# Team Notebook

# Pontificia Universidad Católica de Chile - Bella y Sensual

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#### 1 data structures

#### 1.1 dsu

```
struct Dsu {
   vector<int> p; Dsu(int N = 0) : p(N, -1) {}
   int get(int x) { return p[x] < 0 ? x : get(p[x]); }</pre>
   bool sameSet(int a, int b) { return get(a) == get(b); }
   int size(int x) { return -p[get(x)]; }
   vector<vector<int>> S;
   void unite(int x, int y) {
       if ((x = get(x)) == (y = get(y)))
           return S.push_back({-1});
       if (p[x] > p[y]) swap(x, y);
       S.push_back(\{x, y, p[x], p[y]\});
       p[x] += p[y], p[y] = x;
   void rollback() {
       auto a = S.back(); S.pop_back();
       if (a[0] != -1) p[a[0]] = a[2], p[a[1]] = a[3];
};
```

### 1.2 fenwick-tree

```
int ft[MAXN+1]; // add dimension for multi-d
void upd(int i0, int v){ // add v to i0th element
    for(int i=i0+1;i<=MAXN;i+=i&-i)ft[i]+=v;//+ fors
}
int get(int i0){ // get sum of range [0,i0)
    int r=0; // add fors
    for(int i=i0;i;i-=i&-i)r+=ft[i];
    return r;
}
int get_sum(int i0,int i1){//sum of [i0,i1)
    return get(i1)-get(i0);
}</pre>
```

### 1.3 link-cut-tree

```
const int N_DEL = 0, N_VAL = 0; //delta, value
inline int mOp(int x, int y){return x+y;}//modify
inline int qOp(int lval, int rval){return lval + rval;}//
    query
inline int dOnSeg(int d, int len){return d==N_DEL ? N_DEL :
    d*len;}
```

```
//mostly generic
inline int joinD(int d1, int d2){
 if(d1==N_DEL)return d2;if(d2==N_DEL)return d1;return mOp(
inline int joinVD(int v, int d){return d==N_DEL ? v : mOp(v,
     d):}
struct Node_t{
 int sz, nVal, tVal, d; bool rev;
 Node_t *c[2], *p;
 Node_t(int v) : sz(1), nVal(v), tVal(v), d(N_DEL), rev(0),
       }(0)a
 c[0]=c[1]=0:
 bool isRoot(){return !p || (p->c[0] != this && p->c[1] !=
 void push(){
 if(rev){
   rev=0; swap(c[0], c[1]); fore(x,0,2)if(c[x])c[x]->rev^=1;
 nVal=joinVD(nVal, d); tVal=joinVD(tVal, dOnSeg(d, sz));
 fore(x,0,2)if(c[x])c[x]->d=joinD(c[x]->d, d);
 d=N_DEL;
 }
 void upd();
typedef Node_t* Node;
int getSize(Node r){return r ? r->sz : 0;}
int getPV(Node r){
 return r ? joinVD(r->tVal, dOnSeg(r->d,r->sz)) : N_VAL;}
void Node t::upd(){
 tVal = qOp(qOp(getPV(c[0]), joinVD(nVal, d)), getPV(c[1]))
 sz = 1 + getSize(c[0]) + getSize(c[1]);
void conn(Node c, Node p, int il){if(c)c->p=p;if(il>=0)p->c
    [!ill=c:}
void rotate(Node x){
 Node p = x-p, g = p-p;
 bool gCh=p->isRoot(), isl = x==p->c[0];
 conn(x->c[isl],p,isl); conn(p,x,!isl);
 conn(x,g,gCh?-1:(p==g->c[0])); p->upd();
void spa(Node x){//splay
 while(!x->isRoot()){
 Node p = x-p, g = p-p;
 if(!p->isRoot())g->push();
 p->push(); x->push();
 if(!p-)isRoot())rotate((x=-p-)c[0])==(p=-p-)c[0])? p : x);
 rotate(x);
```

```
x->push(): x->upd():
Node exv(Node x){//expose
 Node last=0:
 for (Node y=x; y; y=y->p) spa(y), y->c[0]=last, y->upd(), last=
 spa(x);
 return last;
void mkR(Node x){exv(x);x->rev^=1;}//makeRoot
Node getR(Node x){exv(x); while(x->c[1])x=x->c[1]; spa(x);
Node lca(Node x, Node y){exv(x); return exv(y);}
bool connected(Node x, Node y){exv(x);exv(y); return x==y?1:
void link(Node x, Node y){mkR(x); x->p=y;}
void cut(Node x, Node y) \{mkR(x); exv(y); y->c[1]->p=0; y->c
    [1]=0;}
Node father(Node x){
 exv(x): Node r=x->c[1]:
 if(!r)return 0;
 while (r->c[0])r=r->c[0];
 return r;
void cut(Node x){ // cuts x from father keeping tree root
 exv(father(x));x->p=0;
int query(Node x, Node y){mkR(x); exv(y); return getPV(y);}
void modifv(Node x. Node v. int d){mkR(x):exv(v):v->d=joinD(
    y->d,d);}
Node lift rec(Node x. int t){
 if(!x)return 0;
 if(t==getSize(x->c[0])){spa(x);return x;}
 if(t<getSize(x->c[0]))return lift_rec(x->c[0],t);
 return lift_rec(x->c[1],t-getSize(x->c[0])-1);
Node lift(Node x, int t) { // t-th ancestor of x (lift(x.1)
    is x's father)
 exv(x);return lift_rec(x,t);}
int depth(Node x){ // distance from x to its tree root
 exv(x);return getSize(x)-1;}
```

## 1.4 persistent-segment-tree-lazy

```
template <class T>
struct Node {
    T x, lz;
    int 1 = -1, r = -1;
};
```

```
template <class T>
struct Pstl {
   int N:
   vector<Node<T>> a:
   vector<int> head;
   T qneut() { return 0; }
   T merge(T 1, T r) { return 1 + r; }
   T uneut() { return 0; }
   T accum(T u, T x) { return u + x; }
   T apply(T x, T lz, int l, int r) { return x + (r - 1) *
   int build(int vl. int vr) {
      if (vr - vl == 1) a.push_back({qneut(), uneut()}); //
            node construction
      else {
          int vm = (vl + vr) / 2, l = build(vl, vm), r =
               build(vm. vr):
          a.push_back({merge(a[1].x, a[r].x), uneut(), 1, r
               }); // query merge
      }
       return a.size() - 1;
   T query(int 1, int r, int v, int v1, int vr, T acc) {
      if (1 >= vr || r <= vl) return qneut();</pre>
           // guery neutral
      if (1 \le vl \&\& r \ge vr) return apply(a[v].x, acc, vl,
            vr): // update op
      acc = accum(acc, a[v].lz);
           // update merge
      int vm = (vl + vr) / 2:
      return merge(query(1, r, a[v].1, v1, vm, acc), query(
           1, r, a[v].r, vm, vr, acc)); // query merge
   int update(int 1, int r, T x, int v, int v1, int vr) {
      if (1 >= vr || r <= vl || r <= 1) return v;
      a.push_back(a[v]);
      v = a.size() - 1:
      if (1 <= v1 && r >= vr) {
          a[v].x = apply(a[v].x, x, vl, vr); // update op
          a[v].lz = accum(a[v].lz, x);  // update merge
      } else {
          int vm = (vl + vr) / 2:
          a[v].1 = update(1, r, x, a[v].1, v1, vm);
          a[v].r = update(1, r, x, a[v].r, vm, vr);
          a[v].x = merge(a[a[v].1].x, a[a[v].r].x); //
               query merge
```

## 1.5 persistent-segment-tree

```
// usage:
// Pst<Node<11>> pst:
// pst = {N};
// int newtime = pst.update(time, index, value);
// Node<ll> result = pst.querv(newtime. left. right):
template <class T>
struct Node {
   T x:
   int 1 = -1, r = -1;
   Node(): x(0) {}
   Node(T x) : x(x) \{ \}
   Node (Node a, Node b, int l = -1, int r = -1): x(a.x + b.
        x), 1(1), r(r) {}
};
template <class U>
struct Pst {
   int N:
    vector<U> a:
    vector<int> head:
   int build(int vl, int vr) {
       if (vr - vl == 1) a.push_back(U());
           int vm = (vl + vr) / 2, l = build(vl, vm),
              r = build(vm, vr);
           a.push_back(U(a[1], a[r], 1, r));
```

```
return a.size() - 1:
   }
   U query(int 1, int r, int v, int v1, int vr) {
       if (1 >= vr || r <= vl) return U();</pre>
       if (1 <= v1 && r >= vr) return a[v]:
       int vm = (v1 + vr) / 2;
       return U(query(1, r, a[v].1, v1, vm),
               query(1, r, a[v].r, vm, vr));
   int update(int i, U x, int v, int vl, int vr) {
       a.push_back(a[v]);
       v = a.size() - 1;
       if (vr - vl == 1) a[v] = x:
       else {
          int vm = (vl + vr) / 2;
          if (i < vm) a[v].l = update(i, x, a[v].l, vl, vm)</pre>
           else a[v].r = update(i, x, a[v].r, vm, vr);
          a[v] = U(a[a[v].1], a[a[v].r], a[v].1, a[v].r);
      }
       return v;
   Pst(int N) : N(N) { head.push_back(build(0, N)); }
   U query(int t, int l, int r) {
       return query(1, r, head[t], 0, N);
   int update(int t, int i, U x) {
       return head.push_back(update(i, x, head[t], 0, N)).
           head.size() - 1;
};
```

## 1.6 rmq-lineal

```
typedef int tf; // O(n) construction, O(1) query
struct rmq{
   int n; tf INF=1e9;//change sign of INF for MAX
   vector<unsigned int> mk; vector<tf> bk,v;
   rmq(){}
   tf op(tf a, tf b){return min(a,b);}//change for maximum
   int f(int x){return x>>5;}
   rmq(vector<tf> &vv):n(SZ(vv)),mk(n),bk(n,INF),v(vv){
      unsigned int lst=0;
      for(int i=0;i<SZ(v);i++,lst<<=1){</pre>
```

```
bk[f(i)]=op(bk[f(i)],v[i]);
          while(lst&&v[i-_builtin_ctz(lst)]>v[i]) lst^=lst
               &-lst: //MIN
          //while(lst&&v[i- builtin ctz(lst)]<v[i]) lst^=
               lst&-lst: //MAX
          mk[i]=++lst:
       for(int k=1,top=f(n);(1<<k)<=top;k++)fore(i,0,top)if(
            i+(1<<k)<=top)
          bk[top*k+i]=op(bk[top*(k-1)+i], bk[top*(k-1)+i]
               +(1<<k-1)]):
   tf get(int st, int en){
       return v[en-31+ builtin clz(mk[en]&((111<<en-st+1)
   tf query(int s, int e){ //[s,e]
       int b1=f(s),b2=f(e),top=f(n);
       if(b1==b2) return get(s,e);
       tf ans=op(get(s,(b1+1)*32-1), get(b2*32.e)); s=(b1+1)
            *32; e=b2*32-1;
       if(s<=e){
          int k=31-__builtin_clz(f(e-s+1));
          ans=op(ans,op(bk[top*k+f(s)],bk[top*k+f(e)-(1<<k)
      }
       return ans:
};
```

## 1.7 segment-tree-2d

```
// #define MAXN 1024 #define op(a,b) (a+b) #define NEUT 0
int n.m: int a[MAXN][MAXN].st[2*MAXN][2*MAXN];
void build(){
   repx(i, 0, n) repx(j, 0, m) st[i+n][j+m] = a[i][j];
   repx(i, 0, n) for(int j = m-1; j; --j)
       st[i+n][j] = op(st[i+n][j<<1], st[i+n][j<<1|1]);
   for(int i = n-1; i; --i) repx(j, 0, 2*m)
       st[i][i] = op(st[i << 1][i], st[i << 1|1][i]):
void upd(int x, int y, int v){
   st[x+n][y+m]=v;
   for(int j = y+m; j > 1; j >>= 1)
       st[x+n][i>>1] = op(st[x+n][i], st[x+n][i^1]):
   for(int i = x+n; i > 1; i >>= 1) for(int j=y+m; j; j>>=1)
       st[i>>1][j] = op(st[i][j], st[i^1][j]);
}
int query(int x0, int x1, int y0, int y1){
```

```
int r=NEUT;
for(int i0=x0+n, i1=x1+n; i0<i1; i0>>=1, i1>>=1){
    int t[4], q = 0;
    if(i0 & 1) t[q++] = i0++;
    if(i1 & 1) t[q++] = --i1;
    repx(k,0,q)
        for(int j0=y0+m, j1=y1+m; j0<j1; j0>>=1,j1>>=1){
        if(j0 & 1) r = op(r, st[t[k]][j0++]);
        if(j1 & 1) r = op(r, st[t[k]][--j1]);
    }
}
return r;
}
```

## 1.8 segment-tree-beats

```
struct Node {
   11 \text{ s. mx1. mx2. mxc. mn1. mn2. mnc. } 1z = 0:
   Node(): s(0), mx1(LLONG_MIN), mx2(LLONG_MIN), mxc(0),
        mn1(LLONG_MAX), mn2(LLONG_MAX), mnc(0) {}
   Node(ll x): s(x), mx1(x), mx2(LLONG MIN), mxc(1), mn1(x)
        , mn2(LLONG_MAX), mnc(1) {}
   Node(const Node &a. const Node &b) {
      // add
      s = a.s + b.s;
      // min
      if (a.mx1 > b.mx1) mx1 = a.mx1, mxc = a.mxc, mx2 =
            max(b.mx1. a.mx2):
       if (a.mx1 < b.mx1) mx1 = b.mx1, mxc = b.mxc, mx2 =
            max(a.mx1, b.mx2);
       if (a.mx1 == b.mx1) mx1 = a.mx1, mxc = a.mxc + b.mxc,
            mx2 = max(a.mx2, b.mx2);
       // max
      if (a.mn1 < b.mn1) mn1 = a.mn1, mnc = a.mnc, mn2 =
            min(b.mn1, a.mn2):
       if (a.mn1 > b.mn1) mn1 = b.mn1, mnc = b.mnc, mn2 =
            min(a.mn1, b.mn2):
       if (a.mn1 == b.mn1) mn1 = a.mn1, mnc = a.mnc + b.mnc,
            mn2 = min(a.mn2, b.mn2):
   }
};
// 0 - indexed / inclusive - inclusive
template <class node>
struct STB {
   vector<node> st; int n;
   void build(int u, int i, int j, vector<node> &arr) {
      if (i == j) {
```

```
st[u] = arr[i]:
       return:
   }
   int m = (i + i) / 2, l = u * 2 + 1, r = u * 2 + 2;
   build(1, i, m, arr), build(r, m + 1, j, arr);
   st[u] = node(st[l], st[r]):
void push_add(int u, int i, int j, ll v) {
   st[u].s += (j - i + 1) * v;
   st[u].mx1 += v, st[u].mn1 += v, st[u].lz += v;
   if (st[u].mx2 != LLONG MIN) st[u].mx2 += v:
   if (st[u].mn2 != LLONG MAX) st[u].mn2 += v:
void push_max(int u, ll v, bool l) { // for min op
   if (v >= st[u].mx1) return:
   st[u].s -= st[u].mx1 * st[u].mxc;
   st[u].mx1 = v:
   st[u].s += st[u].mx1 * st[u].mxc;
   if (1) st[u].mn1 = st[u].mx1:
   else if (v \le st[u].mn1) st[u].mn1 = v:
   else if (v < st[u].mn2) st[u].mn2 = v;
void push_min(int u, ll v, bool l) { // for max op
   if (v <= st[u].mn1) return:</pre>
   st[u].s -= st[u].mn1 * st[u].mnc;
   st[u].mn1 = v;
   st[u].s += st[u].mn1 * st[u].mnc;
   if (1) st[u].mx1 = st[u].mn1:
   else if (v \ge st[u].mx1) st[u].mx1 = v;
   else if (v > st[u].mx2) st[u].mx2 = v:
void push(int u, int i, int j) {
   if (i == i) return:
   // add
   int m = (i + j) / 2, l = u * 2 + 1, r = u * 2 + 2;
   push add(l, i, m, st[u].lz);
   push_add(r, m + 1, j, st[u].lz);
   st[u].lz = 0;
   // min
   push max(1, st[u].mx1, i == m):
   push max(r, st[u].mx1, m + 1 == i):
   // max
   push min(1, st[u].mn1, i == m):
   push_min(r, st[u].mn1, m + 1 == r);
node query(int a, int b, int u, int i, int j) {
   if (b < i || j < a) return node();</pre>
   if (a <= i && j <= b) return st[u];</pre>
   push(u, i, i):
   int m = (i + j) / 2, l = u * 2 + 1, r = u * 2 + 2;
```

```
return node(query(a, b, 1, i, m), query(a, b, r, m +
        1, i));
void update_add(int a, int b, ll v, int u, int i, int i)
   if (b < i | | i < a) return:
   if (a <= i && j <= b) {</pre>
       push_add(u, i, j, v);
       return:
   push(u, i, j);
   int m = (i + j) / 2, l = u * 2 + 1, r = u * 2 + 2;
   update_add(a, b, v, l, i, m);
   update_add(a, b, v, r, m + 1, j);
   st[u] = node(st[1], st[r]);
void update_min(int a, int b, ll v, int u, int i, int j)
   if (b < i \mid | i < a \mid | v >= st[u].mx1) return:
   if (a <= i && i <= b && v > st[u].mx2) {
       push_max(u, v, i == j);
       return:
   push(u, i, j);
   int m = (i + j) / 2, l = u * 2 + 1, r = u * 2 + 2;
   update_min(a, b, v, l, i, m);
   update_min(a, b, v, r, m + 1, j);
   st[u] = node(st[1], st[r]);
void update_max(int a, int b, ll v, int u, int i, int j)
   if (b < i || j < a || v <= st[u].mn1) return;</pre>
   if (a <= i && j <= b && v < st[u].mn2) {</pre>
       push_min(u, v, i == j);
       return:
   push(u, i, j);
   int m = (i + j) / 2, l = u * 2 + 1, r = u * 2 + 2;
   update_max(a, b, v, l, i, m);
   update_max(a, b, v, r, m + 1, j);
   st[u] = node(st[1], st[r]);
STB(vector<node> &v, int N) : n(N), st(N * 4 + 5) { build
    (0, 0, n - 1, v); 
node query(int a, int b) { return query(a, b, 0, 0, n -
    1); }
void update_add(int a, int b, ll v) { update_add(a, b, v,
     0.0.n - 1):
```

### 1.9 segment-tree-lazy

```
template <class T>
struct Stl {
   int n; vector<T> a, b;
   Stl(int n = 0) : n(n), a(4 * n, qneut()),
      b(4 * n. uneut()) {}
   T gneut() { return -2e9: }
   T uneut() { return 0: }
   T merge(T x, T y) { return max(x, y); }
   void upd(int v, T x, int 1, int r)
      \{a[v] += x, b[v] += x; \}
   void push(int v. int vl. int vm. int vr) {
      upd(2 * v, b[v], vl, vm);
      upd(2 * v + 1, b[v], vm, vr);
       b[v] = uneut():
   }
   T query(int 1, int r, int v=1, int vl=0, int vr=1e9) {
       vr = min(vr, n):
      if (1 <= v1 && r >= vr) return a[v];
      if (1 >= vr || r <= vl) return gneut();</pre>
      int vm = (vl + vr) / 2:
      push(v. vl. vm. vr):
      return merge(query(1, r, 2 * v, v1, vm),
          querv(1, r, 2 * v + 1, vm, vr)):
   void update(int l. int r. T x. int v = 1. int vl = 0.
          int vr = 1e9) {
      vr = min(vr, n):
      if (1 >= vr || r <= vl || r <= 1) return;</pre>
      if (1 <= v1 && r >= vr) upd(v, x, v1, vr);
          int vm = (vl + vr) / 2;
          push(v, v1, vm, vr);
          update(1, r, x, 2 * v, v1, vm):
          update(1, r, x, 2 * v + 1, vm, vr);
          a[v] = merge(a[2 * v], a[2 * v + 1]);
      }
```

## 1.10 segment-tree

```
struct St {
    11 neut() { return 0; }
    11 merge(11 x, 11 y) { return x + y; }

    int n; vector<11> a;
    St(int n = 0) : n(n), a(2 * n, neut()) {}

    11 query(int 1, int r) {
        11 x = neut(), y = neut();
        for (1 += n, r += n; 1 < r; 1 /= 2, r /= 2) {
            if (1 & 1) x = merge(x, a[1++]);
            if (r & 1) y = merge(a[--r], y);
        }
        return merge(x, y);
    }

    void update(int i, 11 x) {
        for (a[i += n] = x; i /= 2;)
            a[i] = merge(a[2 * i], a[2 * i + 1]);
    }
};</pre>
```

## 1.11 sparse-table

```
template <class T>
struct Sparse {
   T op(T a, T b) { return max(a, b); }
   vector<vector<T>> st;
   Sparse() {}
   Sparse(vector<T> a) : st{a} {
       int N = st[0].size();
       int npot = N <= 1 ? 1 : 32 - __builtin_clz(N);</pre>
       st.resize(npot):
       repx(i, 1, npot) rep(j, N + 1 - (1 << i))
       st[i].push_back(
          op(st[i-1][j], st[i-1][j+(1 << (i-1))])
      ); // query op
   T query(int 1, int r) { // range must be nonempty!
       int i = 31 - __builtin_clz(r - 1);
       return op(st[i][1], st[i][r - (1 << i)]); // queryop</pre>
```

```
};
```

### 1.12 treap-implicit

```
mt19937 gen(chrono::high_resolution_clock::now().
     time_since_epoch().count());
// 101 Implicit Treap //
struct Node {
    int p, sz = 0, v, acc, l = -1, r = -1:
    Node() : v(0), acc(0) {}
    Node(int x) : p(gen()), sz(1), v(x), acc(x) {}
    void recalc(const Node &a, const Node &b) {
       sz = a.sz + b.sz + 1; acc = v + a.acc + b.acc;
};
template <class node>
struct Treap {
    vector<node> t;
    int n, r = -1;
    node get(int u) { return u != -1 ? t[u] : node(); }
    void recalc(int u){t[u].recalc(get(t[u].l),get(t[u].r));}
    int merge(int 1, int r) {
       if (min(1, r) == -1) return 1 != -1 ? 1 : r:
       int ans = (t[1].p < t[r].p) ? 1 : r;
       if (ans == 1) t[1].r = merge(t[1].r, r), recalc(1);
       if (ans == r) t[r].l = merge(l, t[r].l), recalc(r);
       return ans:
    pii split(int u, int id) {
       if (u == -1) return {-1, -1};
       int szl = get(t[u].1).sz:
       if (szl >= id) {
           pii ans = split(t[u].1, id);
           t[u].1 = ans.ss;
           recalc(u);
           return {ans.ff, u}:
       pii ans = split(t[u].r, id - szl - 1);
       t[u].r = ans.ff:
       recalc(u):
       return {u, ans.ss};
    Treap(vi &v) : n(sz(v)) {
       for (int i=0: i<n: i++) t.eb(v[i]), r = merge(r, i):
};
// Complete Implicit Treap with Lazy propagation //
struct Node {
```

```
int p, sz = 0, v, acc, l = -1, r = -1, par = -1, lzv=0;
   bool lz = false, f = false:
   Node() : v(0), acc(0) {}
   Node(int x) : p(gen()), sz(1), v(x), acc(x) {}
   void recalc(const Node &a, const Node &b) {
      sz = a.sz + b.sz + 1; acc = v + a.acc + b.acc;
   void upd_lazy(int x) { lz = 1, lzv += x; }
   void lazy() {v+=lzv, acc += sz * lzv, lz = 0, lzv = 0; }
   void flip() { swap(1, r), f = 0; }
template <class node> struct Treap {
   vector<node> t;
   int n. r = -1:
   node get(int u) { return u != -1 ? t[u] : node(); }
   void recalc(int u) {
      int 1 = t[u].1. r = t[u].r:
      push(1), push(r), flip(1), flip(r);
      t[u].recalc(get(1), get(r));
   }
   void push(int u) {
       if (u == -1 || !t[u].lz) return;
      int 1 = t[u].1, r = t[u].r;
      if (1 != -1) t[1].upd_lazy(t[u].lzv);
      if (r != -1) t[r].upd_lazy(t[u].lzv);
      t[u].lazy();
   }
   void flip(int u) {
      if (u == -1 || !t[u].f) return;
      int 1 = t[u].1, r = t[u].r:
      if (1 != -1) t[1].f ^= 1;
      if (r != -1) t[r].f ^= 1;
      t[u].flip():
   int merge(int 1, int r) { // (*) = only if parent needed
      if (min(1, r) == -1) return 1 != -1 ? 1 : r:
      push(1), push(r), flip(1), flip(r);
      int ans = (t[1].p < t[r].p) ? 1 : r;
      if (ans == 1) t[1].r = merge(t[1].r, r), recalc(1);
      if (ans == r) t[r].l = merge(l, t[r].l), recalc(r);
      if (t[ans].l != -1) t[t[ans].l].par = ans: // (*)
      if (t[ans].r != -1) t[t[ans].r].par = ans; // (*)
       return ans:
   pii split(int u, int id) {// (*) = only if parent needed
      if (u == -1) return {-1, -1}:
      push(u);
      flip(u);
      int szl = get(t[u].1).sz;
      if (szl >= id) {
```

```
pii ans = split(t[u].1, id):
          if (ans.ss != -1) t[ans.ss].par = u; // (*)
          if (ans.ff != -1) t[ans.ff].par = -1; // (*)
          t[u].1 = ans.ss:
          recalc(u);
          return {ans.ff. u}:
       pii ans = split(t[u].r, id - szl - 1);
       if (ans.ff != -1) t[ans.ff].par = u; // (*)
       if (ans.ss != -1) t[ans.ss].par = -1; // (*)
       t[u].r = ans.ff:
       recalc(u):
       return {u, ans.ss};
   int update(int u, int 1, int r, int v) {
       pii a = split(u, 1), b = split(a.ss, r - 1 + 1);
       t[b.ff].upd lazv(v):
       return merge(a.ff, merge(b.ff, b.ss));
   void print(int u) {
       if (u==-1) return; push(u), flip(u); print(t[u].1);
       cout << t[u].v << ' '; print(t[u].r);</pre>
   Treap(vi &v) : n(sz(v)) {
       for (int i=0; i<n; i++) t.eb(v[i]), r = merge(r, i);</pre>
};
```

#### 1.13 treap

```
typedef struct item *pitem:
struct item {
   int pr,key,cnt; pitem l,r;
   item(int kev):kev(kev).pr(rand()).cnt(1).1(0).r(0) {}
int cnt(pitem t){return t?t->cnt:0;}
void upd_cnt(pitem t){if(t)t->cnt=cnt(t->1)+cnt(t->r)+1;}
void split(pitem t, int key, pitem& l, pitem& r){ // l: <=</pre>
    kev. r: > kev
   if(!t)l=r=0:
   else if(key<t->key)split(t->1,key,1,t->1),r=t;
   else split(t->r,key,t->r,r),l=t;
   upd_cnt(t);
void insert(pitem& t. pitem it){
   if(!t)t=it:
   else if(it->pr>t->pr)split(t,it->key,it->l,it->r),t=it;
   else insert(it->kev<t->kev?t->l:t->r.it):
   upd_cnt(t);
```

```
void merge(pitem& t, pitem l, pitem r){
    if(!l||!r)t=l?l:r;
    else if(l->pr>r->pr)merge(l->r,l->r,r),t=l;
    else merge(r->l,l,r->l),t=r;
    upd_cnt(t);
}

void erase(pitem& t, int key){
    if(t->key==key)merge(t,t->l,t->r);
    else erase(key<t->key?t->l:t->r,key);
    upd_cnt(t);
}
```

## $2 \quad geo2d$

#### 2.1 circle

```
struct C {
   Po; Tr;
   // circle-line intersection, assuming it exists
   // points are sorted along the direction of the line
   pair<P, P> line_inter(L 1) const {
      P c = 1.closest_to(o); T c2 = (c - o).magsq();
      P = 1.d * sqrt(max(r*r - c2, T()) / 1.d.magsq());
      return {c - e, c + e}:
   // check the type of line-circle collision
   // <0: 2 inters, =0: 1 inter, >0: 0 inters
   T line_collide(L 1) const {
      T c2 = (1.closest_to(o) - o).magsq();
       return c2 - r * r:
   // calculates the two intersections between two circles
   // the circles must intersect in one or two points!
   pair<P, P> inter(C h) const {
      P d = h.o - o;
      T c = (r * r - h.r * h.r) / d.magsq();
      return h.line_inter({(1 + c) / 2 * d, d.rot()});
   // check if the given circles intersect
   bool collide(C h) const {
       return (h.o - o).magsq() \le (h.r + r) * (h.r + r);
   // get one of the two tangents that go through the point
   // the point must not be inside the circle
   // a = -1: cw (relative to the circle) tangent
   // a = 1: ccw (relative to the circle) tangent
   P point_tangent(P p, T a) const {
```

```
T c = r * r / p.magsq();
      return o + c*(p-o) - a*sqrt(c*(1-c))*(p-o).rot();
  // get one of the 4 tangents between the two circles
  // a = 1: exterior tangents
  // a = -1: interior tangents (requires no area overlap)
  // b = 1: ccw tangent
  // b = -1: cw tangent
  // the line origin is on this circumference, and the
  // direction is a unit vector towards the other circle
  L tangent(C c, T a, T b) const {
     T dr = a * r - c.r:
     P d = c.o - o:
     P n = (d*dr+b*d.rot()*sqrt(d.magsq()-dr*dr)).unit();
      return {o + n * r, -b * n.rot()};
  // circumcircle of a **non-degenerate** triangle
  static C thru_points(P a, P b, P c) {
     b = b - a, c = c - a:
      P p = (b*c.magsq() - c*b.magsq()).rot() / (b%c*2);
      return {a + p, p.mag()};
  // find the two circles that go through the given point,
  // are tangent to the given line and have radius 'r'
  // the point-line distance must be at most 'r'!
  // the circles are sorted in the direction of the line
  static pair<C, C> thru_point_line_r(P a, L t, T r) {
     P d = t.d.rot().unit():
      if (d * (a - t.o) < 0) d = -d;
      auto p = C(a, r).line_inter(\{t.o + d * r, t.d\});
      return {{p.first, r}, {p.second, r}};
  // find the two circles that go through the given points
  // and have radius 'r'
  // circles sorted by angle from the first point
  // the points must be at most at distance 'r'!
  static pair<C, C> thru_points_r(P a, P b, T r) {
      auto p = C(a, r).line_inter({(a+b)/2, (b-a).rot()});
      return {{p.first, r}, {p.second, r}};
  vector<P> linecol(L 1){
      vector<P> s;P p=1.closest_to(o);double d=(p-o).norm()
      if(d-EPS>r)return s;
      if(abs(d-r)<=EPS){s.pb(p);return s;}</pre>
      d=sqrt(r*r-d*d); s.pb(p+l.pq.unit()*d); s.pb(p-l.pq.
          unit()*d);
      return s:
double intertriangle(P a,P b){ // intersection with oab
```

### 2.2 closest-points

```
// sort by x
11 closest(vector<ii> &p) {
   int n = SZ(p):
   set<ii>> s;
   ll best = 1e18:
   int j = 0;
   fore(i, 0, n) {
       11 d = ceil(sqrt(best));
       while(p[i].fst - p[j].fst >= best)
           s.erase({p[j].snd, p[j].fst}), j++;
       auto it1=s.lower bound({p[i].snd-d.p[i].fst}):
       auto it2=s.upper_bound({p[i].snd+d,p[i].fst});
       for(auto it = it1: it != it2: ++it) {
          11 dx = p[i].fst - it->snd;
          11 dy = p[i].snd - it->fst;
          best = min(best, dx * dx + dy * dy);
       s.insert({p[i].snd, p[i].fst});
   return best:
```

### 2.3 convex-hull

```
// ccw order, excludes collinear points by default
vector<P> chull(vector<P> p) {
   if (p.size() < 3) return p;
   vector<P> r; int m, k = 0;
   sort(p.begin(), p.end(), [](P a, P b) {
```

### 2.4 delaunay

```
typedef __int128_t lll; // if on a 64-bit platform
struct Q {
    Q *rot, *o; P p = {INF, INF}; bool mark;
    P &F() { return r()->p; }
    Q *&r() { return rot->rot: }
    Q *prev() { return rot->o->rot; }
    Q *next() { return r()->prev(); }
};
T cross(P a, P b, P c) { return (b - a) % (c - a); }
bool circ(P p, P a, P b, P c) { // is p in the circumcircle?
    111 p2 = p.magsq(), A = a.magsq() - p2.
       B = b.magsq() - p2, C = c.magsq() - p2;
    return cross(p, a, b) * C + cross(p, b, c) * A + cross(p,
          c. a) * B > 0:
Q *makeEdge(Q *&H, P orig, P dest) {
    Q *r = H ? H : new Q{new Q{new Q{new Q{0}}}};
    H = r -> 0; r -> r() -> r() = r;
    repx(i, 0, 4) r = r->rot, r->p = {INF, INF},
       r->0 = i & 1 ? r : r->r();
    r\rightarrow p = orig; r\rightarrow F() = dest;
    return r;
void splice(Q *a, Q *b) {
    swap(a->o->rot->o, b->o->rot->o); swap(a->o, b->o);
```

```
Q *connect(Q *&H, Q *a, Q *b) {
   Q *q = makeEdge(H, a->F(), b->p);
   splice(q, a->next()); splice(q->r(), b); return q;
pair<Q *, Q *> rec(Q *&H, const vector<P> &s) {
   if (s.size() <= 3) {
       Q *a = makeEdge(H, s[0], s[1]), *b = makeEdge(H, s[0])
            [1], s.back());
       if (s.size() == 2) return {a, a->r()}; splice(a->r(),
       auto side = cross(s[0], s[1], s[2]);
       Q *c = side ? connect(H, b, a) : 0;
       return {side < 0 ? c->r() : a, side < 0 ? c : b->r()
#define J(e) e \rightarrow F(), e \rightarrow p
#define valid(e) (cross(e->F(), J(base)) > 0)
   Q *A, *B, *ra, *rb; int half = s.size() / 2;
   tie(ra, A) = rec(H, {s.begin(), s.end() - half});
   tie(B, rb) = rec(H, {s.begin() + s.size() - half, s.end()
   while ((cross(B->p, J(A)) < 0 \&\& (A = A->next())) | |
          (cross(A->p, J(B)) > 0 && (B = B->r()->o)));
   Q *base = connect(H, B->r(), A);
   if (A->p == ra->p) ra = base->r():
   if (B->p == rb->p) rb = base;
#define DEL(e, init, dir) Q *e = init->dir; \
   if (valid(e)) while (circ(e->dir->F(), J(base), e->F()))
        { \
           Q *t = e->dir; splice(e, e->prev()); \
           splice(e->r(), e->r()->prev()); e->o = H; H = e;
               e = t: \
   for (::) {
       DEL(LC, base->r(), o); DEL(RC, base, prev());
       if (!valid(LC) && !valid(RC)) break;
       if (!valid(LC) || (valid(RC) && circ(J(RC), J(LC))))
            base = connect(H, RC, base->r());
       else base = connect(H, base->r(), LC->r());
   }
   return {ra, rb};
#undef J
#undef valid
#undef DEL
```

```
// there must be no duplicate points
// returns no triangles in the case of all collinear points
// produces counter-clockwise triangles ordered in triples
// maximizes the minimum angle across all triangulations
// the euclidean mst is a subset of these edges
// O(N log N)
vector<P> triangulate(vector<P> pts) {
   sort(pts.begin(), pts.end(), [](P a, P b) {
       return make_pair(a.x, a.y) < make_pair(b.x, b.y);</pre>
   assert(unique(pts.begin(), pts.end()) == pts.end());
   if (pts.size() < 2) return {}:</pre>
   Q *H = 0; Q *e = rec(H, pts).first;
   vector < Q *> q = \{e\}; int qi = 0;
   while (cross(e->o->F(), e->F(), e->p) < 0) e = e->o;
#define ADD
   {
       Q *c = e;
       do {
           c->mark = 1; pts.push_back(c->p); \
           q.push_back(c->r()); c = c->next(); \setminus
       } while (c != e):
   ADD:
   pts.clear();
   while (qi < (int)q.size()) if (!(e = q[qi++])->mark) ADD;
   return pts;
#undef ADD
```

### 2.5 halfplane-intersect

```
while (n \ge 2 \&\& H[i].side(a[0].intersection(a[1])) >
       q.pop_front(), n--;
   if (n > 0 \&\& H[i].parallel(q[n - 1])) {
       if (H[i].d * q[n - 1].d < 0) return {};</pre>
       if (H[i].side(q[n-1].o) > 0) q.pop_back(), n--;
       else continue:
   }
   q.push_back(H[i]), n++;
while (n \ge 3 \&\& a[0].side(a[n - 1].intersection(a[n -
    21)) > 0)
   q.pop_back(), n--;
while (n \ge 3 \&\& q[n - 1].side(q[0].intersection(q[1])) >
   q.pop_front(), n--;
if (n < 3) return {}:
vector<P> ps(n):
rep(i, n) ps[i] = q[i].intersection(q[(i + 1) % n]);
return ps;
```

#### 2.6 line

```
// a segment or an infinite line
// does not handle point segments correctly!
struct L {
   P o, d;
   static L from eq(P ab, T c) {
       return L{ab.rot(), ab * -c / ab.magsq()};
   pair<P, T> line_eq() { return {-d.rot(), d.rot() * o}; }
   // on which side of the line is the point
   // negative: left, positive: right
   T side(P r) const { return (r - o) % d: }
   // returns the intersection coefficient
   // in the range [0, d % r.d]
   // if d % r.d is zero, the lines are parallel
   T inter(L r) const { return (r.o - o) % r.d; }
   // get the single intersection point
   // lines must not be parallel
   P intersection(L r) const {return o+d*inter(r)/(d%r.d);} };
```

```
// check if lines are parallel
bool parallel(L r) const { return abs(d % r.d) <= EPS; }</pre>
// check if segments intersect
bool seg_collide(L r) const {
   Tz = d \% r.d:
   if (abs(z) \le EPS) {
       if (abs(side(r.o)) > EPS) return false;
       T s = (r.o - o) * d, e = s + r.d * d;
       if (s > e) swap(s, e);
       return s <= d * d + EPS && e >= -EPS;
   T s = inter(r), t = -r.inter(*this);
   if (z < 0) s = -s, t = -t, z = -z;
   return s>=-EPS && s<=z+EPS && t>=-EPS && t<=z+EPS:
// full segment intersection
// makes a point segment if the intersection is a point
// however it does not handle point segments as input!
bool seg_inter(L r, L *out) const {
   Tz = d \% r.d:
   if (abs(z) <= EPS) {
       if (abs(side(r.o)) > EPS) return false;
       if (r.d * d < 0) r = \{r.o + r.d, -r.d\};
       P s = o * d < r.o * d ? r.o : o;
       P = (o+d)*d < (r.o+r.d)*d ? o+d : r.o+r.d:
       if (s * d > e * d) return false:
       return *out = {s, e - s}, true;
   T s = inter(r), t = -r.inter(*this);
   if (z < 0) s = -s, t = -t, z = -z;
   if (s>=-EPS && s<=z+EPS && t>=-EPS && t<=z+EPS)
       return *out = \{0 + d * s / z, \{0, 0\}\}, true;
   return false:
// check if the given point is on the segment
bool point_on_seg(P r) const {
   if (abs(side(r)) > EPS) return false;
   if ((r - o) * d < -EPS) return false:</pre>
   if ((r - o - d) * d > EPS) return false;
   return true:
// point in this line that is closest to a given point
P closest_to(P r) const {
   P dr = d.rot(): return r + dr*((o-r)*dr)/d.magsq():
```

#### 2.7 minkowski

```
void reorder_polygon(vector<P> &ps) {
   int pos = 0;
   repx(i, 1, (int)ps.size()) {
       if (ps[i].y < ps[pos].y || (ps[i].y == ps[pos].y &&</pre>
           ps[i].x < ps[pos].x))
           pos = i:
   rotate(ps.begin(), ps.begin() + pos, ps.end());
vector<P> minkowski(vector<P> ps, vector<P> qs) {
   // the first vertex must be the lowest
   reorder_polygon(ps); reorder_polygon(qs);
   ps.push_back(ps[0]); ps.push_back(ps[1]);
   qs.push_back(qs[0]); qs.push_back(qs[1]);
   vector<P> result; int i = 0, j = 0;
   while (i < ps.size() - 2 || j < qs.size() - 2) {</pre>
       result.push back(ps[i] + qs[i]):
       auto z = (ps[i + 1] - ps[i]) % (qs[j + 1] - qs[j]);
       if (z >= 0 && i < ps.size() - 2) ++i;</pre>
       if (z \le 0 \&\& i \le as.size() - 2) ++i:
   return result:
```

### 2.8 point

```
struct P {
   T x, y;
   P(T x, T y) : x(x), y(y) {}
   P() : P(0, 0) \{ \}
   friend ostream &operator<<(ostream &s, const P &r) {</pre>
       return s << r.x << " " << r.v:
   friend istream & operator >> (istream &s, P &r) { return s
        >> r.x >> r.y; }
   P operator+(P r) const { return \{x + r.x, y + r.y\}; \}
   P operator-(P r) const { return {x - r.x, y - r.y}; }
   P operator*(T r) const { return {x * r, y * r}; }
   P operator/(T r) const { return {x / r, y / r}; }
   P operator-() const { return {-x, -y}; }
   friend P operator*(T 1, P r) { return {1 * r.x, 1 * r.v};
   P rot() const { return {-y, x}; }
   T operator*(P r) const { return x * r.x + y * r.y; }
```

```
T operator%(P r) const { return rot() * r; }
   T left(P a. P b) { return (b - a) % (*this - a): }
   T magsq() const { return x * x + v * v; }
   T mag() const { return sqrt(magsq()); }
   P unit() const { return *this / mag(): }
   bool half() const { return abs(v) <= EPS && x < -EPS || v
   T angcmp(P r) const { // like strcmp(this, r)
       int h = (int)half() - r.half();
       return h ? h : r % *this:
   T angcmp_rel(P a, P b) { // like strcmp(a, b)
       Pz = *this:
       int h = z \% a \le 0 \&\& z * a < 0 || z % a < 0;
       h = z \% b \le 0 \&\& z * b < 0 || z \% b < 0:
       return h? h: b % a:
   bool operator==(P r) const { return abs(x - r.x) <= EPS</pre>
        && abs(v - r.v) <= EPS: }
   double angle() const { return atan2(y, x); }
   static P from_angle(double a) { return {cos(a), sin(a)};
        }
};
```

## 2.9 polygon

```
// get TWICE the area of a simple polygon in ccw order
T area2(const vector<P> &p) {
   int n = p.size(); T a = 0;
   rep(i, n) = + = (p[i] - p[0]) \% (p[(i + 1) \% n] - p[i]):
   return a:
}
// checks whether a point is inside a ccw simple polygon
// returns 1 if inside, 0 if on border, -1 if outside
int in_poly(const vector<P> &p, P q) {
   int w = 0:
   rep(i, p.size()) {
       P = p[i], b = p[(i + 1) \% p.size()];
       T k = (b - a) \% (q - a);
       T u = a.v - g.v, v = b.v - g.v;
       if (k > 0 && u < 0 && v >= 0) w++;
       if (k < 0 && v < 0 && u >= 0) w--:
       if (k == 0 && (q - a) * (q - b) <= 0) return 0:
```

```
return w ? 1 : -1:
// check if point in ccw convex polygon, O(log n)
// + if inside, 0 if on border, - if outside
T in convex(const vector<P> &p. P a) {
    int l = 1, h = p.size() - 2; assert(p.size() >= 3);
    while (1 != h) { // collinear points are unsupported!
       int m = (1 + h + 1) / 2:
       if (q.left(p[0], p[m]) >= 0) 1 = m;
       else h = m - 1:
   T in = min(q.left(p[0], p[1]), q.left(p.back(), p[0]));
    return min(in, q.left(p[1], p[1 + 1]));
int extremal(const vector<P> &p. P d) {
    int n = p.size(), l = 0, r = n - 1; assert(n);
    P = 0 = (p[n - 1] - p[0]).rot():
    while (1 < r) { // polygon must be convex
       int m = (1 + r + 1) / 2;
       P = (p[(m + n - 1) \% n] - p[m]).rot();
       if (e0.angcmp_rel(d, e) < 0) r = m - 1;
       else 1 = m:
   }
    return 1;
// square dist of most distant points of a ccw convex
// polygon with NO COLLINEAR POINTS
T callipers(const vector<P> &p) {
    int n = p.size();
   T r = 0:
    for (int i = 0, j = n < 2 ? 0 : 1; <math>i < j; i++) {
       for (;; j = (j + 1) \% n) {
           r = max(r, (p[i] - p[j]).magsq());
           if ((p[(i + 1) % n] - p[i]) % (p[(j + 1) % n] - p
                [i]) <= EPS) break:</pre>
    return r:
P centroid(const vector<P> &p) { // (barycenter)
    P r(0, 0); T t = 0; int n = p.size();
    rep(i, n) {
       r += (p[i] + p[(i+1)\%n]) * (p[i] \% p[(i+1)\%n]);
       t += p[i] \% p[(i+1)\%n];
    return r / t / 3;
```

### 2.10 sweep

```
#include "point.cpp"
// iterate over all pairs of points
// 'op' is called with all ordered pairs of different
    indices '(i, j)'
// additionally, the 'ps' vector is kept sorted by signed
    distance
// to the line formed by 'i' and 'j'
// for example, if the vector from 'i' to 'j' is pointing
// the 'ps' vector is sorted from smallest 'y' to largest 'y
// note that, because the 'ps' vector is sorted by signed
// 'j' is always equal to 'i + 1'
// this means that the amount of points to the left of the
    line is always 'N - i'
template <class OP>
void all pair points(vector<P> &ps. OP op) {
   int N = ps.size();
   sort(ps.begin(), ps.end(), [](P a, P b) {
      return make_pair(a.y, a.x) < make_pair(b.y, b.x);</pre>
   vector<pair<int, int>> ss;
   rep(i, N) rep(j, N) if (i != j) ss.push_back({i, j});
   stable_sort(ss.begin(), ss.end(), [&](auto a, auto b) {
       return (ps[a.second] - ps[a.first]).angle lt(ps[b.
           second] - ps[b.first]);
   vector<int> p(N); rep(i, N) p[i] = i;
   for (auto [i, j] : ss)
```

#### 2.11 theorems

```
// Pick's theorem
// Simple polygon with integer vertices:
// A = I + B / 2 - 1
// A: Area of the polygon
// I: Integer points strictly inside the polygon
// B: Integer points on the boundary of the polygon
```

## 3 graph

## 3.1 artic-bridge-biconn

```
vector<int> g[MAXN]:int n:
struct edge {int u,v,comp;bool bridge;};
vector<edge> e;
void add_edge(int u, int v){
   g[u].pb(e.size());g[v].pb(e.size());
   e.pb((edge){u,v,-1,false});
int D[MAXN],B[MAXN],T;
int nbc; // number of biconnected components
int art[MAXN]; // articulation point iff !=0
stack<int> st: // only for biconnected
void dfs(int u.int pe){
   B[u]=D[u]=T++;
   for(int ne:g[u])if(ne!=pe){
       int v=e[ne].u^e[ne].v^u;
       if(D[v]<0){
          st.push(ne):dfs(v.ne):
          if(B[v]>D[u])e[ne].bridge = true; // bridge
          if(B[v]>=D[u]){
              art[u]++: // articulation
              int last: // start biconnected
              do{last=st.top();st.pop();e[last].comp=nbc;}
              while(last!=ne);
              nbc++; // end biconnected
          B[u]=min(B[u],B[v]);
       else if(D[v]<D[u])st.push(ne),B[u]=min(B[u],D[v]);</pre>
```

```
}
void doit(){
   memset(D,-1,sizeof(D));memset(art,0,sizeof(art));
   nbc=T=0; fore(i,0,n)if(D[i]<0)dfs(i,-1),art[i]--;
}</pre>
```

### 3.2 bellman-ford

```
struct Edge { int u, v; ll w; };
// find distance from source node to all nodes.
// supports negative edge weights.
// returns true if a negative cycle is detected.
11
// time: O(V E)
bool bellman_ford(int N, int s, vector<Edge> &E, vector<11>
    &D, vector<int> &P) {
   P.assign(N, -1), D.assign(N, INF), D[s] = 0;
   rep(i, N - 1) {
       bool f = true:
       rep(ei, E.size()) {
          auto &e = E[ei];
          ll n = D[e.u] + e.w:
          if (D[e.u] < INF && n < D[e.v])
              D[e.v] = n, P[e.v] = ei, f = false;
       if (f) return false;
   return true:
```

#### 3.3 blossom

```
y;else if(!inq[mt[y]])inq[q[++qt]=mt[y]]=true;}}}return -1;}
int aug(int s,int t){int x=t,y,z;while(x>=0){y=ft[x];z=mt[y]};mt[y]=x;mt[x]=y;x=z;}return t>=0;}int edmonds(){int r=0;}
memset(mt,-1,sizeof(mt));rep(x,n)if(mt[x]<0)r+=aug(x,findp(x));return r;}</pre>
```

#### 3.4 chu-liu-minimum-spanning-arborescence

```
//O(n*m) minimum spanning tree in directed graph
//returns -1 if not possible
//included i-th edge if take[i]!=0
typedef int tw; tw INF=111<<30;</pre>
struct edge{int u,v,id;tw len;};
struct ChuLiu{
   int n: vector<edge> e:
   vector<int> inc,dec,take,pre,num,id,vis;
   vector<tw> inw:
   void add_edge(int x, int y, tw w){
       inc.pb(0); dec.pb(0); take.pb(0);
       e.pb(\{x,y,SZ(e),w\});
   ChuLiu(int n):n(n),pre(n),num(n),id(n),vis(n),inw(n){}
   tw doit(int root){
       auto e2=e:
       tw ans=0; int eg=SZ(e)-1,pos=SZ(e)-1;
       while(1){
           fore(i,0,n) inw[i]=INF,id[i]=vis[i]=-1;
           for(auto ed:e2) if(ed.len<inw[ed.v]){</pre>
               inw[ed.v]=ed.len; pre[ed.v]=ed.u;
              num[ed.v]=ed.id;
           inw[root]=0:
           fore(i,0,n) if(inw[i]==INF) return -1;
           int tot=-1:
           fore(i,0,n){
               ans+=inw[i];
              if(i!=root)take[num[i]]++:
              int j=i;
               while(vis[i]!=i&&i!=root&&id[i]<0)vis[i]=i.i=</pre>
                   pre[j];
              if(j!=root&&id[j]<0){</pre>
                  id[i]=++tot:
                  for(int k=pre[j];k!=j;k=pre[k]) id[k]=tot;
           }
           if(tot<0)break;</pre>
           fore(i,0,n) if(id[i]<0)id[i]=++tot;</pre>
           n=tot+1; int j=0;
           fore(i,0,SZ(e2)){
```

```
int v=e2[i].v:
              e2[j].v=id[e2[i].v];
              e2[j].u=id[e2[i].u];
              if(e2[j].v!=e2[j].u){
                  e2[j].len=e2[i].len-inw[v];
                  inc.pb(e2[i].id):
                  dec.pb(num[v]);
                  take.pb(0);
                  e2[j++].id=++pos;
           e2.resize(i):
           root=id[root];
       while(pos>eg){
           if(take[pos]>0) take[inc[pos]]++, take[dec[pos
               11--:
          pos--;
       return ans;
};
```

#### 3.5 dinic

```
// time: 0(E V^2)
       O(E V^(2/3)) / O(E sqrt(E)) unit capacities
       O(E sgrt(V)) (hopcroft-karp) unit networks
//unit network: c in {0,1} & forall v, indeg<=1 or outdeg<=1</pre>
//min-cut: nodes reachable from s in final residual graph
struct Dinic {
   struct Edge { int u, v: ll c, f = 0: }:
   int N, s, t; vector<vector<int>> G;
   vector<Edge> E: vector<int> lvl. ptr:
   Dinic() {}
   Dinic(int N, int s, int t): N(N), s(s), t(t), G(N) {}
   void add_edge(int u, int v, ll c) {
       G[u].push_back(E.size()); E.push_back({u, v, c});
       G[v].push_back(E.size()); E.push_back({v, u, 0});
   11 push(int u, 11 p) {
       if (u == t || p <= 0) return p;</pre>
       while (ptr[u] < G[u].size()) {</pre>
           int ei = G[u][ptr[u]++];
          Edge &e = E[ei];
           if (lvl[e.v] != lvl[u] + 1) continue;
           ll a = push(e.v, min(e.c - e.f, p));
```

```
if (a <= 0) continue;</pre>
           e.f += a, E[ei ^ 1].f -= a; return a;
       return 0:
   }
   11 maxflow() {
       11 f = 0;
       while (true) {
           lvl.assign(N, -1); queue<int> q;
           lvl[s] = 0; q.push(s);
           while (!a.emptv()) {
               int u = q.front(); q.pop();
               for (int ei : G[u]) {
                   Edge &e = E[ei]:
                   if (e.c-e.f<=0||lvl[e.v]!=-1) continue;</pre>
                   lvl[e.v] = lvl[u] + 1; q.push(e.v);
               }
           if (lvl[t] == -1) break:
           ptr.assign(N,0); while(ll ff=push(s,INF))f += ff;
       return f;
};
/* Flujo con demandas (no necesariamente el maximo)
Agregar s' v t' nuevos source and sink
c'(s', v) = sum(d(u, v) \text{ for } u \text{ in } V) \setminus forall \text{ arista } (s', v)
c'(v, t') = sum(d(v, w) \text{ for w in V}) \setminus forall arista (v, t')
c'(u, v) = c(u, v) - d(u, v) \setminus forall arists antiguas
c'(t, s) = INF (el flujo por esta arista es el flujo real)*/
```

#### 3.6 dominator-tree

```
//idom[i]=parent of i in dominator tree with root=rt, or -1
        if not exists
int n,rnk[MAXN],pre[MAXN],anc[MAXN],idom[MAXN],semi[MAXN],
        low[MAXN];
vector<int> g[MAXN],rev[MAXN],dom[MAXN],ord;
void dfspre(int pos){
        rnk[pos]=SZ(ord); ord.pb(pos);
        for(auto x:g[pos]){
            rev[x].pb(pos);
            if(rnk[x]==n) pre[x]=pos,dfspre(x);
        }
}
int eval(int v){
        if(anc[v]<n&&anc[anc[v]]<n){</pre>
```

```
int x=eval(anc[v]):
       if(rnk[semi[low[v]]]>rnk[semi[x]]) low[v]=x;
       anc[v]=anc[anc[v]]:
   return low[v];
void dominators(int rt){
   fore(i.0.n){
       dom[i].clear(); rev[i].clear();
       rnk[i]=pre[i]=anc[i]=idom[i]=n;
       semi[i]=low[i]=i:
   ord.clear(); dfspre(rt);
   for(int i=SZ(ord)-1:i:i--){
       int w=ord[i]:
       for(int v:rev[w]){
          int u=eval(v):
          if(rnk[semi[w]]>rnk[semi[u]])semi[w]=semi[u];
       dom[semi[w]].pb(w); anc[w]=pre[w];
       for(int v:dom[pre[w]]){
          int u=eval(v):
          idom[v]=(rnk[pre[w]]>rnk[semi[u]]?u:pre[w]);
       dom[pre[w]].clear();
   for(int w:ord) if(w!=rt&&idom[w]!=semi[w]) idom[w]=idom[
        idom[w]]:
   fore(i,0,n) if(idom[i]==n)idom[i]=-1;
```

#### 3.7 eulerian

```
// path/tour for directed graphs. uncomment for undirected.
struct Euler {
    struct Edge { int v, rev; };
    vector<vector<Edge>> G; vector<Edge> P;
    Euler(int N = 0) : G(N) {}
    void add_edge(int u, int v) {
        G[u].push_back({v, (int)G[v].size()});
        // G[v].push_back({u, (int)G[u].size() - 1});
    }

    void go(int u) {
        while (G[u].size()) {
            Edge e = G[u].back(); G[u].pop_back();
            // if (e.v == -1) continue;
            // G[e.v][e.rev].v = -1;
            go(e.v); P.push_back(e);
```

```
}

// works ONLY if the vertex degrees are eulerian! check!
vector<Edge> get_path(int u) {
    return P.clear(),go(u),reverse(P.begin(),P.end()),P;
}
```

## 3.8 floyd-warshall

## 3.9 heavy-light

```
struct Hld {
   vector<int> P, H, D, pos, top;
   void init(vector<vector<int>> &G) {
      int N = G.size():
      P.resize(N), H.resize(N), D.resize(N), pos.resize(N),
          top.resize(N);
      D[0] = -1, dfs(G, 0); int t = 0;
      rep(i, N) if (H[P[i]] != i) {
          int j = i;
          while (j != -1)
              \{ top[j] = i, pos[j] = t++; j = H[j]; \}
   }
   int dfs(vector<vector<int>> &G. int i) {
      int w = 1, mw = 0;
      D[i] = D[P[i]] + 1, H[i] = -1;
      for (int c : G[i]) {
          if (c == P[i]) continue;
```

```
P[c] = i: int sw = dfs(G, c): w += sw:
       if (sw > mw) H[i] = c, mw = sw;
   return w:
// visit the log N segments in the path from u to v
template <class OP>
void path(int u, int v, OP op) {
   while (top[u] != top[v]) {
       if (D[top[u]] > D[top[v]]) swap(u, v);
       op(pos[top[v]], pos[v] + 1): v = P[top[v]]:
   if (D[u] > D[v]) swap(u, v);
   op(pos[u], pos[v] + 1); // value on node
   // op(pos[u]+1, pos[v] + 1); // value on edge
// an alternative to 'path' that considers order.
// calls 'op' with an 'l <= r' inclusive-exclusive range.
     and a
// boolean indicating if the query is forwards or
    backwards.
template <class OP>
void path(int u, int v, OP op) {
   int lu = u, lv = v;
   while (top[lu] != top[lv])
       if (D[top[lu]] > D[top[lv]]) lu = P[top[lu]];
       else lv = P[top[lv]];
   int lca = D[lu] > D[lv] ? lv : lu:
   while (top[u] != top[lca])
       op(pos[top[u]], pos[u] + 1, false), u = P[top[u
   if (u != lca) op(pos[lca] + 1, pos[u] + 1, false);
   vector<int> stk:
   while (top[v] != top[lca])
       stk.push_back(v), v = P[top[v]];
   // op(pos[lca], pos[v] + 1, true); // value on node
   op(pos[lca] + 1, pos[v] + 1, true); // value on edge
   reverse(stk.begin(), stk.end());
   for (int w : stk) op(pos[top[w]], pos[w] + 1, true);
// commutative segment tree
template <class T, class S>
void update(S &seg, int i, T val) { seg.update(pos[i],
    val): }
```

```
// commutative segment tree lazy
template <class T, class S>
void update(S &seg, int u, int v, T val) {
    path(u, v, [&](int 1, int r) { seg.update(1, r, val);
        });
}

// commutative (lazy) segment tree
template <class T, class S>
T query(S &seg, int u, int v) {
    T ans = 0;
        // neutral element
    path(u, v, [&](int 1, int r) { ans += seg.query(1, r)
        ; }); // query op
    return ans;
}
};
```

### 3.10 hungarian

```
// find a maximum gain perfect matching in the given
    bipartite complete graph.
// input: gain matrix (G_{xy} = benefit of joining vertex x
    in set X with vertex
// v in set Y).
// output: maximum gain matching in members 'xy[x]' and 'yx[
    v]'.
// runtime: O(N^3)
struct Hungarian {
   int N, qi, root;
   vector<vector<ll>> gain:
   vector<int> xy, yx, p, q, slackx;
   vector<ll> lx. lv. slack:
   vector<bool> S, T;
   void add(int x, int px) {
       S[x] = true, p[x] = px;
       rep(y, N) if (lx[x] + ly[y] - gain[x][y] < slack[y])
          slack[v] = lx[x] + lv[v] - gain[x][v], slackx[v]
   void augment(int x, int y) {
       while (x != -2) {
          yx[y] = x; swap(xy[x], y); x = p[x];
```

```
void improve() {
   S.assign(N, false), T.assign(N, false), p.assign(N,
        -1):
   gi = 0, g.clear();
   rep(x, N) if (xy[x] == -1) {
       q.push_back(root = x), p[x] = -2, S[x] = true;
       break:
   rep(y, N) slack[y] = lx[root] + ly[y] - gain[root][y
        1. slackx[v] = root:
   while (true) {
       while (qi < q.size()) {</pre>
          int x = q[qi++];
          rep(y, N) if (lx[x] + ly[y] == gain[x][y] &&!
               T[v]) {
              if (vx[v] == -1) return augment(x, v):
              T[y] = true, q.push_back(yx[y]), add(yx[y
                  ], x);
       }
       11 d = INF:
       rep(y, N) if (!T[y]) d = min(d, slack[y]);
       rep(x, N) if (S[x]) lx[x] -= d;
       rep(y, N) if (T[y]) ly[y] += d;
       rep(y, N) if (!T[y]) slack[y] -= d;
       rep(y, N) if (!T[y] && slack[y] == 0) {
           if (yx[v] == -1) return augment(slackx[v], v);
          T[v] = true:
          if (!S[yx[y]]) q.push_back(yx[y]), add(yx[y],
               slackx[v]):
      }
   }
Hungarian(vector<vector<11>>> g)
   : N(g.size()), gain(g), xy(N, -1), yx(N, -1), lx(N, -
   lv(N). slack(N). slackx(N) {
   rep(x, N) rep(y, N) lx[x] = max(lx[x], ly[y]);
   rep(i, N) improve();
```

};

#### 3.11 kuhn

```
// get a maximum cardinality matching in a bipartite graph.
// input: adjacency lists.
// output: matching (in 'mt' member).
// runtime: O(V E)
struct Kuhn {
   vector<vector<int>> G:
   int N, size;
   vector<bool> seen;
   vector<int> mt:
   bool visit(int i) {
       if (seen[i]) return false;
       seen[i] = true:
       for (int to : G[i])
           if (mt[to] == -1 || visit(mt[to])) {
              mt[to] = i;
              return true:
       return false;
   Kuhn(vector<vector<int>> adj) : G(adj), N(G.size()), mt(N
        , -1) {
       rep(i, N) {
          seen.assign(N, false);
           size += visit(i);
   }
};
```

#### 3.12 lca

```
// calculates the lowest common ancestor for any two nodes
    in O(log N) time,
// with O(N log N) preprocessing
struct Lca {
    int N, K, t = 0;
    vector<vector<int>> U;
    vector<int>> L, R;

Lca() {}
Lca(vector<vector<int>> &G) : N(G.size()), L(N), R(N) {
        K = N <= 1 ? 0 : 32 - __builtin_clz(N - 1);
        U.resize(K + 1, vector<int>(N));
        visit(G, 0, 0);
        rep(k, K) rep(u, N) U[k + 1][u] = U[k][U[k][u]];
}
```

```
void visit(vector<vector<int>> &G, int u, int p) {
    L[u] = t++, U[0][u] = p;
    for (int v : G[u]) if (v != p) visit(G, v, u);
    R[u] = t++;
}

bool is_anc(int up, int dn) {
    return L[up] <= L[dn] && R[dn] <= R[up];
}

int find(int u, int v) {
    if (is_anc(u, v)) return u;
    if (is_anc(v, u)) return v;
    for (int k = K; k >= 0;)
        if (is_anc(U[k][u], v)) k--;
        else u = U[k][u];
    return U[0][u];
}
```

#### 3.13 maxflow-mincost

```
// time: O(F V E)
                          F is the maximum flow
       O(V E + F E log V) if bellman-ford is replaced by
    johnson
struct Flow {
   struct Edge {
      int u, v;
       11 c, w, f = 0;
   }:
   int N, s, t;
   vector<vector<int>> G:
   vector<Edge> E;
   vector<ll> d, b;
   vector<int> p;
   Flow() {}
   Flow(int N, int s, int t) : N(N), s(s), t(t), G(N) {}
   void add_edge(int u, int v, ll c, ll w) {
       G[u].push_back(E.size());
       E.push_back({u, v, c, w});
      G[v].push back(E.size()):
       E.push_back({v, u, 0, -w});
   // naive distances with bellman-ford: O(V E)
```

```
void calcdists() {
   p.assign(N, -1), d.assign(N, INF), d[s] = 0;
   rep(i, N - 1) rep(ei, E.size()) {
       Edge &e = E[ei]:
       ll n = d[e.u] + e.w;
       if (d[e.u] < INF && e.c - e.f > 0 && n < d[e.v])
           d[e.v] = n, p[e.v] = ei;
   }
// johnsons potentials: O(E log V)
void calcdists() {
   if (b.empty()) {
       b.assign(N, 0);
       // code below only necessary if there are
            negative costs
       rep(i, N - 1) rep(ei, E.size()) {
           Edge &e = E[ei];
          if (e.f < e.c) b[e.v] = min(b[e.v], b[e.u] + e
       }
   p.assign(N, -1), d.assign(N, INF), d[s] = 0;
   priority_queue<pair<11, int>> q;
   q.push({0, s});
   while (!q.empty()) {
       auto [w, u] = q.top();
       g.pop():
       if (d[u] < -w + b[u]) continue;</pre>
       for (int ei : G[u]) {
           auto e = E[ei];
           11 n = d[u] + e.w;
           if (e.f < e.c && n < d[e.v]) {
              d[e.v] = n, p[e.v] = ei;
              q.push({b[e.v] - n, e.v});
       }
   }
   b = d;
11 solve() {
   b.clear():
   11 \text{ ff} = 0;
   while (true) {
       calcdists():
       if (p[t] == -1) break;
       11 f = INF:
```

## 3.14 parallel-dfs

```
struct Tree {
    int n,z[2];
    vector<vector<int>> g;
    vector<int> ex,ey,p,w,f,v[2];
    Tree(int n):g(n), w(n), f(n){}
    void add_edge(int x, int y){
       p.pb(g[x].size());g[x].pb(ex.size());
        ex.pb(x); ey.pb(y);
       p.pb(g[y].size());g[y].pb(ex.size());
        ex.pb(y);ey.pb(x);
    bool go(int k){//returns 1 if it finds new node
       int & x=z[k];
       while(x \ge 0 \& \&
           (w[x] == g[x].size() | |w[x] == g[x].size()-1
           &&(g[x].back()^1)==f[x])
           x=f[x]>=0?ex[f[x]]:-1:
       if(x<0)return false:</pre>
       if((g[x][w[x]]^1)==f[x])w[x]++;
       int e=g[x][w[x]],y=ey[e]; f[y]=e;
       w[x]++; w[y]=0; x=y; v[k].pb(x);
        return true;
    vector<int> erase_edge(int e){
        e*=2;//erases eth edge, returns smaller comp
       int x=ex[e],y=ey[e]; p[g[x].back()]=p[e];
       g[x][p[e]]=g[x].back(); g[x].pop_back();
       p[g[y].back()]=p[e^1]; g[y][p[e^1]]=g[y].back();
       g[y].pop_back();
       f[x]=f[y]=-1; w[x]=w[y]=0; z[0]=x;z[1]=y;
       v[0]=\{x\}:v[1]=\{v\}:
       bool d0=true,d1=true;while(d0&&d1)d0=go(0),d1=go(1);
        return v[1-d1]:
    }
};
```

## 3.15 push-relabel

```
#include "../common.h"
const 11 INF = 1e18:
// maximum flow algorithm.
// to run, use 'maxflow()'.
11
// time: O(V^2 \operatorname{sgrt}(E)) \leq O(V^3)
// memory: 0(V^2)
struct PushRelabel {
   vector<vector<ll>> cap. flow:
   vector<ll> excess:
   vector<int> height;
   PushRelabel() {}
   void resize(int N) { cap.assign(N, vector<11>(N)); }
   // push as much excess flow as possible from u to v.
   void push(int u, int v) {
       11 f = min(excess[u], cap[u][v] - flow[u][v]);
       flow[u][v] += f;
       flow[v][u] -= f:
       excess[v] += f;
       excess[u] -= f;
   // relabel the height of a vertex so that excess flow may
         be pushed.
   void relabel(int u) {
       int d = INT32 MAX:
       rep(v, cap.size()) if (cap[u][v] - flow[u][v] > 0) d
           min(d, height[v]):
       if (d < INF) height[u] = d + 1;</pre>
   // get the maximum flow on the network specified by 'cap'
         with source 's'
   // and sink 't'.
   // node-to-node flows are output to the 'flow' member.
   11 maxflow(int s. int t) {
       int N = cap.size(), M;
       flow.assign(N, vector<ll>(N));
       height.assign(N, 0), height[s] = N;
       excess.assign(N, 0), excess[s] = INF;
```

```
rep(i, N) if (i != s) push(s, i):
vector<int> q;
while (true) {
   // find the highest vertices with excess
    q.clear(), M = 0:
    rep(i, N) {
       if (excess[i] \le 0 \mid | i == s \mid | i == t)
            continue:
       if (height[i] > M) q.clear(), M = height[i];
       if (height[i] >= M) g.push back(i);
    if (q.empty()) break;
    // process vertices
    for (int u : q) {
       bool relab = true;
       rep(v. N) {
           if (excess[u] <= 0) break;</pre>
           if (cap[u][v] - flow[u][v] > 0 && height[u
               ] > height[v])
               push(u, v), relab = false;
       if (relab) {
           relabel(u);
           break;
}
11 f = 0: rep(i, N) f += flow[i][t]: return f:
```

## 3.16 strongly-connected-components

};

```
/* time: O(V + E), memory: O(V)
after building:
   comp = map from vertex to component
        (components are toposorted, root first, leaf last)
   N = number of components
   G = condensation graph (component DAG)
byproducts:
   vgi = transposed graph
   order = reverse topological sort (leaf first, root last)
others:
   vn = number of vertices
   vg = original vertex graph  */
struct Scc {
   int vn, N;
```

```
vector<int> order. comp:
   vector<vector<int>> vg, vgi, G;
   void toposort(int u) {
       if (comp[u]) return:
       comp[u] = -1;
       for (int v : vg[u]) toposort(v);
       order.push_back(u);
   bool carve(int u) {
       if (comp[u] != -1) return false;
       comp[u] = N:
      for (int v : vgi[u]) {
          carve(v):
          if (comp[v] != N) G[comp[v]].push_back(N);
       return true;
   Scc() {}
   Scc(vector<vector<int>> &g)
    : vn(g.size()), vg(g), comp(vn), vgi(vn), G(vn), N(0) {
       rep(u, vn) toposort(u);
       rep(u, vn) for (int v : vg[u]) vgi[v].push_back(u);
       invrep(i, vn) N += carve(order[i]);
};
```

#### 3.17 two-sat

```
// calculate the solvability of a system of logical
    equations, where every equation is of the form 'a or b
// 'neg': get negation of 'u'
// 'then': 'u' implies 'v'
// 'anv': 'u' or 'v'
// 'set': 'u' is true
11
// after 'solve' (O(V+E)) returns true, 'sol' contains one
    possible solution.
// determining all solutions is O(V*E) hard (requires
    computing reachability in a DAG).
struct TwoSat {
   int N; vector<vector<int>> G;
   Scc scc; vector<bool> sol;
   TwoSat(int n) : N(n), G(2 * n), sol(n) {}
   TwoSat() {}
   int neg(int u) { return (u + N) % (2 * N); }
   void then(int u, int v) { G[u].push_back(v), G[neg(v)].
        push_back(neg(u)); }
```

```
void any(int u, int v) { then(neg(u), v); }
void set(int u) { G[neg(u)].push_back(u); }

bool solve() {
    scc = Scc(G);
    rep(u, N) if (scc.comp[u] == scc.comp[neg(u)]) return
        false;
    rep(u, N) sol[u] = (scc.comp[u] > scc.comp[neg(u)]);
    return true;
}
```

## 4 implementation

### 4.1 common-template-and-bit-tricks

```
#pragma GCC optimize("Ofast")
#pragma GCC target("bmi,bmi2,lzcnt,popcnt")
#pragma GCC target("avx,avx2,f16c,fma,sse3,sse3,sse4.1,sse4
     .2")
#include <bits/stdc++.h>
using namespace std;
typedef long long 11:
#define repx(i, a, b) for (int i = a; i < b; i++)
#define rep(i, n) repx(i, 0, n)
#define invrepx(i, a, b) for (int i = b - 1; i \ge a; i--)
#define invrep(i, n) invrepx(i, 0, n)
// Command to check time and memory usage:
       /usr/bin/time -v ./tmp
// See "Maximum resident set size" for max memory used
// Commands for interactive checker:
// mkfifo fifo
11
      (./solution < fifo) | (./interactor > fifo)
// Does not work on the Windows file system, i.e., /mnt/c/
// The special fifo file must be used, otherwise the
// solution will not wait for input and will read EOF
y = x & (x-1) // Turn off rightmost 1bit
y = x & (-x) // Isolate rightmost 1bit
y = x \mid (x-1) // Right propagate rightmost 1bit(fill in 1s)
y = x \mid (x+1) // Turn on rightmost Obit
y = x & (x+1) // Isolate rightmost Obit
// If x is of long type, use __builtin_popcountl(x)
// If x is of long long type, use __builtin_popcountll(x)
// 1. Counts the number of ones(set bits) in an integer.
__builtin_popcount(x)
// 2. Checks the Parity of a number. Returns true(1) if the
// number has odd number of set bits. else it returns
// false(0) for even number of set bits.
```

```
__builtin_parity(x)

// 3. Counts the leading number of zeros of the integer.
__builtin_clz(x)

// 4. Counts the trailing number of zeros of the integer.
__builtin_ctz(x)

// 5. Returns 1 + the index of the least significant 1-bit.
__builtin_ffs(x) // If x == 0, returns 0.

// Iterate over non empty subsets of bitmask
for(int s=m;s;s=(s-1)&m) // Decreasing order
for(int s=0;s=s-m&m;) // Increasing order
```

## 4.2 dp-convex-hull-trick

```
struct Line {
   mutable ll a. b. c:
   bool operator<(Line r) const { return a < r.a; }</pre>
   bool operator<(ll x) const { return c < x; }</pre>
// dynamically insert 'a*x + b' lines and query for maximum
// at any x all operations have complexity O(\log N)
struct LineContainer : multiset<Line, less<>>> {
   11 div(ll a, ll b) {
       return a / b - ((a ^ b) < 0 && a % b):
   bool isect(iterator x, iterator y) {
       if (y == end()) return x->c = INF, 0;
       if (x->a == y->a) x->c = x->b > y->b ? INF : -INF;
       else x->c = div(y->b - x->b, x->a - y->a);
       return x->c >= y->c;
   void add(ll a. ll b) {
       // a *= -1, b *= -1 // for min
       auto z = insert(\{a, b, 0\}), y = z++, x = y;
       while (isect(y, z)) z = erase(z);
       if (x != begin() \&\& isect(--x, y)) isect(x, y = erase)
       while ((y = x) != begin() \&\& (--x)->c >= y->c) isect(
            x, erase(y));
   }
   11 query(11 x) {
       if (empty()) return -INF; // INF for min
       auto 1 = *lower_bound(x);
       return 1.a * x + 1.b;
       // return -l.a * x - l.b: // for min
```

```
,
```

## 4.3 dp-divide-and-conquer

```
// for every index i assign an optimal index j, such that
// cost(i, j) is minimal for every i. the property that if
// i2 >= i1 then j2 >= j1 is exploited (monotonic condition)
// calculate optimal index for all indices in range [1, r)
// knowing that the optimal index for every index in this
// range is within [optl, optr).
// time: O(N log N)
void calc(vector<int> &opt, int 1, int r,int optl,int optr){
   if (1 == r) return:
   int i = (1 + r) / 2;
   11 optc = INF:
   int optj;
   repx(j, optl, optr) {
      11 c = i + j; // cost(i, j)
      if (c < optc) optc = c, optj = j;</pre>
   opt[i] = optj;
   calc(opt, l, i, optl, optj + 1);
   calc(opt, i + 1, r, optj, optr);
```

## 4.4 dynamic-connectivity

```
struct DC {
   int n; Dsu D;
   vector<vector<pair<int, int>>> t;
   DC(int N) : n(N), D(N), t(2 * N) {}
   // add edge p to all times in interval [1, r]
   void upd(int 1, int r, pair<int, int> p) {
       for (1 += n, r += n; 1 < r; 1 >>= 1, r >>= 1) {
          if (1 & 1) t[1++].push_back(p);
          if (r & 1) t[--r].push back(p):
      }
   }
   void process(int u = 1) { // process all queries
       for (auto &e : t[u]) D.unite(e.first, e.second);
       if (u >= n) {
          // do stuff with D at time u - n
       } else process(2 * u), process(2 * u + 1);
       for (auto &e : t[u]) D.rollback();
};
```

#### 4.5 hash-container

```
namespace{//add (#define tmpl template)(#define ty typename)
  tmpl<ty T> size_t mk_h(const T& v){return hash<T>()(v);}
  void h_cmb(size_t& h, const size_t& v)
  { h ~= v + 0x9e3779b9 + (h << 6) + (h >> 2); }
  tmpl<ty T> struct h_ct{size_t operator()(const T& v)const{
  size_t h=0;for(const auto& e:v){h_cmb(h,mk_h(e));}return h;
  }};
}namespace std{//support for pair<T,U>, vector<T> & map<T,U>
  tmpl<ty T, ty U> struct hash<pair<T, U>>{
    size_t operator()(const pair<T,U>& v) const
  {size_t h=mk_h(v.first);h_cmb(h, mk_h(v.second));return h;
  };
  tmpl<ty... T>struct hash<vector<T...>:h_ct<vector<T...>>{};
  tmpl<ty... T>struct hash<map<T...>>:h_ct<map<T...>>{};
}
```

#### 4.6 mo

```
struct Query { int 1, r, idx; };
// answer segment queries using only 'add(i)', 'remove(i)'
    and 'get()'
// functions.
11
// complexity: O((N + Q) * sqrt(N) * F)
// N = length of the full segment
// Q = amount of queries
// F = complexity of the 'add', 'remove' functions
template <class A, class R, class G, class T>
void mo(vector<Query> &queries, vector<T> &ans, A add, R
    remove, G get) {
   int Q = queries.size(), B = (int)sqrt(Q);
   sort(queries.begin(), queries.end(), [&](Query &a, Query
       return make_pair(a.1 / B, a.r) < make_pair(b.1 / B, b</pre>
   });
   ans.resize(0):
   int 1 = 0, r = 0:
   for (auto &q : queries) {
       while (r < q.r) add(r), r++;
       while (1 > q.1) 1--, add(1);
       while (r > q.r) r--, remove(r):
       while (1 < q.1) remove(1), 1++;
       ans[q.idx] = get();
```

#### 4.7 ordered-set

## 4.8 unordered-map

## $5 \mod n$

## 5.1 Linear Diophantine

```
ii extendedEuclid(11 a, 11 b){
    11 x, y; //a*x + b*y = gcd(a,b)
    if (b == 0) return {1, 0};
    auto p = extendedEuclid(b, a%b);
    x = p.second;
    y = p.first - (a/b)*x;
    if(a*x + b*y == -__gcd(a,b)) x=-x, y=-y;
    return {x, y};
}
```

```
pair<ii, ii> diophantine(11 a, 11 b, 11 r){
    //a*x+b*y=r where r is multiple of gcd(a,b);
    1l d = __gcd(a, b);
    a/=d; b/=d; r/=d;
    auto p = extendedEuclid(a, b);
    p.first*=r; p.second*=r;
    assert(a*p.first + b*p.second == r);
    return {p, {-b, a}}; //solutions: p+t*ans.second
}
```

#### 5.2 arithmetic

```
inline int floor_log2(int n)
{ return n <= 1 ? 0 : 31 - __builtin_clz(n); }
inline int ceil_log2(int n)
{ return n <= 1 ? 0 : 32 - __builtin_clz(n - 1); }
inline ll floordiv(ll a, ll b) {return a/b-((a^b)<0&&a%b);}
inline ll ceildiv(ll a, ll b) {return a/b+((a^b)>=0&&a%b);}
```

## 5.3 berlekamp-massey-linear-recurrence

```
vector<int> BM(vector<int> x) {
    vector<int> ls, cur;
    int lf. ld:
    rep(i, x.size()) {
       11 t = 0:
       rep(j, cur.size()) t = (t+x[i-j-1]*(ll)cur[j])%MOD;
       if ((t - x[i]) \% MOD == 0) continue;
       if (!cur.size()) {
           cur.resize(i + 1): lf = i: ld = (t-x[i]) % MOD:
           continue;
       11 k = -(x[i] - t) * bin_exp(1d, MOD - 2) % MOD;
       vector<int> c(i - lf - 1); c.push_back(k);
       rep(j, ls.size()) c.push_back(-ls[j] * k % MOD);
       if (c.size() < cur.size()) c.resize(cur.size());</pre>
       rep(j, cur.size()) c[j] = (c[j] + cur[j]) % MOD;
       if (i - lf + ls.size() >= cur.size())
           ls = cur, lf = i, ld = (t - x[i]) % MOD;
    rep(i, cur.size()) cur[i] = (cur[i] % MOD + MOD) % MOD;
    return cur:
// Linear Recurrence
11 MOD = 998244353:
11 LOG = 60;
```

```
struct LinearRec{
 typedef vector<int> vi:
 int n; vi terms, trans; vector<vi> bin;
 vi add(vi &a. vi &b){
   vi res(n*2+1):
   rep(i.n+1) rep(i.n+1)
       res[i+j]=(res[i+j]*1LL+(ll)a[i]*b[j])%MOD;
   for(int i=2*n; i>n; --i){
      res[i-1-j]=(res[i-1-j]*1LL+(ll)res[i]*trans[j])%MOD;
     res[i]=0:
   res.erase(res.begin()+n+1,res.end());
   return res:
 LinearRec(vi &terms, vi &trans):terms(terms),trans(trans){
   n=trans.size():vi a(n+1):a[1]=1:
   bin.push_back(a);
   repx(i,1,LOG)bin.push back(add(bin[i-1],bin[i-1]));
 int calc(ll k){
   vi a(n+1);a[0]=1;
   rep(i,LOG)if((k>>i)&1)a=add(a,bin[i]);
   rep(i,n)ret=((ll)ret+(ll)a[i+1]*terms[i])%MOD;
   ret = ret%MOD + MOD;
   return ret%MOD:
};
```

#### 5.4 crt

```
pair<11, ll> solve_crt(const vector<pair<11, ll>> &eqs) {
    ll a0 = eqs[0].first, p0 = eqs[0].second;
    repx(i, 1, eqs.size()) {
        ll a1 = eqs[i].first, p1 = eqs[i].second;
        ll k1, k0;
        ll d = ext_gcd(p1, p0, k1, k0);
        a0 -= a1;
        if (a0 % d != 0) return {-1, -1};
        p0 = p0 / d * p1;
        a0 = a0 / d * k1 % p0 * p1 % p0 + a1;
        a0 = (a0 % p0 + p0) % p0;
    }
    return {a0, p0};
}
```

#### 5.5 discrete-log

```
// discrete logarithm log_a(b).
// solve b \hat{x} = a \pmod{M} for the smallest x.
// returns -1 if no solution is found.
//
// time: O(sqrt(M))
11 dlog(ll a, ll b, ll M) {
   11 k = 1, s = 0;
   while (true) {
       11 g = \_gcd(b, M);
       if (g <= 1) break;</pre>
       if (a == k) return s:
       if (a % g != 0) return -1;
       a = g, M = g, s += 1, k = b / g * k % M;
   11 N = sqrt(M) + 1;
   umap<11. 11> r:
   rep(q, N + 1) {
       r[a] = q;
       a = a * b % M:
   ll bN = binexp(b, N, M), bNp = k;
   repx(p, 1, N + 1) {
       bNp = bNp * bN % M;
       if (r.count(bNp)) return N * p - r[bNp] + s;
   return -1:
```

### 5.6 fast-hadamard-transform

```
11 c1[MAXN+9],c2[MAXN+9];//MAXN must be power of 2!
void fht(ll* p, int n, bool inv){
    for(int l=1;2*l<=n;l*=2)for(int i=0;i<n;i+=2*l)fore(j,0,1)
        ){
        ll u=p[i+j],v=p[i+l+j];
        if(!inv)p[i+j]=u+v,p[i+l+j]=u-v; // XOR
        else p[i+j]=(u+v)/2,p[i+l+j]=(u-v)/2;
        //if(!inv)p[i+j]=v,p[i+l+j]=u+v; // AND
        //else p[i+j]=-u+v,p[i+l+j]=u;
        //if(!inv)p[i+j]=u+v,p[i+l+j]=u;
        //else p[i+j]=v,p[i+l+j]=u-v;
    }
}
// like polynomial multiplication, but XORing exponents
// instead of adding them (also ANDing, ORing)</pre>
```

```
vector<11> multiply(vector<11>& p1, vector<11>& p2){
   int n=1<<(32-_builtin_clz(max(SZ(p1),SZ(p2))-1));
   fore(i,0,n)c1[i]=0,c2[i]=0;
   fore(i,0,SZ(p1))c1[i]=p1[i];
   fore(i,0,SZ(p2))c2[i]=p2[i];
   fht(c1,n,false);fht(c2,n,false);
   fore(i,0,n)c1[i]*=c2[i];
   fht(c1,n,true);
   return vector<11>(c1,c1+n);
}
```

#### 5.7 fft

```
using cd = complex<double>;
const double PI = acos(-1):
// compute the DFT of a power-of-two-length sequence.
// if 'inv' is true, computes the inverse DFT.
void fft(vector<cd> &a. bool inv) {
   int N = a.size(), k = 0, b;
   assert(N == 1 << __builtin_ctz(N));</pre>
   repx(i, 1, N) {
      for (b = N >> 1; k & b;) k ^= b, b >>= 1;
       if (i < (k ^= b)) swap(a[i], a[k]);</pre>
   for (int 1 = 2; 1 <= N; 1 <<= 1) {</pre>
       double ang = 2 * PI / 1 * (inv ? -1 : 1);
       cd wl(cos(ang), sin(ang));
       for (int i = 0; i < N; i += 1) {</pre>
           cd w = 1:
           rep(j, 1 / 2) {
              cd u = a[i + j], v = a[i + j + 1 / 2] * w;
              a[i + i] = u + v:
              a[i + j + 1 / 2] = u - v;
              w *= wl:
   }
   if (inv) rep(i, N) a[i] /= N;
const 11 MOD = 998244353, ROOT = 15311432;
// const 11 MOD = 2130706433, ROOT = 1791270792;
// const 11 MOD = 922337203673733529711, ROOT =
    532077456549635698311;
void find_root_of_unity(ll M) {
   11 c = M - 1, k = 0:
   while (c \% 2 == 0) c /= 2, k += 1;
   // find proper divisors of M - 1
   vector<ll> divs:
   for (ll d = 1; d < c; d++) {
```

```
if (d * d > c) break:
       if (c \% d == 0) rep(i, k + 1) divs.push_back(d << i);
   rep(i, k) divs.push_back(c << i);</pre>
   // find any primitive root of M
   11 G = -1:
   repx(g, 2, M) {
       bool ok = true;
       for (ll d : divs) ok &= (binexp(g, d, M) != 1);
       if (ok) {
          G = g;
          break:
      }
   assert(G != -1):
   11 w = binexp(G, c, M);
   cerr << "M = c * 2^k + 1" << endl:
   cerr << " M = " << M << endl;
   cerr << " c = " << c << endl:
   cerr << " k = " << k << endl:
   cerr << " w^(2^k) == 1" << endl;
   cerr << " w = g^{(M-1)/2k} = g^c << endl;
   cerr << " g = " << G << endl;</pre>
   cerr << " w = " << w << endl:
// compute the DFT of a power-of-two-length sequence, modulo
     a special prime
// number with an Nth root of unity, where N is the length
    of the sequence.
void ntt(vector<11> &a. bool inv) {
   vector<ll> wn:
   for (11 p = ROOT; p != 1; p = p * p % MOD) wn.push_back(p
        ):
   int N = a.size(), k = 0, b;
   assert(N == 1 << builtin ctz(N) && N <= 1 << wn.size())
   rep(i, N) a[i] = (a[i] \% MOD + MOD) \% MOD;
   repx(i, 1, N) {
      for (b = N >> 1; k & b;) k ^= b, b >>= 1;
       if (i < (k ^= b)) swap(a[i], a[k]);</pre>
   for (int 1 = 2; 1 <= N; 1 <<= 1) {
      11 wl = wn[wn.size() - builtin ctz(1)]:
       if (inv) wl = multinv(wl, MOD);
      for (int i = 0; i < N; i += 1) {</pre>
          11 w = 1:
          repx(i, 0, 1 / 2) {
              11 u = a[i + j], v = a[i + j + 1 / 2] * w %
              a[i + j] = (u + v) \% MOD;
```

#### 5.8 gauss

```
const double EPS = 1e-9;
// solve a system of equations.
// complexity: O(\min(N, M) * N * M)
// 'a' is a list of rows
// the last value in each row is the result of the equation
// return values:
// 0 -> no solutions
// 1 -> unique solution, stored in 'ans'
// -1 -> infinitely many solutions, one of which is stored
    in 'ans'
// UNTESTED
int gauss(vector<vector<double>> a, vector<double> &ans) {
   int N = a.size(), M = a[0].size() - 1:
   vector<int> where(M, -1):
   for (int j = 0, i = 0; j < M && i < N; j++) {
       int sel = i:
       repx(k, i, N) if (abs(a[k][i]) > abs(a[sel][i])) sel
            = k:
       if (abs(a[sel][i]) < EPS) continue;</pre>
       repx(k, j, M + 1) swap(a[sel][k], a[i][k]);
       where[j] = i;
       rep(k, N) if (k != i) {
          double c = a[k][j] / a[i][j];
          repx(1, j, M + 1) a[k][1] -= a[i][1] * c;
      }
       i++;
   ans.assign(M, 0);
   rep(i, M) if (where[i] != -1) ans[i] = a[where[i]][M] / a
        [where[i]][i]:
   rep(i, N) {
```

```
double sum = 0;
    rep(j, M) sum += ans[j] * a[i][j];
    if (abs(sum - a[i][M]) > EPS) return 0;
}
rep(i, M) if (where[i] == -1) return -1;
return 1;
```

#### 5.9 matrix

```
typedef vector<vector<double>> Mat:
Mat matmul(Mat 1, Mat r) {
   int n = 1.N, m = r.M, p = 1.M; assert(1.M == r.N);
   Mat a(n, vector<double>(m)); // neutral
   rep(i, n) rep(j, m)
      rep(k, p) a[i][j] = a[i][j] + l[i][k] * r[k][j];
   return a:
double reduce(vector<vector<double>> &A) {
   int n = A.size(), m = A[0].size();
   int i = 0, i = 0; double r = 1.:
   while (i < n && j < m) {</pre>
      int 1 = i:
       repx(k, i+1, n) if(abs(A[k][j]) > abs(A[l][j])) l=k;
       if (abs(A[1][i]) < EPS) \{ i++; r = 0.; continue; \}
      if (1 != i) { r = -r; swap(A[i], A[1]); }
      r *= A[i][i];
      for (int k = m - 1; k >= j; k--) A[i][k] /= A[i][j];
       repx(k, 0, n) {
          if (k == i) continue:
          for(int l=m-1;l>=j;l--)A[k][l]-=A[k][j]*A[i][l];
       i++, j++;
   return r; // returns determinant
```

## 5.10 mobius

### 5.11 multiny

```
// a * x + b * y == gcd(a, b)
ll ext_gcd(ll a, ll b, ll &x, ll &y) {
    if (b == 0) { x = 1, y = 0; return a; }
    ll d = ext_gcd(b, a % b, y, x); y -= a / b * x; return d;
}

// inverse exists if and only if a and M are coprime
// if M is prime: multinv(a, M) = (a**(M-2)) % M
ll multinv(ll a, ll M)
{ ll x, y; ext_gcd(a, M, x, y); return x; }

// all modular inverses from 1 to inv.size()-1
void multinv_all(vector<ll> &inv) {
    inv[1] = 1;
    repx(i, 2, inv.size())
        inv[i] = MOD - (MOD / i) * inv[MOD % i] % MOD;
}
```

### 5.12 polar-rho

```
11 mulmod(l1 a, 11 b, 11 m) {
   11 r=a*b-(11)((long double)a*b/m+.5)*m;
   return r<0?r+m:r;</pre>
bool is_prime_prob(ll n, int a){
   if(n==a)return true:
   ll s=0.d=n-1:
   while (d\%2==0)s++, d/=2;
   11 x=expmod(a,d,n):
   if((x==1)||(x+1==n))return true:
   fore(_,0,s-1){
       x=mulmod(x.x.n):
       if(x==1)return false;
       if(x+1==n)return true;
   return false;
bool rabin(ll n){ // true iff n is prime
   if(n==1)return false:
   int ar[]={2,3,5,7,11,13,17,19,23};
   fore(i,0,9)if(!is_prime_prob(n,ar[i]))return false;
   return true:
ll rho(ll n){
   if(!(n&1))return 2;
   11 x=2,y=2,d=1;
   ll c=rand()%n+1;
```

```
while(d==1){
       x=(\text{mulmod}(x,x,n)+c)%n:
       fore(it,0,2) y=(mulmod(y,y,n)+c)%n;
       if(x>=y)d=_gcd(x-y,n);
       else d=__gcd(y-x,n);
    return d==n?rho(n):d;
void fact(ll n, map<ll,int>& f){ //0 (lg n)^3
    if(n==1)return:
    if(rabin(n)){f[n]++:return:}
    11 g=rho(n):fact(g,f):fact(n/g,f):
}
// optimized version: replace rho and fact with the
     following:
const int MAXP=1e6+1; // sieve size
int sv[MAXP]: // sieve
11 add(l1 a, l1 b, l1 m){return (a+=b)<m?a:a-m;}</pre>
11 rho(11 n){
    static ll s[MAXP]:
    while(1){
       11 x=rand()%n,y=x,c=rand()%n;
       11 *px=s,*py=s,v=0,p=1;
       while(1){
           *py++=y=add(mulmod(y,y,n),c,n);
           *py++=y=add(mulmod(y,y,n),c,n);
           if((x=*px++)==y)break;
           11 t=p; p=mulmod(p,abs(y-x),n);
           if(!p)return __gcd(t,n);
           if(++v==26){
               if((p=__gcd(p,n))>1&&p<n)return p;</pre>
           }
       if(v&&(p=_gcd(p,n))>1&&p<n)return p;</pre>
}
void init_sv(){ fore(i,2,MAXP)if(!sv[i])for(11 j=i;j<MAXP;j</pre>
     +=i)sv[j]=i; }
void fact(ll n,map<ll,int>&f){//call init_sv first!
    for(auto&& p:f)while(n%p.fst==0)p.snd++.n/=p.fst:
    if(n \le MAXP) while (n \ge 1) f[sv[n]] ++ , n/=sv[n];
    else if(rabin(n))f[n]++;
    else {ll q=rho(n);fact(q,f);fact(n/q,f);}
```

## 5.13 polynomials

```
typedef int tp; // type of polynomial
```

```
template<class T=tp>
struct poly { // poly<> : 1 variable, poly<poly<>>>: 2
    variables, etc.
   vector<T> c:
   T& operator[](int k){return c[k];}
   poly(vector<T>& c):c(c){}
   poly(initializer_list<T> c):c(c){}
   poly(int k):c(k){}
   poly(){}
   polv operator+(polv<T> o);
   poly operator*(tp k);
   polv operator*(polv o):
   poly operator-(poly<T> o){return *this+(o*-1);}
   T operator()(tp v){
      T sum(0):
       for(int i=c.size()-1;i>=0;--i)sum=sum*v+c[i];
       return sum:
   }
// example: p(x,y)=2*x^2+3*x*y-y+4
// poly<poly<>> p={{4,-1},{0,3},{2}}
// printf("d\n",p(2)(3)) // 27 (p(2,3))
set<tp> roots(poly<> p){ // only for integer polynomials
   while(!p.c.empty()&&!p.c.back())p.c.pop_back();
   if(!p(0))r.insert(0);
   if(p.c.empty())return r;
   tp a0=0,an=abs(p[p.c.size()-1]);
   for(int k=0;!a0;a0=abs(p[k++]));
   vector<tp> ps.qs:
   fore(i,1,sqrt(a0)+1)if(a0%i==0)ps.pb(i),ps.pb(a0/i);
   fore(i,1,sqrt(an)+1)if(an\%i==0)qs.pb(i),qs.pb(an/i);
   for(auto pt:ps)for(auto qt:qs)if(pt%qt==0){
       tp x=pt/qt;
       if(!p(x))r.insert(x);
      if(!p(-x))r.insert(-x):
   }
   return r;
pair<poly<>,tp> ruffini(poly<> p, tp r){ // returns pair (
    result.rem)
   int n=p.c.size()-1;
   vector<tp> b(n):
   b[n-1]=p[n];
   for(int k=n-2; k>=0; --k)b[k]=p[k+1]+r*b[k+1];
   return {poly<>(b),p[0]+r*b[0]};
// only for double polynomials
pair<poly<>,poly<> > polydiv(poly<> p, poly<> q){ // returns
     pair (result,rem)
```

```
int n=p.c.size()-q.c.size()+1:
   vector<tp> b(n):
   for(int k=n-1;k>=0;--k){
       b[k]=p.c.back()/q.c.back();
       fore(i,0,q.c.size())p[i+k]-=b[k]*q[i];
       p.c.pop_back();
   while(!p.c.empty()&&abs(p.c.back()) < EPS)p.c.pop_back();</pre>
   return {poly<>(b),p};
// only for double polynomials
polv<> interpolate(vector<tp> x, vector<tp> v){
   polv<> q={1},S={0};
   for(tp a:x)q=poly<>({-a,1})*q;
   fore(i,0,x.size()){
       poly<> Li=ruffini(q,x[i]).fst;
       Li=Li*(1.0/Li(x[i])); // change for int polynomials
       S=S+Li*v[i];
   return S;
```

## 5.14 primes

```
// counts the divisors of a positive integer in O(\operatorname{sgrt}(n))
ll count divisors(ll x) {
   11 \text{ divs} = 1, i = 2;
   for (11 divs = 1, i = 2; x > 1; i++) {
       if (i * i > x) { divs *= 2; break; }
       for (11 d = divs; x % i == 0; x /= i) divs += d;
   return divs:
// gets the prime factorization of a number in O(sqrt(n))
vector<pair<ll, int>> factorize(ll x) {
   vector<pair<11, int>> f;
   for (11 k = 2: x > 1: k++) {
       if (k * k > x) { f.push_back({x, 1}); break; }
       int n = 0:
       while (x \% k == 0) x /= k, n++;
       if (n > 0) f.push_back(\{k, n\});
   return f;
// iterate over all divisors of a number.
// divisor count upper bound: n^(1.07 / ln ln n)
template <class OP>
void divisors(ll x, OP op) {
   auto facts = factorize(x);
```

```
vector<int> f(facts.size()):
   while (true) {
       11 y = 1;
       rep(i, f.size()) rep(j, f[i]) y *= facts[i].first;
       op(y);
       int i;
       for (i = 0; i < f.size(); i++) {</pre>
           f[i] += 1:
           if (f[i] <= facts[i].second) break;</pre>
           f[i] = 0:
       if (i == f.size()) break;
}
// computes euler totative function phi(x), counting the
// amount of integers in [1, x] that are coprime with x.
// time: O(sqrt(x))
ll phi(ll x) {
   11 phi = 1. k = 2:
   for (; x > 1; k++) {
       if (k * k > x) { phi *= x - 1; break; }
       11 k1 = 1, k0 = 0;
       while (x \% k == 0) x /= k, k0 = k1, k1 *= k;
       phi *= k1 - k0;
   return phi;
}
// test-prime.cpp
// change to int128 if checking numbers over 10^9
bool isprime(ll n) {
   if (n < 2 | | n % 6 % 4 != 1) return n - 2 < 2;
   11 A[] = \{2.325.9375.28178.450775.9780504.1795265022\}:
   ll s = \_builtin\_ctzll(n - 1), d = n >> s;
   for (int a : A) {
       11 p = binexp(a, d, n), i = s;
       while (p != 1 && p != n - 1 && a % n && i--) p = p *
       if (p != n - 1 && i != s) return 0;
   return 1:
```

## 5.15 simplex

/\* Solves a general linear maximization problem: maximize  $c^T x$  subject to  $Ax \le b$ ,  $x \le 0$ . Returns -inf if there is no solution, inf if there are arbitrarily good solutions, or the maximum value of  $c^T x$  otherwise. The

```
input vector is set to an optimal $x$ (or in the unbounded
case, an arbitrary solution fulfilling the constraints).
Numerical stability is not guaranteed. For better
performance, define variables such that x = 0 is viable.
Usage:
vvd A = \{\{1,-1\}, \{-1,1\}, \{-1,-2\}\};
vd b = \{1,1,-4\}, c = \{-1,-1\}, x;
T val = LPSolver(A, b, c).solve(x);
Time: O(NM * \t pivots), where a pivot may be e.g. an edge
relaxation. O(2^n) in the general case.*/
typedef double T://long double, Rational, double + mod<P>...
typedef vector<T> vd:
typedef vector<vd> vvd;
const T eps = 1e-8, inf = 1 / .0;
#define MP make_pair
#define ltj(X) \
   if (s == -1 \mid | MP(X[i], N[i]) < MP(X[s], N[s])) s = i
struct LPSolver {
   int m. n: vector<int> N. B: vvd D:
   LPSolver(const vvd &A.const vd &b.const vd &c) : m(b.size
        ()),n(c.size()),N(n+1),B(m),D(m+2,vd(n+2))
       rep(i, m) rep(j, n) D[i][j] = A[i][j];
       rep(i, m) {
           B[i] = n + i; D[i][n] = -1; D[i][n + 1] = b[i];
       rep(j, n) \{ N[j] = j; D[m][j] = -c[j]; \}
       N[n] = -1; D[m + 1][n] = 1;
   void pivot(int r, int s) {
       T *a = D[r].data(), inv = 1 / a[s]:
       rep(i, m + 2) if (i != r && abs(D[i][s]) > eps) {
           T *b = D[i].data(), inv2 = b[s] * inv;
           repx(j, 0, n + 2) b[j] -= a[j] * inv2;
           b[s] = a[s] * inv2;
       rep(i, n + 2) if (i != s) D[r][i] *= inv:
       rep(i, m + 2) if (i != r) D[i][s] *= -inv;
       D[r][s] = inv:
       swap(B[r], N[s]);
   bool simplex(int phase) {
       int x = m + phase - 1;
       for (::) {
           int s = -1:
           rep(j, n + 1) if (N[j] != -phase) ltj(D[x]);
           if (D[x][s] \ge -eps) return true:
           int r = -1;
           rep(i, m) {
              if (D[i][s] <= eps) continue;</pre>
```

```
if (r == -1 \mid | MP(D[i][n + 1] / D[i][s], B[i])
                    < MP(D[r][n + 1] / D[r][s], B[r])) r = i
           if (r == -1) return false;
           pivot(r. s):
      }
   T solve(vd &x) {
       int r = 0:
       repx(i, 1, m) if (D[i][n + 1] < D[r][n + 1]) r = i;
       if (D[r][n + 1] < -eps) {
           pivot(r, n);
           if (!simplex(2) || D[m + 1][n + 1] < -eps) return</pre>
           rep(i, m) if (B[i] == -1) {
              int s = 0:
              repx(j, 1, n + 1) ltj(D[i]);
              pivot(i, s):
           }
       bool ok = simplex(1);
       x = vd(n);
       rep(i, m) if (B[i] < n) \times [B[i]] = D[i][n + 1];
       return ok ? D[m][n + 1] : inf;
   }
};
```

#### 5.16 theorems-and-formulas

$$n! \sim \sqrt{2\pi n} \left(\frac{n}{e}\right)^n$$

$$\sum_{i=0}^k \binom{n+i}{i} = \binom{n+k+1}{k}$$

$$\binom{n}{k} = \text{perm of } n \text{ elements with } k \text{ cycles}$$

$$\binom{n+1}{k} = n \binom{n}{k} + \binom{n}{k-1}$$

$$\binom{n}{k} = \text{partitions of an } n\text{-element set into } k \text{ parts}$$

$$\binom{n}{k} = \frac{1}{k!} \sum_{i=0}^k (-1)^i \binom{k}{i} (k-i)^n.$$

$$\binom{n+1}{k} = k \binom{n}{k} + \binom{n}{k-1}$$

Integers  $d_1 \geq \cdots \geq d'_n \geq 0$  can be the degree sequence of

a finite simple graph on n vertices 
$$\iff$$
  $d_1 + \dots + d_n$  is even and for every  $k$  in  $1 \le k \le n$  
$$\sum_{i=1}^k d_i \le k(k-1) + \sum_{i=k+1}^n \min(d_i, k)$$

 $a^n = a^{\varphi(m)+n \mod \varphi(m)} \pmod{m}$  if n > lg(m)

Misere Nim: if  $\exists a_i > 1$  then normal nim; else the condition is reversed.

Derangements: Num of permutations of n=0,1,2,... elements without fixed points is 1,0,1,2,9,44,265,1854,14833,... Recurrence:  $D_n=(n-1)(D_{n-1}+D_{n-2})=n*D_{n-1}+(-1)^n$ . Collary: number of permutations with exactly k fixed points  $nCkD_{n-k}$ 

Eulerian numbers: E(n, k) is the number of permutations with exactly k descents  $(i : \pi_i < \pi_{i+1})$ , ascents  $(\pi_i > \pi_{i+1})$  / excedances  $(\pi > i)$  / k + 1 weak excedances  $(\pi \ge i)$ .  $E_{n,k} = (k+1)E_{n-1,k} + (n-k)E_{n-1,k-1}$ 

#### 5.17 theorems

#### Burnside lemma

Tomemos imagenes x en X y operaciones (g: X -> X) en G. Si #g es la cantidad de imagenes que son puntos fijos de g, entonces la cantidad de objetos es '(sum\_{g} in G} #g) / |G|' Es requisito que G tenga la operacion identidad, que toda operacion tenga inversa y que todo par de operaciones tenga su combinacion.

#### Rational root theorem

Las raices racionales de un polinomio de orden n con coeficientes enteros A[i] son de la forma p / q, donde p y q son coprimos, p es divisor de A[0] y q es divisor de A[n]. Notar que si A[0] = 0, cero es raiz, se puede dividir el polinomio por x y aplica nuevamente el teorema.

#### Petersens theorem

Every cubic and bridgeless graph has a perfect matching.

Number of divisors for powers of 10 (0,1) (1,4) (2,12) (3,32) (4,64) (5,128) (6,240) (7,448) (8,768) (9,1344) (10,2304) (11,4032) (12,6720) (13,10752) (14,17280) (15,26880) (16,41472) (17,64512) (18,103680)

Kirchoff Theorem: Sea A la matriz de adyacencia del multi-

grafo (A[u][v] indica la cantidad de aristas entre u y v) Sea D una matriz diagonal tal que D[v][v] es igual al grado de v (considerando auto aristas y multi aristas). Sea L = A - D. Todos los cofactores de L son iguales y equivalen a la cantidad de Spanning Trees del grafo. Un cofactor (i,j) de L es la multiplicacin de  $(-1)^{i}$  + j con el determinant de la matriz al quitar la fila i y la columna j

Prufer Code: Dado un rbol con los nodos indexados: busca la hoja de menor ndice, brrala y anota el ndice del nodo al que estaba conectado. Repite el paso anterior n-2 veces. Lo anterior muestra una biyeccin entre los arreglos de tamao n-2 con elementos en [1, n] y los rboles de n nodos, por lo que hay n^{n-2} spanning trees en un grafo completo. Corolario: Si tenemos k componentes de tamaos  $s1, s2, \ldots, sk$  entonces podemos hacerlos conexos agregando k-1 aristas entre nodos de  $s1*s2*\ldots*sk*n^{k-2}$  formas

#### Combinatoria

Catalan:  $C_{n+1} = sum(C_i*C_{n-i})$  for i \in [0, n]) Catalan:  $C_n = \frac{1}{n+1}*\frac{n}{n}$  Sea  $C_n^k$  las formas de poner n+k pares de parntesis, con los primeros k parntesis abiertos (esto es, hay 2n + 2k carcteres), se tiene que  $C_n^k = \frac{2n+k-1}{(2n+k)/(n*(n+k+1))} * C_{n-1}^k$  Sea  $D_n$  el nmero de permutaciones sin puntos fijos, entoces  $D_n = \frac{(n-1)*(D_{n-1})}{n-1} + D_{n-2}$ ,  $D_n = 1$ ,  $D_n = 0$ 

#### 5.18 tonelli-shanks

```
ll legendre(ll a, ll p) {
   if (a % p == 0) return 0; if (p == 2) return 1;
   return binexp(a, (p - 1) / 2, p);
// sqrt(n) mod p (p must be a prime)
// rnd(a, b) return a random number in [a, b]
ll tonelli shanks(ll n. ll p) {
   if (n == 0) return 0;
   if (legendre(n, p) != 1) return -1; // no existe
   if (p == 2) return 1:
   ll s = builtin ctzll(p - 1):
   11 q = (p - 1LL) >> s, z = rnd(1, p - 1);
   if (s == 1) return binexp(n, (p + 1) / 4LL, p);
   while (legendre(z, p) != p - 1) z = rnd(1, p - 1);
   ll c = binexp(z, q, p), r = binexp(n, (q + 1) / 2, p):
   ll t = binexp(n, q, p), m = s;
   while (t != 1) {
      11 i = 1, ts = (t * t) \% p;
       while (ts != 1) i++, ts = (ts * ts) % p;
```

```
ll b = c;
  repx(_, 0, m - i - 1) b = (b * b) % p;
  r = r*b%p; c = b*b%p; t = t*c%p; m = i;
}
  return r;
}
```

## 6 strings

#### 6.1 aho-corasick

```
struct Vertex {
   int next[26], go[26];
   int p, link = -1, exit = -1, cnt = -1:
   vector<int> leaf:
   char pch;
   Vertex(int p=-1, char ch='$') : p(p), pch(ch) {
       rep(i, 26) next[i] = -1, go[i] = -1;
vector<Vertex> t(1):
void add(string &s. int id) {
   int. v = 0:
   for (char ch : s) {
       int c = ch - 'a':
       if (t[v].next[c] == -1) {
          t[v].next[c] = t.size():
          t.emplace_back(v, ch);
       v = t[v].next[c]:
   t[v].leaf.push_back(id);
int go(int v, char ch);
int get_link(int v) {
   if (t[v].link == -1) {
       if (v == 0 \mid | t[v].p == 0) t[v].link = 0;
       else t[v].link = go(get_link(t[v].p), t[v].pch);
   return t[v].link:
int go(int v, char ch) {
   int c = ch - a:
   if (t[v].go[c] == -1) {
       if (t[v].next[c] != -1) t[v].go[c] = t[v].next[c];
       else t[v].go[c] = v == 0 ? 0 : go(get_link(v), ch);
   }
   return t[v].go[c];
```

## 6.2 debruijn-sequence

```
// Given alphabet [0,k) constructs a cyclic string of length
// k^n that contains every length n string as substr.
vector<int> deBruijnSeq(int k, int n) { // Recursive FKM
    if (k == 1) return {0};
    vector<int> seq, aux(n+1);
    function<void(int,int)> gen = [&](int t, int p) {
        if (t > n) { // +lyndon word of len p
            if (n%p == 0) repx(i,1,p+1) seq.pb(aux[i]);
        } else {
        aux[t] = aux[t-p]; gen(t+1,p);
        while (++aux[t] < k) gen(t+1,t);
      }
    };
    gen(1,1); return seq;
}</pre>
```

#### 6.3 hash

```
const int K = 2;
struct Hash{
    const ll MOD[K] = {999727999, 1070777777};
    const ll P = 1777771;
    vector<1l> h[K], p[K];
    Hash(string &s){
        int n = s.size();
        rep(k, K){
            h[k].resize(n+1, 0);
            p[k].resize(n+1, 1);
        repx(i, 1, n+1){
            h[k][i] = (h[k][i-1]*P + s[i-1]) % MOD[k];
```

```
p[k][i] = (p[k][i-1]*P) % MOD[k];

}

}

vector<11> get(int i, int j){
    vector<11> r(K);
    rep(k, K){
        r[k] = (h[k][j] - h[k][i]*p[k][j-i]) % MOD[k];
        r[k] = (r[k] + MOD[k]) % MOD[k];
} return r;
}
```

#### 6.4 manacher

```
// odd[i]: length of longest palindrome centered at i
// even[i]: ...longest palindrome centered between i and i+1
void manacher(string &s,vector<int> &odd,vector<int> &even){ };
   string t = "$#";
   for(char c: s) t += c + string("#"):
   t += "^":
   int n = t.size();
   vector<int> p(n);
   int 1 = 1, r = 1;
   repx(i, 1, n-1) {
       p[i] = max(0, min(r - i, p[1 + (r - i)]));
       while(t[i - p[i]] == t[i + p[i]]) p[i]++;
       if(i + p[i] > r) l = i - p[i], r = i + p[i];
   repx(i, 2, n-2) {
       if(i%2) even.push_back(p[i]-1);
       else odd.push_back(p[i]-1);
```

## 6.5 palindromic-tree

```
int last:
                  // current node
   EerTree(string &s) : last(0) {
       t.emplace_back(-1); t.emplace_back(0); // root 1 & 2
       rep(i, s.size()) add(i, s); // construct tree
       for(int i = t.size()-1; i > 1; i--)
          t[t[i].link].cnt += t[i].cnt;
   void add(int i, string &s){
                                    // vangrind warning:
       int p=last, c=s[i]-'a';
                                   // i-t[p].len-1 = -1
       while(s[i-t[p].len-1] != s[i]) p = t[p].link;
       if(t[p].to[c]){ last = t[p].to[c]: t[last].cnt++: }
           int q = t[p].link;
           while(s[i-t[q].len-1] != s[i]) q = t[q].link;
          q = max(1, t[q].to[c]);
          last = t[p].to[c] = t.size();
          t.emplace_back(t[p].len + 2, q, i-t[p].len-1);
      }
void main(){
   string s = "abcbab"; EerTree pt(s); // build EerTree
   repx(i, 2, pt.t.size()){// list all distinct palindromes
       repx(j,pt.t[i].i,pt.t[i].i+pt.t[i].len)cout << s[j];</pre>
       cout << " " << pt.t[i].cnt << endl;</pre>
   }
```

## 6.6 prefix-function

```
vector<int> prefix_function(string s) {
   int n = s.size():
   vector<int> pi(n);
   repx(i, 1, n) {
       int j = pi[i-1];
       while (j > 0 \&\& s[i] != s[j])
          i = pi[i-1]:
       if (s[i] == s[i])
           j++;
       pi[i] = j;
   return pi;
vector<vector<int>> aut;
void compute automaton(string s) {
   s += '#':
   int n = s.size();
   vector<int> pi = prefix_function(s);
   aut.assign(n, vector<int>(26));
```

## 6.7 suffix-array

```
// build the suffix array
// suffixes are sorted, with each suffix represented by its
// starting position
vector<int> suffixarray(const string &s) {
   int N = s.size() + 1;//optional: include terminating NUL
   vector<int> p(N), p2(N), c(N), c2(N), cnt(256);
   rep(i, N) cnt[s[i]] += 1:
   repx(b, 1, 256) cnt[b] += cnt[b - 1];
   rep(i, N) p[--cnt[s[i]]] = i;
   repx(i, 1, N) c[p[i]] = c[p[i - 1]] + (s[p[i]] != s[p[i - 1]])
         1]]);
   for (int k = 1: k < N: k <<= 1) {</pre>
       int C = c[p[N - 1]] + 1;
       cnt.assign(C + 1, 0);
       for (int &pi : p) pi = (pi - k + N) % N;
       for (int cl : c) cnt[cl + 1] += 1;
       rep(i, C) cnt[i + 1] += cnt[i];
       rep(i, N) p2[cnt[c[p[i]]]++] = p[i];
       c2[p2[0]] = 0:
       repx(i, 1, N) c2[p2[i]] =
           c2[p2[i-1]] + (c[p2[i]] != c[p2[i-1]] ||
                          c[(p2[i] + k) \% N] != c[(p2[i - 1]
                                + k) % N1):
       swap(c, c2), swap(p, p2);
```

```
p.erase(p.begin()); // optional: erase terminating NUL
   return p;
// build the lcp
// 'lcp[i]' represents the length of the longest common
// prefix between suffix i and suffix i+1 in the suffix
//array 'p'. the last element of 'lcp' is zero by convention
vector<int> makelcp(const string &s, const vector<int> &p) {
   int N = p.size(), k = 0;
   vector<int> r(N), lcp(N);
   rep(i, N) r[p[i]] = i;
   rep(i, N) {
       if (r[i] + 1 \ge N) \{ k = 0; continue; \}
       int j = p[r[i] + 1];
       while (i + k < N \&\& j + k < N \&\& s[i + k] == s[j + k]
           ]) k += 1;
       lcp[r[i]] = k:
       if (k) k -= 1;
   return lcp;
// lexicographically compare the suffixes starting from 'i'
// and 'j', considering only up to 'K' characters.
// 'r' is the inverse suffix array, mapping suffix offsets
// to indices. requires an LCP sparse table.
int lcp_cmp(vector<int> &r, Sparse<int> &lcp, int i, int j,
    int K) {
   if (i == j) return 0;
   int ii = r[i], jj = r[j];
   int 1 = lcp.query(min(ii, jj), max(ii, jj));
   if (1 >= K) return 0;
   return ii < j; ? -1 : 1;
```

#### 6.8 suffix-automaton

```
st[p].next[c] = k:
   if(p == -1) st[k].link = 0:
   else {
       int q = st[p].next[c];
       if(st[p].len + 1 == st[q].len) st[k].link = q;
           int w = sz++; st[w].len = st[p].len + 1;
           st[w].next=st[q].next; st[w].link=st[q].link;
           for(; p!=-1 && st[p].next[c]==q; p=st[p].link)
              st[p].next[c] = w;
           st[a].link=st[k].link = w:
   }
   last = k:
\frac{1}{2} // # states <= 2n-1 && transitions <= 3n-4 (for n > 2)
// Follow link from 'last' to 0, nodes on path are terminal
// # matches = # paths from state to a terminal node
// # substrings = # paths from 0 to any node
// # substrings = sum of (len - len(link)) for all nodes
```

#### 6.9 z-function

```
// i-th element is equal to the greatest number of
// characters starting from the position i that coincide
// with the first characters of s
vector<int> z_function(string s) {
   int n = s.size();
   vector<int> z(n);
   int l = 0, r = 0;
   for(int i = 1; i < n; i++) {
      if(i < r) z[i] = min(r - i, z[i - l]);
      while(i + z[i] < n && s[z[i]] == s[i + z[i]])z[i]++;
      if(i + z[i] > r) {
            l = i;
            r = i + z[i];
      }
   }
   return z;
}
```