Team Notebook

Pontificia Universidad Católica de Chile - Laranjas.clear()

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1 data structures

1.1 dsu

```
struct Dsu {
    vector\langle int \rangle p; Dsu(int N = 0) : p(N, -1) {}
   int get(int x) { return p[x] < 0 ? x : get(p[x]); }</pre>
   bool sameSet(int a, int b) { return get(a) == get(b); }
   int size(int x) { return -p[get(x)]; }
   vector<vector<int>> S;
   void unite(int x, int y) {
       if ((x = get(x)) == (y = get(y)))
           return S.push_back({-1});
       if (p[x] > p[y]) swap(x, y);
       S.push_back(\{x, y, p[x], p[y]\});
       p[x] += p[y], p[y] = x;
   void rollback() {
       auto a = S.back(); S.pop_back();
       if (a[0] != -1) p[a[0]] = a[2], p[a[1]] = a[3];
};
```

1.2 fenwick-tree

```
int ft[MAXN+1]; // add dimension for multi-d
void upd(int i0, int v){ // add v to i0th element
    for(int i=i0+1;i<=MAXN;i+=i&-i)ft[i]+=v;//+ fors
}
int get(int i0){ // get sum of range [0,i0)
    int r=0; // add fors
    for(int i=i0;i;i-=i&-i)r+=ft[i];
    return r;
}
int get_sum(int i0,int i1){//sum of [i0,i1)
    return get(i1)-get(i0);
}</pre>
```

1.3 link-cut-tree

```
//mostly generic
inline int joinD(int d1, int d2){
 if(d1==N_DEL)return d2;if(d2==N_DEL)return d1;return mOp(
inline int joinVD(int v, int d){return d==N_DEL ? v : mOp(v,
     d):}
struct Node_t{
 int sz, nVal, tVal, d; bool rev;
 Node_t *c[2], *p;
 Node_t(int v) : sz(1), nVal(v), tVal(v), d(N_DEL), rev(0),
       }(0)a
 c[0]=c[1]=0:
 bool isRoot(){return !p || (p->c[0] != this && p->c[1] !=
 void push(){
 if(rev){
   rev=0; swap(c[0], c[1]); fore(x,0,2)if(c[x])c[x]->rev^=1;
 nVal=joinVD(nVal, d); tVal=joinVD(tVal, dOnSeg(d, sz));
 fore(x,0,2)if(c[x])c[x]->d=joinD(c[x]->d, d);
 d=N_DEL;
 }
 void upd();
typedef Node_t* Node;
int getSize(Node r){return r ? r->sz : 0;}
int getPV(Node r){
 return r ? joinVD(r->tVal, dOnSeg(r->d,r->sz)) : N_VAL;}
void Node t::upd(){
 tVal = qOp(qOp(getPV(c[0]), joinVD(nVal, d)), getPV(c[1]))
 sz = 1 + getSize(c[0]) + getSize(c[1]);
void conn(Node c, Node p, int il){if(c)c->p=p;if(il>=0)p->c
    [!ill=c:}
void rotate(Node x){
 Node p = x-p, g = p-p;
 bool gCh=p->isRoot(), isl = x==p->c[0];
 conn(x->c[isl],p,isl); conn(p,x,!isl);
 conn(x,g,gCh?-1:(p==g->c[0])); p->upd();
void spa(Node x){//splay
 while(!x->isRoot()){
 Node p = x-p, g = p-p;
 if(!p->isRoot())g->push();
 p->push(); x->push();
 if(!p-)isRoot())rotate((x=-p-)c[0])==(p=-p-)c[0])? p : x);
 rotate(x):
```

```
x->push(): x->upd():
Node exv(Node x){//expose
 Node last=0:
 for (Node y=x; y; y=y->p) spa(y), y->c[0]=last, y->upd(), last=
 spa(x);
 return last;
void mkR(Node x){exv(x);x->rev^=1;}//makeRoot
Node getR(Node x){exv(x); while(x->c[1])x=x->c[1]; spa(x);
Node lca(Node x, Node y){exv(x); return exv(y);}
bool connected(Node x, Node y){exv(x);exv(y); return x==y?1:
void link(Node x, Node y){mkR(x); x->p=y;}
void cut(Node x, Node y) \{mkR(x); exv(y); y->c[1]->p=0; y->c
    [1]=0;}
Node father(Node x){
 exv(x): Node r=x->c[1]:
 if(!r)return 0;
 while (r->c[0])r=r->c[0];
 return r;
void cut(Node x){ // cuts x from father keeping tree root
 exv(father(x));x->p=0;
int query(Node x, Node y){mkR(x); exv(y); return getPV(y);}
void modifv(Node x. Node v. int d){mkR(x):exv(v):v->d=joinD(
    y->d,d);}
Node lift rec(Node x. int t){
 if(!x)return 0;
 if(t==getSize(x->c[0])){spa(x);return x;}
 if(t<getSize(x->c[0]))return lift_rec(x->c[0],t);
 return lift_rec(x->c[1],t-getSize(x->c[0])-1);
Node lift(Node x, int t) { // t-th ancestor of x (lift(x.1)
    is x's father)
 exv(x);return lift_rec(x,t);}
int depth(Node x){ // distance from x to its tree root
 exv(x);return getSize(x)-1;}
```

1.4 persistent-segment-tree-lazy

```
template <class T>
struct Node {
    T x, lz;
    int 1 = -1, r = -1;
};
```

```
template <class T>
struct Pstl {
   int N:
   vector<Node<T>> a:
   vector<int> head;
   T qneut() { return 0; }
   T merge(T 1, T r) { return 1 + r; }
   T uneut() { return 0; }
   T accum(T u, T x) { return u + x; }
   T apply(T x, T lz, int l, int r) { return x + (r - 1) *
   int build(int vl. int vr) {
      if (vr - vl == 1) a.push_back({qneut(), uneut()}); //
            node construction
      else {
          int vm = (vl + vr) / 2, l = build(vl, vm), r =
               build(vm. vr):
          a.push_back({merge(a[1].x, a[r].x), uneut(), 1, r
               }); // query merge
      }
       return a.size() - 1;
   T query(int 1, int r, int v, int v1, int vr, T acc) {
      if (1 >= vr || r <= vl) return qneut();</pre>
           // guery neutral
      if (1 \le vl \&\& r \ge vr) return apply(a[v].x, acc, vl,
            vr): // update op
      acc = accum(acc, a[v].lz);
           // update merge
      int vm = (vl + vr) / 2:
      return merge(query(1, r, a[v].1, v1, vm, acc), query(
           1, r, a[v].r, vm, vr, acc)); // query merge
   int update(int 1, int r, T x, int v, int v1, int vr) {
      if (1 >= vr || r <= vl || r <= 1) return v;
      a.push_back(a[v]);
      v = a.size() - 1:
      if (1 <= v1 && r >= vr) {
          a[v].x = apply(a[v].x, x, vl, vr); // update op
          a[v].lz = accum(a[v].lz, x);  // update merge
      } else {
          int vm = (vl + vr) / 2:
          a[v].1 = update(1, r, x, a[v].1, v1, vm);
          a[v].r = update(1, r, x, a[v].r, vm, vr);
          a[v].x = merge(a[a[v].1].x, a[a[v].r].x); //
               query merge
```

1.5 persistent-segment-tree

```
// usage:
// Pst<Node<11>> pst:
// pst = {N};
// int newtime = pst.update(time, index, value);
// Node<ll> result = pst.querv(newtime. left. right):
template <class T>
struct Node {
   T x:
   int 1 = -1, r = -1;
   Node(): x(0) {}
   Node(T x) : x(x) \{ \}
   Node (Node a, Node b, int l = -1, int r = -1): x(a.x + b.
        x), 1(1), r(r) {}
};
template <class U>
struct Pst {
   int N:
    vector<U> a:
    vector<int> head:
   int build(int vl, int vr) {
       if (vr - vl == 1) a.push_back(U());
           int vm = (vl + vr) / 2, l = build(vl, vm),
              r = build(vm, vr);
           a.push_back(U(a[1], a[r], 1, r));
```

```
return a.size() - 1:
   }
   U query(int 1, int r, int v, int v1, int vr) {
       if (1 >= vr || r <= vl) return U();</pre>
       if (1 <= v1 && r >= vr) return a[v]:
       int vm = (vl + vr) / 2;
       return U(query(1, r, a[v].1, v1, vm),
               query(1, r, a[v].r, vm, vr));
   int update(int i, U x, int v, int vl, int vr) {
       a.push_back(a[v]);
       v = a.size() - 1;
       if (vr - vl == 1) a[v] = x:
       else {
          int vm = (vl + vr) / 2;
          if (i < vm) a[v].l = update(i, x, a[v].l, vl, vm)</pre>
           else a[v].r = update(i, x, a[v].r, vm, vr);
          a[v] = U(a[a[v].1], a[a[v].r], a[v].1, a[v].r);
      }
       return v;
   Pst(int N) : N(N) { head.push_back(build(0, N)); }
   U query(int t, int l, int r) {
       return query(1, r, head[t], 0, N);
   int update(int t, int i, U x) {
       return head.push_back(update(i, x, head[t], 0, N)).
           head.size() - 1;
};
```

1.6 rmq-lineal

```
typedef int tf; // O(n) construction, O(1) query
struct rmq{
   int n; tf INF=1e9;//change sign of INF for MAX
   vector<unsigned int> mk; vector<tf> bk,v;
   rmq(){}
   tf op(tf a, tf b){return min(a,b);}//change for maximum
   int f(int x){return x>>5;}
   rmq(vector<tf> &vv):n(SZ(vv)),mk(n),bk(n,INF),v(vv){
      unsigned int lst=0;
      for(int i=0;i<SZ(v);i++,lst<<=1){</pre>
```

```
bk[f(i)]=op(bk[f(i)],v[i]);
          while(lst&&v[i-_builtin_ctz(lst)]>v[i]) lst^=lst
               &-lst: //MIN
          //while(lst&&v[i- builtin ctz(lst)]<v[i]) lst^=
               lst&-lst: //MAX
          mk[i]=++lst:
       for(int k=1,top=f(n);(1<<k)<=top;k++)fore(i,0,top)if(
            i+(1<<k)<=top)
          bk[top*k+i]=op(bk[top*(k-1)+i], bk[top*(k-1)+i]
               +(1<<k-1)]):
   tf get(int st, int en){
       return v[en-31+ builtin clz(mk[en]&((111<<en-st+1)
   tf query(int s, int e){ //[s,e]
       int b1=f(s),b2=f(e),top=f(n);
       if(b1==b2) return get(s,e);
       tf ans=op(get(s,(b1+1)*32-1), get(b2*32.e)); s=(b1+1)
            *32; e=b2*32-1;
       if(s<=e){
          int k=31-__builtin_clz(f(e-s+1));
          ans=op(ans,op(bk[top*k+f(s)],bk[top*k+f(e)-(1<<k)
      }
       return ans:
};
```

1.7 segment-tree-2d

```
// #define MAXN 1024 #define op(a,b) (a+b) #define NEUT 0
int n.m: int a[MAXN][MAXN].st[2*MAXN][2*MAXN];
void build(){
   repx(i, 0, n) repx(j, 0, m) st[i+n][j+m] = a[i][j];
   repx(i, 0, n) for(int j = m-1; j; --j)
       st[i+n][j] = op(st[i+n][j<<1], st[i+n][j<<1|1]);
   for(int i = n-1; i; --i) repx(j, 0, 2*m)
       st[i][i] = op(st[i << 1][i], st[i << 1|1][i]):
void upd(int x, int y, int v){
   st[x+n][y+m]=v;
   for(int j = y+m; j > 1; j >>= 1)
       st[x+n][i>>1] = op(st[x+n][i], st[x+n][i^1]):
   for(int i = x+n; i > 1; i >>= 1) for(int j=y+m; j; j>>=1)
       st[i>>1][j] = op(st[i][j], st[i^1][j]);
}
int query(int x0, int x1, int y0, int y1){
```

```
int r=NEUT;
for(int i0=x0+n, i1=x1+n; i0<i1; i0>>=1, i1>>=1){
    int t[4], q = 0;
    if(i0 & 1) t[q++] = i0++;
    if(i1 & 1) t[q++] = --i1;
    repx(k,0,q)
        for(int j0=y0+m, j1=y1+m; j0<j1; j0>>=1,j1>>=1){
        if(j0 & 1) r = op(r, st[t[k]][j0++]);
        if(j1 & 1) r = op(r, st[t[k]][--j1]);
    }
}
return r;
}
```

1.8 segment-tree-beats

```
struct Node {
   11 \text{ s. mx1. mx2. mxc. mn1. mn2. mnc. } 1z = 0:
   Node(): s(0), mx1(LLONG_MIN), mx2(LLONG_MIN), mxc(0),
        mn1(LLONG_MAX), mn2(LLONG_MAX), mnc(0) {}
   Node(ll x): s(x), mx1(x), mx2(LLONG MIN), mxc(1), mn1(x)
        , mn2(LLONG_MAX), mnc(1) {}
   Node(const Node &a. const Node &b) {
      // add
      s = a.s + b.s;
      // min
      if (a.mx1 > b.mx1) mx1 = a.mx1, mxc = a.mxc, mx2 =
            max(b.mx1. a.mx2):
       if (a.mx1 < b.mx1) mx1 = b.mx1, mxc = b.mxc, mx2 =
            max(a.mx1, b.mx2);
       if (a.mx1 == b.mx1) mx1 = a.mx1, mxc = a.mxc + b.mxc,
            mx2 = max(a.mx2, b.mx2);
       // max
      if (a.mn1 < b.mn1) mn1 = a.mn1, mnc = a.mnc, mn2 =
            min(b.mn1, a.mn2):
       if (a.mn1 > b.mn1) mn1 = b.mn1, mnc = b.mnc, mn2 =
            min(a.mn1, b.mn2):
       if (a.mn1 == b.mn1) mn1 = a.mn1, mnc = a.mnc + b.mnc,
            mn2 = min(a.mn2, b.mn2):
   }
};
// 0 - indexed / inclusive - inclusive
template <class node>
struct STB {
   vector<node> st; int n;
   void build(int u, int i, int j, vector<node> &arr) {
      if (i == j) {
```

```
st[u] = arr[i]:
       return:
   }
   int m = (i + i) / 2, l = u * 2 + 1, r = u * 2 + 2;
   build(1, i, m, arr), build(r, m + 1, j, arr);
   st[u] = node(st[l], st[r]):
void push_add(int u, int i, int j, ll v) {
   st[u].s += (j - i + 1) * v;
   st[u].mx1 += v, st[u].mn1 += v, st[u].lz += v;
   if (st[u].mx2 != LLONG MIN) st[u].mx2 += v:
   if (st[u].mn2 != LLONG MAX) st[u].mn2 += v:
void push_max(int u, ll v, bool l) { // for min op
   if (v >= st[u].mx1) return:
   st[u].s -= st[u].mx1 * st[u].mxc;
   st[u].mx1 = v:
   st[u].s += st[u].mx1 * st[u].mxc;
   if (1) st[u].mn1 = st[u].mx1:
   else if (v \le st[u].mn1) st[u].mn1 = v:
   else if (v < st[u].mn2) st[u].mn2 = v;
void push_min(int u, ll v, bool l) { // for max op
   if (v <= st[u].mn1) return:</pre>
   st[u].s -= st[u].mn1 * st[u].mnc;
   st[u].mn1 = v;
   st[u].s += st[u].mn1 * st[u].mnc:
   if (1) st[u].mx1 = st[u].mn1:
   else if (v \ge st[u].mx1) st[u].mx1 = v;
   else if (v > st[u].mx2) st[u].mx2 = v:
void push(int u, int i, int j) {
   if (i == i) return:
   // add
   int m = (i + j) / 2, l = u * 2 + 1, r = u * 2 + 2;
   push add(l, i, m, st[u].lz);
   push_add(r, m + 1, j, st[u].lz);
   st[u].lz = 0;
   // min
   push max(1, st[u].mx1, i == m):
   push max(r, st[u].mx1, m + 1 == i):
   // max
   push min(1, st[u].mn1, i == m):
   push_min(r, st[u].mn1, m + 1 == r);
node query(int a, int b, int u, int i, int j) {
   if (b < i || j < a) return node();</pre>
   if (a <= i && j <= b) return st[u];</pre>
   push(u, i, i):
   int m = (i + j) / 2, l = u * 2 + 1, r = u * 2 + 2;
```

```
return node(query(a, b, 1, i, m), query(a, b, r, m +
        1, i));
void update_add(int a, int b, ll v, int u, int i, int i)
   if (b < i | | i < a) return:
   if (a <= i && j <= b) {</pre>
       push_add(u, i, j, v);
       return:
   push(u, i, j);
   int m = (i + j) / 2, l = u * 2 + 1, r = u * 2 + 2;
   update_add(a, b, v, l, i, m);
   update_add(a, b, v, r, m + 1, j);
   st[u] = node(st[1], st[r]);
void update_min(int a, int b, ll v, int u, int i, int j)
   if (b < i \mid | i < a \mid | v >= st[u].mx1) return:
   if (a <= i && i <= b && v > st[u].mx2) {
       push_max(u, v, i == j);
       return:
   push(u, i, j);
   int m = (i + j) / 2, l = u * 2 + 1, r = u * 2 + 2;
   update_min(a, b, v, l, i, m);
   update_min(a, b, v, r, m + 1, j);
   st[u] = node(st[1], st[r]);
void update max(int a. int b. 11 v. int u. int i. int i)
   if (b < i || j < a || v <= st[u].mn1) return;</pre>
   if (a <= i && j <= b && v < st[u].mn2) {</pre>
       push_min(u, v, i == j);
       return:
   push(u, i, j);
   int m = (i + j) / 2, l = u * 2 + 1, r = u * 2 + 2;
   update_max(a, b, v, l, i, m);
   update_max(a, b, v, r, m + 1, j);
   st[u] = node(st[l], st[r]):
STB(vector<node> &v, int N) : n(N), st(N * 4 + 5) { build
    (0, 0, n - 1, v); 
node query(int a, int b) { return query(a, b, 0, 0, n -
    1); }
void update_add(int a, int b, ll v) { update_add(a, b, v,
     0.0.n - 1):
```

1.9 segment-tree-lazy

```
template <class T>
struct Stl {
   int n; vector<T> a, b;
   Stl(int n = 0) : n(n), a(4 * n, qneut()),
      b(4 * n. uneut()) {}
   T gneut() { return -2e9: }
   T uneut() { return 0: }
   T merge(T x, T y) { return max(x, y); }
   void upd(int v, T x, int 1, int r)
      \{a[v] += x, b[v] += x; \}
   void push(int v. int vl. int vm. int vr) {
      upd(2 * v, b[v], vl, vm);
      upd(2 * v + 1, b[v], vm, vr);
       b[v] = uneut():
   }
   T query(int 1, int r, int v=1, int vl=0, int vr=1e9) {
       vr = min(vr, n):
      if (1 <= v1 && r >= vr) return a[v];
      if (1 >= vr || r <= vl) return gneut();</pre>
      int vm = (vl + vr) / 2:
      push(v. vl. vm. vr):
      return merge(query(1, r, 2 * v, v1, vm),
          querv(1, r, 2 * v + 1, vm, vr)):
   void update(int l. int r. T x. int v = 1. int vl = 0.
          int vr = 1e9) {
      vr = min(vr, n):
      if (1 >= vr || r <= vl || r <= 1) return;</pre>
      if (1 <= v1 && r >= vr) upd(v, x, v1, vr);
          int vm = (vl + vr) / 2;
          push(v, v1, vm, vr);
          update(1, r, x, 2 * v, v1, vm):
          update(1, r, x, 2 * v + 1, vm, vr);
          a[v] = merge(a[2 * v], a[2 * v + 1]);
      }
```

1.10 segment-tree

```
struct St {
    ll neut() { return 0; }
    ll merge(ll x, ll y) { return x + y; }

int n; vector<ll> a;
    St(int n = 0) : n(n), a(2 * n, neut()) {}

ll query(int l, int r) {
    ll x = neut(), y = neut();
    for (l += n, r += n; l < r; l /= 2, r /= 2) {
        if (l & 1) x = merge(x, a[l++]);
        if (r & 1) y = merge(a[--r], y);
    }
    return merge(x, y);
}

void update(int i, ll x) {
    for (a[i += n] = x; i /= 2;)
        a[i] = merge(a[2 * i], a[2 * i + 1]);
}
};</pre>
```

1.11 sparse-table

```
template <class T>
struct Sparse {
   T op(T a, T b) { return max(a, b); }
   vector<vector<T>> st;
   Sparse() {}
   Sparse(vector<T> a) : st{a} {
       int N = st[0].size();
       int npot = N <= 1 ? 1 : 32 - __builtin_clz(N);</pre>
       st.resize(npot):
       repx(i, 1, npot) rep(j, N + 1 - (1 << i))
       st[i].push_back(
          op(st[i-1][j], st[i-1][j+(1 << (i-1))])
      ); // query op
   T query(int 1, int r) { // range must be nonempty!
       int i = 31 - __builtin_clz(r - 1);
       return op(st[i][1], st[i][r - (1 << i)]); // queryop</pre>
```

```
};
```

1.12 treap-implicit

```
mt19937 gen(chrono::high_resolution_clock::now().
     time_since_epoch().count());
// 101 Implicit Treap //
struct Node {
    int p, sz = 0, v, acc, l = -1, r = -1:
    Node() : v(0), acc(0) {}
    Node(int x) : p(gen()), sz(1), v(x), acc(x) {}
    void recalc(const Node &a, const Node &b) {
       sz = a.sz + b.sz + 1; acc = v + a.acc + b.acc;
};
template <class node>
struct Treap {
    vector<node> t;
    int n, r = -1;
    node get(int u) { return u != -1 ? t[u] : node(); }
    void recalc(int u){t[u].recalc(get(t[u].l),get(t[u].r));}
    int merge(int 1, int r) {
       if (min(1, r) == -1) return 1 != -1 ? 1 : r:
       int ans = (t[1].p < t[r].p) ? 1 : r;
       if (ans == 1) t[1].r = merge(t[1].r, r), recalc(1);
       if (ans == r) t[r].l = merge(l, t[r].l), recalc(r);
       return ans:
    pii split(int u, int id) {
       if (u == -1) return {-1, -1};
       int szl = get(t[u].1).sz:
       if (szl >= id) {
           pii ans = split(t[u].1, id);
           t[u].1 = ans.ss;
           recalc(u);
           return {ans.ff, u}:
       pii ans = split(t[u].r, id - szl - 1);
       t[u].r = ans.ff:
       recalc(u):
       return {u, ans.ss};
    Treap(vi &v) : n(sz(v)) {
       for (int i=0: i<n: i++) t.eb(v[i]), r = merge(r, i):
};
// Complete Implicit Treap with Lazy propagation //
struct Node {
```

```
int p, sz = 0, v, acc, l = -1, r = -1, par = -1, lzv=0;
   bool lz = false, f = false:
   Node() : v(0), acc(0) {}
   Node(int x) : p(gen()), sz(1), v(x), acc(x) {}
   void recalc(const Node &a, const Node &b) {
      sz = a.sz + b.sz + 1; acc = v + a.acc + b.acc;
   void upd_lazy(int x) { lz = 1, lzv += x; }
   void lazy() {v+=lzv, acc += sz * lzv, lz = 0, lzv = 0; }
   void flip() { swap(1, r), f = 0; }
template <class node> struct Treap {
   vector<node> t;
   int n. r = -1:
   node get(int u) { return u != -1 ? t[u] : node(); }
   void recalc(int u) {
      int 1 = t[u].1. r = t[u].r:
      push(1), push(r), flip(1), flip(r);
      t[u].recalc(get(1), get(r));
   }
   void push(int u) {
       if (u == -1 || !t[u].lz) return;
      int 1 = t[u].1, r = t[u].r;
      if (1 != -1) t[1].upd_lazy(t[u].lzv);
      if (r != -1) t[r].upd_lazy(t[u].lzv);
      t[u].lazy();
   }
   void flip(int u) {
      if (u == -1 || !t[u].f) return;
      int 1 = t[u].1, r = t[u].r:
      if (1 != -1) t[1].f ^= 1;
      if (r != -1) t[r].f ^= 1;
      t[u].flip():
   int merge(int 1, int r) { // (*) = only if parent needed
      if (min(1, r) == -1) return 1 != -1 ? 1 : r:
      push(1), push(r), flip(1), flip(r);
      int ans = (t[1].p < t[r].p) ? 1 : r;
      if (ans == 1) t[1].r = merge(t[1].r, r), recalc(1);
      if (ans == r) t[r].l = merge(l, t[r].l), recalc(r);
      if (t[ans].l != -1) t[t[ans].l].par = ans: // (*)
      if (t[ans].r != -1) t[t[ans].r].par = ans; // (*)
       return ans:
   pii split(int u, int id) {// (*) = only if parent needed
      if (u == -1) return {-1, -1}:
      push(u);
      flip(u);
      int szl = get(t[u].1).sz;
      if (szl >= id) {
```

```
pii ans = split(t[u].1, id):
          if (ans.ss != -1) t[ans.ss].par = u; // (*)
          if (ans.ff != -1) t[ans.ff].par = -1; // (*)
          t[u].1 = ans.ss:
          recalc(u);
          return {ans.ff. u}:
       pii ans = split(t[u].r, id - szl - 1);
       if (ans.ff != -1) t[ans.ff].par = u; // (*)
       if (ans.ss != -1) t[ans.ss].par = -1; // (*)
       t[u].r = ans.ff:
       recalc(u):
       return {u, ans.ss};
   int update(int u, int 1, int r, int v) {
       pii a = split(u, 1), b = split(a.ss, r - 1 + 1);
       t[b.ff].upd lazv(v):
       return merge(a.ff, merge(b.ff, b.ss));
   void print(int u) {
       if (u==-1) return; push(u), flip(u); print(t[u].1);
       cout << t[u].v << ' '; print(t[u].r);</pre>
   Treap(vi &v) : n(sz(v)) {
       for (int i=0; i<n; i++) t.eb(v[i]), r = merge(r, i);</pre>
};
```

1.13 treap

```
typedef struct item *pitem:
struct item {
   int pr,key,cnt; pitem l,r;
   item(int kev):kev(kev).pr(rand()).cnt(1).1(0).r(0) {}
int cnt(pitem t){return t?t->cnt:0;}
void upd_cnt(pitem t){if(t)t->cnt=cnt(t->1)+cnt(t->r)+1;}
void split(pitem t, int key, pitem& l, pitem& r){ // l: <=</pre>
    kev. r: > kev
   if(!t)l=r=0:
   else if(key<t->key)split(t->1,key,1,t->1),r=t;
   else split(t->r,key,t->r,r),l=t;
   upd_cnt(t);
void insert(pitem& t. pitem it){
   if(!t)t=it:
   else if(it->pr>t->pr)split(t,it->key,it->l,it->r),t=it;
   else insert(it->kev<t->kev?t->l:t->r.it):
   upd_cnt(t);
```

```
}
void merge(pitem& t, pitem 1, pitem r){
    if(!1||!r)t=1?1:r;
    else if(1->pr>r->pr)merge(1->r,1->r,r),t=1;
    else merge(r->1,1,r->1),t=r;
    upd_cnt(t);
}
void erase(pitem& t, int key){
    if(t->key==key)merge(t,t->1,t->r);
    else erase(key<t->key?t->1:t->r,key);
    upd_cnt(t);
}
```

$2 \quad geo2d$

2.1 circle

```
struct C {
   Po; Tr;
   // circle-line intersection, assuming it exists
   // points are sorted along the direction of the line
   pair<P, P> line_inter(L 1) const {
      P c = 1.closest_to(o); T c2 = (c - o).magsq();
      P = 1.d * sqrt(max(r*r - c2, T()) / 1.d.magsq());
      return \{c - e, c + e\};
   // check the type of line-circle collision
   // <0: 2 inters, =0: 1 inter, >0: 0 inters
   T line collide(L 1) const {
      T c2 = (1.closest_to(o) - o).magsq();
      return c2 - r * r:
   // calculates the two intersections between two circles
   // the circles must intersect in one or two points!
   pair<P, P> inter(C h) const {
      P d = h.o - o:
      T c = (r * r - h.r * h.r) / d.magsq();
      return h.line_inter({(1 + c) / 2 * d, d.rot()});
   // check if the given circles intersect
   bool collide(C h) const {
      return (h.o - o).magsq() \le (h.r + r) * (h.r + r);
```

```
// get one of the two tangents that go through the point
   // the point must not be inside the circle
   // a = -1: cw (relative to the circle) tangent
   // a = 1: ccw (relative to the circle) tangent
   P point_tangent(P p, T a) const {
      T c = r * r / p.magsq():
       return o + c*(p-o) - a*sqrt(c*(1-c))*(p-o).rot();
   }
   // get one of the 4 tangents between the two circles
   // a = 1: exterior tangents
   // a = -1: interior tangents (requires no area overlap)
   // b = 1: ccw tangent
   // b = -1: cw tangent
   // the line origin is on this circumference, and the
   // direction is a unit vector towards the other circle
   L tangent(C c, T a, T b) const {
      T dr = a * r - c.r;
      P d = c.o - o;
      P n = (d*dr+b*d.rot()*sqrt(d.magsq()-dr*dr)).unit();
       return {o + n * r, -b * n.rot()};
   // circumcircle of a **non-degenerate** triangle
   static C thru_points(P a, P b, P c) {
      b = b - a, c = c - a;
       P p = (b*c.magsq() - c*b.magsq()).rot() / (b%c*2);
       return {a + p, p.mag()};
   // find the two circles that go through the given point,
   // are tangent to the given line and have radius 'r'
   // the point-line distance must be at most 'r'!
   // the circles are sorted in the direction of the line
   static pair<C, C> thru_point_line_r(P a, L t, T r) {
       P d = t.d.rot().unit():
       if (d * (a - t.o) < 0) d = -d;
       auto p = C(a, r).line_inter(\{t.o + d * r, t.d\});
       return {{p.first, r}, {p.second, r}};
   // find the two circles that go through the given points
   // and have radius 'r'
   // circles sorted by angle from the first point
   // the points must be at most at distance 'r'!
   static pair<C, C> thru_points_r(P a, P b, T r) {
       auto p = C(a, r).line_inter(\{(a+b)/2, (b-a).rot()\});
       return {{p.first, r}, {p.second, r}};
   }
};
```

2.2 closest-points

```
// sort by x
11 closest(vector<ii> &p) {
   int n = SZ(p);
   set<ii>> s:
   ll best = 1e18;
   int j = 0;
   fore(i, 0, n) {
      11 d = ceil(sqrt(best));
       while(p[i].fst - p[j].fst >= best)
          s.erase({p[j].snd, p[j].fst}), j++;
       auto it1=s.lower_bound({p[i].snd-d,p[i].fst});
       auto it2=s.upper_bound({p[i].snd+d,p[i].fst});
       for(auto it = it1; it != it2; ++it) {
          ll dx = p[i].fst - it->snd;
          11 dy = p[i].snd - it->fst;
          best = min(best, dx * dx + dy * dy);
       s.insert({p[i].snd, p[i].fst});
   return best:
```

2.3 convex-hull

```
// ccw order, excludes collinear points by default
vector<P> chull(vector<P> p) {
   if (p.size() < 3) return p;</pre>
   vector\langle P \rangle r; int m, k = 0;
   sort(p.begin(), p.end(), [](P a, P b) {
       return a.x != b.x ? a.x < b.x : a.y < b.y; });</pre>
   for (P q : p) { // lower hull
       while (k \ge 2 \&\& r[k - 1].left(r[k - 2], q) \ge 0)
           r.pop_back(), k--; // >= to > to add collinears
       r.push back(q), k++:
   if (k == (int)p.size()) return r;
   r.pop_back(), k--, m = k;
   for (int i = p.size() - 1; i >= 0; --i) { // upper hull
       while (k \ge m+2 \&\& r[k-1].left(r[k-2], p[i]) \ge 0)
           r.pop_back(), k--; // >= to > to add collinears
       r.push_back(p[i]), k++;
   r.pop_back(); return r;
```

2.4 delaunay

```
typedef __int128_t lll; // if on a 64-bit platform
struct Q {
    Q *rot, *o; P p = {INF, INF}; bool mark;
    P &F() { return r()->p; }
    Q *&r() { return rot->rot; }
    Q *prev() { return rot->o->rot; }
    Q *next() { return r()->prev(); }
}:
T cross(P a, P b, P c) { return (b - a) % (c - a); }
bool circ(P p, P a, P b, P c) { // is p in the circumcircle?
    111 p2 = p.magsq(), A = a.magsq() - p2,
       B = b.magsq() - p2, C = c.magsq() - p2;
    return cross(p, a, b) * C + cross(p, b, c) * A + cross(p, b, c)
          c. a) * B > 0:
Q *makeEdge(Q *&H, P orig, P dest) {
    Q *r = H ? H : new Q{new Q{new Q{new Q{0}}}};
    H = r -> 0; r -> r() -> r() = r;
    repx(i, 0, 4) r = r->rot, r->p = {INF, INF},
       r \rightarrow 0 = i & 1 ? r : r \rightarrow r();
    r\rightarrow p = orig; r\rightarrow F() = dest;
    return r;
}
void splice(Q *a, Q *b) {
    swap(a->o->rot->o, b->o->rot->o); swap(a->o, b->o);
0 *connect(Q *&H, Q *a, Q *b) {
    Q *q = makeEdge(H, a->F(), b->p);
    splice(q, a->next()); splice(q->r(), b); return q;
}
pair<Q *, Q *> rec(Q *&H, const vector<P> &s) {
    if (s.size() <= 3) {</pre>
       Q *a = makeEdge(H, s[0], s[1]), *b = makeEdge(H, s[0], s[1])
             [1]. s.back()):
       if (s.size() == 2) return \{a, a->r()\}; splice(a->r(),
        auto side = cross(s[0], s[1], s[2]):
       Q *c = side ? connect(H, b, a) : 0;
       return {side < 0 ? c->r() : a, side < 0 ? c : b->r()
             }:
```

```
#define J(e) e \rightarrow F(), e \rightarrow p
#define valid(e) (cross(e->F(), J(base)) > 0)
   Q *A, *B, *ra, *rb; int half = s.size() / 2;
    tie(ra, A) = rec(H, \{s.begin(), s.end() - half\});
    tie(B, rb) = rec(H, {s.begin() + s.size() - half, s.end()
   while ((cross(B->p, J(A)) < 0 \&\& (A = A->next())) | |
          (cross(A->p, J(B)) > 0 \&\& (B = B->r()->o)));
   Q *base = connect(H, B->r(), A);
   if (A->p == ra->p) ra = base->r();
   if (B->p == rb->p) rb = base:
#define DEL(e, init, dir) Q *e = init->dir; \
    if (valid(e)) while (circ(e->dir->F(), J(base), e->F()))
           Q *t = e->dir; splice(e, e->prev()); \
           splice(e->r(), e->r()->prev()); e->o = H; H = e;
                e = t: \
   for (;;) {
       DEL(LC, base->r(), o); DEL(RC, base, prev());
       if (!valid(LC) && !valid(RC)) break;
       if (!valid(LC) || (valid(RC) && circ(J(RC), J(LC))))
            base = connect(H, RC, base->r());
       else base = connect(H, base->r(), LC->r());
   return {ra, rb}:
#undef J
#undef valid
#undef DEL
// there must be no duplicate points
// returns no triangles in the case of all collinear points
// produces counter-clockwise triangles ordered in triples
// maximizes the minimum angle across all triangulations
// the euclidean mst is a subset of these edges
// O(N log N)
vector<P> triangulate(vector<P> pts) {
    sort(pts.begin(), pts.end(), [](P a, P b) {
       return make_pair(a.x, a.y) < make_pair(b.x, b.y);</pre>
    assert(unique(pts.begin(), pts.end()) == pts.end());
   if (pts.size() < 2) return {};</pre>
   Q *H = 0; Q *e = rec(H, pts).first;
   vector<Q *> q = {e}; int qi = 0;
    while (cross(e->o->F(), e->F(), e->p) < 0) e = e->o;
#define ADD
```

2.5 halfplane-intersect

```
// obtain the convex polygon that results from intersecting
    the given list
// of halfplanes, represented as lines that allow their left
// assumes the halfplane intersection is bounded
vector<P> halfplane intersect(vector<L> &H) {
   L bb(P(-INF, -INF), P(INF, 0));
   rep(k, 4) H.push_back(bb), bb.o = bb.o.rot(), bb.d = bb.d
        .rot():
   sort(begin(H), end(H), [](L a, L b) { return a.d.angcmp(b
        .d) < 0; \});
   deque<L> q; int n = 0;
   rep(i, H.size()) {
       while (n \ge 2 \&\& H[i].side(q[n - 1].intersection(q[n
           -21)) > 0)
           q.pop_back(), n--;
       while (n >= 2 && H[i].side(q[0].intersection(q[1])) >
           q.pop_front(), n--;
       if (n > 0 && H[i].parallel(q[n - 1])) {
          if (H[i].d * q[n - 1].d < 0) return {};</pre>
          if (H[i].side(q[n-1].o) > 0) q.pop_back(), n--;
           else continue:
       7
       q.push_back(H[i]), n++;
   while (n \ge 3 \&\& q[0].side(q[n - 1].intersection(q[n -
        21)) > 0)
       q.pop_back(), n--;
   while (n \ge 3 \&\& q[n - 1].side(q[0].intersection(q[1])) >
       q.pop_front(), n--;
```

```
if (n < 3) return {};

vector<P> ps(n);
rep(i, n) ps[i] = q[i].intersection(q[(i + 1) % n]);
return ps;
```

2.6 line

```
// a segment or an infinite line
// does not handle point segments correctly!
struct L {
   P o, d;
   static L from eq(P ab, T c) {
      return L{ab.rot(), ab * -c / ab.magsq()};
   pair<P, T> line_eq() { return {-d.rot(), d.rot() * o}; }
   // on which side of the line is the point
   // negative: left, positive: right
   T side(P r) const { return (r - o) % d; }
   // returns the intersection coefficient
   // in the range [0, d % r.d]
   // if d % r.d is zero, the lines are parallel
   T inter(L r) const { return (r.o - o) % r.d; }
   // get the single intersection point
   // lines must not be parallel
   P intersection(L r) const {return o+d*inter(r)/(d%r.d):}
   // check if lines are parallel
   bool parallel(L r) const { return abs(d % r.d) <= EPS; }</pre>
   // check if segments intersect
   bool seg collide(L r) const {
      Tz = d \% r.d;
      if (abs(z) \le EPS) {
          if (abs(side(r.o)) > EPS) return false;
          T s = (r.o - o) * d. e = s + r.d * d:
          if (s > e) swap(s, e):
          return s <= d * d + EPS && e >= -EPS;
      T s = inter(r), t = -r.inter(*this):
      if (z < 0) s = -s, t = -t, z = -z:
       return s>=-EPS && s<=z+EPS && t>=-EPS && t<=z+EPS:
```

```
// full segment intersection
// makes a point segment if the intersection is a point
// however it does not handle point segments as input!
bool seg inter(L r, L *out) const {
   Tz = d \% r.d:
   if (abs(z) <= EPS) {
       if (abs(side(r.o)) > EPS) return false;
       if (r.d * d < 0) r = \{r.o + r.d, -r.d\};
       P s = o * d < r.o * d ? r.o : o:
       P = (o+d)*d < (r.o+r.d)*d ? o+d : r.o+r.d;
       if (s * d > e * d) return false:
       return *out = {s. e - s}, true:
   T s = inter(r), t = -r.inter(*this);
   if (z < 0) s = -s. t = -t. z = -z:
   if (s>=-EPS && s<=z+EPS && t>=-EPS && t<=z+EPS)
       return *out = {o + d * s / z, {0, 0}}, true:
   return false;
// check if the given point is on the segment
bool point_on_seg(P r) const {
   if (abs(side(r)) > EPS) return false;
   if ((r - o) * d < -EPS) return false;</pre>
   if ((r - o - d) * d > EPS) return false:
   return true;
// point in this line that is closest to a given point
P closest to(P r) const {
   P dr = d.rot(); return r + dr*((o-r)*dr)/d.magsq();
```

2.7 minkowski

```
void reorder_polygon(vector<P> &ps) {
   int pos = 0;
   repx(i, 1, (int)ps.size()) {
      if (ps[i].y < ps[pos].y || (ps[i].y == ps[pos].y &&
            ps[i].x < ps[pos].x))
      pos = i;
   }
   rotate(ps.begin(), ps.begin() + pos, ps.end());
}
vector<P> minkowski(vector<P> ps, vector<P> qs) {
   // the first vertex must be the lowest
   reorder_polygon(ps); reorder_polygon(qs);
   ps.push_back(ps[0]); ps.push_back(ps[1]);
```

```
qs.push_back(qs[0]); qs.push_back(qs[1]);
vector<P> result; int i = 0, j = 0;
while (i < ps.size() - 2 || j < qs.size() - 2) {
    result.push_back(ps[i] + qs[j]);
    auto z = (ps[i + 1] - ps[i]) % (qs[j + 1] - qs[j]);
    if (z >= 0 && i < ps.size() - 2) ++i;
    if (z <= 0 && j < qs.size() - 2) ++j;
}
return result;
}</pre>
```

2.8 point

```
struct P {
   T x. v:
   P(T x, T y) : x(x), y(y) {}
   P() : P(0, 0) \{ \}
   friend ostream &operator<<(ostream &s, const P &r) {</pre>
      return s << r.x << " " << r.y;
   friend istream &operator>>(istream &s, P &r) { return s
        >> r.x >> r.v: }
   P operator+(P r) const { return \{x + r.x, y + r.y\}; }
   P operator-(P r) const { return {x - r.x, y - r.y}; }
   P operator*(T r) const { return {x * r, y * r}; }
   P operator/(T r) const { return {x / r, v / r}; }
   P operator-() const { return {-x, -y}; }
   friend P operator*(T 1, P r) { return {1 * r.x, 1 * r.y};
   P rot() const { return {-v, x}; }
   T operator*(P r) const { return x * r.x + y * r.y; }
   T operator%(P r) const { return rot() * r; }
   T left(P a, P b) { return (b - a) % (*this - a); }
   T magsq() const { return x * x + y * y; }
   T mag() const { return sqrt(magsq()); }
   P unit() const { return *this / mag(): }
   bool half() const { return abs(y) <= EPS && x < -EPS || y
         < -EPS: }
   T angcmp(P r) const { // like strcmp(this, r)
      int h = (int)half() - r.half();
       return h ? h : r % *this;
   T angcmp_rel(P a, P b) { // like strcmp(a, b)
      Pz = *this:
```

2.9 polygon

```
// get TWICE the area of a simple polygon in ccw order
T area2(const vector<P> &p) {
    int n = p.size(); T a = 0;
    rep(i, n) a += (p[i] - p[0]) \% (p[(i + 1) \% n] - p[i]);
    return a:
}
// checks whether a point is inside a ccw simple polygon
// returns 1 if inside, 0 if on border, -1 if outside
int in_poly(const vector<P> &p, P q) {
    int w = 0:
    rep(i, p.size()) {
       P = p[i], b = p[(i + 1) \% p.size()];
       T k = (b - a) \% (q - a);
       T u = a.y - q.y, v = b.y - q.y;
       if (k > 0 && u < 0 && v >= 0) w++:
       if (k < 0 && v < 0 && u >= 0) w--:
       if (k == 0 && (q - a) * (q - b) <= 0) return 0;
    return w ? 1 : -1;
// check if point in ccw convex polygon, O(log n)
// + if inside, 0 if on border, - if outside
T in_convex(const vector<P> &p, P q) {
    int 1 = 1, h = p.size() - 2; assert(p.size() >= 3);
    while (1 != h) { // collinear points are unsupported!
       int m = (1 + h + 1) / 2;
       if (q.left(p[0], p[m]) >= 0) 1 = m;
       else h = m - 1:
   T in = min(q.left(p[0], p[1]), q.left(p.back(), p[0]));
    return min(in, q.left(p[1], p[1 + 1]));
```

```
int extremal(const vector<P> &p. P d) {
         int n = p.size(), l = 0, r = n - 1; assert(n);
          P = 0 = (p[n - 1] - p[0]).rot();
          while (1 < r) { // polygon must be convex</pre>
                   int m = (1 + r + 1) / 2:
                   P = (p[(m + n - 1) \% n] - p[m]).rot();
                   if (e0.angcmp_rel(d, e) < 0) r = m - 1;
                   else 1 = m:
          }
          return 1:
 // square dist of most distant points of a ccw convex
// polygon with NO COLLINEAR POINTS
T callipers(const vector<P> &p) {
         int n = p.size():
         T r = 0;
          for (int i = 0, i = n < 2 ? 0 : 1; <math>i < j; i++) {
                   for (:: i = (i + 1) \% n) {
                            r = max(r, (p[i] - p[j]).magsq());
                             if ((p[(i + 1) % n] - p[i]) % (p[(j + 1) % n] - p
                                          [i]) <= EPS) break:
          return r;
 P centroid(const vector<P> &p) { // (barycenter)
         P r(0, 0): T t = 0: int n = p.size():
          rep(i, n) {
                   r += (p[i] + p[(i+1)\%n]) * (p[i] \% p[(i+1)\%n]);
                   t += p[i] \% p[(i+1)\%n];
         }
          return r / t / 3:
 // classify collision of a ray inside a ccw polygon vertex.
 // ray is (o, d), vertex is b, previous vertex is a, next is
               С.
 pair <bool, bool > inner collide(P o, P d, P a, P b, P c) {
         T p = (a - o) \% d; // side of previous
         T n = (c - o) \% d:
                                                                // side of next
          T v = (c - b) \% (b - a); // is vertex convex?
          return \{v > 0 ? n < 0 | | (n == 0 \&\& p < 0) : p > 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | | n < 0 | 
                            v > 0 ? p > 0 || (p == 0 \&\& n > 0) : p > 0 || n <
                                            0}:
```

2.10 sweep

```
#include "point.cpp"
// iterate over all pairs of points
// 'op' is called with all ordered pairs of different
    indices '(i, i)'
// additionally, the 'ps' vector is kept sorted by signed
    distance
// to the line formed by 'i' and 'i'
// for example, if the vector from 'i' to 'j' is pointing
// the 'ps' vector is sorted from smallest 'v' to largest 'v
// note that, because the 'ps' vector is sorted by signed
    distance.
// 'i' is always equal to 'i + 1'
// this means that the amount of points to the left of the
    line is always 'N - i'
template <class OP>
void all pair points(vector<P> &ps. OP op) {
   int N = ps.size();
   sort(ps.begin(), ps.end(), [](P a, P b) {
      return make_pair(a.y, a.x) < make_pair(b.y, b.x);</pre>
   vector<pair<int, int>> ss;
   rep(i, N) rep(j, N) if (i != j) ss.push_back({i, j});
   stable_sort(ss.begin(), ss.end(), [&](auto a, auto b) {
      return (ps[a.second] - ps[a.first]).angle_lt(ps[b.
           second] - ps[b.first]);
   vector<int> p(N); rep(i, N) p[i] = i;
   for (auto [i, j] : ss)
      { op(p[i], p[j]); swap(ps[p[i]], ps[p[j]]); swap(p[i
           ], p[i]); }
```

2.11 theorems

```
// Pick's theorem
// Simple polygon with integer vertices:
// A = I + B / 2 - 1
// A: Area of the polygon
// I: Integer points strictly inside the polygon
-// B: Integer points on the boundary of the polygon
```

3 graph

3.1 artic-bridge-biconn

```
vector<int> g[MAXN];int n;
struct edge {int u,v,comp;bool bridge;};
vector<edge> e;
void add_edge(int u, int v){
    g[u].pb(e.size());g[v].pb(e.size());
    e.pb((edge){u.v.-1.false}):
}
int D[MAXN].B[MAXN].T:
int nbc; // number of biconnected components
int art[MAXN]; // articulation point iff !=0
stack<int> st: // only for biconnected
void dfs(int u,int pe){
    B[u]=D[u]=T++;
    for(int ne:g[u])if(ne!=pe){
       int v=e[ne].u^e[ne].v^u;
       if(D[v]<0){
           st.push(ne);dfs(v,ne);
           if(B[v]>D[u])e[ne].bridge = true; // bridge
           if(B[v]>=D[u]){
               art[u]++; // articulation
               int last: // start biconnected
               do{last=st.top();st.pop();e[last].comp=nbc;}
               while(last!=ne);
               nbc++: // end biconnected
           B[u]=min(B[u],B[v]);
       else if(D[v]<D[u])st.push(ne),B[u]=min(B[u],D[v]);</pre>
   }
}
void doit(){
    memset(D,-1.sizeof(D)):memset(art,0.sizeof(art));
    nbc=T=0; fore(i,0,n)if(D[i]<0)dfs(i,-1),art[i]--;</pre>
}
```

3.2 bellman-ford

```
struct Edge { int u, v; ll w; };

// find distance from source node to all nodes.
// supports negative edge weights.
// returns true if a negative cycle is detected.
//
// time: O(V E)
```

3.3 blossom

```
vector<int> g[MAXN];int n,m,mt[MAXN],qh,qt,q[MAXN],ft[MAXN],
bs[MAXN]:bool ing[MAXN].inb[MAXN].inp[MAXN]:int lca(int root
,int x,int y){memset(inp,0,sizeof(inp));while(1){inp[x=bs[x]
]=true;if(x==root)break;x=ft[mt[x]];}while(1){if(inp[y=bs[y])}
])return y;else y=ft[mt[y]];}}void mark(int z,int x){while(
bs[x]!=z){int y=mt[x];inb[bs[x]]=inb[bs[y]]=true;x=ft[y];if(
bs[x]!=z)ft[x]=y;}}void contr(int s,int x,int y){int z=lca(s
,x,y);memset(inb,0,sizeof(inb));mark(z,x);mark(z,y);if(bs[x]
!=z)ft[x]=y;if(bs[y]!=z)ft[y]=x;rep(x,n)if(inb[bs[x]]){bs[x]}
=z;if(!inq[x])inq[q[++qt]=x]=true;}}int findp(int s){memset(
ing.0.sizeof(ing)):memset(ft.-1.sizeof(ft)):rep(i.n)bs[i]=i;
ing[g[qh=qt=0]=s]=true; while(qh<=qt){int x=q[qh++]; for(int y</pre>
:g[x])if(bs[x]!=bs[y]&&mt[x]!=y){if(y==s||mt[y]>=0&&ft[mt[y])}
]>=0)contr(s,x,y);else if(ft[y]<0){ft[y]=x;if(mt[y]<0)return</pre>
v;else if(!ing[mt[v]])ing[g[++qt]=mt[v]]=true;}}}return -1;}
int aug(int s.int t){int x=t.v.z:while(x>=0){v=ft[x]:z=mt[v]
;mt[y]=x;mt[x]=y;x=z;}return t>=0;}int edmonds(){int r=0;
memset(mt,-1,sizeof(mt));rep(x,n)if(mt[x]<0)r+=aug(x,findp(x)</pre>
)):return r:}
```

3.4 chu-liu-minimum-spanning-arborescence

```
//O(n*m) minimum spanning tree in directed graph
//returns -1 if not possible
//included i-th edge if take[i]!=0
typedef int tw; tw INF=111<<30;
struct edge{int u,v,id;tw len;};
struct ChuLiu{
  int n; vector<edge> e;
```

```
vector<int> inc.dec.take.pre.num.id.vis:
vector<tw> inw:
void add_edge(int x, int y, tw w){
   inc.pb(0); dec.pb(0); take.pb(0);
   e.pb(\{x,y,SZ(e),w\});
ChuLiu(int n):n(n),pre(n),num(n),id(n),vis(n),inw(n){}
tw doit(int root){
    auto e2=e:
   tw ans=0; int eg=SZ(e)-1,pos=SZ(e)-1;
   while(1){
       fore(i.0.n) inw[i]=INF.id[i]=vis[i]=-1:
       for(auto ed:e2) if(ed.len<inw[ed.v]){</pre>
           inw[ed.v]=ed.len; pre[ed.v]=ed.u;
           num[ed.v]=ed.id:
       inw[root]=0:
       fore(i,0,n) if(inw[i]==INF) return -1;
       int tot=-1:
       fore(i,0,n){
           ans+=inw[i];
           if(i!=root)take[num[i]]++:
           int j=i;
           while(vis[j]!=i&&j!=root&&id[j]<0)vis[j]=i,j=</pre>
           if(j!=root&&id[j]<0){</pre>
              id[i]=++tot:
               for(int k=pre[j];k!=j;k=pre[k]) id[k]=tot;
       if(tot<0)break:
       fore(i,0,n) if(id[i]<0)id[i]=++tot;</pre>
       n=tot+1: int i=0:
       fore(i,0,SZ(e2)){
           int v=e2[i].v:
           e2[i].v=id[e2[i].v]:
           e2[i].u=id[e2[i].u];
           if(e2[j].v!=e2[j].u){
               e2[j].len=e2[i].len-inw[v];
               inc.pb(e2[i].id);
               dec.pb(num[v]):
               take.pb(0);
               e2[i++].id=++pos:
       e2.resize(i):
       root=id[root];
    while(pos>eg){
```

3.5 dinic

```
// time: 0(E V^2)
       O(E V^{(2/3)}) / O(E \operatorname{sqrt}(E)) unit capacities
       O(E sqrt(V)) (hopcroft-karp) unit networks
//unit network: c in {0,1} & forall v, indeg<=1 or outdeg<=1
//min-cut: nodes reachable from s in final residual graph
struct Dinic {
   struct Edge { int u, v; ll c, f = 0; };
   int N, s, t; vector<vector<int>> G;
   vector<Edge> E; vector<int> lvl, ptr;
   Dinic() {}
   Dinic(int N, int s, int t): N(N), s(s), t(t), G(N) {}
   void add_edge(int u, int v, ll c) {
       G[u].push_back(E.size()); E.push_back({u, v, c});
       G[v].push_back(E.size()); E.push_back({v, u, 0});
   11 push(int u, 11 p) {
       if (u == t || p <= 0) return p;</pre>
       while (ptr[u] < G[u].size()) {</pre>
           int ei = G[u][ptr[u]++];
           Edge &e = E[ei]:
           if (lvl[e.v] != lvl[u] + 1) continue;
           ll a = push(e.v. min(e.c - e.f. p)):
           if (a <= 0) continue;</pre>
           e.f += a, E[ei ^ 1].f -= a; return a;
       return 0;
   11 maxflow() {
       11 f = 0:
       while (true) {
           lvl.assign(N, -1); queue<int> q;
           lvl[s] = 0: a.push(s):
           while (!q.empty()) {
               int u = q.front(); q.pop();
              for (int ei : G[u]) {
                  Edge &e = E[ei];
```

3.6 dominator-tree

```
//idom[i]=parent of i in dominator tree with root=rt. or -1
    if not exists
int n,rnk[MAXN],pre[MAXN],anc[MAXN],idom[MAXN],semi[MAXN],
    low[MAXN]:
vector<int> g[MAXN],rev[MAXN],dom[MAXN],ord;
void dfspre(int pos){
   rnk[pos]=SZ(ord); ord.pb(pos);
   for(auto x:g[pos]){
      rev[x].pb(pos);
      if(rnk[x]==n) pre[x]=pos,dfspre(x);
int eval(int v){
   if (anc[v]<n&&anc[anc[v]]<n){
       int x=eval(anc[v]):
       if(rnk[semi[low[v]]]>rnk[semi[x]]) low[v]=x;
       anc[v]=anc[anc[v]]:
   }
   return low[v]:
void dominators(int rt){
   fore(i.0.n){
       dom[i].clear(); rev[i].clear();
       rnk[i]=pre[i]=anc[i]=idom[i]=n;
       semi[i]=low[i]=i:
   ord.clear(); dfspre(rt);
   for(int i=SZ(ord)-1:i:i--){
      int w=ord[i];
```

```
for(int v:rev[w]){
    int u=eval(v);
    if(rnk[semi[w]]>rnk[semi[u]])semi[w]=semi[u];
}
dom[semi[w]].pb(w); anc[w]=pre[w];
for(int v:dom[pre[w]]){
    int u=eval(v);
    idom[v]=(rnk[pre[w]]>rnk[semi[u]]?u:pre[w]);
}
dom[pre[w]].clear();
}
for(int w:ord) if(w!=rt&&idom[w]!=semi[w]) idom[w]=idom[idom[w]];
fore(i,0,n) if(idom[i]==n)idom[i]=-1;
```

3.7 eulerian-path

```
// Directed version(uncomment commented code for undirected)
struct edge {
   int y; // list<edge>::iterator rev;
   edge(int v):v(v){}
};
list<edge> g[MAXN];
void add_edge(int a, int b){
   g[a].push_front(edge(b));//auto ia=g[a].begin();
    g[b].push_front(edge(a));auto ib=g[b].begin();
11
    ia->rev=ib:ib->rev=ia:
vector<int> p;
void go(int x){
   while(g[x].size()){
       int y=g[x].front().y;//g[y].erase(g[x].front().rev);
       g[x].pop_front(); go(y);
   p.push_back(x);
vector<int> get_path(int x){ // get a path that begins in x
// check that a path exists from x before calling get_path!
   p.clear();go(x);reverse(p.begin(),p.end());
   return p;
```

3.8 floyd-warshall

// calculate distances between every pair of nodes in $O(V^3)$ time.

```
// works with negative edges, but not negative cycles.
void floyd(const vector<vector<pair<11, int>>> &G, vector<
        vector<11>> &D) {
    int N = G.size();
    D.assign(N, vector<11>(N, INF));
    rep(u, N) D[u][u] = 0;
    rep(u, N) for (auto [w, v] : G[u]) D[u][v] = w;
    rep(k, N) rep(u, N) rep(v, N)
        D[u][v] = min(D[u][v], D[u][k] + D[k][v]);
}
```

3.9 heavy-light

```
struct Hld {
   vector<int> P, H, D, pos, top;
   Hld() {}
   void init(vector<vector<int>> &G) {
      int N = G.size();
      P.resize(N), H.resize(N), D.resize(N), pos.resize(N),
          top.resize(N):
      D[0] = -1, dfs(G, 0); int t = 0;
      rep(i, N) if (H[P[i]] != i) {
          int i = i:
          while (j != -1)
              \{ top[j] = i, pos[j] = t++; j = H[j]; \}
   }
   int dfs(vector<vector<int>> &G, int i) {
      int w = 1, mw = 0;
      D[i] = D[P[i]] + 1, H[i] = -1:
      for (int c : G[i]) {
          if (c == P[i]) continue:
          P[c] = i; int sw = dfs(G, c); w += sw;
          if (sw > mw) H[i] = c, mw = sw;
      }
      return w;
   // visit the log N segments in the path from u to v
   template <class OP>
   void path(int u, int v, OP op) {
      while (top[u] != top[v]) {
          if (D[top[u]] > D[top[v]]) swap(u, v);
          op(pos[top[v]], pos[v] + 1); v = P[top[v]];
      if (D[u] > D[v]) swap(u, v);
       op(pos[u], pos[v] + 1); // value on node
```

```
// op(pos[u]+1, pos[v] + 1): // value on edge
// an alternative to 'path' that considers order.
// calls 'op' with an 'l <= r' inclusive-exclusive range,
     and a
// boolean indicating if the query is forwards or
    backwards.
template <class OP>
void path(int u, int v, OP op) {
   int lu = u. lv = v:
   while (top[lu] != top[lv])
       if (D[top[lu]] > D[top[lv]]) lu = P[top[lu]];
       else lv = P[top[lv]];
   int lca = D[lu] > D[lv] ? lv : lu;
   while (top[u] != top[lca])
       op(pos[top[u]], pos[u] + 1, false), u = P[top[u
   if (u != lca) op(pos[lca] + 1, pos[u] + 1, false);
   vector<int> stk:
   while (top[v] != top[lca])
       stk.push_back(v), v = P[top[v]];
   // op(pos[lca], pos[v] + 1, true); // value on node
   op(pos[lca] + 1, pos[v] + 1, true); // value on edge
   reverse(stk.begin(), stk.end()):
   for (int w : stk) op(pos[top[w]], pos[w] + 1, true);
// commutative segment tree
template <class T, class S>
void update(S &seg, int i, T val) { seg.update(pos[i],
    val): }
// commutative segment tree lazy
template <class T, class S>
void update(S &seg, int u, int v, T val) {
   path(u, v, [&](int 1, int r) { seg.update(1, r, val);
         }):
}
// commutative (lazy) segment tree
template <class T, class S>
T query(S &seg, int u, int v) {
   T ans = 0:
        // neutral element
   path(u, v, [\&](int 1, int r) \{ ans += seg.query(1, r) \}
        ; }); // query op
```

```
return ans;
};
```

3.10 hungarian

```
// find a maximum gain perfect matching in the given
    bipartite complete graph.
// input: gain matrix (G_{xy} = benefit of joining vertex x
    in set X with vertex
// y in set Y).
// output: maximum gain matching in members 'xy[x]' and 'yx[
// runtime: O(N^3)
struct Hungarian {
   int N, qi, root;
   vector<vector<ll>>> gain;
   vector<int> xy, yx, p, q, slackx;
   vector<ll> lx, ly, slack;
   vector<bool> S, T;
   void add(int x, int px) {
       S[x] = true, p[x] = px;
       rep(y, N) if (lx[x] + ly[y] - gain[x][y] < slack[y])
          slack[y] = lx[x] + ly[y] - gain[x][y], slackx[y]
   void augment(int x, int y) {
       while (x != -2) {
          yx[y] = x; swap(xy[x], y); x = p[x];
   }
   void improve() {
       S.assign(N, false), T.assign(N, false), p.assign(N,
            -1):
       qi = 0, q.clear();
       rep(x, N) if (xy[x] == -1) {
          q.push_back(root = x), p[x] = -2, S[x] = true;
          break;
       rep(y, N) slack[y] = lx[root] + ly[y] - gain[root][y
           ], slackx[v] = root;
       while (true) {
          while (qi < q.size()) {</pre>
```

```
int x = a[ai++]:
              rep(y, N) if (lx[x] + ly[y] == gain[x][y] &&!
                   T[y]) {
                  if (yx[y] == -1) return augment(x, y);
                  T[y] = true, q.push_back(yx[y]), add(yx[y
                      1. x):
              }
          }
          11 d = INF;
           rep(y, N) if (!T[y]) d = min(d, slack[y]);
           rep(x, N) if (S[x]) lx[x] -= d:
           rep(y, N) if (T[y]) ly[y] += d;
           rep(y, N) if (!T[y]) slack[y] -= d:
           rep(y, N) if (!T[y] && slack[y] == 0) {
              if (yx[y] == -1) return augment(slackx[y], y);
              T[y] = true:
              if (!S[yx[y]]) q.push_back(yx[y]), add(yx[y],
                   slackx[v]):
          }
       }
   }
   Hungarian(vector<vector<ll>>> g)
       : N(g.size()), gain(g), xy(N, -1), yx(N, -1), lx(N, -
       lv(N). slack(N). slackx(N) {
       rep(x, N) rep(y, N) lx[x] = max(lx[x], ly[y]);
       rep(i, N) improve();
};
```

3.11 kuhn

```
// get a maximum cardinality matching in a bipartite graph.
// input: adjacency lists.
// output: matching (in 'mt' member).
// runtime: O(V E)
struct Kuhn {
    vector<vector<int>> G;
    int N, size;
    vector<bool> seen;
    vector<int> mt;

    bool visit(int i) {
        if (seen[i]) return false;
        seen[i] = true;
        for (int to : G[i])
```

3.12 lca

```
// calculates the lowest common ancestor for any two nodes
    in O(log N) time,
// with O(N log N) preprocessing
struct Lca {
   int N, K, t = 0;
   vector<vector<int>> U:
   vector<int> L, R;
   Lca() {}
   Lca(vector<vector<int>> &G) : N(G.size()), L(N), R(N) {
      K = N \le 1 ? 0 : 32 - \_builtin_clz(N - 1);
      U.resize(K + 1, vector<int>(N));
      visit(G, 0, 0):
       rep(k, K) rep(u, N) U[k + 1][u] = U[k][U[k][u]]:
   }
   void visit(vector<vector<int>> &G, int u, int p) {
      L[u] = t++, U[0][u] = p;
      for (int v : G[u]) if (v != p) visit(G, v, u):
      R[u] = t++;
   bool is_anc(int up, int dn) {
       return L[up] <= L[dn] && R[dn] <= R[up];
   }
   int find(int u, int v) {
      if (is_anc(u, v)) return u;
      if (is_anc(v, u)) return v;
      for (int k = K; k \ge 0;)
          if (is_anc(U[k][u], v)) k--;
```

```
else u = U[k][u];
    return U[0][u];
}
```

3.13 maxflow-mincost

```
// time: O(F V E)
                         F is the maximum flow
       O(V E + F E log V) if bellman-ford is replaced by
    johnson
struct Flow {
   struct Edge {
      int u, v;
      11 c, w, f = 0;
   int N, s, t;
   vector<vector<int>> G;
   vector<Edge> E;
   vector<ll> d, b;
   vector<int> p;
   Flow() {}
   Flow(int N, int s, int t) : N(N), s(s), t(t), G(N) {}
   void add_edge(int u, int v, ll c, ll w) {
       G[u].push_back(E.size());
       E.push_back({u, v, c, w});
      G[v].push_back(E.size());
       E.push_back({v, u, 0, -w});
   // naive distances with bellman-ford: O(V E)
   void calcdists() {
       p.assign(N, -1), d.assign(N, INF), d[s] = 0;
      rep(i, N - 1) rep(ei, E.size()) {
          Edge &e = E[ei]:
          ll n = d[e.u] + e.w;
          if (d[e.u] < INF && e.c - e.f > 0 && n < d[e.v])
               d[e.v] = n, p[e.v] = ei;
   // johnsons potentials: O(E log V)
   void calcdists() {
      if (b.empty()) {
          b.assign(N, 0);
          // code below only necessary if there are
               negative costs
```

```
rep(i, N - 1) rep(ei, E.size()) {
           Edge &e = E[ei]:
           if (e.f < e.c) b[e.v] = min(b[e.v], b[e.u] + e
       }
   p.assign(N, -1), d.assign(N, INF), d[s] = 0;
   priority_queue<pair<11, int>> q;
   q.push({0, s});
   while (!q.empty()) {
       auto [w, u] = q.top();
       g.pop():
       if (d[u] < -w + b[u]) continue;</pre>
       for (int ei : G[u]) {
           auto e = E[ei]:
          ll n = d[u] + e.w;
           if (e.f < e.c && n < d[e.v]) {</pre>
              d[e.v] = n, p[e.v] = ei;
              q.push({b[e.v] - n, e.v});
       }
   }
   b = d;
11 solve() {
   b.clear():
   11 ff = 0:
   while (true) {
       calcdists():
       if (p[t] == -1) break;
       11 f = INF:
       for (int cur = t; p[cur] != -1; cur = E[p[cur]].u
          f = min(f, E[p[cur]].c - E[p[cur]].f);
       for (int cur = t; p[cur] != -1; cur = E[p[cur]].u
           E[p[cur]].f += f, E[p[cur] ^ 1].f -= f;
       ff += f:
   }
   return ff;
```

3.14 parallel-dfs

```
struct Tree {
   int n,z[2];
```

};

```
vector<vector<int>> g:
vector<int> ex,ey,p,w,f,v[2];
Tree(int n):g(n), w(n), f(n){}
void add_edge(int x, int y){
   p.pb(g[x].size());g[x].pb(ex.size());
   ex.pb(x); ey.pb(y);
   p.pb(g[y].size());g[y].pb(ex.size());
   ex.pb(y);ey.pb(x);
bool go(int k){//returns 1 if it finds new node
   int& x=z[k]:
   while(x \ge 0 \& \&
       (w[x] == g[x].size() | |w[x] == g[x].size()-1
       &&(g[x].back()^1)==f[x]))
       x=f[x] >= 0?ex[f[x]]:-1;
   if(x<0)return false;</pre>
   if((g[x][w[x]]^1)==f[x])w[x]++;
   int e=g[x][w[x]],y=ey[e]; f[y]=e;
   w[x]++; w[y]=0; x=y; v[k].pb(x);
   return true:
vector<int> erase_edge(int e){
   e*=2;//erases eth edge, returns smaller comp
   int x=ex[e],y=ey[e]; p[g[x].back()]=p[e];
   g[x][p[e]]=g[x].back(); g[x].pop_back();
   p[g[y].back()]=p[e^1]; g[y][p[e^1]]=g[y].back();
   g[y].pop_back();
   f[x]=f[y]=-1; w[x]=w[y]=0; z[0]=x;z[1]=y;
   v[0]={x};v[1]={y};
   bool d0=true.d1=true:while(d0&&d1)d0=go(0).d1=go(1):
   return v[1-d1];
```

3.15 push-relabel

```
#include "../common.h"

const ll INF = 1e18;

// maximum flow algorithm.
// to run, use 'maxflow()'.

//
// time: O(V^2 sqrt(E)) <= O(V^3)
// memory: O(V^2)
struct PushRelabel {
   vector<vector<ll>> cap, flow;
   vector<int> height;
```

```
PushRelabel() {}
void resize(int N) { cap.assign(N, vector<11>(N)); }
// push as much excess flow as possible from u to v.
void push(int u, int v) {
   ll f = min(excess[u], cap[u][v] - flow[u][v]);
   flow[u][v] += f:
   flow[v][u] -= f:
   excess[v] += f;
   excess[u] -= f:
// relabel the height of a vertex so that excess flow may
     be pushed.
void relabel(int u) {
   int d = INT32 MAX:
   rep(v, cap.size()) if (cap[u][v] - flow[u][v] > 0) d
       min(d, height[v]);
   if (d < INF) height[u] = d + 1;</pre>
// get the maximum flow on the network specified by 'cap'
     with source 's'
// and sink 't'.
// node-to-node flows are output to the 'flow' member.
11 maxflow(int s, int t) {
   int N = cap.size(), M;
   flow.assign(N, vector<11>(N));
   height.assign(N, 0), height[s] = N;
   excess.assign(N, 0), excess[s] = INF;
   rep(i, N) if (i != s) push(s, i);
   vector<int> q;
   while (true) {
       // find the highest vertices with excess
       q.clear(), M = 0;
       rep(i, N) {
           if (excess[i] <= 0 || i == s || i == t)</pre>
           if (height[i] > M) q.clear(), M = height[i];
           if (height[i] >= M) g.push back(i):
       if (q.empty()) break;
       // process vertices
       for (int u : q) {
          bool relab = true:
           rep(v, N) {
              if (excess[u] <= 0) break:
```

3.16 strongly-connected-components

```
/* time: O(V + E), memory: O(V)
after building:
   comp = map from vertex to component
         (components are toposorted, root first, leaf last)
   N = number of components
   G = condensation graph (component DAG)
bvproducts:
   vgi = transposed graph
   order = reverse topological sort (leaf first, root last)
   vn = number of vertices
   vg = original vertex graph
struct Scc {
   int vn. N:
   vector<int> order. comp:
   vector<vector<int>> vg, vgi, G;
   void toposort(int u) {
      if (comp[u]) return;
       comp[u] = -1;
       for (int v : vg[u]) toposort(v);
       order.push_back(u);
   bool carve(int u) {
       if (comp[u] != -1) return false;
       comp[u] = N:
      for (int v : vgi[u]) {
          if (comp[v] != N) G[comp[v]].push back(N):
      }
       return true;
   Scc() {}
```

```
Scc(vector<vector<int>> &g)
   : vn(g.size()), vg(g), comp(vn), vgi(vn), G(vn), N(0) {
     rep(u, vn) toposort(u);
     rep(u, vn) for (int v : vg[u]) vgi[v].push_back(u);
     invrep(i, vn) N += carve(order[i]);
}
```

3.17 two-sat

```
// calculate the solvability of a system of logical
    equations, where every equation is of the form 'a or b
// 'neg': get negation of 'u'
// 'then': 'u' implies 'v'
// 'anv': 'u' or 'v'
// 'set': 'u' is true
// after 'solve' (O(V+E)) returns true, 'sol' contains one
    possible solution.
// determining all solutions is O(V*E) hard (requires
     computing reachability in a DAG).
struct TwoSat {
   int N: vector<vector<int>> G:
   Scc scc; vector<bool> sol;
   TwoSat(int n) : N(n), G(2 * n), sol(n) {}
   TwoSat() {}
   int neg(int u) { return (u + N) % (2 * N); }
   void then(int u, int v) { G[u].push_back(v), G[neg(v)].
        push_back(neg(u)); }
   void anv(int u, int v) { then(neg(u), v); }
   void set(int u) { G[neg(u)].push_back(u); }
   bool solve() {
       scc = Scc(G);
       rep(u, N) if (scc.comp[u] == scc.comp[neg(u)]) return
       rep(u, N) sol[u] = (scc.comp[u] > scc.comp[neg(u)]);
       return true:
   }
};
```

4 implementation

4.1 common-template-and-bit-tricks

```
#pragma GCC optimize("Ofast")
#pragma GCC target("bmi,bmi2,lzcnt,popcnt")
#pragma GCC target("avx,avx2,f16c,fma,sse3,sse3,sse4.1,sse4
#include <bits/stdc++.h>
using namespace std:
typedef long long 11;
#define repx(i, a, b) for (int i = a; i < b; i++)
#define rep(i, n) repx(i, 0, n)
#define invrepx(i, a, b) for (int i = b - 1; i \ge a; i--)
#define invrep(i, n) invrepx(i, 0, n)
// Command to check time and memory usage:
      /usr/bin/time -v ./tmp
// See "Maximum resident set size" for max memory used
// Commands for interactive checker:
11
      mkfifo fifo
11
      (./solution < fifo) | (./interactor > fifo)
// Does not work on the Windows file system, i.e., /mnt/c/
// The special fifo file must be used, otherwise the
// solution will not wait for input and will read EOF
v = x & (x-1) // Turn off rightmost 1bit
v = x & (-x) // Isolate rightmost 1bit
y = x \mid (x-1) // Right propagate rightmost 1bit(fill in 1s)
y = x \mid (x+1) // Turn on rightmost Obit
y = x & (x+1) // Isolate rightmost Obit
// If x is of long type, use __builtin_popcountl(x)
// If x is of long long type, use __builtin_popcountll(x)
// 1. Counts the number of ones(set bits) in an integer.
__builtin_popcount(x)
// 2. Checks the Parity of a number. Returns true(1) if the
// number has odd number of set bits, else it returns
// false(0) for even number of set bits.
__builtin_parity(x)
// 3. Counts the leading number of zeros of the integer.
__builtin_clz(x)
// 4. Counts the trailing number of zeros of the integer.
__builtin_ctz(x)
// 5. Returns 1 + the index of the least significant 1-bit.
__builtin_ffs(x) // If x == 0, returns 0.
// Iterate over non empty subsets of bitmask
for(int s=m;s;s=(s-1)&m) // Decreasing order
for(int s=0;s=s-m&m;) // Increasing order
```

4.2 dp-convex-hull-trick

```
struct Line {
    mutable 11 a, b, c;
    bool operator<(Line r) const { return a < r.a; }</pre>
    bool operator<(11 x) const { return c < x; }</pre>
}:
// dynamically insert 'a*x + b' lines and guery for maximum
// at any x all operations have complexity O(\log N)
struct LineContainer : multiset<Line. less<>> {
    11 div(11 a, 11 b) {
       return a / b - ((a ^ b) < 0 && a % b):
    bool isect(iterator x, iterator y) {
       if (y == end()) return x \rightarrow c = INF, 0;
       if (x->a == y->a) x->c = x->b > y->b ? INF : -INF;
       else x->c = div(v->b - x->b, x->a - v->a):
       return x->c >= y->c;
    void add(ll a, ll b) {
       // a *= -1. b *= -1 // for min
       auto z = insert(\{a, b, 0\}), y = z++, x = y;
       while (isect(y, z)) z = erase(z);
       if (x != begin() && isect(--x, y)) isect(x, y = erase
       while ((y = x) != begin() \&\& (--x)->c >= y->c) isect(
            x. erase(v)):
    11 query(11 x) {
       if (empty()) return -INF; // INF for min
       auto 1 = *lower bound(x):
       return 1.a * x + 1.b;
       // return -l.a * x - l.b: // for min
};
```

4.3 dp-divide-and-conquer

```
// for every index i assign an optimal index j, such that // cost(i, j) is minimal for every i. the property that if // i2 >= i1 then j2 >= j1 is exploited (monotonic condition) // calculate optimal index for all indices in range [l, r) // knowing that the optimal index for every index in this // range is within [opt1, optr). // time: O(N log N)
```

```
void calc(vector<int> &opt, int 1, int r,int optl,int optr){
    if (l == r) return;
    int i = (l + r) / 2;
    ll optc = INF;
    int optj;
    repx(j, optl, optr) {
        ll c = i + j; // cost(i, j)
        if (c < optc) optc = c, optj = j;
    }
    opt[i] = optj;
    calc(opt, l, i, optl, optj + 1);
    calc(opt, i + 1, r, optj, optr);
}</pre>
```

4.4 dynamic-connectivity

```
struct DC {
   int n: Dsu D:
   vector<vector<pair<int, int>>> t;
   DC(int N) : n(N), D(N), t(2 * N) {}
   // add edge p to all times in interval [1, r]
   void upd(int 1, int r, pair<int, int> p) {
       for (1 += n, r += n; 1 < r; 1 >>= 1, r >>= 1) {
          if (1 & 1) t[1++].push_back(p);
          if (r & 1) t[--r].push_back(p);
      }
   }
   void process(int u = 1) { // process all queries
       for (auto &e : t[u]) D.unite(e.first, e.second):
      if (u >= n) {
          // do stuff with D at time u - n
      } else process(2 * u), process(2 * u + 1);
       for (auto &e : t[u]) D.rollback();
};
```

4.5 hash-container

```
namespace{//add (#define tmpl template)(#define ty typename)
  tmpl<ty T> size_t mk_h(const T& v){return hash<T>()(v);}
  void h_cmb(size_t& h, const size_t& v)
  { h ^= v + 0x9e3779b9 + (h << 6) + (h >> 2); }
  tmpl<ty T> struct h_ct{size_t operator()(const T& v)const{
  size_t h=0;for(const auto& e:v){h_cmb(h,mk_h(e));}return h;
  }};
}namespace std{//support for pair<T,U>, vector<T> & map<T,U>
  tmpl<ty T, ty U> struct hash<pair<T, U>>{
```

```
size_t operator()(const pair<T,U>& v) const
{size_t h=mk_h(v.first);h_cmb(h, mk_h(v.second));return h;}
};
tmpl<ty... T>struct hash<vector<T...>>:h_ct<vector<T...>>{};
tmpl<ty... T>struct hash<map<T...>>:h_ct<map<T...>>{};
}
```

4.6 mo

```
struct Query { int 1, r, idx; };
// answer segment queries using only 'add(i)', 'remove(i)'
    and 'get()'
// functions.
// complexity: O((N + Q) * sqrt(N) * F)
// N = length of the full segment
// Q = amount of queries
// F = complexity of the 'add', 'remove' functions
template <class A, class R, class G, class T>
void mo(vector<Query> &queries, vector<T> &ans, A add, R
    remove, G get) {
   int Q = queries.size(), B = (int)sqrt(Q);
   sort(queries.begin(), queries.end(), [&](Query &a, Query
       return make_pair(a.1 / B, a.r) < make_pair(b.1 / B, b</pre>
   }):
   ans.resize(Q):
   int 1 = 0, r = 0;
   for (auto &g : gueries) {
       while (r < q.r) add(r), r++;
       while (1 > q.1) 1--, add(1);
       while (r > q.r) r--, remove(r):
       while (1 < q.1) remove(1), 1++;
       ans[q.idx] = get();
   }
```

4.7 ordered-set

```
// find_by_order(i) -> iterator to ith element
// order_of_key(k) -> position (int) of lower_bound of k
```

4.8 unordered-map

5 math

5.1 Linear Diophantine

```
ii extendedEuclid(ll a, ll b){
   11 x, y; //a*x + b*y = gcd(a,b)
   if (b == 0) return {1, 0};
   auto p = extendedEuclid(b, a%b);
   x = p.second;
   y = p.first - (a/b)*x;
   if(a*x + b*y == -_gcd(a,b)) x=-x, y=-y;
   return {x, v}:
pair<ii, ii> diophantine(ll a, ll b, ll r){
   //a*x+b*y=r where r is multiple of gcd(a,b);
   11 d = \_gcd(a, b);
   a/=d; b/=d; r/=d;
   auto p = extendedEuclid(a, b);
   p.first*=r; p.second*=r;
   assert(a*p.first + b*p.second == r);
   return {p, {-b, a}}; //solutions: p+t*ans.second
```

5.2 arithmetic

```
inline int floor_log2(int n)
{ return n <= 1 ? 0 : 31 - __builtin_clz(n); }
inline int ceil_log2(int n)
{ return n <= 1 ? 0 : 32 - __builtin_clz(n - 1); }
inline ll floordiv(ll a, ll b) {return a/b-((a^b)<0&&a%b);}
inline ll ceildiv(ll a, ll b) {return a/b+((a^b)>=0&&a%b);}
```

5.3 berlekamp-massey-linear-recurrence

```
vector<int> BM(vector<int> x) {
   vector<int> ls. cur:
   int lf, ld;
   rep(i, x.size()) {
      11 t = 0:
       rep(j, cur.size()) t = (t+x[i-j-1]*(l1)cur[j])%MOD;
       if ((t - x[i]) \% MOD == 0) continue:
       if (!cur.size()) {
          cur.resize(i + 1); lf = i; ld = (t-x[i]) % MOD;
           continue:
       11 k = -(x[i] - t) * bin_exp(1d, MOD - 2) % MOD;
       vector<int> c(i - lf - 1): c.push back(k):
       rep(j, ls.size()) c.push_back(-ls[j] * k % MOD);
       if (c.size() < cur.size()) c.resize(cur.size());</pre>
       rep(j, cur.size()) c[j] = (c[j] + cur[j]) % MOD;
       if (i - lf + ls.size() >= cur.size())
          ls = cur, lf = i, ld = (t - x[i]) % MOD;
   rep(i, cur.size()) cur[i] = (cur[i] % MOD + MOD) % MOD:
   return cur:
// Linear Recurrence
11 MOD = 998244353;
11 LOG = 60:
struct LinearRec{
 typedef vector<int> vi;
 int n; vi terms, trans; vector<vi> bin;
 vi add(vi &a. vi &b){
   vi res(n*2+1):
   rep(i,n+1) rep(i,n+1)
       res[i+j]=(res[i+j]*1LL+(ll)a[i]*b[j])%MOD;
   for(int i=2*n: i>n: --i){
     rep(i,n)
       res[i-1-j]=(res[i-1-j]*1LL+(ll)res[i]*trans[j])%MOD;
     res[i]=0:
```

```
res.erase(res.begin()+n+1,res.end());
return res;
}
LinearRec(vi &terms, vi &trans):terms(terms),trans(trans){
    n=trans.size();vi a(n+1);a[1]=1;
    bin.push_back(a);
    repx(i,1,LOG)bin.push_back(add(bin[i-1],bin[i-1]));
}
int calc(ll k){
    vi a(n+1);a[0]=1;
    rep(i,LOG)if((k>>i)&1)a=add(a,bin[i]);
    int ret=0;
    rep(i,n)ret=((ll)ret+(ll)a[i+1]*terms[i])%MOD;
    ret = ret%MOD + MOD;
    return ret%MOD;
}
```

5.4 crt

```
pair<11, 11> solve_crt(const vector<pair<11, 11>> &eqs) {
    11 a0 = eqs[0].first, p0 = eqs[0].second;
    repx(i, 1, eqs.size()) {
        11 a1 = eqs[i].first, p1 = eqs[i].second;
        11 k1, k0;
        11 d = ext_gcd(p1, p0, k1, k0);
        a0 -= a1;
        if (a0 % d != 0) return {-1, -1};
        p0 = p0 / d * p1;
        a0 = a0 / d * k1 % p0 * p1 % p0 + a1;
        a0 = (a0 % p0 + p0) % p0;
    }
    return {a0, p0};
}
```

5.5 discrete-log

```
// discrete logarithm log_a(b).
// solve b ^ x = a (mod M) for the smallest x.
// returns -1 if no solution is found.
//
// time: O(sqrt(M))
ll dlog(ll a, ll b, ll M) {
    ll k = 1, s = 0;
    while (true) {
        ll g = __gcd(b, M);
        if (g <= 1) break;</pre>
```

```
if (a == k) return s;
if (a % g != 0) return -1;
a /= g, M /= g, s += 1, k = b / g * k % M;
}
ll N = sqrt(M) + 1;

umap<ll, ll> r;
rep(q, N + 1) {
    r[a] = q;
    a = a * b % M;
}

ll bN = binexp(b, N, M), bNp = k;
repx(p, 1, N + 1) {
    bNp = bNp * bN % M;
    if (r.count(bNp)) return N * p - r[bNp] + s;
}
return -1;
```

5.6 fast-hadamard-transform

```
ll c1[MAXN+9],c2[MAXN+9];//MAXN must be power of 2!
void fht(ll* p, int n, bool inv){
   for(int l=1;2*1<=n;1*=2)for(int i=0;i<n;i+=2*1)fore(j,0,1</pre>
        ){
       11 u=p[i+j],v=p[i+l+j];
       if(!inv)p[i+j]=u+v,p[i+l+j]=u-v; // XOR
       else p[i+j]=(u+v)/2, p[i+l+j]=(u-v)/2;
       //if(!inv)p[i+j]=v,p[i+l+j]=u+v; // AND
       //else p[i+j]=-u+v,p[i+l+j]=u;
       //if(!inv)p[i+j]=u+v,p[i+l+j]=u; // OR
       //else p[i+j]=v,p[i+l+j]=u-v;
// like polynomial multiplication, but XORing exponents
// instead of adding them (also ANDing, ORing)
vector<ll> multiply(vector<ll>& p1, vector<ll>& p2){
   int n=1<<(32-_builtin_clz(max(SZ(p1),SZ(p2))-1));</pre>
   fore(i,0,n)c1[i]=0,c2[i]=0:
   fore(i,0,SZ(p1))c1[i]=p1[i];
   fore(i,0,SZ(p2))c2[i]=p2[i];
   fht(c1,n,false);fht(c2,n,false);
   fore(i,0,n)c1[i]*=c2[i];
   fht(c1.n.true):
   return vector<ll>(c1,c1+n);
```

5.7 fft

```
using cd = complex<double>:
const double PI = acos(-1);
// compute the DFT of a power-of-two-length sequence.
// if 'inv' is true, computes the inverse DFT.
void fft(vector<cd> &a, bool inv) {
   int N = a.size(), k = 0, b;
   assert(N == 1 << builtin ctz(N));</pre>
   repx(i, 1, N) {
       for (b = N >> 1; k \& b;) k ^= b, b >>= 1;
       if (i < (k ^= b)) swap(a[i], a[k]);</pre>
   for (int 1 = 2: 1 <= N: 1 <<= 1) {
       double ang = 2 * PI / 1 * (inv ? -1 : 1);
       cd wl(cos(ang), sin(ang));
       for (int i = 0; i < N; i += 1) {</pre>
           cd w = 1;
           rep(j, 1 / 2) {
              cd u = a[i + j], v = a[i + j + 1 / 2] * w;
              a[i + j] = u + v;
              a[i + i + 1 / 2] = u - v:
              w *= wl:
   if (inv) rep(i, N) a[i] /= N;
const 11 MOD = 998244353, ROOT = 15311432;
// const 11 MOD = 2130706433, ROOT = 1791270792;
// const 11 MOD = 922337203673733529711, ROOT =
    532077456549635698311:
void find root of unitv(ll M) {
   11 c = M - 1, k = 0;
   while (c \% 2 == 0) c /= 2, k += 1:
   // find proper divisors of M - 1
   vector<ll> divs;
   for (11 d = 1: d < c: d++) {
       if (d * d > c) break;
       if (c \% d == 0) rep(i, k + 1) divs.push back(d << i):
   rep(i, k) divs.push_back(c << i);</pre>
   // find any primitive root of M
   11 G = -1;
   repx(g, 2, M) {
       bool ok = true:
       for (ll d : divs) ok &= (binexp(g, d, M) != 1);
       if (ok) {
           G = g:
           break;
```

```
}
   }
   assert(G != -1);
   ll w = binexp(G, c, M):
   cerr << "M = c * 2^k + 1" << endl;
   cerr << " M = " << M << endl:
   cerr << " c = " << c << endl;
   cerr << " k = " << k << endl:
   cerr << " w^(2^k) == 1" << endl;
   cerr << " w = g^{(M-1)/2k} = g^{c} << endl;
   cerr << " g = " << G << endl;
   cerr << " w = " << w << endl:
// compute the DFT of a power-of-two-length sequence, modulo
     a special prime
// number with an Nth root of unity, where N is the length
    of the sequence.
void ntt(vector<ll> &a, bool inv) {
   vector<ll> wn:
   for (11 p = ROOT; p != 1; p = p * p % MOD) wn.push_back(p
       );
   int N = a.size(), k = 0, b;
   assert(N == 1 << __builtin_ctz(N) && N <= 1 << wn.size())
   rep(i, N) a[i] = (a[i] \% MOD + MOD) \% MOD;
   repx(i, 1, N) {
      for (b = N >> 1; k & b;) k ^= b, b >>= 1;
       if (i < (k = b)) swap(a[i], a[k]):
   for (int 1 = 2: 1 <= N: 1 <<= 1) {
      11 wl = wn[wn.size() - __builtin_ctz(1)];
       if (inv) wl = multinv(wl, MOD);
      for (int i = 0: i < N: i += 1) {
          11 w = 1;
          repx(j, 0, 1 / 2) {
              11 u = a[i + i], v = a[i + i + 1 / 2] * w %
              a[i + j] = (u + v) \% MOD;
              a[i + j + 1 / 2] = (u - v + MOD) \% MOD;
              w = w * w1 \% MOD:
      }
   11 q = multinv(N, MOD);
   if (inv) rep(i, N) a[i] = a[i] * q % MOD;
void convolve(vector<cd> &a, vector<cd> b, int n) {
   n = 1 \ll (32 - \_builtin\_clz(2 * n - 1));
   a.resize(n), b.resize(n):
   fft(a, false), fft(b, false);
```

```
rep(i, n) a[i] *= b[i];
fft(a, true);
```

5.8 gauss

```
const double EPS = 1e-9:
// solve a system of equations.
// complexity: O(\min(N, M) * N * M)
// 'a' is a list of rows
// the last value in each row is the result of the equation
// return values:
// 0 -> no solutions
// 1 -> unique solution, stored in 'ans'
// -1 -> infinitely many solutions, one of which is stored
     in 'ans'
// UNTESTED
int gauss(vector<vector<double>> a, vector<double> &ans) {
   int N = a.size(), M = a[0].size() - 1;
   vector<int> where(M. -1):
   for (int j = 0, i = 0; j < M && i < N; j++) {
       int sel = i;
       repx(k, i, N) if (abs(a[k][j]) > abs(a[sel][j])) sel
            = k:
       if (abs(a[sel][j]) < EPS) continue;</pre>
       repx(k, j, M + 1) swap(a[sel][k], a[i][k]);
       where[i] = i;
       rep(k, N) if (k != i) {
          double c = a[k][j] / a[i][j];
           repx(1, j, M + 1) a[k][1] -= a[i][1] * c;
      }
       i++;
   ans.assign(M, 0):
   rep(i, M) if (where[i] != -1) ans[i] = a[where[i]][M] / a
        [where[i]][i];
   rep(i, N) {
       double sum = 0;
       rep(j, M) sum += ans[j] * a[i][j];
       if (abs(sum - a[i][M]) > EPS) return 0;
   rep(i, M) if (where[i] == -1) return -1;
   return 1;
```

```
5.9 matrix
```

```
typedef vector<vector<double>> Mat;
Mat matmul(Mat 1, Mat r) {
   int n = 1.N, m = r.M, p = 1.M; assert(1.M == r.N);
   Mat a(n, vector<double>(m)); // neutral
   rep(i, n) rep(j, m)
       rep(k, p) a[i][j] = a[i][j] + l[i][k] * r[k][j];
   return a;
double reduce(vector<vector<double>> &A) {
   int n = A.size(), m = A[0].size();
   int i = 0, j = 0; double r = 1.;
   while (i < n && j < m) {</pre>
       int 1 = i:
       repx(k, i+1, n) if(abs(A[k][i]) > abs(A[l][i])) l=k;
       if (abs(A[1][j]) < EPS) \{ j++; r = 0.; continue; \}
       if (1 != i) { r = -r; swap(A[i], A[1]); }
       r *= A[i][i]:
       for (int k = m - 1; k >= j; k--) A[i][k] /= A[i][j];
       repx(k, 0, n) {
          if (k == i) continue:
          for(int l=m-1;l>=j;l--)A[k][l]-=A[k][j]*A[i][l];
       i++, j++;
   return r; // returns determinant
```

5.10 mobius

5.11 multiny

```
// a * x + b * y == gcd(a, b)

ll ext_gcd(ll a, ll b, ll &x, ll &y) {
    if (b == 0) { x = 1, y = 0; return a; }
    ll d = ext_gcd(b, a % b, y, x); y -= a / b * x; return d; }

// inverse exists if and only if a and M are coprime
// if M is prime: multinv(a, M) = (a**(M-2)) % M
```

```
11 multinv(11 a, 11 M)
{  ll x, y; ext_gcd(a, M, x, y); return x; }

// all modular inverses from 1 to inv.size()-1
void multinv_all(vector<1l> &inv) {
   inv[1] = 1;
   repx(i, 2, inv.size())
   inv[i] = MOD - (MOD / i) * inv[MOD % i] % MOD;
}
```

5.12 polar-rho

```
11 mulmod(l1 a, 11 b, 11 m) {
   11 r=a*b-(11)((long double)a*b/m+.5)*m;
   return r<0?r+m:r:</pre>
bool is_prime_prob(ll n, int a){
   if(n==a)return true:
   11 s=0, d=n-1;
   while (d\%2==0)s++,d/=2;
   11 x=expmod(a,d,n):
   if((x==1)||(x+1==n))return true;
   fore(_,0,s-1){
       x=mulmod(x.x.n):
       if(x==1)return false;
       if(x+1==n)return true:
   return false;
bool rabin(ll n){ // true iff n is prime
   if(n==1)return false:
   int ar[]={2.3.5.7.11.13.17.19.23}:
   fore(i,0,9)if(!is_prime_prob(n,ar[i]))return false;
   return true:
ll rho(ll n){
   if(!(n&1))return 2:
   11 x=2, y=2, d=1;
   11 c=rand()%n+1:
   while(d==1){
       x=(\text{mulmod}(x,x,n)+c)%n;
       fore(it,0,2) y=(mulmod(y,y,n)+c)%n;
       if(x>=y)d=_gcd(x-y,n);
       else d=__gcd(y-x,n);
   return d==n?rho(n):d;
void fact(ll n, map<ll,int>& f){ //0 (lg n)^3
   if(n==1)return;
```

```
if(rabin(n)){f[n]++:return:}
    11 q=rho(n);fact(q,f);fact(n/q,f);
}
// optimized version: replace rho and fact with the
     following:
const int MAXP=1e6+1: // sieve size
int sv[MAXP]; // sieve
11 add(11 a, 11 b, 11 m){return (a+=b)<m?a:a-m;}</pre>
ll rho(ll n){
    static ll s[MAXP];
    while(1){
       11 x=rand()%n.v=x.c=rand()%n:
       ll *px=s,*py=s,v=0,p=1;
       while(1){
           *py++=y=add(mulmod(y,y,n),c,n);
           *py++=y=add(mulmod(y,y,n),c,n);
           if((x=*px++)==v)break:
           11 t=p; p=mulmod(p,abs(y-x),n);
           if(!p)return __gcd(t,n);
           if(++v==26){
               if((p=__gcd(p,n))>1&&p<n)return p;</pre>
           }
       if(v&&(p=_gcd(p,n))>1&&p<n)return p;</pre>
}
void init_sv(){ fore(i,2,MAXP)if(!sv[i])for(11 j=i;j<MAXP;j</pre>
     +=i)sv[j]=i; }
void fact(ll n.map<ll.int>&f){//call init sv first!
    for(auto&& p:f)while(n%p.fst==0)p.snd++,n/=p.fst;
    if(n<MAXP)while(n>1)f[sv[n]]++,n/=sv[n];
    else if(rabin(n))f[n]++;
    else {ll q=rho(n);fact(q,f);fact(n/q,f);}
```

5.13 polynomials

```
typedef int tp; // type of polynomial
template<class T=tp>
struct poly { // poly<> : 1 variable, poly<poly<>>>: 2
    variables, etc.
    vector<T> c;
    T& operator[](int k){return c[k];}
    poly(vector<T>& c):c(c){}
    poly(initializer_list<T> c):c(c){}
    poly(int k):c(k){}
    poly(){}
    poly operator+(poly<T> o);
```

```
poly operator*(tp k);
   poly operator*(poly o);
   poly operator-(poly<T> o){return *this+(o*-1);}
   T operator()(tp v){
      T sum(0);
       for(int i=c.size()-1:i>=0:--i)sum=sum*v+c[i]:
       return sum:
   }
// example: p(x,y)=2*x^2+3*x*y-y+4
// poly<poly<>> p={{4,-1},{0,3},{2}}
// printf("%d\n",p(2)(3)) // 27 (p(2.3))
set<tp> roots(poly<> p){ // only for integer polynomials
   while(!p.c.empty()&&!p.c.back())p.c.pop_back();
   if(!p(0))r.insert(0);
   if(p.c.empty())return r;
   tp a0=0,an=abs(p[p.c.size()-1]);
   for(int k=0:!a0:a0=abs(p[k++])):
   vector<tp> ps.qs;
   fore(i,1,sqrt(a0)+1) if (a0\%i==0) ps.pb(i),ps.pb(a0/i);
   fore(i,1,sqrt(an)+1)if(an\%i==0)qs.pb(i),qs.pb(an/i);
   for(auto pt:ps)for(auto qt:qs)if(pt%qt==0){
       tp x=pt/qt;
       if(!p(x))r.insert(x);
       if(!p(-x))r.insert(-x);
   return r:
pair<poly<>,tp> ruffini(poly<> p, tp r){ // returns pair (
    result.rem)
   int n=p.c.size()-1;
   vector<tp> b(n);
   b[n-1]=p[n];
   for(int k=n-2; k>=0; --k)b[k]=p[k+1]+r*b[k+1];
   return {polv<>(b).p[0]+r*b[0]}:
// only for double polynomials
pair<poly<>,poly<> > polydiv(poly<> p, poly<> q){ // returns
     pair (result.rem)
   int n=p.c.size()-q.c.size()+1:
   vector<tp> b(n);
   for(int k=n-1:k>=0:--k){
       b[k]=p.c.back()/q.c.back();
       fore(i,0,q.c.size())p[i+k]-=b[k]*q[i];
       p.c.pop_back();
   while(!p.c.empty()&&abs(p.c.back())<EPS)p.c.pop_back();</pre>
   return {poly<>(b),p};
```

```
// only for double polynomials
poly<> interpolate(vector<tp> x, vector<tp> y){
    poly<> q={1},S={0};
    for(tp a:x)q=poly<>({-a,1})*q;
    fore(i,0,x.size()){
        poly<> Li=ruffini(q,x[i]).fst;
        Li=Li*(1.0/Li(x[i])); // change for int polynomials
        S=S+Li*y[i];
    }
    return S;
}
```

5.14 primes

```
// counts the divisors of a positive integer in O(\operatorname{sqrt}(n))
ll count divisors(ll x) {
   11 \text{ divs} = 1, i = 2;
   for (ll divs = 1, i = 2; x > 1; i++) {
       if (i * i > x) { divs *= 2; break; }
       for (11 d = divs; x % i == 0; x /= i) divs += d;
   return divs;
// gets the prime factorization of a number in O(sqrt(n))
vector<pair<11, int>> factorize(11 x) {
   vector<pair<11, int>> f;
   for (ll k = 2; x > 1; k++) {
       if (k * k > x) { f.push_back({x, 1}); break; }
       int n = 0:
       while (x \% k == 0) x /= k, n++;
       if (n > 0) f.push back(\{k, n\}):
   return f;
// iterate over all divisors of a number.
// divisor count upper bound: n^(1.07 / ln ln n)
template <class OP>
void divisors(ll x, OP op) {
   auto facts = factorize(x):
   vector<int> f(facts.size());
   while (true) {
       11 v = 1:
       rep(i, f.size()) rep(j, f[i]) y *= facts[i].first;
       op(y);
       for (i = 0; i < f.size(); i++) {</pre>
           f[i] += 1:
           if (f[i] <= facts[i].second) break;</pre>
```

```
f[i] = 0:
       if (i == f.size()) break;
}
// computes euler totative function phi(x), counting the
// amount of integers in [1, x] that are coprime with x.
// time: O(sqrt(x))
11 phi(11 x) {
   11 phi = 1, k = 2;
    for (: x > 1: k++) {
       if (k * k > x) { phi *= x - 1; break; }
       11 k1 = 1, k0 = 0;
       while (x \% k == 0) x /= k, k0 = k1, k1 *= k;
       phi *= k1 - k0:
    return phi;
}
// test-prime.cpp
// change to __int128 if checking numbers over 10^9
bool isprime(ll n) {
    if (n < 2 | | n % 6 % 4 != 1) return n - 2 < 2;
    11 A[] = \{2,325,9375,28178,450775,9780504,1795265022\};
    ll s = \_builtin\_ctzll(n - 1), d = n >> s;
    for (int a : A) {
       ll p = binexp(a, d, n), i = s;
       while (p != 1 && p != n - 1 && a % n && i--) p = p *
       if (p != n - 1 && i != s) return 0;
    return 1;
}
```

5.15 simplex

/* Solves a general linear maximization problem: maximize \$c^T x\$ subject to \$Ax \le b\$, \$x \ge 0\$. Returns -inf if there is no solution, inf if there are arbitrarily good solutions, or the maximum value of \$c^T x\$ otherwise. The input vector is set to an optimal \$x\$ (or in the unbounded case, an arbitrary solution fulfilling the constraints). Numerical stability is not guaranteed. For better performance, define variables such that \$x = 0\$ is viable. Usage: vvd $A = \{\{1,-1\}, \{-1,1\}, \{-1,-2\}\};$ vd $b = \{1,1,-4\}, c = \{-1,-1\}, x;$ T val = LPSolver(A, b, c).solve(x); Time: $D(NM * \mbox{$\psi$pivots})$, where a pivot may be e.g. an edge relaxation. $D(2^n)$ in the general case.*/

```
typedef double T://long double, Rational, double + mod<P>...
typedef vector<T> vd:
typedef vector<vd> vvd;
const T eps = 1e-8, inf = 1 / .0:
#define MP make_pair
#define lti(X) \
   if (s == -1 \mid | MP(X[i], N[i]) < MP(X[s], N[s])) s = i
struct LPSolver {
   int m, n; vector<int> N, B; vvd D;
   LPSolver(const vvd &A,const vd &b,const vd &c) : m(b.size
        ()),n(c.size()),N(n+1),B(m),D(m+2,vd(n+2))
       rep(i, m) rep(i, n) D[i][i] = A[i][i]:
       rep(i, m) {
          B[i] = n + i; D[i][n] = -1; D[i][n + 1] = b[i];
       rep(j, n) \{ N[j] = j; D[m][j] = -c[j]; \}
       N[n] = -1; D[m + 1][n] = 1;
   }
   void pivot(int r. int s) {
      T *a = D[r].data(), inv = 1 / a[s]:
       rep(i, m + 2) if (i != r && abs(D[i][s]) > eps) {
          T *b = D[i].data(), inv2 = b[s] * inv;
          repx(j, 0, n + 2) b[j] -= a[j] * inv2;
          b[s] = a[s] * inv2:
       rep(j, n + 2) if (j != s) D[r][j] *= inv;
       rep(i, m + 2) if (i != r) D[i][s] *= -inv;
       D[r][s] = inv:
       swap(B[r], N[s]);
   bool simplex(int phase) {
       int x = m + phase - 1;
       for (::) {
          int s = -1;
          rep(j, n + 1) if (N[j] != -phase) ltj(D[x]);
          if (D[x][s] >= -eps) return true:
          int r = -1:
          rep(i, m) {
              if (D[i][s] <= eps) continue;</pre>
              if (r == -1 || MP(D[i][n + 1] / D[i][s], B[i])
                    < MP(D[r][n + 1] / D[r][s], B[r])) r = i
          if (r == -1) return false;
          pivot(r, s);
   }
   T solve(vd &x) {
       int r = 0:
       repx(i, 1, m) if (D[i][n + 1] < D[r][n + 1]) r = i;
```

```
if (D[r][n + 1] < -eps) {
    pivot(r, n);
    if (!simplex(2) || D[m + 1][n + 1] < -eps) return
        -inf;
    rep(i, m) if (B[i] == -1) {
        int s = 0;
        repx(j, 1, n + 1) ltj(D[i]);
        pivot(i, s);
    }
}
bool ok = simplex(1);
    x = vd(n);
    rep(i, m) if (B[i] < n) x[B[i]] = D[i][n + 1];
    return ok ? D[m][n + 1] : inf;
}
};</pre>
```

5.16 theorems-and-formulas

$$n! \sim \sqrt{2\pi n} \left(\frac{n}{e}\right)^n$$

$$\sum_{i=0}^k \binom{n+i}{i} = \binom{n+k+1}{k}$$

 $\begin{bmatrix} n \\ k \end{bmatrix} =$ perm of n elements with k cycles

 $\begin{Bmatrix} n \\ k \end{Bmatrix}$ = partitions of an *n*-element set into *k* parts

$${n \brace k} = \frac{1}{k!} \sum_{i=0}^{k} (-1)^i {k \choose i} (k-i)^n.$$
$${n+1 \brace k} = k {n \brace k} + {n \brace k-1}$$

Integers $d_1 \ge \cdots \ge d_n \ge 0$ can be the degree sequence of a finite simple graph on n vertices \iff

 $d_1 + \cdots + d_n$ is even and for every k in $1 \le k \le n$

$$\sum_{i=1}^{k} d_i \le k(k-1) + \sum_{i=k+1}^{n} \min(d_i, k)$$

 $a^n = a^{\varphi(m)+n \mod \varphi(m)} \pmod{m}$ if n > lg(m)

Misere Nim: if $\exists a_i > 1$ then normal nim; else the condition is reversed.

Derangements: Num of permutations of n = 0, 1, 2, ...

elements without fixed points is $1,0,1,2,9,44,265,1854,14833,.... \ \ \text{Recurrence:}$ $D_n=(n-1)(D_{n-1}+D_{n-2})=n*D_{n-1}+(-1)^n. \ \ \text{Collary:}$ number of permutations with exactly k fixed points $nCkD_{n-k}$

Eulerian numbers: E(n,k) is the number of permutations with exactly k descents $(i: \pi_i < \pi_{i+1})$, ascents $(\pi_i > \pi_{i+1})$ / excedances $(\pi > i)$ / k+1 weak excedances $(\pi \ge i)$. $E_{n,k} = (k+1)E_{n-1,k} + (n-k)E_{n-1,k-1}$

5.17 theorems

Burnside lemma

Tomemos imagenes x en X y operaciones (g: X -> X) en G. Si #g es la cantidad de imagenes que son puntos fijos de g, entonces la cantidad de objetos es '(sum_{g} in G} #g) / |G|' Es requisito que G tenga la operacion identidad, que toda operacion tenga inversa y que todo par de operaciones tenga su combinacion.

Rational root theorem

Las raices racionales de un polinomio de orden n con coeficientes enteros A[i] son de la forma p / q, donde p y q son coprimos, p es divisor de A[0] y q es divisor de A[n]. Notar que si A[0] = 0, cero es raiz, se puede dividir el polinomio por x y aplica nuevamente el teorema.

Petersens theorem

Every cubic and bridgeless graph has a perfect matching.

Number of divisors for powers of 10 (0,1) (1,4) (2,12) (3,32) (4,64) (5,128) (6,240) (7,448) (8,768) (9,1344) (10,2304) (11,4032) (12,6720) (13,10752) (14,17280) (15,26880) (16,41472) (17,64512) (18,103680)

Kirchoff Theorem: Sea A la matriz de adyacencia del multigrafo (A[u][v] indica la cantidad de aristas entre u y v) Sea D una matriz diagonal tal que D[v][v] es igual al grado de v (considerando auto aristas y multi aristas). Sea L = A - D. Todos los cofactores de L son iguales y equivalen a la cantidad de Spanning Trees del grafo. Un cofactor (i,j) de L es la multiplicacin de $(-1)^{i}$ + j con el determinant de la matriz al quitar la fila i y la columna j

Prufer Code: Dado un rbol con los nodos indexados: busca la hoja de menor ndice, brrala y anota el ndice del nodo al que estaba conectado. Repite el paso anterior n-2 veces. Lo

anterior muestra una biyeccin entre los arreglos de tamao n-2 con elementos en [1, n] y los rboles de n nodos, por lo que hay n^{n-2} spanning trees en un grafo completo. Corolario: Si tenemos k componentes de tamaos $s1,s2,\ldots,sk$ entonces podemos hacerlos conexos agregando k-1 aristas entre nodos de $s1*s2*\ldots*sk*n^{k-2}$ formas

Combinatoria Catalan: $C_{n+1} = sum(C_i*C_{n-i})$ for i \in [0, n]) Catalan: $C_n = \frac{1}{n+1}*\frac{0}{n}$ Sea C_n^k las formas de poner n+k pares de parntesis, con los primeros k parntesis abiertos (esto es, hay 2n + 2k carcteres), se tiene que $C_n^k = \frac{(2n+k-1)*(2n+k)}{(n*(n+k+1))} * C_{n-1}^k$ Sea D_n el nmero de permutaciones sin puntos fijos, entoces $D_n = (n-1)*(D_{n-1} + D_{n-2})$, $D_0 = 1$, $D_1 = 0$

5.18 tonelli-shanks

```
ll legendre(ll a, ll p) {
    if (a % p == 0) return 0; if (p == 2) return 1;
   return binexp(a, (p-1) / 2, p);
// sqrt(n) mod p (p must be a prime)
// rnd(a, b) return a random number in [a, b]
ll tonelli shanks(ll n. ll p) {
    if (n == 0) return 0;
    if (legendre(n, p) != 1) return -1; // no existe
   if (p == 2) return 1:
   ll s = __builtin_ctzll(p - 1);
   11 q = (p - 1LL) >> s, z = rnd(1, p - 1);
    if (s == 1) return binexp(n, (p + 1) / 4LL, p);
    while (legendre(z, p) != p - 1) z = rnd(1, p - 1);
   ll c = binexp(z, q, p), r = binexp(n, (q + 1) / 2, p):
   ll t = binexp(n, q, p), m = s;
   while (t != 1) {
       11 i = 1, ts = (t * t) \% p;
       while (ts != 1) i++, ts = (ts * ts) % p;
       11 b = c:
       repx(_, 0, m - i - 1) b = (b * b) \% p;
       r = r*b\%p; c = b*b\%p; t = t*c\%p; m = i;
    return r;
```

6 strings

6.1 aho-corasick

```
struct Vertex {
   int next[26], go[26];
   int p, link = -1, exit = -1, cnt = -1;
   vector<int> leaf:
   char pch:
   Vertex(int p=-1, char ch='$') : p(p), pch(ch) {
      rep(i, 26) next[i] = -1, go[i] = -1:
};
vector<Vertex> t(1):
void add(string &s, int id) {
   int v = 0:
   for (char ch : s) {
       int c = ch - 'a':
       if (t[v].next[c] == -1) {
          t[v].next[c] = t.size();
          t.emplace_back(v, ch);
       v = t[v].next[c];
   t[v].leaf.push_back(id);
int go(int v. char ch);
int get_link(int v) {
   if (t[v].link == -1) {
       if (v == 0 || t[v].p == 0) t[v].link = 0;
       else t[v].link = go(get_link(t[v].p), t[v].pch);
   return t[v].link:
int go(int v. char ch) {
   int c = ch - 'a':
   if (t[v].go[c] == -1) {
       if (t[v].next[c] != -1) t[v].go[c] = t[v].next[c];
       else t[v].go[c] = v == 0 ? 0 : go(get_link(v), ch);
   return t[v].go[c];
int next_match(int v){ // Optional
   if(t[v].exit == -1){
       if(t[get_link(v)].leaf.size())t[v].exit=get_link(v);
       else t[v].exit = v==0 ? 0 : next match(get link(v));
   return t[v].exit;
int cnt_matches(int v){ // Optional
```

```
if(t[v].cnt == -1)
    t[v].cnt = v == 0 ? 0 : t[v].leaf.size() +
        cnt_matches(get_link(v));
return t[v].cnt;
```

6.2 debruijn-sequence

```
// Given alphabet [0,k) constructs a cyclic string of length
// k^n that contains every length n string as substr.
vector<int> deBruijnSeq(int k, int n) { // Recursive FKM
   if (k == 1) return {0};
   vector<int> seq, aux(n+1);
   function<void(int,int)> gen = [&](int t, int p) {
      if (t > n) { // +lyndon word of len p
         if (n%p == 0) repx(i,1,p+1) seq.pb(aux[i]);
      } else {
        aux[t] = aux[t-p]; gen(t+1,p);
        while (++aux[t] < k) gen(t+1,t);
      }
   };
   gen(1,1); return seq;
}</pre>
```

6.3 hash

```
const int K = 2:
struct Hash{
   const 11 MOD[K] = {999727999, 1070777777};
   const 11 P = 1777771:
   vector<ll> h[K], p[K];
   Hash(string &s){
      int n = s.size();
      rep(k, K){}
          h[k].resize(n+1, 0):
          p[k].resize(n+1, 1);
          repx(i, 1, n+1){
              h[k][i] = (h[k][i-1]*P + s[i-1]) % MOD[k];
              p[k][i] = (p[k][i-1]*P) \% MOD[k];
          }
      }
   vector<ll> get(int i, int j){
      vector<ll> r(K);
      rep(k, K){
          r[k] = (h[k][j] - h[k][i]*p[k][j-i]) % MOD[k];
          r[k] = (r[k] + MOD[k]) \% MOD[k];
```

```
} return r;
}
```

6.4 manacher

```
// odd[i]: length of longest palindrome centered at i
// even[i]: ...longest palindrome centered between i and i+1
void manacher(string &s,vector<int> &odd,vector<int> &even){ };
   string t = "$#":
   for(char c: s) t += c + string("#");
   t += "^":
   int n = t.size():
   vector<int> p(n);
   int 1 = 1, r = 1:
   repx(i, 1, n-1) {
      p[i] = max(0, min(r - i, p[1 + (r - i)]));
       while(t[i - p[i]] == t[i + p[i]]) p[i]++;
       if(i + p[i] > r) l = i - p[i], r = i + p[i];
   repx(i, 2, n-2) {
      if(i%2) even.push_back(p[i]-1);
       else odd.push_back(p[i]-1);
   }
```

6.5 palindromic-tree

```
struct Node { // (*) = Optional
   int len: // length of substring
   int to[26]; // insertion edge for all characters a-z
   int link; // maximum palindromic suffix
   int i; // (*) start index of current Node
   int cnt; // (*) # of occurrences of this substring
   Node(int len, int link=0, int i=0, int cnt=1): len(len).
   link(link), i(i), cnt(cnt) {memset(to, 0, sizeof(to));}
}: struct EerTree { // Palindromic Tree
   vector<Node> t: // tree (max size of tree is n+2)
   int last:
                 // current node
   EerTree(string &s) : last(0) {
       t.emplace_back(-1); t.emplace_back(0); // root 1 & 2
       rep(i, s.size()) add(i, s); // construct tree
       for(int i = t.size()-1: i > 1: i--)
          t[t[i].link].cnt += t[i].cnt;
   void add(int i, string &s){
                                   // vangrind warning:
       int p=last, c=s[i]-'a';
                                   // i-t[p].len-1 = -1
```

```
while(s[i-t[p].len-1] != s[i]) p = t[p].link;
    if(t[p].to[c]){ last = t[p].to[c]; t[last].cnt++; }
    else{
        int q = t[p].link;
        while(s[i-t[q].len-1] != s[i]) q = t[q].link;
        q = max(1, t[q].to[c]);
        last = t[p].to[c] = t.size();
        t.emplace_back(t[p].len + 2, q, i-t[p].len-1);
    }
};
void main(){
    string s = "abcbab"; EerTree pt(s); // build EerTree
    repx(i, 2, pt.t.size()){// list all distinct palindromes
        repx(j,pt.t[i].i,pt.t[i].i+pt.t[i].len)cout << s[j];
        cout << " " << pt.t[i].cnt << endl;
}</pre>
```

6.6 prefix-function

```
vector<int> prefix_function(string s) {
   int n = s.size():
   vector<int> pi(n):
   repx(i, 1, n) {
       int j = pi[i-1];
       while (j > 0 \&\& s[i] != s[j])
          j = pi[j-1];
       if (s[i] == s[j])
           j++;
       pi[i] = j;
   return pi;
vector<vector<int>> aut;
void compute_automaton(string s) {
   s += '#':
   int n = s.size();
   vector<int> pi = prefix_function(s);
   aut.assign(n, vector<int>(26));
   rep(i, n) {
      rep(c, 26) {
          int j = i;
          while (j > 0 \&\& 'a' + c != s[j])
              j = pi[j-1];
           if ('a' + c == s[i])
              j++;
           aut[i][c] = j;
```

```
}    }    // k = n - pi[n - 1]; if k divides n, then the string can be    // aprtitioned into blocks of length k otherwise there is no    // effective compression and the answer is n.
```

6.7 suffix-array

```
// build the suffix array
// suffixes are sorted, with each suffix represented by its
// starting position
vector<int> suffixarray(const string &s) {
   int N = s.size() + 1;//optional: include terminating NUL
   vector<int> p(N), p2(N), c(N), c2(N), cnt(256);
   rep(i, N) cnt[s[i]] += 1;
   repx(b, 1, 256) cnt[b] += cnt[b - 1];
   rep(i, N) p[--cnt[s[i]]] = i;
   repx(i, 1, N) c[p[i]] = c[p[i - 1]] + (s[p[i]] != s[p[i - 1]])
         1]]):
   for (int k = 1; k < N; k <<= 1) {</pre>
       int C = c[p[N - 1]] + 1;
       cnt.assign(C + 1, 0);
       for (int &pi : p) pi = (pi - k + N) % N;
       for (int cl : c) cnt[cl + 1] += 1:
       rep(i, C) cnt[i + 1] += cnt[i];
       rep(i, N) p2[cnt[c[p[i]]]++] = p[i];
       c2[p2[0]] = 0;
       repx(i, 1, N) c2[p2[i]] =
           c2[p2[i-1]] + (c[p2[i]] != c[p2[i-1]] ||
                          c[(p2[i] + k) \% N] != c[(p2[i - 1]
                                + k) % N1):
       swap(c, c2), swap(p, p2);
   p.erase(p.begin()); // optional: erase terminating NUL
   return p:
// build the lcp
// 'lcp[i]' represents the length of the longest common
// prefix between suffix i and suffix i+1 in the suffix
```

```
//array 'p'. the last element of 'lcp' is zero by convention
vector<int> makelcp(const string &s, const vector<int> &p) {
   int N = p.size(), k = 0;
   vector<int> r(N), lcp(N);
   rep(i, N) r[p[i]] = i;
   rep(i, N) {
      if (r[i] + 1 >= N) { k = 0; continue; }
      int j = p[r[i] + 1];
       while (i + k < N && j + k < N && s[i + k] == s[j + k]
      lcp[r[i]] = k:
      if (k) k -= 1:
   return lcp;
// lexicographically compare the suffixes starting from 'i'
// and 'j', considering only up to 'K' characters.
// 'r' is the inverse suffix array, mapping suffix offsets
// to indices, requires an LCP sparse table.
int lcp_cmp(vector<int> &r, Sparse<int> &lcp, int i, int j,
    int K) {
   if (i == j) return 0;
   int ii = r[i], jj = r[j];
   int 1 = lcp.query(min(ii, jj), max(ii, jj));
   if (1 >= K) return 0;
   return ii < jj ? -1 : 1;
```

6.8 suffix-automaton

```
int a = st[p].next[c]:
       if(st[p].len + 1 == st[q].len) st[k].link = q;
       else {
           int w = sz++; st[w].len = st[p].len + 1;
           st[w].next=st[q].next; st[w].link=st[q].link;
          for(; p!=-1 && st[p].next[c]==q; p=st[p].link)
              st[p].next[c] = w;
          st[q].link=st[k].link = w;
      }
   }
   last = k:
} // # states <= 2n-1 && transitions <= 3n-4 (for n > 2)
// Follow link from 'last' to 0, nodes on path are terminal
// # matches = # paths from state to a terminal node
// # substrings = # paths from 0 to any node
// # substrings = sum of (len - len(link)) for all nodes
```

6.9 z-function

```
// i-th element is equal to the greatest number of
// characters starting from the position i that coincide
// with the first characters of s
vector<int> z_function(string s) {
   int n = s.size();
   vector<int> z(n);
   int l = 0, r = 0;
   for(int i = 1; i < n; i++) {
      if(i < r) z[i] = min(r - i, z[i - l]);
      while(i + z[i] < n && s[z[i]] == s[i + z[i]])z[i]++;
      if(i + z[i] > r) {
        l = i;
        r = i + z[i];
      }
   }
   return z;
}
```