Team Notebook

October 15, 2023

Contents			-		4.6 minkowski	9	6.5	dynamic-connectivity	17
					4.7 point	9	6.6	mo	17
1	\mathbf{Str}	ngs	2		4.8 polygon	9	6.7	ordered-set	
	1.1	Manacher	2		4.9 sweep	10	6.8	persistent-segment-tree-lazy	
	1.2	aho-corasick	2			10	6.9	persistent-segment-tree	
	1.3	hash	2) segment-tree-lazy	
	1.4	hash2d	3	5	graph	10		segment-tree	
	1.5	palindromic-tree	3		5.1 bellman-ford	10		2 sparse-table	
	1.6	prefix-function	4		5.2 blossom	10		3 unordered-map	
	1.7	suffix-array	4		$5.3 \text{dinic} \dots \dots \dots \dots \dots$	10	0.1.	o unordered-map	16
	1.8	suffix-automaton	4		5.4 floyd-warshall	11 ,	7 im	orimible	1(
	1.9	z-function	5		5.5 heavy-light	11	11111	5111111616	Τ.
					5.6 hungarian	$12 \mid \xi$	3 ma	h	19
2	con	nmon	6		5.6 hungarian	- 1 (8 ma 8 1		19 19
		nmon	6		-	12	8.1	arithmetic	
	$d\mathbf{p}$		6 6		5.7 kuhn	12 12 13	8.1 8.2	arithmetic	19
		convex-hull-trick	6 6		5.7 kuhn 5.8 lca 5.9 maxflow-mincost 5.10 push-relabel	12 12 13 13	8.1 8.2 8.3	arithmetic	19 19
	$d\mathbf{p}$		6 6 6		5.7 kuhn	12 12 13 13	8.1 8.2 8.3 8.4	arithmetic	19 19 19
3	dp 3.1 3.2	convex-hull-trick	6 6 6		5.7 kuhn 5.8 lca 5.9 maxflow-mincost 5.10 push-relabel	12 12 13 13 14	8.1 8.2 8.3 8.4 8.5	arithmetic	19 19 19 20
3	dp 3.1	convex-hull-trick	6 6 6 6		5.7 kuhn 5.8 lca 5.9 maxflow-mincost 5.10 push-relabel 5.11 strongly-connected-components	12 12 13 13 14	8.1 8.2 8.3 8.4 8.5 8.6	arithmetic	19 19 19 20 21
3	dp 3.1 3.2	convex-hull-trick	6 6 6 6 6	6	5.7 kuhn 5.8 lca 5.9 maxflow-mincost 5.10 push-relabel 5.11 strongly-connected-components 5.12 two-sat	12 12 13 13 14	8.1 8.2 8.3 8.4 8.5 8.6 8.7	arithmetic	19 19 19 20 21 21
3	dp 3.1 3.2	convex-hull-trick divide-and-conquer	6 6 6 6 6 7	6	5.7 kuhn 5.8 lca 5.9 maxflow-mincost 5.10 push-relabel 5.11 strongly-connected-components 5.12 two-sat	12 12 13 13 14 14 14	8.1 8.2 8.3 8.4 8.5 8.6 8.7	arithmetic crt discrete-log fft gauss matrix mobius mod mod create crt mod crt	19 19 20 21 21 21
3	dp 3.1 3.2 geo 4.1	convex-hull-trick	6 6 6 6 6 7 7	6	5.7 kuhn 5.8 lca 5.9 maxflow-mincost 5.10 push-relabel 5.11 strongly-connected-components 5.12 two-sat implementation	12 12 13 13 14 14 14 14	8.1 8.2 8.3 8.4 8.5 8.6 8.7	arithmetic	19 19 20 21 21 21
3	dp 3.1 3.2 geo 4.1 4.2	convex-hull-trick	6 6 6 6 7 7 8	6	5.7 kuhn 5.8 lca 5.9 maxflow-mincost 5.10 push-relabel 5.11 strongly-connected-components 5.12 two-sat implementation 6.1 SegmentTreeBeats	12 12 13 13 14 14 14 14 14 15	8.1 8.2 8.3 8.4 8.5 8.6 8.7 8.8 8.9	arithmetic crt discrete-log fft gauss matrix mobius mod mod create crt mod crt	19 19 19 20 21 21 21 22

1 Strings

1.1 Manacher

```
// odd[i]: length of longest palindrome centered at i
// even[i]: ...longest palindrome centered between i and i+1
void manacher(string &s,vector<int> &odd,vector<int> &even){
   string t = "$#";
   for(char c: s) t += c + string("#");
   t += "^":
   int n = t.size();
   vector<int> p(n):
   int 1 = 1, r = 1;
   repx(i, 1, n-1) {
       p[i] = max(0, min(r - i, p[1 + (r - i)]));
       while(t[i - p[i]] == t[i + p[i]]) p[i]++;
       if(i + p[i] > r) l = i - p[i], r = i + p[i];
   repx(i, 2, n-2) {
       if(i%2) even.push_back(p[i]-1);
       else odd.push_back(p[i]-1);
}
```

1.2 aho-corasick

```
const int K = 26:
struct Vertex {
    int next[K];
    int leaf = 0:
    int leaf_id = -1;
    int p = -1;
    char pch:
    int link = -1;
    int exit = -1:
    int cnt = -1:
    int go[K];
    Vertex(int p=-1, char ch='$') : p(p), pch(ch) {
       fill(begin(next), end(next), -1);
       fill(begin(go), end(go), -1);
   }
};
vector<Vertex> t(1):
void add(string &s, int id) {
    int. v = 0:
    for (char ch : s) {
```

```
int c = ch - 'a':
      if (t[v].next[c] == -1) {
          t[v].next[c] = t.size();
          t.emplace_back(v, ch);
      v = t[v].next[c]:
   t[v].leaf++;
   t[v].leaf_id = id;
int go(int v. char ch):
int get_link(int v) {
   if (t[v].link == -1) {
      if (v == 0 || t[v].p == 0)
          t[v].link = 0:
          t[v].link = go(get link(t[v].p), t[v].pch):
   }
   return t[v].link;
int go(int v, char ch) {
   int c = ch - 'a':
   if (t[v].go[c] == -1) {
      if (t[v].next[c] != -1)
          t[v].go[c] = t[v].next[c]:
          t[v].go[c] = v == 0 ? 0 : go(get_link(v), ch);
   return t[v].go[c];
int next match(int v){
   if(t[v].exit == -1)
      if(t[get_link(v)].leaf)
          t[v].exit = get_link(v);
          t[v].exit = v==0 ? 0 : next match(get link(v)):
   return t[v].exit:
int cnt matches(int v){
   if(t[v].cnt == -1)
      t[v].cnt = v == 0 ? 0 : t[v].leaf + cnt_matches(
           get link(v)):
   return t[v].cnt;
```

1.3 hash

```
// compute substring hashes in O(1).
// hashes are compatible between different strings.
struct Hash {
   11 HMOD:
   int N:
   vector<int> h:
   vector<int> p;
   Hash() {}
   // O(N)
   Hash(const string \&s. 11 HMOD = 1000003931)
       : N(s.size() + 1), HMOD(HMOD_), p(N), h(N) {
       static const 11 P =
           chrono::steady_clock::now().time_since_epoch().
               count() % (1 << 29);
       p[0] = 1;
       rep(i, N - 1) p[i + 1] = p[i] * P % HMOD;
       rep(i, N-1) h[i+1] = (h[i] + (ll)s[i] * p[i]) %
   }
   pair<11, int> get(int i, int j) { return {(h[j] - h[i] +
        HMOD) % HMOD, il: }
   bool cmp(pair<11, int> x0, pair<11, int> x1) {
       int d = x0.second - x1.second:
       11 &lo = d < 0 ? x0.first : x1.first;</pre>
      lo = lo * p[abs(d)] % HMOD;
       return x0.first == x1.first;
   }
};
// compute hashes in multiple prime modulos simultaneously,
// to reduce the chance of collisions.
struct HashM {
   int N:
   vector<Hash> sub:
   HashM() {}
   // O(K N)
   HashM(const string &s, const vector<11> &mods) : N(mods.
        size()), sub(N) {
       rep(i, N) sub[i] = Hash(s, mods[i]);
```

1.4 hash2d

```
using Hash = pair<11, int>;
struct Block {
   int x0, y0, x1, y1;
struct Hash2d {
   11 HMOD;
   int W. H:
   vector<int> h;
   vector<int> p;
   Hash2d() {}
   Hash2d(const string &s, int W_, int H_, 11 HMOD_ =
        1000003931)
       : W(W_ + 1), H(H_ + 1), HMOD(HMOD_) {
       static const 11 P =
          chrono::steady_clock::now().time_since_epoch().
               count() % (1 << 29);
      p.resize(W * H):
      p[0] = 1;
      rep(i, W * H - 1) p[i + 1] = p[i] * P % HMOD;
      h.assign(W * H, 0);
      repx(v, 1, H) repx(x, 1, W) {
```

```
ll c = (ll)s[(y - 1) * (W - 1) + x - 1] * p[y * W
          h[y * W + x] = (HMOD + h[y * W + x - 1] + h[(y -
                         h[(y - 1) * W + x - 1] + c) %
      }
   }
   bool isout(Block s) {
       return s.x0 < 0 || s.x0 >= W || s.x1 < 0 || s.x1 >= W
             | | s.v0 < 0 | |
             s.y0 >= H \mid \mid s.y1 < 0 \mid \mid s.y1 >= H;
   }
   Hash get(Block s) {
       return \{(2 * HMOD + h[s.v1 * W + s.x1] - h[s.v1 * W +
             s.x01 -
               h[s.v0 * W + s.x1] + h[s.v0 * W + s.x0]) %
              s.v0 * W + s.x0;
   }
   bool cmp(Hash x0, Hash x1) {
       int d = x0.second - x1.second;
       11 &lo = d < 0 ? x0.first : x1.first;</pre>
       lo = lo * p[abs(d)] % HMOD;
       return x0.first == x1.first:
   }
}:
struct Hash2dM {
   int N:
   vector<Hash2d> sub;
   Hash2dM(const string &s, int W, int H, const vector<11> &
        mods)
       : N(mods.size()), sub(N) {
       rep(i, N) sub[i] = Hash2d(s, W, H, mods[i]):
   bool isout(Block s) { return sub[0].isout(s): }
   vector<Hash> get(Block s) {
       vector<Hash> hs(N):
       rep(i, N) hs[i] = sub[i].get(s);
       return hs:
   }
```

1.5 palindromic-tree

```
struct Node { // (*) = Optional
int len; // length of substring
int edge[26];// insertion edge for all characters a-z
int link: // the Maximum Palindromic Suffix Node for the
     current Node
           // (*) start index of current Node
int i:
   int cnt = 1; // (*) # of occurrences of this substring
   Node(){ fill(begin(edge), end(edge), -1); }
}:
struct EerTree { // Palindromic Tree
   vector<Node> t; // tree
   int curr:
                 // current node
   EerTree(string &s) {
       t.resize(2);
       t.reserve(s.size()+2):// (*) max size of tree
       t[0].len = -1:
                          // root 1
       t[0].link = 0;
       t[1].len = 0:
                          // root 2
      t[1].link = 0;
       curr = 1:
       rep(i, s.size()) insert(i, s); // construct tree
       // (*) calculate number of occurrences of each node
      for(int i = t.size()-1; i > 1; i--)
          t[t[i].link].cnt += t[i].cnt;
   void insert(int i, string &s) {
       int tmp = curr;
       while (i - t[tmp].len < 1 \mid | s[i] != s[i-t[tmp].len
           -17)
          tmp = t[tmp].link;
```

```
if(t[tmp].edge[s[i]-'a'] != -1){
          curr = t[tmp].edge[s[i]-'a']; // already exists
          t[curr].cnt++:
                             // (*) increase cnt
          return:
       // create new node
       curr = t[tmp].edge[s[i]-'a'] = t.size();
       t.emplace_back();
      t[curr].len = t[tmp].len + 2; // set length
       t[curr].i = i - t[curr].len + 1://(*)set start index
      if (t[curr].len == 1) {
                                     // set suffix link
          t[curr].link = 1:
      } else {
          tmp = t[tmp].link:
          while (i-t[tmp].len < 1 \mid | s[i] != s[i-t[tmp].len
              tmp = t[tmp].link;
          t[curr].link = t[tmp].edge[s[i]-'a'];
   }
void main(){
string s = "abcbab";
   EerTree pt(s);
                         // construct palindromic tree
repx(i, 2, pt.t.size()) // list all distinct palindromes
 cout << i-1 << ") ":
 repx(j, pt.t[i].i, pt.t[i].i + pt.t[i].len)
  cout << s[i];
 cout << " " << pt.t[i].cnt << endl;</pre>
```

1.6 prefix-function

```
return pi:
vector<vector<int>> aut:
void compute_automaton(string s) {
   s += '#';
   int n = s.size():
   vector<int> pi = prefix_function(s);
   aut.assign(n, vector<int>(26));
   rep(i, n) {
      rep(c, 26) {
          int i = i:
          while (i > 0 \&\& 'a' + c != s[i])
              j = pi[j-1];
          if ('a' + c == s[j])
              j++;
          aut[i][c] = j;
      }
   }
// k = n - pi[n - 1]; if k divides n, then the string can be
// aprtitioned into blocks of length k otherwise there is no
// effective compression and the answer is n.
```

1.7 suffix-array

```
// build the suffix array
// suffixes are sorted, with each suffix represented by its
// starting position
vector<int> suffixarray(const string &s) {
   int N = s.size() + 1;//optional: include terminating NUL
   vector<int> p(N), p2(N), c(N), c2(N), cnt(256);
   rep(i, N) cnt[s[i]] += 1:
   repx(b, 1, 256) cnt[b] += cnt[b - 1];
   rep(i, N) p[--cnt[s[i]]] = i;
   repx(i, 1, N) c[p[i]] = c[p[i - 1]] + (s[p[i]] != s[p[i]])
         1]]);
   for (int k = 1: k < N: k <<= 1) {
      int C = c[p[N - 1]] + 1;
       cnt.assign(C + 1, 0);
       for (int &pi : p) pi = (pi - k + N) % N;
       for (int cl : c) cnt[cl + 1] += 1;
      rep(i, C) cnt[i + 1] += cnt[i];
       rep(i, N) p2[cnt[c[p[i]]]++] = p[i];
       c2[p2[0]] = 0;
       repx(i, 1, N) c2[p2[i]] =
          c2[p2[i-1]] + (c[p2[i]] != c[p2[i-1]] ||
                          c[(p2[i] + k) \% N] != c[(p2[i - 1]
                                + k) % N]);
       swap(c, c2), swap(p, p2);
```

```
p.erase(p.begin()); // optional: erase terminating NUL
// build the lcp
// 'lcp[i]' represents the length of the longest common
// prefix between suffix i and suffix i+1 in the suffix
//array 'p'. the last element of 'lcp' is zero by convention
vector<int> makelcp(const string &s, const vector<int> &p) {
   int N = p.size(), k = 0;
   vector<int> r(N), lcp(N):
   rep(i, N) r[p[i]] = i:
   rep(i, N) {
       if (r[i] + 1 \ge N) \{ k = 0; continue; \}
       int j = p[r[i] + 1];
       while (i + k < N \&\& j + k < N \&\& s[i + k] == s[j + k]
           1) k += 1:
       lcp[r[i]] = k;
       if (k) k -= 1:
   return lcp;
// lexicographically compare the suffixes starting from 'i'
// and 'j', considering only up to 'K' characters.
// 'r' is the inverse suffix array, mapping suffix offsets
// to indices. requires an LCP sparse table.
int lcp_cmp(vector<int> &r, Sparse<int> &lcp, int i, int j,
    int K) {
   if (i == j) return 0;
   int ii = r[i]. ii = r[i]:
   int 1 = lcp.query(min(ii, jj), max(ii, jj));
   if (1 >= K) return 0;
   return ii < jj ? -1 : 1;</pre>
```

1.8 suffix-automaton

```
struct SuffixAutomaton {
    // edges[i]: the labeled edges from node i
    vector<map<char,int>> edges;
    // link[i]: the suffix link of i
    vector<int> link;
    // length[i]: len of the longest string in the ith class
    vector<int> length;
    // cnt[i]: # occurrences of each string in the ith class
    vector<int> cnt;
    // paths[i]: # of paths on the automaton starting from i
    vector<int> paths;
    // terminal[i]: true if i is a terminal state
```

```
vector<bool> terminal:
   vector<int> first_pos;
   vector<int> last_pos;
   //the index of the equivalence class of the whole string
   int last:
   SuffixAutomaton(string s) {
       edges.push_back(map<char,int>());
      link.push_back(-1);
      length.push_back(0);
      last = 0:
      rep(i, s.size()) { // construct r
          edges.push_back(map<char,int>());
          length.push_back(i+1);
          link.push_back(0);
          int r = edges.size() - 1;
          int p = last;
          // add edges to r and find p with link to q
          while(p >= 0 && !edges[p].count(s[i])) {
              edges[p][s[i]] = r;
              p = link[p];
          }
          if(p != -1) {
              int q = edges[p][s[i]];
              if(length[p] + 1 == length[q]) {
//we don't have to split q, just set the correct suffix link
                 link[r] = a:
              } else { // we have to split, add q'
                 // copy edges of q
                  edges.push_back(edges[q]);
                 length.push_back(length[p] + 1);
                 // copy parent of q
                 link.push_back(link[q]);
                 int qq = edges.size()-1;
                 // add qq as the new parent of q and r
                 link[q] = qq;
                 link[r] = qq;
          // move short classes polling to q to poll to q'
                 while(p >= 0 && edges[p][s[i]] == q) {
                     edges[p][s[i]] = qq;
                     p = link[p];
                 }
              }
          last = r:
   /* ----- Optional ----- */
      // mark terminal nodes
       terminal.assign(edges.size(), false);
```

```
int p = last:
   while (p > 0) {
       terminal[p] = true;
       p = link[p];
   // precompute match count
   cnt.assign(edges.size(), -1);
   cnt matches(0):
   //precompute # of paths (substr) starting from state
   paths.assign(edges.size(), -1);
   cnt paths(0):
   first_pos.assign(edges.size(), -1);
   get_first_pos(0);
   last_pos.assign(edges.size(), -1);
   get_last_pos(0);
int cnt matches(int state) {
   if(cnt[state] != -1) return cnt[state];
   int ans = terminal[state];
   for(auto edge : edges[state])
       ans += cnt_matches(edge.second);
   return cnt[state] = ans:
}
int cnt paths(int state) {
   if(paths[state] != -1) return paths[state];
   // without repetition (counts different substrings)
   int ans = state == 0 ? 0 : 1:
   // with repetition
// int ans = state == 0 ? 0 : cnt[state];
   for(auto edge : edges[state])
       ans += cnt_paths(edge.second);
   return paths[state] = ans:
int get_first_pos(int state) {
   if(first_pos[state] != -1) return first_pos[state];
   int ans = 0:
   for(auto edge : edges[state])
       ans = max(ans, get_first_pos(edge.second)+1);
   return first_pos[state] = ans;
int get_last_pos(int state) {
   if(last_pos[state] != -1) return last_pos[state];
   int ans = terminal[state] ? 0 : INT_MAX;//fix
   for(auto edge : edges[state])
```

```
ans = min(ans, get_last_pos(edge.second)+1);
       return last_pos[state] = ans;
   7
   string get_k_substring(int k) { // 0-indexed
       string ans:
       int state = 0;
       while(true){
          // without repetition (counts different substrs)
          int curr = state == 0 ? 0 : 1;
          // with repetition
       // int curr = state == 0 ? 0 : cnt[state]:
          if(curr > k) return ans;
          k -= curr:
          for(auto edge : edges[state]) {
              if(paths[edge.second] <= k) {</pre>
                  k -= paths[edge.second];
              } else {
                  ans += edge.first:
                  state = edge.second;
                  break;
          }
   }
};
```

1.9 z-function

```
// i-th element is equal to the greatest number of
// characters starting from the position i that coincide
// with the first characters of s
vector<int> z_function(string s) {
    int n = s.size();
    vector<int> z(n);
    int l = 0, r = 0;
    for(int i = 1; i < n; i++) {
        if(i < r) z[i] = min(r - i, z[i - 1]);
        while(i + z[i] < n && s[z[i]] == s[i + z[i]])z[i]++;
        if(i + z[i] > r) {
            l = i;
            r = i + z[i];
        }
    }
    return z;
}
```

2 common

$3 ext{ dp}$

3.1 convex-hull-trick

```
struct Line {
    mutable 11 a, b, c;
    bool operator<(Line r) const { return a < r.a; }</pre>
    bool operator<(11 x) const { return c < x; }</pre>
};
// dynamically insert 'a*x + b' lines and guery for maximum
// at any x all operations have complexity O(\log N)
struct LineContainer : multiset<Line, less<>>> {
    11 div(ll a, ll b) {
       return a / b - ((a ^ b) < 0 && a % b);
    bool isect(iterator x, iterator v) {
       if (y == end()) return x \rightarrow c = INF, 0;
       if (x->a == y->a) x->c = x->b > y->b? INF : -INF;
       else x->c = div(y->b - x->b, x->a - y->a);
       return x->c >= y->c;
    void add(ll a. ll b) {
       // a *= -1, b *= -1 // for min
       auto z = insert(\{a, b, 0\}), y = z++, x = y;
       while (isect(y, z)) z = erase(z);
       if (x != begin() \&\& isect(--x, v)) isect(x, v = erase)
       while ((y = x) != begin() && (--x)->c >= y->c) isect(
            x. erase(v)):
```

```
11 query(11 x) {
    if (empty()) return -INF; // INF for min
    auto 1 = *lower_bound(x);
    return 1.a * x + 1.b;
    // return -l.a * x - 1.b; // for min
}
};
```

3.2 divide-and-conquer

```
// for every index i assign an optimal index j, such that
    cost(i, j) is
// minimal for every i. the property that if i2 >= i1 then
// exploited (monotonic condition).
// calculate optimal index for all indices in range [1, r)
    knowing that
// the optimal index for every index in this range is within
     [optl, optr).
// time: O(N log N)
void calc(vector<int> &opt, int 1, int r, int optl, int optr
   if (1 == r) return:
   int i = (1 + r) / 2;
   11 optc = INF;
   int opti:
   repx(j, optl, optr) {
      11 c = i + j; // cost(i, j)
       if (c < optc) optc = c, optj = j;</pre>
   opt[i] = opti:
   calc(opt, 1, i, optl, optj + 1);
   calc(opt, i + 1, r, optj, optr);
```

4 geo2d

4.1 circle

```
struct C {
   P o; T r;

   C(P o, T r) : o(o), r(r) {}
   C() : C(P(), T()) {}
```

```
// intersects the circle with a line, assuming they
// results are sorted with respect to the direction of
     the line
pair<P, P> line_inter(L 1) const {
   P c = 1.closest to(o):
   T c2 = (c - o).magsq();
   P = sqrt(max(r * r - c2, T())) * 1.d.unit();
   return {c - e, c + e};
// checks whether the given line collides with the circle
// negative: 2 intersections
// zero: 1 intersection
// positive: 0 intersections
T line_collide(L 1) const {
   T c2 = (1.closest_to(o) - o).magsq();
   return c2 - r * r;
// calculates the two intersections between two circles
// the circles must intersect in one or two points!
pair<P, P> inter(C h) const {
   P d = h.o - o;
   T c = (r * r - h.r * h.r) / d.magsq();
   return h.line_inter(\{(1 + c) / 2 * d, d.rot()\});
// check if the given circles intersect
bool collide(C h) const {
   return (h.o - o).magsq() \le (h.r + r) * (h.r + r);
// get one of the two tangents that cross through the
// the point must not be inside the circle
// a = -1: cw (relative to the circle) tangent
// a = 1: ccw (relative to the circle) tangent
P point_tangent(P p, T a) const {
   T c = r * r / p.magsq();
   return o + c * (p - o) - a * sqrt(c * (1 - c)) * (p - c)
         o).rot():
// get one of the 4 tangents between the two circles
// a = 1: exterior tangents
// a = -1: interior tangents (requires no area overlap)
// b = 1: ccw tangent
// b = -1: cw tangent
```

```
// the line origin is on this circumference, and the
    direction
// is a unit vector towards the other circle
L tangent(C c, T a, T b) const {
   T dr = a * r - c.r;
   P d = c.o - o:
   P n = (d * dr + b * d.rot() * sqrt(d.magsq() - dr *
        dr)).unit():
   return {o + n * r, -b * n.rot()};
// find the circumcircle of the given **non-degenerate**
    triangle
static C thru_points(P a, P b, P c) {
   L 1((a + b) / 2, (b - a).rot());
   P p = 1.intersection(L((a + c) / 2, (c - a).rot()));
   return {p, (p - a).mag()};
// find the two circles that go through the given point,
    are tangent
// to the given line and have radius 'r'
// the point-line distance must be at most 'r'!
// the circles are sorted in the direction of the line
static pair<C, C> thru_point_line_r(P a, L t, T r) {
   P d = t.d.rot().unit();
   if (d * (a - t.o) < 0) d = -d:
   auto p = C(a, r).line inter(\{t.o + d * r, t.d\}):
   return {{p.first, r}, {p.second, r}};
// find the two circles that go through the given points
    and have
// radius 'r'
// the circles are sorted by angle with respect to the
    first point
// the points must be at most at distance 'r'!
static pair<C, C> thru_points_r(P a, P b, T r) {
   return {{p.first, r}, {p.second, r}};
```

convex-hull

};

```
// get the convex hull with the least amount of vertices for
     the given set
// of points
```

```
// probably misbehaves if points are not all distinct!
vector<P> convex_hull(vector<P> &ps) {
   int N = ps.size(), n = 0, k = 0;
   if (N <= 2) return ps:
   rep(i, N) if (make_pair(ps[i].v, ps[i].x) < make_pair(ps[</pre>
        kl.v. ps[kl.x)) k = i:
   swap(ps[k], ps[0]);
   sort(++ps.begin(), ps.end(), [&](P 1, P r) {
       T x = (r - 1) \% (ps[0] - 1), d = (r - 1) * (ps[0] - 1)
       return x > 0 \mid | x == 0 && d < 0:
   }):
   vector<P> H:
   for (P p : ps) {
       while (n \ge 2 \&\& (H[n-1] - p) \% (H[n-2] - p) >=
            0) H.pop back(), n--:
       H.push_back(p), n++;
   return H:
```

4.3 delaunay

```
typedef __int128_t lll; // if on a 64-bit platform
                                                          struct Q {
                                                             Q *rot, *o; P p = {INF, INF}; bool mark;
                                                             P &F() { return r()->p; }
                                                             Q *&r() { return rot->rot; }
                                                             Q *prev() { return rot->o->rot; }
                                                             Q *next() { return r()->prev(): }
                                                         };
                                                         T cross(P a, P b, P c) { return (b - a) % (c - a); }
auto p = C(a, r).line_inter({(a + b) / 2, (b - a).rot | bool circ(P p, P a, P b, P c) { // is p in the circumcircle?
                                                             111 p2 = p.magsq(), A = a.magsq() - p2,
                                                                 B = b.magsq() - p2, C = c.magsq() - p2;
                                                             return cross(p, a, b) * C + cross(p, b, c) * A + cross(p, b, c)
                                                                   c. a) * B > 0:
                                                         Q *makeEdge(Q *&H, P orig, P dest) {
                                                             Q *r = H ? H : new Q{new Q{new Q{new Q{0}}}}:
                                                             H = r -> 0; r -> r() -> r() = r;
                                                             repx(i, 0, 4) r = r->rot, r->p = {INF, INF},
                                                                 r->0 = i & 1 ? r : r->r():
                                                             r\rightarrow p = orig; r\rightarrow F() = dest;
```

```
return r:
void splice(Q *a, Q *b) {
   swap(a->o->rot->o, b->o->rot->o); swap(a->o, b->o);
Q *connect(Q *&H, Q *a, Q *b) {
   Q *q = makeEdge(H, a->F(), b->p);
   splice(q, a->next()); splice(q->r(), b); return q;
pair<Q *, Q *> rec(Q *&H, const vector<P> &s) {
   if (s.size() <= 3) {</pre>
       Q *a = makeEdge(H, s[0], s[1]), *b = makeEdge(H, s[0], s[1])
            [1], s.back());
       if (s.size() == 2) return \{a, a->r()\}; splice(a->r(),
       auto side = cross(s[0], s[1], s[2]):
       Q *c = side ? connect(H, b, a) : 0:
       return {side < 0 ? c->r() : a, side < 0 ? c : b->r()
   }
#define J(e) e \rightarrow F(), e \rightarrow p
#define valid(e) (cross(e->F(), J(base)) > 0)
   Q *A, *B, *ra, *rb; int half = s.size() / 2;
   tie(ra, A) = rec(H, {s.begin(), s.end() - half});
   tie(B, rb) = rec(H, {s.begin() + s.size() - half, s.end()
        }):
   while ((cross(B->p, J(A)) < 0 \&\& (A = A->next())) | |
          (cross(A->p, J(B)) > 0 \&\& (B = B->r()->o)));
   Q *base = connect(H, B->r(), A):
   if (A->p == ra->p) ra = base->r();
   if (B->p == rb->p) rb = base:
#define DEL(e, init, dir) Q *e = init->dir; \
   if (valid(e)) while (circ(e->dir->F(), J(base), e->F()))
           Q *t = e->dir; splice(e, e->prev()); \
           splice(e->r(), e->r()->prev()); e->o = H; H = e;
                e = t: \
       }
   for (;;) {
       DEL(LC, base->r(), o); DEL(RC, base, prev());
       if (!valid(LC) && !valid(RC)) break;
       if (!valid(LC) || (valid(RC) && circ(J(RC), J(LC))))
           base = connect(H, RC, base->r()):
       else base = connect(H, base->r(), LC->r());
```

```
return {ra, rb}:
#undef J
#undef valid
#undef DEL
// there must be no duplicate points
// returns no triangles in the case of all collinear points
// produces counter-clockwise triangles ordered in triples
// maximizes the minimum angle across all triangulations
// the euclidean mst is a subset of these edges
// O(N log N)
vector<P> triangulate(vector<P> pts) {
   sort(pts.begin(), pts.end(), [](P a, P b) {
       return make_pair(a.x, a.y) < make_pair(b.x, b.y);</pre>
   assert(unique(pts.begin(), pts.end()) == pts.end());
   if (pts.size() < 2) return {};</pre>
   Q *H = 0; Q *e = rec(H, pts).first;
   vector<Q *> a = {e}: int ai = 0:
   while (cross(e->o->F(), e->F(), e->p) < 0) e = e->o;
#define ADD
   ſ
       0 *c = e:
       do {
           c->mark = 1; pts.push_back(c->p); \
           q.push_back(c->r()); c = c->next(); \setminus
       } while (c != e):
   ADD:
   pts.clear();
   while (qi < (int)q.size()) if (!(e = q[qi++])->mark) ADD;
   return pts;
#undef ADD
```

4.4 halfplane-intersect

```
// obtain the convex polygon that results from intersecting
    the given list
// of halfplanes, represented as lines that allow their left
    side
// assumes the halfplane intersection is bounded
vector<P> halfplane_intersect(vector<L> &H) {
    L bb(P(-INF, -INF), P(INF, 0));
    rep(k, 4) H.push_back(bb), bb.o = bb.o.rot(), bb.d = bb.d
    .rot();
```

```
sort(begin(H), end(H), [](L a, L b) { return a.d.angcmp(b
     .d) < 0: }):
deque<L> q; int n = 0;
rep(i, H.size()) {
   while (n \ge 2 \&\& H[i].side(q[n - 1].intersection(q[n
        -21)) > 0)
       q.pop_back(), n--;
   while (n \ge 2 \&\& H[i].side(q[0].intersection(q[1])) >
       q.pop_front(), n--;
   if (n > 0 \&\& H[i].parallel(q[n - 1])) {
       if (H[i].d * a[n - 1].d < 0) return {}:
       if (H[i].side(q[n - 1].o) > 0) q.pop_back(), n--;
       else continue:
   }
   q.push_back(H[i]), n++;
while (n \ge 3 \&\& a[0].side(a[n - 1].intersection(a[n -
    21)) > 0)
    g.pop_back(), n--;
while (n \ge 3 \&\& q[n - 1].side(q[0].intersection(q[1])) >
   q.pop_front(), n--;
if (n < 3) return {};</pre>
vector<P> ps(n);
rep(i, n) ps[i] = q[i].intersection(q[(i + 1) % n]);
return ps;
```

4.5 line

```
// a segment or an infinite line
// does not handle point segments correctly!
struct L {
   P o, d;
   L() : o(), d() {}
   L(P o, P d) : o(o), d(d) {}

   L(P ab, T c) : d(ab.rot()), o(ab * -c / ab.magsq()) {}
   pair<P, T> line_eq() { return {-d.rot(), d.rot() * o}; }

   // returns a number indicating which side of the line the point is in
   // negative: left, positive: right
   T side(P r) const { return (r - o) % d; }

   // returns the intersection coefficient
```

```
// in the range [0, d % r.d]
// if d % r.d is zero, the lines are parallel
T inter(L r) const { return (r.o - o) % r.d; }
// get the single intersection point
// lines must not be parallel
P intersection(L r) const { return o + d * inter(r) / (d
    % r.d): }
// check if lines are parallel
bool parallel(L r) const { return abs(d % r.d) <= EPS: }</pre>
// check if segments intersect
bool seg_collide(L r) const {
   Tz = d \% r.d:
   if (abs(z) <= EPS) {
       if (abs(side(r.o)) > EPS) return false;
       T s = (r.o - o) * d, e = s + r.d * d;
       if (s > e) swap(s, e):
       return s <= d * d + EPS && e >= -EPS:
   T s = inter(r), t = -r.inter(*this);
   if (z < 0) s = -s, t = -t, z = -z;
   return s >= -EPS && s <= z + EPS && t >= -EPS && t <=
         z + EPS:
}
// full segment intersection
// produces a point segment if the intersection is a
// however it **does not** handle point segments as input
bool seg_inter(L r, L *out) const {
   T z = d \% r.d;
   if (abs(z) <= EPS) {
       if (abs(side(r.o)) > EPS) return false;
       if (r.d * d < 0) r = \{r.o + r.d, -r.d\};
       P s = o * d < r.o * d ? r.o : o;
       P = (o + d) * d < (r.o + r.d) * d ? o + d : r.o
             + r.d:
       if (s * d > e * d) return false;
       return *out = {s, e - s}, true;
   }
   T s = inter(r), t = -r.inter(*this);
   if (z < 0) s = -s, t = -t, z = -z;
   if (s >= -EPS && s <= z + EPS && t >= -EPS && t <= z
       return *out = {o + d * s / z. {0, 0}}, true:
   return false:
```

4.6 minkowski

```
void reorder polygon(vector<P> &ps) {
   int pos = 0:
   repx(i, 1, (int)ps.size()) {
       if (ps[i].y < ps[pos].y || (ps[i].y == ps[pos].y &&</pre>
           ps[i].x < ps[pos].x)
          pos = i:
   rotate(ps.begin(), ps.begin() + pos, ps.end());
vector<P> minkowski(vector<P> ps. vector<P> qs) {
   // the first vertex must be the lowest
   reorder_polygon(ps); reorder_polygon(qs);
   ps.push_back(ps[0]); ps.push_back(ps[1]);
   qs.push_back(qs[0]); qs.push_back(qs[1]);
   vector<P> result: int i = 0, i = 0:
   while (i < ps.size() - 2 || j < qs.size() - 2) {
       result.push_back(ps[i] + qs[j]);
       auto z = (ps[i + 1] - ps[i]) \% (qs[j + 1] - qs[j]);
       if (z \ge 0 \&\& i < ps.size() - 2) ++i;
       if (z <= 0 && j < qs.size() - 2) ++j;</pre>
   }
   return result;
```

4.7 point

```
struct P {
   T x, y;
```

```
P(T x, T y) : x(x), y(y) {}
P() : P(0, 0) \{ \}
friend ostream &operator<<(ostream &s, const P &r) {</pre>
   return s << r.x << " " << r.y;
friend istream & operator >> (istream &s, P &r) { return s
    >> r.x >> r.v: }
P operator+(P r) const { return \{x + r.x, y + r.y\}; }
P operator-(P r) const { return {x - r.x, y - r.y}; }
P operator*(T r) const { return {x * r, v * r}: }
P operator/(T r) const { return {x / r, y / r}; }
P operator-() const { return {-x, -y}; }
friend P operator*(T 1, P r) { return {1 * r.x, 1 * r.y};
P rot() const { return {-y, x}; }
T operator*(P r) const { return x * r.x + v * r.v: }
T operator%(P r) const { return rot() * r; }
T magsq() const { return x * x + y * y; }
T mag() const { return sqrt(magsq()); }
P unit() const { return *this / mag(); }
bool half() const { return abs(y) <= EPS && x < -EPS || y | struct InConvex {
T angcmp(P r) const {
   int h = (int)half() - r.half();
   return h ? h : r % *this:
bool operator==(P r) const { return abs(x - r.x) <= EPS</pre>
     && abs(y - r.y) <= EPS; }
double angle() const { return atan2(y, x); }
static P from_angle(double a) { return {cos(a), sin(a)};
    }
```

4.8 polygon

```
// get the area of a simple polygon in ccw order
T area(const vector<P> &ps) {
   int N = ps.size();
   T a = 0;
   rep(i, N) a += (ps[i] - ps[0]) % (ps[(i + 1) % N] - ps[i]);
   return a / 2;
```

```
// checks whether a point is inside a simple polygon
// returns -1 if inside, 0 if on border, 1 if outside
// O(N)
int in polv(const vector<P> &ps. P p) {
   int N = ps.size(), w = 0;
   rep(i, N) {
       P s = ps[i] - p, e = ps[(i + 1) \% N] - p;
       if (s == P()) return 0;
       if (s.v == 0 \&\& e.v == 0) {
          if (\min(s.x. e.x) \le 0 \&\& 0 \le \max(s.x. e.x))
               return 0:
      } else {
           bool b = s.v < 0;
          if (b != (e.v < 0)) {
              Tz = s \% e: if (z == 0) return 0:
              if (b == (z > 0)) w += b ? 1 : -1;
      }
   return w ? -1 : 1:
// check if a point is in a convex polygon
   vector<P> ps:
   T 11. 1h. rl. rh:
   int N, m;
   // preprocess polygon
   // O(N)
   InConvex(const vector<P> &p) : ps(p), N(ps.size()), m(0)
       assert(N >= 2):
       rep(i, N) if (ps[i].x < ps[m].x) m = i;
       rotate(ps.begin(), ps.begin() + m, ps.end());
       rep(i, N) if (ps[i].x > ps[m].x) m = i;
      11 = 1h = ps[0].v, r1 = rh = ps[m].v;
       for (P p : ps) {
          if (p.x == ps[0].x) 11 = min(11, p.y), 1h = max(
               lh, p.v);
          if (p.x == ps[m].x) rl = min(rl, p.y), rh = max(
               rh, p.y);
      }
   InConvex() {}
   // check if point belongs in polygon
   // returns -1 if inside. 0 if on border. 1 if outside
```

```
// O(log N)
   int in_poly(P p) {
       if (p.x < ps[0].x || p.x > ps[m].x) return 1;
       if (p.x == ps[0].x) return p.y < 11 || p.y > 1h;
       if (p.x == ps[m].x) return p.y < rl \mid \mid p.y > rh;
       int r = upper_bound(ps.begin(), ps.begin() + m, p,
           [](Pa, Pb) { return a.x < b.x; }) - ps.begin();
       Tz = (ps[r-1] - ps[r]) \% (p - ps[r]); if (z >= 0)
            return !!z;
       r = upper_bound(ps.begin() + m, ps.end(), p,
           [](P a, P b) { return a.x > b.x; }) - ps.begin();
       z = (ps[r - 1] - ps[r \% N]) \% (p - ps[r \% N]);
       if (z >= 0) return !!z; return -1;
};
// classify collision of a ray inside a ccw polygon vertex.
// ray is (o, d), vertex is b, previous vertex is a, next is
pair<bool, bool> inner_collide(P o, P d, P a, P b, P c) {
   T p = (a - o) \% d; // side of previous
                        // side of next
   T n = (c - o) \% d:
   T v = (c - b) \% (b - a); // is vertex convex?
   return \{v > 0 ? n < 0 | l (n == 0 && p < 0) : p > 0 | l n < l//
          v > 0 ? p > 0 || (p == 0 && n > 0) : p > 0 || n <
```

4.9 sweep

```
#include "point.cpp"
// iterate over all pairs of points
// 'op' is called with all ordered pairs of different
    indices '(i, j)'
// additionally, the 'ps' vector is kept sorted by signed
     distance
// to the line formed by 'i' and 'i'
// for example, if the vector from 'i' to 'j' is pointing
     right.
// the 'ps' vector is sorted from smallest 'y' to largest 'y
// note that, because the 'ps' vector is sorted by signed
    distance.
// 'j' is always equal to 'i + 1'
// this means that the amount of points to the left of the
    line is always 'N - i'
template <class OP>
```

4.10 theorems

```
// Pick's theorem
// Simple polygon with integer vertices:
// A = I + B / 2 - 1
// A: Area of the polygon
// I: Integer points strictly inside the polygon
// B: Integer points on the boundary of the polygon
```

5 graph

5.1 bellman-ford

```
struct Edge { int u, v; ll w; };

// find distance from source node to all nodes.
// supports negative edge weights.
// returns true if a negative cycle is detected.
//
// time: O(V E)
bool bellman_ford(int N, int s, vector<Edge> &E, vector<ll> &D, vector<int> &P) {
    P.assign(N, -1), D.assign(N, INF), D[s] = 0;
    rep(i, N - 1) {
        bool f = true;
        rep(ei, E.size()) {
            auto &e = E[ei];
            ll n = D[e.u] + e.w;
            if (D[e.v] < INF && n < D[e.v])</pre>
```

```
D[e.v] = n, P[e.v] = ei, f = false;
}
if (f) return false;
}
return true;
}
```

5.2 blossom

vector<int> g[MAXN];int n,m,mt[MAXN],qh,qt,q[MAXN],ft[MAXN], bs[MAXN];bool inq[MAXN],inb[MAXN],inp[MAXN];int lca(int root ,int x,int y){memset(inp,0,sizeof(inp));while(1){inp[x=bs[x]]=true:if(x==root)break:x=ft[mt[x]]:}while(1){if(inp[v=bs[v]]])return y;else y=ft[mt[y]];}}void mark(int z,int x){while(bs[x]!=z){int y=mt[x];inb[bs[x]]=inb[bs[y]]=true;x=ft[y];if(bs[x]!=z)ft[x]=y;}}void contr(int s,int x,int y){int z=lca(s ,x,y);memset(inb,0,sizeof(inb));mark(z,x);mark(z,y);if(bs[x] $!=z)ft[x]=y;if(bs[y]!=z)ft[y]=x;rep(x,n)if(inb[bs[x]]){bs[x]}$ =z;if(!inq[x])inq[q[++qt]=x]=true;}}int findp(int s){memset(ing.O.sizeof(ing)):memset(ft.-1.sizeof(ft)):rep(i.n)bs[i]=i: inq[q[qh=qt=0]=s]=true; while(qh<=qt){int x=q[qh++]; for(int y</pre> $:g[x])if(bs[x]!=bs[y]\&\&mt[x]!=y)\{if(y==s||mt[y]>=0\&\&ft[mt[y]])$]>=0)contr(s,x,y);else if(ft[y]<0){ft[y]=x;if(mt[y]<0)return} v;else if(!ing[mt[v]])ing[g[++qt]=mt[v]]=true;}}}return -1;} int aug(int s,int t){int x=t,y,z;while(x>=0){y=ft[x];z=mt[y]} ;mt[y]=x;mt[x]=y;x=z;}return t>=0;}int edmonds(){int r=0; memset(mt,-1,sizeof(mt));rep(x,n)if(mt[x]<0)r+=aug(x,findp(x))</pre>)):return r:}

5.3 dinic

```
struct Edge { int u, v: 11 c, f = 0: }:
// maximum flow algorithm.
// time: 0(E V^2)
//
        O(E V^{(2/3)}) / O(E \operatorname{sqrt}(E)) unit capacities
//
        O(E sart(V))
                                     unit networks (hopcroft-
    karp)
// unit network: c in {0, 1} and forall v, len(incoming(v))
    <= 1 or len(outgoing(v)) <= 1
// min-cut: find all nodes reachable from the source in the
    residual graph
struct Dinic {
   int N, s, t; vector<vector<int>> G;
   vector<Edge> E; vector<int> lvl, ptr;
   Dinic() {}
   Dinic(int N, int s, int t): N(N), s(s), t(t), G(N) {}
```

```
void add_edge(int u, int v, ll c) {
       G[u].push_back(E.size()); E.push_back({u, v, c});
       G[v].push_back(E.size()); E.push_back({v, u, 0});
   11 push(int u, 11 p) {
       if (u == t || p <= 0) return p;</pre>
       while (ptr[u] < G[u].size()) {</pre>
           int ei = G[u][ptr[u]++];
           Edge &e = E[ei]:
           if (lvl[e.v] != lvl[u] + 1) continue:
           ll a = push(e.v, min(e.c - e.f, p));
           if (a <= 0) continue; e.f += a, E[ei ^ 1].f -= a;</pre>
       }
       return 0:
   11 maxflow() {
       11 f = 0;
       while (true) {
           lvl.assign(N, -1); queue < int > q; <math>lvl[s] = 0, q.
                push(s);
           while (!q.empty()) {
               int u = q.front(); q.pop();
              for (int ei : G[u]) {
                  Edge &e = E[ei]:
                  if (e.c - e.f <= 0 || lvl[e.v] != -1)</pre>
                       continue:
                  lvl[e.v] = lvl[u] + 1, q.push(e.v);
           }
           if (lvl[t] == -1) break;
           ptr.assign(N, 0): while (11 ff = push(s, INF)) f
       }
       return f;
};
```

5.4 floyd-warshall

```
// calculate distances between every pair of nodes in O(V^3)
time.
// works with negative edges, but not negative cycles.
void floyd(const vector<vector<pair<11, int>>> &G, vector<
    vector<11>>> &D) {
```

```
int N = G.size();
D.assign(N, vector<11>(N, INF));
rep(u, N) D[u][u] = 0;
rep(u, N) for (auto [w, v] : G[u]) D[u][v] = w;
rep(k, N) rep(u, N) rep(v, N)
        D[u][v] = min(D[u][v], D[u][k] + D[k][v]);
}
```

5.5 heavy-light

```
struct Hld {
   vector<int> P, H, D, pos, top;
   void init(vector<vector<int>> &G) {
      int N = G.size():
      P.resize(N), H.resize(N), D.resize(N), pos.resize(N),
          top.resize(N):
      D[0] = -1, dfs(G, 0); int t = 0;
      rep(i, N) if (H[P[i]] != i) {
          int i = i:
          while (j != -1)
             \{ top[j] = i, pos[j] = t++; j = H[j]; \}
      7
   }
   int dfs(vector<vector<int>> &G, int i) {
      int w = 1, mw = 0:
      D[i] = D[P[i]] + 1, H[i] = -1;
      for (int c : G[i]) {
          if (c == P[i]) continue;
          P[c] = i: int sw = dfs(G, c): w += sw:
          if (sw > mw) H[i] = c, mw = sw;
      }
      return w;
   // visit the log N segments in the path from u to v
   template <class OP>
   void path(int u, int v, OP op) {
      while (top[u] != top[v]) {
          if (D[top[u]] > D[top[v]]) swap(u, v);
          op(pos[top[v]], pos[v] + 1); v = P[top[v]];
      if (D[u] > D[v]) swap(u, v):
      op(pos[u], pos[v] + 1); // value on node
      // op(pos[u]+1, pos[v] + 1); // value on edge
   }
```

```
// an alternative to 'path' that considers order.
   // calls 'op' with an 'l <= r' inclusive-exclusive range,</pre>
         and a
   // boolean indicating if the query is forwards or
        backwards.
   template <class OP>
   void path(int u, int v, OP op) {
       int lu = u, lv = v;
       while (top[lu] != top[lv])
          if (D[top[lu]] > D[top[lv]]) lu = P[top[lu]];
          else lv = P[top[lv]]:
       int lca = D[lu] > D[lv] ? lv : lu:
       while (top[u] != top[lca])
          op(pos[top[u]], pos[u] + 1, false), u = P[top[u
       if (u != lca) op(pos[lca] + 1, pos[u] + 1, false):
       vector<int> stk:
       while (top[v] != top[lca])
          stk.push_back(v), v = P[top[v]];
       // op(pos[lca], pos[v] + 1, true); // value on node
       op(pos[lca] + 1, pos[v] + 1, true); // value on edge
       reverse(stk.begin(), stk.end());
      for (int w : stk) op(pos[top[w]], pos[w] + 1, true);
   // commutative segment tree
   template <class T, class S>
   void update(S &seg, int i, T val) { seg.update(pos[i],
        val): }
   // commutative segment tree lazy
   template <class T, class S>
   void update(S &seg, int u, int v, T val) {
      path(u, v, [&](int 1, int r) { seg.update(1, r, val);
            }):
   }
   // commutative (lazv) segment tree
   template <class T, class S>
   T query(S &seg, int u, int v) {
      T ans = 0:
           // neutral element
       path(u, v, [\&](int 1, int r) \{ ans += seg.query(1, r) \}
           ; }); // query op
      return ans:
   }
};
```

5.6 hungarian

```
// find a maximum gain perfect matching in the given
    bipartite complete graph.
// input: gain matrix (G_{xy} = benefit of joining vertex x
    in set X with vertex
// v in set Y).
// output: maximum gain matching in members 'xy[x]' and 'yx[
    y]'.
// runtime: O(N^3)
struct Hungarian {
   int N, qi, root;
   vector<vector<ll>>> gain;
   vector<int> xy, yx, p, q, slackx;
   vector<ll> lx. lv. slack:
   vector<bool> S, T;
   void add(int x, int px) {
      S[x] = true, p[x] = px;
      rep(y, N) if (lx[x] + ly[y] - gain[x][y] < slack[y])
          slack[y] = lx[x] + ly[y] - gain[x][y], slackx[y]
      }
   void augment(int x, int y) {
       while (x != -2) {
          yx[y] = x; swap(xy[x], y); x = p[x];
      }
   void improve() {
       S.assign(N, false), T.assign(N, false), p.assign(N,
           -1):
       qi = 0, q.clear();
       rep(x, N) if (xy[x] == -1) {
          q.push_back(root = x), p[x] = -2, S[x] = true;
          break:
      rep(v, N) slack[v] = lx[root] + lv[v] - gain[root][v
           ], slackx[y] = root;
       while (true) {
          while (gi < g.size()) {</pre>
              int x = q[qi++];
              rep(y, N) if (lx[x] + ly[y] == gain[x][y] &&!
                  T[v]) {
                  if (yx[y] == -1) return augment(x, y);
```

```
T[y] = true, q.push_back(yx[y]), add(yx[y
                      ], x);
              }
          11 d = INF:
          rep(y, N) if (!T[y]) d = min(d, slack[y]);
          rep(x, N) if (S[x]) lx[x] -= d;
          rep(y, N) if (T[y]) ly[y] += d;
          rep(y, N) if (!T[y]) slack[y] -= d;
          rep(v, N) if (!T[v] && slack[v] == 0) {
              if (yx[y] == -1) return augment(slackx[y], y);
              T[v] = true:
              if (!S[yx[y]]) q.push_back(yx[y]), add(yx[y],
                   slackx[v]);
          }
      }
   }
   Hungarian(vector<vector<ll>> g)
       : N(g.size()), gain(g), xy(N, -1), yx(N, -1), lx(N, -
           INF).
       ly(N), slack(N), slackx(N) {
       rep(x, N) rep(y, N) lx[x] = max(lx[x], ly[y]);
       rep(i, N) improve();
   }
};
```

5.7 kuhn

```
// get a maximum cardinality matching in a bipartite graph.
// input: adjacency lists.
// output: matching (in 'mt' member).
// runtime: O(V E)
struct Kuhn {
   int N. size:
   vector<vector<int>> G;
   vector<bool> seen:
   vector<int> mt:
   bool visit(int i) {
       if (seen[i]) return false;
       seen[i] = true:
       for (int to : G[i])
           if (mt[to] == -1 || visit(mt[to])) {
              mt[to] = i:
              return true;
```

5.8 lca

```
// calculates the lowest common ancestor for any two nodes
    in O(log N) time,
// with O(N log N) preprocessing
struct Lca {
   int N, K, t = 0;
   vector<vector<int>> U;
   vector<int> L. R:
   Lca() {}
   Lca(vector<vector<int>> &G) : N(G.size()), L(N), R(N) {
      K = N \le 1 ? 0 : 32 - \_builtin_clz(N - 1);
      U.resize(K + 1, vector<int>(N));
       visit(G, 0, 0);
       rep(k, K) rep(u, N) U[k + 1][u] = U[k][U[k][u]];
   void visit(vector<vector<int>> &G, int u, int p) {
      L[u] = t++, U[0][u] = p:
      for (int v : G[u]) if (v != p) visit(G, v, u);
      R[u] = t++:
   }
   bool is anc(int up, int dn) {
       return L[up] <= L[dn] && R[dn] <= R[up];</pre>
   int find(int u, int v) {
       if (is_anc(u, v)) return u;
      if (is_anc(v, u)) return v;
      for (int k = K; k >= 0;)
          if (is anc(U[k][u], v)) k--:
          else u = U[k][u];
       return U[0][u];
   }
```

F is the maximum flow

5.9 maxflow-mincost

// time: O(F V E)

```
O(V E + F E log V) if bellman-ford is replaced by
    johnson
struct Flow {
   struct Edge {
      int u, v;
      11 c, w, f = 0;
   }:
   int N. s. t:
   vector<vector<int>> G;
   vector<Edge> E;
   vector<ll> d, b;
   vector<int> p;
   Flow() {}
   Flow(int N, int s, int t) : N(N), s(s), t(t), G(N) {}
   void add_edge(int u, int v, ll c, ll w) {
      G[u].push_back(E.size());
       E.push_back({u, v, c, w});
      G[v].push_back(E.size());
      E.push_back(\{v, u, 0, -w\});
   // naive distances with bellman-ford: O(V E)
   void calcdists() {
      p.assign(N, -1), d.assign(N, INF), d[s] = 0;
      rep(i, N - 1) rep(ei, E.size()) {
          Edge &e = E[ei];
          11 n = d[e.u] + e.w;
          if (d[e.u] < INF && e.c - e.f > 0 && n < d[e.v])
               d[e.v] = n, p[e.v] = ei;
      }
   }
   // johnsons potentials: O(E log V)
   void calcdists() {
      if (b.empty()) {
          b.assign(N. 0):
          // code below only necessary if there are
               negative costs
          rep(i, N - 1) rep(ei, E.size()) {
              Edge &e = E[ei];
```

```
if (e.f < e.c) b[e.v] = min(b[e.v], b[e.u] + e // maximum flow algorithm.
          }
       p.assign(N, -1), d.assign(N, INF), d[s] = 0;
       priority_queue<pair<11, int>> q;
       q.push({0, s});
       while (!q.empty()) {
           auto [w, u] = q.top();
           q.pop();
           if (d[u] < -w + b[u]) continue:
           for (int ei : G[u]) {
              auto e = E[ei];
              ll n = d[u] + e.w:
              if (e.f < e.c && n < d[e.v]) {</pre>
                  d[e.v] = n, p[e.v] = ei;
                  q.push({b[e.v] - n, e.v});
              }
           }
       }
       b = d;
   ll solve() {
       b.clear();
       11 \text{ ff} = 0;
       while (true) {
           calcdists():
           if (p[t] == -1) break;
           for (int cur = t; p[cur] != -1; cur = E[p[cur]].u
              f = min(f, E[p[cur]].c - E[p[cur]].f);
           for (int cur = t; p[cur] != -1; cur = E[p[cur]].u
              E[p[cur]].f += f, E[p[cur] ^ 1].f -= f;
          ff += f:
       return ff:
   }
};
```

5.10 push-relabel

```
#include "../common.h"
const ll INF = 1e18;
```

```
// to run. use 'maxflow()'.
11
// time: O(V^2 \operatorname{sart}(E)) \leq O(V^3)
// memory: 0(V^2)
struct PushRelabel {
   vector<vector<11>> cap, flow;
   vector<ll> excess;
   vector<int> height;
   PushRelabel() {}
   void resize(int N) { cap.assign(N, vector<11>(N)); }
   // push as much excess flow as possible from u to v.
   void push(int u, int v) {
       11 f = min(excess[u], cap[u][v] - flow[u][v]);
       flow[u][v] += f:
       flow[v][u] -= f;
       excess[v] += f:
       excess[u] -= f:
   // relabel the height of a vertex so that excess flow may
         be pushed.
   void relabel(int u) {
       int d = INT32 MAX:
       rep(v, cap.size()) if (cap[u][v] - flow[u][v] > 0) d
           min(d, height[v]);
       if (d < INF) height[u] = d + 1:</pre>
   // get the maximum flow on the network specified by 'cap'
         with source 's'
   // and sink 't'.
   // node-to-node flows are output to the 'flow' member.
   11 maxflow(int s, int t) {
       int N = cap.size(), M;
       flow.assign(N, vector<11>(N));
       height.assign(N, 0), height[s] = N;
       excess.assign(N. 0). excess[s] = INF:
       rep(i, N) if (i != s) push(s, i);
       vector<int> q;
       while (true) {
          // find the highest vertices with excess
           q.clear(), M = 0;
          rep(i, N) {
              if (excess[i] <= 0 || i == s || i == t)</pre>
                   continue:
```

```
if (height[i] > M) q.clear(), M = height[i];
                 if (height[i] >= M) q.push_back(i);
            }
             if (q.empty()) break;
            // process vertices
             for (int u : a) {
                bool relab = true;
                rep(v, N) {
                     if (excess[u] <= 0) break;</pre>
                     \label{eq:cap_u} \begin{subarray}{ll} if $(cap[u][v] - flow[u][v] > 0 && height[u] \end{subarray}
                          1 > height[v])
                         push(u, v), relab = false:
                }
                 if (relab) {
                     relabel(u):
                     break:
            }
        11 f = 0; rep(i, N) f += flow[i][t]; return f;
};
```

5.11 strongly-connected-components

```
// compute strongly connected components.
// time: O(V + E), memory: O(V)
// after building:
// comp = map from vertex to component (components are
    toposorted, root first, leaf last)
// N = number of components
// G = condensation graph (component DAG)
// byproducts:
// vgi = transposed graph
// order = reverse topological sort (leaf first, root last)
//
// others:
// vn = number of vertices
// vg = original vertex graph
struct Scc {
   int vn, N;
   vector<int> order. comp:
   vector<vector<int>> vg, vgi, G;
   void toposort(int u) {
       if (comp[u]) return;
```

```
comp[u] = -1:
       for (int v : vg[u]) toposort(v);
       order.push_back(u);
   bool carve(int u) {
       if (comp[u] != -1) return false;
       comp[u] = N;
       for (int v : vgi[u]) {
          carve(v);
          if (comp[v] != N) G[comp[v]].push_back(N);
       return true;
   Scc(vector<vector<int>> &g) : vn(g.size()), vg(g), comp(
        vn), vgi(vn), G(vn), N(0) {
       rep(u, vn) toposort(u):
       rep(u, vn) for (int v : vg[u]) vgi[v].push_back(u);
       invrep(i, vn) N += carve(order[i]);
};
```

5.12 two-sat

```
// calculate the solvability of a system of logical
    equations, where every equation is of the form 'a or b
// 'neg': get negation of 'u'
// 'then': 'u' implies 'v'
// 'any': 'u' or 'v'
// 'set': 'u' is true
// after 'solve' (O(V+E)) returns true, 'sol' contains one
    possible solution.
// determining all solutions is O(V*E) hard (requires
    computing reachability in a DAG).
struct TwoSat {
   int N; vector<vector<int>> G;
   Scc scc; vector<bool> sol;
   TwoSat(int n) : N(n), G(2 * n), sol(n) {}
   TwoSat() {}
   int neg(int u) { return (u + N) % (2 * N); }
   void then(int u, int v) { G[u].push_back(v), G[neg(v)].
        push_back(neg(u)); }
   void any(int u, int v) { then(neg(u), v); }
   void set(int u) { G[neg(u)].push_back(u); }
```

f 6 implementation

6.1 SegmentTreeBeats

```
struct Node {
   11 \text{ s, } mx1, mx2, mxc, mn1, mn2, mnc, }1z = 0;
   Node(): s(0), mx1(LLONG_MIN), mx2(LLONG_MIN), mxc(0),
        mn1(LLONG_MAX), mn2(LLONG_MAX), mnc(0) {}
   Node(ll x): s(x), mx1(x), mx2(LLONG_MIN), mxc(1), mn1(x)
        . mn2(LLONG MAX). mnc(1) {}
   Node(const Node &a, const Node &b) {
       // add
       s = a.s + b.s:
       // min
       if (a.mx1 > b.mx1) mx1 = a.mx1, mxc = a.mxc, mx2 =
           max(b.mx1, a.mx2);
       if (a.mx1 < b.mx1) mx1 = b.mx1, mxc = b.mxc, mx2 =
            max(a.mx1, b.mx2):
       if (a.mx1 == b.mx1) mx1 = a.mx1, mxc = a.mxc + b.mxc,
             mx2 = max(a.mx2, b.mx2):
       if (a.mn1 < b.mn1) mn1 = a.mn1, mnc = a.mnc, mn2 =</pre>
           min(b.mn1, a.mn2):
       if (a.mn1 > b.mn1) mn1 = b.mn1, mnc = b.mnc, mn2 =
            min(a.mn1, b.mn2);
       if (a.mn1 == b.mn1) mn1 = a.mn1, mnc = a.mnc + b.mnc,
             mn2 = min(a.mn2, b.mn2);
   }
};
// 0 - indexed / inclusive - inclusive
template <class node>
struct STB {
   vector<node> st: int n:
   void build(int u, int i, int j, vector<node> &arr) {
       if (i == j) {
          st[u] = arr[i];
```

```
return:
   int m = (i + j) / 2, l = u * 2 + 1, r = u * 2 + 2;
   build(l, i, m, arr), build(r, m + 1, i, arr);
   st[u] = node(st[1], st[r]);
void push_add(int u, int i, int j, ll v) {
   st[u].s += (j - i + 1) * v;
   st[u].mx1 += v, st[u].mn1 += v, st[u].lz += v;
   if (st[u].mx2 != LLONG_MIN) st[u].mx2 += v:
   if (st[u].mn2 != LLONG MAX) st[u].mn2 += v:
void push_max(int u, ll v, bool l) { // for min op
   if (v >= st[u].mx1) return;
   st[u].s -= st[u].mx1 * st[u].mxc:
   st[u].mx1 = v;
   st[u].s += st[u].mx1 * st[u].mxc:
   if (1) st[u].mn1 = st[u].mx1;
   else if (v \le st[u].mn1) st[u].mn1 = v:
   else if (v < st[u].mn2) st[u].mn2 = v;
void push_min(int u, ll v, bool l) { // for max op
   if (v <= st[u].mn1) return;</pre>
   st[u].s -= st[u].mn1 * st[u].mnc;
   st[u].mn1 = v;
   st[u].s += st[u].mn1 * st[u].mnc;
   if (1) st[u].mx1 = st[u].mn1:
   else if (v \ge st[u].mx1) st[u].mx1 = v:
   else if (v > st[u].mx2) st[u].mx2 = v;
void push(int u, int i, int j) {
   if (i == j) return;
   // add
   int m = (i + j) / 2, l = u * 2 + 1, r = u * 2 + 2;
   push_add(1, i, m, st[u].lz);
   push add(r, m + 1, i, st[u].lz):
   st[u].lz = 0:
   // min
   push_max(1, st[u].mx1, i == m);
   push_max(r, st[u].mx1, m + 1 == j);
   push_min(1, st[u].mn1, i == m);
   push min(r, st[u].mn1, m + 1 == r):
node query(int a, int b, int u, int i, int j) {
   if (b < i | | i < a) return node();</pre>
   if (a <= i && j <= b) return st[u];</pre>
   push(u, i, j);
   int m = (i + j) / 2, l = u * 2 + 1, r = u * 2 + 2;
```

```
return node(query(a, b, 1, i, m), query(a, b, r, m +
        1, i));
}
void update add(int a. int b. 11 v. int u. int i. int i)
   if (b < i | | i < a) return:
   if (a <= i && j <= b) {</pre>
       push_add(u, i, j, v);
       return:
   push(u, i, i):
   int m = (i + i) / 2, l = u * 2 + 1, r = u * 2 + 2;
   update_add(a, b, v, l, i, m);
   update_add(a, b, v, r, m + 1, j);
   st[u] = node(st[l], st[r]):
void update min(int a, int b, ll v, int u, int i, int i)
   if (b < i \mid | i < a \mid | v >= st[u].mx1) return:
   if (a <= i && i <= b && v > st[u].mx2) {
       push_max(u, v, i == j);
       return:
   push(u, i, j);
   int m = (i + j) / 2, l = u * 2 + 1, r = u * 2 + 2;
   update_min(a, b, v, l, i, m);
   update_min(a, b, v, r, m + 1, j);
   st[u] = node(st[l], st[r]):
void update max(int a, int b, ll v, int u, int i, int i)
   if (b < i || j < a || v <= st[u].mn1) return;</pre>
   if (a <= i && j <= b && v < st[u].mn2) {</pre>
       push_min(u, v, i == j);
       return:
   push(u, i, j);
   int m = (i + j) / 2, l = u * 2 + 1, r = u * 2 + 2;
   update_max(a, b, v, 1, i, m);
   update_max(a, b, v, r, m + 1, j);
   st[u] = node(st[1], st[r]);
}
STB(vector<node> &v, int N) : n(N), st(N * 4 + 5) { build
     (0, 0, n - 1, v); 
node querv(int a, int b) { return querv(a, b, 0, 0, n -
    1); }
void update_add(int a, int b, ll v) { update_add(a, b, v,
     0, 0, n - 1); }
```

6.2 Treap

```
mt19937 gen(chrono::high_resolution_clock::now().
    time since epoch().count());
// 101 Implicit Treap //
struct Node {
   int p, sz = 0, v, acc, l = -1, r = -1:
   Node(): v(0), acc(0) {}
   Node(int x) : p(gen()), sz(1), v(x), acc(x) {}
   void recalc(const Node &a. const Node &b) {
       sz = a.sz + b.sz + 1;
       acc = v + a.acc + b.acc;
};
template <class node>
struct Treap {
   vector<node> t:
   int n, r = -1;
   node get(int u) { return u != -1 ? t[u] : node(); }
   void recalc(int u) { t[u].recalc(get(t[u].1), get(t[u].r)
       ): }
   int merge(int 1, int r) {
       if (min(1, r) == -1) return 1 != -1 ? 1 : r;
       int ans = (t[1].p < t[r].p) ? 1 : r:
       if (ans == 1) t[1].r = merge(t[1].r, r), recalc(1);
       if (ans == r) t[r].l = merge(l, t[r].l), recalc(r);
       return ans:
   pii split(int u, int id) {
       if (u == -1) return {-1, -1};
       int szl = get(t[u].1).sz:
      if (szl >= id) {
          pii ans = split(t[u].1, id);
          t[u].1 = ans.ss;
          recalc(u):
          return {ans.ff, u};
       pii ans = split(t[u].r, id - szl - 1);
       t[u].r = ans.ff;
```

```
recalc(u):
       return {u, ans.ss};
   Treap(vi &v) : n(sz(v)) {
       for (int i = 0; i < n; i++) t.eb(v[i]), r = merge(r)
}:
// Complete Implicit Treap with Lazy propagation //
struct Node {
    int p, sz = 0, v, acc, l = -1, r = -1, par = -1, lzv = 0;
    bool lz = false, f = false;
    Node() : v(0), acc(0) {}
    Node(int x) : p(gen()), sz(1), v(x), acc(x) {}
    void recalc(const Node &a, const Node &b) {
       sz = a.sz + b.sz + 1:
       acc = v + a.acc + b.acc:
    void upd_lazy(int x) { lz = 1, lzv += x; }
    void lazy() { v += lzv, acc += sz * lzv, lz = 0, lzv = 0;
    void flip() { swap(1, r), f = 0; }
};
template <class node>
struct Treap {
    vector<node> t:
    int n, r = -1;
    node get(int u) { return u != -1 ? t[u] : node(): }
    void recalc(int u) {
       int 1 = t[u].1, r = t[u].r:
       push(1). push(r). flip(1). flip(r):
       t[u].recalc(get(1), get(r));
    void push(int u) {
       if (u == -1 || !t[u].lz) return;
       int 1 = t[u].1, r = t[u].r:
       if (1 != -1) t[1].upd_lazy(t[u].lzv);
       if (r != -1) t[r].upd_lazy(t[u].lzv);
       t[u].lazy();
    void flip(int u) {
       if (u == -1 || !t[u].f) return;
       int 1 = t[u].1, r = t[u].r:
       if (1 != -1) t[1].f ^= 1:
       if (r != -1) t[r].f ^= 1;
```

```
t[u].flip():
}
int merge(int 1, int r) {
   if (min(l, r) == -1) return l != -1 ? l : r:
   push(1), push(r), flip(1), flip(r);
   int ans = (t[1].p < t[r].p) ? 1 : r:</pre>
   if (ans == 1) t[1].r = merge(t[1].r, r), recalc(1);
   if (ans == r) t[r].l = merge(l, t[r].l), recalc(r);
   if (t[ans].l != -1) t[t[ans].l].par = ans; // only if
         parent needed
   if (t[ans].r!= -1) t[t[ans].r].par = ans: // only if
         parent needed
   return ans;
pii split(int u, int id) {
   if (u == -1) return {-1, -1};
   push(u):
   flip(u);
   int szl = get(t[u].1).sz:
   if (szl >= id) {
       pii ans = split(t[u].1, id);
       if (ans.ss != -1) t[ans.ss].par = u; // only if
            parent needed
       if (ans.ff !=-1) t[ans.ff].par = -1; // only if
            parent needed
       t[u].1 = ans.ss;
       recalc(u):
       return {ans.ff. u}:
   pii ans = split(t[u].r, id - szl - 1);
   if (ans.ff != -1) t[ans.ff].par = u; // only if
        parent needed
   if (ans.ss != -1) t[ans.ss].par = -1; // only if
        parent needed
   t[u].r = ans.ff:
   recalc(u):
   return {u, ans.ss};
int update(int u, int 1, int r, int v) {
   pii a = split(u, 1), b = split(a.ss, r - 1 + 1);
   t[b.ff].upd lazv(v):
   return merge(a.ff, merge(b.ff, b.ss));
void print(int u) {
   if (u == -1) return;
   push(u), flip(u);
   print(t[u].1);
   cout << t[u].v << ' ';
   print(t[u].r);
```

6.3 bit-tricks

```
y = x & (x-1) // Turn off rightmost 1bit
y = x & (-x) // Isolate rightmost 1bit
y = x \mid (x-1) // Right propagate rightmost 1bit(fill in 1s)
v = x \mid (x+1) // Turn on rightmost Obit
y = ~x & (x+1) // Isolate rightmost Obit
// If x is of long type, use __builtin_popcountl(x)
// If x is of long long type, use __builtin_popcountll(x)
// 1. Counts the number of ones(set bits) in an integer.
__builtin_popcount(x)
// 2. Checks the Parity of a number. Returns true(1) if the
// number has odd number of set bits, else it returns
// false(0) for even number of set bits.
__builtin_parity(x)
// 3. Counts the leading number of zeros of the integer.
__builtin_clz(x)
// 4. Counts the trailing number of zeros of the integer.
__builtin_ctz(x)
// 5. Returns 1 + the index of the least significant 1-bit.
__builtin_ffs(x) // If x == 0, returns 0.
// Iterate over non empty subsets of bitmask
for(int s=m;s;s=(s-1)&m) // Decreasing order
for (int s=0:s=s-m&m:) // Increasing order
```

6.4 dsu

```
void rollback() {
    auto a = S.back(); S.pop_back();
    if (a[0] != -1) p[a[0]] = a[2], p[a[1]] = a[3];
};
```

6.5 dynamic-connectivity

```
struct DC {
   int n: Dsu D:
   vector<vector<pair<int, int>>> t;
   DC(int N) : n(N), D(N), t(2 * N) {}
   // add edge p to all times in interval [1, r]
   void upd(int 1, int r, pair<int, int> p) {
       for (1 += n, r += n; 1 < r; 1 >>= 1, r >>= 1) {
           if (1 & 1) t[1++].push_back(p);
           if (r & 1) t[--r].push_back(p);
   void process(int u = 1) { // process all queries
       for (auto &e : t[u]) D.unite(e.first, e.second):
       if (u >= n) {
           // do stuff with D at time u - n
       } else process(2 * u), process(2 * u + 1);
       for (auto &e : t[u]) D.rollback():
};
```

6.6 mo

6.7 ordered-set

6.8 persistent-segment-tree-lazy

```
int build(int vl. int vr) {
   if (vr - vl == 1) a.push_back({gneut(), uneut()}); //
         node construction
   else {
       int vm = (vl + vr) / 2, l = build(vl, vm), r =
            build(vm, vr):
       a.push_back({merge(a[1].x, a[r].x), uneut(), 1, r
            }); // query merge
   return a.size() - 1;
T query(int 1, int r, int v, int v1, int vr, T acc) {
   if (1 >= vr || r <= vl) return qneut();</pre>
        // query neutral
   if (1 <= v1 && r >= vr) return apply(a[v].x, acc, v1,
         vr): // update op
   acc = accum(acc, a[v].lz);
        // update merge
   int vm = (v1 + vr) / 2:
   return merge(query(1, r, a[v].1, v1, vm, acc), query(
        1, r, a[v].r, vm, vr, acc)); // query merge
int update(int 1, int r, T x, int v, int v1, int vr) {
   if (1 >= vr || r <= vl || r <= 1) return v;</pre>
   a.push_back(a[v]);
   v = a.size() - 1:
   if (1 <= v1 && r >= vr) {
       a[v].x = apply(a[v].x, x, vl, vr); // update op
       a[v].lz = accum(a[v].lz, x); // update merge
       int vm = (vl + vr) / 2:
       a[v].1 = update(1, r, x, a[v].1, v1, vm);
       a[v].r = update(1, r, x, a[v].r, vm, vr);
       a[v].x = merge(a[a[v].1].x, a[a[v].r].x); //
            query merge
   }
    return v;
7
Pst1() {}
Pstl(int N) : N(N) { head.push back(build(0, N)); }
T query(int t, int 1, int r) {
   return query(1, r, head[t], 0, N, uneut()); // update
         neutral
int update(int t, int 1, int r, T x) {
```

6.9 persistent-segment-tree

// usage:

```
// Pst<Node<11>> pst;
// pst = {N}:
// int newtime = pst.update(time, index, value);
// Node<11> result = pst.query(newtime, left, right);
template <class T>
struct Node {
   T x:
   int 1 = -1, r = -1;
   Node() : x(0) {}
   Node(T x) : x(x) \{ \}
   Node (Node a. Node b. int l = -1, int r = -1) : x(a.x + b.
        x), 1(1), r(r) {}
};
template <class U>
struct Pst {
   int N:
   vector<U> a:
   vector<int> head;
   int build(int vl. int vr) {
       if (vr - vl == 1) a.push back(U()):
       else {
           int vm = (vl + vr) / 2, l = build(vl, vm),
              r = build(vm, vr);
           a.push_back(U(a[1], a[r], 1, r));
       return a.size() - 1;
   U query(int 1, int r, int v, int v1, int vr) {
       if (1 >= vr || r <= vl) return U():
       if (1 <= v1 && r >= vr) return a[v];
       int vm = (v1 + vr) / 2;
       return U(query(1, r, a[v].1, v1, vm),
               query(1, r, a[v].r, vm, vr));
   }
   int update(int i, U x, int v, int vl, int vr) {
```

```
a.push back(a[v]):
   v = a.size() - 1:
   if (vr - vl == 1) a[v] = x:
       int vm = (v1 + vr) / 2;
       if (i < vm) a[v].l = update(i, x, a[v].l, vl, vm)</pre>
       else a[v].r = update(i, x, a[v].r, vm, vr);
       a[v] = U(a[a[v].1], a[a[v].r], a[v].1, a[v].r);
   return v:
}
Pst() {}
Pst(int N) : N(N) { head.push_back(build(0, N)); }
U querv(int t, int 1, int r) {
   return query(1, r, head[t], 0, N);
int update(int t, int i, U x) {
   return head.push_back(update(i, x, head[t], 0, N)),
        head.size() - 1:
```

6.10 segment-tree-lazy

```
template <class T>
struct Stl {
   int n;
   vector<T> a. b:
   T qneut() { return -2e9; }
   T uneut() { return 0: }
   T merge(T x, T y) { return max(x, y); }
   void upd(int v, T x, int 1, int r)
      \{ a[v] += x, b[v] += x; \}
   Stl(int n = 0) : n(n), a(4 * n, qneut()),
      b(4 * n. uneut()) {}
   void push(int v, int vl, int vm, int vr) {
      upd(2 * v, b[v], vl, vm);
      upd(2 * v + 1, b[v], vm, vr);
      b[v] = uneut():
   }
   T query(int 1, int r, int v=1, int vl=0, int vr=1e9) {
      vr = min(vr. n):
```

```
if (1 \le v1 \&\& r \ge vr) return a[v]:
       if (1 >= vr || r <= vl) return gneut();</pre>
       int vm = (v1 + vr) / 2;
       push(v. vl. vm. vr):
       return merge(query(1, r, 2 * v, v1, vm),
           querv(1, r, 2 * v + 1, vm, vr)):
   void update(int 1, int r, T x, int v = 1, int vl = 0,
           int vr = 1e9) {
       vr = min(vr, n):
       if (1 >= vr || r <= vl || r <= 1) return;</pre>
       if (1 <= v1 && r >= vr) upd(v, x, v1, vr);
           int vm = (v1 + vr) / 2:
           push(v, vl, vm, vr);
           update(1, r, x, 2 * v, v1, vm);
           update(1, r, x, 2 * v + 1, vm, vr);
           a[v] = merge(a[2 * v], a[2 * v + 1]):
};
```

6.11 segment-tree

```
struct St {
   int n;
   vector<1l> a;

ll neut() { return 0; }
   ll merge(ll x, ll y) { return x + y; }

St(int n = 0) : n(n), a(2 * n, neut()) {}

ll query(int l, int r) {
    ll x = neut(), y = neut();
    for (1 += n, r += n; 1 < r; 1 /= 2, r /= 2) {
        if (1 & 1) x = merge(x, a[1++]);
        if (r & 1) y = merge(a[--r], y);
    }

   return merge(x, y);
}

void update(int i, ll x) {
   for (a[i += n] = x; i /= 2;)
        a[i] = merge(a[2 * i], a[2 * i + 1]);
};</pre>
```

6.12 sparse-table

```
template <class T>
struct Sparse {
   T op(T a, T b) { return max(a, b); }
   vector<vector<T>> st;
   Sparse() {}
   Sparse(vector<T> a) : st{a} {
       int N = st[0].size();
       int npot = N <= 1 ? 1 : 32 - __builtin_clz(N);</pre>
       st.resize(npot);
       repx(i, 1, npot) rep(j, N + 1 - (1 << i))
       st[i].push back(
           op(st[i-1][j], st[i-1][j+(1 << (i-1))])
       ); // query op
   T querv(int 1, int r) { // range must be nonempty!
       int i = 31 - \_builtin\_clz(r - 1);
       return op(st[i][1], st[i][r - (1 << i)]); // queryop</pre>
};
```

6.13 unordered-map

```
// hackproof rng
static mt19937 rng(chrono::steady_clock::now().
    time_since_epoch().count());
// deterministic rng
uint64 t splitmix64(uint64 t *x) {
   uint64_t z = (*x += 0x9e3779b97f4a7c15);
   z = (z ^ (z >> 30)) * 0xbf58476d1ce4e5b9:
   z = (z ^ (z >> 27)) * 0x94d049bb133111eb;
   return z^(z >> 31);
}
// hackproof unordered map hash
struct Hash {
    size_t operator()(const 11 &x) const {
       static const uint64 t RAND =
           chrono::steady_clock::now().time_since_epoch().
               count();
       uint64 t z = x + RAND + 0x9e3779b97f4a7c15:
       z = (z ^ (z >> 30)) * 0xbf58476d1ce4e5b9;
       z = (z ^ (z >> 27)) * 0x94d049bb133111eb;
       return z^(z >> 31):
```

```
};

// hackproof unordered_map
template <class T, class U>
using umap = unordered_map<T, U, Hash>;

// hackproof unordered_set
template <class T>
using uset = unordered_set<T, Hash>;
```

7 imprimible

8 math

8.1 arithmetic

```
inline int floor_log2(int n) { return n <= 1 ? 0 : 31 -</pre>
     builtin clz(n): }
inline int ceil_log2(int n) { return n <= 1 ? 0 : 32 -</pre>
     builtin clz(n - 1): }
inline 11 floordiv(11 a, 11 b) {
   return a / b - ((a ^ b) < 0 && a % b);
inline ll ceildiv(ll a, ll b) {
   return a / b + ((a ^ b) >= 0 && a % b):
ll binexp(ll a, ll e) {
   ll res = 1: // neutral element
   while (e) {
       if (e & 1) res = res * a; // multiplication
                               // multiplication
       a = a * a:
       e >>= 1:
   return res;
```

8.2 crt

```
pair<11, 11> solve_crt(const vector<pair<11, 11>> &eqs) {
    11 a0 = eqs[0].first, p0 = eqs[0].second;
    repx(i, 1, eqs.size()) {
        11 a1 = eqs[i].first, p1 = eqs[i].second;
        11 k1, k0;
        11 d = ext_gcd(p1, p0, k1, k0);
```

```
a0 -= a1;

if (a0 % d != 0) return {-1, -1};

p0 = p0 / d * p1;

a0 = a0 / d * k1 % p0 * p1 % p0 + a1;

a0 = (a0 % p0 + p0) % p0;

}

return {a0, p0};
```

8.3 discrete-log

```
// discrete logarithm log a(b).
// solve b \hat{x} = a \pmod{M} for the smallest x.
// returns -1 if no solution is found.
// time: O(sqrt(M))
11 dlog(ll a, ll b, ll M) {
   11 k = 1. s = 0:
   while (true) {
       11 g = \_gcd(b, M);
       if (g <= 1) break;</pre>
       if (a == k) return s;
       if (a % g != 0) return -1;
       a /= g, M /= g, s += 1, k = b / g * k % M;
   ll N = sqrt(M) + 1;
   umap<11. 11> r:
   rep(q, N + 1) {
       r[a] = q;
       a = a * b % M:
   ll bN = binexp(b, N, M), bNp = k;
   repx(p, 1, N + 1) {
       bNp = bNp * bN % M;
       if (r.count(bNp)) return N * p - r[bNp] + s;
   return -1;
```

8.4 fft.

```
using cd = complex<double>;
const double PI = acos(-1);

// compute the DFT of a power-of-two-length sequence.
```

```
// if 'inv' is true, computes the inverse DFT.
void fft(vector<cd> &a, bool inv) {
   int N = a.size(), k = 0, b;
   assert(N == 1 << builtin ctz(N)):</pre>
   repx(i, 1, N) {
       for (b = N >> 1; k & b;) k ^= b, b >>= 1;
       if (i < (k ^= b)) swap(a[i], a[k]);</pre>
   for (int 1 = 2: 1 <= N: 1 <<= 1) {
       double ang = 2 * PI / 1 * (inv ? -1 : 1):
       cd wl(cos(ang), sin(ang));
       for (int i = 0: i < N: i += 1) {</pre>
          cd w = 1:
           rep(i, 1 / 2) {
              cd u = a[i + j], v = a[i + j + 1 / 2] * w;
              a[i + j] = u + v;
              a[i + i + 1 / 2] = u - v:
              w *= wl:
          }
       }
   if (inv) rep(i, N) a[i] /= N;
const 11 MOD = 998244353, ROOT = 15311432:
// const 11 MOD = 2130706433, ROOT = 1791270792;
// const 11 MOD = 922337203673733529711, ROOT =
    532077456549635698311:
void find root of unitv(ll M) {
   11 c = M - 1, k = 0;
   while (c \% 2 == 0) c /= 2, k += 1:
   // find proper divisors of M - 1
   vector<ll> divs:
   for (11 d = 1; d < c; d++) {
       if (d * d > c) break:
       if (c \% d == 0) rep(i, k + 1) divs.push back(d << i):
   rep(i, k) divs.push back(c << i):
   // find any primitive root of M
   11 G = -1:
   repx(g, 2, M) {
       bool ok = true:
       for (ll d : divs) ok &= (binexp(g, d, M) != 1);
       if (ok) {
```

```
G = g;
          break;
      }
   assert(G != -1);
   ll w = binexp(G, c, M);
   cerr << "M = c * 2^k + 1" << endl;
   cerr << " M = " << M << endl:
   cerr << " c = " << c << endl;
   cerr << " k = " << k << endl:
   cerr << " w^(2^k) == 1" << endl:
   cerr << " w = g^{(M-1)/2k} = g^c << endl;
   cerr << " g = " << G << endl;</pre>
   cerr << " w = " << w << endl:
// compute the DFT of a power-of-two-length sequence, modulo
     a special prime
// number with an Nth root of unity, where N is the length
    of the sequence.
void ntt(vector<11> &a. bool inv) {
   vector<11> wn:
   for (11 p = ROOT; p != 1; p = p * p % MOD) wn.push_back(p
   int N = a.size(), k = 0, b;
   assert(N == 1 << __builtin_ctz(N) && N <= 1 << wn.size())
   rep(i, N) a[i] = (a[i] % MOD + MOD) % MOD:
   repx(i, 1, N) {
      for (b = N >> 1: k & b:) k ^= b, b >>= 1:
       if (i < (k ^= b)) swap(a[i], a[k]);</pre>
   for (int 1 = 2: 1 <= N: 1 <<= 1) {
      11 wl = wn[wn.size() - __builtin_ctz(1)];
      if (inv) wl = multinv(wl, MOD);
      for (int i = 0: i < N: i += 1) {
          11 w = 1:
          repx(i, 0, 1 / 2)  {
             11 u = a[i + j], v = a[i + j + 1 / 2] * w %
              a[i + j] = (u + v) \% MOD;
              a[i + j + 1 / 2] = (u - v + MOD) \% MOD;
              w = w * w1 \% MOD:
          }
```

8.5 gauss

```
const double EPS = 1e-9;
// solve a system of equations.
// complexity: O(\min(N, M) * N * M)
11
// 'a' is a list of rows
// the last value in each row is the result of the equation
// return values:
// 0 -> no solutions
// 1 -> unique solution, stored in 'ans'
// -1 -> infinitely many solutions, one of which is stored
    in 'ans'
// UNTESTED
int gauss(vector<vector<double>> a, vector<double> &ans) {
   int N = a.size(), M = a[0].size() - 1:
   vector<int> where(M, -1);
   for (int j = 0, i = 0; j < M && i < N; j++) {
       int sel = i:
       repx(k, i, N) if (abs(a[k][i]) > abs(a[sel][i])) sel
       if (abs(a[sel][i]) < EPS) continue;</pre>
       repx(k, j, M + 1) swap(a[sel][k], a[i][k]);
       where[i] = i:
       rep(k, N) if (k != i) {
           double c = a[k][j] / a[i][j];
           repx(1, j, M + 1) a[k][1] -= a[i][1] * c;
       }
       i++;
   ans.assign(M, 0);
```

8.6 matrix

```
using T = 11:
struct Mat {
   int N. M:
   vector<vector<T>> v:
   Mat(int n, int m) : N(n), M(m), v(N, vector<T>(M)) {}
   Mat(int n) : Mat(n, n) { rep(i, N) v[i][i] = 1; }
   vector<T> &operator[](int i) { return v[i]; }
   Mat operator*(Mat &r) {
       assert(M == r.N):
      int n = N, m = r.M, p = M;
      Mat a(n, m);
      rep(i, n) rep(j, m) {
                                                       11
          a[i][j] = T();
               neutral
          rep(k, p) a[i][j] = a[i][j] + v[i][k] * r[k][j];
               // mul, add
      }
       return a;
   Mat binexp(ll e) {
       assert(N == M):
      Mat a = *this, res(N); // neutral
      while (e) {
          if (e & 1) res = res * a: // mul
                                  // mul
          a = a * a;
          e >>= 1;
      }
       return res;
   friend ostream &operator<<(ostream &s, Mat &a) {</pre>
```

```
rep(i, a.N) {
      rep(j, a.M) s << a[i][j] << " ";
      s << endl;
    }
    return s;
}</pre>
```

8.7 mobius

```
short mu[MAXN] = {0,1};
void mobius(){
  repx(i,1,MAXN)if(mu[i])for(int j=i+i;j<MAXN;j+=i)mu[j]-=mu[
        i];
}</pre>
```

$8.8 \mod$

```
11 binexp(ll a, ll e, ll M) {
   assert(e >= 0):
   ll res = 1 % M:
    while (e) {
       if (e & 1) res = res * a % M;
       a = a * a % M:
       e >>= 1:
   }
   return res:
ll multinv(ll a, ll M) { return binexp(a, M - 2, M); }
// calculate gcd(a, b).
// also, calculate x and y such that:
// a * x + b * y == gcd(a, b)
// time: O(log min(a, b))
// (ignoring complexity of arithmetic)
ll ext_gcd(ll a, ll b, ll &x, ll &y) {
   if (b == 0) {
       x = 1, y = 0;
       return a;
   ll d = ext_gcd(b, a \% b, y, x);
   v = a / b * x;
   return d:
```

```
// compute inverse with any M.
// a and M must be coprime for inverse to exist!
11 multinv_euc(l1 a, l1 M) {
   11 x. v:
   ext_gcd(a, M, x, y);
   return x:
// multiply two big numbers (~10^18) under a large modulo,
    without resorting to
// bigints.
11 bigmul(l1 x, l1 v, l1 M) {
   11 z = 0;
   while (v) {
       if (y \& 1) z = (z + x) \% M;
       x = (x << 1) \% M, y >>= 1;
   return z;
// all modular inverses from 1 to inv.size()-1
void multinv all(vector<ll> &inv) {
   inv[1] = 1;
   repx(i, 2, inv.size())
       inv[i] = MOD - (MOD / i) * inv[MOD % i] % MOD;
// change to int128 if checking numbers over 10^9
bool isprime(ll n) {
   if (n < 2 | | n % 6 % 4 != 1) return n - 2 < 2:
   11 A[] = \{2,325,9375,28178,450775,9780504,1795265022\};
   ll s = \_builtin\_ctzll(n - 1), d = n >> s;
   for (int a : A) {
      11 p = binexp(a, d, n), i = s;
       while (p != 1 && p != n - 1 && a % n && i--) p = p *
           р % n:
       if (p != n - 1 && i != s) return 0;
   return 1;
struct Mod {
   int a:
   static const int M = 1e9 + 7;
   Mod(ll aa) : a((aa % M + M) % M) {}
   Mod operator+(Mod rhs) const { return (a + rhs.a) % M; }
   Mod operator-(Mod rhs) const { return (a - rhs.a + M) % M
```

```
Mod operator-() const { return Mod(0) - *this; }
   Mod operator*(Mod rhs) const { return (11)a * rhs.a % M;
   Mod operator+=(Mod rhs) { return *this = *this + rhs; }
   Mod operator==(Mod rhs) { return *this = *this - rhs; }
   Mod operator*=(Mod rhs) { return *this = *this * rhs; }
   Mod bigmul(ll big) const { return ::bigmul(a, big, M); }
   Mod binexp(ll e) const { return ::binexp(a, e, M); }
   // Mod multinv() const { return ::multinv(a, M): } //
        prime M
   Mod multinv() const { return ::multinv_euc(a, M); } //
        possibly composite M
};
// dynamic modulus
struct DMod {
   int a. M:
   DMod(ll aa, ll m) : M(m), a((aa % m + m) % m) {}
   DMod operator+(DMod rhs) const { return {(a + rhs.a) % M,
   DMod operator-(DMod rhs) const { return {(a - rhs.a + M)
        % M. M}: }
   DMod operator-() const { return DMod(0, M) - *this; }
   DMod operator*(DMod rhs) const { return {(11)a * rhs.a %
        M. M:: }
   DMod operator+=(DMod rhs) { return *this = *this + rhs; }
   DMod operator-=(DMod rhs) { return *this = *this - rhs; }
   DMod operator*=(DMod rhs) { return *this = *this * rhs; }
   DMod bigmul(ll big) const { return {::bigmul(a, big, M),
        M}: }
   DMod binexp(ll e) const { return {::binexp(a, e, M), M};
   DMod multinv() const { return {::multinv(a, M), M}: } //
   // DMod multinv() const { return {::multinv_euc(a, M), M
        }; } // possibly composite M
};
```

8.9 primes

```
// counts the divisors of a positive integer in O(\operatorname{sqrt}(n))
```

```
ll count divisors(ll x) {
    11 \text{ divs} = 1. i = 2:
    for (11 divs = 1, i = 2; x > 1; i++) {
        if (i * i > x) {
           divs *= 2:
           break:
        for (11 d = divs; x % i == 0; x /= i) divs += d;
    return divs;
// gets the prime factorization of a number in O(sqrt(n))
vector<pair<11, int>> factorize(11 x) {
    vector<pair<11, int>> f;
    for (11 k = 2; x > 1; k++) {
       if (k * k > x) {
           f.push_back({x, 1});
           break:
       int n = 0;
        while (x \% k == 0) x /= k, n++;
        if (n > 0) f.push_back(\{k, n\});
    return f;
 // iterate over all divisors of a number.
// divisor count upper bound: n^(1.07 / ln ln n)
template <class OP>
void divisors(ll x, OP op) {
    auto facts = factorize(x):
    vector<int> f(facts.size());
    while (true) {
       11 v = 1:
        rep(i, f.size()) rep(j, f[i]) y *= facts[i].first;
        op(y);
        int i:
        for (i = 0: i < f.size(): i++) {</pre>
           f[i] += 1:
           if (f[i] <= facts[i].second) break:</pre>
           f[i] = 0;
        if (i == f.size()) break:
    }
```

```
// computes euler totative function phi(x), counting the
    amount of integers in
// [1, x] that are coprime with x.
// time: O(sqrt(x))
ll phi(ll x) {
   11 phi = 1, k = 2;
   for (; x > 1; k++) {
       if (k * k > x) {
          phi *= x - 1;
           break:
       }
      11 k1 = 1, k0 = 0;
       while (x \% k == 0) x /= k, k0 = k1, k1 *= k;
       phi *= k1 - k0:
   return phi;
// isprime is in mod.cpp
```

8.10 simplex

```
// Solves a general linear maximization problem: maximize $c
     ^T x$ subject to $Ax \le b$, $x \ge 0$.
// Returns -inf if there is no solution, inf if there are
    arbitrarily good solutions, or the maximum value of $c^
    T x$ otherwise.
// The input vector is set to an optimal $x$ (or in the
    unbounded case, an arbitrary solution fulfilling the
    constraints).
// Numerical stability is not guaranteed. For better
    performance, define variables such that x = 0 is
    viable.
// Usage:
// vvd A = \{\{1,-1\}, \{-1,1\}, \{-1,-2\}\};
// vd b = \{1.1.-4\}, c = \{-1.-1\}, x:
// T val = LPSolver(A, b, c).solve(x);
// Time: O(NM * \t ), where a pivot may be e.g. an edge
     relaxation. O(2^n) in the general case.
#include "../common.h"
typedef double T; // long double, Rational, double + mod<P
typedef vector<T> vd;
typedef vector<vd> vvd;
const T eps = 1e-8, inf = 1 / .0;
```

```
#define MP make pair
#define ltj(X) \
   if (s == -1 \mid | MP(X[j], N[j]) < MP(X[s], N[s])) s = j
struct LPSolver {
   int m. n:
   vector<int> N, B;
   vvd D:
   LPSolver(const vvd &A, const vd &b, const vd &c) : m(b.
        size()), n(c.size()), N(n + 1), B(m), D(m + 2), vd(n)
        + 2)) {
       rep(i, m) rep(j, n) D[i][j] = A[i][j];
       rep(i, m) {
          B[i] = n + i;
          D[i][n] = -1;
          D[i][n + 1] = b[i]:
       rep(j, n) {
          N[i] = i:
          D[m][i] = -c[i];
       N[n] = -1;
       D[m + 1][n] = 1;
   void pivot(int r, int s) {
      T *a = D[r].data(), inv = 1 / a[s]:
       rep(i, m + 2) if (i != r \&\& abs(D[i][s]) > eps) {
          T *b = D[i].data(), inv2 = b[s] * inv:
          repx(j, 0, n + 2) b[j] -= a[j] * inv2;
          b[s] = a[s] * inv2;
       rep(j, n + 2) if (j != s) D[r][j] *= inv;
       rep(i, m + 2) if (i != r) D[i][s] *= -inv;
       D[r][s] = inv:
       swap(B[r], N[s]);
```

```
bool simplex(int phase) {
       int x = m + phase - 1;
       for (;;) {
           int s = -1:
           rep(j, n + 1) if (N[j] != -phase) ltj(D[x]);
           if (D[x][s] >= -eps) return true:
           int r = -1;
           rep(i, m) {
              if (D[i][s] <= eps) continue;</pre>
              if (r == -1 \mid | MP(D[i][n + 1] / D[i][s], B[i])
                    < MP(D[r][n + 1] / D[r][s], B[r])) r = i
           if (r == -1) return false:
           pivot(r, s);
   }
   T solve(vd &x) {
       int r = 0:
       repx(i, 1, m) if (D[i][n + 1] < D[r][n + 1]) r = i;
       if (D[r][n + 1] < -eps) {
           pivot(r, n);
           if (!simplex(2) || D[m + 1][n + 1] < -eps) return</pre>
           rep(i, m) if (B[i] == -1) {
              int s = 0:
              repx(j, 1, n + 1) ltj(D[i]);
              pivot(i, s);
           }
       bool ok = simplex(1);
       x = vd(n):
       rep(i, m) if (B[i] < n) x[B[i]] = D[i][n + 1];
       return ok ? D[m][n + 1] : inf:
   }
};
```

8.11 theorems

```
Burnside lemma
For a set X, with members x in X, and a group G, with
operations g in G, where g(x): X \rightarrow X.
F g is the set of x which are fixed points of g (ie.
\{x \text{ in } X / g(x) = x \}).
The number of orbits (connected components in the
graph formed by assigning each x a node and
a directed edge between x and g(x) for every g) is
called M.
M = the average of the fixed points of all
g = (|F_g1| + |F_g2| + ... + |F_gn|) / |G|
If x are images and g are simmetries, then M
corresponds to the amount of objects, |G|
corresponds to the amount of simmetries, and F_g
corresponds to the amount of simmetrical images under
the simmetry g.
Rational root theorem
All rational roots of the polynomials with integer
a0 * x^0 + a1 * x^1 + a2 * x^2 + ... + an * x^n = 0
If these roots are represented as p / q, with p and q
     coprime,
- p is an integer factor of a0
- q is an integer factor of an
Note that if a0 = 0, then x = 0 is a root, the polynomial
    can be divided by x and the theorem
applies once again.
Petersen's theorem
Every cubic and bridgeless graph has a perfect matching.
Number of divisors for powers of 10
(0,1) (1,4) (2,12) (3,32) (4,64) (5,128) (6,240) (7,448)
(8,768) (9,1344) (10,2304) (11,4032) (12,6720) (13,10752)
(14,17280) (15,26880) (16,41472) (17,64512) (18,103680)
```