# Team Notebook

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# 1 Strings

#### 1.1 Manacher

```
#include<bits/stdc++.h>
using namespace std;
#define rep(i, n) for (int i = 0; i < (int)n; i++)
#define repx(i, a, b) for (int i = (int)a; i < (int)b; i++)</pre>
// odd[i] : length of the longest palindrome centered at i
// even[i] : length of the longest palindrome centered
    between i and i+1
void manacher(string &s, vector<int> &odd, vector<int> &even
    ) {
   string t = "$#":
   for(char c: s)
       t += c + string("#");
   t += "^":
   int n = t.size();
   vector<int> p(n);
   int 1 = 1, r = 1;
   repx(i, 1, n-1) {
       p[i] = max(0, min(r - i, p[1 + (r - i)]));
       while(t[i - p[i]] == t[i + p[i]]) {
          p[i]++;
       if(i + p[i] > r) {
          1 = i - p[i], r = i + p[i];
      }
   repx(i, 2, n-2) {
       if(i%2) even.push_back(p[i]-1);
       else odd.push_back(p[i]-1);
```

#### 1.2 aho-corasick

```
#include "../common.h"

const int K = 26;
struct Vertex {
   int next[K];
   int leaf = 0;
   int leaf_id = -1;
   int p = -1;
   char pch;
```

```
int link = -1:
   int exit = -1:
   int cnt = -1;
   int go[K];
   Vertex(int p=-1, char ch='$') : p(p), pch(ch) {
       fill(begin(next), end(next), -1);
       fill(begin(go), end(go), -1);
};
vector<Vertex> t(1):
void add(string &s, int id) {
   for (char ch : s) {
       int c = ch - 'a':
       if (t[v].next[c] == -1) {
           t[v].next[c] = t.size():
           t.emplace_back(v, ch);
       v = t[v].next[c];
   t[v].leaf++:
   t[v].leaf_id = id;
int go(int v. char ch):
int get_link(int v) {
   if (t[v].link == -1) {
       if (v == 0 || t[v].p == 0)
           t[v].link = 0:
           t[v].link = go(get_link(t[v].p), t[v].pch);
   return t[v].link;
int go(int v, char ch) {
   int c = ch - 'a':
   if (t[v].go[c] == -1) {
       if (t[v].next[c] != -1)
           t[v].go[c] = t[v].next[c];
           t[v].go[c] = v == 0 ? 0 : go(get_link(v), ch);
   return t[v].go[c];
```

#### 1.3 hash

```
// compute substring hashes in O(1).
// hashes are compatible between different strings.
struct Hash {
   11 HMOD:
   int N;
   vector<int> h;
   vector<int> p;
   Hash() {}
   // O(N)
   Hash(const string &s, 11 HMOD_ = 1000003931)
       : N(s.size() + 1), HMOD(HMOD_), p(N), h(N) {
       static const 11 P =
          chrono::steady_clock::now().time_since_epoch().
               count() % (1 << 29):
      p[0] = 1;
       rep(i, N - 1) p[i + 1] = p[i] * P % HMOD;
       rep(i, N-1) h[i+1] = (h[i] + (ll)s[i] * p[i]) %
   }
   pair<ll, int> get(int i, int j) { return {(h[j] - h[i] +
        HMOD) % HMOD, i}; }
   bool cmp(pair<11, int> x0, pair<11, int> x1) {
       int d = x0.second - x1.second;
```

```
ll &lo = d < 0 ? x0.first : x1.first:
       lo = lo * p[abs(d)] % HMOD;
       return x0.first == x1.first;
};
// compute hashes in multiple prime modulos simultaneously,
     to reduce the chance
// of collisions.
struct HashM {
    int N:
    vector<Hash> sub:
    HashM() {}
    // D(K N)
    HashM(const string &s, const vector<11> &mods) : N(mods.
        size()), sub(N) {
       rep(i, N) sub[i] = Hash(s, mods[i]);
    // O(K)
    vector<pair<11, int>> get(int i, int j) {
       vector<pair<11, int>> hs(N);
       rep(k, N) hs[k] = sub[k].get(i, j);
       return hs;
    bool cmp(const vector<pair<11, int>> &x0, const vector<</pre>
        pair<11, int>> &x1) {
       rep(i, N) if (!sub[i].cmp(x0[i], x1[i])) return false
       return true;
    bool cmp(int i0, int j0, int i1, int j1) {
       rep(i, N) if (!sub[i].cmp(sub[i].get(i0, j0),
                               sub[i].get(i1, j1))) return
                                    false:
       return true;
};
```

### 1.4 hash2d

```
using Hash = pair<11, int>;
struct Block {
   int x0, y0, x1, y1;
};
```

```
struct Hash2d {
   11 HMOD:
   int W. H:
   vector<int> h;
   vector<int> p:
   Hash2d() \{ \}
    Hash2d(const string &s, int W_, int H_, 11 HMOD_ =
        1000003931)
       : W(W + 1), H(H + 1), HMOD(HMOD) {
       static const 11 P =
           chrono::steady_clock::now().time_since_epoch().
                count() % (1 << 29):
       p.resize(W * H):
       p[0] = 1;
       rep(i, W * H - 1) p[i + 1] = p[i] * P % HMOD;
       h.assign(W * H, 0);
       repx(y, 1, H) repx(x, 1, W) {
           ll c = (ll)s[(y - 1) * (W - 1) + x - 1] * p[y * W
                 + x] % HMOD;
           h[y * W + x] = (HMOD + h[y * W + x - 1] + h[(y - x)]
               1) * W + xl -
                         h[(y-1) * W + x - 1] + c) %
    bool isout(Block s) {
       return s.x0 < 0 || s.x0 >= W || s.x1 < 0 || s.x1 >= W
             || s.y0 < 0 ||
              s.y0 >= H || s.y1 < 0 || s.y1 >= H;
   }
   Hash get(Block s) {
       return \{(2 * HMOD + h[s.v1 * W + s.x1] - h[s.v1 * W +
               h[s.y0 * W + s.x1] + h[s.y0 * W + s.x0]) %
              s.v0 * W + s.x0:
   }
    bool cmp(Hash x0, Hash x1) {
       int d = x0.second - x1.second;
       11 &lo = d < 0 ? x0.first : x1.first;</pre>
       lo = lo * p[abs(d)] % HMOD;
       return x0.first == x1.first;
};
```

```
struct Hash2dM {
   int N:
   vector<Hash2d> sub:
   Hash2dM() {}
   Hash2dM(const string &s. int W. int H. const vector<11> &
       : N(mods.size()), sub(N) {
       rep(i, N) sub[i] = Hash2d(s, W, H, mods[i]);
   bool isout(Block s) { return sub[0].isout(s): }
   vector<Hash> get(Block s) {
       vector<Hash> hs(N);
       rep(i, N) hs[i] = sub[i].get(s);
       return hs:
   }
   bool cmp(const vector<Hash> &x0, const vector<Hash> &x1)
       rep(i, N) if (!sub[i].cmp(x0[i], x1[i])) return false
       return true;
   bool cmp(Block s0, Block s1) {
       rep(i, N) if (!sub[i].cmp(sub[i].get(s0), sub[i].get(
           s1))) return false;
       return true:
   }
};
const vector<11> HMOD = {1000002649, 1000000933, 1000003787,
     1000002173}:
```

# 1.5 palindromic-tree

};

```
Node(){ fill(begin(edge), end(edge), -1): }
struct EerTree { // Palindromic Tree
   vector<Node> t; // tree
   int curr:
              // current node
   EerTree(string &s) {
       t.resize(2):
       t.reserve(s.size()+2); // (optional) maximum size of
           tree
       t[0].len = -1:
                           // root 1
      t[0].link = 0;
      t[1].len = 0:
                           // root 2
      t[1].link = 0;
      curr = 1;
      rep(i, s.size()) insert(i, s): // construct tree
      // (optional) calculate number of occurrences of each
       for(int i = t.size()-1; i > 1; i--)
          t[t[i].link].cnt += t[i].cnt;
   void insert(int i, string &s) {
      int tmp = curr;
       while (i - t[tmp].len < 1 \mid | s[i] != s[i-t[tmp].len
           -11)
          tmp = t[tmp].link;
      if(t[tmp].edge[s[i]-'a'] != -1){
          curr = t[tmp].edge[s[i]-'a']; // node already
               exists
          t[curr].cnt++;
                                         // (optional)
               increase cnt
          return:
      }
       curr = t[tmp].edge[s[i]-'a'] = t.size(); // create
           new node
       t.emplace_back();
      t[curr].len = t[tmp].len + 2:
                                         // set length
      t[curr].i = i - t[curr].len + 1; // (optional) set
            start index
       if (t[curr].len == 1) {
                                         // set suffix link
          t[curr].link = 1:
      } else {
          tmp = t[tmp].link;
```

```
while (i-t[tmp].len < 1 || s[i] != s[i-t[tmp].len
              tmp = t[tmp].link;
          t[curr].link = t[tmp].edge[s[i]-'a'];
};
int main()
 string s = "abcbab":
   EerTree pt(s): // construct palindromic tree
repx(i, 2, pt.t.size()) // list all distinct palindromes
 cout << i-1 << ") ":
 repx(j, pt.t[i].i, pt.t[i].i + pt.t[i].len)
  cout << s[i]:
 cout << " " << pt.t[i].cnt << endl;</pre>
return 0;
```

# 1.6 prefix-function

```
#include "../common.h"
vector<int> prefix_function(string s) {
   int n = s.size();
   vector<int> pi(n):
   repx(i, 1, n) {
      int j = pi[i-1];
      while (j > 0 \&\& s[i] != s[j])
          j = pi[j-1];
       if (s[i] == s[j])
          j++:
       pi[i] = j;
   return pi;
vector<vector<int>> aut;
void compute_automaton(string s) {
   s += '#':
   int n = s.size();
   vector<int> pi = prefix_function(s);
   aut.assign(n, vector<int>(26));
   rep(i, n) {
```

```
rep(c, 26) {
           int j = i;
           while (j > 0 \&\& 'a' + c != s[j])
              j = pi[j-1];
          if ('a' + c == s[i])
              j++:
           aut[i][c] = j;
       }
   }
// k = n - pi[n - 1]
// if k divides n, then the string can be aprtitioned into
    blocks of length k
// otherwise there is no effective compression and the
    answer is n.
```

### suffix-array-martin

```
// build the suffix array
// suffixes are sorted, with each suffix represented by its
    starting position
vector<int> suffixarray(const string &s) {
   int N = s.size() + 1: // optional: include terminating
   vector\langle int \rangle p(N), p2(N), c(N), c2(N), cnt(256);
   rep(i, N) cnt[s[i]] += 1;
   repx(b, 1, 256) cnt[b] += cnt[b - 1];
   rep(i, N) p[--cnt[s[i]]] = i;
   repx(i, 1, N) c[p[i]] = c[p[i - 1]] + (s[p[i]] != s[p[i - 1]])
         111):
   for (int k = 1: k < N: k <<= 1) {
       int C = c[p[N - 1]] + 1;
       cnt.assign(C + 1, 0):
       for (int &pi : p) pi = (pi - k + N) % N;
       for (int cl : c) cnt[cl + 1] += 1;
       rep(i, C) cnt[i + 1] += cnt[i];
       rep(i, N) p2[cnt[c[p[i]]]++] = p[i];
       c2[p2[0]] = 0;
       repx(i, 1, N) c2[p2[i]] =
          c2[p2[i-1]] + (c[p2[i]] != c[p2[i-1]] ||
                          c[(p2[i] + k) \% N] != c[(p2[i - 1]
                                + k) % N]);
       swap(c, c2), swap(p, p2);
   p.erase(p.begin()); // optional: erase terminating NUL
   return p;
```

```
// build the lcp
// 'lcp[i]' represents the length of the longest common
    prefix between suffix i
// and suffix i+1 in the suffix array 'p'. the last element
    of 'lcp' is zero by
// convention
vector<int> makelcp(const string &s, const vector<int> &p) {
   int N = p.size(), k = 0;
   vector<int> r(N), lcp(N);
   rep(i, N) r[p[i]] = i;
   rep(i, N) {
       if (r[i] + 1 >= N) {
          k = 0:
           continue:
       int j = p[r[i] + 1];
       while (i + k < N \&\& i + k < N \&\& s[i + k] == s[i + k]
           ]) k += 1;
       lcp[r[i]] = k:
       if (k) k -= 1:
   return lcp;
// lexicographically compare the suffixes starting from 'i'
    and 'i',
// considering only up to 'K' characters.
// 'r' is the inverse suffix array, mapping suffix offsets
     to indices.
// requires an LCP sparse table.
int lcp_cmp(vector<int> &r, Sparse<int> &lcp, int i, int j,
    int K) {
   if (i == j) return 0;
   int ii = r[i], jj = r[j];
   int 1 = lcp.query(min(ii, jj), max(ii, jj));
   if (1 >= K) return 0:
   return ii < j; ? -1 : 1;
}
```

## 1.8 suffix-array

```
#include "../common.h"

struct SuffixArray {
  int n; vector<int> C, R, R_, sa, sa_, lcp;
  inline int gr(int i) { return i < n ? R[i] : 0; } // sort
      suffixes</pre>
```

```
//inline int gr(int i) { return R[i%n]: } // sort
    cvclic shifts
void csort(int maxv, int k) {
   C.assign(maxv + 1, 0); rep(i, n) C[gr(i + k)]++;
   repx(i, 1, maxv + 1) C[i] += C[i - 1];
   for (int i = n - 1; i \ge 0; i--) sa [--C[gr(sa[i] + k]]
       )]] = sa[i];
   sa.swap(sa_);
void getSA(vector<int>& s) {
   R = R = sa = sa = vector < int > (n): rep(i, n) sa[i] =
   sort(sa.begin(), sa.end(), [&s](int i, int j) {
        return s[i] < s[j]; });</pre>
   int r = R[sa[0]] = 1;
   repx(i, 1, n) R[sa[i]] = (s[sa[i]] != s[sa[i - 1]]) ?
   for (int h = 1; h < n && r < n; h <<= 1) {
       csort(r, h): csort(r, 0): r = R [sa[0]] = 1:
       repx(i, 1, n) {
          if (R[sa[i]] != R[sa[i - 1]] || gr(sa[i] + h)
               != gr(sa[i - 1] + h)) r++:
          R_[sa[i]] = r:
       } R.swap(R_);
void getLCP(vector<int> &s) {// only works with suffixes
     (not cyclic shifts)
   lcp.assign(n, 0); int k = 0;
   rep(i, n) {
       int r = R[i] - 1;
       if (r == n - 1) \{ k = 0; continue; \}
       int i = sa[r + 1]:
       while (i + k < n &  i + k < n &  s[i + k] == s[i]
           + kl) k++:
       lcp[r] = k: if (k) k--:
SuffixArray(vector<int> &s) { n = s.size(); getSA(s);
    getLCP(s); constructLCP(); }
/* ----- */
vector<vector<int>> T:
void constructLCP() {
   T.assign(LOG2(n)+1, lcp);
   for(int k = 1: (1<<k) <= n: ++k)
       for(int i = 0; i + (1<<k) <= n; ++i)</pre>
          T[k][i] = min(T[k-1][i],T[k-1][i+(1<<(k-1))]):
}
```

```
// get LCP of suffix starting at i and suffix starting at
   int queryLCP(int i, int j) {
       if(i == i) return n-i;
       i = R[i]-1; j = R[j]-1;
       if(i > j) swap(i, j);
       11 k = LOG2(j-i);
       return min(T[k][i],T[k][j-(1<<k)]);</pre>
   // compare substring of length len1 starting at i
   // with substring of length len2 starting at j
   bool cmp(int i, int len1, int i, int len2) {
       if(queryLCP(i, j) >= min(len1, len2))
           return (len1 < len2):</pre>
           return (R[i] < R[j]);</pre>
};
vector<int> suffix array:
vector<vector<int>> C;
int n:
void sort_cyclic_shifts(string s) {
   s += "$":
   n = s.size();
   const int alphabet = 256;
   vector<int> p(n), c(n), cnt(max(alphabet, n), 0);
   for (int i = 0; i < n; i++)
       cnt[s[i]]++:
   for (int i = 1; i < alphabet; i++)</pre>
       cnt[i] += cnt[i-1];
   for (int i = 0: i < n: i++)
       p[--cnt[s[i]]] = i;
   c[p[0]] = 0;
   int classes = 1:
   for (int i = 1; i < n; i++) {</pre>
       if (s[p[i]] != s[p[i-1]])
           classes++;
       c[p[i]] = classes - 1;
   C.emplace_back(c.begin(), c.end());
   vector<int> pn(n), cn(n);
   for (int h = 0; (1 << h) < n; ++h) {
       for (int i = 0; i < n; i++) {</pre>
           pn[i] = p[i] - (1 << h);
           if (pn[i] < 0)</pre>
              pn[i] += n;
       fill(cnt.begin(), cnt.begin() + classes, 0);
```

```
for (int i = 0: i < n: i++)</pre>
           cnt[c[pn[i]]]++:
       for (int i = 1; i < classes; i++)</pre>
           cnt[i] += cnt[i-1]:
       for (int i = n-1; i >= 0; i--)
           p[--cnt[c[pn[i]]]] = pn[i]:
       cn[p[0]] = 0;
       classes = 1;
       for (int i = 1; i < n; i++) {</pre>
           pair < int, int > cur = \{c[p[i]], c[(p[i] + (1 << h))\}
                ) % nl}:
           pair<int, int> prev = \{c[p[i-1]], c[(p[i-1] + (1
                << h)) % n]};
           if (cur != prev)
               ++classes:
           cn[p[i]] = classes - 1;
       c.swap(cn);
       C.emplace back(c.begin(), c.end());
    p.erase(p.begin());
    suffix_array = p;
}
vector<int> lcp_construction(string &s, vector<int> &p) {
    int n = s.size():
    vector<int> rank(n):
    rep(i, n) rank[p[i]] = i:
    int k = 0:
    vector<int> lcp(n-1, 0);
    rep(i, n) {
       if (rank[i] == n - 1) {
           k = 0;
           continue:
       int j = p[rank[i] + 1];
       while (i + k < n \&\& j + k < n \&\& s[i+k] == s[j+k])
           k++;
       lcp[rank[i]] = k;
       if (k)
           k--;
    return lcp;
bool compare1(int i, int j, int l) {
    int k = LOG2(1):
    pair<int, int> a = \{C[k][i], C[k][(i+1-(1 << k))\%n]\};
    pair<int, int> b = \{C[k][j], C[k][(j+1-(1 << k))\%n]\};
```

```
return a >= b:
bool compare2(int i, int j, int l) {
   int k = LOG2(1);
   pair<int, int> a = \{C[k][i], C[k][(i+1-(1 << k))\%n]\};
   pair<int, int> b = \{C[k][j], C[k][(j+1-(1 << k))\%n]\};
   return a <= b:
pair<int.int> find(int i, int len)
   int 1 = 0, r = suffix_array.size()-1;
   while(1 != r)
       int mid = (1+r)/2:
       if(compare1(suffix_array[mid], i, len))
           r = mid:
       else
          l = mid+1;
   int left = 1;
   1 = 0, r = suffix_array.size()-1;
   while(1 != r)
       int mid = (1+r+1)/2:
       if(compare2(suffix_array[mid], i, len))
          1 = mid:
       else
          r = mid-1;
   int right = 1;
   if(!compare1(suffix_array[left], i, len)) return {-1,-1};
   if(!compare2(suffix_array[right], i, len)) return
        {-1,-1}:
   if(left > right) return {-1,-1};
   return {left, right};
```

### 1.9 suffix-automaton

```
#include "../common.h"
struct SuffixAutomaton {
```

```
vector<map<char,int>> edges; // edges[i] : the labeled
    edges from node i
vector<int> link:
                          // link[i] : the suffix link
    of i
vector<int> length;
                          // length[i] : the length of
    the longest string in the ith class
                          // cnt[i] : number of
vector<int> cnt;
    occurrences of each string in the ith class
                          // paths[i] : number of paths
vector<int> paths;
    on the automaton starting from i
                         // terminal[i] : true if i is
vector<bool> terminal:
    a terminal state
vector<int> first_pos;
vector<int> last_pos;
                          // the index of the
int last:
    equivalence class of the whole string
SuffixAutomaton(string s) {
   edges.push back(map<char.int>()):
   link.push back(-1):
   length.push_back(0);
   last = 0:
   rep(i, s.size()) { // construct r
       edges.push_back(map<char,int>());
       length.push_back(i+1);
       link.push_back(0);
       int r = edges.size() - 1:
       int p = last; // add edges to r and find p with
           link to a
       while(p >= 0 && !edges[p].count(s[i])) {
           edges[p][s[i]] = r;
          p = link[p];
       if(p != -1) {
           int q = edges[p][s[i]];
           if(length[p] + 1 == length[q]) {
              link[r] = q; // we do not have to split q,
                    just set the correct suffix link
          } else { // we have to split, add q'
              edges.push_back(edges[q]); // copy edges
              length.push_back(length[p] + 1);
              link.push_back(link[q]); // copy parent of
              int qq = edges.size()-1;
              link[q] = qq; // add qq as the new parent
                   of a and r
              link[r] = qq;
```

```
while(p >= 0 && edges[p][s[i]] == q) { //
                  move short classes polling to q to
                  poll to q'
                 edges[p][s[i]] = qq;
                 p = link[p];
          }
       last = r;
/* ----- Optional ----- */
   // mark terminal nodes
   terminal.assign(edges.size(), false);
   int p = last;
   while(p > 0) {
      terminal[p] = true;
       p = link[p]:
   }
   // precompute match count
   cnt.assign(edges.size(), -1);
   cnt_matches(0);
   // precompute number of paths (substrings) starting
        from state
   paths.assign(edges.size(), -1);
   cnt_paths(0);
   first_pos.assign(edges.size(), -1);
   get_first_pos(0);
   last_pos.assign(edges.size(), -1);
   get_last_pos(0);
int cnt matches(int state) {
   if(cnt[state] != -1) return cnt[state];
   int ans = terminal[state]:
   for(auto edge : edges[state])
       ans += cnt_matches(edge.second);
   return cnt[state] = ans:
int cnt_paths(int state) {
   if(paths[state] != -1) return paths[state];
   int ans = state == 0 ? 0 : 1; // without repetition ( );
        counts diferent substrings)
```

```
// int ans = state == 0 ? 0 : cnt[state]: // with
    repetition
   for(auto edge : edges[state])
       ans += cnt_paths(edge.second);
   return paths[state] = ans;
int get_first_pos(int state) {
   if(first_pos[state] != -1) return first_pos[state];
   int ans = 0;
   for(auto edge : edges[state])
       ans = max(ans, get first pos(edge.second)+1):
   return first_pos[state] = ans;
int get_last_pos(int state) {
   if(last_pos[state] != -1) return last_pos[state];
   int ans = terminal[state] ? 0 : INT_MAX;//fix
   for(auto edge : edges[state])
       ans = min(ans, get_last_pos(edge.second)+1);
   return last_pos[state] = ans;
string get_k_substring(int k) // 0-indexed
   string ans;
   int state = 0;
   while(true)
       int curr = state == 0 ? 0 : 1: // without
           repetition (counts different substrings)
   // int curr = state == 0 ? 0 : cnt[state]; // with
       repetition
       if(curr > k) return ans;
       k -= curr:
       for(auto edge : edges[state]) {
          if(paths[edge.second] <= k) {</pre>
              k -= paths[edge.second];
          } else {
              ans += edge.first:
              state = edge.second;
              break:
          }
```

#### 1.10 z-function

```
#include "../common.h"
// i-th element is equal to the greatest number of
    characters starting
// from the position i that coincide with the first
    characters of s
vector<int> z_function(string s) {
   int n = s.size():
   vector<int> z(n);
   int 1 = 0, r = 0:
   for(int i = 1: i < n: i++) {</pre>
       if(i < r) {
          z[i] = min(r - i, z[i - 1]);
       while(i + z[i] < n && s[z[i]] == s[i + z[i]]) {
          z[i]++:
       if(i + z[i] > r) {
          1 = i:
          r = i + z[i];
       }
   }
   return z;
```

# $2 ext{ dp}$

#### 2.1 convex-hull-trick

```
struct Line {
    mutable ll a, b, c;

    bool operator<(Line r) const { return a < r.a; }
    bool operator<(ll x) const { return c < x; }
};

// dynamically insert 'a*x + b' lines and query for maximum
    at any x

// all operations have complexity O(log N)
struct LineContainer : multiset<Line, less<>> {

    ll div(ll a, ll b) {
        return a / b - ((a ^ b) < 0 && a % b);
    }
}</pre>
```

```
bool isect(iterator x, iterator v) {
       if (y == end()) return x \rightarrow c = INF. 0:
       if (x->a == y->a) x->c = x->b > y->b ? INF : -INF;
       else x\rightarrow c = div(y\rightarrow b - x\rightarrow b, x\rightarrow a - y\rightarrow a);
       return x->c >= y->c;
   void add(ll a, ll b) {
       // a *= -1, b *= -1 // for min
       auto z = insert(\{a, b, 0\}), y = z++, x = y;
       while (isect(y, z)) z = erase(z);
       if (x != begin() && isect(--x, v)) isect(x, v = erase
       while ((y = x) != begin() \&\& (--x)->c >= y->c) isect(
            x. erase(v)):
   11 query(11 x) {
       if (empty()) return -INF; // INF for min
       auto 1 = *lower bound(x):
       return 1.a * x + 1.b;
       // return -l.a * x - l.b: // for min
};
```

# 2.2 divide-and-conquer

```
// for every index i assign an optimal index j, such that
    cost(i, i) is
// minimal for every i. the property that if i2 >= i1 then
    j2 >= j1 is
// exploited (monotonic condition).
// calculate optimal index for all indices in range [1, r)
     knowing that
// the optimal index for every index in this range is within
      [optl, optr).
// time: O(N log N)
void calc(vector<int> &opt, int 1, int r, int opt1, int optr
    ) {
   if (1 == r) return;
   int i = (1 + r) / 2:
   11 optc = INF;
   int optj;
   repx(j, optl, optr) {
       11 c = i + i; // cost(i, i)
       if (c < optc) optc = c, optj = j;</pre>
   opt[i] = optj;
   calc(opt, 1, i, optl, optj + 1);
```

```
calc(opt, i + 1, r, optj, optr);
}
```

# 3 geo2d

### 3.1 circle

```
struct C {
   Po: Tr:
   C(P \ o, T \ r) : o(o), r(r) \{\}
   C() : C(P(), T()) \{ \}
   // intersects the circle with a line, assuming they
        intersect
   // results are sorted with respect to the direction of
        the line
   pair<P, P> line_inter(L 1) const {
      P c = 1.closest_to(o);
      T c2 = (c - o).magsq():
      P = sqrt(max(r * r - c2, T())) * 1.d.unit();
      return {c - e, c + e}:
   // checks whether the given line collides with the circle
   // negative: 2 intersections
   // zero: 1 intersection
   // positive: 0 intersections
   T line_collide(L 1) const {
      T c2 = (1.closest_to(o) - o).magsq();
      return c2 - r * r:
   }
   // calculates the two intersections between two circles
   // the circles must intersect in one or two points!
   pair<P. P> inter(C h) const {
      P d = h.o - o;
      T c = (r * r - h.r * h.r) / d.magsq();
      return h.line_inter({(1 + c) / 2 * d, d.rot()});
   }
   // check if the given circles intersect
   bool collide(C h) const {
      return (h.o - o).magsq() \le (h.r + r) * (h.r + r):
   // get one of the two tangents that cross through the
```

```
// the point must not be inside the circle
// a = -1: cw (relative to the circle) tangent
// a = 1: ccw (relative to the circle) tangent
P point_tangent(P p, T a) const {
   T c = r * r / p.magsq();
   return o + c * (p - o) - a * sqrt(c * (1 - c)) * (p -
         o).rot():
}
// get one of the 4 tangents between the two circles
// a = 1: exterior tangents
// a = -1: interior tangents (requires no area overlap)
// b = 1: ccw tangent
// b = -1: cw tangent
// the line origin is on this circumference, and the
     direction
// is a unit vector towards the other circle
L tangent(C c, T a, T b) const {
   T dr = a * r - c.r:
   P d = c.o - o:
   P n = (d * dr + b * d.rot() * sqrt(d.magsq() - dr *
        dr)).unit();
   return {o + n * r, -b * n.rot()};
// find the circumcircle of the given **non-degenerate**
     triangle
static C thru_points(P a, P b, P c) {
   L 1((a + b) / 2, (b - a).rot());
   P p = 1.intersection(L((a + c) / 2, (c - a).rot())):
   return {p, (p - a).mag()};
// find the two circles that go through the given point,
     are tangent
// to the given line and have radius 'r'
// the point-line distance must be at most 'r'!
// the circles are sorted in the direction of the line
static pair<C, C> thru_point_line_r(P a, L t, T r) {
   P d = t.d.rot().unit():
   if (d * (a - t.o) < 0) d = -d:
   auto p = C(a, r).line_inter(\{t.o + d * r, t.d\});
   return {{p.first, r}, {p.second, r}};
// find the two circles that go through the given points
     and have
// radius 'r'
// the circles are sorted by angle with respect to the
     first point
```

```
// the points must be at most at distance 'r'!
   static pair<C, C> thru_points_r(P a, P b, T r) {
       return {{p.first, r}, {p.second, r}};
};
```

#### convex-hull

```
// get the convex hull with the least amount of vertices for
      the given set
// of points
// probably misbehaves if points are not all distinct!
vector<P> convex_hull(vector<P> &ps) {
   int N = ps.size(), n = 0, k = 0;
   if (N <= 2) return ps;</pre>
   rep(i, N) if (make_pair(ps[i].y, ps[i].x) < make_pair(ps[</pre>
        k].y, ps[k].x)) k = i;
   swap(ps[k], ps[0]):
   sort(++ps.begin(), ps.end(), [&](P 1, P r) {
       T x = (r - 1) \% (ps[0] - 1), d = (r - 1) * (ps[0] - 1)
       return x > 0 \mid | x == 0 && d < 0;
   }):
   vector<P> H:
   for (P p : ps) {
       while (n \ge 2 \&\& (H[n - 1] - p) \% (H[n - 2] - p) >=
            0) H.pop_back(), n--;
       H.push back(p), n++:
   return H;
```

### delaunay

```
typedef __int128_t lll; // if on a 64-bit platform
struct O {
    Q *rot, *o; P p = {INF, INF}; bool mark;
    P &F() { return r()->p; }
    Q *&r() { return rot->rot: }
    Q *prev() { return rot->o->rot; }
    Q *next() { return r()->prev(); }
};
```

```
T cross(P a, P b, P c) { return (b - a) % (c - a): }
auto p = C(a, r).line_inter({(a + b) / 2, (b - a).rot | bool circ(P p, P a, P b, P c) { // is p in the circumcircle?
                                                                 111 p2 = p.magsq(), A = a.magsq() - p2,
                                                                     B = b.magsq() - p2, C = c.magsq() - p2;
                                                                 return cross(p, a, b) * C + cross(p, b, c) * A + cross(p, b, c)
                                                                        c, a) * B > 0;
                                                             Q *makeEdge(Q *&H, P orig, P dest) {
                                                                 Q *r = H ? H : new Q{new Q{new Q{new Q{0}}}}:
                                                                 H = r -> 0: r -> r() -> r() = r:
                                                                 repx(i, 0, 4) r = r->rot, r->p = {INF, INF},
                                                                     r->0 = i & 1 ? r : r->r();
                                                                 r\rightarrow p = orig; r\rightarrow F() = dest;
                                                                 return r;
                                                              void splice(0 *a, 0 *b) {
                                                                  swap(a->o->rot->o, b->o->rot->o); swap(a->o, b->o);
                                                             Q *connect(Q *&H, Q *a, Q *b) {
                                                                  Q *q = makeEdge(H, a->F(), b->p);
                                                                  splice(q, a->next()); splice(q->r(), b); return q;
                                                             pair<0 *. 0 *> rec(0 *&H. const vector<P> &s) {
                                                                 if (s.size() <= 3) {</pre>
                                                                     Q *a = makeEdge(H, s[0], s[1]), *b = makeEdge(H, s[0])
                                                                           [1], s.back());
                                                                     if (s.size() == 2) return \{a, a\rightarrow r()\}; splice(a\rightarrow r(), a\rightarrow r());
                                                                     auto side = cross(s[0], s[1], s[2]);
                                                                     Q *c = side ? connect(H, b, a) : 0;
                                                                     return \{\text{side} < 0 ? c \rightarrow r() : a. \text{side} < 0 ? c : b \rightarrow r()\}
                                                                 }
                                                              #define J(e) e \rightarrow F(), e \rightarrow p
                                                              #define valid(e) (cross(e->F(), J(base)) > 0)
                                                                  Q *A, *B, *ra, *rb; int half = s.size() / 2;
                                                                 tie(ra, A) = rec(H, {s.begin(), s.end() - half});
                                                                 tie(B, rb) = rec(H, {s.begin() + s.size() - half, s.end()
                                                                      }):
                                                                  while ((cross(B\rightarrow p, J(A)) < 0 \&\& (A = A\rightarrow next())) | |
                                                                        (cross(A->p, J(B)) > 0 && (B = B->r()->o)));
                                                                 Q *base = connect(H, B->r(), A);
                                                                 if (A->p == ra->p) ra = base->r():
                                                                  if (B->p == rb->p) rb = base:
```

```
#define DEL(e. init. dir) Q *e = init->dir: \
   if (valid(e)) while (circ(e->dir->F(), J(base), e->F()))
           Q *t = e->dir; splice(e, e->prev()); \
          splice(e->r(), e->r()->prev()); e->o = H; H = e;
       }
   for (;;) {
       DEL(LC, base->r(), o); DEL(RC, base, prev());
       if (!valid(LC) && !valid(RC)) break;
       if (!valid(LC) || (valid(RC) && circ(J(RC), J(LC))))
            base = connect(H, RC, base->r());
       else base = connect(H, base->r(), LC->r());
   return {ra, rb};
#undef J
#undef valid
#undef DEL
// there must be no duplicate points
// returns no triangles in the case of all collinear points
// produces counter-clockwise triangles ordered in triples
// maximizes the minimum angle across all triangulations
// the euclidean mst is a subset of these edges
// O(N log N)
vector<P> triangulate(vector<P> pts) {
   sort(pts.begin(), pts.end(), [](P a, P b) {
       return make_pair(a.x, a.y) < make_pair(b.x, b.y);</pre>
   assert(unique(pts.begin(), pts.end()) == pts.end());
   if (pts.size() < 2) return {};</pre>
   Q *H = 0; Q *e = rec(H, pts).first;
   vector < Q *> q = \{e\}; int qi = 0;
   while (cross(e->o->F(), e->F(), e->p) < 0) e = e->o:
#define ADD
   {
       Q *c = e;
       do {
           c->mark = 1; pts.push back(c->p); \
          q.push_back(c->r()); c = c->next(); \
       } while (c != e):
   ADD:
   while (qi < (int)q.size()) if (!(e = q[qi++])->mark) ADD;
   return pts:
#undef ADD
```

### 3.4 halfplane-intersect

```
// obtain the convex polygon that results from intersecting
     the given list
// of halfplanes, represented as lines that allow their left
      side
// assumes the halfplane intersection is bounded
vector<P> halfplane intersect(vector<L> &H) {
   L bb(P(-INF, -INF), P(INF, 0));
   rep(k, 4) H.push_back(bb), bb.o = bb.o.rot(), bb.d = bb.d
        .rot():
   sort(begin(H), end(H), [](L a, L b) { return a.d.angcmp(b
        .d) < 0: }:
   deque<L> q; int n = 0;
   rep(i, H.size()) {
       while (n \ge 2 \&\& H[i].side(q[n - 1].intersection(q[n - 1]))
            -21)) > 0)
           g.pop back(), n--:
       while (n >= 2 && H[i].side(q[0].intersection(q[1])) >
           q.pop_front(), n--;
       if (n > 0 \&\& H[i].parallel(q[n - 1])) {
           if (H[i].d * q[n - 1].d < 0) return {};</pre>
           if (H[i].side(q[n - 1].o) > 0) q.pop_back(), n--;
           else continue:
       }
       q.push_back(H[i]), n++;
   while (n \ge 3 \&\& q[0].side(q[n - 1].intersection(q[n -
        21)) > 0)
       q.pop_back(), n--;
   while (n \ge 3 \&\& q[n - 1].side(q[0].intersection(q[1])) >
       q.pop_front(), n--;
   if (n < 3) return {}:</pre>
   vector<P> ps(n);
   rep(i, n) ps[i] = q[i].intersection(q[(i + 1) % n]);
   return ps;
```

### 3.5 line

```
// a segment or an infinite line
// does not handle point segments correctly!
struct L {
   P o, d;
```

```
L() : o(), d()  {}
L(P \ o, P \ d) : o(o), d(d) \{ \}
L(P ab, T c) : d(ab.rot()), o(ab * -c / ab.magsq()) {}
pair<P, T> line_eq() { return {-d.rot(), d.rot() * o}; }
// returns a number indicating which side of the line the
     point is in
// negative: left, positive: right
T side(P r) const { return (r - o) % d; }
// returns the intersection coefficient
// in the range [0, d % r.d]
// if d % r.d is zero, the lines are parallel
T inter(L r) const { return (r.o - o) % r.d: }
// get the single intersection point
// lines must not be parallel
P intersection(L r) const { return o + d * inter(r) / (d
// check if lines are parallel
bool parallel(L r) const { return abs(d % r.d) <= EPS; }</pre>
// check if segments intersect
bool seg_collide(L r) const {
   Tz = d \% r.d:
   if (abs(z) <= EPS) {
       if (abs(side(r.o)) > EPS) return false;
       T s = (r.o - o) * d. e = s + r.d * d:
       if (s > e) swap(s, e);
       return s <= d * d + EPS && e >= -EPS;
   T s = inter(r), t = -r.inter(*this);
   if (z < 0) s = -s, t = -t, z = -z:
   return s >= -EPS && s <= z + EPS && t >= -EPS && t <=
         z + EPS:
}
// full segment intersection
// produces a point segment if the intersection is a
// however it **does not** handle point segments as input
bool seg_inter(L r, L *out) const {
   T z = d \% r.d;
   if (abs(z) <= EPS) {
       if (abs(side(r.o)) > EPS) return false;
       if (r.d * d < 0) r = \{r.o + r.d, -r.d\}:
       Ps = o * d < r.o * d ? r.o : o:
```

```
P = (o + d) * d < (r.o + r.d) * d ? o + d : r.o
          if (s * d > e * d) return false:
          return *out = {s. e - s}. true:
      }
      T s = inter(r), t = -r.inter(*this):
       if (z < 0) s = -s, t = -t, z = -z;
       if (s \ge -EPS \&\& s \le z + EPS \&\& t \ge -EPS \&\& t \le z
          return *out = \{0 + d * s / z, \{0, 0\}\}, true;
       return false:
   // check if the given point is on the segment
   bool point_on_seg(P r) const {
       if (abs(side(r)) > EPS) return false;
       if ((r - o) * d < -EPS) return false:
       if ((r - o - d) * d > EPS) return false;
       return true:
   }
   // get the point in this line that is closest to a given
        point
   P closest to(P r) const {
       P dr = d.rot(); return r + (o - r) * dr * dr / d.
            magsq();
}:
```

### 3.6 minkowski

```
void reorder_polygon(vector<P> &ps) {
   int pos = 0;
   repx(i, 1, (int)ps.size()) {
      if (ps[i].y < ps[pos].y || (ps[i].y == ps[pos].y &&
            ps[i].x < ps[pos].x))
      pos = i;
   }
   rotate(ps.begin(), ps.begin() + pos, ps.end());
}

vector<P> minkowski(vector<P> ps, vector<P> qs) {
      // the first vertex must be the lowest
      reorder_polygon(ps); reorder_polygon(qs);
      ps.push_back(ps[0]); ps.push_back(ps[1]);
      qs.push_back(qs[0]); qs.push_back(qs[1]);
      vector<P> result; int i = 0, j = 0;
      while (i < ps.size() - 2 || j < qs.size() - 2) {
        result.push_back(ps[i] + qs[j]);
      auto z = (ps[i + 1] - ps[i]) % (qs[j + 1] - qs[j]);
}</pre>
```

```
if (z \ge 0 \&\& i < ps.size() - 2) ++i:
   if (z \le 0 \&\& j \le gs.size() - 2) ++j;
return result:
```

### 3.7 point

```
struct P {
    T x. v:
    P(T x, T y) : x(x), y(y) {}
    P() : P(0, 0) {}
    friend ostream &operator<<(ostream &s, const P &r) {</pre>
       return s << r.x << " " << r.v:
    friend istream &operator>>(istream &s, P &r) { return s
        >> r.x >> r.v: }
    P operator+(P r) const { return \{x + r.x, y + r.y\}; }
    P operator-(P r) const { return {x - r.x, y - r.y}; }
    P operator*(T r) const { return \{x * r, y * r\}; \}
    P operator/(T r) const { return {x / r, y / r}; }
    P operator-() const { return {-x, -v}; }
    friend P operator*(T 1, P r) { return {1 * r.x, 1 * r.y};
         }
    P rot() const { return {-v, x}; }
    T operator*(P r) const { return x * r.x + y * r.y; }
    T operator%(P r) const { return rot() * r; }
    T magsq() const { return x * x + v * v; }
    T mag() const { return sqrt(magsq()); }
    P unit() const { return *this / mag(): }
    bool half() const { return abs(y) <= EPS && x < -EPS | | y | // check if a point is in a convex polygon
         < -EPS: }
    T angcmp(P r) const {
       int h = (int)half() - r.half();
       return h ? h : r % *this:
    bool operator==(P r) const { return abs(x - r.x) <= EPS</pre>
         && abs(v - r.v) <= EPS; }
    double angle() const { return atan2(v, x); }
    static P from_angle(double a) { return {cos(a), sin(a)};
};
```

### 3.8 polygon

```
// get the area of a simple polygon in ccw order
T area(const vector<P> &ps) {
   int N = ps.size():
   T a = 0:
   rep(i, N) a += (ps[i] - ps[0]) % (ps[(i + 1) % N] - ps[i]
   return a / 2;
// checks whether a point is inside a simple polygon
// returns -1 if inside, 0 if on border, 1 if outside
int in_poly(const vector<P> &ps, P p) {
   int N = ps.size(), w = 0:
   rep(i, N) {
      P = ps[i] - p, e = ps[(i + 1) % N] - p;
       if (s == P()) return 0:
       if (s.v == 0 \&\& e.v == 0) {
           if (\min(s.x. e.x) \le 0 \&\& 0 \le \max(s.x. e.x))
               return 0:
       } else {
           bool b = s.v < 0:
           if (b != (e.v < 0)) {
              T z = s \% e; if (z == 0) return 0;
              if (b == (z > 0)) w += b ? 1 : -1:
   }
   return w ? -1 : 1;
struct InConvex {
   vector<P> ps;
   T 11. 1h. rl. rh:
   int N. m:
   // preprocess polygon
   // O(N)
   InConvex(const vector<P> &p) : ps(p), N(ps.size()), m(0)
       assert(N >= 2);
       rep(i, N) if (ps[i].x < ps[m].x) m = i;
       rotate(ps.begin(), ps.begin() + m, ps.end());
       rep(i, N) if (ps[i].x > ps[m].x) m = i;
```

```
11 = 1h = ps[0].v. rl = rh = ps[m].v:
       for (P p : ps) {
           if (p.x == ps[0].x) 11 = min(11, p.y), 1h = max(
           if (p.x == ps[m].x) rl = min(rl, p.y), rh = max(
               rh. p.v):
      }
   InConvex() {}
   // check if point belongs in polygon
   // returns -1 if inside. 0 if on border. 1 if outside
   // O(log N)
   int in_poly(P p) {
       if (p.x < ps[0].x \mid\mid p.x > ps[m].x) return 1;
       if (p.x == ps[0].x) return p.y < 11 || p.y > 1h;
       if (p.x == ps[m].x) return p.y < rl || p.y > rh;
       int r = upper_bound(ps.begin(), ps.begin() + m, p,
           \lceil \rceil (P a, P b) \mid return a.x < b.x: \} - ps.begin():
       Tz = (ps[r - 1] - ps[r]) \% (p - ps[r]); if (z >= 0)
            return !!z;
       r = upper_bound(ps.begin() + m, ps.end(), p,
           [](P a, P b) { return a.x > b.x; }) - ps.begin();
       z = (ps[r - 1] - ps[r \% N]) \% (p - ps[r \% N]);
       if (z \ge 0) return !!z: return -1:
   }
};
// classify collision of a ray inside a ccw polygon vertex.
// ray is (o, d), vertex is b, previous vertex is a, next is
pair<bool, bool> inner_collide(P o, P d, P a, P b, P c) {
   T p = (a - o) \% d: // side of previous
   T n = (c - o) \% d;
                        // side of next
   T v = (c - b) \% (b - a); // is vertex convex?
   return \{v > 0 ? n < 0 | | (n == 0 && p < 0) : p > 0 | | n 
          v > 0 ? p > 0 || (p == 0 \&\& n > 0) : p > 0 || n <
                0};
```

#### 3.9 sweep

```
#include "point.cpp"
// iterate over all pairs of points
// 'op' is called with all ordered pairs of different
    indices '(i, i)'
```

```
// additionally, the 'ps' vector is kept sorted by signed
// to the line formed by 'i' and 'j'
// for example, if the vector from 'i' to 'j' is pointing
// the 'ps' vector is sorted from smallest 'v' to largest 'v
// note that, because the 'ps' vector is sorted by signed
     distance.
// 'j' is always equal to 'i + 1'
// this means that the amount of points to the left of the
    line is always 'N - i'
template <class OP>
void all_pair_points(vector<P> &ps, OP op) {
   int N = ps.size();
   sort(ps.begin(), ps.end(), [](P a, P b) {
       return make_pair(a.y, a.x) < make_pair(b.y, b.x);</pre>
   });
   vector<pair<int, int>> ss:
   rep(i, N) rep(j, N) if (i != j) ss.push_back({i, j});
   stable_sort(ss.begin(), ss.end(), [&](auto a, auto b) {
       return (ps[a.second] - ps[a.first]).angle_lt(ps[b.
            second] - ps[b.first]);
   vector<int> p(N); rep(i, N) p[i] = i;
   for (auto [i, j] : ss)
       { op(p[i], p[j]); swap(ps[p[i]], ps[p[j]]); swap(p[i
           ], p[i]); }
```

#### 3.10 theorems

```
// Pick's theorem
// Simple polygon with integer vertices:
// A = I + B / 2 - 1
// A: Area of the polygon
// I: Integer points strictly inside the polygon
// B: Integer points on the boundary of the polygon
```

# 4 graph

### 4.1 bellman-ford

```
struct Edge { int u, v; ll w; };
// find distance from source node to all nodes.
```

```
// supports negative edge weights.
// returns true if a negative cycle is detected.
11
// time: 0(V E)
bool bellman_ford(int N, int s, vector<Edge> &E, vector<11>
    &D. vector<int> &P) {
   P.assign(N, -1), D.assign(N, INF), D[s] = 0;
   rep(i, N - 1) {
       bool f = true:
       rep(ei, E.size()) {
          auto &e = E[ei]:
          ll n = D[e.u] + e.w:
          if (D[e.u] < INF && n < D[e.v])
              D[e.v] = n, P[e.v] = ei, f = false;
       if (f) return false;
   return true;
```

#### 4.2 dinic

```
struct Edge { int u, v; ll c, f = 0; };
// maximum flow algorithm.
// time: 0(E V^2)
//
       O(E \ V^{(2/3)}) / O(E \ sqrt(E)) unit capacities
11
                                    unit networks (hopcroft-
// unit network: c in {0, 1} and forall v, len(incoming(v))
    <= 1 or len(outgoing(v)) <= 1
// min-cut: find all nodes reachable from the source in the
    residual graph
struct Dinic {
   int N, s, t; vector<vector<int>> G;
   vector<Edge> E; vector<int> lvl, ptr;
   Dinic(int N, int s, int t): N(N), s(s), t(t), G(N) {}
   void add_edge(int u, int v, ll c) {
       G[u].push_back(E.size()); E.push_back({u, v, c});
       G[v].push_back(E.size()); E.push_back({v, u, 0});
   }
   11 push(int u, 11 p) {
      if (u == t || p <= 0) return p;</pre>
       while (ptr[u] < G[u].size()) {</pre>
          int ei = G[u][ptr[u]++];
           Edge &e = E[ei];
```

```
if (lvl[e.v] != lvl[u] + 1) continue;
          ll a = push(e.v, min(e.c - e.f, p));
          if (a <= 0) continue; e.f += a, E[ei ^ 1].f -= a;</pre>
          return a:
      }
       return 0:
   }
   ll maxflow() {
      11 f = 0:
       while (true) {
          // bfs to build levels
          lvl.assign(N, -1); queue < int > q; lvl[s] = 0, q.
               push(s);
           while (!q.empty()) {
              int u = q.front(); q.pop();
              for (int ei : G[u]) {
                  Edge &e = E[ei];
                  if (e.c - e.f <= 0 || lvl[e.v] != -1)
                  lvl[e.v] = lvl[u] + 1, q.push(e.v);
          }
           if (lvl[t] == -1) break;
          // dfs to find blocking flow
          ptr.assign(N, 0); while (11 ff = push(s, INF)) f
               += ff:
      }
       return f:
   }
};
```

# 4.3 floyd-warshall

### 4.4 heavy-light

```
struct Hld {
   vector<int> P, H, D, pos, top;
   H14() {}
   void init(vector<vector<int>> &G) {
      int N = G.size():
      P.resize(N), H.resize(N), D.resize(N), pos.resize(N),
          top.resize(N);
      D[0] = -1, dfs(G, 0); int t = 0;
      rep(i, N) if (H[P[i]] != i) {
          int j = i;
          while (j != -1)
              \{ top[i] = i, pos[i] = t++; j = H[i]; \}
   }
   int dfs(vector<vector<int>> &G. int i) {
      int w = 1, mw = 0;
      D[i] = D[P[i]] + 1, H[i] = -1;
      for (int c : G[i]) {
          if (c == P[i]) continue;
          P[c] = i; int sw = dfs(G, c); w += sw;
          if (sw > mw) H[i] = c. mw = sw:
      }
      return w:
   // visit the log N segments in the path from u to v
   template <class OP>
   void path(int u, int v, OP op) {
      while (top[u] != top[v]) {
          if (D[top[u]] > D[top[v]]) swap(u, v);
          op(pos[top[v]], pos[v] + 1); v = P[top[v]];
      if (D[u] > D[v]) swap(u, v);
       op(pos[u], pos[v] + 1); // value on node
      // op(pos[u]+1, pos[v] + 1); // value on edge
   // an alternative to 'path' that considers order.
   // calls 'op' with an 'l <= r' inclusive-exclusive range,</pre>
         and a
   // boolean indicating if the query is forwards or
        backwards.
   template <class OP>
   void path(int u, int v, OP op) {
      int lu = u. lv = v:
      while (top[lu] != top[lv])
```

```
if (D[top[lu]] > D[top[lv]]) lu = P[top[lu]];
           else lv = P[top[lv]];
       int lca = D[lu] > D[lv] ? lv : lu;
       while (top[u] != top[lca])
           op(pos[top[u]], pos[u] + 1, false), u = P[top[u]]
       if (u != lca) op(pos[lca] + 1, pos[u] + 1, false);
       vector<int> stk:
       while (top[v] != top[lca])
           stk.push back(v), v = P[top[v]]:
       // op(pos[lca], pos[v] + 1, true); // value on node
       op(pos[lca] + 1, pos[v] + 1, true); // value on edge
       reverse(stk.begin(), stk.end());
       for (int w : stk) op(pos[top[w]], pos[w] + 1, true);
   }
   // commutative segment tree
   template <class T, class S>
   void update(S &seg, int i, T val) { seg.update(pos[i],
        val): }
   // commutative segment tree lazy
   template <class T, class S>
   void update(S &seg, int u, int v, T val) {
       path(u, v, [&](int 1, int r) { seg.update(1, r, val);
             });
   }
   // commutative (lazy) segment tree
   template <class T. class S>
   T query(S &seg, int u, int v) {
      T ans = 0:
            // neutral element
       path(u, v, [\&](int 1, int r) \{ ans += seg.query(1, r) \}
           ; }); // query op
       return ans;
}:
```

### 4.5 hungarian

```
// find a maximum gain perfect matching in the given
    bipartite complete graph.
// input: gain matrix (G_{xy} = benefit of joining vertex x
    in set X with vertex
// y in set Y).
```

```
// output: maximum gain matching in members 'xv[x]' and 'vx[
    vΑ.
// runtime: O(N^3)
struct Hungarian {
   int N, qi, root;
   vector<vector<ll>>> gain;
   vector<int> xy, yx, p, q, slackx;
   vector<ll> lx, ly, slack;
   vector<bool> S. T:
   void add(int x, int px) {
       S[x] = true, p[x] = px:
       rep(y, N) if (lx[x] + ly[y] - gain[x][y] < slack[y])
           slack[v] = lx[x] + ly[v] - gain[x][v], slackx[v]
      }
   }
   void augment(int x, int y) {
       while (x != -2) {
          yx[y] = x; swap(xy[x], y); x = p[x];
   }
   void improve() {
       S.assign(N, false), T.assign(N, false), p.assign(N,
           -1):
       qi = 0, q.clear();
       rep(x, N) if (xv[x] == -1) {
          q.push_back(root = x), p[x] = -2, S[x] = true;
       rep(y, N) slack[y] = lx[root] + ly[y] - gain[root][y
           1. slackx[v] = root:
       while (true) {
           while (qi < q.size()) {</pre>
              int x = q[qi++];
              rep(y, N) if (lx[x] + ly[y] == gain[x][y] &&!
                   T[v]) {
                  if (vx[v] == -1) return augment(x, v);
                  T[y] = true, q.push_back(yx[y]), add(yx[y
                      1. x):
              }
          }
          11 d = INF:
          rep(y, N) if (!T[y]) d = min(d, slack[y]);
          rep(x, N) if (S[x]) lx[x] -= d:
```

### 4.6 kuhn

```
// get a maximum cardinality matching in a bipartite graph.
// input: adjacency lists.
// output: matching (in 'mt' member).
// runtime: O(V E)
struct Kuhn {
   int N. size:
   vector<vector<int>> G;
   vector<bool> seen:
   vector<int> mt:
   bool visit(int i) {
       if (seen[i]) return false;
       seen[i] = true;
       for (int to : G[i])
          if (mt[to] == -1 || visit(mt[to])) {
              mt[to] = i:
              return true:
          }
       return false;
   Kuhn(vector<vector<int>> adj) : G(adj), N(G.size()), mt(N
        , -1) {
       rep(i, N) {
          seen.assign(N, false);
           size += visit(i);
```

```
}
};
```

# 4.7 lca

```
// calculates the lowest common ancestor for any two nodes
    in O(log N) time,
// with O(N log N) preprocessing
struct Lca {
   int N, K, t = 0;
   vector<vector<int>> U:
   vector<int> L. R:
   Lca() {}
   Lca(vector<vector<int>> &G) : N(G.size()), L(N), R(N) {
       K = N \le 1 ? 0 : 32 - \_builtin_clz(N - 1);
      U.resize(K + 1, vector<int>(N));
       visit(G, 0, 0);
       rep(k, K) rep(u, N) U[k + 1][u] = U[k][U[k][u]]:
   }
   void visit(vector<vector<int>> &G, int u, int p) {
       L[u] = t++, U[0][u] = p;
       for (int v : G[u]) if (v != p) visit(G, v, u);
       R[u] = t++:
   }
   bool is_anc(int up, int dn) {
       return L[up] <= L[dn] && R[dn] <= R[up];
   int find(int u, int v) {
       if (is anc(u, v)) return u:
       if (is_anc(v, u)) return v;
       for (int k = K: k >= 0:)
          if (is_anc(U[k][u], v)) k--;
           else u = U[k][u];
       return U[0][u];
   }
};
```

# 4.8 maxflow-mincost

```
// untested
#include "../common.h"
```

```
const 11 INF = 1e18:
struct Edge {
   int u, v;
   11 c. w. f = 0:
// find the minimum-cost flow among all maximum-flow flows.
// time: 0(F V E)
                          F is the maximum flow
     O(V E + F E log V) if bellman-ford is replaced by
    johnson
struct Flow {
   int N. s. t:
   vector<vector<int>> G;
   vector<Edge> E:
   vector<ll> d;
   vector<int> p:
   Flow() {}
   Flow(int N, int s, int t) : N(N), s(s), t(t), G(N) {}
   void add_edge(int u, int v, ll c, ll w) {
       G[u].push_back(E.size());
       E.push_back({u, v, c, w});
       G[v].push_back(E.size());
       E.push back({v. u. 0. -w}):
   }
   void calcdists() {
       // replace bellman-ford with johnson for better time
       d.assign(N, INF);
       p.assign(N, -1);
       d[s] = 0:
       rep(i, N - 1) rep(ei, E.size()) {
          Edge &e = E[ei]:
          ll n = d[e.u] + e.w:
          if (d[e.u] < INF && e.c - e.f > 0 && n < d[e.v])
               d[e.v] = n, p[e.v] = ei;
      }
   }
   11 maxflow() {
       11 \text{ ff} = 0:
       while (true) {
          calcdists();
          if (p[t] == -1) break;
          11 f = INF:
```

```
int cur = t;
while (p[cur] != -1) {
    Edge &e = E[p[cur]];
    f = min(f, e.c - e.f);
    cur = e.u;
}

int cur = t;
while (p[cur] != -1) {
    E[p[cur]].f += f;
    E[p[cur] ^ 1].f -= f;
}

ff += f;
}
return ff;
}
```

### 4.9 push-relabel

```
#include "../common.h"
const ll INF = 1e18:
// maximum flow algorithm.
// to run, use 'maxflow()'.
// time: O(V^2 \operatorname{sqrt}(E)) \leq O(V^3)
// memory: 0(V^2)
struct PushRelabel {
   vector<vector<ll>> cap. flow:
   vector<11> excess;
   vector<int> height;
   PushRelabel() {}
   void resize(int N) { cap.assign(N, vector<11>(N)); }
   // push as much excess flow as possible from u to v.
   void push(int u, int v) {
       11 f = min(excess[u], cap[u][v] - flow[u][v]);
       flow[u][v] += f:
       flow[v][u] -= f;
       excess[v] += f;
       excess[u] -= f:
   // relabel the height of a vertex so that excess flow may
         be pushed.
```

```
void relabel(int u) {
   int d = INT32 MAX:
   rep(v, cap.size()) if (cap[u][v] - flow[u][v] > 0) d
       min(d, height[v]);
   if (d < INF) height[u] = d + 1:</pre>
// get the maximum flow on the network specified by 'cap'
     with source 's'
// and sink 't'.
// node-to-node flows are output to the 'flow' member.
11 maxflow(int s, int t) {
   int N = cap.size(), M;
   flow.assign(N, vector<ll>(N));
   height.assign(N, 0), height[s] = N;
   excess.assign(N, 0), excess[s] = INF;
   rep(i, N) if (i != s) push(s, i);
   vector<int> a:
   while (true) {
       // find the highest vertices with excess
       q.clear(), M = 0;
       rep(i, N) {
           if (excess[i] <= 0 || i == s || i == t)</pre>
               continue:
          if (height[i] > M) q.clear(), M = height[i];
           if (height[i] >= M) q.push back(i):
       if (q.empty()) break;
       // process vertices
       for (int u : q) {
          bool relab = true:
          rep(v, N) {
              if (excess[u] <= 0) break;</pre>
              if (cap[u][v] - flow[u][v] > 0 && height[u]
                   ] > height[v])
                  push(u, v), relab = false;
          if (relab) {
              relabel(u):
              break;
   11 f = 0; rep(i, N) f += flow[i][t]; return f;
```

### 4.10 strongly-connected-components

```
// compute strongly connected components.
// time: O(V + E), memory: O(V)
//
// after building:
// comp = map from vertex to component (components are
    toposorted, root first, leaf last)
// N = number of components
// G = condensation graph (component DAG)
11
// byproducts:
// vgi = transposed graph
// order = reverse topological sort (leaf first, root last)
// others:
// vn = number of vertices
// vg = original vertex graph
struct Scc {
   int vn. N:
   vector<int> order, comp;
   vector<vector<int>> vg, vgi, G;
   void toposort(int u) {
       if (comp[u]) return:
       comp[u] = -1;
       for (int v : vg[u]) toposort(v);
       order.push back(u):
   bool carve(int u) {
       if (comp[u] != -1) return false;
       comp[u] = N;
       for (int v : vgi[u]) {
           carve(v):
           if (comp[v] != N) G[comp[v]].push_back(N);
       return true:
   Scc(vector<vector<int>>> &g) : vn(g.size()), vg(g), comp(
        vn), vgi(vn), G(vn), N(0) {
       rep(u, vn) toposort(u);
       rep(u, vn) for (int v : vg[u]) vgi[v].push_back(u);
       invrep(i, vn) N += carve(order[i]);
};
```

#### 4.11 two-sat

```
// calculate the solvability of a system of logical
    equations, where every equation is of the form 'a or b
// 'neg': get negation of 'u'
// 'then': 'u' implies 'v'
// 'anv': 'u' or 'v'
// 'set': 'u' is true
// after 'solve' (O(V+E)) returns true. 'sol' contains one
    possible solution.
// determining all solutions is O(V*E) hard (requires
    computing reachability in a DAG).
struct TwoSat {
   int N; vector<vector<int>> G;
   Scc scc; vector<bool> sol;
   TwoSat(int n) : N(n), G(2 * n), sol(n) {}
   TwoSat() {}
   int neg(int u) { return (u + N) % (2 * N); }
   void then(int u, int v) { G[u].push_back(v), G[neg(v)].
        push_back(neg(u)); }
   void any(int u, int v) { then(neg(u), v); }
   void set(int u) { G[neg(u)].push_back(u); }
   bool solve() {
       scc = Scc(G);
       rep(u, N) if (scc.comp[u] == scc.comp[neg(u)]) return
       rep(u, N) sol[u] = (scc.comp[u] > scc.comp[neg(u)]);
       return true:
};
```

# 5 implementation

# 5.1 SegmentTreeBeats

```
if (a.mx1 > b.mx1) mx1 = a.mx1, mxc = a.mxc, mx2 =
            max(b.mx1, a.mx2):
       if (a.mx1 < b.mx1) mx1 = b.mx1, mxc = b.mxc, mx2 =
            max(a.mx1, b.mx2);
       if (a.mx1 == b.mx1) mx1 = a.mx1, mxc = a.mxc + b.mxc.
            mx2 = max(a.mx2, b.mx2);
      if (a.mn1 < b.mn1) mn1 = a.mn1, mnc = a.mnc, mn2 =
            min(b.mn1, a.mn2);
      if (a.mn1 > b.mn1) mn1 = b.mn1, mnc = b.mnc, mn2 =
            min(a.mn1, b.mn2):
       if (a.mn1 == b.mn1) mn1 = a.mn1, mnc = a.mnc + b.mnc,
            mn2 = min(a.mn2, b.mn2):
   }
};
// 0 - indexed / inclusive - inclusive
template <class node>
struct STB {
   vector<node> st; int n;
   void build(int u, int i, int j, vector<node> &arr) {
      if (i == j) {
          st[u] = arr[i];
          return:
       int m = (i + i) / 2, l = u * 2 + 1, r = u * 2 + 2;
       build(1, i, m, arr), build(r, m + 1, j, arr);
       st[u] = node(st[l], st[r]):
   void push_add(int u, int i, int j, ll v) {
       st[u].s += (i - i + 1) * v:
       st[u].mx1 += v, st[u].mn1 += v, st[u].lz += v;
       if (st[u].mx2 != LLONG MIN) st[u].mx2 += v:
       if (st[u].mn2 != LLONG MAX) st[u].mn2 += v:
   void push_max(int u, ll v, bool l) { // for min op
      if (v >= st[u].mx1) return;
      st[u].s -= st[u].mx1 * st[u].mxc:
       st[u].mx1 = v:
       st[u].s += st[u].mx1 * st[u].mxc;
      if (1) st[u].mn1 = st[u].mx1:
       else if (v <= st[u].mn1) st[u].mn1 = v;</pre>
       else if (v < st[u].mn2) st[u].mn2 = v;
   void push_min(int u, ll v, bool l) { // for max op
       if (v <= st[u].mn1) return:</pre>
       st[u].s -= st[u].mn1 * st[u].mnc:
       st[u].mn1 = v:
```

```
st[u].s += st[u].mn1 * st[u].mnc:
   if (1) st[u].mx1 = st[u].mn1:
   else if (v \ge st[u].mx1) st[u].mx1 = v;
   else if (v > st[u].mx2) st[u].mx2 = v:
void push(int u. int i. int i) {
   if (i == j) return;
   // add
   int m = (i + j) / 2, l = u * 2 + 1, r = u * 2 + 2;
   push_add(1, i, m, st[u].lz);
   push add(r. m + 1. i. st[u].lz):
   st[u].lz = 0:
   // min
    push_max(1, st[u].mx1, i == m);
   push_max(r, st[u].mx1, m + 1 == j);
   // max
   push_min(1, st[u].mn1, i == m);
   push_min(r, st[u].mn1, m + 1 == r);
node query(int a, int b, int u, int i, int j) {
   if (b < i || j < a) return node();</pre>
   if (a <= i && j <= b) return st[u];</pre>
   push(u, i, j);
   int m = (i + j) / 2, l = u * 2 + 1, r = u * 2 + 2;
   return node(query(a, b, 1, i, m), query(a, b, r, m +
        1, j));
void update add(int a, int b, ll v, int u, int i, int i)
   if (b < i || j < a) return;</pre>
   if (a <= i && j <= b) {</pre>
       push_add(u, i, j, v);
       return:
   push(u, i, j);
   int m = (i + i) / 2, l = u * 2 + 1, r = u * 2 + 2:
   update_add(a, b, v, l, i, m);
   update_add(a, b, v, r, m + 1, j);
   st[u] = node(st[1], st[r]);
void update min(int a, int b, ll v, int u, int i, int i)
   if (b < i || j < a || v >= st[u].mx1) return;
   if (a <= i && j <= b && v > st[u].mx2) {
       push_max(u, v, i == j);
       return:
   push(u, i, j);
   int m = (i + j) / 2, l = u * 2 + 1, r = u * 2 + 2;
    update_min(a, b, v, l, i, m);
```

```
update_min(a, b, v, r, m + 1, j);
       st[u] = node(st[1], st[r]);
   void update_max(int a, int b, ll v, int u, int i, int j)
       if (b < i || i < a || v <= st[u].mn1) return;</pre>
       if (a <= i && j <= b && v < st[u].mn2) {</pre>
          push_min(u, v, i == j);
          return:
       push(u, i, j);
       int m = (i + j) / 2, l = u * 2 + 1, r = u * 2 + 2;
       update_max(a, b, v, l, i, m);
       update_max(a, b, v, r, m + 1, j);
       st[u] = node(st[1], st[r]);
   STB(vector<node> &v, int N) : n(N), st(N * 4 + 5) { build
        (0, 0, n - 1, v): }
   node query(int a, int b) { return query(a, b, 0, 0, n -
        1); }
   void update_add(int a, int b, ll v) { update_add(a, b, v,
         0, 0, n - 1); }
   void update_min(int a, int b, ll v) { update_min(a, b, v,
         0.0.n - 1): 
   void update_max(int a, int b, ll v) { update_max(a, b, v,
         0.0.n - 1): 
};
```

### 5.2 Treap

```
#include "../Template.cpp"

mt19937 gen(chrono::high_resolution_clock::now().
    time_since_epoch().count());

// 101 Implicit Treap //

struct Node {
    int p, sz = 0, v, acc, l = -1, r = -1;
    Node() : v(0), acc(0) {}
    Node(int x) : p(gen()), sz(1), v(x), acc(x) {}
    void recalc(const Node &a, const Node &b) {
        sz = a.sz + b.sz + 1;
        acc = v + a.acc + b.acc;
    }
};

template <class node>
```

```
struct Treap {
   vector<node> t:
   int n. r = -1:
   node get(int u) { return u != -1 ? t[u] : node(); }
   void recalc(int u) { t[u].recalc(get(t[u].l), get(t[u].r)
   int merge(int 1, int r) {
       if (min(1, r) == -1) return 1 != -1 ? 1 : r;
       int ans = (t[1].p < t[r].p) ? 1 : r;
       if (ans == 1) t[1].r = merge(t[1].r, r), recalc(1);
       if (ans == r) t[r].l = merge(l, t[r].l), recalc(r):
       return ans;
   pii split(int u, int id) {
       if (u == -1) return {-1, -1};
       int szl = get(t[u].1).sz:
      if (szl >= id) {
          pii ans = split(t[u].1, id);
          t[u].1 = ans.ss:
          recalc(u);
          return {ans.ff. u}:
      pii ans = split(t[u].r, id - szl - 1);
       t[u].r = ans.ff;
       recalc(u):
       return {u. ans.ss}:
   }
   Treap(vi &v) : n(sz(v)) {
      for (int i = 0; i < n; i++) t.eb(v[i]), r = merge(r, v)
   }
};
// Complete Implicit Treap with Lazy propagation //
struct Node {
   int p, sz = 0, v, acc, l = -1, r = -1, par = -1, lzv = 0;
   bool lz = false, f = false:
   Node(): v(0), acc(0) {}
   Node(int x) : p(gen()), sz(1), v(x), acc(x) {}
   void recalc(const Node &a. const Node &b) {
      sz = a.sz + b.sz + 1;
      acc = v + a.acc + b.acc;
   void upd_lazv(int x) { lz = 1, lzv += x; }
   void lazy() { v += lzv, acc += sz * lzv, lz = 0, lzv = 0;
   void flip() { swap(1, r), f = 0; }
```

```
};
template <class node>
struct Treap {
   vector<node> t;
   int n. r = -1:
   node get(int u) { return u != -1 ? t[u] : node(); }
   void recalc(int u) {
       int 1 = t[u].1, r = t[u].r;
       push(1), push(r), flip(1), flip(r);
       t[u].recalc(get(1), get(r)):
   void push(int u) {
       if (u == -1 || !t[u].lz) return;
       int 1 = t[u].1, r = t[u].r;
       if (1 != -1) t[1].upd_lazy(t[u].lzv);
       if (r != -1) t[r].upd_lazy(t[u].lzv);
       t[u].lazv():
   }
   void flip(int u) {
       if (u == -1 || !t[u].f) return;
       int 1 = t[u].1, r = t[u].r;
       if (1 != -1) t[1].f ^= 1;
       if (r != -1) t[r].f ^= 1;
       t[u].flip();
   int merge(int 1, int r) {
       if (min(1, r) == -1) return 1 != -1 ? 1 : r;
       push(1), push(r), flip(1), flip(r);
       int ans = (t[1].p < t[r].p) ? 1 : r;</pre>
       if (ans == 1) t[1].r = merge(t[1].r, r), recalc(1);
       if (ans == r) t[r].l = merge(l, t[r].l), recalc(r);
       if (t[ans].l != -1) t[t[ans].l].par = ans; // only if
             parent needed
       if (t[ans].r != -1) t[t[ans].r].par = ans: // only if
            parent needed
       return ans;
   pii split(int u, int id) {
       if (u == -1) return {-1, -1}:
       push(u);
      flip(u):
       int szl = get(t[u].1).sz;
       if (szl >= id) {
          pii ans = split(t[u].1, id):
          if (ans.ss != -1) t[ans.ss].par = u; // only if
               parent needed
          if (ans.ff != -1) t[ans.ff].par = -1; // only if
               parent needed
```

```
t[u].1 = ans.ss:
       recalc(u):
       return {ans.ff, u};
   pii ans = split(t[u].r, id - szl - 1);
   if (ans.ff != -1) t[ans.ff].par = u; // only if
        parent needed
   if (ans.ss != -1) t[ans.ss].par = -1; // only if
        parent needed
   t[u].r = ans.ff;
   recalc(u):
   return {u. ans.ss}:
int update(int u, int 1, int r, int v) {
   pii a = split(u, 1), b = split(a.ss, r - 1 + 1);
   t[b.ff].upd_lazy(v);
   return merge(a.ff, merge(b.ff, b.ss));
void print(int u) {
   if (u == -1) return:
   push(u), flip(u);
   print(t[u].1);
   cout << t[u].v << ' ';
   print(t[u].r);
Treap(vi &v) : n(sz(v)) {
   for (int i = 0; i < n; i++) t.eb(v[i]), r = merge(r)
        i);
```

### 5.3 dsu

};

```
struct Dsu {
   vector<int> p, r;

   // initialize the disjoint-set-union to all unitary sets
   void reset(int N) {
      p.resize(N), r.assign(N, 0);
      rep(i, N) p[i] = i;
   }

   // find the leader node corresponding to node 'i'
   int find(int i) {
      if (p[i] != i) p[i] = find(p[i]);
      return p[i];
   }
```

```
// perform union on the two sets that 'i' and 'j' belong
    to

void unite(int i, int j) {
        i = find(i), j = find(j);
        if (i == j) return;
        if (r[i] > r[j]) swap(i, j);
        if (r[i] == r[j]) r[j] += 1;
        p[i] = j;
    }
};
```

#### 5.4 mo

```
struct Query { int 1, r, idx; };
// answer segment queries using only 'add(i)', 'remove(i)'
    and 'get()'
// functions.
// complexity: O((N + Q) * sqrt(N) * F)
// N = length of the full segment
// Q = amount of queries
// F = complexity of the 'add', 'remove' functions
template <class A, class R, class G, class T>
void mo(vector<Query> &queries, vector<T> &ans. A add. R
    remove, G get) {
   int Q = queries.size(), B = (int)sqrt(Q);
   sort(queries.begin(), queries.end(), [&](Query &a, Query
       return make_pair(a.1 / B, a.r) < make_pair(b.1 / B, b</pre>
            .r):
   }):
   ans.resize(Q);
   int 1 = 0, r = 0;
   for (auto &q : queries) {
       while (r < q.r) add(r), r++;
       while (1 > q.1) 1--, add(1);
       while (r > q.r) r--, remove(r);
       while (1 < q.1) remove(1), 1++;</pre>
       ans[q.idx] = get();
   }
```

### 5.5 ordered-set

#include <ext/pb\_ds/assoc\_container.hpp>

```
#include <ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;
typedef tree<int, null_type, less<int>, rb_tree_tag,
            tree_order_statistics_node_update>
   ordered set:
int main() {
   ordered set p:
   p.insert(5); p.insert(2); p.insert(6); p.insert(4); // 0(
   // value at 3rd index in sorted array. O(log n). Output:
   cout << "Value at 3rd index: " << *p.find_by_order(3) <<</pre>
   // index of number 6. O(log n). Output: 3
   cout << "Index of number 6: " << p.order of kev(6) <<</pre>
        endl:
   // number 7 not in the set but it will show the index
   // number if it was there in sorted array. Output: 4
   cout << "Index of number 7:" << p.order_of_key(7) << end1</pre>
   // number of elements in the range [3, 10)
   cout << p.order_of_key(10) - p.order_of_key(3) << endl;</pre>
```

# 5.6 persistent-segment-tree-lazy

```
template <class T>
struct Node {
   T x. lz:
   int 1 = -1, r = -1;
template <class T>
struct Pstl {
   int N;
   vector<Node<T>> a:
   vector<int> head:
   T qneut() { return 0; }
   T merge(T 1, T r) { return 1 + r; }
   T uneut() { return 0; }
   T accum(T u, T x) { return u + x: }
   T apply(T x, T lz, int l, int r) { return x + (r - 1) *
        1z: }
   int build(int vl, int vr) {
```

```
if (vr - vl == 1) a.push back({gneut(), uneut()}): //
         node construction
   else {
       int vm = (vl + vr) / 2, l = build(vl, vm), r =
           build(vm, vr);
       a.push_back({merge(a[1].x, a[r].x), uneut(), 1, r
           }); // query merge
   }
   return a.size() - 1:
T querv(int 1, int r, int v, int v1, int vr, T acc) {
   if (1 >= vr || r <= vl) return gneut();</pre>
        // query neutral
   if (1 <= v1 && r >= vr) return apply(a[v].x, acc, vl, | struct Node {
         vr); // update op
   acc = accum(acc, a[v].lz):
        // update merge
   int vm = (vl + vr) / 2:
   return merge(query(1, r, a[v].1, v1, vm, acc), query(
        1, r, a[v].r, vm, vr, acc)); // query merge
int update(int 1, int r, T x, int v, int v1, int vr) {
   if (1 >= vr || r <= vl || r <= 1) return v:
   a.push_back(a[v]);
   v = a.size() - 1:
   if (1 <= v1 && r >= vr) {
       a[v].x = apply(a[v].x, x, vl, vr); // update op
       a[v].lz = accum(a[v].lz, x); // update merge
   } else {
       int vm = (vl + vr) / 2;
       a[v].1 = update(1, r, x, a[v].1, v1, vm);
       a[v].r = update(1, r, x, a[v].r, vm, vr);
       a[v].x = merge(a[a[v].1].x, a[a[v].r].x); //
           query merge
   }
   return v;
Pstl(int N) : N(N) { head.push_back(build(0, N)); }
T query(int t, int 1, int r) {
   return query(1, r, head[t], 0, N, uneut()); // update
         neutral
int update(int t, int 1, int r, T x) {
   return head.push_back(update(1, r, x, head[t], 0, N))
        . head.size() - 1:
```

# 5.7 persistent-segment-tree

// Pst<Node<11>> pst;

// pst = {N}:

```
// int newtime = pst.update(time, index, value);
// Node<11> result = pst.querv(newtime, left, right):
template <class T>
   int 1 = -1, r = -1:
   Node(): x(0) {}
   Node(T x) : x(x) \{ \}
   Node (Node a, Node b, int 1 = -1, int r = -1) : x(a.x + b.
        x), 1(1), r(r) {}
};
template <class U>
struct Pst {
   int N;
   vector<U> a;
    vector<int> head;
   int build(int vl, int vr) {
       if (vr - vl == 1) a.push_back(U()); // node
            construction
           int vm = (vl + vr) / 2, l = build(vl, vm), r =
                build(vm. vr):
           a.push_back(U(a[1], a[r], 1, r)); // query merge
       return a.size() - 1:
   }
    U query(int 1, int r, int v, int v1, int vr) {
       if (1 >= vr || r <= vl) return U(); // query neutral</pre>
       if (1 <= v1 && r >= vr) return a[v]:
       int vm = (v1 + vr) / 2;
       return U(query(1, r, a[v].1, v1, vm), query(1, r, a[v
            1.r. vm. vr)): // querv merge
   }
    int update(int i, U x, int v, int vl, int vr) {
       a.push_back(a[v]);
```

```
v = a.size() - 1:
       if (vr - vl == 1) a[v] = x; // update op
       else {
          int vm = (vl + vr) / 2:
          if (i < vm) a[v].l = update(i, x, a[v].l, vl, vm)</pre>
          else a[v].r = update(i, x, a[v].r, vm, vr);
          a[v] = U(a[a[v].1], a[a[v].r], a[v].1, a[v].r);
               // query merge
      }
       return v:
   Pst() {}
   Pst(int N) : N(N) { head.push_back(build(0, N)); }
   U querv(int t, int 1, int r) {
       return query(1, r, head[t], 0, N);
   int update(int t, int i, U x) {
       return head.push_back(update(i, x, head[t], 0, N)),
           head.size() - 1:
   }
};
```

### 5.8 segment-tree-lazy

```
// 0-based. inclusive-exclusive
// usage:
// St13<11> a;
// a = {N}:
template <class T>
struct Stl {
   // immediate, lazy
   vector<pair<T, T>> a;
   T qneutral() { return 0; }
   T merge(T 1, T r) { return 1 + r: }
   T uneutral() { return 0: }
   void update(pair<T, T> &u, T val, int l, int r) { u.first
         += val * (r - 1), u.second += val; }
   St1() {}
   Stl(int N) : a(4 * N, {gneutral(), uneutral()}) {} //
        node neutral
   void push(int v, int vl, int vm, int vr) {
       update(a[2 * v], a[v].second, v1, vm); // node update
```

```
update(a[2 * v + 1], a[v].second, vm, vr): // node
       a[v].second = uneutral();
                                              // update
            neutral
   // query for range [1, r)
   T query(int 1, int r, int v = 1, int vl = 0, int vr = -1)
       if (vr == -1) vr = a.size() / 4;
       if (1 <= v1 && r >= vr) return a[v].first: // guerv
       if (1 >= vr || r <= vl) return qneutral(); // query</pre>
       int vm = (vl + vr) / 2:
       push(v, vl, vm, vr);
       return merge(query(1, r, 2 * v, v1, vm), query(1, r,
            2 * v + 1, vm, vr)); // item merge
   // update range [1, r) using val
   void update(int 1, int r, T val, int v = 1, int vl = 0,
        int vr = -1) {
       if (vr == -1) vr = a.size() / 4;
       if (1 >= vr || r <= vl || r <= 1) return;</pre>
       if (1 \le v1 \&\& r \ge vr) update(a[v], val, vl, vr); //
             node update
       else {
           int vm = (v1 + vr) / 2;
           push(v. vl. vm. vr):
           update(1, r, val, 2 * v, vl, vm);
           update(1, r, val, 2 * v + 1, vm, vr);
           a[v].first = merge(a[2 * v].first, a[2 * v + 1].
               first); // node merge
       }
};
```

### 5.9 segment-tree

```
// usage:
// St<Node<ll>> st;
// st = {N};
// st.update(index, new_value);
// Node<ll> result = st.query(left, right);

template <class T>
struct Node {
   T x;
```

```
Node(): x(0) {}
   Node(T x) : x(x) \{ \}
   Node(Node a, Node b) : x(a.x + b.x) {}
template <class U>
struct St {
   vector<U> a;
   St() {}
   St(int N) : a(4 * N, U()) {} // node neutral
   // query for range [1, r)
   U query(int 1, int r, int v = 1, int vl = 0, int vr = -1)
       if (vr == -1) vr = a.size() / 4;
       if (1 <= v1 && r >= vr) return a[v]: // item
           construction
       int vm = (vl + vr) / 2:
       if (1 >= vr || r <= vl) return U();</pre>
                                       // item neutral
       return U(query(1, r, 2 * v, v1, vm), query(1, r, 2 *
           v + 1, vm, vr)); // item merge
   }
   // set element i to val
   void update(int i, U val, int v = 1, int vl = 0, int vr =
         -1) {
       if (vr == -1) vr = a.size() / 4;
      if (vr - vl == 1) a[v] = val: // item update
       else {
          int vm = (vl + vr) / 2;
          if (i < vm) update(i, val, 2 * v, vl, vm);</pre>
          else update(i, val, 2 * v + 1, vm, vr);
          a[v] = U(a[2 * v], a[2 * v + 1]); // node merge
      }
   }
```

# 5.10 sparse-table

```
Sparse(int N) : st{vector<T>(N)} {}
   T &operator[](int i) { return st[0][i]; }
   // O(N log N) time, O(N log N) memory
   void init() {
       int N = st[0].size();
       int npot = N <= 1 ? 1 : 32 - __builtin_clz(N);</pre>
       st.resize(npot);
       repx(i, 1, npot) rep(j, N + 1 - (1 << i)) st[i].
            push back(
              op(st[i - 1][i], st[i - 1][i + (1 << (i - 1))]
       ); // query op
   // query maximum in the range [1, r) in O(1) time
   // range must be nonempty!
   T querv(int 1. int r) {
       int i = 31 - __builtin_clz(r - 1);
       return op(st[i][1], st[i][r - (1 << i)]); // query op</pre>
};
```

# 5.11 unordered-map

```
// hackproof rng
static mt19937 rng(chrono::steady_clock::now().
    time_since_epoch().count());
// deterministic rng
uint64 t splitmix64(uint64 t *x) {
   uint64_t z = (*x += 0x9e3779b97f4a7c15);
   z = (z ^ (z >> 30)) * 0xbf58476d1ce4e5b9:
   z = (z ^ (z >> 27)) * 0x94d049bb133111eb;
   return z ^ (z >> 31);
// hackproof unordered map hash
struct Hash {
   size_t operator()(const 11 &x) const {
       static const uint64_t RAND =
          chrono::steady_clock::now().time_since_epoch().
               count();
       uint64 t z = x + RAND + 0x9e3779b97f4a7c15:
       z = (z ^ (z >> 30)) * 0xbf58476d1ce4e5b9;
       z = (z \hat{z} > 27) * 0x94d049bb133111eb;
       return z^(z >> 31):
```

```
};

// hackproof unordered_map
template <class T, class U>
using umap = unordered_map<T, U, Hash>;

// hackproof unordered_set
template <class T>
using uset = unordered_set<T, Hash>;
```

# 6 imprimible

### 7 math

### 7.1 arithmetic

```
// floor(log2(n)) without precision loss
inline int floor_log2(int n) { return n <= 1 ? 0 : 31 -</pre>
     __builtin_clz(n): }
// ceil(log2(n)) without precision loss
inline int ceil_log2(int n) { return n <= 1 ? 0 : 32 -</pre>
     __builtin_clz(n - 1): }
inline ll floordiv(ll a, ll b) {
   return a / b - ((a ^ b) < 0 && a % b);
inline 11 ceildiv(11 a, 11 b) {
   return a / b + ((a ^ b) >= 0 && a % b);
// a^e through binary exponentiation.
ll binexp(ll a, ll e) {
   ll res = 1: // neutral element
   while (e) {
       if (e & 1) res = res * a; // multiplication
                              // multiplication
       a = a * a;
       e >>= 1;
   return res;
```

### 7.2 crt

```
pair<11, 11> solve_crt(const vector<pair<11, 11>> &eqs) {
```

```
1l a0 = eqs[0].first, p0 = eqs[0].second;
repx(i, 1, eqs.size()) {
    1l a1 = eqs[i].first, p1 = eqs[i].second;
    1l k1, k0;
    1l d = ext_gcd(p1, p0, k1, k0);
    a0 -= a1;
    if (a0 % d != 0) return {-1, -1};
    p0 = p0 / d * p1;
    a0 = a0 / d * k1 % p0 * p1 % p0 + a1;
    a0 = (a0 % p0 + p0) % p0;
}
return {a0, p0};
```

# 7.3 discrete-log

```
// discrete logarithm log a(b).
// solve b \hat{x} = a \pmod{M} for the smallest x.
// returns -1 if no solution is found.
// time: O(sqrt(M))
11 dlog(ll a, ll b, ll M) {
   11 k = 1, s = 0;
   while (true) {
       ll g = \_gcd(b, M);
       if (g <= 1) break;</pre>
       if (a == k) return s:
       if (a % g != 0) return -1;
       a = g, M = g, s += 1, k = b / g * k % M;
   11 N = sqrt(M) + 1;
   umap<11, 11> r:
   rep(q, N + 1) {
       r[a] = q;
       a = a * b % M;
   ll bN = binexp(b, N, M), bNp = k;
   repx(p, 1, N + 1) {
       bNp = bNp * bN % M;
       if (r.count(bNp)) return N * p - r[bNp] + s;
   return -1;
```

# 7.4 fft

```
using cd = complex<double>;
const double PI = acos(-1);
// compute the DFT of a power-of-two-length sequence.
// if 'inv' is true, computes the inverse DFT.
void fft(vector<cd> &a. bool inv) {
   int N = a.size(), k = 0, b;
   assert(N == 1 << builtin ctz(N));</pre>
   repx(i, 1, N) {
       for (b = N >> 1: k & b:) k ^= b, b >>= 1:
       if (i < (k ^= b)) swap(a[i], a[k]);</pre>
   for (int 1 = 2; 1 <= N; 1 <<= 1) {
       double ang = 2 * PI / 1 * (inv ? -1 : 1):
       cd wl(cos(ang), sin(ang));
       for (int i = 0; i < N; i += 1) {</pre>
           cd w = 1:
           rep(i, 1 / 2) {
              cd u = a[i + j], v = a[i + j + 1 / 2] * w;
              a[i + i] = u + v:
              a[i + j + 1 / 2] = u - v;
              w *= wl:
      }
   }
   if (inv) rep(i, N) a[i] /= N:
const 11 MOD = 998244353, ROOT = 15311432:
// const 11 MOD = 2130706433, ROOT = 1791270792;
// const 11 MOD = 922337203673733529711, ROOT =
    532077456549635698311;
void find_root_of_unity(ll M) {
   11 c = M - 1, k = 0:
   while (c \% 2 == 0) c /= 2, k += 1;
   // find proper divisors of M - 1
   vector<ll> divs:
   for (11 d = 1: d < c: d++) {
       if (d * d > c) break;
       if (c % d == 0) rep(i, k + 1) divs.push_back(d << i);</pre>
   rep(i, k) divs.push_back(c << i);</pre>
   // find any primitive root of M
   11 G = -1;
```

```
repx(g, 2, M) {
       bool ok = true:
       for (ll d : divs) ok &= (binexp(g, d, M) != 1);
       if (ok) {
           G = g;
           break;
   assert(G != -1);
   ll w = binexp(G, c, M):
   cerr << "M = c * 2^k + 1" << endl:
   cerr << " M = " << M << endl;
    cerr << " c = " << c << endl:
    cerr << " k = " << k << endl:
    cerr << " w^(2^k) == 1" << endl;
   cerr << " w = g^{(M-1)/2^k} = g^c" << endl;
   cerr << " g = " << G << endl;</pre>
   cerr << " w = " << w << endl:
}
// compute the DFT of a power-of-two-length sequence, modulo
     a special prime
// number with an Nth root of unity, where N is the length
    of the sequence.
void ntt(vector<ll> &a, bool inv) {
   vector<ll> wn:
   for (ll p = ROOT; p != 1; p = p * p % MOD) wn.push back(p \frac{1}{7} 'a' is a list of rows
        );
   int N = a.size(), k = 0, b;
   assert(N == 1 << \_builtin\_ctz(N) && N <= 1 << wn.size()) | // 1 -> unique solution, stored in 'ans'
   rep(i, N) a[i] = (a[i] \% MOD + MOD) \% MOD;
   repx(i, 1, N) {
       for (b = N >> 1; k & b;) k ^= b, b >>= 1;
       if (i < (k ^= b)) swap(a[i], a[k]);</pre>
   for (int 1 = 2: 1 <= N: 1 <<= 1) {
       11 wl = wn[wn.size() - __builtin_ctz(1)];
       if (inv) wl = multinv(wl, MOD);
       for (int i = 0; i < N; i += 1) {</pre>
          11 w = 1:
           repx(i, 0, 1 / 2) {
              11 u = a[i + j], v = a[i + j + 1 / 2] * w %
              a[i + j] = (u + v) \% MOD;
```

```
a[i + i + 1 / 2] = (u - v + MOD) \% MOD:
              w = w * w1 \% MOD:
   }
   11 q = multinv(N, MOD);
   if (inv) rep(i, N) a[i] = a[i] * q % MOD;
void convolve(vector<cd> &a. vector<cd> b. int n) {
   n = 1 \ll (32 - builtin clz(2 * n - 1)):
   a.resize(n), b.resize(n);
   fft(a, false), fft(b, false);
   rep(i, n) a[i] *= b[i];
   fft(a, true);
```

### **7.5** gauss

```
const double EPS = 1e-9:
// solve a system of equations.
// complexity: O(min(N, M) * N * M)
// the last value in each row is the result of the equation
// return values:
// 0 -> no solutions
// -1 -> infinitely many solutions, one of which is stored
// UNTESTED
int gauss(vector<vector<double>> a, vector<double> &ans) {
   int N = a.size(). M = a[0].size() - 1:
   vector<int> where(M, -1):
   for (int j = 0, i = 0; j < M && i < N; j++) {
       int sel = i:
       repx(k, i, N) if (abs(a[k][j]) > abs(a[sel][j])) sel
       if (abs(a[sel][j]) < EPS) continue;</pre>
       repx(k, j, M + 1) swap(a[sel][k], a[i][k]);
       where[j] = i;
       rep(k, N) if (k != i) {
          double c = a[k][j] / a[i][j];
          repx(1, j, M + 1) a[k][1] -= a[i][1] * c;
```

```
i++:
}
ans.assign(M, 0);
rep(i, M) if (where[i] != -1) ans[i] = a[where[i]][M] / a
     [where[i]][i]:
rep(i, N) {
   double sum = 0;
   rep(j, M) sum += ans[j] * a[i][j];
   if (abs(sum - a[i][M]) > EPS) return 0;
rep(i, M) if (where[i] == -1) return -1;
return 1:
```

#### 7.6 matrix

```
using T = 11;
struct Mat {
   int N. M:
   vector<vector<T>> v;
   Mat(int n, int m) : N(n), M(m), v(N, vector<T>(M)) {}
   Mat(int n) : Mat(n, n) { rep(i, N) v[i][i] = 1; }
   vector<T> &operator[](int i) { return v[i]; }
   Mat operator*(Mat &r) {
       assert(M == r.N);
      int n = N, m = r.M, p = M;
      Mat a(n. m):
      rep(i, n) rep(j, m) {
          a[i][i] = T():
                                                       //
               neutral
          rep(k, p) a[i][j] = a[i][j] + v[i][k] * r[k][j];
               // mul. add
      }
       return a:
   Mat binexp(ll e) {
       assert(N == M);
      Mat a = *this, res(N); // neutral
          if (e & 1) res = res * a; // mul
          a = a * a:
                                  // miil
          e >>= 1:
```

```
return res;
}

friend ostream & operator << (ostream & s, Mat & a) {
    rep(i, a.N) {
        rep(j, a.M) s << a[i][j] << " ";
        s << endl;
    }
    return s;
}
</pre>
```

### 7.7 mobius

```
short mu[MAXN] = {0,1};
void mobius(){
  fore(i,1,MAXN)if(mu[i])for(int j=i+i;j<MAXN;j+=i)mu[j]-=mu[
         i];
}</pre>
```

### 7.8 mod

```
11 binexp(ll a, ll e, ll M) {
   assert(e >= 0):
   ll res = 1 \% M:
   while (e) {
       if (e & 1) res = res * a % M:
       a = a * a % M;
       e >>= 1:
   return res;
11 multinv(11 a, 11 M) { return binexp(a, M - 2, M); }
// calculate gcd(a, b).
// also, calculate x and v such that:
// a * x + b * v == gcd(a, b)
// time: O(log min(a, b))
// (ignoring complexity of arithmetic)
11 ext_gcd(ll a, ll b, ll &x, ll &y) {
   if (b == 0) {
       x = 1, y = 0;
       return a:
   ll d = ext_gcd(b, a \% b, y, x);
```

```
y -= a / b * x:
   return d:
// compute inverse with any M.
// a and M must be coprime for inverse to exist!
11 multinv_euc(11 a, 11 M) {
   11 x, v;
   ext_gcd(a, M, x, y);
   return x;
// multiply two big numbers (~10^18) under a large modulo,
     without resorting to
// bigints.
11 bigmul(ll x, ll v, ll M) {
   11 z = 0:
   while (y) {
       if (v \& 1) z = (z + x) \% M:
       x = (x << 1) \% M, v >>= 1:
   }
   return z:
// all modular inverses from 1 to inv.size()-1
void multinv_all(vector<ll> &inv) {
   inv[1] = 1:
   repx(i, 2, inv.size())
       inv[i] = MOD - (MOD / i) * inv[MOD % i] % MOD;
// change to __int128 if checking numbers over 10^9
bool isprime(ll n) {
   if (n < 2 | | n % 6 % 4 != 1) return n - 2 < 2;
   11 A\Gamma1 = {2, 325, 9375, 28178, 450775, 9780504,
        1795265022}:
   11 s = builtin ctzll(n - 1), d = n >> s:
   for (int a : A) {
       11 p = binexp(a, d, n), i = s;
       while (p != 1 && p != n - 1 && a % n && i--) p = p *
            р % n:
       if (p != n - 1 && i != s) return 0;
   return 1:
struct Mod {
   int a:
   static const int M = 1e9 + 7:
```

```
Mod(11 aa) : a((aa % M + M) % M) {}
   Mod operator+(Mod rhs) const { return (a + rhs.a) % M; }
   Mod operator-(Mod rhs) const { return (a - rhs.a + M) % M
        ; }
   Mod operator-() const { return Mod(0) - *this: }
   Mod operator*(Mod rhs) const { return (11)a * rhs.a % M;
   Mod operator+=(Mod rhs) { return *this = *this + rhs; }
   Mod operator == (Mod rhs) { return *this = *this - rhs: }
   Mod operator*=(Mod rhs) { return *this = *this * rhs: }
   Mod bigmul(ll big) const { return ::bigmul(a, big, M); }
   Mod binexp(ll e) const { return ::binexp(a, e, M); }
   // Mod multinv() const { return ::multinv(a, M): } //
        prime M
   Mod multinv() const { return ::multinv_euc(a, M); } //
        possibly composite M
};
// dynamic modulus
struct DMod {
   int a. M:
   DMod(11 aa. 11 m) : M(m), a((aa % m + m) % m) {}
   DMod operator+(DMod rhs) const { return {(a + rhs.a) % M,
   DMod operator-(DMod rhs) const { return {(a - rhs.a + M)
        % M, M}; }
   DMod operator-() const { return DMod(0, M) - *this: }
   DMod operator*(DMod rhs) const { return {(11)a * rhs.a %
        M. M: }
   DMod operator+=(DMod rhs) { return *this = *this + rhs: }
   DMod operator==(DMod rhs) { return *this = *this - rhs; }
   DMod operator*=(DMod rhs) { return *this = *this * rhs; }
   DMod bigmul(ll big) const { return {::bigmul(a, big, M).
   DMod binexp(ll e) const { return {::binexp(a, e, M), M};
   DMod multinv() const { return {::multinv(a, M), M}: } //
   // DMod multinv() const { return {::multinv euc(a, M), M
        }; } // possibly composite M
};
```

# 7.9 primes

```
// counts the divisors of a positive integer in O(\operatorname{sqrt}(n))
11 count_divisors(11 x) {
   11 \text{ divs} = 1, i = 2;
   for (11 divs = 1, i = 2; x > 1; i++) {
       if (i * i > x) {
           divs *= 2:
           break:
       for (11 d = divs: x % i == 0: x /= i) divs += d:
   return divs;
}
// gets the prime factorization of a number in O(sqrt(n))
vector<pair<ll, int>> factorize(ll x) {
   vector<pair<11, int>> f;
   for (11 k = 2; x > 1; k++) {
       if (k * k > x) {
           f.push_back(\{x, 1\});
           break:
       int n = 0;
       while (x \% k == 0) x /= k, n++:
       if (n > 0) f.push_back(\{k, n\});
   return f;
}
// iterate over all divisors of a number.
// divisor count upper bound: n^(1.07 / ln ln n)
template <class OP>
void divisors(ll x. OP op) {
   auto facts = factorize(x);
   vector<int> f(facts.size());
   while (true) {
       11 v = 1;
       rep(i, f.size()) rep(j, f[i]) y *= facts[i].first;
       op(y);
       int i:
       for (i = 0; i < f.size(); i++) {</pre>
           if (f[i] <= facts[i].second) break:</pre>
           f[i] = 0;
       if (i == f.size()) break;
```

```
// computes euler totative function phi(x), counting the
     amount of integers in
// [1, x] that are coprime with x.
// time: O(sqrt(x))
ll phi(ll x) {
   11 phi = 1. k = 2:
   for (; x > 1; k++) {
       if (k * k > x) {
           phi *= x - 1:
           break;
       11 k1 = 1, k0 = 0;
       while (x \% k == 0) x /= k, k0 = k1, k1 *= k;
       phi *= k1 - k0:
   return phi:
// computes primality up to N.
// considers 0 and 1 prime.
// O(N log N)
void sieve(int N, vector<bool> &prime) {
   prime.assign(N + 1, true);
   repx(n, 2, N + 1) if (prime[n]) for (int k = 2 * n; k <=
        N: k += n) prime[k] = false:
```

# $7.10 \quad \text{simplex}$

```
// Solves a general linear maximization problem: maximize $c
     ^T x$ subject to $Ax \le b$, $x \ge 0$.
// Returns -inf if there is no solution, inf if there are
     arbitrarily good solutions, or the maximum value of $c^
     T x$ otherwise.
// The input vector is set to an optimal $x$ (or in the
     unbounded case, an arbitrary solution fulfilling the
// Numerical stability is not guaranteed. For better
     performance, define variables such that x = 0 is
     viable.
// Usage:
// vvd A = \{\{1,-1\}, \{-1,1\}, \{-1,-2\}\}:
// vd b = \{1,1,-4\}, c = \{-1,-1\}, x;
// T val = LPSolver(A, b, c).solve(x);
// Time: O(NM * \#pivots), where a pivot may be e.g. an edge
     relaxation. O(2^n) in the general case.
```

```
#include "../common.h"
typedef double T; // long double, Rational, double + mod<P
typedef vector<T> vd:
typedef vector<vd> vvd;
const T eps = 1e-8, inf = 1 / .0;
#define MP make_pair
#define lti(X) \
   if (s == -1 \mid | MP(X[j], N[j]) < MP(X[s], N[s])) s = j
struct LPSolver {
   int m. n:
   vector<int> N, B;
   vvd D:
   LPSolver(const vvd &A. const vd &b. const vd &c) : m(b.
        size()), n(c.size()), N(n + 1), B(m), D(m + 2), vd(n)
        + 2)) {
       rep(i, m) rep(j, n) D[i][j] = A[i][j];
       rep(i, m) {
          B[i] = n + i;
          D[i][n] = -1;
          D[i][n + 1] = b[i];
       rep(j, n) {
          N[j] = j;
          D[m][i] = -c[i];
       }
       N[n] = -1;
       D[m + 1][n] = 1:
   void pivot(int r. int s) {
       T *a = D[r].data(), inv = 1 / a[s];
       rep(i, m + 2) if (i != r && abs(D[i][s]) > eps) {
          T *b = D[i].data(), inv2 = b[s] * inv;
          repx(j, 0, n + 2) b[j] -= a[j] * inv2;
          b[s] = a[s] * inv2:
       rep(j, n + 2) if (j != s) D[r][j] *= inv;
       rep(i, m + 2) if (i != r) D[i][s] *= -inv;
       D[r][s] = inv;
       swap(B[r], N[s]);
   bool simplex(int phase) {
       int x = m + phase - 1;
```

```
for (::) {
       int s = -1:
       rep(j, n + 1) if (N[j] != -phase) ltj(D[x]);
       if (D[x][s] >= -eps) return true;
       int r = -1;
       rep(i, m) {
          if (D[i][s] <= eps) continue;</pre>
           if (r == -1 || MP(D[i][n + 1] / D[i][s], B[i])
                < MP(D[r][n + 1] / D[r][s], B[r])) r = i
       if (r == -1) return false:
       pivot(r, s);
T solve(vd &x) {
   int r = 0;
   repx(i, 1, m) if (D[i][n + 1] < D[r][n + 1]) r = i;
   if (D[r][n + 1] < -eps) {
       pivot(r, n);
       if (!simplex(2) || D[m + 1][n + 1] < -eps) return</pre>
             -inf;
       rep(i, m) if (B[i] == -1) {
           int s = 0;
          repx(j, 1, n + 1) ltj(D[i]);
           pivot(i, s);
       }
```

```
}
bool ok = simplex(1);
    x = vd(n);
    rep(i, m) if (B[i] < n) x[B[i]] = D[i][n + 1];
    return ok ? D[m][n + 1] : inf;
}
};</pre>
```

### 7.11 theorems

```
// Burnside lemma
      For a set X, with members x in X, and a group G, with
    operations g in G, where g(x): X \rightarrow X.
     F_g is the set of x which are fixed points of g (ie. {
     x in X / g(x) = x }).
      The number of orbits (connected components in the
    graph formed by assigning each x a node and
     a directed edge between x and g(x) for every g) is
      M = the average of the fixed points of all g = (|F_g1|
     + |F_g2| + ... + |F_gn| ) / |G|
11
11
      If x are images and g are simmetries, then M
    corresponds to the amount of objects, |G|
      corresponds to the amount of simmetries, and F_g
    corresponds to the amount of simmetrical
```

```
images under the simmetry g.
11
// Rational root theorem
     All rational roots of the polynomials with integer
    coefficients:
//
11
      a0 * x^0 + a1 * x^1 + a2 * x^2 + ... + an * x^n = 0
11
11
     If these roots are represented as p / q, with p and q
11
      - p is an integer factor of a0
11
     - q is an integer factor of an
//
      Note that if a0 = 0, then x = 0 is a root, the
    polynomial can be divided by x and the theorem
      applies once again.
11
// Legendre's formula
     Considering a prime p, the largest power p^k that
    divides n! is given by:
//
      k = floor(n/p) + floor(n/p^2) + floor(n/p^3) + ...
11
11
     Which can be computed in O(log n / log p) time
```