## Smooth Sorting

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softmax: 
$$\mathbb{R}^n \to [0,1]^n$$
, softmax $(v)_i = \frac{\exp(v_i)}{\sum_{j=1}^n \exp(v_j)}$ 

softmax:  $\mathbb{R}^n \to [0,1]^n$ , softmax $(v)_i = \frac{\exp(v_i)}{\sum_{j=1}^n \exp(v_j)}$ Gives smoothed version of argmax, which returns vector with 1 in position of largest element and 0 elsewhere Entries have sum 1: may be interpreted as probability that given element of v is the largest Want to generalize to  $\operatorname{soft}_k(v)$  for  $k \in \{1, \ldots, n\} = [n]$ : probability that  $v_i$  is the  $k^{\text{th}}$ -largest Desired properties:

- 1.  $soft_1 = softmax$
- 2.  $\operatorname{soft}_{n+1-k}(v) = \operatorname{soft}_k(-v)$
- 3. The largest entry of soft<sub>k</sub>(v) is in the same place as the  $k^{\text{th}}$ -largest entry of v
- 4.  $\sum_{i=1}^{n} \operatorname{soft}_{k}(v)_{j} = 1$  (exactly one element of v is at position k)
- 5.  $\sum_{k=1}^{n} \operatorname{soft}_{k}(v)_{j} = 1$  ( $v_{j}$  is in exactly one position)
- 6.  $\operatorname{soft}_k$  is smooth
- 7.  $\operatorname{soft}_k(v) = \operatorname{soft}_k(u)$  if there is  $c \in \mathbb{R}$  such that for all  $i, u_i = v_i + c$ , i.e.  $soft_k$  does not change when adding a constant to all elements of the input

Define  $S \in \mathbb{R}^{n \times n}$  by  $S_{jk} = \operatorname{soft}_k(v)_j$ 

Properties 4 and 5: S is doubly stochastic

Birkhoff's theorem:  $S = \sum_{\pi} P(\pi) M_{\pi}$  over permutations  $\pi$  of [n],  $M_{\pi}$  has 1 in position  $(\pi(i), i)$  and 0 elsewhere,  $P(\pi) \in [0, 1]$ ,  $\sum_{\pi} P(\pi) = 1$  Interpretation:  $P(\pi) = \text{probability of } v_{\pi(k)}$  being  $k^{\text{th}}$ -largest for all k

Justification:  $S_{jk} = \sum_{\pi(k)=j} P(\pi)$ :  $v_j$  is  $k^{\text{th}}$ -largest iff a permutation where  $v_j$  is  $k^{\text{th}}$ -largest is chosen Reframe problem: define a distribution over permutations, get 4 and 5 for free

Equivalently, define procedure to pick a random permutation: "smooth" sorting

First attempt: property 1 may be rewritten  $P(\pi(1) = j) = \operatorname{softmax}(v)_i$ Can ensure by picking  $\pi(1) = a$  first, then permuting the remaining n-1 elements recursively:

- Define  $\sigma: [n] \setminus \{a\} \leftrightarrow [n-1]$  by  $\sigma(i) = i$  if i < a and i-1 if i > a: map between indices
- Let  $v' \in \mathbb{R}^{n-1}$  where  $v'_i = v_{\sigma(i)}$  (i.e. v' is v with  $v_a$  removed)
- Find permutation  $\pi'$  of [n-1] based on v'
- Define  $\pi(i) = a$  if i = 1 and  $\sigma^{-1}(\pi'(i-1))$  otherwise

Base case n=0: identity permutation

Since softmax satisfies independence of irrelevant alternatives (IIA), this is equivalent to:

- Repeatedly pick  $i \in [n]$  with probability softmax $(v)_i$
- If already seen, discard; else, add to list

• Once list has length  $n, \pi(k) =$  element at  $k^{\text{th}}$  position

However, this procedure does not satisfy property 2:

- Let n = 3,  $v_1 = 2$ ,  $v_2 = 1$ ,  $v_3 = 0$
- Then  $soft_3(v)_3 = soft_1(-v)_3 = \frac{\exp(0)}{\exp(-2) + \exp(-1) + \exp(0)} \approx 0.67$
- $\bullet$  Procedure picks  $v_3$  last iff either  $(v_1$  first, then  $v_2)$  or  $(v_2$  first, then  $v_1)$
- Probability =  $\frac{\exp(2)}{\exp(2) + \exp(1) + \exp(0)} \frac{\exp(1)}{\exp(1) + \exp(0)} + \frac{\exp(1)}{\exp(2) + \exp(1) + \exp(0)} \frac{\exp(2)}{\exp(2) + \exp(0)} \approx 0.70$

New procedure: should act on v "symmetrically" to satisfy property 2

Idea: pick  $\pi(1) = a$  and  $\pi(n) = b$  at once, then find middle recursively using v without  $v_a$  or  $v_b$ 

(Base cases are now n = 0 or 1: still identity)

First attempt: pick a by softmax(v) and b by softmax(-v) independently Equivalently:

- Define matrix E by  $E_{ab} = \exp(v_a v_b)$
- Pick (a,b) with probability proportional to  $E_{ab}$ , i.e. equal to  $E_{ab}/(\text{sum of entries of }E)$

E stores non-normalized joint distribution of (a,b)

Non-normalized marginal distribution of a given by  $\sum_b E_{ab} = \exp(v_a) \sum_b \exp(-v_b)$ 

Proportional to  $\exp(v_a)$ , which is proportional to softmax $(v)_a$ , so property 1 satisfied

Problem: a could equal b since  $E_{aa} = \exp(v_a - v_a) = 1 \neq 0$ : can't have  $\pi(1) = \pi(n)$ 

Need to fix E by making diagonal 0 without changing marginal distributions (ratios of row/column sums) First attempt:

- Just zero out diagonal
- Equivalently, try to pick a and b again with the same probabilities until  $a \neq b$
- Sum of each row decreases by 1, but since rows can have different sums, their ratio can change
- Thus marginal distribution of a not preserved (same for b)
- Scaling E to restore sum of entries doesn't help since ratios of row/column sums don't change

## Second attempt:

- Zero out diagonal, add  $\frac{1}{n-1}$  to all other entries
- Equivalently, pick a and b as usual, but if a = b, pick a distinct pair uniformly instead
- Row/column sums decrease by 1 in 1 spot and increase by  $\frac{1}{n-1}$  in n-1 spots: no change
- $\bullet$  Thus marginal distributions of a and b preserved: success!

Recall this is a single iteration: overall process finds  $\pi(1)$  and  $\pi(n)$ , then  $\pi(2)$  and  $\pi(n-1)$ , ... Use probabilities of picking each  $\pi$  to define soft<sub>k</sub> as described by Birkhoff's thm.

Satisfies properties 1, 2, 4, 5, 6, and 7 by virtue of its construction

However, procedure suggests that closed form of soft<sub>k</sub> must case on whether k < n/2

(Also expensive by naive method, but not currently a concern)