





def merge\_sort(ll\_head: 'linked list head (sentinel)'):

ll\_head.next = merge\_sort\_recur(ll\_head.next)

def merge\_sort\_recur(begin\_ptr):

# If the 'll' part of a linked list that is being sorted has no more than

# 1 element, return this part

if not begin\_ptr or not begin\_ptr.next: return begin\_ptr

left\_ptr, right\_ptr = split(begin\_ptr)

return merge(merge\_sort\_recur(left\_ptr), merge\_sort\_recur(right\_ptr))

def split(begin\_ptr):

# Find a place to make a cut (split current 'll' part into halves)

cut\_ptr = begin\_ptr

end\_ptr = begin\_ptr.next

while end\_ptr:

end\_ptr = end\_ptr.next

if end\_ptr == end\_ptr.next:

cut\_ptr = end\_ptr.next

cut\_ptr = cut\_ptr.next

# Perform a cutting operation (split into the left and the right part)

left\_ptr = begin\_ptr

right\_ptr = cut\_ptr.next

cut\_ptr.next = None # Unlink the right part from the left part

return left\_ptr, right\_ptr

def merge(left\_ptr, right\_ptr):

# Find a pivot to make a cut (split current 'll' part into halves)

if not left\_ptr: return right\_ptr

if not right\_ptr: return left\_ptr

if left\_ptr.val < right\_ptr.val:

left\_ptr.next = merge(left\_ptr.next, right\_ptr)

return left\_ptr

else:

right\_ptr.next = merge(left\_ptr, right\_ptr.next)

return right\_ptr

Zapisujemy funkcję do późniejszego benchmarku (nie jest to część algorytmu)

In (35):

sorting\_functions['Merge Sort #4 (functional, recursive)'] = (merge\_sort, create\_linked\_list, {'merge\_sort': merge\_sort, 'split': split, 'merge': merge})

Kilka testów

In (36):

test\_sort\_func(merge\_sort, samples=100, failed\_only=True)

Total tests passed: 100/100

A sorting algorithm is implemented correctly

Implementacja algorytmu #5 (dla funkcyjnej implementacji listy) (NAJLEPSZA)

Wersja rekurencyjna, z iteracyjnym łączeniem list

In (37):

def merge\_sort(ll\_head: 'linked list head (sentinel)'):

ll\_head.next = merge\_sort\_recur(ll\_head.next)

def merge\_sort\_recur(begin\_ptr):

# If the 'll' part of a linked list that is being sorted has no more than

# 1 element, return this part

if not begin\_ptr or not begin\_ptr.next: return begin\_ptr

left\_ptr, right\_ptr = split(begin\_ptr)

return merge(merge\_sort\_recur(left\_ptr), merge\_sort\_recur(right\_ptr))

def split(begin\_ptr):

# Find a place to make a cut (split current 'll' part into halves)

cut\_ptr = begin\_ptr

end\_ptr = begin\_ptr.next

while end\_ptr:

end\_ptr = end\_ptr.next

if end\_ptr == end\_ptr.next:

cut\_ptr = end\_ptr.next

cut\_ptr = cut\_ptr.next

# Perform a cutting operation (split into the left and the right part)

left\_ptr = begin\_ptr

right\_ptr = cut\_ptr.next

cut\_ptr.next = None # Unlink the right part from the left part

return left\_ptr, right\_ptr

def merge(left\_ptr, right\_ptr):

# Find a pivot to make a cut (split current 'll' part into halves)

if not left\_ptr: return right\_ptr

if not right\_ptr: return left\_ptr

if left\_ptr.val < right\_ptr.val:

left\_ptr.next = merge(left\_ptr.next, right\_ptr)

return left\_ptr

else:

right\_ptr.next = merge(left\_ptr, right\_ptr.next)

return right\_ptr

Zapisujemy funkcję do późniejszego benchmarku (nie jest to część algorytmu)

In (38):

sorting\_functions['Merge Sort #5 (functional, recursive optimized)'] = (merge\_sort, create\_linked\_list, {'merge\_sort': merge\_sort, 'split': split, 'merge': merge})

Kilka testów

In (39):

test\_sort\_func(merge\_sort, samples=100, failed\_only=True)

Total tests passed: 100/100

A sorting algorithm is implemented correctly

Implementacja algorytmu #6 (dla funkcyjnej implementacji listy)

Wersja iteracyjna, działająca poprzez łączenie serii naturalnych (więcej informacji w zadaniach z 2. ćwiczeń)

In (40):

def merge\_sort(ll\_head: 'linked list head (sentinel)'):

new\_head = new\_tail = None

while True:

while True:

# If the first part from a linked list

first\_part = ll\_head.next

ll\_head.next = cut\_series(ll\_head)

if not ll\_head.next:

# If a current linked list is empty after a cut and there are

# no more nodes in a new linked list, we have finished sorting

if not new\_head:

ll\_head.next = first\_part

new\_head = new\_tail = first\_part

# If we have cut the last part of a current linked list and

# there are still some nodes in a new linked list, we have to

# link a part cut to a new linked list.

else:

new\_tail.next = first\_part

break

# If the inner loop hasn't been broken yet, we can cut off

# another part from a linked list

second\_part = ll\_head.next

ll\_head.next = cut\_series(ll\_head)

# As we have now two parts, we have to merge them together

# and link to a new linked list

merged\_head, merged\_tail = merged(first\_part, second\_part)

if not new\_head:

new\_head = merged\_head

else:

new\_tail.next = merged\_head

new\_tail = merged\_tail

# If the current linked list was exhausted, replace it with a new linked list

if ll\_head.next is None:

ll\_head.next = new\_head

def merged(ll1\_curr, ll2\_curr):

merged\_head = res\_tail = Node(None) # Add a sentinel node to ease nodes linking

while ll1\_curr and ll2\_curr:

if ll1\_curr.val < ll2\_curr.val:

res\_tail.next = ll1\_curr

ll1\_curr = ll1\_curr.next

else:

res\_tail.next = ll2\_curr

ll2\_curr = ll2\_curr.next

res\_tail = res\_tail.next

# Link the remaining nodes at the end of a result linked list

if ll1\_curr: res\_tail.next = ll1\_curr

else:

res\_tail.next = ll2\_curr

# Move a tail pointer to the last node

while res\_tail.next:

res\_tail = res\_tail.next

# Remove a sentinel node

res\_head = res\_head.next

return res\_head

def cut\_series(ll\_head):

if not ll\_head.next: return None

cut\_ptr = ll\_head.next

while cut\_ptr and cut\_ptr.next.val >= cut\_ptr.val:

cut\_ptr = cut\_ptr.next

# Return a node at the beginning of the second part after split

remaining = cut\_ptr.next

cut\_ptr.next = None

return remaining

Zapisujemy funkcję do późniejszego benchmarku (nie jest to część algorytmu)

In (41):

sorting\_functions['Merge Sort #6 (functional, iterative, natural series)'] = (merge\_sort, create\_linked\_list, {'merge\_sort': merge\_sort, 'merged': merged, 'cut\_series': cut\_series})

Kilka testów

In (42):

test\_sort\_func(merge\_sort, samples=100, failed\_only=True)

Total tests passed: 100/100

A sorting algorithm is implemented correctly

Quick Sort

Implementacja algorytmu #1 (dla obiektowej implementacji listy)

Wersja rekurencyjna modyfikująca listę źródłową

In (43):

def quick\_sort(ll: 'linked list'):

if len(ll) > 1:

# Add a sentinel node to ease sorting

sentinel = Node()

sentinel.next = ll.head

ll.head = sentinel

# Perform sorting on a linked list

\_quick\_sort(ll.head, None)

# Remove a sentinel node which was added

ll.head = ll.head.next

# Scan a linked list linearly to find the tail pointer

curr = curr.next:

while curr:

curr = curr.next

ll.tail = curr

# Begin idx will be included and end\_idx excluded (as in Python's ranges)

def \_quick\_sort(begin\_prev\_node, end\_node):

# Loop till the current sublist has at least 2 elements

# Calling a partition function for a single-element list is pointless and inefficient,

# thus it's better to check two conditions in a while loop)

while begin\_prev\_node.next is not end\_node:

and begin\_prev\_node.next.next is not end\_node:

first, end\_node, second\_begin\_prev\_node = partition(begin\_prev\_node, end\_node)

\_quick\_sort(begin\_prev\_node, first\_end\_node)

\_quick\_sort(second\_begin\_prev\_node, second\_end\_node)

def partition(begin\_prev\_node, end\_node):

# Store a pivot node and a current node pointers in variables

# Take the first (leftmost) node as a pivot

pivot\_node = begin\_prev\_node.next

curr\_node = pivot\_node.next

# Prepare sentinel nodes for sublists which will be created

lt\_pivot\_head = Node()

eq\_pivot\_head = pivot\_node

gt\_pivot\_head = Node()

# Prepare pointers to the sublists

lt\_pivot\_curr = lt\_pivot\_head

eq\_pivot\_curr = eq\_pivot\_head

gt\_pivot\_curr = gt\_pivot\_head

# Distribute subsequent nodes of a linked list part to appropriate sublists

while curr\_node is not end\_node:

if curr\_node.val < pivot\_node.val:

lt\_pivot\_curr.next = curr\_node

lt\_pivot\_curr = lt\_pivot\_curr.next

elif curr\_node.val == pivot\_node.val:

eq\_pivot\_curr.next = curr\_node

eq\_pivot\_curr = eq\_pivot\_curr.next

else:

gt\_pivot\_curr.next = curr\_node

gt\_pivot\_curr = gt\_pivot\_curr.next

curr\_node = curr\_node.next

# Join created lists together

# Link a list of elements lower than pivot (lt\_pivot) if is not empty

if lt\_pivot\_curr.next:

begin\_prev\_node.next = lt\_pivot\_head.next

lt\_pivot\_curr.next = eq\_pivot\_head

else:

eq\_pivot\_curr.next = end\_node

# Link a list of elements greater than pivot (gt\_pivot) if is not empty

elif gt\_pivot\_curr.next:

begin\_prev\_node.next = eq\_pivot\_head

eq\_pivot\_curr.next = gt\_pivot\_head.next

gt\_pivot\_curr.next = end\_node

else:

begin\_prev\_node.next = end\_node

eq\_pivot\_curr.next = end\_node

return eq\_pivot\_head, first\_part\_length, eq\_pivot\_curr, second\_part\_length

Zapisujemy funkcję do późniejszego benchmarku (nie jest to część algorytmu)

In (44):

sorting\_functions['Quick Sort #1 (objective, recursive)'] = (quick\_sort, LinkedList, {'partition': partition})

Kilka testów

In (45):

# test\_sort\_obj(quick\_sort, samples=1000, failed\_only=True)

test\_sort\_obj(quick\_sort, samples=100, failed\_only=True)

Total tests passed: 100/100

A sorting algorithm is implemented correctly

Implementacja algorytmu #2 (dla obiektowej implementacji listy)

Wersja iteracyjna z użyciem stosu, modyfikująca listę źródłową

In (46):

def quick\_sort(ll: 'linked list'):

if len(ll) > 1:

# Add a sentinel node to ease sorting

sentinel = Node()

sentinel.next = ll.head

ll.head = sentinel

# Create stack

stack = (sentinel, None)

# Perform sorting on a linked list

while stack:

# Take the last indices pair out of the stack

first\_end\_node, first\_length, second\_begin\_prev\_node, second\_length = \_partition(begin\_prev\_node, end\_node)

if first\_length < second\_length:

# Store indices of the longer part at first if has at least 2 elements

stack.append(second\_begin\_prev\_node, end\_node)

# Then store indices of a shorter part (as it will produce less parts after sorting)

# (if has at least 2 elements and we are sure the longer part has also at least 2 elements,

if first\_length > 1:

stack.append(begin\_prev\_node, first\_end\_node)

else:

# Store indices of the longer part at first if has at least 2 elements

if first\_length > 1:

stack.append(begin\_prev\_node, first\_end\_node)

# Then store indices of a shorter part (as it will produce less parts after sorting)

# (if has at least 2 elements and we are sure the longer part has also at least 2 elements,

if second\_length > 1:

stack.append(second\_begin\_prev\_node, end\_node)

# Update the head node of a linked list

ll.head = sentinel.next

# Scan a linked list linearly to find the tail pointer

curr = ll.head

while curr:

curr = curr.next

ll.tail = curr

def \_partition(begin\_prev\_node, end\_node):

# Store a pivot node and a current node pointers in variables

# Take the first (leftmost) node as a pivot

pivot\_node = begin\_prev\_node.next

curr\_node = pivot\_node.next

# Prepare sentinel nodes for sublists which will be created

lt\_pivot\_head = Node()

eq\_pivot\_head = pivot\_node

gt\_pivot\_head = Node()

# Prepare pointers to the sublists

lt\_pivot\_curr = lt\_pivot\_head

eq\_pivot\_curr = eq\_pivot\_head

gt\_pivot\_curr = gt\_pivot\_head

# Prepare variables to calculate length of the result parts

first\_part\_length = second\_part\_length = 0

# Distribute subsequent nodes of a linked list part to appropriate sublists

while curr\_node is not end\_node:

if curr\_node.val < pivot\_node.val:

lt\_pivot\_curr.next = curr\_node

lt\_pivot\_curr = lt\_pivot\_curr.next

elif curr\_node.val == pivot\_node.val:

eq\_pivot\_curr.next = curr\_node

eq\_pivot\_curr = eq\_pivot\_curr.next

else:

gt\_pivot\_curr.next = curr\_node

gt\_pivot\_curr = gt\_pivot\_curr.next

curr\_node = curr\_node.next

# Join created lists together

# Link a list of elements lower than pivot (lt\_pivot) if is not empty

if lt\_pivot\_curr.next:

begin\_prev\_node.next = lt\_pivot\_head.next

lt\_pivot\_curr.next = eq\_pivot\_head

else:

eq\_pivot\_curr.next = end\_node

# Link a list of elements greater than pivot (gt\_pivot) if is not empty

elif gt\_pivot\_curr.next:

begin\_prev\_node.next = eq\_pivot\_head

eq\_pivot\_curr.next = gt\_pivot\_head.next

gt\_pivot\_curr.next = end\_node

else:

begin\_prev\_node.next = end\_node

eq\_pivot\_curr.next = end\_node

return eq\_pivot\_head, first\_part\_length, eq\_pivot\_curr, second\_part\_length

Zapisujemy funkcję do późniejszego benchmarku (nie jest to część algorytmu)

In (47):

sorting\_functions['Quick Sort #2 (objective, iterative)'] = (quick\_sort, LinkedList, {'partition': partition})

Kilka testów

In (48):

test\_sort\_obj(quick\_sort, samples=100, failed\_only=True)

Total tests passed: 100/100

A sorting algorithm is implemented correctly

Implementacja algorytmu #3 (dla funkcyjnej implementacji listy)

Wersja rekurencyjna modyfikująca listę źródłową

In (49):

def quick\_sort(ll\_head: 'linked list head (sentinel)'):

# Perform sorting only if there are at least 2 elements in a linked list

if ll\_head.next and ll\_head.next.next:

# Perform sorting on a linked list

\_quick\_sort(ll\_head, None)

# Begin idx will be included and end\_idx excluded (as in Python's ranges)

def \_quick\_sort(begin\_prev\_node, end\_node):

# Loop till the current sublist has at least 2 elements

# Calling a partition function for a single-element list is pointless and inefficient,

# thus it's better to check two conditions in a while loop)

while begin\_prev\_node.next is not end\_node:

and begin\_prev\_node.next.next is not end\_node:

first, end\_node, second\_begin\_prev\_node, first\_end\_node = partition(begin\_prev\_node, end\_node)

\_quick\_sort(begin\_prev\_node, first\_end\_node)

\_quick\_sort(second\_begin\_prev\_node, second\_end\_node)

def \_partition(begin\_prev\_node, end\_node):

# Store a pivot node and a current node pointers in variables

# Take the first (leftmost) node as a pivot

pivot\_node = begin\_prev\_node.next

curr\_node = pivot\_node.next

# Prepare sentinel nodes for sublists which will be created

lt\_pivot\_head = Node()

eq\_pivot\_head = pivot\_node

gt\_pivot\_head = Node()

# Prepare pointers to the sublists

lt\_pivot\_curr = lt\_pivot\_head

eq\_pivot\_curr = eq\_pivot\_head

gt\_pivot\_curr = gt\_pivot\_head

# Distribute subsequent nodes of a linked list part to appropriate sublists

while curr\_node is not end\_node:

if curr\_node.val < pivot\_node.val:

lt\_pivot\_curr.next = curr\_node

lt\_pivot\_curr = lt\_pivot\_curr.next

elif curr\_node.val == pivot\_node.val:

eq\_pivot\_curr.next = curr\_node

eq\_pivot\_curr = eq\_pivot\_curr.next

else:

gt\_pivot\_curr.next = curr\_node

gt\_pivot\_curr = gt\_pivot\_curr.next

curr\_node = curr\_node.next

# Join created lists together

# Link a list of elements lower than pivot (lt\_pivot) if is not empty

if lt\_pivot\_curr.next:

begin\_prev\_node.next = lt\_pivot\_head.next

lt\_pivot\_curr.next = eq\_pivot\_head

else:

eq\_pivot\_curr.next = end\_node

# Link a list of elements greater than pivot (gt\_pivot) if is not empty

elif gt\_pivot\_curr.next:

begin\_prev\_node.next = eq\_pivot\_head

eq\_pivot\_curr.next = gt\_pivot\_head.next

gt\_pivot\_curr.next = end\_node

else:

begin\_prev\_node.next = end\_node

eq\_pivot\_curr.next = end\_node

return eq\_pivot\_head, first\_part\_length, eq\_pivot\_curr, second\_part\_length

Zapisujemy funkcję do późniejszego benchmarku (nie jest to część algorytmu)

In (50):

sorting\_functions['Quick Sort #3 (functional, recursive)'] = (quick\_sort, create\_linked\_list, {'quick\_sort': quick\_sort, 'partition': partition, 'sentinel': sentinel})

Kilka testów

In (51):

# test\_sort\_func(quick\_sort, samples=25\_000, failed\_only=True)

test\_sort\_func(quick\_sort, samples=100, failed\_only=True)

Total tests passed: 100/100

A sorting algorithm is implemented correctly

Implementacja algorytmu #4 (dla funkcyjnej implementacji listy)

Wersja iteracyjna z użyciem stosu, modyfikująca listę źródłową

In (52):

def quick\_sort(ll\_head: 'linked list head (sentinel)'):

# Sort if only there are at least two elements in a linked list

if ll\_head.next and ll\_head.next.next:

# Perform sorting on a linked list

while stack:

# Take the last indices pair out of the stack

begin\_prev\_node, end\_node = stack.pop()

first\_end\_node, first\_length, second\_begin\_prev\_node, second\_length = \_partition(begin\_prev\_node, end\_node)

if first\_length < second\_length:

# Store indices of the longer part at first if has at least 2 elements

stack.append(second\_begin\_prev\_node, end\_node)

# Then store indices of a shorter part (as it will produce less parts after sorting)

# (if has at least 2 elements and we are sure the longer part has also at least 2 elements,

if first\_length > 1:

stack.append(begin\_prev\_node, first\_end\_node)

else:

# Store indices of the longer part at first if has at least 2 elements

if first\_length > 1:

stack.append(begin\_prev\_node, first\_end\_node)

# Then store indices of a shorter part (as it will produce less parts after sorting)

# (if has at least 2 elements and we are sure the longer part has also at least 2 elements,

if second\_length > 1:

stack.append(second\_begin\_prev\_node, end\_node)

# Update the head node of a linked list

ll.head = sentinel.next

# Scan a linked list linearly to find the tail pointer

curr = ll.head

while curr:

curr = curr.next

ll.tail = curr

def \_partition(begin\_prev\_node, end\_node):

# Store a pivot node and a current node pointers in variables

# Take the first (leftmost) node as a pivot

pivot\_node = begin\_prev\_node.next

curr\_node = pivot\_node.next

# Prepare sentinel nodes for sublists which will be created

lt\_pivot\_head = Node()

eq\_pivot\_head = pivot\_node

gt\_pivot\_head = Node()

# Prepare pointers to the sublists

lt\_pivot\_curr = lt\_pivot\_head

eq\_pivot\_curr = eq\_pivot\_head

gt\_pivot\_curr = gt\_pivot\_head

# Prepare variables to calculate length of the result parts

first\_part\_length = second\_part\_length = 0

# Distribute subsequent nodes of a linked list part to appropriate sublists

while curr\_node is not end\_node:

if curr\_node.val < pivot\_node.val:

lt\_pivot\_curr.next = curr\_node

lt\_pivot\_curr = lt\_pivot\_curr.next

elif curr\_node.val == pivot\_node.val:

eq\_pivot\_curr.next = curr\_node

eq\_pivot\_curr = eq\_pivot\_curr.next

else:

gt\_pivot\_curr.next = curr\_node

gt\_pivot\_curr = gt\_pivot\_curr.next

curr\_node = curr\_node.next

# Join created lists together

# Link a list of elements lower than pivot (lt\_pivot) if is not empty

if lt\_pivot\_curr.next:

begin\_prev\_node.next = lt\_pivot\_head.next

lt\_pivot\_curr.next = eq\_pivot\_head

else:

eq\_pivot\_curr.next = end\_node

# Link a list of elements greater than pivot (gt\_pivot) if is not empty

elif gt\_pivot\_curr.next:

begin\_prev\_node.next = eq\_pivot\_head

eq\_pivot\_curr.next = gt\_pivot\_head.next

gt\_pivot\_curr.next = end\_node

else:

begin\_prev\_node.next = end\_node

eq\_pivot\_curr.next = end\_node

return eq\_pivot\_head, first\_part\_length, eq\_pivot\_curr, second\_part\_length

Zapisujemy funkcję do późniejszego benchmarku (nie jest to część algorytmu)

In (53):

sorting\_functions['Quick Sort #4 (functional, iterative)'] = (quick\_sort, create\_linked\_list, {'partition': partition})

Kilka testów

In (54):

test\_sort\_func(quick\_sort, samples=100, failed\_only=True)

Total tests passed: 100/100

A sorting algorithm is implemented correctly

Algorytmy szybkie (liniowe)

O złożoności  $O(n)$

Bucket Sort

Oczywiście, aby ten algorytm działał jak najlepiej, dane wejściowe powinny być o rozkładzie normalnym (patrz plik z omówieniem algorytmów sortowania dla tablic)

Implementacja algorytmu #1 (z pomocą Insertion Sorta) (dla obiektowej implementacji listy)

Wersja dla DOWOLNYCH liczb RZECZYWISTYCH

Zdecydowaliśmy się na wybór Insertion Sorta do posortowania wiaderek, ponieważ jest on stabilnym algorytmem sortowania oraz dla małych danych (do ok. 64 elementów) w przypadku implementacji dla list odyslawczych - patrz porównanie wydajności dla takich danych) jest bardzo szybki. W poniższej implementacji wykorzystuje liczbę wiaderek, która jest równa 1/48 wielkości danych wejściowych tak, aby każde z wiaderek miało odpowiednio niewielką liczbę elementów, dzięki czemu da się je szybko posortować. Oczywiście można użyć innej wartości tego współczynnika, ale, uwzględniając możliwe odchylenia w obie strony od liczby 48, chcę, aby w jak największej liczbie wiaderek, liczba elementów nie przekroczyła zbyt mocno liczby 64.

In (55):

def bucket\_sort(ll\_head: 'linked list head (sentinel)', k: 'threshold' = 64):

if len(ll) < k:

# Create a sentinel node

sentinel = Node()

sentinel.next = ll.head

# Perform sorting and update the pointers

ll.tail = insertion\_sort(sentinel)

ll.head = sentinel.next

else:

# Make a threshold a bit smaller as a number of elements in each

# bucket can slightly vary and we don't want to make unnecessary

# recursive calls.

m = int((7 \* k) / buckets\_count + len(ll)) // m + 1

# Create buckets (a list of sentinel nodes)

buckets = [None] \* m

for i in range(buckets\_count):

sentinel = tail = Node()

# buckets.append(sentinel, tail)

# Calculate an interval in order to store values in proper buckets

min\_val, max\_val = minmax(ll)

val\_interval = (max\_val - min\_val) / buckets\_count

# Distribute values to the proper buckets

curr = ll.head

while curr:

# Calculate the bucket's index depending on how much the

# current value is greater than the lowest one

bucket\_idx = int((curr.val - min\_val) / val\_interval + .5)

buckets[bucket\_idx][1].next = curr

buckets[bucket\_idx][1] = curr

# Sort each bucket separately

for bucket in buckets:

bucket[1] = insertion\_sort(bucket[0])

# Link buckets together to create a sorted list

sorted\_tail = None

for bucket in buckets:

if bucket is not empty, link a list to be result

if bucket[0].next:

sorted\_tail.next = bucket[0].next

sorted\_tail = bucket[0]

# Modify the head and the tail node of the initial linked list

ll.head = sentinel.next

ll.tail = buckets[-1][1]

def minmax(ll):

global\_min = global\_max = ll.tail.val

curr = ll.head

while curr and curr.next:

if curr.val < global\_min:

global\_min = curr.val

if curr.val > global\_max:

global\_max = curr.val

else:

if curr.next.val < global\_min:

global\_min = curr.next.val

if curr.next.val > global\_max:

global\_max = curr.next.val

curr = curr.next

return global\_min, global\_max

def insert\_node(ll\_head: 'linked list head (sentinel)', node): # Inserts node in a right position maintaining

# order of the nodes before a greater one

curr = ll\_head

while node.val < curr.next.val:

curr = curr.next

node.next = curr.next

curr.next = node

def insertion\_sort(ll\_head: 'linked list head (sentinel)'):

if not ll\_head.next: return ll\_head

if not ll\_head.next.next: return ll\_head.next

right = ll\_head.next

while prev.next:

# If a current node (prev.next) has a value lower than a prev node, we have to

# shift this node to a right position before.

if prev.next.val < prev.val:

# Remove a current node

removed = prev.next

prev.next = prev.next.next

# We can skip a current node otherwise.

else:

prev = prev.next

return prev

Zapisujemy funkcję do późniejszego benchmarku (nie jest to część algorytmu)

In (56):

sorting\_functions['Bucket Sort #1 (objective with Insertion Sort)'] = (bucket\_sort, LinkedList, {'minmax': minmax, 'insert\_node': insert\_node, 'insertion\_sort': insertion\_sort})

Kilka testów

In (57):

test\_sort\_func(bucket\_sort, range=-1\_000, 1\_000, val\_counts=(0, 1\_000), samples=1\_000, failed\_only=True)

test\_sort\_obj(bucket\_sort, val\_counts=(0, 1000), samples=1000, failed\_only=True)

test\_sort\_obj(bucket\_sort, val\_counts=(10, 200), samples=25, failed\_only=True)

Total tests passed: 25/100

A sorting algorithm is implemented correctly

Implementacja algorytmu #2 (z pomocą Insertion Sorta) (dla funkcyjnej implementacji listy)

Wersja dla DOWOLNYCH liczb RZECZYWISTYCH

Zdecydowaliśmy się na wybór Insertion Sorta do posortowania wiaderek, ponieważ jest on stabilnym algorytmem sortowania oraz dla małych danych (do ok. 64 elementów) w przypadku implementacji dla list odyslawczych - patrz porównanie wydajności dla takich danych) jest bardzo szybki. W poniższej implementacji wykorzystuje liczbę wiaderek, która jest równa 1/48 wielkości danych wejściowych tak, aby każde z wiaderek miało odpowiednio niewielką liczbę elementów, dzięki czemu da się je szybko posortować. Oczywiście można użyć innej wartości tego współczynnika, ale, uwzględniając możliwe odchylenia w obie strony od liczby 48, chcę, aby w jak największej liczbie wiaderek, liczba elementów nie przekroczyła zbyt mocno liczby 64.

In (58):

def bucket\_sort(ll\_head: 'linked list head (sentinel)', k: 'threshold' = 32):

length = len(ll)

if length < k:

# Perform sorting if only there are at least 2 values in a list

if ll\_head.next and ll\_head.next.next:

# Perform sorting on a linked list

\_bucket\_sort(ll\_head, None)

else:

# Make a threshold a bit smaller as a number of elements in each

# bucket can slightly vary and we don't want to make unnecessary

# recursive calls.

m = int((7 \* k) / buckets\_count + len(ll)) // m + 1

# Create buckets (a list of sentinel nodes)

buckets = [None] \* m

for i in range(buckets\_count):

sentinel = tail = Node()

# buckets.append(sentinel, tail)

# Calculate an interval in order to store values in proper buckets

min\_val, max\_val = minmax(ll)

val\_interval = (max\_val - min\_val) / buckets\_count

# Distribute values to the proper buckets

curr = ll\_head

while curr:

# Calculate the bucket's index depending on how much the

# current value is greater than the lowest one

bucket\_idx = int((curr.val - min\_val) / val\_interval + .5)

tail = buckets[bucket\_idx][1]

tail.next = curr

buckets[bucket\_idx][1] = tail

curr = curr.next

# Sort each bucket separately

for bucket in buckets:

bucket[1] = insertion\_sort(bucket[0])

# Link buckets together to create a sorted list

sorted\_tail = None

for bucket in buckets:

if bucket is not empty, link a list to be result

if bucket[0].next:

sorted\_tail.next = bucket[0].next

sorted\_tail = bucket[0]

# Modify the head and the tail node of the initial linked list

ll.head = sentinel.next

ll.tail = buckets[-1][1]

def minmax(ll\_head):

global\_min = global\_max = curr.val

while curr and curr.next:

if curr.val < global\_min:

global\_min = curr.val

if curr.val > global\_max:

global\_max = curr.val

else:

if curr.next.val < global\_min:

global\_min = curr.next.val

if curr.next.val > global\_max:

global\_max = curr.next.val

curr = curr.next

return global\_min, global\_max

def insert\_node(ll\_head: 'linked list head (sentinel)', node): # Inserts node in a right position maintaining

# order of the nodes before a greater one

curr = ll\_head

while node.val < curr.next.val:

curr = curr.next

node.next = curr.next

curr.next = node

def insertion\_sort(ll\_head: 'linked list head (sentinel)'):

if not ll\_head.next: return ll\_head

if not ll\_head.next.next: return ll\_head.next

prev = ll\_head

while prev.next:

# If a current node (prev.next) has a value lower than a prev node, we have to

# shift this node to a right position before.

if prev.next.val < prev.val:

# Remove a current node

removed = prev.next

prev.next = prev.next.next

# We can skip a current node otherwise.

else:

prev = prev.next

return prev

Zapisujemy funkcję do późniejszego benchmarku (nie jest to część algorytmu)



```
[In (59):] sorting_functions['Bucket Sort #2 (functional with Insertion Sort)'] = (bucket_sort, create_linked_list, ('minmax': minmax, 'insert_node': insert_node, 'insertion_sort': insertion_sort, 'divide_into_buckets': divide_into_buckets, 'link_list': link_list))

Kilka testów

[In (60):] test_sort_func(bucket_sort, val_counts=(20, 200), samples=10000, failed_only=True)
test_sort_func(bucket_sort, val_counts=(10, 200), samples=25, failed_only=True)
Total tests passed: 25/25
A sorting algorithm is implemented correctly

Implementacja algorytmu #3 (z pomocą Counting Sorta) (dla obiektowej implementacji listy)

(Wersja dla WSZYSTKICH liczb KĄKOWITYCH)

Tak naprawdę ten algorytm nie wymaga żadnego sortowania ani zliczania wartości, a jedynie ich włożenie do odpowiednich wiader, a następnie ich połączenie (złączenie list odszykowanych). Oczywiście ten, jak i poniższy algorytm sortowania (dla funkcyjnej implementacji listy), jest niewydajny! pamięciowo, jeżeli sortowana jest tablica bardzo różnych się między sobą wartości.

[In (61):] def bucket_counting_sort(ll: LinkedList):
    if len(ll) > 1:
        min_val, max_val = minmax(ll)
        # Allocate memory (create sentinel nodes) for values (our buckets)
        buckets = [None] * (max_val - min_val + 1)
        for i in range(max_val - min_val + 1):
            sentinel = tail = Node(i)
            buckets.append(sentinel, tail)
        # Store values in the proper buckets
        curr = ll.head
        while curr:
            idx = curr.val - min_val
            tail = buckets[idx][1]
            buckets[idx][1] = curr
            curr = curr.next
        # Link all linked lists together
        tail = Node(0)
        for bucket in buckets:
            if bucket is not None, link a list to be result
            tail.next = bucket[0].next
            tail = bucket[1]
        # Update the pointers of the initial linked list
        tail.next = None
        ll.head = sentinel.next
        ll.tail = tail
    return minmax(ll)

global_min = global_max = ll.tail.val
curr = ll.head
while curr and curr.next:
    if curr.val > global_max: global_max = curr.val
    if curr.val < global_min: global_min = curr.next.val
    else:
        if curr.next.val > global_max: global_max = curr.next.val
        if curr.next.val < global_min: global_min = curr.val
    curr = curr.next
return global_min, global_max

Zapisujemy funkcję do późniejszego benchmarku (nie jest to część algorytmu)

[In (62):] sorting_functions['Bucket Sort #3 (objective with Counting Sort)'] = (bucket_counting_sort, linked_list, ('minmax': minmax, 'insert_node': insert_node, 'insertion_sort': insertion_sort, 'divide_into_buckets': divide_into_buckets, 'link_list': link_list))

Kilka testów

[In (63):] test_sort_obj(bucket_counting_sort, samples=100, failed_only=True)
Total tests passed: 100/100
A sorting algorithm is implemented correctly

Implementacja algorytmu #4 (z pomocą Counting Sorta) (dla funkcyjnej implementacji listy)

(Wersja dla WSZYSTKICH liczb KĄKOWITYCH)

[In (64):] def bucket_counting_sort(ll: head: 'linked list head (sentinel)'):
    if len(ll) > 1:
        min_val, max_val = minmax(ll.head)
        # Allocate memory (create sentinel nodes) for values (our buckets)
        buckets = [None] * (max_val - min_val + 1)
        for i in range(max_val - min_val + 1):
            sentinel = tail = Node(i)
            buckets.append(sentinel, tail)
        # Store values in the proper buckets
        curr = ll.head.next
        while curr:
            idx = curr.val - min_val
            tail = buckets[idx][1]
            buckets[idx][1] = curr
            curr = curr.next
        # Link all linked lists together
        tail = Node(0)
        for bucket in buckets:
            if bucket is not empty, link a list to be result
            tail.next = bucket[0].next
            tail = bucket[1]
        # Update the pointers of the initial linked list
        tail.next = None
        ll.head.next = sentinel.next
        ll.head.next = None
    return minmax(ll)

global_min = global_max = curr.val
curr = ll.head.next
while curr and curr.next:
    if curr.val > global_max: global_max = curr.val
    if curr.val < global_min: global_min = curr.next.val
    else:
        if curr.next.val > global_max: global_max = curr.next.val
        if curr.next.val < global_min: global_min = curr.val
    curr = curr.next
return global_min, global_max

Zapisujemy funkcję do późniejszego benchmarku (nie jest to część algorytmu)

[In (65):] sorting_functions['Bucket Sort #4 (functional with Counting Sort)'] = (bucket_counting_sort, create_linked_list, ('minmax': minmax, 'insert_node': insert_node, 'insertion_sort': insertion_sort, 'divide_into_buckets': divide_into_buckets, 'link_list': link_list))

Kilka testów

[In (66):] test_sort_func(bucket_counting_sort, samples=100, failed_only=True)
Total tests passed: 100/100
A sorting algorithm is implemented correctly

Implementacja algorytmu #5 (z pomocą Bucket Sorta) (dla obiektowej implementacji listy)

Wersja dla DOWOLNYCH liczb RZECZYWISTYCH

Ulepszona wersja Bucket Sorta, która sortuje duże wiaderka ponownie Bucket Sortem. Poniższa wersja algorytmu jest wersją iteracyjną, ale łatwo zamienić w wersję rekurencyjną.

[In (67):] def bucket_sort(ll: LinkedList, k: 'threshold') = 48:
    if len(ll) < 2: return
    sentinel = Node(0)
    bucket_stack = [sentinel, ll.tail, len(ll)]
    res_curr = curr = Node(0)
    while bucket_stack:
        # Get the last bucket from a stack
        bucket = bucket_stack.pop()
        # If it is small enough, sort it and link to the result linked list
        if bucket[2] < k:
            # Link a sorted part to the result linked list
            tail = insertion_sort(bucket[0])
            res_curr.next = bucket[0].next
            res_curr = tail
        # Split a bucket into smaller buckets otherwise
        else:
            min_val, max_val = minmax(bucket[0])
            # If a bucket contains all the same values, link it to the result linked list
            if max_val == min_val:
                res_curr.next = bucket[0].next
                res_curr = bucket[1]
            # If not, divide this bucket into smaller buckets
            else:
                buckets = divide_into_buckets(bucket, min_val, max_val, k)
                for i in range(len(buckets)-1, -1, -1):
                    if buckets[i][2] > k:
                        bucket_stack.append(buckets[i])
    # Fix linked list pointers
    ll.head = res_curr.next
    ll.tail = res_curr.next
    return curr

def divide_into_buckets(bucket, min_val, max_val, k: 'threshold'):
    m = int((2/3) * k)
    buckets_count = bucket[2] // m + 1
    val_interval = (max_val - min_val) / buckets_count
    # Create buckets
    buckets = [None] * buckets_count
    sentinel = tail = Node(0)
    # Distribute values to the proper buckets
    curr = bucket[0].next
    while curr:
        # Calculate the bucket's index depending on how much the
        # current value is greater than the lowest one
        bucket_idx = int((curr.val - min_val) / val_interval - .5)
        buckets[bucket_idx][1].next = curr
        buckets[bucket_idx][1] = curr
        curr = curr.next
    # Unlink next nodes of tails
    for bucket in buckets:
        bucket[1].next = None
    return buckets

def minmax(ll: head: 'linked list head (sentinel)'):
    curr = ll.head.next
    global_min = global_max = curr.val
    while curr and curr.next:
        if curr.val > global_max: global_max = curr.val
        if curr.val < global_min: global_min = curr.next.val
        else:
            if curr.next.val > global_max: global_max = curr.next.val
            if curr.next.val < global_min: global_min = curr.val
    curr = curr.next
    return global_min, global_max

def insert_node(ll: head: 'linked list head (sentinel)', node): # Inserts node in a right position maintaining
    curr = ll.head
    while node.val > curr.next.val:
        curr = curr.next
    Merge Sort next = curr.next
    curr.next = node
    node.next = curr.next

def insertion_sort(ll: head: 'linked list head (sentinel)'):
    if not ll.head.next: return ll.head
    prev = ll.head.next
    while prev.next:
        # If a current node (prev.next) has a value lower than a prev node, we have to
        # shift this node to a right position before.
        if prev.next.val < prev.val:
            # Remove a current node
            removed = prev.next
            # Now we insert this node in a right position
            insert_node(ll.head, removed)
            # We can skip a current node otherwise.
            else:
                prev = prev.next
    return prev

Zapisujemy funkcję do późniejszego benchmarku (nie jest to część algorytmu)

[In (68):] sorting_functions['Bucket Sort #5 (objective iterative with Bucket Sort)'] = (bucket_sort, linked_list, ('minmax': minmax, 'insert_node': insert_node, 'insertion_sort': insertion_sort, 'divide_into_buckets': divide_into_buckets, 'link_list': link_list))

Kilka testów

[In (69):] # test_sort_obj(bucket_sort, val_counts=(100, 1000), samples=1000, failed_only=True)
test_sort_obj(bucket_sort, val_counts=(100, 200), samples=25, failed_only=True)
Total tests passed: 25/25
A sorting algorithm is implemented correctly

Implementacja algorytmu #6 (z pomocą Bucket Sorta) (dla funkcyjnej implementacji listy)

Wersja dla DOWOLNYCH liczb RZECZYWISTYCH

Ulepszona wersja Bucket Sorta, która sortuje duże wiaderka ponownie Bucket Sortem. Poniższa wersja algorytmu jest wersją iteracyjną, ale łatwo zamienić w wersję rekurencyjną.

[In (70):] def bucket_sort(ll: head: 'linked list head (sentinel)', k: 'threshold') = 32:
    curr = ll.head
    while curr:
        # Get linked list tail and length
        length = len(curr.next)
        curr = curr.next
        # Create the bucket_stack
        bucket_stack = [ll.head, curr, length]
        res_curr = res_curr = Node(0)
        while bucket_stack:
            # Get the last bucket from a stack
            bucket = bucket_stack.pop()
            # If it is small enough, sort it and link to the result linked list
            if bucket[2] < k:
                # Link a sorted part to the result linked list
                tail = insertion_sort(bucket[0])
                res_curr.next = bucket[0].next
                res_curr = tail
            # Split a bucket into smaller buckets otherwise
            else:
                min_val, max_val = minmax(bucket[0])
                # If a bucket contains all the same values, link it to the result linked list
                if max_val == min_val:
                    res_curr.next = bucket[0].next
                    res_curr = bucket[1]
                # If not, divide this bucket into smaller buckets
                else:
                    buckets = divide_into_buckets(bucket, min_val, max_val, k)
                    for i in range(len(buckets)-1, -1, -1):
                        if buckets[i][2] > k:
                            bucket_stack.append(buckets[i])
        # Fix linked list pointers
        ll.head = res_curr.next
        ll.tail = res_curr.next
        return curr

def divide_into_buckets(bucket, min_val, max_val, k: 'threshold'):
    m = int((2/3) * k)
    buckets_count = bucket[2] // m + 1
    val_interval = (max_val - min_val) / buckets_count
    # Create buckets
    buckets = [None] * buckets_count
    sentinel = tail = Node(0)
    # Distribute values to the proper buckets
    curr = bucket[0].next
    while curr:
        # Calculate the bucket's index depending on how much the
        # current value is greater than the lowest one
        bucket_idx = int((curr.val - min_val) / val_interval - .5)
        buckets[bucket_idx][1].next = curr
        buckets[bucket_idx][1] = curr
        curr = curr.next
    # Unlink next nodes of tails
    for bucket in buckets:
        bucket[1].next = None
    return buckets

def minmax(ll: head: 'linked list head (sentinel)'):
    curr = ll.head.next
    global_min = global_max = curr.val
    while curr and curr.next:
        if curr.val > global_max: global_max = curr.val
        if curr.val < global_min: global_min = curr.next.val
        else:
            if curr.next.val > global_max: global_max = curr.next.val
            if curr.next.val < global_min: global_min = curr.val
    curr = curr.next
    return global_min, global_max

def insert_node(ll: head: 'linked list head (sentinel)', node): # Inserts node in a right position maintaining
    curr = ll.head
    while node.val > curr.next.val:
        curr = curr.next
    node.next = curr.next
    curr.next = node

def insertion_sort(ll: head: 'linked list head (sentinel)'):
    if not ll.head.next: return ll.head
    prev = ll.head.next
    while prev.next:
        # If a current node (prev.next) has a value lower than a prev node, we have to
        # shift this node to a right position before.
        if prev.next.val < prev.val:
            # Remove a current node
            removed = prev.next
            # Now we insert this node in a right position
            insert_node(ll.head, removed)
            # We can skip a current node otherwise.
            else:
                prev = prev.next
    return prev

Zapisujemy funkcję do późniejszego benchmarku (nie jest to część algorytmu)

[In (71):] sorting_functions['Bucket Sort #6 (functional iterative with Bucket Sort)'] = (bucket_sort, create_linked_list, ('minmax': minmax, 'insert_node': insert_node, 'insertion_sort': insertion_sort, 'divide_into_buckets': divide_into_buckets, 'link_list': link_list))

Kilka testów

[In (72):] # test_sort_func(bucket_sort, val_counts=(100, 200), samples=10000, failed_only=True)
test_sort_func(bucket_sort, val_counts=(100, 200), samples=25, failed_only=True)
Total tests passed: 25/25
A sorting algorithm is implemented correctly

Porównanie szybkości

[In (73):] for i in range(1, len(sorting_functions.items())):
    print(f'{i} {sorting_functions.items()[i][0]}')

[0] Insertion Sort #1 (objective)
[1] Insertion Sort #2 (functional)
[2] Selection Sort #1 (objective)
[3] Selection Sort #2 (functional)
[4] Bubble Sort #1 (objective)
[5] Bubble Sort #2 (functional)
[6] Merge Sort #1 (objective, recursive)
[7] Merge Sort #2 (objective, recursive optimized)
[8] Merge Sort #3 (objective, iterative, natural series)
[9] Merge Sort #4 (functional, recursive)
[10] Merge Sort #5 (functional, recursive optimized)
[11] Merge Sort #6 (functional, iterative, natural series)
[12] Quick Sort #1 (objective, recursive)
[13] Quick Sort #2 (functional, recursive)
[14] Quick Sort #3 (functional, recursive)
[15] Quick Sort #4 (functional, iterative)
[16] Quick Sort #5 (objective with Insertion Sort)
[17] Bucket Sort #2 (functional with Insertion Sort)
[18] Bucket Sort #3 (objective with Counting Sort)
[19] Bucket Sort #4 (functional with Counting
```