Team Notebook

April 10, 2024

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1 common

1.1 common

```
#pragma GCC optimize("Ofast")
#pragma GCC target("bmi,bmi2,lzcnt,popcnt")
#pragma GCC target("avx,avx2,f16c,fma,sse3,sse3,sse4.1,sse4
#include <bits/stdc++.h>
using namespace std;
typedef long long 11;
#define repx(i, a, b) for (int i = a; i < b; i++)
#define rep(i, n) repx(i, 0, n)
#define invrepx(i, a, b) for (int i = b - 1; i \ge a; i--)
#define invrep(i, n) invrepx(i, 0, n)
// Command to check time and memory usage:
       /usr/bin/time -v ./tmp
// See "Maximum resident set size" for max memory used
// Commands for interactive checker:
       mkfifo fifo
       (./solution < fifo) | (./interactor > fifo)
// Does not work on the Windows file system. i.e., /mnt/c/
// The special fifo file must be used, otherwise the
// solution will not wait for input and will read EOF
```

2 data structures

2.1 Treap

```
mt19937 gen(chrono::high_resolution_clock::now().
    time_since_epoch().count());

// 101 Implicit Treap //

struct Node {
    int p, sz = 0, v, acc, l = -1, r = -1;
    Node(): v(0), acc(0) {}
    Node(int x): p(gen()), sz(1), v(x), acc(x) {}
    void recalc(const Node &a, const Node &b) {
        sz = a.sz + b.sz + 1;
        acc = v + a.acc + b.acc;
    }
};

template <class node>
struct Treap {
    vector<node> t;
```

```
int n. r = -1:
   node get(int u) { return u != -1 ? t[u] : node(); }
    void recalc(int u) { t[u].recalc(get(t[u].1), get(t[u].r)
        ); }
    int merge(int 1. int r) {
       if (min(1, r) == -1) return 1 != -1 ? 1 : r;
       int ans = (t[1].p < t[r].p) ? 1 : r;
       if (ans == 1) t[1].r = merge(t[1].r, r), recalc(1);
       if (ans == r) t[r].l = merge(l, t[r].l), recalc(r);
       return ans:
   pii split(int u, int id) {
       if (u == -1) return {-1, -1};
       int szl = get(t[u].1).sz;
       if (szl >= id) {
           pii ans = split(t[u].1, id);
           t[u].1 = ans.ss;
           recalc(u):
           return {ans.ff, u};
       pii ans = split(t[u].r, id - szl - 1);
       t[u].r = ans.ff;
       recalc(u):
       return {u, ans.ss};
   Treap(vi &v) : n(sz(v)) {
       for (int i = 0; i < n; i++) t.eb(v[i]), r = merge(r, left)
   }
};
// Complete Implicit Treap with Lazy propagation //
   int p, sz = 0, v, acc, l = -1, r = -1, par = -1, lzv = 0;
   bool lz = false, f = false;
   Node() : v(0), acc(0) {}
   Node(int x) : p(gen()), sz(1), v(x), acc(x) {}
   void recalc(const Node &a. const Node &b) {
       sz = a.sz + b.sz + 1;
       acc = v + a.acc + b.acc:
   void upd_lazy(int x) { lz = 1, lzv += x; }
   void lazy() { v += lzv, acc += sz * lzv, lz = 0, lzv = 0;
   void flip() { swap(1, r), f = 0; }
};
```

```
template <class node>
struct Treap {
   vector<node> t;
   int n. r = -1:
   node get(int u) { return u != -1 ? t[u] : node(): }
   void recalc(int u) {
       int 1 = t[u].1, r = t[u].r;
       push(1), push(r), flip(1), flip(r);
      t[u].recalc(get(1), get(r));
   void push(int u) {
       if (u == -1 || !t[u].lz) return;
       int 1 = t[u].1, r = t[u].r;
      if (1 != -1) t[1].upd_lazy(t[u].lzv);
      if (r != -1) t[r].upd_lazy(t[u].lzv);
      t[u].lazv():
   }
   void flip(int u) {
       if (u == -1 || !t[u].f) return:
      int 1 = t[u].1, r = t[u].r;
      if (1 != -1) t[1].f ^= 1:
      if (r != -1) t[r].f ^= 1;
       t[u].flip();
   }
   int merge(int 1, int r) {
       if (min(l, r) == -1) return l != -1 ? l : r;
       push(1), push(r), flip(1), flip(r);
       int ans = (t[1].p < t[r].p) ? 1 : r;
       if (ans == 1) t[1].r = merge(t[1].r, r), recalc(1);
       if (ans == r) t[r].l = merge(l, t[r].l), recalc(r);
       if (t[ans].l != -1) t[t[ans].l].par = ans; // only if
            parent needed
       if (t[ans].r != -1) t[t[ans].r].par = ans; // only if
            parent needed
       return ans:
   pii split(int u, int id) {
       if (u == -1) return \{-1, -1\};
       push(u):
      flip(u):
       int szl = get(t[u].1).sz;
       if (szl >= id) {
          pii ans = split(t[u].1, id);
          if (ans.ss != -1) t[ans.ss].par = u; // only if
               parent needed
          if (ans.ff != -1) t[ans.ff].par = -1; // only if
               parent needed
          t[u].1 = ans.ss:
          recalc(u):
```

```
return {ans.ff. u}:
      pii ans = split(t[u].r, id - szl - 1);
       if (ans.ff != -1) t[ans.ff].par = u; // only if
            parent needed
      if (ans.ss != -1) t[ans.ss].par = -1; // only if
            parent needed
      t[u].r = ans.ff;
       recalc(u):
       return {u, ans.ss};
   int update(int u, int 1, int r, int v) {
      pii a = split(u, 1), b = split(a.ss, r - 1 + 1);
      t[b.ff].upd_lazy(v);
       return merge(a.ff, merge(b.ff, b.ss));
   void print(int u) {
      if (u == -1) return;
      push(u), flip(u):
      print(t[u].1):
      cout << t[u].v << ' ';
      print(t[u].r);
   Treap(vi &v) : n(sz(v)) {
      for (int i = 0; i < n; i++) t.eb(v[i]), r = merge(r, left)
   }
};
```

2.2 dsu

```
struct Dsu {
    vector<int> p; Dsu(int N = 0) : p(N, -1) {}
    int get(int x) { return p[x] < 0 ? x : get(p[x]); }
    bool sameSet(int a, int b) { return get(a) == get(b); }
    int size(int x) { return -p[get(x)]; }
    vector<vector<int>> S;
    void unite(int x, int y) {
        if ((x = get(x)) == (y = get(y)))
            return S.push_back({-1});
        if (p[x] > p[y]) swap(x, y);
        S.push_back({x, y, p[x], p[y]});
        p[x] += p[y], p[y] = x;
    }
    void rollback() {
        auto a = S.back(); S.pop_back();
        if (a[0] != -1) p[a[0]] = a[2], p[a[1]] = a[3];
    }
}
```

```
2.3 fenwick-tree
```

```
int ft[MAXN+1]; // add dimension for multi-d
void upd(int i0, int v){ // add v to i0th element
for(int i=i0+1;i<=MAXN;i+=i&-i)ft[i]+=v;//+ fors
}
int get(int i0){ // get sum of range [0,i0)
int r=0; // add fors
for(int i=i0;i;i-=i&-i)r+=ft[i];
return r;
}
int get_sum(int i0,int i1){//sum of [i0,i1)
return get(i1)-get(i0);</pre>
```

2.4 link-cut-tree

```
const int N_DEL = 0, N_VAL = 0; //delta, value
inline int mOp(int x, int y){return x+y;}//modify
inline int gOp(int lval, int rval){return lval + rval;}//
inline int dOnSeg(int d. int len){return d==N DEL ? N DEL :
    d*len;}
//mostly generic
inline int joinD(int d1, int d2){
 if(d1==N_DEL)return d2;if(d2==N_DEL)return d1;return m0p(
      d1. d2):}
inline int joinVD(int v. int d){return d==N DEL ? v : mOp(v.
     d):}
struct Node t{
 int sz, nVal, tVal, d; bool rev;
 Node_t *c[2], *p;
 Node_t(int v) : sz(1), nVal(v), tVal(v), d(N_DEL), rev(0),
       }(0)q
 c[0]=c[1]=0:
 }
 bool isRoot(){return !p || (p->c[0] != this && p->c[1] !=
      this):}
 void push(){
 if(rev){
   rev=0: swap(c[0], c[1]): fore(x,0,2)if(c[x])c[x]->rev^=1:
 nVal=joinVD(nVal, d); tVal=joinVD(tVal, dOnSeg(d, sz));
 fore(x,0,2)if(c[x])c[x]->d=joinD(c[x]->d, d);
 d=N DEL:
```

```
void upd();
typedef Node_t* Node;
int getSize(Node r){return r ? r->sz : 0;}
int getPV(Node r){
 return r ? joinVD(r->tVal, dOnSeg(r->d,r->sz)) : N_VAL;}
void Node_t::upd(){
 tVal = qOp(qOp(getPV(c[0]), joinVD(nVal, d)), getPV(c[1]))
  sz = 1 + getSize(c[0]) + getSize(c[1]):
void conn(Node c, Node p, int il){if(c)c->p=p;if(il>=0)p->c
     [!ill=c:}
void rotate(Node x){
 Node p = x-p, g = p-p;
 bool gCh=p->isRoot(), isl = x==p->c[0];
  conn(x->c[isl],p,isl); conn(p,x,!isl);
  conn(x,g,gCh?-1:(p==g->c[0])); p->upd();
void spa(Node x){//splay
 while(!x->isRoot()){
 Node p = x->p, g = p->p;
  if(!p->isRoot())g->push();
  p->push(); x->push();
 if(!p-)isRoot())rotate((x=-p-)c[0])==(p=-p-)c[0])? p : x);
  rotate(x):
 x->push(); x->upd();
Node exv(Node x){//expose
 Node last=0;
  for(Node y=x; y; y=y->p)spa(y),y->c[0]=last,y->upd(),last=
  spa(x);
 return last:
void mkR(Node x){exv(x);x->rev^=1;}//makeRoot
Node getR(Node x){exv(x); while(x->c[1])x=x->c[1]; spa(x);
    return x:}
Node lca(Node x. Node v){exv(x): return exv(v):}
bool connected(Node x, Node y){exv(x);exv(y); return x==y?1:
     x - p! = 0:
void link(Node x, Node y){mkR(x); x->p=y;}
void cut(Node x, Node y){mkR(x); exv(y); y\rightarrow c[1]\rightarrow p=0; y\rightarrow c
     [1]=0:}
Node father(Node x){
 exv(x): Node r=x->c[1]:
if(!r)return 0:
 while (r->c[0])r=r->c[0];
```

2.5 persistent-segment-tree-lazy

```
template <class T>
struct Node {
   T x. lz:
    int 1 = -1, r = -1;
};
template <class T>
struct Pstl {
    int N;
    vector<Node<T>> a;
    vector<int> head:
    T aneut() { return 0: }
    T merge(T 1, T r) { return 1 + r; }
    T uneut() { return 0; }
    T accum(T u, T x) { return u + x; }
    T apply(T x, T lz, int l, int r) { return x + (r - 1) *
        lz: }
    int build(int vl, int vr) {
       if (vr - vl == 1) a.push_back({qneut(), uneut()}); // }:
             node construction
       else {
           int vm = (vl + vr) / 2, l = build(vl, vm), r =
               build(vm, vr);
           a.push_back({merge(a[1].x, a[r].x), uneut(), 1, r
                }); // query merge
```

```
return a.size() - 1;
}
T query(int 1, int r, int v, int v1, int vr, T acc) {
   if (1 >= vr || r <= vl) return gneut();</pre>
        // guery neutral
   if (1 <= vl && r >= vr) return apply(a[v].x, acc, vl,
         vr); // update op
   acc = accum(acc, a[v].lz);
        // update merge
   int vm = (v1 + vr) / 2;
   return merge(query(1, r, a[v].1, v1, vm, acc), query(
        1, r, a[v].r, vm, vr, acc)); // query merge
int update(int 1, int r, T x, int v, int v1, int vr) {
   if (1 >= vr || r <= vl || r <= 1) return v:
   a.push_back(a[v]);
   v = a.size() - 1:
   if (1 <= v1 && r >= vr) {
       a[v].x = apply(a[v].x, x, vl, vr); // update op
       a[v].lz = accum(a[v].lz, x); // update merge
   } else {
       int vm = (vl + vr) / 2;
       a[v].1 = update(1, r, x, a[v].1, vl, vm);
       a[v].r = update(1, r, x, a[v].r, vm, vr);
       a[v].x = merge(a[a[v].1].x, a[a[v].r].x); //
            query merge
   return v;
}
Pst1() {}
Pstl(int N) : N(N) { head.push_back(build(0, N)); }
T query(int t, int 1, int r) {
   return query(1, r, head[t], 0, N, uneut()); // update
         neutral
int update(int t, int 1, int r, T x) {
   return head.push back(update(1, r, x, head[t], 0, N))
        , head.size() - 1;
```

2.6 persistent-segment-tree

```
// usage:
// Pst<Node<11>> pst;
```

```
// pst = {N}:
// int newtime = pst.update(time, index, value);
// Node<ll> result = pst.query(newtime, left, right);
template <class T>
struct Node {
   T x;
   int 1 = -1, r = -1;
   Node() : x(0) {}
   Node(T x) : x(x) \{ \}
   Node (Node a, Node b, int l = -1, int r = -1) : x(a.x + b.
        x), 1(1), r(r) {}
};
template <class U>
struct Pst {
   int N:
   vector<U> a:
   vector<int> head:
   int build(int vl. int vr) {
       if (vr - vl == 1) a.push_back(U());
           int vm = (v1 + vr) / 2, 1 = build(v1, vm),
              r = build(vm, vr);
          a.push_back(U(a[1], a[r], 1, r));
       return a.size() - 1;
   U query(int 1, int r, int v, int v1, int vr) {
       if (1 >= vr || r <= vl) return U();</pre>
       if (1 <= v1 && r >= vr) return a[v];
       int vm = (vl + vr) / 2:
       return U(query(1, r, a[v].1, v1, vm),
               query(1, r, a[v].r, vm, vr));
   }
   int update(int i, U x, int v, int vl, int vr) {
       a.push back(a[v]):
       v = a.size() - 1;
       if (vr - vl == 1) a[v] = x:
       else {
           int vm = (vl + vr) / 2;
          if (i < vm) a[v].l = update(i, x, a[v].l, vl, vm)</pre>
           else a[v].r = update(i, x, a[v].r, vm, vr);
           a[v] = U(a[a[v].1], a[a[v].r], a[v].1, a[v].r);
       }
```

2.7 rmq-lineal

```
typedef int tf; // O(n) construction, O(1) query
struct rmq{
int n; tf INF=1e9;//change sign of INF for MAX
vector<unsigned int> mk: vector<tf> bk.v:
tf op(tf a, tf b){return min(a,b);}//change for maximum
int f(int x){return x>>5:}
rmq(vector<tf> &vv):n(SZ(vv)),mk(n),bk(n,INF),v(vv){
 unsigned int 1st=0:
 for(int i=0;i<SZ(v);i++,lst<<=1){</pre>
  bk[f(i)]=op(bk[f(i)],v[i]);
  while(lst&&v[i-__builtin_ctz(lst)]>v[i]) lst^=lst&-lst;
  //while(lst&&v[i- builtin ctz(lst)]<v[i]) lst^=lst&-lst:
        //MAX
  mk[i]=++lst;
 for(int k=1,top=f(n);(1<<k)<=top;k++)fore(i,0,top)if(i</pre>
      +(1<<k)<=top)
  bk[top*k+i]=op(bk[top*(k-1)+i], bk[top*(k-1)+i+(1<< k-1)])
tf get(int st, int en){
 return v[en-31+__builtin_clz(mk[en]&((111<<en-st+1)-1))];</pre>
tf query(int s, int e){ //[s,e]
 int b1=f(s),b2=f(e),top=f(n);
 if(b1==b2) return get(s,e):
 tf ans=op(get(s,(b1+1)*32-1), get(b2*32,e)); s=(b1+1)*32;
      e=b2*32-1:
 if(s \le e)
  int k=31-_builtin_clz(f(e-s+1));
```

```
ans=op(ans,op(bk[top*k+f(s)],bk[top*k+f(e)-(1<<k)+1]));
}
return ans;
}
};</pre>
```

2.8 segment-tree-2d

```
// #define MAXN 1024 #define op(a,b) (a+b) #define NEUT 0
int n,m; int a[MAXN][MAXN],st[2*MAXN][2*MAXN];
void build(){
repx(i, 0, n) repx(j, 0, m) st[i+n][j+m] = a[i][j];
repx(i, 0, n) for(int i = m-1; i; --i)
 st[i+n][j] = op(st[i+n][j<<1], st[i+n][j<<1|1]);
for(int i = n-1; i; --i) repx(j, 0, 2*m)
 st[i][j] = op(st[i << 1][j], st[i << 1|1][j]);
void upd(int x, int v, int v){
st[x+n][y+m]=v;
for(int i = v+m: i > 1: i >>= 1)
      st[x+n][i>>1] = op(st[x+n][i], st[x+n][i^1]):
for(int i = x+n; i > 1; i >>= 1) for(int j=y+m; j; j>>=1)
 st[i>>1][j] = op(st[i][j], st[i^1][j]);
int query(int x0, int x1, int y0, int y1){
int r=NEUT:
for(int i0=x0+n, i1=x1+n; i0<i1; i0>>=1, i1>>=1){
 int t[4], q = 0;
 if(i0 \& 1) t[q++] = i0++:
 if(i1 \& 1) t[q++] = --i1;
 repx(k,0,q)
          for(int j0=y0+m, j1=y1+m; j0<j1; j0>>=1,j1>>=1){
              if(j0 \& 1) r = op(r, st[t[k]][j0++]);
              if(j1 \& 1) r = op(r, st[t[k]][--j1]);
return r;
```

2.9 segment-tree-beats

```
Node(const Node &a. const Node &b) {
       // add
       s = a.s + b.s;
       // min
       if (a.mx1 > b.mx1) mx1 = a.mx1, mxc = a.mxc, mx2 =
           max(b.mx1, a.mx2):
       if (a.mx1 < b.mx1) mx1 = b.mx1, mxc = b.mxc, mx2 =
            max(a.mx1. b.mx2):
       if (a.mx1 == b.mx1) mx1 = a.mx1, mxc = a.mxc + b.mxc
             mx2 = max(a.mx2, b.mx2);
       if (a.mn1 < b.mn1) mn1 = a.mn1, mnc = a.mnc, mn2 =</pre>
            min(b.mn1, a.mn2):
       if (a.mn1 > b.mn1) mn1 = b.mn1, mnc = b.mnc, mn2 =
            min(a.mn1, b.mn2):
       if (a.mn1 == b.mn1) mn1 = a.mn1, mnc = a.mnc + b.mnc
            mn2 = min(a.mn2, b.mn2):
   }
}:
// 0 - indexed / inclusive - inclusive
template <class node>
struct STB {
   vector<node> st; int n;
   void build(int u, int i, int j, vector<node> &arr) {
       if (i == i) {
          st[u] = arr[i]:
          return;
       int m = (i + j) / 2, l = u * 2 + 1, r = u * 2 + 2;
       build(l, i, m, arr), build(r, m + 1, j, arr);
       st[u] = node(st[l], st[r]):
   void push_add(int u, int i, int j, ll v) {
       st[u].s += (i - i + 1) * v:
       st[u].mx1 += v, st[u].mn1 += v, st[u].lz += v;
       if (st[u].mx2 != LLONG_MIN) st[u].mx2 += v;
       if (st[u].mn2 != LLONG_MAX) st[u].mn2 += v;
   void push max(int u, 11 v, bool 1) { // for min op
       if (v >= st[u].mx1) return;
       st[u].s -= st[u].mx1 * st[u].mxc:
       st[u].mx1 = v:
       st[u].s += st[u].mx1 * st[u].mxc;
       if (1) st[u].mn1 = st[u].mx1:
       else if (v <= st[u].mn1) st[u].mn1 = v;</pre>
       else if (v < st[u].mn2) st[u].mn2 = v:
   void push_min(int u, ll v, bool l) { // for max op
```

```
if (v <= st[u].mn1) return:</pre>
   st[u].s -= st[u].mn1 * st[u].mnc;
   st[u].mn1 = v;
   st[u].s += st[u].mn1 * st[u].mnc:
   if (1) st[u].mx1 = st[u].mn1;
   else if (v \ge st[u].mx1) st[u].mx1 = v:
   else if (v > st[u].mx2) st[u].mx2 = v;
void push(int u, int i, int j) {
   if (i == j) return;
   // add
   int m = (i + j) / 2, l = u * 2 + 1, r = u * 2 + 2;
   push_add(1, i, m, st[u].lz);
   push_add(r, m + 1, j, st[u].lz);
   st[u].lz = 0:
   // min
   push max(1. st[u].mx1. i == m):
   push_max(r, st[u].mx1, m + 1 == j);
   // max
   push min(l, st[u].mn1, i == m):
   push_min(r, st[u].mn1, m + 1 == r);
node query(int a, int b, int u, int i, int j) {
   if (b < i || j < a) return node();</pre>
   if (a <= i && j <= b) return st[u];</pre>
   push(u, i, j);
   int m = (i + j) / 2, l = u * 2 + 1, r = u * 2 + 2;
   return node(querv(a, b, l, i, m), querv(a, b, r, m +
        1, j));
void update_add(int a, int b, ll v, int u, int i, int j) };
   if (b < i | | i < a) return:
   if (a <= i && j <= b) {
       push_add(u, i, j, v);
       return:
   push(u, i, j);
   int m = (i + j) / 2, l = u * 2 + 1, r = u * 2 + 2;
   update_add(a, b, v, l, i, m);
   update add(a, b, v, r, m + 1, i):
   st[u] = node(st[1], st[r]);
void update_min(int a, int b, ll v, int u, int i, int j)
   if (b < i | | i < a | | v >= st[u].mx1) return:
   if (a <= i && j <= b && v > st[u].mx2) {
       push_max(u, v, i == j);
       return;
   }
```

```
push(u, i, i):
   int m = (i + i) / 2, 1 = u * 2 + 1, r = u * 2 + 2;
   update_min(a, b, v, l, i, m);
   update_min(a, b, v, r, m + 1, j);
   st[u] = node(st[1], st[r]);
void update_max(int a, int b, ll v, int u, int i, int j)
   if (b < i || j < a || v <= st[u].mn1) return;</pre>
   if (a <= i && j <= b && v < st[u].mn2) {</pre>
       push min(u, v, i == i):
       return:
   }
   push(u, i, j);
   int m = (i + j) / 2, l = u * 2 + 1, r = u * 2 + 2;
   update_max(a, b, v, l, i, m);
   update_max(a, b, v, r, m + 1, j);
   st[u] = node(st[1], st[r]);
STB(vector<node> &v, int N) : n(N), st(N * 4 + 5) { build
     (0, 0, n - 1, v): }
node query(int a, int b) { return query(a, b, 0, 0, n -
void update_add(int a, int b, ll v) { update_add(a, b, v,
     0, 0, n - 1); }
void update_min(int a, int b, ll v) { update_min(a, b, v, | };
     0.0.n - 1): 
void update_max(int a, int b, ll v) { update_max(a, b, v,
     0.0.n - 1):
```

2.10 segment-tree-lazy

```
template <class T>
struct St1 {
    int n; vector<T> a, b;
    St1(int n = 0) : n(n), a(4 * n, qneut()),
        b(4 * n, uneut()) {}

    T qneut() { return -2e9; }
    T uneut() { return 0; }
    T merge(T x, T y) { return max(x, y); }
    void upd(int v, T x, int l, int r)
        { a[v] += x, b[v] += x; }

    void push(int v, int v1, int vm, int vr) {
        upd(2 * v, b[v], v1, vm);
        upd(2 * v + 1, b[v], vm, vr);
    }
}
```

```
b[v] = uneut():
T querv(int 1, int r, int v=1, int v1=0, int vr=1e9) {
   vr = min(vr, n);
   if (1 <= v1 && r >= vr) return a[v];
   if (1 >= vr || r <= vl) return gneut();</pre>
   int vm = (vl + vr) / 2:
    push(v, v1, vm, vr);
   return merge(query(1, r, 2 * v, v1, vm),
       querv(1, r, 2 * v + 1, vm, vr)):
void update(int 1, int r, T x, int v = 1, int vl = 0,
       int vr = 1e9) {
   vr = min(vr, n);
   if (1 >= vr || r <= vl || r <= 1) return:
   if (1 \le v1 \&\& r \ge vr) upd(v, x, vl, vr);
       int vm = (vl + vr) / 2:
       push(v, v1, vm, vr);
       update(1, r, x, 2 * v, v1, vm);
       update(1, r, x, 2 * v + 1, vm, vr);
       a[v] = merge(a[2 * v], a[2 * v + 1]);
   }
```

2.11 segment-tree

```
struct St {
    ll neut() { return 0; }
    ll merge(ll x, ll y) { return x + y; }

    int n; vector<ll> a;
    St(int n = 0) : n(n), a(2 * n, neut()) {}

    ll query(int l, int r) {
        ll x = neut(), y = neut();
        for (1 += n, r += n; 1 < r; 1 /= 2, r /= 2) {
            if (1 & 1) x = merge(x, a[l++]);
            if (r & 1) y = merge(a[--r], y);
        }

    return merge(x, y);
    }

void update(int i, ll x) {
    for (a[i += n] = x; i /= 2;)
        a[i] = merge(a[2 * i], a[2 * i + 1]);</pre>
```

```
}
};
```

2.12 sparse-table

```
template <class T>
struct Sparse {
   T op(T a, T b) { return max(a, b); }
   vector<vector<T>> st:
   Sparse() {}
   Sparse(vector<T> a) : st{a} {
       int N = st[0].size();
       int npot = N <= 1 ? 1 : 32 - __builtin_clz(N);</pre>
       st.resize(npot);
       repx(i, 1, npot) rep(j, N + 1 - (1 << i))
       st[i].push_back(
           op(st[i-1][j], st[i-1][j+(1 << (i-1))])
       ): // guery op
   T query(int 1, int r) { // range must be nonempty!
       int i = 31 - __builtin_clz(r - 1);
       return op(st[i][1], st[i][r - (1 << i)]); // queryop</pre>
};
```

$3 ext{ dp}$

3.1 convex-hull-trick

```
struct Line {
   mutable ll a, b, c;

  bool operator<(Line r) const { return a < r.a; }
  bool operator<(ll x) const { return c < x; }
};

// dynamically insert 'a*x + b' lines and query for maximum
// at any x all operations have complexity O(log N)
struct LineContainer : multiset<Line, less<>>> {
   ll div(ll a, ll b) {
      return a / b - ((a ^ b) < 0 && a % b);
   }

  bool isect(iterator x, iterator y) {
    if (y == end()) return x->c = INF, 0;
```

```
if (x->a == v->a) x->c = x->b > v->b? INF : -INF:
       else x->c = div(y->b - x->b, x->a - y->a);
       return x->c >= y->c;
   void add(ll a. ll b) {
       // a *= -1, b *= -1 // for min
       auto z = insert(\{a, b, 0\}), y = z++, x = y;
       while (isect(y, z)) z = erase(z);
       if (x != begin() && isect(--x, y)) isect(x, y = erase
       while ((v = x) != begin() \&\& (--x)->c >= v->c) isect(
            x, erase(y));
   }
   11 query(11 x) {
       if (empty()) return -INF; // INF for min
       auto 1 = *lower_bound(x);
       return 1.a * x + 1.b:
       // return -l.a * x - l.b: // for min
   }
};
```

3.2 divide-and-conquer

```
// for every index i assign an optimal index j, such that
// cost(i, j) is minimal for every i. the property that if
// i2 >= i1 then j2 >= j1 is exploited (monotonic condition)
// calculate optimal index for all indices in range [1, r)
// knowing that the optimal index for every index in this
// range is within [optl, optr).
// time: O(N log N)
void calc(vector<int> &opt, int 1, int r,int optl,int optr){
   if (1 == r) return:
   int i = (1 + r) / 2:
   11 optc = INF;
   int optj;
   repx(j, optl, optr) {
      11 c = i + j; // cost(i, j)
       if (c < optc) optc = c, optj = j;</pre>
   opt[i] = optj;
   calc(opt, 1, i, optl, optj + 1);
   calc(opt, i + 1, r, optj, optr);
```

$4 \quad \text{geo2d}$

4.1 circle

```
struct C {
   Po: Tr:
   // circle-line intersection, assuming it exists
   // points are sorted along the direction of the line
   pair<P, P> line_inter(L 1) const {
      P c = 1.closest to(o): T c2 = (c - o).magsq():
      P = 1.d * sqrt(max(r*r - c2, T()) / 1.d.magsq());
      return {c - e, c + e};
   // check the type of line-circle collision
   // <0: 2 inters, =0: 1 inter, >0: 0 inters
   T line_collide(L 1) const {
      T c2 = (1.closest_to(o) - o).magsq();
       return c2 - r * r;
   }
   // calculates the two intersections between two circles
   // the circles must intersect in one or two points!
   pair<P. P> inter(C h) const {
      P d = h.o - o;
      T c = (r * r - h.r * h.r) / d.magsq();
       return h.line_inter({(1 + c) / 2 * d, d.rot()});
   // check if the given circles intersect
   bool collide(C h) const {
       return (h.o - o).magsq() \le (h.r + r) * (h.r + r):
   }
   // get one of the two tangents that go through the point
   // the point must not be inside the circle
   // a = -1: cw (relative to the circle) tangent
   // a = 1: ccw (relative to the circle) tangent
   P point_tangent(P p, T a) const {
      T c = r * r / p.magsq();
       return o + c*(p-o) - a*sqrt(c*(1-c))*(p-o).rot();
   // get one of the 4 tangents between the two circles
   // a = 1: exterior tangents
   // a = -1: interior tangents (requires no area overlap)
   // b = 1: ccw tangent
   // b = -1: cw tangent
   // the line origin is on this circumference, and the
```

```
// direction is a unit vector towards the other circle
L tangent(C c, T a, T b) const {
   T dr = a * r - c.r:
   P d = c.o - o:
   P n = (d*dr+b*d.rot()*sqrt(d.magsq()-dr*dr)).unit();
   return {o + n * r. -b * n.rot()}:
// circumcircle of a **non-degenerate** triangle
static C thru_points(P a, P b, P c) {
   b = b - a, c = c - a:
   P p = (b*c.magsq() - c*b.magsq()).rot() / (b%c*2):
   return {a + p, p.mag()};
// find the two circles that go through the given point,
// are tangent to the given line and have radius 'r'
// the point-line distance must be at most 'r'!
// the circles are sorted in the direction of the line
static pair<C, C> thru_point_line_r(P a, L t, T r) {
   P d = t.d.rot().unit();
   if (d * (a - t.o) < 0) d = -d:
   auto p = C(a, r).line_inter(\{t.o + d * r, t.d\});
   return {{p.first, r}, {p.second, r}};
// find the two circles that go through the given points
// and have radius 'r'
// circles sorted by angle from the first point
// the points must be at most at distance 'r'!
static pair<C, C> thru_points_r(P a, P b, T r) {
   auto p = C(a, r).line_inter({(a+b)/2, (b-a).rot()});
   return {{p.first, r}, {p.second, r}};
```

4.2 closest-points

};

```
// sort by x

11 closest(vector<ii> &p) {
    int n = SZ(p);
    set<ii> s;
    st = 1e18;
    int j = 0;
    fore(i, 0, n) {
        11 d = ceil(sqrt(best));
        while(p[i].fst - p[j].fst >= best)
            s.erase({p[j].snd, p[j].fst}), j++;
        auto it1=s.lower_bound({p[i].snd-d,p[i].fst});
    }
}
```

```
auto it2=s.upper_bound({p[i].snd+d,p[i].fst});
    for(auto it = it1; it != it2; ++it) {
        11 dx = p[i].fst - it->snd;
        11 dy = p[i].snd - it->fst;
        best = min(best, dx * dx + dy * dy);
    }
    s.insert({p[i].snd, p[i].fst});
}
return best;
}
```

4.3 convex-hull

```
// ccw order, excludes collinear points by default
vector<P> chull(vector<P> p) {
   if (p.size() < 3) return p;</pre>
   vectorP r; int m, k = 0;
   sort(p.begin(), p.end(), [](P a, P b) {
       return a.x != b.x ? a.x < b.x : a.y < b.y; });</pre>
   for (P a : p) { // lower hull
       while (k \ge 2 \&\& r[k - 1].left(r[k - 2], a) \ge 0)
          r.pop_back(), k--; // >= to > to add collinears
      r.push_back(q), k++;
   if (k == (int)p.size()) return r;
   r.pop_back(), k--, m = k;
   for (int i = p.size() - 1; i >= 0; --i) { // upper hull
       while (k \ge m+2 \&\& r[k-1].left(r[k-2], p[i]) \ge 0)
          r.pop back(), k--: // >= to > to add collinears
       r.push_back(p[i]), k++;
   r.pop_back(); return r;
```

4.4 delaunay

```
typedef __int128_t lll; // if on a 64-bit platform

struct Q {
    Q *rot, *o; P p = {INF, INF}; bool mark;
    P &F() { return r()->p; }
    Q *&r() { return rot->rot; }
    Q *prev() { return rot->o->rot; }
    Q *next() { return r()->prev(); }
};
T cross(P a, P b, P c) { return (b - a) % (c - a); }
```

```
bool circ(P p, P a, P b, P c) { // is p in the circumcircle?
   111 p2 = p.magsq(), A = a.magsq() - p2,
       B = b.magsq() - p2, C = c.magsq() - p2;
   return cross(p, a, b) * C + cross(p, b, c) * A + cross(p, b, c)
         c. a) * B > 0:
Q *makeEdge(Q *&H, P orig, P dest) {
   Q *r = H ? H : new Q{new Q{new Q{0}}};
   H = r -> 0: r -> r() -> r() = r:
   repx(i, 0, 4) r = r->rot, r->p = {INF, INF}.
       r->0 = i & 1 ? r : r->r();
   r\rightarrow p = orig; r\rightarrow F() = dest;
   return r:
void splice(Q *a, Q *b) {
   swap(a->o->rot->o, b->o->rot->o); swap(a->o, b->o);
Q *connect(Q *&H, Q *a, Q *b) {
   Q *q = makeEdge(H, a->F(), b->p);
   splice(q, a->next()); splice(q->r(), b); return q;
pair<Q *, Q *> rec(Q *&H, const vector<P> &s) {
   if (s.size() <= 3) {</pre>
       Q *a = makeEdge(H, s[0], s[1]), *b = makeEdge(H, s[0], s[1])
             [1]. s.back()):
       if (s.size() == 2) return \{a, a->r()\}; splice(a->r(),
       auto side = cross(s[0], s[1], s[2]):
       Q *c = side ? connect(H, b, a) : 0;
       return \{\text{side} < 0 ? c \rightarrow r() : a. \text{side} < 0 ? c : b \rightarrow r()\}
            }:
#define J(e) e \rightarrow F(), e \rightarrow p
#define valid(e) (cross(e->F(), J(base)) > 0)
   Q *A. *B. *ra. *rb: int half = s.size() / 2:
   tie(ra, A) = rec(H, {s.begin(), s.end() - half});
   tie(B, rb) = rec(H, {s.begin() + s.size() - half, s.end()
   while ((cross(B->p, J(A)) < 0 && (A = A->next())) ||
          (cross(A->p, J(B)) > 0 \&\& (B = B->r()->o)));
   Q *base = connect(H, B->r(), A);
   if (A->p == ra->p) ra = base->r();
   if (B->p == rb->p) rb = base;
```

```
#define DEL(e, init, dir) Q *e = init->dir: \
   if (valid(e)) while (circ(e->dir->F(), J(base), e->F()))
        { \
           Q *t = e->dir; splice(e, e->prev()); \
           splice(e->r(), e->r()->prev()); e->o = H; H = e;
               e = t: \
   for (;;) {
       DEL(LC, base->r(), o); DEL(RC, base, prev());
       if (!valid(LC) && !valid(RC)) break;
       if (!valid(LC) || (valid(RC) && circ(J(RC), J(LC))))
            base = connect(H, RC, base->r());
       else base = connect(H, base->r(), LC->r());
   return {ra, rb};
#undef J
#undef valid
#undef DEL
// there must be no duplicate points
// returns no triangles in the case of all collinear points
// produces counter-clockwise triangles ordered in triples
// maximizes the minimum angle across all triangulations
// the euclidean mst is a subset of these edges
// O(N log N)
vector<P> triangulate(vector<P> pts) {
   sort(pts.begin(), pts.end(), [](P a, P b) {
       return make_pair(a.x, a.y) < make_pair(b.x, b.y);</pre>
   }):
   assert(unique(pts.begin(), pts.end()) == pts.end());
   if (pts.size() < 2) return {};</pre>
   Q *H = 0; Q *e = rec(H, pts).first;
   vector<Q *> q = {e}; int qi = 0;
   while (cross(e->o->F(), e->F(), e->p) < 0) e = e->o:
#define ADD
   ł
       Q *c = e:
       do {
           c->mark = 1; pts.push_back(c->p); \
           g.push back(c->r()): c = c->next(): \setminus
       } while (c != e);
   ADD;
   pts.clear();
   while (qi < (int)q.size()) if (!(e = q[qi++])->mark) ADD;
   return pts;
#undef ADD
```

4.5 halfplane-intersect

```
// obtain the convex polygon that results from intersecting
    the given list
// of halfplanes, represented as lines that allow their left
// assumes the halfplane intersection is bounded
vector<P> halfplane intersect(vector<L> &H) {
   L bb(P(-INF, -INF), P(INF, 0));
   rep(k, 4) H.push_back(bb), bb.o = bb.o.rot(), bb.d = bb.d
        .rot():
   sort(begin(H), end(H), [](L a, L b) { return a.d.angcmp(b
        .d) < 0: }):
   deque<L> q; int n = 0;
   rep(i, H.size()) {
       while (n \ge 2 \&\& H[i].side(q[n - 1].intersection(q[n - 1]))
           -21)) > 0)
           g.pop back(), n--:
       while (n >= 2 && H[i].side(q[0].intersection(q[1])) >
           q.pop_front(), n--;
       if (n > 0 \&\& H[i].parallel(q[n - 1])) {
           if (H[i].d * q[n - 1].d < 0) return {};</pre>
           if (H[i].side(q[n - 1].o) > 0) q.pop_back(), n--;
           else continue:
       q.push_back(H[i]), n++;
   while (n \ge 3 \&\& q[0].side(q[n - 1].intersection(q[n -
        21)) > 0)
       q.pop_back(), n--;
   while (n \ge 3 \&\& q[n - 1].side(q[0].intersection(q[1])) >
       q.pop_front(), n--;
   if (n < 3) return {}:</pre>
   vector<P> ps(n);
   rep(i, n) ps[i] = q[i].intersection(q[(i + 1) % n]);
   return ps;
```

4.6 line

```
// a segment or an infinite line
// does not handle point segments correctly!
struct L {
   P o, d;
```

```
static L from_eq(P ab, T c) {
   return L{ab.rot(), ab * -c / ab.magsq()};
pair<P, T> line_eq() { return {-d.rot(), d.rot() * o}; }
// on which side of the line is the point
// negative: left, positive: right
T side(P r) const { return (r - o) % d; }
// returns the intersection coefficient
// in the range [0, d % r.d]
// if d % r.d is zero, the lines are parallel
T inter(L r) const { return (r.o - o) % r.d; }
// get the single intersection point
// lines must not be parallel
P intersection(L r) const {return o+d*inter(r)/(d%r.d);}
// check if lines are parallel
bool parallel(L r) const { return abs(d % r.d) <= EPS; }</pre>
// check if segments intersect
bool seg_collide(L r) const {
   Tz = d \% r.d:
   if (abs(z) \le EPS) {
       if (abs(side(r.o)) > EPS) return false;
       T s = (r.o - o) * d. e = s + r.d * d:
       if (s > e) swap(s, e);
       return s <= d * d + EPS && e >= -EPS:
   T s = inter(r), t = -r.inter(*this);
   if (z < 0) s = -s, t = -t, z = -z;
   return s>=-EPS && s<=z+EPS && t>=-EPS && t<=z+EPS;
// full segment intersection
// makes a point segment if the intersection is a point
// however it does not handle point segments as input!
bool seg_inter(L r, L *out) const {
   Tz = d \% r.d:
   if (abs(z) <= EPS) {
       if (abs(side(r.o)) > EPS) return false;
       if (r.d * d < 0) r = \{r.o + r.d, -r.d\};
       P s = o * d < r.o * d ? r.o : o;
       P = (o+d)*d < (r.o+r.d)*d ? o+d : r.o+r.d:
       if (s * d > e * d) return false;
       return *out = {s. e - s}, true:
   T s = inter(r), t = -r.inter(*this);
```

```
if (z < 0) s = -s, t = -t, z = -z:
       if (s>=-EPS \&\& s<=z+EPS \&\& t>=-EPS \&\& t<=z+EPS)
           return *out = \{0 + d * s / z, \{0, 0\}\}, true;
       return false:
    // check if the given point is on the segment
    bool point_on_seg(P r) const {
       if (abs(side(r)) > EPS) return false;
       if ((r - o) * d < -EPS) return false;</pre>
       if ((r - o - d) * d > EPS) return false:
       return true:
    }
    // point in this line that is closest to a given point
    P closest_to(P r) const {
       P dr = d.rot(): return r + dr*((o-r)*dr)/d.magsq():
   }
};
```

4.7 minkowski

```
void reorder_polygon(vector<P> &ps) {
   int pos = 0:
   repx(i, 1, (int)ps.size()) {
       if (ps[i].v < ps[pos].v || (ps[i].v == ps[pos].v &&</pre>
           ps[i].x < ps[pos].x))
           pos = i;
   rotate(ps.begin(), ps.begin() + pos, ps.end());
vector<P> minkowski(vector<P> ps, vector<P> qs) {
   // the first vertex must be the lowest
   reorder_polygon(ps); reorder_polygon(qs);
   ps.push back(ps[0]): ps.push back(ps[1]):
   qs.push_back(qs[0]); qs.push_back(qs[1]);
   vector<P> result; int i = 0, j = 0;
   while (i < ps.size() - 2 || j < qs.size() - 2) {</pre>
       result.push_back(ps[i] + qs[j]);
       auto z = (ps[i + 1] - ps[i]) % (qs[j + 1] - qs[j]);
       if (z >= 0 && i < ps.size() - 2) ++i;</pre>
       if (z \le 0 \&\& j \le gs.size() - 2) ++j;
   return result;
```

4.8 point

```
struct P {
   T x, y;
   P(T x, T y) : x(x), y(y) {}
   P() : P(0, 0) {}
   friend ostream &operator<<(ostream &s, const P &r) {</pre>
       return s << r.x << " " << r.v:
   friend istream & operator >> (istream &s, P &r) { return s
        >> r.x >> r.v: }
   P operator+(P r) const { return {x + r.x, y + r.y}; }
   P operator-(P r) const { return {x - r.x. y - r.y}: }
   P operator*(T r) const { return {x * r, y * r}; }
   P operator/(T r) const { return {x / r, y / r}; }
   P operator-() const { return {-x, -v}; }
   friend P operator*(T 1, P r) { return {1 * r.x, 1 * r.y};
   P rot() const { return {-v, x}: }
   T operator*(P r) const { return x * r.x + y * r.y; }
   T operator%(P r) const { return rot() * r; }
   T left(P a, P b) { return (b - a) % (*this - a); }
   T magsq() const { return x * x + y * y; }
   T mag() const { return sqrt(magsq()); }
   P unit() const { return *this / mag(); }
   bool half() const { return abs(y) <= EPS && x < -EPS || y
         < -EPS: }
   T angcmp(P r) const { // like strcmp(this, r)
       int h = (int)half() - r.half();
       return h ? h : r % *this;
   T angcmp_rel(P a, P b) { // like strcmp(a, b)
      Pz = *this:
       int h = z \% a \le 0 \&\& z * a \le 0 || z \% a \le 0;
      h = z \% b \le 0 \&\& z * b \le 0 || z \% b \le 0:
      return h ? h : b % a:
   bool operator==(P r) const { return abs(x - r.x) <= EPS</pre>
        && abs(y - r.y) <= EPS; }
   double angle() const { return atan2(y, x); }
   static P from_angle(double a) { return {cos(a), sin(a)};
       }
```

4.9 polygon

```
// get TWICE the area of a simple polygon in ccw order
T area2(const vector<P> &p) {
   int n = p.size(); T a = 0;
   rep(i, n) a += (p[i] - p[0]) % (p[(i + 1) % n] - p[i]);
   return a;
// checks whether a point is inside a ccw simple polygon
// returns 1 if inside. 0 if on border. -1 if outside
int in_poly(const vector<P> &p, P q) {
   int w = 0;
   rep(i, p.size()) {
      P = p[i], b = p[(i + 1) \% p.size()];
      T k = (b - a) \% (q - a):
      T u = a.y - q.y, v = b.y - q.y;
       if (k > 0 && u < 0 && v >= 0) w++;
       if (k < 0 && v < 0 && u >= 0) w--:
       if (k == 0 \&\& (q - a) * (q - b) <= 0) return 0;
   return w ? 1 : -1:
// check if point in ccw convex polygon. O(log n)
// + if inside, 0 if on border, - if outside
T in_convex(const vector<P> &p, P q) {
   int 1 = 1, h = p.size() - 2; assert(p.size() >= 3);
   while (1 != h) { // collinear points are unsupported!
       int m = (1 + h + 1) / 2:
       if (q.left(p[0], p[m]) >= 0) 1 = m;
       else h = m - 1:
   T in = min(q.left(p[0], p[1]), q.left(p.back(), p[0]));
   return min(in, q.left(p[1], p[1 + 1]));
int extremal(const vector<P> &p. P d) {
   int n = p.size(), 1 = 0, r = n - 1; assert(n);
   P = 0 = (p[n - 1] - p[0]).rot();
   while (1 < r)  { // polygon must be convex
       int m = (1 + r + 1) / 2:
      P = (p[(m + n - 1) \% n] - p[m]).rot();
      if (e0.angcmp_rel(d, e) < 0) r = m - 1;
       else 1 = m:
   }
   return 1;
// square dist of most distant points of a ccw convex
```

```
// polygon with NO COLLINEAR POINTS
T callipers(const vector<P> &p) {
   int n = p.size();
   T r = 0:
   for (int i = 0, j = n < 2 ? 0 : 1; <math>i < j; i++) {
       for (;; j = (j + 1) \% n) {
           r = max(r, (p[i] - p[j]).magsq());
           if ((p[(i + 1) % n] - p[i]) % (p[(j + 1) % n] - p
                [i]) <= EPS) break:</pre>
   return r:
}
P centroid(const vector<P> &p) { // (barycenter)
   P r(0, 0); T t = 0; int n = p.size();
   rep(i, n) {
       r += (p[i] + p[(i+1)\%n]) * (p[i] \% p[(i+1)\%n]);
       t += p[i] \% p[(i+1)\%n];
   return r / t / 3;
// classify collision of a ray inside a ccw polygon vertex.
// ray is (o, d), vertex is b, previous vertex is a, next is
pair<bool, bool> inner_collide(P o, P d, P a, P b, P c) {
   T p = (a - o) \% d: // side of previous
   T n = (c - o) \% d;
                        // side of next
   T v = (c - b) \% (b - a); // is vertex convex?
   return {v > 0 ? n < 0 || (n == 0 && p < 0) : p > 0 || n < //
           v > 0 ? p > 0 || (p == 0 && n > 0) : p > 0 || n < //
}
```

4.10 sweep

```
#include "point.cpp"

// iterate over all pairs of points

// 'op' is called with all ordered pairs of different
    indices '(i, j)'

// additionally, the 'ps' vector is kept sorted by signed
    distance

// to the line formed by 'i' and 'j'

// for example, if the vector from 'i' to 'j' is pointing
    right,
```

```
// the 'ps' vector is sorted from smallest 'y' to largest 'y
// note that, because the 'ps' vector is sorted by signed
    distance.
// 'j' is always equal to 'i + 1'
// this means that the amount of points to the left of the
    line is always 'N - i'
template <class OP>
void all_pair_points(vector<P> &ps, OP op) {
   int N = ps.size();
   sort(ps.begin(), ps.end(), [](P a, P b) {
      return make pair(a.v. a.x) < make pair(b.v. b.x):
   });
   vector<pair<int, int>> ss;
   rep(i, N) rep(j, N) if (i != j) ss.push_back({i, j});
   stable_sort(ss.begin(), ss.end(), [&](auto a, auto b) {
       return (ps[a.second] - ps[a.first]).angle_lt(ps[b.
           second] - ps[b.first]);
   vector<int> p(N); rep(i, N) p[i] = i;
   for (auto [i, j] : ss)
      { op(p[i], p[j]); swap(ps[p[i]], ps[p[j]]); swap(p[i
           ], p[i]); }
```

4.11 theorems

```
// Pick's theorem
// Simple polygon with integer vertices:
// A = I + B / 2 - 1
// A: Area of the polygon
// I: Integer points strictly inside the polygon
// B: Integer points on the boundary of the polygon
```

5 graph

5.1 artic-bridge-biconn

```
vector<int> g[MAXN]; int n;
struct edge {int u,v,comp; bool bridge;};
vector<edge> e;
void add_edge(int u, int v){
  g[u].pb(e.size()); g[v].pb(e.size());
        e.pb((edge){u,v,-1,false});
}
int D[MAXN],B[MAXN],T;
```

```
int nbc; // number of biconnected components
int art[MAXN]; // articulation point iff !=0
stack<int> st; // only for biconnected
void dfs(int u.int pe){
B[u]=D[u]=T++;
for(int ne:g[u])if(ne!=pe){
 int v=e[ne].u^e[ne].v^u;
 if(D[v]<0){
  st.push(ne);dfs(v,ne);
  if(B[v]>D[u])e[ne].bridge = true; // bridge
  if(B[v]>=D[u]){
   art[u]++; // articulation
   int last; // start biconnected
   do{last=st.top();st.pop();e[last].comp=nbc;}
              while(last!=ne);
   nbc++; // end biconnected
  B[u]=min(B[u],B[v]);
 else if(D[v]<D[u])st.push(ne),B[u]=min(B[u],D[v]);</pre>
void doit(){
memset(D,-1,sizeof(D));memset(art,0,sizeof(art));
nbc=T=0; fore(i,0,n)if(D[i]<0)dfs(i,-1),art[i]--;</pre>
```

5.2 bellman-ford

```
struct Edge { int u, v; ll w; };
// find distance from source node to all nodes.
// supports negative edge weights.
// returns true if a negative cycle is detected.
//
// time: O(V E)
bool bellman ford(int N. int s. vector<Edge> &E. vector<11>
    &D, vector<int> &P) {
   P.assign(N, -1), D.assign(N, INF), D[s] = 0;
   rep(i, N - 1) {
       bool f = true:
       rep(ei, E.size()) {
           auto &e = E[ei];
           ll n = D[e.u] + e.w;
           if (D[e.u] < INF && n < D[e.v])</pre>
              D[e.v] = n, P[e.v] = ei, f = false;
       if (f) return false:
```

```
return true;
}
```

5.3 blossom

```
vector<int> g[MAXN];int n,m,mt[MAXN],qh,qt,q[MAXN],ft[MAXN],
 bs[MAXN]:bool ing[MAXN].inb[MAXN].inp[MAXN]:int lca(int root
 ,int x,int y){memset(inp,0,sizeof(inp));while(1){inp[x=bs[x]
]=true;if(x==root)break;x=ft[mt[x]];}while(1){if(inp[y=bs[y]
])return y;else y=ft[mt[y]];}}void mark(int z,int x){while(
bs[x]!=z){int y=mt[x];inb[bs[x]]=inb[bs[y]]=true;x=ft[y];if(
bs[x]!=z)ft[x]=y;}}void contr(int s,int x,int y){int z=lca(s
 .x.v):memset(inb.0.sizeof(inb)):mark(z.x):mark(z.v):if(bs[x]
 !=z)ft[x]=y;if(bs[y]!=z)ft[y]=x;rep(x,n)if(inb[bs[x]]){bs[x]}
 =z;if(!inq[x])inq[q[++qt]=x]=true;}}int findp(int s){memset(
 inq,0,sizeof(inq));memset(ft,-1,sizeof(ft));rep(i,n)bs[i]=i;
 ing[g[qh=qt=0]=s]=true; while(qh<=qt){int x=q[qh++]; for(int y</pre>
 g[x] = g[x]  if g[x] = bs[y]  & g[x] = g[x]  if g[x] = g[x
]>=0)contr(s,x,y);else if(ft[y]<0){ft[y]=x;if(mt[y]<0)return}
y;else if(!inq[mt[y]])inq[q[++qt]=mt[y]]=true;}}}return -1;}
int aug(int s.int t){int x=t.v.z:while(x>=0){v=ft[x]:z=mt[v]}
 ;mt[v]=x;mt[x]=v;x=z;}return t>=0;}int edmonds(){int r=0;
memset(mt,-1,sizeof(mt));rep(x,n)if(mt[x]<0)r+=aug(x,findp(x</pre>
)):return r:}
```

5.4 chu-liu-minimum-spanning-arborescence

```
//O(n*m) minimum spanning tree in directed graph
//returns -1 if not possible
//included i-th edge if take[i]!=0
typedef int tw; tw INF=111<<30;</pre>
struct edge{int u.v.id:tw len:}:
struct ChuLiu{
int n; vector<edge> e;
vector<int> inc,dec,take,pre,num,id,vis;
vector<tw> inw;
void add_edge(int x, int y, tw w){
 inc.pb(0); dec.pb(0); take.pb(0);
 e.pb(\{x,y,SZ(e),w\});
ChuLiu(int n):n(n),pre(n),num(n),id(n),vis(n),inw(n){}
tw doit(int root){
 auto e2=e:
 tw ans=0; int eg=SZ(e)-1,pos=SZ(e)-1;
  fore(i,0,n) inw[i]=INF,id[i]=vis[i]=-1;
  for(auto ed:e2) if(ed.len<inw[ed.v]){</pre>
```

```
inw[ed.v]=ed.len: pre[ed.v]=ed.u:
  num[ed.v]=ed.id:
 inw[root]=0;
 fore(i,0,n) if(inw[i]==INF) return -1;
 int tot=-1:
 fore(i,0,n){
  ans+=inw[i]:
  if(i!=root)take[num[i]]++:
  while(vis[i]!=i&&i!=root&&id[i]<0)vis[i]=i.i=pre[i]:</pre>
  if(i!=root&&id[i]<0){</pre>
  id[j]=++tot:
  for(int k=pre[j];k!=j;k=pre[k]) id[k]=tot;
 if(tot<0)break:</pre>
 fore(i,0,n) if(id[i]<0)id[i]=++tot;</pre>
 n=tot+1; int j=0;
 fore(i,0,SZ(e2)){
  int v=e2[i].v;
  e2[j].v=id[e2[i].v];
  e2[j].u=id[e2[i].u];
  if(e2[j].v!=e2[j].u){
  e2[j].len=e2[i].len-inw[v];
  inc.pb(e2[i].id);
  dec.pb(num[v]);
   take.pb(0):
  e2[j++].id=++pos;
 e2.resize(j);
 root=id[root]:
while(pos>eg){
if(take[pos]>0) take[inc[pos]]++. take[dec[pos]]--:
return ans;
```

5.5 dinic

```
// time: O(E V^2)

// O(E V^(2/3)) / O(E sqrt(E)) unit capacities

// O(E sqrt(V)) (hopcroft-karp) unit networks

//unit network: c in {0,1} & forall v, indeg<=1 or outdeg<=1

//min-cut: nodes reachable from s in final residual graph
```

```
struct Dinic {
   struct Edge { int u, v; ll c, f = 0; };
   int N, s, t; vector<vector<int>> G;
   vector<Edge> E: vector<int> lvl. ptr:
   Dinic() {}
   Dinic(int N, int s, int t): N(N), s(s), t(t), G(N) {}
   void add_edge(int u, int v, 11 c) {
       G[u].push_back(E.size()); E.push_back({u, v, c});
       G[v].push_back(E.size()); E.push_back({v, u, 0});
   11 push(int u, ll p) {
       if (u == t || p <= 0) return p;</pre>
       while (ptr[u] < G[u].size()) {</pre>
           int ei = G[u][ptr[u]++];
           Edge &e = E[ei]:
           if (lvl[e.v] != lvl[u] + 1) continue;
           ll a = push(e.v. min(e.c - e.f. p)):
           if (a <= 0) continue:</pre>
           e.f += a, E[ei ^ 1].f -= a; return a;
       return 0;
   ll maxflow() {
      11 f = 0:
       while (true) {
           lvl.assign(N, -1); queue<int> q;
          lvl[s] = 0; q.push(s);
           while (!q.empty()) {
              int u = q.front(); q.pop();
              for (int ei : G[u]) {
                  Edge &e = E[ei];
                  if (e.c-e.f<=0||lvl[e.v]!=-1) continue;</pre>
                  lvl[e.v] = lvl[u] + 1; q.push(e.v);
           }
           if (lvl[t] == -1) break;
           ptr.assign(N,0); while(ll ff=push(s,INF))f += ff;
       }
       return f;
};
/* Flujo con demandas (no necesariamente el maximo)
Agregar s' y t' nuevos source and sink
c'(s', v) = sum(d(u, v) \text{ for } u \text{ in } V) \setminus sum(sin v)
c'(v, t') = sum(d(v, w) \text{ for w in V}) \setminus forall arista (v, t')
```

```
c'(t, s) = INF (el flujo por esta arista es el flujo real)*/
```

5.6 dominator-tree

```
//idom[i]=parent of i in dominator tree with root=rt, or -1
int n,rnk[MAXN],pre[MAXN],anc[MAXN],idom[MAXN],semi[MAXN],
vector<int> g[MAXN],rev[MAXN],dom[MAXN],ord;
void dfspre(int pos){
rnk[pos]=SZ(ord); ord.pb(pos);
for(auto x:g[pos]){
 rev[x].pb(pos);
 if(rnk[x]==n) pre[x]=pos.dfspre(x):
}
int eval(int v){
if(anc[v]<n&&anc[anc[v]]<n){</pre>
 int x=eval(anc[v]):
 if(rnk[semi[low[v]]]>rnk[semi[x]]) low[v]=x:
 anc[v]=anc[anc[v]]:
return low[v]:
void dominators(int rt){
fore(i,0,n){
 dom[i].clear(): rev[i].clear():
 rnk[i]=pre[i]=anc[i]=idom[i]=n:
 semi[i]=low[i]=i;
}
ord.clear(); dfspre(rt);
for(int i=SZ(ord)-1;i;i--){
 int w=ord[i]:
 for(int v:rev[w]){
  int u=eval(v):
  if(rnk[semi[w]]>rnk[semi[u]])semi[w]=semi[u]:
 dom[semi[w]].pb(w): anc[w]=pre[w]:
 for(int v:dom[pre[w]]){
  int u=eval(v);
  idom[v]=(rnk[pre[w]]>rnk[semi[u]]?u:pre[w]);
 dom[pre[w]].clear():
for(int w:ord) if(w!=rt&&idom[w]!=semi[w]) idom[w]=idom[
     idom[w]]:
fore(i,0,n) if(idom[i]==n)idom[i]=-1;
```

5.7 eulerian-path

```
// Directed version(uncomment commented code for undirected)
struct edge {
int v; // list<edge>::iterator rev;
edge(int v):v(v){}
list<edge> g[MAXN];
void add edge(int a, int b){
g[a].push_front(edge(b));//auto ia=g[a].begin();
// g[b].push_front(edge(a));auto ib=g[b].begin();
// ia->rev=ib:ib->rev=ia:
vector<int> p;
void go(int x){
while(g[x].size()){
 int y=g[x].front().y;//g[y].erase(g[x].front().rev);
 g[x].pop_front(); go(y);
p.push_back(x);
vector<int> get_path(int x){ // get a path that begins in x
// check that a path exists from x before calling get_path!
p.clear();go(x);reverse(p.begin(),p.end());
return p;
```

5.8 floyd-warshall

5.9 heavy-light

```
struct Hld {
   vector<int> P, H, D, pos, top;
```

```
Hld() {}
void init(vector<vector<int>> &G) {
    int N = G.size():
   P.resize(N), H.resize(N), D.resize(N), pos.resize(N),
       top.resize(N);
   D[0] = -1, dfs(G, 0); int t = 0:
   rep(i, N) if (H[P[i]] != i) {
       int j = i;
       while (j != -1)
           \{ top[i] = i, pos[i] = t++; j = H[i]; \}
}
int dfs(vector<vector<int>> &G. int i) {
   int w = 1. mw = 0:
   D[i] = D[P[i]] + 1, H[i] = -1;
   for (int c : G[i]) {
       if (c == P[i]) continue;
       P[c] = i: int sw = dfs(G, c): w += sw:
       if (sw > mw) H[i] = c. mw = sw:
   }
   return w:
// visit the log N segments in the path from u to v
template <class OP>
void path(int u, int v, OP op) {
   while (top[u] != top[v]) {
       if (D[top[u]] > D[top[v]]) swap(u, v);
       op(pos[top[v]], pos[v] + 1); v = P[top[v]];
    if (D[u] > D[v]) swap(u, v);
    op(pos[u], pos[v] + 1); // value on node
    // op(pos[u]+1, pos[v] + 1); // value on edge
// an alternative to 'path' that considers order.
// calls 'op' with an 'l <= r' inclusive-exclusive range,</pre>
// boolean indicating if the query is forwards or
     backwards.
template <class OP>
void path(int u, int v, OP op) {
   int lu = u. lv = v:
   while (top[lu] != top[lv])
       if (D[top[lu]] > D[top[lv]]) lu = P[top[lu]];
       else lv = P[top[lv]];
   int lca = D[lu] > D[lv] ? lv : lu:
    while (top[u] != top[lca])
```

```
op(pos[top[u]], pos[u] + 1, false), u = P[top[u]
       if (u != lca) op(pos[lca] + 1, pos[u] + 1, false);
       vector<int> stk;
       while (top[v] != top[lca])
          stk.push_back(v), v = P[top[v]];
       // op(pos[lca], pos[v] + 1, true); // value on node
       op(pos[lca] + 1, pos[v] + 1, true); // value on edge
       reverse(stk.begin(), stk.end());
       for (int w : stk) op(pos[top[w]], pos[w] + 1, true):
   // commutative segment tree
   template <class T, class S>
   void update(S &seg, int i, T val) { seg.update(pos[i],
        val); }
   // commutative segment tree lazy
   template <class T, class S>
   void update(S &seg, int u, int v, T val) {
       path(u, v, [&](int 1, int r) { seg.update(1, r, val);
            }):
   // commutative (lazy) segment tree
   template <class T, class S>
   T query(S &seg, int u, int v) {
      T ans = 0:
           // neutral element
       path(u, v, [&](int 1, int r) { ans += seg.query(1, r)
           ; }); // query op
       return ans;
};
```

5.10 hungarian

```
// find a maximum gain perfect matching in the given
    bipartite complete graph.
// input: gain matrix (G_{xy} = benefit of joining vertex x
    in set X with vertex
// y in set Y).
// output: maximum gain matching in members 'xy[x]' and 'yx[
    y]'.
// runtime: O(N^3)
struct Hungarian {
    int N, qi, root;
```

```
vector<vector<ll>>> gain:
vector<int> xy, yx, p, q, slackx;
vector<ll> lx, ly, slack;
vector<bool> S. T:
void add(int x, int px) {
   S[x] = true, p[x] = px;
   rep(y, N) if (lx[x] + ly[y] - gain[x][y] < slack[y])
       slack[y] = lx[x] + ly[y] - gain[x][y], slackx[y]
   }
}
void augment(int x, int y) {
   while (x != -2) {
       yx[y] = x; swap(xy[x], y); x = p[x];
}
void improve() {
   S.assign(N, false), T.assign(N, false), p.assign(N,
        -1):
   qi = 0, q.clear();
   rep(x, N) if (xy[x] == -1) {
       q.push_back(root = x), p[x] = -2, S[x] = true;
   rep(v, N) slack[v] = lx[root] + lv[v] - gain[root][v
        1. slackx[v] = root:
   while (true) {
       while (qi < q.size()) {</pre>
          int x = q[qi++];
          rep(y, N) if (lx[x] + ly[y] == gain[x][y] &&!
               T[v]) {
              if (vx[v] == -1) return augment(x, y);
              T[y] = true, q.push_back(yx[y]), add(yx[y
                   ], x);
          }
       }
       11 d = INF;
       rep(y, N) if (!T[y]) d = min(d, slack[y]);
       rep(x, N) if (S[x]) lx[x] -= d;
       rep(y, N) if (T[y]) ly[y] += d;
       rep(y, N) if (!T[y]) slack[y] -= d;
       rep(y, N) if (!T[y] && slack[y] == 0) {
          if (yx[y] == -1) return augment(slackx[y], y); };
```

5.11 kuhn

```
// get a maximum cardinality matching in a bipartite graph.
// input: adjacency lists.
// output: matching (in 'mt' member).
// runtime: O(V E)
struct Kuhn {
   vector<vector<int>> G:
   int N, size;
   vector<bool> seen:
   vector<int> mt:
   bool visit(int i) {
       if (seen[i]) return false;
       seen[i] = true:
       for (int to : G[i])
          if (mt[to] == -1 || visit(mt[to])) {
              mt[to] = i:
              return true:
       return false:
   Kuhn(vector<vector<int>> adj) : G(adj), N(G.size()), mt(N
        , -1) {
      rep(i, N) {
          seen.assign(N, false);
           size += visit(i):
      }
```

5.12 lca

```
// calculates the lowest common ancestor for any two nodes
     in O(log N) time,
// with O(N log N) preprocessing
struct Lca {
   int N, K, t = 0;
   vector<vector<int>> U:
   vector<int> L, R;
   Lca() {}
   Lca(vector<vector<int>> &G) : N(G.size()), L(N), R(N) {
       K = N \le 1 ? 0 : 32 - \_builtin_clz(N - 1);
       U.resize(K + 1, vector<int>(N));
       visit(G, 0, 0);
       rep(k, K) rep(u, N) U[k + 1][u] = U[k][U[k][u]];
   void visit(vector<vector<int>> &G. int u. int p) {
       L[u] = t++, U[0][u] = p;
       for (int v : G[u]) if (v != p) visit(G, v, u);
       R[u] = t++:
   }
   bool is_anc(int up, int dn) {
       return L[up] <= L[dn] && R[dn] <= R[up];</pre>
   int find(int u, int v) {
       if (is anc(u, v)) return u:
       if (is_anc(v, u)) return v;
       for (int k = K; k \ge 0;)
           if (is_anc(U[k][u], v)) k--;
           else u = U[k][u];
       return U[0][u]:
};
```

5.13 maxflow-mincost

```
// time: O(F V E) F is the maximum flow
// O(V E + F E log V) if bellman-ford is replaced by
    johnson
struct Flow {
    struct Edge {
        int u, v;
        ll c, w, f = 0;
    };
```

```
int N. s. t:
vector<vector<int>> G:
vector<Edge> E;
vector<ll> d. b:
vector<int> p;
Flow() {}
Flow(int N, int s, int t) : N(N), s(s), t(t), G(N) {}
void add_edge(int u, int v, ll c, ll w) {
   G[u].push back(E.size()):
   E.push back({u, v, c, w}):
   G[v].push_back(E.size());
   E.push_back({v, u, 0, -w});
}
// naive distances with bellman-ford: O(V E)
void calcdists() {
   p.assign(N, -1), d.assign(N, INF), d[s] = 0;
   rep(i, N - 1) rep(ei, E.size()) {
       Edge &e = E[ei];
       ll n = d[e.u] + e.w:
       if (d[e.u] < INF && e.c - e.f > 0 && n < d[e.v])
           d[e.v] = n, p[e.v] = ei;
   }
}
// johnsons potentials: O(E log V)
void calcdists() {
   if (b.emptv()) {
       b.assign(N, 0);
       // code below only necessary if there are
           negative costs
       rep(i, N - 1) rep(ei, E.size()) {
          Edge &e = E[ei]:
          if (e.f < e.c) b[e.v] = min(b[e.v], b[e.u] + e
   p.assign(N, -1), d.assign(N, INF), d[s] = 0;
   priority queue<pair<11. int>> q:
   q.push({0, s});
   while (!a.emptv()) {
       auto [w, u] = q.top();
       q.pop();
       if (d[u] < -w + b[u]) continue:
       for (int ei : G[u]) {
          auto e = E[ei]:
          11 n = d[u] + e.w:
          if (e.f < e.c && n < d[e.v]) {</pre>
```

```
d[e.v] = n, p[e.v] = ei:
                  q.push({b[e.v] - n, e.v});
          }
      }
       b = d:
   }
   ll solve() {
       b.clear();
       11 \text{ ff} = 0:
       while (true) {
          calcdists();
          if (p[t] == -1) break;
           for (int cur = t; p[cur] != -1; cur = E[p[cur]].u
              f = min(f, E[p[cur]].c - E[p[cur]].f);
           for (int cur = t: p[cur] != -1: cur = E[p[cur]].u
              E[p[cur]].f += f, E[p[cur] ^ 1].f -= f;
          ff += f;
       }
       return ff;
   }
};
```

5.14 parallel-dfs

```
struct Tree {
int n.z[2]:
vector<vector<int>> g;
vector<int> ex,ey,p,w,f,v[2];
Tree(int n):g(n), w(n), f(n){}
void add_edge(int x, int y){
 p.pb(g[x].size());g[x].pb(ex.size());
       ex.pb(x);ey.pb(y);
 p.pb(g[y].size());g[y].pb(ex.size());
       ex.pb(y);ey.pb(x);
bool go(int k){//returns 1 if it finds new node
 int & x=z[k];
 while(x \ge 0 \& \&
  (w[x]==g[x].size()||w[x]==g[x].size()-1
          &&(g[x].back()^1)==f[x]))
  x=f[x]>=0?ex[f[x]]:-1;
 if(x<0)return false:</pre>
 if((g[x][w[x]]^1)==f[x])w[x]++;
```

```
int e=g[x][w[x]],y=ey[e]; f[y]=e;
      w[x]++; w[y]=0; x=y; v[k].pb(x);
 return true;
 vector<int> erase_edge(int e){
 e*=2;//erases eth edge, returns smaller comp
 int x=ex[e],y=ey[e]; p[g[x].back()]=p[e];
 g[x][p[e]]=g[x].back(); g[x].pop_back();
 p[g[y].back()]=p[e^1]; g[y][p[e^1]]=g[y].back();
       g[y].pop_back();
 f[x]=f[y]=-1; w[x]=w[y]=0; z[0]=x;z[1]=y;
      v[0]=\{x\}:v[1]=\{v\}:
 bool d0=true,d1=true;while(d0&&d1)d0=go(0),d1=go(1);
 return v[1-d1]:
}
};
```

5.15 push-relabel

```
#include "../common.h"
const 11 INF = 1e18:
// maximum flow algorithm.
// to run, use 'maxflow()'.
// time: O(V^2 \operatorname{sqrt}(E)) \leq O(V^3)
// memory: □(V^2)
struct PushRelabel {
   vector<vector<ll>> cap, flow;
   vector<ll> excess:
   vector<int> height:
   PushRelabel() {}
   void resize(int N) { cap.assign(N, vector<11>(N)); }
   // push as much excess flow as possible from u to v.
   void push(int u, int v) {
       11 f = min(excess[u], cap[u][v] - flow[u][v]);
       flow[u][v] += f:
       flow[v][u] -= f;
       excess[v] += f;
       excess[u] -= f;
   // relabel the height of a vertex so that excess flow may
         be pushed.
   void relabel(int u) {
       int d = INT32_MAX;
```

```
rep(v, cap.size()) if (cap[u][v] - flow[u][v] > 0) d | 5.16 strongly-connected-components
       min(d, height[v]);
   if (d < INF) height[u] = d + 1;</pre>
}
// get the maximum flow on the network specified by 'cap'
     with source 's'
// and sink 't.'.
// node-to-node flows are output to the 'flow' member.
11 maxflow(int s. int t) {
   int N = cap.size(). M:
   flow.assign(N, vector<11>(N));
   height.assign(N, 0), height[s] = N;
   excess.assign(N, 0), excess[s] = INF;
   rep(i, N) if (i != s) push(s, i);
   vector<int> q;
   while (true) {
       // find the highest vertices with excess
       q.clear(), M = 0;
       rep(i, N) {
          if (excess[i] <= 0 || i == s || i == t)</pre>
          if (height[i] > M) q.clear(), M = height[i];
          if (height[i] >= M) q.push_back(i);
       if (q.empty()) break;
       // process vertices
       for (int u : a) {
          bool relab = true;
          rep(v, N) {
              if (excess[u] <= 0) break;</pre>
              if (cap[u][v] - flow[u][v] > 0 && height[u]
                  ] > height[v])
                  push(u, v), relab = false:
          if (relab) {
              relabel(u);
              break:
   11 f = 0; rep(i, N) f += flow[i][t]; return f;
```

```
/* time: O(V + E), memory: O(V)
after building:
   comp = map from vertex to component
        (components are toposorted, root first, leaf last)
   N = number of components
   G = condensation graph (component DAG)
byproducts:
   vgi = transposed graph
   order = reverse topological sort (leaf first, root last)
   vn = number of vertices
   vg = original vertex graph
struct Scc {
   int vn. N:
   vector<int> order, comp;
   vector<vector<int>> vg, vgi, G;
   void toposort(int u) {
       if (comp[u]) return;
       comp[u] = -1;
      for (int v : vg[u]) toposort(v);
       order.push_back(u);
   bool carve(int u) {
       if (comp[u] != -1) return false;
       comp[u] = N:
      for (int v : vgi[u]) {
          carve(v):
          if (comp[v] != N) G[comp[v]].push_back(N);
      }
       return true;
   }
   Scc() {}
   Scc(vector<vector<int>> &g)
    : vn(g.size()), vg(g), comp(vn), vgi(vn), G(vn), N(0) {
      rep(u, vn) toposort(u);
       rep(u, vn) for (int v : vg[u]) vgi[v].push_back(u);
       invrep(i, vn) N += carve(order[i]);
};
```

5.17 two-sat

```
// calculate the solvability of a system of logical
    equations, where every equation is of the form 'a or b
// 'neg': get negation of 'u'
// 'then': 'u' implies 'v'
```

```
// 'anv': 'u' or 'v'
// 'set': 'u' is true
11
// after 'solve' (O(V+E)) returns true, 'sol' contains one
    possible solution.
// determining all solutions is O(V*E) hard (requires
    computing reachability in a DAG).
struct TwoSat {
   int N; vector<vector<int>> G;
   Scc scc; vector<bool> sol;
   TwoSat(int n): N(n), G(2 * n), sol(n) {}
   TwoSat() {}
   int neg(int u) { return (u + N) % (2 * N); }
   void then(int u, int v) { G[u].push_back(v), G[neg(v)].
        push_back(neg(u)); }
   void any(int u, int v) { then(neg(u), v); }
   void set(int u) { G[neg(u)].push_back(u); }
   bool solve() {
       scc = Scc(G);
       rep(u, N) if (scc.comp[u] == scc.comp[neg(u)]) return
       rep(u, N) sol[u] = (scc.comp[u] > scc.comp[neg(u)]);
       return true:
};
```

6 implementation

6.1 bit-tricks

```
v = x & (x-1) // Turn off rightmost 1bit
y = x & (-x) // Isolate rightmost 1bit
y = x \mid (x-1) // Right propagate rightmost 1bit(fill in 1s)
y = x \mid (x+1) // Turn on rightmost Obit
y = x & (x+1) // Isolate rightmost Obit
// If x is of long type, use __builtin_popcountl(x)
// If x is of long long type, use __builtin_popcountll(x)
// 1. Counts the number of ones(set bits) in an integer.
__builtin_popcount(x)
// 2. Checks the Parity of a number. Returns true(1) if the
// number has odd number of set bits, else it returns
// false(0) for even number of set bits.
__builtin_parity(x)
// 3. Counts the leading number of zeros of the integer.
__builtin_clz(x)
// 4. Counts the trailing number of zeros of the integer.
```

```
__builtin_ctz(x)
// 5. Returns 1 + the index of the least significant 1-bit.
__builtin_ffs(x) // If x == 0, returns 0.
// Iterate over non empty subsets of bitmask
for(int s=m;s;s=(s-1)&m) // Decreasing order
for (int s=0;s=s-m&m;) // Increasing order
```

6.2 dynamic-connectivity

```
struct DC {
   int n; Dsu D;
   vector<vector<pair<int, int>>> t;
   DC(int N) : n(N), D(N), t(2 * N) {}
   // add edge p to all times in interval [1, r]
   void upd(int 1, int r, pair<int, int> p) {
      for (1 += n, r += n; 1 < r; 1 >>= 1, r >>= 1) {
          if (1 & 1) t[1++].push back(p):
          if (r & 1) t[--r].push_back(p);
      }
   void process(int u = 1) { // process all queries
       for (auto &e : t[u]) D.unite(e.first, e.second):
      if (u \ge n) {
          // do stuff with D at time u - n
      } else process(2 * u), process(2 * u + 1);
       for (auto &e : t[u]) D.rollback();
   }
};
```

6.3 hash-container

```
namespace{//add (#define tmpl template)(#define ty typename)
  tmpl<ty T> size_t mk_h(const T& v){return hash<T>()(v);}
  void h_cmb(size_t& h, const size_t& v)
  { h ^= v + 0x9e3779b9 + (h << 6) + (h >> 2); }
  tmpl<ty T> struct h_ct{size_t operator()(const T& v)const{
  size_t h=0;for(const auto& e:v){h_cmb(h,mk_h(e));}return h;
  }};
}namespace std{//support for pair<T,U>, vector<T> & map<T,U>
  tmpl<ty T, ty U> struct hash<pair<T, U>>{
    size_t operator()(const pair<T,U>& v) const
  {size_t operator()(const pair<T,U>& v) const
  {size_t h=mk_h(v.first);h_cmb(h, mk_h(v.second));return h;}
  };

tmpl<ty... T>struct hash<vector<T...>>:h_ct<vector<T...>>{};
  tmpl<ty... T>struct hash<map<T...>>:h_ct<map<T...>>{};
```

6.4 mo

```
struct Querv { int 1. r. idx: }:
// answer segment queries using only 'add(i)', 'remove(i)'
    and 'get()'
// functions.
//
// complexity: O((N + Q) * sqrt(N) * F)
// N = length of the full segment
// Q = amount of queries
// F = complexity of the 'add', 'remove' functions
template <class A, class R, class G, class T>
void mo(vector<Query> &queries, vector<T> &ans, A add, R
    remove, G get) {
   int Q = queries.size(), B = (int)sqrt(Q):
   sort(queries.begin(), queries.end(), [&](Query &a, Query
       return make_pair(a.1 / B, a.r) < make_pair(b.1 / B, b</pre>
   }):
   ans.resize(Q);
   int 1 = 0, r = 0:
   for (auto &g : gueries) {
       while (r < q.r) add(r), r++;
       while (1 > q.1) 1--, add(1):
       while (r > q.r) r--, remove(r);
       while (1 < q.1) remove(1), 1++:
       ans[q.idx] = get();
```

6.5 ordered-set

6.6 unordered-map

7 imprimible

8 math

8.1 Linear Diophantine

```
ii extendedEuclid(ll a, ll b){
11 x, y; //a*x + b*y = gcd(a,b)
if (b == 0) return {1, 0}:
auto p = extendedEuclid(b, a%b);
x = p.second;
y = p.first - (a/b)*x;
if(a*x + b*y == -_gcd(a,b)) x=-x, y=-y;
return {x, y};
pair<ii, ii> diophantine(ll a, ll b, ll r){
//a*x+b*y=r where r is multiple of gcd(a,b);
11 d = \_gcd(a, b);
a/=d: b/=d: r/=d:
auto p = extendedEuclid(a, b);
p.first*=r; p.second*=r;
assert(a*p.first + b*p.second == r);
return {p, {-b, a}}; //solutions: p+t*ans.second
```

8.2 arithmetic

```
inline int floor_log2(int n)
```

```
{ return n <= 1 ? 0 : 31 - __builtin_clz(n); }
inline int ceil_log2(int n)
{ return n <= 1 ? 0 : 32 - __builtin_clz(n - 1); }
inline ll floordiv(ll a, ll b) {return a/b-((a^b)<0&&a%b);}
inline ll ceildiv(ll a, ll b) {return a/b+((a^b)>=0&&a%b);}
```

8.3 berlekamp-massey-linear-recurrence

```
vector<int> BM(vector<int> x) {
   vector<int> ls. cur:
   int lf, ld;
   rep(i, x.size()) {
       11 t = 0:
       rep(j, cur.size()) t = (t+x[i-j-1]*(ll)cur[j])%MOD;
       if ((t - x[i]) \% MOD == 0) continue:
       if (!cur.size()) {
           cur.resize(i + 1); lf = i; ld = (t-x[i]) % MOD;
           continue:
       11 k = -(x[i] - t) * bin_exp(1d, MOD - 2) % MOD;
       vector<int> c(i - lf - 1); c.push_back(k);
       rep(j, ls.size()) c.push_back(-ls[j] * k % MOD);
       if (c.size() < cur.size()) c.resize(cur.size());</pre>
       rep(j, cur.size()) c[j] = (c[j] + cur[j]) % MOD;
       if (i - lf + ls.size() >= cur.size())
           ls = cur, lf = i, ld = (t - x[i]) % MOD:
   rep(i, cur.size()) cur[i] = (cur[i] % MOD + MOD) % MOD;
   return cur;
// Linear Recurrence
11 \text{ MOD} = 998244353:
11 LOG = 60:
struct LinearRec{
 typedef vector<int> vi;
 int n: vi terms, trans: vector<vi> bin:
 vi add(vi &a, vi &b){
   vi res(n*2+1):
   rep(i,n+1) rep(j,n+1)
       res[i+j]=(res[i+j]*1LL+(ll)a[i]*b[j])%MOD;
   for(int i=2*n; i>n; --i){
     rep(j,n)
       res[i-1-j]=(res[i-1-j]*1LL+(ll)res[i]*trans[j])%MOD;
   res.erase(res.begin()+n+1,res.end());
   return res:
```

```
LinearRec(vi &terms, vi &trans):terms(terms),trans(trans){
    n=trans.size();vi a(n+1);a[1]=1;
    bin.push_back(a);
    repx(i,1,LOG)bin.push_back(add(bin[i-1],bin[i-1]));
}
int calc(ll k){
    vi a(n+1);a[0]=1;
    rep(i,LOG)if((k>>i)&1)a=add(a,bin[i]);
    int ret=0;
    rep(i,n)ret=((l1)ret+(l1)a[i+1]*terms[i])%MOD;
    ret = ret%MOD + MOD;
    return ret%MOD;
}
```

8.4 crt

```
pair<11, 11> solve_crt(const vector<pair<11, 11>> &eqs) {
    11 a0 = eqs[0].first, p0 = eqs[0].second;
    repx(i, 1, eqs.size()) {
        11 a1 = eqs[i].first, p1 = eqs[i].second;
        11 k1, k0;
        11 d = ext_gcd(p1, p0, k1, k0);
        a0 -= a1;
        if (a0 % d != 0) return {-1, -1};
        p0 = p0 / d * p1;
        a0 = a0 / d * k1 % p0 * p1 % p0 + a1;
        a0 = (a0 % p0 + p0) % p0;
    }
    return {a0, p0};
}
```

8.5 discrete-log

```
// discrete logarithm log_a(b).
// solve b ^ x = a (mod M) for the smallest x.
// returns -1 if no solution is found.
//
// time: 0(sqrt(M))
ll dlog(ll a, ll b, ll M) {
    ll k = 1, s = 0;
    while (true) {
        ll g = __gcd(b, M);
        if (g <= 1) break;
        if (a == k) return s;
        if (a % g != 0) return -1;
        a /= g, M /= g, s += 1, k = b / g * k % M;</pre>
```

```
}
11 N = sqrt(M) + 1;

umap<ll, ll> r;
rep(q, N + 1) {
    r[a] = q;
    a = a * b % M;
}

ll bN = binexp(b, N, M), bNp = k;
repx(p, 1, N + 1) {
    bNp = bNp * bN % M;
    if (r.count(bNp)) return N * p - r[bNp] + s;
}
return -1;
```

8.6 fast-hadamard-transform

```
ll c1[MAXN+9].c2[MAXN+9]://MAXN must be power of 2!
void fht(ll* p, int n, bool inv){
for(int l=1;2*1<=n;1*=2)for(int i=0;i<n;i+=2*1)fore(j,0,1){</pre>
 11 u=p[i+j],v=p[i+l+j];
 if(!inv)p[i+j]=u+v,p[i+l+j]=u-v; // XOR
 else p[i+j]=(u+v)/2, p[i+l+j]=(u-v)/2;
 //if(!inv)p[i+j]=v,p[i+l+j]=u+v; // AND
 //else p[i+j]=-u+v,p[i+l+j]=u;
 //if(!inv)p[i+j]=u+v,p[i+l+j]=u; // OR
 //else p[i+i]=v.p[i+l+i]=u-v:
// like polynomial multiplication, but XORing exponents
// instead of adding them (also ANDing, ORing)
vector<ll> multiply(vector<ll>& p1, vector<ll>& p2){
int n=1 << (32-\_builtin\_clz(max(SZ(p1),SZ(p2))-1));
fore(i,0,n)c1[i]=0,c2[i]=0:
fore(i,0,SZ(p1))c1[i]=p1[i];
fore(i,0,SZ(p2))c2[i]=p2[i];
fht(c1,n,false);fht(c2,n,false);
fore(i,0,n)c1[i]*=c2[i];
fht(c1,n,true);
return vector<ll>(c1.c1+n):
```

8.7 ff

```
using cd = complex<double>;
```

```
const double PI = acos(-1):
// compute the DFT of a power-of-two-length sequence.
// if 'inv' is true, computes the inverse DFT.
void fft(vector<cd> &a, bool inv) {
   int N = a.size(), k = 0, b;
   assert(N == 1 << __builtin_ctz(N));</pre>
   repx(i, 1, N) {
      for (b = N >> 1; k & b;) k ^= b, b >>= 1;
       if (i < (k ^= b)) swap(a[i], a[k]);</pre>
   for (int 1 = 2; 1 <= N; 1 <<= 1) {</pre>
       double ang = 2 * PI / 1 * (inv ? -1 : 1);
       cd wl(cos(ang), sin(ang));
       for (int i = 0: i < N: i += 1) {
           cd w = 1;
          rep(j, 1 / 2) {
              cd u = a[i + j], v = a[i + j + 1 / 2] * w;
              a[i + j] = u + v;
              a[i + j + 1 / 2] = u - v;
              w *= w1;
   }
   if (inv) rep(i, N) a[i] /= N:
const 11 MOD = 998244353, ROOT = 15311432;
// const 11 MOD = 2130706433, ROOT = 1791270792;
// const 11 MOD = 922337203673733529711. ROOT =
    532077456549635698311;
void find root of unitv(ll M) {
   11 c = M - 1. k = 0:
   while (c \% 2 == 0) c /= 2, k += 1;
   // find proper divisors of M - 1
   vector<ll> divs:
   for (11 d = 1; d < c; d++) {
      if (d * d > c) break:
       if (c % d == 0) rep(i, k + 1) divs.push_back(d << i);</pre>
   rep(i, k) divs.push_back(c << i);</pre>
   // find any primitive root of M
   11 G = -1:
   repx(g, 2, M) {
```

```
bool ok = true:
      for (ll d : divs) ok &= (binexp(g, d, M) != 1);
      if (ok) {
          G = g:
          break;
      }
   }
   assert(G != -1);
   ll w = binexp(G, c, M);
   cerr << "M = c * 2^k + 1" << endl:
   cerr << " M = " << M << endl;
   cerr << " c = " << c << endl;
   cerr << " k = " << k << endl;
   cerr << " w^(2^k) == 1" << endl;
   cerr << " w = g^{(M-1)/2k} = g^c << endl;
   cerr << " g = " << G << endl;
   cerr << " w = " << w << endl:
// compute the DFT of a power-of-two-length sequence, modulo
     a special prime
// number with an Nth root of unity, where N is the length
    of the sequence.
void ntt(vector<ll>> &a, bool inv) {
   vector<ll> wn;
   for (11 p = ROOT; p != 1; p = p * p % MOD) wn.push_back(p
        ):
   int N = a.size(), k = 0, b:
   assert(N == 1 << \_builtin_ctz(N) && N <= 1 << wn.size())
   rep(i, N) a[i] = (a[i] % MOD + MOD) % MOD;
   repx(i, 1, N) {
      for (b = N >> 1; k & b;) k ^= b, b >>= 1;
       if (i < (k ^= b)) swap(a[i], a[k]);</pre>
   for (int 1 = 2: 1 <= N: 1 <<= 1) {
      11 wl = wn[wn.size() - __builtin_ctz(1)];
       if (inv) wl = multinv(wl, MOD);
      for (int i = 0; i < N; i += 1) {</pre>
          11 w = 1:
          repx(j, 0, 1 / 2) {
              11 u = a[i + j], v = a[i + j + 1 / 2] * w %
                   MOD:
              a[i + j] = (u + v) \% MOD;
              a[i + j + 1 / 2] = (u - v + MOD) \% MOD;
```

```
w = w * wl % MOD;
}

}

ll q = multinv(N, MOD);
if (inv) rep(i, N) a[i] = a[i] * q % MOD;
}

void convolve(vector<cd> &a, vector<cd> b, int n) {
    n = 1 << (32 - __builtin_clz(2 * n - 1));
    a.resize(n), b.resize(n);
    fft(a, false), fft(b, false);
    rep(i, n) a[i] *= b[i];
    fft(a, true);
}</pre>
```

8.8 gauss

```
const double EPS = 1e-9;
// solve a system of equations.
// complexity: O(\min(N, M) * N * M)
// 'a' is a list of rows
// the last value in each row is the result of the equation
// return values:
// 0 -> no solutions
// 1 -> unique solution, stored in 'ans'
// -1 -> infinitely many solutions, one of which is stored
    in 'ans'
// UNTESTED
int gauss(vector<vector<double>> a, vector<double> &ans) {
   int N = a.size(), M = a[0].size() - 1:
   vector<int> where(M, -1);
   for (int j = 0, i = 0; j < M && i < N; j++) {
       int sel = i;
       repx(k, i, N) if (abs(a[k][j]) > abs(a[sel][j])) sel
            = k:
       if (abs(a[sel][j]) < EPS) continue;</pre>
       repx(k, j, M + 1) swap(a[sel][k], a[i][k]);
       where[j] = i;
       rep(k, N) if (k != i) {
          double c = a[k][j] / a[i][j];
          repx(1, j, M + 1) a[k][1] -= a[i][1] * c;
      }
       i++;
```

8.9 matrix

```
typedef vector<vector<double>> Mat;
Mat matmul(Mat 1, Mat r) {
   int n = 1.N, m = r.M, p = 1.M; assert(1.M == r.N);
   Mat a(n, vector<double>(m)); // neutral
   rep(i, n) rep(j, m)
      rep(k, p) a[i][j] = a[i][j] + l[i][k] * r[k][j];
   return a;
double reduce(vector<vector<double>> &A) {
   int n = A.size(), m = A[0].size():
   int i = 0, j = 0; double r = 1.;
   while (i < n && j < m) {</pre>
      int 1 = i:
       repx(k, i+1, n) if(abs(A[k][j]) > abs(A[l][j])) l=k;
       if (abs(A[1][j]) < EPS) { j++; r = 0.; continue; }</pre>
      if (1 != i) \{ r = -r : swap(A[i], A[1]) : \}
      r *= A[i][i];
      for (int k = m - 1; k >= j; k--) A[i][k] /= A[i][j];
       repx(k, 0, n) {
          if (k == i) continue;
          for(int l=m-1;l>=j;l--)A[k][l]-=A[k][j]*A[i][l];
       i++, j++;
   return r; // returns determinant
```

8.10 mobius

8.11 multiny

```
// a * x + b * y == gcd(a, b)
11 ext_gcd(11 a, 11 b, 11 &x, 11 &y) {
    if (b == 0) { x = 1, y = 0; return a; }
    11 d = ext_gcd(b, a % b, y, x); y -= a / b * x; return d;
}

// inverse exists if and only if a and M are coprime
// if M is prime: multinv(a, M) = (a**(M-2)) % M
11 multinv(11 a, 11 M)
{ 11 x, y; ext_gcd(a, M, x, y); return x; }

// all modular inverses from 1 to inv.size()-1
void multinv_all(vector<11> &inv) {
    inv[1] = 1;
    repx(i, 2, inv.size())
        inv[i] = MOD - (MOD / i) * inv[MOD % i] % MOD;
}
```

8.12 polar-rho

```
ll mulmod(ll a, ll b, ll m) {
    ll r=a*b-(ll)((long double)a*b/m+.5)*m;
    return r<0?r+m:r;
}
bool is_prime_prob(ll n, int a){
    if(n==a)return true;
    ll s=0,d=n-1;
    while(d%2==0)s++,d/=2;
    ll x=expmod(a,d,n);
    if((x==1)||(x+1==n))return true;
    fore(_,0,s-1){
        x=mulmod(x,x,n);
        if(x=1)return false;
        if(x+1==n)return true;
}
return false;
}
bool rabin(ll n){ // true iff n is prime</pre>
```

```
if(n==1)return false:
int ar[]={2,3,5,7,11,13,17,19,23};
fore(i,0,9)if(!is_prime_prob(n,ar[i]))return false;
return true:
7
ll rho(ll n){
if(!(n&1))return 2;
11 x=2, y=2, d=1;
11 c=rand()%n+1;
while(d==1){
 x=(mulmod(x.x.n)+c)%n:
 fore(it.0.2) v=(mulmod(v,v,n)+c)%n:
 if(x>=y)d=_gcd(x-y,n);
 else d=__gcd(y-x,n);
}
return d==n?rho(n):d;
void fact(ll n, map<ll,int>& f){ //0 (lg n)^3
if(n==1)return:
if(rabin(n)){f[n]++:return:}
11 g=rho(n);fact(q,f);fact(n/q,f);
// optimized version: replace rho and fact with the
    following:
const int MAXP=1e6+1; // sieve size
int sv[MAXP]: // sieve
11 add(l1 a, l1 b, l1 m){return (a+=b)<m?a:a-m;}</pre>
ll rho(ll n){
static ll s[MAXP];
while(1){
 11 x=rand()%n,y=x,c=rand()%n;
 ll *px=s,*py=s,v=0,p=1;
 while(1){
  *py++=y=add(mulmod(y,y,n),c,n);
  *py++=y=add(mulmod(y,y,n),c,n);
  if((x=*px++)==v)break:
  11 t=p; p=mulmod(p,abs(y-x),n);
  if(!p)return __gcd(t,n);
  if(++v==26){
   if((p=_gcd(p,n))>1&&p<n)return p;</pre>
   v=0:
  }
 if(v&&(p=_gcd(p,n))>1&&p<n)return p;</pre>
void init_sv(){ fore(i,2,MAXP)if(!sv[i])for(ll j=i;j<MAXP;j</pre>
    +=i)sv[i]=i: }
void fact(ll n,map<ll,int>&f){//call init_sv first!
for(auto&& p:f)while(n%p.fst==0)p.snd++,n/=p.fst;
```

```
if(n<MAXP)while(n>1)f[sv[n]]++,n/=sv[n];
else if(rabin(n))f[n]++;
else {ll q=rho(n);fact(q,f);fact(n/q,f);}
}
```

8.13 polynomials

```
typedef int tp; // type of polynomial
template < class T=tp>
struct polv { // polv<> : 1 variable, polv<polv<>>: 2
    variables, etc.
vector<T> c;
T& operator[](int k){return c[k];}
poly(vector<T>& c):c(c){}
poly(initializer_list<T> c):c(c){}
poly(int k):c(k){}
polv(){}
poly operator+(poly<T> o);
poly operator*(tp k);
polv operator*(polv o);
poly operator-(poly<T> o){return *this+(o*-1);}
T operator()(tp v){
 T sum(0):
 for(int i=c.size()-1:i>=0:--i)sum=sum*v+c[i]:
 return sum;
// example: p(x,y)=2*x^2+3*x*y-y+4
// poly<poly<>> p={{4,-1},{0,3},{2}}
// printf("d\n",p(2)(3)) // 27 (p(2,3))
set<tp> roots(poly<> p){ // only for integer polynomials
while(!p.c.empty()&&!p.c.back())p.c.pop_back();
if(!p(0))r.insert(0);
if(p.c.empty())return r;
tp a0=0,an=abs(p[p.c.size()-1]);
for(int k=0;!a0;a0=abs(p[k++]));
vector<tp> ps,qs;
fore(i,1,sqrt(a0)+1)if(a0%i==0)ps.pb(i),ps.pb(a0/i);
fore(i,1,sqrt(an)+1)if(an\%i==0)qs.pb(i),qs.pb(an/i);
for(auto pt:ps)for(auto qt:qs)if(pt%qt==0){
 tp x=pt/at:
 if(!p(x))r.insert(x);
 if(!p(-x))r.insert(-x);
return r;
pair<poly<>,tp> ruffini(poly<> p, tp r){ // returns pair (
    result, rem)
```

```
int n=p.c.size()-1:
vector<tp> b(n);
b[n-1]=p[n];
for(int k=n-2;k>=0;--k)b[k]=p[k+1]+r*b[k+1];
return {poly<>(b),p[0]+r*b[0]};
// only for double polynomials
pair<poly<>,poly<> > polydiv(poly<> p, poly<> q){ // returns
     pair (result,rem)
int n=p.c.size()-q.c.size()+1;
vector<tp> b(n):
for(int k=n-1:k>=0:--k){
 b[k]=p.c.back()/q.c.back();
 fore(i,0,q.c.size())p[i+k]-=b[k]*q[i];
 p.c.pop_back();
while(!p.c.empty()&&abs(p.c.back())<EPS)p.c.pop_back();</pre>
return {poly<>(b),p};
// only for double polynomials
poly<> interpolate(vector<tp> x, vector<tp> y){
poly<> q={1},S={0};
for(tp a:x)q=poly<>({-a,1})*q;
fore(i,0,x.size()){
 poly<> Li=ruffini(q,x[i]).fst;
 Li=Li*(1.0/Li(x[i])); // change for int polynomials
 S=S+Li*y[i];
}
return S;
```

8.14 primes

```
// counts the divisors of a positive integer in O(sqrt(n))
ll count_divisors(ll x) {
    ll divs = 1, i = 2;
    for (ll divs = 1, i = 2; x > 1; i++) {
        if (i * i > x) { divs *= 2; break; }
        for (ll d = divs; x % i == 0; x /= i) divs += d;
    }
    return divs;
}

// gets the prime factorization of a number in O(sqrt(n))
vector<pair<ll, int>> factorize(ll x) {
    vector<pair<ll, int>> f;
    for (ll k = 2; x > 1; k++) {
        if (k * k > x) { f.push_back({x, 1}); break; }
        int n = 0;
```

```
while (x \% k == 0) x /= k, n++:
       if (n > 0) f.push_back(\{k, n\});
    return f:
}
// iterate over all divisors of a number.
// divisor count upper bound: n^(1.07 / ln ln n)
template <class OP>
void divisors(11 x, OP op) {
    auto facts = factorize(x):
    vector<int> f(facts.size()):
    while (true) {
       11 v = 1:
       rep(i, f.size()) rep(j, f[i]) v *= facts[i].first;
       (v)qo
       int i;
       for (i = 0: i < f.size(): i++) {</pre>
          f[i] += 1:
          if (f[i] <= facts[i].second) break;</pre>
           f[i] = 0:
       }
       if (i == f.size()) break;
// computes euler totative function phi(x), counting the
// amount of integers in [1, x] that are coprime with x.
// time: O(sart(x))
11 phi(11 x) {
   11 \text{ phi} = 1, k = 2;
    for (: x > 1: k++) {
       if (k * k > x) { phi *= x - 1; break; }
       11 k1 = 1, k0 = 0:
       while (x \% k == 0) x /= k, k0 = k1, k1 *= k:
       phi *= k1 - k0:
    return phi;
```

8.15 simplex

```
// The input vector is set to an optimal $x$ (or in the
    unbounded case, an arbitrary solution fulfilling the
    constraints).
// Numerical stability is not guaranteed. For better
    performance, define variables such that x = 0 is
// Usage:
// vvd A = \{\{1,-1\}, \{-1,1\}, \{-1,-2\}\};
// vd b = \{1,1,-4\}, c = \{-1,-1\}, x;
// T val = LPSolver(A, b, c).solve(x);
// Time: O(NM * \#pivots), where a pivot may be e.g. an edge
     relaxation. O(2^n) in the general case.
#include "../common.h"
typedef double T; // long double, Rational, double + mod<P
typedef vector<T> vd:
typedef vector<vd> vvd:
const T eps = 1e-8, inf = 1 / .0;
#define MP make_pair
#define ltj(X) \
   if (s == -1 || MP(X[i], N[i]) < MP(X[s], N[s])) s = i
struct LPSolver {
   int m. n:
   vector<int> N. B:
   vvd D:
   LPSolver(const vvd &A, const vd &b, const vd &c) : m(b.
        size()), n(c.size()), N(n + 1), B(m), D(m + 2), vd(n)
        + 2)) {
       rep(i, m) rep(j, n) D[i][j] = A[i][j];
       rep(i, m) {
          B[i] = n + i;
          D[i][n] = -1:
          D[i][n + 1] = b[i];
       rep(j, n) {
          N[i] = i:
          D[m][i] = -c[i];
      }
       N[n] = -1:
       D[m + 1][n] = 1;
   void pivot(int r, int s) {
      T *a = D[r].data(), inv = 1 / a[s]:
       rep(i, m + 2) if (i != r \&\& abs(D[i][s]) > eps) {
```

```
T *b = D[i].data(), inv2 = b[s] * inv:
          repx(j, 0, n + 2) b[j] -= a[j] * inv2;
          b[s] = a[s] * inv2;
       rep(j, n + 2) if (j != s) D[r][j] *= inv;
       rep(i, m + 2) if (i != r) D[i][s] *= -inv:
       D[r][s] = inv;
       swap(B[r], N[s]);
   bool simplex(int phase) {
       int x = m + phase - 1:
      for (;;) {
          int s = -1:
          rep(j, n + 1) if (N[j] != -phase) ltj(D[x]);
          if (D[x][s] >= -eps) return true;
          int r = -1:
          rep(i, m) {
              if (D[i][s] <= eps) continue:</pre>
              if (r == -1 || MP(D[i][n + 1] / D[i][s], B[i])
                    < MP(D[r][n + 1] / D[r][s], B[r])) r = i
          if (r == -1) return false:
          pivot(r, s);
      }
   }
   T solve(vd &x) {
       int r = 0:
       repx(i, 1, m) if (D[i][n + 1] < D[r][n + 1]) r = i;
       if (D[r][n + 1] < -eps) {
          pivot(r, n):
          if (!simplex(2) || D[m + 1][n + 1] < -eps) return
          rep(i, m) if (B[i] == -1) {
              int s = 0:
              repx(j, 1, n + 1) ltj(D[i]);
              pivot(i, s);
       bool ok = simplex(1);
       x = vd(n):
       rep(i, m) if (B[i] < n) \times [B[i]] = D[i][n + 1];
       return ok ? D[m][n + 1] : inf;
};
```

8.16 test-prime

8.17 theorems

Burnside lemma

Tomemos imagenes x en X y operaciones (g: X -> X) en G. Si #g es la cantidad de imagenes que son puntos fijos de g, entonces la cantidad de objetos es '(sum_{g in G} #g) / |G|' Es requisito que G tenga la operacion identidad, que toda operacion tenga inversa y que todo par de operaciones tenga su combinacion.

Rational root theorem

Las raices racionales de un polinomio de orden n con coeficientes enteros A[i] son de la forma p / q, donde p y q son coprimos, p es divisor de A[0] y q es divisor de A[n]. Notar que si A[0] = 0, cero es raiz, se puede dividir el polinomio por x y aplica nuevamente el teorema.

Petersens theorem

Every cubic and bridgeless graph has a perfect matching.

Number of divisors for powers of 10 (0,1) (1,4) (2,12) (3,32) (4,64) (5,128) (6,240) (7,448) (8,768) (9,1344) (10,2304) (11,4032) (12,6720) (13,10752) (14,17280) (15,26880) (16,41472) (17,64512) (18,103680)

Kirchoff Theorem: Sea A la matriz de adyacencia del multigrafo (A[u][v] indica la cantidad de aristas entre u y v) Sea D una matriz diagonal tal que D[v][v] es igual al grado de v (considerando auto aristas y multi aristas). Sea L = A - D. Todos los cofactores de L son iguales y equivalen a la cantidad de Spanning Trees del grafo. Un cofactor (i,j) de L es la multiplicacin de $(-1)^{i}$ + j con el determinant de la matriz al quitar la fila i y la columna j

Prufer Code: Dado un rbol con los nodos indexados: busca la hoja de menor ndice, brrala y anota el ndice del nodo al que estaba conectado. Repite el paso anterior n-2 veces. Lo anterior muestra una biyeccin entre los arreglos de tamao n-2 con elementos en [1, n] y los rboles de n nodos, por lo que hay n^{n-2} spanning trees en un grafo completo. Corolario: Si tenemos k componentes de tamaos n-20 spanning trees en un grafo completo. Corolario: Si tenemos k componentes de tamaos n-20 spanning trees en un grafo completo. Corolario: Si tenemos k componentes de tamaos n-20 spanning trees en un grafo completo. Corolario: Si tenemos k componentes de tamaos n-20 formas entre nodos de n-20 formas n-20 formas n-21 formas n-22 formas n-23 formas n-24 formas n-25 formas n-25 formas n-26 formas n-26 formas n-27 formas n-28 formas n-29 formas formas n-29 formas n-29 formas f

Catalan: C_n = \frac{1}{n+1}*\binom{2n, n} Sea C_n^k las formas de poner n+k pares de parntesis, con los primeros k parntesis abiertos (esto es, hay 2n + 2k carcteres), se tiene que C_n^k = (2n+k-1)*(2n+k)/(n*(n+k+1)) * C_{n-1}^k Sea D_n el nmero de permutaciones sin puntos fijos, entoces D n = (n-1)*(D {n-1} + D {n-2}), D 0 = 1, D 1 = 0

8.18 tonelli-shanks

```
ll legendre(ll a, ll p) {
    if (a % p == 0) return 0; if (p == 2) return 1;
    return binexp(a, (p-1) / 2, p);
// sqrt(n) mod p (p must be a prime)
// rnd(a, b) return a random number in [a, b]
ll tonelli shanks(ll n. ll p) {
    if (n == 0) return 0;
    if (legendre(n, p) != 1) return -1; // no existe
    if (p == 2) return 1;
    ll s = __builtin_ctzll(p - 1);
    11 q = (p - 1LL) >> s, z = rnd(1, p - 1);
    if (s == 1) return binexp(n, (p + 1) / 4LL, p);
    while (legendre(z, p) != p - 1) z = rnd(1, p - 1);
    ll c = binexp(z, q, p), r = binexp(n, (q + 1) / 2, p):
    ll t = binexp(n, q, p), m = s;
    while (t != 1) {
       11 i = 1, ts = (t * t) % p:
       while (ts != 1) i++, ts = (ts * ts) % p;
       11 b = c:
       repx(_, 0, m - i - 1) b = (b * b) \% p;
       r = r*b\%p; c = b*b\%p; t = t*c\%p; m = i;
    return r;
```

9 strings

9.1 Manacher

```
// odd[i]: length of longest palindrome centered at i
// even[i]: ...longest palindrome centered between i and i+1
void manacher(string &s,vector<int> &odd,vector<int> &even){
   string t = "$#";
   for(char c: s) t += c + string("#");
   t += "^":
   int n = t.size();
   vector<int> p(n);
   int 1 = 1, r = 1;
   repx(i, 1, n-1) {
      p[i] = max(0, min(r - i, p[1 + (r - i)]));
       while(t[i - p[i]] == t[i + p[i]]) p[i]++;
       if(i + p[i] > r) l = i - p[i], r = i + p[i];
   repx(i, 2, n-2) {
       if(i%2) even.push_back(p[i]-1);
       else odd.push_back(p[i]-1);
```

9.2 aho-corasick

```
struct Vertex {
   int next[26], go[26];
   int p, link = -1, exit = -1, cnt = -1;
   vector<int> leaf:
   Vertex(int p=-1, char ch='$') : p(p), pch(ch) {
       rep(i, 26) next[i] = -1, go[i] = -1:
vector<Vertex> t(1):
void add(string &s, int id) {
   int v = 0:
   for (char ch : s) {
      int c = ch - 'a';
      if (t[v].next[c] == -1) {
          t[v].next[c] = t.size();
          t.emplace_back(v, ch);
       v = t[v].next[c];
   t[v].leaf.push_back(id);
```

```
int go(int v. char ch):
int get_link(int v) {
   if (t[v].link == -1) {
       if (v == 0 | | t[v].p == 0) t[v].link = 0:
       else t[v].link = go(get_link(t[v].p), t[v].pch);
   return t[v].link;
}
int go(int v, char ch) {
   int c = ch - 'a';
   if (t[v].go[c] == -1) {
       if (t[v].next[c] != -1) t[v].go[c] = t[v].next[c]:
       else t[v].go[c] = v == 0 ? 0 : go(get_link(v), ch);
   return t[v].go[c];
int next match(int v){ // Optional
   if(t[v].exit == -1){
       if(t[get link(v)].leaf.size())t[v].exit=get link(v):
       else t[v].exit = v==0 ? 0 : next match(get link(v)):
   return t[v].exit:
int cnt_matches(int v){ // Optional
   if(t[v].cnt == -1)
       t[v].cnt = v == 0 ? 0 : t[v].leaf.size() +
            cnt_matches(get_link(v));
   return t[v].cnt:
```

9.3 hash

```
const int K = 2;
struct Hash{
   const 11 MOD[K] = {999727999, 1070777777};
   const 11 P = 1777771;
   vector<ll> h[K], p[K]:
   Hash(string &s){
      int n = s.size():
      rep(k, K){
          h[k].resize(n+1, 0);
          p[k].resize(n+1, 1);
          repx(i, 1, n+1){
              h[k][i] = (h[k][i-1]*P + s[i-1]) % MOD[k];
              p[k][i] = (p[k][i-1]*P) % MOD[k];
          }
      }
   vector<ll> get(int i, int j){
```

```
vector<ll> r(K);
rep(k, K){
    r[k] = (h[k][j] - h[k][i]*p[k][j-i]) % MOD[k];
    r[k] = (r[k] + MOD[k]) % MOD[k];
} return r;
}
};
```

9.4 palindromic-tree

```
struct Node { // (*) = Optional
 int len; // length of substring
 int to[26]; // insertion edge for all characters a-z
 int link; // maximum palindromic suffix
           // (*) start index of current Node
   int cnt: // (*) # of occurrences of this substring
   Node(int len, int link=0, int i=0, int cnt=1): len(len),
   link(link), i(i), cnt(cnt) {memset(to, 0, sizeof(to));}
}: struct EerTree { // Palindromic Tree
   vector<Node> t: // tree (max size of tree is n+2)
   int last:
                // current node
   EerTree(string &s) : last(0) {
       t.emplace_back(-1); t.emplace_back(0); // root 1 & 2
       rep(i, s.size()) add(i, s): // construct tree
       for(int i = t.size()-1; i > 1; i--)
          t[t[i].link].cnt += t[i].cnt:
   void add(int i, string &s){
                                   // vangrind warning:
       int p=last, c=s[i]-'a';
                                 // i-t[p].len-1 = -1
       while(s[i-t[p].len-1] != s[i]) p = t[p].link;
       if(t[p].to[c]){ last = t[p].to[c]; t[last].cnt++; }
       else{
          int q = t[p].link;
          while(s[i-t[q].len-1] != s[i]) q = t[q].link;
          q = max(1, t[q].to[c]);
          last = t[p].to[c] = t.size();
          t.emplace back(t[p].len + 2, q, i-t[p].len-1):
   }
}:
 string s = "abcbab"; EerTree pt(s); // build EerTree
 repx(i, 2, pt.t.size()){// list all distinct palindromes
 repx(j,pt.t[i].i,pt.t[i].i+pt.t[i].len)cout << s[j];
 cout << " " << pt.t[i].cnt << endl;</pre>
```

9.5 prefix-function

```
vector<int> prefix_function(string s) {
   int n = s.size();
   vector<int> pi(n);
   repx(i, 1, n) {
       int j = pi[i-1];
       while (j > 0 \&\& s[i] != s[i])
          j = pi[j-1];
       if (s[i] == s[i])
          j++:
       pi[i] = j;
   return pi;
vector<vector<int>> aut:
void compute_automaton(string s) {
   s += '#':
   int n = s.size():
   vector<int> pi = prefix_function(s);
   aut.assign(n, vector<int>(26));
   rep(i, n) {
      rep(c, 26) {
          int i = i:
          while (j > 0 \&\& 'a' + c != s[j])
              j = pi[j-1];
          if ('a' + c == s[j])
              j++;
          aut[i][c] = j;
      }
// k = n - pi[n - 1]; if k divides n, then the string can be
// aprtitioned into blocks of length k otherwise there is no
// effective compression and the answer is n.
```

9.6 suffix-array

```
for (int k = 1: k < N: k <<= 1) {</pre>
       int C = c[p[N - 1]] + 1:
       cnt.assign(C + 1, 0);
       for (int &pi : p) pi = (pi - k + N) % N;
       for (int cl : c) cnt[cl + 1] += 1;
       rep(i, C) cnt[i + 1] += cnt[i];
       rep(i, N) p2[cnt[c[p[i]]]++] = p[i];
       c2[p2[0]] = 0;
       repx(i, 1, N) c2[p2[i]] =
           c2[p2[i-1]] + (c[p2[i]] != c[p2[i-1]] ||
                          c[(p2[i] + k) \% N] != c[(p2[i - 1]
                                + k) % N]):
       swap(c, c2), swap(p, p2);
   p.erase(p.begin()); // optional: erase terminating NUL
   return p;
// build the lcp
// 'lcp[i]' represents the length of the longest common
// prefix between suffix i and suffix i+1 in the suffix
//array 'p'. the last element of 'lcp' is zero by convention
vector<int> makelcp(const string &s, const vector<int> &p) {
   int N = p.size(), k = 0;
   vector<int> r(N), lcp(N);
   rep(i, N) r[p[i]] = i;
   rep(i, N) {
       if (r[i] + 1 >= N) { k = 0; continue; }
       int i = p[r[i] + 1]:
       while (i + k < N \&\& j + k < N \&\& s[i + k] == s[j + k]
            1) k += 1:
       lcp[r[i]] = k;
       if (k) k -= 1;
```

```
return lcp;
}
// lexicographically compare the suffixes starting from 'i'
// and 'j', considering only up to 'K' characters.
// 'r' is the inverse suffix array, mapping suffix offsets
// to indices. requires an LCP sparse table.
int lcp_cmp(vector<int> &r, Sparse<int> &lcp, int i, int j,
    int K) {
    if (i == j) return 0;
    int ii = r[i], jj = r[j];
    int l = lcp.query(min(ii, jj), max(ii, jj));
    if (l >= K) return 0;
    return ii < jj ? -1 : 1;
}</pre>
```

9.7 suffix-automaton

```
st[p].next[c] = w;
st[q].link=st[k].link = w;
}
last = k;
} // # states <= 2n-1 && transitions <= 3n-4 (for n > 2)
// Follow link from 'last' to 0, nodes on path are terminal
// # matches = # paths from state to a terminal node
// # substrings = # paths from 0 to any node
// # substrings = sum of (len - len(link)) for all nodes
```

9.8 z-function

```
// i-th element is equal to the greatest number of
// characters starting from the position i that coincide
// with the first characters of s
vector<int> z_function(string s) {
   int n = s.size();
   vector<int> z(n);
   int l = 0, r = 0;
   for(int i = 1; i < n; i++) {
      if(i < r) z[i] = min(r - i, z[i - 1]);
      while(i + z[i] < n && s[z[i]] == s[i + z[i]])z[i]++;
      if(i + z[i] > r) {
        l = i;
        r = i + z[i];
      }
   }
   return z;
}
```