Perfect forms and perfect Delaunay polytopes

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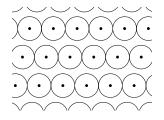
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I. Lattices, packings

and coverings

Lattice packings

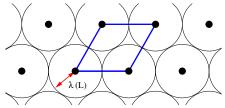
- ▶ A lattice $L \subset \mathbb{R}^n$ is a set of the form $L = \mathbb{Z}v_1 + \cdots + \mathbb{Z}v_n$.
- ▶ A packing is a family of balls $B_n(x_i, r)$, $i \in I$ of the same radius r and center x_i such that their interiors are disjoint.



▶ If L is a lattice, the lattice packing is the packing defined by taking the maximal value of $\alpha > 0$ such that $L + B_n(0, \alpha)$ is a packing.

Density of lattice packings

► Take the lattice packing defined by a lattice *L*:



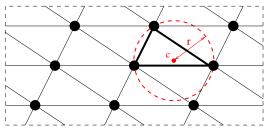
The packing density has the expression

$$\delta(L) = \frac{\lambda(L)^n \operatorname{vol}(B_n(0,1))}{\det L} \quad \text{with} \quad \lambda(L) = \frac{1}{2} \min_{v \in L - \{0\}} ||v||,$$

 $\operatorname{vol}(B_n(0,1))$ the volume of the unit ball $B_n(0,1)$ and $\det L$ the volume of an unit cell.

Empty sphere and Delaunay polytopes

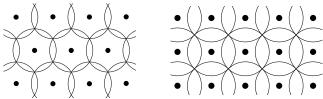
- ▶ Definition: A sphere S(c, r) of center c and radius r in an n-dimensional lattice L is said to be an empty sphere if:
 - (i) $||v-c|| \ge r$ for all $v \in L$,
 - (ii) the set $S(c, r) \cap L$ contains n + 1 affinely independent points.
- ▶ Definition: A Delaunay polytope P in a lattice L is a polytope, whose vertex-set is $L \cap S(c, r)$.



ightharpoonup Delaunay polytopes define a tesselation of the Euclidean space \mathbb{R}^n

Lattice covering

▶ For a lattice L we define the covering radius $\mu(L)$ to be the smallest r such that the family of balls $v + B_n(0, r)$ for $v \in L$ cover \mathbb{R}^n .



The covering density has the expression

$$\Theta(L) = \frac{\mu(L)^n \operatorname{vol}(B_n(0,1))}{\det(L)} \ge 1$$

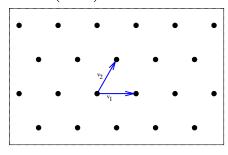
with $\mu(L)$ being the largest radius of Delaunay polytopes

▶ The only general method for computing $\Theta(L)$ is to compute all Delaunay polytopes of L.

II. Gram matrix formalism

Gram matrix and lattices

- ▶ Denote by S^n the vector space of real symmetric $n \times n$ matrices and $S^n_{>0}$ the convex cone of real symmetric positive definite $n \times n$ matrices.
- ▶ Take a basis $(v_1, ..., v_n)$ of a lattice L and associate to it the Gram matrix $G_{\mathbf{v}} = (\langle v_i, v_j \rangle)_{1 \leq i,j \leq n} \in S^n_{>0}$.
- lacktriangle Example: take the hexagonal lattice generated by $v_1=(1,0)$ and $v_2=\left(rac{1}{2},rac{\sqrt{3}}{2}
 ight)$



$$G_{\mathbf{v}} = \frac{1}{2} \left(\begin{array}{cc} 2 & 1 \\ 1 & 2 \end{array} \right)$$

Isometric lattices

▶ Take a basis $(v_1, ..., v_n)$ of a lattice L with $v_i = (v_{i,1}, ..., v_{i,n}) \in \mathbb{R}^n$ and write the matrix

$$V = \left(\begin{array}{ccc} v_{1,1} & \dots & v_{n,1} \\ \vdots & \ddots & \vdots \\ v_{1,n} & \dots & v_{n,n} \end{array}\right)$$

and $G_{\mathbf{v}} = V^T V$.

The matrix $G_{\mathbf{v}}$ is defined by $\frac{n(n+1)}{2}$ variables as opposed to n^2 for the basis V.

- ▶ If $M \in S_{>0}^n$, then there exists V such that $M = V^T V$ (Gram Schmidt orthonormalization)
- ▶ If $M = V_1^T V_1 = V_2^T V_2$, then $V_1 = OV_2$ with $O^T O = I_n$ (i.e. O corresponds to an isometry of \mathbb{R}^n).
- Also if L is a lattice of \mathbb{R}^n with basis \mathbf{v} and u an isometry of \mathbb{R}^n , then $G_{\mathbf{v}} = G_{u(\mathbf{v})}$.

Arithmetic minimum

▶ The arithmetic minimum of $A \in S_{>0}^n$ is

$$\min(A) = \min_{x \in \mathbb{Z}^n - \{0\}} x^T A x$$

▶ The minimal vector set of $A \in S_{>0}^n$ is

$$Min(A) = \left\{ x \in \mathbb{Z}^n \mid x^T A x = min(A) \right\}$$

- ▶ Both min(A) and Min(A) can be computed using some programs (for example sv by Vallentin)
- ▶ The matrix $A_{hex} = \begin{pmatrix} 2 & 1 \\ 1 & 2 \end{pmatrix}$ has

$$Min(A_{hex}) = \{\pm(1,0), \pm(0,1), \pm(1,-1)\}.$$

Reexpression of previous definitions

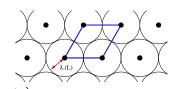
▶ Take a lattice $L = \mathbb{Z}v_1 + \cdots + \mathbb{Z}v_n$. If $x \in L$,

$$x = x_1 v_1 + \cdots + x_n v_n$$
 with $x_i \in \mathbb{Z}$

we associate to it the column vector $X = \begin{pmatrix} x_1 \\ \vdots \\ x_n \end{pmatrix}$

• We get $||x||^2 = X^T G_{\mathbf{v}} X$ and

$$\det L = \sqrt{\det G_{\mathbf{v}}} \text{ and } \lambda(L) = \frac{1}{2} \sqrt{\min(G_{\mathbf{v}})}$$



▶ For $A_{hex} = \begin{pmatrix} 2 & 1 \\ 1 & 2 \end{pmatrix}$, det $A_{hex} = 3$ and min $(A_{hex}) = 2$

Changing basis

▶ If **v** and **v**' are two basis of a lattice *L* then V' = VP with $P \in GL_n(\mathbb{Z})$. This implies

$$G_{v'} = V'^T V' = (VP)^T VP = P^T \{V^T V\}P = P^T G_{v}P$$

▶ If $A, B \in S_{>0}^n$, they are called arithmetically equivalent if there is at least one $P \in GL_n(\mathbb{Z})$ such that

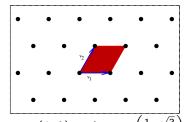
$$A = P^T B P$$

- ▶ Lattices up to isometric equivalence correspond to $S_{>0}^n$ up to arithmetic equivalence.
- ► In practice, Plesken/Souvignier wrote a program isom for testing arithmetic equivalence and a program autom for computing automorphism group of lattices.

 All such programs take Gram matrices as input.

An example of equivalence

▶ Take the hexagonal lattice and two basis in it.



$$v_1=(1,0) \text{ and } v_2=\left(\frac{1}{2},\frac{\sqrt{3}}{2}\right) \qquad v_1'=\left(\frac{5}{2},\frac{\sqrt{3}}{2}\right) \text{ and } v_2'=(-1,0)$$

▶ One has
$$v_1' = 2v_1 + v_2$$
, $v_2' = -v_1$ and $P = \begin{pmatrix} 2 & -1 \\ 1 & 0 \end{pmatrix}$

$$G_{\mathbf{v}} = \begin{pmatrix} 1 & \frac{1}{2} \\ \frac{1}{2} & 1 \end{pmatrix}$$
 and $G_{\mathbf{v}'} = \begin{pmatrix} 7 & -\frac{5}{2} \\ -\frac{5}{2} & 1 \end{pmatrix} = P^T G_{\mathbf{v}} P$

Root lattices

Let us take the lattice

$$A_n = \left\{ x \in \mathbb{Z}^{n+1} \text{ s.t. } \sum_{i=1}^{n+1} x_i = 0 \right\}$$

If we take the basis $v_i = e_{i+1} - e_i$ then we get the Gram matrix $A = (a_{ij})_{1 \le i,j \le n}$ with $a_{i,i} = 2$, $a_{i,i+1} = a_{i+1,i} = -1$ and $a_{i,j} = 0$ otherwise.

Let us take the lattice

$$D_n = \left\{ x \in \mathbb{Z}^n \text{ s.t. } \sum_{i=1}^n x_i \equiv 0 \pmod{2} \right\}$$

For the basis $v_1 = e_1 + e_2$, $v_2 = e_1 - e_2$, $v_i = e_i - e_{i-1}$ we get

$$G_{v} = \left(\begin{array}{ccccccc} 2 & 0 & -1 & 0 & \dots & 0 \\ 0 & 2 & -1 & 0 & \dots & 0 \\ -1 & -1 & 2 & -1 & \dots & 0 \\ 0 & 0 & -1 & 2 & \ddots & 0 \\ \vdots & 0 & \ddots & \ddots & \ddots & -1 \\ 0 & \dots & \dots & \dots & -1 & 2 \end{array}\right)$$

III. Perfect and eutactic forms

Hermite function

▶ If $A \in S_{>0}^n$ then the arithmetic minimum is

$$\min(A) = \min_{x \in \mathbb{Z}^n - \{0\}} x^T A x$$

and the set of minimal vectors is

$$Min(A) = \left\{ x \in \mathbb{Z}^n : x^T A x = min(A) \right\}$$

▶ The Hermite function on the space $S_{>0}^n$ is

$$\gamma(A) = \frac{\min(A)}{(\det A)^{1/n}}$$

▶ The density of the lattice packing *L* associated to *A* is

$$\delta(L) = \sqrt{\gamma(A)^n} \frac{\operatorname{vol}(B_n(0,1))}{2^n}$$

► Finding lattice packings with highest packing density is the same as maximizing the Hermite function.

Perfect forms

A form A is extreme if there is a neighborhood V of A in $S_{>0}^n$ such that

If
$$B \in V$$
 with $B \neq \lambda A$ then $\gamma(B) < \gamma(A)$

▶ A matrix $A \in S_{>0}^n$ is perfect (Korkine & Zolotarev, 1873) if the equation

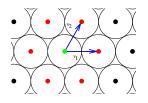
$$B \in S^n$$
 and $x^T B x = \min(A)$ for all $x \in \min(A)$

implies B = A.

- ► Theorem: (Korkine & Zolotarev, 1873) If a form is extreme then it is perfect.
- Perfect forms are rational forms.
- ▶ If A is perfect then $\gamma(A)^n$ is rational.

A perfect form

▶ $A_{hex} = \begin{pmatrix} 2 & 1 \\ 1 & 2 \end{pmatrix}$ corresponds to the lattice:



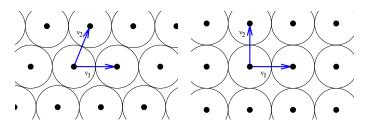
▶ If $B = \begin{pmatrix} a & b \\ b & c \end{pmatrix}$ satisfies to $x^T B x = \min(A_{hex})$ for $x \in \text{Min}(A_{hex}) = \{\pm(1,0), \pm(0,1), \pm(1,-1)\}$, then:

$$a = 2$$
, $b = 2$ and $a - 2c + b = 2$

which implies $B = A_{hex}$. A_{hex} is perfect.

A non-perfect form

- $A_{sqr} = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}$ has $Min(A_{sqr}) = \{\pm(0,1), \pm(1,0)\}.$
- ▶ See below lattices L_B , L_{sqr} associated to matrices $B, A_{sqr} \in S^2_{>0}$ with $Min(B) = Min(A_{sqr})$:



Eutactic forms

▶ A form $A \in S_{>0}^n$ is eutactic (Voronoi, 1908) if there exist $\lambda_V > 0$ such that

$$A^{-1} = \sum_{v \in \mathsf{Min}(A)} \lambda_v v v^T$$

- ▶ Theorem: (Voronoi, 1908) A form *A* is extreme if and only if it is perfect and eutactic.
- ► Theorem: (Ash, 1977)
 - (i) If A is not an eutactic form then it is topologically ordinary point for γ
 - (ii) If A is an eutactic form then it is a critical but topologically non-degenerate point for γ .
 - (ii) γ is a topological Morse function.

Examples of perfect forms

▶ The root lattice are all perfect:

Name	Min	Min	det	Aut
A_n	$e_i - e_j$	2n(n+1)	n+1	2(n+1)!
D_n	$\pm e_i \pm e_j$	4n(n-1)	4	2 ⁿ n!
E ₆	complex	72	3	103680
E ₇	complex	126	2	2903040
E ₈	complex	240	1	696729600

- Another remarkable lattice is the Leech lattice of dimension 24.
 - Every vector v has $||v||^2 \ge 4$ and $\det Leech = 1$.
 - ► There are 196280 shortest vectors (maximal number in dimension 24)
 - Its automorphism group quotiented by $\pm Id_{24}$ is the sporadic simple group Co_0
 - It plays a significant role in modular form theory and Lorentzian lattice theory.

Known results on lattice packing density maximization

dim.	Nr. of perfect forms	Absolute maximum	
		of γ realized by	
2	1 (Lagrange)	A_2	
3	1 (Gauss)	A_3	
4	2 (Korkine & Zolotarev)	D_4	
5	3 (Korkine & Zolotarev)	D_5	
6	7 (Barnes)	E ₆ (Blichfeldt)	
7	33 (Jaquet)	E ₇ (Blichfeldt)	
8	10916 (DSV)	E ₈ (Blichfeldt)	
9	≥500000	Λ_9 ?	
24	?	Leech (Cohn & Kumar)	

Remarks

- ► The enumeration of perfect forms is done with the Voronoi algorithm.
- ► The solution in dimension 24 was obtained by different methods.

and the Voronoi algorithm

IV. Ryshkov polyhedron

The Ryshkov polyhedron

▶ The Ryshkov polyhedron R_n is defined as

$$R_n = \left\{ A \in S^n \text{ s.t. } x^T A x \ge 1 \text{ for all } x \in \mathbb{Z}^n - \{0\} \right\}$$

- ▶ The cone is invariant under the action of $GL_n(\mathbb{Z})$.
- ▶ The cone is locally polyhedral, i.e. for a given $A \in R_n$

$$\left\{x \in \mathbb{Z}^n \text{ s.t. } x^T A x = 1\right\}$$

is finite

- ▶ Vertices of R_n correspond to perfect forms.
- ▶ For a form $A \in R_n$ we define the local cone

$$Loc(A) = \left\{ Q \in S^n \text{ s.t. } x^T Q x \ge 0 \text{ if } x^T A x = 1 \right\}$$

The Voronoi algorithm

- ▶ Find a perfect form (say A_n), insert it to the list \mathcal{L} as undone.
- Iterate
 - For every undone perfect form A in L, compute the local cone Loc(A) and then its extreme rays.
 - For every extreme ray r of Loc(A) realize the flipping, i.e. compute the adjacent perfect form $A' = A + \alpha r$.
 - ▶ If A' is not equivalent to a form in \mathcal{L} , then we insert it into \mathcal{L} as undone.
- Finish when all perfect domains have been treated.

The subalgorithms are:

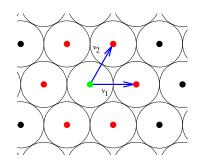
- Find the extreme rays of the local cone Loc(A) (use cdd or Irs or any other program)
- For any extreme ray r of Loc(A) find the adjacent perfect form A' in the Ryshkov polyhedron R_n
- ► Test equivalence of perfect forms using autom

Flipping on an edge I

$$\mathsf{Min}(A_{hex}) = \{\pm(1,0), \pm(0,1), \pm(1,-1)\}$$

with

$$A_{hex}=\left(egin{array}{cc} 1 & 1/2 \ 1/2 & 1 \end{array}
ight) \ \ {
m and} \ \ D=\left(egin{array}{cc} 0 & -1 \ -1 & 0 \end{array}
ight)$$





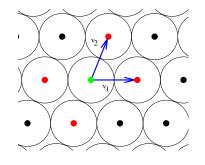
Ahex

Flipping on an edge II

$$\mathsf{Min}(B) = \{\pm(1,0), \pm(0,1)\}$$

with

$$B = \begin{pmatrix} 1 & 1/4 \\ 1/4 & 1 \end{pmatrix} = A_{hex} + D/4$$



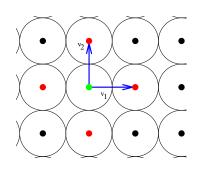


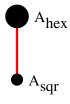
Flipping on an edge III

$$Min(A_{sqr}) = \{\pm(1,0),\pm(0,1)\}$$

with

$$A_{sqr} = \left(egin{array}{cc} 1 & 0 \\ 0 & 1 \end{array}
ight) = A_{hex} + D/2$$



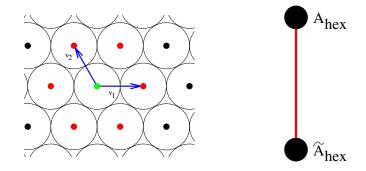


Flipping on an edge IV

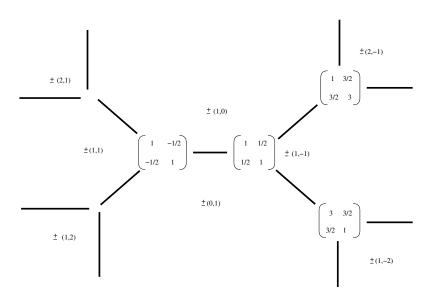
$$Min(\tilde{A}_{hex}) = \{\pm(1,0), \pm(0,1), \pm(1,1)\}$$

with

$$\widetilde{A}_{hex} = \left(egin{array}{cc} 1 & -1/2 \ -1/2 & 1 \end{array}
ight) = A_{hex} + D$$



The Ryshkov polyhedron R_2



Well rounded forms and retract

- A form Q is said to be well rounded if it admits vectors v₁, ..., vn such that
 - (v_1, \ldots, v_n) form a basis of \mathbb{R}^n
 - \triangleright v_1, \ldots, v_n are shortest vectors.
 - $P Q[v_1] = \cdots = Q[v_n].$
- ▶ Well rounded forms correspond to bounded faces of R_n .
- ▶ Every form can be continuously deformed to a well rounded form and this defines a retracting homotopy of R_n onto a polyhedral complex WR_n of dimension $\frac{n(n-1)}{2}$.
- ▶ Every face of WR_n has finite stabilizer, hence we can use it for computing the homology of $GL_n(\mathbb{Z})$ and other arithmetic groups.
- Actually, in term of dimension, we cannot do better:
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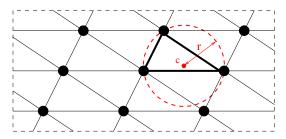
V. The lattice covering problem

Empty sphere and Delaunay polytopes

A sphere S(c, r) of radius r and center c in an n-dimensional lattice L is said to be an empty sphere if:

- (i) $||v-c|| \ge r$ for all $v \in L$,
- (ii) the set $S(c,r) \cap L$ contains n+1 affinely independent points.

A Delaunay polytope P in a lattice L is a polytope, whose vertex-set is $L \cap S(c, r)$.



Equalities and inequalities

- ▶ Take $M = G_v$ with $v = (v_1, ..., v_n)$ a basis of lattice L.
- ▶ If $V = (w_1, ..., w_N)$ with $w_i \in \mathbb{Z}^n$ are the vertices of a Delaunay polytope of empty sphere S(c, r) then:

$$||w_i - c|| = r$$
 i.e. $w_i^T M w_i - 2 w_i^T M c + c^T M c = r^2$

Substracting one obtains

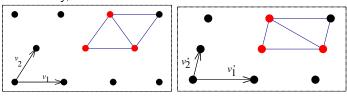
$$\{w_{i}^{T}Mw_{i} - w_{j}^{T}Mw_{j}\} - 2\{w_{i}^{T} - w_{j}^{T}\}Mc = 0$$

- ▶ Inverting matrices, one obtains $Mc = \psi(M)$ with ψ linear and so one gets linear equalities on M.
- ▶ Similarly $||w c|| \ge r$ translates into linear inequalities on M: Take $V = (v_0, \ldots, v_n)$ a simplex $(v_i \in \mathbb{Z}^n)$, $w \in \mathbb{Z}^n$. If one writes $w = \sum_{i=0}^n \lambda_i v_i$ with $1 = \sum_{i=0}^n \lambda_i$, then one has

$$||w-c|| \ge r \Leftrightarrow w^T M w - \sum_{i=0}^n \lambda_i v_i^T M v_i \ge 0$$

Iso-Delaunay domains

- ▶ Take a lattice L and select a basis v_1, \ldots, v_n .
- ▶ We want to assign the Delaunay polytopes of a lattice. Geometrically, this means that



are part of the same iso-Delaunay domain.

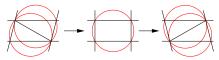
▶ An iso-Delaunay domain is the assignment of Delaunay polytopes, so it is also the assignment of the Voronoi polytope of the lattice.

Primitive iso-Delaunay

- ▶ If one takes a generic matrix M in $S_{>0}^n$, then all its Delaunay are simplices and so no linear equality are implied on M.
- ► Hence the corresponding iso-Delaunay domain is of dimension $\frac{n(n+1)}{2}$, they are called primitive

Equivalence and enumeration

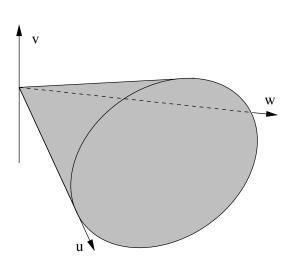
- ▶ The group $GL_n(\mathbb{Z})$ acts on $S_{>0}^n$ by arithmetic equivalence and preserve the primitive iso-Delaunay domains.
- Voronoi proved that after this action, there is a finite number of primitive iso-Delaunay domains.
- ▶ Bistellar flipping creates one iso-Delaunay from a given iso-Delaunay domain and a facet of the domain. In dim. 2:



- ► Enumerating primitive iso-Delaunay domains is done classicaly:
 - Find one primitive iso-Delaunay domain.
 - ▶ Find the adjacent ones and reduce by arithmetic equivalence.
- ▶ This is very similar to the Voronoi algorithm for perfect forms.

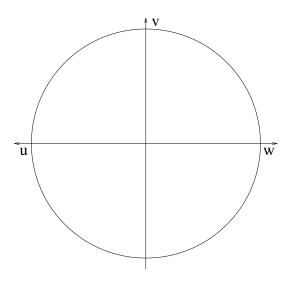
The partition of $S^2_{>0} \subset \mathbb{R}^3$ I

If $q(x,y) = ux^2 + 2vxy + wy^2$ then $q \in S_{>0}^2$ if and only if $v^2 < uw$ and u > 0.



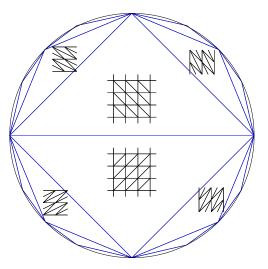
The partition of $S^2_{>0} \subset \mathbb{R}^3$ II

We cut by the plane $\mathrm{u}+\mathrm{w}=1$ and get a circle representation.



The partition of $S^2_{>0}\subset \mathbb{R}^3$ III

Primitive iso-Delaunay domains in $S_{>0}^2$:



Optimization problem

- The lattice covering problem is to find a lattice covering of minimal density.
- ► Thm. Given an iso-Delaunay domain *LT*, there exist a unique lattice, which minimize the covering density over *LT*.
- ► The effective lattice is obtained by solving a semidefinite programming problem, so no exact solution, but approximate solutions available at any precision.
- ► The local maxima that are found are defined by algebraic integers.
- See for more details
 - A. Schürmann and F. Vallentin, *Computational approaches to lattice packing and covering problems*, Discrete & Compututational Geometry **35** (2006) 73–116.
 - A. Schürmann, Computational geometry of positive definite quadratic forms, University Lecture Notes, AMS.

Known results on covering density minimization

dim.	Best covering	nr of iso-Delaunay
2	A ₂ (Kershner)	1 (Voronoi)
3	A_3^* (Bambah)	1(Voronoi)
4	A ₄ (Delone & Ryshkov)	3(Voronoi)
5	A ₅ (Ryshkov & Baranovski)	222(Engel)
6	L_6 (conj. Vallentin)?	?
7	L ₇ (conj. Schürmann & Vallentin)?	?
24	Leech (conj.)?	?

- It turn out that the lattice of minimal covering density are unique for $n \le 5$
- In general the best lattice coverings are expected to be non-rational and with low symmetry.
- ▶ But experimentations seemed to indicate that E₆ is a local covering maxima.

VI. Quadratic functions

and the Erdahl cone

The Erdahl cone

▶ Denote by $E_2(n)$ the vector space of degree 2 polynomial functions on \mathbb{R}^n . We write $f \in E_2(n)$ in the form

$$f(x) = a_f + b_f \cdot x + Q_f[x]$$

with $a_f \in \mathbb{R}$, $b_f \in \mathbb{R}^n$ and Q_f a $n \times n$ symmetric matrix

▶ The Erdahl cone is defined as

$$Erdahl(n) = \{ f \in E_2(n) \text{ such that } f(x) \ge 0 \text{ for } x \in \mathbb{Z}^n \}$$

- It is a convex cone, which is non-polyhedral since defined by an infinity of inequalities.
- ▶ The group acting on Erdahl(n) is $AGL_n(\mathbb{Z})$, i.e. the group of affine integral transformations

$$x \mapsto b + Px$$
 for $b \in \mathbb{Z}^n$ and $P \in GL_n(\mathbb{Z})$

Scalar product

▶ Definition: If $f, g \in E_2(n)$, then:

$$\langle f,g \rangle = a_f a_g + \langle b_f,b_g \rangle + \langle Q_f,Q_g \rangle$$

- ▶ Definition: For $v \in \mathbb{Z}^n$, define $ev_v(x) = (1 + v \cdot x)^2$.
- We have

$$\langle f, ev_v \rangle = f(v)$$

- Thus finding the rays of Erdahl(n) is a dual description problem with an infinity of inequalities and infinite group acting on it.
- ▶ If $f \in Erdahl(n)$ then Q_f is positive semidefinite.
- ▶ Definition: We also define

$$Erdahl_{>0}(n) = \{ f \in Erdahl(n) : Q_f \text{ positive definite} \}$$

Relation with Delaunay polytope

▶ If *D* is a Delaunay polytope of a lattice $L = \mathbb{Z}v_1 + \cdots + \mathbb{Z}v_n$ of empty sphere S(c, r) then we define the function

$$f_{D,\mathbf{v}}: \mathbb{Z}^n \to \mathbb{R}$$

$$x = (x_1, \dots, x_n) \mapsto \|\sum_{i=1}^n x_i v_i - c\|^2 - r^2$$

Clearly $f_{D,\mathbf{v}} \in Erdahl_{>0}(n)$.

- ► The perfection rank of a Delaunay polytope is the dimension of the face it defines in *Erdahl(n)*.
- ▶ Definition: If $f \in Erdahl(n)$ then

$$Z(f) = \{ v \in \mathbb{Z}^n : f(v) = 0 \}$$

▶ Theorem: If $f \in Erdahl(n)$ then there exist a lattice L_f and a lattice L' containing a Delaunay polytope D_f such that

$$Z(f) = D_f + L_f$$

▶ We have dim L' + dim $L_f \le n$. In case of equality Z(f) is called a Delaunay polyhedra.

Perfect Delaunay polytopes/polyhedra

▶ Definition: If *D* is a *n*-dimensional Delaunay polyhedra then we define

$$\mathsf{Dom}_{\mathbf{v}} \ D = \sum_{v\mathbf{v} \in D} \mathbb{R}_{+} e v_{v}$$

- We have $\langle f_{D,\mathbf{v}}, \mathsf{Dom}_{\mathbf{v}} \ D \rangle = 0$.
- ▶ Definition: D is perfect if Dom D is of dimension $\binom{n+2}{2} 1$ that is if the perfection rank is 1.
- ▶ This implies that f_D is an extreme rays of Erdahl(n) and f_D is rational.
- A perfect *n*-dimensional Delaunay polytope has at least $\binom{n+2}{2}-1$ vertices. There is only one way to embed it as a Delaunay polytope of a lattice.
- Perfect Delaunay polytopes are remarkable and rare objects that we want to enumerate.

Perfect Delaunay polytope

► There is a finite number of them in each dimension *n*. Known results:

suits.		
dim.	perfect Delaunay	authors
1	$[0,1]$ in $\mathbb Z$	
2	Ø	
3	Ø	
4	Ø	
5	Ø	↑ (Deza, Laurent & Grishukhin)
6	2 ₂₁ in E ₆	(Deza & Dutour)
7	3 ₂₁ in E ₇	
	and ER_7 in $L(ER_7)$	
8	≥ 27	(Dutour Sikiric & Rybnikov)
9	≥ 100000	(Dutour Sikiric)

- ▶ Theorem: There exist perfect Delaunay polytopes D such that $\mathbb{Z}D \neq \mathbb{Z}^n$.
- ► Theorem: There exist lattices with several perfect Delaunay polytopes.
- ▶ Theorem: For $n \ge 6$ there exist a perfect Delaunay polytope with exactly $\binom{n+2}{2} 1$ vertices.

Extreme rays of Erdahl(n)

▶ Definition: If $f \in Erdahl_{>0}(n)$ then we define

$$Dom f = \sum_{v \in Z(f)} \mathbb{R}_+ ev_v$$

- We have $\langle f, \mathsf{Dom} \ f \rangle = 0$.
- ▶ Erdahl, 1992: The extreme rays of Erdahl(n) are
 - (a) The constant function 1.
 - (b) The functions

$$(a_1x_1+\cdots+a_nx_n+\beta)^2$$

with (a_1, \ldots, a_n) not collinear to an integral vector.

- (c) The functions f such that Z(f) is a perfect Delaunay polyhedra.
- Note that if $f \in Erdahl(n)$ with Z(f) a Delaunay polyhedra, then there exist a lattice L' of dimension $k \leq n$, a Delaunay polytope D of L', a basis \mathbf{v}' of L' and a function $\phi \in \mathsf{AGL}_n(\mathbb{Z})$ such that

$$f \circ \phi(x_1,\ldots,x_n) = f_{D,\mathbf{v}'}(x_1,\ldots,x_k)$$

Delaunay polyhedra retract

- For a function $f \in Erdahl(n)$ a proper decomposition is a pair (g,h) with f = g + h, $g \in Erdahl(n)$ and $h(x) \ge 0$ for $x \in \mathbb{R}^n$.
- ▶ Lemma: For a proper decomposition we have

$$Vect Z(f) + Ker Q_f \subset Ker Q_h$$

and there exist a proper decomposition with equality.

- ▶ Fix an integral complement L' of $Vect\ Z(f) + Ker\ Q_f$. A proper decomposition is called extremal if det $Q_h|_{L'}$ is maximal among all proper decompositions.
- ▶ Theorem: For $f \in Erdahl(n)$, there exist a unique extremal decomposition. For it we have that Z(g) is a delaunay polyhedra.
- ► Conjecture: The decomposition depends continuously on $f \in Erdahl(n)$.
- ▶ On the other hand in a neighborhood of $f \in Erdahl(n)$ we can have an infinity of Delaunay polyhedra.

Voronoi algorithm on the Erdahl cone

► From a given n-dimensional perfect Delaunay polytope Q of form f we can define the local cone

$$Loc(f) = \{g \in E_2(n) \text{ s.t. } g(x) \ge 0 \text{ for } x \in Z(f)\}$$

- ► The flipping algorithm finds the adjacent quadratic perfect form *g* from a given perfect form *f*.
- ▶ The problem is Erdahl(n) is not locally polyhedral, i.e. the rank of g can be lower than n.
- ▶ The technique is to use a recursive algorithm for realizing the enumeration. We start form $[0,1] \times \mathbb{R}^{n-1}$ and by subdivizion reach $[0,1]^n$ (its local cone is the cut cone CUT_{n+1} occuring in combinatorial optimization).

VII. Covering maxima, pessima and their characterization

Eutacticity

▶ If $f \in Erdahl_{>0}(n)$ then define μ_f and c_f such that

$$f(x) = Q_f[x - c_f] - \mu_f$$

Then define

$$u_f(x) = (1 + c_f \cdot x)^2 + \frac{\mu_f}{n} Q_f^{-1}[x]$$

- ▶ Definition: $f \in Erdahl_{>0}(n)$ is eutactic if u_f is in the relative interior of Dom f.
- ▶ Definition: Take a Delaunay polytope P for a quadratic form Q of center c_P and square radius μ_P . P is called eutactic if there are $\alpha_V > 0$ so that

$$\begin{cases} 1 &= \sum_{v \in \text{vert } P} \alpha_v, \\ 0 &= \sum_{v \in \text{vert } P} \alpha_v(v - c_P), \\ \frac{\mu_P}{n} Q^{-1} &= \sum_{v \in \text{vert } P} \alpha_v(v - c_P)(v - c_P)^T. \end{cases}$$

Covering maxima

- A given lattice L is called a covering maxima if for any lattice L' near L we have $\Theta(L') < \Theta(L)$.
- ▶ Theorem: The following are equivalent:
 - ▶ L is a covering maxima
 - Every Delaunay polytope of maximal circumradius is perfect and eutactic.
- ▶ The following are perfect Delaunay polytope:

name	# vertices	\mid $\#$ orbits Delaunay polytopes \mid
E ₆	27	1
E_7	56	2
ER ₇	35	4
O_{10}	160	6
BW_{16}	512	4
O_{23}	94208	5
Λ_{23}	47104	709

▶ Theorem: For any $n \ge 6$ there exist one lattice $L(DS_n)$ which is a covering maxima.

There is only one perfect Delaunay polytope $P(DS_n)$ of maximal radius in $L(DS_n)$.

The infinite series

- ▶ For n even $P(DS_n)$ is defined as the lamination over D_{n-1} of
 - one vertex
 - ▶ the half cube $\frac{1}{2}H_{n-1}$
 - the cross polytope CP_{n-1}

For n = 6, it is E_6 .

- ▶ For *n* odd as the lamination over D_{n-1} of
 - ightharpoonup the cross polytope CP_{n-1}
 - ▶ the half cube $\frac{1}{2}H_{n-1}$
 - ▶ the cross polytope CP_{n-1}

For n = 7, it is E_7 .

- ▶ Conjecture: The lattice DS_n has the following properties:
 - L(DS_n) has the maximum covering density among all covering maxima
 - ▶ Among all perfect Delaunay polytopes, $P(DS_n)$ has
 - maximum number of vertices
 - maximum volume

If true this would imply Minkovski conjecture.

Pessimum and Morse function property

- ▶ For a lattice L let us denote $D_{crit}(L)$ the space of direction d of deformation of L such that Θ increases in the direction d.
- ▶ Definition: A lattice L is said to be a covering pessimum if the space D_{crit} is of measures 0.
- ► Theorem: If a lattice L has all its Delaunay polytopes of maximum circumradius are eutactic and are not simplices then Q is a pessimum.

name	# vertices	# orbits Delaunay polytopes
\mathbb{Z}^n	2 ⁿ	1
D_4	8	1
$D_n \ (n \geq 5)$	2^{n-1}	2
E ₆ *	9	1
E ₇ *	16	1
E ₈	16	2
K_{12}	81	4

▶ Theorem: The covering density function $Q \mapsto \Theta(Q)$ is a topological Morse function if and only if $n \leq 3$.

THANK YOU