

# Mohammad Hassaan Ejaz

22i - 2434

SE - B

## Question 1:

### Approach and Data Structures Used:

#### 1. Heap Data Structure:

A **d-ary heap** is a tree-based data structure where each node can have up to **d** children..

#### 2. Key Operations:

- **Insert Operation (ins):**
  - After insertion at the leaf, the element "floats up" to maintain the heap property.
- **Extract Operation (extract):**
  - Removes and returns root (either min or max).
  - The last element in the heap replaces the root, and then the heap is maintained by pushing the element down to satisfy heap order.
- **Heapify:**
  - Uses heapifyUp or heapifyDown methods to maintain heap structure.

#### 3. Implementation Details:

- The heap is implemented as an array where:
  - Root is first and has up to **d** children and similarly to each node.

- The index of a node's parent is calculated by the formula  $(i-1)/d$  where  $i$  is the current node's index.
  - The children of a node at index  $i$  are located at indices  $d*i + 1$ ,  $d*i + 2$ , ...,  $d*i + d$ .
  - **Float Up:** when a new element is added,. The element goes up until heap property is satisfied (min or max)
  - **Sink Down:** The element at the root is replaced by the last element in the heap, and the new root goes down till heap is maintained.
  - **Heap Type Change:** Supports both min-heap and max-heap based on input
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### Code Explanation:

- The HeapX class implements the d-ary heap. It is generic and works with any data type  $T$  (although integers are preferred).
- `add()` inserts an element into the heap, `popTop()` removes and returns the root, and `show()` prints the heap in a level-wise format.
- The `flipType()` method allows switching between min-heap and max-heap

## Question 2:

### Finite Automaton:

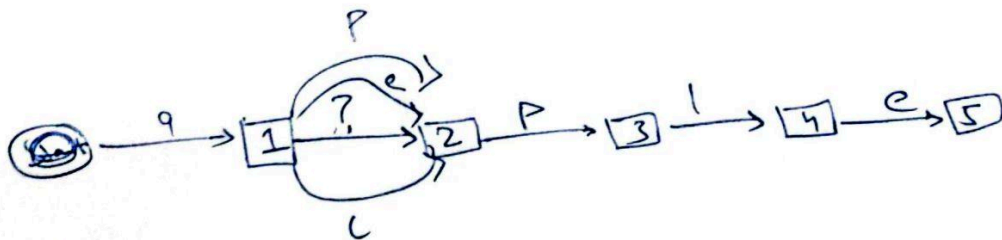
- **Wildcard (?):** The  $?$  wildcard matches any single character. This enhancement allows patterns like `a?ple` to match both "apple" and "ample".
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## 2. Transition Table and Diagram

## QUESTION 2

Pattern : apple

	Input	Next	Regex	Suffix
0	a	1	"a"	?ple
1	?	2	"a"	ple
2	p	3	"ap"	le
3	l	4	"apl"	e
4	e	5	"apple"	"
5	Finish		"apple"	"

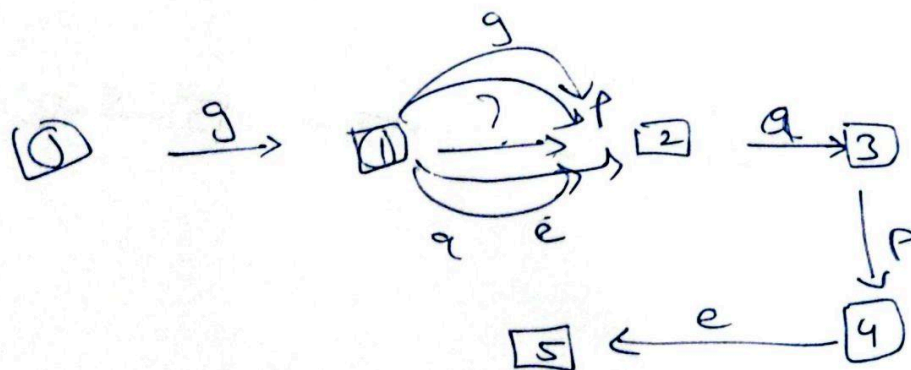




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Pattern g?ape

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State	Input	Next State	Prefix	Suffix
0	g	1	"	ape
1	?	2	g	ape
2	a	3	g?	pe
3	p	4	g?a	e
4	e	5	g?ape	"
5	"		g!ape	



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## 4. Time Complexity Analysis

- **Preprocessing (Building Automaton):**  $O(m * n)$  where  $m$  is pattern length and  $n$  is the number of characters (256 ASCII).
  - **Search Time:**  $O(L * m)$  where  $L$  is the length of the text and  $m$  is the pattern length.
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## 5. Code Explanation

- **FAEngine Class:** Builds the automaton's transition table and searches the text.
  - **Construct Method:** Creates the table based on the pattern.
  - **Scan Method:** Scans the text and finds matches using the automaton.
  - **Lowercase Conversion:** Ensures case-insensitive matching.
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## Conclusion

The enhanced FA algorithm efficiently handles wildcard characters and matches multiple patterns across text with optimal time complexity.

## Question 3:

### Top K Gardens:

#### Steps

1. **Max-Heap** creation — this takes  $O(n)$  time.

2. **Pop top k** gardens from heap — each pop is  **$O(\log n)$** , so total is  **$O(k \log n)$** .
3. Return the registration IDs from popped gardens.

### Time Complexity

$O(n + k \log n)$

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### Problem B: Gardens Above a Score x (Max-Heap)

#### What's the Task?

Now the gardens are in a **Max-Heap**. We need to return all gardens with a score **strictly greater than x**.

#### How to Do It?

- Start at the top (index 0).
- If the current score  $> x$ :
  - Save its ID.
  - Recursively check the left and right children.
- If score  $\leq x$ : **✗** Skip the whole subtree.

#### Pseudocode

PROCEDURE gardensAboveX

BEGIN

INPUT:

; heap is a structure array/list in which each has a current score  
data member

List[1..n] heap

Number x, index

OUTPUT: number of gardens above a particular garden



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IF index >= n THEN
    return List [1 .. n] ; empty list
ENDIF

current INTEGER

SET current = heap[index]

IF current.score <= x THEN
    return List[1 .. n] ; empty list i.e. skipping this branch
ENDIF

SET result = [current.registration_id]

left Integer, right Integer

SET left = 2 * index + 1

SET right = 2 * index + 2

SET result = result + gardensAboveX(heap, x, left) +
gardensAboveX(heap, x, right)

OUTPUT result

END

```

### Time Complexity

$O(n_x)$  — where  $n_x$  is the number of gardens with score  $> x$ .

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### Quick Notes

### Useful Tricks

- **Problem A:** Heap is fast for top-k.
- **Problem B:** Heap helps you skip bad branches quickly.

### Edge Cases

- If  $k > \text{total gardens}$ : just return all.
- If all scores  $\leq x$ : return empty.
- Empty list? Still works — just return nothing