

## Crypto-mid QP - Mid semester examination answers

Cryptography (Birla Institute of Technology and Science, Pilani)

## Birla Institute of Technology & Science – Pilani Hyderabad Campus 2<sup>nd</sup> Semester 2019-2020

Cryptography (BITS F463) – Mid Sem Test (Regular)

Date: 05.03.2020 Weightage: 30% Duration: 1 hr. 30 min. Type: Closed Book

**Instructions:** Answer all questions. Answer all parts of a question consecutively. No of pages: 2

Q1. (a) Consider the Playfair cipher. Given the keyword KEYWORD, encrypt the message "Why, don't you?" (b) Consider the Hill cipher encryption of the form

$$\begin{bmatrix} y1 \\ y2 \end{bmatrix} = \begin{bmatrix} a & b \\ c & d \end{bmatrix} \begin{bmatrix} x1 \\ x2 \end{bmatrix} \mod 26.$$

Encipher the word MISSISSIPPI, with K as the padding character and given a = 22, b = 13, c = 11 and d = 5.

- (c) Given the plaintext MICHIGAN TECHNOLOGICAL UNIVERSITY and the keyword as HOUGHTON, find the outcome of Vigenere cipher encryption.
- (d)(i) Consider the rail fence cipher with key = 4. Find the ciphertext for the plaintext "This is a secret message".
- (ii) Decrypt the message **DEZCDZGAOIG** with depth = 3 using rail fence decryption.
- (e) Monoalphabetic substitution cipher can be constructed using alphabet mixing via columnar transposition. The steps are (i) The letters from the keyword form the headings of the columns, and the remaining letters of the alphabet fill in order in the rows below (ii) Mixing is achieved by transcribing columns (i.e., transcribing columns left-to-right gives the substitution). If the keyword is CORNELL, use this scheme to encrypt the message "FAR ABOVE CAYUGA'S WATERS".

(4+4+4+4+4=20 marks)

- Q2. (a) In cryptography congruence relations play a major role. Solve the following using congruences. (i) 5<sup>-1</sup> mod 7 (ii)  $23x \equiv -9 \mod 60$  (iii)  $5x \equiv 1 \mod 12$  (iv)  $x \equiv 13/8 \mod 29$  (v) Show that the number 345,546,711 is divisible by 9 using congruences.
- (b) Consider the finite field  $GF(2^m)$ . Given the polynomials  $P1=(x^2+x+1)$  and P2=(x+1), show that the results of P1+P2and P1-P2 will be the same if we perform modulo 2 arithmetic.
- (c) Let us assume that the equation xyz = 1 holds in a group G. (i) Does it follow that yzx = 1?(ii) Does it follow that yxz = 1? Reason your answer.
- (d) Let G be the set of all  $2\times 2$  matrices (as shown below) where a, b, c, d are integers modulo 2, such that ad  $-bc \neq 0$ . Use matrix multiplication as the operation in G to show that G is a group of order 6. Also show the elements.

(e) This problem pertains to block cipher modes of operations and their security aspects. Match every mode of operation with its correct security feature in the following table.

Mode	Security aspect	
1. ECB	A. encryption would not be resistant to chosen-plain text attacks if IV is known	
2. CBC	B. we should encrypt ciphertext data from the previous round	
3. CFB	C. repetition of encrypting the initialization vector may produce the same state that has occurred before	
4. OFB	D. nonce plays the same role as initialization vector	
5. CTR	E. Created cipher text is not blurred	

(4+4+4+4+4=20 marks)

<b>Q3.</b> (a)(i) Consider the AES implementation. Find $w[0]$ , $w[1]$ , $w[2]$ , $w[3]$ and $w[5]$ given $w[4] = (E2, 32, FC, F1)$ .
Given the 128 bit key (in hex): 54 68 61 74 73 20 6D 79 20 4B 75 6E 67 20 46 75.
(ii) Consider the 128-bit key implementation of Advanced Encryption Standard. The state matrix after the AddRoundKey
operation in the 10 <sup>th</sup> round is as shown below. What is the ciphertext?

29 57 40 1A C3 14 22 02 50 20 99 D7 5F F6 B3 3A

(b) Consider the following simple implementation of DES. Let the plaintext be the string  $0010\ 1000$ . Let the 10 bit key be 1100011110. The final result of the encryption is  $1000\ 1010$ . The IP and IP<sup>-1</sup> functions are as defined below. Find the respective outputs of the transformations IP and IP<sup>-1</sup>.

IP Table	IP <sup>-1</sup> Table
26314857	41357286

(c) Consider a snapshot of the 16<sup>th</sup> round of the standard DES implementation. Fill in the blanks in the following table:

 $L_{16} = 0100\ 0011\ 0100\ 0010\ 0011\ 0010\ 0011\ 0100$   $R_{16} = 0000\ 1010\ 0100\ 1100\ 1101\ 1001\ 1001\ 0101$ Outcome of final permutation:  $IP^{I} = \underline{\hspace{1cm}}$ Let the final ciphertext C = 85E813540F0AB405

(d) Fill in the blanks in the following table on 3DES implementation with 2keys.

Triple-DES is just DES with two \_\_\_\_bit keys applied. Given a plaintext message, the first key is used to \_\_\_\_the message. The second key is used to \_\_\_\_ the encrypted message. The twice-scrambled message is then \_\_\_\_ again with the \_\_\_\_ key to yield the final cipher text. This three-step procedure is called triple-DES.

Triple-DES is just DES done three times with two keys used in a particular order. (Triple-DES can also be done with three separate keys instead of only two. In either case the resultant key space is about  $2^{112}$ .)

(e) The following table shows the steps in a generic stream cipher system. Fill in the blanks.

Start with a secret key called ("\_\_\_\_\_"); Generate a \_\_\_\_\_. The i-th bit/byte of keying stream is a function of the \_\_\_\_\_ and the first \_\_\_\_\_ ciphertext bits. Combine the stream with the \_\_\_\_\_ to produce the ciphertext (typically by XOR)

(4+4+4+4+4=20 marks)