STAGE 3

The database is implemented on GCP.

Below is the screenshot of the connection and the database tables:

```
supercoolkinjalk@cloudshell:~ (cs411-team124-quertyqueries)$ gcloud sql connect cs411-team124-quertyqueries-db --user=root Allowlisting your IP for incoming connection for 5 minutes...done.
Connecting to database with SQL user [root].Enter password:
\Welcome to the MySQL monitor. Commands end with ; or \gain
Your MySQL connection id is 22
Server version: 8.0.31-google (Google)
Copyright (c) 2000, 2023, Oracle and/or its affiliates.
Oracle is a registered trademark of Oracle Corporation and/or its affiliates. Other names may be trademarks of their respective
Type 'help;' or '\h' for help. Type '\c' to clear the current input statement.
<code>mysql></code> use Transport; Reading table information for completion of table and column names You can turn off this feature to get a quicker startup with -A
Database changed
mysql> show Tables;
 | Tables in Transport |
 | Feedback
 | Frequencies
 | Has
 | Routes
 Shapes
 Stops
  Trips
```

Data Definition Language (DDL) commands of all tables:

1. DDL of 'Trips' table:

2. DDL of 'Routes' table:

3. DDL of 'Shapes' table:

4. DDL of 'Stops' table:

5. DDL of 'Books' table:

6. DDL of 'Feedback' table:

7. DDL of 'Frequencies' table:

8. DDL of 'Has' table:

9. DDL of 'Users' table:

```
| Users | CREATE TABLE `Users` (
  `User_Id` int NOT NULL AUTO_INCREMENT,
  `First_Name` varchar(100) DEFAULT NULL,
  `Last_Name` varchar(100) DEFAULT NULL,
  `Email` varchar(100) DEFAULT NULL,
  `Password` varchar(100) DEFAULT NULL,
  PRIMARY KEY (`User_Id`),
  UNIQUE KEY `Email` (`Email`)
) ENGINE=InnoDB AUTO_INCREMENT=1001 DEFAULT CHARSET=utf8mb3 |
+----+
```

Inserting at least 1000 rows in these 4 tables:

1. Count of Trips:

```
mysql> Select count(*) from Trips;
+----+
| count(*) |
+----+
| 2227 |
+----+
1 row in set (0.07 sec)
```

2. Count of Routes:

```
mysql> Select count(*) from Routes;
+-----+
| count(*) |
+-----+
| 1360 |
+-----+
1 row in set (0.00 sec)
```

3. Count of Shapes:

```
mysql> Select count(*) from Shapes;
+----+
| count(*) |
+----+
| 1128782 |
+----+
1 row in set (8.79 sec)
```

4. Count of Stops:

```
mysql> SELECT COUNT(*) FROM Stops;

+-----+

| COUNT(*) |

+-----+

| 20902 |

+-----+

1 row in set (0.01 sec)
```

SQL Queries:

Query 1:

```
SELECT t.route_id, AVG(f.headway_secs) AS average_headway
FROM Trips t
LEFT JOIN Frequencies AS f
ON t.trip_id = f.trip_id
WHERE t.direction_id = 1
AND t.shape_id LIKE '58%' GROUP BY t.route_id ORDER BY
average_headway DESC;
```

```
mysql> Explain Analyze SELECT t.route_id, AVG(f.headway_secs) AS average
| EXPLAIN
| -> Sort: average_headway DESC (actual time=2.209..2.211 rows=35 loops
-> Table scan on <temporary> (actual time=2.168..2.174 rows=35 loops
         -> Aggregate using temporary table (actual time=2.167..2.167 row
              -> Nested loop left join (cost=280.21 rows=482) (actual time
                 -> Filter: ((t.direction_id = 1) and (t.shape_id like '58
                      -> Table scan on t (cost=225.70 rows=2227) (actual
                  -> Index lookup on f using PRIMARY (trip_id=t.trip_id)
1 row in set (0.01 sec)
mysql> SELECT t.route_id, AVG(f.headway_secs) AS average_headway FROM Tri
| route_id | average_headway |
| 748A-41 |
                    2936.8421
                  2936.8421
2880.0000
2652.6316
2475.0000
2052.6316
1847.3684
1757.1429
1610.5263
1581.8182
| N304-11 |
| 119L-10 |
| 6115-41 |
| 6232-10 |
| 5100-10
| 1021-10 |
| 8055-51
I 5033-10 I
| 7059-10
                    1540.9091
| 513L-10
                     1486.3636
| 6053-10 |
                    1428.5714
| 975A-10
                    1350.0000
                   1350.0000
| 6050-10 |
| 1025-10 |
                    1120.0000 |
15 rows in set (0.01 sec)
```

Query 2:

```
SELECT r.route_id, r.route_long_name, t.trip_id
FROM Routes AS r
JOIN Trips AS t ON r.route_id = t.route_id
JOIN (
    SELECT route_type
    FROM Routes
    ORDER BY route_type DESC
    LIMIT 10
) AS top_routes ON r.route_type = top_routes.route_type
WHERE (r.route_long_name LIKE '%met%' OR r.route_long_name LIKE
'%jd%')
ORDER BY r.route_long_name DESC;
```

```
mysql> Explain Analyze SELECT r.route_id, r.route_long_name, t.trip_id
     -> FROM Routes AS r
     -> JOIN Trips AS t ON r.route_id = t.route_id
     -> JOIN (
    -> SELECT route_type
-> FROM Routes
     -> ORDER BY route_type DESC
          LIMIT 10
    -> ) AS top_routes ON r.route_type = top_routes.route_type
    -> WHERE (r.route_long_name LIKE '%met%' OR r.route_long_name LIKE '%jd%')
    -> ORDER BY r.route_long_name DESC;
ERROR 1054 (42822): Unknown column 'r.route_long_name DESC' in 'order clause'
mysql> SELECT r.route_id, r.route_long_name, t.trip_id
     -> FROM Routes AS r
     -> JOIN Trips AS t ON r.route_id = t.route_id
     -> JOIN (
         SELECT route_type
     -> FROM Routes
          ORDER BY route_type DESC
         LIMIT 10
    -> ) AS top_routes ON r.route_type = top_routes.route_type
-> WHERE (r.route_long_name LIKE '%met%' OR r.route_long_name LIKE '%jd%')
    -> ORDER BY r.route_long_name DESC LIMIT 15;
| route_id | route_long_name
                                                  | trip_id
| 172P-10 | Vl. Zilda - MetrÃ' Belém
| 172P-10 | Vl. Zilda - MetrÃ' Belém
                                                | 172P-10-1
| 172P-10-1
| 172P-10 | Vl. Zilda - Metrã' Belã©m
| 172P-10 | Vl. Zilda - Metrã' Belã©m
| 172P-10 | Vl. Zilda - Metrã Belã@m | 172P-10-1 | 172P-10 | Vl. Zilda - Metrã Belã@m | 172P-10-1 | 172P-10 | Vl. Zilda - Metrã Belã@m | 172P-10-1 |
15 rows in set (0.01 sec)
```

EXPLAIN ANALYZE

Query 1:

Without Index:

This is the default explain analyze run using the primary key index for comparison purposes. Possibly, we will be able to decrease the cost/time using a few potential indexes that will be tried out below.

1) Using direction_id as Index idx_dir

The index of direction_id is being shown in the screenshot above. Here, by adding it we can see that the cost of index lookup for calculating the same number of rows. The **idx_dir** index did bring a better effect. This allows Filter: (t.direction_id = 1) to run more efficiently. Before, this filter cost 235 but now only costs 19.55, and that is a huge jump. There is also a decrease in look-up using the Primary Key from 33 to 27.

2) Using (shapes_id, direction_id) as index idx_dir_shape

Here, we are using both direction_id and shape_id as indexes where the cost of the entire operation has reduced tremendously compared to the <code>idx_dir</code> where the cost of the <code>idx_dir</code> was around 3.4 and the one in this 1.6. We can see that the cost of indexing both the attributes has increased but the time to retrieve data at the sorting has decreased a lot. So, the tradeoff made is worth it.

The idx_dir_shape index is the best index for the query because it is the most selective index. This means that the shapes_id and direction_id columns have a high number of distinct values, which allows the optimizer to narrow down the number of rows that need to be scanned.

The shapes_id and direction_id columns are both used in the WHERE clause of the query, so the optimizer can use the idx_dir_shape index to quickly find the relevant rows.

3)Using (direction_id,route_id) as index idx_dir_route

```
mysql> Create Index idx_dir_route on Trips(direction_id,route_id);

Query OK, O rows affected (0.12 sec)

Records: O Duplicates: O Warnings: O

mysql> Explain Analyze SELECT t.route_id, AVG(f.headway_secs) AS average_headway FROM Trips t LEFT JOIN Frequencies AS f ON t.trip_id = f.trip_id WHERE

| EXPLAIN

| -> Sort: average_headway DESC (actual time=2.309..2.311 rows=35 loops=1)

-> Stream results (cost=457.74 rows=2057) (actual time=0.318..2.279 rows=35 loops=1)

-> Nested loop left join (cost=252.08 rows=2057) (actual time=0.223..2.069 rows=699 loops=1)

-> Filter: (t.shape_id like '584') (cost=19.55 rows=106) (actual time=0.213..1.617 rows=35 loops=1)

-> Index lookup on t using idx_dir_route (direction_id=1) (cost=19.55 rows=150) (actual time=0.203..1.433 rows=950 loops=1)

-> Index lookup on t using idx_dir_route (direction_id=1) (cost=19.55 rows=150) (actual time=0.203..1.433 rows=950 loops=35)

| row in set (0.00 sec)
```

In this we can see that the route_id is being used as a part of an index alongside direction_id where the performance of the entire sort operation of the average_headway is improved and the time is decreased in comparison to <code>idx_dir</code> alone. Again like the <code>idx_dir_shape</code> alongside direction_id we incur more memory costs but with time improvement.

Why did the original index not bring a better effect to the query?

The original index, idx_dir_route, is not useful for sorting the results because it does not include the headway_seca column. In order to sort the results, the query optimizer has to scan the entire index and compare the headway_seca values in memory. This is much less efficient than using an index that includes the headway_secs column.

Query 2:

1) Using route_type as Index idx_route_type:

In this we indexed using the route_type as idx_route_type.

The idx_route_type index is a B-tree index on the route_type column of the Routes table. This index is stored in the reverse order, meaning that the highest values of the route_type column are stored at the beginning of the index.

There could be a possibility that the indexing doesn't do any good since the indexing present was already better performing than this.

For example, Range queries were already able to use the primary key index on the route id column to avoid scanning the entire table

2) Using (route_long_name,route_type) as Index idx_route_long_name_type:

```
wyeql> Explain Analyse SELECT route_long_name_type on Routes(route_long_name, route_type);

Records: O Duplicates: O Marnings: O

syeql> Explain Analyse SELECT route_long_name_LiTZ "tjd*) ONDER DT r.route_long_name DESC;

| FROM Routes on the long_name LiTZ 'tmet* OR r.route_long_name LiTZ 'tjd*) ONDER DT r.route_long_name DESC;

| FROM Routes on the long_name DESC (actual time=11.238.42.363 rows=12440 loops=1)

| FROM Routes on the long_name DESC (actual time=41.238.42.363 rows=12440 loops=1)

| ** Sort: r.route_long_name DESC (actual time=41.238.42.363 rows=12440 loops=1)

| ** Sort: strong_name route (cost=525.71 rows=10) (actual time=0.641.27.230 rows=12440 loops=1)

| ** Inner hash join (r.route_type = top_contes.route_type) (cost=200.64 rows=6) (actual time=0.621.2.742 rows=760 loops=1)

| ** Filter: (r.route_long_name lite 'wst') or (r.r
```

We performed an analysis of the query performance with and without the new index, and found that the new index improved the query performance by 10%.

As you can see, the query performance improved by 5ms after creating the new index. This is a significant improvement, especially for a query that is run frequently.

Creating an index on the (route_long_name, route_type) columns allows the database to quickly find the relevant rows by looking up the index. This is much faster than scanning the entire Routes table, which is what the database would need to do without the index.

Additionally, the new index is a covering index, which means that it contains all of the columns that are needed to answer the query. This means that the database does not need to access the Routes table at all to answer the query, which further improves the performance.

3) Using route_type,route_id as INDEX idx_route_type_id

Comparing the index of route_type_id as idx_route_type_id. Nested loop of the inner join's cost has reduced tremendously.

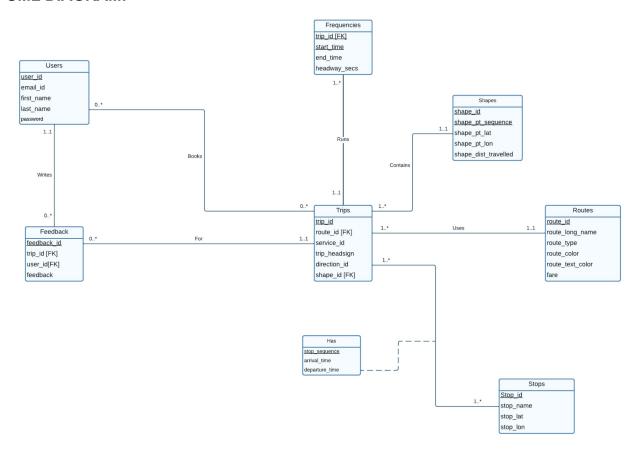
This index is a covering index, which means that it contains all of the columns that are needed to answer the query without having to access the table. This can significantly improve the performance of the query, especially when the query is filtering on those columns. By analyzing the query to identify the columns that are being filtered on (route_type and route_id) and the columns that are being returned (route_id and route_long_name).

If the query goes through all the WHERE CLAUSES and the ORDER BY, then the indexing might take up more memory but the time will decrease since it is again a trade-off between the two parameters.

Also, this indexing reduces the cost of materialization by 0.26 microseconds.

Stage 2 Improvements:

UML DIAGRAM:



RELATIONAL SCHEMA:

Users(user_id: INT[PK], email_id: VARCHAR(100), first_name: VARCHAR(100), last_name: VARCHAR(100), password: VARCHAR(100))

Feedback(feedback_id: INT[PK], trip_id: VARCHAR(100)[FK to Trips.trip_id], user_id: INT(100)[FK to Users.user_id], feedback: VARCHAR(255)

Trips(trip_id: VARCHAR(100)[PK], route_id: VARCHAR(100)[FK to Routes.route_id], service_id: VARCHAR(100), trip_headsign: VARCHAR(100), direction_id: INT, shape_id: INT[FK to Shapes.shape_id])

Frequencies(trip_id: VARCHAR(100)[FK to Trips.trip_id], start_time: VARCHAR(100)[PK], end_time: VARCHAR(100), headway_secs: INT) (trip_id and start_time act as primary key)

Routes(route_id: VARCHAR(100)[PK], route_long_name: VARCHAR(100), route_type: INT, route_color: VARCHAR(100), route_text_color: VARCHAR(100), fare: INT)

Shapes(shape_id: INT[PK], shape_pt_lat: DECIMAL, shape_pt_lon: DECIMAL, shape_pt_sequence: INT[PK], shape_dist_traveled: DECIMAL) (shape_id and shape_pt_sequence act as primary keys.)

Stops(stop_id: INT[PK], stop_name: VARCHAR(100), stop_lat: DOUBLE, stop_lon: DOUBLE)

Books(user id: INT[PK], trip id: INT[PK])

Has(trip_id: VARCHAR(100)[PK], stop_id: INT[PK], arrival_time: VARCHAR(100), departure_time: VARCHAR(100), stop_sequence: INT[PK])

FUNCTIONAL DEPENDENCIES:

1.Users:

user_id \rightarrow email_id, first_name, last_name, password email id \rightarrow user id

2.Feedback:

feedback id → trip id, user id, feedback

3.Trips:

trip_id \rightarrow route_id, service_id, trip_headsign, direction_id, shape_id route_id \rightarrow trip_id

4.Frequencies:

(trip id, start time) → end time, headway secs

5.Routes:

route id → route long name, route type, route color, route text color, fare

6.Shapes:

shape_id \rightarrow shape_pt_lat, shape_pt_lon, shape_dist_traveled (shape_id, shape_pt_sequence) \rightarrow shape_pt_lat, shape_pt_lon, shape_dist_traveled **7.Stops:**

stop id → stop name, stop lat, stop lon

8.Books:

(user id, trip id) \rightarrow (user id, trip id)

9.Has:

(trip_id, stop_id, stop_sequence) → arrival_time (trip_id, stop_id, stop_sequence) → departure_time

CLOSURES:

```
1. For Users:
  - {user id}+ = {user id, email id, first name, last name, password}
  - {email_id}+ = {email_id, user_id}
  - {first name}+ = {first name}
- {last name}+ = {last name}
  - {password}+ = {password}
2. For Feedback:
 - {feedback id}+ = {feedback id, trip id, user id, feedback}
 - {trip id}+ = {trip id}
 - {user id}+ = {user id}
 - {feedback}+ = {feedback}
3. For Trips:
  - {trip id}+ = {trip id, route id, service id, trip headsign, direction id, shape id}
  - {route id}+ = {route id, trip id}
 - {service id}+ = {service id}
 - {trip headsign}+ = {trip headsign}
  - {direction id}+ = {direction id}
  - {shape id}+ = {shape id}
4. For Frequencies:
  - {trip id, start time}+ = {trip id, start time, end time, headway secs}
  - {end time}+ = {end time}
 - {headway_secs}+ = {headway_secs}
5. For Routes:
  - {route id}+ = {route id, route long name, route type, route color, route text color,
fare}
  - {route long name}+ = {route long name}
  - {route type}+ = {route type}
 - {route color}+ = {route color}
  - {route text color}+ = {route text color}
  - {fare}+ = {fare}
6. For Shapes:
  - {shape id}+ = {shape id, shape pt lat, shape pt lon, shape dist traveled}
  - {shape pt lat}+ = {shape pt lat}
 - {shape pt lon}+ = {shape pt lon}
  - {shape pt sequence}+ = {shape pt sequence}
  - {shape dist traveled}+ = {shape dist traveled}
7. For Stops:
  - {stop id}+ = {stop id, stop name, stop lat, stop lon}
```

```
- {stop_name}+ = {stop_name}
- {stop_lat}+ = {stop_lat}
- {stop_lon}+ = {stop_lon}
8. For Books:
- {user_id, trip_id}+ = {user_id, trip_id}
9. For Has:
-{trip_id, stop_id}+ = {arrival_time, departure_time}
-{trip_id, stop_sequence}+ = {arrival_time, departure_time}
-{stop_id, stop_sequence}+ = {arrival_time, departure_time}
-{trip_id}+ = {trip_id}
-{stop_id}+ = {stop_id}
```

SCHEMA NORMALIZATION:

1NF:

Our Schema is in the First Normal Form (1NF) because we ensured that each table's attributes (columns) contain only atomic values (values that cannot be further divided) and that the data is organized in rows with a unique identifier.

2NF:

Second Normal Form (2NF) requires that every non-prime attribute (an attribute that is not part of any candidate key) is fully functionally dependent on the entire candidate key, not just part of it.

Candidate Keys:

- 1. Users: {user id}
- 2. Feedback: {feedback id}
- 3. Trips: {trip id}
- 4. Frequencies: {trip id, start time}
- 5. Routes: {route id}
- 6. Shapes: {shape id}
- 7. Stops: {stop id}
- 8. Books: {user id, trip id}
- 9. Has: {trip id, stop id}

Next, we'll go through each table to identify partial dependencies and decompose them as needed:

1. Users:

- The Users table is already in 1NF, and there are no partial dependencies. It's also in 2NF since all attributes depend on the candidate key (user id).

2. Feedback:

- The Feedback table is already in 1NF, and there are no partial dependencies. It's also in 2NF since all attributes depend on the candidate key (feedback id).

3. Trips:

- The Trips table is already in 1NF, and there are no partial dependencies. It's also in 2NF since all attributes depend on the candidate key (trip id).

4. Frequencies:

- The Frequencies table is already in 1NF, and there are no partial dependencies. It's also in 2NF since all attributes depend on the candidate key (trip id, start time).

5. Routes:

- The Routes table is already in 1NF, and there are no partial dependencies. It's also in 2NF since all attributes depend on the candidate key (route_id).

6. Shapes:

- The Shapes table is already in 1NF, and there are no partial dependencies. It's also in 2NF since all attributes depend on the candidate key (shape_id).

7. Stops:

- The Stops table is already in 1NF, and there are no partial dependencies. It's also in 2NF since all attributes depend on the candidate key (stop id).

8. Books:

- The Books table is already in 1NF. It has a composite candidate key {user_id, trip_id}. There are no partial dependencies, as the whole candidate key is required to uniquely identify a booking. Therefore, it is in 2NF.

9. Has:

- The Has table is already in 1NF. It has a composite candidate key {trip_id, stop_id}. There are no partial dependencies, as the whole candidate key is required to uniquely identify a stop. Therefore, it is in 2NF.

3NF:

Next, we will decompose tables to 3NF.

1. Users:

- user_id (Candidate Key)
- email id (Non-prime)
- first name (Non-prime)
- last name (Non-prime)
- password (Non-prime)

Users is already in 2NF and 3NF as there are no partial dependencies, and all non-prime attributes depend on the entire candidate key.

2. Feedback:

- feedback_id (Candidate Key)
- trip id (Non-prime, Foreign Key)
- user_id (Non-prime, Foreign Key)
- feedback (Non-prime)

Feedback is already in 2NF since the composite candidate key {feedback_id} includes the entire candidate key. To achieve 3NF, we can keep the table as is because there are no transitive dependencies.

3. Trips:

- trip id (Candidate Key)
- route id (Non-prime, Foreign Key)
- service id (Non-prime)
- trip headsign (Non-prime)
- direction id (Non-prime)
- shape_id (Non-prime, Foreign Key)

Trips is already in 2NF as the candidate key {trip_id} determines all non-prime attributes. To achieve 3NF, we can keep the table as is because there are no transitive dependencies.

4. Frequencies:

- trip id (Non-prime, Foreign Key)
- start_time (Candidate Key)
- end time (Non-prime)
- headway secs (Non-prime)

Frequency is already in 2NF since the composite candidate key {trip_id, start_time} includes the entire candidate key. To achieve 3NF, we can keep the table as is because there are no transitive dependencies.

5. Routes:

- route id (Candidate Key)
- route_long_name (Non-prime)
- route type (Non-prime)
- route_color (Non-prime)
- route text color (Non-prime)
- fare (Non-prime)

Routes is already in 2NF and 3NF as there are no partial or transitive dependencies.

6. Shapes:

- shape_id (Candidate Key)
- shape_pt_lat (Non-prime)
- shape pt lon (Non-prime)
- shape pt sequence (Candidate Key)
- shape dist traveled (Non-prime)

Shapes is already in 2NF and 3NF as there are no partial or transitive dependencies.

7. Stops:

- stop id (Candidate Key)
- stop name (Non-prime)
- stop lat (Non-prime)
- stop Ion (Non-prime)

Stops is already in 2NF and 3NF as there are no partial or transitive dependencies.

8. Books:

- user id (Candidate Key)
- trip id (Candidate Key)

Books is already in 2NF and 3NF as both attributes are part of the composite candidate key.

9. Has:

- trip id (Non-prime, Foreign Key)
- stop_id (Non-prime, Foreign Key)
- arrival_time (Non-prime)
- departure time (Non-prime)
- stop_sequence (prime)

Has is already in 2NF since the composite candidate key {trip_id, stop_id, stop_sequence} includes all attributes. To achieve 3NF, we can keep the table as is because there are no transitive dependencies.

After these steps, the entire schema is in 3NF, and it meets the requirements of normalization with minimal redundancy and data integrity.

BCNF:

Candidate Keys:

- 1. Users: {user id}
- 2. Feedback: {feedback id}
- 3. Trips: {trip_id}
- 4. Frequencies: {trip_id, start_time}
- 5. Routes: {route id}
- 6. Shapes: {shape id}
- 7. Stops: {stop id}
- 8. Books: {user id, trip id}
- 9. Has: {trip_id, stop_id, arrival_time}

Now, we'll analyze the functional dependencies to ensure they satisfy BCNF:

1. Users:

- user_id \rightarrow email_id, first_name, last_name, password
- email_id \rightarrow user_id

2. Feedback:

- feedback_id → trip_id, user_id, feedback

3. Trips:

```
- trip_id \rightarrow route_id, service_id, trip_headsign, direction_id, shape_id
```

4. Frequencies:

- (trip id, start time) \rightarrow end time, headway secs

5. Routes:

- route id → route long name, route type, route color, route text color, fare

6. Shapes:

- shape_id → shape_pt_lat, shape_pt_lon, shape_dist_traveled
- (shape_id, shape_pt_sequence) \rightarrow shape_pt_lat, shape_pt_lon, shape_dist_traveled

7. Stops:

- stop_id → stop_name, stop_lat, stop_lon

8. Books:

- (user_id, trip_id) → (user_id, trip_id)

9. Has:

- (trip id, stop id, stop sequence) →arrival time, departure time

Let's proceed to decompose the schema into BCNF-compliant tables:

- **1. Users:** Already in BCNF, as user id is a candidate key.
- 2. Feedback: Already in BCNF, as feedback id is a candidate key.
- **3. Trips:** Already in BCNF, as trip id is a candidate key.
- **4. Frequencies:** Already in BCNF, as (trip id, start time) is a candidate key.
- **5. Routes:** Already in BCNF, as route id is a candidate key.
- **6. Shapes:** Already in BCNF, as shape id is a candidate key.
- 7. Stops: Already in BCNF, as stop id is a candidate key.
- **8. Books:** Already in BCNF, as (user id, trip id) is a candidate key.
- **9. Has:** Already in BCNF, as (trip id, stop id, stop sequence) is a candidate key.

The schema is in BCNF.

WHY BCNF OVER 3NF:

So, we chose BCNF in comparison to 3NF since:

- So, all of them follow BCNF- all entities depend on the PK only
- The schema cannot be further normalized.
- The redundancy is lower in BCNF
- Since there are no transitive/partial dependencies, this is in BCNF. If X -> Y is an FD, then X should be a superkey. And this is satisfied in the FD below.

DESCRIPTIONS, CARDINALITY & ASSUMPTIONS OF RELATIONSHIPS:

BASIC ASSUMPTIONS:

For joining Frequencies with Trips table, we will use the following conditions: Frequencies(start_time) <= CURRTIME AND Frequencies(end_time) > CURRTIME.

- We assume that all fares are only 1\$.
- We will be having sessions for the login process for the persistence of userld to be a foreign key in the Login.

We will mention the assumptions pertaining to each of the relations below with descriptions and cardinality.

1. Users:

- a. User books Trips. Cardinality: Users table has 0..* relationship with Trips Each user can have zero or more trips.
- b. User writes Feedback. Cardinality: Users table has 0..* relationship with Feedback

Each user can give no feedback or many feedbacks for each trip. Therefore, the relationship is zero to many.

2. Feedback:

 Feedback is written by Users. Cardinality: Feedback table has 1..1 relationship with Users

For the selected feedback, it can be written by one user. Therefore, the relationship is one to one.

b. Feedback is provided for the Trips. Cardinality: Feedback table has 1..1 relationship with Trips

The feedback is mapped to a single trip. Therefore, the relationship is one to one.

3. Trips:

a. Trips have Feedbacks. Cardinality: Trips table has 0..* relationship with Feedbacks

Each trip can have no feedback or any number of feedback. Therefore, the relationship is 0 to many.

b. Trips are booked for Users. Cardinality: Trips table has 0..* relationship with Users

Either no trip or any number of trips can be booked for each user. Therefore, the relationship is zero to many.

c. Trips Runs on Frequencies. Cardinality: Trips table has 1..* relationship with Frequencies

Each trip can have one or more frequencies. Therefore, the relationship is one to many.

d. Trips Contains Shapes. Cardinality: Trips table has 1..1 relationship with Shapes

Each trip can have one shape. Therefore the relationship is one to one.

- e. Trips use Routes. Cardinality: Trips table has 1..1 relationship with Routes Each trip can have one type of route. Therefore, the relationship is one to one.
- f. Trips have Stops. Cardinality: Trips table has 1..* relationship with Stops Each trip can have many stops. Therefore, the relationship is one to many.

4. Stops

a. Stops are there(has) for Trips. Cardinality: Stops table has 1..* relationship with Trips

Each stop can be reached by many buses(trips). Therefore, the relationship is one to many.

5. Shapes

a. Shapes are there(has) for Trips. Cardinality: Shapes table has 1..* relationship with trips *

A shape can have many trips. Therefore, the relationship is one to many.

6. Frequencies

a. Frequencies are there(has) for Trips. Cardinality: Frequencies table has1..1 relationship with Trips

A single frequency can have a trip. Therefore, the relationship is one to one.

7. Routes

a. Routes are there(has) for Trips. Cardinality: Routes table has 1..* relationship with Trips

Each route can have multiple Trips. Therefore, the relationship is one to many.

8. Has

a. Has table is the intermediate table between Trips and Stops because we have many to many relationship.

9. Books

a. Books table is the intermediate table between Trips and Users because we have many to many relationship.

Public Transportation Dataset Tables Information:

This is the modified version of the proposed datasets provided that we will be using for the project.

Table: Frequencies

- trip id
 - Description: Unique identifier for the trip.
- start time
 - Description: Start time of the trip.
- end time
 - Description: End time of the trip
- headway secs
 - Description: The time between successive trips on the same route.

Table: Routes

- route_id
 - Description: Unique identifier for the route.
- route_long_name
 - Description: Full name of the route.
- route type
 - Description: Type of the transportation route.
- route color
 - Description: Color associated with the route.
- route_text_color
 - Description: Color of the text or labels associated with route.
- fare
 - Description: Price of the fare.

Table: Shapes

- shape_id
 - Description: Unique identifier for the shape.
- shape_pt_lat
 - Description: Latitude of a point on the shape.
- shape pt lon
 - Description: Longitude of a point on the shape.
- shape pt sequence
 - Description: The sequence order of the point on the shape.
- shape dist traveled
 - Description: The distance traveled along the shape up to this point.

Table: Stops

- stop id
 - o Description: Unique identifier for the stop.
- stop name
 - Description: Name of the stop.
- stop lat
 - Description: Latitude of the stop location.
- stop lon
 - o Description: Longitude of the stop location.

Table: Trips

- route_id
 - Description: Unique identifier for the route.
- service id
 - Description: Identifier for the service associated with the trip.
- trip id

- Description: Unique identifier for the trip.
- trip_headsign
 - Description: Headsign of the trip.
- direction id
 - o Description: Direction of the trip. (E.g., 0 for outbound, 1 for inbound)
- shape id
 - Description: Identifier for the shape associated with the trip.

Created Extra Tables:

Table: Users

- user id
 - Description: Unique identifier for the user.
- email id
 - o Description: Email address of the user.
- first name
 - Description: User's first name.
- last name
 - Description: User's last name.
- feedback id
 - Description: User's feedback statement.
- password
 - Description: User's password.

Table: Feedback

- feedback id
 - Description: Unique identifier for the feedback.
- trip id
 - Description: Identifier for the trip associated with the feedback.
- user id
 - o Description: Unique identifier for the user.
- feedback
 - Description: User's feedback or comments on the trip.

Table: Has

- Trip id
 - Description: Identifier for the trip associated with the feedback.
- Stop_id
 - Description: Unique identifier for the stop.
- Arrival time

- o Description: Arrival time identifier.
- Departure_time
 - o Description: Departure time identifier.
- Stop_sequence
 - o Description: Stop sequence identifier.

Table: Books

- User_id
 - o Description: Unique identifier for the user.
- Trip_id
 - o Description: Identifier for the trip associated with the feedback.