# **Message Passing Interface**

Basic functions

### **Outline**

- · Introduction to the Message Passing Interface
  - Parallel programming model of MPI
  - "MPI for Dummies": the 6 basic instructions.
    - » How to run a MPI program.
  - More avanced MPI:
    - » Collective communication
    - » Non-blocking communication

## **Message Passing Interface (MPI)**

- · MPI is a library for Parallel Programming
  - Extends C or Fortran (bindings for C++).
  - Provides abstract datatypes and functions for communication.
- MPI is the de-facto Message Passing Standard
- · More than 180 functionalities
  - Point-to-point and colective comm, non-blocking com, abstract datatypes, DMA, logical organization of the processes, dynamicity, etc...
  - "MPI is as simple as using 6 functions and as complicated as a user wishes to make it."
- · Two main open source distributions:
  - MPICH www-unix.mcs.anl.gov/mpi/mpich
  - OpenMPI (LAM-MPI): www.open-mpi.org

### **Development of MPI**

 November 1992 First draft of MPI 1 November 1993 Second draft of MPI 1.0

MPI 1.0 · June 1994 · June 1995 **MPI 1.1** MPI 1.2 November 1997 · November 1997 MPI 2.0

· Late 1998 Partial implementation of MPI 2.0

• 2000 Most of MPI-2 available

• 2005 All major MPI distributions include MPI

### A few references on MPI

- http://www.mpi-forum.org, for the norm MPI.
- International conference: EuroPVM-MPI (LNCS, Springer)
- · Books:
  - Gropp, William et al., Using MPI, MIT Press.
  - Gropp, William et al., Using MPI-2, MIT Press.
  - Snir, M. et al., Dongarra, J., MPI: The Complete Reference.

# The MPI paradigm

- · Each one of the p processes run the same binary program
  - Single program, Multiple Data paradigm.
  - In "basic" MPI you launch the p processes at the start of the program, and all the p processes must run until the end.
- Each process is identified by a unique rank (a number between 0 and p-1).
- · Based on the rank, each process can:
  - Execute tests (if... then) to run those parts of the program that are relevant;
    - » (Advanced use): nothing prevents the processes to launch
  - essages to/from any other given process.
    - » There are many types of possible messages.

### MPI in 6 words

- 1) MPI\_Init( &argc, &argv) // No comment.
- 2) MPI\_Comm\_rank(&r, communicator)

// returns the rank in the var. int 'r'

3) MPI\_Comm\_size(&p, communicator)

// returns the number of processes in the var. int 'p'

- 4) MPI\_Send( /\* a bunch of args \*/)
- 5) MPI\_Recv( /\* almost the same bunch of args \*/)
- 6) MPI\_Finalize() // No comment

The basic communicator is MPI\_COMM\_WORLD.

# What is a MPI message?

- Look at the MPI\_send function: int MPI\_Send(void\*, int, MPI\_Datatype, int, int, MPI\_Comm).
- · Typical call: MPI\_Send(&work, 1, MPI\_INT, dest, WORKTAG, MPI\_COMM\_WORLD);
- It sends the content of a buffer from the current process to the receiver 'dest' process.
- The buffer is defined by the 3 first arguments:
  - Work (void\*): pointer to the memory area where the data are found.
     1, MPI\_INT: number and basic type of the data (almost sizeof()!)
- A message is identified by a tag (see WORKTAG).
  - The tag must be the same in the Recv and Send.
  - The type is irrelevant to the matching algorithm between sender and receiver.

### MPI\_Recv

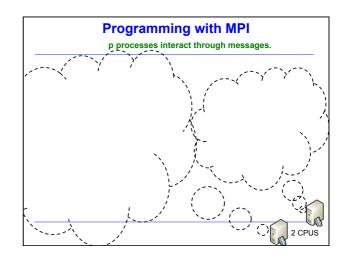
- · Profile of the call:
  - int MPI\_Recv(void\*, int, MPI\_Datatype, int, int, MPI\_Comm, MPI\_Status\*)
- Typical use:

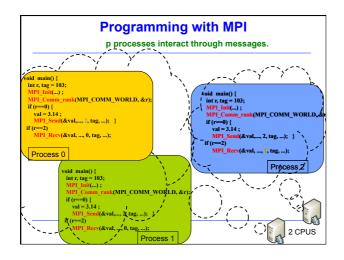
MPI\_Status\* s;

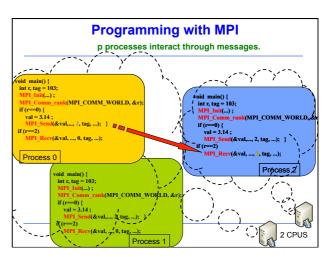
int d, TAG = 103;

MPI\_Recv(&d, 1, MPI\_INT, source, TAG, MPI\_COMM\_WORLD, s);

- 'source' is the rank of the sender proc., TAG is the tag of
- This call is blocking.
  - When the process executes the next instruction, d contains the data that was expected.







## Whole example

```
void main() {
    int p, r, tag = 103;
    MPI_Status stat;
    double val;
    MPI_Init(&argc, &argv);
    MPI_Comm_rank(MPI_COMM_WORLD, &r);
    if (r==0) {
        printf("Processor 0 sends a message to 1\n");
        val = 3.14;
        MPI_Send(&val, 1, MPI_DOUBLE, 1, tag, MPI_COMM_WORLD); }
    else {
        printf("Processor 1 receives a message from 0\n");
        MPI_Recv(&val, 1, MPI_DOUBLE, 0, tag, MPI_COMM_WORLD, &stat);
        printf("I received the value: %.2lf\n", valor); }
}
```

### A few observations

- MPI is much more powerfull than trivial Master/Slave programming.
  - A frequent template: par ranks processos perform something, impar ranks processes run something else.
- · The tags must be predefined by the programmer.
  - The set (source, tag, dest) identifies the message, so be carefull with casts.
- · The buffer is a contiguous memory area
  - Non-contiguous data must be serialized before being sent.
- A message has a fixed size, which must be known before issuing a MPI recv
  - It is frequent to have to send 2 messages:
    - » First the size (1 int),
    - » Then the "real" message ('size' elements).

## **Blocking vs. Non-blocking communication**

- MPI\_Recv(&x,...) is blocking.
- MPI\_Send(&x,...) is "non-blocking"
  - But x is copied into na internal buffer.
  - Send is non-blocking... Until the internal buffer gets full!
- Non-blocking variants: MPI\_Irecv() and MPI\_Isend()
  - Same args as Recv/Send, with one extra of type MPI\_Request.
  - The MPI Request enables the testing of the completion of the non-blocking communication.
- Explicitly bufferized versions of Send/recv: MPI\_Bsend, MPI\_Brecv().

## Test & Wait non-blocking comm.

- MPI\_Test(MPI\_Request\* req, int\* flag, MPI\_Status\* stat)
  - Sets 'flag' to 0 or 1, depending of 'req'
  - You have to test 'flag' afterwards (if (flag)...)
- MPI\_Wait(MPI\_Request\* req, MPI\_Status\* stat)
  - Waits until the completion of the non-blocking comm.
- Key for High-Performance: computation/ communication overlap.
  - » Launch a non-blocking communication (e.g recv),
  - » (in a loop) run all you can run, without having received the data, and test regularly for the reception.
  - » If the loop ends up, then block with a Wait.

# That's all folk!

· See you next week!