# Team Notebook

# February 11, 2024

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# 1 Strings

#### 1.1 Manacher

```
// odd[i]: length of longest palindrome centered at i
// even[i]: ...longest palindrome centered between i and i+1
string t = "$#";
   for(char c: s) t += c + string("#");
   t += "^":
   int n = t.size();
   vector<int> p(n);
   int 1 = 1, r = 1;
   repx(i, 1, n-1) {
      p[i] = max(0, min(r - i, p[1 + (r - i)]));
      while(t[i - p[i]] == t[i + p[i]]) p[i]++;
      if(i + p[i] > r) l = i - p[i], r = i + p[i];
   repx(i, 2, n-2) {
      if(i%2) even.push_back(p[i]-1);
      else odd.push_back(p[i]-1);
}
```

#### 1.2 aho-corasick

```
struct Vertex {
    int next[26], go[26];
    int p, link = -1, exit = -1, cnt = -1;
    vector<int> leaf:
    char pch:
    Vertex(int p=-1, char ch='$') : p(p), pch(ch) {
       rep(i, 26) next[i] = -1, go[i] = -1:
};
vector<Vertex> t(1):
void add(string &s, int id) {
    int v = 0:
   for (char ch : s) {
       int c = ch - a:
       if (t[v].next[c] == -1) {
           t[v].next[c] = t.size();
           t.emplace_back(v, ch);
       v = t[v].next[c];
    t[v].leaf.push_back(id);
```

```
int go(int v. char ch);
int get_link(int v) {
   if (t[v].link == -1) {
       if (v == 0 || t[v].p == 0) t[v].link = 0;
       else t[v].link = go(get_link(t[v].p), t[v].pch);
   return t[v].link;
int go(int v, char ch) {
   int c = ch - 'a';
   if (t[v].go[c] == -1) {
      if (t[v].next[c] != -1) t[v].go[c] = t[v].next[c];
       else t[v].go[c] = v == 0 ? 0 : go(get_link(v), ch);
   return t[v].go[c];
int next_match(int v){ // Optional
   if(t[v].exit == -1){
      if(t[get link(v)].leaf.size())t[v].exit=get link(v):
       else t[v].exit = v==0 ? 0 : next match(get link(v)):
   }
   return t[v].exit:
int cnt_matches(int v){ // Optional
   if(t[v].cnt == -1)
       t[v].cnt = v == 0 ? 0 : t[v].leaf.size() +
           cnt_matches(get_link(v));
   return t[v].cnt;
```

#### 1.3 hash

```
const int K = 2;
struct Hash{
    const ll MOD[K] = {999727999, 1070777777};
    const ll P = 1777771;
    vector<1l> h[K], p[K];
    Hash(string &s){
        int n = s.size();
        rep(k, K){
            h[k].resize(n+1, 0);
            p[k].resize(n+1, 1);
            repx(i, 1, n+1){
                h[k][i] = (h[k][i-1]*P + s[i-1]) % MOD[k];
            p[k][i] = (p[k][i-1]*P) % MOD[k];
        }
    }
}
vector<1l> get(int i, int j){
```

```
vector<ll> r(K);
rep(k, K){
    r[k] = (h[k][j] - h[k][i]*p[k][j-i]) % MOD[k];
    r[k] = (r[k] + MOD[k]) % MOD[k];
} return r;
}
};
```

#### 1.4 hash2d

```
using Hash = pair<11, int>;
struct Block {
   int x0, y0, x1, y1;
struct Hash2d {
   11 HMOD:
   int W, H;
   vector<int> h;
   vector<int> p;
   Hash2d() {}
   Hash2d(const string &s, int W_, int H_, 11 HMOD_ =
        1000003931)
       : W(W_ + 1), H(H_ + 1), HMOD(HMOD_) {
       static const 11 P =
           chrono::steady_clock::now().time_since_epoch().
                count() % (1 << 29);
       p.resize(W * H);
       p[0] = 1;
       rep(i, W * H - 1) p[i + 1] = p[i] * P % HMOD;
       h.assign(W * H, 0);
       repx(y, 1, H) repx(x, 1, W) {
          ll c = (ll)s[(y - 1) * (W - 1) + x - 1] * p[y * W
                + x] % HMOD;
          h[v * W + x] = (HMOD + h[v * W + x - 1] + h[(v -
               1) * W + x] -
                         h[(y-1)*W+x-1]+c)%
   bool isout(Block s) {
       return s.x0 < 0 \mid | s.x0 >= W \mid | s.x1 < 0 \mid | s.x1 >= W
             | | s.y0 < 0 | |
             s.y0 >= H \mid \mid s.y1 < 0 \mid \mid s.y1 >= H;
   }
```

```
Hash get(Block s) {
       return {(2 * HMOD + h[s.y1 * W + s.x1] - h[s.y1 * W +
             s.x01 -
               h[s.y0 * W + s.x1] + h[s.y0 * W + s.x0]) %
                  HMOD.
              s.v0 * W + s.x0:
    bool cmp(Hash x0, Hash x1) {
       int d = x0.second - x1.second;
       ll &lo = d < 0 ? x0.first : x1.first:
       lo = lo * p[abs(d)] % HMOD:
       return x0.first == x1.first;
};
struct Hash2dM {
    int N:
    vector<Hash2d> sub:
    Hash2dM() {}
    Hash2dM(const string &s, int W, int H, const vector<11> &
        mods)
       : N(mods.size()), sub(N) {
       rep(i, N) sub[i] = Hash2d(s, W, H, mods[i]);
    bool isout(Block s) { return sub[0].isout(s): }
    vector<Hash> get(Block s) {
       vector<Hash> hs(N);
       rep(i, N) hs[i] = sub[i].get(s);
       return hs:
    bool cmp(const vector<Hash> &x0. const vector<Hash> &x1)
       rep(i, N) if (!sub[i].cmp(x0[i], x1[i])) return false | };
       return true:
    bool cmp(Block s0, Block s1) {
       rep(i, N) if (!sub[i].cmp(sub[i].get(s0), sub[i].get(
            s1))) return false;
       return true:
};
```

# 1.5 palindromic-tree

```
struct Node { // (*) = Optional
int len: // length of substring
int to[26]; // insertion edge for all characters a-z
int link; // maximum palindromic suffix
int i:
           // (*) start index of current Node
   int cnt; // (*) # of occurrences of this substring
   Node(int len, int link=0, int i=0, int cnt=1): len(len).
   link(link), i(i), cnt(cnt) {memset(to, 0, sizeof(to));}
struct EerTree { // Palindromic Tree
   vector<Node> t; // tree (max size of tree is n+2)
                // current node
   EerTree(string &s) : last(0) {
      t.emplace_back(-1); t.emplace_back(0); // root 1 & 2
      rep(i, s.size()) add(i, s): // construct tree
      for(int i = t.size()-1: i > 1: i--)
          t[t[i].link].cnt += t[i].cnt;
   void add(int i, string &s){
       int p=last, c=s[i]-'a';
       while(s[i-t[p].len-1] != s[i]) p = t[p].link;
      if(t[p].to[c]){ last = t[p].to[c]; t[last].cnt++; }
      elsef
          int q = t[p].link:
          while(s[i-t[q].len-1] != s[i]) q = t[q].link;
          q = max(1, t[q].to[c]);
          last = t[p].to[c] = t.size();
          t.emplace_back(t[p].len + 2, q, i-t[p].len-1);
   }
string s = "abcbab"; EerTree pt(s); // build EerTree
repx(i, 2, pt.t.size()){// list all distinct palindromes
 repx(j,pt.t[i].i,pt.t[i].i+pt.t[i].len)cout << s[j];
 cout << " " << pt.t[i].cnt << endl;</pre>
```

# 1.6 prefix-function

```
vector<int> prefix_function(string s) {
```

```
int n = s.size():
   vector<int> pi(n):
   repx(i, 1, n) {
       int j = pi[i-1];
       while (j > 0 \&\& s[i] != s[j])
          i = pi[i-1]:
       if (s[i] == s[i])
           j++;
       pi[i] = j;
   return pi:
vector<vector<int>> aut:
void compute automaton(string s) {
   s += '#':
   int n = s.size();
   vector<int> pi = prefix_function(s);
   aut.assign(n, vector<int>(26));
   rep(i, n) {
       rep(c, 26) {
          int j = i;
          while (j > 0 \&\& 'a' + c != s[j])
              j = pi[j-1];
           if ('a' + c == s[i])
              j++;
           aut[i][c] = j;
   }
// k = n - pi[n - 1]; if k divides n, then the string can be
// aprtitioned into blocks of length k otherwise there is no
// effective compression and the answer is n.
```

## 1.7 suffix-array

```
for (int &pi : p) pi = (pi - k + N) % N;
       for (int cl : c) cnt[cl + 1] += 1;
       rep(i, C) cnt[i + 1] += cnt[i];
       rep(i, N) p2[cnt[c[p[i]]]++] = p[i];
       c2[p2[0]] = 0;
       repx(i, 1, N) c2[p2[i]] =
           c2[p2[i-1]] + (c[p2[i]] != c[p2[i-1]] ||
                          c[(p2[i] + k) \% N] != c[(p2[i - 1]
                                + k) % N1):
       swap(c, c2), swap(p, p2);
   p.erase(p.begin()); // optional: erase terminating NUL
   return p;
// build the lcp
// 'lcp[i]' represents the length of the longest common
// prefix between suffix i and suffix i+1 in the suffix
//array 'p'. the last element of 'lcp' is zero by convention
vector<int> makelcp(const string &s. const vector<int> &p) {
   int N = p.size(), k = 0;
   vector<int> r(N), lcp(N);
   rep(i, N) r[p[i]] = i;
   rep(i, N) {
       if (r[i] + 1 >= N) { k = 0; continue; }
       int j = p[r[i] + 1];
       while (i + k < N \&\& j + k < N \&\& s[i + k] == s[j + k]
           1) k += 1:
       lcp[r[i]] = k:
       if (k) k -= 1;
   return lcp;
// lexicographically compare the suffixes starting from 'i'
// and 'j', considering only up to 'K' characters.
// 'r' is the inverse suffix array, mapping suffix offsets
// to indices. requires an LCP sparse table.
int lcp_cmp(vector<int> &r, Sparse<int> &lcp, int i, int j,
    int K) {
   if (i == j) return 0;
   int ii = r[i], jj = r[j];
   int 1 = lcp.query(min(ii, jj), max(ii, jj));
   if (1 >= K) return 0;
   return ii < jj ? -1 : 1;</pre>
```

#### 1.8 suffix-automaton

```
struct SuffixAutomaton {
  vector<map<char,int>> edges;
```

```
vector<int> link, len, cnt, paths, pos:
vector<bool> terminal;
int last; // idx of the eq. class of the whole string
SuffixAutomaton(string s) :last(0) {
   edges.push_back({});
   link.push back(-1):
   len.push_back(0);
   rep(i, s.size()) {
       edges.push_back({});
       len.push_back(i+1);
       link.push_back(0);
       int r = len.size() - 1, p = last:
       while(p >= 0 && !edges[p].count(s[i])) {
          edges[p][s[i]] = r;
          p = link[p];
       if(p != -1) {
          int q = edges[p][s[i]];
          if(len[p] + 1 == len[q]) link[r] = q;
              edges.push_back(edges[q]);
              len.push_back(len[p] + 1);
              link.push_back(link[q]);
              int qq= link[q] = link[r] =len.size()-1;
              while(p >= 0 && edges[p][s[i]] == q){
                  edges[p][s[i]] = qq;
                  p = link[p];
          }
       last = r;
   } /* ----- Optional ----- */
   terminal.assign(len.size(), 0);
   for(int p = last; p > 0; p = link[p]) terminal[p]=1;
   cnt.assign(len.size(), -1); cnt_matches(0);
   //precompute # of paths (substr) starting from state
   paths.assign(len.size(), -1); cnt_paths(0);
   pos.assign(len.size(), -1); get_pos(0);
int cnt matches(int state) {
   if(cnt[state] != -1) return cnt[state]:
   int ans = terminal[state];
   for(auto edge : edges[state])
       ans += cnt_matches(edge.second);
   return cnt[state] = ans;
int cnt_paths(int state) {
   if(paths[state] != -1) return paths[state];
   int ans = state != 0:
                              // without repetitions
// int ans = state == 0 ? 0 : cnt[state]; // with rep.
```

```
for(auto edge : edges[state])
          ans += cnt_paths(edge.second);
      return paths[state] = ans;
   int get_pos(int state) { // gets first pos
       if(pos[state] != -1) return pos[state];
                                  // max->first_pos
       int ans = 0:
// int ans = terminal[state] ? 0 : 1e9;// min->last_pos
      for(auto edge : edges[state])
          ans = max(ans, get_pos(edge.second)+1); //or min
       return pos[state] = ans:
   string get_k_substring(int k) { // 0-indexed
       string ans; int state = 0;
       while(1){
           int curr = state != 0; // without repetition
       // int curr = state == 0 ? 0 : cnt[state]: // with
          if(curr > k) return ans;
          k -= curr:
          for(auto edge : edges[state]) {
              if(paths[edge.second] <= k) {</pre>
                 k -= paths[edge.second];
              } else {
                 ans += edge.first;
                 state = edge.second;
                 break:
          }
   }
};
```

# 1.9 z-function

```
// i-th element is equal to the greatest number of
// characters starting from the position i that coincide
// with the first characters of s
vector<int> z_function(string s) {
   int n = s.size();
   vector<int> z(n);
   int l = 0, r = 0;
   for(int i = 1; i < n; i++) {
      if(i < r) z[i] = min(r - i, z[i - 1]);
      while(i + z[i] < n && s[z[i]] == s[i + z[i]])z[i]++;
      if(i + z[i] > r) {
            l = i;
            r = i + z[i];
      }
}
```

```
return z;
}
```

#### 2 common

#### 2.1 common

### 2.2 debug

```
// Command to check time and memory usage:
/usr/bin/time -v ./tmp
// See "Maximum resident set size" for max memory used
// Commands for interactive checker:
mkfifo fifo
(./solution < fifo) | (./interactor > fifo)
// Does not work on the Windows file system, i.e., /mnt/c/
// The special fifo file must be used, otherwise the
// solution will not wait for input and will read EOF
```

# $3 ext{ dp}$

#### 3.1 convex-hull-trick

```
struct Line {
   mutable ll a, b, c;

bool operator<(Line r) const { return a < r.a; }
   bool operator<(ll x) const { return c < x; }
};

// dynamically insert 'a*x + b' lines and query for maximum
// at any x all operations have complexity O(log N)</pre>
```

```
struct LineContainer : multiset<Line, less<>> {
   11 div(11 a, 11 b) {
      return a / b - ((a ^ b) < 0 && a % b);
   bool isect(iterator x, iterator v) {
       if (y == end()) return x->c = INF, 0;
      if (x->a == y->a) x->c = x->b > y->b ? INF : -INF;
       else x->c = div(y->b - x->b, x->a - y->a);
       return x->c >= y->c;
   }
   void add(ll a, ll b) {
       // a *= -1, b *= -1 // for min
       auto z = insert(\{a, b, 0\}), y = z++, x = y;
       while (isect(y, z)) z = erase(z);
      if (x != begin() && isect(--x, y)) isect(x, y = erase
       while ((v = x) != begin() \&\& (--x)->c >= v->c) isect(
           x. erase(v)):
   11 query(11 x) {
      if (empty()) return -INF; // INF for min
       auto 1 = *lower_bound(x);
      return 1.a * x + 1.b;
       // return -l.a * x - l.b: // for min
   }
};
```

# 3.2 divide-and-conquer

```
// for every index i assign an optimal index j, such that
// cost(i, j) is minimal for every i. the property that if
// i2 >= i1 then j2 >= j1 is exploited (monotonic condition)
// calculate optimal index for all indices in range [1, r)
// knowing that the optimal index for every index in this
// range is within [optl, optr).
// time: O(N log N)
void calc(vector<int> &opt, int 1, int r,int opt1,int optr){
   if (1 == r) return:
   int i = (1 + r) / 2:
   11 optc = INF;
   int optj;
   repx(j, optl, optr) {
      11 c = i + j; // cost(i, j)
       if (c < optc) optc = c, optj = j;</pre>
   }
    opt[i] = optj;
```

```
calc(opt, 1, i, opt1, optj + 1);
calc(opt, i + 1, r, optj, optr);
}
```

# $4 \quad \text{geo2d}$

#### 4.1 circle

```
struct C {
   Po; Tr;
   // circle-line intersection, assuming it exists
   // points are sorted along the direction of the line
   pair<P. P> line inter(L 1) const {
      P c = 1.closest_to(o); T c2 = (c - o).magsq();
      P = 1.d * sqrt(max(r*r - c2, T()) / 1.d.magsq());
      return {c - e, c + e}:
   // check the type of line-circle collision
   // <0: 2 inters, =0: 1 inter, >0: 0 inters
   T line collide(L 1) const {
      T c2 = (1.closest to(o) - o).magsq():
       return c2 - r * r;
   // calculates the two intersections between two circles
   // the circles must intersect in one or two points!
   pair<P, P> inter(C h) const {
      P d = h.o - o:
      T c = (r * r - h.r * h.r) / d.magsq();
      return h.line_inter(\{(1 + c) / 2 * d, d.rot()\});
   // check if the given circles intersect
   bool collide(C h) const {
       return (h.o - o).magsq() \le (h.r + r) * (h.r + r);
   // get one of the two tangents that go through the point
   // the point must not be inside the circle
   // a = -1: cw (relative to the circle) tangent
   // a = 1: ccw (relative to the circle) tangent
   P point tangent(P p. T a) const {
      Tc = r * r / p.magsq();
       return o + c*(p-o) - a*sqrt(c*(1-c))*(p-o).rot();
```

```
// get one of the 4 tangents between the two circles
// a = 1: exterior tangents
// a = -1: interior tangents (requires no area overlap)
// b = 1: ccw tangent
// b = -1: cw tangent
// the line origin is on this circumference, and the
// direction is a unit vector towards the other circle
L tangent(C c, T a, T b) const {
   T dr = a * r - c.r:
   P d = c.o - o;
   P n = (d*dr+b*d.rot()*sqrt(d.magsq()-dr*dr)).unit();
   return {o + n * r. -b * n.rot()}:
// circumcircle of a **non-degenerate** triangle
static C thru_points(P a, P b, P c) {
   b = b - a, c = c - a:
   P p = (b*c.magsq() - c*b.magsq()).rot() / (b%c*2);
   return {a + p, p.mag()};
// find the two circles that go through the given point,
// are tangent to the given line and have radius 'r'
// the point-line distance must be at most 'r'!
// the circles are sorted in the direction of the line
static pair<C, C> thru_point_line_r(P a, L t, T r) {
   P d = t.d.rot().unit():
   if (d * (a - t.o) < 0) d = -d:
   auto p = C(a, r).line_inter(\{t.o + d * r, t.d\});
   return {{p.first, r}, {p.second, r}};
// find the two circles that go through the given points
// and have radius 'r'
// circles sorted by angle from the first point
// the points must be at most at distance 'r'!
static pair<C, C> thru_points_r(P a, P b, T r) {
   auto p = C(a, r).line_inter({(a+b)/2, (b-a).rot()}); }
   return {{p.first, r}, {p.second, r}};
```

#### 4.2 convex-hull

};

```
// ccw order, excludes collinear points by default
vector<P> chull(vector<P> p) {
   if (p.size() < 3) return p;
   vector<P> r; int m, k = 0;
   sort(p.begin(), p.end(), [](P a, P b) {
```

```
return a.x != b.x ? a.x < b.x : a.y < b.y; });
for (P q : p) { // lower hull
    while (k >= 2 && r[k - 1].left(r[k - 2], q) >= 0)
        r.pop_back(), k--; // >= to > to add collinears
    r.push_back(q), k++;
}
if (k == (int)p.size()) return r;
r.pop_back(), k--, m = k;
for (int i = p.size() - 1; i >= 0; --i) { // upper hull
    while (k >= m+2 && r[k-1].left(r[k-2], p[i]) >= 0)
        r.pop_back(), k--; // >= to > to add collinears
    r.push_back(p[i]), k++;
}
r.pop_back(); return r;
```

# 4.3 delaunay

```
typedef __int128_t lll; // if on a 64-bit platform
struct 0 {
   Q *rot, *o; P p = {INF, INF}; bool mark;
   P &F() { return r()->p; }
   0 *&r() { return rot->rot: }
   Q *prev() { return rot->o->rot; }
   Q *next() { return r()->prev(); }
T cross(P a, P b, P c) { return (b - a) % (c - a); }
bool circ(P p, P a, P b, P c) { // is p in the circumcircle?
   111 p2 = p.magsq(), A = a.magsq() - p2.
      B = b.magsq() - p2, C = c.magsq() - p2;
   return cross(p. a. b) * C + cross(p. b. c) * A + cross(p.
         c. a) * B > 0:
Q *makeEdge(Q *&H, P orig, P dest) {
   Q *r = H ? H : new Q{new Q{new Q{new Q{0}}}}:
   H = r -> 0: r -> r() -> r() = r:
   repx(i, 0, 4) r = r->rot, r->p = {INF, INF},
      r->0 = i & 1 ? r : r->r();
   r\rightarrow p = orig; r\rightarrow F() = dest;
   return r;
void splice(Q *a, Q *b) {
   swap(a->o->rot->o, b->o->rot->o); swap(a->o, b->o);
```

```
Q * connect(Q * \& H. Q * a. Q * b) {
   Q *q = makeEdge(H, a->F(), b->p);
   splice(q, a->next()); splice(q->r(), b); return q;
pair<Q *, Q *> rec(Q *&H, const vector<P> &s) {
   if (s.size() <= 3) {</pre>
       Q *a = makeEdge(H, s[0], s[1]), *b = makeEdge(H, s[0])
            [1], s.back());
       if (s.size() == 2) return \{a, a->r()\}: splice(a->r(), a->r()\}:
       auto side = cross(s[0], s[1], s[2]);
       Q *c = side ? connect(H, b, a) : 0;
       return {side < 0 ? c->r() : a, side < 0 ? c : b->r()
            };
#define J(e) e \rightarrow F(), e \rightarrow p
#define valid(e) (cross(e->F(), J(base)) > 0)
   Q *A, *B, *ra, *rb; int half = s.size() / 2;
   tie(ra, A) = rec(H, {s.begin(), s.end() - half});
   tie(B, rb) = rec(H, {s.begin() + s.size() - half, s.end()
   while ((cross(B->p, J(A)) < 0 \&\& (A = A->next())) | |
          (cross(A->p, J(B)) > 0 && (B = B->r()->o)));
   Q *base = connect(H, B->r(), A);
   if (A->p == ra->p) ra = base->r():
   if (B->p == rb->p) rb = base;
#define DEL(e, init, dir) Q *e = init->dir; \
   if (valid(e)) while (circ(e->dir->F(), J(base), e->F()))
        { \
           Q *t = e->dir; splice(e, e->prev()); \
           splice(e->r(), e->r()->prev()); e->o = H; H = e;
                e = t: \
       }
   for (;;) {
       DEL(LC, base->r(), o); DEL(RC, base, prev());
       if (!valid(LC) && !valid(RC)) break;
       if (!valid(LC) || (valid(RC) && circ(J(RC), J(LC))))
            base = connect(H, RC, base->r());
       else base = connect(H, base->r(), LC->r());
   }
   return {ra, rb};
#undef J
#undef valid
#undef DEL
```

```
// there must be no duplicate points
// returns no triangles in the case of all collinear points
// produces counter-clockwise triangles ordered in triples
// maximizes the minimum angle across all triangulations
// the euclidean mst is a subset of these edges
// O(N log N)
vector<P> triangulate(vector<P> pts) {
   sort(pts.begin(), pts.end(), [](P a, P b) {
       return make_pair(a.x, a.y) < make_pair(b.x, b.y);</pre>
   assert(unique(pts.begin(), pts.end()) == pts.end());
   if (pts.size() < 2) return {};</pre>
   Q *H = 0; Q *e = rec(H, pts).first;
   vector < Q *> q = \{e\}; int qi = 0;
   while (cross(e->o->F(), e->F(), e->p) < 0) e = e->o;
#define ADD
   {
       Q *c = e;
       do {
           c->mark = 1; pts.push_back(c->p); \
           q.push_back(c->r()); c = c->next(); \
       } while (c != e):
   ADD:
   pts.clear();
   while (qi < (int)q.size()) if (!(e = q[qi++])->mark) ADD;
   return pts:
#undef ADD
```

# 4.4 halfplane-intersect

```
while (n >= 2 && H[i].side(q[0].intersection(q[1])) >
       q.pop_front(), n--;
   if (n > 0 && H[i].parallel(q[n - 1])) {
       if (H[i].d * q[n - 1].d < 0) return {};</pre>
       if (H[i].side(q[n - 1].o) > 0) q.pop_back(), n--;
       else continue:
   q.push_back(H[i]), n++;
while (n \ge 3 \&\& a[0].side(a[n - 1].intersection(a[n -
    21)) > 0)
   q.pop_back(), n--;
while (n \ge 3 \&\& q[n - 1].side(q[0].intersection(q[1])) >
   q.pop_front(), n--;
if (n < 3) return {};</pre>
vector<P> ps(n):
rep(i, n) ps[i] = q[i].intersection(q[(i + 1) % n]);
return ps:
```

#### 4.5 line

```
// a segment or an infinite line
// does not handle point segments correctly!
struct L {
   P o, d;
   static L from eq(P ab, T c) {
      return L{ab.rot(), ab * -c / ab.magsq()};
   pair<P, T> line_eq() { return {-d.rot(), d.rot() * o}; }
   // on which side of the line is the point
   // negative: left, positive: right
   T side(P r) const { return (r - o) % d: }
   // returns the intersection coefficient
   // in the range [0, d % r.d]
   // if d % r.d is zero, the lines are parallel
   T inter(L r) const { return (r.o - o) % r.d; }
   // get the single intersection point
   // lines must not be parallel
   P intersection(L r) const {return o+d*inter(r)/(d%r.d);}
```

```
// check if lines are parallel
bool parallel(L r) const { return abs(d % r.d) <= EPS; }</pre>
// check if segments intersect
bool seg_collide(L r) const {
   Tz = d \% r.d:
   if (abs(z) <= EPS) {
       if (abs(side(r.o)) > EPS) return false;
       T s = (r.o - o) * d, e = s + r.d * d;
       if (s > e) swap(s, e);
       return s <= d * d + EPS && e >= -EPS:
   T s = inter(r), t = -r.inter(*this);
   if (z < 0) s = -s, t = -t, z = -z;
   return s>=-EPS && s<=z+EPS && t>=-EPS && t<=z+EPS:
// full segment intersection
// makes a point segment if the intersection is a point
// however it does not handle point segments as input!
bool seg_inter(L r, L *out) const {
   Tz = d \% r.d:
   if (abs(z) <= EPS) {
       if (abs(side(r.o)) > EPS) return false;
       if (r.d * d < 0) r = \{r.o + r.d, -r.d\};
       P s = o * d < r.o * d ? r.o : o;
       P = (o+d)*d < (r.o+r.d)*d ? o+d : r.o+r.d:
       if (s * d > e * d) return false:
       return *out = {s, e - s}, true;
   T s = inter(r), t = -r.inter(*this);
   if (z < 0) s = -s, t = -t, z = -z;
   if (s>=-EPS && s<=z+EPS && t>=-EPS && t<=z+EPS)
       return *out = \{0 + d * s / z, \{0, 0\}\}, true;
   return false:
// check if the given point is on the segment
bool point_on_seg(P r) const {
   if (abs(side(r)) > EPS) return false;
   if ((r - o) * d < -EPS) return false:
   if ((r - o - d) * d > EPS) return false;
   return true:
// point in this line that is closest to a given point
P closest_to(P r) const {
   P dr = d.rot(); return r + dr*((o-r)*dr)/d.magsq();
```

#### 4.6 minkowski

```
void reorder_polygon(vector<P> &ps) {
    int pos = 0;
    repx(i, 1, (int)ps.size()) {
       if (ps[i].y < ps[pos].y || (ps[i].y == ps[pos].v &&</pre>
            ps[i].x < ps[pos].x))
           pos = i:
    rotate(ps.begin(), ps.begin() + pos, ps.end());
}
vector<P> minkowski(vector<P> ps, vector<P> qs) {
    // the first vertex must be the lowest
    reorder_polygon(ps); reorder_polygon(qs);
    ps.push_back(ps[0]); ps.push_back(ps[1]);
    qs.push_back(qs[0]); qs.push_back(qs[1]);
    vector<P> result; int i = 0, j = 0;
    while (i < ps.size() - 2 || j < qs.size() - 2) {
       result.push back(ps[i] + qs[i]):
       auto z = (ps[i + 1] - ps[i]) \% (qs[j + 1] - qs[j]);
       if (z >= 0 && i < ps.size() - 2) ++i;</pre>
       if (z <= 0 && j < qs.size() - 2) ++j;</pre>
    return result:
```

#### 4.7 point

```
struct P {
   T x, v;
   P(T x, T y) : x(x), y(y) {}
   P() : P(0, 0) \{ \}
   friend ostream &operator<<(ostream &s. const P &r) {</pre>
       return s << r.x << " " << r.v:
   friend istream &operator>>(istream &s, P &r) { return s
        >> r.x >> r.y; }
   P operator+(P r) const { return \{x + r.x, y + r.y\}; \}
   P operator-(P r) const { return {x - r.x, y - r.y}; }
   P operator*(T r) const { return {x * r, y * r}; }
   P operator/(T r) const { return {x / r, y / r}; }
   P operator-() const { return {-x, -y}; }
   friend P operator*(T 1, P r) { return {1 * r.x. 1 * r.v};
         }
   P rot() const { return {-y, x}; }
   T operator*(P r) const { return x * r.x + y * r.y; }
```

```
T operator%(P r) const { return rot() * r; }
   T left(P a. P b) { return (b - a) % (*this - a): }
   T magsq() const { return x * x + v * v: }
   T mag() const { return sqrt(magsq()); }
   P unit() const { return *this / mag(): }
   bool half() const { return abs(y) <= EPS && x < -EPS || y
   T angcmp(P r) const { // like strcmp(this, r)
       int h = (int)half() - r.half();
       return h ? h : r % *this:
   T angcmp_rel(P a, P b) { // like strcmp(a, b)
       Pz = *this:
       int h = z % a <= 0 && z * a < 0 || z % a < 0;
       h = z \% b \le 0 \&\& z * b \le 0 || z \% b \le 0:
       return h? h: b % a;
   bool operator==(P r) const { return abs(x - r.x) <= EPS</pre>
        && abs(v - r.v) <= EPS: }
   double angle() const { return atan2(v, x); }
   static P from_angle(double a) { return {cos(a), sin(a)};
        }
};
```

# 4.8 polygon

```
// get TWICE the area of a simple polygon in ccw order
T area2(const vector<P> &p) {
   int n = p.size(); T a = 0;
   rep(i, n) = + = (p[i] - p[0]) \% (p[(i + 1) \% n] - p[i]):
   return a:
// checks whether a point is inside a ccw simple polygon
// returns 1 if inside, 0 if on border, -1 if outside
int in_poly(const vector<P> &p, P q) {
   int w = 0:
   rep(i, p.size()) {
       P = p[i], b = p[(i + 1) \% p.size()];
       T k = (b - a) \% (q - a);
       T u = a.y - q.y, v = b.y - q.y;
       if (k > 0 \&\& u < 0 \&\& v >= 0) w++;
       if (k < 0 && v < 0 && u >= 0) w--:
       if (k == 0 \&\& (q - a) * (q - b) <= 0) return 0;
```

```
return w ? 1 : -1:
// check if point in ccw convex polygon, O(log n)
// + if inside, 0 if on border, - if outside
T in convex(const vector<P> &p. P q) {
   int 1 = 1, h = p.size() - 2; assert(p.size() >= 3);
   while (1 != h) { // collinear points are unsupported!
       int m = (1 + h + 1) / 2:
       if (q.left(p[0], p[m]) >= 0) 1 = m;
       else h = m - 1:
   T in = min(q.left(p[0], p[1]), q.left(p.back(), p[0]));
   return min(in, q.left(p[1], p[1 + 1]));
int extremal(const vector<P> &p. P d) {
   int n = p.size(), 1 = 0, r = n - 1; assert(n);
   P = 0 = (p[n - 1] - p[0]).rot();
   while (1 < r) { // polygon must be convex
       int m = (1 + r + 1) / 2;
       P = (p[(m + n - 1) \% n] - p[m]).rot();
       if (e0.angcmp_rel(d, e) < 0) r = m - 1;
       else 1 = m:
   }
   return 1;
// square dist of most distant points of a ccw convex
// polygon with NO COLLINEAR POINTS
T callipers(const vector<P> &p) {
   int n = p.size():
   T r = 0:
   for (int i = 0, j = n < 2 ? 0 : 1; <math>i < j; i++) {
       for (;; j = (j + 1) % n) {
          r = max(r, (p[i] - p[j]).magsq());
          if ((p[(i + 1) % n] - p[i]) % (p[(j + 1) % n] - p
                [i]) <= EPS) break;</pre>
       }
   return r:
P centroid(const vector<P> &p) { // (barycenter)
   P r(0, 0); T t = 0; int n = p.size();
   rep(i, n) {
       r += (p[i] + p[(i+1)\%n]) * (p[i] \% p[(i+1)\%n]);
       t += p[i] \% p[(i+1)\%n];
   return r / t / 3;
```

# 4.9 sweep

```
#include "point.cpp"
// iterate over all pairs of points
// 'op' is called with all ordered pairs of different
    indices '(i, i)'
// additionally, the 'ps' vector is kept sorted by signed
    distance
// to the line formed by 'i' and 'j'
// for example, if the vector from 'i' to 'j' is pointing
// the 'ps' vector is sorted from smallest 'y' to largest 'y
// note that, because the 'ps' vector is sorted by signed
// 'j' is always equal to 'i + 1'
// this means that the amount of points to the left of the
    line is always 'N - i'
template <class OP>
void all pair points(vector<P> &ps. OP op) {
   int N = ps.size();
   sort(ps.begin(), ps.end(), [](P a, P b) {
       return make_pair(a.y, a.x) < make_pair(b.y, b.x);</pre>
   }):
   vector<pair<int, int>> ss;
   rep(i, N) rep(j, N) if (i != j) ss.push_back({i, j});
   stable_sort(ss.begin(), ss.end(), [&](auto a, auto b) {
       return (ps[a.second] - ps[a.first]).angle lt(ps[b.
            second] - ps[b.first]);
   vector<int> p(N); rep(i, N) p[i] = i;
   for (auto [i, j] : ss)
```

```
{ op(p[i], p[j]); swap(ps[p[i]], ps[p[j]]); swap(p[i ], p[j]); }
```

#### 4.10 theorems

```
// Pick's theorem
// Simple polygon with integer vertices:
// A = I + B / 2 - 1
// A: Area of the polygon
// I: Integer points strictly inside the polygon
// B: Integer points on the boundary of the polygon
```

# 5 graph

# 5.1 bellman-ford

```
struct Edge { int u, v; ll w; };
// find distance from source node to all nodes.
// supports negative edge weights.
// returns true if a negative cycle is detected.
11
// time: O(V E)
bool bellman_ford(int N, int s, vector<Edge> &E, vector<11>
    &D. vector<int> &P) {
   P.assign(N, -1), D.assign(N, INF), D[s] = 0;
   rep(i. N - 1) {
       bool f = true:
       rep(ei, E.size()) {
           auto &e = E[ei]:
           ll n = D[e.u] + e.w;
           if (D[e.u] < INF && n < D[e.v])</pre>
              D[e.v] = n, P[e.v] = ei, f = false;
       if (f) return false:
   }
   return true;
```

### 5.2 blossom

```
vector<int> g[MAXN];int n,m,mt[MAXN],qh,qt,q[MAXN],ft[MAXN],
bs[MAXN];bool inq[MAXN],inb[MAXN],inp[MAXN];int lca(int root
```

```
,int x,int y){memset(inp,0,sizeof(inp));while(1){inp[x=bs[x]
]=true;if(x==root)break;x=ft[mt[x]];}while(1){if(inp[y=bs[y])}
])return y;else y=ft[mt[y]];}}void mark(int z,int x){while(
bs[x]!=z){int y=mt[x];inb[bs[x]]=inb[bs[y]]=true;x=ft[y];if(
bs[x]!=z)ft[x]=y;}}void contr(int s,int x,int y){int z=lca(s
.x.v):memset(inb.0.sizeof(inb)):mark(z.x):mark(z.v):if(bs[x]
!=z)ft[x]=y;if(bs[y]!=z)ft[y]=x;rep(x,n)if(inb[bs[x]]){bs[x]}
=z;if(!ing[x])ing[g[++gt]=x]=true;}}int findp(int s){memset(
inq,0,sizeof(inq));memset(ft,-1,sizeof(ft));rep(i,n)bs[i]=i;
inq[q[qh=qt=0]=s]=true; while(qh<=qt){int x=q[qh++]; for(int y</pre>
g[x] = g[x]  if (bs[x]!=bs[y]&&mt[x]!=y){if(y==s||mt[y]>=0&&ft[mt[y])}
>=0)contr(s.x.v):else if(ft[v]<0){ft[v]=x:if(mt[v]<0)return</pre>
v;else if(!ing[mt[v]])ing[g[++qt]=mt[v]]=true;}}}return -1;}
int aug(int s,int t){int x=t,y,z;while(x>=0){y=ft[x];z=mt[y]}
;mt[y]=x;mt[x]=y;x=z;}return t>=0;}int edmonds(){int r=0;
memset(mt,-1,sizeof(mt));rep(x,n)if(mt[x]<0)r+=aug(x,findp(x</pre>
));return r;}
```

#### 5.3 dinic

```
// time: 0(E V^2)
       O(E V^(2/3)) / O(E sqrt(E)) unit capacities
11
      O(E sqrt(V)) (hopcroft-karp) unit networks
//unit network: c in {0,1} & forall v, indeg<=1 or outdeg<=1
//min-cut: nodes reachable from s in final residual graph
struct Dinic {
   struct Edge { int u, v; ll c, f = 0; };
   int N. s. t: vector<vector<int>> G:
   vector<Edge> E; vector<int> lvl, ptr;
   Dinic() {}
   Dinic(int N, int s, int t): N(N), s(s), t(t), G(N) {}
   void add_edge(int u, int v, 11 c) {
       G[u].push back(E.size()): E.push back({u, v, c}):
       G[v].push_back(E.size()); E.push_back({v, u, 0});
   11 push(int u, 11 p) {
       if (u == t || p <= 0) return p;</pre>
       while (ptr[u] < G[u].size()) {</pre>
           int ei = G[u][ptr[u]++];
           Edge &e = E[ei];
           if (lvl[e.v] != lvl[u] + 1) continue;
           ll a = push(e.v, min(e.c - e.f, p));
           if (a <= 0) continue:</pre>
           e.f += a, E[ei ^ 1].f -= a; return a;
       }
       return 0:
```

```
11 maxflow() {
       11 f = 0;
        while (true) {
           lvl.assign(N, -1); queue<int> q;
           lvl[s] = 0; q.push(s);
            while (!q.empty()) {
               int u = q.front(); q.pop();
               for (int ei : G[u]) {
                   Edge &e = E[ei];
                   if (e.c-e.f<=0||lvl[e.v]!=-1) continue;</pre>
                   lvl[e.v] = lvl[u] + 1; q.push(e.v);
               }
            if (lvl[t] == -1) break;
            ptr.assign(N,0); while(ll ff=push(s,INF))f += ff;
        return f;
};
/* Flujo con demandas (no necesariamente el maximo)
Agregar s' v t' nuevos source and sink
c'(s', v) = sum(d(u, v) \text{ for } u \text{ in } V) \setminus forall \text{ arista } (s', v)
c'(v, t') = sum(d(v, w) \text{ for w in V}) \setminus forall arista (v, t')
c'(u, v) = c(u, v) - d(u, v) \forall aristas antiguas
c'(t, s) = INF (el flujo por esta arista es el flujo real)*/
```

# 5.4 floyd-warshall

# 5.5 heavy-light

```
struct Hld {
  vector<int> P, H, D, pos, top;
```

```
Hld() {}
void init(vector<vector<int>> &G) {
   int N = G.size():
   P.resize(N), H.resize(N), D.resize(N), pos.resize(N),
       top.resize(N):
   D[0] = -1, dfs(G, 0); int t = 0;
   rep(i, N) if (H[P[i]] != i) {
       int j = i;
       while (j != -1)
          \{ top[j] = i, pos[j] = t++; j = H[j]; \}
   }
}
int dfs(vector<vector<int>> &G, int i) {
   int w = 1, mw = 0;
   D[i] = D[P[i]] + 1, H[i] = -1;
   for (int c : G[i]) {
       if (c == P[i]) continue:
       P[c] = i: int sw = dfs(G, c): w += sw:
       if (sw > mw) H[i] = c, mw = sw;
   return w;
// visit the log N segments in the path from \boldsymbol{u} to \boldsymbol{v}
template <class OP>
void path(int u, int v, OP op) {
   while (top[u] != top[v]) {
       if (D[top[u]] > D[top[v]]) swap(u, v);
       op(pos[top[v]], pos[v] + 1); v = P[top[v]];
   if (D[u] > D[v]) swap(u, v);
   op(pos[u], pos[v] + 1); // value on node
   // op(pos[u]+1, pos[v] + 1); // value on edge
// an alternative to 'path' that considers order.
// calls 'op' with an 'l <= r' inclusive-exclusive range,
     and a
// boolean indicating if the query is forwards or
    backwards.
template <class OP>
void path(int u, int v, OP op) {
   int lu = u, lv = v;
   while (top[lu] != top[lv])
       if (D[top[lu]] > D[top[lv]]) lu = P[top[lu]];
       else lv = P[top[lv]]:
   int lca = D[lu] > D[lv] ? lv : lu:
```

```
while (top[u] != top[lca])
          op(pos[top[u]], pos[u] + 1, false), u = P[top[u
       if (u != lca) op(pos[lca] + 1, pos[u] + 1, false);
       vector<int> stk:
       while (top[v] != top[lca])
          stk.push_back(v), v = P[top[v]];
       // op(pos[lca], pos[v] + 1, true); // value on node
       op(pos[lca] + 1, pos[v] + 1, true); // value on edge
       reverse(stk.begin(), stk.end()):
       for (int w : stk) op(pos[top[w]], pos[w] + 1, true);
   // commutative segment tree
   template <class T, class S>
   void update(S &seg, int i, T val) { seg.update(pos[i],
        val): }
   // commutative segment tree lazy
   template <class T, class S>
   void update(S &seg, int u, int v, T val) {
       path(u, v, [&](int 1, int r) { seg.update(1, r, val);
   // commutative (lazv) segment tree
   template <class T, class S>
   T query(S &seg, int u, int v) {
      T ans = 0:
           // neutral element
       path(u, v, [\&](int 1, int r) \{ ans += seg.query(1, r) \}
           ; }); // query op
       return ans:
   }
};
```

# 5.6 hungarian

```
// find a maximum gain perfect matching in the given
    bipartite complete graph.
// input: gain matrix (G_{xy} = benefit of joining vertex x
    in set X with vertex
// y in set Y).
// output: maximum gain matching in members 'xy[x]' and 'yx[
    y]'.
// runtime: O(N^3)
struct Hungarian {
```

```
int N. gi. root:
vector<vector<ll>>> gain;
vector<int> xy, yx, p, q, slackx;
vector<ll> lx, ly, slack;
vector<bool> S, T;
void add(int x, int px) {
   S[x] = true, p[x] = px;
   rep(y, N) if (lx[x] + ly[y] - gain[x][y] < slack[v])
       slack[y] = lx[x] + ly[y] - gain[x][y], slackx[y]
void augment(int x, int y) {
   while (x != -2) {
       yx[y] = x; swap(xy[x], y); x = p[x];
}
void improve() {
   S.assign(N, false), T.assign(N, false), p.assign(N,
        -1):
   qi = 0, q.clear();
   rep(x, N) if (xy[x] == -1) {
       q.push_back(root = x), p[x] = -2, S[x] = true;
       break:
   rep(y, N) slack[y] = lx[root] + ly[y] - gain[root][y
        ], slackx[y] = root;
   while (true) {
       while (qi < q.size()) {</pre>
          int x = q[qi++];
          rep(y, N) if (lx[x] + ly[y] == gain[x][y] &&!
              if (yx[y] == -1) return augment(x, y);
              T[y] = true, q.push_back(yx[y]), add(yx[y])
          }
       }
       11 d = INF:
       rep(y, N) if (!T[y]) d = min(d, slack[y]);
       rep(x, N) if (S[x]) lx[x] -= d;
       rep(v, N) if (T[v]) lv[v] += d;
       rep(y, N) if (!T[y]) slack[y] -= d;
       rep(y, N) if (!T[y] && slack[y] == 0) {
```

```
if (yx[y] == -1) return augment(slackx[y], y); | 5.8 | lca
              if (!S[yx[y]]) q.push_back(yx[y]), add(yx[y],
                   slackx[v]):
      }
   }
   Hungarian(vector<vector<11>> g)
       : N(g.size()), gain(g), xy(N, -1), yx(N, -1), lx(N, -
      lv(N). slack(N). slackx(N) {
       rep(x, N) rep(y, N) lx[x] = max(lx[x], ly[y]);
       rep(i, N) improve();
};
```

#### 5.7 kuhn

```
// get a maximum cardinality matching in a bipartite graph.
// input: adjacency lists.
// output: matching (in 'mt' member).
// runtime: O(V E)
struct Kuhn {
   int N. size:
   vector<vector<int>> G;
   vector<bool> seen:
   vector<int> mt:
   bool visit(int i) {
       if (seen[i]) return false;
       seen[i] = true:
       for (int to : G[i])
           if (mt[to] == -1 || visit(mt[to])) {
              mt[to] = i:
              return true:
       return false:
   Kuhn(vector<vector<int>> adj) : G(adj), N(G.size()), mt(N // time: O(F V E)
        , -1) {
       rep(i, N) {
          seen.assign(N, false);
           size += visit(i):
   }
};
```

```
// calculates the lowest common ancestor for any two nodes
    in O(log N) time,
// with O(N log N) preprocessing
struct Lca {
   int N, K, t = 0;
   vector<vector<int>> U:
   vector<int> L, R;
   Lca() {}
   Lca(vector<vector<int>> &G) : N(G.size()), L(N), R(N) {
       K = N \le 1 ? 0 : 32 - \_builtin_clz(N - 1);
       U.resize(K + 1, vector<int>(N));
       visit(G, 0, 0);
       rep(k, K) rep(u, N) U[k + 1][u] = U[k][U[k][u]];
   void visit(vector<vector<int>> &G. int u. int p) {
      L[u] = t++, U[0][u] = p;
       for (int v : G[u]) if (v != p) visit(G, v, u);
       R[u] = t++:
   }
   bool is_anc(int up, int dn) {
       return L[up] <= L[dn] && R[dn] <= R[up];</pre>
   int find(int u, int v) {
       if (is anc(u, v)) return u:
       if (is_anc(v, u)) return v;
       for (int k = K; k \ge 0;)
          if (is_anc(U[k][u], v)) k--;
          else u = U[k][u];
       return U[0][u]:
};
```

#### 5.9 maxflow-mincost

```
F is the maximum flow
       O(V E + F E log V) if bellman-ford is replaced by
    johnson
struct Flow {
   struct Edge {
      int u, v;
      11 c, w, f = 0;
   }:
```

```
int N. s. t:
vector<vector<int>> G:
vector<Edge> E;
vector<ll> d. b:
vector<int> p;
Flow() {}
Flow(int N, int s, int t) : N(N), s(s), t(t), G(N) {}
void add_edge(int u, int v, ll c, ll w) {
   G[u].push_back(E.size());
   E.push back({u, v, c, w}):
   G[v].push_back(E.size());
   E.push_back({v, u, 0, -w});
// naive distances with bellman-ford: O(V E)
void calcdists() {
   p.assign(N, -1), d.assign(N, INF), d[s] = 0:
   rep(i, N - 1) rep(ei, E.size()) {
       Edge &e = E[ei];
       ll n = d[e.u] + e.w:
       if (d[e.u] < INF && e.c - e.f > 0 && n < d[e.v])
           d[e.v] = n, p[e.v] = ei;
   }
// johnsons potentials: O(E log V)
void calcdists() {
   if (b.emptv()) {
      b.assign(N, 0);
       // code below only necessary if there are
           negative costs
       rep(i, N - 1) rep(ei, E.size()) {
          Edge &e = E[ei];
          if (e.f < e.c) b[e.v] = min(b[e.v], b[e.u] + e
       }
   p.assign(N, -1), d.assign(N, INF), d[s] = 0;
   priority queue<pair<11, int>> q:
   q.push({0, s});
   while (!a.emptv()) {
       auto [w, u] = q.top();
       q.pop();
       if (d[u] < -w + b[u]) continue:
       for (int ei : G[u]) {
          auto e = E[ei]:
          ll n = d[u] + e.w;
          if (e.f < e.c \&\& n < d[e.v]) {
```

```
d[e.v] = n, p[e.v] = ei;
                  q.push({b[e.v] - n, e.v});
          }
      }
       b = d:
   }
   11 solve() {
       b.clear():
      11 ff = 0:
       while (true) {
          calcdists();
          if (p[t] == -1) break;
          for (int cur = t: p[cur] != -1: cur = E[p[cur]].u
              f = min(f, E[p[cur]].c - E[p[cur]].f);
          for (int cur = t: p[cur] != -1: cur = E[p[cur]].u
              E[p[cur]].f += f, E[p[cur] ^ 1].f -= f;
          ff += f;
      return ff;
   }
};
```

# 5.10 push-relabel

```
#include "../common.h"

const ll INF = 1e18;

// maximum flow algorithm.
// to run, use 'maxflow()'.

//
// time: O(V^2 sqrt(E)) <= O(V^3)
// memory: O(V^2)
struct PushRelabel {
    vector<vector<ll>> cap, flow;
    vector<ll> excess;
    vector<int> height;

PushRelabel() {}
    void resize(int N) { cap.assign(N, vector<ll>(N)); }

// push as much excess flow as possible from u to v.
    void push(int u, int v) {
```

```
11 f = min(excess[u], cap[u][v] - flow[u][v]);
   flow[u][v] += f:
   flow[v][u] -= f;
   excess[v] += f:
   excess[u] -= f;
// relabel the height of a vertex so that excess flow may
     be pushed.
void relabel(int u) {
   int d = INT32 MAX:
   rep(v, cap.size()) if (cap[u][v] - flow[u][v] > 0) d
       min(d, height[v]);
   if (d < INF) height[u] = d + 1;</pre>
// get the maximum flow on the network specified by 'cap'
     with source 's'
// and sink 't'.
// node-to-node flows are output to the 'flow' member.
11 maxflow(int s. int t) {
   int N = cap.size(), M;
   flow.assign(N, vector<11>(N));
   height.assign(N, 0), height[s] = N;
   excess.assign(N, 0), excess[s] = INF;
   rep(i, N) if (i != s) push(s, i);
   vector<int> q;
   while (true) {
       // find the highest vertices with excess
       q.clear(), M = 0;
       rep(i, N) {
          if (excess[i] <= 0 || i == s || i == t)</pre>
                continue:
           if (height[i] > M) q.clear(), M = height[i];
           if (height[i] >= M) q.push_back(i);
       if (q.empty()) break;
       // process vertices
       for (int u : a) {
           bool relab = true;
           rep(v. N) {
              if (excess[u] <= 0) break;</pre>
              if (cap[u][v] - flow[u][v] > 0 && height[u]
                   1 > height[v])
                  push(u, v), relab = false;
           if (relab) {
              relabel(u):
```

```
break;
}
}

ll f = 0; rep(i, N) f += flow[i][t]; return f;
}
};
```

#### 5.11 strongly-connected-components

```
// compute strongly connected components.
// time: O(V + E), memory: O(V)
// after building:
// comp = map from vertex to component (components are
    toposorted, root first, leaf last)
// N = number of components
// G = condensation graph (component DAG)
// byproducts:
// vgi = transposed graph
// order = reverse topological sort (leaf first, root last)
// others:
// vn = number of vertices
// vg = original vertex graph
struct Scc {
   int vn. N:
   vector<int> order, comp;
   vector<vector<int>> vg, vgi, G;
   void toposort(int u) {
       if (comp[u]) return;
       comp[u] = -1;
       for (int v : vg[u]) toposort(v);
       order.push_back(u);
   bool carve(int u) {
       if (comp[u] != -1) return false;
       comp[u] = N;
       for (int v : vgi[u]) {
          if (comp[v] != N) G[comp[v]].push back(N):
      }
       return true;
```

```
Scc() {}
Scc(vector<vector<int>> &g) : vn(g.size()), vg(g), comp(
     vn), vgi(vn), G(vn), N(0) {
    rep(u, vn) toposort(u);
    rep(u, vn) for (int v : vg[u]) vgi[v].push_back(u);
    invrep(i, vn) N += carve(order[i]);
}
};
```

#### 5.12 two-sat

```
// calculate the solvability of a system of logical
    equations, where every equation is of the form 'a or b
// 'neg': get negation of 'u'
// 'then': 'u' implies 'v'
// 'anv': 'u' or 'v'
// 'set': 'u' is true
// after 'solve' (O(V+E)) returns true. 'sol' contains one
    possible solution.
// determining all solutions is O(V*E) hard (requires
     computing reachability in a DAG).
struct TwoSat {
   int N: vector<vector<int>> G:
   Scc scc; vector<bool> sol;
   TwoSat(int n) : N(n), G(2 * n), sol(n) {}
   TwoSat() {}
   int neg(int u) { return (u + N) % (2 * N); }
   void then(int u, int v) { G[u].push_back(v), G[neg(v)].
        push_back(neg(u)); }
   void anv(int u, int v) { then(neg(u), v); }
   void set(int u) { G[neg(u)].push_back(u); }
   bool solve() {
       scc = Scc(G);
       rep(u, N) if (scc.comp[u] == scc.comp[neg(u)]) return
       rep(u, N) sol[u] = (scc.comp[u] > scc.comp[neg(u)]);
       return true:
   }
};
```

# 6 implementation

# 6.1 SegmentTreeBeats

```
struct Node {
   11 \text{ s, mx1, mx2, mxc, mn1, mn2, mnc, } 1z = 0;
   Node(): s(0), mx1(LLONG_MIN), mx2(LLONG_MIN), mxc(0),
        mn1(LLONG_MAX), mn2(LLONG_MAX), mnc(0) {}
   Node(ll x): s(x), mx1(x), mx2(LLONG MIN), mxc(1), mn1(x)
        , mn2(LLONG_MAX), mnc(1) {}
   Node(const Node &a. const Node &b) {
       // add
       s = a.s + b.s;
       // min
       if (a.mx1 > b.mx1) mx1 = a.mx1, mxc = a.mxc, mx2 =
           max(b.mx1, a.mx2):
       if (a.mx1 < b.mx1) mx1 = b.mx1, mxc = b.mxc, mx2 =
            max(a.mx1, b.mx2);
       if (a.mx1 == b.mx1) mx1 = a.mx1, mxc = a.mxc + b.mxc,
             mx2 = max(a.mx2, b.mx2);
       // max
       if (a.mn1 < b.mn1) mn1 = a.mn1, mnc = a.mnc, mn2 =
            min(b.mn1, a.mn2);
       if (a.mn1 > b.mn1) mn1 = b.mn1, mnc = b.mnc, mn2 =
            min(a.mn1, b.mn2):
       if (a.mn1 == b.mn1) mn1 = a.mn1, mnc = a.mnc + b.mnc,
            mn2 = min(a.mn2, b.mn2):
};
// 0 - indexed / inclusive - inclusive
template <class node>
struct STB {
   vector<node> st; int n;
   void build(int u, int i, int j, vector<node> &arr) {
       if (i == j) {
          st[u] = arr[i]:
          return;
       int m = (i + j) / 2, l = u * 2 + 1, r = u * 2 + 2;
       build(1, i, m, arr), build(r, m + 1, j, arr);
       st[u] = node(st[1], st[r]);
   void push_add(int u, int i, int j, ll v) {
       st[u].s += (i - i + 1) * v:
       st[u].mx1 += v, st[u].mn1 += v, st[u].lz += v;
       if (st[u].mx2 != LLONG_MIN) st[u].mx2 += v;
       if (st[u].mn2 != LLONG MAX) st[u].mn2 += v:
```

```
void push max(int u, 11 v, bool 1) { // for min op
   if (v \ge st[u].mx1) return:
   st[u].s -= st[u].mx1 * st[u].mxc;
   st[u].mx1 = v:
   st[u].s += st[u].mx1 * st[u].mxc;
   if (1) st[u].mn1 = st[u].mx1:
   else if (v <= st[u].mn1) st[u].mn1 = v;</pre>
   else if (v < st[u].mn2) st[u].mn2 = v;
void push_min(int u, 11 v, bool 1) { // for max op
   if (v <= st[u].mn1) return:</pre>
   st[u].s -= st[u].mn1 * st[u].mnc:
   st[u].mn1 = v;
   st[u].s += st[u].mn1 * st[u].mnc:
   if (1) st[u].mx1 = st[u].mn1:
   else if (v \ge st[u].mx1) st[u].mx1 = v;
   else if (v > st[u].mx2) st[u].mx2 = v:
void push(int u, int i, int i) {
   if (i == i) return:
   // add
   int m = (i + j) / 2, l = u * 2 + 1, r = u * 2 + 2;
   push_add(1, i, m, st[u].lz);
   push_add(r, m + 1, j, st[u].lz);
   st[u].lz = 0;
   // min
   push_max(1, st[u].mx1, i == m);
   push max(r, st[u].mx1, m + 1 == i):
   // max
   push min(l, st[u].mn1, i == m):
   push_min(r, st[u].mn1, m + 1 == r);
node query(int a, int b, int u, int i, int j) {
   if (b < i || j < a) return node();</pre>
   if (a <= i && j <= b) return st[u];</pre>
   push(u, i, j);
   int m = (i + j) / 2, l = u * 2 + 1, r = u * 2 + 2;
   return node(query(a, b, 1, i, m), query(a, b, r, m +
        1, i));
void update add(int a, int b, ll v, int u, int i, int i) }:
   if (b < i || j < a) return;</pre>
   if (a <= i && j <= b) {</pre>
       push_add(u, i, j, v);
       return:
   push(u, i, j);
   int m = (i + j) / 2, l = u * 2 + 1, r = u * 2 + 2;
   update_add(a, b, v, l, i, m);
```

```
update add(a, b, v, r, m + 1, i):
   st[u] = node(st[1], st[r]):
void update min(int a, int b, ll v, int u, int i, int i)
   if (b < i \mid | i < a \mid | v >= st[u].mx1) return:
   if (a <= i && j <= b && v > st[u].mx2) {
       push_max(u, v, i == j);
       return:
   push(u, i, i):
   int m = (i + i) / 2, l = u * 2 + 1, r = u * 2 + 2:
   update_min(a, b, v, l, i, m);
   update_min(a, b, v, r, m + 1, j);
   st[u] = node(st[l], st[r]):
void update max(int a, int b, ll v, int u, int i, int i)
   if (b < i || i < a || v <= st[u].mn1) return;</pre>
   if (a <= i && i <= b && v < st[u].mn2) {
       push_min(u, v, i == j);
       return:
   push(u, i, j);
   int m = (i + j) / 2, l = u * 2 + 1, r = u * 2 + 2;
   update_max(a, b, v, l, i, m);
   update_max(a, b, v, r, m + 1, j);
   st[u] = node(st[1], st[r]):
}
STB(vector<node> &v, int N) : n(N), st(N * 4 + 5) { build
     (0, 0, n - 1, v); 
node query(int a, int b) { return query(a, b, 0, 0, n -
    1); }
void update add(int a, int b, ll v) { update add(a, b, v,
     0.0.n - 1): 
void update_min(int a, int b, ll v) { update_min(a, b, v,
     0.0.n - 1): 
void update_max(int a, int b, ll v) { update_max(a, b, v,
     0.0.n - 1):
```

# 6.2 Treap

```
struct Node {
   int p, sz = 0, v, acc, l = -1, r = -1;
   Node(): v(0), acc(0) {}
   Node(int x) : p(gen()), sz(1), v(x), acc(x) {}
   void recalc(const Node &a, const Node &b) {
      sz = a.sz + b.sz + 1:
       acc = v + a.acc + b.acc;
   }
};
template <class node>
struct Treap {
   vector<node> t;
   int n. r = -1:
   node get(int u) { return u != -1 ? t[u] : node(); }
   void recalc(int u) { t[u].recalc(get(t[u].l), get(t[u].r)
        ); }
   int merge(int 1, int r) {
       if (min(1, r) == -1) return 1 != -1 ? 1 : r:
       int ans = (t[1].p < t[r].p) ? 1 : r;
       if (ans == 1) t[1].r = merge(t[1].r, r), recalc(1);
       if (ans == r) t[r].l = merge(l, t[r].l), recalc(r);
       return ans:
   pii split(int u, int id) {
       if (u == -1) return {-1, -1};
       int szl = get(t[u].1).sz:
      if (szl >= id) {
          pii ans = split(t[u].1, id):
          t[u].1 = ans.ss:
          recalc(u);
          return {ans.ff. u}:
       pii ans = split(t[u].r. id - szl - 1):
      t[u].r = ans.ff:
       recalc(u):
       return {u, ans.ss};
   Treap(vi &v) : n(sz(v)) {
      for (int i = 0; i < n; i++) t.eb(v[i]), r = merge(r, v[i])
           i):
};
// Complete Implicit Treap with Lazy propagation //
   int p, sz = 0, v, acc, l = -1, r = -1, par = -1, lzv = 0;
```

```
bool lz = false, f = false:
   Node(): v(0), acc(0) {}
   Node(int x) : p(gen()), sz(1), v(x), acc(x) {}
   void recalc(const Node &a. const Node &b) {
       sz = a.sz + b.sz + 1:
       acc = v + a.acc + b.acc:
   void upd_lazv(int x) { lz = 1, lzv += x; }
   void lazy() { v += lzv, acc += sz * lzv, lz = 0, lzv = 0;
   void flip() { swap(1, r), f = 0; }
}:
template <class node>
struct Treap {
   vector<node> t;
   int n. r = -1:
   node get(int u) { return u != -1 ? t[u] : node(): }
   void recalc(int u) {
       int 1 = t[u].1, r = t[u].r;
       push(1), push(r), flip(1), flip(r);
       t[u].recalc(get(1), get(r));
   void push(int u) {
       if (u == -1 || !t[u].lz) return;
       int 1 = t[u].1. r = t[u].r:
       if (1 != -1) t[1].upd_lazy(t[u].lzv);
       if (r != -1) t[r].upd_lazy(t[u].lzv);
       t[u].lazv():
    void flip(int u) {
       if (u == -1 || !t[u].f) return:
       int 1 = t[u].1, r = t[u].r;
       if (1 != -1) t[1].f ^= 1;
       if (r != -1) t[r].f ^= 1;
       t[u].flip():
   int merge(int 1, int r) {
       if (min(1, r) == -1) return 1 != -1 ? 1 : r:
       push(1). push(r). flip(1). flip(r):
       int ans = (t[1].p < t[r].p) ? 1 : r;
       if (ans == 1) t[1].r = merge(t[1].r, r), recalc(1);
       if (ans == r) t[r].l = merge(l, t[r].l), recalc(r);
       if (t[ans].l != -1) t[t[ans].l].par = ans; // only if
             parent needed
       if (t[ans].r != -1) t[t[ans].r].par = ans; // only if
            parent needed
       return ans;
```

```
pii split(int u, int id) {
      if (u == -1) return \{-1, -1\}:
       push(u):
      flip(u):
       int szl = get(t[u].1).sz;
      if (szl >= id) {
          pii ans = split(t[u].1, id);
          if (ans.ss != -1) t[ans.ss].par = u; // only if
               parent needed
          if (ans.ff != -1) t[ans.ff].par = -1; // only if
               parent needed
          t[u].1 = ans.ss:
          recalc(u);
          return {ans.ff, u};
      pii ans = split(t[u].r, id - szl - 1);
      if (ans.ff != -1) t[ans.ff].par = u: // only if
            parent needed
       if (ans.ss != -1) t[ans.ss].par = -1; // only if
            parent needed
      t[u].r = ans.ff;
       recalc(u):
       return {u, ans.ss};
   int update(int u, int 1, int r, int v) {
       pii a = split(u, l), b = split(a.ss, r - l + 1);
       t[b.ff].upd_lazy(v);
       return merge(a.ff, merge(b.ff, b.ss));
   void print(int u) {
       if (u == -1) return;
       push(u), flip(u);
      print(t[u].1):
      cout << t[u].v << ' ';
       print(t[u].r):
   Treap(vi &v) : n(sz(v)) {
      for (int i = 0; i < n; i++) t.eb(v[i]), r = merge(r,</pre>
   }
};
```

# 6.3 bit-tricks

```
y = x & (x-1) // Turn off rightmost 1bit
y = x & (-x) // Isolate rightmost 1bit
y = x | (x-1) // Right propagate rightmost 1bit(fill in 1s)
y = x | (x+1) // Turn on rightmost 0bit
```

```
v = x & (x+1) // Isolate rightmost Obit
// If x is of long type, use __builtin_popcountl(x)
// If x is of long long type, use __builtin_popcountll(x)
// 1. Counts the number of ones(set bits) in an integer.
__builtin_popcount(x)
// 2. Checks the Parity of a number. Returns true(1) if the
// number has odd number of set bits, else it returns
// false(0) for even number of set bits.
__builtin_parity(x)
// 3. Counts the leading number of zeros of the integer.
__builtin_clz(x)
// 4. Counts the trailing number of zeros of the integer.
__builtin_ctz(x)
// 5. Returns 1 + the index of the least significant 1-bit.
__builtin_ffs(x) // If x == 0, returns 0.
// Iterate over non empty subsets of bitmask
for(int s=m;s;s=(s-1)&m) // Decreasing order
for (int s=0;s=s-m&m;) // Increasing order
```

#### 6.4 dsu

```
struct Dsu {
   vector<int> p; Dsu() {} Dsu(int N) : p(N, -1) {}
   int get(int x) { return p[x] < 0 ? x : get(p[x]); }
   bool sameSet(int a, int b) { return get(a) == get(b); }
   int size(int x) { return -p[get(x)]; }
   vector<vector<int>> S:
   void unite(int x, int y) {
       if ((x = get(x)) == (v = get(v))) { S.push back(\{-1\})
           ; return; }
      if (p[x] > p[y]) swap(x, y);
      S.push_back(\{x, y, p[x], p[y]\});
      p[x] += p[y], p[y] = x;
   void rollback() {
       auto a = S.back(); S.pop_back();
       if (a[0] != -1) p[a[0]] = a[2], p[a[1]] = a[3];
};
```

# 6.5 dynamic-connectivity

```
struct DC {
  int n; Dsu D;
  vector<vector<pair<int, int>>> t;
  DC(int N) : n(N), D(N), t(2 * N) {}
  // add edge p to all times in interval [1, r]
```

```
void upd(int 1, int r, pair<int, int> p) {
    for (1 += n, r += n; 1 < r; 1 >>= 1, r >>= 1) {
        if (1 & 1) t[1++].push_back(p);
        if (r & 1) t[--r].push_back(p);
    }
}
void process(int u = 1) { // process all queries
    for (auto &e : t[u]) D.unite(e.first, e.second);
    if (u >= n) {
        // do stuff with D at time u - n
    } else process(2 * u), process(2 * u + 1);
    for (auto &e : t[u]) D.rollback();
}
};
```

#### 6.6 hash-container

```
namespace{//add (#define tmpl template)(#define ty typename)
  tmpl<ty T> size_t mk_h(const T& v){return hash<T>()(v);}
  void h_cmb(size_t& h, const size_t& v)
  { h ^= v + 0x9e3779b9 + (h << 6) + (h >> 2); }
  tmpl<ty T> struct h_ct{size_t operator()(const T& v)const{
  size_t h=0;for(const auto& e:v){h_cmb(h,mk_h(e));}return h; }};
}namespace std{//support for pair<T,U>, vector<T> & map<T,U>
  tmpl<ty T, ty U> struct hash<pair<T, U>>{
    size_t operator()(const pair<T,U>& v) const
  {size_t h=mk_h(v.first);h_cmb(h, mk_h(v.second));return h;};
};
tmpl<ty... T>struct hash<vector<T...>:h_ct<vector<T...>>{};
tmpl<ty... T>struct hash<map<T...>>:h_ct<map<T...>>{};
}
```

#### 6.7 mo

#### 6.8 ordered-set

# 6.9 persistent-segment-tree-lazy

```
template <class T>
struct Node {
    T x, lz;
    int l = -1, r = -1;
};

template <class T>
struct Pstl {
    int N;
    vector<Node<T>> a;
    vector<int> head;

    T qneut() { return 0; }
    T merge(T l, T r) { return l + r; }
    T uneut() { return 0; }
    T accum(T u, T x) { return u + x; }
}
```

```
T apply(T x, T lz, int l, int r) { return x + (r - 1) *
     1z: }
int build(int vl. int vr) {
    if (vr - vl == 1) a.push_back({qneut(), uneut()}); //
         node construction
    else {
       int vm = (vl + vr) / 2, l = build(vl, vm), r =
            build(vm, vr):
       a.push_back({merge(a[1].x, a[r].x), uneut(), 1, r
            }): // querv merge
   return a.size() - 1;
T query(int 1, int r, int v, int v1, int vr, T acc) {
   if (1 >= vr || r <= vl) return gneut();</pre>
        // query neutral
   if (1 <= v1 && r >= vr) return apply(a[v].x. acc. vl.
         vr): // update op
   acc = accum(acc, a[v].lz);
        // update merge
   int vm = (vl + vr) / 2;
   return merge(query(1, r, a[v].1, v1, vm, acc), query(
        1, r, a[v].r, vm, vr, acc)); // query merge
int update(int 1, int r, T x, int v, int v1, int vr) {
   if (1 >= vr || r <= vl || r <= 1) return v;</pre>
   a.push back(a[v]):
   v = a.size() - 1:
   if (1 <= v1 && r >= vr) {
       a[v].x = apply(a[v].x, x, vl, vr); // update op
       a[v].lz = accum(a[v].lz, x); // update merge
   } else {
       int vm = (vl + vr) / 2:
       a[v].1 = update(1, r, x, a[v].1, v1, vm);
       a[v].r = update(1, r, x, a[v].r, vm, vr);
       a[v].x = merge(a[a[v].1].x, a[a[v].r].x); //
            query merge
   }
    return v;
Pst1() {}
Pstl(int N) : N(N) { head.push_back(build(0, N)); }
T query(int t, int 1, int r) {
   return query(1, r, head[t], 0, N, uneut()); // update
         neutral
```

# 6.10 persistent-segment-tree

```
// usage:
// Pst<Node<11>> pst;
// pst = {N};
// int newtime = pst.update(time, index, value);
// Node<ll> result = pst.query(newtime, left, right);
template <class T>
struct Node {
   T x:
    int 1 = -1, r = -1;
    Node(): x(0) {}
    Node(T x) : x(x) \{ \}
    Node (Node a, Node b, int l = -1, int r = -1) : x(a.x + b.
        x), 1(1), r(r) {}
};
template <class U>
struct Pst {
    int N:
    vector<U> a;
    vector<int> head:
    int build(int vl, int vr) {
       if (vr - vl == 1) a.push_back(U());
       else {
           int vm = (vl + vr) / 2, l = build(vl, vm),
              r = build(vm, vr):
           a.push_back(U(a[1], a[r], 1, r));
       return a.size() - 1:
    U query(int 1, int r, int v, int v1, int vr) {
       if (1 >= vr || r <= vl) return U();</pre>
       if (1 <= v1 && r >= vr) return a[v]:
       int vm = (v1 + vr) / 2;
       return U(query(1, r, a[v].1, v1, vm),
               query(1, r, a[v].r, vm, vr));
```

# 6.11 segment-tree-lazy

```
template <class T>
struct Stl {
   int n:
   vector<T> a, b;
   T qneut() { return -2e9; }
   T uneut() { return 0; }
   T merge(T x, T y) { return max(x, y); }
   void upd(int v, T x, int 1, int r)
      \{ a[v] += x, b[v] += x; \}
   Stl(int n = 0) : n(n), a(4 * n, qneut()),
      b(4 * n, uneut()) {}
   void push(int v, int vl, int vm, int vr) {
      upd(2 * v. b[v]. vl. vm):
      upd(2 * v + 1, b[v], vm, vr);
       b[v] = uneut();
   }
```

```
T querv(int 1, int r, int v=1, int v1=0, int vr=1e9) {
       vr = min(vr, n):
       if (1 <= v1 && r >= vr) return a[v];
       if (1 >= vr || r <= vl) return qneut();</pre>
       int vm = (v1 + vr) / 2;
       push(v. vl. vm. vr):
       return merge(query(1, r, 2 * v, v1, vm),
          query(1, r, 2 * v + 1, vm, vr));
   void update(int 1, int r, T x, int v = 1, int vl = 0,
          int vr = 1e9) {
       vr = min(vr, n);
       if (1 >= vr || r <= vl || r <= 1) return;</pre>
       if (1 <= vl && r >= vr) upd(v, x, vl, vr);
          int vm = (vl + vr) / 2:
          push(v, v1, vm, vr);
          update(1, r, x, 2 * v, vl, vm):
          update(1, r, x, 2 * v + 1, vm, vr);
          a[v] = merge(a[2 * v], a[2 * v + 1]);
   }
};
```

# 6.12 segment-tree

```
struct St {
    ll neut() { return 0; }
    ll merge(ll x, ll y) { return x + y; }

    int n; vector<ll> a;
    St(int n = 0) : n(n), a(2 * n, neut()) {}

    ll query(int l, int r) {
        ll x = neut(), y = neut();
        for (1 += n, r += n; 1 < r; 1 /= 2, r /= 2) {
            if (1 & 1) x = merge(x, a[1++]);
            if (r & 1) y = merge(a[--r], y);
        }
        return merge(x, y);
    }

    void update(int i, ll x) {
        for (a[i += n] = x; i /= 2;)
            a[i] = merge(a[2 * i], a[2 * i + 1]);
    }
};</pre>
```

#### 6.13 sparse-table

```
template <class T>
struct Sparse {
   T op(T a, T b) { return max(a, b); }
   vector<vector<T>> st:
   Sparse() {}
   Sparse(vector<T> a) : st{a} {
       int N = st[0].size();
       int npot = N <= 1 ? 1 : 32 - __builtin_clz(N);</pre>
       st.resize(npot):
       repx(i, 1, npot) rep(i, N + 1 - (1 << i))
       st[i].push_back(
           op(st[i - 1][i], st[i - 1][i + (1 << (i - 1))])
       ); // query op
   T query(int 1, int r) { // range must be nonempty!
       int i = 31 - builtin clz(r - 1):
       return op(st[i][1], st[i][r - (1 << i)]); // queryop</pre>
   }
};
```

# 6.14 unordered-map

# 7 imprimible

# 8 math

# 8.1 Linear Diophantine

```
ii extendedEuclid(ll a, ll b){
11 x, y; //a*x + b*y = gcd(a,b)
if (b == 0) return {1, 0};
auto p = extendedEuclid(b, a%b):
x = p.second:
y = p.first - (a/b)*x;
if(a*x + b*y == -_gcd(a,b)) x=-x, y=-y;
return {x, y};
pair<ii, ii> diophantine(ll a, ll b, ll r){
//a*x+b*y=r where r is multiple of gcd(a,b);
11 d = _{gcd(a, b)};
a/=d: b/=d: r/=d:
auto p = extendedEuclid(a, b);
p.first*=r: p.second*=r:
assert(a*p.first + b*p.second == r);
return {p, {-b, a}}; //solutions: p+t*ans.second
```

#### 8.2 arithmetic

```
inline int floor_log2(int n) { return n <= 1 ? 0 : 31 -</pre>
    __builtin_clz(n); }
inline int ceil log2(int n) { return n <= 1 ? 0 : 32 -
    __builtin_clz(n - 1); }
inline 11 floordiv(11 a, 11 b) {
   return a / b - ((a ^ b) < 0 && a % b):
inline 11 ceildiv(11 a, 11 b) {
   return a / b + ((a ^ b) >= 0 && a % b);
ll binexp(ll a, ll e) {
   ll res = 1: // neutral element
   while (e) {
      if (e & 1) res = res * a; // multiplication
                               // multiplication
      a = a * a:
       e >>= 1:
   }
   return res:
```

#### 8.3 crt

```
pair<11, ll> solve_crt(const vector<pair<11, ll>> &eqs) {
    ll a0 = eqs[0].first, p0 = eqs[0].second;
    repx(i, 1, eqs.size()) {
        ll a1 = eqs[i].first, p1 = eqs[i].second;
        ll k1, k0;
        ll d = ext_gcd(p1, p0, k1, k0);
        a0 -= a1;
        if (a0 % d != 0) return {-1, -1};
        p0 = p0 / d * p1;
        a0 = a0 / d * k1 % p0 * p1 % p0 + a1;
        a0 = (a0 % p0 + p0) % p0;
    }
    return {a0, p0};
}
```

# 8.4 discrete-log

```
// discrete logarithm log a(b).
// solve b \hat{x} = a \pmod{M} for the smallest x.
// returns -1 if no solution is found.
//
// time: O(sqrt(M))
11 dlog(ll a, ll b, ll M) {
   11 k = 1, s = 0;
   while (true) {
       11 g = \_gcd(b, M);
       if (g <= 1) break;</pre>
       if (a == k) return s:
       if (a % g != 0) return -1;
       a \neq g, M \neq g, s += 1, k = b \neq g * k % M;
   ll N = sqrt(M) + 1;
   umap<11. 11> r:
   rep(q, N + 1) {
       r[a] = q;
       a = a * b % M;
   ll bN = binexp(b, N, M), bNp = k;
   repx(p, 1, N + 1) {
       bNp = bNp * bN % M;
       if (r.count(bNp)) return N * p - r[bNp] + s;
   return -1;
```

#### 8.5 fft

```
using cd = complex<double>:
const double PI = acos(-1);
// compute the DFT of a power-of-two-length sequence.
// if 'inv' is true, computes the inverse DFT.
void fft(vector<cd> &a, bool inv) {
   int N = a.size(), k = 0, b:
   assert(N == 1 << __builtin_ctz(N));</pre>
   repx(i, 1, N) {
       for (b = N >> 1; k & b;) k ^= b, b >>= 1;
       if (i < (k ^= b)) swap(a[i], a[k]);</pre>
   for (int 1 = 2; 1 <= N; 1 <<= 1) {</pre>
       double ang = 2 * PI / 1 * (inv ? -1 : 1);
       cd wl(cos(ang), sin(ang));
       for (int i = 0; i < N; i += 1) {</pre>
          cd w = 1;
           rep(j, 1 / 2) {
               cd u = a[i + j], v = a[i + j + 1 / 2] * w;
              a[i + j] = u + v;
              a[i + j + 1 / 2] = u - v;
              w *= w1;
          }
       }
   if (inv) rep(i, N) a[i] /= N;
const 11 MOD = 998244353, ROOT = 15311432;
// const 11 MOD = 2130706433, ROOT = 1791270792:
// const 11 MOD = 922337203673733529711, ROOT =
    532077456549635698311;
void find_root_of_unity(11 M) {
   11 c = M - 1, k = 0:
   while (c \% 2 == 0) c /= 2, k += 1:
   // find proper divisors of M - 1
   vector<ll> divs;
   for (11 d = 1; d < c; d++) {</pre>
       if (d * d > c) break:
       if (c \% d == 0) rep(i, k + 1) divs.push_back(d << i);
   rep(i, k) divs.push_back(c << i);
```

```
// find any primitive root of M
   11 G = -1:
   repx(g, 2, M) {
      bool ok = true:
      for (ll d : divs) ok &= (binexp(g, d, M) != 1);
      if (ok) {
          G = g;
          break;
      }
   }
   assert(G != -1):
   ll w = binexp(G, c, M);
   cerr << "M = c * 2^k + 1" << endl;
   cerr << " M = " << M << endl;
   cerr << " c = " << c << endl;
   cerr << " k = " << k << endl:
   cerr << " w^(2^k) == 1" << endl;
   cerr << " w = g^{(M-1)/2k} = g^c << endl;
   cerr << " g = " << G << endl;</pre>
   cerr << " w = " << w << endl;
// compute the DFT of a power-of-two-length sequence, modulo
     a special prime
// number with an Nth root of unity, where N is the length
    of the sequence.
void ntt(vector<11> &a. bool inv) {
   vector<ll> wn;
   for (11 p = ROOT; p != 1; p = p * p % MOD) wn.push_back(p
       );
   int N = a.size(), k = 0, b:
   assert(N == 1 << \_builtin_ctz(N) && N <= 1 << wn.size()) // 0 -> no solutions
   rep(i, N) a[i] = (a[i] \% MOD + MOD) \% MOD:
   repx(i, 1, N) {
      for (b = N >> 1; k & b;) k ^= b, b >>= 1;
      if (i < (k ^= b)) swap(a[i], a[k]);</pre>
   for (int 1 = 2: 1 <= N: 1 <<= 1) {
      11 wl = wn[wn.size() - __builtin_ctz(1)];
      if (inv) wl = multinv(wl, MOD);
      for (int i = 0; i < N; i += 1) {</pre>
          11 w = 1:
          repx(j, 0, 1 / 2) {
```

```
11 u = a[i + i], v = a[i + i + 1 / 2] * w %
              a[i + j] = (u + v) \% MOD;
              a[i + j + 1 / 2] = (u - v + MOD) \% MOD;
              w = w * wl % MOD;
      }
   }
   11 q = multinv(N, MOD);
   if (inv) rep(i, N) a[i] = a[i] * q % MOD;
void convolve(vector<cd> &a, vector<cd> b, int n) {
   n = 1 \ll (32 - builtin clz(2 * n - 1)):
   a.resize(n), b.resize(n);
   fft(a, false), fft(b, false);
   rep(i, n) a[i] *= b[i];
   fft(a, true):
```

#### 8.6 gauss

```
const double EPS = 1e-9:
// solve a system of equations.
// complexity: O(\min(N, M) * N * M)
// 'a' is a list of rows
// the last value in each row is the result of the equation
// return values:
// 1 -> unique solution, stored in 'ans'
// -1 -> infinitely many solutions, one of which is stored
// UNTESTED
int gauss(vector<vector<double>> a, vector<double> &ans) {
   int N = a.size(), M = a[0].size() - 1;
   vector<int> where(M, -1);
   for (int j = 0, i = 0; j < M && i < N; j++) {
       int sel = i:
       repx(k, i, N) if (abs(a[k][j]) > abs(a[sel][j])) sel
            = k:
       if (abs(a[sel][i]) < EPS) continue;</pre>
       repx(k, j, M + 1) swap(a[sel][k], a[i][k]);
       where[j] = i;
       rep(k, N) if (k != i) {
```

```
double c = a[k][j] / a[i][j];
    repx(l, j, M + 1) a[k][l] -= a[i][l] * c;
}
i++;
}
ans.assign(M, 0);
rep(i, M) if (where[i] != -1) ans[i] = a[where[i]][M] / a
    [where[i]][i];
rep(i, N) {
    double sum = 0;
    rep(j, M) sum += ans[j] * a[i][j];
    if (abs(sum - a[i][M]) > EPS) return 0;
}
rep(i, M) if (where[i] == -1) return -1;
return 1;
```

#### 8.7 matrix

```
typedef vector<vector<double>> Mat;
Mat matmul(Mat 1, Mat r) {
   int n = 1.N, m = r.M, p = 1.M; assert(1.M == r.N);
   Mat a(n, vector<double>(m)): // neutral
   rep(i, n) rep(j, m)
       rep(k, p) a[i][i] = a[i][i] + 1[i][k] * r[k][i]:
   return a:
}
double reduce(vector<vector<double>> &A) {
   int n = A.size(), m = A[0].size();
   int i = 0, j = 0; double r = 1.;
   while (i < n && j < m) {</pre>
       int 1 = i:
       repx(k, i+1, n) if(abs(A[k][j]) > abs(A[l][j])) l=k;
       if (abs(A[1][i]) < EPS) \{ i++; r = 0.; continue; \}
       if (1 != i) \{ r = -r : swap(A[i], A[1]) : \}
       for (int k = m - 1; k \ge j; k--) A[i][k] /= A[i][j];
       repx(k, 0, n) {
           if (k == i) continue;
           for(int l=m-1;l>=j;l--)A[k][l]-=A[k][j]*A[i][l];
       }
       i++, j++;
   return r; // returns determinant
```

#### 8.8 mobius

```
short mu[MAXN] = {0,1};
void mobius(){
  repx(i,1,MAXN)if(mu[i])for(int j=i+i;j<MAXN;j+=i)mu[j]-=mu[
   i];
}</pre>
```

```
8.9 mod
11 binexp(11 a, 11 e, 11 M) {
    assert(e >= 0);
    ll res = 1 % M:
    while (e) {
       if (e & 1) res = res * a % M:
       a = a * a % M:
       e >>= 1:
    }
    return res;
11 multinv(ll a, ll M) { return binexp(a, M - 2, M); }
// calculate gcd(a, b).
// also, calculate x and y such that:
// a * x + b * v == gcd(a, b)
11
// time: O(log min(a, b))
// (ignoring complexity of arithmetic)
ll ext_gcd(ll a, ll b, ll &x, ll &y) {
    if (b == 0) {
       x = 1, v = 0:
       return a;
    11 d = ext_gcd(b, a % b, y, x);
    y = a / b * x;
    return d:
// compute inverse with any M.
// a and M must be coprime for inverse to exist!
ll multinv euc(ll a. ll M) {
    11 x, v;
    ext_gcd(a, M, x, y);
    return x:
// multiply two big numbers (~10^18) under a large modulo,
     without resorting to
```

```
// bigints.
11 bigmul(11 x, 11 y, 11 M) {
   11 z = 0:
   while (v) {
       if (y \& 1) z = (z + x) \% M;
       x = (x << 1) \% M, v >>= 1:
   return z;
// all modular inverses from 1 to inv.size()-1
void multinv all(vector<ll> &inv) {
   inv[1] = 1;
   repx(i, 2, inv.size())
       inv[i] = MOD - (MOD / i) * inv[MOD % i] % MOD:
// change to __int128 if checking numbers over 10^9
bool isprime(ll n) {
   if (n < 2 | | n % 6 % 4 != 1) return n - 2 < 2:
   11 A[] = \{2,325,9375,28178,450775,9780504,1795265022\};
   ll s = builtin ctzll(n - 1), d = n >> s:
   for (int a : A) {
      11 p = binexp(a, d, n), i = s;
       while (p != 1 && p != n - 1 && a % n && i--) p = p *
           p % n;
       if (p != n - 1 && i != s) return 0;
   return 1;
struct Mod {
   static const int M = 1e9 + 7;
   Mod(11 aa) : a((aa % M + M) % M) {}
   Mod operator+(Mod rhs) const { return (a + rhs.a) % M; }
   Mod operator-(Mod rhs) const { return (a - rhs.a + M) % M
        : }
   Mod operator-() const { return Mod(0) - *this: }
   Mod operator*(Mod rhs) const { return (11)a * rhs.a % M;
   Mod operator+=(Mod rhs) { return *this = *this + rhs; }
   Mod operator = (Mod rhs) { return *this = *this - rhs: }
   Mod operator*=(Mod rhs) { return *this = *this * rhs; }
   Mod bigmul(ll big) const { return ::bigmul(a, big, M); }
```

```
Mod binexp(ll e) const { return ::binexp(a, e, M); }
   // Mod multinv() const { return ::multinv(a. M): } //
        prime M
   Mod multinv() const { return ::multinv euc(a, M): } //
        possibly composite M
}:
// dynamic modulus
struct DMod {
   int a, M;
   DMod(11 aa. 11 m) : M(m). a((aa % m + m) % m) {}
   DMod operator+(DMod rhs) const { return {(a + rhs.a) % M,
   DMod operator-(DMod rhs) const { return {(a - rhs.a + M)
        % M. M}: }
   DMod operator-() const { return DMod(0, M) - *this; }
   DMod operator*(DMod rhs) const { return {(11)a * rhs.a %
   DMod operator+=(DMod rhs) { return *this = *this + rhs; }
   DMod operator-=(DMod rhs) { return *this = *this - rhs; }
   DMod operator*=(DMod rhs) { return *this = *this * rhs; }
   DMod bigmul(ll big) const { return {::bigmul(a, big, M),
        M}: }
   DMod binexp(ll e) const { return {::binexp(a, e, M), M};
   DMod multinv() const { return {::multinv(a, M), M}; } //
        prime M
   // DMod multinv() const { return {::multinv_euc(a, M), M
        }; } // possibly composite M
};
```

# 8.10 primes

```
// counts the divisors of a positive integer in O(sqrt(n))
ll count_divisors(ll x) {
    ll divs = 1, i = 2;
    for (ll divs = 1, i = 2; x > 1; i++) {
        if (i * i > x) {
            divs *= 2;
            break;
        }
        for (ll d = divs; x % i == 0; x /= i) divs += d;
    }
    return divs;
```

```
// gets the prime factorization of a number in O(sqrt(n))
vector<pair<11, int>> factorize(11 x) {
    vector<pair<11, int>> f;
    for (11 k = 2: x > 1: k++) {
       if (k * k > x) {
           f.push_back(\{x, 1\});
           break:
       int n = 0:
       while (x \% k == 0) x /= k, n++:
       if (n > 0) f.push_back(\{k, n\});
    return f:
// iterate over all divisors of a number.
// divisor count upper bound: n^(1.07 / ln ln n)
template <class OP>
void divisors(ll x, OP op) {
    auto facts = factorize(x);
    vector<int> f(facts.size());
    while (true) {
       11 v = 1:
       rep(i, f.size()) rep(j, f[i]) y *= facts[i].first;
       op(y);
       int i:
       for (i = 0; i < f.size(); i++) {</pre>
           f[i] += 1:
           if (f[i] <= facts[i].second) break;</pre>
           f[i] = 0;
       }
       if (i == f.size()) break;
// computes euler totative function phi(x), counting the
     amount of integers in
// [1, x] that are coprime with x.
11
// time: O(sqrt(x))
11 phi(11 x) {
   11 phi = 1, k = 2:
   for (; x > 1; k++) {
       if (k * k > x) {
           phi *= x - 1:
           break:
```

```
}
ll k1 = 1, k0 = 0;
while (x % k == 0) x /= k, k0 = k1, k1 *= k;
phi *= k1 - k0;
}
return phi;
}
// isprime is in mod.cpp
```

#### 8.11 simplex

```
// Solves a general linear maximization problem: maximize $c
     ^T x$ subject to $Ax \le b$, $x \ge 0$.
// Returns -inf if there is no solution, inf if there are
    arbitrarily good solutions, or the maximum value of $c^
    T x$ otherwise.
// The input vector is set to an optimal $x$ (or in the
    unbounded case, an arbitrary solution fulfilling the
    constraints).
// Numerical stability is not guaranteed. For better
    performance, define variables such that x = 0 is
    viable.
// Usage:
// vvd A = \{\{1,-1\}, \{-1,1\}, \{-1,-2\}\};
// vd b = \{1.1.-4\}, c = \{-1.-1\}, x:
// T val = LPSolver(A, b, c).solve(x);
// Time: O(NM * \t pivots), where a pivot may be e.g. an edge
     relaxation. O(2^n) in the general case.
#include "../common.h"
typedef double T; // long double, Rational, double + mod<P
    >...
typedef vector<T> vd;
typedef vector<vd> vvd;
const T eps = 1e-8, inf = 1 / .0;
#define MP make_pair
#define lti(X) \
   if (s == -1 \mid | MP(X[i], N[i]) < MP(X[s], N[s])) s = i
struct LPSolver {
   int m, n;
   vector<int> N. B:
   vvd D:
   LPSolver(const vvd &A, const vd &b, const vd &c) : m(b.
        size()), n(c.size()), N(n + 1), B(m), D(m + 2), vd(n)
```

```
+ 2)) {
   rep(i, m) rep(j, n) D[i][j] = A[i][j];
   rep(i, m) {
      B[i] = n + i:
       D[i][n] = -1;
       D[i][n + 1] = b[i]:
   rep(j, n) {
       N[j] = j;
       D[m][i] = -c[i];
   N[n] = -1:
   D[m + 1][n] = 1;
void pivot(int r, int s) {
   T *a = D[r].data(), inv = 1 / a[s]:
   rep(i, m + 2) if (i != r \&\& abs(D[i][s]) > eps) {
      T *b = D[i].data(), inv2 = b[s] * inv:
       repx(i, 0, n + 2) b[i] -= a[i] * inv2:
       b[s] = a[s] * inv2;
   rep(j, n + 2) if (j != s) D[r][j] *= inv;
   rep(i, m + 2) if (i != r) D[i][s] *= -inv;
   D[r][s] = inv;
   swap(B[r], N[s]);
bool simplex(int phase) {
   int x = m + phase - 1:
   for (;;) {
       int s = -1;
       rep(j, n + 1) if (N[j] != -phase) ltj(D[x]);
      if (D[x][s] >= -eps) return true;
       int r = -1:
       rep(i, m) {
          if (D[i][s] <= eps) continue;</pre>
          if (r == -1 || MP(D[i][n + 1] / D[i][s], B[i])
                < MP(D[r][n + 1] / D[r][s], B[r])) r = i
       if (r == -1) return false;
       pivot(r. s):
T solve(vd &x) {
   int r = 0:
   repx(i, 1, m) if (D[i][n + 1] < D[r][n + 1]) r = i;
   if (D[r][n + 1] < -eps) {
```

#### 8.12 theorems

#### Burnside lemma

Tomemos imagenes x en X y operaciones (g: X -> X) en G. Si #g es la cantidad de imagenes que son puntos fijos de g, entonces la cantidad de objetos es '(sum\_{g} in G} #g) / |G|' Es requisito que G tenga la operacion identidad, que toda operacion tenga inversa y que todo par de operaciones tenga su combinacion.

#### Rational root theorem

Las raices racionales de un polinomio de orden n con coeficientes enteros A[i] son de la forma p / q, donde p y q son coprimos, p es divisor de A[0] y q es divisor de A[n]. Notar que si A[0] = 0, cero es raiz, se puede dividir el polinomio por x y aplica nuevamente el teorema.

#### Petersens theorem

Every cubic and bridgeless graph has a perfect matching.

```
Number of divisors for powers of 10 (0,1) (1,4) (2,12) (3,32) (4,64) (5,128) (6,240) (7,448) (8,768) (9,1344) (10,2304) (11,4032) (12,6720) (13,10752) (14,17280) (15,26880) (16,41472) (17,64512) (18,103680)
```

Kirchoff Theorem: Sea A la matriz de adyacencia del multigrafo (A[u][v] indica la cantidad de aristas entre u y v) Sea D una matriz diagonal tal que D[v][v] es igual al grado de v (considerando auto aristas y multi aristas). Sea L = A - D. Todos los cofactores de L son iguales y equivalen a la cantidad de Spanning Trees del grafo. Un cofactor (i,j) de L es la multiplicacin de  $(-1)^{i}$  + j con el determinant de la matriz al quitar la fila i y la columna j

Prufer Code: Dado un rbol con los nodos indexados: busca la hoja de menor ndice, brrala y anota el ndice del nodo al que estaba conectado. Repite el paso anterior n-2 veces. Lo anterior muestra una biveccin entre los arreglos de tamao n-2 con elementos en [1, n] v los rboles de n nodos, por lo que hay n^{n-2} spanning trees en un grafo completo. Corolario: Si tenemos k componentes de tamaos s1,s2,...,sk entonces podemos hacerlos conexos agregando k-1 aristas entre nodos de s1\*s2\*...\*sk\*n^{k-2} formas Combinatoria Catalan:  $C_{n+1} = sum(C_i*C_{n-i})$  for  $i \in [0, n]$ Catalan:  $C_n = \frac{1}{n+1}*\frac{2n}{n}$ Sea C\_n^k las formas de poner n+k pares de parntesis, con los primeros k parntesis abiertos (esto es, hay 2n + 2k carcteres), se tiene que  $C_n^k = (2n+k-1)*(2n+k)/(n*(n+k+1)) * C_{n-1}^k$ Sea D\_n el nmero de permutaciones sin puntos fijos, entoces  $D n = (n-1)*(D \{n-1\} + D \{n-2\}), D 0 = 1, D 1 = 0$ 

#### 8.13 tonelli-shanks

```
ll legendre(ll a, ll p) {
   if (a % p == 0) return 0; if (p == 2) return 1;
   return binexp(a, (p - 1) / 2, p);
// sqrt(n) mod p (p must be a prime)
// rnd(a, b) return a random number in [a, b]
11 tonelli_shanks(ll n, ll p) {
   if (n == 0) return 0:
   if (legendre(n, p) != 1) return -1; // no existe
   if (p == 2) return 1;
   ll s = builtin ctzll(p - 1):
   11 q = (p - 1LL) >> s, z = rnd(1, p - 1);
   if (s == 1) return binexp(n, (p + 1) / 4LL, p);
   while (legendre(z, p) != p - 1) z = rnd(1, p - 1);
   11 c = binexp(z, q, p), r = binexp(n, (q + 1) / 2, p);
   11 t = binexp(n, q, p), m = s;
   while (t != 1) {
      11 i = 1, ts = (t * t) \% p;
       while (ts != 1) i++, ts = (ts * ts) % p:
      repx(_, 0, m - i - 1) b = (b * b) \% p;
      r = r*b\%p; c = b*b\%p; t = t*c\%p; m = i;
   }
   return r:
```