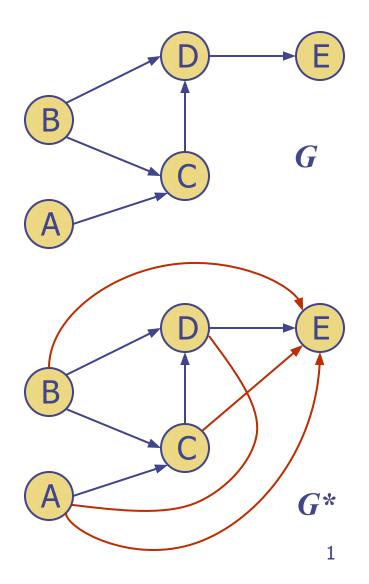
Transitive Closure

- Given a digraph G, the transitive closure of G is the digraph G^* such that
 - G* has the same vertices
 as G
 - if G has a directed path from u to v ($u \neq v$), G* has a directed edge from u to v
- The transitive closure provides reachability information about a digraph



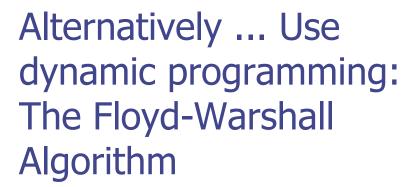
Computing the Transitive Closure

IWW.GENIUS.

We can performDFS starting at each vertex

O(n(n+m))

If there's a way to get from A to B and from B to C, then there's a way to get from A to C.



Floyd-Warshall Transitive Closure

- □ Idea #1: Number the vertices 1, 2, ..., n.
- Idea #2: Consider paths that use only vertices numbered 1, 2, ..., k, as intermediate vertices:

Uses only vertices numbered 1,...,k (add this edge if it's not already in)

Uses only vertices numbered 1,...,k-1

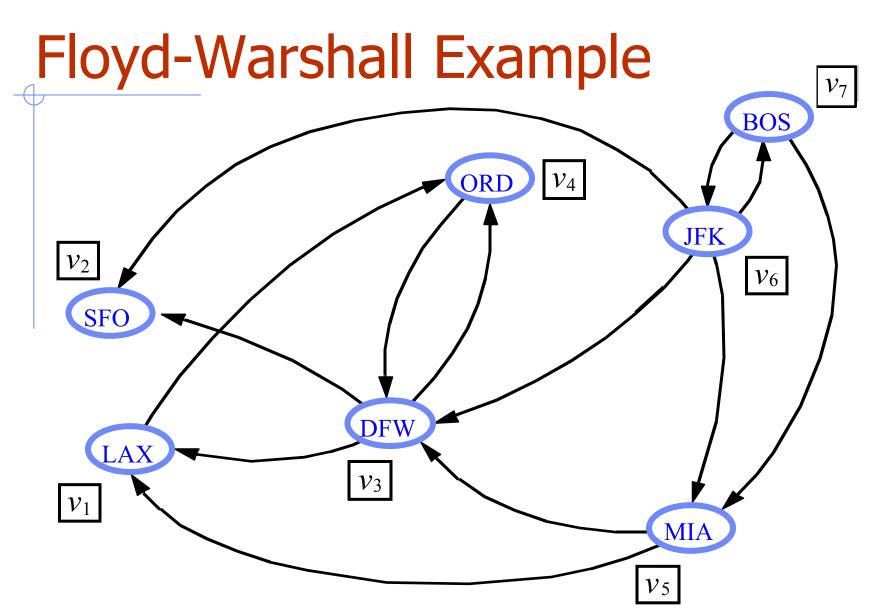
Uses only vertices numbered 1,...,k-1





- \neg Number vertices $v_1, ..., v_n$
- \Box Compute digraphs $G_0, ..., G_n$
 - $\mathbf{G}_0 = \mathbf{G}$
 - G_k has directed edge (v_i, v_j) if G has a directed path from v_i to v_j with intermediate vertices in $\{v_1, ..., v_k\}$
- □ We have that $G_n = G^*$
- In phase k, digraph G_k is computed from G_{k-1}
- Running time: $O(n^3)$, assuming areAdjacent is O(1) (e.g., adjacency matrix)

```
Algorithm FloydWarshall(G)
      Input digraph G
      Output transitive closure G^* of G
      i \leftarrow 1
      for all v \in G.vertices()
             denote v as v,
             i \leftarrow i + 1
      G_0 \leftarrow G
      for k \leftarrow 1 to n do
             G_k \leftarrow G_{k-1}
              for i \leftarrow 1 to n \ (i \neq k) do
                   for j \leftarrow 1 to n \ (j \neq i, k) do
                          if G_{k-1}.areAdjacent(v_i, v_k)
_{1}.areAdjacent(v_{k}, v_{j})
\neg G_{k}.areAdjacent(v, v_{i})
```



Floyd-Warshall, Iteration 1 BOS JFK v_6 DFW V_3 MIA v_5

Floyd-Warshall, Iteration 2 BOS JFK v_6 DFW V_3 MIA v_5

Floyd-Warshall, Iteration 3 BOS ORD JFK v_6 DFW V_3 MIA

