CS 496: Homework 4

Due: 7 March, 11:55pm

1 Assignment Policies

Collaboration Policy. Homework will be done individually: each student must hand in their own answers. It is acceptable for students to collaborate in understanding the material but not in solving the problems or programming. Use of the Internet is allowed, but should not include searching for existing solutions.

Under absolutely no circumstances code can be exchanged between students. Excerpts of code presented in class can be used.

Assignments from previous offerings of the course must not be re-used. Violations will be penalized appropriately.

2 Assignment

This assignment consists in extending REC to allow for mutually recursive function definitions. The resulting language will be called REC-M. It modifies the concrete syntax for letrec as follows. The production

```
<Expression> ::= letrec <Identifier>(<Identifier>) = <Expression> in <Expression> in REC is replaced with:
```

```
<Expression> ::= letrec { <Identifier>( <Identifier>) = <Expression>} + in <Expression>
```

in REC-M. The expression $\{ < ldentifier > (< ldentifier >) = < Expression > \}^+$ above means that there may be 1 or more declarations. Here is an example of a valid program in REC-M:

Evaluating that expression should produce the result NumVal 1, meaning that 99 is indeed odd. If we replace 99 in the code above with 98 and evaluate the resulting expression, this time we should get NumVal 0 as a result. This is correct since 98 is not an odd number.

Note that the above expression is not syntactically valid in REC. To see this, try running it in the interpreter for REC.

3 Implementing REC-M

To facilitate the process of implementing REC-M a stub has been provided for you in Canvas. This stub has been obtained by taking the interpreter for REC and applying some changes. Here is a summary of the changes:

1. The parser.mly file has been updated so that the parser is capable of parsing expressions such as

Here is the result of parsing it:

Note that Letrec now has two arguments, Letrec of (dec list)*expr, where dec is a new type with the constructor Dec of (string*string*expr). Each dec in the first parameter of Letrec represents one of the function declarations in the first part of the letrec expression.

2. The environment datatype has been updated by creating a new sum type with the syntax

Instead of maintaining the environment as a list and calling functions to add to it when needed, you will construct a new env type with the necessary constructor, where the final parameter is the environment that you are extending.

You will have to update

1. lookup in the file ds.ml by implementing the find_dec function. It currently reads as follows:

This function utilizes OCaml's option type which has the syntax

```
type 'a option = Some of 'a | None
```

This can be seen as an alternative to error handling with "null" checks used in Java. You can return None if you didn't find what you were looking for or Some x if you did. The result can be unpacked by the caller, as seen in the following excerpt from lookup

```
...
2 | LetrecEnv (decs, saved_env) ->
    let dec = find_dec decs id in
4 match dec with
    | None -> lookup saved_env id
6 | Some (Ast.Dec(name, var, body)) -> ProcVal(var, body, env)
```

2. eval in the file interp.ml. The case that must be updated is this one:

```
...
2 | Letrec(decs, e2) ->
    (* TODO evaluate e2 with a new LetrecEnv *)
4 failwith "Implement me"
```

4 Submission instructions

Submit a file named HW4_<SURNAME>.zip through Canvas. Include only the supporting files uploaded into Canvas but where interp.ml and ds.ml have been completed, as described in this document. Please write your name in the source code using comments.