

# Notes on Type Theory

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Base types, sum/product/function types, dependent types, inductive types, universes

Base types:  $\perp$  AKA false,  $\top$  AKA true, bool

$\#:\top$ ,  $0:\text{bool}$ ,  $1:\text{bool}$  ( $\perp$  has no terms),  $\perp:\text{Type}$ ,  $\top:\text{Type}$ ,  $\text{bool}:\text{Type}$

Dependent types:  $\sum$  AKA  $\exists$ ,  $\prod$  AKA  $\forall$   
 $\exists$

$$A + B \qquad A \vee B \qquad (1)$$

$$A \times B \qquad A \wedge B \qquad (2)$$

$$B^A \qquad A \rightarrow B \qquad (3)$$

(1): The union type of types A and B has terms which are either of type A or of type B. For example, if A had 2 terms (and thus stored one bit of information) and B had 3 terms (one trit), there would be  $2 + 3 = 5$  possible values of the union.

(2): The pair type of A and B has terms which contain both a term of type A and a term of type B. Using the previous examples of A and B, there would be  $2 \times 3 = 6$  possible pairs.

(3): The function type that goes from A to B has terms which associate exactly one (non-unique) element of B with each element of A. Since for each of the 2 terms of A (input), there are 3 choices for the corresponding element of B (output), there would be  $3^2 = 9$  possible functions.

$$A \times (B + C) = (A \times B) + (A \times C) \qquad A \wedge (B \vee C) \equiv (A \wedge B) \vee (A \wedge C) \qquad (4)$$

$$C^{A+B} = C^A \times C^B \qquad (A \vee B) \rightarrow C \equiv (A \rightarrow C) \wedge (B \rightarrow C) \qquad (5)$$

$$(B \times C)^A = B^A \times C^A \qquad A \rightarrow (B \wedge C) \equiv (A \rightarrow B) \wedge (A \rightarrow C) \qquad (6)$$

$$C^{A \times B} = (C^B)^A \qquad (A \wedge B) \rightarrow C \equiv A \rightarrow (B \rightarrow C) \qquad (7)$$

$$A + 0 = A \qquad A \vee F \equiv A \qquad (8)$$

$$A \times 1 = A \qquad A \wedge T \equiv A \qquad (9)$$