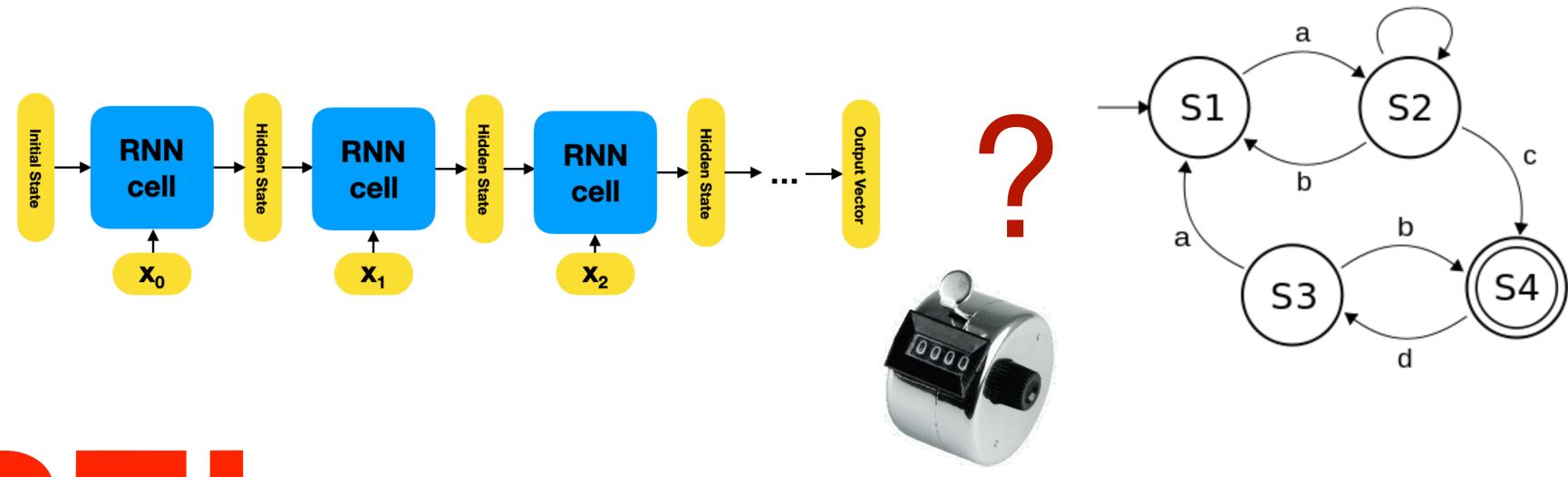
## Theoretical Properties of RNNs

#### Gail Weiss







- The state machine relation
- Theoretical Power of RNNs: Infinite Precision and Time
- Theoretical Power of RNNs: Finite Precision and Time

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  - A Completely Made Up Architecture (Failed RNN 😔)
- Augmentations (Stack RNNs)

The state machine relation

Simple state machines...

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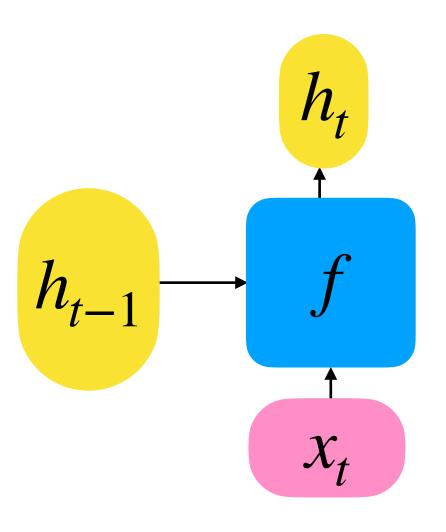
Simple state machines...

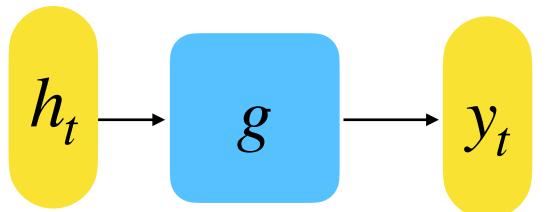
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Turing complete!?

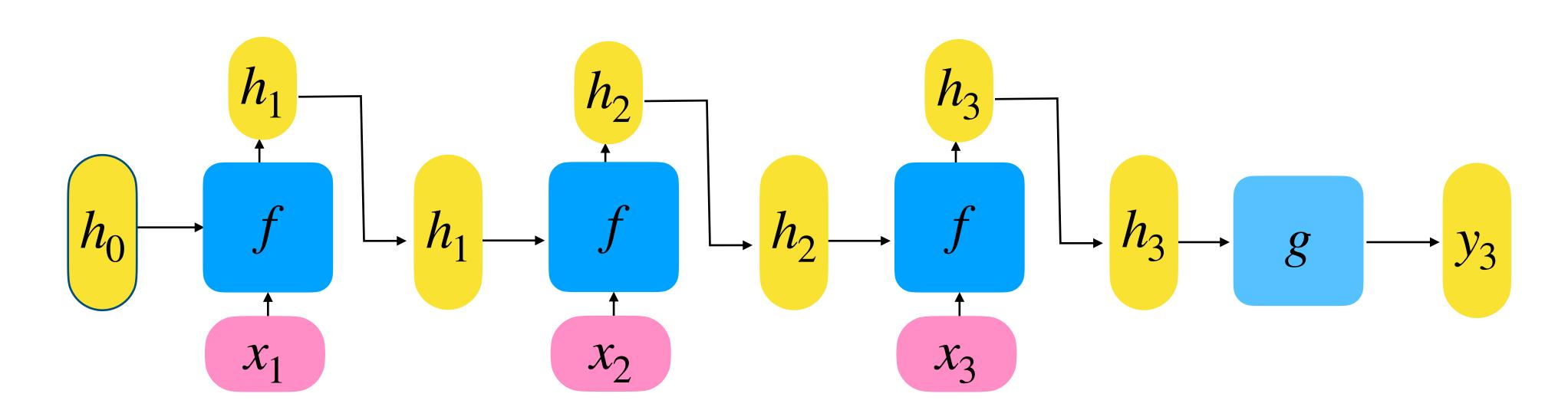
- Theoretical Power of RNNs: Finite Precision and Time Oh, state machines... with extras
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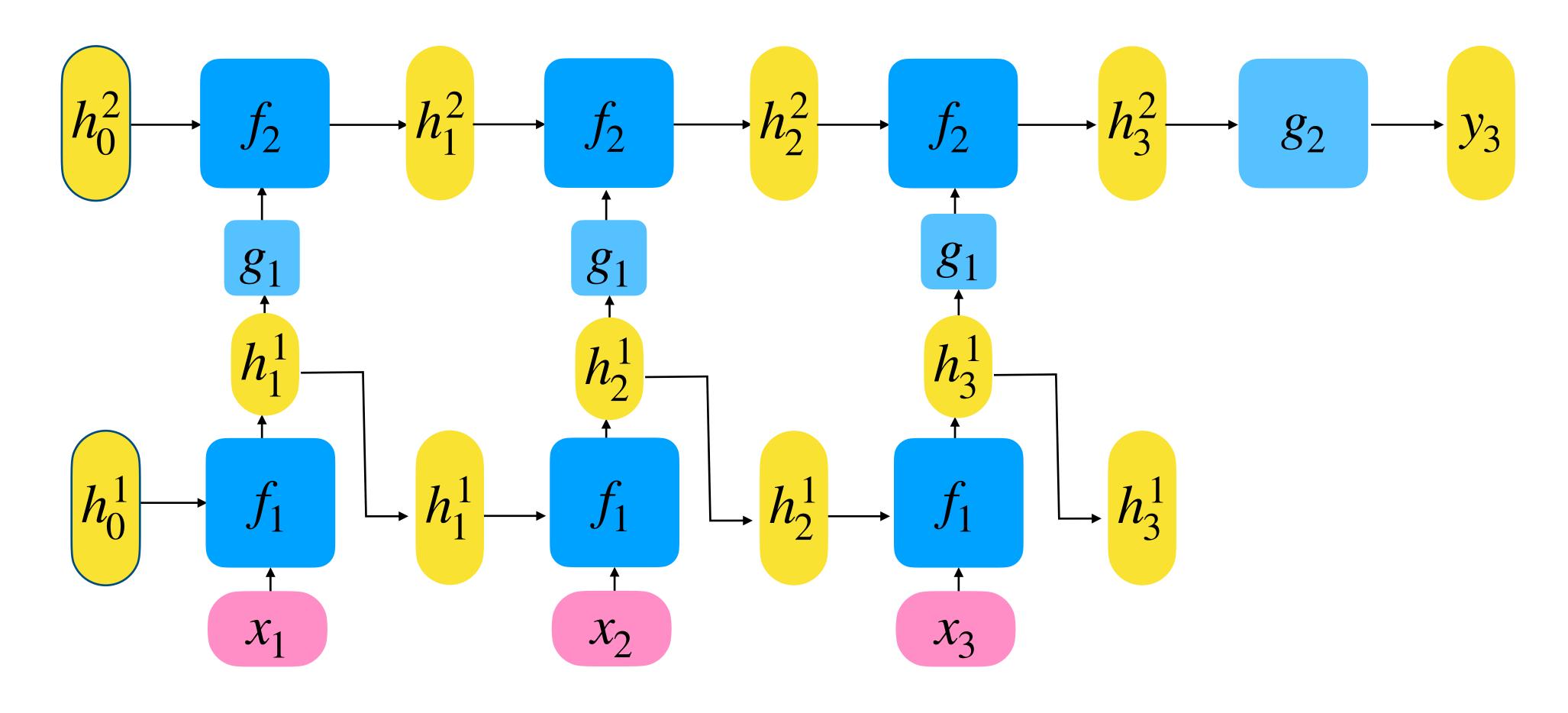


 $h_t \in \mathbb{R}^d$  for some fixed dimension d



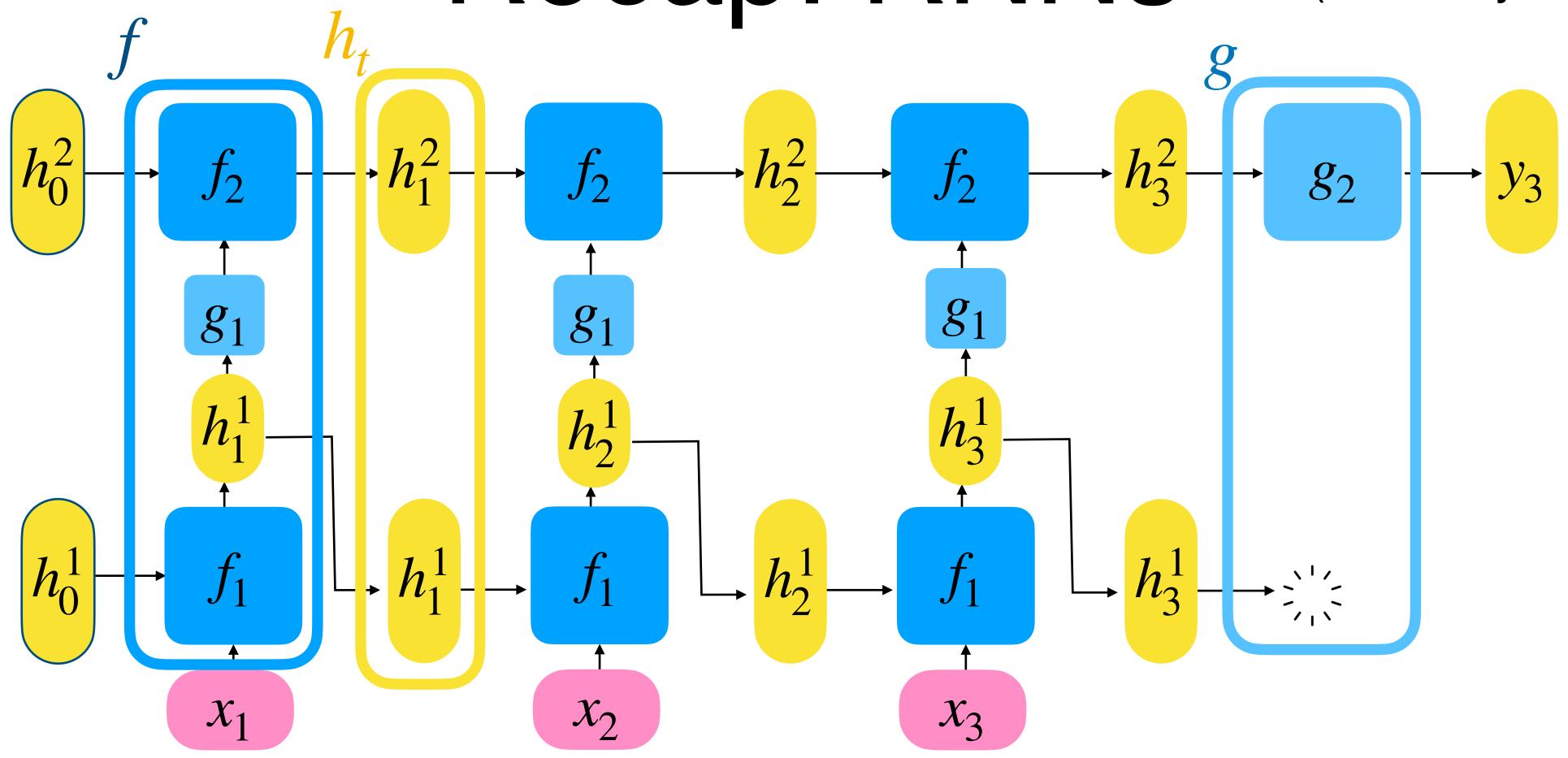
$$R = (h_0, f, g) \qquad \begin{aligned} h_t &= f(h_{t-1}, x_t) \\ y_t &= g(h_t) \end{aligned} \qquad h_t \in \mathbb{R}^d \text{ for some fixed dimension } d$$

(Multi-Layer RNNs)

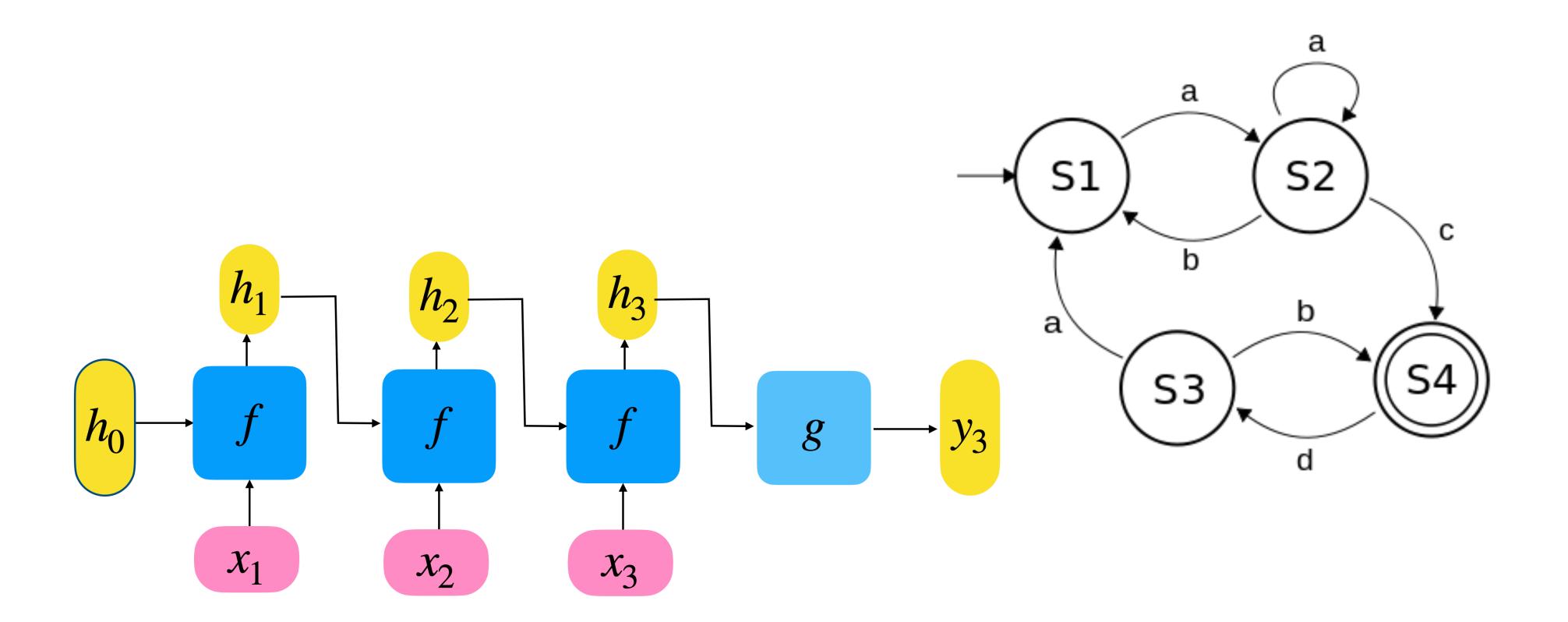


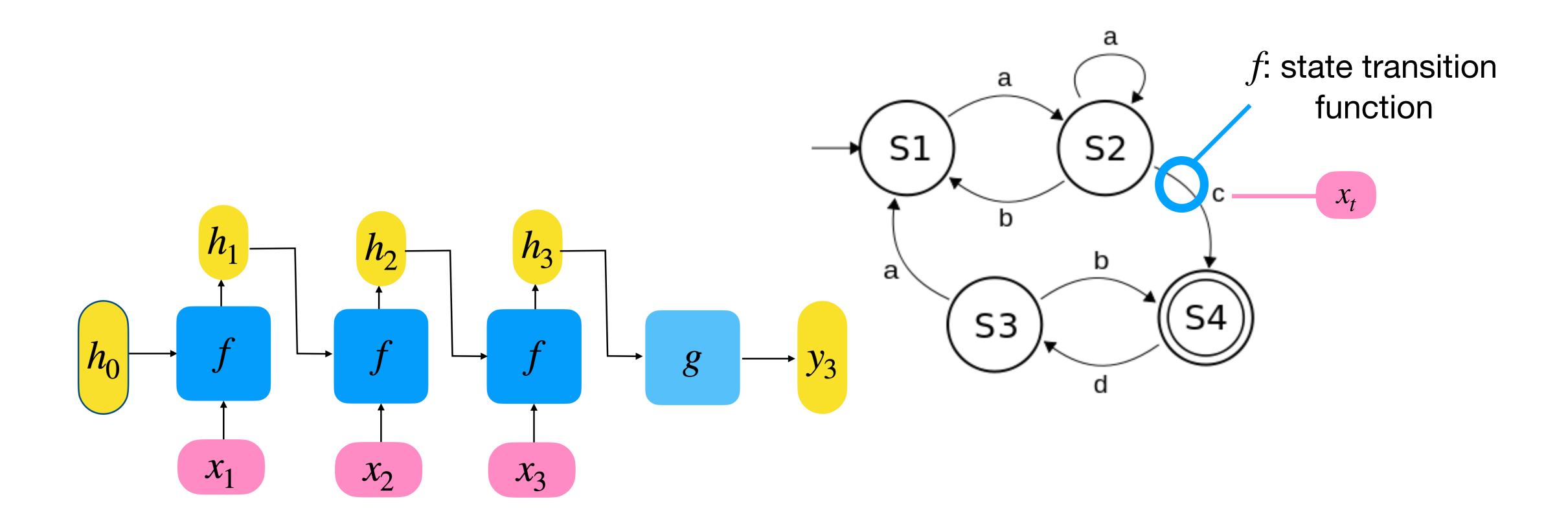
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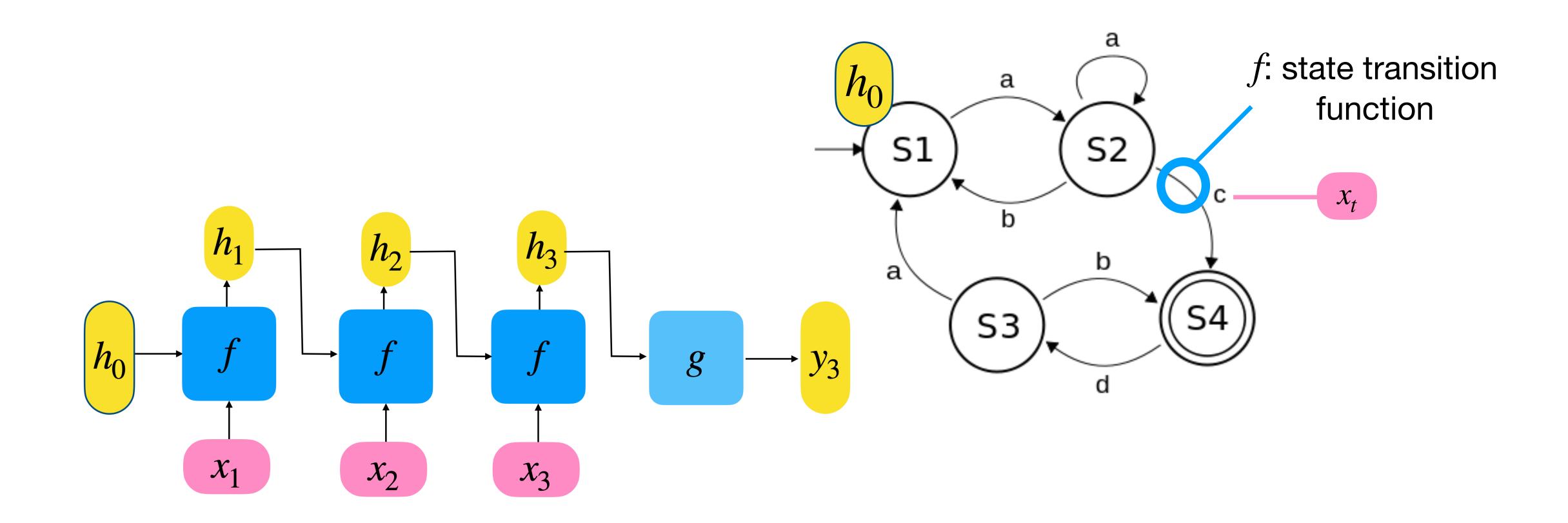
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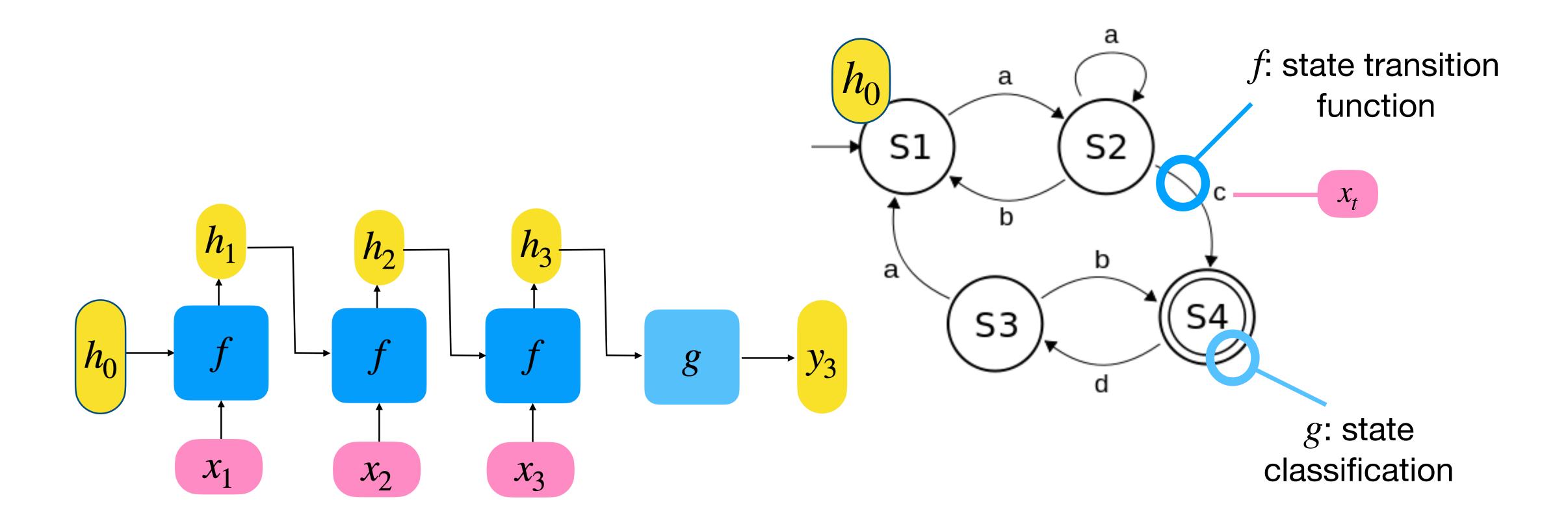


$$h_t \in \mathbb{R}^d$$
 for some fixed dimension  $d$ 

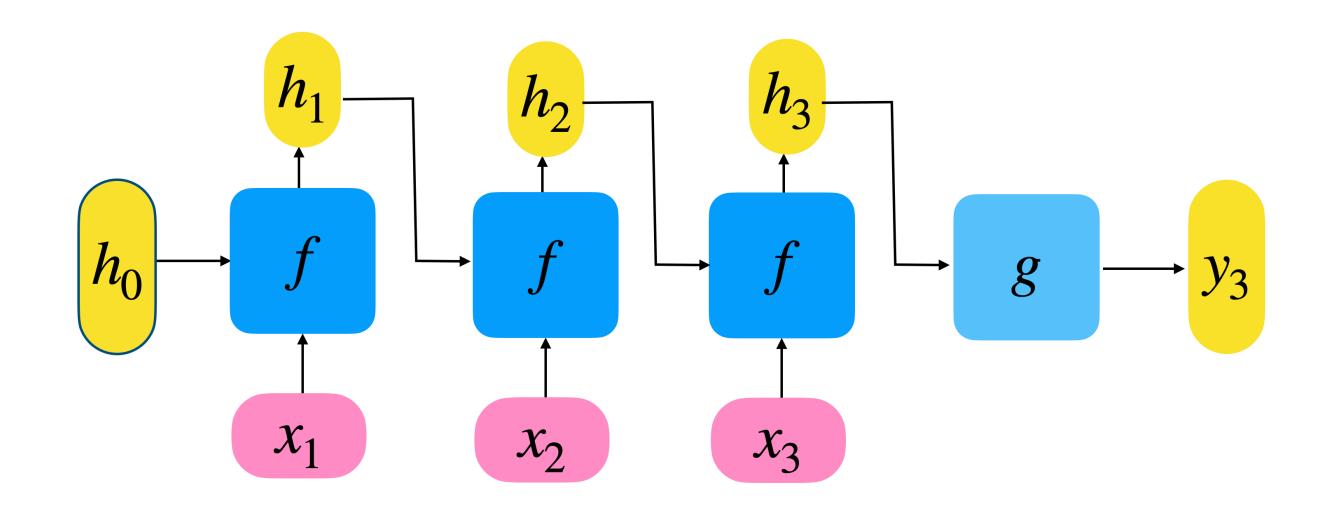






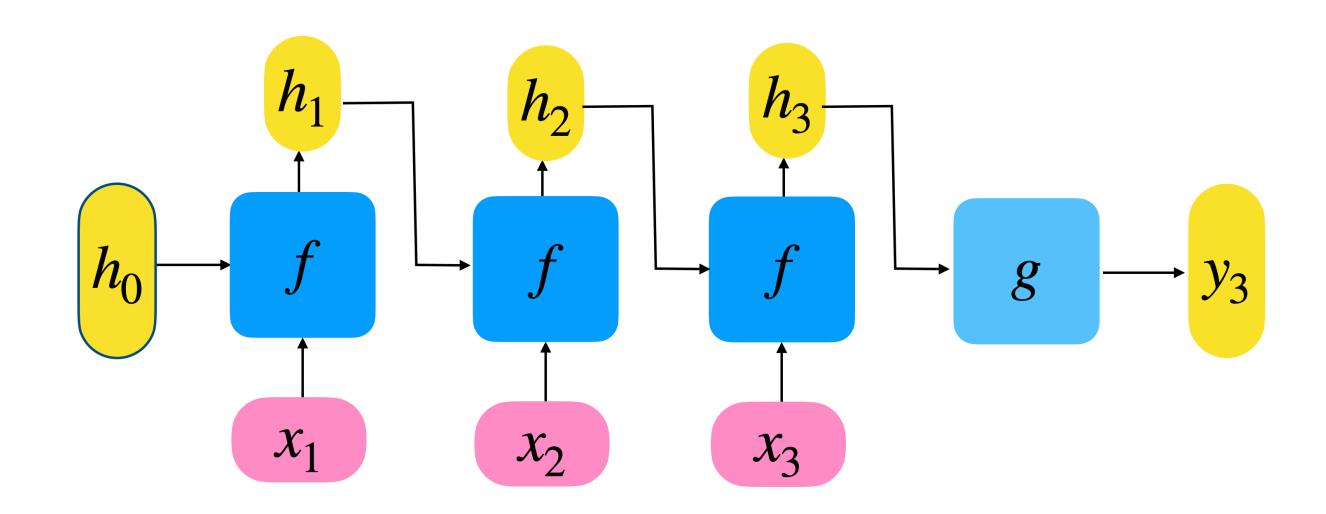


Several works use this parallel as inspiration to extract deterministic finite automata from trained RNNs!



Wait, are RNNs just deterministic finite automata?:/

### What's in a state?



Wait, are RNNs just deterministic finite automata?:/

Let's assume  $h_t \in \mathbb{Q}^d$  for some fixed d, and  $f: \mathbb{Q}^d \to \mathbb{Q}^d$  perfectly precise

- 1. How big, potentially, is the set of reachable states?
- 2. How strong, potentially, is such an RNN?

A simple RNN, with (1) rational states and weights and (2) clipped ReLU (piecewise linear sigmoid), is Turing complete. Siegelmann and Sonntag, 1992

$$h_{t+1} = \sigma(Wx_t + Uh_t + b)$$

$$\sigma(x) = \begin{cases} 0 & : x \le 0 \\ x & : 0 < x < 1 \\ 1 & : 1 \le x \end{cases}$$

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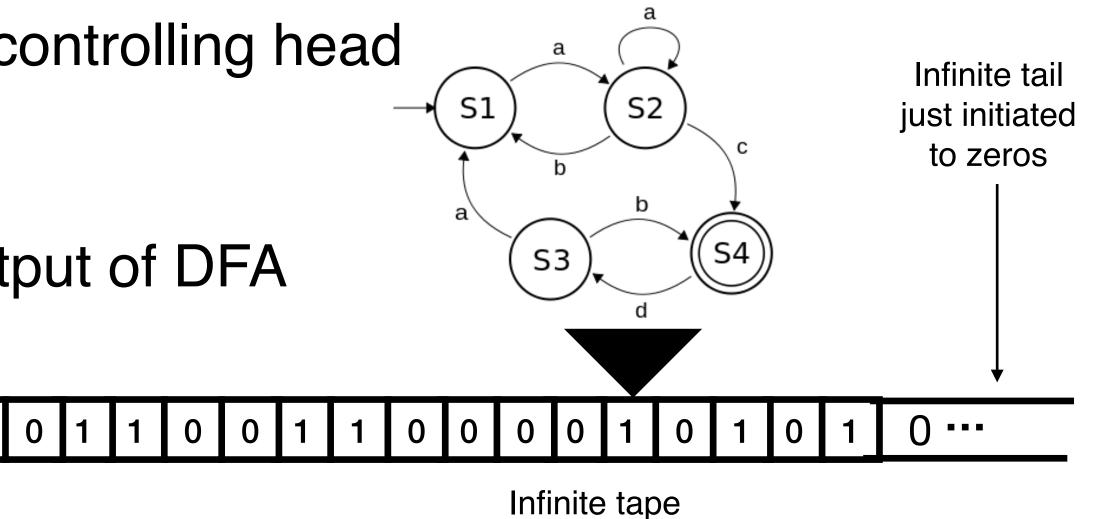
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Theory time! Buckle up

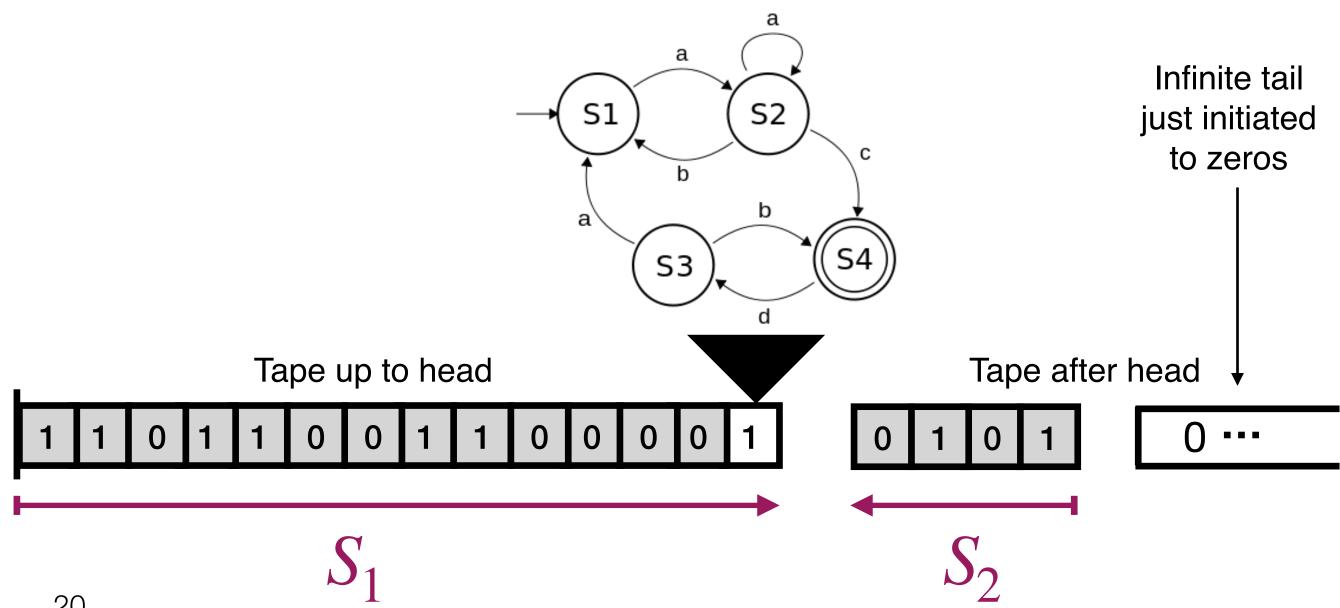
#### Recall: Turing Machines (most basic version)

- Infinite tape (infinite in one direction), binary (for simplicity)
- Head moving on and reading from/writing to tape
- Controller DFA responding to tape and controlling head
  - Current head value is input to DFA
  - Head movements and writes are output of DFA



The tape can be described with two stacks:

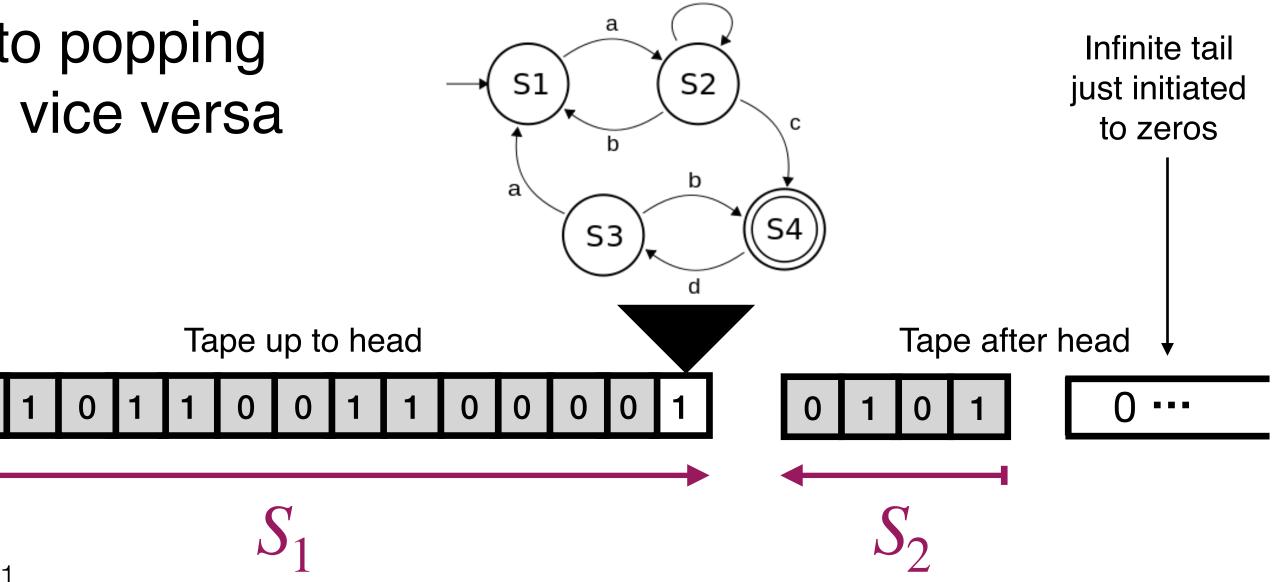
- One from the bottom of the tape, up to the head  $(S_1)$
- One from the highest location written to in the tape, down to the head  $(S_2)$



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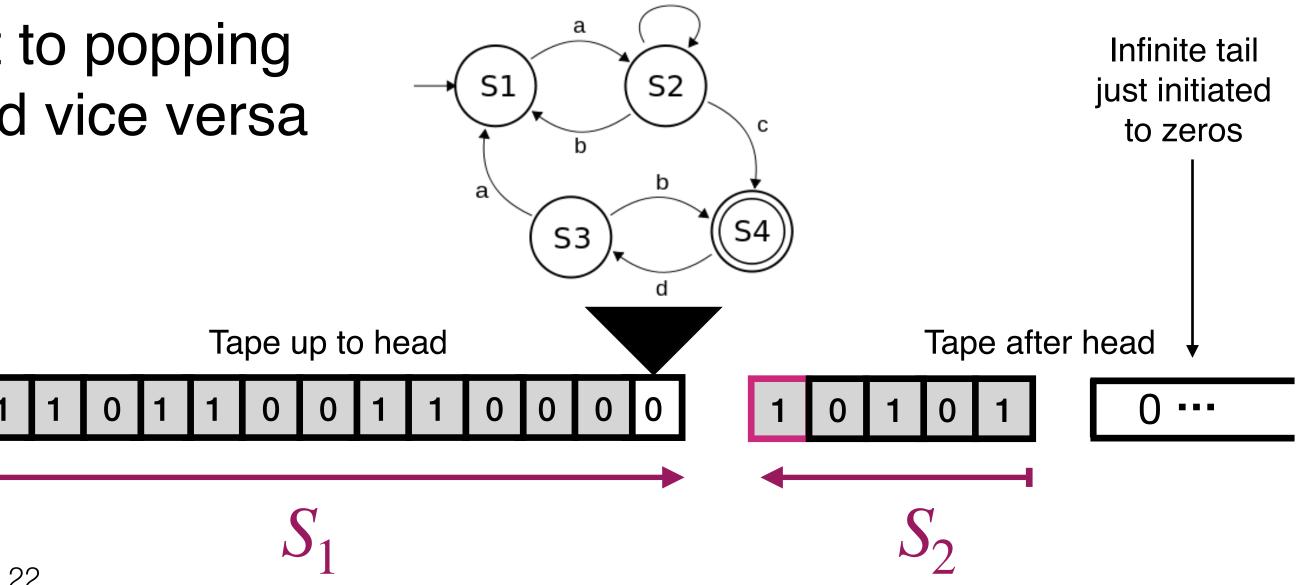
Moving left or right on the tape is equivalent to popping from one stack and pushing to the other, and vice versa



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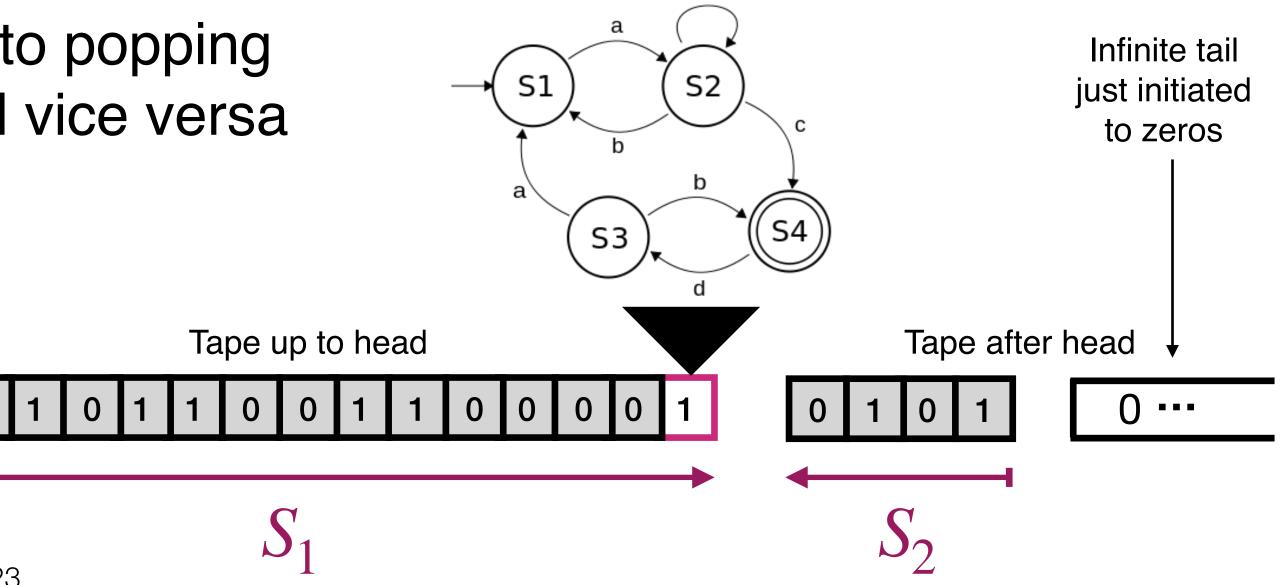
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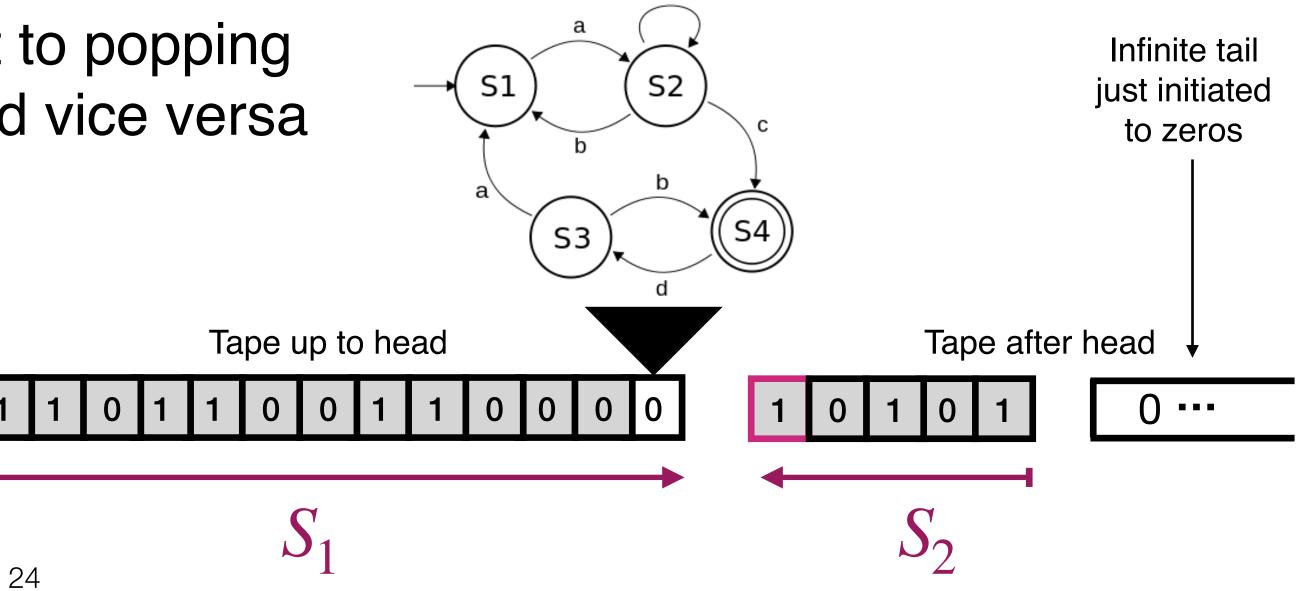
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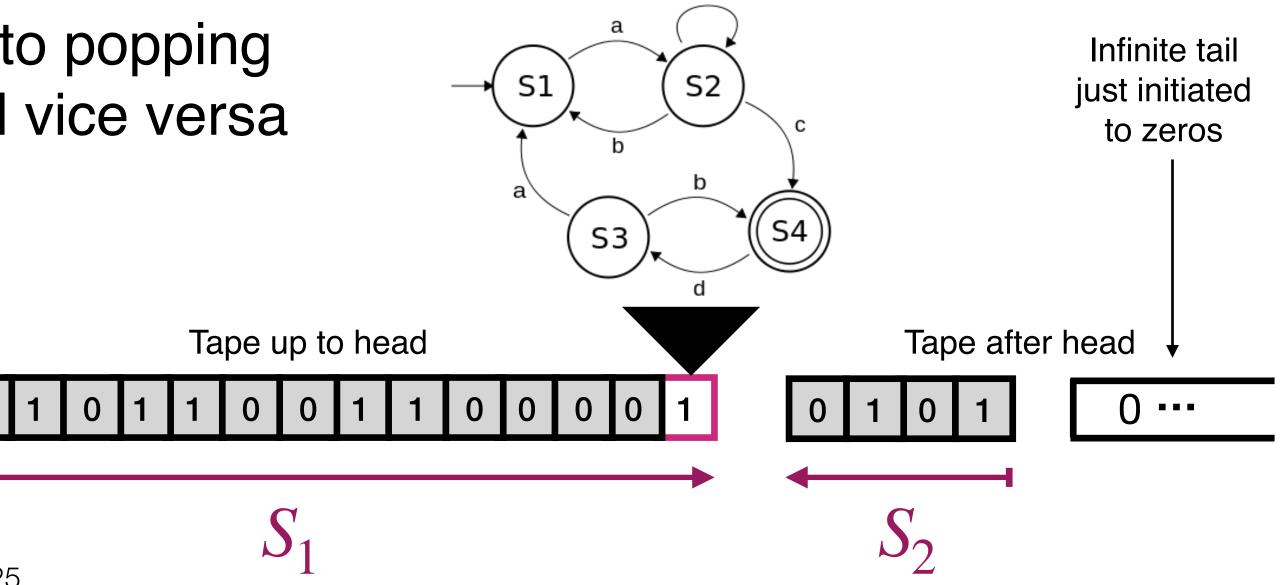
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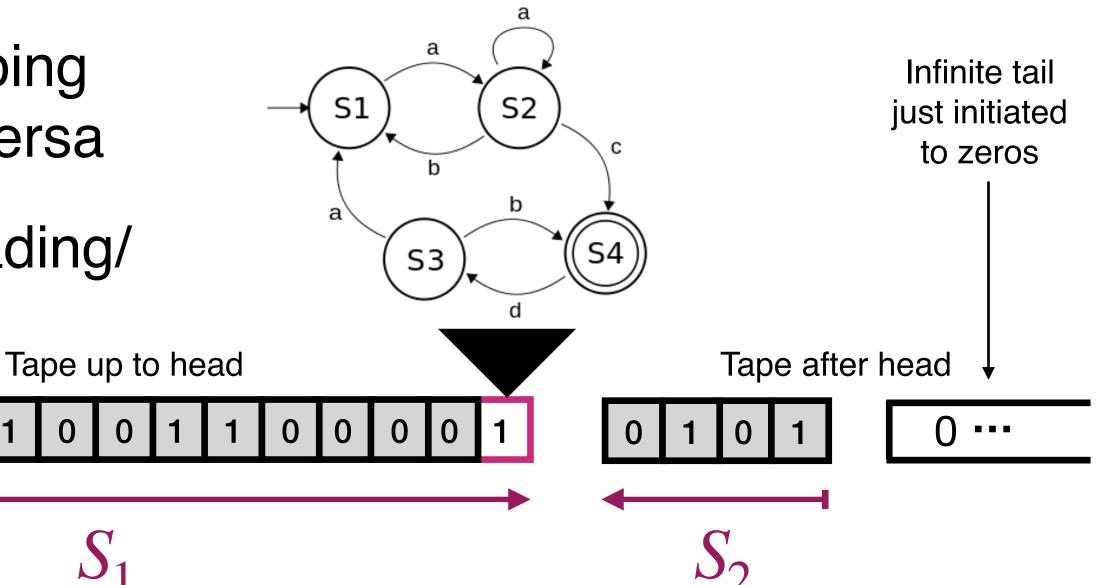


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Reading from/writing to the tape is equivalent to reading/writing the top value in  $S_1$ 



27

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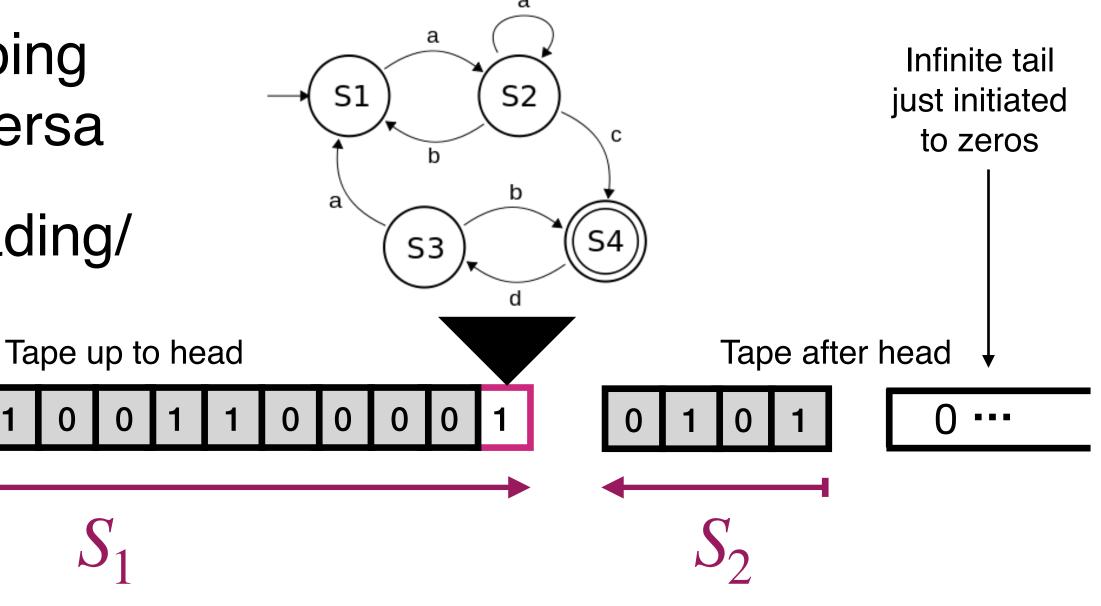
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The Turing completeness proof hinges on simulating  $S_1$  and  $S_2$ 



The stacks  $S_1$  and  $S_2$  can be encoded with a pair of fractions,  $q_1$  and  $q_2$ 

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- To recover the top value z from non-empty  $q_i$ , we compute  $z=\sigma(4q_i-2)$

$$= \sigma(q_i^{prev} + 2z - 1)$$

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- To recover the top value z from non-empty  $q_i$ , we compute  $z=\sigma(4q_i-2)$
- To pop the (recovered) top value z from  $q_i$ , we compute  $q_i^\prime = 4q_i 2z 1$

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- Both stacks are initially empty, and for this  $q_1$  and  $q_2$  are initiated to 0
- To push  $z\in\{0,1\}$  into  $q_i$ , we compute:  $q_i'=\frac{q_i}{4}+\frac{z}{2}+\frac{1}{4}$  Why add 1/4? And why compute  $q_i/4$  instead of  $q_i/2$ ?
- To recover the top value z from non-empty  $q_i$ , we compute  $z = \sigma(4q_i 2)$
- To pop the (recovered) top value z from  $q_i$ , we compute  $q_i' = 4q_i 2z 1$

We can simulate an entire Turing machine in our RNN by:

- 1. Using specific dimensions in the state  $h_t$  for:
  - 1. The stacks,  $q_1$  and  $q_2$
  - 2. The current head values, z, and value to be written, z', and
  - 3. The controller states
- 2. Padding the input, to allow the RNN time at each 'actual' input token to compute the full Turing machine update (read, decide updates, write, move)

So RNNs are Turing complete!

But we also just said they were plain state machines...

(Well, we didn't say they were *finite* state machines)

Still, where's the truth?

Do you see any practical problems with the stack fractions  $q_1$  and  $q_2$  ?

$$q_i' = \frac{q_i}{4} + \frac{z}{2} + \frac{1}{4}$$

push 
$$z$$
 to  $S_i$ 

$$q_i' = 4q_i - 2z - 1$$

pop 
$$z$$
 from  $S_i$ 

$$z = \sigma(4q_i - 2)$$

read from top of  $S_i$ 

# High Level: Turing Completeness of RNNs

Do you see any practical problems with the stack fractions  $q_1$  and  $q_2$ ?

$$q_i' = \frac{q_i}{4} + \frac{z}{2} + \frac{1}{4}$$

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$$z = \sigma(4q_i - 2)$$

push z to  $S_i$ 

pop z from  $S_i$ 

read from top of  $S_i$ 

Do we normally provide RNNs with heavily padded input?

Do we normally allow RNNs arbitrary additional time to 'finish' a computation before we read the output?

# High Level: Turing Completeness of RNNs

Do you see any practical problems with the stack fractions  $q_1$  and  $q_2$  ?

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$$z = \sigma(4q_i - 2)$$

read from top of  $S_i$ 

Do we normally provid

Do we normally allow RN computation

This result does not say much about how we use RNNs in practice (finite precision floating point calculations, bounded inference time)

### Practical Power of RNNs

Instead of checking their maximum potential, let's consider what different RNN mechanisms are really set up for...

GRU LSTM

Teaser: these are not the same...

#### GRU

$$z_{t} = \sigma(W^{z}x_{t} + U^{z}h_{t-1} + b^{z})$$

$$r_{t} = \sigma(W^{r}x_{t} + U^{r}h_{t-1} + b^{r})$$

$$\tilde{h}_{t} = \tanh(W^{h}x_{t} + U^{h}(r_{t} \circ h_{t-1}) + b^{h})$$

$$h_{t} = z_{t} \circ h_{t-1} + (1 - z_{t}) \circ \tilde{h}_{t}$$

$$f_t = \sigma(W^f x_t + U^f h_{t-1} + b^f)$$

$$i_t = \sigma(W^i x_t + U^i h_{t-1} + b^i)$$

$$o_t = \sigma(W^o x_t + U^o h_{t-1} + b^o)$$

$$\tilde{c}_t = \tanh(W^c x_t + U^c h_{t-1} + b^c)$$

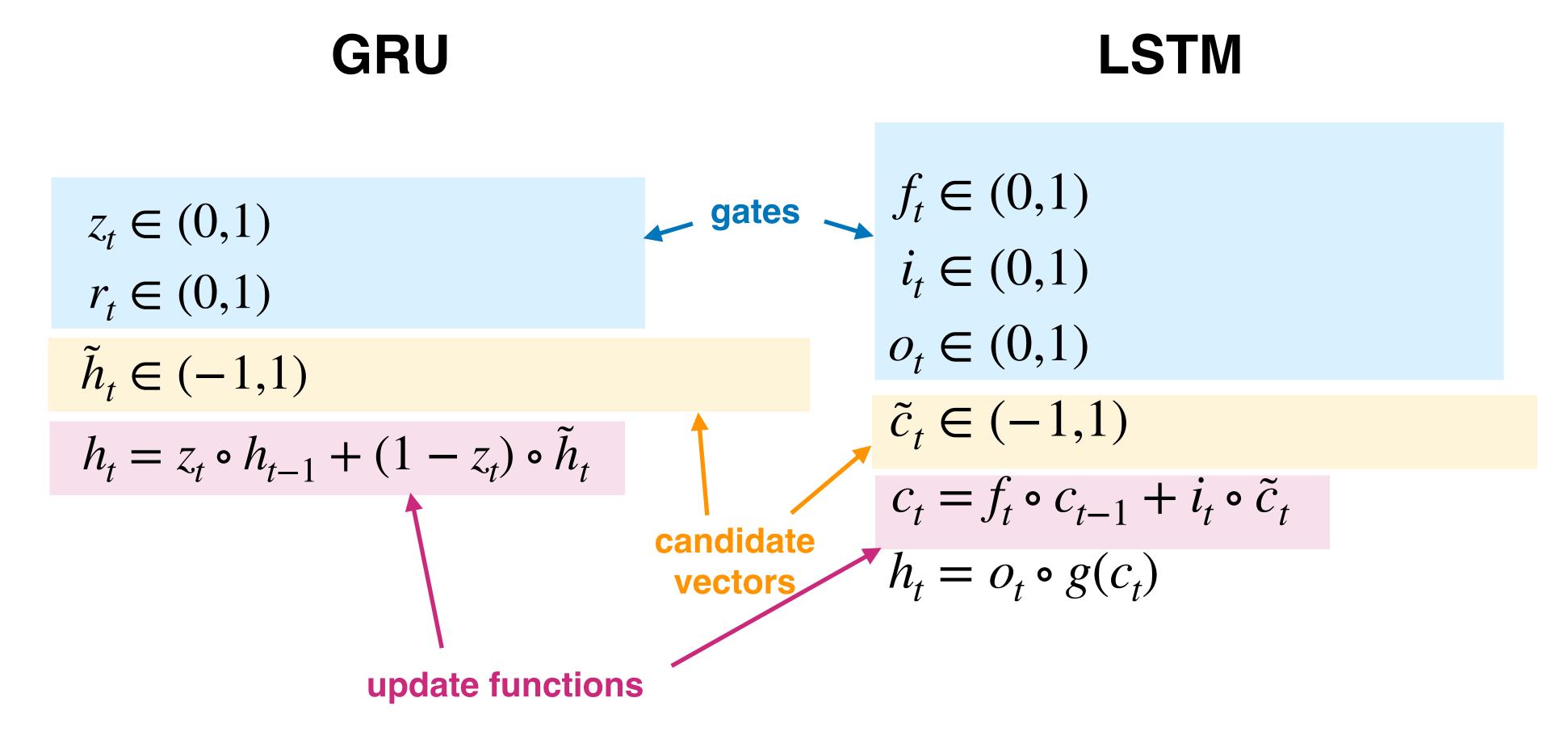
$$c_t = f_t \circ c_{t-1} + i_t \circ \tilde{c}_t$$

$$h_t = o_t \circ g(c_t)$$

#### **GRU**

$$\begin{aligned} z_t &= \sigma(W^z x_t + U^z h_{t-1} + b^z) \\ r_t &= \sigma(W^r x_t + U^r h_{t-1} + b^r) \end{aligned} \qquad \begin{aligned} f_t &= \sigma(W^f x_t + U^f h_{t-1} + b^f) \\ i_t &= \sigma(W^i x_t + U^i h_{t-1} + b^i) \\ \tilde{h}_t &= \tanh(W^h x_t + U^h (r_t \circ h_{t-1}) + b^h) \end{aligned} \qquad \begin{aligned} i_t &= \sigma(W^i x_t + U^i h_{t-1} + b^i) \\ o_t &= \sigma(W^o x_t + U^o h_{t-1} + b^o) \\ \tilde{c}_t &= \tanh(W^c x_t + U^c h_{t-1} + b^c) \\ \tilde{c}_t &= \tanh(W^c x_t + U^c h_{t-1} + b^c) \end{aligned} \\ c_t &= f_t \circ c_{t-1} + i_t \circ \tilde{c}_t \\ h_t &= o_t \circ g(c_t) \end{aligned}$$

#### GRU **LSTM** $\int_{t} \mathbf{gates} \int_{t} f_{t} \in (0,1)$ $i_{t} \in (0,1)$ $z_t \in (0,1)$ $r_t \in (0,1)$ $o_t \in (0,1)$ $\tilde{h}_t = \tanh(W^h x_t + U^h (r_t \circ h_{t-1}) + b^h)$ $\tilde{c}_t = \tanh(W^c x_t + U^c h_{t-1} + b^c)$ $h_t = z_t \circ h_{t-1} + (1 - z_t) \circ \tilde{h}_t$ $c_t = f_t \circ c_{t-1} + i_t \circ \tilde{c}_t$ $h_t = o_t \circ g(c_t)$ update functions



#### GRU

# $z_t \in (0,1)$ $r_t \in (0,1)$ Bounded! $\tilde{h}_t \in (-1,1)$ $h_t = z_t \circ h_{t-1} + (1-z_t) \circ \tilde{h}_t$ Interpolation

```
f_t \in (0,1)
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#### GRU

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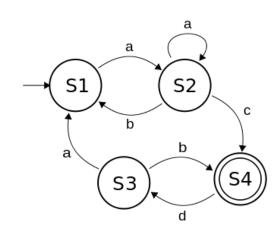
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$$f_t \in (0,1)$$
 reset/keep, then...

$$i_t \in (0,1)$$
 stay/step, by...

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 subtracting/adding 1

$$c_t = f_t \circ c_{t-1} + i_t \circ \tilde{c}_t$$

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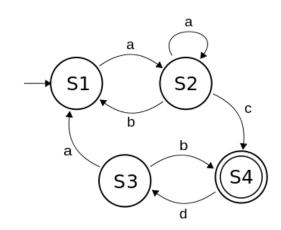
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#### **LSTM**

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Counts!

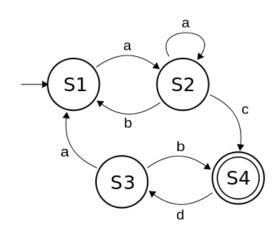
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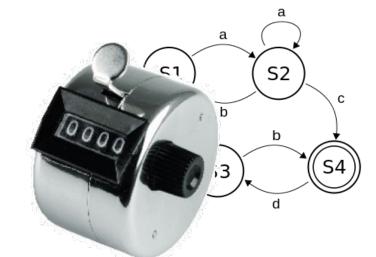
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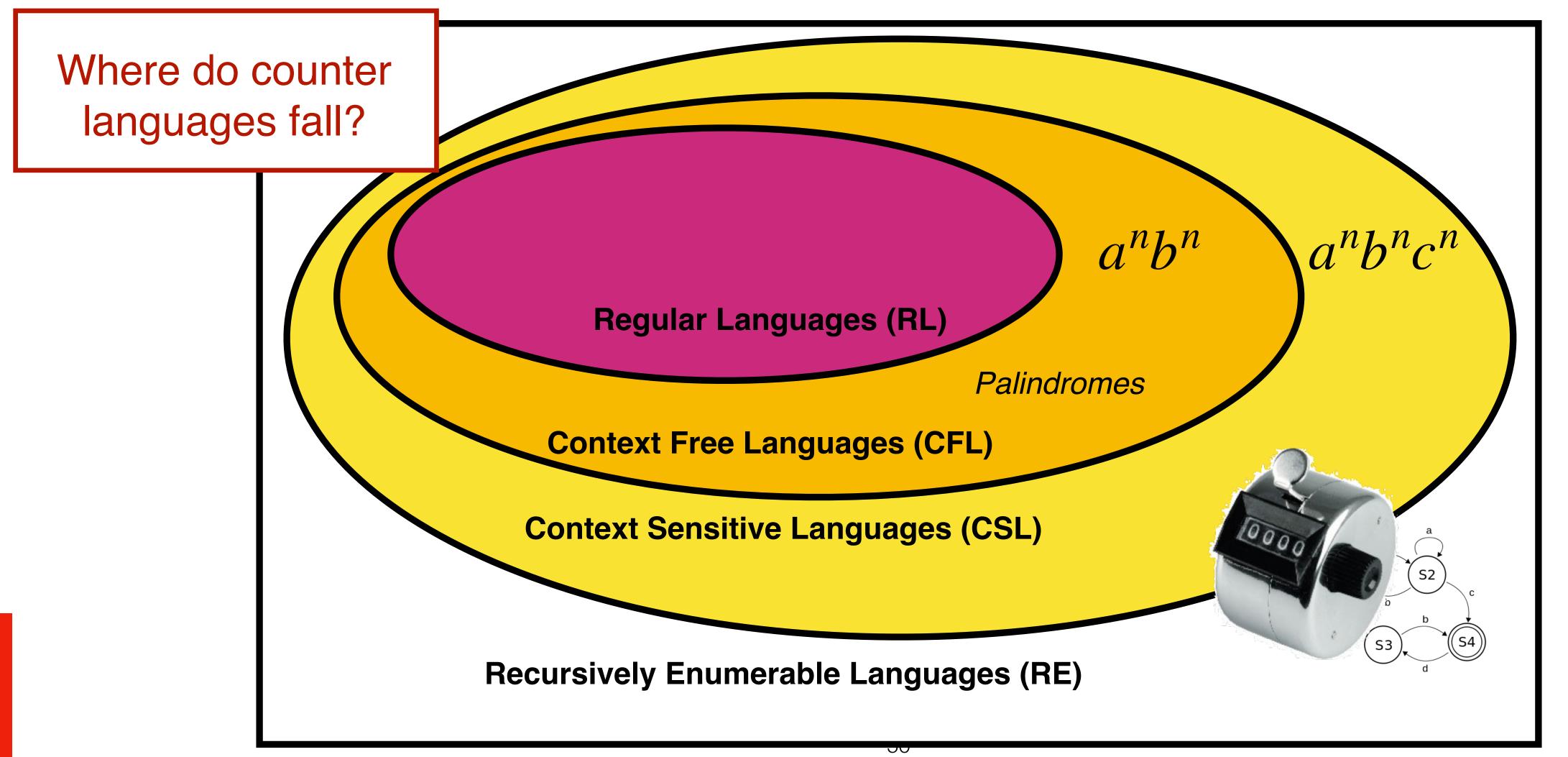
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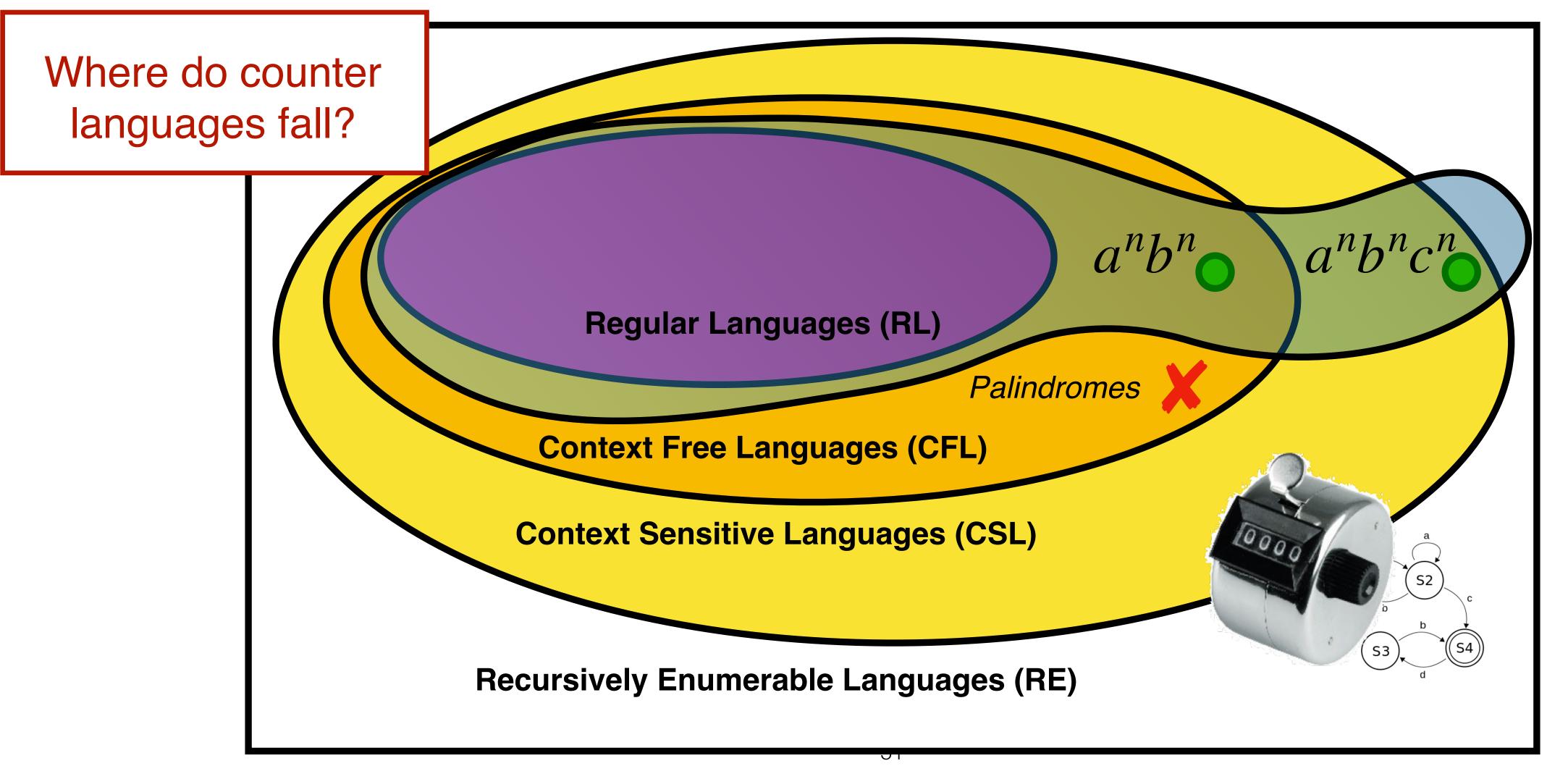


Counts!

# Aside: The Chomsky Hierarchy and Counting

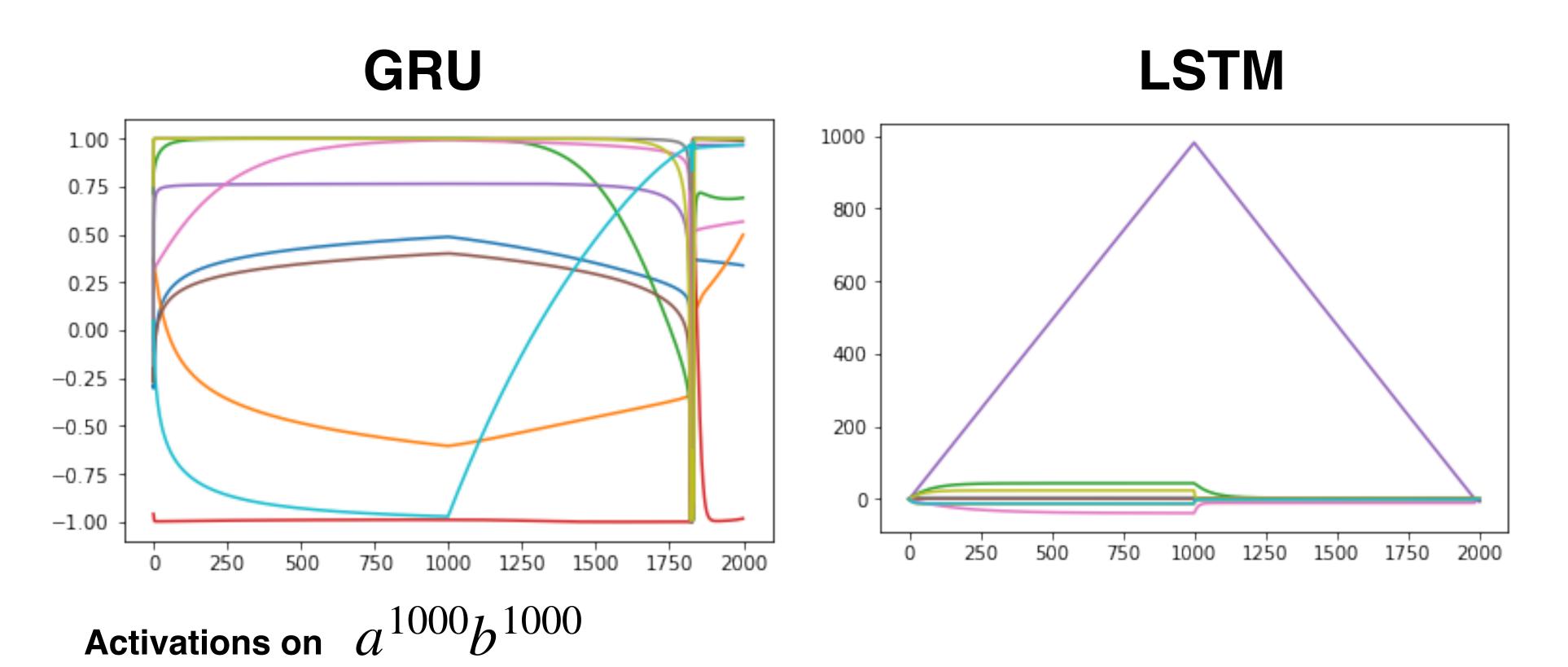


# Aside: The Chomsky Hierarchy and Counting



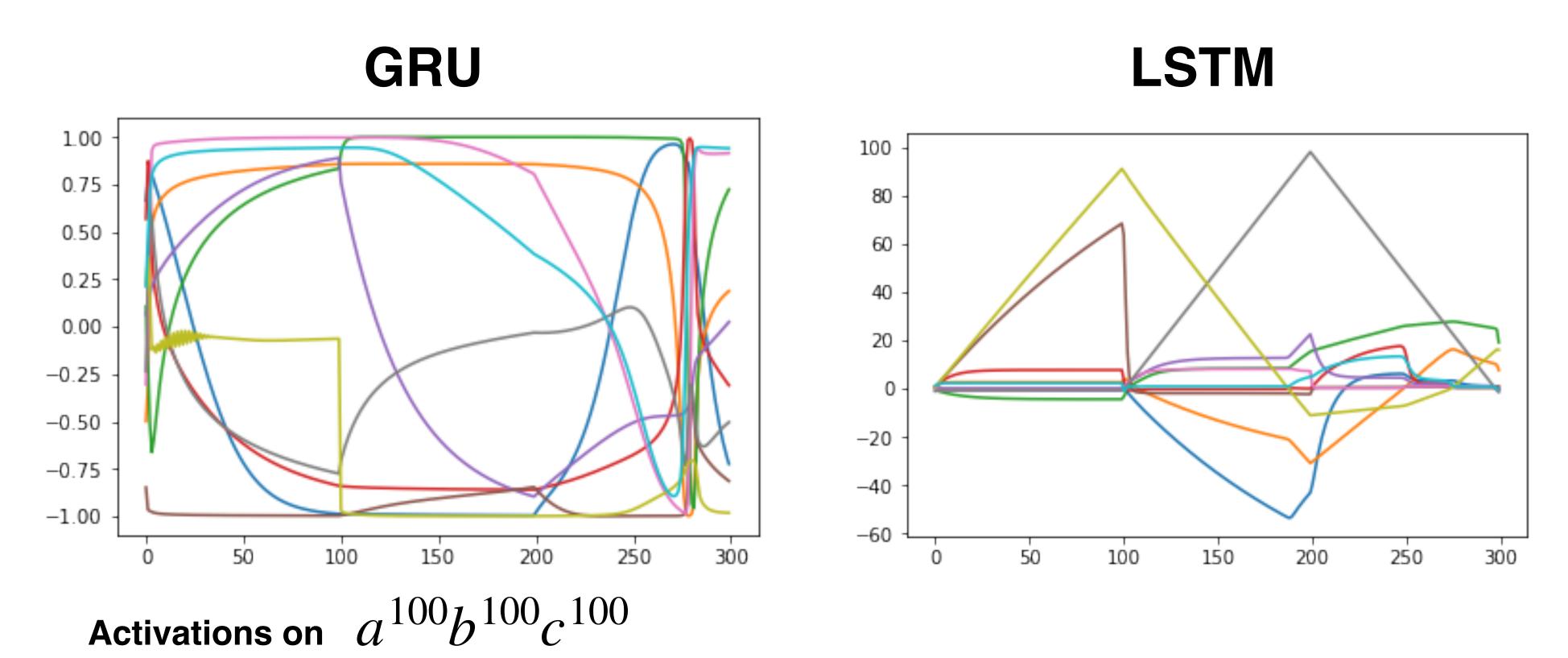
What will happen if we train GRUs and LSTMs on counting-based tasks, like  $a^nb^n$  and  $a^nb^nc^n$ ?

What kind of counters do these tasks need?



**Trained**  $a^nb^n$  , (on positive examples up to length 100)

GRU begins failing at length 39



**Trained**  $a^n b^n c^n$ , (on positive examples up to length 100)

GRU begins failing at length 9

#### Conclusion:

LSTMs can dedicate any single dimension to counting, and can learn to do so when necessary

GRUs cannot, and seem to struggle on counting tasks overall

Are LSTMs always better than GRUs?

# Other RNN Architectures

Let's have a little fun

#### **QRNNs (Quasi-Recurrent Neural Networks)**

A proposed efficient alternative to LSTMs

Idea: state update is function of last k input tokens - without current state

(This way on large sequence, heavy calculation of all state updates can be computed in parallel, and only the updates themselves need to be chained)

High level: 
$$x^{t,k} = (x_t, x_{t-1}, \dots, x_{t+1-k})$$
 
$$z_t, f_t, o_t, i_t = w(x^{t,k})$$
 
$$c_t = f_t \odot c_{t-1} + i_t \odot z_t$$
 
$$h_t = o_t \odot c_t$$

Could count - like an LSTM!

Can it be as strong as an LSTM?

#### **QRNNs (Quasi-Recurrent Neural Networks)**

High level: 
$$x^{t,k} = (x_t, x_{t-1}, \dots, x_{t+1-k})$$
 
$$z_t, f_t, o_t, i_t = w(x^{t,k})$$
 
$$c_t = f_t \odot c_{t-1} + i_t \odot z_t$$
 
$$h_t = o_t \odot c_t$$

Could count - like an LSTM!

But, counter updates won't be state-dependent

#### **QRNNs (Quasi-Recurrent Neural Networks)**

High level: 
$$z^{t,k} = (x_t, x_{t-1}, \dots, x_{t+1-k})$$
 
$$z_t, f_t, o_t, i_t = w(x^{t,k})$$
 Could count - like an LSTM! 
$$c_t = f_t \odot c_{t-1} + i_t \odot z_t$$
 But, counter updates won't be state-dependent

For languages like  $a^nb^n$ , can do all the needed counting without 'self awareness'

- RNN classifier function g can finish the job

#### **QRNNs (Quasi-Recurrent Neural Networks)**

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Does this mean a QRNN is as powerful as an LSTM?

#### **QRNNs (Quasi-Recurrent Neural Networks)**

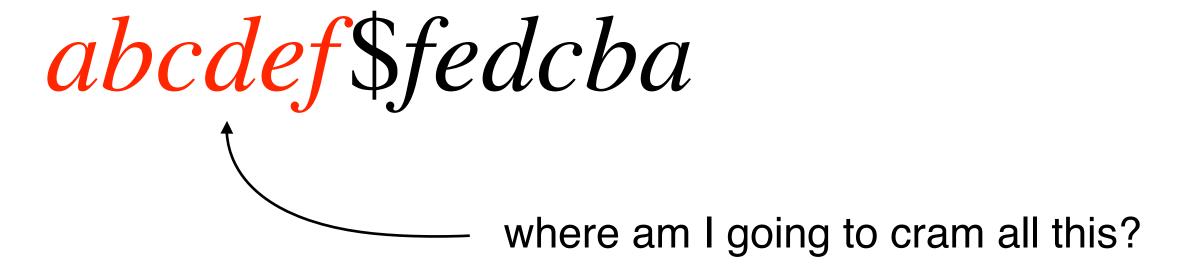
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 But, counter updates won't be state-dependent

#### Some counter languages require state/counter-dependent updates!

Consider 
$$L = \{a^n b^n w | n \ge 1, w \in \{a, b\}^*\}$$

Need to know to stop updating the counters after reading  $a^nb^n$ !

#### **Motivation:**

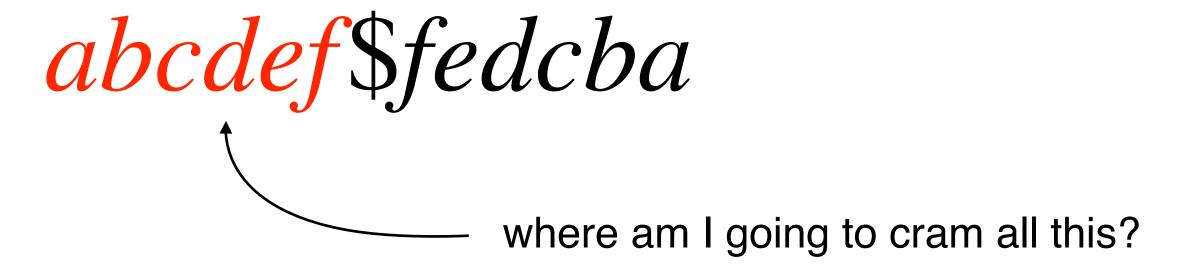


LSTMs can count! But that's not that strong

State values increase at most linearly with input sequence length: state "space" increases only polynomially

Need much more space to recognise things like palindromes (which requires holding entire prefix)

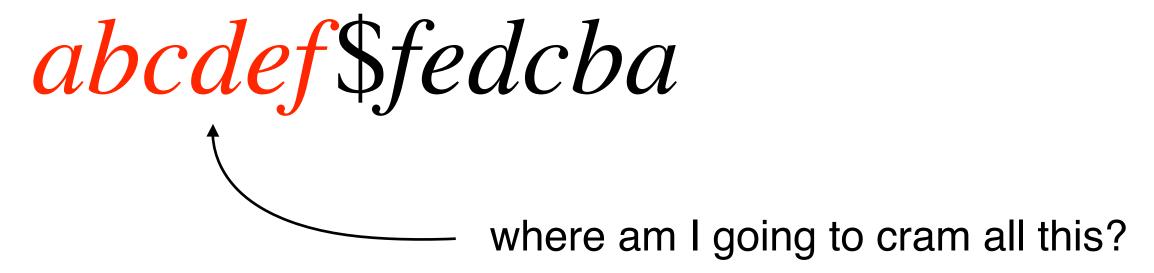
#### **Motivation:**



Need state space to grow much faster...

What if I allow state values to double??

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#### **Doubling-LSTM**

$$f_t^1 = \sigma(W^f x_t + U^f h_{t-1} + b^f)$$

$$f_t^2 = \sigma(W^f x_t + U^f h_{t-1} + b^f)$$

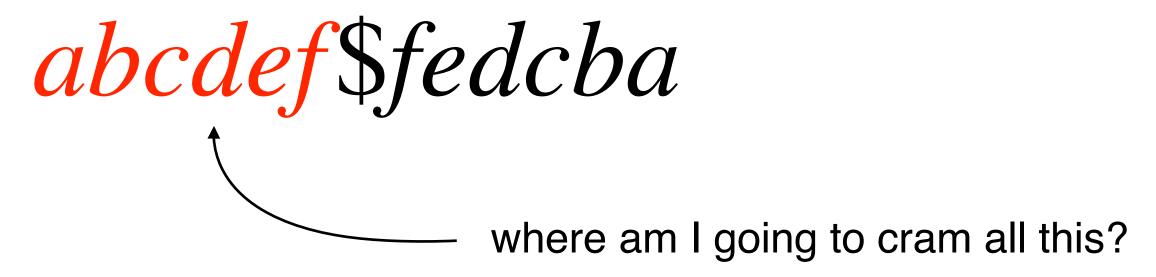
$$f_t = f_t^1 + f_t^2$$

$$\vdots$$

$$c_t = f_t \circ c_{t-1} + i_t \circ \tilde{c}_t$$

$$h_t = o_t \circ g(c_t)$$

#### **Motivation:**



Now the RNN can implement reset ( $f_t = 0$ ), keep ( $f_t = 1$ ) and double ( $f_t = 2$ )

If we also assume it manages to stabilise on halve ( $f_t = 1/2$ ), it effectively has a simple stack to push and pop binary tokens to/from!

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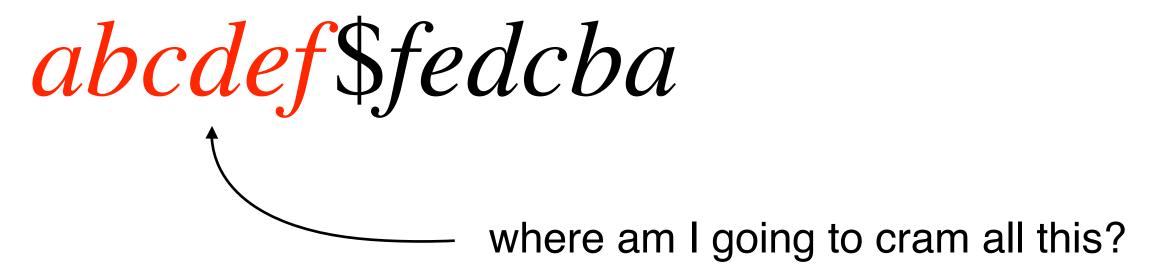
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Did it work?

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#### Did it work?

It turns out doubling the values in RNN states quickly gets too big for training...

# RNNs with Augmentations

The doubling LSTM was a failure... but getting a stack doesn't have to be doomed

Several works propose augmenting an RNN with an external stack

(Other works propose adding attention, but you've already seen that :))

The stack is not encoded as some horrifically large number

Instead an actual stack, with differentiable push, pop, and read operations, is carefully designed

(That's the hard part)

The RNN learns to interact with the stack to solve tasks

## Stack RNNs

#### 1. Stratification (layer cake stack)

Items in stack have a "thickness" - layers of a cake

RNN computes a distribution d over push, pop, do nothing

Push, pop, and nothing all(!) happen, each with thickness according to d

When read width covers several stack layers: they are averaged (by thickness)





Reality



# Stack RNNs

#### 2. Superposition (Try everything stack)

Items in stack get pushed, popped, or left alone completely - like a true stack

RNN still computes a distribution d over push, pop, do nothing

Push, pop, and nothing all(!) happen, each in different stack - all weighted by d

Creates polynomial (!) number of potential stacks over time, d is composed with the existing distribution

Reading from stack takes weighted average of stack tops



# Stack RNNs

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#### (3. Nondeterminism)

New: similar to superposition, but achieved with simulation of a nondeterministic pushdown automaton



# References

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