Team Notebook

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1 Strings

1.1 Manacher

```
// odd[i]: length of longest palindrome centered at i
// even[i]: ...longest palindrome centered between i and i+1
void manacher(string &s,vector<int> &odd,vector<int> &even){
   string t = "$#";
   for(char c: s) t += c + string("#");
   t += "^":
   int n = t.size();
   vector<int> p(n):
   int 1 = 1, r = 1;
   repx(i, 1, n-1) {
       p[i] = max(0, min(r - i, p[1 + (r - i)]));
       while(t[i - p[i]] == t[i + p[i]]) p[i]++;
       if(i + p[i] > r) l = i - p[i], r = i + p[i];
   repx(i, 2, n-2) {
       if(i%2) even.push_back(p[i]-1);
       else odd.push_back(p[i]-1);
}
```

1.2 aho-corasick

```
const int K = 26:
struct Vertex {
    int next[K];
    int leaf = 0:
    int leaf_id = -1;
    int p = -1;
    char pch:
    int link = -1;
    int exit = -1:
    int cnt = -1:
    int go[K];
    Vertex(int p=-1, char ch='$') : p(p), pch(ch) {
       fill(begin(next), end(next), -1);
       fill(begin(go), end(go), -1);
   }
};
vector<Vertex> t(1):
void add(string &s, int id) {
    int. v = 0:
    for (char ch : s) {
```

```
int c = ch - 'a':
      if (t[v].next[c] == -1) {
          t[v].next[c] = t.size();
          t.emplace_back(v, ch);
      v = t[v].next[c]:
   t[v].leaf++;
   t[v].leaf_id = id;
int go(int v. char ch):
int get_link(int v) {
   if (t[v].link == -1) {
      if (v == 0 || t[v].p == 0)
          t[v].link = 0:
          t[v].link = go(get link(t[v].p), t[v].pch):
   }
   return t[v].link;
int go(int v, char ch) {
   int c = ch - 'a':
   if (t[v].go[c] == -1) {
      if (t[v].next[c] != -1)
          t[v].go[c] = t[v].next[c]:
          t[v].go[c] = v == 0 ? 0 : go(get_link(v), ch);
   return t[v].go[c];
int next match(int v){
   if(t[v].exit == -1)
      if(t[get_link(v)].leaf)
          t[v].exit = get_link(v);
          t[v].exit = v==0 ? 0 : next match(get link(v)):
   return t[v].exit:
int cnt matches(int v){
   if(t[v].cnt == -1)
      t[v].cnt = v == 0 ? 0 : t[v].leaf + cnt_matches(
           get link(v)):
   return t[v].cnt;
```

1.3 hash

```
// compute substring hashes in O(1).
// hashes are compatible between different strings.
struct Hash {
   11 HMOD:
   int N:
   vector<int> h:
   vector<int> p;
   Hash() {}
   // O(N)
   Hash(const string \&s. 11 HMOD = 1000003931)
       : N(s.size() + 1), HMOD(HMOD_), p(N), h(N) {
       static const 11 P =
           chrono::steady_clock::now().time_since_epoch().
               count() % (1 << 29);
       p[0] = 1;
       rep(i, N - 1) p[i + 1] = p[i] * P % HMOD;
       rep(i, N-1) h[i+1] = (h[i] + (ll)s[i] * p[i]) %
   }
   pair<11, int> get(int i, int j) { return {(h[j] - h[i] +
        HMOD) % HMOD, il: }
   bool cmp(pair<11, int> x0, pair<11, int> x1) {
       int d = x0.second - x1.second:
       11 &lo = d < 0 ? x0.first : x1.first;</pre>
      lo = lo * p[abs(d)] % HMOD;
       return x0.first == x1.first;
   }
};
// compute hashes in multiple prime modulos simultaneously,
// to reduce the chance of collisions.
struct HashM {
   int N:
   vector<Hash> sub:
   HashM() {}
   // O(K N)
   HashM(const string &s, const vector<11> &mods) : N(mods.
        size()), sub(N) {
       rep(i, N) sub[i] = Hash(s, mods[i]);
```

1.4 hash2d

```
using Hash = pair<11, int>;
struct Block {
   int x0, y0, x1, y1;
struct Hash2d {
   11 HMOD;
   int W. H:
   vector<int> h;
   vector<int> p;
   Hash2d() {}
   Hash2d(const string &s, int W_, int H_, 11 HMOD_ =
        1000003931)
       : W(W_ + 1), H(H_ + 1), HMOD(HMOD_) {
       static const 11 P =
          chrono::steady_clock::now().time_since_epoch().
               count() % (1 << 29);
      p.resize(W * H):
      p[0] = 1;
      rep(i, W * H - 1) p[i + 1] = p[i] * P % HMOD;
      h.assign(W * H, 0);
      repx(v, 1, H) repx(x, 1, W) {
```

```
ll c = (ll)s[(y - 1) * (W - 1) + x - 1] * p[y * W
          h[y * W + x] = (HMOD + h[y * W + x - 1] + h[(y -
                         h[(y - 1) * W + x - 1] + c) %
      }
   }
   bool isout(Block s) {
       return s.x0 < 0 || s.x0 >= W || s.x1 < 0 || s.x1 >= W
             | | s.v0 < 0 | |
             s.y0 >= H \mid \mid s.y1 < 0 \mid \mid s.y1 >= H;
   }
   Hash get(Block s) {
       return \{(2 * HMOD + h[s.v1 * W + s.x1] - h[s.v1 * W +
             s.x01 -
               h[s.v0 * W + s.x1] + h[s.v0 * W + s.x0]) %
              s.v0 * W + s.x0;
   }
   bool cmp(Hash x0, Hash x1) {
       int d = x0.second - x1.second;
       11 &lo = d < 0 ? x0.first : x1.first;</pre>
       lo = lo * p[abs(d)] % HMOD;
       return x0.first == x1.first:
   }
}:
struct Hash2dM {
   int N:
   vector<Hash2d> sub;
   Hash2dM(const string &s, int W, int H, const vector<11> &
        mods)
       : N(mods.size()), sub(N) {
       rep(i, N) sub[i] = Hash2d(s, W, H, mods[i]):
   bool isout(Block s) { return sub[0].isout(s): }
   vector<Hash> get(Block s) {
       vector<Hash> hs(N):
       rep(i, N) hs[i] = sub[i].get(s);
       return hs:
   }
```

1.5 palindromic-tree

```
struct Node { // (*) = Optional
int len; // length of substring
int edge[26];// insertion edge for all characters a-z
int link: // the Maximum Palindromic Suffix Node for the
     current Node
           // (*) start index of current Node
int i:
   int cnt = 1; // (*) # of occurrences of this substring
   Node(){ fill(begin(edge), end(edge), -1); }
}:
struct EerTree { // Palindromic Tree
   vector<Node> t; // tree
   int curr:
                 // current node
   EerTree(string &s) {
       t.resize(2);
       t.reserve(s.size()+2):// (*) max size of tree
       t[0].len = -1:
                          // root 1
       t[0].link = 0;
       t[1].len = 0:
                          // root 2
      t[1].link = 0;
       curr = 1:
       rep(i, s.size()) insert(i, s); // construct tree
       // (*) calculate number of occurrences of each node
      for(int i = t.size()-1; i > 1; i--)
          t[t[i].link].cnt += t[i].cnt;
   void insert(int i, string &s) {
       int tmp = curr;
       while (i - t[tmp].len < 1 \mid | s[i] != s[i-t[tmp].len
           -17)
          tmp = t[tmp].link;
```

```
if(t[tmp].edge[s[i]-'a'] != -1){
          curr = t[tmp].edge[s[i]-'a']; // already exists
          t[curr].cnt++:
                             // (*) increase cnt
          return:
       // create new node
       curr = t[tmp].edge[s[i]-'a'] = t.size();
       t.emplace_back();
      t[curr].len = t[tmp].len + 2; // set length
       t[curr].i = i - t[curr].len + 1://(*)set start index
      if (t[curr].len == 1) {
                                     // set suffix link
          t[curr].link = 1:
      } else {
          tmp = t[tmp].link:
          while (i-t[tmp].len < 1 \mid | s[i] != s[i-t[tmp].len
              tmp = t[tmp].link;
          t[curr].link = t[tmp].edge[s[i]-'a'];
   }
int main(){
string s = "abcbab";
   EerTree pt(s);
                        // construct palindromic tree
repx(i, 2, pt.t.size()) // list all distinct palindromes
 cout << i-1 << ") ":
 repx(j, pt.t[i].i, pt.t[i].i + pt.t[i].len)
  cout << s[i];
 cout << " " << pt.t[i].cnt << endl:</pre>
return 0;
```

1.6 prefix-function

```
return pi;
vector<vector<int>> aut:
void compute_automaton(string s) {
   s += '#':
   int n = s.size();
   vector<int> pi = prefix_function(s);
   aut.assign(n, vector<int>(26));
   rep(i, n) {
      rep(c, 26) {
          int i = i:
          while (j > 0 \&\& 'a' + c != s[j])
              j = pi[j-1];
          if ('a' + c == s[i])
              j++;
          aut[i][c] = j;
// k = n - pi[n - 1]; if k divides n, then the string can be
// aprtitioned into blocks of length k otherwise there is no
// effective compression and the answer is n.
```

1.7 suffix-array-martin

```
// build the suffix array
// suffixes are sorted, with each suffix represented by its
// starting position
vector<int> suffixarray(const string &s) {
   int N = s.size() + 1;//optional: include terminating NUL
   vector\langle int \rangle p(N), p2(N), c(N), c2(N), cnt(256);
   rep(i, N) cnt[s[i]] += 1;
   repx(b, 1, 256) cnt[b] += cnt[b - 1];
   rep(i, N) p[--cnt[s[i]]] = i;
   repx(i, 1, N) c[p[i]] = c[p[i - 1]] + (s[p[i]] != s[p[i - 1]])
         1]]):
   for (int k = 1; k < N; k <<= 1) {</pre>
       int C = c[p[N - 1]] + 1;
       cnt.assign(C + 1, 0);
       for (int &pi : p) pi = (pi - k + N) % N;
       for (int cl : c) cnt[cl + 1] += 1;
       rep(i, C) cnt[i + 1] += cnt[i];
       rep(i, N) p2[cnt[c[p[i]]]++] = p[i];
       c2[p2[0]] = 0;
       repx(i, 1, N) c2[p2[i]] =
           c2[p2[i-1]] + (c[p2[i]] != c[p2[i-1]] ||
                          c[(p2[i] + k) \% N] != c[(p2[i - 1]
                                + k) % N]);
```

```
swap(c, c2), swap(p, p2);
   p.erase(p.begin()); // optional: erase terminating NUL
   return p:
// build the lcp
// 'lcp[i]' represents the length of the longest common
// prefix between suffix i and suffix i+1 in the suffix
//array 'p'. the last element of 'lcp' is zero by convention
vector<int> makelcp(const string &s, const vector<int> &p) {
   int N = p.size(), k = 0:
   vector<int> r(N), lcp(N);
   rep(i, N) r[p[i]] = i;
   rep(i, N) {
       if (r[i] + 1 \ge N) \{ k = 0; continue; \}
       int j = p[r[i] + 1];
       while (i + k < N \&\& i + k < N \&\& s[i + k] == s[i + k]
           ]) k += 1;
       lcp[r[i]] = k:
       if (k) k -= 1:
   return lcp;
// lexicographically compare the suffixes starting from 'i'
// and 'j', considering only up to 'K' characters.
// 'r' is the inverse suffix array, mapping suffix offsets
// to indices. requires an LCP sparse table.
int lcp cmp(vector<int> &r. Sparse<int> &lcp. int i. int i.
    int K) {
   if (i == i) return 0:
   int ii = r[i], jj = r[j];
   int 1 = lcp.query(min(ii, jj), max(ii, jj));
   if (1 >= K) return 0:
   return ii < j; ? -1 : 1;
```

1.8 suffix-automaton

```
struct SuffixAutomaton {
    // edges[i]: the labeled edges from node i
    vector<map<char,int>> edges;
    // link[i]: the suffix link of i
    vector<int> link;
    // length[i]: len of the longest string in the ith class
    vector<int> length;
    // cnt[i]: # occurrences of each string in the ith class
    vector<int> cnt;
    // paths[i]: # of paths on the automaton starting from i
    vector<int> paths;
```

```
// terminal[i]: true if i is a terminal state
   vector<bool> terminal:
   vector<int> first_pos;
   vector<int> last_pos;
   //the index of the equivalence class of the whole string
   int last:
   SuffixAutomaton(string s) {
      edges.push_back(map<char,int>());
      link.push_back(-1);
      length.push_back(0);
      last = 0:
      rep(i, s.size()) { // construct r
          edges.push_back(map<char,int>());
          length.push_back(i+1);
          link.push back(0):
          int r = edges.size() - 1;
          int p = last:
          // add edges to r and find p with link to q
          while(p >= 0 && !edges[p].count(s[i])) {
             edges[p][s[i]] = r;
             p = link[p];
          if(p != -1) {
             int q = edges[p][s[i]];
             if(length[p] + 1 == length[q]) {
//we don't have to split q, just set the correct suffix link
                 link[r] = q;
             } else { // we have to split, add g'
                 // copy edges of q
                 edges.push_back(edges[q]);
                 length.push_back(length[p] + 1);
                 // copy parent of q
                 link.push_back(link[q]);
                 int ag = edges.size()-1:
                 // add qq as the new parent of q and r
                 link[q] = qq;
                 link[r] = qq;
          // move short classes polling to q to poll to q'
                 while(p >= 0 && edges[p][s[i]] == q) {
                     edges[p][s[i]] = qq;
                    p = link[p];
                 }
             }
          last = r;
   /* ----- */
      // mark terminal nodes
```

```
terminal.assign(edges.size(), false):
   int p = last:
   while(p > 0) {
       terminal[p] = true:
       p = link[p];
   // precompute match count
   cnt.assign(edges.size(), -1);
   cnt matches(0):
   //precompute # of paths (substr) starting from state
   paths.assign(edges.size(), -1);
   cnt paths(0):
   first_pos.assign(edges.size(), -1);
   get_first_pos(0);
   last_pos.assign(edges.size(), -1);
   get_last_pos(0);
int cnt_matches(int state) {
   if(cnt[state] != -1) return cnt[state]:
   int ans = terminal[state];
   for(auto edge : edges[state])
       ans += cnt_matches(edge.second);
   return cnt[state] = ans:
int cnt_paths(int state) {
   if(paths[state] != -1) return paths[state];
   // without repetition (counts different substrings)
   int ans = state == 0 ? 0 : 1;
   // with repetition
// int ans = state == 0 ? 0 : cnt[state];
   for(auto edge : edges[state])
       ans += cnt paths(edge.second):
   return paths[state] = ans;
int get_first_pos(int state) {
   if(first pos[state] != -1) return first pos[state]:
   int ans = 0:
   for(auto edge : edges[state])
       ans = max(ans, get_first_pos(edge.second)+1);
   return first_pos[state] = ans;
int get_last_pos(int state) {
   if(last_pos[state] != -1) return last_pos[state];
   int ans = terminal[state] ? 0 : INT_MAX;//fix
```

```
for(auto edge : edges[state])
          ans = min(ans, get_last_pos(edge.second)+1);
      return last_pos[state] = ans;
   string get k substring(int k) { // 0-indexed
       string ans;
       int state = 0;
       while(true){
          // without repetition (counts different substrs)
          int curr = state == 0 ? 0 : 1:
          // with repetition
      // int curr = state == 0 ? 0 : cnt[state];
          if(curr > k) return ans:
          k -= curr:
          for(auto edge : edges[state]) {
              if(paths[edge.second] <= k) {</pre>
                  k -= paths[edge.second];
              } else {
                  ans += edge.first;
                  state = edge.second;
                  break:
      }
   }
};
```

1.9 z-function

```
// i-th element is equal to the greatest number of
// characters starting from the position i that coincide
// with the first characters of s
vector<int> z_function(string s) {
   int n = s.size():
   vector<int> z(n):
   int 1 = 0, r = 0:
   for(int i = 1; i < n; i++) {</pre>
       if(i < r) z[i] = min(r - i, z[i - 1]);
       while(i + z[i] < n && s[z[i]] == s[i + z[i]])z[i]++:
       if(i + z[i] > r) {
          1 = i:
          r = i + z[i];
       }
   }
   return z;
```

2 dp

2.1 convex-hull-trick

```
struct Line {
    mutable 11 a. b. c:
    bool operator<(Line r) const { return a < r.a; }</pre>
    bool operator<(ll x) const { return c < x; }</pre>
};
// dynamically insert 'a*x + b' lines and query for maximum
// all operations have complexity O(log N)
struct LineContainer : multiset<Line, less<>>> {
    11 div(ll a, ll b) {
       return a / b - ((a ^ b) < 0 && a % b);
    bool isect(iterator x, iterator y) {
       if (y == end()) return x->c = INF, 0;
       if (x->a == y->a) x->c = x->b > y->b? INF : -INF;
       else x->c = div(y->b - x->b, x->a - y->a);
       return x->c >= y->c;
    void add(ll a. ll b) {
       // a *= -1, b *= -1 // for min
       auto z = insert(\{a, b, 0\}), y = z++, x = y;
       while (isect(y, z)) z = erase(z);
       if (x != begin() \&\& isect(--x, y)) isect(x, y = erase)
       while ((y = x) != begin() \&\& (--x)->c >= y->c) isect(
            x. erase(v)):
    11 querv(ll x) {
       if (empty()) return -INF; // INF for min
       auto 1 = *lower_bound(x);
       return 1.a * x + 1.b:
       // return -l.a * x - l.b; // for min
   }
};
```

2.2 divide-and-conquer

```
// for every index i assign an optimal index i, such that
// minimal for every i. the property that if i2 >= i1 then
    i2 >= i1 is
// exploited (monotonic condition).
// calculate optimal index for all indices in range [1, r)
    knowing that
// the optimal index for every index in this range is within
     [optl, optr).
// time: O(N log N)
void calc(vector<int> &opt, int 1, int r, int optl, int optr
   if (1 == r) return;
   int i = (1 + r) / 2:
   11 optc = INF;
   int optj;
   repx(j, optl, optr) {
      11 c = i + j; // cost(i, j)
       if (c < optc) optc = c, optj = j;</pre>
   opt[i] = optj;
   calc(opt, 1, i, optl, optj + 1);
   calc(opt, i + 1, r, optj, optr);
```

$3 \quad geo2d$

3.1 circle

```
struct C {
   P o; T r;

C(P o, T r) : o(o), r(r) {}
C() : C(P(), T()) {}

// intersects the circle with a line, assuming they intersect
// results are sorted with respect to the direction of the line
pair<P, P> line_inter(L 1) const {
   P c = 1.closest_to(o);
   T c2 = (c - o).magsq();
   P e = sqrt(max(r * r - c2, T())) * 1.d.unit();
   return {c - e, c + e};
}

// checks whether the given line collides with the circle
// negative: 2 intersections
```

```
// zero: 1 intersection
// positive: 0 intersections
T line_collide(L 1) const {
   T c2 = (1.closest_to(o) - o).magsq();
   return c2 - r * r;
// calculates the two intersections between two circles
// the circles must intersect in one or two points!
pair<P, P> inter(C h) const {
   P d = h.o - o:
   T c = (r * r - h.r * h.r) / d.magsq():
   return h.line_inter(\{(1 + c) / 2 * d, d.rot()\});
// check if the given circles intersect
bool collide(C h) const {
   return (h.o - o).magsq() \le (h.r + r) * (h.r + r);
// get one of the two tangents that cross through the
// the point must not be inside the circle
// a = -1: cw (relative to the circle) tangent
// a = 1: ccw (relative to the circle) tangent
P point_tangent(P p, T a) const {
   T c = r * r / p.magsq();
   return o + c * (p - o) - a * sqrt(c * (1 - c)) * (p -
         o).rot();
// get one of the 4 tangents between the two circles
// a = 1: exterior tangents
// a = -1: interior tangents (requires no area overlap)
// b = 1: ccw tangent
// b = -1: cw tangent
// the line origin is on this circumference, and the
    direction
// is a unit vector towards the other circle
L tangent(C c, T a, T b) const {
   T dr = a * r - c.r:
   P d = c.o - o;
   P n = (d * dr + b * d.rot() * sqrt(d.magsq() - dr *
        dr)).unit():
   return {o + n * r, -b * n.rot()};
// find the circumcircle of the given **non-degenerate**
static C thru_points(P a, P b, P c) {
```

```
L 1((a + b) / 2, (b - a).rot()):
      P p = 1.intersection(L((a + c) / 2, (c - a).rot()));
      return {p, (p - a).mag()};
   // find the two circles that go through the given point.
        are tangent
   // to the given line and have radius 'r'
   // the point-line distance must be at most 'r'!
   // the circles are sorted in the direction of the line
   static pair<C, C> thru_point_line_r(P a, L t, T r) {
      P d = t.d.rot().unit():
      if (d * (a - t.o) < 0) d = -d;
       auto p = C(a, r).line_inter(\{t.o + d * r, t.d\});
      return {{p.first, r}, {p.second, r}};
   // find the two circles that go through the given points
        and have
   // radius 'r'
   // the circles are sorted by angle with respect to the
        first point
   // the points must be at most at distance 'r'!
   static pair<C, C> thru_points_r(P a, P b, T r) {
       auto p = C(a, r).line_inter({(a + b) / 2, (b - a).rot}
           ()}):
      return {{p.first, r}, {p.second, r}};
};
```

3.2 convex-hull

3.3 delaunay

```
typedef __int128_t lll; // if on a 64-bit platform
struct 0 {
    Q *rot, *o; P p = {INF, INF}; bool mark;
   P &F() { return r()->p; }
    Q *&r() { return rot->rot; }
   Q *prev() { return rot->o->rot; }
    Q *next() { return r()->prev(); }
T cross(P a, P b, P c) { return (b - a) % (c - a): }
bool circ(P p, P a, P b, P c) { // is p in the circumcircle?
   111 p2 = p.magsq(), A = a.magsq() - p2.
       B = b.magsq() - p2, C = c.magsq() - p2;
   return cross(p, a, b) * C + cross(p, b, c) * A + cross(p, b, c)
         c. a) * B > 0:
Q *makeEdge(Q *&H, P orig, P dest) {
   Q *r = H ? H : new Q{new Q{new Q{0}}};
   H = r -> 0: r -> r() -> r() = r:
   repx(i, 0, 4) r = r->rot, r->p = {INF, INF},
       r->0 = i & 1 ? r : r->r():
   r\rightarrow p = orig; r\rightarrow F() = dest;
   return r;
void splice(0 *a, 0 *b) {
    swap(a->o->rot->o, b->o->rot->o); swap(a->o, b->o);
Q *connect(Q *&H, Q *a, Q *b) {
   Q *q = makeEdge(H, a->F(), b->p);
    splice(q, a->next()): splice(q->r(), b): return q:
pair<Q *, Q *> rec(Q *&H, const vector<P> &s) {
  if (s.size() <= 3) {</pre>
```

```
Q *a = makeEdge(H, s[0], s[1]), *b = makeEdge(H, s[1])
            [1]. s.back()):
       if (s.size() == 2) return \{a, a->r()\}; splice(a->r(),
       auto side = cross(s[0], s[1], s[2]);
       Q *c = side ? connect(H, b, a) : 0:
       return {side < 0 ? c->r() : a, side < 0 ? c : b->r()
            }:
   7
#define J(e) e \rightarrow F(), e \rightarrow p
#define valid(e) (cross(e->F(), J(base)) > 0)
   Q *A, *B, *ra, *rb; int half = s.size() / 2;
   tie(ra, A) = rec(H, {s.begin(), s.end() - half});
   tie(B, rb) = rec(H, {s.begin() + s.size() - half, s.end()
   while ((cross(B->p, J(A)) < 0 \&\& (A = A->next())) | |
          (cross(A->p, J(B)) > 0 \&\& (B = B->r()->o)));
   Q *base = connect(H, B->r(), A):
   if (A->p == ra->p) ra = base->r():
   if (B->p == rb->p) rb = base;
#define DEL(e, init, dir) Q *e = init->dir; \
   if (valid(e)) while (circ(e->dir->F(), J(base), e->F()))
           Q *t = e->dir; splice(e, e->prev()); \
           splice(e->r(), e->r()->prev()); e->o = H; H = e;
                e = t: \
       }
   for (::) {
       DEL(LC, base->r(), o); DEL(RC, base, prev());
       if (!valid(LC) && !valid(RC)) break;
       if (!valid(LC) || (valid(RC) && circ(J(RC), J(LC))))
            base = connect(H, RC, base->r());
       else base = connect(H, base->r(), LC->r());
   return {ra. rb}:
#undef J
#undef valid
#undef DEL
// there must be no duplicate points
// returns no triangles in the case of all collinear points
// produces counter-clockwise triangles ordered in triples
// maximizes the minimum angle across all triangulations
// the euclidean mst is a subset of these edges
// O(N log N)
vector<P> triangulate(vector<P> pts) {
   sort(pts.begin(), pts.end(), [](P a, P b) {
```

```
return make pair(a.x. a.v) < make pair(b.x. b.v):
   assert(unique(pts.begin(), pts.end()) == pts.end());
   if (pts.size() < 2) return {};</pre>
   Q *H = 0; Q *e = rec(H, pts).first;
   vector < Q *> q = \{e\}; int qi = 0;
   while (cross(e->o->F(), e->F(), e->p) < 0) e = e->o;
#define ADD
   {
       Q *c = e;
          c->mark = 1: pts.push back(c->p): \
          q.push_back(c->r()); c = c->next(); \
      } while (c != e):
   ADD;
   pts.clear();
   while (qi < (int)q.size()) if (!(e = q[qi++])->mark) ADD;
   return pts:
#undef ADD
```

3.4 halfplane-intersect

```
// obtain the convex polygon that results from intersecting
     the given list
// of halfplanes, represented as lines that allow their left
// assumes the halfplane intersection is bounded
vector<P> halfplane_intersect(vector<L> &H) {
   L bb(P(-INF, -INF), P(INF, 0));
   rep(k, 4) H.push back(bb), bb.o = bb.o.rot(), bb.d = bb.d
        .rot();
   sort(begin(H), end(H), [](L a, L b) { return a.d.angcmp(b
         .d) < 0; \});
   deque < L > q: int n = 0:
   rep(i, H.size()) {
       while (n \ge 2 \&\& H[i].side(q[n - 1].intersection(q[n + 1]))
           -21)) > 0)
           q.pop_back(), n--;
       while (n \ge 2 \&\& H[i].side(q[0].intersection(q[1])) >
           q.pop_front(), n--;
       if (n > 0 && H[i].parallel(g[n - 1])) {
           if (H[i].d * q[n - 1].d < 0) return {};
           if (H[i].side(q[n-1].o) > 0) q.pop_back(), n--;
           else continue:
```

```
q.push_back(H[i]), n++;
}

while (n >= 3 && q[0].side(q[n - 1].intersection(q[n - 2])) > 0)
    q.pop_back(), n--;
while (n >= 3 && q[n - 1].side(q[0].intersection(q[1])) > 0)
    q.pop_front(), n--;
if (n < 3) return {};

vector<P> ps(n);
rep(i, n) ps[i] = q[i].intersection(q[(i + 1) % n]);
return ps;
}
```

3.5 line

```
// a segment or an infinite line
// does not handle point segments correctly!
struct L {
   P o. d:
   L() : o(), d() {}
   L(P o, P d) : o(o), d(d) {}
   L(P ab, T c) : d(ab.rot()), o(ab * -c / ab.magsq()) {}
   pair<P, T> line_eq() { return {-d.rot(), d.rot() * o}; }
   // returns a number indicating which side of the line the
         point is in
   // negative: left, positive: right
   T side(P r) const { return (r - o) % d: }
   // returns the intersection coefficient
   // in the range [0, d % r.d]
   // if d % r.d is zero, the lines are parallel
   T inter(L r) const { return (r.o - o) % r.d: }
   // get the single intersection point
   // lines must not be parallel
   P intersection(L r) const { return o + d * inter(r) / (d
        % r.d): }
   // check if lines are parallel
   bool parallel(L r) const { return abs(d % r.d) <= EPS: }</pre>
   // check if segments intersect
   bool seg_collide(L r) const {
      Tz = d \% r.d;
```

```
if (abs(z) <= EPS) {
       if (abs(side(r.o)) > EPS) return false;
       T s = (r.o - o) * d, e = s + r.d * d;
       if (s > e) swap(s, e);
       return s <= d * d + EPS && e >= -EPS;
   T s = inter(r), t = -r.inter(*this);
   if (z < 0) s = -s. t = -t. z = -z:
   return s >= -EPS && s <= z + EPS && t >= -EPS && t <=
         z + EPS:
// full segment intersection
// produces a point segment if the intersection is a
// however it **does not** handle point segments as input
bool seg_inter(L r, L *out) const {
   Tz = d \% r.d:
   if (abs(z) <= EPS) {
       if (abs(side(r.o)) > EPS) return false;
       if (r.d * d < 0) r = \{r.o + r.d, -r.d\};
       P s = o * d < r.o * d ? r.o : o;
       P = (o + d) * d < (r.o + r.d) * d ? o + d : r.o
       if (s * d > e * d) return false;
       return *out = {s, e - s}, true;
   T s = inter(r), t = -r.inter(*this);
   if (z < 0) s = -s, t = -t, z = -z;
   if (s \ge -EPS \&\& s \le z + EPS \&\& t \ge -EPS \&\& t \le z
       return *out = {o + d * s / z, {0, 0}}, true:
   return false;
// check if the given point is on the segment
bool point_on_seg(P r) const {
   if (abs(side(r)) > EPS) return false;
   if ((r - o) * d < -EPS) return false;</pre>
   if ((r - o - d) * d > EPS) return false:
   return true:
// get the point in this line that is closest to a given
P closest_to(P r) const {
   P dr = d.rot(): return r + (o - r) * dr * dr / d.
        magsq();
```

```
};
```

3.6 minkowski

```
void reorder_polygon(vector<P> &ps) {
   int pos = 0;
   repx(i, 1, (int)ps.size()) {
       if (ps[i].y < ps[pos].y || (ps[i].y == ps[pos].y &&</pre>
            ps[i].x < ps[pos].x)
           pos = i:
   rotate(ps.begin(), ps.begin() + pos, ps.end());
}
vector<P> minkowski(vector<P> ps, vector<P> qs) {
   // the first vertex must be the lowest
   reorder_polygon(ps); reorder_polygon(qs);
   ps.push_back(ps[0]); ps.push_back(ps[1]);
   qs.push_back(qs[0]); qs.push_back(qs[1]);
   vector<P> result; int i = 0, j = 0;
   while (i < ps.size() - 2 || j < qs.size() - 2) {</pre>
       result.push_back(ps[i] + qs[j]);
       auto z = (ps[i + 1] - ps[i]) \% (qs[j + 1] - qs[j]);
       if (z \ge 0 \&\& i < ps.size() - 2) ++i;
       if (z <= 0 && i < as.size() - 2) ++i:
   return result:
```

3.7 point

```
struct P {
    T x, y;
    P(T x, T y) : x(x), y(y) {}
    P() : P(0, 0) {}

    friend ostream & operator << (ostream & s, const P & r) {
        return s << r.x << " " << r.y;
    }
    friend istream & operator >> (istream & s, P & r) { return s
        >> r.x >> r.y; }

P operator + (P r) const { return {x + r.x, y + r.y}; }
    P operator - (P r) const { return {x - r.x, y - r.y}; }
    P operator / (T r) const { return {x * r, y * r}; }
    P operator - () const { return {x / r, y / r}; }
    P operator - () const { return {x / r, y / r}; }
}
```

```
friend P operator*(T 1, P r) { return {1 * r.x, 1 * r.v};
P rot() const { return {-y, x}; }
T operator*(P r) const { return x * r.x + y * r.y; }
T operator%(P r) const { return rot() * r: }
T magsq() const { return x * x + y * y; }
T mag() const { return sqrt(magsq()); }
P unit() const { return *this / mag(); }
bool half() const { return abs(v) <= EPS && x < -EPS | | v | struct InConvex {
     < -EPS: }
T angcmp(P r) const {
   int h = (int)half() - r.half();
   return h ? h : r % *this;
bool operator==(P r) const { return abs(x - r.x) <= EPS</pre>
    && abs(v - r.v) <= EPS: }
double angle() const { return atan2(y, x); }
static P from_angle(double a) { return {cos(a), sin(a)};
```

3.8 polygon

```
// get the area of a simple polygon in ccw order
T area(const vector<P> &ps) {
   int N = ps.size();
   rep(i, N) a += (ps[i] - ps[0]) % (ps[(i + 1) % N] - ps[i]
        1):
   return a / 2;
// checks whether a point is inside a simple polygon
// returns -1 if inside, 0 if on border, 1 if outside
int in_poly(const vector<P> &ps, P p) {
   int N = ps.size(), w = 0;
   rep(i, N) {
       P = ps[i] - p, e = ps[(i + 1) % N] - p;
       if (s == P()) return 0:
       if (s.y == 0 && e.y == 0) {
           if (\min(s.x. e.x) \le 0 \&\& 0 \le \max(s.x. e.x))
               return 0:
       } else {
```

```
bool b = s.v < 0:
          if (b != (e.v < 0)) {
             T z = s \% e; if (z == 0) return 0;
             if (b == (z > 0)) w += b ? 1 : -1:
   return w ? -1 : 1;
// check if a point is in a convex polygon
   vector<P> ps;
   T 11, 1h, rl, rh;
   int N. m:
   // preprocess polygon
   // O(N)
   InConvex(const vector<P> &p) : ps(p), N(ps.size()), m(0)
       assert(N >= 2);
      rep(i, N) if (ps[i].x < ps[m].x) m = i;
      rotate(ps.begin(), ps.begin() + m, ps.end());
      rep(i, N) if (ps[i].x > ps[m].x) m = i;
      11 = 1h = ps[0].y, r1 = rh = ps[m].y;
      for (P p : ps) {
          if (p.x == ps[0].x) 11 = min(11, p.y), 1h = max(
               lh. p.v):
          if (p.x == ps[m].x) rl = min(rl, p.y), rh = max(
      }
   InConvex() {}
   // check if point belongs in polygon
   // returns -1 if inside. 0 if on border. 1 if outside
   // O(log N)
   int in_poly(P p) {
      if (p.x < ps[0].x || p.x > ps[m].x) return 1;
      if (p.x == ps[0].x) return p.y < 11 || p.y > 1h;
      if (p.x == ps[m].x) return p.v < rl || p.v > rh:
      int r = upper_bound(ps.begin(), ps.begin() + m, p,
          [](P a, P b) { return a.x < b.x; }) - ps.begin();
      Tz = (ps[r-1] - ps[r]) \% (p - ps[r]); if (z >= 0)
           return !!z;
      r = upper_bound(ps.begin() + m, ps.end(), p,
          [](P a, P b) { return a.x > b.x; }) - ps.begin();
      z = (ps[r - 1] - ps[r \% N]) \% (p - ps[r \% N]);
      if (z >= 0) return !!z; return -1;
```

3.9 sweep

```
#include "point.cpp"
// iterate over all pairs of points
// 'op' is called with all ordered pairs of different
    indices '(i, j)'
// additionally, the 'ps' vector is kept sorted by signed
    distance
// to the line formed by 'i' and 'j'
// for example, if the vector from 'i' to 'j' is pointing
// the 'ps' vector is sorted from smallest 'y' to largest 'y
// note that, because the 'ps' vector is sorted by signed
// 'j' is always equal to 'i + 1'
// this means that the amount of points to the left of the
    line is always 'N - i'
template <class OP>
void all_pair_points(vector<P> &ps, OP op) {
   int N = ps.size();
   sort(ps.begin(), ps.end(), [](P a, P b) {
      return make_pair(a.y, a.x) < make_pair(b.y, b.x);</pre>
   }):
   vector<pair<int, int>> ss;
   rep(i, N) rep(j, N) if (i != j) ss.push_back({i, j});
   stable_sort(ss.begin(), ss.end(), [&](auto a, auto b) {
       return (ps[a.second] - ps[a.first]).angle lt(ps[b.
            second] - ps[b.first]);
   vector<int> p(N); rep(i, N) p[i] = i;
   for (auto [i, j] : ss)
```

3.10 theorems

```
// Pick's theorem
// Simple polygon with integer vertices:
// A = I + B / 2 - 1
// A: Area of the polygon
// I: Integer points strictly inside the polygon
// B: Integer points on the boundary of the polygon
```

4 graph

4.1 bellman-ford

```
struct Edge { int u, v; ll w; };
// find distance from source node to all nodes.
// supports negative edge weights.
// returns true if a negative cycle is detected.
11
// time: O(V E)
bool bellman_ford(int N, int s, vector<Edge> &E, vector<11>
     &D. vector<int> &P) {
   P.assign(N, -1), D.assign(N, INF), D[s] = 0;
   rep(i. N - 1) {
       bool f = true:
       rep(ei, E.size()) {
           auto &e = E[ei]:
           ll n = D[e.u] + e.w;
           if (D[e.u] < INF && n < D[e.v])</pre>
              D[e.v] = n, P[e.v] = ei, f = false;
       if (f) return false;
   }
   return true;
```

4.2 dinic

```
struct Edge { int u, v; ll c, f = 0; };
```

```
// maximum flow algorithm.
// time: O(E V^2)
11
        O(E V^{(2/3)}) / O(E sqrt(E)) unit capacities
11
        O(E sart(V))
                                    unit networks (hopcroft-
    karp)
// unit network: c in {0, 1} and forall v, len(incoming(v))
    <= 1 or len(outgoing(v)) <= 1
// min-cut: find all nodes reachable from the source in the
    residual graph
struct Dinic {
   int N. s. t: vector<vector<int>> G:
   vector<Edge> E: vector<int> lvl. ptr:
   Dinic() {}
   Dinic(int N, int s, int t) : N(N), s(s), t(t), G(N) {}
   void add_edge(int u, int v, ll c) {
       G[u].push_back(E.size()); E.push_back({u, v, c});
       G[v].push_back(E.size()); E.push_back({v, u, 0});
   11 push(int u, 11 p) {
       if (u == t || p <= 0) return p;</pre>
       while (ptr[u] < G[u].size()) {</pre>
           int ei = G[u][ptr[u]++];
           Edge &e = E[ei];
           if (lvl[e.v] != lvl[u] + 1) continue;
           11 a = push(e.v, min(e.c - e.f, p));
           if (a <= 0) continue: e.f += a. E[ei ^ 1].f -= a:
           return a;
       }
       return 0;
   11 maxflow() {
       11 f = 0:
       while (true) {
           // bfs to build levels
           lvl.assign(N, -1); queue<int> q; lvl[s] = 0, q.
                push(s);
           while (!q.empty()) {
              int u = q.front(); q.pop();
              for (int ei : G[u]) {
                  Edge &e = E[ei]:
                  if (e.c - e.f <= 0 || lvl[e.v] != -1)</pre>
                       continue;
                  lvl[e.v] = lvl[u] + 1, q.push(e.v);
              }
           if (lvl[t] == -1) break:
```

4.3 floyd-warshall

4.4 heavy-light

```
struct Hld {
   vector<int> P, H, D, pos, top;
   Hld() {}
   void init(vector<vector<int>> &G) {
      int N = G.size();
      P.resize(N), H.resize(N), D.resize(N), pos.resize(N),
          top.resize(N);
      D[0] = -1, dfs(G, 0); int t = 0;
      rep(i, N) if (H[P[i]] != i) {
          int j = i;
          while (i != -1)
              \{ top[j] = i, pos[j] = t++; j = H[j]; \}
   }
   int dfs(vector<vector<int>> &G, int i) {
      int w = 1, mw = 0:
      D[i] = D[P[i]] + 1, H[i] = -1;
      for (int c : G[i]) {
          if (c == P[i]) continue;
          P[c] = i; int sw = dfs(G, c); w += sw;
```

```
if (sw > mw) H[i] = c. mw = sw:
   }
   return w;
// visit the log N segments in the path from u to v
template <class OP>
void path(int u, int v, OP op) {
   while (top[u] != top[v]) {
       if (D[top[u]] > D[top[v]]) swap(u, v);
       op(pos[top[v]], pos[v] + 1); v = P[top[v]];
   if (D[u] > D[v]) swap(u, v);
   op(pos[u], pos[v] + 1); // value on node
   // op(pos[u]+1, pos[v] + 1); // value on edge
// an alternative to 'path' that considers order.
// calls 'op' with an 'l <= r' inclusive-exclusive range.
// boolean indicating if the query is forwards or
    backwards.
template <class OP>
void path(int u, int v, OP op) {
   int lu = u, lv = v;
   while (top[lu] != top[lv])
       if (D[top[lu]] > D[top[lv]]) lu = P[top[lu]];
       else lv = P[top[lv]]:
   int lca = D[lu] > D[lv] ? lv : lu;
   while (top[u] != top[lca])
       op(pos[top[u]], pos[u] + 1, false), u = P[top[u
           11:
   if (u != lca) op(pos[lca] + 1, pos[u] + 1, false);
   vector<int> stk:
   while (top[v] != top[lca])
       stk.push_back(v), v = P[top[v]];
   // op(pos[lca], pos[v] + 1, true); // value on node
   op(pos[lca] + 1, pos[v] + 1, true); // value on edge
   reverse(stk.begin(), stk.end());
   for (int w : stk) op(pos[top[w]], pos[w] + 1, true);
// commutative segment tree
template <class T, class S>
void update(S &seg, int i, T val) { seg.update(pos[i],
    val): }
```

```
// commutative segment tree lazy
template <class T, class S>
void update(S &seg, int u, int v, T val) {
    path(u, v, [&](int l, int r) { seg.update(l, r, val);
        });
}

// commutative (lazy) segment tree
template <class T, class S>
T query(S &seg, int u, int v) {
    T ans = 0;
        // neutral element
    path(u, v, [&](int l, int r) { ans += seg.query(l, r)
        ; }); // query op
    return ans;
}
};
```

4.5 hungarian

```
// find a maximum gain perfect matching in the given
    bipartite complete graph.
// input: gain matrix (G_{xy}) = benefit of joining vertex x
    in set X with vertex
// y in set Y).
// output: maximum gain matching in members 'xy[x]' and 'yx[
    v]'.
// runtime: O(N^3)
struct Hungarian {
   int N, qi, root;
   vector<vector<ll>>> gain;
   vector<int> xy, yx, p, q, slackx;
   vector<ll> lx, ly, slack;
   vector<bool> S. T:
   void add(int x, int px) {
      S[x] = true, p[x] = px;
       rep(y, N) if (lx[x] + ly[y] - gain[x][y] < slack[y])
          slack[y] = lx[x] + ly[y] - gain[x][y], slackx[y]
               = x:
   }
   void augment(int x, int v) {
       while (x != -2) {
          yx[y] = x; swap(xy[x], y); x = p[x];
      }
```

```
void improve() {
   S.assign(N, false), T.assign(N, false), p.assign(N,
        -1):
   qi = 0, q.clear();
   rep(x, N) if (xv[x] == -1) {
       q.push_back(root = x), p[x] = -2, S[x] = true;
       break:
   rep(y, N) slack[y] = lx[root] + ly[y] - gain[root][y
        1. slackx[v] = root:
   while (true) {
       while (qi < q.size()) {</pre>
          int x = q[qi++];
          rep(v, N) if (lx[x] + lv[v] == gain[x][v] &&!
              if (yx[y] == -1) return augment(x, y);
              T[v] = true, g.push back(vx[v]), add(vx[v]
       }
       11 d = INF:
       rep(y, N) if (!T[y]) d = min(d, slack[y]);
       rep(x, N) if (S[x]) lx[x] -= d;
       rep(y, N) if (T[y]) ly[y] += d;
       rep(v, N) if (!T[v]) slack[v] -= d:
       rep(y, N) if (!T[y] && slack[y] == 0) {
          if (yx[y] == -1) return augment(slackx[y], y);
          T[v] = true;
          if (!S[yx[y]]) q.push_back(yx[y]), add(yx[y],
               slackx[v]);
       }
   }
Hungarian(vector<vector<11>> g)
   : N(g.size()), gain(g), xy(N, -1), yx(N, -1), lx(N, -1)
        INF).
   ly(N), slack(N), slackx(N) {
   rep(x, N) rep(y, N) lx[x] = max(lx[x], ly[y]);
   rep(i, N) improve();
```

```
kuhn
```

};

```
// get a maximum cardinality matching in a bipartite graph.
// input: adjacency lists.
// output: matching (in 'mt' member).
// runtime: O(V E)
struct Kuhn {
   int N. size:
   vector<vector<int>> G;
   vector<bool> seen:
   vector<int> mt;
   bool visit(int i) {
       if (seen[i]) return false;
       seen[i] = true:
       for (int to : G[i])
          if (mt[to] == -1 || visit(mt[to])) {
              mt[to] = i:
              return true;
       return false;
   Kuhn(vector<vector<int>> adj) : G(adj), N(G.size()), mt(N | 4.8 maxflow-mincost
        . -1) {
       rep(i, N) {
          seen.assign(N, false);
          size += visit(i);
      }
   }
};
```

4.7 lca

```
// calculates the lowest common ancestor for any two nodes
    in O(log N) time,
// with O(N log N) preprocessing
struct Lca {
   int N, K, t = 0;
   vector<vector<int>> U:
   vector<int> L. R:
   Lca() {}
   Lca(vector<vector<int>> &G) : N(G.size()), L(N), R(N) {
      K = N \le 1 ? 0 : 32 - \_builtin_clz(N - 1);
      U.resize(K + 1, vector<int>(N));
      visit(G, 0, 0);
      rep(k, K) rep(u, N) U[k + 1][u] = U[k][U[k][u]];
```

```
void visit(vector<vector<int>> &G. int u. int p) {
      L[u] = t++, U[0][u] = p;
      for (int v : G[u]) if (v != p) visit(G, v, u);
      R[u] = t++:
   }
   bool is_anc(int up, int dn) {
       return L[up] <= L[dn] && R[dn] <= R[up];
   int find(int u, int v) {
       if (is anc(u, v)) return u:
       if (is_anc(v, u)) return v;
      for (int k = K; k \ge 0;)
          if (is_anc(U[k][u], v)) k--;
          else u = U[k][u];
      return U[0][u]:
   }
};
```

```
// time: O(F V E)
                          F is the maximum flow
       O(V E + F E log V) if bellman-ford is replaced by
    johnson
struct Flow {
   struct Edge {
       int u, v;
       11 c, w, f = 0;
   int N. s. t:
   vector<vector<int>> G;
   vector<Edge> E;
   vector<ll> d, b;
   vector<int> p;
   Flow() {}
   Flow(int N, int s, int t): N(N), s(s), t(t), G(N) {}
   void add_edge(int u, int v, ll c, ll w) {
      G[u].push_back(E.size());
       E.push_back({u, v, c, w});
       G[v].push_back(E.size());
       E.push back({v. u. 0. -w}):
   // naive distances with bellman-ford: O(V E)
   void calcdists() {
```

```
p.assign(N, -1), d.assign(N, INF), d[s] = 0;
   rep(i, N - 1) rep(ei, E.size()) {
       Edge &e = E[ei];
       ll n = d[e.u] + e.w:
       if (d[e.u] < INF && e.c - e.f > 0 && n < d[e.v])
            d[e.v] = n, p[e.v] = ei:
}
// johnsons potentials: O(E log V)
void calcdists() {
   if (b.emptv()) {
       b.assign(N, 0);
       // code below only necessary if there are
           negative costs
       rep(i, N - 1) rep(ei, E.size()) {
          Edge &e = E[ei]:
          if (e.f < e.c) b[e.v] = min(b[e.v], b[e.u] + e
       }
   p.assign(N, -1), d.assign(N, INF), d[s] = 0;
   priority_queue<pair<11, int>> q;
   q.push({0, s});
   while (!q.empty()) {
       auto [w, u] = q.top();
       q.pop();
       if (d[u] < -w + b[u]) continue:
       for (int ei : G[u]) {
          auto e = E[ei]:
          ll n = d[u] + e.w;
          if (e.f < e.c && n < d[e.v]) {</pre>
              d[e.v] = n, p[e.v] = ei;
              q.push({b[e.v] - n, e.v});
       }
   b = d:
ll solve() {
   b.clear();
   11 ff = 0:
   while (true) {
       calcdists();
       if (p[t] == -1) break:
       11 f = INF:
       for (int cur = t; p[cur] != -1; cur = E[p[cur]].u
```

4.9 push-relabel

```
#include "../common.h"
const ll INF = 1e18;
// maximum flow algorithm.
// to run, use 'maxflow()'.
// time: O(V^2 \text{ sqrt}(E)) \leq O(V^3)
// memory: 0(V^2)
struct PushRelabel {
    vector<vector<ll>> cap. flow:
    vector<ll> excess;
    vector<int> height;
    PushRelabel() {}
    void resize(int N) { cap.assign(N, vector<ll>(N)); }
    // push as much excess flow as possible from u to v.
    void push(int u, int v) {
       11 f = min(excess[u], cap[u][v] - flow[u][v]);
       flow[u][v] += f:
       flow[v][u] -= f;
       excess[v] += f;
       excess[u] -= f:
    }
    // relabel the height of a vertex so that excess flow may
          be pushed.
    void relabel(int u) {
       int d = INT32_MAX;
       rep(v, cap.size()) if (cap[u][v] - flow[u][v] > 0) d
           min(d, height[v]);
       if (d < INF) height[u] = d + 1;</pre>
```

```
// get the maximum flow on the network specified by 'cap'
     with source 's'
// and sink 't'.
// node-to-node flows are output to the 'flow' member.
11 maxflow(int s, int t) {
   int N = cap.size(), M:
   flow.assign(N, vector<11>(N));
   height.assign(N, 0), height[s] = N;
    excess.assign(N, 0), excess[s] = INF;
   rep(i, N) if (i != s) push(s, i);
   vector<int> a:
    while (true) {
       // find the highest vertices with excess
       q.clear(), M = 0;
       rep(i, N) {
           if (excess[i] <= 0 || i == s || i == t)</pre>
                continue:
           if (height[i] > M) q.clear(), M = height[i];
           if (height[i] >= M) q.push_back(i);
       if (q.empty()) break;
       // process vertices
       for (int u : q) {
           bool relab = true;
           rep(v, N) {
              if (excess[u] <= 0) break;</pre>
              if (cap[u][v] - flow[u][v] > 0 && height[u]
                   ] > height[v])
                  push(u, v), relab = false:
           if (relab) {
              relabel(u):
              break;
   }
   11 f = 0; rep(i, N) f += flow[i][t]; return f;
```

4.10 strongly-connected-components

};

```
// compute strongly connected components.
// time: O(V + E), memory: O(V)
//
// after building:
```

```
// comp = map from vertex to component (components are
     toposorted, root first, leaf last)
// N = number of components
// G = condensation graph (component DAG)
// byproducts:
// vgi = transposed graph
// order = reverse topological sort (leaf first, root last)
// others:
// vn = number of vertices
// vg = original vertex graph
struct Scc {
   int vn, N;
   vector<int> order, comp;
   vector<vector<int>> vg, vgi, G;
   void toposort(int u) {
       if (comp[u]) return:
       comp[u] = -1:
       for (int v : vg[u]) toposort(v);
       order.push_back(u);
   bool carve(int u) {
       if (comp[u] != -1) return false;
       comp[u] = N;
       for (int v : vgi[u]) {
          carve(v);
          if (comp[v] != N) G[comp[v]].push back(N):
       }
       return true;
   Scc() {}
   Scc(vector<vector<int>> &g) : vn(g.size()), vg(g), comp(
        vn), vgi(vn), G(vn), N(0) {
       rep(u, vn) toposort(u);
       rep(u, vn) for (int v : vg[u]) vgi[v].push_back(u);
       invrep(i, vn) N += carve(order[i]):
};
```

4.11 two-sat

```
// calculate the solvability of a system of logical
    equations, where every equation is of the form 'a or b
    '.
// 'neg': get negation of 'u'
```

```
'then': 'u' implies 'v'
// 'anv': 'u' or 'v'
// 'set': 'u' is true
// after 'solve' (O(V+E)) returns true, 'sol' contains one
    possible solution.
// determining all solutions is O(V*E) hard (requires
    computing reachability in a DAG).
struct TwoSat {
   int N; vector<vector<int>> G;
   Scc scc: vector<bool> sol:
   TwoSat(int n): N(n), G(2 * n), sol(n) {}
   TwoSat() {}
   int neg(int u) { return (u + N) \% (2 * N); }
   void then(int u, int v) { G[u].push_back(v), G[neg(v)].
        push back(neg(u)): }
   void any(int u, int v) { then(neg(u), v); }
   void set(int u) { G[neg(u)].push back(u): }
   bool solve() {
      scc = Scc(G):
      rep(u, N) if (scc.comp[u] == scc.comp[neg(u)]) return
       rep(u, N) sol[u] = (scc.comp[u] > scc.comp[neg(u)]);
       return true;
};
```

5 implementation

5.1 SegmentTreeBeats

```
if (a.mx1 == b.mx1) mx1 = a.mx1, mxc = a.mxc + b.mxc
             mx2 = max(a.mx2. b.mx2):
       // max
       if (a.mn1 < b.mn1) mn1 = a.mn1, mnc = a.mnc, mn2 =
            min(b.mn1, a.mn2);
       if (a.mn1 > b.mn1) mn1 = b.mn1, mnc = b.mnc, mn2 =
           min(a.mn1, b.mn2);
       if (a.mn1 == b.mn1) mn1 = a.mn1, mnc = a.mnc + b.mnc,
            mn2 = min(a.mn2, b.mn2):
   }
};
// 0 - indexed / inclusive - inclusive
template <class node>
struct STB {
   vector<node> st; int n;
   void build(int u, int i, int j, vector<node> &arr) {
       if (i == i) {
          st[u] = arr[i]:
          return;
       int m = (i + j) / 2, l = u * 2 + 1, r = u * 2 + 2;
       build(l, i, m, arr), build(r, m + 1, j, arr);
       st[u] = node(st[1], st[r]);
   void push_add(int u, int i, int j, ll v) {
       st[u].s += (i - i + 1) * v:
       st[u].mx1 += v, st[u].mn1 += v, st[u].lz += v;
       if (st[u].mx2 != LLONG MIN) st[u].mx2 += v:
       if (st[u].mn2 != LLONG_MAX) st[u].mn2 += v;
   void push_max(int u, ll v, bool l) { // for min op
       if (v >= st[u].mx1) return;
       st[u].s -= st[u].mx1 * st[u].mxc:
       st[u].mx1 = v:
       st[u].s += st[u].mx1 * st[u].mxc;
       if (1) st[u].mn1 = st[u].mx1:
       else if (v <= st[u].mn1) st[u].mn1 = v;</pre>
       else if (v < st[u].mn2) st[u].mn2 = v;
   void push_min(int u, ll v, bool l) { // for max op
       if (v <= st[u].mn1) return;</pre>
       st[u].s -= st[u].mn1 * st[u].mnc;
       st[u].mn1 = v:
       st[u].s += st[u].mn1 * st[u].mnc:
       if (1) st[u].mx1 = st[u].mn1;
       else if (v \ge st[u].mx1) st[u].mx1 = v:
       else if (v > st[u].mx2) st[u].mx2 = v:
```

```
void push(int u, int i, int i) {
   if (i == j) return;
   // add
   int m = (i + j) / 2, l = u * 2 + 1, r = u * 2 + 2;
   push_add(1, i, m, st[u].lz);
   push add(r. m + 1, i, st[u].lz):
   st[u].lz = 0;
   // min
   push_max(1, st[u].mx1, i == m);
   push_max(r, st[u].mx1, m + 1 == j);
   // max
   push min(l, st[u].mn1, i == m):
   push_min(r, st[u].mn1, m + 1 == r);
node query(int a, int b, int u, int i, int j) {
   if (b < i || j < a) return node();</pre>
   if (a <= i && j <= b) return st[u];</pre>
   push(u, i, j);
   int m = (i + j) / 2, l = u * 2 + 1, r = u * 2 + 2;
   return node(query(a, b, 1, i, m), query(a, b, r, m +
        1, i));
void update_add(int a, int b, ll v, int u, int i, int j)
   if (b < i || j < a) return;</pre>
   if (a <= i && j <= b) {</pre>
       push_add(u, i, j, v);
       return:
   push(u, i, i):
   int m = (i + j) / 2, 1 = u * 2 + 1, r = u * 2 + 2;
   update_add(a, b, v, l, i, m);
   update_add(a, b, v, r, m + 1, j);
   st[u] = node(st[1], st[r]);
void update min(int a, int b, ll v, int u, int i, int i)
   if (b < i || j < a || v >= st[u].mx1) return;
   if (a <= i && j <= b && v > st[u].mx2) {
       push_max(u, v, i == j);
       return:
   push(u, i, j);
   int m = (i + j) / 2, l = u * 2 + 1, r = u * 2 + 2;
   update_min(a, b, v, l, i, m);
   update min(a, b, v, r, m + 1, i):
   st[u] = node(st[1], st[r]);
void update_max(int a, int b, ll v, int u, int i, int j)
```

```
if (b < i || i < a || v <= st[u].mn1) return;</pre>
   if (a <= i && j <= b && v < st[u].mn2) {</pre>
       push_min(u, v, i == j);
       return:
   push(u, i, i):
   int m = (i + j) / 2, l = u * 2 + 1, r = u * 2 + 2;
   update_max(a, b, v, 1, i, m);
   update_max(a, b, v, r, m + 1, j);
   st[u] = node(st[1], st[r]);
}
STB(vector<node> &v, int N) : n(N), st(N * 4 + 5) { build
     (0, 0, n - 1, v): }
node query(int a, int b) { return query(a, b, 0, 0, n -
    1): }
void update add(int a, int b, ll v) { update add(a, b, v,
     0.0.n - 1): 
void update min(int a, int b, ll v) { update min(a, b, v.
     0.0.n - 1): 
void update_max(int a, int b, ll v) { update_max(a, b, v,
     0.0.n - 1):
```

5.2 Treap

```
#include "../Template.cpp"
mt19937 gen(chrono::high_resolution_clock::now().
     time_since_epoch().count());
// 101 Implicit Treap //
struct Node {
    int p, sz = 0, v, acc, l = -1, r = -1;
   Node() : v(0), acc(0) {}
   Node(int x) : p(gen()), sz(1), v(x), acc(x) {}
   void recalc(const Node &a, const Node &b) {
       sz = a.sz + b.sz + 1:
       acc = v + a.acc + b.acc:
};
template <class node>
struct Treap {
   vector<node> t;
   int n, r = -1;
   node get(int u) { return u != -1 ? t[u] : node(); }
```

```
void recalc(int u) { t[u].recalc(get(t[u].l), get(t[u].r)
   int merge(int 1, int r) {
       if (min(l, r) == -1) return l != -1 ? l : r:
       int ans = (t[1].p < t[r].p) ? 1 : r;
       if (ans == 1) t[1].r = merge(t[1].r, r), recalc(1):
       if (ans == r) t[r].1 = merge(1, t[r].1), recalc(r);
       return ans:
   pii split(int u, int id) {
       if (u == -1) return {-1, -1}:
       int szl = get(t[u].1).sz:
       if (szl >= id) {
          pii ans = split(t[u].1, id);
          t[u].1 = ans.ss:
          recalc(u);
          return {ans.ff, u}:
       pii ans = split(t[u].r. id - szl - 1):
       t[u].r = ans.ff:
       recalc(u);
       return {u, ans.ss};
   }
   Treap(vi &v) : n(sz(v)) {
      for (int i = 0; i < n; i++) t.eb(v[i]), r = merge(r, v[i])
   }
};
// Complete Implicit Treap with Lazy propagation //
struct Node {
   int p, sz = 0, v, acc, l = -1, r = -1, par = -1, lzv = 0;
   bool lz = false, f = false:
   Node(): v(0), acc(0) {}
   Node(int x) : p(gen()), sz(1), v(x), acc(x) {}
   void recalc(const Node &a, const Node &b) {
       sz = a.sz + b.sz + 1;
       acc = v + a.acc + b.acc:
   void upd_lazy(int x) { lz = 1, lzv += x; }
   void lazv() { v += lzv, acc += sz * lzv, lz = 0, lzv = 0;
   void flip() { swap(1, r), f = 0; }
template <class node>
struct Treap {
   vector<node> t:
```

```
int n. r = -1:
node get(int u) { return u != -1 ? t[u] : node(); }
void recalc(int u) {
   int 1 = t[u].1, r = t[u].r;
   push(1), push(r), flip(1), flip(r);
   t[u].recalc(get(1), get(r));
void push(int u) {
   if (u == -1 || !t[u].lz) return;
   int 1 = t[u].1, r = t[u].r:
   if (1 != -1) t[1].upd lazv(t[u].lzv):
   if (r != -1) t[r].upd_lazy(t[u].lzv);
   t[u].lazy();
void flip(int u) {
   if (u == -1 || !t[u].f) return:
   int 1 = t[u].1, r = t[u].r;
   if (1 != -1) t[1].f ^= 1:
   if (r != -1) t[r].f ^= 1:
   t[u].flip();
int merge(int 1, int r) {
   if (min(1, r) == -1) return 1 != -1 ? 1 : r;
   push(1), push(r), flip(1), flip(r);
   int ans = (t[1].p < t[r].p) ? 1 : r;</pre>
   if (ans == 1) t[1].r = merge(t[1].r, r), recalc(1);
   if (ans == r) t[r].l = merge(l, t[r].l), recalc(r):
   if (t[ans].l != -1) t[t[ans].l].par = ans; // only if
         parent needed
   if (t[ans].r != -1) t[t[ans].r].par = ans; // only if
         parent needed
   return ans:
pii split(int u, int id) {
   if (u == -1) return {-1, -1};
   push(u):
   flip(u);
   int szl = get(t[u].1).sz;
   if (szl >= id) {
       pii ans = split(t[u].l. id):
       if (ans.ss != -1) t[ans.ss].par = u; // only if
           parent needed
       if (ans.ff != -1) t[ans.ff].par = -1; // only if
           parent needed
       t[u].1 = ans.ss:
       recalc(u);
       return {ans.ff. u}:
   pii ans = split(t[u].r, id - szl - 1);
```

```
if (ans.ff != -1) t[ans.ff].par = u: // only if
            parent needed
       if (ans.ss != -1) t[ans.ss].par = -1; // only if
            parent needed
       t[u].r = ans.ff;
       recalc(u):
       return {u, ans.ss};
   int update(int u, int 1, int r, int v) {
       pii a = split(u, 1), b = split(a.ss, r - 1 + 1);
       t[b.ff].upd lazv(v):
       return merge(a.ff, merge(b.ff, b.ss));
   }
   void print(int u) {
       if (u == -1) return:
       push(u), flip(u);
       print(t[u].1):
       cout << t[u].v << ' ';
       print(t[u].r):
   }
   Treap(vi &v) : n(sz(v)) {
       for (int i = 0; i < n; i++) t.eb(v[i]), r = merge(r, left)
   }
};
```

5.3 dsu

```
struct Dsu {
    vector<int> p, r;

    // initialize the disjoint-set-union to all unitary sets
    void reset(int N) {
        p.resize(N), r.assign(N, 0);
        rep(i, N) p[i] = i;
    }

    // find the leader node corresponding to node 'i'
    int find(int i) {
        if (p[i] != i) p[i] = find(p[i]);
        return p[i];
    }

    // perform union on the two sets that 'i' and 'j' belong
        to
    void unite(int i, int j) {
        i = find(i), j = find(j);
        if (i == j) return;
```

```
if (r[i] > r[j]) swap(i, j);
    if (r[i] == r[j]) r[j] += 1;
    p[i] = j;
};
```

5.4 mo

```
struct Query { int 1, r, idx; };
// answer segment queries using only 'add(i)', 'remove(i)'
    and 'get()'
// functions.
// complexity: O((N + Q) * sqrt(N) * F)
// N = length of the full segment
// Q = amount of queries
// F = complexity of the 'add', 'remove' functions
template <class A, class R, class G, class T>
void mo(vector<Query> &queries, vector<T> &ans, A add, R
    remove, G get) {
   int Q = queries.size(), B = (int)sqrt(Q);
   sort(queries.begin(), queries.end(), [&](Query &a, Query
       return make_pair(a.1 / B, a.r) < make_pair(b.1 / B, b</pre>
   }):
   ans.resize(Q):
   int 1 = 0, r = 0;
   for (auto &g : gueries) {
       while (r < q.r) add(r), r++;
       while (1 > q.1) 1--, add(1);
       while (r > q.r) r--, remove(r):
       while (1 < q.1) remove(1), 1++;
       ans[q.idx] = get();
   }
```

5.5 ordered-set

```
ordered set:
int main() {
   ordered set p:
   p.insert(5); p.insert(2); p.insert(6); p.insert(4); // 0(
        log n)
    // value at 3rd index in sorted array. O(log n). Output:
   cout << "Value at 3rd index: " << *p.find_by_order(3) <<</pre>
    // index of number 6. O(log n). Output: 3
   cout << "Index of number 6: " << p.order of kev(6) <<</pre>
        endl:
   // number 7 not in the set but it will show the index
   // number if it was there in sorted array. Output: 4
   cout << "Index of number 7:" << p.order_of_key(7) << endl</pre>
   // number of elements in the range [3, 10)
    cout << p.order of kev(10) - p.order of kev(3) << endl:</pre>
```

persistent-segment-tree-lazy

```
template <class T>
struct Node {
   T x, lz;
   int 1 = -1, r = -1;
}:
template <class T>
struct Pstl {
   int N:
   vector<Node<T>> a:
   vector<int> head:
   T aneut() { return 0: }
   T merge(T 1, T r) { return 1 + r; }
   T uneut() { return 0: }
   T accum(T u, T x) { return u + x; }
   T apply(T x, T lz, int l, int r) { return x + (r - 1) *
        lz: }
   int build(int vl, int vr) {
       if (vr - vl == 1) a.push_back({qneut(), uneut()}); // }:
             node construction
       else {
           int vm = (vl + vr) / 2, l = build(vl, vm), r =
               build(vm, vr);
```

```
a.push_back({merge(a[1].x, a[r].x), uneut(), 1, r | 5.7 persistent-segment-tree
           }): // guerv merge
   return a.size() - 1:
T query(int 1, int r, int v, int v1, int vr, T acc) {
   if (1 >= vr || r <= vl) return gneut();</pre>
        // query neutral
   if (1 \le vl \&\& r \ge vr) return apply(a[v].x, acc, vl,
         vr): // update op
   acc = accum(acc, a[v].lz):
        // update merge
   int vm = (vl + vr) / 2:
   return merge(query(1, r, a[v].1, v1, vm, acc), query(
        1, r, a[v].r, vm, vr, acc)); // query merge
int update(int 1, int r, T x, int v, int v1, int vr) {
   if (1 >= vr || r <= vl || r <= 1) return v:
   a.push_back(a[v]);
   v = a.size() - 1:
   if (1 <= v1 && r >= vr) {
       a[v].x = apply(a[v].x, x, vl, vr); // update op
       a[v].lz = accum(a[v].lz, x); // update merge
   } else {
       int vm = (vl + vr) / 2:
       a[v].1 = update(1, r, x, a[v].1, v1, vm);
       a[v].r = update(1, r, x, a[v].r, vm, vr);
       a[v].x = merge(a[a[v].1].x, a[a[v].r].x); //
            query merge
   return v;
}
Pstl(int N) : N(N) { head.push_back(build(0, N)); }
T query(int t, int 1, int r) {
   return query(1, r, head[t], 0, N, uneut()); // update
         neutral
int update(int t, int 1, int r, T x) {
   return head.push_back(update(1, r, x, head[t], 0, N))
        , head.size() - 1;
```

```
// usage:
// Pst<Node<11>> pst;
// pst = {N};
// int newtime = pst.update(time, index, value);
// Node<ll> result = pst.query(newtime, left, right);
template <class T>
struct Node {
   T x:
   int 1 = -1, r = -1;
   Node(): x(0) {}
   Node(T x) : x(x) \{ \}
   Node (Node a, Node b, int l = -1, int r = -1) : x(a.x + b.
        x), 1(1), r(r) {}
};
template <class U>
struct Pst {
   int N:
   vector<U> a:
   vector<int> head:
   int build(int vl, int vr) {
       if (vr - vl == 1) a.push_back(U()); // node
            construction
       else {
           int vm = (vl + vr) / 2, l = build(vl, vm), r =
               build(vm, vr);
          a.push_back(U(a[1], a[r], 1, r)); // query merge
       return a.size() - 1;
   U query(int 1, int r, int v, int v1, int vr) {
       if (1 >= vr || r <= vl) return U(); // query neutral</pre>
       if (1 <= v1 && r >= vr) return a[v];
       int vm = (vl + vr) / 2:
       return U(query(1, r, a[v].1, v1, vm), query(1, r, a[v
            ].r, vm, vr)); // query merge
   int update(int i, U x, int v, int vl, int vr) {
       a.push back(a[v]):
       v = a.size() - 1;
       if (vr - vl == 1) a[v] = x; // update op
       else {
          int vm = (vl + vr) / 2;
```

5.8 segment-tree-lazy

```
template <class T>
struct Stl {
   int n:
   vector<T> a, b;
   T qneut() { return -2e9; }
   T uneut() { return 0; }
   T merge(T x, T y) { return max(x, y); }
   void upd(int v. T x. int l. int r)
      \{ a[v] += x, b[v] += x; \}
   Stl(int n = 0) : n(n), a(4 * n, qneut()),
      b(4 * n, uneut()) {}
   void push(int v, int vl, int vm, int vr) {
       upd(2 * v, b[v], vl, vm);
       upd(2 * v + 1, b[v], vm, vr);
      b[v] = uneut():
   T query(int 1, int r, int v=1, int vl=0, int vr=1e9) {
       vr = min(vr, n):
      if (1 <= v1 && r >= vr) return a[v];
      if (1 >= vr || r <= vl) return qneut();</pre>
      int vm = (vl + vr) / 2:
      push(v, v1, vm, vr);
```

5.9 segment-tree

```
struct St {
   int n;
   vector<1l> a;

   ll neut() { return 0; }
   ll merge(ll x, ll y) { return x + y; }

   St(int n = 0) : n(n), a(2 * n, neut()) {}

   ll query(int l, int r) {
      ll x = neut(), y = neut();
      for (l += n, r += n; l < r; l /= 2, r /= 2) {
        if (l & 1) x = merge(x, a[l++]);
        if (r & 1) y = merge(a[--r], y);
      }

      return merge(x, y);
   }

   void update(int i, ll x) {
      for (a[i += n] = x; i /= 2;)
        a[i] = merge(a[2 * i], a[2 * i + 1]);
   }
};</pre>
```

5.10 sparse-table

```
// handle immutable range maximum queries (or any idempotent
     query) in O(1)
template <class T>
struct Sparse {
   T op(T a, T b) { return max(a, b); }
   vector<vector<T>> st;
   Sparse() {}
   Sparse(int N) : st{vector<T>(N)} {}
   T &operator[](int i) { return st[0][i]; }
   // O(N log N) time, O(N log N) memory
   void init() {
       int N = st[0].size();
       int npot = N <= 1 ? 1 : 32 - __builtin_clz(N);</pre>
       st.resize(npot);
       repx(i, 1, npot) rep(j, N + 1 - (1 << i)) st[i].
           push_back(
              op(st[i-1][i], st[i-1][i+(1 << (i-1))
       ): // query op
   // query maximum in the range [1, r) in O(1) time
   // range must be nonempty!
   T query(int 1, int r) {
       int i = 31 - __builtin_clz(r - 1);
      return op(st[i][l], st[i][r - (1 << i)]); // query op</pre>
   }
};
```

5.11 unordered-map

6 imprimible

7 math

7.1 arithmetic

```
// floor(log2(n)) without precision loss
inline int floor_log2(int n) { return n <= 1 ? 0 : 31 -</pre>
     __builtin_clz(n); }
// ceil(log2(n)) without precision loss
inline int ceil log2(int n) { return n <= 1 ? 0 : 32 -</pre>
     __builtin_clz(n - 1); }
inline ll floordiv(ll a, ll b) {
   return a / b - ((a ^ b) < 0 && a % b);
inline 11 ceildiv(11 a, 11 b) {
   return a / b + ((a ^ b) >= 0 && a % b):
}
// a^e through binary exponentiation.
ll binexp(ll a, ll e) {
   ll res = 1: // neutral element
   while (e) {
       if (e & 1) res = res * a; // multiplication
                               // multiplication
       a = a * a:
       e >>= 1:
```

```
}
return res;
}
```

7.2 crt

```
pair<11, ll> solve_crt(const vector<pair<11, ll>> &eqs) {
    ll a0 = eqs[0].first, p0 = eqs[0].second;
    repx(i, 1, eqs.size()) {
        ll a1 = eqs[i].first, p1 = eqs[i].second;
        ll k1, k0;
        ll d = ext_gcd(p1, p0, k1, k0);
        a0 -= a1;
        if (a0 % d != 0) return {-1, -1};
        p0 = p0 / d * p1;
        a0 = a0 / d * k1 % p0 * p1 % p0 + a1;
        a0 = (a0 % p0 + p0) % p0;
    }
    return {a0, p0};
}
```

7.3 discrete-log

```
// discrete logarithm log_a(b).
// solve b \hat{x} = a \pmod{M} for the smallest x.
// returns -1 if no solution is found.
// time: O(sqrt(M))
11 dlog(ll a, ll b, ll M) {
    11 k = 1, s = 0:
    while (true) {
       ll g = \_gcd(b, M);
       if (g <= 1) break;</pre>
       if (a == k) return s;
       if (a % g != 0) return -1;
       a = g, M = g, s += 1, k = b / g * k % M;
    11 N = sqrt(M) + 1;
    umap<11, 11> r;
    rep(q, N + 1) {
       r[a] = q;
       a = a * b % M;
    ll bN = binexp(b, N, M), bNp = k;
    repx(p, 1, N + 1) {
```

```
bNp = bNp * bN % M;
    if (r.count(bNp)) return N * p - r[bNp] + s;
}
return -1;
}
```

7.4 fft

```
using cd = complex<double>;
const double PI = acos(-1):
// compute the DFT of a power-of-two-length sequence.
// if 'inv' is true, computes the inverse DFT.
void fft(vector<cd> &a, bool inv) {
   int N = a.size(), k = 0, b:
   assert(N == 1 << __builtin_ctz(N));</pre>
   repx(i, 1, N) {
       for (b = N >> 1; k & b;) k ^= b, b >>= 1;
       if (i < (k ^= b)) swap(a[i], a[k]);</pre>
   for (int 1 = 2; 1 <= N; 1 <<= 1) {</pre>
       double ang = 2 * PI / 1 * (inv ? -1 : 1):
       cd wl(cos(ang), sin(ang));
       for (int i = 0: i < N: i += 1) {
           cd w = 1;
           rep(j, 1 / 2) {
              cd u = a[i + j], v = a[i + j + 1 / 2] * w;
              a[i + j] = u + v;
              a[i + j + 1 / 2] = u - v;
              w *= wl:
      }
   }
   if (inv) rep(i, N) a[i] /= N;
const 11 MOD = 998244353, ROOT = 15311432;
// const 11 MOD = 2130706433. ROOT = 1791270792:
// const 11 MOD = 922337203673733529711, ROOT =
    532077456549635698311;
void find root of unitv(ll M) {
   11 c = M - 1, k = 0;
   while (c \% 2 == 0) c /= 2, k += 1;
   // find proper divisors of M - 1
```

```
vector<ll> divs:
   for (11 d = 1: d < c: d++) {
       if (d * d > c) break;
       if (c \% d == 0) rep(i, k + 1) divs.push back(d << i):
   rep(i, k) divs.push back(c << i):
   // find any primitive root of M
   11 G = -1;
   repx(g, 2, M) {
       bool ok = true:
       for (ll d : divs) ok &= (binexp(g, d, M) != 1):
       if (ok) {
          G = g;
          break:
       }
   assert(G != -1);
   ll w = binexp(G, c, M);
   cerr << "M = c * 2^k + 1" << endl;
   cerr << " M = " << M << endl:
   cerr << " c = " << c << endl;
   cerr << " k = " << k << endl:
   cerr << " w^(2^k) == 1" << endl:
   cerr << " w = g^{(M-1)/2k} = g^c << endl;
   cerr << " g = " << G << endl;</pre>
   cerr << " w = " << w << endl:
// compute the DFT of a power-of-two-length sequence, modulo
     a special prime
// number with an Nth root of unity, where N is the length
     of the sequence.
void ntt(vector<ll> &a. bool inv) {
   vector<ll> wn:
   for (11 p = ROOT; p != 1; p = p * p % MOD) wn.push_back(p
        );
   int N = a.size(), k = 0, b:
   assert(N == 1 << __builtin_ctz(N) && N <= 1 << wn.size()) \frac{1}{0} 0 -> no solutions
   rep(i, N) a[i] = (a[i] \% MOD + MOD) \% MOD:
   repx(i, 1, N) {
       for (b = N >> 1: k & b:) k ^= b, b >>= 1:
       if (i < (k ^= b)) swap(a[i], a[k]);</pre>
   for (int 1 = 2; 1 <= N; 1 <<= 1) {
```

```
11 wl = wn[wn.size() - builtin ctz(1)]:
      if (inv) wl = multinv(wl. MOD):
       for (int i = 0: i < N: i += 1) {</pre>
          11 w = 1:
          repx(i, 0, 1 / 2)  {
              11 u = a[i + j], v = a[i + j + 1 / 2] * w %
              a[i + j] = (u + v) \% MOD;
              a[i + j + 1 / 2] = (u - v + MOD) \% MOD;
              w = w * w1 % MOD:
      }
   11 q = multinv(N, MOD);
   if (inv) rep(i, N) a[i] = a[i] * q % MOD:
void convolve(vector<cd> &a. vector<cd> b. int n) {
   n = 1 \ll (32 - \_builtin\_clz(2 * n - 1));
   a.resize(n), b.resize(n);
   fft(a, false), fft(b, false);
   rep(i, n) a[i] *= b[i];
   fft(a, true);
```

gauss

```
const double EPS = 1e-9:
// solve a system of equations.
// complexity: O(min(N, M) * N * M)
// 'a' is a list of rows
// the last value in each row is the result of the equation
// return values:
// 1 -> unique solution, stored in 'ans'
// -1 -> infinitely many solutions, one of which is stored
    in 'ans'
// UNTESTED
int gauss(vector<vector<double>> a, vector<double> &ans) {
   int N = a.size(), M = a[0].size() - 1;
   vector<int> where(M, -1);
   for (int j = 0, i = 0; j < M && i < N; j++) {
      int sel = i:
```

```
repx(k, i, N) if (abs(a[k][i]) > abs(a[sel][i])) sel
   if (abs(a[sel][j]) < EPS) continue;</pre>
   repx(k, j, M + 1) swap(a[sel][k], a[i][k]);
    where[i] = i;
    rep(k, N) if (k != i) {
       double c = a[k][j] / a[i][j];
       repx(1, j, M + 1) a[k][1] -= a[i][1] * c;
   }
   i++:
}
ans.assign(M, 0);
rep(i, M) if (where[i] != -1) ans[i] = a[where[i]][M] / a
     [where[i]][i];
rep(i, N) {
   double sum = 0;
   rep(i, M) sum += ans[i] * a[i][i]:
   if (abs(sum - a[i][M]) > EPS) return 0:
rep(i, M) if (where[i] == -1) return -1;
return 1:
```

7.6 matrix

```
using T = 11;
struct Mat {
   int N. M:
   vector<vector<T>> v:
   Mat(int n, int m) : N(n), M(m), v(N, vector<T>(M)) {}
   Mat(int n) : Mat(n, n) { rep(i, N) v[i][i] = 1; }
   vector<T> &operator[](int i) { return v[i]: }
   Mat operator*(Mat &r) {
       assert(M == r.N):
       int n = N, m = r.M, p = M;
      Mat a(n. m):
      rep(i, n) rep(j, m) {
          a[i][j] = T();
                                                       //
               neutral
          rep(k, p) a[i][j] = a[i][j] + v[i][k] * r[k][j];
               // mul. add
      }
       return a;
```

```
Mat binexp(ll e) {
       assert(N == M):
       Mat a = *this, res(N); // neutral
       while (e) {
          if (e & 1) res = res * a; // mul
                                  // mul
          a = a * a:
          e >>= 1:
       }
       return res:
   friend ostream &operator<<(ostream &s. Mat &a) {</pre>
       rep(i, a.N) {
          rep(j, a.M) s << a[i][j] << " ";
          s << endl:
       }
       return s:
};
```

7.7 mobius

7.8 mod

```
11 binexp(ll a, ll e, ll M) {
    assert(e >= 0);
    ll res = 1 % M;
    while (e) {
        if (e & 1) res = res * a % M;
        a = a * a % M;
        e >>= 1;
    }
    return res;
}

ll multinv(ll a, ll M) { return binexp(a, M - 2, M); }

// calculate gcd(a, b).
// also, calculate x and y such that:
```

```
// a * x + b * v == gcd(a, b)
// time: O(log min(a, b))
// (ignoring complexity of arithmetic)
ll ext_gcd(ll a, ll b, ll &x, ll &y) {
   if (b == 0) {
      x = 1, y = 0;
      return a;
   ll d = ext_gcd(b, a \% b, y, x);
   v = a / b * x:
   return d:
// compute inverse with any M.
// a and M must be coprime for inverse to exist!
ll multinv euc(ll a. ll M) {
   11 x, y;
   ext gcd(a, M, x, v):
   return x:
// multiply two big numbers (~10^18) under a large modulo,
     without resorting to
// bigints.
11 bigmul(11 x, 11 v, 11 M) {
   11 z = 0:
    while (v) {
      if (y \& 1) z = (z + x) \% M;
      x = (x << 1) \% M, v >>= 1:
   }
   return z;
// all modular inverses from 1 to inv.size()-1
void multinv all(vector<ll> &inv) {
   inv[1] = 1:
   repx(i, 2, inv.size())
       inv[i] = MOD - (MOD / i) * inv[MOD % i] % MOD;
// change to __int128 if checking numbers over 10^9
bool isprime(ll n) {
   if (n < 2 | | n % 6 % 4 != 1) return n - 2 < 2;
   1795265022}:
   ll s = \_builtin\_ctzll(n - 1), d = n >> s;
   for (int a : A) {
      ll p = binexp(a, d, n), i = s:
```

```
while (p != 1 && p != n - 1 && a % n && i--) p = p *
       if (p != n - 1 && i != s) return 0:
   return 1;
struct Mod {
   int a:
   static const int M = 1e9 + 7;
   Mod(11 aa) : a((aa % M + M) % M) {}
   Mod operator+(Mod rhs) const { return (a + rhs.a) % M; }
   Mod operator-(Mod rhs) const { return (a - rhs.a + M) % M
   Mod operator-() const { return Mod(0) - *this: }
   Mod operator*(Mod rhs) const { return (11)a * rhs.a % M;
   Mod operator+=(Mod rhs) { return *this = *this + rhs; }
   Mod operator==(Mod rhs) { return *this = *this - rhs; }
   Mod operator*=(Mod rhs) { return *this = *this * rhs; }
   Mod bigmul(ll big) const { return ::bigmul(a, big, M); }
   Mod binexp(ll e) const { return ::binexp(a, e, M); }
   // Mod multinv() const { return ::multinv(a, M): } //
   Mod multinv() const { return ::multinv euc(a, M): } //
        possibly composite M
};
// dynamic modulus
struct DMod {
   int a. M:
   DMod(ll aa, ll m) : M(m), a((aa % m + m) % m) {}
   DMod operator+(DMod rhs) const { return {(a + rhs.a) % M.
   DMod operator-(DMod rhs) const { return {(a - rhs.a + M)
        % M. M}: }
   DMod operator-() const { return DMod(0, M) - *this; }
   DMod operator*(DMod rhs) const { return {(11)a * rhs.a %
        M. M}: }
   DMod operator+=(DMod rhs) { return *this = *this + rhs; }
   DMod operator == (DMod rhs) { return *this = *this - rhs: }
   DMod operator*=(DMod rhs) { return *this = *this * rhs: }
```

7.9 primes

```
// counts the divisors of a positive integer in O(\operatorname{sqrt}(n))
11 count divisors(11 x) {
   11 \text{ divs} = 1, i = 2;
   for (11 divs = 1, i = 2; x > 1; i++) {
       if (i * i > x) {
           divs *= 2:
           break:
        for (11 d = divs: x % i == 0: x /= i) divs += d:
    return divs;
}
// gets the prime factorization of a number in O(\operatorname{sqrt}(n))
vector<pair<11, int>> factorize(11 x) {
    vector<pair<11, int>> f;
    for (11^{-}k = 2; x > 1; k++) {
       if (k * k > x) {
           f.push_back(\{x, 1\});
           break:
       int n = 0;
       while (x \% k == 0) x /= k, n++:
       if (n > 0) f.push_back(\{k, n\});
    return f:
// iterate over all divisors of a number.
// divisor count upper bound: n^(1.07 / ln ln n)
template <class OP>
void divisors(ll x, OP op) {
    auto facts = factorize(x):
    vector<int> f(facts.size());
```

```
while (true) {
       11 v = 1:
       rep(i, f.size()) rep(j, f[i]) y *= facts[i].first;
       op(y);
       int i:
       for (i = 0; i < f.size(); i++) {</pre>
           f[i] += 1:
           if (f[i] <= facts[i].second) break;</pre>
       if (i == f.size()) break:
   }
// computes euler totative function phi(x), counting the
     amount of integers in
// [1, x] that are coprime with x.
// time: O(sart(x))
ll phi(ll x) {
   ll phi = 1, k = 2;
   for (; x > 1; k++) {
       if (k * k > x) {
           phi *= x - 1;
           break:
       11 k1 = 1, k0 = 0:
       while (x \% k == 0) x /= k, k0 = k1, k1 *= k;
       phi *= k1 - k0;
   return phi;
// computes primality up to N.
// considers 0 and 1 prime.
// O(N log N)
void sieve(int N, vector<bool> &prime) {
   prime.assign(N + 1, true);
   repx(n, 2, N + 1) if (prime[n]) for (int k = 2 * n; k <=
        N: k += n) prime[k] = false:
```

$7.10 \quad \text{simplex}$

```
T x$ otherwise.
// The input vector is set to an optimal $x$ (or in the
    unbounded case, an arbitrary solution fulfilling the
    constraints).
// Numerical stability is not guaranteed. For better
    performance, define variables such that x = 0 is
// Usage:
// vvd A = \{\{1,-1\}, \{-1,1\}, \{-1,-2\}\};
// vd b = \{1,1,-4\}, c = \{-1,-1\}, x;
// T val = LPSolver(A, b, c).solve(x):
// Time: O(NM * \#pivots), where a pivot may be e.g. an edge
     relaxation. O(2^n) in the general case.
#include "../common.h"
typedef double T; // long double, Rational, double + mod<P
    >...
typedef vector<T> vd:
tvpedef vector<vd> vvd:
const T eps = 1e-8, inf = 1 / .0;
#define MP make_pair
#define ltj(X) \
   if (s == -1 \mid | MP(X[i], N[i]) < MP(X[s], N[s])) s = i
struct LPSolver {
   int m. n:
   vector<int> N, B;
   vvd D:
   LPSolver(const vvd &A, const vd &b, const vd &c) : m(b.
        size()), n(c.size()), N(n + 1), B(m), D(m + 2), vd(n + 1)
        + 2)) {
       rep(i, m) rep(j, n) D[i][j] = A[i][j];
       rep(i, m) {
           B[i] = n + i;
           D[i][n] = -1;
           D[i][n + 1] = b[i];
       rep(j, n) {
           N[j] = j;
           D[m][i] = -c[i];
       }
       N[n] = -1;
       D[m + 1][n] = 1:
   void pivot(int r, int s) {
       T *a = D[r].data(). inv = 1 / a[s]:
```

```
rep(i, m + 2) if (i != r && abs(D[i][s]) > eps) {
       T *b = D[i].data(), inv2 = b[s] * inv;
       repx(j, 0, n + 2) b[j] -= a[j] * inv2;
       b[s] = a[s] * inv2:
   rep(i, n + 2) if (i != s) D[r][i] *= inv:
   rep(i, m + 2) if (i != r) D[i][s] *= -inv;
   D[r][s] = inv;
   swap(B[r], N[s]);
bool simplex(int phase) {
   int x = m + phase - 1;
   for (;;) {
       int s = -1:
       rep(j, n + 1) if (N[j] != -phase) ltj(D[x]);
       if (D[x][s] >= -eps) return true;
       int r = -1;
       rep(i, m) {
          if (D[i][s] <= eps) continue;</pre>
          if (r == -1 || MP(D[i][n + 1] / D[i][s], B[i])
                < MP(D[r][n + 1] / D[r][s], B[r])) r = i
       if (r == -1) return false;
       pivot(r, s);
T solve(vd &x) {
   int r = 0:
   repx(i, 1, m) if (D[i][n + 1] < D[r][n + 1]) r = i;
```

7.11 theorems

```
// Burnside lemma

//

// For a set X, with members x in X, and a group G, with operations g in G, where g(x): X -> X.

// F_g is the set of x which are fixed points of g (ie. { x in X / g(x) = x }).

// The number of orbits (connected components in the graph formed by assigning each x a node and

// a directed edge between x and g(x) for every g) is called M.

// M = the average of the fixed points of all g = (|F_g1| + |F_g2| + ... + |F_gn|) / |G|
```

```
If x are images and g are simmetries, then M
    corresponds to the amount of objects, |G|
      corresponds to the amount of simmetries, and F_g
    corresponds to the amount of simmetrical
     images under the simmetry g.
11
// Rational root theorem
     All rational roots of the polynomials with integer
    coefficients:
11
11
      a0 * x^0 + a1 * x^1 + a2 * x^2 + ... + an * x^n = 0
//
11
     If these roots are represented as p / q, with p and q
11
      - p is an integer factor of a0
11
      - q is an integer factor of an
11
11
     Note that if a0 = 0, then x = 0 is a root, the
    polynomial can be divided by x and the theorem
      applies once again.
11
// Legendre's formula
//
      Considering a prime p, the largest power p^k that
    divides n! is given by:
//
11
     k = floor(n/p) + floor(n/p^2) + floor(n/p^3) + ...
11
//
     Which can be computed in O(log n / log p) time
```