffffe4a4 → "SHELL=/bin/bash"
$rsp : 0x00007fffffffe148 → 0x00007ffffff7ded0b3 → <__libc_start_main+2
$rbp : 0x0
$rsi : 0x00007fffffffe238 → 0x00007fffffffe477 → "/home/griffith/mcc_
$rdi : 0x1
$rip : 0x0000555555555149 → <main+0> endbr64
$r8 : 0x0
$r9 : 0x00007ffff7fe0d50 → endbr64
$r10 : 0xf
$r11 : 0x2
$r12 : 0x0000555555555060 → <_start+0> endbr64
$r13 : 0x00007fffffffe230 → 0x0000000000000001
$r14 : 0x0
$r15 : 0x0
```

Computer Memory Hierarchy



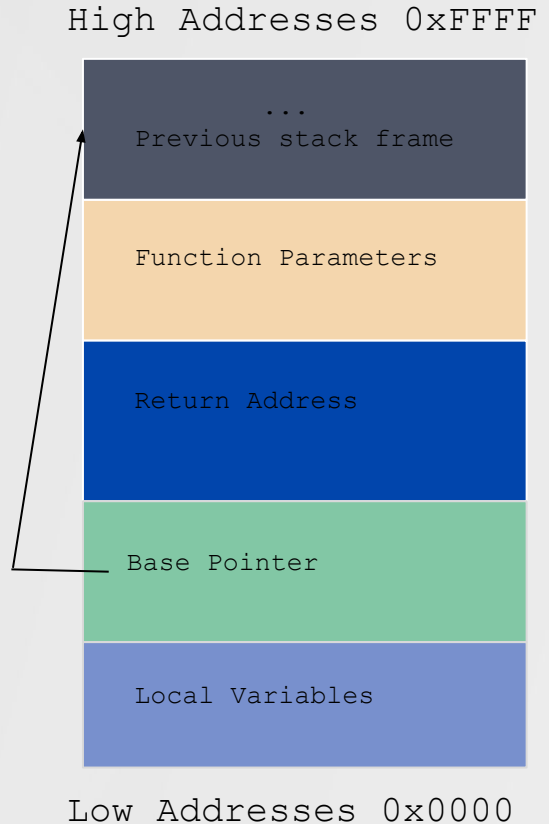
Common Instructions

- **Mov S, D - Move data from source to destination**
- **Add S, D - add source to destination**
- **Sub S, D - subtract source from destination**
- **and/or/xor S, D - bitwise operations on destination by source**
- **Jmp $Label$ - jump to execute at $label$**
- **Jg/Jl $Label$ - jump if greater/less than**

Nice x64 cheatsheet: https://cs.brown.edu/courses/cs033/docs/guides/x64_cheatsheet.pdf

Stack Frame

- What if we want to store something in memory but it can't fit in a register? Put it on the stack!
- What exactly goes on stack varies by architecture but typically includes at least return address, base pointer, and some local variables.
- Each function gets its own stack frame set up for itself when it gets called
- The return address points to the code that will execute after this frame is closed
 - **Hmmm, that sounds interesting**
- Base pointer holds the address of the top of the stack frame and is used to reference the local variables



C to Assembly

```
#include<stdio.h>

int main(){
    puts("hello world!\n");
}
```

Disassembled

```
1000: 4c 5a c3                test    %eax,%eax
1001: 74 00                   je      1003 <register_tm_clones+0x38>
1002: ff e0                   jmpq    *%rax
1003: 66 0f 1f 44 00 00       nopw    0x0(%rax,%rax,1)
1004: c3                      retq
1005: 0f 1f 80 00 00 00       nopl    0x0(%rax)

0000000000001100: <_do_global_ctors_aux>
1100: f3 0f 1e fa             endbr64
1101: 80 3d 05 2f 00 00 00    cmpb    $0x0,0x2f05(%rip)           # 4010 <__TMC_END__>
1102: 75 2b                   jne     1138 <_do_global_ctors_aux+0x38>
1103: 55                      push    %rbp
1104: 48 83 3d e2 2e 00 00    cmpq    $0x0,0x2ee2(%rip)           # 3ff8 <__cxa_finalize@GLIBC_2.2.5>
1105: 00                      je
1106: 48 89 e5                mov     %rsp,%rbp
1107: 74 0c                   je      1127 <_do_global_ctors_aux+0x27>
1108: 48 8b 3d e6 2e 00 00    mov     0x2ee6(%rip),%rdi           # 4008 <__dso_handle>
1109: e8 19 ff ff ff         callq   1040 <__cxa_finalize@plt>
1110: e0 64 ff ff ff         callq   1000 <register_tm_clones>
1111: c5 05 dd 2e 00 00 01    movb    $0x1,0x2edd(%rip)           # 4010 <__TMC_END__>
1112: 5d                      pop     %rbp
1113: c3                      retq
1114: 0f 1f 00                nopl    (%rax)
1115: c3                      retq
1116: 0f 1f 80 00 00 00       nopl    0x0(%rax)

0000000000001140: <frame_dummy>
1140: f3 0f 1e fa             endbr64
1141: e9 77 ff ff ff         jmpq    10c0 <register_tm_clones>

0000000000001149: <main>
1149: f3 0f 1e fa             endbr64
1150: 55                      push    %rbp
1151: 48 89 e5                mov     %rsp,%rbp
1152: 48 8d 3d ac 0e 00 00    lea     0xaeac(%rip),%rdi           # 2004 <_IO_stdin_used+0x4>
1153: e8 f3 fe ff ff         callq   1050 <puts@plt>
1154: b8 00 00 00 00         mov     $0x0,%eax
1155: 5d                      pop     %rbp
1156: c3                      retq
1157: 66 2e 0f 1f 84 00 00    nopw    %cs:0x0(%rax,%rax,1)
1158: 00 00 00                nopw
1159: 66 90                   xchgb   %ax,%ax

0000000000001170: <__libc_csu_init>
1170: f3 0f 1e fa             endbr64
1171: 4c 57                   test    %eax,%eax
1172: 4c 8d 3b 3b 2c 00 00    lea     0xc23b(%rip),%r15           # 3db8 <__frame_dummy_init_array_entry>
1173: 41 56                   push    %r14
1174: 41 56                   push    %r14
1175: 49 89 d6                mov     %rdx,%r14
1176: 41 55                   push    %r13
1177: 49 89 f5                mov     %rsi,%r13
1178: 41 54                   push    %r12
1179: 41 89 fc                mov     %edi,%r12d
1180: 55                      push    %rbp
1181: 48 8d 2d 2c 2c 00 00    lea     0xc22c(%rip),%rbp           # 3dc0 <_do_global_ctors_aux_fini_array_entry>
1182: 53                      push    %ebx
```


C to Assembly cont.

Dump of assembler code for function main:

```
0x00000000000001149 <+0>:    endbr64
0x0000000000000114d <+4>:    push    rbp
0x0000000000000114e <+5>:    mov     rbp, rsp
0x00000000000001151 <+8>:    lea     rdi, [rip+0xeac]          # 0x2004
0x00000000000001158 <+15>:   call    0x1050 <puts@plt>
0x0000000000000115d <+20>:   mov     eax, 0x0
0x00000000000001162 <+25>:   pop     rbp
0x00000000000001163 <+26>:   ret
```

End of assembler dump.

Save previous frame pointer

Frame pointer is also the top
of our stack

0x2004

Move string "Hello world!"
into rdi register

Call puts function

Close stack frame and return
execution to previous location

Loops in Assembly

```
#include<stdio.h>
```

```
int main(){  
    int i;  
    for(i = 0; i < 10; i++){  
        puts("hello\n");  
    }  
}
```

Disassembled

Dump of assembler code for function main:

```
0x0000000000001149 <+0>:    endbr64  
0x000000000000114d <+4>:    push   rbp  
0x000000000000114e <+5>:    mov    rbp, rsp  
0x0000000000001151 <+8>:    sub    rsp, 0x10  
0x0000000000001155 <+12>:   mov     DWORD PTR [rbp-0x4], 0x0  
0x000000000000115c <+19>:   jmp     0x116e <main+37>  
0x000000000000115e <+21>:   lea     rdi, [rip+0xe9f]          # 0x2004  
0x0000000000001165 <+28>:   call    0x1050 <puts@plt>  
0x000000000000116a <+33>:   add     DWORD PTR [rbp-0x4], 0x1  
0x000000000000116e <+37>:   cmp     DWORD PTR [rbp-0x4], 0x9  
0x0000000000001172 <+41>:   jle     0x115e <main+21>  
0x0000000000001174 <+43>:   mov     eax, 0x0  
0x0000000000001179 <+48>:   leave  
0x000000000000117a <+49>:   ret
```

End of assembler dump.

Loops in Assembly cont.

Dump of assembler code for function main:

```
0x00000000000001149 <+0>:      endbr64
0x0000000000000114d <+4>:      push    rbp
0x0000000000000114e <+5>:      mov     rbp, rsp
0x00000000000001151 <+8>:      sub     rsp, 0x10
0x00000000000001155 <+12>:     mov     DWORD PTR [rbp-0x4], 0x0
0x0000000000000115c <+19>:     jmp     0x116e <main+37>
0x0000000000000115e <+21>:     lea     rdi, [rip+0xe9f]
0x00000000000001165 <+28>:     call    0x1050 <puts@plt>
0x0000000000000116a <+33>:     add     DWORD PTR [rbp-0x4], 0x1
0x0000000000000116e <+37>:     cmp     DWORD PTR [rbp-0x4], 0x9
0x00000000000001172 <+41>:     jle     0x115e <main+21>
0x00000000000001174 <+43>:     mov     eax, 0x0
0x00000000000001179 <+48>:     leave
0x0000000000000117a <+49>:     ret
```

Annotations:

- Make room for local variable "i" (points to `sub rsp, 0x10`)
- Initialize i to 0 (points to `mov DWORD PTR [rbp-0x4], 0x0`)
- Jump to comparison (points to `jmp 0x116e <main+37>`)
- Increment i (points to `add DWORD PTR [rbp-0x4], 0x1`)
- Continue loop if i < 9 (points to `jle 0x115e <main+21>`)

End of assembler dump.



Assembly Conclusion

- Assembly sucks, wouldn't it be great if there was a better way of pulling apart programs?
- Learning assembly is crucial for understanding the most low-level operations of a computer

Surely there must be a better way! This is the 21st century!

There is.



The background is a light gray with an abstract geometric pattern. On the left side, there is a dense network of thin gray lines connecting small dark gray dots, forming a complex web. Scattered across the entire background are various thin-lined triangles of different sizes and orientations. Some triangles are solid gray, while others are just outlines. In the top right corner, there are several small, faint circles or dots.

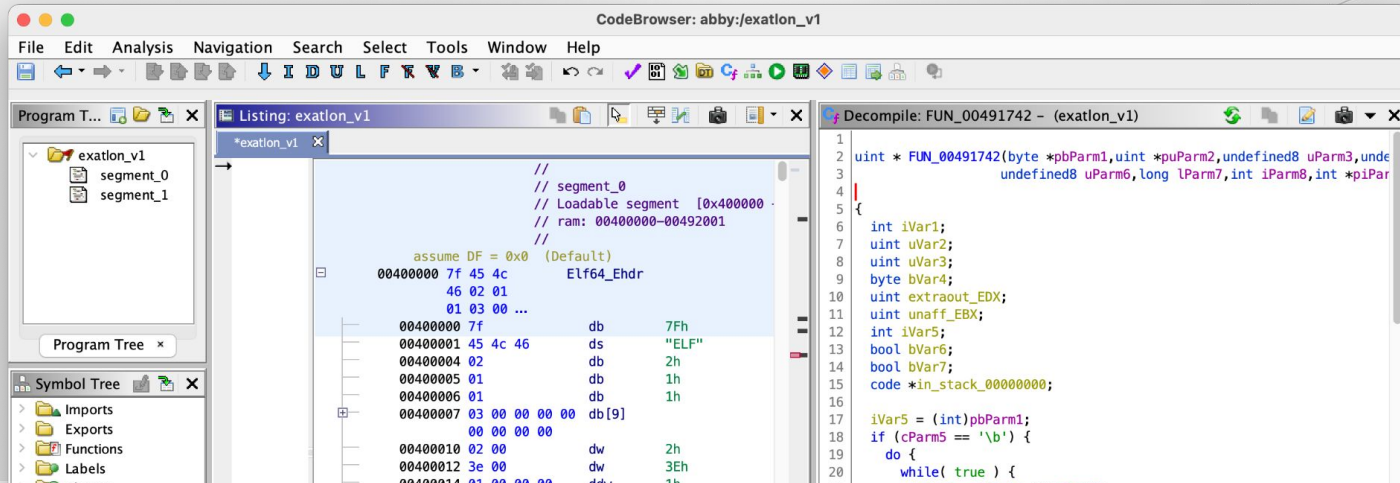
REAL Reverse Engineering

Time to bring out the big guns.

What is Reverse Engineering?

- Taking apart software without any source code and trying to figure out how it works
- A very useful and in-demand skill to have in the security field
- If we can pull apart a program and find vulnerabilities and exploits, so could evil hackers!

How can we use what we know to find exploits? Can we automate exploit discovery?



Decompilers

- Best invention of the century
- Analyze binaries and shows you the disassembly, but also shows you a nice view of what it **thinks** the source code may have looked like
- MUCH easier to read than assembly, but not always accurate

 Decompile: main – (hello_world)



```
1 |
2 | undefined8 main(void)
3 |
4 | {
5 |     _puts("Hello world!\n");
6 |     return 0;
7 | }
8 |
```


Decompilers cont

- Comes in many flavors
- You should start with Ghidra because it's free and open source



**IDA
Pro**

https://www.ninefx.com/img/products/ida_pro.png

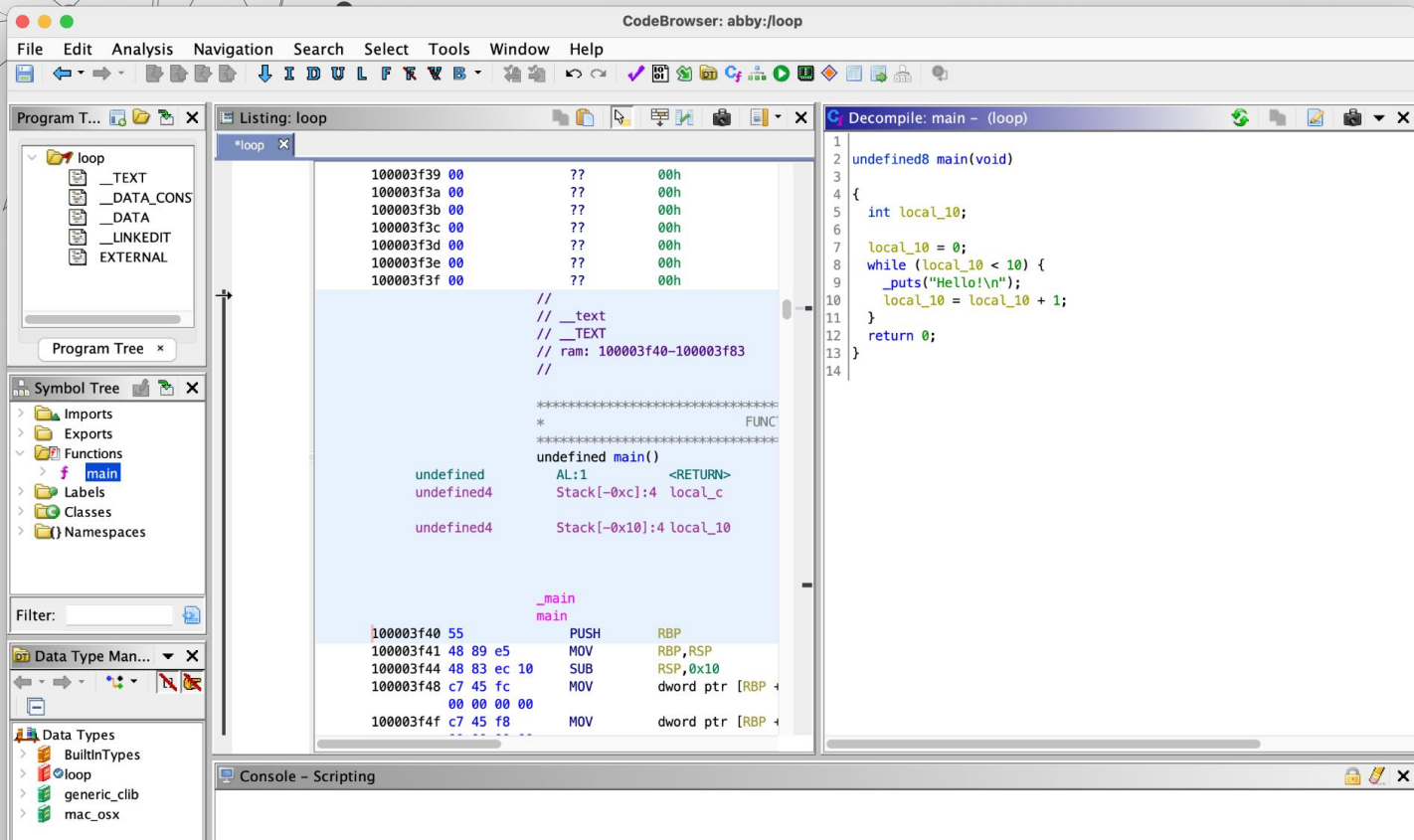


https://upload.wikimedia.org/wikipedia/commons/a/a3/Ghidra_Logo.png



<https://binary.ninja/ico/apple-icon-180x180.png>

Much better!



GDB

- Viewing the decompilation is a form of **static analysis**
- What if we want to view what a program is doing while it's running? We want **dynamic analysis**!
- And the best tool for that is the GNU Debugger (GDB)!
 - If you want to be really cool, use the GEF extension to get the most out of GDB
- From the GDB website:

GDB can do four main kinds of things (plus other things in support of these) to help you catch bugs in the act:

- **Start your program, specifying anything that might affect its behavior.**
- **Make your program stop on specified conditions.**
- **Examine what has happened, when your program has stopped.**
- **Change things in your program, so you can experiment with correcting the effects of one bug and go on to learn about another.**

You saw screenshots of GDB with GEF on previous slides!



GDB w/ GEF

```
gef> [ Legend: Modified register | Code | Heap | Stack | String ]

registers
$rax : 0x0000555555555149 → <main+0> endbr64
$rbx : 0x0000555555555170 → <_libc_csu_init+0> endbr64
$rcx : 0x0000555555555170 → <_libc_csu_init+0> endbr64
$rdx : 0x00007fffffffe278 → 0x00007fffffffe555 → "SHELL=/bin/bash"
$rsp : 0x00007fffffffe178 → 0x00007ffff7ded0b3 → <_libc_start_main+243> mov edi, eax
$rbp : 0x0
$rsi : 0x00007fffffffe268 → 0x00007fffffffe531 → "/home/griffith/assembly/hello_world"
$rdi : 0x1
$rip : 0x0000555555555149 → <main+0> endbr64
$r8 : 0x0
$r9 : 0x00007ffff7fe0d50 → endbr64
$r10 : 0x0
$r11 : 0x0
$r12 : 0x0000555555555060 → <_start+0> endbr64
$r13 : 0x00007fffffffe260 → 0x0000000000000001
$r14 : 0x0
$r15 : 0x0
$eflags: [ZERO carry PARITY adjust sign trap INTERRUPT direction overflow resume virtualx86 identification]
$cs: 0x0033 $ss: 0x002b $ds: 0x0000 $es: 0x0000 $fs: 0x0000 $gs: 0x0000

stack
0x00007fffffffe178 | +0x0000: 0x00007ffff7ded0b3 → <_libc_start_main+243> mov edi, eax ← $rsp
0x00007fffffffe180 | +0x0008: 0x00007ffff7ffc620 → 0x0005081200000000
0x00007fffffffe188 | +0x0010: 0x00007fffffffe268 → 0x00007fffffffe531 → "/home/griffith/assembly/hello_world"
0x00007fffffffe190 | +0x0018: 0x0000000010000000
0x00007fffffffe198 | +0x0020: 0x0000555555555149 → <main+0> endbr64
0x00007fffffffe1a0 | +0x0028: 0x0000555555555170 → <_libc_csu_init+0> endbr64
0x00007fffffffe1a8 | +0x0030: 0xefd2b5c60c0187b7
0x00007fffffffe1b0 | +0x0038: 0x0000555555555060 → <_start+0> endbr64

code:x86:64
0x555555555139 <_do_global_dtors_aux+57> nop     DWORD PTR [rax+0x0]
0x555555555140 <frame_dummy+0> endbr64
0x555555555144 <frame_dummy+4> jmp     0x5555555550c0 <register_tm_clones>
→ 0x555555555149 <main+0> endbr64
0x55555555514d <main+4> push    rbp
0x55555555514e <main+5> mov     rbp, rsp
0x555555555151 <main+8> lea     rdi, [rip+0xeac] # 0x555555556004
0x555555555158 <main+15> call    0x555555555050 <puts@plt>
0x55555555515d <main+20> mov     eax, 0x0

threads
[#0] Id 1, Name: "hello_world", stopped 0x555555555149 in main (), reason: BREAKPOINT

trace
[#0] 0x555555555149 → main()

gef>
```

Concolic Analysis

- Super modern technology allows us to combine both static and dynamic execution to create: concolic analysis
- Angr is a concolic analysis framework that uses the Z3 theorem prover to essentially simulate the execution of a binary, giving us the ability to analyze maximum code coverage and use symbolic execution to automate exploit discovery

This could be its own talk (or 4 credit course) so check out [their documentation](#) if you're interested.

If you flex your Angr ability in an interview, jaws will hit the floor.



https://angr.io/img/angry_face.png

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The background of the slide features a complex, abstract geometric pattern. It consists of numerous thin, light gray lines that connect various points, creating a network-like structure. Some of these points are highlighted as larger, solid dark gray circles, while others are smaller dots. The overall effect is a modern, tech-inspired aesthetic. The text 'LIVE DEMO' is centered in the upper half of the slide, rendered in a large, bold, dark gray sans-serif font.

LIVE DEMO

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