

Let Σ be some alphabet.

Prove that:

$$\forall S, T \in \mathcal{P}(\Sigma^*), S \neq \emptyset \wedge S \subseteq T \wedge (\forall t \in T, \exists s \in S, s \leq^s t) \longrightarrow \\ \Sigma^* T \Sigma^* = \Sigma^* S \Sigma^*.$$

Let $S, T \in \mathcal{P}(\Sigma^*)$ be such that $S \neq \emptyset$, $S \subseteq T$ and $\forall t \in T, \exists s \in S, s \leq^s t$.

(We must show that $\Sigma^* T \Sigma^* = \Sigma^* S \Sigma^*$)

First we'll show that: $\Sigma^* T \Sigma^* \subseteq \Sigma^* S \Sigma^*$

Let $w \in \Sigma^* T \Sigma^*$, By definition of $\Sigma^* T \Sigma^*$ we get that

$$\exists w_1, w_2 \in \Sigma^*, t \in T, w = w_1 t w_2.$$

Now since $t \in T$ we get that $\exists s \in S, s \leq^s t$, By definition of the \leq^s relation we get that $\exists \psi, \chi \in \Sigma^*, t = \psi s \chi$, Thus $w = w_1 \psi s \chi w_2$. Lets denote $u_1 = w_1 \psi$ and $u_2 = \chi w_2$, We get that $w = u_1 s u_2$, Since $u_1, u_2 \in \Sigma^*$ and $s \in S$ we get that $u_1 s u_2 \in \Sigma^* S \Sigma^*$, and so $w \in \Sigma^* S \Sigma^*$. Thus $\Sigma^* T \Sigma^* \subseteq \Sigma^* S \Sigma^*$

Now we'll show that: $\Sigma^* S \Sigma^* \subseteq \Sigma^* T \Sigma^*$

Let $w \in \Sigma^* S \Sigma^*$, By definition of $\Sigma^* S \Sigma^*$ we get that

$$\exists w_1, w_2 \in \Sigma^*, s \in S, w = w_1 s w_2.$$

Now since $s \in S$ and $S \subseteq T$ we get that $s \in T$ and we get that

$$\exists w_1, w_2 \in \Sigma^*, s \in T, w = w_1 s w_2. \text{ Now by definition of } \Sigma^* T \Sigma^*$$

we get that $w \in \Sigma^* T \Sigma^*$. Thus $\Sigma^* S \Sigma^* \subseteq \Sigma^* T \Sigma^*$

From the two set inclusions we get that $\Sigma^* T \Sigma^* = \Sigma^* S \Sigma^*$ as was to be shown.

Q.E.D.
