



**UNIVERZITET U SARAJEVU**  
**ELEKTROTEHNIČKI FAKULTET SARAJEVO**

# **ZADATAK 1 – LV5**

## **OPERACIONA ISTRAŽIVANJA**

**Studenti: Mašović Haris, Muminović Amir**

**Indeks: 1689/17993, 1661/17744**

**Odsjek: Računarstvo i Informatika**

**Datum:**

**04.02.2020**

**Potpisi:**

---

## Zadatak 1

Brod od 10 tona treba natovariti s tri vrste robe. U tabeli je data jedinična težina u tonama za svaku vrstu robe, raspoloživa količina i cijena po utovarenom komadu robe u hiljadama KM. Kako treba natovariti brod da se dobije maksimalan profit od utovarene robe i koja je suma u pitanju? Ukoliko trgovac bude prinuđen da prebaci robu koristeći dva broda od 4 i 6 tona kako će tada izgledati optimalni utovar i hoće li doći do promjene u profitu? Koristiti dinamičko programiranje prilikom rješavanja ovog problema. Ukoliko ima više rješenja dovoljno je navesti samo jedno uz napomenu da nije jedinstveno.

Vrsta robe ( $i$ )	1	2	3
Težina ( $w_i$ ) [t]	2	3	1
Količina ( $t_i$ )	4	3	5
Cijena ( $c_i$ ) [hiljadu KM]	31	35	20

### Rješenje:

Matematički model problema:

$$\arg \max Z = 31x_1 + 35x_2 + 20x_3$$

$$po. \quad 2x_1 + 3x_2 + x_3 \leq 10$$

$$0 \leq x_1 \leq 4$$

$$0 \leq x_2 \leq 3$$

$$0 \leq x_3 \leq 5$$

$$x_1 \in Z, x_2 \in Z, x_3 \in Z$$

Rekurzivna jednačina Bellmana (raspisana po iteracijama):

$$z(v, i) = \max \left\{ c_i x_i + z(v - w_i x_i, i - 1) \mid x_i = 0, \dots, \min \left\{ t_i, \left\lfloor \frac{v}{w_i} \right\rfloor \right\} \right\}, i = 1, 2, 3$$

$$z(v, 0) = 0$$

$$z(v, 1) = \max \left\{ c_1 x_1 + z(v - w_1 x_1, 0) \mid x_1 = 0, \dots, \min \left\{ t_1, \left\lfloor \frac{v}{w_1} \right\rfloor \right\} \right\}$$

$$= \max \left\{ 31x_1 + z(v - 2x_1, 0) \mid x_1 = 0, \dots, \min \left\{ 4, \left\lfloor \frac{v}{2} \right\rfloor \right\} \right\}$$

$$= \max \left\{ 31x_1 \mid x_1 = 0, \dots, \min \left\{ 4, \left\lfloor \frac{v}{2} \right\rfloor \right\} \right\}$$

$$z(v, 2) = \max \left\{ c_2 x_2 + z(v - w_2 x_2, 1) \mid x_2 = 0, \dots, \min \left\{ t_2, \left\lfloor \frac{v}{w_2} \right\rfloor \right\} \right\}$$

$$= \max \left\{ 35x_2 + z(v - 3x_2, 1) \mid x_2 = 0, \dots, \min \left\{ 3, \left\lfloor \frac{v}{3} \right\rfloor \right\} \right\}$$

$$z(v, 2) = \max \left\{ c_2 x_2 + z(v - w_2 x_2, 1) \mid x_2 = 0, \dots, \min \left\{ t_2, \left\lfloor \frac{v}{w_2} \right\rfloor \right\} \right\}$$

$$= \max \left\{ 35x_2 + z(v - 3x_2, 1) \mid x_2 = 0, \dots, \min \left\{ 3, \left\lfloor \frac{v}{3} \right\rfloor \right\} \right\}$$

$$z(v, 3) = \max \left\{ c_3 x_3 + z(v - w_3 x_3, 2) \mid x_3 = 0, \dots, \min \left\{ t_3, \left\lfloor \frac{v}{w_3} \right\rfloor \right\} \right\}$$

$$= \max \left\{ 20x_3 + z(v - x_3, 2) \mid x_3 = 0, \dots, \min \left\{ 5, \left\lfloor \frac{v}{1} \right\rfloor \right\} \right\}$$

Za i = 1:

$$z(v, 1) = \max \left\{ 31x_1 \mid x_1 = 0, \dots, \min \left\{ 4, \left\lfloor \frac{v}{2} \right\rfloor \right\} \right\} = 31 \min \left\{ 4, \left\lfloor \frac{v}{2} \right\rfloor \right\}$$

v	x	Z(v,1)
0	0	0
1	0	0
2	1	31
3	1	31
4	2	62
5	2	62
6	3	93
7	3	93
8	4	124
9	4	124
10	4	124

Za i = 2:

$$z(v, 2) = \max \left\{ 35x_2 + z(v - 3x_2, 1) \mid x_2 = 0, \dots, \min \left\{ 3, \left\lfloor \frac{v}{3} \right\rfloor \right\} \right\}$$

v	$(x_2, v - 3x_2)$	$35x_2 + z(v - 3x_2, 1)$	Z(v,2)
0	(0,0)	0	0
1	(0,1)	0 + z(1,1)=0	0
2	(0,2)	0+z(2,1)=31	31
3	(0,3), (1,0)	31, 35	35
4	(0,4), (1,1)	62, 35	62
5	(0,5), (1,2)	62, 66	66
6	(0,6), (1,3), (2,0)	93, 66, 70	93
7	(0,7), (1,4), (2,1)	93, 97, 70	97
8	(0,8), (1,5), (2,2)	124, 97, 101	124
9	(0,9), (1,6), (2,3), (3,0)	124, 97, 101, 105	124
10	(0,10), (1,7), (2,4), (3,1)	124, 128, 70+62, 105	132

Za  $i = 3$ :

$$z(v, 3) = \max \left\{ 20x_3 + z(v - x_3, 2) \mid x_3 = 0, \dots, \min \left\{ 5, \left\lfloor \frac{v}{1} \right\rfloor \right\} \right\}$$

v	$(x_3, v - x_3)$	$20x_3 + z(v - x_3, 2)$	$Z(v, 2)$
0	(0,0)	0	0
1	(0,1), (1,0)	0, 20	20
2	(0,2), (1,1), (2,0)	31, 20, 40	40
3	(0,3), (1,2), (2,1), (3,0)	35, 51, 40, 60	60
4	(0,4), (1,3), (2,2), (3,1), (4,0)	62, 55, 71, 60, 80	80
5	(0,5), (1,4), (2,3), (3,2), (4,1), (5,0)	66, 82, 75, 91, 80, 100	100
6	(0,6), (1,5), (2,4), (3,3), (4,2), (5,1)	93, 86, 102, 95, 111, 100	111
7	(0,7), (1,6), (2,5), (3,4), (4,3), (5,2)	97, 113, 106, 122, 115, 131	131
8	(0,8), (1,7), (2,6), (3,5), (4,4), (5,3)	124, 117, 133, 126, 142, 135	142
9	(0,9), (1,8), (2,7), (3,6), (4,5), (5,4)	124, 144, 137, 153, 146, 162	162
10	(0,10), (1,9), (2,8), (3,7), (4,6), (5,5)	132, 144, 164, 157, 173, 166	173

Optimalna vrijednost funkcije cilja  $Z = 173$ . Optimalna strategija:

- Za  $i=3$  i  $v = 10$  maksimum za  $x_3 = 4$ , preostaje kapacitet  $10-4 = 6$ .
- Za  $i=2$  i  $v = 6$  maksimum za  $x_2 = 0$  što ostavlja ranac neizmjenjenog kapaciteta.
- Za  $i = 1$  i  $v = 6$  maksimum za  $x_1 = 3$  preostaje kapacitet ranca  $v=0$  (u potpunosti popunjen ranac).

Optimalno rješenje:

$$x_1 = 3, x_2 = 0, x_3 = 4, Z = 173$$

Za dva broda od 4 i 6 tona imamo sljedeće. Za brod od 4 tone:

Za  $i = 1$ :

$$z(v, 1) = \max \left\{ 31x_1 \mid x_1 = 0, \dots, \min \left\{ 4, \left\lfloor \frac{v}{2} \right\rfloor \right\} \right\} = 31 \min \left\{ 4, \left\lfloor \frac{v}{2} \right\rfloor \right\}$$

v	x	$Z(v, 1)$
0	0	0
1	0	0
2	1	31
3	1	31
4	2	62

Za  $i = 2$ :

$$z(v, 2) = \max \left\{ 35x_2 + z(v - 3x_2, 1) \mid x_3 = 0, \dots, \min \left\{ 3, \left\lfloor \frac{v}{3} \right\rfloor \right\} \right\}$$

v	$(x_2, v - 3x_2)$	$35x_2 + z(v - 3x_2, 1)$	$Z(v, 2)$
0	(0,0)	0	0
1	(0,1)	$0 + z(1,1)=0$	0
2	(0,2)	$0 + z(2,1)=31$	31
3	(0,3), (1,0)	31, 35	35
4	(0,4), (1,1)	62, 35	62

Za  $i = 3$ :

$$z(v, 3) = \max \left\{ 20x_3 + z(v - x_3, 2) \mid x_3 = 0, \dots, \min \left\{ 5, \left\lfloor \frac{v}{1} \right\rfloor \right\} \right\}$$

v	$(x_3, v - x_3)$	$20x_3 + z(v - x_3, 2)$	$Z(v, 2)$
0	(0,0)	0	0
1	(0,1), (1,0)	0, 20	20
2	(0,2), (1,1), (2,0)	31, 20, 40	40
3	(0,3), (1,2), (2,1), (3,0)	35, 51, 40, 60	60
4	(0,4), (1,3), (2,2), (3,1), (4,0)	62, 55, 71, 60, 80	80

Optimalna vrijednost funkcije cilja  $Z = 80$ . Optimalna strategija:

- Za  $i=3$  i  $v = 4$  maksimum je postignut za  $x_3 = 4$ , kapacitet ranca  $v=0$  (u potpunosti popunjen ranac).

Optimalno rješenje za brod kapaciteta 4 glasi:

$$x_1 = 0, x_2 = 0, x_3 = 4, Z = 80$$

Za brod od 6 tona imamo:

Za  $i = 1$ :

$$z(v, 1) = \max \left\{ 31x_1 \mid x_1 = 0, \dots, \min \left\{ 4, \left\lfloor \frac{v}{2} \right\rfloor \right\} \right\} = 31 \min \left\{ 4, \left\lfloor \frac{v}{2} \right\rfloor \right\}$$

v	x	$Z(v, 1)$
0	0	0
1	0	0
2	1	31
3	1	31
4	2	62
5	2	62
6	3	93

Za  $i = 2$ :

$$z(v, 2) = \max \left\{ 35x_2 + z(v - 3x_2, 1) \mid x_3 = 0, \dots, \min \left\{ 3, \left\lfloor \frac{v}{3} \right\rfloor \right\} \right\}$$

v	$(x_2, v - 3x_2)$	$35x_2 + z(v - 3x_2, 1)$	$Z(v, 2)$
0	(0,0)	0	0
1	(0,1)	$0 + z(1,1)=0$	0
2	(0,2)	$0 + z(2,1)=31$	31
3	(0,3), (1,0)	31, 35	35
4	(0,4), (1,1)	62, 35	62
5	(0,5), (1,2)	62, 66	66
6	(0,6), (1,3), (2,0)	93, 66, 70	93

Za  $i = 3$ :

$$z(v, 3) = \max \left\{ 20x_3 + z(v - x_3, 2) \mid x_3 = 0, \dots, \min \left\{ 1, \left\lfloor \frac{v}{1} \right\rfloor \right\} \right\}$$

v	$(x_3, v - x_3)$	$20x_3 + z(v - x_3, 2)$	$Z(v, 2)$
0	(0,0)	0	0
1	(0,1), (1,0)	0, 20	20
2	(0,2), (1,1)	31, 20	31
3	(0,3), (1,2)	35, 51	51
4	(0,4), (1,3)	62, 55	62
5	(0,5), (1,4)	66, 82	82
6	(0,6), (1,5)	93, 86	93

Optimalna vrijednost funkcije cilja  $Z = 93$ .

Optimalna strategija:

- Za  $i=3$  i  $v=6$  maksimum je postignut za  $x_3 = 0$  te kapacitet ostaje nepromjenjen
- Za  $i=2$  i  $v=6$  maksimum je postignut za  $x_2 = 0$  te kapacitet ostaje nepromjenjen
- Za  $i=1$   $v=6$  maksimum je postignut za  $x_1 = 3$  nakon čega je kapacitet ranca  $v=0$  (u potpunosti popunjen ranac).

Optimalno rješenje za brod kapaciteta 6 glasi:

$$x_1 = 3, x_2 = 0, x_3 = 0, Z = 93$$

Vidimo da smo opet postigli isto rješenje kao i u slučaju kada imamo samo jedan brod kapaciteta 10 jer je  $93+80=173$  te smo uzimali iste predmete. Rješenje nije jedinstveno.