



# COMPILER DESIGN

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## CHAPTER-4

### Syntax-directed definitions



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## Syntax Directed Definitions

- We can associate information with a language construct by attaching attributes to the grammar symbols.
- A syntax directed definition specifies the values of attributes by associating semantic rules with the grammar productions.

Production  
 $E \rightarrow E1 + T$

Semantic Rule  
 $E.code = E1.code = T.code$





## Syntax Directed Definitions

- A SDD is a context free grammar with attributes and rules.
- Attributes are associated with grammar symbols and rules with productions.
- Attributes may be of many kinds: numbers, types, table references, strings, etc.



## Syntax Directed Definitions

Synthesized attributes:

A synthesized attribute at node N is defined only in terms of attribute values of children of N and at N it.

Inherited attributes:

An inherited attribute at node N is defined only in terms of attribute values at N's parent, N itself and N's siblings



## Synthesized attributes

- A synthesized attribute at node N is defined only in terms of attribute values of children of N and at N it.
- Each of the non-terminals has a single synthesized attribute, called val.
- An SDD that involves only synthesized attributes is called S-attributed.
- Each rule computes an attribute for the non-terminal at the head of a production from attributes taken from the body of the production.



## SDD for expression grammar with Synthesized attribute

### Production

- 1)  $L \rightarrow E \ n$
- 2)  $E \rightarrow E1 + T$
- 3)  $E \rightarrow T$
- 4)  $T \rightarrow T1 * F$
- 5)  $T \rightarrow F$
- 6)  $F \rightarrow (E)$
- 7)  $F \rightarrow \text{digit}$

### Semantic Rules

$L.val = E.val$   
 $E.val = E1.val + T.val$   
 $E.val = T.val$   
 $T.val = T1.val * F.val$   
 $T.val = F.val$   
 $F.val = E.val$   
 $F.val = \text{digit.lexval}$



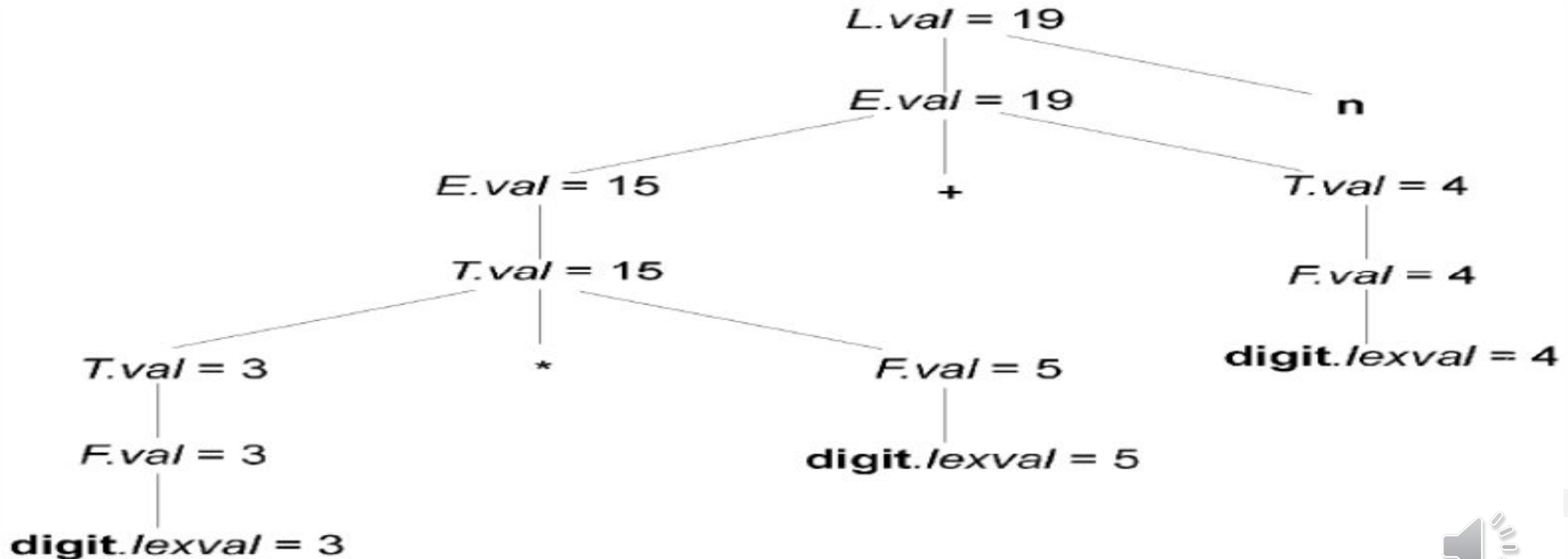


## Evaluating an SDD at the Nodes of a Parse Tree

- A parse tree, showing the value(s) of its attribute(s) is called an annotated parse tree.
- With synthesized attributes, evaluate attributes in bottom-up order.



## Annotated Parse Tree for $3*5+4$



## Inherited Attributes

- An INHERITED ATTRIBUTE for a non-terminal B at a parse-tree node N is defined by a semantic rule associated with the production at the parent of N. The production must have B as a symbol in its body.
- An inherited attribute at node N is defined only in terms of attribute values at N's parent, N itself, and N's siblings



## SDD for expression grammar with inherited grammar

PRODUCTION	SEMANTIC RULE
$D \rightarrow TL$	$L.in := T.type$
$T \rightarrow \text{int}$	$T.type := \text{integer}$
$T \rightarrow \text{real}$	$T.type := \text{real}$
$L \rightarrow L_1, \text{id}$	$L_1.in := L.in; \text{addtype}(\text{id.entry}, L.in)$
$L \rightarrow \text{id}$	$\text{addtype}(\text{id.entry}, L.in)$



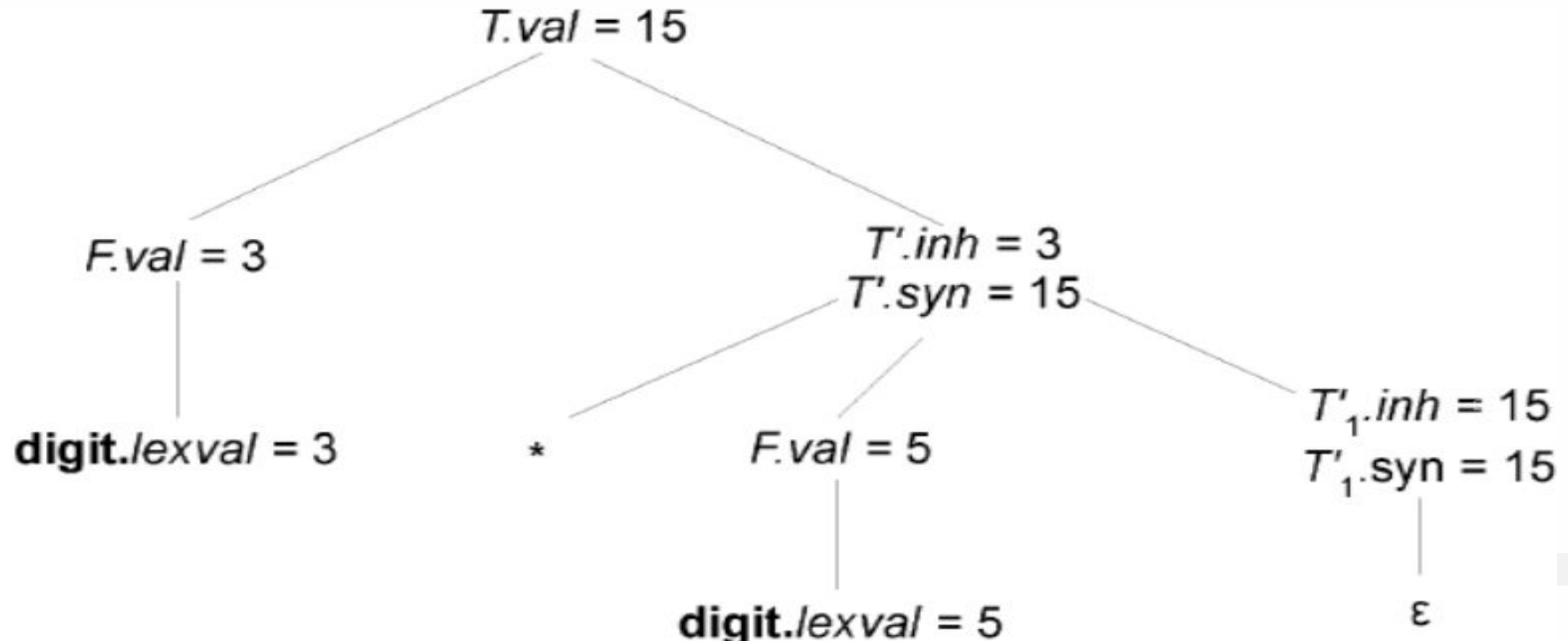
## SDD for expression grammar with inherited grammar

Production	Semantic Rules
$T \rightarrow F T'$	$T'.inh = F.val$ $T.val = T'.syn$
$T' \rightarrow *F T'_1$	$T'_1.inh = T'.inh \times F.val$ $T'.syn = T'_1.syn$
$T' \rightarrow \epsilon$	$T'.syn = T'.inh$
$F \rightarrow \mathbf{digit}$	$F.val = \mathbf{digit.lexval}$





## Annotated Parse Tree for $3*5$



## Evaluation Orders for SDD's

- Dependency graphs" tool for determining an evaluation order for the attribute instances in a given parse tree.
- annotated parse tree shows the values of attributes, a dependency graph helps to determine how those values can be computed.



## Dependency graphs

- Edges express constraints implied by the semantic rules.
- Each attribute is associated to a node
- If a semantic rule associated with a production  $p$  defines the value of synthesized attribute  $A.b$  in terms of the value of  $X.c$ , then graph has an edge from  $X.c$  to  $A.b$ .
- If a semantic rule associated with a production  $p$  defines the value of inherited attribute  $B.c$  in terms of value of  $X.a$ , then graph has an edge



## Topological Sort

- A dependency graph characterizes the possible order in which we can evaluate the attributes at various nodes of a parse tree.
- If there is an edge from node M to N, then attribute corresponding to M first be evaluated before evaluating N.
- Thus the allowable orders of evaluation are  $N_1, N_2, \dots, N_k$  such that if there is an edge from  $N_i$  to  $N_j$  then  $i < j$ .



## Topological Sort

- Such an ordering embeds a directed graph into a linear order, and is called a topological sort of the graph.
- If there is any cycle in the graph, then there are no topological sorts.





## S-Attributed

- If every attribute is synthesized.
- S-attributed SDD can be evaluated in bottom up order of the nodes of the parse tree.





## L-Attributed

- Synthesized, or Inherited, but with the rules limited as follows. Suppose that there is a production  $A \rightarrow X_1 X_2 \dots X_n$ , and that there is an inherited attribute  $X_i$  computed by a rule associated with this production.
- Inherited attributes associated with the head  $A$ .
- Either inherited or synthesized attributes associated with the occurrences of symbols  $X_1$ ,  $X_2$ , ...,  $X_{i-1}$  located to the left of  $X_i$ .
- Inherited or synthesized attributes associated with this occurrence of  $X_i$  itself, but only in such a way that there are no cycles in a dependency graph of attributes of this  $X_i$ .



## Recursive Descent Parser

- Top-down parsing strategy, for LL(1) grammars.
- One procedure per nonterminal.
- Stack contents embedded in recursive call sequence
- Each procedure “commits” to one production, based on the next input symbol, and the select sets.
- It uses back- tracking. So, it is not usefull.

