

Algorithm Design and Analysis

Not Strong Enough

March 7, 2020

1 Homework

1: Problem 4

[A Dynamic Programming Algorithm for the Binomial Coefficient] Using pseudocode, write a dynamic programming algorithm computing $\binom{n}{k}$. Implement it in python! What is its running time in terms of n and k ? Would you say your algorithm is efficient? Why or why not?

Algorithm 1 Calculate Binomial Coefficient Using DP

Create a 2-dimension array G

for $i = 0$ to n **do**

$G[i][0] = 1$

end for

for $i = 0$ to k **do**

$G[i][i] = 1$

end for

for $i = 1$ to n **do**

for $j = 1$ to $\min(i-1, k)$ **do**

$G[i][j] = G[i-1][j] + G[i-1][j-1]$

end for

end for

return $G[n][k]$

```
def calc_dp(n,k):
    arr = [[0 for i in range(k+1)] for j in range(n+1)]
    for i in range(n+1):
        arr[i][0]=1
    for i in range(k+1):
        arr[i][i]=1
```

```

for i in range(1,n+1):
    for j in range(1,min(i-1,k)+1):
        arr[i][j]=arr[i-1][j]+arr[i-1][j-1]
return arr[n][k]

```

Complexity analysis:

When we traverse the array and visit `arr[i][j]`, we actually perform the add operation $\binom{i}{j} = \binom{i-1}{j} + \binom{i-1}{j-1}$. So the operation costs $O(\log \binom{i}{j})$ time. Using the Stirling Formula, we can estimate $\log \binom{i}{j} \approx (i + \frac{1}{2}) \log i - (i - j + \frac{1}{2}) \log(i - j) - (j + \frac{1}{2}) \log j = O(i)$. The number of nodes we visit is $O(k \cdot (2n - k)/2) = O(kn)$. Since we will only visit `arr[i][j]` once, we can estimate the total complexity as below: the upper bound is $O(kn \cdot n) = O(kn^2)$, the lower bound is $\Omega(kn)$. The algorithm is efficient, because there is no redundant calculation.

▮ **2: Problem 5**

[Binomial Coefficient modulo 2] Suppose we are only interested in whether $\binom{n}{k}$ is even or odd, i.e., we want to compute $\binom{n}{k} \bmod 2$. You could do this by computing $\binom{n}{k}$ using dynamic programming and then taking the result modulo 2. What is the running time? Would you say this algorithm is efficient? Why or why not?

The running time is the same as Problem 4. It's not efficient, because the cost of addition is still large. However, we can do modulo operation after every addition, so that the complexity of addition can be reduced to $O(1)$.