



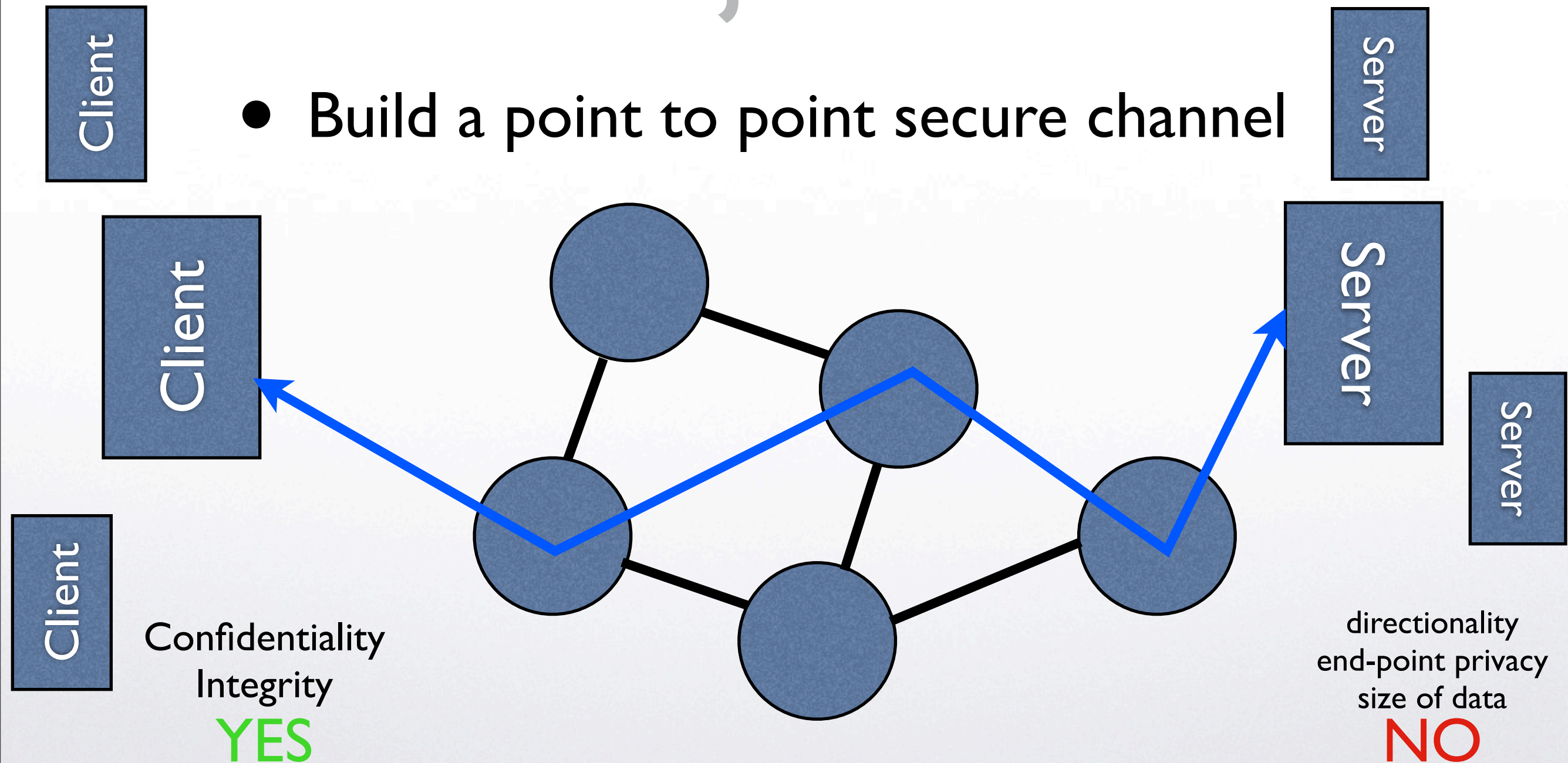
SSL/TLS/X.509

Aggelos Kiayias



Objective

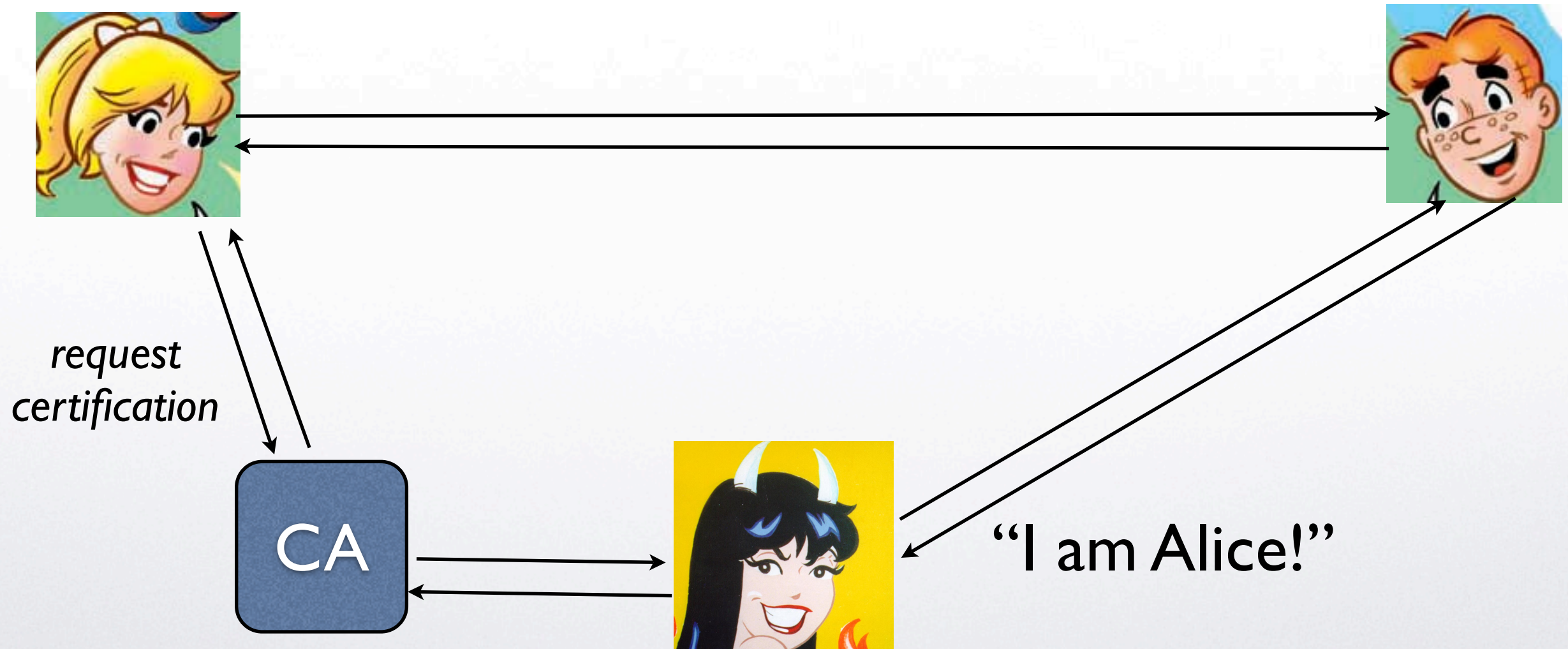
- Build a point to point secure channel





Identification Problem

“I am Alice!”





SSL/TLS

- Secure Sockets Layer: Developed by Netscape.
- SSL Version 3 released in 1996.
- Substituted by TLS in 1999: (Transport Layer Security). Standardized by IETF. published as RFC 2246.
- On top of TCP/IP, below application protocols.



TLS

- The Record protocol
 - encapsulates higher level protocols
- The Handshake protocol
 - enables authentication and session key establishment.



TLS Connections Parameters

- Connection end (client, server)
- Bulk encryption algorithm
- MAC algorithm
- Compression algorithm
- Master secret (48 bytes)
- Client random (32 byte)
- Server random (32 byte)

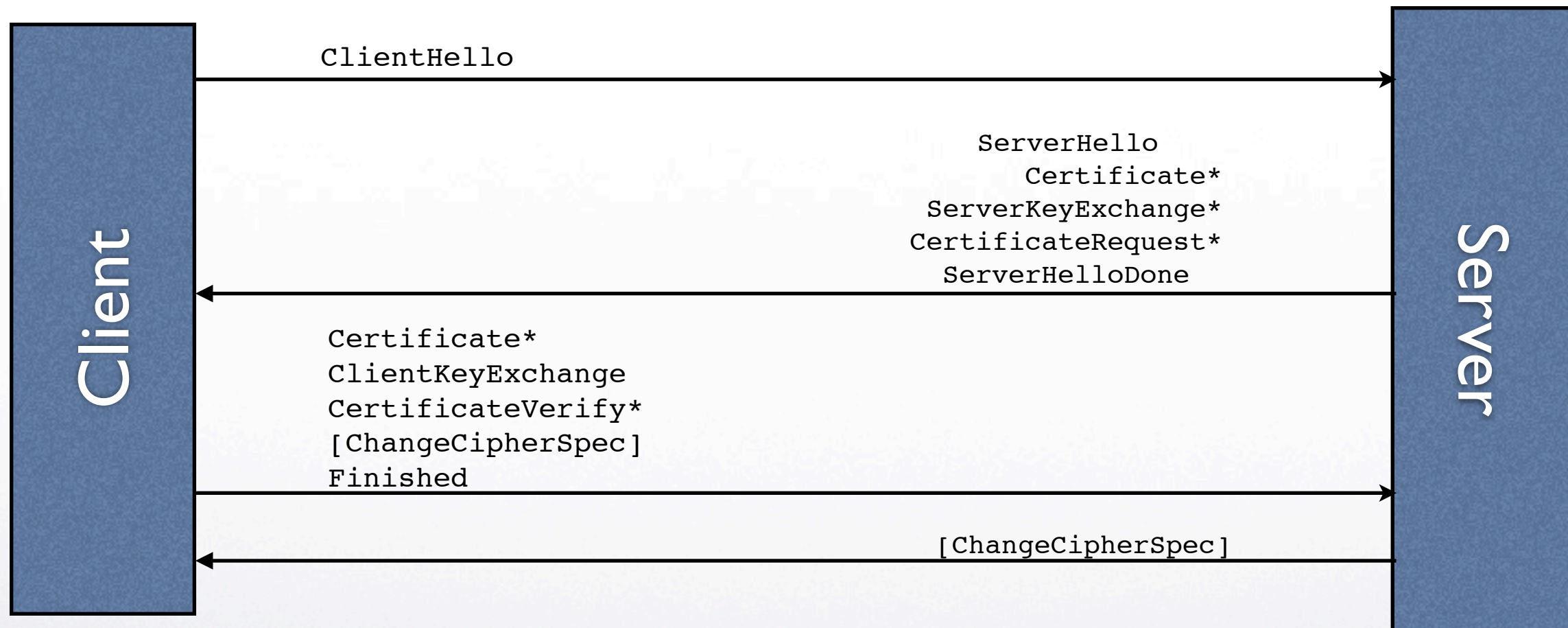


TLS Handshake

- Identities can be authenticated using public-key means. Authentication is optional but usually required at least for server side.
- Session-key is generated by public-key techniques.
- Special protocol modifications allow robustness against modification and man-in-the-middle attacks.



TLS Handshake



* = optional or situation-dependent



TLS Handshake, II

- The `ClientHello` message:
 - Random 28 byte string.
 - Client Timestamp
 - SessionID array (can be empty / it is used for resuming a previous session)
 - Ciphersuite list.
 - Compression list.

```
struct {  
    ProtocolVersion client_version;  
    Random random;  
    SessionID session_id;  
    CipherSuite cipher_suites<2..2^16-1>;  
    CompressionMethod compression_methods<1..2^8-1>;  
} ClientHello;
```



Example - ClientHello

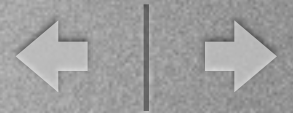
```
ClientVersion 3,1
ClientRandom[32]
SessionID: None (new session)
Suggested Cipher Suites:
  TLS_RSA_WITH_3DES_EDE_CBC_SHA
  TLS_RSA_WITH_DES_CBC_SHA
Suggested Compression Algorithm: NONE
```

format:

TLS_[handshake specs]_WITH_[record specs]

RSA for key exchange
3DES symmetric cipher in EDE mode
CBC encryption mode
SHA = hash algorithm

as above but DES is used.
this is not as secure
(will only use a 56 bit key as opposed to
168 bits for 3DES)



TLS Handshake III

- The ServerHello message:

taking client's suggestions into account

highest version
strongest protocol

Structure of this message:

```
struct {  
    ProtocolVersion server_version;  
    Random random;  
    SessionID session_id;  
    CipherSuite cipher_suite;  
    CompressionMethod compression_method;  
} ServerHello;
```

of the same format as of client's

one choice

will either match
client's suggestion (resume)
indicate a new session_id
or will be NULL



Example - ServerHello

```
Version 3,1  
ServerRandom[32]  
SessionID:  
bd608869f0c629767ea7e3ebf7a63bdcffb0ef58b1b941e6b0c044acb6820a77  
Use Cipher Suite:  
TLS_RSA_WITH_3DES_EDE_CBC_SHA  
Compression Algorithm: NONE
```




TLS Handshake IV

- The Certificate message: the server will provide its certificate to authenticate its name on the public-key it provides to the client.

```
opaque ASN.1Cert<1..224-1>;
```

```
struct {  
    ASN.1Cert certificate_list<0..224-1>;  
} Certificate;
```

chain of certificates

The end entity certificate's public key (and associated restrictions) MUST be compatible with the selected key exchange algorithm.



server [certificate]

<u>Key Exchange Alg.</u>	<u>Required Certificate Key Type</u>
RSA RSA_PSK	RSA public key; the certificate MUST allow the key to be used for encryption (the keyEncipherment bit MUST be set if the key usage extension is present).
DHE_RSA ECDHE_RSA	RSA public key; the certificate MUST allow the key to be used for signing (the digitalSignature bit MUST be set if the key usage extension is present) with the signature scheme and hash algorithm that will be employed in the server key exchange message. Note: ECDHE_RSA is defined in [TLSECC].
DHE_DSS	DSA public key; the certificate MUST allow the key to be used for signing with the hash algorithm that will be employed in the server key exchange message.
DH_DSS DH_RSA	Diffie-Hellman public key; the keyAgreement bit MUST be set if the key usage extension is present.



[serverkeyexchange]

- `ServerKeyExchange` : will contain public key information in case the information in the `Certificate` message is not sufficient (or this message has not been provided at all).



[serverkeyexchange]

required when TLS
follows:

- DHE_DSS
- DHE_RSA
- DH_anon

In this case it will
contain $\langle p, g, g^x \bmod p \rangle$

signed with DSS or RSA signing
algorithm with the public-key
that was provided in the [certificate]
except for DH_anon where no signature is given

illegal when TLS
follows:

- RSA
- DH_DSS
- DH_RSA

in these cases
the [certificate]
has enough info
for the key-exchange
to complete



[certificaterequest]

- `CertificateRequest` : will prompt the client to send a certificate to authenticate itself (typically not used).



TLS Handshake, V

- The `ClientKeyExchange` message

In case of RSA, client encrypts a random 46 byte string with the public-key of the server

In case of DH, client prepares his public-key according to the exchange.

```
struct {  
    select (KeyExchangeAlgorithm) {  
        case rsa: EncryptedPreMasterSecret;  
        case diffie_hellman: ClientDiffieHellmanPublic;  
    } exchange_keys;  
} ClientKeyExchange;
```




TLS Handshake, VI

Client Authentication

- `CertificateRequest` : server will use this message to request a certificate-based authentication from the client.
- `Certificate` : response to a `CertificateRequest` message. This will be sent before `ClientKeyExchange`
- `CertificateVerify` : if client's certificate has signing capability, this message will be a digital signature of all handshake messages so far. This will be sent after `ClientKeyExchange`



TLS Handshake VII

The Master-Secret (48 bytes):

```
master_secret = PRF(pre_master_secret, "master secret",  
                    ClientHello.random + ServerHello.random)
```

Computed depending on key exchange method

The Finished message:

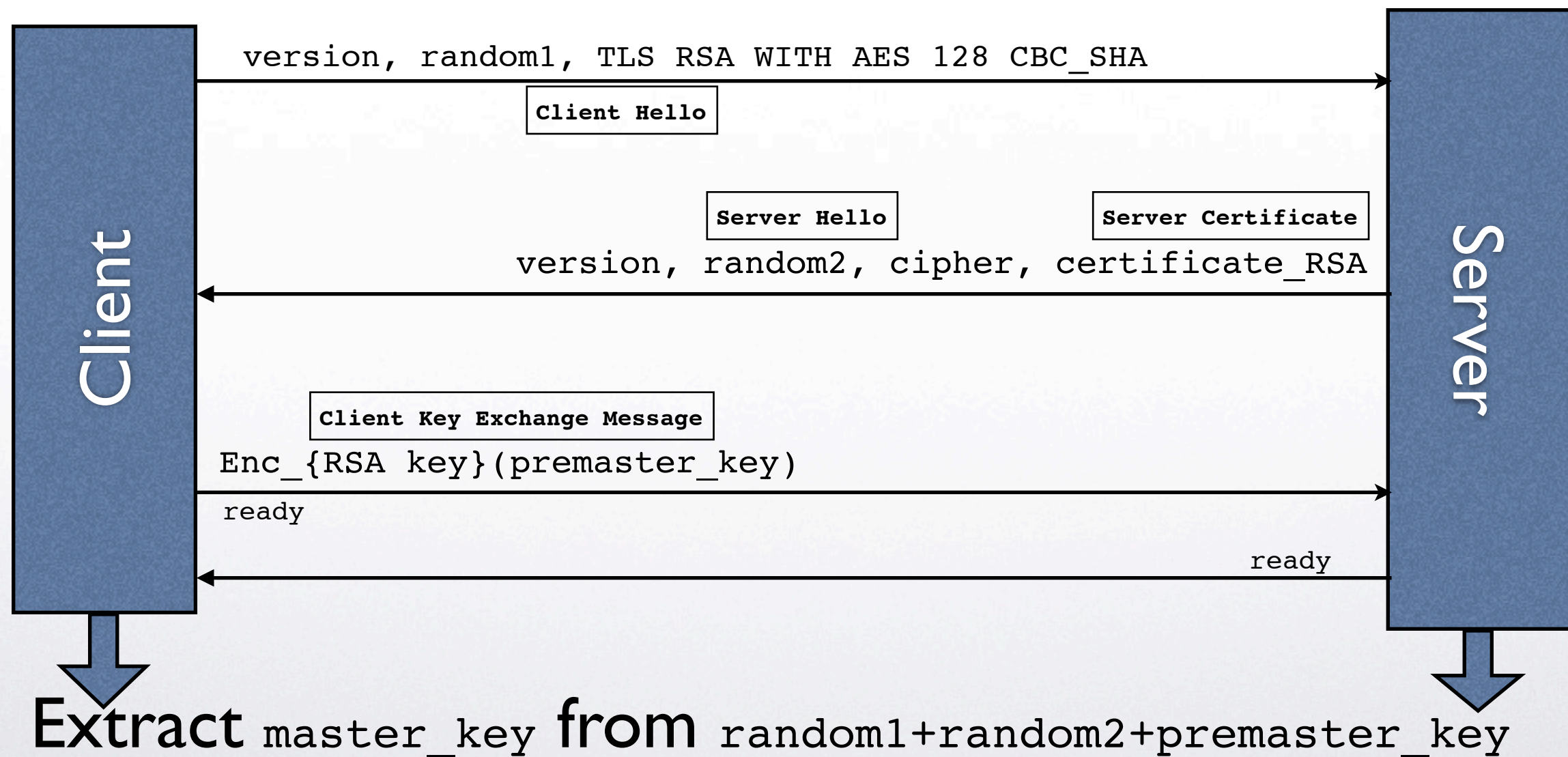
```
PRF(master_secret, finished_label, MD5(handshake_messages) +  
    SHA-1(handshake_messages)) [0..11];
```

Pseudorandom Function based on SHA-1 xor MD5



TLS Handshake, VIII

EXAMPLE





Summary :TLS Authentication

- Anonymous (neither client nor server provide certificates) - rare.
- Server-side - common.
- Both client and server - rare.



Alert sub-protocol

- RFC 2246.
- When things go wrong an alert is generated and either session ends or the recipient is given the opportunity to continue.

bad_certificate

certificate_expired

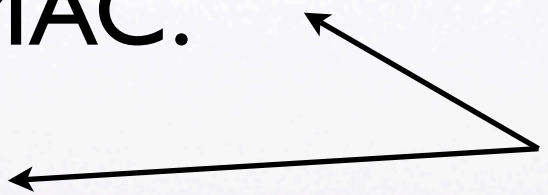
handshake_failure

unsupported_certificate

unknown_ca



Record Protocol

- takes messages to be transmitted and
 - fragments data into manageable blocks.
 - (optionally) compresses data.
 - applies a MAC.
 - encrypts.
 - transmits while maintaining proper order.
- Key derived from
master_key and prf
- 



Certificates

- Solving the authentication problem:

?

version, random2, cipher, SignCA(PK), PK

How do you know that a certain public key you receive is really coming from the source?



Server



What is a certificate?



What is a certificate?

- A digital signature on a public-key + other info.



What is a certificate?

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- What should be contained into the 'info'?



What is a certificate?

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- What should be contained into the 'info'?
- Who is the signer?



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- A digital signature on a public-key + other info.
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- How do you verify the signature?



What is a certificate?

- A digital signature on a public-key + other info.
- What should be contained into the 'info'?
- Who is the signer?
- How do you verify the signature?
- How do you know that you have the correct signer's public-key?



Certification

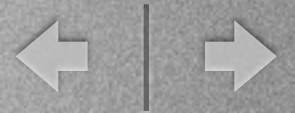
- Hierarchical certification structure.
 - Root certification authorities sign public-keys of lower-level authorities and so on till a public-key used for a session-key exchange is created.
- The web-of-trust (PGP).
 - “I sign your key you sign mine.”



X.509 Certificates

- Internet standard since 1988.
- Hierarchical.

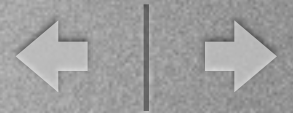
<http://www.ietf.org/rfc/rfc3280.txt>



Structure of X.509 Certificates

Version
Serial Number
Algorithm / Parameters
Issuer
Period of Validity: not before date not after date
Subject
Algorithm/ Parameters/ Key
x509v3 extensions
...
<i>Signature</i>

**X.509
does not
specify
cryptographic
algorithms**



Example

Certificate:

Data:

Version: 3 (0x2)

Serial Number:

13:86:35:4d:1d:3f:06:f2:c1:f9:65:05:d5:90:1c:62

Signature Algorithm: sha1WithRSAEncryption

Issuer: C=US, O=VISA, OU=Visa International Service Association, CN=Visa eCommerce Root

Validity

Not Before: Jun 26 02:18:36 2002 GMT

Not After : Jun 24 00:16:12 2022 GMT

Subject: C=US, O=VISA, OU=Visa International Service Association, CN=Visa eCommerce Root

Subject Public Key Info:

Public Key Algorithm: rsaEncryption

RSA Public Key: (2048 bit)

Modulus (2048 bit):

00:af:57:de:56:1e:6e:a1:da:60:b1:94:27:cb:17:
db:07:3f:80:85:4f:c8:9c:b6:d0:f4:6f:4f:cf:99:
d8:e1:db:c2:48:5c:3a:ac:39:33:c7:1f:6a:8b:26:
3d:2b:35:f5:48:b1:91:c1:02:4e:04:96:91:7b:b0:
33:f0:b1:14:4e:11:6f:b5:40:af:1b:45:a5:4a:ef:
7e:b6:ac:f2:a0:1f:58:3f:12:46:60:3c:8d:a1:e0:
7d:cf:57:3e:33:1e:fb:47:f1:aa:15:97:07:55:66:
a5:b5:2d:2e:d8:80:59:b2:a7:0d:b7:46:ec:21:63:

Hierarchical name spec

C=Country, L=Location, O=Organization,
OU=Organization Unit, CN=Common Name

**A self-signed
root certificate**



Example continued

4a:ea:db:df:72:38:8c:f3:96:bd:f1:17:bc:d2:ba:
3b:45:5a:c6:a7:f6:c6:17:8b:01:9d:fc:19:a8:2a:
83:16:b8:3a:48:fe:4e:3e:a0:ab:06:19:e9:53:f3:
80:13:07:ed:2d:bf:3f:0a:3c:55:20:39:2c:2c:00:
69:74:95:4a:bc:20:b2:a9:79:e5:18:89:91:a8:dc:
1c:4d:ef:bb:7e:37:0b:5d:fe:39:a5:88:52:8c:00:
6c:ec:18:7c:41:bd:f6:8b:75:77:ba:60:9d:84:e7:
fe:2d

Exponent: 65537 (0x10001)

X509v3 extensions:

X509v3 Basic Constraints: critical
CA:TRUE

X509v3 Key Usage: critical
Certificate Sign, CRL Sign

X509v3 Subject Key Identifier:

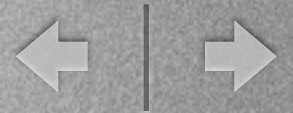
15:38:83:0F:3F:2C:3F:70:33:1E:CD:46:FE:07:8C:20:E0:D7:C3:B7

Signature Algorithm: sha1WithRSAEncryption

5f:f1:41:7d:7c:5c:08:b9:2b:e0:d5:92:47:fa:67:5c:a5:13:
c3:03:21:9b:2b:4c:89:46:cf:59:4d:c9:fe:a5:40:b6:63:cd:

.....

note:
public-exponent,
fixed / small





Binary Encoding

- The elements of certificate are encoded using DER (distinguished encoding rules)

Items are stored in TLV format : triplets <Type,Length,Value>

```
1.      30 23          ; SEQUENCE (23 Bytes)
2.      | 31 0f          ; SET (f Bytes)
3.      | | 30 0d          ; SEQUENCE (d Bytes)
4.      | | | 06 03          ; OBJECT_ID (3 Bytes)
5.      | | | | 55 04 03
6.      | | | | | 2.5.4.3 Common Name (CN)
7.      | | | | 13 06          ; PRINTABLE_STRING (6 Bytes)
8.      | | | | | 54 65 73 74 43 4e          ; TestCN
9.      | | | | | "TestCN"
10.     | | 31 10          ; SET (10 Bytes)
11.     | | | 30 0e          ; SEQUENCE (e Bytes)
12.     | | | | 06 03          ; OBJECT_ID (3 Bytes)
13.     | | | | | 55 04 0a
14.     | | | | | 2.5.4.10 Organization (O)
15.     | | | | 13 07          ; PRINTABLE_STRING (7 Bytes)
16.     | | | | | 54 65 73 74 4f 72 67          ; TestOrg
17.     | | | | | "TestOrg"
```

example from <http://msdn.microsoft.com>



Parsing DER encoding

```
openssl asn1parse -in certforstar.engr.uconn.edu.pem
```

```
0:d=0  hl=4 l=1375 cons: SEQUENCE
4:d=1  hl=4 l=1095 cons: SEQUENCE
8:d=2  hl=2 l=  3 cons: cont [ 0 ]
10:d=3  hl=2 l=  1 prim: INTEGER           :02
13:d=2  hl=2 l=  7 prim: INTEGER           :2B906692AF6A3F
22:d=2  hl=2 l= 13 cons: SEQUENCE
24:d=3  hl=2 l=  9 prim: OBJECT             :sha1WithRSAEncryption
35:d=3  hl=2 l=  0 prim: NULL
37:d=2  hl=3 l=202 cons: SEQUENCE
40:d=3  hl=2 l= 11 cons: SET
42:d=4  hl=2 l=  9 cons: SEQUENCE
44:d=5  hl=2 l=  3 prim: OBJECT             :countryName
49:d=5  hl=2 l=  2 prim: PRINTABLESTRING    :US
53:d=3  hl=2 l= 16 cons: SET
55:d=4  hl=2 l= 14 cons: SEQUENCE
57:d=5  hl=2 l=  3 prim: OBJECT             :stateOrProvinceName
62:d=5  hl=2 l=  7 prim: PRINTABLESTRING    :Arizona
71:d=3  hl=2 l= 19 cons: SET
73:d=4  hl=2 l= 17 cons: SEQUENCE
75:d=5  hl=2 l=  3 prim: OBJECT             :localityName
80:d=5  hl=2 l= 10 prim: PRINTABLESTRING    :Scottsdale
92:d=3  hl=2 l= 26 cons: SET
94:d=4  hl=2 l= 24 cons: SEQUENCE
96:d=5  hl=2 l=  3 prim: OBJECT             :organizationName
101:d=5  hl=2 l= 17 prim: PRINTABLESTRING    :GoDaddy.com, Inc.
120:d=3  hl=2 l= 51 cons: SET
122:d=4  hl=2 l= 49 cons: SEQUENCE
124:d=5  hl=2 l=  3 prim: OBJECT             :organizationalUnitName
129:d=5  hl=2 l= 42 prim: PRINTABLESTRING    :http://certificates.godaddy.com/repository
173:d=3  hl=2 l= 48 cons: SET
175:d=4  hl=2 l= 46 cons: SEQUENCE
```

FORMAT of each line

offset:

d=depth

hl=header length

l=content length

... followed by type of encoding

primitive - the contents represent the final payload.

constructed - the contents are other TLV values.

.....



Hash and Sign Paradigm

- Sign is a cryptographic operation that operates on fixed length inputs.
- Hashing allows the signing of arbitrary input lengths.
- What must be the required properties of the hash function?



Available Signature Algorithms

- Common choices:
 - sha1withRSA
 - sha1withDSA
 - (previously md5withRSA etc. md5 is now totally broken)



HTTP Secure

- Combination of HTTP protocol and TLS.
- Integration with browser :
 - CA public-keys come within the browser.
 - Browser issues warning if certificate is invalid or not trusted. Key points : (1) CommonName in certificate should match web-site DNS name. (2) Wildcards are allowed.
 - Typically only server side authentication is performed.



Colliding Certificates

- Based on collision attacks against hash functions the feasibility of producing MD5-based colliding certificates was demonstrated.
- What does this mean for SSL/TLS?

Arjen Lenstra and Xiaoyun Wang and Benne de Weger

Colliding X.509 Certificates

<http://eprint.iacr.org/2005/067>



Colliding Certificates

- In 2008 directed collisions were produced for SSL certificates => CA hacks feasible.

<http://www.win.tue.nl/hashclash/rogue-ca/>



Revoking Certificates

- What happens when a certified public-key is compromised?
- E.g., secret key is leaked and published online?
- What if a previously trusted entity is proved to be untrustworthy?



Microsoft Security Bulletin MS01-017

Erroneous VeriSign-Issued Digital Certificates Pose Spoofing Hazard

Originally posted: March 22, 2001

Updated: June 23, 2003

Summary

Who should read this bulletin:

All customers using Microsoft® products.

Impact of vulnerability:

Attacker could digitally sign code using the name "Microsoft Corporation".

Recommendation:

All customers should install the update discussed below.

<http://www.microsoft.com/technet/security/bulletin/MS01-017.msp>



Microsoft Security Bulletin MS01-017

- **Technical description:**

- In mid-March 2001, VeriSign, Inc., advised Microsoft that on January 29 and 30, 2001, it issued two VeriSign Class 3 code-signing digital certificates to an individual who fraudulently claimed to be a Microsoft employee. The common name assigned to both certificates is "Microsoft Corporation". The ability to sign executable content using keys that purport to belong to Microsoft would clearly be advantageous to an attacker who wished to convince users to allow the content to run.
- The certificates could be used to sign programs, ActiveX controls, Office macros, and other executable content. Of these, signed ActiveX controls and Office macros would pose the greatest risk, because the attack scenarios involving them would be the most straightforward. Both ActiveX controls and Word documents can be delivered via either web pages or HTML mails. ActiveX controls can be automatically invoked via script, and Word documents can be automatically opened via script unless the user has applied the [Office Document Open Confirmation Tool](#).

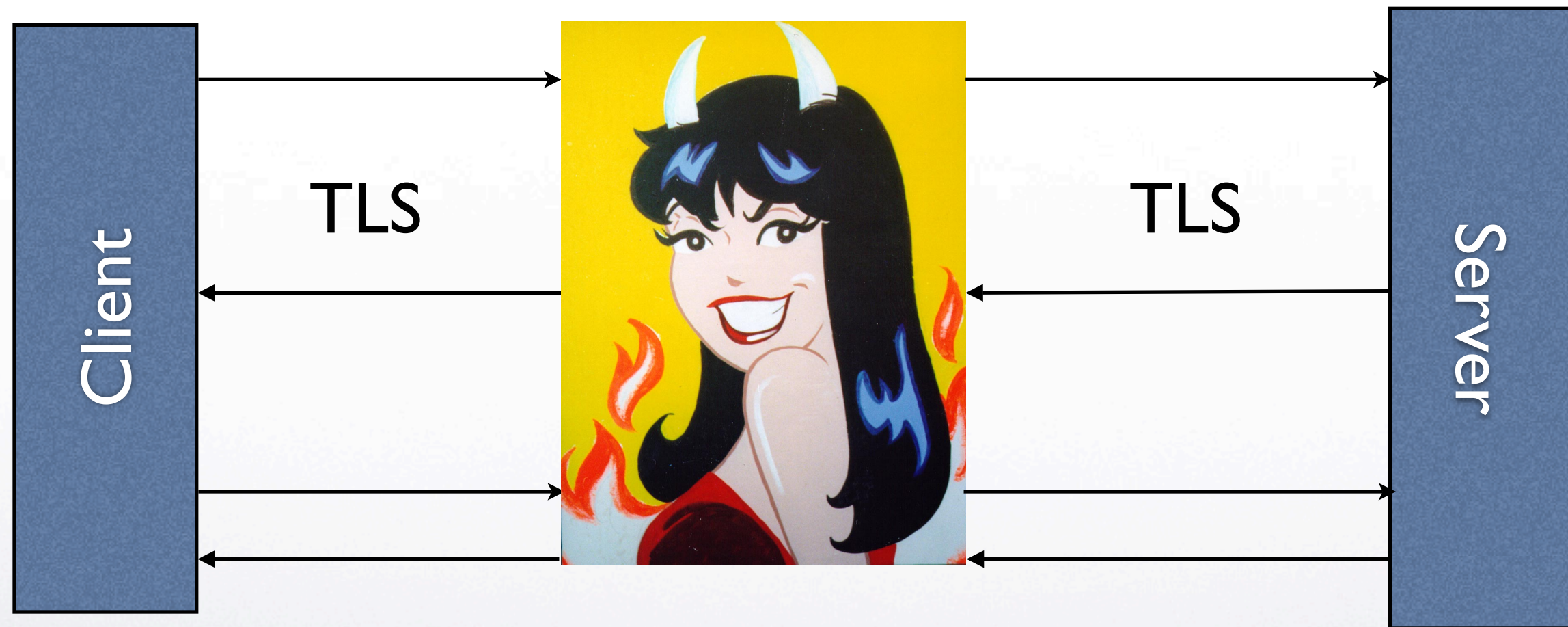


Fraudulent Certificates

- The attacker could use the certified public-key and sign code as “Microsoft Corporation”
- How to revoke such certificates?



“man-in-the-middle”



In *anonymous mode* TLS Authentication is totally vulnerable to a man-in-the-middle attack.



Version Rollback Attacks

(a man in the middle attack)

- When a client and server rollback to a previous version of a protocol despite both being capable.
- Can be a problem when client and server negotiate a version in the clear.
- Related example : rollback to DH_anon



certificate verification flaws

that have been demonstrated due to bad implementation

- Check validity of certificate but don't check common name.
- Check common name but fail to verify the whole chain of certificates.
- Check everything but allow non-CA signing certificates to be used for issuing certificates.



CA hacks

- Comodo and DigiNotar in 2011
- "several dozens of SSL certificates" issued fraudulently.



Certificate Revocation Lists

- A certificate revocation list (CRL) contains the serial numbers of valid certificates that should not be accepted.
- A CRL must be issued and signed by the CA that originally issued the revoked certificate.
- From where can you get a CRL?

try <http://crl.verisign.com/>



Usability aspects of Certificate Management

- From an internet discussion board:

subject: HELP

for some unknown reason, as of today, every file that I access,
every program that runs, every folder in explorer.exe that I open wants
to connect to some address at crl.verisign.com 198.49.161.200

I have done 3 restores as far back as last week.....no difference.....

I have search the registry and the hard drives.....

I have ran a virus scanner and a trojan detector.....

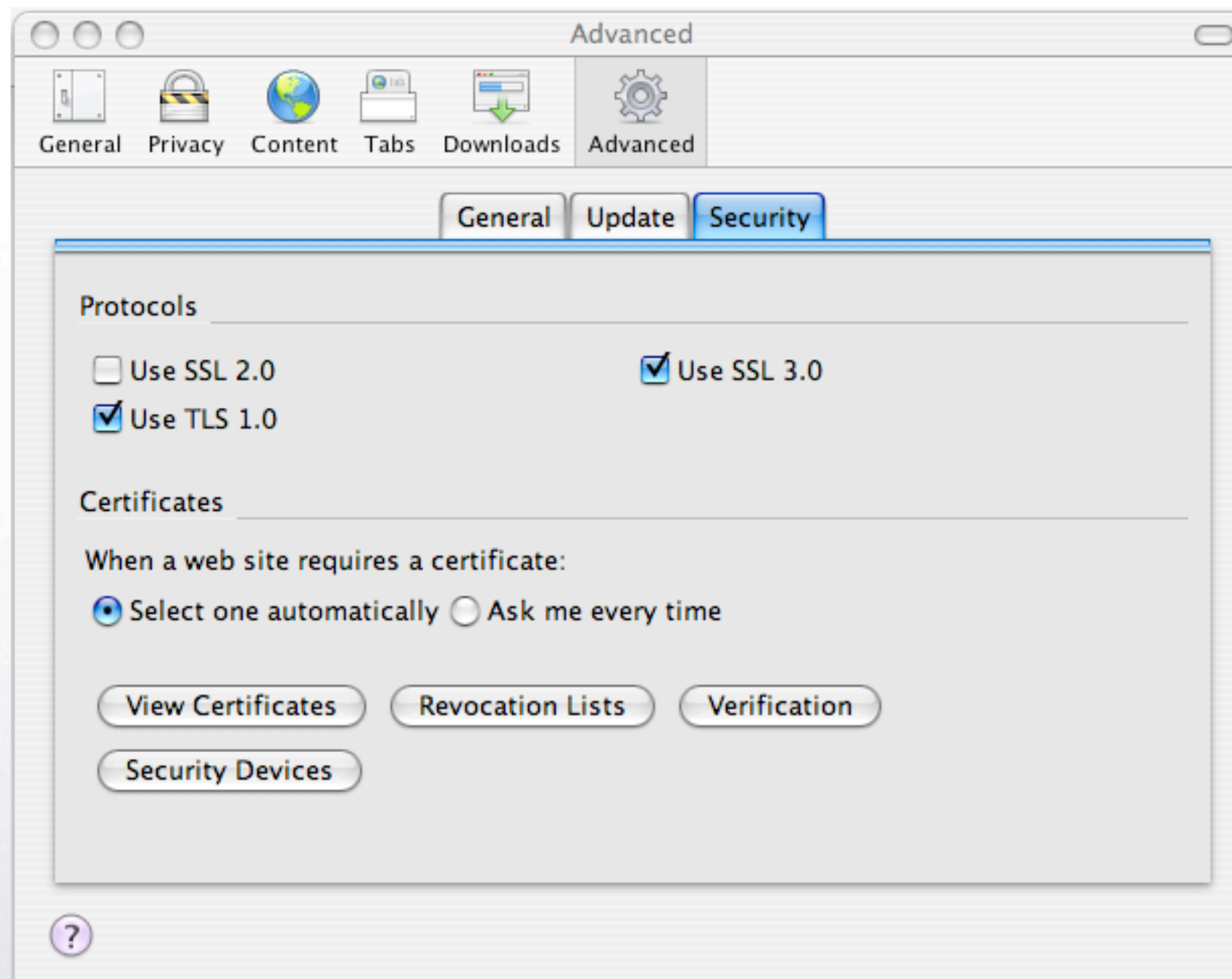


Distributing CRLs

- Certificates include inside the extension fields the URL of a **CRL distribution point**.
- Your system will be responsible to collect such CRLs.



A quick look at Firefox





OCSP

- Online Certificate Status Protocol.
- Idea: obviate the need for CRLs by verifying certificates online with the CA
- Has various advantages over CRLs.
- but requires CA to be online.
- privacy issues?



PKI

- Public-key Infrastructure
 - A structure that allows entities communicating over an insecure network to form authenticated and secure communication channels.
- Basic tool: digital certification of public-keys.
- Essential ingredients: certificate issuing criteria, trusted distribution, certificate revocation and expiration.



CCA Attacks

- In public-key encryption based key exchange protocols, the server acts as a decryption box.
- Even though it may not return the plaintext it may leak some information about it (e.g., through an error message).
- Such leakage has been exploited for attacks.



Bleichenbacher's PKCS #1 Attack

- Applicable to SSL v3.0
- A chosen ciphertext attack:
 - goal compute $c^{\{d\}} \bmod n$ for a given c (i.e., a decryption).
 - By exploiting only error-messages in SSL protocol!!
 - Many queries needed but possible ($>300,000$).