

Cryptomorphic Descriptions of Matroids

Adriel Matei, Béla Schneider, Javier Vela, Juš Kocutar

Presenting: Javier, Juš

June 9, 2023

1 Motivation

- 1 Motivation
- 2 What is a matroid?

- 1 Motivation
- 2 What is a matroid?
- 3 Linear algebra and graph theory examples

- 1 Motivation
- 2 What is a matroid?
- 3 Linear algebra and graph theory examples
- 4 Matroids defined by independent sets and bases

- 1 Motivation
- 2 What is a matroid?
- 3 Linear algebra and graph theory examples
- 4 Matroids defined by independent sets and bases
- 5 Matroids in terms of circuits

- 1 Motivation
- 2 What is a matroid?
- 3 Linear algebra and graph theory examples
- 4 Matroids defined by independent sets and bases
- 5 Matroids in terms of circuits
- 6 Rank function and closure

- 1 Motivation
- 2 What is a matroid?
- 3 Linear algebra and graph theory examples
- 4 Matroids defined by independent sets and bases
- 5 Matroids in terms of circuits
- 6 Rank function and closure
- 7 Duality

Abstracting Independence - Motivation

- Notions of independence seem to independently appear in different branches of mathematics (linear algebra, graph theory, etc.)

Abstracting Independence - Motivation

- Notions of independence seem to independently appear in different branches of mathematics (linear algebra, graph theory, etc.)
- These notions share fundamental common properties (e.g. taking an element out of an independent set still yields an independent set)

Abstracting Independence - Motivation

- Notions of independence seem to independently appear in different branches of mathematics (linear algebra, graph theory, etc.)
- These notions share fundamental common properties (e.g. taking an element out of an independent set still yields an independent set)
- Abstracting fundamental properties gives us deeper understanding of the topic (e.g. by developing different ways to visualise the concept of independence)

What is a Matroid?

A matroid is a *structure* that abstracts the notion of *linear independence*. It can be defined in many different ways, such as using:

What is a Matroid?

A matroid is a *structure* that abstracts the notion of *linear independence*. It can be defined in many different ways, such as using:

- Independent sets
- Bases
- Circuits
- and many more...

What is a Matroid?

A matroid is a *structure* that abstracts the notion of *linear independence*. It can be defined in many different ways, such as using:

- Independent sets
- Bases
- Circuits
- and many more...

Due to all the possible characterizations this structure is useful to relate different areas in mathematics such as linear algebra and graph theory.

Linear Independence, Graph Cycles

$$A = \begin{matrix} & \begin{matrix} 1 & 2 & 3 & 4 & 5 \end{matrix} \\ \begin{pmatrix} 1 & 0 & 0 & 0 & 1 \\ 0 & 1 & 1 & 0 & 1 \end{pmatrix} & E = \{1, 2, 3, 4, 5\} \end{matrix}$$

Here, E is the set corresponding to the column vectors of A .

Linear Independence, Graph Cycles

$$A = \begin{matrix} & \begin{matrix} 1 & 2 & 3 & 4 & 5 \end{matrix} \\ \begin{pmatrix} 1 & 0 & 0 & 0 & 1 \\ 0 & 1 & 1 & 0 & 1 \end{pmatrix} & E = \{1, 2, 3, 4, 5\} \end{matrix}$$

Here, E is the set corresponding to the column vectors of A . The linearly independent sets are

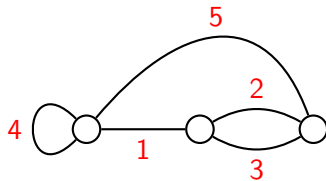
$\{1, 5\}, \{2, 5\}, \{3, 5\}, \{1, 2\}, \{1, 3\}, \{1\}, \{2\}, \{3\}, \{5\}$.

Linear Independence, Graph Cycles

$$A = \begin{pmatrix} \overset{1}{1} & \overset{2}{0} & \overset{3}{0} & \overset{4}{0} & \overset{5}{1} \\ 0 & 1 & 1 & 0 & 1 \end{pmatrix} \quad E = \{\overset{1}{1}, \overset{2}{2}, \overset{3}{3}, \overset{4}{4}, \overset{5}{5}\}$$

Here, E is the set corresponding to the column vectors of A . The linearly independent sets are

$\{\overset{1}{1}, \overset{5}{5}\}, \{\overset{2}{2}, \overset{5}{5}\}, \{\overset{3}{3}, \overset{5}{5}\}, \{\overset{1}{1}, \overset{2}{2}\}, \{\overset{1}{1}, \overset{3}{3}\}, \{\overset{1}{1}\}, \{\overset{2}{2}\}, \{\overset{3}{3}\}, \{\overset{5}{5}\}.$



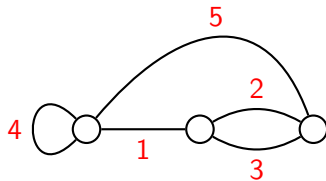
Here, E is the set of edges.

Linear Independence, Graph Cycles

$$A = \begin{pmatrix} \overset{1}{1} & \overset{2}{0} & \overset{3}{0} & \overset{4}{0} & \overset{5}{1} \\ 0 & 1 & 1 & 0 & 1 \end{pmatrix} \quad E = \{\overset{1}{1}, \overset{2}{2}, \overset{3}{3}, \overset{4}{4}, \overset{5}{5}\}$$

Here, E is the set corresponding to the column vectors of A . The linearly independent sets are

$\{\overset{1}{1}, \overset{5}{5}\}, \{\overset{2}{2}, \overset{5}{5}\}, \{\overset{3}{3}, \overset{5}{5}\}, \{\overset{1}{1}, \overset{2}{2}\}, \{\overset{1}{1}, \overset{3}{3}\}, \{\overset{1}{1}\}, \{\overset{2}{2}\}, \{\overset{3}{3}\}, \{\overset{5}{5}\}.$



Here, E is the set of edges. The cycles of this graph are $\{\overset{1}{1}, \overset{2}{2}, \overset{5}{5}\}, \{\overset{1}{1}, \overset{3}{3}, \overset{5}{5}\}, \{\overset{2}{2}, \overset{3}{3}\}, \{\overset{4}{4}\}.$ The independent sets are just the sets of edges that do not contain a cycle.

Independent Sets, Bases

Our setup - finite set E , collection of its subsets \mathcal{I} satisfying:

Independent Sets, Bases

Our setup - finite set E , collection of its subsets \mathcal{I} satisfying:

- 1 $\emptyset \in \mathcal{I}$.

Independent Sets, Bases

Our setup - finite set E , collection of its subsets \mathcal{I} satisfying:

- 1 $\emptyset \in \mathcal{I}$.
- 2 If $I \in \mathcal{I}$ and $J \subset I$, then $J \in \mathcal{I}$.

Independent Sets, Bases

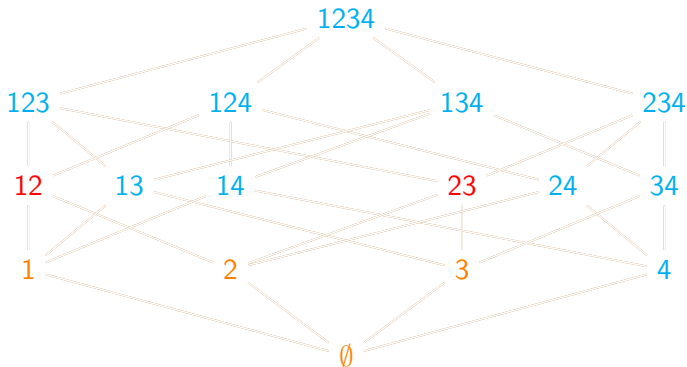
Our setup - finite set E , collection of its subsets \mathcal{I} satisfying:

- 1 $\emptyset \in \mathcal{I}$.
- 2 If $I \in \mathcal{I}$ and $J \subset I$, then $J \in \mathcal{I}$.
- 3 If $J, I \in \mathcal{I}$ and $|J| < |I|$, then there is $e \in I - J$ such that $J \cup e \in \mathcal{I}$.

Independent Sets, Bases

Our setup - finite set E , collection of its subsets \mathcal{I} satisfying:

- 1 $\emptyset \in \mathcal{I}$.
- 2 If $I \in \mathcal{I}$ and $J \subset I$, then $J \in \mathcal{I}$.
- 3 If $J, I \in \mathcal{I}$ and $|J| < |I|$, then there is $e \in I - J$ such that $J \cup e \in \mathcal{I}$.



Independent set, Basis, Dependent set

Circuits — motivation

- Cycles are central to graph theory.

Circuits — motivation

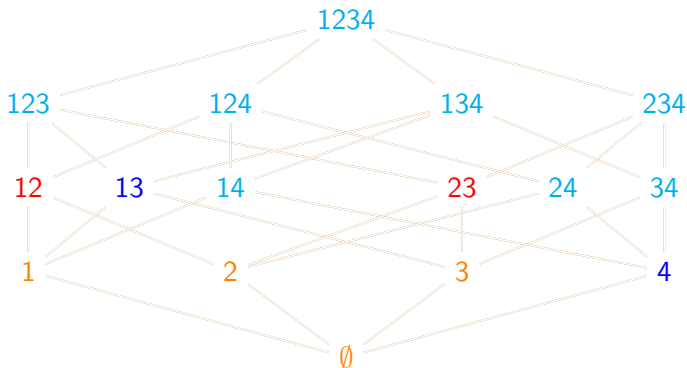
- Cycles are central to graph theory.
- Intuitively, spanning trees and (the absence) of cycles are intimately connected.

Circuits — motivation

- Cycles are central to graph theory.
- Intuitively, spanning trees and (the absence) of cycles are intimately connected.
- It turns out this idea can be used to describe general matroids using a generalized notion of cycles — *Circuits*.

Circuits — motivation

- Cycles are central to graph theory.
- Intuitively, spanning trees and (the absence) of cycles are intimately connected.
- It turns out this idea can be used to describe general matroids using a generalized notion of cycles — *Circuits*.



Independent set, Basis, Circuit, Dependent set

Circuits — Definition

- *Circuits* are minimal dependent sets

Circuits — Definition

- *Circuits* are minimal dependent sets
- We can formally define a matroid given its ground set E and a set of circuits \mathcal{C} :

Circuits — Definition

- *Circuits* are minimal dependent sets
- We can formally define a matroid given its ground set E and a set of circuits \mathcal{C} :
- ① $\emptyset \notin \mathcal{C}$ (the empty set cannot be dependent)

Circuits — Definition

- *Circuits* are minimal dependent sets
- We can formally define a matroid given its ground set E and a set of circuits \mathcal{C} :
 - 1 $\emptyset \notin \mathcal{C}$ (the empty set cannot be dependent)
 - 2 If $C_1, C_2 \in \mathcal{C}$ and $C_1 \subseteq C_2$ then $C_1 = C_2$ (otherwise C_2 would clearly not be minimal)

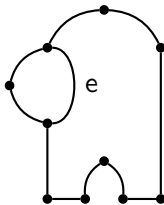
Circuits — Definition

- *Circuits* are minimal dependent sets
- We can formally define a matroid given its ground set E and a set of circuits \mathcal{C} :
 - 1 $\emptyset \notin \mathcal{C}$ (the empty set cannot be dependent)
 - 2 If $C_1, C_2 \in \mathcal{C}$ and $C_1 \subseteq C_2$ then $C_1 = C_2$ (otherwise C_2 would clearly not be minimal)
 - 3 If $C_1, C_2 \in \mathcal{C}$ are distinct circuits and $e \in C_1 \cap C_2$ then a circuit $C_3 \in \mathcal{C}$ exists with $C_3 \subseteq (C_1 \cup C_2) - e$

Circuits — Definition

- *Circuits* are minimal dependent sets
- We can formally define a matroid given its ground set E and a set of circuits \mathcal{C} :
 - 1 $\emptyset \notin \mathcal{C}$ (the empty set cannot be dependent)
 - 2 If $C_1, C_2 \in \mathcal{C}$ and $C_1 \subseteq C_2$ then $C_1 = C_2$ (otherwise C_2 would clearly not be minimal)
 - 3 If $C_1, C_2 \in \mathcal{C}$ are distinct circuits and $e \in C_1 \cap C_2$ then a circuit $C_3 \in \mathcal{C}$ exists with $C_3 \subseteq (C_1 \cup C_2) - e$

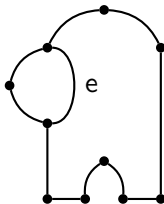
(a) Graph



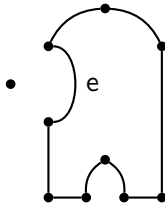
Circuits — Definition

- *Circuits* are minimal dependent sets
- We can formally define a matroid given its ground set E and a set of circuits \mathcal{C} :
 - 1 $\emptyset \notin \mathcal{C}$ (the empty set cannot be dependent)
 - 2 If $C_1, C_2 \in \mathcal{C}$ and $C_1 \subseteq C_2$ then $C_1 = C_2$ (otherwise C_2 would clearly not be minimal)
 - 3 If $C_1, C_2 \in \mathcal{C}$ are distinct circuits and $e \in C_1 \cap C_2$ then a circuit $C_3 \in \mathcal{C}$ exists with $C_3 \subseteq (C_1 \cup C_2) - e$

(a) Graph

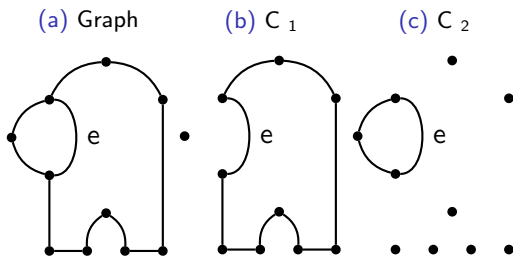


(b) C_1



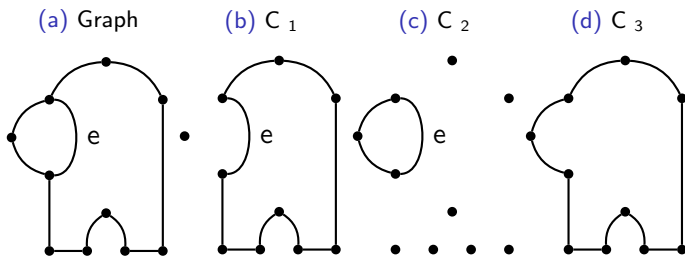
Circuits — Definition

- *Circuits* are minimal dependent sets
- We can formally define a matroid given its ground set E and a set of circuits \mathcal{C} :
 - 1 $\emptyset \notin \mathcal{C}$ (the empty set cannot be dependent)
 - 2 If $C_1, C_2 \in \mathcal{C}$ and $C_1 \subseteq C_2$ then $C_1 = C_2$ (otherwise C_2 would clearly not be minimal)
 - 3 If $C_1, C_2 \in \mathcal{C}$ are distinct circuits and $e \in C_1 \cap C_2$ then a circuit $C_3 \in \mathcal{C}$ exists with $C_3 \subseteq (C_1 \cup C_2) - e$



Circuits — Definition

- *Circuits* are minimal dependent sets
- We can formally define a matroid given its ground set E and a set of circuits \mathcal{C} :
 - 1 $\emptyset \notin \mathcal{C}$ (the empty set cannot be dependent)
 - 2 If $C_1, C_2 \in \mathcal{C}$ and $C_1 \subseteq C_2$ then $C_1 = C_2$ (otherwise C_2 would clearly not be minimal)
 - 3 If $C_1, C_2 \in \mathcal{C}$ are distinct circuits and $e \in C_1 \cap C_2$ then a circuit $C_3 \in \mathcal{C}$ exists with $C_3 \subseteq (C_1 \cup C_2) - e$



Rank Function, Closure

- The notion of *rank* from linear algebra can also be generalized to arbitrary matroids.

- The notion of *rank* from linear algebra can also be generalized to arbitrary matroids. The rank of a set S , denoted $r(S)$, measures the size of the largest independent sets contained in S . For example, the rank of the ground set E is the size of a basis of the matroid (since all bases are of the same size).

Rank Function, Closure

- The notion of *rank* from linear algebra can also be generalized to arbitrary matroids. The rank of a set S , denoted $r(S)$, measures the size of the largest independent sets contained in S . For example, the rank of the ground set E is the size of a basis of the matroid (since all bases are of the same size).
- The closure of a set S , denoted $\text{cl}(S)$, is the maximal superset of S with equal rank. In other words:

Rank Function, Closure

- The notion of *rank* from linear algebra can also be generalized to arbitrary matroids. The rank of a set S , denoted $r(S)$, measures the size of the largest independent sets contained in S . For example, the rank of the ground set E is the size of a basis of the matroid (since all bases are of the same size).
- The closure of a set S , denoted $\text{cl}(S)$, is the maximal superset of S with equal rank. In other words:

$$\text{cl}(S) = \{x \in E : r(S \cup x) = r(S)\}.$$

Rank Function, Closure — examples

1 Graphic matroids:

Rank Function, Closure — examples

- 1 Graphic matroids:
 - the rank function measures the size of the spanning forest.

Rank Function, Closure — examples

1 Graphic matroids:

- the rank function measures the size of the spanning forest.
- the closure is the maximal superset one can obtain by only adding edges when they introduce cycles.

Rank Function, Closure — examples

1 Graphic matroids:

- the rank function measures the size of the spanning forest.
- the closure is the maximal superset one can obtain by only adding edges when they introduce cycles.

2 Vector matroids:

Rank Function, Closure — examples

1 Graphic matroids:

- the rank function measures the size of the spanning forest.
- the closure is the maximal superset one can obtain by only adding edges when they introduce cycles.

2 Vector matroids:

- the rank function measures the rank of the submatrix, so the dimension of the span.

Rank Function, Closure — examples

1 Graphic matroids:

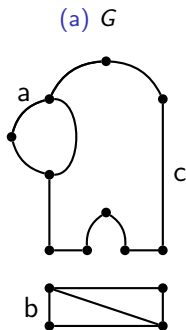
- the rank function measures the size of the spanning forest.
- the closure is the maximal superset one can obtain by only adding edges when they introduce cycles.

2 Vector matroids:

- the rank function measures the rank of the submatrix, so the dimension of the span.
- the closure operator constructs the span, intersected with E .

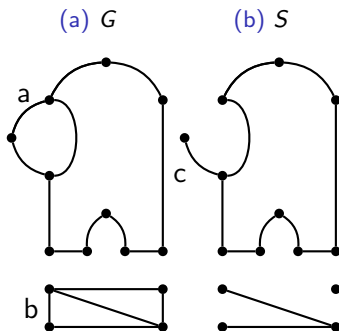
Rank Function, Closure Operator — worked example

Consider the matroid induced by the graph G , and the set S defined below.



Rank Function, Closure Operator — worked example

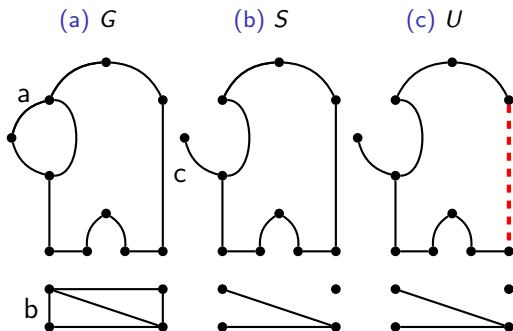
Consider the matroid induced by the graph G , and the set S defined below.



Rank Function, Closure Operator — worked example

Consider the matroid induced by the graph G , and the set S defined below.

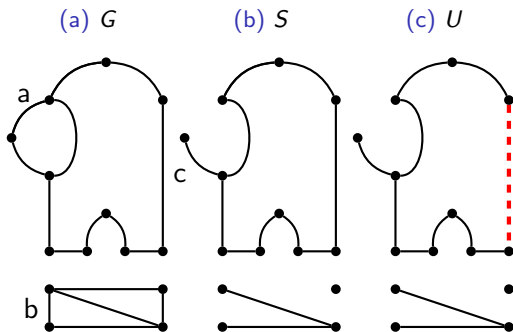
- U is a spanning forest (we removed c to get rid of a cycle)



Rank Function, Closure Operator — worked example

Consider the matroid induced by the graph G , and the set S defined below.

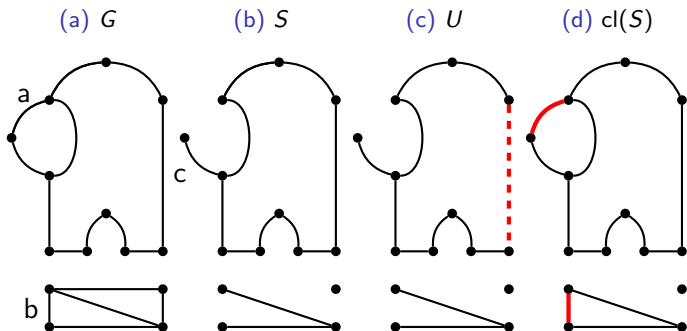
- U is a spanning forest (we removed c to get rid of a cycle)
- The rank is the size of the spanning forest, i.e.
 $r(S) = |U| = 11$.



Rank Function, Closure Operator — worked example

Consider the matroid induced by the graph G , and the set S defined below.

- U is a spanning forest (we removed c to get rid of a cycle)
- The rank is the size of the spanning forest, i.e.
 $r(S) = |U| = 11$.
- The only edges which introduce cycles when added to S are a and b , so $\text{cl}(S) = S + a + b$.



Duality - Definition

Given a matroid M , the dual of M is a matroid denoted by M^* having the same ground set E , such that the bases of M^* are precisely the complements of bases of M .

Duality - Definition

Given a matroid M , the dual of M is a matroid denoted by M^* having the same ground set E , such that the bases of M^* are precisely the complements of bases of M . Formally:

- let M be a matroid on E with \mathcal{B} as its collection of bases, and:

$$\mathcal{B}^*(M) = \{E(M) - B : B \in \mathcal{B}(M)\}.$$

- Then the dual matroid M^* is a matroid having $\mathcal{B}^*(M)$ as its collection of bases.

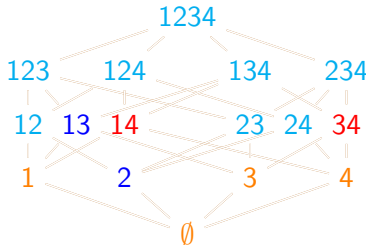
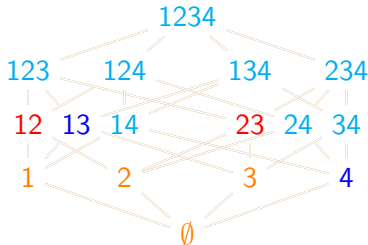
Duality - Definition

Given a matroid M , the dual of M is a matroid denoted by M^* having the same ground set E , such that the bases of M^* are precisely the complements of bases of M . Formally:

- let M be a matroid on E with \mathcal{B} as its collection of bases, and:

$$\mathcal{B}^*(M) = \{E(M) - B : B \in \mathcal{B}(M)\}.$$

- Then the dual matroid M^* is a matroid having $\mathcal{B}^*(M)$ as its collection of bases.



Independent set, Basis, Circuit, Dependent set

Notation: The bases of M^* are called cobases of M

- $(M^*)^* = M$

Notation: The bases of M^* are called cobases of M

- $(M^*)^* = M$
- Let M be a matroid in a set E and suppose $X \subseteq E$. Then

Notation: The bases of M^* are called cobases of M

- $(M^*)^* = M$
- Let M be a matroid in a set E and suppose $X \subseteq E$. Then
 - ① X is *independent* if and only if $E - X$ is *cospanning*.

Notation: The bases of M^* are called cobases of M

- $(M^*)^* = M$
- Let M be a matroid in a set E and suppose $X \subseteq E$. Then
 - 1 X is *independent* if and only if $E - X$ is *cospanning*.
 - 2 X is *spanning* if and only if $E - X$ is *coindependent*.

Notation: The bases of M^* are called cobases of M

- $(M^*)^* = M$
- Let M be a matroid in a set E and suppose $X \subseteq E$. Then
 - ① X is *independent* if and only if $E - X$ is *cospanning*.
 - ② X is *spanning* if and only if $E - X$ is *coindependent*.
- If M is *representable* over the field F , then M^* is also representable over F .

Notation: The bases of M^* are called cobases of M

- $(M^*)^* = M$
- Let M be a matroid in a set E and suppose $X \subseteq E$. Then
 - ① X is *independent* if and only if $E - X$ is *cospanning*.
 - ② X is *spanning* if and only if $E - X$ is *coindependent*.
- If M is *representable* over the field F , then M^* is also representable over F .
- $r(M) + r^*(M) = |E(M)| = |E(M^*)|$

Notation: The bases of M^* are called cobases of M

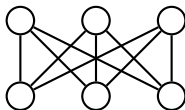
- $(M^*)^* = M$
- Let M be a matroid in a set E and suppose $X \subseteq E$. Then
 - 1 X is *independent* if and only if $E - X$ is *cospanning*.
 - 2 X is *spanning* if and only if $E - X$ is *coindependent*.
- If M is *representable* over the field F , then M^* is also representable over F .
- $r(M) + r^*(M) = |E(M)| = |E(M^*)|$
- $r^*(X) = r(E - X) + |X| - r(M)$

Duality - Examples

Example 1: Let us consider the bipartite graph $K_{3,3}$.

Duality - Examples

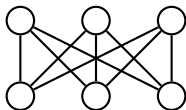
Example 1: Let us consider the bipartite graph $K_{3,3}$.



We can see it has 9 edges, so the ground set is $E = \{1, 2, 3, \dots, 9\}$.

Duality - Examples

Example 1: Let us consider the bipartite graph $K_{3,3}$.



We can see it has 9 edges, so the ground set is $E = \{1, 2, 3, \dots, 9\}$.

The matrix associated to the matroid $M(K_{3,3})$ is the following:

$$B = \begin{bmatrix} 1 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 1 \\ 0 & 1 & 0 & 0 & 0 & 1 & 1 & 0 & 1 \\ 0 & 0 & 1 & 0 & 0 & 1 & 1 & 1 & 1 \\ 0 & 0 & 0 & 1 & 0 & 0 & 1 & 1 & 1 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 1 & 1 \end{bmatrix}$$

Theorem

Let M be the vector matroid of the matrix $[I_r|D]$ where the columns of this matrix are labelled, in order, e_1, e_2, \dots, e_n and $1 \leq r < n$.

Theorem

Let M be the vector matroid of the matrix $[I_r|D]$ where the columns of this matrix are labelled, in order, e_1, e_2, \dots, e_n and $1 \leq r < n$. Then M^ is the vector matroid of $[-D^T|I_{n-r}]$ where its columns are also labelled e_1, e_2, \dots, e_n in that order.*

Duality - Examples

Theorem

Let M be the vector matroid of the matrix $[I_r | D]$ where the columns of this matrix are labelled, in order, e_1, e_2, \dots, e_n and $1 \leq r < n$. Then M^ is the vector matroid of $[-D^T | I_{n-r}]$ where its columns are also labelled e_1, e_2, \dots, e_n in that order.*

So, the dual matroid $M^*(K_{3,3})$ is represented by the following matrix:

$$B^* = \begin{bmatrix} 1 & 1 & 1 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 1 & 1 & 1 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 & 1 & 0 & 0 & 1 & 0 \\ 1 & 1 & 1 & 1 & 1 & 0 & 0 & 0 & 1 \end{bmatrix}$$

Duality - Examples

So we have,

$$B = \begin{bmatrix} 1 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 1 \\ 0 & 1 & 0 & 0 & 0 & 1 & 1 & 0 & 1 \\ 0 & 0 & 1 & 0 & 0 & 1 & 1 & 1 & 1 \\ 0 & 0 & 0 & 1 & 0 & 0 & 1 & 1 & 1 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 1 & 1 \end{bmatrix} \quad B^* = \begin{bmatrix} 1 & 1 & 1 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 1 & 1 & 1 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 & 1 & 0 & 0 & 1 & 0 \\ 1 & 1 & 1 & 1 & 1 & 0 & 0 & 0 & 1 \end{bmatrix}$$

Duality - Examples

So we have,

$$B = \begin{bmatrix} 1 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 1 \\ 0 & 1 & 0 & 0 & 0 & 1 & 1 & 0 & 1 \\ 0 & 0 & 1 & 0 & 0 & 1 & 1 & 1 & 1 \\ 0 & 0 & 0 & 1 & 0 & 0 & 1 & 1 & 1 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 1 & 1 \end{bmatrix} \quad B^* = \begin{bmatrix} 1 & 1 & 1 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 1 & 1 & 1 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 & 1 & 0 & 0 & 1 & 0 \\ 1 & 1 & 1 & 1 & 1 & 0 & 0 & 0 & 1 \end{bmatrix}$$

Where, $r(M) = 5$ and $r^*(M) = 4$, and $|E(M)| = 9$. Hence, we see that $r(M) + r^*(M) = |E(M)|$ holds.

Duality - Examples

So we have,

$$B = \begin{bmatrix} 1 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 1 \\ 0 & 1 & 0 & 0 & 0 & 1 & 1 & 0 & 1 \\ 0 & 0 & 1 & 0 & 0 & 1 & 1 & 1 & 1 \\ 0 & 0 & 0 & 1 & 0 & 0 & 1 & 1 & 1 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 1 & 1 \end{bmatrix} \quad B^* = \begin{bmatrix} 1 & 1 & 1 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 1 & 1 & 1 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 & 1 & 0 & 0 & 1 & 0 \\ 1 & 1 & 1 & 1 & 1 & 0 & 0 & 0 & 1 \end{bmatrix}$$

Where, $r(M) = 5$ and $r^*(M) = 4$, and $|E(M)| = 9$. Hence, we see that $r(M) + r^*(M) = |E(M)|$ holds.

Also, in this example, both matroids are vector matroids over the field \mathbb{F}_2 .

Duality - Examples

So we have,

$$B = \begin{bmatrix} 1 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 1 \\ 0 & 1 & 0 & 0 & 0 & 1 & 1 & 0 & 1 \\ 0 & 0 & 1 & 0 & 0 & 1 & 1 & 1 & 1 \\ 0 & 0 & 0 & 1 & 0 & 0 & 1 & 1 & 1 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 1 & 1 \end{bmatrix} \quad B^* = \begin{bmatrix} 1 & 1 & 1 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 1 & 1 & 1 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 & 1 & 0 & 0 & 1 & 0 \\ 1 & 1 & 1 & 1 & 1 & 0 & 0 & 0 & 1 \end{bmatrix}$$

Where, $r(M) = 5$ and $r^*(M) = 4$, and $|E(M)| = 9$. Hence, we see that $r(M) + r^*(M) = |E(M)|$ holds.

Also, in this example, both matroids are vector matroids over the field \mathbb{F}_2 . However, although $M(K_{3,3})$ is a graphic matroid, its dual, $M^*(K_{3,3})$ is not graphic. This results points to the following:

Duality - Examples

So we have,

$$B = \begin{bmatrix} 1 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 1 \\ 0 & 1 & 0 & 0 & 0 & 1 & 1 & 0 & 1 \\ 0 & 0 & 1 & 0 & 0 & 1 & 1 & 1 & 1 \\ 0 & 0 & 0 & 1 & 0 & 0 & 1 & 1 & 1 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 1 & 1 \end{bmatrix} \quad B^* = \begin{bmatrix} 1 & 1 & 1 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 1 & 1 & 1 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 & 1 & 0 & 0 & 1 & 0 \\ 1 & 1 & 1 & 1 & 1 & 0 & 0 & 0 & 1 \end{bmatrix}$$

Where, $r(M) = 5$ and $r^*(M) = 4$, and $|E(M)| = 9$. Hence, we see that $r(M) + r^*(M) = |E(M)|$ holds.

Also, in this example, both matroids are vector matroids over the field \mathbb{F}_2 . However, although $M(K_{3,3})$ is a graphic matroid, its dual, $M^*(K_{3,3})$ is not graphic. This results points to the following:

Theorem

The dual of a graphic matroid is itself graphic if and only if the underlying graph is planar.

To summarize

To summarize

- We abstracted linear independence with three axioms.

To summarize

- We abstracted linear independence with three axioms.
- This abstract independence turns out to have connections to other fields too, like graph theory.

To summarize

- We abstracted linear independence with three axioms.
- This abstract independence turns out to have connections to other fields too, like graph theory.
- Each matroid has circuits, bases, rank function and closure operator, each with their unique properties that determine them from the reverse direction.

To summarize

- We abstracted linear independence with three axioms.
- This abstract independence turns out to have connections to other fields too, like graph theory.
- Each matroid has circuits, bases, rank function and closure operator, each with their unique properties that determine them from the reverse direction.
- Every matroid has an "opposite matroid" called its dual, that is formed using the complements of the bases of the original matroid.

- J. G. Oxley. Matroid theory. Oxford Science Publications. Oxford University Press, USA, 1993.

- J. G. Oxley. Matroid theory. Oxford Science Publications. Oxford University Press, USA, 1993.
- G. Gordon and J. McNulty. Matroids: A geometric introduction. Cambridge University Press, 2021.

Thank You!

For can now give an example for the concepts given until now.

Example:

Now, we will formulate another important point, the so called, orthogonality, which refers to the link between circuits and cocircuits.

Proposition: For a given matroid M , let C be a circuit and C^* be a cocircuit. Then, $|C \cap C^*| \neq 1$.

Proof.

We will prove this by contradiction. Suppose $C \cap C^* = \{x\}$, for some $x \in E$, this is, there exists an element in the intersection such that the cardinality will be 1. Now consider the hyperplane $H = E - C^*$, and recall that the closure $cl(H) = H$. We can observe that by the way in which it is defined if $x \in C^*$, then $x \notin H$, but $C - x \subseteq H$. Moreover, we have that $x \in cl(C - x)$ hence $x \in cl(H) = H$. Then $x \notin C^*$, which is a contradiction. So, $C \cap C^* \neq \{x\}$, which implies $|C \cap C^*| \neq 1$. □