



⑤ Var $x: \text{Int}$; ⑥ Var $x, y: \text{Int}$; ⑦ Var $x, y: \text{Int}$; ⑧ Var $x, y: \text{Int}$; ⑨ El ⑦ se puede escribir como el ⑧ para tener una asignación Multiple
 $[\sigma_0: (x \mapsto 4)]$ $[\sigma_0: (x \mapsto 2, y \mapsto 5)]$ $[\sigma_0: (x \mapsto 2, y \mapsto 5)]$ $[\sigma_0: (x \mapsto 2, y \mapsto 5)]$
 $x := 5;$ $x := x + y;$ $y := y + x;$ $y, x := y + x, x + y;$
 $[\sigma_1: (x \mapsto 5)]$ $[\sigma_1: (x \mapsto 7, y \mapsto 5)]$ $[\sigma_1: (x \mapsto 2, y \mapsto 10)]$ $[\sigma_1: (y \mapsto 10, x \mapsto 7)]$
 $x := y + y;$ $x := x + y;$
 $[\sigma_2: (x \mapsto 7, y \mapsto 10)]$ $[\sigma_2: (x \mapsto 12, y \mapsto 10)]$

⑩ Var $x, y: \text{Int}$; ⑪ Var $x, y: \text{Int}$; ⑫ Var $x, y: \text{Int}$; ⑬ El ⑩ programa no podría ejecutarse ya que no hay ningún caso que cumpla $x=y$
 $[\sigma_0: (x \mapsto 3, y \mapsto 1)]$ $[\sigma_0: (x \mapsto 100, y \mapsto 1)]$ $[\sigma_0: (x \mapsto 1, y \mapsto 1)]$ $[\sigma_0: (x \mapsto 1, y \mapsto 1)]$
 if $x \geq y \rightarrow$ ~~if $x \geq y \rightarrow$~~ if $x \geq y \rightarrow$ ~~if $x \geq y \rightarrow$~~
 $[\sigma_1: (x \mapsto 3, y \mapsto 1)]$ ~~$[\sigma_1:]$~~ $[\sigma_1: (x \mapsto 1, y \mapsto 1)]$ ~~$[\sigma_1:]$~~
 $x := 0;$ ~~$x := 0;$~~ $x := 0;$ ~~$x := 0;$~~
 $[\sigma_2: (x \mapsto 0, y \mapsto 1)]$ ~~$[\sigma_2:]$~~ $[\sigma_2: (x \mapsto 0, y \mapsto 1)]$ ~~$[\sigma_2:]$~~
~~if $x \leq y \rightarrow$~~ if $x \leq y$ ~~if $x \leq y \rightarrow$~~ if $x \leq y \rightarrow$
 ~~$[\sigma_1:]$~~ $[\sigma_1: (x \mapsto 100, y \mapsto 1)]$ ~~$[\sigma_1:]$~~ $[\sigma_1: (x \mapsto 1, y \mapsto 1)]$
 ~~$x := 2;$~~ $x := 2;$ ~~$x := 2;$~~ $x := 2;$
 ~~$[\sigma_2:]$~~ $[\sigma_2: (x \mapsto 2, y \mapsto 1)]$ ~~$[\sigma_2:]$~~ $[\sigma_2: (x \mapsto 2, y \mapsto 1)]$
 f_i f_i f_i f_i
 $[\sigma_3: (x \mapsto 0, y \mapsto 1)]$ $[\sigma_3: (x \mapsto 2, y \mapsto 1)]$ $[\sigma_3: (x \mapsto 0, y \mapsto 1)]$ $[\sigma_3: (x \mapsto 2, y \mapsto 1)]$

⑭ Var $i: \text{Int}$; ⑮ Var $i: \text{Int}$; ⑯ Var $i: \text{Int}$; ⑰ Var $i: \text{Int}$
 $[\sigma_0: (i \mapsto 4)]$ $[\sigma_0: (i \mapsto 400)]$ $[\sigma_0: (i \mapsto 4)]$ $[\sigma_0: (i \mapsto 0)]$
 do $i \neq 0 \rightarrow$ do $i \neq 0 \rightarrow$ do $i < 0 \rightarrow$ do $i \leq 0 \rightarrow$
 $[\sigma_1: (i \mapsto 4), \sigma_1^1: (i \mapsto 3), \sigma_1^2: (i \mapsto 2), \sigma_1^3: (i \mapsto 1)]$ $[\sigma_1: (i \mapsto 400)]$ ~~$[\sigma_1:]$~~ $[\sigma_1: \dots, \dots]$
 $i := i - 1$ $i := 0$ ~~$i := i - 1$~~ $i := i - 1;$
 $[\sigma_2: (i \mapsto 3), \sigma_2^1: (i \mapsto 2), \sigma_2^2: (i \mapsto 1), \sigma_2^3: (i \mapsto 0)]$ $[\sigma_2: (i \mapsto 0)]$ ~~$[\sigma_2:]$~~ $[\sigma_2: \dots, \dots]$
 od od od od
 $[\sigma_3: (i \mapsto 0)]$ $[\sigma_3: (i \mapsto 0)]$ $[\sigma_3: (i \mapsto 4)]$ ~~$[\sigma_3:]$~~ *programa no termina nunca*

⑱ Var $r: \text{Int}$ ⑲
 $[\sigma_0: (r \mapsto 3)]$
 do $r \neq 0 \rightarrow$
 if $r < 0 \rightarrow$
 ~~$[\sigma_1:]$~~
 ~~$r := r + 1;$~~
 ~~$[\sigma_2:]$~~
 if $r > 0 \rightarrow$
 $[\sigma_3: (r \mapsto 3), \sigma_3^1: (r \mapsto 2), \sigma_3^2: (r \mapsto 1)]$
 $r := r - 1;$
 $[\sigma_4: (r \mapsto 2), \sigma_4^1: (r \mapsto 1), \sigma_4^2: (r \mapsto 0)]$
 f_i
 $[\sigma_5: (r \mapsto 2), \sigma_5^1: (r \mapsto 1), \sigma_5^2: (r \mapsto 0)]$
 od
 $[\sigma_6: (r \mapsto 0)]$

⑦ ⑥

```

{N ≥ 0} S {r = N!}
var r, x: Int
r := 1;
x := N;
do x > 0 →
  r := r * x;
  x := x - 1;
od

{N ≥ 0} S {r = N!}
var r, x: Int
r := 1;
x := 1;
do x ≤ N →
  r := r * x
  x := x + 1
od

```

⑧ $[\sigma_0: (r \mapsto 1, x \mapsto 5)]$

```

do x > 0 →
   $[\sigma_1^1: (r \mapsto 1, x \mapsto 5), \sigma_1^2: (r \mapsto 5, x \mapsto 4), \sigma_1^3: (r \mapsto 20, x \mapsto 3), \sigma_1^4: (r \mapsto 60, x \mapsto 2), \sigma_1^5: (r \mapsto 120, x \mapsto 1)]$ 
  r := r * x;
  x := x - 1;
   $[\sigma_2^1: (r \mapsto 5, x \mapsto 4), \sigma_2^2: (r \mapsto 20, x \mapsto 3), \sigma_2^3: (r \mapsto 60, x \mapsto 2), \sigma_2^4: (r \mapsto 120, x \mapsto 1), \sigma_2^5: (r \mapsto 120, x \mapsto 0)]$ 
od
 $[\sigma_3: (r \mapsto 120, x \mapsto 0)]$ 

```

$[\sigma_0: (r \mapsto 1, x \mapsto 1)]$

```

do x ≤ 5 →
   $[\sigma_1^1: (r \mapsto 1, x \mapsto 1), \sigma_1^2: (r \mapsto 1, x \mapsto 2), \sigma_1^3: (r \mapsto 2, x \mapsto 3), \sigma_1^4: (r \mapsto 6, x \mapsto 4), \sigma_1^5: (r \mapsto 24, x \mapsto 5)]$ 
  r := r * x
  x := x + 1
   $[\sigma_2^1: (r \mapsto 1, x \mapsto 2), \sigma_2^2: (r \mapsto 2, x \mapsto 3), \sigma_2^3: (r \mapsto 6, x \mapsto 4), \sigma_2^4: (r \mapsto 24, x \mapsto 5), \sigma_2^5: (r \mapsto 120, x \mapsto 6)]$ 
od
 $[\sigma_3: r \mapsto 120, x \mapsto 6]$ 

```

⑨ ⑩ Const A: array [0,4] of Int;

```

var i, s: Int;
 $[\sigma_0: (i \mapsto 3, s \mapsto 5, A \mapsto [2 \ 10 \ 10 \ -1])]$ 
i, s := 0, 0;
 $[\sigma_0]$ 
do i < 4 →
   $[\sigma_1^0, \dots, \sigma_1^4]$ 
  s, i := s + A[i], i + 1
   $[\sigma_2^0, \dots, \sigma_2^4]$ 
od
 $[\sigma_3]$ 

```

⑨ ⑩ agrega una p:= s/i cuando termine el do
 ⑥ agrega una condicion en el do i % 2 == 0

Estado	i	s
σ_0	-3	5
σ_0^0	0	0
σ_1^0	0	0
σ_2^0	2	1
σ_1^1	2	1
σ_2^1	12	2
σ_1^2	12	2
σ_2^2	22	3
σ_1^3	22	3
σ_2^3	21	4
σ_3	21	4

⑩ ⑩ Const A: array [0,4];

```

var i, s, sm3;
i, s, sm3 := 0, 0, 0
 $[\sigma_0]$ 

```

```

do i < N
   $[\sigma_1^0, \dots, \sigma_1^4]$ 
  if A[i] > 0
    s = s + A[i]
  fi
   $[\sigma_2^0, \dots, \sigma_2^4]$ 
  if A[i] % 3 == 0
    sm3 = sm3 + A[i]
  fi
  i = i + 1
   $[\sigma_3^0, \dots, \sigma_3^4]$ 
od
 $[\sigma_4]$ 

```

⑥ Const A: array [0,4] of Int;

```

var i, c: Int;
 $[\sigma_0: (i \mapsto 3, c \mapsto 12, A \mapsto [12 \ -9 \ 10 \ -1])]$ 
i, c := 0, 0;
 $[\sigma_0^1]$ 
do i < 4 →
   $[\sigma_1^0, \dots, \sigma_1^4]$ 
  if A[i] > 0 →
    c := c + 1
  fi
   $[\sigma_2^0, \dots, \sigma_2^4]$ 
  i := i + 1
   $[\sigma_3^0, \dots, \sigma_3^4]$ 
od
 $[\sigma_4]$ 

```

Estado	i	A[i]	c
σ_0	3	-1	12
σ_0^1	0	12	0
σ_0^0	0	12	1
σ_0^3	1	-9	1
σ_1^1	1	-9	1
σ_1^2	1	-9	1
σ_2^3	2	10	1
σ_2^1	2	10	1
σ_2^2	2	10	2
σ_2^3	3	-1	2
σ_3^1	3	-1	2
σ_3^2	3	-1	2
σ_3^3	4	/	2
σ_4	4	/	2

⑥ N=5, A= [12 -9 10 0 -1]

Estado	i	A[i]	s	sm3	Estado	i	A[i]	s	sm3
σ_0	0	12	0	0	σ_4^1	4	-1	22	3
σ_1^0	0	12	0	0	σ_2^4	4	-1	22	3
σ_2^0	0	12	12	0	σ_3^1	5	-1	22	3
σ_3^0	1	-9	12	12	σ_4	5	/	22	3
σ_1^1	1	-9	12	12					
σ_2^1	1	-9	12	12					
σ_3^1	2	10	12	3					
σ_1^2	2	10	12	3					
σ_2^2	2	10	22	3					
σ_3^2	3	0	22	3					
σ_4^2	3	0	22	3					
σ_2^3	3	0	22	3					
σ_3^3	4	-1	22	3					

11 12

16

12 12 Var $x: \text{Int};$
 $\{x > 0\}$
 $x := x * x$
 $\{True\}$

$\{x > 0\} \subseteq \{True\}$
 $\equiv \{TH, WP\}$

$x > 0 \Rightarrow wp.(x := x * x). True$

$\equiv \{Asignacion\}$

$x > 0 \Rightarrow True$

$\equiv \{Logica\}$

$True$

16 Var $x: \text{Int};$
 $\{x \neq 100\}$
 $x := x * x$
 $\{x \geq 0\}$

$\{x \neq 100\} \subseteq \{x \geq 0\}$
 $\equiv \{TH, WP\}$

$(x \neq 100) \Rightarrow wp.(x := x * x). (x \geq 0)$

$\equiv \{Asignacion\}$

$(x \neq 100) \Rightarrow (x * x \geq 0)$

$\equiv \{Logica\}$

$(x \neq 100) \Rightarrow True$

$\equiv \{Logica\}$

$True$

16 Var $x: \text{Int};$
 $\{x > 0 \wedge x < 100\}$
 $x := x * x$
 $\{x \geq 0\}$

$\{x > 0 \wedge x < 100\} \subseteq \{x \geq 0\}$
 $\equiv \{TH, WP\}$

$(x > 0 \wedge x < 100) \Rightarrow wp.(x := x * x). (x \geq 0)$

$\equiv \{Asignacion\}$

$(x > 0 \wedge x < 100) \Rightarrow x * x \geq 0$

$\equiv \{Logica\}$

$(x > 0 \wedge x < 100) \Rightarrow True$

$\equiv \{Logica\}$

$True$

16 Var $x: \text{Int};$
 $\{x > 0\}$
 $x := x * x$
 $\{x < 0\}$

$\{x > 0\} \subseteq \{x < 0\}$
 $\equiv \{TH, WP\}$

$(x > 0) \Rightarrow wp.(x := x * x). (x < 0)$

$\equiv \{Asignacion\}$

$(x > 0) \Rightarrow (x * x < 0)$

$\equiv \{Logica\}$

$(x > 0) \Rightarrow False$

No es valida la TH, $(x > 0) \Rightarrow False$ no es valida, si $x = 1$ la formula es falsa.

16 Var $x: \text{Int};$
 $\{x < 0\}$
 $x := x * x$
 $\{x < 0\}$

$\{x < 0\} \subseteq \{x < 0\}$
 $\equiv \{TH, WP\}$

$(x < 0) \Rightarrow wp.(x := x * x). (x < 0)$

$\equiv \{Asignacion\}$

$(x < 0) \Rightarrow (x * x < 0)$

16 $\{True\} \subseteq \{x \geq 0\}$
 $\equiv \{TH, WP\}$
 $True \Rightarrow wp.(x := x * x). (x \geq 0)$
 $\equiv \{Asignacion\}$
 $True \Rightarrow (x * x \geq 0)$
 $\equiv \{Logica\}$
 $True \Rightarrow True$
 $\equiv \{Logica\}$
 $True$

13

14) a) Var $x, y: \text{Num};$
 $\{ \}$
 $x := x + y$
 $\{x = 6 \wedge y = 5\}$

$\{P\} S \{x=6 \wedge y=5\}$
 $\equiv \{TH, WP\}$
 $P \Rightarrow wp.(x := x + y).(x=6 \wedge y=5)$
 $\equiv \{Assignment\}$
 $P \Rightarrow (x+y=6 \wedge y=5)$
 $\equiv \{Leibniz\}$
 $P \Rightarrow (x+5=6 \wedge y=5)$
 $\equiv \{Arithmetic\}$
 $P \Rightarrow (x=1 \wedge y=5)$

WP es $x=1 \wedge y=5$
 $\{x=1 \wedge y=5\} S \{x=6 \wedge y=5\}$

d) Var $x, y: \text{Num};$
 $\{ \}$
 $x, y := y, x$
 $\{x=6 \wedge y=5\}$

$\{P\} S \{x=B \wedge y=A\}$
 $\equiv \{TH, WP\}$
 $P \Rightarrow wp.(x, y := y, x)(x=B \wedge y=A)$
 $\equiv \{Assignment\}$
 $P \Rightarrow (y=B \wedge x=A)$

La WP es $y=B \wedge x=A$
 $\{y=B \wedge x=A\} S \{x=B \wedge y=A\}$

f) $\{y=1\}$
 if $x \geq y \rightarrow$
 $\{ \}$
 $x := 0$
 $\{ \}$
 else $x \leq y \rightarrow$
 $\{ \}$
 $x := 2$
 $\{ \}$

f:
 $\{(x=0 \vee x=2) \wedge y=1\}$

$\{P\} S; T \{(x=0 \vee x=2) \wedge y=1\}$

b) $\{ \}$
 $x := 8$
 $\{x=8\}$

$\{P\} S \{x=8\}$
 $\equiv \{TH, WP\}$
 $P \Rightarrow wp.(x := 8).(x=8)$
 $\equiv \{Assignment\}$
 $P \Rightarrow 8=8$
 $\equiv \{Logica\}$
 $P \Rightarrow \text{True}$

WP es True
 $\{\text{True}\} S \{x=8\}$

c) $\{ \}$
 $x := 8$
 $\{x=7\}$

$\{P\} S \{x=7\}$
 $\equiv \{TH, WP\}$
 $P \Rightarrow wp.(x := 8).(x=7)$
 $\equiv \{Assignment\}$
 $P \Rightarrow (8=7)$
 $\equiv \{Logica\}$
 $P \Rightarrow \text{False}$

La WP es False
 $\{\text{False}\} S \{x=8\}$

e) Var $x, y, a: \text{Num};$
 $\{ \} \rightarrow \{y=B \wedge x=A\}$
 $a, x := x, y;$
 $\{ \} \longrightarrow \{x=B \wedge a=A\}$
 $y := a$
 $\{x=B \wedge y=A\}$

$\{P\} S; T \{x=B \wedge y=A\}$
 $\equiv \{Sec, Comp o Concat\}$
 $P \Rightarrow wp.(a, x := x, y; y := a).(x=B \wedge y=A)$
 $\equiv \{Def\}$
 $P \Rightarrow wp.(a, x := x, y).wp.(y := a).(x=B \wedge y=A)$
 $\equiv \{Assignment\}$
 $P \Rightarrow wp.(a, x := x, y).(x=B \wedge a=A)$
 $\equiv \{Assignment\}$
 $P \Rightarrow y=B \wedge x=A$

La WP es $y=B \wedge x=A$
 $\{y=B \wedge x=A\} S; T \{x=B \wedge y=A\}$

$((x > y) \vee (x \leq y)) \wedge ((x > y) \Rightarrow wp.(x:=0).((x=0 \vee x=2) \wedge y=1) \wedge ((x \leq y) \Rightarrow wp.(x:=2).((x=0 \vee x=2) \wedge y=1))$

$\equiv \{WP\}$
 $((x > y) \vee (x \leq y)) \wedge ((x > y) \Rightarrow ((0=0 \vee 0=2) \wedge y=1)) \wedge ((x \leq y) \Rightarrow ((2=0 \vee 2=2) \wedge y=1))$

$\equiv \{Logica\}$
 $((x > y) \vee (x \leq y)) \wedge ((x > y) \Rightarrow \text{True} \wedge y=1) \wedge ((x \leq y) \Rightarrow \text{True} \wedge y=1)$

$\equiv \{Logica\}$
 $((x > y) \vee (x \leq y)) \wedge ((x > y) \Rightarrow y=1) \wedge ((x \leq y) \Rightarrow y=1)$
 $\equiv \{Characterization\}$
 $((x > y) \vee (x \leq y)) \wedge (\neg(x > y) \vee y=1) \wedge (\neg(x \leq y) \vee y=1)$
 $\equiv \{destrucción\}$
 $((x > y) \vee (x \leq y)) \wedge (\neg(x > y) \wedge \neg(x \leq y)) \vee y=1$

$\equiv \{Logica\}$
 $\text{True} \wedge ((x \leq y) \wedge (x > y)) \vee y=1$
 $\equiv \{Logica\} \rightarrow y=1$

13) ② $\{True\}$

if $x \geq 1 \longrightarrow x := x+1$

\square $x \leq 1 \longrightarrow x := x-1$

f;

$\{x \neq 1\}$

$$True \Rightarrow ((x \geq 1) \vee (x \leq 1)) \wedge \begin{array}{l} \{x \geq 1 \wedge True\}. (x := x+1). \{x \neq 1\} \\ \wedge \\ \{x \leq 1 \wedge True\}. (x := x-1). \{x \neq 1\} \end{array}$$

$$\begin{array}{ll} x \geq 1 \Rightarrow wp.(x := x+1). (x \neq 1) & x \leq 1 \Rightarrow wp.(x := x-1). (x \neq 1) \\ \equiv \{wp \text{ y } Assign\} & \equiv \{wp \text{ y } Assign\} \\ x \geq 1 \Rightarrow x+1 \neq 1 & x \leq 1 \Rightarrow x-1 \neq 1 \end{array}$$

$$\begin{array}{l} ((x \geq 1) \vee (x \leq 1)) \wedge (x \geq 1 \Rightarrow x+1 \neq 1) \wedge (x \leq 1 \Rightarrow x-1 \neq 1) \\ \equiv \{logical\} \\ (x \geq 1 \Rightarrow x+1 \neq 1) \wedge (x \leq 1 \Rightarrow x-1 \neq 1) \\ \equiv \{logical\} \\ True \wedge True \\ \equiv \{Conj\} \\ True \end{array}$$

(16) a) La post condiccion indica que la posiccion k del arreglo es la sumatoria de todos los elementos previos a k , hasta $N-1$.
El programa no es valido ya que nunca se incrementa i .

b) La post condiccion indica que la posiccion k del arreglo es la sumatoria de todos los elementos previos a k , hasta $N-1$.
Valida las anotaciones

c) La post condiccion indica que la posiccion k del arreglo es la sumatoria de todos los elementos previos a k , hasta $N-1$.
No valida ya que la posiccion -1 del array no existe (numero aleatorio)

d) La postcondiccion comprueba que existe un elemento en la posiccion k del arreglo y que esto implica que la posiccion $i = km$ tambien es igual a E .
La anotacion no es valida

Faltan Extras (17, 18, 19, 20) y (14 f), (15)