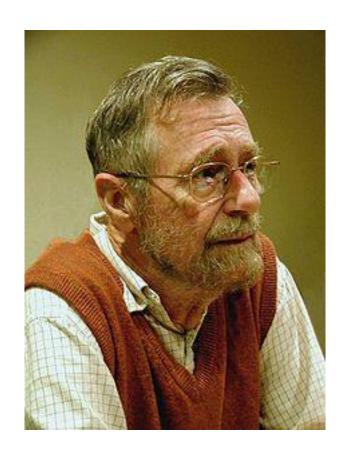
DIJKSTRA'S ALGORITHM

By Laksman Veeravagu and Luis Barrera

THE AUTHOR: EDSGER WYBE DIJKSTRA

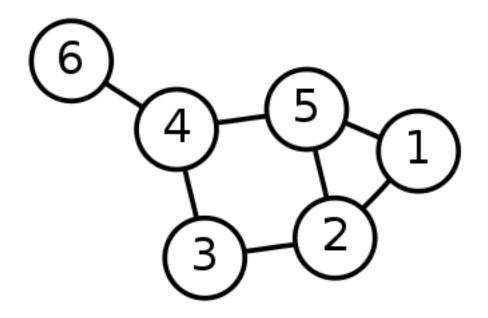


"Computer Science is no more about computers than astronomy is about telescopes."

http://www.cs.utexas.edu/~EWD/

SINGLE-SOURCE SHORTEST PATH PROBLEM

<u>Single-Source Shortest Path Problem</u> - The problem of finding shortest paths from a source vertex *v* to all other vertices in the graph.



DIJKSTRA'S ALGORITHM

<u>Dijkstra's algorithm</u> - is a solution to the single-source shortest path problem in graph theory.

Works on both directed and undirected graphs. However, all edges must have nonnegative weights.

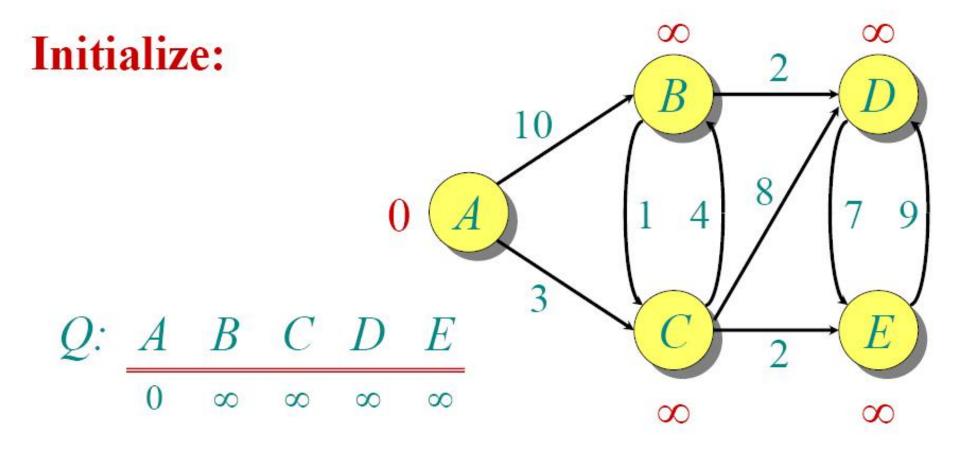
Approach: Greedy

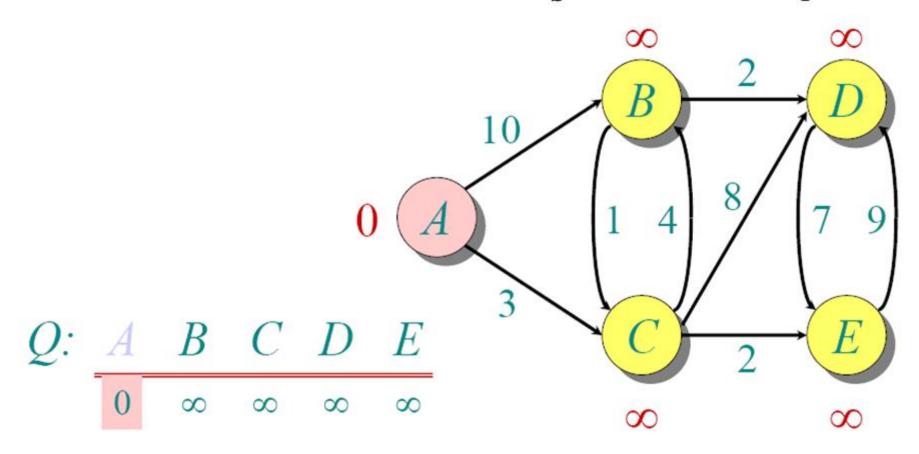
Input: Weighted graph G={E,V} and source vertex *v*∈V, such that all edge weights are nonnegative

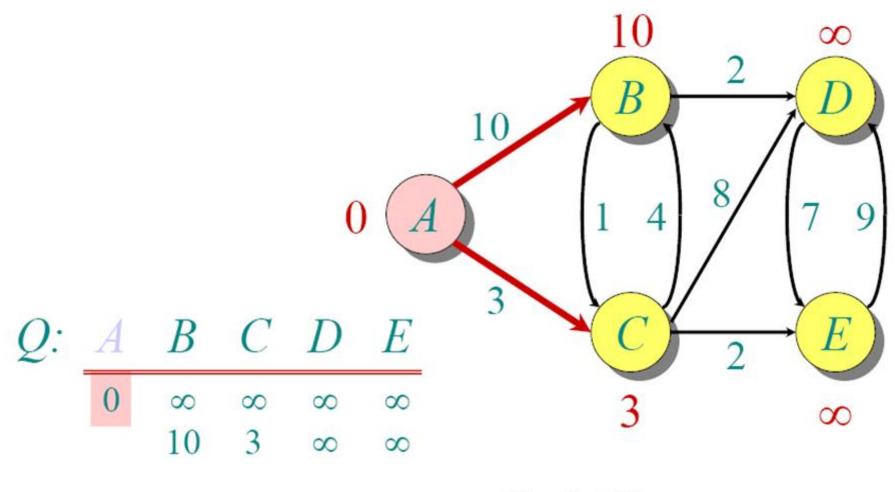
Output: Lengths of shortest paths (or the shortest paths themselves) from a given source vertex *v*∈V to all other vertices

DIJKSTRA'S ALGORITHM - PSEUDOCODE

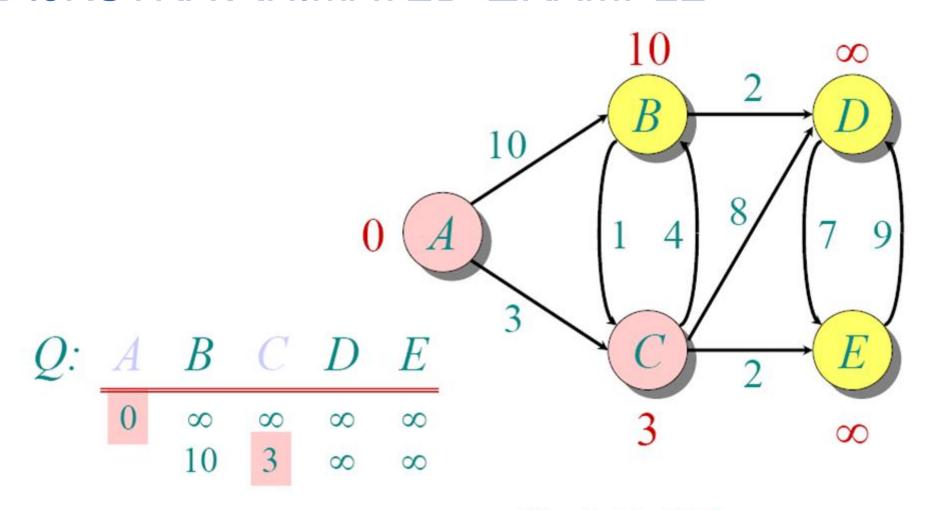
```
dist[s] \leftarrow o
                                      (distance to source vertex is zero)
for all v \in V - \{s\}
     do dist[v] \leftarrow \infty
                                     (set all other distances to infinity)
                                      (S, the set of visited vertices is initially empty)
S←Ø
                                      (Q, the queue initially contains all vertices)
Q←V
while Q ≠Ø
                                      (while the queue is not empty)
do u \leftarrow mindistance(Q,dist)
                                      (select the element of Q with the min. distance)
                                      (add u to list of visited vertices)
    S \leftarrow S \cup \{u\}
    for all v \in neighbors[u]
         do if dist[v] > dist[u] + w(u, v)
                                                        (if new shortest path found)
                then d[v] \leftarrow d[u] + w(u, v)
                                                        (set new value of shortest path)
                   (if desired, add traceback code)
return dist
```



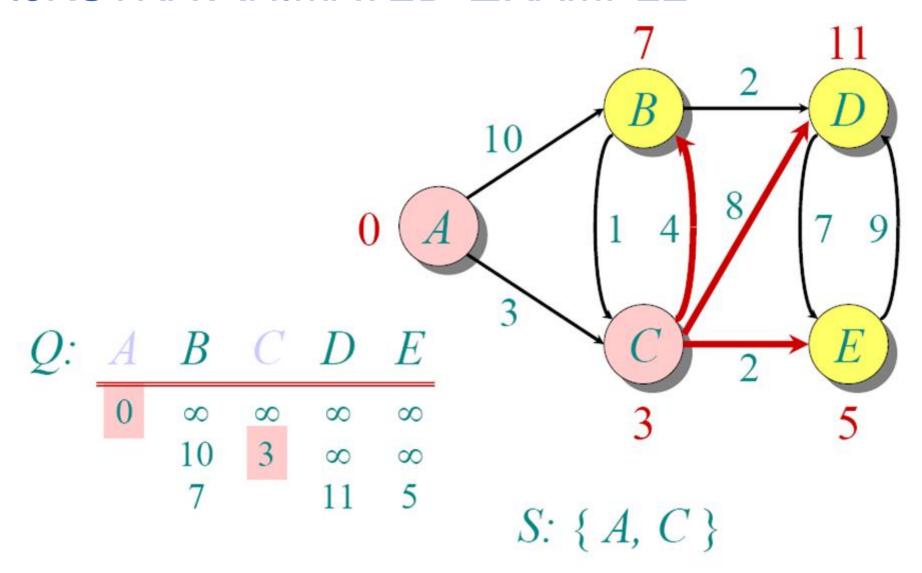


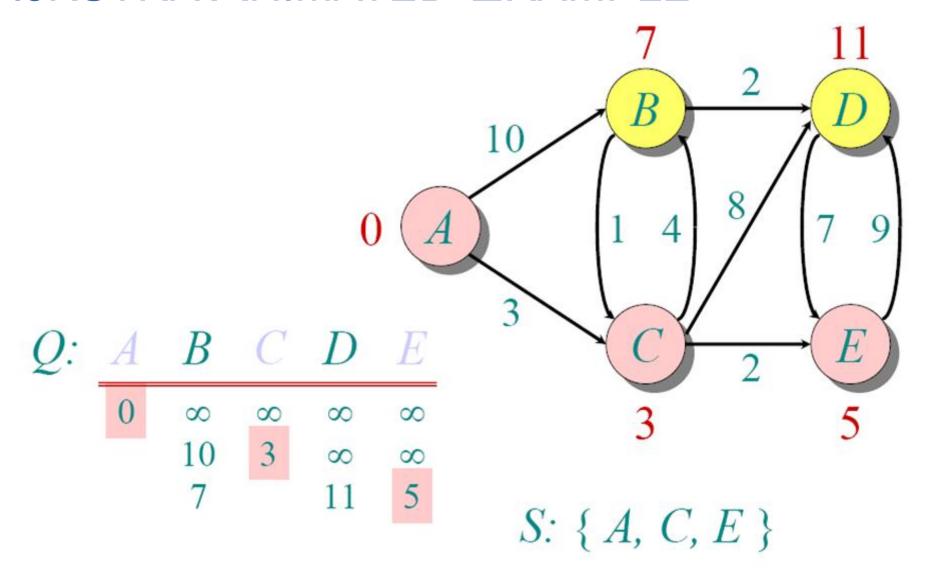


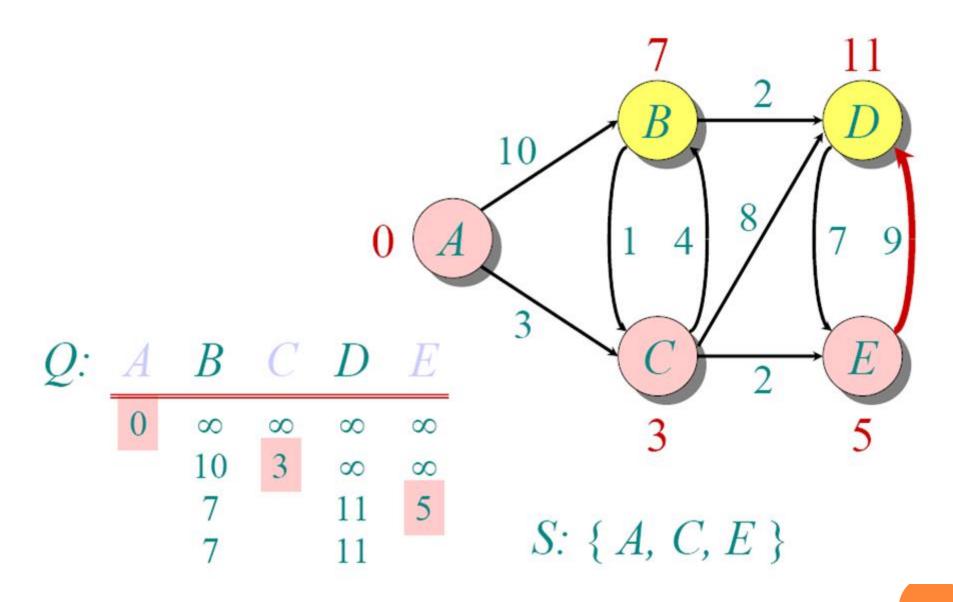
S: { A }

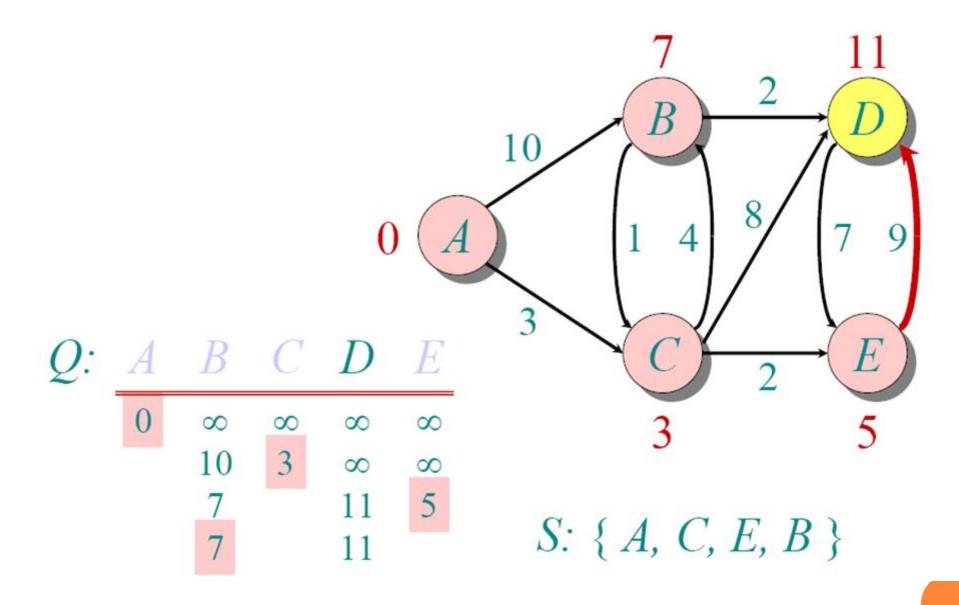


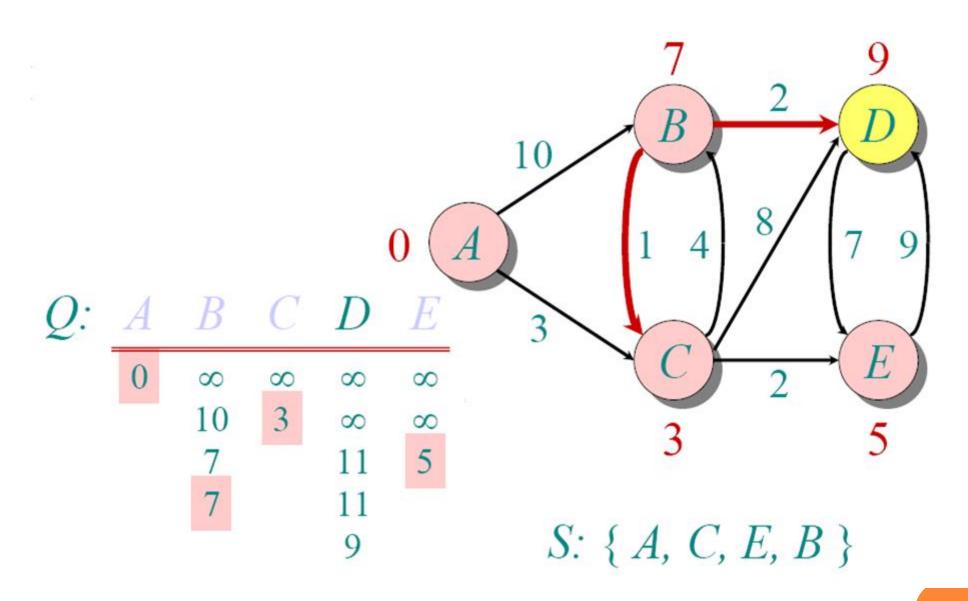
S: { A, C }

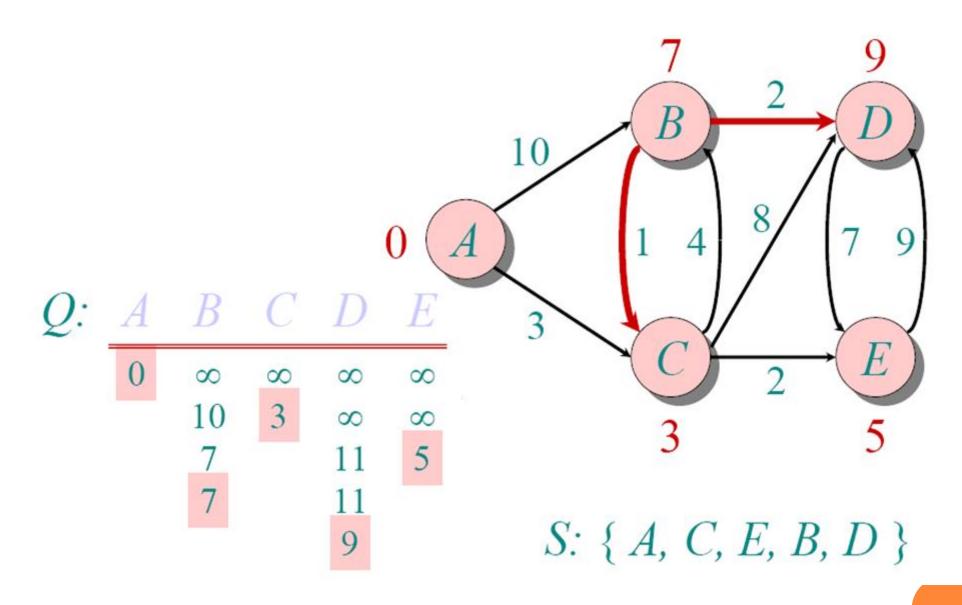












IMPLEMENTATIONS AND RUNNING TIMES

The simplest implementation is to store vertices in an array or linked list. This will produce a running time of

$$O(|V|^2 + |E|)$$

For sparse graphs, or graphs with very few edges and many nodes, it can be implemented more efficiently storing the graph in an adjacency list using a binary heap or priority queue. This will produce a running time of

$$O((|E|+|V|)\log |V|)$$

DIJKSTRA'S ALGORITHM - WHY IT WORKS

- As with all greedy algorithms, we need to make sure that it is a correct algorithm (e.g., it *always* returns the right solution if it is given correct input).
- A formal proof would take longer than this presentation, but we can understand how the argument works intuitively.
- If you can't sleep unless you see a proof, see the second reference or ask us where you can find it.

DIJKSTRA'S ALGORITHM - WHY IT WORKS

- To understand how it works, we'll go over the previous example again. However, we need two mathematical results first:
- **Lemma 1**: Triangle inequality If $\delta(u,v)$ is the shortest path length between u and v, $\delta(u,v) \leq \delta(u,x) + \delta(x,v)$
- Lemma 2: The subpath of any shortest path is itself a shortest path.
- The key is to understand why we can claim that anytime we put a new vertex in S, we can say that we already know the shortest path to it.
- Now, back to the example...

DIJKSTRA'S ALGORITHM - WHY USE IT?

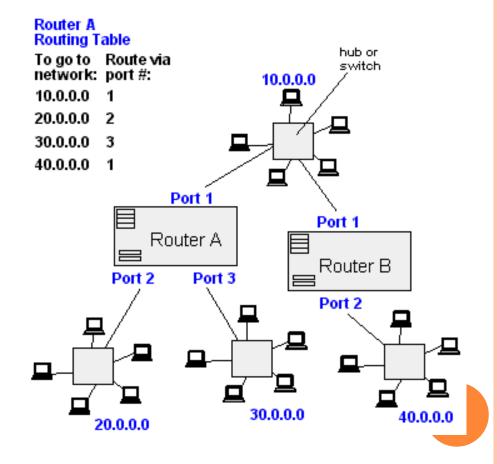
- As mentioned, Dijkstra's algorithm calculates the shortest path to every vertex.
- However, it is about as computationally expensive to calculate the shortest path from vertex u to every vertex using Dijkstra's as it is to calculate the shortest path to some particular vertex v.
- Therefore, anytime we want to know the optimal path to some other vertex from a determined origin, we can use Dijkstra's algorithm.

APPLICATIONS OF DIJKSTRA'S ALGORITHM

- Traffic Information Systems are most prominent use
- Mapping (Map Quest, Google Maps)
- Routing Systems

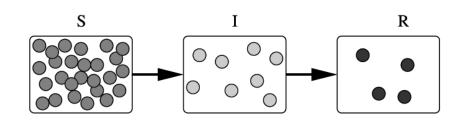
From Computer Desktop Encyclopedia

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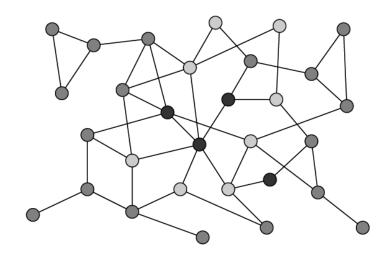


APPLICATIONS OF DIJKSTRA'S ALGORITHM

- One particularly relevant this week: epidemiology
- Prof. Lauren Meyers (Biology Dept.) uses networks to model the spread of infectious diseases and design prevention and response strategies.
- Vertices represent individuals, and edges their possible contacts. It is useful to calculate how a particular individual is connected to others.
- Knowing the shortest path lengths to other individuals can be a relevant indicator of the potential of a particular individual to infect others.



Network



REFERENCES

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 E. W. Dijkstra. (1959) A Note on Two Problems in Connection with Graphs. Numerische Mathematik, 1. 269-271.
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- Meyers, L.A. (2007) Contact network epidemiology: Bond percolation applied to infectious disease prediction and control. *Bulletin of the American Mathematical Society* **44**: 63-86.
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