┌ 1: t8

On expection, how many comparisions will it make to use QUICKSELECT to find the minimum of the array.

Using the analysis in exercise 9(see below)2, we can get that

$$\sum_{1 \le i < j \le n} B_{i,j,k} = \sum_{i=1}^{k-1} \sum_{j=i+1}^{k} B_{i,j,k} + \sum_{i=k}^{n-1} \sum_{j=i+1}^{n} B_{i,j,k}$$

$$= k - H_k + n - (k+1) - H_{n-k+1}$$

$$= 1 - H_1 + n - 2 - H_n$$

$$= n - 2 - H_n$$

$$= n - \log(n) - 2 - o(1)$$

┌_{2: t9}

Derive a formula for $\mathbb{E}_{\pi}[C(\pi,k)]$, up to additive terms of order o(n). You might want to introduce $\kappa = k/n$.

• case $i < j \le k$:

$$\sum_{i=1}^{k-1} \sum_{j=i+1}^{k} B_{i,j,k} = \sum_{i=1}^{k-1} \frac{1}{k-i+1} * (k-i)$$

$$= \sum_{i=1}^{k-1} (1 - \frac{1}{k-i+1}) = k - (1 + \frac{1}{2} + \frac{1}{3} + \dots + \frac{1}{k})$$

$$= k - H_k.$$

• case $k \le i < j$:

$$\sum_{i=k}^{n-1} \sum_{j=i+1}^{n} B_{i,j,k} = \sum_{j=k+1}^{n} n \sum_{i=k}^{j-1} B_{i,j,k}$$

$$= \sum_{j=k+1}^{n} \frac{j-k+1}{*} (j-k)$$

$$= \sum_{j=k+1}^{n} (1 - \frac{1}{j-k+1})$$

$$= n - (k+1) - (1 + \frac{1}{2} + \frac{1}{3} + \dots + \frac{1}{n-k+1})$$

$$= n - (k+1) - H_{n-k+1}.$$

• case i < k < j:

$$\begin{split} \sum_{i=1}^{k-1} \sum_{j=k+1}^{n} \frac{1}{j-i+1} &< \sum_{i=1}^{k-1} \int_{k+1}^{n+1} \frac{1}{j-i} \mathrm{d}j \\ &= \sum_{i=1}^{k-1} \ln(n+1-i) - \ln(k+1-i) \\ &< \int_{i=0}^{k-1} \ln(n+1-i) - \ln(k-i) \mathrm{d}i \\ &= (n+1) \ln(n+1) - k \ln k - (n-k+2) \ln(n-k+2). \end{split}$$

On the other hand, similarly we can have

$$\sum_{i=1}^{k-1} \sum_{j=k+1}^{n} \frac{1}{j-i+1} > -(-k+n+2)\ln(-k+n+2) - (k+2)\ln(k+2) + (n+1)\ln(n+1) + 3\ln(3).$$

The difference between the upper bound and the lower bound is o(n), thus the inaccuracy is o(n). To write this formula more simply, we have

$$\sum_{i=1}^{k-1} \sum_{j=k+1}^{n} \frac{1}{j-i+1} = n \ln n - k \ln k - (n-k) \ln(n-k) + o(n).$$

Summing three terms up,

$$\mathbb{E}_{\pi}[C(\pi,k)] = n - 1 - H_k - H_{n-k+1} + n \ln n - k \ln k - (n-k) \ln(n-k) + o(n).$$