Team Notebook

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1 Strings

1.1 Manacher

```
// odd[i]: length of longest palindrome centered at i
// even[i]: ...longest palindrome centered between i and i+1
string t = "$#";
   for(char c: s) t += c + string("#");
   t += "^":
   int n = t.size();
   vector<int> p(n);
   int 1 = 1, r = 1;
   repx(i, 1, n-1) {
      p[i] = max(0, min(r - i, p[1 + (r - i)]));
      while(t[i - p[i]] == t[i + p[i]]) p[i]++;
      if(i + p[i] > r) l = i - p[i], r = i + p[i];
   repx(i, 2, n-2) {
      if(i%2) even.push_back(p[i]-1);
      else odd.push_back(p[i]-1);
}
```

1.2 aho-corasick

```
struct Vertex {
    int next[26], go[26];
    int p, link = -1, exit = -1, cnt = -1;
    vector<int> leaf:
    char pch:
    Vertex(int p=-1, char ch='$') : p(p), pch(ch) {
       rep(i, 26) next[i] = -1, go[i] = -1:
};
vector<Vertex> t(1):
void add(string &s, int id) {
    int v = 0:
   for (char ch : s) {
       int c = ch - a:
       if (t[v].next[c] == -1) {
           t[v].next[c] = t.size();
           t.emplace_back(v, ch);
       v = t[v].next[c];
    t[v].leaf.push_back(id);
```

```
int go(int v. char ch);
int get_link(int v) {
   if (t[v].link == -1) {
       if (v == 0 || t[v].p == 0) t[v].link = 0;
       else t[v].link = go(get_link(t[v].p), t[v].pch);
   return t[v].link;
int go(int v, char ch) {
   int c = ch - 'a';
   if (t[v].go[c] == -1) {
      if (t[v].next[c] != -1) t[v].go[c] = t[v].next[c];
       else t[v].go[c] = v == 0 ? 0 : go(get_link(v), ch);
   return t[v].go[c];
int next_match(int v){ // Optional
   if(t[v].exit == -1){
      if(t[get link(v)].leaf.size())t[v].exit=get link(v):
       else t[v].exit = v==0 ? 0 : next match(get link(v)):
   }
   return t[v].exit:
int cnt_matches(int v){ // Optional
   if(t[v].cnt == -1)
       t[v].cnt = v == 0 ? 0 : t[v].leaf.size() +
           cnt_matches(get_link(v));
   return t[v].cnt;
```

1.3 hash

```
struct Hash{
   const int K = 2;
   const ll MOD[K] = {999727999, 1070777777};
   const ll P = 1777771;
   vector<ll> h[K], p[K];
   Hash(string &s){
      int n = s.size();
      rep(k, K){
        h[k].resize(n+1, 0);
        p[k].resize(n+1, 1);
        repx(i, 1, n+1){
            h[k][i] = (h[k][i-1]*P + s[i-1]) % MOD[k];
            p[k][i] = (p[k][i-1]*P) % MOD[k];
        }
   }
   }
   vector<ll> get(int i, int j){
```

```
vector<ll> r(K);
rep(k, K){
    r[k] = (h[k][j] - h[k][i]*p[k][j-i]) % MOD[k];
    r[k] = (r[k] + MOD[k]) % MOD[k];
} return r;
}
};
```

1.4 hash2d

```
using Hash = pair<11, int>;
struct Block {
   int x0, y0, x1, y1;
struct Hash2d {
   11 HMOD:
   int W, H;
   vector<int> h;
   vector<int> p;
   Hash2d() {}
   Hash2d(const string &s, int W_, int H_, 11 HMOD_ =
        1000003931)
       : W(W_ + 1), H(H_ + 1), HMOD(HMOD_) {
       static const 11 P =
           chrono::steady_clock::now().time_since_epoch().
                count() % (1 << 29);
       p.resize(W * H);
       p[0] = 1;
       rep(i, W * H - 1) p[i + 1] = p[i] * P % HMOD;
       h.assign(W * H, 0);
       repx(y, 1, H) repx(x, 1, W) {
          ll c = (ll)s[(y - 1) * (W - 1) + x - 1] * p[y * W
                + x] % HMOD;
          h[v * W + x] = (HMOD + h[v * W + x - 1] + h[(v -
               1) * W + x] -
                         h[(y-1)*W+x-1]+c)%
   bool isout(Block s) {
       return s.x0 < 0 \mid | s.x0 >= W \mid | s.x1 < 0 \mid | s.x1 >= W
             | | s.y0 < 0 | |
             s.y0 >= H \mid \mid s.y1 < 0 \mid \mid s.y1 >= H;
   }
```

```
Hash get(Block s) {
       return {(2 * HMOD + h[s.y1 * W + s.x1] - h[s.y1 * W +
             s.x01 -
               h[s.y0 * W + s.x1] + h[s.y0 * W + s.x0]) %
                  HMOD.
              s.v0 * W + s.x0:
    bool cmp(Hash x0, Hash x1) {
       int d = x0.second - x1.second;
       ll &lo = d < 0 ? x0.first : x1.first:
       lo = lo * p[abs(d)] % HMOD:
       return x0.first == x1.first;
};
struct Hash2dM {
    int N:
    vector<Hash2d> sub:
    Hash2dM() {}
    Hash2dM(const string &s, int W, int H, const vector<11> &
        mods)
       : N(mods.size()), sub(N) {
       rep(i, N) sub[i] = Hash2d(s, W, H, mods[i]);
    bool isout(Block s) { return sub[0].isout(s): }
    vector<Hash> get(Block s) {
       vector<Hash> hs(N);
       rep(i, N) hs[i] = sub[i].get(s);
       return hs:
    bool cmp(const vector<Hash> &x0. const vector<Hash> &x1)
       rep(i, N) if (!sub[i].cmp(x0[i], x1[i])) return false | };
       return true:
    bool cmp(Block s0, Block s1) {
       rep(i, N) if (!sub[i].cmp(sub[i].get(s0), sub[i].get(
            s1))) return false;
       return true:
};
```

1.5 palindromic-tree

```
struct Node { // (*) = Optional
int len: // length of substring
int to[26]; // insertion edge for all characters a-z
int link; // maximum palindromic suffix
int i:
           // (*) start index of current Node
   int cnt; // (*) # of occurrences of this substring
   Node(int len, int link=0, int i=0, int cnt=1): len(len).
   link(link), i(i), cnt(cnt) {memset(to, 0, sizeof(to));}
struct EerTree { // Palindromic Tree
   vector<Node> t; // tree (max size of tree is n+2)
                // current node
   EerTree(string &s) : last(0) {
      t.emplace_back(-1); t.emplace_back(0); // root 1 & 2
      rep(i, s.size()) add(i, s): // construct tree
      for(int i = t.size()-1: i > 1: i--)
          t[t[i].link].cnt += t[i].cnt;
   void add(int i, string &s){
       int p=last, c=s[i]-'a';
       while(s[i-t[p].len-1] != s[i]) p = t[p].link;
      if(t[p].to[c]){ last = t[p].to[c]; t[last].cnt++; }
      elsef
          int q = t[p].link:
          while(s[i-t[q].len-1] != s[i]) q = t[q].link;
          q = max(1, t[q].to[c]);
          last = t[p].to[c] = t.size();
          t.emplace_back(t[p].len + 2, q, i-t[p].len-1);
   }
string s = "abcbab"; EerTree pt(s); // build EerTree
repx(i, 2, pt.t.size()){// list all distinct palindromes
 repx(j,pt.t[i].i,pt.t[i].i+pt.t[i].len)cout << s[j];
 cout << " " << pt.t[i].cnt << endl;</pre>
```

1.6 prefix-function

```
vector<int> prefix_function(string s) {
```

```
int n = s.size():
   vector<int> pi(n):
   repx(i, 1, n) {
       int j = pi[i-1];
       while (j > 0 \&\& s[i] != s[j])
          i = pi[i-1]:
       if (s[i] == s[i])
           j++;
       pi[i] = j;
   return pi:
vector<vector<int>> aut:
void compute automaton(string s) {
   s += '#':
   int n = s.size();
   vector<int> pi = prefix_function(s);
   aut.assign(n, vector<int>(26));
   rep(i, n) {
       rep(c, 26) {
          int j = i;
          while (j > 0 \&\& 'a' + c != s[j])
              j = pi[j-1];
           if ('a' + c == s[i])
              j++;
           aut[i][c] = j;
   }
// k = n - pi[n - 1]; if k divides n, then the string can be
// aprtitioned into blocks of length k otherwise there is no
// effective compression and the answer is n.
```

1.7 suffix-array

```
for (int &pi : p) pi = (pi - k + N) % N;
       for (int cl : c) cnt[cl + 1] += 1;
       rep(i, C) cnt[i + 1] += cnt[i];
       rep(i, N) p2[cnt[c[p[i]]]++] = p[i];
       c2[p2[0]] = 0;
       repx(i, 1, N) c2[p2[i]] =
           c2[p2[i-1]] + (c[p2[i]] != c[p2[i-1]] ||
                          c[(p2[i] + k) \% N] != c[(p2[i - 1]
                                + k) % N1):
       swap(c, c2), swap(p, p2);
   p.erase(p.begin()); // optional: erase terminating NUL
   return p;
// build the lcp
// 'lcp[i]' represents the length of the longest common
// prefix between suffix i and suffix i+1 in the suffix
//array 'p'. the last element of 'lcp' is zero by convention
vector<int> makelcp(const string &s. const vector<int> &p) {
   int N = p.size(), k = 0;
   vector<int> r(N), lcp(N);
   rep(i, N) r[p[i]] = i;
   rep(i, N) {
       if (r[i] + 1 >= N) { k = 0; continue; }
       int j = p[r[i] + 1];
       while (i + k < N \&\& j + k < N \&\& s[i + k] == s[j + k]
           1) k += 1:
       lcp[r[i]] = k:
       if (k) k -= 1;
   return lcp;
// lexicographically compare the suffixes starting from 'i'
// and 'j', considering only up to 'K' characters.
// 'r' is the inverse suffix array, mapping suffix offsets
// to indices. requires an LCP sparse table.
int lcp_cmp(vector<int> &r, Sparse<int> &lcp, int i, int j,
    int K) {
   if (i == j) return 0;
   int ii = r[i], jj = r[j];
   int 1 = lcp.query(min(ii, jj), max(ii, jj));
   if (1 >= K) return 0;
   return ii < jj ? -1 : 1;</pre>
```

1.8 suffix-automaton

```
struct SuffixAutomaton {
  vector<map<char,int>> edges;
```

```
vector<int> link, len, cnt, paths, pos:
vector<bool> terminal;
int last; // idx of the eq. class of the whole string
SuffixAutomaton(string s) :last(0) {
   edges.push_back({});
   link.push back(-1):
   len.push_back(0);
   rep(i, s.size()) {
       edges.push_back({});
       len.push_back(i+1);
       link.push_back(0);
       int r = len.size() - 1, p = last:
       while(p >= 0 && !edges[p].count(s[i])) {
          edges[p][s[i]] = r;
          p = link[p];
       if(p != -1) {
          int q = edges[p][s[i]];
          if(len[p] + 1 == len[q]) link[r] = q;
              edges.push_back(edges[q]);
              len.push_back(len[p] + 1);
              link.push_back(link[q]);
              int qq= link[q] = link[r] =len.size()-1;
              while(p >= 0 && edges[p][s[i]] == q){
                  edges[p][s[i]] = qq;
                  p = link[p];
          }
       last = r;
   } /* ----- Optional ----- */
   terminal.assign(len.size(), 0);
   for(int p = last; p > 0; p = link[p]) terminal[p]=1;
   cnt.assign(len.size(), -1); cnt_matches(0);
   //precompute # of paths (substr) starting from state
   paths.assign(len.size(), -1); cnt_paths(0);
   pos.assign(len.size(), -1); get_pos(0);
int cnt matches(int state) {
   if(cnt[state] != -1) return cnt[state]:
   int ans = terminal[state];
   for(auto edge : edges[state])
       ans += cnt_matches(edge.second);
   return cnt[state] = ans;
int cnt_paths(int state) {
   if(paths[state] != -1) return paths[state];
   int ans = state != 0:
                              // without repetitions
// int ans = state == 0 ? 0 : cnt[state]; // with rep.
```

```
for(auto edge : edges[state])
          ans += cnt_paths(edge.second);
      return paths[state] = ans;
   int get_pos(int state) { // gets first pos
       if(pos[state] != -1) return pos[state];
                                  // max->first_pos
       int ans = 0:
// int ans = terminal[state] ? 0 : 1e9;// min->last_pos
      for(auto edge : edges[state])
          ans = max(ans, get_pos(edge.second)+1); //or min
       return pos[state] = ans:
   string get_k_substring(int k) { // 0-indexed
       string ans; int state = 0;
       while(1){
           int curr = state != 0; // without repetition
       // int curr = state == 0 ? 0 : cnt[state]: // with
          if(curr > k) return ans;
          k -= curr:
          for(auto edge : edges[state]) {
              if(paths[edge.second] <= k) {</pre>
                 k -= paths[edge.second];
              } else {
                 ans += edge.first;
                 state = edge.second;
                 break:
          }
   }
};
```

1.9 z-function

```
// i-th element is equal to the greatest number of
// characters starting from the position i that coincide
// with the first characters of s
vector<int> z_function(string s) {
   int n = s.size();
   vector<int> z(n);
   int l = 0, r = 0;
   for(int i = 1; i < n; i++) {
      if(i < r) z[i] = min(r - i, z[i - 1]);
      while(i + z[i] < n && s[z[i]] == s[i + z[i]])z[i]++;
      if(i + z[i] > r) {
            l = i;
            r = i + z[i];
      }
}
```

```
return z;
}
```

2 common

$3 ext{ dp}$

3.1 convex-hull-trick

```
struct Line {
    mutable 11 a, b, c;
    bool operator<(Line r) const { return a < r.a; }</pre>
    bool operator<(11 x) const { return c < x; }</pre>
}:
// dynamically insert 'a*x + b' lines and query for maximum
// at any x all operations have complexity O(log N)
struct LineContainer : multiset<Line. less<>> {
    11 div(11 a. 11 b) {
       return a / b - ((a ^ b) < 0 && a % b);
    bool isect(iterator x, iterator v) {
       if (v == end()) return x \rightarrow c = INF. 0:
       if (x->a == y->a) x->c = x->b > y->b ? INF : -INF;
       else x->c = div(y->b - x->b, x->a - y->a);
       return x->c >= y->c;
    void add(ll a, ll b) {
       // a *= -1, b *= -1 // for min
       auto z = insert(\{a, b, 0\}), y = z++, x = y;
       while (isect(y, z)) z = erase(z);
```

3.2 divide-and-conquer

```
// for every index i assign an optimal index j, such that
    cost(i, i) is
// minimal for every i. the property that if i2 >= i1 then
    i2 >= i1 is
// exploited (monotonic condition).
// calculate optimal index for all indices in range [1, r)
    knowing that
// the optimal index for every index in this range is within
     [optl, optr).
// time: O(N log N)
void calc(vector<int> &opt, int 1, int r, int opt1, int optr
    ) {
   if (1 == r) return:
   int i = (1 + r) / 2;
   11 optc = INF;
   int opti:
   repx(j, optl, optr) {
      11 c = i + j; // cost(i, j)
      if (c < optc) optc = c, optj = j;</pre>
   opt[i] = opti:
   calc(opt, 1, i, optl, optj + 1);
   calc(opt, i + 1, r, optj, optr);
```

$4 \quad \mathbf{geo2d}$

4.1 circle

```
struct C {
```

```
Po: Tr:
C(P o, T r) : o(o), r(r) {}
C() : C(P(), T()) \{ \}
// intersects the circle with a line. assuming they
// results are sorted with respect to the direction of
     the line
pair<P, P> line_inter(L 1) const {
   P c = 1.closest to(o):
   T c2 = (c - o).magsq():
   P = sqrt(max(r * r - c2, T())) * 1.d.unit();
   return {c - e, c + e};
}
// checks whether the given line collides with the circle
// negative: 2 intersections
// zero: 1 intersection
// positive: 0 intersections
T line_collide(L 1) const {
   T c2 = (1.closest_to(o) - o).magsq();
   return c2 - r * r;
// calculates the two intersections between two circles
// the circles must intersect in one or two points!
pair<P. P> inter(C h) const {
   P d = h.o - o;
   T c = (r * r - h.r * h.r) / d.magsq();
   return h.line_inter({(1 + c) / 2 * d, d.rot()});
// check if the given circles intersect
bool collide(C h) const {
   return (h.o - o).magsq() \le (h.r + r) * (h.r + r):
// get one of the two tangents that cross through the
// the point must not be inside the circle
// a = -1: cw (relative to the circle) tangent
// a = 1: ccw (relative to the circle) tangent
P point_tangent(P p, T a) const {
   T c = r * r / p.magsq();
   return o + c * (p - o) - a * sqrt(c * (1 - c)) * (p -
         o).rot();
}
```

// get one of the 4 tangents between the two circles

```
// a = 1: exterior tangents
// a = -1: interior tangents (requires no area overlap)
// b = 1: ccw tangent
// b = -1: cw tangent
// the line origin is on this circumference, and the
    direction
// is a unit vector towards the other circle
L tangent(C c, T a, T b) const {
   T dr = a * r - c.r:
   P d = c.o - o;
   P n = (d * dr + b * d.rot() * sart(d.magsq() - dr *
        dr)).unit();
   return {o + n * r, -b * n.rot()};
// find the circumcircle of the given **non-degenerate**
    triangle
static C thru_points(P a, P b, P c) {
   L 1((a + b) / 2, (b - a).rot()):
   P p = 1.intersection(L((a + c) / 2, (c - a).rot()));
   return {p, (p - a).mag()};
// find the two circles that go through the given point,
    are tangent
// to the given line and have radius 'r'
// the point-line distance must be at most 'r'!
// the circles are sorted in the direction of the line
static pair<C, C> thru_point_line_r(P a, L t, T r) {
   P d = t.d.rot().unit():
   if (d * (a - t.o) < 0) d = -d;
   auto p = C(a, r).line_inter(\{t.o + d * r, t.d\});
   return {{p.first, r}, {p.second, r}};
// find the two circles that go through the given points
// radius 'r'
// the circles are sorted by angle with respect to the
    first point
// the points must be at most at distance 'r'!
static pair<C, C> thru_points_r(P a, P b, T r) {
   auto p = C(a, r).line inter(\{(a + b) / 2, (b - a).rot
   return {{p.first, r}, {p.second, r}};
```

};

4.2 convex-hull

```
// ccw order, excludes collinear points by default
// adding collinears duplicates points
vector<P> chull(vector<P> p) {
   if (p.size() < 3) return p;</pre>
   vector < P > r; int k = 0;
   sort(p.begin(), p.end(), [](P a, P b) {
       return a.x != b.x ? a.x < b.x : a.v < b.y; });</pre>
   for (P q : p) { // lower hull
       while (k \ge 2 \&\& r[k - 1].left(r[k - 2], a) \ge 0)
           r.pop_back(), k--; // >= to > to add collinears
       r.push back(q), k++:
   r.pop_back(), k--; int m = k;
   for (int i = p.size() - 1; i >= 0; --i) { // upper hull
       while (k \ge m+2 \&\& r[k-1].left(r[k-2], p[i]) \ge 0)
           r.pop_back(), k--; // >= to > to add collinears
       r.push back(p[i]), k++;
   r.pop_back(); return r;
```

4.3 delaunay

```
typedef __int128_t lll; // if on a 64-bit platform
struct Q {
   Q *rot, *o; P p = {INF, INF}; bool mark;
   P &F() { return r()->p; }
    Q *&r() { return rot->rot: }
   Q *prev() { return rot->o->rot: }
    Q *next() { return r()->prev(); }
T cross(P a, P b, P c) { return (b - a) % (c - a); }
bool circ(P p, P a, P b, P c) { // is p in the circumcircle?
   111 p2 = p.magsq(), A = a.magsq() - p2,
       B = b.magsq() - p2, C = c.magsq() - p2;
    return cross(p, a, b) * C + cross(p, b, c) * A + cross(p, b, c)
         c. a) * B > 0:
0 *makeEdge(0 *&H, P orig, P dest) {
   Q *r = H ? H : new Q{new Q{new Q{new Q{0}}}};
   H = r -> 0; r -> r() -> r() = r;
   repx(i, 0, 4) r = r->rot, r->p = {INF, INF},
       r \rightarrow 0 = i \& 1 ? r : r \rightarrow r();
```

```
r\rightarrow p = orig: r\rightarrow F() = dest:
   return r:
void splice(Q *a, Q *b) {
   swap(a->o->rot->o, b->o->rot->o): swap(a->o, b->o):
Q *connect(Q *&H, Q *a, Q *b) {
   Q *q = makeEdge(H, a->F(), b->p);
   splice(q, a->next()); splice(q->r(), b); return q;
pair<Q *. Q *> rec(Q *&H. const vector<P> &s) {
   if (s.size() <= 3) {</pre>
       Q *a = makeEdge(H, s[0], s[1]), *b = makeEdge(H, s[0])
            [1]. s.back()):
       if (s.size() == 2) return \{a, a->r()\}; splice(a->r(),
       auto side = cross(s[0], s[1], s[2]):
       Q *c = side ? connect(H, b, a) : 0;
       return \{\text{side} < 0 ? c > r() : a. \text{side} < 0 ? c : b > r()\}
            }:
   }
#define J(e) e \rightarrow F(), e \rightarrow p
#define valid(e) (cross(e->F(), J(base)) > 0)
   0 *A. *B. *ra. *rb: int half = s.size() / 2:
   tie(ra, A) = rec(H, {s.begin(), s.end() - half});
   tie(B, rb) = rec(H, {s.begin() + s.size() - half, s.end()
        }):
   while ((cross(B->p, J(A)) < 0 \&\& (A = A->next())) | |
          (cross(A->p, J(B)) > 0 \&\& (B = B->r()->o)));
   Q *base = connect(H, B->r(), A);
   if (A->p == ra->p) ra = base->r();
   if (B\rightarrow p == rb\rightarrow p) rb = base:
#define DEL(e, init, dir) Q *e = init->dir; \
   if (valid(e)) while (circ(e->dir->F(), J(base), e->F()))
           Q *t = e->dir: splice(e, e->prev()): \
           splice(e->r(), e->r()->prev()); e->o = H; H = e;
                e = t: \
       }
   for (;;) {
       DEL(LC, base->r(), o); DEL(RC, base, prev());
       if (!valid(LC) && !valid(RC)) break;
       if (!valid(LC) || (valid(RC) && circ(J(RC), J(LC))))
            base = connect(H, RC, base->r()):
       else base = connect(H. base->r(). LC->r()):
```

```
return {ra, rb};
#undef J
#undef valid
#undef DEL
// there must be no duplicate points
// returns no triangles in the case of all collinear points
// produces counter-clockwise triangles ordered in triples
// maximizes the minimum angle across all triangulations
// the euclidean mst is a subset of these edges
// O(N log N)
vector<P> triangulate(vector<P> pts) {
   sort(pts.begin(), pts.end(), [](P a, P b) {
       return make_pair(a.x, a.y) < make_pair(b.x, b.y);</pre>
   assert(unique(pts.begin(), pts.end()) == pts.end());
   if (pts.size() < 2) return {};</pre>
   Q *H = 0; Q *e = rec(H, pts).first;
   vector<Q *> q = {e}; int qi = 0;
   while (cross(e->o->F(), e->F(), e->p) < 0) e = e->o;
#define ADD
   ł
       0 *c = e:
           c->mark = 1; pts.push_back(c->p); \
           q.push_back(c->r()); c = c->next(); \setminus
       } while (c != e);
   ADD;
   pts.clear();
   while (qi < (int)q.size()) if (!(e = q[qi++])->mark) ADD;
   return pts;
#undef ADD
```

4.4 halfplane-intersect

```
sort(begin(H), end(H), [](L a, L b) { return a.d.angcmp(b
     .d) < 0: }):
deque<L> q; int n = 0;
rep(i, H.size()) {
   while (n \ge 2 \&\& H[i].side(q[n - 1].intersection(q[n
        -21)) > 0)
       q.pop_back(), n--;
   while (n \ge 2 \&\& H[i].side(q[0].intersection(q[1])) >
       q.pop_front(), n--;
   if (n > 0 \&\& H[i].parallel(q[n - 1])) {
       if (H[i].d * a[n - 1].d < 0) return {}:
       if (H[i].side(q[n - 1].o) > 0) q.pop_back(), n--;
       else continue:
   q.push_back(H[i]), n++;
while (n \ge 3 \&\& a[0].side(a[n - 1].intersection(a[n -
    21)) > 0)
    g.pop_back(), n--;
while (n \ge 3 \&\& q[n - 1].side(q[0].intersection(q[1])) >
   q.pop_front(), n--;
if (n < 3) return {};</pre>
vector<P> ps(n);
rep(i, n) ps[i] = q[i].intersection(q[(i + 1) % n]);
return ps;
```

4.5 line

```
// a segment or an infinite line
// does not handle point segments correctly!
struct L {
   P o, d;
   L() : o(), d() {}
   L(P o, P d) : o(o), d(d) {}

   L(P ab, T c) : d(ab.rot()), o(ab * -c / ab.magsq()) {}
   pair<P, T> line_eq() { return {-d.rot(), d.rot() * o}; }

   // returns a number indicating which side of the line the point is in
   // negative: left, positive: right
   T side(P r) const { return (r - o) % d; }

   // returns the intersection coefficient
```

```
// in the range [0, d % r.d]
// if d % r.d is zero, the lines are parallel
T inter(L r) const { return (r.o - o) % r.d; }
// get the single intersection point
// lines must not be parallel
P intersection(L r) const { return o + d * inter(r) / (d
    % r.d): }
// check if lines are parallel
bool parallel(L r) const { return abs(d % r.d) <= EPS: }</pre>
// check if segments intersect
bool seg_collide(L r) const {
   Tz = d \% r.d;
   if (abs(z) <= EPS) {
       if (abs(side(r.o)) > EPS) return false;
       T s = (r.o - o) * d, e = s + r.d * d;
       if (s > e) swap(s, e):
       return s <= d * d + EPS && e >= -EPS:
   T s = inter(r), t = -r.inter(*this);
   if (z < 0) s = -s, t = -t, z = -z;
   return s >= -EPS && s <= z + EPS && t >= -EPS && t <=
         z + EPS:
}
// full segment intersection
// produces a point segment if the intersection is a
// however it **does not** handle point segments as input
bool seg_inter(L r, L *out) const {
   T z = d \% r.d;
   if (abs(z) <= EPS) {
       if (abs(side(r.o)) > EPS) return false;
       if (r.d * d < 0) r = \{r.o + r.d, -r.d\};
       P s = o * d < r.o * d ? r.o : o;
       P = (o + d) * d < (r.o + r.d) * d ? o + d : r.o
             + r.d:
       if (s * d > e * d) return false;
       return *out = {s, e - s}, true;
   }
   T s = inter(r), t = -r.inter(*this);
   if (z < 0) s = -s, t = -t, z = -z;
   if (s >= -EPS && s <= z + EPS && t >= -EPS && t <= z
       return *out = {o + d * s / z, {0, 0}}, true:
   return false:
```

```
// check if the given point is on the segment
bool point_on_seg(P r) const {
    if (abs(side(r)) > EPS) return false;
    if ((r - o) * d < -EPS) return false;
    if ((r - o - d) * d > EPS) return false;
    return true;
}

// get the point in this line that is closest to a given point
P closest_to(P r) const {
    P dr = d.rot(); return r + (o - r) * dr * dr / d.
        magsq();
}
```

4.6 minkowski

```
void reorder polygon(vector<P> &ps) {
   int pos = 0:
   repx(i, 1, (int)ps.size()) {
       if (ps[i].y < ps[pos].y || (ps[i].y == ps[pos].y &&</pre>
           ps[i].x < ps[pos].x)
          pos = i:
   rotate(ps.begin(), ps.begin() + pos, ps.end());
vector<P> minkowski(vector<P> ps. vector<P> qs) {
   // the first vertex must be the lowest
   reorder_polygon(ps); reorder_polygon(qs);
   ps.push_back(ps[0]); ps.push_back(ps[1]);
   qs.push_back(qs[0]); qs.push_back(qs[1]);
   vector<P> result: int i = 0, i = 0:
   while (i < ps.size() - 2 || j < qs.size() - 2) {
       result.push_back(ps[i] + qs[j]);
       auto z = (ps[i + 1] - ps[i]) \% (qs[j + 1] - qs[j]);
       if (z \ge 0 \&\& i < ps.size() - 2) ++i;
       if (z <= 0 && j < qs.size() - 2) ++j;</pre>
   }
   return result;
```

4.7 point

```
struct P {
   T x, y;
```

```
P(T x, T y) : x(x), y(y) {}
   P() : P(0, 0) \{ \}
   friend ostream &operator<<(ostream &s, const P &r) {</pre>
       return s << r.x << " " << r.y;
   friend istream &operator>>(istream &s, P &r) { return s
        >> r.x >> r.v: }
   P operator+(P r) const { return \{x + r.x, y + r.y\}; }
   P operator-(P r) const { return {x - r.x, y - r.y}; }
   P operator*(T r) const { return {x * r, v * r}: }
   P operator/(T r) const { return {x / r, y / r}; }
   P operator-() const { return {-x, -y}; }
   friend P operator*(T 1, P r) { return {1 * r.x, 1 * r.y};
   P rot() const { return {-y, x}; }
   T operator*(P r) const { return x * r.x + v * r.v: }
   T operator%(P r) const { return rot() * r: }
   T left(P a, P b) { return (b - a) % (*this - a); }
   T magsq() const { return x * x + y * y; }
   T mag() const { return sqrt(magsq()); }
   P unit() const { return *this / mag(); }
   bool half() const { return abs(y) <= EPS && x < -EPS || y
         < -EPS: }
   T angcmp(P r) const { // like strcmp(this, r)
       int h = (int)half() - r.half();
       return h ? h : r % *this;
   T angcmp_rel(P a, P b) { // like strcmp(a, b)
      Pz = *this;
       int h = z \% a \le 0 & z * a \le 0 | | z \% a \le 0:
      h = z \% b \le 0 \&\& z * b < 0 || z \% b < 0:
       return h ? h : b % a:
   }
   bool operator==(P r) const { return abs(x - r.x) <= EPS</pre>
        && abs(v - r.v) <= EPS: }
   double angle() const { return atan2(v, x): }
   static P from_angle(double a) { return {cos(a), sin(a)};
}:
```

4.8 polygon

```
// get TWICE the area of a simple polygon in ccw order
T area2(const vector<P> &p) {
   int n = p.size(): T a = 0:
   rep(i, n) a += (p[i] - p[0]) % (p[(i + 1) % n] - p[i]);
   return a:
// checks whether a point is inside a ccw simple polygon
// returns 1 if inside, 0 if on border, -1 if outside
int in_poly(const vector<P> &p, P q) {
   int. w = 0:
   rep(i, p.size()) {
       P = p[i], b = p[(i + 1) \% p.size()];
       T k = (b - a) \% (q - a);
       T u = a.y - q.y, v = b.y - q.y;
       if (k > 0 \&\& u < 0 \&\& v >= 0) w++:
       if (k < 0 \&\& v < 0 \&\& u >= 0) w--;
       if (k == 0 && (q - a) * (q - b) <= 0) return 0;
   return w ? 1 : -1;
// check if point in ccw convex polygon, O(log n)
// + if inside, 0 if on border, - if outside
T in_convex(const vector<P> &p, P q) {
   int 1 = 1, h = p.size() - 2; assert(p.size() >= 3);
   while (1 != h) { // collinear points are unsupported!
       int m = (1 + h + 1) / 2;
       if (a.left(p[0], p[m]) >= 0) 1 = m:
       else h = m - 1:
   T in = min(q.left(p[0], p[1]), q.left(p.back(), p[0]));
   return min(in, q.left(p[1], p[1 + 1]));
int extremal(const vector<P> &p, P d) {
   int n = p.size(), l = 0, r = n - 1; assert(n);
   P = 0 = (p[n - 1] - p[0]).rot():
   while (1 < r) { // polygon must be convex
       int m = (1 + r + 1) / 2:
       P = (p[(m + n - 1) \% n] - p[m]).rot();
       if (e0.angcmp_rel(d, e) < 0) r = m - 1:
       else 1 = m:
   return 1:
// square dist of most distant points of a ccw convex
// polygon with NO COLLINEAR POINTS
T callipers(const vector<P> &p) {
```

```
int n = p.size();
   T r = 0:
   for (int i = 0, j = n < 2 ? 0 : 1; <math>i < j; i++) {
       for (;; j = (j + 1) % n) {
          r = max(r, (p[i] - p[j]).magsq());
           if ((p[(i + 1) % n] - p[i]) % (p[(j + 1) % n] - p
                [i]) <= EPS) break;</pre>
   return r;
P centroid(const vector<P> &p) { // (barycenter)
   P r(0, 0); T t = 0; int n = p.size();
   rep(i, n) {
       r += (p[i] + p[(i+1)\%n]) * (p[i] \% p[(i+1)\%n]);
       t += p[i] \% p[(i+1)\%n];
   return r / t / 3:
}
// classify collision of a ray inside a ccw polygon vertex.
// ray is (o, d), vertex is b, previous vertex is a, next is
pair<bool, bool> inner_collide(P o, P d, P a, P b, P c) {
   T p = (a - o) \% d; // side of previous
   T n = (c - o) \% d:
                         // side of next
   T v = (c - b) \% (b - a); // is vertex convex?
   return {v > 0 ? n < 0 || (n == 0 && p < 0) : p > 0 || n < //
          v > 0 ? p > 0 || (p == 0 && n > 0) : p > 0 || n <
```

4.9 sweep

```
#include "point.cpp"

// iterate over all pairs of points

// 'op' is called with all ordered pairs of different
    indices '(i, j)'

// additionally, the 'ps' vector is kept sorted by signed
    distance

// to the line formed by 'i' and 'j'

// for example, if the vector from 'i' to 'j' is pointing
    right,

// the 'ps' vector is sorted from smallest 'y' to largest 'y
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```

```
// note that, because the 'ps' vector is sorted by signed
// 'j' is always equal to 'i + 1'
// this means that the amount of points to the left of the
    line is always 'N - i'
template <class OP>
void all_pair_points(vector<P> &ps, OP op) {
   int N = ps.size():
   sort(ps.begin(), ps.end(), [](P a, P b) {
       return make_pair(a.v, a.x) < make_pair(b.v, b.x);</pre>
   vector<pair<int, int>> ss:
   rep(i, N) rep(j, N) if (i != j) ss.push_back({i, j});
   stable_sort(ss.begin(), ss.end(), [&](auto a, auto b) {
       return (ps[a.second] - ps[a.first]).angle_lt(ps[b.
            second] - ps[b.first]);
   }):
   vector<int> p(N); rep(i, N) p[i] = i;
   for (auto [i, i] : ss)
       { op(p[i], p[j]); swap(ps[p[i]], ps[p[j]]); swap(p[i
           ], p[i]); }
```

4.10 theorems

```
// Pick's theorem

// Simple polygon with integer vertices:

// A = I + B / 2 - 1

// A: Area of the polygon

// I: Integer points strictly inside the polygon

// B: Integer points on the boundary of the polygon
```

5 graph

5.1 bellman-ford

```
rep(i, N - 1) {
    bool f = true;
    rep(ei, E.size()) {
        auto &e = E[ei];
        ll n = D[e.u] + e.w;
        if (D[e.u] < INF && n < D[e.v])
            D[e.v] = n, P[e.v] = ei, f = false;
    }
    if (f) return false;
}
return true;
}</pre>
```

5.2 blossom

vector<int> g[MAXN];int n,m,mt[MAXN],qh,qt,q[MAXN],ft[MAXN], bs[MAXN];bool ing[MAXN],inb[MAXN],inp[MAXN];int lca(int root ,int x,int y){memset(inp,0,sizeof(inp));while(1){inp[x=bs[x]]=true;if(x==root)break;x=ft[mt[x]];}while(1){if(inp[y=bs[y])}])return v:else v=ft[mt[v]]:}}void mark(int z.int x){while(bs[x]!=z){int y=mt[x];inb[bs[x]]=inb[bs[y]]=true;x=ft[y];if(bs[x]!=z)ft[x]=y;}}void contr(int s,int x,int y){int z=lca(s ,x,y);memset(inb,0,sizeof(inb));mark(z,x);mark(z,y);if(bs[x] $!=z)ft[x]=y;if(bs[y]!=z)ft[y]=x;rep(x,n)if(inb[bs[x]]){bs[x]}$ =z;if(!ing[x])ing[q[++qt]=x]=true;}}int findp(int s){memset(inq,0,sizeof(inq));memset(ft,-1,sizeof(ft));rep(i,n)bs[i]=i; inq[q[qh=qt=0]=s]=true; while(qh<=qt){int x=q[qh++]; for(int y</pre> g[x] = g[x] if bs[x] = bs[y] & mt[x] = y {if $y = s \mid mt[y] > 0 & ft[mt[y] > 0$ $1 \ge 0$ contr(s.x.v):else if(ft[v]<0){ft[v]=x:if(mt[v]<0)return y;else if(!inq[mt[y]])inq[q[++qt]=mt[y]]=true;}}}return -1;} int aug(int s,int t){int x=t,y,z;while(x>=0){y=ft[x];z=mt[y]} ;mt[y]=x;mt[x]=y;x=z;}return t>=0;}int edmonds(){int r=0; memset(mt,-1,sizeof(mt));rep(x,n)if(mt[x]<0)r+=aug(x,findp(x</pre>)):return r:}

5.3 dinic

```
// time: 0(E V^2)
// 0(E V^(2/3)) / 0(E sqrt(E)) unit capacities
// 0(E sqrt(V)) (hopcroft-karp) unit networks
//unit network: c in {0,1} & forall v, indeg<=1 or outdeg<=1
//min-cut: nodes reachable from s in final residual graph
struct Dinic {
   struct Edge { int u, v; ll c, f = 0; };
   int N, s, t; vector<vector<int>> G;
   vector<Edge> E; vector<int> lvl, ptr;
   Dinic() {}
```

```
Dinic(int N, int s, int t): N(N), s(s), t(t), G(N) {}
    void add_edge(int u, int v, ll c) {
       G[u].push_back(E.size()); E.push_back({u, v, c});
       G[v].push_back(E.size()); E.push_back({v, u, 0});
    11 push(int u, ll p) {
       if (u == t || p <= 0) return p;</pre>
       while (ptr[u] < G[u].size()) {</pre>
           int ei = G[u][ptr[u]++];
           Edge &e = E[ei]:
           if (lvl[e.v] != lvl[u] + 1) continue;
           11 a = push(e.v, min(e.c - e.f, p));
           if (a <= 0) continue:
           e.f += a, E[ei ^ 1].f -= a; return a;
       }
       return 0;
    11 maxflow() {
       11 f = 0:
       while (true) {
           lvl.assign(N, -1); queue<int> q;
           lvl[s] = 0; q.push(s);
           while (!q.empty()) {
               int u = q.front(); q.pop();
               for (int ei : G[u]) {
                   Edge &e = E[ei];
                   if (e.c-e.f<=0||lvl[e.v]!=-1) continue;</pre>
                   lvl[e.v] = lvl[u] + 1; q.push(e.v);
           }
           if (lvl[t] == -1) break;
           ptr.assign(N,0); while(ll ff=push(s,INF))f += ff;
       }
       return f:
   }
};
/* Fluio con demandas (no necesariamente el maximo)
Agregar s' y t' nuevos source and sink
c'(s', v) = sum(d(u, v) \text{ for } u \text{ in } V) \setminus sum(sim v)
c'(v, t') = sum(d(v, w) \text{ for w in V}) \setminus forall arista (v, t')
c'(u, v) = c(u, v) - d(u, v) \forall aristas antiguas
c'(t, s) = INF (el flujo por esta arista es el flujo real)*/
```

5.4 floyd-warshall

5.5 heavy-light

```
struct Hld {
   vector<int> P, H, D, pos, top;
   Hld() {}
   void init(vector<vector<int>> &G) {
       int N = G.size();
       P.resize(N), H.resize(N), D.resize(N), pos.resize(N),
           top.resize(N):
       D[0] = -1, dfs(G, 0); int t = 0;
       rep(i, N) if (H[P[i]] != i) {
          int j = i;
           while (j != -1)
              \{ top[j] = i, pos[j] = t++; j = H[j]; \}
      }
   }
   int dfs(vector<vector<int>> &G, int i) {
       int w = 1, mw = 0:
       D[i] = D[P[i]] + 1, H[i] = -1;
       for (int c : G[i]) {
           if (c == P[i]) continue:
          P[c] = i; int sw = dfs(G, c); w += sw;
           if (sw > mw) H[i] = c. mw = sw:
      7
       return w;
   // visit the log N segments in the path from \boldsymbol{u} to \boldsymbol{v}
   template <class OP>
   void path(int u, int v, OP op) {
       while (top[u] != top[v]) {
          if (D[top[u]] > D[top[v]]) swap(u, v);
           op(pos[top[v]], pos[v] + 1); v = P[top[v]];
```

```
if (D[u] > D[v]) swap(u, v);
   op(pos[u], pos[v] + 1); // value on node
   // op(pos[u]+1, pos[v] + 1); // value on edge
// an alternative to 'path' that considers order.
// calls 'op' with an 'l <= r' inclusive-exclusive range,</pre>
// boolean indicating if the query is forwards or
     backwards.
template <class OP>
void path(int u, int v, OP op) {
   int lu = u. lv = v:
   while (top[lu] != top[lv])
       if (D[top[lu]] > D[top[lv]]) lu = P[top[lu]];
       else lv = P[top[lv]];
   int lca = D[lu] > D[lv] ? lv : lu;
   while (top[u] != top[lca])
       op(pos[top[u]], pos[u] + 1, false), u = P[top[u]]
   if (u != lca) op(pos[lca] + 1, pos[u] + 1, false);
   vector<int> stk:
   while (top[v] != top[lca])
       stk.push_back(v), v = P[top[v]];
   // op(pos[lca], pos[v] + 1, true); // value on node
   op(pos[lca] + 1, pos[v] + 1, true); // value on edge
   reverse(stk.begin(), stk.end());
   for (int w : stk) op(pos[top[w]], pos[w] + 1, true);
// commutative segment tree
template <class T. class S>
void update(S &seg, int i, T val) { seg.update(pos[i],
    val): }
// commutative segment tree lazy
template <class T. class S>
void update(S &seg, int u, int v, T val) {
   path(u, v, [&](int 1, int r) { seg.update(1, r, val);
         });
// commutative (lazy) segment tree
template <class T, class S>
T query(S &seg, int u, int v) {
```

5.6 hungarian

```
// find a maximum gain perfect matching in the given
    bipartite complete graph.
// input: gain matrix (G_{xy} = benefit of joining vertex x
    in set X with vertex
// v in set Y).
// output: maximum gain matching in members 'xy[x]' and 'yx[
// runtime: O(N^3)
struct Hungarian {
   int N, qi, root;
   vector<vector<ll>> gain;
   vector<int> xy, yx, p, q, slackx;
   vector<ll> lx, ly, slack;
   vector<bool> S, T;
   void add(int x, int px) {
       S[x] = true, p[x] = px;
       rep(y, N) if (lx[x] + ly[y] - gain[x][y] < slack[y])
           slack[y] = lx[x] + ly[y] - gain[x][y], slackx[y]
   void augment(int x, int y) {
       while (x != -2) {
          yx[y] = x; swap(xy[x], y); x = p[x];
   }
   void improve() {
       S.assign(N, false), T.assign(N, false), p.assign(N,
            -1);
       qi = 0, q.clear();
       rep(x, N) if (xy[x] == -1) {
          q.push_back(root = x), p[x] = -2, S[x] = true;
           break:
       }
```

```
rep(y, N) slack[y] = lx[root] + ly[y] - gain[root][y
        ], slackx[v] = root;
   while (true) {
       while (qi < q.size()) {</pre>
           int x = a[ai++]:
           rep(y, N) if (lx[x] + ly[y] == gain[x][y] &&!
               T[v]) {
              if (yx[y] == -1) return augment(x, y);
              T[y] = true, q.push_back(yx[y]), add(yx[y
                   1. x):
          }
       }
       11 d = INF:
       rep(y, N) if (!T[y]) d = min(d, slack[y]);
       rep(x, N) if (S[x]) lx[x] -= d;
       rep(y, N) if (T[y]) ly[y] += d;
       rep(y, N) if (!T[y]) slack[y] -= d;
       rep(y, N) if (!T[y] && slack[y] == 0) {
           if (yx[y] == -1) return augment(slackx[y], y);
          T[v] = true;
           if (!S[yx[y]]) q.push_back(yx[y]), add(yx[y],
                slackx[y]);
   }
}
Hungarian(vector<vector<11>>> g)
   : N(g.size()), gain(g), xy(N, -1), yx(N, -1), lx(N, -
   ly(N), slack(N), slackx(N) {
   rep(x, N) rep(y, N) lx[x] = max(lx[x], ly[y]);
   rep(i, N) improve();
}
```

5.7 kuhn

};

```
// get a maximum cardinality matching in a bipartite graph.
// input: adjacency lists.
// output: matching (in 'mt' member).
// runtime: O(V E)
struct Kuhn {
   int N, size;
   vector<vector<int>> G;
   vector<bool> seen;
   vector<int> mt;
```

```
bool visit(int i) {
       if (seen[i]) return false;
       seen[i] = true:
       for (int to : G[i])
          if (mt[to] == -1 || visit(mt[to])) {
              mt[to] = i;
              return true;
       return false;
   Kuhn(vector<vector<int>> adj) : G(adj), N(G.size()), mt(N
        , -1) {
       rep(i, N) {
          seen.assign(N, false);
          size += visit(i):
      }
   }
};
```

5.8 lca

```
// calculates the lowest common ancestor for any two nodes
    in O(log N) time,
// with O(N log N) preprocessing
struct Lca {
   int N, K, t = 0;
   vector<vector<int>> U;
   vector<int> L, R;
   Lca(vector<vector<int>> &G) : N(G.size()), L(N), R(N) {
      K = N \le 1 ? 0 : 32 - builtin clz(N - 1):
      U.resize(K + 1, vector<int>(N));
       visit(G, 0, 0);
       rep(k, K) rep(u, N) U[k + 1][u] = U[k][U[k][u]];
   }
   void visit(vector<vector<int>> &G, int u, int p) {
      L[u] = t++, U[0][u] = p;
      for (int v : G[u]) if (v != p) visit(G, v, u);
      R[u] = t++;
   bool is_anc(int up, int dn) {
       return L[up] <= L[dn] && R[dn] <= R[up];</pre>
```

```
int find(int u, int v) {
    if (is_anc(u, v)) return u;
    if (is_anc(v, u)) return v;
    for (int k = K; k >= 0;)
        if (is_anc(U[k][u], v)) k--;
        else u = U[k][u];
    return U[0][u];
}
```

5.9 maxflow-mincost

```
// time: O(F V E)
                         F is the maximum flow
       O(V E + F E log V) if bellman-ford is replaced by
    johnson
struct Flow {
   struct Edge {
      int u, v;
      11 c, w, f = 0;
   int N, s, t;
   vector<vector<int>> G;
   vector<Edge> E:
   vector<ll> d, b;
   vector<int> p;
   Flow() {}
   Flow(int N, int s, int t) : N(N), s(s), t(t), G(N) {}
   void add_edge(int u, int v, ll c, ll w) {
      G[u].push back(E.size()):
      E.push_back({u, v, c, w});
      G[v].push back(E.size());
      E.push_back({v, u, 0, -w});
   // naive distances with bellman-ford: O(V E)
   void calcdists() {
      p.assign(N, -1), d.assign(N, INF), d[s] = 0;
      rep(i, N - 1) rep(ei, E.size()) {
          Edge &e = E[ei]:
          ll n = d[e.u] + e.w;
          if (d[e.u] < INF && e.c - e.f > 0 && n < d[e.v])
               d[e.v] = n, p[e.v] = ei:
   // johnsons potentials: O(E log V)
```

```
void calcdists() {
       if (b.empty()) {
           b.assign(N, 0);
           // code below only necessary if there are
               negative costs
           rep(i, N - 1) rep(ei, E.size()) {
              Edge &e = E[ei];
              if (e.f < e.c) b[e.v] = min(b[e.v], b[e.u] + e
           }
       p.assign(N. -1), d.assign(N. INF), d[s] = 0:
       priority_queue<pair<11, int>> q;
       q.push({0, s});
       while (!q.empty()) {
           auto [w, u] = q.top();
           q.pop();
           if (d[u] < -w + b[u]) continue;</pre>
           for (int ei : G[u]) {
              auto e = E[ei]:
              ll n = d[u] + e.w;
              if (e.f < e.c && n < d[e.v]) {</pre>
                  d[e.v] = n, p[e.v] = ei;
                  q.push({b[e.v] - n, e.v});
           }
       }
       b = d:
   }
   ll solve() {
       b.clear();
       11 ff = 0:
       while (true) {
           calcdists():
           if (p[t] == -1) break:
           11 f = INF:
           for (int cur = t; p[cur] != -1; cur = E[p[cur]].u
              f = min(f, E[p[cur]].c - E[p[cur]].f);
           for (int cur = t; p[cur] != -1; cur = E[p[cur]].u
              E[p[cur]].f += f, E[p[cur] ^ 1].f -= f;
           ff += f;
       return ff;
   }
};
```

5.10 push-relabel

```
#include "../common.h"
const 11 INF = 1e18;
// maximum flow algorithm.
// to run, use 'maxflow()'.
11
// time: O(V^2 \operatorname{sqrt}(E)) \leq O(V^3)
// memory: 0(V^2)
struct PushRelabel {
   vector<vector<11>> cap, flow;
   vector<ll> excess:
   vector<int> height;
   PushRelabel() {}
   void resize(int N) { cap.assign(N, vector<11>(N)); }
   // push as much excess flow as possible from u to v.
   void push(int u, int v) {
       11 f = min(excess[u], cap[u][v] - flow[u][v]);
       flow[u][v] += f:
       flow[v][u] -= f:
       excess[v] += f:
       excess[u] -= f;
   // relabel the height of a vertex so that excess flow may
         be pushed.
   void relabel(int u) {
       int d = INT32 MAX:
       rep(v, cap.size()) if (cap[u][v] - flow[u][v] > 0) d
           min(d, height[v]):
       if (d < INF) height[u] = d + 1;</pre>
   // get the maximum flow on the network specified by 'cap'
         with source 's'
   // and sink 't.'.
   // node-to-node flows are output to the 'flow' member.
   11 maxflow(int s, int t) {
       int N = cap.size(), M;
       flow.assign(N, vector<ll>(N));
       height.assign(N, 0), height[s] = N:
       excess.assign(N, 0), excess[s] = INF;
       rep(i, N) if (i != s) push(s, i);
       vector<int> q;
```

```
while (true) {
           // find the highest vertices with excess
           q.clear(), M = 0;
           rep(i, N) {
              if (excess[i] <= 0 || i == s || i == t)</pre>
                    continue:
              if (height[i] > M) q.clear(), M = height[i];
              if (height[i] >= M) q.push_back(i);
           if (q.empty()) break;
           // process vertices
           for (int u : a) {
              bool relab = true;
              rep(v. N) {
                  if (excess[u] <= 0) break;</pre>
                  if (cap[u][v] - flow[u][v] > 0 && height[u]
                       1 > height[v])
                      push(u, v), relab = false;
              if (relab) {
                  relabel(u);
                  break:
       }
       11 f = 0; rep(i, N) f += flow[i][t]; return f;
};
```

5.11 strongly-connected-components

```
// compute strongly connected components.
// time: O(V + E), memory: O(V)
//
//
//
after building:
// comp = map from vertex to component (components are toposorted, root first, leaf last)
// N = number of components
// G = condensation graph (component DAG)
//
// byproducts:
// vgi = transposed graph
// order = reverse topological sort (leaf first, root last)
//
// others:
// vn = number of vertices
// vg = original vertex graph
struct Scc {
```

```
int vn. N:
   vector<int> order, comp;
   vector<vector<int>> vg, vgi, G;
   void toposort(int u) {
       if (comp[u]) return:
       comp[u] = -1;
       for (int v : vg[u]) toposort(v);
       order.push_back(u);
   bool carve(int u) {
       if (comp[u] != -1) return false;
       comp[u] = N;
       for (int v : vgi[u]) {
          carve(v);
          if (comp[v] != N) G[comp[v]].push_back(N);
       return true:
   }
   Scc() {}
   Scc(vector<vector<int>> &g) : vn(g.size()), vg(g), comp(
        vn), vgi(vn), G(vn), N(0) {
       rep(u, vn) toposort(u);
       rep(u, vn) for (int v : vg[u]) vgi[v].push_back(u);
       invrep(i, vn) N += carve(order[i]);
   }
};
```

5.12 two-sat

6 implementation

6.1 SegmentTreeBeats

```
struct Node {
   11 \text{ s, mx1, mx2, mxc, mn1, mn2, mnc, } 1z = 0;
   Node(): s(0). mx1(LLONG MIN). mx2(LLONG MIN). mxc(0).
        mn1(LLONG_MAX), mn2(LLONG_MAX), mnc(0) {}
   Node(ll x): s(x), mx1(x), mx2(LLONG_MIN), mxc(1), mn1(x)
        , mn2(LLONG_MAX), mnc(1) {}
   Node(const Node &a. const Node &b) {
       // add
       s = a.s + b.s;
       // min
       if (a.mx1 > b.mx1) mx1 = a.mx1, mxc = a.mxc, mx2 =
            max(b.mx1, a.mx2);
       if (a.mx1 < b.mx1) mx1 = b.mx1. mxc = b.mxc. mx2 =
            max(a.mx1, b.mx2);
       if (a.mx1 == b.mx1) mx1 = a.mx1, mxc = a.mxc + b.mxc,
             mx2 = max(a.mx2, b.mx2):
       // max
       if (a.mn1 < b.mn1) mn1 = a.mn1, mnc = a.mnc, mn2 =
            min(b.mn1, a.mn2):
       if (a.mn1 > b.mn1) mn1 = b.mn1, mnc = b.mnc, mn2 =
            min(a.mn1, b.mn2):
       if (a.mn1 == b.mn1) mn1 = a.mn1, mnc = a.mnc + b.mnc,
             mn2 = min(a.mn2, b.mn2);
   }
};
// 0 - indexed / inclusive - inclusive
template <class node>
```

```
struct STB {
   vector<node> st; int n;
   void build(int u, int i, int i, vector<node> &arr) {
      if (i == i) {
          st[u] = arr[i]:
          return;
      int m = (i + j) / 2, l = u * 2 + 1, r = u * 2 + 2;
      build(l, i, m, arr), build(r, m + 1, j, arr);
      st[u] = node(st[1], st[r]);
   void push_add(int u, int i, int j, ll v) {
      st[u].s += (j - i + 1) * v;
      st[u].mx1 += v, st[u].mn1 += v, st[u].lz += v;
      if (st[u].mx2 != LLONG_MIN) st[u].mx2 += v;
      if (st[u].mn2 != LLONG MAX) st[u].mn2 += v:
   void push max(int u, 11 v, bool 1) { // for min op
      if (v >= st[u].mx1) return;
      st[u].s -= st[u].mx1 * st[u].mxc;
      st[u].mx1 = v:
      st[u].s += st[u].mx1 * st[u].mxc;
      if (1) st[u].mn1 = st[u].mx1:
       else if (v <= st[u].mn1) st[u].mn1 = v;</pre>
       else if (v < st[u].mn2) st[u].mn2 = v;
   void push min(int u, 11 v, bool 1) { // for max op
      if (v <= st[u].mn1) return;</pre>
      st[u].s -= st[u].mn1 * st[u].mnc:
      st[u].mn1 = v;
      st[u].s += st[u].mn1 * st[u].mnc;
      if (1) st[u].mx1 = st[u].mn1:
      else if (v \ge st[u].mx1) st[u].mx1 = v;
       else if (v > st[u].mx2) st[u].mx2 = v:
   void push(int u, int i, int j) {
      if (i == j) return;
      // add
      int m = (i + j) / 2, l = u * 2 + 1, r = u * 2 + 2;
      push add(l, i, m, st[u].lz):
      push_add(r, m + 1, j, st[u].lz);
      st[u].lz = 0:
      // min
       push_max(1, st[u].mx1, i == m);
      push_max(r, st[u].mx1, m + 1 == j);
      // max
      push_min(1, st[u].mn1, i == m);
       push min(r, st[u].mn1, m + 1 == r):
```

```
node guerv(int a, int b, int u, int i, int i) {
   if (b < i | | i < a) return node():
   if (a <= i && j <= b) return st[u];</pre>
   push(u, i, i):
   int m = (i + j) / 2, l = u * 2 + 1, r = u * 2 + 2;
   return node(querv(a, b, l, i, m), querv(a, b, r, m +
}
void update_add(int a, int b, ll v, int u, int i, int j)
   if (b < i | | i < a) return:
   if (a <= i && i <= b) {
       push_add(u, i, j, v);
       return:
   push(u, i, j);
   int m = (i + j) / 2, l = u * 2 + 1, r = u * 2 + 2;
   update_add(a, b, v, l, i, m);
   update add(a, b, v, r, m + 1, i):
   st[u] = node(st[1], st[r]);
void update min(int a, int b, ll v, int u, int i, int i)
   if (b < i || j < a || v >= st[u].mx1) return;
   if (a <= i && j <= b && v > st[u].mx2) {
       push_max(u, v, i == j);
       return:
   push(u, i, j);
   int m = (i + j) / 2, l = u * 2 + 1, r = u * 2 + 2;
   update_min(a, b, v, l, i, m);
   update_min(a, b, v, r, m + 1, j);
   st[u] = node(st[l], st[r]):
void update_max(int a, int b, ll v, int u, int i, int j)
   if (b < i || j < a || v <= st[u].mn1) return;</pre>
   if (a <= i && j <= b && v < st[u].mn2) {</pre>
       push_min(u, v, i == j);
       return:
   push(u, i, j);
   int m = (i + j) / 2, l = u * 2 + 1, r = u * 2 + 2;
   update_max(a, b, v, l, i, m);
   update_max(a, b, v, r, m + 1, j);
   st[u] = node(st[1], st[r]);
STB(vector<node> &v, int N) : n(N), st(N * 4 + 5) { build
     (0, 0, n - 1, v);
```

6.2 Treap

```
mt19937 gen(chrono::high_resolution_clock::now().
    time_since_epoch().count());
// 101 Implicit Treap //
struct Node {
   int p, sz = 0, v, acc, l = -1, r = -1;
   Node() : v(0), acc(0) {}
   Node(int x) : p(gen()), sz(1), v(x), acc(x) {}
   void recalc(const Node &a, const Node &b) {
       sz = a.sz + b.sz + 1:
       acc = v + a.acc + b.acc:
   }
};
template <class node>
struct Treap {
   vector<node> t;
   int n. r = -1:
   node get(int u) { return u != -1 ? t[u] : node(); }
   void recalc(int u) { t[u].recalc(get(t[u].l), get(t[u].r)
        ): }
   int merge(int 1, int r) {
       if (min(l, r) == -1) return l != -1 ? l : r:
       int ans = (t[1].p < t[r].p) ? 1 : r;
       if (ans == 1) t[1].r = merge(t[1].r, r), recalc(1);
       if (ans == r) t[r].l = merge(l, t[r].l), recalc(r):
       return ans:
   pii split(int u, int id) {
       if (u == -1) return {-1, -1};
       int szl = get(t[u].1).sz:
       if (szl >= id) {
          pii ans = split(t[u].1, id);
          t[u].1 = ans.ss:
           recalc(u);
```

```
return {ans.ff. u}:
       pii ans = split(t[u].r, id - szl - 1);
       t[u].r = ans.ff:
       recalc(u);
       return {u. ans.ss}:
    Treap(vi &v) : n(sz(v)) {
       for (int i = 0; i < n; i++) t.eb(v[i]), r = merge(r, left)
            i):
};
// Complete Implicit Treap with Lazy propagation //
struct Node {
    int p, sz = 0, v, acc, l = -1, r = -1, par = -1, lzv = 0;
    bool lz = false, f = false:
    Node(): v(0), acc(0) {}
    Node(int x): p(gen()), sz(1), v(x), acc(x) {}
    void recalc(const Node &a. const Node &b) {
       sz = a.sz + b.sz + 1:
       acc = v + a.acc + b.acc:
    void upd_lazy(int x) { lz = 1, lzv += x; }
    void lazy() { v += lzv, acc += sz * lzv, lz = 0, lzv = 0;
    void flip() { swap(1, r), f = 0; }
}:
template <class node>
struct Treap {
    vector<node> t;
    int n. r = -1:
    node get(int u) { return u != -1 ? t[u] : node(); }
    void recalc(int u) {
       int 1 = t[u].1, r = t[u].r;
       push(1), push(r), flip(1), flip(r);
       t[u].recalc(get(1), get(r)):
    void push(int u) {
       if (u == -1 || !t[u].lz) return;
       int 1 = t[u].1, r = t[u].r;
       if (1 != -1) t[1].upd lazv(t[u].lzv):
       if (r != -1) t[r].upd_lazy(t[u].lzv);
       t[u].lazv():
    void flip(int u) {
```

```
if (u == -1 || !t[u].f) return;
   int 1 = t[u].1. r = t[u].r:
   if (1 != -1) t[1].f ^= 1;
   if (r != -1) t[r].f ^= 1:
   t[u].flip();
int merge(int 1, int r) {
   if (min(1, r) == -1) return 1 != -1 ? 1 : r;
   push(1), push(r), flip(1), flip(r);
   int ans = (t[1].p < t[r].p) ? 1 : r;
   if (ans == 1) t[1].r = merge(t[1].r, r), recalc(1);
   if (ans == r) t[r].l = merge(l, t[r].l), recalc(r):
   if (t[ans].l != -1) t[t[ans].l].par = ans; // only if
         parent needed
   if (t[ans].r != -1) t[t[ans].r].par = ans; // only if
         parent needed
   return ans:
pii split(int u, int id) {
   if (u == -1) return {-1, -1}:
   push(u);
   flip(u):
   int szl = get(t[u].1).sz;
   if (szl >= id) {
       pii ans = split(t[u].1, id);
       if (ans.ss != -1) t[ans.ss].par = u; // only if
           parent needed
       if (ans.ff != -1) t[ans.ff].par = -1; // only if
            parent needed
       t[u].1 = ans.ss:
       recalc(u):
       return {ans.ff, u};
   pii ans = split(t[u].r, id - szl - 1);
   if (ans.ff != -1) t[ans.ff].par = u; // only if
        parent needed
   if (ans.ss != -1) t[ans.ss].par = -1; // only if
        parent needed
   t[u].r = ans.ff;
   recalc(u):
   return {u. ans.ss}:
int update(int u, int 1, int r, int v) {
   pii a = split(u, l), b = split(a.ss, r - l + 1);
   t[b.ff].upd_lazy(v);
   return merge(a.ff, merge(b.ff, b.ss));
void print(int u) {
   if (u == -1) return:
   push(u), flip(u);
```

```
print(t[u].1);
    cout << t[u].v << ' ';
    print(t[u].r);
}

Treap(vi &v) : n(sz(v)) {
    for (int i = 0; i < n; i++) t.eb(v[i]), r = merge(r, i);
}
};</pre>
```

6.3 bit-tricks

```
y = x & (x-1) // Turn off rightmost 1bit
y = x & (-x) // Isolate rightmost 1bit
y = x \mid (x-1) // Right propagate rightmost 1bit(fill in 1s)
y = x \mid (x+1) // Turn on rightmost Obit
v = ~x & (x+1) // Isolate rightmost Obit
// If x is of long type, use __builtin_popcountl(x)
// If x is of long long type, use __builtin_popcountll(x)
// 1. Counts the number of ones(set bits) in an integer.
__builtin_popcount(x)
// 2. Checks the Parity of a number. Returns true(1) if the
// number has odd number of set bits, else it returns
// false(0) for even number of set bits.
__builtin_parity(x)
// 3. Counts the leading number of zeros of the integer.
__builtin_clz(x)
// 4. Counts the trailing number of zeros of the integer.
__builtin_ctz(x)
// 5. Returns 1 + the index of the least significant 1-bit.
__builtin_ffs(x) // If x == 0, returns 0.
// Iterate over non empty subsets of bitmask
for(int s=m:s:s=(s-1)&m) // Decreasing order
for (int s=0;s=s-m&m;) // Increasing order
```

6.4 dsu

```
struct Dsu {
  vector<int> p; Dsu() {} Dsu(int N) : p(N, -1) {}
  int get(int x) { return p[x] < 0 ? x : get(p[x]); }
  bool sameSet(int a, int b) { return get(a) == get(b); }
  int size(int x) { return -p[get(x)]; }
  vector<vector<int>> S;
  void unite(int x, int y) {
    if ((x = get(x)) == (y = get(y))) { S.push_back({-1})}
      ; return; }
```

```
if (p[x] > p[y]) swap(x, y);
    S.push_back({x, y, p[x], p[y]});
    p[x] += p[y], p[y] = x;
}
void rollback() {
    auto a = S.back(); S.pop_back();
    if (a[0] != -1) p[a[0]] = a[2], p[a[1]] = a[3];
}
};
```

6.5 dynamic-connectivity

```
struct DC {
   int n; Dsu D;
   vector<vector<pair<int, int>>> t;
   DC(int N) : n(N), D(N), t(2 * N) {}
   // add edge p to all times in interval [1, r]
   void upd(int 1, int r, pair<int, int> p) {
       for (1 += n, r += n; 1 < r; 1 >>= 1, r >>= 1) {
          if (1 & 1) t[1++].push back(p):
           if (r & 1) t[--r].push back(p):
    void process(int u = 1) { // process all queries
       for (auto &e : t[u]) D.unite(e.first. e.second):
       if (n >= n) {
           // do stuff with D at time u - n
       } else process(2 * u), process(2 * u + 1);
       for (auto &e : t[u]) D.rollback():
};
```

6.6 hash-container

```
namespace{//add (#define tmpl template)(#define ty typename)
  tmpl<ty T> size_t mk_h(const T& v){return hash<T>()(v);}
  void h_cmb(size_t& h, const size_t& v)
  { h ^= v + 0x9e3779b9 + (h << 6) + (h >> 2); }
  tmpl<ty T> struct h_ct{size_t operator()(const T& v)const{
  size_t h=0;for(const auto& e:v){h_cmb(h,mk_h(e));}return h;
  }};
}namespace std{//support for pair<T,U>, vector<T> & map<T,U>
  tmpl<ty T, ty U> struct hash<pair<T, U>>{
    size_t operator()(const pair<T,U>& v) const
  {size_t h=mk_h(v.first);h_cmb(h, mk_h(v.second));return h;}
  };
tmpl<ty... T>struct hash<vector<T...>:h_ct<vector<T...>>{};
```

```
tmpl<ty... T>struct hash<map<T...>>:h_ct<map<T...>>{}; }
```

6.7 mo

```
struct Query { int 1, r, idx; };
// answer segment queries using only 'add(i)', 'remove(i)'
    and 'get()'
// functions.
// complexity: O((N + Q) * sqrt(N) * F)
// N = length of the full segment
// Q = amount of gueries
// F = complexity of the 'add', 'remove' functions
template <class A, class R, class G, class T>
void mo(vector<Query> &queries, vector<T> &ans, A add, R
    remove, G get) {
   int Q = queries.size(), B = (int)sqrt(Q);
   sort(queries.begin(), queries.end(), [&](Query &a, Query
       return make_pair(a.1 / B, a.r) < make_pair(b.1 / B, b</pre>
            .r):
   });
   ans.resize(0):
   int 1 = 0, r = 0;
   for (auto &q : queries) {
       while (r < q.r) add(r), r++;
       while (1 > q.1) 1--, add(1);
       while (r > q.r) r--, remove(r);
       while (1 < q.1) remove(1), 1++;</pre>
       ans[q.idx] = get();
   }
```

6.8 ordered-set

6.9 persistent-segment-tree-lazy

```
template <class T>
struct Node {
   T x, lz;
   int 1 = -1, r = -1:
template <class T>
struct Pstl {
   int N:
   vector<Node<T>> a;
   vector<int> head;
   T qneut() { return 0; }
   T merge(T 1, T r) { return 1 + r: }
   T uneut() { return 0; }
   T accum(T u, T x) { return u + x; }
   T apply(T x, T lz, int l, int r) { return x + (r - 1) *
        lz; }
   int build(int vl. int vr) {
       if (vr - vl == 1) a.push_back({qneut(), uneut()}); //
            node construction
          int vm = (vl + vr) / 2, l = build(vl, vm), r =
               build(vm, vr);
          a.push_back({merge(a[1].x, a[r].x), uneut(), 1, r
               }); // query merge
       return a.size() - 1;
   T query(int 1, int r, int v, int v1, int vr, T acc) {
       if (1 >= vr || r <= vl) return gneut();</pre>
           // query neutral
       if (1 <= v1 && r >= vr) return apply(a[v].x, acc, v1,
            vr): // update op
       acc = accum(acc, a[v].lz);
           // update merge
       int vm = (vl + vr) / 2:
       return merge(query(1, r, a[v].1, v1, vm, acc), query(
           1, r, a[v].r, vm, vr, acc)); // query merge
   }
   int update(int 1, int r, T x, int v, int v1, int vr) {
      if (1 >= vr || r <= vl || r <= 1) return v;
       a.push_back(a[v]);
       v = a.size() - 1:
       if (1 <= v1 && r >= vr) {
```

```
a[v].x = apply(a[v].x, x, vl, vr); // update op
          a[v].lz = accum(a[v].lz. x): // update merge
      } else {
          int vm = (vl + vr) / 2:
          a[v].1 = update(1, r, x, a[v].1, v1, vm);
          a[v].r = update(1, r, x, a[v].r, vm, vr);
          a[v].x = merge(a[a[v].1].x, a[a[v].r].x); //
               query merge
      }
       return v;
   Pst1() {}
   Pstl(int N) : N(N) { head.push_back(build(0, N)); }
   T query(int t, int 1, int r) {
       return query(1, r, head[t], 0, N, uneut()); // update
            neutral
   int update(int t, int 1, int r, T x) {
       return head.push_back(update(1, r, x, head[t], 0, N))
            . head.size() - 1:
};
```

6.10 persistent-segment-tree

```
// usage:
// Pst<Node<11>> pst;
// pst = {N};
// int newtime = pst.update(time, index, value);
// Node<ll> result = pst.querv(newtime, left, right);
template <class T>
struct Node {
   T x;
    int 1 = -1, r = -1:
    Node(): x(0) {}
    Node(T x) : x(x) \{ \}
    Node (Node a, Node b, int l = -1, int r = -1) : x(a.x + b.
        x), 1(1), r(r) {}
};
template <class U>
struct Pst {
    int N:
    vector<U> a:
    vector<int> head;
```

```
int build(int vl. int vr) {
   if (vr - vl == 1) a.push_back(U());
       int vm = (vl + vr) / 2, l = build(vl, vm),
          r = build(vm, vr):
       a.push_back(U(a[1], a[r], 1, r));
   return a.size() - 1:
}
U querv(int 1, int r, int v, int v1, int vr) {
   if (1 >= vr || r <= vl) return U();</pre>
   if (1 <= v1 && r >= vr) return a[v]:
   int vm = (vl + vr) / 2:
   return U(query(1, r, a[v].1, v1, vm),
           query(1, r, a[v].r, vm, vr));
}
int update(int i, U x, int v, int vl, int vr) {
   a.push_back(a[v]);
   v = a.size() - 1:
   if (vr - vl == 1) a[v] = x;
   else {
       int vm = (vl + vr) / 2:
       if (i < vm) a[v].l = update(i, x, a[v].l, vl, vm)</pre>
       else a[v].r = update(i, x, a[v].r, vm, vr);
       a[v] = U(a[a[v].1], a[a[v].r], a[v].1, a[v].r);
   }
   return v;
}
Pst(int N) : N(N) { head.push_back(build(0, N)); }
U query(int t, int 1, int r) {
   return query(1, r, head[t], 0, N);
int update(int t, int i, U x) {
   return head.push back(update(i, x, head[t], 0, N)).
        head.size() - 1;
```

6.11 segment-tree-lazy

```
template <class T>
struct Stl {
```

```
int n:
   vector<T> a, b;
   T qneut() { return -2e9; }
   T uneut() { return 0; }
   T merge(T x, T y) { return max(x, y); }
   void upd(int v, T x, int 1, int r)
      \{ a[v] += x, b[v] += x; \}
   Stl(int n = 0) : n(n), a(4 * n, qneut()),
       b(4 * n. uneut()) {}
   void push(int v, int vl, int vm, int vr) {
       upd(2 * v, b[v], vl, vm);
       upd(2 * v + 1, b[v], vm, vr);
      b[v] = uneut();
   T querv(int 1, int r, int v=1, int v1=0, int vr=1e9) {
       vr = min(vr, n):
       if (1 <= v1 && r >= vr) return a[v];
       if (1 >= vr || r <= vl) return qneut();</pre>
       int vm = (vl + vr) / 2;
       push(v, v1, vm, vr);
       return merge(query(1, r, 2 * v, v1, vm),
          query(1, r, 2 * v + 1, vm, vr));
   void update(int 1, int r, T x, int v = 1, int vl = 0,
          int vr = 1e9) {
       vr = min(vr, n):
       if (1 >= vr || r <= vl || r <= l) return;</pre>
       if (1 <= v1 && r >= vr) upd(v, x, v1, vr);
       else {
          int vm = (vl + vr) / 2:
          push(v. vl. vm. vr):
          update(1, r, x, 2 * v, v1, vm);
          update(1, r, x, 2 * v + 1, vm, vr);
           a[v] = merge(a[2 * v], a[2 * v + 1]);
};
```

6.12 segment-tree

```
struct St {
    ll neut() { return 0; }
    ll merge(ll x, ll y) { return x + y; }
```

```
int n; vector<ll> a;
St(int n = 0) : n(n), a(2 * n, neut()) {}

ll query(int l, int r) {
    ll x = neut(), y = neut();
    for (1 += n, r += n; 1 < r; 1 /= 2, r /= 2) {
        if (1 & 1) x = merge(x, a[1++]);
        if (r & 1) y = merge(a[--r], y);
    }
    return merge(x, y);
}

void update(int i, ll x) {
    for (a[i += n] = x; i /= 2;)
        a[i] = merge(a[2 * i], a[2 * i + 1]);
};</pre>
```

6.13 sparse-table

```
template <class T>
struct Sparse {
   T op(T a, T b) { return max(a, b); }
   vector<vector<T>> st;
   Sparse() {}
   Sparse(vector<T> a) : st{a} {
       int N = st[0].size():
       int npot = N <= 1 ? 1 : 32 - __builtin_clz(N);</pre>
       st.resize(npot);
       repx(i, 1, npot) rep(j, N + 1 - (1 << i))
       st[i].push_back(
           op(st[i-1][j], st[i-1][j+(1 << (i-1))])
       ); // query op
   T query(int 1, int r) { // range must be nonempty!
       int i = 31 - __builtin_clz(r - 1);
       return op(st[i][l], st[i][r - (1 << i)]); // queryop</pre>
};
```

6.14 unordered-map

7 imprimible

8 math

8.1 arithmetic

```
inline int floor_log2(int n) { return n <= 1 ? 0 : 31 -</pre>
     __builtin_clz(n); }
inline int ceil_log2(int n) { return n <= 1 ? 0 : 32 -</pre>
     __builtin_clz(n - 1); }
inline ll floordiv(ll a, ll b) {
    return a / b - ((a \hat{b}) < 0 & a \% b):
inline 11 ceildiv(11 a, 11 b) {
    return a / b + ((a ^ b) >= 0 && a % b);
ll binexp(ll a, ll e) {
    ll res = 1: // neutral element
    while (e) {
       if (e & 1) res = res * a; // multiplication
       a = a * a:
                                // multiplication
       e >>= 1;
    }
    return res:
```

8.2 crt

```
pair<11, 11> solve_crt(const vector<pair<11, 11>> &eqs) {
    11 a0 = eqs[0].first, p0 = eqs[0].second;
```

```
repx(i, 1, eqs.size()) {
    ll a1 = eqs[i].first, p1 = eqs[i].second;
    ll k1, k0;
    ll d = ext_gcd(p1, p0, k1, k0);
    a0 -= a1;
    if (a0 % d != 0) return {-1, -1};
    p0 = p0 / d * p1;
    a0 = a0 / d * k1 % p0 * p1 % p0 + a1;
    a0 = (a0 % p0 + p0) % p0;
}
return {a0, p0};
```

8.3 discrete-log

```
// discrete logarithm log_a(b).
// solve b \hat{x} = a \pmod{M} for the smallest x.
// returns -1 if no solution is found.
11
// time: O(sqrt(M))
11 dlog(ll a, ll b, ll M) {
   11 k = 1, s = 0;
   while (true) {
       11 g = \_gcd(b, M);
       if (g <= 1) break;</pre>
       if (a == k) return s:
       if (a % g != 0) return -1;
       a \neq g, M \neq g, s += 1, k = b \neq g * k % M;
   ll N = sqrt(M) + 1;
   umap<11, 11> r;
   rep(q, N + 1) {
       r[a] = q;
       a = a * b % M;
   ll bN = binexp(b, N, M), bNp = k;
   repx(p, 1, N + 1) {
       bNp = bNp * bN % M;
       if (r.count(bNp)) return N * p - r[bNp] + s:
   return -1;
```

8.4 fft

```
using cd = complex<double>;
const double PI = acos(-1);
// compute the DFT of a power-of-two-length sequence.
// if 'inv' is true, computes the inverse DFT.
void fft(vector<cd> &a. bool inv) {
   int N = a.size(), k = 0, b;
   assert(N == 1 << builtin ctz(N)):</pre>
   repx(i, 1, N) {
       for (b = N >> 1: k & b:) k ^= b, b >>= 1:
       if (i < (k ^= b)) swap(a[i], a[k]);</pre>
   for (int 1 = 2; 1 <= N; 1 <<= 1) {</pre>
       double ang = 2 * PI / 1 * (inv ? -1 : 1):
       cd wl(cos(ang), sin(ang));
       for (int i = 0; i < N; i += 1) {</pre>
           cd w = 1:
           rep(j, 1 / 2) {
              cd u = a[i + j], v = a[i + j + 1 / 2] * w;
              a[i + i] = u + v:
              a[i + j + 1 / 2] = u - v;
              w *= wl:
       }
   if (inv) rep(i, N) a[i] /= N:
const 11 MOD = 998244353, ROOT = 15311432:
// const ll MOD = 2130706433, ROOT = 1791270792;
// const 11 MOD = 922337203673733529711. ROOT =
     532077456549635698311:
void find root of unitv(ll M) {
   11 c = M - 1, k = 0:
   while (c \% 2 == 0) c /= 2, k += 1;
   // find proper divisors of M - 1
   vector<ll> divs:
   for (ll d = 1: d < c: d++) {
       if (d * d > c) break;
       if (c \% d == 0) rep(i, k + 1) divs.push back(d << i):
   rep(i, k) divs.push_back(c << i);</pre>
   // find any primitive root of M
   11 G = -1;
```

```
repx(g, 2, M) {
      bool ok = true:
      for (ll d : divs) ok &= (binexp(g, d, M) != 1);
      if (ok) {
          G = g;
          break:
      }
   assert(G != -1);
   11 w = binexp(G, c, M):
   cerr << "M = c * 2^k + 1" << endl:
   cerr << " M = " << M << endl;
   cerr << " c = " << c << endl:
   cerr << " k = " << k << endl:
   cerr << " w^(2^k) == 1" << endl;
   cerr << " w = g^{(M-1)/2^k} = g^c << endl:
   cerr << " g = " << G << endl;
   cerr << " w = " << w << endl:
// compute the DFT of a power-of-two-length sequence, modulo
     a special prime
// number with an Nth root of unity, where N is the length
    of the sequence.
void ntt(vector<ll> &a, bool inv) {
   vector<ll> wn:
   for (11 p = ROOT: p != 1: p = p * p % MOD) wn.push back(p
   int N = a.size(), k = 0, b;
   assert(N == 1 << __builtin_ctz(N) && N <= 1 << wn.size()) // 1 -> unique solution, stored in 'ans'
   rep(i, N) a[i] = (a[i] \% MOD + MOD) \% MOD;
   repx(i, 1, N) {
      for (b = N >> 1; k & b;) k ^= b, b >>= 1;
      if (i < (k ^= b)) swap(a[i], a[k]);</pre>
   for (int 1 = 2: 1 <= N: 1 <<= 1) {
      11 wl = wn[wn.size() - __builtin_ctz(1)];
      if (inv) wl = multinv(wl, MOD);
      for (int i = 0; i < N; i += 1) {</pre>
          11 w = 1:
          repx(i, 0, 1 / 2) {
             11 u = a[i + j], v = a[i + j + 1 / 2] * w %
              a[i + j] = (u + v) \% MOD;
```

```
a[i + i + 1 / 2] = (u - v + MOD) \% MOD:
             w = w * w1 % MOD:
      }
   }
   11 q = multinv(N, MOD);
   if (inv) rep(i, N) a[i] = a[i] * q % MOD;
void convolve(vector<cd> &a. vector<cd> b. int n) {
   n = 1 \ll (32 - builtin clz(2 * n - 1)):
   a.resize(n), b.resize(n);
   fft(a, false), fft(b, false);
   rep(i, n) a[i] *= b[i];
   fft(a, true);
```

8.5 gauss

```
const double EPS = 1e-9:
// solve a system of equations.
// complexity: O(min(N, M) * N * M)
// 'a' is a list of rows
// the last value in each row is the result of the equation
// return values:
// 0 -> no solutions
// -1 -> infinitely many solutions, one of which is stored
// UNTESTED
int gauss(vector<vector<double>> a, vector<double> &ans) {
   int N = a.size(). M = a[0].size() - 1:
   vector<int> where(M, -1):
   for (int j = 0, i = 0; j < M && i < N; <math>j++) {
       int sel = i:
       repx(k, i, N) if (abs(a[k][j]) > abs(a[sel][j])) sel
       if (abs(a[sel][j]) < EPS) continue;</pre>
       repx(k, j, M + 1) swap(a[sel][k], a[i][k]);
       where[i] = i;
       rep(k, N) if (k != i) {
           double c = a[k][j] / a[i][j];
           repx(1, j, M + 1) a[k][1] -= a[i][1] * c;
```

8.6 matrix

```
typedef vector<vector<double>> Mat:
Mat matmul(Mat 1. Mat r) {
   int n = 1.N, m = r.M, p = 1.M; assert(1.M == r.N);
   Mat a(n, vector<double>(m)); // neutral
   rep(i, n) rep(j, m)
       rep(k, p) a[i][j] = a[i][j] + l[i][k] * r[k][j];
   return a:
}
double reduce(vector<vector<double>> &A) {
   int n = A.size(), m = A[0].size();
   int i = 0, j = 0; double r = 1.;
   while (i < n && j < m) {</pre>
       int 1 = i:
       repx(k, i+1, n) if(abs(A[k][j]) > abs(A[l][j])) l=k;
       if (abs(A[1][j]) < EPS) { j++; r = 0.; continue; }</pre>
       if (1 != i) { r = -r; swap(A[i], A[1]); }
       for (int k = m - 1; k \ge j; k--) A[i][k] /= A[i][j];
       repx(k, 0, n) {
           if (k == i) continue;
           for(int l=m-1;l>=j;l--)A[k][l]-=A[k][j]*A[i][l];
       }
       i++, j++;
   return r; // returns determinant
```

8.7 mobius

```
short mu[MAXN] = {0,1};
void mobius(){
  repx(i,1,MAXN)if(mu[i])for(int j=i+i;j<MAXN;j+=i)mu[j]-=mu[
         i];
}</pre>
```

8.8 mod

```
11 binexp(ll a, ll e, ll M) {
   assert(e >= 0);
   ll res = 1 % M;
    while (e) {
       if (e & 1) res = res * a % M;
       a = a * a % M:
       e >>= 1:
   }
   return res:
11 multinv(11 a, 11 M) { return binexp(a, M - 2, M): }
// calculate gcd(a, b).
// also, calculate x and v such that:
// a * x + b * y == gcd(a, b)
//
// time: O(log min(a, b))
// (ignoring complexity of arithmetic)
ll ext_gcd(ll a, ll b, ll &x, ll &y) {
   if (b == 0) {
       x = 1, y = 0;
       return a:
   11 d = ext gcd(b, a % b, v, x):
   y = a / b * x;
   return d:
// compute inverse with any M.
// a and M must be coprime for inverse to exist!
11 multinv_euc(ll a, ll M) {
   11 x, y;
   ext_gcd(a, M, x, y);
   return x;
// multiply two big numbers (~10^18) under a large modulo,
    without resorting to
// bigints.
```

```
ll bigmul(ll x. ll v. ll M) {
   11 z = 0:
   while (y) {
      if (v \& 1) z = (z + x) \% M:
      x = (x << 1) \% M, y >>= 1;
   return z;
// all modular inverses from 1 to inv.size()-1
void multinv all(vector<ll> &inv) {
   inv[1] = 1:
   repx(i, 2, inv.size())
       inv[i] = MOD - (MOD / i) * inv[MOD % i] % MOD:
// change to __int128 if checking numbers over 10^9
bool isprime(ll n) {
   if (n < 2 | | n % 6 % 4 != 1) return n - 2 < 2:
   11 A[] = \{2.325.9375.28178.450775.9780504.1795265022\}:
   ll s = \_builtin\_ctzll(n - 1), d = n >> s;
   for (int a : A) {
      11 p = binexp(a, d, n), i = s;
       while (p != 1 && p != n - 1 && a % n && i--) p = p *
       if (p != n - 1 && i != s) return 0;
   return 1:
struct Mod {
   int a;
   static const int M = 1e9 + 7:
   Mod(11 aa) : a((aa % M + M) % M) {}
   Mod operator+(Mod rhs) const { return (a + rhs.a) % M; }
   Mod operator-(Mod rhs) const { return (a - rhs.a + M) % M
   Mod operator-() const { return Mod(0) - *this; }
   Mod operator*(Mod rhs) const { return (11)a * rhs.a % M;
   Mod operator+=(Mod rhs) { return *this = *this + rhs; }
   Mod operator == (Mod rhs) { return *this = *this - rhs; }
   Mod operator*=(Mod rhs) { return *this = *this * rhs: }
   Mod bigmul(ll big) const { return ::bigmul(a, big, M); }
   Mod binexp(ll e) const { return ::binexp(a, e, M); }
```

```
// Mod multinv() const { return ::multinv(a, M): } //
    Mod multinv() const { return ::multinv_euc(a, M); } //
        possibly composite M
};
// dynamic modulus
struct DMod {
    int a, M;
    DMod(11 aa. 11 m) : M(m). a((aa % m + m) % m) {}
    DMod operator+(DMod rhs) const { return {(a + rhs.a) % M,
    DMod operator-(DMod rhs) const { return {(a - rhs.a + M)
        % M, M}; }
    DMod operator-() const { return DMod(0, M) - *this; }
    DMod operator*(DMod rhs) const { return {(11)a * rhs.a %
        M. M:: }
    DMod operator+=(DMod rhs) { return *this = *this + rhs; } void divisors(11 x, OP op) {
    DMod operator==(DMod rhs) { return *this = *this - rhs; }
    DMod operator*=(DMod rhs) { return *this = *this * rhs; }
    DMod bigmul(ll big) const { return {::bigmul(a, big, M),
        M}; }
    DMod binexp(ll e) const { return {::binexp(a, e, M), M}:
    DMod multinv() const { return {::multinv(a, M), M}; } //
        prime M
    // DMod multinv() const { return {::multinv_euc(a, M), M
        }; } // possibly composite M
```

8.9 primes

```
// counts the divisors of a positive integer in O(sqrt(n))
11 count_divisors(11 x) {
    11 divs = 1, i = 2;
    for (11 divs = 1, i = 2; x > 1; i++) {
        if (i * i > x) {
            divs *= 2;
            break;
        }
        for (11 d = divs; x % i == 0; x /= i) divs += d;
    }
    return divs;
}
```

```
// gets the prime factorization of a number in O(\operatorname{sqrt}(n))
vector<pair<11, int>> factorize(11 x) {
    vector<pair<11. int>> f:
    for (ll k = 2; x > 1; k++) {
       if (k * k > x) {
           f.push_back({x, 1});
           break;
       int n = 0;
       while (x \% k == 0) x /= k, n++:
       if (n > 0) f.push back(\{k, n\}):
    return f:
// iterate over all divisors of a number.
// divisor count upper bound: n^(1.07 / ln ln n)
template <class OP>
    auto facts = factorize(x):
    vector<int> f(facts.size());
    while (true) {
       11 y = 1;
       rep(i, f.size()) rep(j, f[i]) y *= facts[i].first;
       op(y);
       for (i = 0: i < f.size(): i++) {</pre>
           f[i] += 1:
           if (f[i] <= facts[i].second) break;</pre>
           f[i] = 0:
       if (i == f.size()) break;
    }
// computes euler totative function phi(x), counting the
     amount of integers in
// [1. x] that are coprime with x.
11
// time: O(sart(x))
11 phi(11 x) {
   11 phi = 1, k = 2;
    for (: x > 1: k++) {
       if (k * k > x) {
           phi *= x - 1:
           break;
       }
```

```
ll k1 = 1, k0 = 0;
while (x % k == 0) x /= k, k0 = k1, k1 *= k;
    phi *= k1 - k0;
}
return phi;
}
// isprime is in mod.cpp
```

8.10 simplex

```
// Solves a general linear maximization problem: maximize $c
    ^T x$ subject to $Ax \le b$, $x \ge 0$.
// Returns -inf if there is no solution, inf if there are
    arbitrarily good solutions, or the maximum value of $c^
    T x$ otherwise.
// The input vector is set to an optimal $x$ (or in the
    unbounded case, an arbitrary solution fulfilling the
    constraints).
// Numerical stability is not guaranteed. For better
    performance, define variables such that x = 0 is
// Usage:
// vvd A = \{\{1,-1\}, \{-1,1\}, \{-1,-2\}\}\}:
// vd b = \{1,1,-4\}, c = \{-1,-1\}, x;
// T val = LPSolver(A, b, c).solve(x);
// Time: O(NM * \t ), where a pivot may be e.g. an edge
     relaxation. O(2^n) in the general case.
#include "../common.h"
typedef double T: // long double, Rational, double + mod<P
typedef vector<T> vd:
typedef vector<vd> vvd;
const T eps = 1e-8, inf = 1 / .0:
#define MP make_pair
#define ltj(X) \
   if (s == -1 \mid | MP(X[i], N[i]) < MP(X[s], N[s])) s = i
struct LPSolver {
   int m, n;
   vector<int> N, B;
   vvd D:
   LPSolver(const vvd &A, const vd &b, const vd &c) : m(b.
        size()), n(c.size()), N(n + 1), B(m), D(m + 2), vd(n)
        + 2)) {
```

```
rep(i, m) rep(i, n) D[i][i] = A[i][i]:
   rep(i, m) {
       B[i] = n + i;
      D[i][n] = -1:
       D[i][n + 1] = b[i];
   rep(j, n) {
       N[i] = i;
       D[m][i] = -c[i];
   N[n] = -1:
   D[m + 1][n] = 1:
void pivot(int r, int s) {
   T *a = D[r].data(), inv = 1 / a[s];
   rep(i, m + 2) if (i != r && abs(D[i][s]) > eps) {
      T *b = D[i].data(), inv2 = b[s] * inv;
       repx(i, 0, n + 2) b[i] -= a[i] * inv2:
      b[s] = a[s] * inv2:
   rep(j, n + 2) if (j != s) D[r][j] *= inv;
   rep(i, m + 2) if (i != r) D[i][s] *= -inv;
   D[r][s] = inv:
   swap(B[r], N[s]);
bool simplex(int phase) {
   int x = m + phase - 1;
   for (::) {
       int s = -1;
       rep(j, n + 1) if (N[j] != -phase) ltj(D[x]);
      if (D[x][s] >= -eps) return true:
       int r = -1;
       rep(i, m) {
          if (D[i][s] <= eps) continue;</pre>
          if (r == -1 || MP(D[i][n + 1] / D[i][s], B[i])
                < MP(D[r][n + 1] / D[r][s], B[r])) r = i
       if (r == -1) return false:
       pivot(r, s);
}
T solve(vd &x) {
   int r = 0;
   repx(i, 1, m) if (D[i][n + 1] < D[r][n + 1]) r = i;
   if (D[r][n + 1] < -eps) {
       pivot(r. n):
```

8.11 theorems

Burnside lemma

Tomemos imagenes x en X y operaciones (g: X -> X) en G. Si #g es la cantidad de imagenes que son puntos fijos de g, entonces la cantidad de objetos es '(sum_{g in G} #g) / |G|' Es requisito que G tenga la operacion identidad, que toda operacion tenga inversa y que todo par de operaciones tenga su combinacion.

Rational root theorem

Las raices racionales de un polinomio de orden n con coeficientes enteros A[i] son de la forma p / q, donde p y q son coprimos, p es divisor de A[0] y q es divisor de A[n]. Notar que si A[0] = 0, cero es raiz, se puede dividir el polinomio por x y aplica nuevamente el teorema.

Petersens theorem

Every cubic and bridgeless graph has a perfect matching.

Number of divisors for powers of 10 (0,1) (1,4) (2,12) (3,32) (4,64) (5,128) (6,240) (7,448) (8,768) (9,1344) (10,2304) (11,4032) (12,6720) (13,10752) (14,17280) (15,26880) (16,41472) (17,64512) (18,103680)

Kirchoff Theorem: Sea A la matriz de adyacencia del multigrafo (A[u][v] indica la cantidad de aristas entre u y v) Sea D una matriz diagonal tal que D[v][v] es igual al grado de v (considerando auto aristas y multi aristas). Sea L = A - D. Todos los cofactores de L son iguales y equivalen a la cantidad de Spanning Trees del grafo. Un cofactor (i,j) de L es la multiplicacin de $(-1)^{i}$ + j con el determinant de la matriz al quitar la fila i y la columna j

Prufer Code: Dado un rbol con los nodos indexados: busca la hoja de menor ndice, brrala y anota el ndice del nodo al que estaba conectado. Repite el paso anterior n-2 veces. Lo anterior muestra una biyeccin entre los arreglos de tamao n-2 con elementos en [1, n] y los rboles de n nodos, por lo que hav n^{n-2} spanning trees en un grafo completo. Corolario: Si tenemos k componentes de tamaos s1,s2,...,sk entonces podemos hacerlos conexos agregando k-1 aristas entre nodos de s1*s2*...*sk*n^{k-2} formas Combinatoria Catalan: $C \{n+1\} = sum(C i*C \{n-i\} for i \in [0, n])$ Catalan: $C_n = \frac{1}{n+1}*\frac{2n}{n}$ Sea C_n^k las formas de poner n+k pares de parntesis, con los primeros k parntesis abiertos (esto es, hay 2n + 2k carcteres), se tiene que $C n^k = (2n+k-1)*(2n+k)/(n*(n+k+1)) * C {n-1}^k$ Sea D_n el nmero de permutaciones sin puntos fijos, entoces $D n = (n-1)*(D \{n-1\} + D \{n-2\}), D 0 = 1, D 1 = 0$

8.12 tonelli-shanks

```
ll legendre(ll a. ll p) {
   if (a % p == 0) return 0; if (p == 2) return 1;
   return binexp(a, (p-1) / 2, p);
// sgrt(n) mod p (p must be a prime)
// rnd(a, b) return a random number in [a, b]
ll tonelli_shanks(ll n, ll p) {
   if (n == 0) return 0:
   if (legendre(n, p) != 1) return -1; // no existe
   if (p == 2) return 1:
   ll s = __builtin_ctzll(p - 1);
   11 q = (p - 1LL) >> s, z = rnd(1, p - 1);
   if (s == 1) return binexp(n, (p + 1) / 4LL, p);
   while (legendre(z, p) != p - 1) z = rnd(1, p - 1);
   ll c = binexp(z, q, p), r = binexp(n, (q + 1) / 2, p);
   ll t = binexp(n, q, p), m = s;
   while (t != 1) {
       11 i = 1, ts = (t * t) % p:
       while (ts != 1) i++, ts = (ts * ts) % p;
      11 b = c:
       repx(_, 0, m - i - 1) b = (b * b) % p;
      r = r*b\%p; c = b*b\%p; t = t*c\%p; m = i;
   return r:
```