Team Notebook

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1 Strings

1.1 AhoCorasick

```
#include<bits/stdc++.h>
using namespace std;
const int K = 26:
struct Vertex {
   int next[K]:
   int leaf = 0:
   int leaf_id = -1;
   int p = -1;
   char pch;
   int link = -1;
   int exit = -1:
   int cnt = -1;
   int go[K];
   Vertex(int p=-1, char ch='$') : p(p), pch(ch) {
       fill(begin(next), end(next), -1);
       fill(begin(go), end(go), -1);
};
vector<Vertex> t(1);
void add(string &s, int id) {
   int v = 0:
   for (char ch : s) {
       int c = ch - 'a';
       if (t[v].next[c] == -1) {
           t[v].next[c] = t.size():
           t.emplace_back(v, ch);
       v = t[v].next[c];
   t[v].leaf++:
    t[v].leaf_id = id;
int go(int v, char ch);
int get_link(int v) {
   if (t[v].link == -1) {
       if (v == 0 || t[v].p == 0)
           t[v].link = 0;
       else
           t[v].link = go(get_link(t[v].p), t[v].pch);
```

```
return t[v].link:
int go(int v, char ch) {
   int c = ch - 'a';
   if (t[v].go[c] == -1) {
      if (t[v].next[c] != -1)
          t[v].go[c] = t[v].next[c];
          t[v].go[c] = v == 0 ? 0 : go(get_link(v), ch);
   return t[v].go[c];
int next match(int v)
   if(t[v].exit == -1)
       if(t[get link(v)].leaf)
          t[v].exit = get_link(v);
          t[v].exit = v == 0 ? 0 : next match(get link(v)):
   return t[v].exit;
int cnt matches(int v)
   if(t[v].cnt == -1)
       t[v].cnt = v == 0 ? 0 : t[v].leaf + cnt matches(
           get_link(v));
   return t[v].cnt:
```

1.2 Manacher

```
#include<bits/stdc++.h>
using namespace std;
#define rep(i, n) for (int i = 0; i < (int)n; i++)
#define repx(i, a, b) for (int i = (int)a; i < (int)b; i++)

// odd[i] : length of the longest palindrome centered at i
// even[i] : length of the longest palindrome centered
    between i and i+1
void manacher(string &s, vector<int> &odd, vector<int> &even
    ) {
    string t = "$#";
    for(char c: s)
        t += c + string("#");
```

```
t += "^";
int n = t.size();
vector<int> p(n);
int l = 1, r = 1;
repx(i, 1, n-1) {
    p[i] = max(0, min(r - i, p[1 + (r - i)]));
    while(t[i - p[i]] == t[i + p[i]]) {
        p[i]++;
    }
    if(i + p[i] > r) {
        l = i - p[i], r = i + p[i];
    }
}
repx(i, 2, n-2) {
    if(i%2) even.push_back(p[i]-1);
    else odd.push_back(p[i]-1);
}
```

1.3 PalindromicTree

```
#include<bits/stdc++.h>
using namespace std;
#define rep(i, n) for (int i = 0; i < (int)n; i++)
#define repx(i, a, b) for (int i = (int)a; i < (int)b; i++)
struct Node {
int len:
                  // length of substring
                  // insertion edge for all characters a-z
int edge[26];
int link;
                  // the Maximum Palindromic Suffix Node
     for the current Node
                  // (optional) start index of current Node
   int cnt = 1;
                     // (optional) number of occurrences of
        this substring
   Node(){ fill(begin(edge), end(edge), -1); }
};
struct EerTree { // Palindromic Tree
   vector<Node> t: // tree
   int curr:
              // current node
   EerTree(string &s) {
       t.resize(2);
       t.reserve(s.size()+2); // (optional) maximum size of
       t[0].len = -1;
                           // root 1
       t[0].link = 0;
       t[1].len = 0:
                           // root 2
       t[1].link = 0;
```

```
curr = 1:
       rep(i, s.size()) insert(i, s); // construct tree
       // (optional) calculate number of occurrences of each
             node
       for(int i = t.size()-1: i > 1: i--)
           t[t[i].link].cnt += t[i].cnt;
   }
   void insert(int i, string &s) {
       int tmp = curr:
       while (i - t[tmp].len < 1 \mid | s[i] != s[i-t[tmp].len
            -17)
           tmp = t[tmp].link;
       if(t[tmp].edge[s[i]-'a'] != -1){
           curr = t[tmp].edge[s[i]-'a']: // node already
               exists
           t[curr].cnt++:
                                         // (optional)
               increase cnt
           return:
       }
       curr = t[tmp].edge[s[i]-'a'] = t.size(); // create
            new node
       t.emplace_back();
       t[curr].len = t[tmp].len + 2:
                                         // set length
       t[curr].i = i - t[curr].len + 1; // (optional) set
            start index
       if (t[curr].len == 1) {
                                         // set suffix link
           t[curr].link = 1:
       } else {
           tmp = t[tmp].link;
           while (i-t[tmp].len < 1 || s[i] != s[i-t[tmp].len
              tmp = t[tmp].link;
           t[curr].link = t[tmp].edge[s[i]-'a'];
};
int main()
 string s = "abcbab":
   EerTree pt(s);  // construct palindromic tree
 repx(i, 2, pt.t.size()) // list all distinct palindromes
 cout << i-1 << ") ":
```

```
repx(j, pt.t[i].i, pt.t[i].i + pt.t[i].len)
  cout << s[j];
  cout << " " << pt.t[i].cnt << endl;
}
return 0;
}</pre>
```

1.4 PrefixFunction

```
#include<bits/stdc++.h>
using namespace std:
#define rep(i, n) for (int i = 0; i < (int)n; i++)
#define repx(i, a, b) for (int i = (int)a; i < (int)b; i++)
vector<int> prefix_function(string s) {
   int n = s.size():
   vector<int> pi(n);
   repx(i, 1, n) {
       int j = pi[i-1];
       while (j > 0 \&\& s[i] != s[j])
          j = pi[j-1];
       if (s[i] == s[j])
          j++;
       pi[i] = j;
   return pi;
vector<vector<int>> aut;
void compute_automaton(string s) {
   s += '#':
   int n = s.size():
   vector<int> pi = prefix_function(s);
   aut.assign(n, vector<int>(26));
   rep(i, n) {
      rep(c, 26) {
          int j = i;
          while (j > 0 \&\& 'a' + c != s[j])
              j = pi[j-1];
          if ('a' + c == s[j])
              j++;
          aut[i][c] = j;
      }
   }
```

1.5 SuffixArray

```
#include<bits/stdc++.h>
using namespace std;
#define LOG2(X) ((unsigned) (8*sizeof (unsigned long long) -
     __builtin_clzll((X)) - 1))
#define rep(i, n) for (int i = 0; i < (int)n; i++)
#define repx(i, a, b) for (int i = (int)a; i < (int)b; i++)</pre>
struct SuffixArray {
   int n: vector<int> C. R. R. sa. sa. lcp:
   inline int gr(int i) { return i < n ? R[i] : 0; } // sort</pre>
         suffixes
   //inline int gr(int i) { return R[i%n]; } // sort
        cyclic shifts
   void csort(int maxv. int k) {
       C.assign(maxv + 1, 0); rep(i, n) C[gr(i + k)]++;
       repx(i, 1, maxv + 1) C[i] += C[i - 1];
       for (int i = n - 1; i \ge 0; i--) sa_{-}[--C[gr(sa[i] + k]]
           )]] = sa[i];
       sa.swap(sa_);
   void getSA(vector<int>& s) {
       R = R_{=} = sa = sa_{=} = vector < int > (n); rep(i, n) sa[i] =
       sort(sa.begin(), sa.end(), [&s](int i, int j) {
            return s[i] < s[j]; });</pre>
       int r = R[sa[0]] = 1;
       repx(i, 1, n) R[sa[i]] = (s[sa[i]] != s[sa[i - 1]]) ?
             ++r : r:
       for (int h = 1; h < n && r < n; h <<= 1) {
          csort(r, h): csort(r, 0): r = R [sa[0]] = 1:
          repx(i, 1, n) {
              if (R[sa[i]] != R[sa[i - 1]] || gr(sa[i] + h)
                   != gr(sa[i - 1] + h)) r++:
              R_{sa}[i] = r;
          } R.swap(R_);
   void getLCP(vector<int> &s) {// only works with suffixes
        (not cyclic shifts)
       lcp.assign(n, 0); int k = 0;
       rep(i, n) {
          int r = R[i] - 1;
          if (r == n - 1) \{ k = 0; continue; \}
          int i = sa[r + 1]:
          while (i + k < n \&\& j + k < n \&\& s[i + k] == s[j]
               + kl) k++:
          lcp[r] = k: if (k) k--:
```

```
SuffixArray(vector<int> &s) { n = s.size(); getSA(s);
        getLCP(s); constructLCP(); }
    /* ----- */
    vector<vector<int>> T:
    void constructLCP() {
       T.assign(LOG2(n)+1, lcp);
       for(int k = 1; (1<<k) <= n; ++k)
           for(int i = 0; i + (1<<k) <= n; ++i)</pre>
              T[k][i] = min(T[k-1][i],T[k-1][i+(1<<(k-1))]):
    // get LCP of suffix starting at i and suffix starting at
    int queryLCP(int i, int j) {
       if(i == j) return n-i;
       i = R[i]-1; j = R[j]-1;
       if(i > j) swap(i, j);
       11 k = LOG2(i-i):
       return min(T[k][i],T[k][j-(1<<k)]);</pre>
    // compare substring of length len1 starting at i
    // with substring of length len2 starting at j
    bool cmp(int i, int len1, int j, int len2) {
       if(queryLCP(i, j) >= min(len1, len2))
           return (len1 < len2);</pre>
           return (R[i] < R[j]);</pre>
}:
vector<int> suffix_array;
vector<vector<int>> C:
int n;
void sort_cyclic_shifts(string s) {
    s += "$":
    n = s.size();
    const int alphabet = 256;
    vector<int> p(n), c(n), cnt(max(alphabet, n), 0);
    for (int i = 0: i < n: i++)</pre>
       cnt[s[i]]++;
    for (int i = 1; i < alphabet; i++)</pre>
       cnt[i] += cnt[i-1];
    for (int i = 0; i < n; i++)</pre>
       p[--cnt[s[i]]] = i:
    c[p[0]] = 0;
    int classes = 1:
    for (int i = 1; i < n; i++) {</pre>
       if (s[p[i]] != s[p[i-1]])
```

```
classes++:
       c[p[i]] = classes - 1;
   C.emplace_back(c.begin(), c.end());
   vector<int> pn(n), cn(n);
   for (int h = 0; (1 << h) < n; ++h) {
       for (int i = 0; i < n; i++) {</pre>
          pn[i] = p[i] - (1 << h);
          if (pn[i] < 0)
              pn[i] += n;
       fill(cnt.begin(), cnt.begin() + classes, 0):
      for (int i = 0; i < n; i++)</pre>
           cnt[c[pn[i]]]++;
       for (int i = 1; i < classes; i++)</pre>
          cnt[i] += cnt[i-1];
       for (int i = n-1: i >= 0: i--)
          p[--cnt[c[pn[i]]]] = pn[i];
       cn[p[0]] = 0:
       classes = 1:
       for (int i = 1; i < n; i++) {</pre>
          pair < int, int > cur = \{c[p[i]], c[(p[i] + (1 << h))\}
          pair<int, int> prev = {c[p[i-1]], c[(p[i-1] + (1
               << h)) % nl}:
          if (cur != prev)
              ++classes:
          cn[p[i]] = classes - 1:
       c.swap(cn):
       C.emplace_back(c.begin(), c.end());
   p.erase(p.begin());
   suffix_array = p;
vector<int> lcp_construction(string &s, vector<int> &p) {
   int n = s.size():
   vector<int> rank(n);
   rep(i, n) rank[p[i]] = i;
   int k = 0;
   vector < int > lcp(n-1, 0):
   rep(i, n) {
      if (rank[i] == n - 1) {
          k = 0:
          continue;
      int j = p[rank[i] + 1];
      while (i + k < n \&\& j + k < n \&\& s[i+k] == s[j+k])
```

```
k++:
       lcp[rank[i]] = k;
       if (k)
          k--:
   return lcp:
bool compare1(int i, int j, int l) {
   int k = LOG2(1);
   pair<int, int> a = \{C[k][i], C[k][(i+l-(1 << k))\%n]\};
   pair<int, int> b = \{C[k][i], C[k][(i+1-(1 << k))\%n]\}:
   return a >= b;
bool compare2(int i, int j, int 1) {
   int k = LOG2(1):
   pair<int, int> a = \{C[k][i], C[k][(i+l-(1 << k))\%n]\};
   pair<int, int> b = \{C[k][j], C[k][(j+1-(1 << k))%n]\};
   return a <= b:
pair<int,int> find(int i, int len)
   int 1 = 0, r = suffix_array.size()-1;
   while(1 != r)
       int mid = (1+r)/2;
       if(compare1(suffix arrav[mid], i, len))
          r = mid:
       else
          1 = mid+1:
   int left = 1:
   1 = 0, r = suffix_array.size()-1;
   while(1 != r)
       int mid = (1+r+1)/2:
       if(compare2(suffix arrav[mid], i, len))
          1 = mid:
       else
          r = mid-1:
   int right = 1:
   if(!compare1(suffix_array[left], i, len)) return {-1,-1};
   if(!compare2(suffix_array[right], i, len)) return
        {-1,-1}:
```

```
if(left > right) return {-1,-1};
return {left, right};
```

1.6 SuffixAutomaton

```
#include<bits/stdc++.h>
using namespace std;
#define rep(i, n) for (int i = 0; i < (int)n; i++)
#define repx(i, a, b) for (int i = (int)a; i < (int)b; i++)
struct SuffixAutomaton {
   vector<map<char,int>> edges; // edges[i] : the labeled
        edges from node i
   vector<int> link:
                             // link[i] : the suffix link
       of i
   vector<int> length;
                             // length[i] : the length of
        the longest string in the ith class
   vector<int> cnt;
                             // cnt[i] : number of
       occurrences of each string in the ith class
   vector<int> paths;
                             // paths[i] : number of paths
       on the automaton starting from i
                             // terminal[i] : true if i is
   vector<br/>terminal:
       a terminal state
   vector<int> first_pos;
   vector<int> last_pos;
                             // the index of the
   int last:
        equivalence class of the whole string
   SuffixAutomaton(string s) {
       edges.push back(map<char.int>()):
      link.push_back(-1);
      length.push back(0):
      last = 0:
      rep(i, s.size()) { // construct r
          edges.push_back(map<char,int>());
          length.push_back(i+1);
          link.push back(0):
          int r = edges.size() - 1:
          int p = last; // add edges to r and find p with
               link to q
          while(p >= 0 && !edges[p].count(s[i])) {
              edges[p][s[i]] = r:
              p = link[p];
          if(p != -1) {
              int q = edges[p][s[i]];
```

```
if(length[p] + 1 == length[a]) {
              link[r] = q; // we do not have to split q,
                   just set the correct suffix link
          } else { // we have to split, add q'
              edges.push_back(edges[q]); // copy edges
                  of a
              length.push_back(length[p] + 1);
              link.push_back(link[q]); // copy parent of
              int qq = edges.size()-1;
              link[q] = qq; // add qq as the new parent
                  of a and r
              link[r] = qq;
              while(p >= 0 && edges[p][s[i]] == q) { //
                  move short classes polling to q to
                  poll to q'
                 edges[p][s[i]] = qq;
                 p = link[p];
          }
       }
       last = r:
/* ----- */
   // mark terminal nodes
   terminal.assign(edges.size(), false):
   int p = last;
   while(p > 0) {
       terminal[p] = true;
       p = link[p];
   }
   // precompute match count
   cnt.assign(edges.size(), -1);
   cnt matches(0):
   // precompute number of paths (substrings) starting
        from state
   paths.assign(edges.size(), -1);
   cnt_paths(0);
   first_pos.assign(edges.size(), -1);
   get_first_pos(0);
   last_pos.assign(edges.size(), -1);
   get_last_pos(0);
}
```

```
int cnt matches(int state) {
    if(cnt[state] != -1) return cnt[state]:
    int ans = terminal[state];
   for(auto edge : edges[state])
       ans += cnt_matches(edge.second);
   return cnt[state] = ans:
int cnt_paths(int state) {
    if(paths[state] != -1) return paths[state];
    int ans = state == 0 ? 0 : 1; // without repetition (
        counts diferent substrings)
// int ans = state == 0 ? 0 : cnt[state]; // with
     repetition
   for(auto edge : edges[state])
       ans += cnt_paths(edge.second);
   return paths[state] = ans:
}
int get_first_pos(int state) {
   if(first_pos[state] != -1) return first_pos[state];
    int ans = 0:
   for(auto edge : edges[state])
       ans = max(ans, get_first_pos(edge.second)+1);
   return first_pos[state] = ans;
}
int get last pos(int state) {
   if(last_pos[state] != -1) return last_pos[state];
   int ans = terminal[state] ? 0 : INT MAX://fix
   for(auto edge : edges[state])
       ans = min(ans, get_last_pos(edge.second)+1);
   return last_pos[state] = ans;
string get k substring(int k) // 0-indexed
    string ans;
   int state = 0;
   while(true)
       int curr = state == 0 ? 0 : 1; // without
            repetition (counts different substrings)
   // int curr = state == 0 ? 0 : cnt[state]; // with
        repetition
       if(curr > k) return ans:
       k -= curr;
       for(auto edge : edges[state]) {
           if(paths[edge.second] <= k) {</pre>
```

```
k -= paths[edge.second];
} else {
    ans += edge.first;
    state = edge.second;
    break;
}
}
}
}
```

1.7 hash

```
// compute substring hashes in O(1).
// hashes are compatible between different strings.
struct Hash {
   11 HMOD;
   int N:
   vector<int> h;
   vector<int> p;
   Hash() {}
   // O(N)
   Hash(const string &s, 11 HMOD_ = 1000003931)
       : N(s.size() + 1), HMOD(HMOD_), p(N), h(N) {
       static const 11 P =
           chrono::steady_clock::now().time_since_epoch().
               count() % (1 << 29);
       p[0] = 1:
       rep(i, N - 1) p[i + 1] = p[i] * P % HMOD;
       rep(i, N-1) h[i+1] = (h[i] + (ll)s[i] * p[i]) %
   pair<ll, int> get(int i, int j) { return {(h[j] - h[i] +
        HMOD) % HMOD, i}: }
   bool cmp(pair<11, int> x0, pair<11, int> x1) {
       int d = x0.second - x1.second:
       ll &lo = d < 0 ? x0.first : x1.first:
       lo = lo * p[abs(d)] % HMOD;
       return x0.first == x1.first;
}:
// compute hashes in multiple prime modulos simultaneously,
     to reduce the chance
// of collisions.
```

```
struct HashM {
   int N:
   vector<Hash> sub:
   HashM() {}
   // O(K N)
   HashM(const string &s, const vector<11> &mods) : N(mods.
        size()), sub(N) {
       rep(i, N) sub[i] = Hash(s, mods[i]);
   // O(K)
   vector<pair<11, int>> get(int i, int j) {
       vector<pair<11, int>> hs(N);
       rep(k, N) hs[k] = sub[k].get(i, j);
       return hs;
   }
   bool cmp(const vector<pair<11, int>> &x0, const vector<
        pair<11. int>> &x1) {
       rep(i, N) if (!sub[i].cmp(x0[i], x1[i])) return false
       return true;
   }
   bool cmp(int i0, int j0, int i1, int j1) {
       rep(i, N) if (!sub[i].cmp(sub[i].get(i0, j0),
                               sub[i].get(i1, j1))) return
                                    false;
       return true:
   }
};
#ifndef NOMAIN_HASH
int main() {
   const vector<ll> HMOD = {1000001237, 1000003931};
   // 01234567890123456789012
   string s = "abracadabra abracadabra";
   HashM h(s, HMOD):
   rep(i0, s.size() + 1) repx(i0, i0, s.size() + 1) rep(i1, i0, s.size() + 1)
        s.size() + 1)
       repx(i1, i1, s.size() + 1) {
       bool eq = h.cmp(h.get(i0, j0), h.get(i1, j1));
       bool eq2 = s.substr(i0, j0 - i0) == s.substr(i1, j1 -
             i1):
       if (eq != eq2) {
           cout << " hash says strings \"" << s.substr(i0,</pre>
               i0 - i0) << "\" "
```

1.8 hash2d

```
using Hash = pair<11, int>;
struct Block {
   int x0, y0, x1, y1;
};
struct Hash2d {
   11 HMOD:
   int W, H;
   vector<int> h:
   vector<int> p;
   Hash2d() {}
   Hash2d(const string &s, int W_, int H_, 11 HMOD_ =
        1000003931)
       : W(W_ + 1), H(H_ + 1), HMOD(HMOD_) {
       static const 11 P =
          chrono::steady_clock::now().time_since_epoch().
               count() % (1 << 29):
       p.resize(W * H);
       p[0] = 1:
       rep(i, W * H - 1) p[i + 1] = p[i] * P % HMOD;
      h.assign(W * H, 0);
       repx(y, 1, H) repx(x, 1, W) {
          ll c = (ll)s[(v - 1) * (W - 1) + x - 1] * p[v * W
                + xl % HMOD:
          h[y * W + x] = (HMOD + h[y * W + x - 1] + h[(y -
               1) * W + x] -
                        h[(y-1) * W + x - 1] + c) %
                        HMOD:
   7
   bool isout(Block s) {
       return s.x0 < 0 || s.x0 >= W || s.x1 < 0 || s.x1 >= W
            | | s.y0 < 0 | |
```

```
s.v0 >= H || s.v1 < 0 || s.v1 >= H:
   }
    Hash get(Block s) {
       return \{(2 * HMOD + h[s.v1 * W + s.x1] - h[s.v1 * W +
             s.x01 -
               h[s.y0 * W + s.x1] + h[s.y0 * W + s.x0]) %
                  HMOD.
              s.v0 * W + s.x0:
    bool cmp(Hash x0, Hash x1) {
       int d = x0.second - x1.second:
       ll &lo = d < 0 ? x0.first : x1.first:
       lo = lo * p[abs(d)] % HMOD;
       return x0.first == x1.first;
};
struct Hash2dM {
    int N;
    vector<Hash2d> sub:
    Hash2dM() {}
    Hash2dM(const string &s, int W, int H, const vector<11> &
        mods)
       : N(mods.size()), sub(N) {
       rep(i, N) sub[i] = Hash2d(s, W, H, mods[i]):
    bool isout(Block s) { return sub[0].isout(s); }
    vector<Hash> get(Block s) {
       vector<Hash> hs(N);
       rep(i, N) hs[i] = sub[i].get(s):
       return hs:
    bool cmp(const vector<Hash> &x0, const vector<Hash> &x1)
       rep(i, N) if (!sub[i].cmp(x0[i], x1[i])) return false
       return true:
    bool cmp(Block s0, Block s1) {
            s1))) return false:
       return true:
```

```
}:
#ifndef NOMAIN HASH2D
const vector<ll> HMOD = {1000002649, 1000000933, 1000003787,
      1000002173}:
int main() {}
#endif
```

1.9 sufarr

```
// build the suffix array
                                                        // suffixes are sorted, with each suffix represented by its
                                                             starting position
                                                        vector<int> suffixarray(const string &s) {
                                                            int N = s.size() + 1; // optional: include terminating
                                                            vector\langle int \rangle p(N), p2(N), c(N), c2(N), cnt(256);
                                                            rep(i, N) cnt[s[i]] += 1:
                                                            repx(b, 1, 256) cnt[b] += cnt[b - 1];
                                                            rep(i, N) p[--cnt[s[i]]] = i;
                                                            repx(i, 1, N) c[p[i]] = c[p[i - 1]] + (s[p[i]] != s[p[i - 1]])
                                                                 1]]);
                                                            for (int k = 1: k < N: k <<= 1) {
                                                               int C = c[p[N - 1]] + 1;
                                                               cnt.assign(C + 1, 0);
                                                               for (int &pi : p) pi = (pi - k + N) % N;
                                                               for (int cl : c) cnt[cl + 1] += 1;
                                                               rep(i, C) cnt[i + 1] += cnt[i];
                                                               rep(i, N) p2[cnt[c[p[i]]]++] = p[i]:
                                                               c2[p2[0]] = 0;
                                                               repx(i, 1, N) c2[p2[i]] =
                                                                   c2[p2[i-1]] + (c[p2[i]] != c[p2[i-1]] ||
                                                                                  c[(p2[i] + k) \% N] != c[(p2[i - 1]
                                                                                        + k) % N]):
                                                                swap(c, c2), swap(p, p2);
                                                            p.erase(p.begin()); // optional: erase terminating NUL
                                                            return p;
                                                        // build the lcp
rep(i, N) if (!sub[i].cmp(sub[i].get(s0), sub[i].get( | // 'lcp[i]' represents the length of the longest common
                                                             prefix between suffix i
                                                        // and suffix i+1 in the suffix array 'p'. the last element
                                                             of 'lcp' is zero by
                                                        // convention
```

```
vector<int> makelcp(const string &s. const vector<int> &p) {
   int N = p.size(), k = 0:
   vector<int> r(N), lcp(N);
   rep(i, N) r[p[i]] = i;
   rep(i, N) {
       if (r[i] + 1 >= N) {
          k = 0:
           continue;
       int j = p[r[i] + 1];
       while (i + k < N \&\& i + k < N \&\& s[i + k] == s[i + k]
           1) k += 1:
       lcp[r[i]] = k;
       if (k) k -= 1:
   return lcp;
#ifndef NOMAIN SUFARR
void test(const string &s) {
   cout << "suffix array for string \"" << s << "\" (length</pre>
        " << s.size()
        << "):" << endl:
   vector<int> sa = suffixarray(s);
   vector<int> lcp = makelcp(s, sa);
   rep(i, sa.size()) {
       int i = sa[i]:
       if (i > 0) cout << " " << lcp[i - 1] << endl;</pre>
       cout << " \"" << s.substr(j) << "\"" << endl:
   }
int main() {
   test("hello"):
   test("abracadabra"):
#endif
```

dp

2.1 convex-hull-trick

```
struct Line {
   mutable ll a, b, c;
   bool operator<(Line r) const { return a < r.a; }</pre>
```

```
bool operator<(11 x) const { return c < x: }</pre>
}:
// dynamically insert 'a*x + b' lines and guery for maximum
     at any x
// all operations have complexity O(log N)
struct LineContainer : multiset<Line, less<>>> {
    11 div(ll a, ll b) {
       return a / b - ((a ^ b) < 0 && a % b);
    bool isect(iterator x, iterator y) {
       if (y == end()) return x->c = INF, 0;
       if (x->a == y->a) x->c = x->b > y->b? INF : -INF;
       else x->c = div(y->b - x->b, x->a - y->a);
       return x->c >= v->c:
    void add(ll a. ll b) {
       // a *= -1, b *= -1 // for min
       auto z = insert(\{a, b, 0\}), y = z++, x = y;
       while (isect(y, z)) z = erase(z);
       if (x != begin() && isect(--x, y)) isect(x, y = erase
       while ((y = x) != begin() && (--x)->c >= y->c) isect(
            x. erase(v)):
   11 guerv(ll x) {
       if (empty()) return -INF; // INF for min
       auto 1 = *lower_bound(x):
       return 1.a * x + 1.b:
       // return -l.a * x - l.b; // for min
};
```

2.2 divide-and-conquer

```
// for every index i assign an optimal index j, such that
    cost(i, j) is
// minimal for every i. the property that if i2 >= i1 then
        j2 >= j1 is
// exploited (monotonic condition).
// calculate optimal index for all indices in range [1, r)
        knowing that
// the optimal index for every index in this range is within
        [opt1, optr).
// time: O(N log N)
```

```
void calc(vector<int> &opt, int 1, int r, int opt1, int optr
    ) {
    if (1 == r) return;
    int i = (1 + r) / 2;
    ll optc = INF;
    int optj;
    repx(j, opt1, optr) {
        ll c = i + j; // cost(i, j)
        if (c < optc) optc = c, optj = j;
    }
    opt[i] = optj;
    calc(opt, 1, i, opt1, optj + 1);
    calc(opt, i + 1, r, optj, optr);
}</pre>
```

$3 \quad \text{geo2d}$

3.1 circle

```
struct C {
   Po; Tr;
   C(P \ o, T \ r) : o(o), r(r) \{\}
   C() : C(P(), T()) {}
   // intersects the circle with a line, assuming they
   // results are sorted with respect to the direction of
        the line
   pair<P. P> line inter(L 1) const {
      P c = 1.closest to(o):
      T c2 = (c - o).magsq();
      P = sart(max(r * r - c2, T())) * 1.d.unit():
      return {c - e, c + e}:
   // checks whether the given line collides with the circle
   // negative: 2 intersections
   // zero: 1 intersection
   // positive: 0 intersections
   T line collide(L 1) const {
      T c2 = (1.closest_to(o) - o).magsq();
      return c2 - r * r;
   // calculates the two intersections between two circles
   // the circles must intersect in one or two points!
   pair<P, P> inter(C h) const {
```

```
P d = h.o - o:
   T c = (r * r - h.r * h.r) / d.magsq();
   return h.line_inter({(1 + c) / 2 * d, d.rot()});
// check if the given circles intersect
bool collide(C h) const {
   return (h.o - o).magsq() \le (h.r + r) * (h.r + r);
// get one of the two tangents that cross through the
// the point must not be inside the circle
// a = -1: cw (relative to the circle) tangent
// a = 1: ccw (relative to the circle) tangent
P point_tangent(P p, T a) const {
   T c = r * r / p.magsq();
   return o + c * (p - o) - a * sqrt(c * (1 - c)) * (p -
         o).rot():
// get one of the 4 tangents between the two circles
// a = 1: exterior tangents
// a = -1: interior tangents (requires no area overlap)
// b = 1: ccw tangent
// b = -1: cw tangent
// the line origin is on this circumference, and the
     direction
// is a unit vector towards the other circle
L tangent(C c, T a, T b) const {
   T dr = a * r - c.r:
   P d = c.o - o;
   P n = (d * dr + b * d.rot() * sqrt(d.magsq() - dr *
        dr)).unit();
   return {o + n * r. -b * n.rot()}:
// find the circumcircle of the given **non-degenerate**
     triangle
static C thru_points(P a, P b, P c) {
   L 1((a + b) / 2, (b - a).rot()):
   P p = 1.intersection(L((a + c) / 2, (c - a).rot()));
   return {p, (p - a).mag()}:
// find the two circles that go through the given point,
     are tangent
// to the given line and have radius 'r'
// the point-line distance must be at most 'r'!
// the circles are sorted in the direction of the line
```

3.2 convex-hull

```
// get the convex hull with the least amount of vertices for
      the given set
// of points
// probably misbehaves if points are not all distinct!
vector<P> convex_hull(vector<P> &ps) {
   int N = ps.size(), n = 0, k = 0;
   if (N <= 2) return ps;</pre>
   rep(i, N) if (make_pair(ps[i].y, ps[i].x) < make_pair(ps[])</pre>
        k].y, ps[k].x)) k = i;
   swap(ps[k], ps[0]);
   sort(++ps.begin(), ps.end(), [&](P 1, P r) {
       T x = (r - 1) / (ps[0] - 1), d = (r - 1) * (ps[0] - 1)
       return x > 0 \mid | x == 0 && d < 0:
   }):
   vector<P> H:
   for (P p : ps) {
       while (n \ge 2 \&\& (H[n-1] - p) / (H[n-2] - p) >=
            0) H.pop_back(), n--;
       H.push_back(p), n++;
   return H;
```

3.3 delaunay

```
typedef int128 t lll: // if on a 64-bit platform
    Q *rot, *o; P p = {INF, INF}; bool mark;
    P &F() { return r()->p; }
    Q *&r() { return rot->rot; }
    Q *prev() { return rot->o->rot; }
    Q *next() { return r()->prev(); }
}:
T cross(P a, P b, P c) { return (b - a) / (c - a); }
bool circ(P p, P a, P b, P c) { // is p in the circumcircle?
    111 p2 = p.magsq(), A = a.magsq() - p2,
        B = b.magsq() - p2, C = c.magsq() - p2;
    return cross(p, a, b) * C + cross(p, b, c) * A + cross(p, b, c)
          c. a) * B > 0:
Q *makeEdge(Q *&H, P orig, P dest) {
    Q *r = H ? H : new Q{new Q{new Q{0}}};
    H = r -> 0: r -> r() -> r() = r:
    repx(i, 0, 4) r = r->rot, r->p = {INF, INF}.
        r \rightarrow 0 = i \& 1 ? r : r \rightarrow r();
    r->p = orig; r->F() = dest;
    return r;
void splice(Q *a, Q *b) {
    swap(a\rightarrow o\rightarrow rot\rightarrow o, b\rightarrow o\rightarrow rot\rightarrow o): swap(a\rightarrow o, b\rightarrow o):
0 *connect(0 *&H. 0 *a. 0 *b) {
    Q *q = makeEdge(H, a->F(), b->p);
    splice(q, a->next()); splice(q->r(), b); return q;
pair<Q *. Q *> rec(Q *&H. const vector<P> &s) {
    if (s.size() <= 3) {</pre>
        Q *a = makeEdge(H, s[0], s[1]), *b = makeEdge(H, s[0], s[1])
              [1], s.back()):
        if (s.size() == 2) return \{a, a\rightarrow r()\}; splice(a\rightarrow r(), a\rightarrow r());
        auto side = cross(s[0], s[1], s[2]):
        Q *c = side ? connect(H, b, a) : 0;
        return {side < 0 ? c->r() : a, side < 0 ? c : b->r()
             }:
```

```
#define J(e) e \rightarrow F(), e \rightarrow p
#define valid(e) (cross(e->F(), J(base)) > 0)
   Q *A, *B, *ra, *rb: int half = s.size() / 2:
   tie(ra, A) = rec(H, {s.begin(), s.end() - half});
   tie(B, rb) = rec(H, {s.begin() + s.size() - half, s.end()
   while ((cross(B\rightarrow p, J(A)) < 0 \&\& (A = A\rightarrow next())) | |
          (cross(A->p, J(B)) > 0 \&\& (B = B->r()->o)));
   Q *base = connect(H, B->r(), A);
   if (A->p == ra->p) ra = base->r();
   if (B->p == rb->p) rb = base:
#define DEL(e, init, dir) Q *e = init->dir: \
    if (valid(e)) while (circ(e->dir->F(), J(base), e->F()))
           Q *t = e->dir; splice(e, e->prev()); \
           splice(e->r(), e->r()->prev()); e->o = H; H = e;
                e = t: \
   for (;;) {
       DEL(LC, base->r(), o); DEL(RC, base, prev());
       if (!valid(LC) && !valid(RC)) break;
       if (!valid(LC) || (valid(RC) && circ(J(RC), J(LC))))
            base = connect(H, RC, base->r());
       else base = connect(H, base->r(), LC->r());
   return {ra, rb}:
#undef J
#undef valid
#undef DEL
// there must be no duplicate points
// returns no triangles in the case of all collinear points
// produces counter-clockwise triangles ordered in triples
// maximizes the minimum angle across all triangulations
// the euclidean mst is a subset of these edges
// O(N log N)
vector<P> triangulate(vector<P> pts) {
   sort(pts.begin(), pts.end(), [](P a, P b) {
       return make_pair(a.x, a.y) < make_pair(b.x, b.y);</pre>
   assert(unique(pts.begin(), pts.end()) == pts.end());
   if (pts.size() < 2) return {};</pre>
   Q *H = 0: Q *e = rec(H. pts).first:
   vector<Q *> q = {e}; int qi = 0;
   while (cross(e->o->F(), e->F(), e->p) < 0) e = e->o;
#define ADD
```

3.4 halfplane-intersect

```
// obtain the convex polygon that results from intersecting
    the given list
// of halfplanes, represented as lines that allow their left
// assumes the halfplane intersection is bounded
vector<P> halfplane intersect(vector<L> &H) {
   L bb(P(-INF, -INF), P(INF, 0));
   rep(k, 4) H.push_back(bb), bb.o = bb.o.rot(), bb.d = bb.d
        .rot():
   sort(begin(H), end(H), [](L a, L b) { return a.d.angcmp(b
        .d) < 0: }):
   deque<L> q; int n = 0;
   rep(i, H.size()) {
       while (n \ge 2 \&\& H[i].side(q[n - 1].intersection(q[n - 1]))
            -21)) > 0)
           g.pop back(), n--:
       while (n >= 2 && H[i].side(q[0].intersection(q[1])) >
           q.pop_front(), n--;
       if (n > 0 \&\& H[i].parallel(q[n - 1])) {
          if (H[i].d * q[n - 1].d < 0) return {};</pre>
          if (H[i].side(q[n-1].o) > 0) q.pop_back(), n--;
           else continue:
      }
       q.push_back(H[i]), n++;
   while (n \ge 3 \&\& q[0].side(q[n - 1].intersection(q[n -
        21)) > 0)
       q.pop_back(), n--;
   while (n \ge 3 \&\& q[n - 1].side(q[0].intersection(q[1])) >
       q.pop_front(), n--;
```

```
if (n < 3) return {};

vector<P> ps(n);
 rep(i, n) ps[i] = q[i].intersection(q[(i + 1) % n]);
 return ps;
}
```

3.5 line

```
// a segment or an infinite line
// does not handle point segments correctly!
struct L {
   P o. d:
   L() : o(), d()  {}
   L(P o, P d) : o(o), d(d) {}
   L(P ab, T c) : d(ab.rot()), o(ab * -c / ab.magsq()) {}
   pair<P, T> line_eq() { return {-d.rot(), d.rot() * o}; }
   // returns a number indicating which side of the line the
         point is in
   // negative: left, positive: right
   T side(P r) const { return (r - o) / d: }
   // returns the intersection coefficient
   // in the range [0, d / r.d]
   // if d / r.d is zero, the lines are parallel
   T inter(L r) const { return (r.o - o) / r.d; }
   // get the single intersection point
   // lines must not be parallel
   P intersection(L r) const { return o + d * inter(r) / (d
        / r.d): }
   // check if lines are parallel
   bool parallel(L r) const { return abs(d / r.d) <= EPS; }</pre>
   // check if segments intersect
   bool seg collide(L r) const {
      Tz = d / r.d:
      if (abs(z) <= EPS) {
          if (abs(side(r.o)) > EPS) return false;
          T s = (r.o - o) * d, e = s + r.d * d;
          if (s > e) swap(s, e):
          return s <= d * d + EPS && e >= -EPS;
      T s = inter(r), t = -r.inter(*this);
       if (z < 0) s = -s, t = -t, z = -z;
```

```
return s \geq= -EPS && s \leq= z + EPS && t \geq= -EPS && t \leq=
             z + EPS:
   }
   // full segment intersection
   // produces a point segment if the intersection is a
   // however it **does not** handle point segments as input
   bool seg_inter(L r, L *out) const {
      Tz = d / r.d:
       if (abs(z) <= EPS) {
          if (abs(side(r.o)) > EPS) return false;
          if (r.d * d < 0) r = \{r.o + r.d, -r.d\};
          P s = o * d < r.o * d ? r.o : o:
          P = (o + d) * d < (r.o + r.d) * d ? o + d : r.o
          if (s * d > e * d) return false;
          return *out = L(s, e - s), true;
      T s = inter(r), t = -r.inter(*this);
       if (z < 0) s = -s, t = -t, z = -z;
       if (s >= -EPS && s <= z + EPS && t >= -EPS && t <= z
          return *out = L(o + d * s / z, P()), true;
       return false:
   // check if the given point is on the segment
   bool point_on_seg(P r) const {
       if (abs(side(r)) > EPS) return false;
       if ((r - o) * d < -EPS) return false;</pre>
       if ((r - o - d) * d > EPS) return false:
       return true;
   // get the point in this line that is closest to a given
        point
   P closest_to(P r) const {
       P dr = d.rot(): return r + (o - r) * dr * dr / d.
           magsq();
}:
```

3.6 point

```
struct P {
    T x, y;
    P(T x, T y) : x(x), y(y) {}
```

```
P() : P(0, 0) \{ \}
   friend ostream &operator<<(ostream &s, const P &r) {</pre>
       return s << r.x << " " << r.v:
   friend istream & operator >> (istream &s. P &r) { return s
        >> r.x >> r.y; }
   P operator+(P r) const { return \{x + r.x, y + r.y\}; \}
   P operator-(P r) const { return {x - r.x, y - r.y}; }
   P operator*(T r) const { return {x * r, y * r}; }
   P operator/(T r) const { return {x / r, v / r}: }
   P operator-() const { return {-x, -v}; }
   friend P operator*(T 1, P r) { return {1 * r.x, 1 * r.y};
   P rot() const { return {-v, x}; }
   T operator*(P r) const { return x * r.x + y * r.y; }
   T operator/(P r) const { return rot() * r: }
   T magsq() const { return x * x + y * y; }
   T mag() const { return sqrt(magsq()); }
   P unit() const { return *this / mag(); }
   bool half() const { return abs(y) <= EPS && x < -EPS || y | struct InConvex {
         < -EPS: }
   T angcmp(P r) const {
       int h = (int)half() - r.half();
       return h ? h : r / *this;
   bool operator==(P r) const { return abs(x - r.x) <= EPS</pre>
        && abs(v - r.v) <= EPS: }
   double angle() const { return atan2(y, x); }
   static P from_angle(double a) { return {cos(a), sin(a)};
};
```

3.7 polygon

```
// get the area of a simple polygon in ccw order
T area(const vector<P> &ps) {
   int N = ps.size();
   T a = 0;
   rep(i, N) a += (ps[i] - ps[0]) / (ps[(i + 1) % N] - ps[i
        ]);
   return a / 2;
}
```

```
// checks whether a point is inside a simple polygon
// returns -1 if inside, 0 if on border, 1 if outside
int in_poly(const vector<P> &ps, P p) {
   int N = ps.size(). w = 0:
   rep(i, N) {
       P = ps[i] - p, e = ps[(i + 1) \% N] - p;
       if (s == P()) return 0:
       if (s.y == 0 && e.y == 0) {
           if (\min(s.x. e.x) \le 0 \&\& 0 \le \max(s.x. e.x))
               return 0:
       } else {
           bool b = s.v < 0:
           if (b != (e.v < 0)) {
              Tz = s / e; if (z == 0) return 0;
              if (b == (z > 0)) w += b ? 1 : -1:
    return w ? -1 : 1;
// check if a point is in a convex polygon
   vector<P> ps;
   T 11. 1h. rl. rh:
   int N. m:
   // preprocess polygon
   InConvex(const vector<P> &p) : ps(p), N(ps.size()), m(0)
        {
       assert(N >= 2);
       rep(i, N) if (ps[i].x < ps[m].x) m = i;
       rotate(ps.begin(), ps.begin() + m, ps.end()):
       rep(i, N) if (ps[i].x > ps[m].x) m = i;
       11 = 1h = ps[0].y, r1 = rh = ps[m].y;
       for (P p : ps) {
           if (p.x == ps[0].x) 11 = min(11, p.y), 1h = max(
           if (p.x == ps[m].x) rl = min(rl, p.y), rh = max(
               rh. p.v):
       }
    InConvex() {}
   // check if point belongs in polygon
   // returns -1 if inside, 0 if on border, 1 if outside
   // O(log N)
```

```
int in_poly(P p) {
    if (p.x < ps[0].x || p.x > ps[m].x) return 1;
    if (p.x == ps[0].x) return p.y < 11 || p.y > 1h;
    if (p.x == ps[m].x) return p.y < rl || p.y > rh;
    int r = upper_bound(ps.begin(), ps.begin() + m, p,
        [](P a, P b) { return a.x < b.x; }) - ps.begin();
    T z = (ps[r - 1] - ps[r]) / (p - ps[r]); if (z >= 0)
        return !!z;
    r = upper_bound(ps.begin() + m, ps.end(), p,
        [](P a, P b) { return a.x > b.x; }) - ps.begin();
    z = (ps[r - 1] - ps[r % N]) / (p - ps[r % N]);
    if (z >= 0) return !!z; return -1;
}
```

11

3.8 sweep

```
#include "point.cpp"
// iterate over all pairs of points
// 'op' is called with all ordered pairs of different
    indices '(i, j)'
// additionally, the 'ps' vector is kept sorted by signed
    distance
// to the line formed by 'i' and 'i'
// for example, if the vector from 'i' to 'j' is pointing
// the 'ps' vector is sorted from smallest 'y' to largest 'y
// note that, because the 'ps' vector is sorted by signed
// 'j' is always equal to 'i + 1'
// this means that the amount of points to the left of the
    line is always 'N - i'
template <class OP>
void all pair points(vector<P> &ps. OP op) {
   int N = ps.size();
   sort(ps.begin(), ps.end(), [](P a, P b) {
      return make_pair(a.y, a.x) < make_pair(b.y, b.x);</pre>
   vector<pair<int, int>> ss:
   rep(i, N) rep(j, N) if (i != j) ss.push_back({i, j});
   stable_sort(ss.begin(), ss.end(), [&](auto a, auto b) {
       return (ps[a.second] - ps[a.first]).angle lt(ps[b.
           second] - ps[b.first]);
   vector<int> p(N); rep(i, N) p[i] = i;
   for (auto [i, j] : ss)
```

3.9 theorems

```
// Pick's theorem
// Simple polygon with integer vertices:
// A = I + B / 2 - 1
// A: Area of the polygon
// I: Integer points strictly inside the polygon
// B: Integer points on the boundary of the polygon
```

4 graph

4.1 bellman-ford

```
struct Edge { int u, v; ll w; };
// find distance from source node to all nodes.
// supports negative edge weights.
// returns true if a negative cycle is detected.
//
// time: 0(V E)
bool bellman_ford(int N, int s, vector<Edge> &E, vector<11>
    &D. vector<int> &P) {
   P.assign(N, -1), D.assign(N, INF), D[s] = 0;
   rep(i. N - 1) {
       bool f = true:
       rep(ei, E.size()) {
           auto &e = E[ei]:
          ll n = D[e.u] + e.w;
           if (D[e.u] < INF && n < D[e.v])</pre>
              D[e.v] = n, P[e.v] = ei, f = false;
       if (f) return false;
   }
   return true;
```

4.2 dinic

```
struct Edge { int u, v; ll c, f = 0; };
```

```
// maximum flow algorithm.
 // time: 0(E V^2)
 11
         O(E V^(2/3)) / O(E sqrt(E)) unit capacities
 11
         O(E sart(V))
                                     unit networks (hopcroft-
     karp)
 // unit network: c in {0, 1} and forall v. len(incoming(v))
      <= 1 or len(outgoing(v)) <= 1
 // min-cut: find all nodes reachable from the source in the
      residual graph
 struct Dinic {
    int N. s. t: vector<vector<int>> G:
    vector<Edge> E: vector<int> lvl. ptr:
    Dinic() {}
    Dinic(int N, int s, int t) : N(N), s(s), t(t), G(N) {}
    void add_edge(int u, int v, 11 c) {
        G[u].push_back(E.size()); E.push_back({u, v, c});
        G[v].push_back(E.size()); E.push_back({v, u, 0});
    11 push(int u, 11 p) {
        if (u == t || p <= 0) return p;</pre>
        while (ptr[u] < G[u].size()) {</pre>
            int ei = G[u][ptr[u]++];
            Edge &e = E[ei];
            if (lvl[e.v] != lvl[u] + 1) continue;
            11 a = push(e.v, min(e.c - e.f, p));
            if (a <= 0) continue; e.f += a, E[ei ^ 1].f -= a;
            return a;
        }
        return 0;
    }
    11 maxflow() {
        11 f = 0:
        while (true) {
            // bfs to build levels
            lvl.assign(N, -1); queue < int > q; lvl[s] = 0, q.
                 push(s);
            while (!q.empty()) {
               int u = q.front(); q.pop();
               for (int ei : G[u]) {
                   Edge &e = E[ei]:
                   if (e.c - e.f <= 0 || lvl[e.v] != -1)</pre>
                        continue;
                   lvl[e.v] = lvl[u] + 1, q.push(e.v);
               }
            if (lvl[t] == -1) break:
```

4.3 floyd-warshall

```
// O(V^3) time and O(V^2) memory.

// requires an NxN array to store results.

// works with negative edges, but not negative cycles.

void floyd(const vector<vector<pair<int, ll>>> &G, vector<
    vector<ll>>> &dists) {
    int N = G.size();
    rep(i, N) rep(j, N) dists[i][j] = i == j ? 0 : INF;
    rep(i, N) for (auto edge : G[i]) dists[i][edge.first] =
        edge.second;
    rep(k, N) rep(i, N) rep(j, N)
        dists[i][j] = min(dists[i][j], dists[i][k] + dists[k
        ][j]);
}
```

4.4 heavy-light

```
struct Hld {
   vector<int> P, H, D, pos, top;
   void init(vector<vector<int>> &G) {
       int N = G.size();
      P.resize(N), H.resize(N), D.resize(N), pos.resize(N),
          top.resize(N);
       D[0] = -1, dfs(G, 0); int t = 0;
      rep(i, N) if (H[P[i]] != i) {
          int j = i;
          while (i != -1)
              \{ top[i] = i, pos[i] = t++; j = H[i]; \}
   }
   int dfs(vector<vector<int>> &G. int i) {
      int w = 1, mw = 0;
       D[i] = D[P[i]] + 1, H[i] = -1;
      for (int c : G[i]) {
          if (c == P[i]) continue;
```

```
P[c] = i: int sw = dfs(G, c): w += sw:
          if (sw > mw) H[i] = c, mw = sw;
      }
       return w;
   template <class OP>
   void path(int u, int v, OP op) {
       while (top[u] != top[v]) {
          if (D[top[u]] > D[top[v]]) swap(u, v);
          op(pos[top[v]], pos[v] + 1); v = P[top[v]];
      if (D[u] > D[v]) swap(u, v);
       op(pos[u], pos[v] + 1); // value on vertex
      // op(pos[u]+1, pos[v] + 1); // value on path
   // segment tree
   template <class T. class S>
   void update(S &seg, int i, T val) {
       seg.update(pos[i], val);
   // segment tree lazy
   template <class T, class S>
   void update(S &seg, int u, int v, T val) {
      path(u, v, [&](int 1, int r) { seg.update(1, r, val);
            }):
   }
   template <class T, class S>
   T query(S &seg, int u, int v) {
      T ans = 0;
           // neutral element
      path(u, v, [\&](int 1, int r) \{ ans += seg.query(1, r) \}
           : }): // guerv op
       return ans:
};
```

4.5 hungarian

```
// find a maximum gain perfect matching in the given
    bipartite complete graph.
// input: gain matrix (G_{xy} = benefit of joining vertex x
    in set X with vertex
// y in set Y).
// output: maximum gain matching in members 'xy[x]' and 'yx[
    y]'.
```

```
// runtime: O(N^3)
struct Hungarian {
   int N, qi, root;
   vector<vector<ll>>> gain;
   vector<int> xy, yx, p, q, slackx;
   vector<ll> lx. lv. slack:
   vector<bool> S, T;
   void add(int x, int px) {
       S[x] = true, p[x] = px;
       rep(y, N) if (lx[x] + ly[y] - gain[x][y] < slack[y])
          slack[y] = lx[x] + ly[y] - gain[x][y], slackx[y]
               = x:
      }
   }
   void augment(int x, int y) {
       while (x != -2) {
          yx[y] = x; swap(xy[x], y); x = p[x];
   }
   void improve() {
       S.assign(N, false), T.assign(N, false), p.assign(N,
            -1);
       qi = 0, q.clear();
       rep(x, N) if (xv[x] == -1) {
          q.push_back(root = x), p[x] = -2, S[x] = true;
          break:
       rep(y, N) slack[y] = lx[root] + ly[y] - gain[root][y
           1. slackx[v] = root:
       while (true) {
          while (ai < a.size()) {</pre>
              int x = q[qi++];
              rep(y, N) if (lx[x] + ly[y] == gain[x][y] &&!
                  T[v]) {
                  if (yx[y] == -1) return augment(x, y);
                  T[y] = true, q.push_back(yx[y]), add(yx[y
                      ], x);
              }
          }
          11 d = INF:
          rep(y, N) if (!T[y]) d = min(d, slack[y]);
          rep(x, N) if (S[x]) lx[x] -= d;
          rep(y, N) if (T[y]) ly[y] += d;
          rep(y, N) if (!T[y]) slack[y] -= d;
```

13

4.6 kuhn

```
// get a maximum cardinality matching in a bipartite graph.
// input: adjacency lists.
// output: matching (in 'mt' member).
// runtime: O(V E)
struct Kuhn {
   int N, size;
   vector<vector<int>> G:
   vector<bool> seen:
   vector<int> mt;
   bool visit(int i) {
       if (seen[i]) return false;
       seen[i] = true:
       for (int to : G[i])
          if (mt[to] == -1 || visit(mt[to])) {
              mt[to] = i:
              return true;
       return false;
   Kuhn(vector<vector<int>> adj) : G(adj), N(G.size()), mt(N
        , -1) {
       rep(i, N) {
          seen.assign(N, false);
           size += visit(i):
      }
```

};

4.7 lca

```
// calculates the lowest common ancestor for any two nodes
     in O(log N) time,
// with O(N log N) preprocessing
struct Lca {
   int L:
   vector<vector<int>> up;
   vector<pair<int, int>> time;
   Lca() {}
   void init(const vector<vector<int>> &G) {
       int N = G.size(); L = N <= 1 ? 0 : 32 - __builtin_clz // time: O(F V E)</pre>
            (N - 1);
       up.resize(L + 1); rep(l, L + 1) up[l].resize(N);
       time.resize(N); int t = 0; visit(G, 0, 0, t);
       rep(1, L) rep(i, N) up[1 + 1][i] = up[1][up[1][i]];
   void visit(const vector<vector<int>> &G, int i, int p,
        int &t) {
       up[0][i] = p;
       time[i].first = t++:
       for (int edge : G[i]) {
           if (edge == p) continue;
           visit(G, edge, i, t);
       time[i].second = t++:
   bool is_anc(int up, int dn) {
       return time[up].first <= time[dn].first &&
              time[dn].second <= time[up].second;</pre>
   int get(int i, int i) {
       if (is_anc(i, j)) return i;
       if (is_anc(j, i)) return j;
       int 1 = L:
       while (1 >= 0) {
           if (is_anc(up[1][i], j)) 1--;
           else i = up[1][i];
       }
       return up[0][i];
};
```

4.8 maxflow-mincost

```
// untested
#include "../common.h"
const ll INF = 1e18;
struct Edge {
   int u, v;
   11 c. w. f = 0:
// find the minimum-cost flow among all maximum-flow flows.
                          F is the maximum flow
       O(V E + F E log V) if bellman-ford is replaced by
    johnson
struct Flow {
   int N, s, t;
   vector<vector<int>> G;
   vector<Edge> E:
   vector<ll> d;
   vector<int> p;
   Flow() {}
   Flow(int N, int s, int t) : N(N), s(s), t(t), G(N) {}
   void add_edge(int u, int v, ll c, ll w) {
      G[u].push_back(E.size());
       E.push_back({u, v, c, w});
       G[v].push_back(E.size());
       E.push back(\{v, u, 0, -w\}):
   }
   void calcdists() {
       // replace bellman-ford with johnson for better time
       d.assign(N, INF);
      p.assign(N, -1);
      d[s] = 0:
       rep(i, N - 1) rep(ei, E.size()) {
          Edge &e = E[ei]:
          ll n = d[e.u] + e.w:
          if (d[e.u] < INF && e.c - e.f > 0 && n < d[e.v])
               d[e.v] = n, p[e.v] = ei;
      }
   }
   ll maxflow() {
      11 \text{ ff} = 0;
```

```
while (true) {
          calcdists():
          if (p[t] == -1) break;
          11 f = INF;
          int cur = t:
          while (p[cur] != -1) {
              Edge &e = E[p[cur]];
              f = min(f, e.c - e.f);
              cur = e.u;
          int cur = t:
          while (p[cur] != -1) {
              E[p[cur]].f += f;
              E[p[cur] ^ 1].f -= f;
          ff += f:
       }
       return ff;
};
```

4.9 push-relabel

```
#include "../common.h"
const 11 INF = 1e18;
// maximum flow algorithm.
// to run, use 'maxflow()'.
//
// time: O(V^2 \operatorname{sart}(E)) \leq O(V^3)
// memory: 0(V^2)
struct PushRelabel {
   vector<vector<ll>> cap. flow:
   vector<ll> excess;
   vector<int> height;
   PushRelabel() {}
   void resize(int N) { cap.assign(N, vector<11>(N)); }
   // push as much excess flow as possible from u to v.
   void push(int u, int v) {
       ll f = min(excess[u], cap[u][v] - flow[u][v]);
       flow[u][v] += f;
       flow[v][u] -= f:
       excess[v] += f;
```

```
excess[u] -= f:
// relabel the height of a vertex so that excess flow may :
     be pushed.
void relabel(int u) {
   int d = INT32_MAX;
   rep(v, cap.size()) if (cap[u][v] - flow[u][v] > 0) d
       min(d, height[v]);
   if (d < INF) height[u] = d + 1;</pre>
// get the maximum flow on the network specified by 'cap'
     with source 's'
// and sink 't'.
// node-to-node flows are output to the 'flow' member.
11 maxflow(int s, int t) {
   int N = cap.size(), M:
   flow.assign(N, vector<ll>(N)):
   height.assign(N, 0), height[s] = N;
   excess.assign(N, 0), excess[s] = INF;
   rep(i, N) if (i != s) push(s, i);
   vector<int> q;
   while (true) {
       // find the highest vertices with excess
       q.clear(), M = 0:
       rep(i, N) {
           if (excess[i] <= 0 || i == s || i == t)</pre>
                continue;
           if (height[i] > M) q.clear(), M = height[i];
           if (height[i] >= M) q.push_back(i);
       if (q.empty()) break;
       // process vertices
       for (int u : q) {
          bool relab = true;
           rep(v, N) {
              if (excess[u] <= 0) break:</pre>
              if (cap[u][v] - flow[u][v] > 0 && height[u]
                   ] > height[v])
                  push(u, v), relab = false:
           if (relab) {
              relabel(u):
              break;
       }
   }
```

```
ll f = 0; rep(i, N) f += flow[i][t]; return f;
};
```

4.10 strongly-connected-components

```
// compute strongly connected components.
// time: O(V + E), memory: O(V)
// after building:
// comp = map from vertex to component (components are
    toposorted, root first, leaf last)
// N = number of components
// G = condensation graph (component DAG)
11
// byproducts:
// vgi = transposed graph
// order = reverse topological sort (leaf first, root last)
// others:
// vn = number of vertices
// vg = original vertex graph
struct Scc {
   int vn, N;
   vector<int> order, comp;
   vector<vector<int>> vg, vgi, G;
   void toposort(int u) {
       if (comp[u]) return;
       comp[u] = -1;
       for (int v : vg[u]) toposort(v);
       order.push_back(u);
   bool carve(int u) {
       if (comp[u] != -1) return false:
       comp[u] = N;
       for (int v : vgi[u]) {
           carve(v):
           if (comp[v] != N) G[comp[v]].push_back(N);
       return true;
   }
   Scc(vector<vector<int>> &g) : vn(g.size()), vg(g), comp(
        vn), vgi(vn), G(vn), N(0) {
       rep(u, vn) toposort(u);
```

```
rep(u, vn) for (int v : vg[u]) vgi[v].push_back(u);
    invrep(i, vn) N += carve(order[i]);
};
```

4.11 two-sat

```
// calculate the solvability of a system of logical
    equations, where every equation is of the form 'a or b
// 'neg': get negation of 'u'
// 'then': 'u' implies 'v'
// 'any': 'u' or 'v'
// 'set': 'u' is true
//
// after 'solve' (O(V+E)) returns true, 'sol' contains one
    possible solution.
// determining all solutions is O(V*E) hard (requires
    computing reachability in a DAG).
struct TwoSat {
   int N: vector<vector<int>> G:
   Scc scc; vector<bool> sol;
   TwoSat(int n) : N(n), G(2 * n), sol(n) {}
   TwoSat() {}
   int neg(int u) { return (u + N) \% (2 * N); }
   void then(int u, int v) { G[u].push back(v), G[neg(v)].
        push_back(neg(u)); }
   void any(int u, int v) { then(neg(u), v); }
   void set(int u) { G[neg(u)].push_back(u); }
   bool solve() {
       scc = Scc(G);
       rep(u, N) if (scc.comp[u] == scc.comp[neg(u)]) return
       rep(u, N) sol[u] = (scc.comp[u] > scc.comp[neg(u)]);
       return true:
   }
};
```

5 implementation

5.1 dsu

```
struct Dsu {
  vector<int> p, r;
```

```
// initialize the disjoint-set-union to all unitary sets
   void reset(int N) {
       p.resize(N), r.assign(N, 0);
       rep(i, N) p[i] = i;
   // find the leader node corresponding to node 'i'
   int find(int i) {
       if (p[i] != i) p[i] = find(p[i]);
       return p[i]:
   // perform union on the two sets that 'i' and 'j' belong
   void unite(int i, int j) {
      i = find(i), j = find(j);
       if (i == j) return;
       if (r[i] > r[i]) swap(i, i):
       if (r[i] == r[j]) r[j] += 1;
       p[i] = j;
};
```

$5.2 \quad \text{mo}$

```
struct Query { int 1, r, idx; };
// answer segment queries using only 'add(i)', 'remove(i)'
    and 'get()'
// functions.
// complexity: O((N + Q) * sqrt(N) * F)
// N = length of the full segment
// Q = amount of queries
// F = complexity of the 'add', 'remove' functions
template <class A. class R. class G. class T>
void mo(vector<Query> &queries, vector<T> &ans, A add, R
    remove, G get) {
   int Q = queries.size(), B = (int)sqrt(Q);
   sort(queries.begin(), queries.end(), [&](Query &a, Query
       return make_pair(a.1 / B, a.r) < make_pair(b.1 / B, b</pre>
            .r);
   }):
   ans.resize(Q);
   int 1 = 0, r = 0:
   for (auto &q : queries) {
```

```
while (r < q.r) add(r), r++;
    while (l > q.l) l--, add(l);
    while (r > q.r) r--, remove(r);
    while (l < q.l) remove(l), l++;
    ans[q.idx] = get();
}</pre>
```

5.3 persistent-segment-tree-lazy

```
template <class T>
struct Node {
   T x. lz:
   int 1 = -1, r = -1;
template <class T>
struct Pstl {
   int N;
   vector<Node<T>> a;
   vector<int> head:
   T qneut() { return 0; }
   T merge(T 1, T r) { return 1 + r; }
   T uneut() { return 0; }
   T accum(T u, T x) { return u + x; }
   T apply(T x, T lz, int l, int r) { return x + (r - 1) *
   int build(int vl, int vr) {
       if (vr - vl == 1) a.push_back({qneut(), uneut()}); //
            node construction
       else {
           int vm = (vl + vr) / 2, l = build(vl, vm), r =
               build(vm. vr):
           a.push_back({merge(a[1].x, a[r].x), uneut(), 1, r
               }): // guerv merge
       return a.size() - 1:
   }
   T query(int 1, int r, int v, int v1, int vr, T acc) {
       if (1 >= vr || r <= vl) return gneut();</pre>
            // query neutral
       if (1 \le v1 \&\& r \ge vr) return apply(a[v].x. acc. vl.
             vr); // update op
       acc = accum(acc, a[v].lz);
            // update merge
       int vm = (vl + vr) / 2;
```

```
return merge(query(1, r, a[v].1, v1, vm, acc), query(
        1, r, a[v].r, vm, vr, acc)); // query merge
}
int update(int 1, int r, T x, int v, int v1, int vr) {
   if (1 >= vr || r <= vl || r <= 1) return v:
   a.push_back(a[v]);
   v = a.size() - 1:
   if (1 <= v1 && r >= vr) {
       a[v].x = apply(a[v].x, x, vl, vr); // update op
       a[v].lz = accum(a[v].lz, x); // update merge
   } else {
       int vm = (vl + vr) / 2;
       a[v].1 = update(1, r, x, a[v].1, v1, vm);
       a[v].r = update(1, r, x, a[v].r, vm, vr);
       a[v].x = merge(a[a[v].1].x, a[a[v].r].x); //
            query merge
   }
   return v:
7
Pst1() {}
Pstl(int N) : N(N) { head.push_back(build(0, N)); }
T query(int t, int 1, int r) {
   return query(1, r, head[t], 0, N, uneut()); // update
         neutral
int update(int t, int 1, int r, T x) {
   return head.push back(update(1, r, x, head[t], 0, N))
        , head.size() - 1;
```

5.4 persistent-segment-tree

```
// usage:
// Pst<Node<11>> pst;
// pst = {N};
// int newtime = pst.update(time, index, value);
// Node<11> result = pst.query(newtime, left, right);

template <class T>
struct Node {
   T x;
   int 1 = -1, r = -1;

   Node() : x(0) {}
   Node(T x) : x(x) {}
```

```
Node(Node a, Node b, int l = -1, int r = -1): x(a.x + b.
        x), 1(1), r(r) {}
};
template <class U>
struct Pst {
    int N:
    vector<U> a:
    vector<int> head:
    int build(int vl. int vr) {
       if (vr - vl == 1) a.push back(U()): // node
            construction
       else {
           int vm = (vl + vr) / 2, l = build(vl, vm), r =
               build(vm, vr);
           a.push_back(U(a[1], a[r], 1, r)); // query merge
       return a.size() - 1:
    U query(int 1, int r, int v, int v1, int vr) {
       if (1 >= vr || r <= vl) return U(); // query neutral</pre>
       if (1 <= v1 && r >= vr) return a[v]:
       int vm = (vl + vr) / 2:
       return U(query(1, r, a[v].1, v1, vm), query(1, r, a[v
            ].r, vm, vr)); // query merge
   }
    int update(int i. U x. int v. int vl. int vr) {
       a.push_back(a[v]);
       v = a.size() - 1;
       if (vr - vl == 1) a[v] = x: // update op
       else {
           int vm = (vl + vr) / 2:
           if (i < vm) a[v].l = update(i, x, a[v].l, vl, vm)</pre>
           else a[v].r = update(i, x, a[v].r, vm, vr);
           a[v] = U(a[a[v].1], a[a[v].r], a[v].1, a[v].r);
                // query merge
       }
       return v;
    Pst() {}
    Pst(int N) : N(N) { head.push_back(build(0, N)); }
    U query(int t, int 1, int r) {
       return query(1, r, head[t], 0, N);
```

5.5 segment-tree-lazy

```
// 0-based, inclusive-exclusive
// usage:
// St13<11> a;
// a = {N};
template <class T>
struct St13 {
   // immediate, lazv
   vector<pair<T, T>> a;
   T gneutral() { return 0: }
   T merge(T 1, T r) { return 1 + r; }
   T uneutral() { return 0; }
   void update(pair<T, T> &u, T val, int 1, int r) { u.first
         += val * (r - 1). u.second += val: }
   St13() {}
   Stl3(int N) : a(4 * N, {gneutral(), uneutral()}) {} //
        node neutral
   void push(int v, int vl, int vm, int vr) {
       update(a[2 * v], a[v].second, vl, vm); // node update
       update(a[2 * v + 1], a[v].second, vm, vr); // node
            update
       a[v].second = uneutral():
                                              // update
            neutral
   // query for range [1, r)
   T querv(int 1, int r, int v = 1, int vl = 0, int vr = -1)
       if (vr == -1) vr = a.size() / 4:
       if (1 <= v1 && r >= vr) return a[v].first; // query
       if (1 >= vr || r <= vl) return qneutral(); // query</pre>
            neutral
       int vm = (vl + vr) / 2;
       push(v. vl. vm. vr):
       return merge(query(1, r, 2 * v, v1, vm), query(1, r,
            2 * v + 1, vm, vr)); // item merge
   }
```

```
// update range [1, r) using val
   void update(int 1, int r, T val, int v = 1, int vl = 0,
        int vr = -1) {
       if (vr == -1) vr = a.size() / 4;
       if (1 >= vr || r <= vl || r <= l) return;</pre>
       if (1 \le vl \&\& r \ge vr) update(a[v], val, vl, vr): //
             node update
       else {
          int vm = (vl + vr) / 2:
          push(v, v1, vm, vr);
          update(1, r, val, 2 * v, vl, vm):
          update(1, r, val, 2 * v + 1, vm, vr):
          a[v].first = merge(a[2 * v].first, a[2 * v + 1].
               first): // node merge
      }
   }
}:
struct Node {
   ll x. lazv:
   Node() : x(neutral()), lazy(0) {} // query neutral,
        update neutral
   Node(ll x_-) : Node() { x = x_-; }
   Node(Node &1, Node &r) : Node() { refresh(1, r); } //
        node merge construction
   void refresh(Node &1, Node &r) { x = merge(1.x, r.x); }
        // node merge
   void update(ll val. int l. int r) { x += val * (r - 1).
        lazy += val; } // update-guery, update accumulate
   11 take() {
       11 z = 0: // update neutral
       swap(lazy, z);
       return z:
   11 query() { return x; }
   static ll neutral() { return 0; }
                                             // query
   static ll merge(ll l, ll r) { return l + r; } // query
}:
template <class T, class Node>
struct Stl {
   vector<Node> node;
   void reset(int N) { node.assign(4 * N, {}); } // node
        neutral
```

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```
void build(const vector\langle T \rangle &a. int v = 1. int vl = 0. int
   node.resize(4 * a.size()), vr = vr == -1 ? node.size
        () / 4 : vr:
   if (vr - vl == 1) {
       node[v] = {a[vl]}; // node construction
       return:
   int vm = (vl + vr) / 2;
   build(a, 2 * v, v1, vm);
   build(a, 2 * v + 1, vm, vr):
   node[v] = {node[2 * v], node[2 * v + 1]}; // node
        merge construction
void push(int v, int vl, int vm, int vr) {
   T lazy = node[v].take();
                                     // update neutral
   node[2 * v].update(lazv, vl, vm): // node update
   node[2 * v + 1].update(lazy, vm, vr); // node update
// query for range [1, r)
T query(int 1, int r, int v = 1, int vl = 0, int vr = -1)
   if (vr == -1) vr = node.size() / 4;
   if (1 <= vl && r >= vr) return node[v].query(); //
        query op
   if (1 >= vr || r <= vl) return Node::neutral(); //</pre>
        query neutral
   int vm = (vl + vr) / 2;
   push(v, v1, vm, vr);
   return Node::merge(query(1, r, 2 * v, v1, vm), query(
        1, r, 2 * v + 1, vm, vr)); // item merge
// update range [1, r) using val
void update(int 1, int r, T val, int v = 1, int vl = 0,
    int vr = -1) {
   if (vr == -1) vr = node.size() / 4:
   if (1 >= vr || r <= vl || r <= 1) return:
   if (1 <= vl && r >= vr) node[v].update(val, vl, vr);
        // node update
   else {
       int vm = (vl + vr) / 2;
       push(v. vl. vm. vr):
       update(1, r, val, 2 * v, vl, vm);
       update(1, r, val, 2 * v + 1, vm, vr);
       node[v].refresh(node[2 * v], node[2 * v + 1]); //
            node merge
```

```
}
};
```

5.6 segment-tree

```
// usage:
// St<Node<11>> st:
// st = {N}:
// st.update(index. new value):
// Node<ll> result = st.query(left, right);
template <class T>
struct Node {
   T x:
   Node(): x(0) {}
   Node(T x) : x(x) \{ \}
   Node(Node a. Node b) : x(a.x + b.x) {}
};
template <class U>
struct St {
   vector<U> a:
   St() {}
   St(int N) : a(4 * N, U()) {} // node neutral
   // query for range [1, r)
   U query(int 1, int r, int v = 1, int vl = 0, int vr = -1)
       if (vr == -1) vr = a.size() / 4:
       if (1 <= v1 && r >= vr) return a[v]: // item
            construction
       int vm = (v1 + vr) / 2;
       if (1 >= vr || r <= vl) return U();</pre>
                                        // item neutral
       return U(query(1, r, 2 * v, v1, vm), query(1, r, 2 *
            v + 1, vm, vr)); // item merge
   // set element i to val
   void update(int i, U val, int v = 1, int vl = 0, int vr =
         -1) {
       if (vr == -1) vr = a.size() / 4;
       if (vr - vl == 1) a[v] = val; // item update
       else {
           int vm = (v1 + vr) / 2;
           if (i < vm) update(i, val, 2 * v, vl, vm);</pre>
           else update(i, val, 2 * v + 1, vm, vr);
```

```
a[v] = U(a[2 * v], a[2 * v + 1]); // node merge
}
};
```

5.7 sparse-table

```
// handle immutable range maximum queries (or any idempotent
     query) in O(1)
template <class T>
struct Sparse {
   vector<vector<T>> st:
   T op(T a, T b) { return max(a, b); }
   Sparse() {}
   void reset(int N) { st = {vector<T>(N)}: }
   void set(int i, T val) { st[0][i] = val; }
   // O(N log N) time
   // O(N log N) memory
   void init() {
       int N = st[0].size();
       int npot = N <= 1 ? 1 : 32 - __builtin_clz(N);</pre>
       st.resize(npot);
       repx(i, 1, npot) rep(j, N + 1 - (1 << i)) st[i].
           push back(
          op(st[i-1][i], st[i-1][i+(1 << (i-1))]));
                // query op
   // query maximum in the range [1, r) in O(1) time
   // range must be nonempty!
   T query(int 1, int r) {
       int i = 31 - __builtin_clz(r - 1);
       return op(st[i][1], st[i][r - (1 << i)]); // query op</pre>
};
```

5.8 unordered-map

```
// hackproof rng
static mt19937 rng(chrono::steady_clock::now().
    time_since_epoch().count());
// deterministic rng
```

```
uint64 t splitmix64(uint64 t *x) {
   uint64 t z = (*x += 0x9e3779b97f4a7c15):
   z = (z ^ (z >> 30)) * 0xbf58476d1ce4e5b9;
   z = (z ^ (z >> 27)) * 0x94d049bb133111eb:
   return z^(z >> 31);
}
// hackproof unordered map hash
struct Hash {
   size_t operator()(const 11 &x) const {
       static const uint64 t RAND =
           chrono::steadv clock::now().time since epoch().
               count();
       uint64_t z = x + RAND + 0x9e3779b97f4a7c15;
       z = (z ^ (z >> 30)) * 0xbf58476d1ce4e5b9:
       z = (z \hat{z} > 27) * 0x94d049bb133111eb;
       return z (z >> 31):
}:
// hackproof unordered_map
template <class T, class U>
using umap = unordered_map<T, U, Hash>;
// hackproof unordered_set
template <class T>
using uset = unordered_set<T, Hash>;
```

6 imprimible

7 math

7.1 arithmetic

```
// floor(log2(n)) without precision loss
inline int floor_log2(int n) { return n <= 1 ? 0 : 31 -
    __builtin_clz(n); }
// ceil(log2(n)) without precision loss
inline int ceil_log2(int n) { return n <= 1 ? 0 : 32 -
    __builtin_clz(n - 1); }
inline ll floordiv(ll a, ll b) {
    return a / b - ((a ^ b) < 0 && a % b);
}
inline ll ceildiv(ll a, ll b) {
    return a / b + ((a ^ b) >= 0 && a % b);
}
```

```
}

// a^e through binary exponentiation.

ll binexp(ll a, ll e) {
    ll res = 1; // neutral element
    while (e) {
        if (e & 1) res = res * a; // multiplication
            a = a * a; // multiplication
        e >>= 1;
    }
    return res;
}
```

7.2 crt

```
pair<11, 11> solve_crt(const vector<pair<11, 11>> &eqs) {
    11 a0 = eqs[0].first, p0 = eqs[0].second;
    repx(i, 1, eqs.size()) {
        11 a1 = eqs[i].first, p1 = eqs[i].second;
        11 k1, k0;
        11 d = ext_gcd(p1, p0, k1, k0);
        a0 -= a1;
        if (a0 % d != 0) return {-1, -1};
        p0 = p0 / d * p1;
        a0 = a0 / d * k1 % p0 * p1 % p0 + a1;
        a0 = (a0 % p0 + p0) % p0;
    }
    return {a0, p0};
}
```

7.3 discrete-log

```
// discrete logarithm log_a(b).
// solve b ^ x = a (mod M) for the smallest x.
// returns -1 if no solution is found.
//
// time: O(sqrt(M))
ll dlog(ll a, ll b, ll M) {
    ll k = 1, s = 0;
    while (true) {
        ll g = __gcd(b, M);
        if (g <= 1) break;
        if (a == k) return s;
        if (a % g != 0) return -1;
        a /= g, M /= g, s += 1, k = b / g * k % M;
}
ll N = sqrt(M) + 1;</pre>
```

```
umap<11, 11> r;
rep(q, N + 1) {
    r[a] = q;
    a = a * b % M;
}

ll bN = binexp(b, N, M), bNp = k;
repx(p, 1, N + 1) {
    bNp = bNp * bN % M;
    if (r.count(bNp)) return N * p - r[bNp] + s;
}
return -1;
```

7.4 gauss

```
const double EPS = 1e-9:
// solve a system of equations.
// complexity: O(min(N, M) * N * M)
//
// 'a' is a list of rows
// the last value in each row is the result of the equation
// return values:
// 0 -> no solutions
// 1 -> unique solution, stored in 'ans'
// -1 -> infinitely many solutions, one of which is stored
    in 'ans'
// UNTESTED
int gauss(vector<vector<double>> a, vector<double> &ans) {
   int N = a.size(), M = a[0].size() - 1:
   vector<int> where(M, -1):
   for (int j = 0, i = 0; j < M && i < N; j++) {
       int sel = i;
       repx(k, i, N) if (abs(a[k][j]) > abs(a[sel][j])) sel
            = k;
       if (abs(a[sel][j]) < EPS) continue;</pre>
       repx(k, j, M + 1) swap(a[sel][k], a[i][k]);
       where[i] = i;
       rep(k, N) if (k != i) {
          double c = a[k][j] / a[i][j];
          repx(1, j, M + 1) a[k][1] -= a[i][1] * c;
       }
       i++:
   }
```

7.5 matrix

```
using T = 11;
struct Mat {
   int N. M:
   vector<vector<T>> v;
   Mat(int n, int m) : N(n), M(m), v(N, vector<T>(M)) {}
   Mat(int n) : Mat(n, n) { rep(i, N) v[i][i] = 1; }
   vector<T> &operator[](int i) { return v[i]; }
   Mat operator*(Mat &r) {
      assert(M == r.N);
      int n = N, m = r.M, p = M;
      Mat a(n. m):
      rep(i, n) rep(j, m) {
          a[i][j] = T();
               neutral
          rep(k, p) a[i][k] = a[i][j] + v[i][k] * r[k][j];
               // mul. add
      }
       return a;
   Mat binexp(ll e) {
      assert(N == M):
      Mat a = *this, res(N); // neutral
      while (e) {
          if (e & 1) res = res * a; // mul
          a = a * a;
                                  // mul
          e >>= 1:
      }
       return res;
```

```
friend ostream &operator<<(ostream &s, Mat &a) {
    rep(i, a.N) {
        rep(j, a.M) s << a[i][j] << " ";
        s << endl;
    }
    return s;
}</pre>
```

7.6 mod

```
11 binexp(ll a, ll e, ll M) {
    assert(e >= 0):
   ll res = 1 % M:
    while (e) {
       if (e & 1) res = res * a % M;
       a = a * a % M:
       e >>= 1;
   }
   return res:
ll multinv(ll a, ll M) { return binexp(a, M - 2, M); }
// calculate gcd(a, b).
// also, calculate x and y such that:
// a * x + b * y == gcd(a, b)
// time: O(log min(a, b))
// (ignoring complexity of arithmetic)
ll ext gcd(ll a, ll b, ll &x, ll &v) {
   if (b == 0) {
       x = 1, y = 0;
       return a;
   11 d = ext_gcd(b, a \% b, y, x);
   y = a / b * x;
   return d:
// compute inverse with any M.
// a and M must be coprime for inverse to exist!
11 multinv_euc(ll a, ll M) {
   11 x. v:
   ext_gcd(a, M, x, y);
   return x:
```

```
// multiply two big numbers (~10^18) under a large modulo.
    without resorting to
// bigints.
11 bigmul(l1 x, l1 y, l1 M) {
   11 z = 0:
   while (v) {
       if (y \& 1) z = (z + x) \% M;
      x = (x << 1) \% M, y >>= 1;
   return z;
// all modular inverses from 1 to inv.size()-1
void multinv all(vector<ll> &inv) {
   inv[1] = 1:
   repx(i, 2, inv.size())
       inv[i] = MOD - (MOD / i) * inv[MOD % i] % MOD:
struct Mod {
   int a;
   static const int M = 1e9 + 7;
   Mod(11 aa) : a((aa % M + M) % M) {}
   Mod operator+(Mod rhs) const { return (a + rhs.a) % M; }
   Mod operator-(Mod rhs) const { return (a - rhs.a + M) % M
        : }
   Mod operator-() const { return Mod(0) - *this; }
   Mod operator*(Mod rhs) const { return (11)a * rhs.a % M:
   Mod operator+=(Mod rhs) { return *this = *this + rhs: }
   Mod operator = (Mod rhs) { return *this = *this - rhs; }
   Mod operator*=(Mod rhs) { return *this = *this * rhs; }
   Mod bigmul(ll big) const { return ::bigmul(a, big, M); }
   Mod binexp(ll e) const { return ::binexp(a, e, M); }
   // Mod multinv() const { return ::multinv(a, M); } //
   Mod multinv() const { return ::multinv_euc(a, M); } //
        possibly composite M
};
// dvnamic modulus
struct DMod {
   int a. M:
   DMod(11 aa, 11 m) : M(m), a((aa % m + m) % m) {}
```

```
DMod operator+(DMod rhs) const { return {(a + rhs.a) % M,
         M}: }
   DMod operator-(DMod rhs) const { return {(a - rhs.a + M)
        % M. M}: }
   DMod operator-() const { return DMod(0, M) - *this: }
   DMod operator*(DMod rhs) const { return {(11)a * rhs.a %
        M, M}; }
   DMod operator+=(DMod rhs) { return *this = *this + rhs; }
   DMod operator==(DMod rhs) { return *this = *this - rhs; }
   DMod operator*=(DMod rhs) { return *this = *this * rhs: }
   DMod bigmul(ll big) const { return {::bigmul(a, big, M),
        M}: }
   DMod binexp(ll e) const { return {::binexp(a, e, M), M}; }
   DMod multinv() const { return {::multinv(a, M), M}: } //
   // DMod multinv() const { return {::multinv_euc(a, M), M
       }; } // possibly composite M
};
```

7.7 poly

```
using cd = complex<double>;
const double PI = acos(-1):
// compute the DFT of a power-of-two-length sequence.
// if 'inv' is true, computes the inverse DFT.
// the DFT of a polynomial A(x) = A0 + A1*x + A2*x^2 + ... +
      An*x^n is the array
// of the polynomial A evaluated in all nths roots of unity:
      \Gamma A(w0), A(w1),
// A(w2), ..., A(wn-1)], where w0 = 1 and w1 is the nth
     principal root of unity.
void fft(vector<cd> &a, bool inv) {
    int N = a.size(), k = 0;
    assert(N == 1 << __builtin_ctz(N));</pre>
    rep(i, N) {
       int b = N \gg 1;
       while (k \& b) k = b, b >>= 1:
       k ^= b:
       if (i < k) swap(a[i], a[k]);</pre>
```

```
for (int 1 = 2: 1 <= N: 1 <<= 1) {
       double ang = 2 * PI / 1 * (inv ? -1 : 1);
       cd wl(cos(ang), sin(ang));
       for (int i = 0: i < N: i += 1) {</pre>
           cd w(1);
           repx(i, 0, 1 / 2)  {
              cd u = a[i + j], v = a[i + j + 1 / 2] * w;
              a[i + j] = u + v;
              a[i + j + 1 / 2] = u - v;
              w *= w1;
      }
   }
   if (inv)
       for (cd &x : a) x \neq N;
const 11 MOD = 7340033, ROOT = 5, ROOTPOW = 1 << 20;</pre>
void find_root_of_unity(l1 M) {
   11 c = M - 1, k = 0:
   while (c \% 2 == 0) c /= 2, k += 1;
   // find proper divisors of M - 1
   vector<int> divs;
   repx(d, 1, c) {
       if (d * d > c) break:
       if (c % d == 0) rep(i, k + 1) divs.push_back(d << i);
   rep(i, k) divs.push_back(c << i);</pre>
   // find any primitive root of M
   11 G = -1;
   repx(g, 2, M) {
       bool ok = true:
       for (int d : divs) ok &= (binexp(g, d, M) != 1);
       if (ok) {
           G = g;
          break;
       }
   }
   assert(G != -1):
   ll w = binexp(G, c, M);
   cerr << M << " = c * 2^k + 1" << endl:
   cerr << " c = " << c << endl;
   cerr << " k = " << k << endl:
   cerr << w^(2^k) == 1 << endl:
   cerr << " w = " << w << endl:
```

```
// compute the DFT of a power-of-two-length sequence, modulo
     a special prime
// number with principal root.
// the modulus _must_ be a prime number with an Nth root of
    unity, where N is a
// power of two. the FFT can only be performed on arrays of
void ntt(vector<ll> &a. bool inv) {
   int N = a.size(), k = 0:
   assert(N == 1 << __builtin_ctz(N) && N <= ROOTPOW);</pre>
   rep(i, N) a[i] = (a[i] \% MOD + MOD) \% MOD;
   repx(i, 1, N) {
       int b = N \gg 1:
       while (k & b) k ^= b, b >>= 1;
      k ^= b:
       if (i < k) swap(a[i], a[k]);</pre>
   for (int 1 = 2; 1 <= N; 1 <<= 1) {
       11 wl = inv ? multinv(ROOT, MOD) : ROOT;
       for (ll i = ROOTPOW; i > 1; i >>= 1) wl = wl * wl %
       for (int i = 0: i < N: i += 1) {
          11 w = 1:
          repx(i, 0, 1 / 2) {
              11 u = a[i + i], v = a[i + i + 1 / 2] * w %
              a[i + j] = (u + v) \% MOD;
              a[i + i + 1 / 2] = (u - v + MOD) \% MOD:
              w = w * wl % MOD;
      }
   11 ninv = multinv(N, MOD);
   if (inv)
       for (11 &x : a) x = x * ninv % MOD:
void convolve(vector<ll> &a, vector<ll> b, int n) {
   n = 1 \ll (32 - \_builtin\_clz(2 * n - 1));
   a.resize(n), b.resize(n):
   ntt(a, false), ntt(b, false);
   rep(i, n) a[i] *= b[i];
   ntt(a, true), ntt(b, true);
```

```
using T = 11:
T pmul(T a, T b) { return a * b % MOD; }
T padd(T a, T b) { return (a + b) % MOD; }
T psub(T a, T b) { return (a - b + MOD) % MOD; }
T pinv(T a) { return multinv(a, MOD): }
struct Poly {
   vector<T> a;
   Polv() {}
   Polv(T c) : a(c) { trim(): }
   Poly(vector<T> c) : a(c) { trim(); }
   void trim() {
       while (!a.empty() && a.back() == 0) a.pop_back();
   int deg() const { return a.empty() ? -1000000 : a.size()
        - 1: }
   Polv sub(int 1, int r) const {
       r = min(r, (int)a.size()), l = min(l, r);
       return vector<T>(a.begin() + 1, a.begin() + r);
   Poly trunc(int n) const { return sub(0, n); }
   Poly shl(int n) const {
       Poly out = *this;
       out.a.insert(out.a.begin(), n, 0);
       return out:
   Polv rev(int n. bool r = false) const {
       Polv out(*this);
       if (r) out.a.resize(max(n, (int)a.size()));
       reverse(out.a.begin(), out.a.end());
       return out.trunc(n);
   Poly &operator+=(const Poly &rhs) {
       auto &b = rhs.a:
       a.resize(max(a.size(), b.size()));
       rep(i, b.size()) a[i] = padd(a[i], b[i]); // add
       trim():
       return *this;
   Poly &operator-=(const Poly &rhs) {
       auto &b = rhs.a;
       a.resize(max(a.size(), b.size())):
       rep(i, b.size()) a[i] = psub(a[i], b[i]); // sub
       trim():
       return *this:
```

```
Polv &operator *= (const Polv &rhs) {
   int n = deg() + rhs.deg() + 1;
   if (n <= 0) return *this = Poly();</pre>
   n = 1 \ll (n \ll 1?0:32 - builtin clz(n - 1)):
   vector<T> b = rhs.a:
   a.resize(n), b.resize(n):
   ntt(a, false), ntt(b, false);
                                          // fft
   rep(i, a.size()) a[i] = pmul(a[i], b[i]); // mul
   ntt(a, true), trim();
                                          // invfft
   return *this;
Polv inv(int n) const {
   assert(deg() >= 0);
   Poly ans = pinv(a[0]); // inverse
   int b = 1:
   while (b < n) {
       Polv C = (ans * trunc(2 * b)).sub(b, 2 * b):
       ans -= (ans * C).trunc(b).shl(b);
       b *= 2:
   return ans.trunc(n);
Poly operator+(const Poly &rhs) const { return Poly(*this
    ) += rhs: }
Poly operator-(const Poly &rhs) const { return Poly(*this
     ) -= rhs: }
Poly operator*(const Poly &rhs) const { return Poly(*this
    ) *= rhs; }
pair<Poly, Poly> divmod(const Poly &b) const {
   if (deg() < b.deg()) return {Poly(), *this};</pre>
   int d = deg() - b.deg() + 1:
   Poly D = (rev(d) * b.rev(d).inv(d)).trunc(d).rev(d,
        true):
   return {D, *this - D * b}:
Poly operator/(const Poly &b) const { return divmod(b).
Poly operator%(const Poly &b) const { return divmod(b).
Poly &operator/=(const Poly &b) { return *this = divmod(b
    ).first: }
Poly &operator%=(const Poly &b) { return *this = divmod(b
    ).second: }
T \text{ eval}(T x)  {
   T v = 0:
   invrep(i, a.size()) y = padd(pmul(y, x), a[i]); //
```

```
return v:
   }
   Poly &build(vector<Poly> &tree, vector<T> &x, int v, int
        1. int r) {
       if (1 == r) return tree[v] = vectorT - x[1], 1};
       int m = (1 + r) / 2:
       return tree[v] = build(tree, x, 2 * v, 1, m) *
                      build(tree, x, 2 * v + 1, m + 1, r);
   void subeval(vector<Poly> &tree, vector<T> &x, vector<T>
        &v. int v. int 1.
               int r) {
       if (1 == r) {
          y[1] = eval(x[1]);
          return:
       int m = (1 + r) / 2:
       (*this % tree[2 * v]).subeval(tree, x, y, 2 * v, 1, m
       (*this % tree[2 * v + 1]).subeval(tree, x, v, 2 * v +
            1, m + 1, r);
   // evaluate m points in O(k (\log k)^2) with k = \max(n, m)
   vector<T> multieval(vector<T> &x) {
       int N = x.size();
       if (deg() < 0) return vector<T>(N, 0);
       vector<Polv> tree(4 * N):
      build(tree, x, 1, 0, N - 1);
       vector<T> v(N):
       subeval(tree, x, y, 1, 0, N - 1);
       return v;
   }
};
```

7.8 primes

```
// counts the divisors of a positive integer in O(sqrt(n))
ll count_divisors(ll x) {
    ll divs = 1, i = 2;
    for (ll divs = 1, i = 2; x > 1; i++) {
        if (i * i > x) {
            divs *= 2;
            break;
        }
        for (ll d = divs; x % i == 0; x /= i) divs += d;
    }
    return divs;
}
```

```
// gets the prime factorization of a number in O(\operatorname{sqrt}(n))
vector<pair<11, int>> factorize(11 x) {
   vector<pair<11. int>> f:
   for (11 k = 2; x > 1; k++) {
       if (k * k > x) {
           f.push_back({x, 1});
           break:
       int n = 0;
       while (x \% k == 0) x /= k, n++:
       if (n > 0) f.push back(\{k, n\}):
   return f;
// iterate over all divisors of a number.
11
// divisor count upper bound: n^(1.07 / ln ln n)
template <class OP>
void divisors(11 x, OP op) {
   auto facts = factorize(x):
   vector<int> f(facts.size());
   while (true) {
       11 y = 1;
       rep(i, f.size()) rep(j, f[i]) y *= facts[i].first;
       op(y);
       int i;
       for (i = 0: i < f.size(): i++) {</pre>
           f[i] += 1;
           if (f[i] <= facts[i].second) break;</pre>
       if (i == f.size()) break;
```

```
// computes euler totative function phi(x), counting the
    amount of integers in
// [1, x] that are coprime with x.
// time: O(sqrt(x))
11 phi(11 x) {
   11 phi = 1, k = 2;
   for (; x > 1; k++) {
       if (k * k > x) {
          phi *= x - 1;
           break:
      11 k1 = 1, k0 = 0;
       while (x \% k == 0) x /= k, k0 = k1, k1 *= k;
       phi *= k1 - k0:
   return phi;
// computes primality up to N.
// considers 0 and 1 prime.
// O(N log N)
void sieve(int N, vector<bool> &prime) {
   prime.assign(N + 1, true);
   repx(n, 2, N + 1) if (prime[n]) for (int k = 2 * n; k <=
        N; k += n) prime[k] = false;
```

7.9 theorems

```
// Burnside lemma
//
// For a set X, with members x in X, and a group G, with
operations g in G, where g(x): X -> X.
//
// F_g is the set of x which are fixed points of g (ie. {
    x in X / g(x) = x }).
```

```
The number of orbits (connected components in the
    graph formed by assigning each x a node and
    a directed edge between x and g(x) for every g) is
     M = the average of the fixed points of all g = (|F_g1|)
     + |F g2| + ... + |F gn|) / |G|
11
   If x are images and g are simmetries, then M
    corresponds to the amount of objects, |G|
     corresponds to the amount of simmetries, and F_g
    corresponds to the amount of simmetrical
     images under the simmetry g.
11
// Rational root theorem
     All rational roots of the polynomials with integer
    coefficients:
11
      a0 * x^0 + a1 * x^1 + a2 * x^2 + ... + an * x^n = 0
11
//
     If these roots are represented as p / q, with p and q
     - p is an integer factor of a0
11
     - q is an integer factor of an
11
     Note that if a0 = 0, then x = 0 is a root, the
    polynomial can be divided by x and the theorem
     applies once again.
11
// Legendre's formula
     Considering a prime p, the largest power p^k that
    divides n! is given by:
     k = floor(n/p) + floor(n/p^2) + floor(n/p^3) + ...
      Which can be computed in O(\log n / \log p) time
```