

①

RIGHT-ROTATE(T, x)

 $y = x.left$
 $x.left = y.right$

if $y.right \neq NIL$

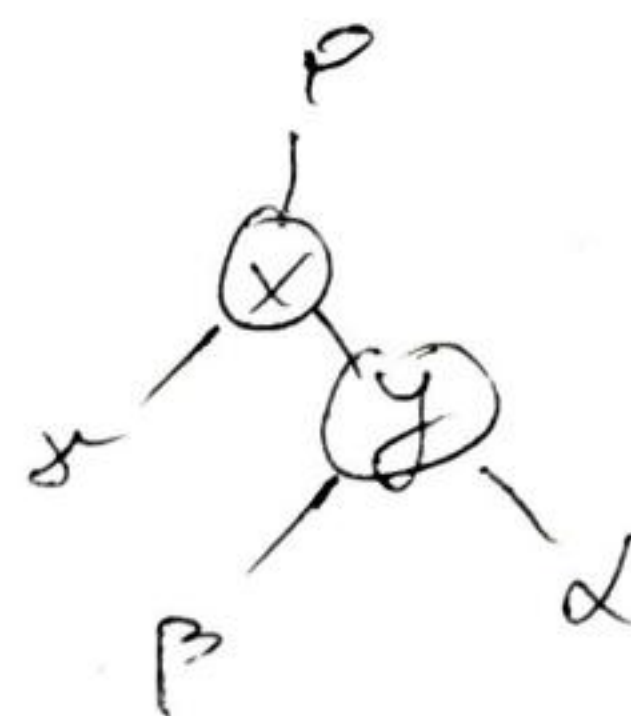
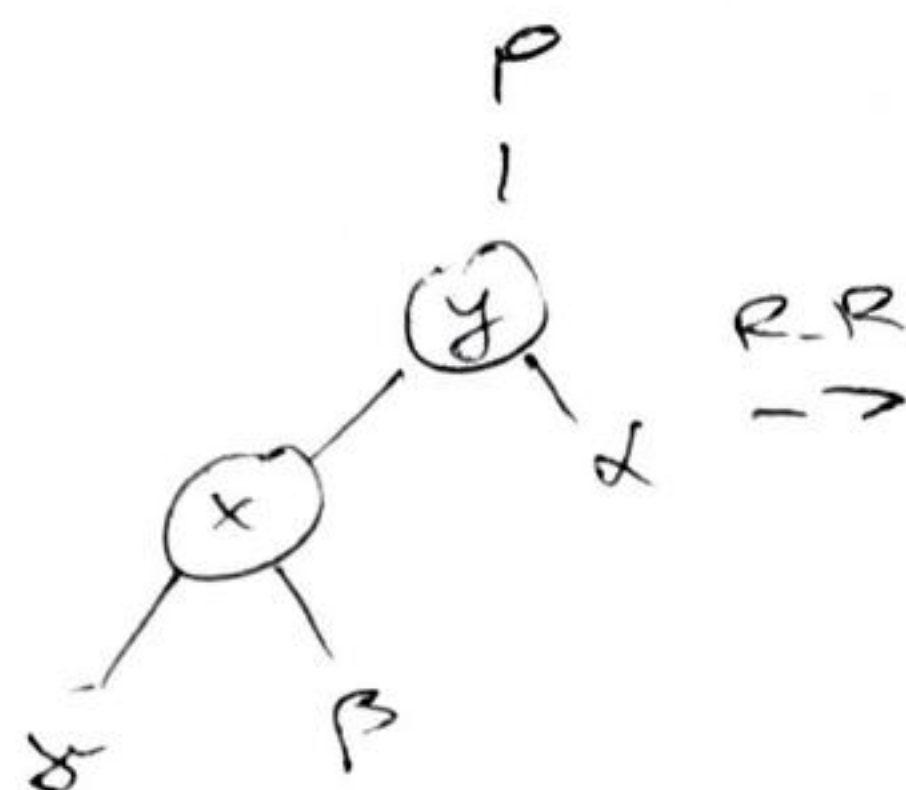
| $y.right.p = x$

if $x.p == NIL$

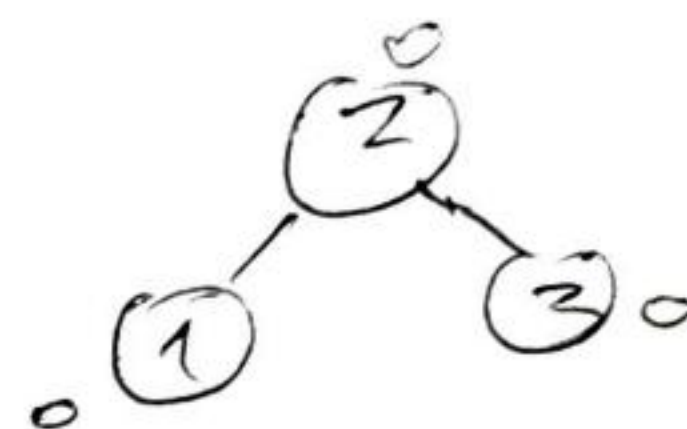
| $T.root = y$

else if $x == x.p.right$

| $x.p.right = y$

else $x.p.left = y$
 $y.right = x$
 $y.p = x.p$
 $x.p = y$


Primjerice:



②

~~INSERT~~
INSERT(T, x, key)

if $x == NIL$

| $x.key = key$

| $x.h = 1$

else if $key < x.key$

| $x.left = INSERT(T, x.left, key)$

| $x.left.p = x$

else if $key > x.key$

| $x.right = INSERT(T, x.right, key)$

| $x.right.p = x$

else return x
 $x.h = 1 + \max(x.left.h, x.right.h)$
 $balance = BALANCE_OF_NODE(x)$

if $balance > 1$ && $key < x.left.key$

| ~~return~~ return RIGHT-ROTATE(T, x)

if $balance < -1$ && $key > x.right.key$

| return LEFT-ROTATE(T, x)

if $balance > 1$ && $key > x.left.key$

| $x.left = LEFT_ROTATE(T, x.left)$

| return RIGHT-ROTATE(T, x)

if $balance < -1$ && $key < x.right.key$

| $x.right = RIGHT_ROTATE(T, x.right)$

| return LEFT-ROTATE(T, x)

return x

ANALIZA:

UBACIVANJE U BST KOD
BALANSIRANOG STABLA (AVL) JE
 $O(\lg n)$ JER JE VISINA STABLA
UVJEK $O(\lg n)$.

SVI IF-STATEMENTI SU OČITO $O(1)$
PREOSTAJE SAMO LEFT/RIGHT
ROTATE ŠTO SMO POKAZALI DA
JE U KONSTANTNOM VREMENU $O(1)$
 \Rightarrow UKUPNA VREMENSKA SLOŽENOST
IZNOSI $O(\lg n)$.