# Team Notebook

## July 29, 2023

| Contents |  |          |   | 3.4                      | halfplane-intersect | 10                   | 5                     | .2 Treap  | 17                         |
|----------|--|----------|---|--------------------------|---------------------|----------------------|-----------------------|---|----------------------------|
|          |  |          |   | 3.5                      | line                | 10                   | 5                     | .3 dsu  | 18                         |
| 1        | Strings  | <b>2</b> |   | 3.6                      | point               | 11                   | 5                     | .4 mo   | 18                         |
|          | 1.1 Manacher   | 2        |   | 3.7                      | polygon             | 11                   | 5                     | .5 persistent-segment-tree-lazy                       | 18                         |
|          | 1.2 aho-corasick   | 2        |   | 3.8                      | sweep               | 11                   | 5                     | .6 persistent-segment-tree                            | 19                         |
|          | 1.3 hash   | 2        |   | 3.9                      | theorems            | 12                   | 5                     | .7 segment-tree-lazy                                  | 19                         |
|          | 1.4 hash2d   | 3        |   |                          |                     |                      | 5                     | .8 segment-tree                                       | 19                         |
|          | 1.5 palindromic-tree   | 4        | 4 | grap                     | ph                  | 12                   | 5                     | .9 sparse-table                                       | 20                         |
|          | 1.6 prefix-function  | 4        |   | 4.1                      | bellman-ford        | 12                   |                       | .10 unordered-map                                     |                            |
|          | 1.7 suffix-array-martin  | 5        |   | 4.2                      | dinic               | 12                   |                       |   |                            |
|          | 1.8 suffix-array   | 5        |   | 4.3                      | floyd-warshall      | 12                   | 6 i                   | mprimible   | 20                         |
|          | 1.9 suffix-automaton   | 6        |   | 4.4                      | heavy-light         | 13                   |                       |   |                            |
|          | 1.10 z-function  | 7        |   | 4.5                      | hungarian           | 13                   |                       | nath  | 20                         |
|          | 1.10 Z-Tunetion  |          |   |                          |                     |                      |                       | .1 arithmetic $.$ $.$ $.$ $.$ $.$ $.$ $.$ $.$ $.$ $.$ | 20                         |
|          |  |          |   | 4.6                      | kuhn                | 13                   | 7                     |   |                            |
| 2        | dp   | 8        |   | $4.6 \\ 4.7$             | kuhn                |                      | 7                     | .2 crt  |                            |
| 2        |  | 8        |   | 4.6<br>4.7<br>4.8        |                     | 14                   | 7                     |   |                            |
| 2        | dp   |          |   | 4.6<br>4.7<br>4.8<br>4.9 | lca                 | 14<br>14             | 7<br>7                | .2 crt  | 21                         |
|          | dp 2.1 convex-hull-trick   |          |   | 4.9                      | lca                 | 14<br>14<br>15       | 7<br>7<br>7           | .2 crt  | 21<br>21                   |
|          | <b>dp</b> 2.1 convex-hull-trick  |          |   | 4.9<br>4.10              | lca                 | 14<br>14<br>15<br>15 | 7<br>7<br>7<br>7      | .2 crt  | 21<br>21<br>21             |
|          | dp         2.1 convex-hull-trick          2.2 divide-and-conquer          geo2d       3.1 circle |          |   | 4.9<br>4.10              | lca                 | 14<br>14<br>15<br>15 | 7<br>7<br>7<br>7      | .2 crt  | 21<br>21<br>21<br>22       |
|          | dp         2.1 convex-hull-trick          2.2 divide-and-conquer          geo2d                  |          |   | 4.9<br>4.10<br>4.11      | lca                 | 14<br>14<br>15<br>15 | 7<br>7<br>7<br>7<br>7 | .2 crt  | 21<br>21<br>21<br>22<br>22 |

## 1 Strings

#### 1.1 Manacher

```
#include<bits/stdc++.h>
using namespace std;
#define rep(i, n) for (int i = 0; i < (int)n; i++)
#define repx(i, a, b) for (int i = (int)a; i < (int)b; i++)</pre>
// odd[i] : length of the longest palindrome centered at i
// even[i] : length of the longest palindrome centered
    between i and i+1
void manacher(string &s, vector<int> &odd, vector<int> &even
    ) {
   string t = "$#":
   for(char c: s)
       t += c + string("#");
   t += "^":
   int n = t.size();
   vector<int> p(n);
   int 1 = 1, r = 1;
   repx(i, 1, n-1) {
       p[i] = max(0, min(r - i, p[1 + (r - i)]));
       while(t[i - p[i]] == t[i + p[i]]) {
          p[i]++;
       if(i + p[i] > r) {
          1 = i - p[i], r = i + p[i];
       }
   repx(i, 2, n-2) {
       if(i%2) even.push_back(p[i]-1);
       else odd.push_back(p[i]-1);
```

#### 1.2 aho-corasick

```
#include "../common.h"

const int K = 26;
struct Vertex {
   int next[K];
   int leaf = 0;
   int leaf_id = -1;
   int p = -1;
   char pch;
```

```
int link = -1:
   int exit = -1:
   int cnt = -1;
   int go[K];
   Vertex(int p=-1, char ch='$') : p(p), pch(ch) {
       fill(begin(next), end(next), -1);
       fill(begin(go), end(go), -1);
};
vector<Vertex> t(1):
void add(string &s, int id) {
   for (char ch : s) {
       int c = ch - 'a':
       if (t[v].next[c] == -1) {
           t[v].next[c] = t.size():
           t.emplace_back(v, ch);
       v = t[v].next[c];
   t[v].leaf++:
   t[v].leaf_id = id;
int go(int v. char ch):
int get_link(int v) {
   if (t[v].link == -1) {
       if (v == 0 || t[v].p == 0)
           t[v].link = 0:
           t[v].link = go(get_link(t[v].p), t[v].pch);
   return t[v].link;
int go(int v, char ch) {
   int c = ch - 'a':
   if (t[v].go[c] == -1) {
       if (t[v].next[c] != -1)
           t[v].go[c] = t[v].next[c];
           t[v].go[c] = v == 0 ? 0 : go(get_link(v), ch);
   return t[v].go[c];
```

#### 1.3 hash

```
// compute substring hashes in O(1).
// hashes are compatible between different strings.
struct Hash {
   11 HMOD:
   int N;
   vector<int> h;
   vector<int> p;
   Hash() {}
   // O(N)
   Hash(const string &s, 11 HMOD_ = 1000003931)
       : N(s.size() + 1), HMOD(HMOD_), p(N), h(N) {
       static const 11 P =
          chrono::steady_clock::now().time_since_epoch().
               count() % (1 << 29):
      p[0] = 1;
       rep(i, N - 1) p[i + 1] = p[i] * P % HMOD;
       rep(i, N-1) h[i+1] = (h[i] + (ll)s[i] * p[i]) %
   }
   pair<ll, int> get(int i, int j) { return {(h[j] - h[i] +
        HMOD) % HMOD, i}; }
   bool cmp(pair<11, int> x0, pair<11, int> x1) {
       int d = x0.second - x1.second;
```

```
11 &lo = d < 0 ? x0.first : x1.first:
       lo = lo * p[abs(d)] % HMOD:
       return x0.first == x1.first:
};
// compute hashes in multiple prime modulos simultaneously,
     to reduce the chance
// of collisions.
struct HashM {
    int N:
    vector<Hash> sub:
    HashM() {}
    // D(K N)
    HashM(const string &s, const vector<11> &mods) : N(mods.
        size()), sub(N) {
       rep(i, N) sub[i] = Hash(s, mods[i]);
    // O(K)
    vector<pair<11, int>> get(int i, int j) {
       vector<pair<11, int>> hs(N);
       rep(k, N) hs[k] = sub[k].get(i, j);
       return hs;
    bool cmp(const vector<pair<11, int>> &x0, const vector<
        pair<11, int>> &x1) {
       rep(i, N) if (!sub[i].cmp(x0[i], x1[i])) return false |};
       return true;
    bool cmp(int i0, int j0, int i1, int j1) {
       rep(i, N) if (!sub[i].cmp(sub[i].get(i0, j0),
                               sub[i].get(i1, j1))) return
                                   false:
       return true;
}:
#ifndef NOMAIN HASH
int main() {
    const vector<11> HMOD = {1000001237, 1000003931};
    // 01234567890123456789012
    string s = "abracadabra abracadabra";
    HashM h(s, HMOD):
```

```
rep(i0, s.size() + 1) repx(i0, i0, s.size() + 1) rep(i1, s.size() + 1) rep(i1, s.size() + 1) rep(i1, s.size() + 1) repx(i0, i0, s.size() + 1) repx(i0, i0, s.size() + 1) repx(i1, s.s
                                           s.size() + 1)
                                     repx(j1, i1, s.size() + 1) {
                                      bool eq = h.cmp(h.get(i0, j0), h.get(i1, j1));
                                     bool eq2 = s.substr(i0, j0 - i0) == s.substr(i1, j1 -
                                                                        i1):
                                      if (eq != eq2) {
                                                            cout << " hash says strings \"" << s.substr(i0,</pre>
                                                                                      j0 - i0) << "\" "
                                                                                      << (eq ? "==" : "!=") << " \"" << s.substr(i1
                                                                                                                 , i1 - i1)
                                                                                      << "\" but in reality they are " << (eg2 ? "
                                                                                                                ==" : "!=")
                                                                                      << endl:
                                    }
#endif
```

#### 1.4 hash2d

```
using Hash = pair<11, int>;
struct Block {
   int x0, y0, x1, y1;
struct Hash2d {
   11 HMOD:
   int W. H:
   vector<int> h;
   vector<int> p;
   Hash2d() {}
   Hash2d(const string &s, int W , int H , 11 HMOD =
       1000003931)
      : W(W + 1), H(H + 1), HMOD(HMOD) {
      static const 11 P =
          chrono::steady_clock::now().time_since_epoch().
               count() % (1 << 29):
      p.resize(W * H);
      p[0] = 1;
      rep(i, W * H - 1) p[i + 1] = p[i] * P % HMOD:
      h.assign(W * H, 0);
      repx(v, 1, H) repx(x, 1, W) {
          11 c = (11)s[(y - 1) * (W - 1) + x - 1] * p[y * W
                + x1 % HMOD:
```

```
h[y * W + x] = (HMOD + h[y * W + x - 1] + h[(y -
               1) * W + x] -
                         h[(y-1) * W + x - 1] + c) %
                        HMOD:
   bool isout(Block s) {
       return s.x0 < 0 || s.x0 >= W || s.x1 < 0 || s.x1 >= W
            | | s.y0 < 0 | |
             s.v0 >= H || s.v1 < 0 || s.v1 >= H:
   Hash get(Block s) {
       return \{(2 * HMOD + h[s.v1 * W + s.x1] - h[s.v1 * W +
               h[s.v0 * W + s.x1] + h[s.v0 * W + s.x0]) %
                  HMOD.
              s.v0 * W + s.x0:
   bool cmp(Hash x0, Hash x1) {
       int d = x0.second - x1.second;
       11 &lo = d < 0 ? x0.first : x1.first;</pre>
      lo = lo * p[abs(d)] % HMOD;
       return x0.first == x1.first;
}:
struct Hash2dM {
   int N:
   vector<Hash2d> sub;
   Hash2dM() {}
   Hash2dM(const string &s, int W, int H, const vector<11> &
        mods)
       : N(mods.size()), sub(N) {
       rep(i, N) sub[i] = Hash2d(s, W, H, mods[i]);
   bool isout(Block s) { return sub[0].isout(s): }
   vector<Hash> get(Block s) {
       vector<Hash> hs(N):
       rep(i, N) hs[i] = sub[i].get(s);
       return hs:
   bool cmp(const vector<Hash> &x0, const vector<Hash> &x1)
```

## 1.5 palindromic-tree

```
#include "../common.h"
struct Node {
 int len:
                   // length of substring
                  // insertion edge for all characters a-z
 int edge[26];
                  // the Maximum Palindromic Suffix Node
 int link;
     for the current Node
                  // (optional) start index of current Node
    int cnt = 1;
                     // (optional) number of occurrences of
        this substring
    Node(){ fill(begin(edge), end(edge), -1); }
};
struct EerTree { // Palindromic Tree
    vector<Node> t: // tree
    int curr: // current node
    EerTree(string &s) {
       t.resize(2);
       t.reserve(s.size()+2); // (optional) maximum size of
       t[0].len = -1;
                           // root 1
       t[0].link = 0;
       t[1].len = 0:
                           // root 2
       t[1].link = 0;
```

```
rep(i, s.size()) insert(i, s); // construct tree
       // (optional) calculate number of occurrences of each
            node
       for(int i = t.size()-1: i > 1: i--)
          t[t[i].link].cnt += t[i].cnt;
   }
   void insert(int i, string &s) {
       int tmp = curr:
       while (i - t[tmp].len < 1 \mid | s[i] != s[i-t[tmp].len
           -17)
          tmp = t[tmp].link;
      if(t[tmp].edge[s[i]-'a'] != -1){
          curr = t[tmp].edge[s[i]-'a']; // node already
               exists
          t[curr].cnt++:
                                         // (optional)
               increase cnt
          return;
       curr = t[tmp].edge[s[i]-'a'] = t.size(); // create
           new node
       t.emplace_back();
       t[curr].len = t[tmp].len + 2;
                                        // set length
       t[curr].i = i - t[curr].len + 1; // (optional) set
           start index
       if (t[curr].len == 1) {
                                         // set suffix link
          t[curr].link = 1:
      } else {
          tmp = t[tmp].link;
          while (i-t[tmp].len < 1 | s[i] != s[i-t[tmp].len
              tmp = t[tmp].link;
          t[curr].link = t[tmp].edge[s[i]-'a'];
   }
};
int main()
string s = "abcbab":
   EerTree pt(s); // construct palindromic tree
repx(i, 2, pt.t.size()) // list all distinct palindromes
 cout << i-1 << ") ":
```

```
repx(j, pt.t[i].i, pt.t[i].i + pt.t[i].len)
  cout << s[j];
  cout << " " << pt.t[i].cnt << endl;
}
return 0;
}</pre>
```

## 1.6 prefix-function

```
#include "../common.h"
vector<int> prefix_function(string s) {
   int n = s.size():
   vector<int> pi(n);
   repx(i, 1, n) {
       int j = pi[i-1];
       while (j > 0 \&\& s[i] != s[j])
          j = pi[j-1];
       if (s[i] == s[j])
           j++;
       pi[i] = j;
   return pi;
vector<vector<int>> aut:
void compute_automaton(string s) {
   s += '#':
   int n = s.size():
   vector<int> pi = prefix_function(s);
   aut.assign(n, vector<int>(26)):
   rep(i, n) {
       rep(c, 26) {
          int i = i:
          while (j > 0 \&\& 'a' + c != s[j])
              j = pi[j-1];
           if ('a' + c == s[j])
              j++;
          aut[i][c] = j;
       }
// k = n - pi[n - 1]
// if k divides n, then the string can be aprtitioned into
    blocks of length k
```

```
// otherwise there is no effective compression and the answer is {\tt n.}
```

### 1.7 suffix-array-martin

```
// build the suffix array
// suffixes are sorted, with each suffix represented by its
     starting position
vector<int> suffixarray(const string &s) {
   int N = s.size() + 1: // optional: include terminating
   vector\langle int \rangle p(N), p2(N), c(N), c2(N), cnt(256);
   rep(i, N) cnt[s[i]] += 1;
   repx(b, 1, 256) cnt[b] += cnt[b - 1];
   rep(i, N) p[--cnt[s[i]]] = i;
   repx(i, 1, N) c[p[i]] = c[p[i - 1]] + (s[p[i]] != s[p[i - 1]])
         1]]):
   for (int k = 1: k < N: k <<= 1) {</pre>
       int C = c[p[N - 1]] + 1;
       cnt.assign(C + 1, 0);
       for (int &pi : p) pi = (pi - k + N) % N;
       for (int cl : c) cnt[cl + 1] += 1;
       rep(i, C) cnt[i + 1] += cnt[i];
       rep(i, N) p2[cnt[c[p[i]]]++] = p[i];
       c2[p2[0]] = 0;
       repx(i, 1, N) c2[p2[i]] =
           c2[p2[i-1]] + (c[p2[i]] != c[p2[i-1]] ||
                          c[(p2[i] + k) \% N] != c[(p2[i - 1]
                                + k) % N]):
       swap(c, c2), swap(p, p2);
   p.erase(p.begin()); // optional: erase terminating NUL
   return p;
// build the lcp
// 'lcp[i]' represents the length of the longest common
    prefix between suffix i
// and suffix i+1 in the suffix array 'p'. the last element
    of 'lcp' is zero by
// convention
vector<int> makelcp(const string &s, const vector<int> &p) {
   int N = p.size(), k = 0;
   vector<int> r(N), lcp(N);
   rep(i, N) r[p[i]] = i:
   rep(i, N) {
       if (r[i] + 1 >= N) {
          k = 0:
           continue;
```

```
int j = p[r[i] + 1];
       while (i + k < N \&\& j + k < N \&\& s[i + k] == s[j + k]
            1) k += 1:
       lcp[r[i]] = k;
       if (k) k -= 1:
   return lcp;
#ifndef NOMAIN SUFARR
void test(const string &s) {
   cout << "suffix array for string \"" << s << "\" (length</pre>
        " << s.size()
        << "):" << endl;
   vector<int> sa = suffixarrav(s):
   vector<int> lcp = makelcp(s, sa);
   rep(i, sa.size()) {
       int i = sa[i]:
       if (i > 0) cout << " " << lcp[i - 1] << endl;</pre>
       cout << " \"" << s.substr(j) << "\"" << endl:</pre>
   }
int main() {
   test("hello"):
   test("abracadabra"):
#endif
```

## 1.8 suffix-array

```
void getSA(vector<int>& s) {
   R = R_{-} = sa = sa_{-} = vector < int > (n); rep(i, n) sa[i] =
   sort(sa.begin(), sa.end(), [&s](int i, int j) {
        return s[i] < s[i]: }):</pre>
   int r = R[sa[0]] = 1;
   repx(i, 1, n) R[sa[i]] = (s[sa[i]] != s[sa[i - 1]]) ?
         ++r : r:
   for (int h = 1; h < n && r < n; h <<= 1) {
       csort(r, h); csort(r, 0); r = R_[sa[0]] = 1;
       repx(i, 1, n) {
          if (R[sa[i]] != R[sa[i - 1]] || gr(sa[i] + h)
                != gr(sa[i - 1] + h)) r++;
          R_{sa}[sa[i]] = r;
       } R.swap(R_);
void getLCP(vector<int> &s) {// only works with suffixes
     (not cyclic shifts)
   lcp.assign(n, 0); int k = 0;
   rep(i, n) {
       int r = R[i] - 1;
       if (r == n - 1) \{ k = 0; continue; \}
       int j = sa[r + 1];
       while (i + k < n \&\& j + k < n \&\& s[i + k] == s[j]
            + kl) k++:
       lcp[r] = k: if (k) k--:
   }
SuffixArray(vector<int> &s) { n = s.size(); getSA(s);
     getLCP(s); constructLCP(); }
/* ----- */
vector<vector<int>> T:
void constructLCP() {
   T.assign(LOG2(n)+1, lcp);
   for(int k = 1; (1<<k) <= n; ++k)
       for(int i = 0; i + (1<<k) <= n; ++i)</pre>
          T[k][i] = min(T[k-1][i], T[k-1][i+(1<<(k-1))]);
// get LCP of suffix starting at i and suffix starting at
int queryLCP(int i, int j) {
   if(i == j) return n-i;
   i = R[i]-1; i = R[i]-1:
   if(i > j) swap(i, j);
   11 k = LOG2(i-i):
   return min(T[k][i],T[k][j-(1<<k)]);</pre>
```

```
// compare substring of length len1 starting at i
    // with substring of length len2 starting at j
    bool cmp(int i, int len1, int j, int len2) {
       if(queryLCP(i, j) >= min(len1, len2))
           return (len1 < len2);</pre>
           return (R[i] < R[j]);</pre>
}:
vector<int> suffix array:
vector<vector<int>> C:
int n;
void sort_cyclic_shifts(string s) {
    s += "$";
    n = s.size():
    const int alphabet = 256;
    vector<int> p(n), c(n), cnt(max(alphabet, n), 0);
    for (int i = 0: i < n: i++)</pre>
       cnt[s[i]]++;
    for (int i = 1; i < alphabet; i++)</pre>
       cnt[i] += cnt[i-1];
    for (int i = 0; i < n; i++)</pre>
       p[--cnt[s[i]]] = i;
    c[p[0]] = 0;
    int classes = 1:
    for (int i = 1: i < n: i++) {</pre>
       if (s[p[i]] != s[p[i-1]])
           classes++:
       c[p[i]] = classes - 1;
    C.emplace_back(c.begin(), c.end());
    vector<int> pn(n), cn(n);
    for (int h = 0; (1 << h) < n; ++h) {
       for (int i = 0: i < n: i++) {</pre>
           pn[i] = p[i] - (1 << h);
           if (pn[i] < 0)</pre>
               pn[i] += n;
       fill(cnt.begin(), cnt.begin() + classes, 0);
       for (int i = 0; i < n; i++)</pre>
           cnt[c[pn[i]]]++:
       for (int i = 1; i < classes; i++)</pre>
           cnt[i] += cnt[i-1];
       for (int i = n-1: i >= 0: i--)
           p[--cnt[c[pn[i]]]] = pn[i];
       cn[p[0]] = 0;
       classes = 1:
       for (int i = 1; i < n; i++) {</pre>
```

```
pair<int, int> cur = \{c[p[i]], c[(p[i] + (1 << h))\}
          << h)) % nl}:
          if (cur != prev)
             ++classes:
          cn[p[i]] = classes - 1;
      c.swap(cn);
      C.emplace_back(c.begin(), c.end());
   p.erase(p.begin()):
   suffix_array = p;
vector<int> lcp_construction(string &s, vector<int> &p) {
   int n = s.size():
   vector<int> rank(n):
   rep(i, n) rank[p[i]] = i:
   int k = 0;
   vector<int> lcp(n-1, 0);
   rep(i, n) {
      if (rank[i] == n - 1) {
          k = 0:
          continue;
      int i = p[rank[i] + 1]:
      while (i + k < n \&\& j + k < n \&\& s[i+k] == s[j+k])
      lcp[rank[i]] = k;
      if (k)
          k--:
   }
   return lcp;
bool compare1(int i, int j, int l) {
   int k = LOG2(1);
   pair<int, int> a = \{C[k][i], C[k][(i+1-(1 << k))\%n]\};
   pair<int. int> b = \{C[k][i], C[k][(i+1-(1 << k))\%n]\}:
   return a >= b:
bool compare2(int i, int j, int l) {
   int k = LOG2(1):
   pair<int, int> a = \{C[k][i], C[k][(i+l-(1 << k))\%n]\};
   pair<int, int> b = \{C[k][j], C[k][(j+1-(1 << k))\%n]\};
   return a <= b:
```

```
pair<int,int> find(int i, int len)
   int 1 = 0, r = suffix_array.size()-1;
   while(1 != r)
      int mid = (1+r)/2;
      if(compare1(suffix_array[mid], i, len))
          r = mid;
       else
          1 = mid+1:
   int left = 1:
   1 = 0, r = suffix_array.size()-1;
   while(1 != r)
      int mid = (1+r+1)/2:
      if(compare2(suffix_array[mid], i, len))
          1 = mid;
      else
          r = mid-1;
   int right = 1;
   if(!compare1(suffix_array[left], i, len)) return {-1,-1};
   if(!compare2(suffix array[right], i, len)) return
        \{-1,-1\};
   if(left > right) return {-1,-1};
   return {left, right};
```

#### 1.9 suffix-automaton

```
#include "../common.h"

struct SuffixAutomaton {
   vector<map<char,int>> edges; // edges[i] : the labeled
        edges from node i
   vector<int> link; // link[i] : the suffix link
        of i
   vector<int> length; // length[i] : the length of
        the longest string in the ith class
   vector<int> cnt; // cnt[i] : number of
        occurrences of each string in the ith class
```

```
vector<int> paths:
                          // paths[i] : number of paths
    on the automaton starting from i
vector<bool> terminal:
                         // terminal[i] : true if i is
    a terminal state
vector<int> first_pos;
vector<int> last pos:
int last;
                         // the index of the
    equivalence class of the whole string
SuffixAutomaton(string s) {
   edges.push_back(map<char,int>());
   link.push back(-1):
   length.push_back(0);
   last = 0:
   rep(i, s.size()) { // construct r
       edges.push_back(map<char,int>());
       length.push_back(i+1);
       link.push back(0):
       int r = edges.size() - 1:
       int p = last; // add edges to r and find p with
           link to a
       while(p >= 0 && !edges[p].count(s[i])) {
          edges[p][s[i]] = r;
          p = link[p];
       if(p != -1) {
          int a = edges[p][s[i]]:
          if(length[p] + 1 == length[q]) {
              link[r] = q; // we do not have to split q,
                   just set the correct suffix link
          } else { // we have to split, add q'
              edges.push_back(edges[q]); // copy edges
              length.push_back(length[p] + 1);
              link.push back(link[q]): // copy parent of
              int qq = edges.size()-1;
              link[q] = qq; // add qq as the new parent
                   of q and r
              link[r] = aa:
              while(p >= 0 && edges[p][s[i]] == q) { //
                  move short classes polling to q to
                  poll to q'
                  edges[p][s[i]] = qq;
                 p = link[p];
              }
       }
       last = r;
```

```
/* ----- Optional ----- */
   // mark terminal nodes
   terminal.assign(edges.size(), false):
   int p = last;
   while(p > 0) {
       terminal[p] = true;
       p = link[p];
   // precompute match count
   cnt.assign(edges.size(), -1);
   cnt matches(0):
   // precompute number of paths (substrings) starting
        from state
   paths.assign(edges.size(), -1);
   cnt paths(0):
   first_pos.assign(edges.size(), -1);
   get_first_pos(0);
   last_pos.assign(edges.size(), -1);
   get_last_pos(0);
int cnt_matches(int state) {
   if(cnt[state] != -1) return cnt[state]:
   int ans = terminal[state];
   for(auto edge : edges[state])
       ans += cnt_matches(edge.second);
   return cnt[state] = ans;
int cnt_paths(int state) {
   if(paths[state] != -1) return paths[state];
   int ans = state == 0 ? 0 : 1; // without repetition (
        counts diferent substrings)
// int ans = state == 0 ? 0 : cnt[state]: // with
    repetition
   for(auto edge : edges[state])
       ans += cnt_paths(edge.second);
   return paths[state] = ans;
int get_first_pos(int state) {
   if(first_pos[state] != -1) return first_pos[state];
   int ans = 0:
```

```
for(auto edge : edges[state])
          ans = max(ans, get_first_pos(edge.second)+1);
       return first_pos[state] = ans;
   int get_last_pos(int state) {
       if(last_pos[state] != -1) return last_pos[state];
       int ans = terminal[state] ? 0 : INT_MAX;//fix
       for(auto edge : edges[state])
          ans = min(ans, get_last_pos(edge.second)+1);
       return last pos[state] = ans:
   string get_k_substring(int k) // 0-indexed
       string ans;
       int state = 0:
       while(true)
          int curr = state == 0 ? 0 : 1: // without
               repetition (counts different substrings)
       // int curr = state == 0 ? 0 : cnt[state]: // with
           repetition
          if(curr > k) return ans:
          k -= curr:
          for(auto edge : edges[state]) {
              if(paths[edge.second] <= k) {</pre>
                  k -= paths[edge.second];
              } else {
                  ans += edge.first;
                  state = edge.second;
                  break:
          }
      }
};
```

#### 1.10 z-function

```
int n = s.size();
vector<int> z(n);
int 1 = 0, r = 0;
for(int i = 1; i < n; i++) {
    if(i < r) {
        z[i] = min(r - i, z[i - 1]);
    }
    while(i + z[i] < n && s[z[i]] == s[i + z[i]]) {
        z[i]++;
    }
    if(i + z[i] > r) {
        1 = i;
        r = i + z[i];
    }
}
return z;
}
```

## $2 ext{ dp}$

#### 2.1 convex-hull-trick

```
struct Line {
    mutable 11 a, b, c;
    bool operator<(Line r) const { return a < r.a; }</pre>
    bool operator<(ll x) const { return c < x; }</pre>
}:
// dynamically insert 'a*x + b' lines and query for maximum
// all operations have complexity O(log N)
struct LineContainer : multiset<Line. less<>>> {
    11 div(ll a, ll b) {
       return a / b - ((a ^ b) < 0 && a % b):
    bool isect(iterator x, iterator y) {
       if (y == end()) return x \rightarrow c = INF, 0;
       if (x->a == y->a) x->c = x->b > y->b ? INF : -INF;
       else x->c = div(y->b - x->b, x->a - y->a);
       return x->c >= y->c;
    void add(ll a, ll b) {
       // a *= -1, b *= -1 // for min
       auto z = insert(\{a, b, 0\}), y = z++, x = y;
```

## 2.2 divide-and-conquer

```
// for every index i assign an optimal index j, such that
    cost(i, j) is
// minimal for every i. the property that if i2 >= i1 then
    i2 >= i1 is
// exploited (monotonic condition).
// calculate optimal index for all indices in range [1, r)
    knowing that
// the optimal index for every index in this range is within
     [optl, optr).
// time: O(N log N)
void calc(vector<int> &opt, int 1, int r, int optl, int optr
    ) {
   if (1 == r) return:
   int i = (1 + r) / 2:
   11 optc = INF:
   int optj;
   repx(j, optl, optr) {
      11 c = i + j; // cost(i, j)
      if (c < optc) optc = c, optj = j;</pre>
   opt[i] = optj;
   calc(opt, 1, i, optl, optj + 1);
   calc(opt, i + 1, r, optj, optr);
```

- $3 \quad \text{geo2d}$
- 3.1 circle

```
struct C {
   Po; Tr;
   C(P \circ, T r) : o(o), r(r) {}
   C() : C(P(), T()) \{ \}
   // intersects the circle with a line, assuming they
   // results are sorted with respect to the direction of
        the line
   pair<P. P> line inter(L 1) const {
       P c = 1.closest to(o):
      T c2 = (c - o).magsq();
       P = sqrt(max(r * r - c2, T())) * 1.d.unit();
       return {c - e, c + e};
   }
   // checks whether the given line collides with the circle
   // negative: 2 intersections
   // zero: 1 intersection
   // positive: 0 intersections
   T line collide(L 1) const {
      T c2 = (1.closest_to(o) - o).magsq();
       return c2 - r * r:
   // calculates the two intersections between two circles
   // the circles must intersect in one or two points!
   pair<P. P> inter(C h) const {
       P d = h.o - o:
      T c = (r * r - h.r * h.r) / d.magsq();
       return h.line_inter(\{(1 + c) / 2 * d, d.rot()\});
   // check if the given circles intersect
   bool collide(C h) const {
       return (h.o - o).magsq() \le (h.r + r) * (h.r + r);
   // get one of the two tangents that cross through the
   // the point must not be inside the circle
   // a = -1: cw (relative to the circle) tangent
   // a = 1: ccw (relative to the circle) tangent
   P point_tangent(P p, T a) const {
       T c = r * r / p.magsq():
       return o + c * (p - o) - a * sqrt(c * (1 - c)) * (p - c)
             o).rot():
   }
```

```
// get one of the 4 tangents between the two circles
// a = 1: exterior tangents
// a = -1: interior tangents (requires no area overlap)
// b = 1: ccw tangent
// b = -1: cw tangent
// the line origin is on this circumference, and the
    direction
// is a unit vector towards the other circle
L tangent(C c, T a, T b) const {
   T dr = a * r - c.r;
   P d = c.o - o:
   P n = (d * dr + b * d.rot() * sqrt(d.magsq() - dr *
        dr)).unit():
   return {o + n * r, -b * n.rot()};
// find the circumcircle of the given **non-degenerate**
    triangle
static C thru points(P a, P b, P c) {
   L 1((a + b) / 2, (b - a).rot()):
   P p = 1.intersection(L((a + c) / 2, (c - a).rot()));
   return {p, (p - a).mag()};
// find the two circles that go through the given point,
    are tangent
// to the given line and have radius 'r'
// the point-line distance must be at most 'r'!
// the circles are sorted in the direction of the line
static pair<C. C> thru point line r(P a, L t, T r) {
   P d = t.d.rot().unit();
   if (d * (a - t.o) < 0) d = -d;
   auto p = C(a, r).line_inter(\{t.o + d * r, t.d\});
   return {{p.first, r}, {p.second, r}};
// find the two circles that go through the given points
    and have
// radius 'r'
// the circles are sorted by angle with respect to the
// the points must be at most at distance 'r'!
static pair < C. C > thru points r(P a. P b. T r) {
   return {{p.first, r}, {p.second, r}};
```

};

#### 3.2 convex-hull

```
// get the convex hull with the least amount of vertices for
      the given set
// of points
// probably misbehaves if points are not all distinct!
vector<P> convex_hull(vector<P> &ps) {
   int N = ps.size(), n = 0, k = 0;
   if (N <= 2) return ps;</pre>
   rep(i, N) if (make_pair(ps[i].v, ps[i].x) < make_pair(ps[</pre>
        k].y, ps[k].x)) k = i;
   swap(ps[k], ps[0]);
   sort(++ps.begin(), ps.end(), [&](P 1, P r) {
       T x = (r - 1) / (ps[0] - 1), d = (r - 1) * (ps[0] - 1)
       return x > 0 \mid | x == 0 && d < 0:
   }):
   vector<P> H:
   for (P p : ps) {
       while (n \ge 2 \&\& (H[n-1] - p) / (H[n-2] - p) >=
            0) H.pop back(), n--:
       H.push_back(p), n++;
   return H;
```

## 3.3 delaunay

```
Q *r = H ? H : new Q{new Q{new Q{new Q{0}}}}:
   H = r - > 0: r - > r() - > r() = r:
   repx(i, 0, 4) r = r->rot, r->p = {INF, INF},
       r->0 = i & 1 ? r : r->r():
   r\rightarrow p = orig; r\rightarrow F() = dest;
   return r:
void splice(Q *a, Q *b) {
   swap(a->o->rot->o, b->o->rot->o); swap(a->o, b->o);
Q *connect(Q *&H, Q *a, Q *b) {
   Q *q = makeEdge(H, a->F(), b->p);
   splice(q, a->next()); splice(q->r(), b); return q;
pair<Q *, Q *> rec(Q *&H, const vector<P> &s) {
   if (s.size() <= 3) {</pre>
       Q *a = makeEdge(H, s[0], s[1]), *b = makeEdge(H, s[0])
            [1], s.back());
       if (s.size() == 2) return \{a, a->r()\}; splice(a->r(),
       auto side = cross(s[0], s[1], s[2]);
       Q *c = side ? connect(H, b, a) : 0;
       return {side < 0 ? c->r() : a, side < 0 ? c : b->r()
   }
#define J(e) e->F(), e->p
#define valid(e) (cross(e->F(), J(base)) > 0)
   Q *A, *B, *ra, *rb; int half = s.size() / 2;
   tie(ra, A) = rec(H, {s.begin(), s.end() - half});
   tie(B, rb) = rec(H, {s.begin() + s.size() - half, s.end()
        }):
   while ((cross(B->p, J(A)) < 0 && (A = A->next())) | |
          (cross(A->p, J(B)) > 0 \&\& (B = B->r()->o)));
   Q *base = connect(H, B->r(), A);
   if (A->p == ra->p) ra = base->r();
   if (B->p == rb->p) rb = base:
#define DEL(e, init, dir) Q *e = init->dir; \
   if (valid(e)) while (circ(e->dir->F(), J(base), e->F()))
           Q *t = e->dir; splice(e, e->prev()); \
           splice(e->r(), e->r()->prev()); e->o = H; H = e;
               e = t; \
   for (::) {
       DEL(LC, base->r(), o); DEL(RC, base, prev());
```

```
if (!valid(LC) && !valid(RC)) break;
       if (!valid(LC) || (valid(RC) && circ(J(RC), J(LC))))
            base = connect(H, RC, base->r());
       else base = connect(H, base->r(), LC->r());
   }
   return {ra, rb}:
#undef J
#undef valid
#undef DEL
// there must be no duplicate points
// returns no triangles in the case of all collinear points
// produces counter-clockwise triangles ordered in triples
// maximizes the minimum angle across all triangulations
// the euclidean mst is a subset of these edges
// O(N log N)
vector<P> triangulate(vector<P> pts) {
   sort(pts.begin(), pts.end(), [](P a, P b) {
       return make_pair(a.x, a.y) < make_pair(b.x, b.y);</pre>
   assert(unique(pts.begin(), pts.end()) == pts.end());
   if (pts.size() < 2) return {};</pre>
   Q *H = 0; Q *e = rec(H, pts).first;
   vector < Q *> q = \{e\}; int qi = 0;
   while (cross(e->o->F(), e->F(), e->p) < 0) e = e->o;
#define ADD
   {
       0 *c = e:
       do {
           c->mark = 1; pts.push_back(c->p); \
           q.push_back(c->r()); c = c->next(); \
       } while (c != e):
   ADD:
   pts.clear():
   while (qi < (int)q.size()) if (!(e = q[qi++])->mark) ADD;
   return pts;
#undef ADD
```

## 3.4 halfplane-intersect

```
// obtain the convex polygon that results from intersecting
    the given list
// of halfplanes, represented as lines that allow their left
    side
// assumes the halfplane intersection is bounded
vector<P> halfplane_intersect(vector<L> &H) {
```

```
L bb(P(-INF, -INF), P(INF, 0)):
rep(k, 4) H.push_back(bb), bb.o = bb.o.rot(), bb.d = bb.d
     .rot();
sort(begin(H), end(H), [](L a, L b) { return a.d.angcmp(b
     .d) < 0: }):
deque < L > q; int n = 0;
rep(i, H.size()) {
   while (n >= 2 && H[i].side(q[n - 1].intersection(q[n
        -21)) > 0)
       g.pop back(), n--:
   while (n \ge 2 \&\& H[i].side(a[0].intersection(a[1])) >
       q.pop_front(), n--;
   if (n > 0 && H[i].parallel(q[n - 1])) {
       if (H[i].d * q[n - 1].d < 0) return {};</pre>
       if (H[i].side(q[n-1].o) > 0) q.pop_back(), n--;
       else continue;
   q.push_back(H[i]), n++;
while (n \ge 3 \&\& q[0].side(q[n - 1].intersection(q[n -
    21)) > 0)
    q.pop_back(), n--;
while (n \ge 3 \&\& q[n - 1].side(q[0].intersection(q[1])) >
   q.pop_front(), n--;
if (n < 3) return {};</pre>
vector<P> ps(n);
rep(i, n) ps[i] = q[i].intersection(q[(i + 1) % n]);
return ps;
```

#### 3.5 line

```
// a segment or an infinite line
// does not handle point segments correctly!
struct L {
  P o, d;
  L() : o(), d() {}
  L(P o, P d) : o(o), d(d) {}

  L(P ab, T c) : d(ab.rot()), o(ab * -c / ab.magsq()) {}
  pair<P, T> line_eq() { return {-d.rot(), d.rot() * o}; }

// returns a number indicating which side of the line the point is in
```

```
// negative: left, positive: right
T side(P r) const { return (r - o) / d: }
// returns the intersection coefficient
// in the range [0, d / r.d]
// if d / r.d is zero, the lines are parallel
T inter(L r) const { return (r.o - o) / r.d; }
// get the single intersection point
// lines must not be parallel
P intersection(L r) const { return o + d * inter(r) / (d
     / r.d): }
// check if lines are parallel
bool parallel(L r) const { return abs(d / r.d) <= EPS; }</pre>
// check if segments intersect
bool seg_collide(L r) const {
   Tz = d / r.d:
   if (abs(z) <= EPS) {
       if (abs(side(r.o)) > EPS) return false;
       T s = (r.o - o) * d, e = s + r.d * d;
       if (s > e) swap(s, e);
       return s \leq d * d + EPS && e \geq -EPS:
   T s = inter(r), t = -r.inter(*this);
   if (z < 0) s = -s, t = -t, z = -z;
   return s >= -EPS && s <= z + EPS && t >= -EPS && t <=
         z + EPS;
}
// full segment intersection
// produces a point segment if the intersection is a
// however it **does not** handle point segments as input
bool seg_inter(L r, L *out) const {
   Tz = d / r.d:
   if (abs(z) <= EPS) {
       if (abs(side(r.o)) > EPS) return false;
       if (r.d * d < 0) r = \{r.o + r.d, -r.d\}:
       P s = o * d < r.o * d ? r.o : o;
       P = (o + d) * d < (r.o + r.d) * d ? o + d : r.o
             + r.d:
       if (s * d > e * d) return false;
       return *out = L(s, e - s), true;
   T s = inter(r), t = -r.inter(*this):
   if (z < 0) s = -s, t = -t, z = -z;
```

```
if (s \ge -EPS \&\& s \le z + EPS \&\& t \ge -EPS \&\& t \le z
            + EPS)
           return *out = L(o + d * s / z, P()), true;
       return false:
   // check if the given point is on the segment
   bool point_on_seg(P r) const {
       if (abs(side(r)) > EPS) return false:
       if ((r - o) * d < -EPS) return false;</pre>
       if ((r - o - d) * d > EPS) return false:
       return true:
   }
   // get the point in this line that is closest to a given
        point
   P closest to(P r) const {
       P dr = d.rot(); return r + (o - r) * dr * dr / d.
            magsq():
   }
};
```

### 3.6 point

```
struct P {
   T x, y;
   P(T x, T y) : x(x), y(y) {}
   P() : P(0, 0) \{ \}
   friend ostream &operator<<(ostream &s, const P &r) {</pre>
      return s << r.x << " " << r.y;
   friend istream & operator >> (istream &s, P &r) { return s
        >> r.x >> r.v: }
   P operator+(P r) const { return \{x + r.x, y + r.y\}; }
   P operator-(P r) const { return {x - r.x, y - r.y}; }
   P operator*(T r) const { return {x * r, y * r}; }
   P operator/(T r) const { return {x / r, y / r}; }
   P operator-() const { return {-x, -y}; }
   friend P operator*(T 1, P r) { return {1 * r.x, 1 * r.y};
   P rot() const { return {-y, x}; }
   T operator*(P r) const { return x * r.x + v * r.v: }
   T operator/(P r) const { return rot() * r; }
   T magsq() const { return x * x + y * y; }
   T mag() const { return sqrt(magsq()); }
```

## 3.7 polygon

```
// get the area of a simple polygon in ccw order
T area(const vector<P> &ps) {
   int N = ps.size();
   Ta = 0:
   rep(i, N) a += (ps[i] - ps[0]) / (ps[(i + 1) % N] - ps[i]
   return a / 2:
// checks whether a point is inside a simple polygon
// returns -1 if inside, 0 if on border, 1 if outside
int in_poly(const vector<P> &ps, P p) {
   int N = ps.size(), w = 0;
   rep(i, N) {
       P s = ps[i] - p, e = ps[(i + 1) \% N] - p;
       if (s == P()) return 0;
       if (s.v == 0 \&\& e.v == 0) {
           if (\min(s.x, e.x) \le 0 \&\& 0 \le \max(s.x, e.x))
               return 0:
       } else {
           bool b = s.y < 0;
           if (b != (e.y < 0)) {
              Tz = s / e; if (z == 0) return 0;
              if (b == (z > 0)) w += b ? 1 : -1:
       }
   return w ? -1 : 1:
```

```
// check if a point is in a convex polygon
   vector<P> ps:
   T 11, 1h, rl, rh;
   int N. m:
   // preprocess polygon
   // O(N)
   InConvex(const vector<P> &p) : ps(p), N(ps.size()), m(0)
       assert(N >= 2):
       rep(i, N) if (ps[i].x < ps[m].x) m = i;
       rotate(ps.begin(), ps.begin() + m, ps.end());
       rep(i, N) if (ps[i].x > ps[m].x) m = i;
      11 = 1h = ps[0].v, r1 = rh = ps[m].v;
      for (P p : ps) {
          if (p.x == ps[0].x) 11 = min(11, p.y), 1h = max(
          if (p.x == ps[m].x) rl = min(rl, p.y), rh = max(
               rh, p.v);
      }
   InConvex() {}
   // check if point belongs in polygon
   // returns -1 if inside. 0 if on border. 1 if outside
   // O(log N)
   int in_poly(P p) {
       if (p.x < ps[0].x || p.x > ps[m].x) return 1;
       if (p.x == ps[0].x) return p.y < 11 || p.y > 1h;
       if (p.x == ps[m].x) return p.y < rl \mid \mid p.y > rh;
       int r = upper_bound(ps.begin(), ps.begin() + m, p,
           [](Pa, Pb) \{ return a.x < b.x; \}) - ps.begin();
      Tz = (ps[r-1] - ps[r]) / (p - ps[r]); if (z >= 0)
           return !!z:
       r = upper_bound(ps.begin() + m, ps.end(), p,
          [](P a, P b) { return a.x > b.x; }) - ps.begin();
       z = (ps[r - 1] - ps[r \% N]) / (p - ps[r \% N]);
       if (z \ge 0) return !!z: return -1:
   }
};
```

#### 3.8 sweep

```
#include "point.cpp"

// iterate over all pairs of points
```

```
// 'op' is called with all ordered pairs of different
    indices '(i, i)'
// additionally, the 'ps' vector is kept sorted by signed
// to the line formed by 'i' and 'j'
// for example, if the vector from 'i' to 'i' is pointing
// the 'ps' vector is sorted from smallest 'y' to largest 'y
// note that, because the 'ps' vector is sorted by signed
    distance.
// 'i' is always equal to 'i + 1'
// this means that the amount of points to the left of the
    line is always 'N - i'
template <class OP>
void all_pair_points(vector<P> &ps, OP op) {
   int N = ps.size():
   sort(ps.begin(), ps.end(), [](P a, P b) {
       return make pair(a.v. a.x) < make pair(b.v. b.x):
   vector<pair<int, int>> ss;
   rep(i, N) rep(j, N) if (i != j) ss.push_back({i, j});
   stable_sort(ss.begin(), ss.end(), [&](auto a, auto b) {
       return (ps[a.second] - ps[a.first]).angle_lt(ps[b.
            second] - ps[b.first]);
   vector<int> p(N); rep(i, N) p[i] = i;
   for (auto [i, j] : ss)
       { op(p[i], p[j]); swap(ps[p[i]], ps[p[j]]); swap(p[i
           ], p[i]); }
```

#### 3.9 theorems

```
// Pick's theorem
// Simple polygon with integer vertices:
// A = I + B / 2 - 1
// A: Area of the polygon
// I: Integer points strictly inside the polygon
// B: Integer points on the boundary of the polygon
```

## 4 graph

### 4.1 bellman-ford

```
struct Edge { int u, v; ll w; };
```

```
// find distance from source node to all nodes.
// supports negative edge weights.
// returns true if a negative cycle is detected.
// time: O(V E)
bool bellman_ford(int N, int s, vector<Edge> &E, vector<11>
    &D. vector<int> &P) {
   P.assign(N, -1), D.assign(N, INF), D[s] = 0;
   rep(i, N - 1) {
       bool f = true:
       rep(ei, E.size()) {
           auto &e = E[ei];
           ll n = D[e.u] + e.w;
          if (D[e.u] < INF && n < D[e.v])</pre>
              D[e.v] = n, P[e.v] = ei, f = false;
       if (f) return false;
   return true:
```

#### 4.2 dinic

```
struct Edge { int u, v; ll c, f = 0; };
// maximum flow algorithm.
// time: 0(E V^2)
11
        O(E V^(2/3)) / O(E sqrt(E)) unit capacities
11
        O(E sqrt(V))
                                    unit networks (hopcroft-
// unit network: c in {0, 1} and forall v. len(incoming(v))
    <= 1 or len(outgoing(v)) <= 1
// min-cut: find all nodes reachable from the source in the
    residual graph
struct Dinic {
   int N. s. t: vector<vector<int>> G:
   vector<Edge> E; vector<int> lvl, ptr;
   Dinic() {}
   Dinic(int N, int s, int t) : N(N), s(s), t(t), G(N) {}
   void add_edge(int u, int v, ll c) {
       G[u].push_back(E.size()); E.push_back({u, v, c});
       G[v].push_back(E.size()); E.push_back({v, u, 0});
   11 push(int u, 11 p) {
       if (u == t || p <= 0) return p;</pre>
       while (ptr[u] < G[u].size()) {</pre>
```

```
int ei = G[u][ptr[u]++];
       Edge &e = E[ei]:
       if (lvl[e.v] != lvl[u] + 1) continue;
       ll a = push(e.v. min(e.c - e.f. p)):
       if (a <= 0) continue; e.f += a, E[ei ^ 1].f -= a;</pre>
       return a:
   }
    return 0;
11 maxflow() {
   11 f = 0:
    while (true) {
       // bfs to build levels
       lvl.assign(N, -1); queue < int > q; lvl[s] = 0, q.
            push(s);
       while (!q.empty()) {
           int u = q.front(); q.pop();
           for (int ei : G[u]) {
               Edge &e = E[ei]:
               if (e.c - e.f <= 0 || lvl[e.v] != -1)</pre>
                   continue:
               lvl[e.v] = lvl[u] + 1, q.push(e.v);
       }
       if (lvl[t] == -1) break;
       // dfs to find blocking flow
       ptr.assign(N, 0); while (11 ff = push(s, INF)) f
            += ff:
   }
    return f;
```

## 4.3 floyd-warshall

4.4 heavy-light

```
struct Hld {
   vector<int> P, H, D, pos, top;
   Hld() {}
   void init(vector<vector<int>> &G) {
      int N = G.size():
      P.resize(N), H.resize(N), D.resize(N), pos.resize(N),
          top.resize(N);
      D[0] = -1, dfs(G, 0); int t = 0;
      rep(i, N) if (H[P[i]] != i) {
          int i = i:
          while (j != -1)
             \{ top[i] = i, pos[i] = t++; i = H[i]; \}
      }
   }
   int dfs(vector<vector<int>> &G. int i) {
      int w = 1. mw = 0:
      D[i] = D[P[i]] + 1, H[i] = -1;
      for (int c : G[i]) {
          if (c == P[i]) continue;
          P[c] = i: int sw = dfs(G, c): w += sw:
          if (sw > mw) H[i] = c, mw = sw;
      }
       return w;
   template <class OP>
   void path(int u, int v, OP op) {
      while (top[u] != top[v]) {
          if (D[top[u]] > D[top[v]]) swap(u, v);
          op(pos[top[v]], pos[v] + 1); v = P[top[v]];
      if (D[u] > D[v]) swap(u, v);
       op(pos[u], pos[v] + 1); // value on vertex
      // op(pos[u]+1, pos[v] + 1); // value on path
   // segment tree
   template <class T, class S>
   void update(S &seg. int i. T val) {
      seg.update(pos[i], val);
   // segment tree lazy
```

```
template <class T, class S>
void update(S &seg, int u, int v, T val) {
    path(u, v, [&](int 1, int r) { seg.update(1, r, val);
      });
}

template <class T, class S>
T query(S &seg, int u, int v) {
    T ans = 0;
      // neutral element
    path(u, v, [&](int 1, int r) { ans += seg.query(1, r)
      ; }); // query op
    return ans;
}
};
```

## 4.5 hungarian

```
// find a maximum gain perfect matching in the given
    bipartite complete graph.
// input: gain matrix (G_{xy} = benefit of joining vertex x
    in set X with vertex
// v in set Y).
// output: maximum gain matching in members 'xy[x]' and 'yx[
    v]'.
// runtime: O(N^3)
struct Hungarian {
   int N, qi, root;
   vector<vector<ll>>> gain;
   vector<int> xy, yx, p, q, slackx;
   vector<ll> lx, ly, slack;
   vector<bool> S. T:
   void add(int x, int px) {
      S[x] = true, p[x] = px;
       rep(y, N) if (lx[x] + ly[y] - gain[x][y] < slack[y])
          slack[y] = lx[x] + ly[y] - gain[x][y], slackx[y]
               = x:
      }
   }
   void augment(int x, int y) {
       while (x != -2) {
          yx[y] = x; swap(xy[x], y); x = p[x];
   }
   void improve() {
```

```
S.assign(N, false), T.assign(N, false), p.assign(N,
   qi = 0, q.clear();
   rep(x, N) if (xv[x] == -1) {
       q.push_back(root = x), p[x] = -2, S[x] = true;
       break:
   rep(v, N) slack[v] = lx[root] + lv[v] - gain[root][v
        ], slackx[y] = root;
    while (true) {
       while (ai < a.size()) {</pre>
           int x = q[qi++];
           rep(y, N) if (lx[x] + ly[y] == gain[x][y] &&!
               T[v]) {
               if (yx[y] == -1) return augment(x, y);
              T[y] = true, q.push_back(yx[y]), add(yx[y
                   ], x);
          }
       }
       11 d = INF:
       rep(y, N) if (!T[y]) d = min(d, slack[y]);
       rep(x, N) if (S[x]) lx[x] -= d;
       rep(y, N) if (T[y]) ly[y] += d;
       rep(y, N) if (!T[y]) slack[y] -= d;
       rep(y, N) if (!T[y] && slack[y] == 0) {
          if (yx[y] == -1) return augment(slackx[y], y);
          T[v] = true:
           if (!S[yx[y]]) q.push_back(yx[y]), add(yx[y],
               slackx[v]);
       }
   }
}
Hungarian(vector<vector<ll>>> g)
   : N(g.size()), gain(g), xy(N, -1), yx(N, -1), lx(N, -
   ly(N), slack(N), slackx(N) {
   rep(x, N) rep(y, N) lx[x] = max(lx[x], ly[y]);
   rep(i, N) improve();
```

### 4.6 kuhn

};

// get a maximum cardinality matching in a bipartite graph.
// input: adjacency lists.

```
// output: matching (in 'mt' member).
// runtime: O(V E)
struct Kuhn {
   int N. size:
   vector<vector<int>> G;
   vector<bool> seen:
   vector<int> mt;
   bool visit(int i) {
       if (seen[i]) return false;
       seen[i] = true:
       for (int to : G[i])
          if (mt[to] == -1 || visit(mt[to])) {
              mt[to] = i:
              return true:
          }
       return false:
   Kuhn(vector<vector<int>> adj) : G(adj), N(G.size()), mt(N
        , -1) {
       rep(i, N) {
          seen.assign(N, false);
           size += visit(i):
       }
};
```

#### 4.7 lca

```
// calculates the lowest common ancestor for any two nodes
    in O(log N) time.
// with O(N log N) preprocessing
struct Lca {
   int L:
   vector<vector<int>> up;
   vector<pair<int, int>> time;
   Lca() {}
   void init(const vector<vector<int>> &G) {
       int N = G.size(); L = N <= 1 ? 0 : 32 - __builtin_clz</pre>
            (N - 1):
       up.resize(L + 1); rep(l, L + 1) up[l].resize(N);
       time.resize(N); int t = 0; visit(G, 0, 0, t);
       rep(1, L) rep(i, N) up[1 + 1][i] = up[1][up[1][i]]:
   void visit(const vector<vector<int>> &G, int i, int p,
        int &t) {
```

```
up[0][i] = p;
       time[i].first = t++:
       for (int edge : G[i]) {
          if (edge == p) continue;
          visit(G, edge, i, t);
       time[i].second = t++;
   }
   bool is_anc(int up, int dn) {
       return time[up].first <= time[dn].first &&
             time[dn].second <= time[up].second:</pre>
   }
   int get(int i, int j) {
       if (is_anc(i, j)) return i;
       if (is_anc(j, i)) return j;
      int 1 = L:
       while (1 >= 0) {
          if (is_anc(up[1][i], j)) 1--;
          else i = up[1][i];
       return up[0][i];
};
```

## 4.8 maxflow-mincost

```
// untested
#include "../common.h"
const 11 INF = 1e18;
struct Edge {
   int u, v;
   11 c. w. f = 0:
// find the minimum-cost flow among all maximum-flow flows.
// time: O(F V E)
                          F is the maximum flow
        O(V E + F E log V) if bellman-ford is replaced by
    johnson
struct Flow {
   int N, s, t;
   vector<vector<int>> G;
   vector<Edge> E;
   vector<ll> d:
```

```
vector<int> p:
   Flow() {}
   Flow(int N, int s, int t) : N(N), s(s), t(t), G(N) {}
   void add edge(int u, int v, ll c, ll w) {
      G[u].push_back(E.size());
       E.push_back({u, v, c, w});
      G[v].push_back(E.size());
       E.push_back({v, u, 0, -w});
   void calcdists() {
       // replace bellman-ford with johnson for better time
       d.assign(N, INF);
      p.assign(N, -1);
      d[s] = 0:
       rep(i, N - 1) rep(ei, E.size()) {
          Edge &e = E[ei]:
          ll n = d[e.u] + e.w:
          if (d[e.u] < INF && e.c - e.f > 0 && n < d[e.v])
               d[e.v] = n, p[e.v] = ei;
      }
   }
   ll maxflow() {
      11 ff = 0:
       while (true) {
          calcdists();
          if (p[t] == -1) break:
          11 f = INF;
          int cur = t:
          while (p[cur] != -1) {
              Edge &e = E[p[cur]];
              f = min(f, e.c - e.f):
              cur = e.u:
          int cur = t:
          while (p[cur] != -1) {
              E[p[cur]].f += f;
              E[p[cur] ^ 1].f -= f;
          ff += f:
      }
      return ff:
   }
};
```

### 4.9 push-relabel

```
#include "../common.h"
const 11 INF = 1e18;
// maximum flow algorithm.
// to run. use 'maxflow()'.
// time: O(V^2 \operatorname{sqrt}(E)) \leq O(V^3)
// memory: 0(V^2)
struct PushRelabel {
   vector<vector<ll>> cap, flow;
   vector<ll> excess:
   vector<int> height;
   PushRelabel() {}
   void resize(int N) { cap.assign(N, vector<11>(N)); }
   // push as much excess flow as possible from u to v.
   void push(int u, int v) {
       11 f = min(excess[u], cap[u][v] - flow[u][v]);
       flow[u][v] += f:
       flow[v][u] -= f:
       excess[v] += f:
       excess[u] -= f;
   // relabel the height of a vertex so that excess flow may
         be pushed.
   void relabel(int u) {
       int d = INT32 MAX:
       rep(v, cap.size()) if (cap[u][v] - flow[u][v] > 0) d
           min(d, height[v]):
       if (d < INF) height[u] = d + 1;</pre>
   // get the maximum flow on the network specified by 'cap'
         with source 's'
   // and sink 't'.
    // node-to-node flows are output to the 'flow' member.
   11 maxflow(int s, int t) {
       int N = cap.size(), M;
       flow.assign(N, vector<ll>(N));
       height.assign(N, 0), height[s] = N:
       excess.assign(N, 0), excess[s] = INF;
       rep(i, N) if (i != s) push(s, i);
       vector<int> q;
```

```
while (true) {
          // find the highest vertices with excess
          q.clear(), M = 0;
          rep(i, N) {
              if (excess[i] <= 0 || i == s || i == t)</pre>
                   continue:
              if (height[i] > M) q.clear(), M = height[i];
              if (height[i] >= M) q.push_back(i);
          if (q.empty()) break;
          // process vertices
          for (int u : a) {
              bool relab = true;
              rep(v. N) {
                  if (excess[u] <= 0) break;</pre>
                  if (cap[u][v] - flow[u][v] > 0 && height[u]
                      1 > height[v])
                      push(u, v), relab = false;
              if (relab) {
                  relabel(u);
                  break:
      11 f = 0; rep(i, N) f += flow[i][t]; return f;
   }
};
```

## ${\bf 4.10}\quad {\bf strongly\text{-}connected\text{-}components}$

```
// compute strongly connected components.
// time: O(V + E), memory: O(V)
//
// after building:
// comp = map from vertex to component (components are toposorted, root first, leaf last)
// N = number of components
// G = condensation graph (component DAG)
//
// byproducts:
// vgi = transposed graph
// order = reverse topological sort (leaf first, root last)
//
// others:
// vn = number of vertices
// vg = original vertex graph
struct Scc {
```

```
int vn. N:
   vector<int> order, comp;
   vector<vector<int>> vg, vgi, G;
   void toposort(int u) {
       if (comp[u]) return:
       comp[u] = -1;
       for (int v : vg[u]) toposort(v);
       order.push_back(u);
   bool carve(int u) {
       if (comp[u] != -1) return false;
       comp[u] = N;
       for (int v : vgi[u]) {
          carve(v);
          if (comp[v] != N) G[comp[v]].push_back(N);
      }
       return true:
   }
   Scc() {}
   Scc(vector<vector<int>> &g) : vn(g.size()), vg(g), comp(
        vn), vgi(vn), G(vn), N(0) {
       rep(u, vn) toposort(u);
       rep(u, vn) for (int v : vg[u]) vgi[v].push_back(u);
       invrep(i, vn) N += carve(order[i]);
   }
};
```

#### 4.11 two-sat

```
// calculate the solvability of a system of logical
    equations, where every equation is of the form 'a or b
// 'neg': get negation of 'u'
// 'then': 'u' implies 'v'
// 'any': 'u' or 'v'
// 'set': 'u' is true
11
// after 'solve' (O(V+E)) returns true, 'sol' contains one
    possible solution.
// determining all solutions is O(V*E) hard (requires
    computing reachability in a DAG).
struct TwoSat {
   int N; vector<vector<int>> G;
   Scc scc; vector<bool> sol;
   TwoSat(int n) : N(n), G(2 * n), sol(n) {}
   TwoSat() {}
```

## 5 implementation

## 5.1 SegmentTreeBeats

```
struct Node {
   11 \text{ s. mx1. mx2. mxc. mn1. mn2. mnc. } 1z = 0:
   Node(): s(0). mx1(LLONG MIN). mx2(LLONG MIN). mxc(0).
        mn1(LLONG_MAX), mn2(LLONG_MAX), mnc(0) {}
   Node(ll x) : s(x), mx1(x), mx2(LLONG_MIN), mxc(1), mn1(x)
        , mn2(LLONG_MAX), mnc(1) {}
   Node(const Node &a. const Node &b) {
       // add
       s = a.s + b.s;
       // min
       if (a.mx1 > b.mx1) mx1 = a.mx1, mxc = a.mxc, mx2 =
            max(b.mx1. a.mx2):
       if (a.mx1 < b.mx1) mx1 = b.mx1, mxc = b.mxc, mx2 =
            max(a.mx1. b.mx2):
       if (a.mx1 == b.mx1) mx1 = a.mx1, mxc = a.mxc + b.mxc,
             mx2 = max(a.mx2, b.mx2):
       // max
       if (a.mn1 < b.mn1) mn1 = a.mn1, mnc = a.mnc, mn2 =
            min(b.mn1, a.mn2):
       if (a.mn1 > b.mn1) mn1 = b.mn1, mnc = b.mnc, mn2 =
            min(a.mn1, b.mn2):
       if (a.mn1 == b.mn1) mn1 = a.mn1, mnc = a.mnc + b.mnc,
             mn2 = min(a.mn2, b.mn2);
};
// 0 - indexed / inclusive - inclusive
template <class node>
```

```
struct STB {
   vector<node> st; int n;
   void build(int u. int i. int i. vector<node> &arr) {
      if (i == i) {
          st[u] = arr[i]:
          return;
       int m = (i + j) / 2, l = u * 2 + 1, r = u * 2 + 2;
       build(1, i, m, arr), build(r, m + 1, j, arr);
       st[u] = node(st[l], st[r]):
   void push_add(int u, int i, int j, ll v) {
       st[u].s += (j - i + 1) * v;
       st[u].mx1 += v, st[u].mn1 += v, st[u].lz += v;
      if (st[u].mx2 != LLONG_MIN) st[u].mx2 += v;
       if (st[u].mn2 != LLONG MAX) st[u].mn2 += v:
   void push max(int u, ll v, bool l) { // for min op
       if (v >= st[u].mx1) return:
       st[u].s -= st[u].mx1 * st[u].mxc;
      st[n].mx1 = v:
       st[u].s += st[u].mx1 * st[u].mxc;
       if (1) st[u].mn1 = st[u].mx1:
       else if (v <= st[u].mn1) st[u].mn1 = v;</pre>
       else if (v < st[u].mn2) st[u].mn2 = v;
   void push_min(int u, ll v, bool l) { // for max op
       if (v <= st[u].mn1) return;</pre>
       st[u].s -= st[u].mn1 * st[u].mnc:
       st[u].mn1 = v:
       st[u].s += st[u].mn1 * st[u].mnc;
       if (1) st[u].mx1 = st[u].mn1:
       else if (v \ge st[u].mx1) st[u].mx1 = v;
       else if (v > st[u].mx2) st[u].mx2 = v:
   void push(int u, int i, int j) {
       if (i == j) return;
       // add
       int m = (i + j) / 2, l = u * 2 + 1, r = u * 2 + 2;
       push add(l, i, m, st[u].lz):
      push_add(r, m + 1, j, st[u].lz);
      st[u].lz = 0:
       push_max(1, st[u].mx1, i == m);
       push_max(r, st[u].mx1, m + 1 == j);
       push_min(1, st[u].mn1, i == m);
       push_min(r, st[u].mn1, m + 1 == r);
```

```
node query(int a, int b, int u, int i, int j) {
   if (b < i || i < a) return node():</pre>
   if (a <= i && j <= b) return st[u];</pre>
   push(u, i, i):
   int m = (i + j) / 2, l = u * 2 + 1, r = u * 2 + 2;
   return node(querv(a, b, l, i, m), querv(a, b, r, m +
void update_add(int a, int b, ll v, int u, int i, int j)
   if (b < i | | i < a) return:
   if (a <= i && i <= b) {
       push_add(u, i, j, v);
       return:
   push(u, i, j);
   int m = (i + j) / 2, l = u * 2 + 1, r = u * 2 + 2;
   update_add(a, b, v, l, i, m);
   update add(a, b, v, r, m + 1, i):
   st[u] = node(st[l], st[r]):
void update min(int a, int b, ll v, int u, int i, int i)
   if (b < i || j < a || v >= st[u].mx1) return;
   if (a <= i && j <= b && v > st[u].mx2) {
       push_max(u, v, i == j);
       return:
   push(u, i, j);
   int m = (i + i) / 2, l = u * 2 + 1, r = u * 2 + 2;
   update_min(a, b, v, 1, i, m);
   update_min(a, b, v, r, m + 1, j);
   st[u] = node(st[l], st[r]);
void update_max(int a, int b, ll v, int u, int i, int j)
   if (b < i || j < a || v <= st[u].mn1) return;</pre>
   if (a <= i && j <= b && v < st[u].mn2) {</pre>
       push_min(u, v, i == j);
       return:
   push(u, i, j);
   int m = (i + i) / 2, l = u * 2 + 1, r = u * 2 + 2:
   update_max(a, b, v, l, i, m);
   update_max(a, b, v, r, m + 1, j);
   st[u] = node(st[l], st[r]):
STB(vector<node> &v, int N) : n(N), st(N * 4 + 5) { build
     (0, 0, n - 1, v);
```

## 5.2 Treap

```
#include "../Template.cpp"
mt19937 gen(chrono::high resolution clock::now().
    time_since_epoch().count());
// 101 Implicit Treap //
struct Node
   int p, sz = 0, v, acc, l = -1, r = -1;
   Node() : v(0), acc(0) {}
   Node(int x): p(gen()), sz(1), v(x), acc(x) {}
   void recalc(const Node &a, const Node &b)
       sz = a.sz + b.sz + 1;
       acc = v + a.acc + b.acc;
};
template <class node>
struct Treap
   vector<node> t; int n, r = -1;
   node get(int u) { return u != -1 ? t[u] : node(): }
   void recalc(int u) { t[u].recalc(get(t[u].l), get(t[u].r) }
        ): }
   int merge(int 1, int r)
       if (min(1, r) == -1) return 1 != -1 ? 1 : r:
       int ans = (t[1].p < t[r].p) ? 1 : r;
       if (ans == 1) t[1].r = merge(t[1].r, r), recalc(1);
       if (ans == r) t[r].l = merge(l, t[r].l), recalc(r):
       return ans;
    pii split(int u, int id)
```

```
if (u == -1) return {-1, -1}:
       int szl = get(t[u].1).sz;
       if (szl >= id)
           pii ans = split(t[u].1, id);
           t[u].1 = ans.ss: recalc(u):
           return {ans.ff, u};
       pii ans = split(t[u].r, id - szl - 1);
       t[u].r = ans.ff; recalc(u);
       return {u, ans.ss}:
   Treap(vi &v) : n(sz(v))
   { for (int i = 0; i < n; i++) t.eb(v[i]), r = merge(r, i)
        ; }
}:
// Complete Implicit Treap with Lazy propagation //
struct Node
   int p, sz = 0, v, acc, l = -1, r = -1, par = -1, lzv = 0;
   bool lz = false, f = false;
   Node() : v(0), acc(0) {}
   Node(int x): p(gen()), sz(1), v(x), acc(x) {}
   void recalc(const Node &a. const Node &b)
       sz = a.sz + b.sz + 1;
       acc = v + a.acc + b.acc:
   void upd_lazv(int x) { lz = 1, lzv += x; }
   void lazy() { v += 1zv, acc += sz * 1zv, 1z = 0, 1zv = 0;
   void flip() { swap(1, r), f = 0; }
template <class node>
struct Treap
   vector<node> t: int n. r = -1:
   node get(int u) { return u != -1 ? t[u] : node(): }
   void recalc(int u)
       int l = t[u].l. r = t[u].r:
       push(1); push(r); flip(1); flip(r);
       t[u].recalc(get(1), get(r));
   void push(int u)
```

```
if (u == -1 || !t[u].lz) return:
   int 1 = t[u].1, r = t[u].r;
   if (1 != -1) t[1].upd_lazy(t[u].lzv);
   if (r != -1) t[r].upd_lazy(t[u].lzv);
   t[u].lazv():
}
void flip(int u)
   if (u == -1 || !t[u].f) return;
   int l = t[u].l. r = t[u].r:
   if (1 != -1) t[1].f ^= 1;
   if (r != -1) t[r].f ^= 1;
   t[u].flip();
int merge(int 1, int r)
    if (min(1, r) == -1) return 1 != -1 ? 1 : r;
    push(1): push(r): flip(1): flip(r):
   int ans = (t[1].p < t[r].p) ? 1 : r;</pre>
   if (ans == 1) t[1].r = merge(t[1].r, r), recalc(1);
    if (ans == r) t[r].l = merge(l, t[r].l), recalc(r);
    if (t[ans].1 != -1) t[t[ans].1].par = ans; // only if
         parent needed
    if (t[ans].r != -1) t[t[ans].r].par = ans; // only if
         parent needed
   return ans:
pii split(int u, int id)
    if (u == -1) return {-1, -1};
    push(u); flip(u);
   int szl = get(t[u].1).sz;
   if (szl >= id)
       pii ans = split(t[u].1, id):
       if (ans.ss != -1) t[ans.ss].par = u; // only if
            parent needed
       if (ans.ff != -1) t[ans.ff].par = -1; // only if
            parent needed
       t[u].1 = ans.ss: recalc(u):
       return {ans.ff, u};
   pii ans = split(t[u].r, id - szl - 1);
    if (ans.ff != -1) t[ans.ff].par = u; // only if
        parent needed
    if (ans.ss != -1) t[ans.ss].par = -1; // only if
        parent needed
   t[u].r = ans.ff; recalc(u);
    return {u, ans.ss};
```

```
}
int update(int u, int l, int r, int v)
{
    pii a = split(u, l), b = split(a.ss, r - l + 1);
        t[b.ff].upd_lazy(v);
    return merge(a.ff, merge(b.ff, b.ss));
}
void print(int u)
{
    if (u == -1) return;
    push(u); flip(u);
    print(t[u].1);
    cout << t[u].v << ' ';
    print(t[u].r);
}
Treap(vi &v) : n(sz(v))
{ for (int i = 0; i < n; i++) t.eb(v[i]), r = merge(r, i)
    ; }
};</pre>
```

#### 5.3 dsu

```
struct Dsu {
   vector<int> p, r:
   // initialize the disjoint-set-union to all unitary sets
   void reset(int N) {
       p.resize(N), r.assign(N, 0);
       rep(i, N) p[i] = i;
   // find the leader node corresponding to node 'i'
    int find(int i) {
       if (p[i] != i) p[i] = find(p[i]);
       return p[i];
   // perform union on the two sets that 'i' and 'j' belong
   void unite(int i, int j) {
       i = find(i), j = find(j);
       if (i == j) return;
       if (r[i] > r[j]) swap(i, j);
       if (r[i] == r[j]) r[j] += 1;
       p[i] = j;
};
```

#### 5.4 mo

```
struct Query { int 1, r, idx; };
// answer segment queries using only 'add(i)', 'remove(i)'
    and 'get()'
// functions.
// complexity: O((N + Q) * sqrt(N) * F)
// N = length of the full segment
// Q = amount of gueries
// F = complexity of the 'add', 'remove' functions
template <class A. class R. class G. class T>
void mo(vector<Query> &queries, vector<T> &ans, A add, R
    remove, G get) {
   int Q = queries.size(), B = (int)sqrt(Q);
   sort(queries.begin(), queries.end(), [&](Query &a, Query
       return make pair(a.1 / B. a.r) < make pair(b.1 / B. b.
   }):
   ans.resize(0):
   int 1 = 0, r = 0:
   for (auto &q : queries) {
       while (r < q.r) add(r), r++;
       while (1 > q.1) 1--, add(1);
       while (r > q.r) r--, remove(r);
       while (1 < q.1) remove(1), 1++;</pre>
       ans[a.idx] = get():
   }
```

## 5.5 persistent-segment-tree-lazy

```
template <class T>
struct Node {
    T x, lz;
    int l = -1, r = -1;
};

template <class T>
struct Pstl {
    int N;
    vector<Node<T>> a;
    vector<int> head;

    T qneut() { return 0; }
    T merge(T l, T r) { return l + r; }
```

```
T uneut() { return 0: }
T accum(T u, T x) { return u + x; }
T apply(T x, T lz, int l, int r) { return x + (r - 1) *
     1z: }
int build(int vl. int vr) {
   if (vr - vl == 1) a.push_back({qneut(), uneut()}); //
         node construction
   else {
       int vm = (vl + vr) / 2, l = build(vl, vm), r =
            build(vm, vr):
       a.push back({merge(a[1].x.a[r].x), uneut(), 1, r
            }); // query merge
   return a.size() - 1:
T query(int 1, int r, int v, int v1, int vr, T acc) {
   if (1 >= vr || r <= vl) return gneut();</pre>
        // guery neutral
   if (1 <= v1 && r >= vr) return apply(a[v].x, acc, v1,
         vr): // update op
    acc = accum(acc, a[v].lz);
        // update merge
   int vm = (vl + vr) / 2:
   return merge(query(1, r, a[v].1, v1, vm, acc), query(
        1, r, a[v].r, vm, vr, acc)); // query merge
int update(int 1, int r, T x, int v, int v1, int vr) {
   if (1 \ge vr \mid | r \le vl \mid | r \le 1) return v:
   a.push_back(a[v]);
   v = a.size() - 1:
   if (1 <= v1 && r >= vr) {
       a[v].x = apply(a[v].x, x, vl, vr); // update op
       a[v].lz = accum(a[v].lz, x): // update merge
   } else {
       int vm = (vl + vr) / 2:
       a[v].1 = update(1, r, x, a[v].1, v1, vm);
       a[v].r = update(1, r, x, a[v].r, vm, vr);
       a[v].x = merge(a[a[v].1].x, a[a[v].r].x); //
            query merge
   }
    return v;
Pstl() {}
Pstl(int N) : N(N) { head.push_back(build(0, N)); }
T query(int t, int 1, int r) {
```

## 5.6 persistent-segment-tree

```
// usage:
// Pst<Node<11>> pst;
// pst = {N};
// int newtime = pst.update(time, index, value):
// Node<ll> result = pst.query(newtime, left, right);
template <class T>
struct Node {
    T x;
    int 1 = -1, r = -1:
    Node(): x(0) {}
    Node(T x) : x(x)  {}
    Node (Node a, Node b, int l = -1, int r = -1) : x(a.x + b.
        x), 1(1), r(r) {}
};
template <class U>
struct Pst {
    int N:
    vector<U> a:
    vector<int> head;
    int build(int vl, int vr) {
       if (vr - vl == 1) a.push_back(U()); // node
            construction
       else {
           int vm = (vl + vr) / 2, l = build(vl, vm), r =
                build(vm. vr):
           a.push_back(U(a[1], a[r], 1, r)); // query merge
       return a.size() - 1;
    U query(int 1, int r, int v, int v1, int vr) {
       if (1 >= vr || r <= vl) return U(); // query neutral</pre>
       if (1 <= v1 && r >= vr) return a[v]:
       int vm = (vl + vr) / 2:
```

```
return U(query(1, r, a[v].1, v1, vm), query(1, r, a[v
           ].r, vm, vr)); // query merge
   }
   int update(int i, U x, int v, int vl, int vr) {
       a.push back(a[v]):
       v = a.size() - 1;
      if (vr - vl == 1) a[v] = x; // update op
          int vm = (vl + vr) / 2;
          if (i < vm) a[v].l = update(i, x, a[v].l, vl, vm)</pre>
          else a[v].r = update(i, x, a[v].r, vm, vr);
          a[v] = U(a[a[v].1], a[a[v].r], a[v].1, a[v].r);
               // query merge
       return v:
   }
   Pst() {}
   Pst(int N) : N(N) { head.push_back(build(0, N)); }
   U query(int t, int 1, int r) {
       return query(1, r, head[t], 0, N);
   int update(int t, int i, U x) {
       return head.push_back(update(i, x, head[t], 0, N)),
           head.size() - 1:
   }
}:
```

## |5.7 segment-tree-lazy

```
Stl(int N) : a(4 * N, {gneutral(), uneutral()}) {} //
        node neutral
   void push(int v, int vl, int vm, int vr) {
       update(a[2 * v], a[v].second, vl, vm); // node update
       update(a[2 * v + 1], a[v].second, vm, vr); // node
       a[v].second = uneutral();
                                              // update
           neutral
   // guerv for range [1, r)
   T query(int 1, int r, int v = 1, int vl = 0, int vr = -1)
       if (vr == -1) vr = a.size() / 4;
       if (1 <= vl && r >= vr) return a[v].first; // query
       if (1 >= vr || r <= vl) return qneutral(); // query</pre>
           neutral
       int vm = (vl + vr) / 2:
       push(v, v1, vm, vr);
       return merge(query(1, r, 2 * v, v1, vm), query(1, r,
           2 * v + 1, vm, vr)); // item merge
   // update range [1, r) using val
   void update(int 1, int r, T val, int v = 1, int vl = 0,
        int vr = -1) {
       if (vr == -1) vr = a.size() / 4;
       if (1 >= vr || r <= vl || r <= 1) return:
       if (1 <= vl && r >= vr) update(a[v], val, vl, vr); //
             node update
       else {
          int vm = (v1 + vr) / 2;
          push(v. vl. vm. vr):
          update(1, r, val, 2 * v, vl, vm):
          update(1, r, val, 2 * v + 1, vm, vr);
          a[v].first = merge(a[2 * v].first, a[2 * v + 1].
               first); // node merge
      }
   }
};
```

## 5.8 segment-tree

```
// usage:

// St<Node<11>> st;

// st = {N};

// st.update(index, new_value);
```

```
// Node<ll> result = st.querv(left, right):
template <class T>
struct Node {
   T x;
   Node(): x(0) {}
   Node(T x) : x(x) \{ \}
   Node(Node a, Node b) : x(a.x + b.x) {}
};
template <class U>
struct St {
   vector<U> a;
   St() {}
   St(int N) : a(4 * N, U()) {} // node neutral
   // query for range [1, r)
   U querv(int 1, int r, int v = 1, int vl = 0, int vr = -1)
       if (vr == -1) vr = a.size() / 4;
       if (1 \le v1 \&\& r \ge vr) return a[v]: // item
            construction
       int vm = (vl + vr) / 2;
       if (1 >= vr || r <= vl) return U();</pre>
                                        // item neutral
       return U(query(1, r, 2 * v, v1, vm), query(1, r, 2 *
            v + 1, vm, vr)): // item merge
   }
   // set element i to val
   void update(int i, U val, int v = 1, int vl = 0, int vr =
         -1) {
       if (vr == -1) vr = a.size() / 4;
       if (vr - vl == 1) a[v] = val: // item update
       else {
           int vm = (vl + vr) / 2;
           if (i < vm) update(i, val, 2 * v, vl, vm);</pre>
           else update(i, val, 2 * v + 1, vm, vr);
           a[v] = U(a[2 * v], a[2 * v + 1]); // node merge
       }
};
```

## 5.9 sparse-table

```
// handle immutable range maximum queries (or any idempotent
    query) in O(1)
template <class T>
    struct Hash {
    size_t ope
    static
```

```
struct Sparse {
   vector<vector<T>> st;
   T op(T a, T b) { return max(a, b); }
   Sparse() {}
   void reset(int N) { st = {vector<T>(N)}; }
   void set(int i, T val) { st[0][i] = val; }
   // O(N log N) time
   // O(N log N) memory
   void init() {
       int N = st[0].size();
       int npot = N <= 1 ? 1 : 32 - __builtin_clz(N);</pre>
       st.resize(npot);
       repx(i, 1, npot) rep(j, N + 1 - (1 << i)) st[i].
            push_back(
          op(st[i-1][j], st[i-1][j+(1 << (i-1))]);
                // query op
   }
   // query maximum in the range [1, r) in O(1) time
   // range must be nonempty!
   T query(int 1, int r) {
       int i = 31 - \_builtin\_clz(r - 1);
       return op(st[i][1], st[i][r - (1 << i)]); // query op</pre>
   }
};
```

## 5.10 unordered-map

```
// hackproof rng
static mt19937 rng(chrono::steady_clock::now().
    time_since_epoch().count());

// deterministic rng
uint64_t splitmix64(uint64_t *x) {
    uint64_t z = (*x += 0x9e3779b97f4a7c15);
    z = (z ^ (z >> 30)) * 0xbf58476d1ce4e5b9;
    z = (z ^ (z >> 27)) * 0x94d049bb133111eb;
    return z ^ (z >> 31);
}

// hackproof unordered map hash
struct Hash {
    size_t operator()(const l1 &x) const {
        static const uint64_t RAND =
```

20

## 6 imprimible

### 7 math

#### 7.1 arithmetic

```
// floor(log2(n)) without precision loss
inline int floor_log2(int n) { return n <= 1 ? 0 : 31 -</pre>
     builtin_clz(n); }
// ceil(log2(n)) without precision loss
inline int ceil_log2(int n) { return n <= 1 ? 0 : 32 -</pre>
    builtin clz(n - 1): }
inline 11 floordiv(11 a, 11 b) {
   return a / b - ((a ^ b) < 0 && a % b):
inline 11 ceildiv(11 a, 11 b) {
   return a / b + ((a ^ b) >= 0 && a % b);
// a^e through binary exponentiation.
ll binexp(ll a, ll e) {
   11 res = 1; // neutral element
       if (e & 1) res = res * a: // multiplication
                               // multiplication
       a = a * a:
       e >>= 1:
   }
   return res;
```

-

#### 7.2 crt

## 7.3 discrete-log

```
// discrete logarithm log_a(b).
// solve b \hat{x} = a \pmod{M} for the smallest x.
// returns -1 if no solution is found.
// time: O(sqrt(M))
11 dlog(ll a, ll b, ll M) {
   11 k = 1, s = 0;
   while (true) {
       11 g = \_gcd(b, M);
       if (g <= 1) break;</pre>
       if (a == k) return s:
       if (a % g != 0) return -1;
       a = g, M = g, s += 1, k = b / g * k % M;
   ll N = sqrt(M) + 1;
   umap<11, 11> r;
   rep(q, N + 1) {
       r[a] = q;
       a = a * b % M;
   ll bN = binexp(b, N, M), bNp = k;
   repx(p, 1, N + 1) {
       bNp = bNp * bN % M;
       if (r.count(bNp)) return N * p - r[bNp] + s;
```

```
}
return -1;
```

#### **7.4** gauss

```
const double EPS = 1e-9:
// solve a system of equations.
// complexity: O(min(N, M) * N * M)
11
// 'a' is a list of rows
// the last value in each row is the result of the equation
// return values:
// 0 -> no solutions
// 1 -> unique solution, stored in 'ans'
// -1 -> infinitely many solutions, one of which is stored
    in 'ans'
// UNTESTED
int gauss(vector<vector<double>> a, vector<double> &ans) {
   int N = a.size(), M = a[0].size() - 1:
   vector<int> where(M, -1);
   for (int j = 0, i = 0; j < M && i < N; j++) {
       int sel = i;
       repx(k, i, N) if (abs(a[k][j]) > abs(a[sel][j])) sel
       if (abs(a[sel][j]) < EPS) continue;</pre>
       repx(k, j, M + 1) swap(a[sel][k], a[i][k]);
       where[i] = i;
       rep(k, N) if (k != i) {
          double c = a[k][i] / a[i][i];
          repx(1, j, M + 1) a[k][1] -= a[i][1] * c;
       i++;
   }
   ans.assign(M, 0);
   rep(i, M) if (where[i] != -1) ans[i] = a[where[i]][M] / a
        [where[i]][i]:
   rep(i, N) {
       double sum = 0;
       rep(j, M) sum += ans[j] * a[i][j];
       if (abs(sum - a[i][M]) > EPS) return 0:
   }
   rep(i, M) if (where[i] == -1) return -1;
   return 1;
```

#### 7.5 matrix

```
using T = 11;
struct Mat {
   int N. M:
   vector<vector<T>> v;
   Mat(int n, int m) : N(n), M(m), v(N, vector<T>(M)) {}
   Mat(int n) : Mat(n, n) { rep(i, N) v[i][i] = 1; }
   vector<T> &operator[](int i) { return v[i]; }
   Mat operator*(Mat &r) {
       assert(M == r.N);
       int n = N, m = r.M, p = M:
      Mat a(n. m):
      rep(i, n) rep(j, m) {
          a[i][j] = T();
                                                       11
               neutral
          rep(k, p) a[i][k] = a[i][j] + v[i][k] * r[k][j];
               // mul. add
      }
       return a:
   Mat binexp(ll e) {
       assert(N == M);
      Mat a = *this, res(N); // neutral
      while (e) {
          if (e & 1) res = res * a; // mul
          a = a * a:
                                  // mul
          e >>= 1:
      }
       return res:
   friend ostream &operator<<(ostream &s, Mat &a) {</pre>
      rep(i, a.N) {
          rep(j, a.M) s << a[i][j] << " ";
          s << endl:
      }
      return s:
   }
};
```

#### 7.6 mod

```
11 binexp(ll a, ll e, ll M) {
   assert(e >= 0);
   ll res = 1 % M;
   while (e) {
       if (e & 1) res = res * a % M;
       a = a * a % M:
       e >>= 1:
   return res:
}
11 multinv(11 a, 11 M) { return binexp(a, M - 2, M); }
// calculate gcd(a, b).
// also, calculate x and y such that:
// a * x + b * y == gcd(a, b)
11
// time: O(log min(a, b))
// (ignoring complexity of arithmetic)
ll ext gcd(ll a, ll b, ll &x, ll &v) {
   if (b == 0) {
       x = 1, y = 0;
       return a:
   11 d = ext_gcd(b, a % b, y, x);
   v = a / b * x;
   return d:
// compute inverse with any M.
// a and M must be coprime for inverse to exist!
11 multinv_euc(ll a, ll M) {
   11 x. v:
   ext_gcd(a, M, x, y);
   return x;
}
// multiply two big numbers (~10^18) under a large modulo.
    without resorting to
// bigints.
11 bigmul(11 x, 11 y, 11 M) {
   11 z = 0;
   while (y) {
       if (v \& 1) z = (z + x) \% M:
       x = (x << 1) \% M, y >>= 1;
   return z;
```

```
// all modular inverses from 1 to inv.size()-1
void multinv all(vector<ll> &inv) {
   inv[1] = 1:
   repx(i, 2, inv.size())
      inv[i] = MOD - (MOD / i) * inv[MOD % i] % MOD:
struct Mod {
   int a:
   static const int M = 1e9 + 7:
   Mod(11 aa) : a((aa % M + M) % M) {}
   Mod operator+(Mod rhs) const { return (a + rhs.a) % M; }
   Mod operator-(Mod rhs) const { return (a - rhs.a + M) % M
   Mod operator-() const { return Mod(0) - *this; }
   Mod operator*(Mod rhs) const { return (11)a * rhs.a % M:
   Mod operator+=(Mod rhs) { return *this = *this + rhs; }
   Mod operator = (Mod rhs) { return *this = *this - rhs; }
   Mod operator*=(Mod rhs) { return *this = *this * rhs; }
   Mod bigmul(ll big) const { return ::bigmul(a, big, M); }
   Mod binexp(ll e) const { return ::binexp(a, e, M): }
   // Mod multinv() const { return ::multinv(a, M); } //
        prime M
   Mod multinv() const { return ::multinv_euc(a, M); } //
        possibly composite M
}:
// dvnamic modulus
struct DMod {
   int a. M:
   DMod(ll aa, ll m) : M(m), a((aa % m + m) % m) {}
   DMod operator+(DMod rhs) const { return {(a + rhs.a) % M.
   DMod operator-(DMod rhs) const { return {(a - rhs.a + M)
        % M. M}: }
   DMod operator-() const { return DMod(0, M) - *this; }
   DMod operator*(DMod rhs) const { return {(11)a * rhs.a %
        M, M}; }
   DMod operator+=(DMod rhs) { return *this = *this + rhs; }
   DMod operator = (DMod rhs) { return *this = *this - rhs; }
```

### 7.7 poly

```
using cd = complex<double>;
const double PI = acos(-1);
// compute the DFT of a power-of-two-length sequence.
// if 'inv' is true, computes the inverse DFT.
11
// the DFT of a polynomial A(x) = A0 + A1*x + A2*x^2 + ... +
     An*x^n is the array
// of the polynomial A evaluated in all nths roots of unity:
     [A(w0), A(w1),
// A(w2), ..., A(wn-1)], where w0 = 1 and w1 is the nth
    principal root of unity.
void fft(vector<cd> &a. bool inv) {
   int N = a.size(), k = 0:
   assert(N == 1 << __builtin_ctz(N));</pre>
   rep(i, N) {
       int b = N \gg 1;
       while (k \& b) k = b, b >>= 1:
       if (i < k) swap(a[i], a[k]);</pre>
   for (int 1 = 2: 1 <= N: 1 <<= 1) {
       double ang = 2 * PI / 1 * (inv ? -1 : 1):
       cd wl(cos(ang), sin(ang)):
      for (int i = 0: i < N: i += 1) {
          cd w(1);
          repx(j, 0, 1 / 2) {
              cd u = a[i + i], v = a[i + i + 1 / 2] * w:
              a[i + j] = u + v;
              a[i + j + 1 / 2] = u - v;
              w *= wl:
```

```
if (inv)
       for (cd &x : a) x \neq N;
const 11 MOD = 7340033, ROOT = 5, ROOTPOW = 1 << 20;</pre>
void find_root_of_unity(ll M) {
   11 c = M - 1, k = 0:
   while (c \% 2 == 0) c /= 2, k += 1:
   // find proper divisors of M - 1
   vector<int> divs:
   repx(d, 1, c) {
       if (d * d > c) break:
       if (c % d == 0) rep(i, k + 1) divs.push_back(d << i);</pre>
   rep(i, k) divs.push_back(c << i);</pre>
   // find any primitive root of M
   11 G = -1;
   repx(g, 2, M) {
       bool ok = true;
       for (int d : divs) ok &= (binexp(g, d, M) != 1);
       if (ok) {
           G = g:
           break;
       }
   assert(G != -1);
   ll w = binexp(G, c, M);
    cerr << M << " = c * 2^k + 1" << endl;
   cerr << " c = " << c << endl:
    cerr << " k = " << k << endl:
   cerr << "w^(2^k) == 1" << endl:
   cerr << " w = " << w << endl;
}
// compute the DFT of a power-of-two-length sequence, modulo
      a special prime
// number with principal root.
// the modulus must be a prime number with an Nth root of
     unity, where N is a
// power of two. the FFT can only be performed on arrays of
void ntt(vector<ll> &a. bool inv) {
```

```
int N = a.size(), k = 0:
    assert(N == 1 << __builtin_ctz(N) && N <= ROOTPOW);</pre>
    rep(i, N) a[i] = (a[i] % MOD + MOD) % MOD;
    repx(i, 1, N) {
       int b = N \gg 1:
        while (k \& b) k = b, b >>= 1;
       k ^= b:
       if (i < k) swap(a[i], a[k]);</pre>
    for (int 1 = 2: 1 <= N: 1 <<= 1) {
       11 wl = inv ? multinv(ROOT, MOD) : ROOT;
        for (ll i = ROOTPOW: i > 1: i >>= 1) wl = wl * wl %
        for (int i = 0; i < N; i += 1) {</pre>
           11 w = 1:
            repx(j, 0, 1 / 2) {
               11 u = a[i + i], v = a[i + i + 1 / 2] * w %
               a[i + j] = (u + v) \% MOD;
               a[i + j + 1 / 2] = (u - v + MOD) \% MOD;
               w = w * w1 \% MOD;
       }
    11 ninv = multinv(N, MOD);
    if (inv)
        for (11 &x : a) x = x * ninv % MOD:
 void convolve(vector<11> &a. vector<11> b. int n) {
    n = 1 \ll (32 - \_builtin\_clz(2 * n - 1));
    a.resize(n), b.resize(n):
    ntt(a, false), ntt(b, false);
    rep(i, n) a[i] *= b[i];
    ntt(a, true), ntt(b, true);
 using T = 11:
T pmul(T a, T b) { return a * b % MOD; }
T padd(T a, T b) { return (a + b) % MOD; }
T psub(T a, T b) { return (a - b + MOD) % MOD; }
 T pinv(T a) { return multinv(a, MOD); }
 struct Poly {
    vector<T> a:
    Poly() {}
```

```
Polv(T c) : a(c) { trim(): }
Poly(vector<T> c) : a(c) { trim(); }
void trim() {
    while (!a.empty() && a.back() == 0) a.pop_back();
int deg() const { return a.empty() ? -1000000 : a.size()
     - 1: }
Poly sub(int 1, int r) const {
   r = min(r, (int)a.size()), l = min(l, r);
   return vector<T>(a.begin() + 1, a.begin() + r);
Poly trunc(int n) const { return sub(0, n); }
Poly shl(int n) const {
    Polv out = *this:
    out.a.insert(out.a.begin(), n, 0);
   return out:
Polv rev(int n. bool r = false) const {
   Polv out(*this):
    if (r) out.a.resize(max(n, (int)a.size()));
   reverse(out.a.begin(), out.a.end());
   return out.trunc(n);
Poly &operator+=(const Poly &rhs) {
   auto &b = rhs.a:
   a.resize(max(a.size(), b.size()));
   rep(i, b.size()) a[i] = padd(a[i], b[i]); // add
   trim():
   return *this;
Poly & operator -= (const Poly &rhs) {
   auto &b = rhs.a;
   a.resize(max(a.size(), b.size()));
   rep(i, b.size()) a[i] = psub(a[i], b[i]); // sub
   trim():
   return *this:
Poly &operator*=(const Poly &rhs) {
   int n = deg() + rhs.deg() + 1:
   if (n <= 0) return *this = Poly();</pre>
   n = 1 \ll (n \ll 1?0:32 - builtin clz(n - 1)):
   vector<T> b = rhs.a:
   a.resize(n), b.resize(n);
   ntt(a, false), ntt(b, false);
   rep(i, a.size()) a[i] = pmul(a[i], b[i]); // mul
   ntt(a, true), trim();
                                         // invfft
   return *this:
```

```
Polv inv(int n) const {
   assert(deg() >= 0):
   Poly ans = pinv(a[0]); // inverse
   int b = 1:
   while (b < n) {
       Polv C = (ans * trunc(2 * b)).sub(b, 2 * b):
       ans -= (ans * C).trunc(b).shl(b);
      b *= 2:
   }
   return ans.trunc(n);
Poly operator+(const Poly &rhs) const { return Poly(*this
    ) += rhs: }
Poly operator-(const Poly &rhs) const { return Poly(*this
    ) -= rhs: }
Poly operator*(const Poly &rhs) const { return Poly(*this
    ) *= rhs: }
pair<Polv. Polv> divmod(const Polv &b) const {
   if (deg() < b.deg()) return {Poly(), *this};</pre>
   int d = deg() - b.deg() + 1;
   Poly D = (rev(d) * b.rev(d).inv(d)).trunc(d).rev(d,
   return \{D, *this - D * b\};
Poly operator/(const Poly &b) const { return divmod(b).
Poly operator%(const Poly &b) const { return divmod(b).
    second: }
Poly &operator/=(const Poly &b) { return *this = divmod(b
    ).first: }
Poly &operator%=(const Poly &b) { return *this = divmod(b
    ).second; }
T eval(T x) {
   T v = 0:
   invrep(i, a.size()) y = padd(pmul(y, x), a[i]); //
        add, mul
   return y;
Poly &build(vector<Poly> &tree, vector<T> &x, int v, int
    1. int r) {
   if (1 == r) return tree[v] = vector<T>{-x[1], 1};
   int m = (1 + r) / 2;
   return tree[v] = build(tree, x, 2 * v, 1, m) *
                  build(tree, x, 2 * v + 1, m + 1, r);
void subeval(vector<Poly> &tree, vector<T> &x, vector<T>
    &v, int v, int 1,
```

```
int r) {
       if (1 == r) {
          y[1] = eval(x[1]);
          return:
       int m = (1 + r) / 2:
       (*this % tree[2 * v]).subeval(tree, x, y, 2 * v, 1, m
       (*this % tree[2 * v + 1]).subeval(tree, x, v, 2 * v +
            1. m + 1. r):
   // evaluate m points in O(k (log k)^2) with k = max(n, m)
   vector<T> multieval(vector<T> &x) {
       int N = x.size():
       if (deg() < 0) return vector<T>(N, 0);
       vector<Polv> tree(4 * N):
       build(tree, x, 1, 0, N - 1);
       vector<T> y(N);
       subeval(tree, x, y, 1, 0, N - 1);
       return v;
   }
};
```

### 7.8 primes

```
// counts the divisors of a positive integer in O(\operatorname{sqrt}(n))
11 count divisors(11 x) {
    11 \text{ divs} = 1, i = 2:
    for (ll divs = 1, i = 2; x > 1; i++) {
       if (i * i > x) {
           divs *= 2:
           break;
       for (11 d = divs; x % i == 0; x /= i) divs += d;
    return divs:
// gets the prime factorization of a number in O(sqrt(n))
vector<pair<11, int>> factorize(11 x) {
    vector<pair<11, int>> f;
    for (11 k = 2; x > 1; k++) {
       if (k * k > x) {
           f.push back(\{x, 1\}):
           break;
       int n = 0:
       while (x \% k == 0) x /= k, n++;
```

```
if (n > 0) f.push back(\{k, n\}):
   return f;
// iterate over all divisors of a number.
// divisor count upper bound: n^(1.07 / ln ln n)
template <class OP>
void divisors(ll x, OP op) {
   auto facts = factorize(x):
   vector<int> f(facts.size()):
   while (true) {
       11 v = 1:
       rep(i, f.size()) rep(j, f[i]) v *= facts[i].first;
       (v)qo
       int i;
       for (i = 0: i < f.size(): i++) {</pre>
           f[i] += 1:
           if (f[i] <= facts[i].second) break;</pre>
       if (i == f.size()) break;
// computes euler totative function phi(x), counting the
    amount of integers in
// [1, x] that are coprime with x.
11
// time: O(sqrt(x))
11 phi(11 x) {
   11 \text{ phi} = 1, k = 2:
   for (; x > 1; k++) {
       if (k * k > x) {
           phi *= x - 1:
           break:
       }
       11 k1 = 1, k0 = 0;
       while (x \% k == 0) x /= k, k0 = k1, k1 *= k:
       phi *= k1 - k0:
   return phi;
// computes primality up to N.
// considers 0 and 1 prime.
// O(N log N)
void sieve(int N, vector<bool> &prime) {
   prime.assign(N + 1, true);
```

#### 7.9 theorems

```
// Burnside lemma
//

// For a set X, with members x in X, and a group G, with operations g in G, where g(x): X -> X.

// F_g is the set of x which are fixed points of g (ie. { x in X / g(x) = x }).

// The number of orbits (connected components in the graph formed by assigning each x a node and a directed edge between x and g(x) for every g) is called M.
```

```
M = the average of the fixed points of all g = (|F_g1|) / |
      + |F_g2| + ... + |F_gn|) / |G|
11
      If x are images and g are simmetries, then M
    corresponds to the amount of objects, |G|
      corresponds to the amount of simmetries, and F_g
    corresponds to the amount of simmetrical
      images under the simmetry g.
11
// Rational root theorem
11
      All rational roots of the polynomials with integer
    coefficients:
//
      a0 * x^0 + a1 * x^1 + a2 * x^2 + ... + an * x^n = 0
//
11
      If these roots are represented as p / q, with p and q
    coprime,
```

```
- p is an integer factor of a0
      - q is an integer factor of an
//
11
      Note that if a0 = 0, then x = 0 is a root, the
    polynomial can be divided by x and the theorem
      applies once again.
11
// Legendre's formula
//
11
     Considering a prime p, the largest power p^k that
    divides n! is given by:
11
//
     k = floor(n/p) + floor(n/p^2) + floor(n/p^3) + ...
//
11
      Which can be computed in O(log n / log p) time
```