

EJERCICIO 10

a. b. $T(n) = \begin{cases} 2 & n=1 \\ T(n-1)+n & , n \geq 2 \end{cases}$

Suponemos que $n \geq 2$:

$T(n) =$
Paso 1: $\downarrow T(n-1)+n$

$T(n) =$
Paso 2: $\downarrow T(n-2) + (n-1) + n$

$T(n) =$
Paso 3: $\downarrow T(n-3) + (n-2) + (n-1) + n$

$T(n) =$
Paso i: $\downarrow T(n-i) + (n-i+1) + (n-i+2) + \dots + (n-i+i)$

Paso n-1: $T(n) = 2 + \sum_{i=2}^n i$

$T(n) = 2 + \sum_{i=2}^n i - 1 = 2 + \frac{n \cdot (n+1)}{2} - 1 = 1 + \frac{n^2 + n}{2} = O(n^2)$

Existen constantes $c > 0$ y N_0 tales que:

$1 + \frac{n^2 + n}{2} \leq c \cdot n^2$ para todo $n \geq n_0$, $n \in \mathbb{N}$

$1 \leq c_1 \cdot n^2$	$(n^2 + n) : 2 \leq c_2 n^2$
$1/n^2 \leq c_1$	$\frac{n+1}{2n} \leq c_2$
$n_1 = 1$	$n_2 = 1$
$c_1 = 1$	$c_2 = 1$

$1 + \frac{n^2 + n}{2} \leq (c_1 + c_2) \cdot n^2$

$T(n) \leq (1+1) \cdot n^2$

$T(n) \leq c \cdot n^2$

$T(n) \leq O(n^2)$, con $c=2$ para todo $n \geq n_0$, con $n_0 = 1$

$$T(n) = \begin{cases} 2 & n=1 \\ T(n-1) + n/2 & n \geq 2 \end{cases}$$

Paso 1: $T(n-1) + n/2$

Paso 2: $T(n) = T(n-2) + \frac{n-1}{2} + n/2$

Paso 3: $T(n) = T(n-3) + \frac{n-2}{2} + \frac{n-1}{2} + n/2$

Paso i: $T(n) = T(n-i) + \frac{(n-i+1)}{2} + \frac{(n-i+2)}{2} + \dots + \frac{n}{2}$

$n-i=1$
 $i=n-1$

Reemplazo i:

$$T(n) = T(1) + \frac{(n-(n-1)+1)}{2} + \frac{(n-(n-1)+2)}{2} + \dots + \frac{n}{2}$$

$$T(n) = 2 + \frac{1}{2} + \frac{3}{2} + \frac{5}{2} + \dots + \frac{n}{2}$$

$$T(n) = 2 + \sum_{i=2}^n i \cdot \frac{1}{2}$$

$$T(n) = 2 + \frac{1}{2} \cdot \left(\frac{n(n+1)}{2} - 1 \right) = 2 + \frac{1}{2} \left(\frac{n^2 + n - 1}{2} \right) = 2 + \frac{n^2}{4} + \frac{n}{4} - \frac{1}{2} = \frac{n^2}{4} + \frac{n}{4} + \frac{3}{4} \therefore O(n^2)$$

Existen constantes $c > 0$ y No tales que:

$$\frac{n^2}{4} + \frac{n}{4} + \frac{3}{4} \leq c \cdot n^2 \text{ para todo } n \geq n_0 \quad n \in \mathbb{N}$$

PRIMER TÉRMINO

$$\frac{n^2}{4} \leq c_1 n^2$$

$$\frac{n^2}{4} : n^2 \leq c_1$$

$$\frac{1}{4} \leq c_1$$

$$n_1 = 0 \quad c_1 = 1$$

SEGUNDO TÉRMINO

$$\frac{n}{4} \leq c_2 n^2$$

$$1/4n \leq c_2$$

$$n_2 \geq 1$$

$$c_2 = 1$$

TERCER TÉRMINO:

$$\frac{3}{4} \leq c_3 n_3$$

$$\frac{3}{4n} \leq c_3$$

$$n_3 = 1$$

$$c_3 = 1$$

$$\frac{n^2}{4} + \frac{n}{4} + \frac{3}{4} \leq (c_1 + c_2 + c_3) \cdot n^2$$

$$T(n) \leq (1+1+1) \cdot n^2$$

$$T(n) \leq c \cdot n^2$$

$$T(n) \leq O(n^2) \text{ con } c=3 \text{ para todo } n \geq n_0, \text{ con } n_0=1$$

$$T(n) = \begin{cases} 1 & n=1 \\ 2T\left(\frac{n}{4}\right) + \sqrt{n} & , n \geq 2 \end{cases}$$

$$\text{Paso 1: } T(n) = 2T(n/4) + \sqrt{n}$$

$$\text{Paso 2: } T(n) = 2[2T(n/16) + \sqrt{n/4}] + \sqrt{n}$$

$$T(n) = 4T(n/16) + 2\sqrt{n/4} + \sqrt{n} = 4T(n/16) + 2\sqrt{n}$$

$$\text{Paso 3: } T(n) = 4[2T(n/48) + \sqrt{n/16}] + 2\sqrt{n}$$

$$T(n) = 8T(n/48) + 4\sqrt{n/16} + 2\sqrt{n} = 8T(n/48) + 3\sqrt{n}$$

$$\text{Paso } i: T(n) = 2^i T(n/4^i) + i \cdot \sqrt{n} \quad \rightarrow n/4^i = 1 \rightarrow \log_4 n = i$$

$$\text{Reemplazo } i: T(n) = 2^{\log_4 n} \cdot T(n/4^{\log_4 n}) + \log_4 n \cdot \sqrt{n}$$

$$T(n) = 2^{\log_4 n} \cdot T(1) + \log_4 n \cdot \sqrt{n} = \sqrt{n} \cdot 1 + \log_4 n \cdot \sqrt{n} = O(\log_4(n) \cdot \sqrt{n})$$

Existen constantes $c > 0$ y N_0 tales que:

$$\sqrt{n} + \log_4(n) \cdot \sqrt{n} \leq c \cdot \log_4(n) \cdot \sqrt{n} \text{ para todo } n \geq n_0, n \in \mathbb{N}$$

$$\sqrt{n} \leq c_1 \cdot \log_4(n) \cdot \sqrt{n}$$

$$\sqrt{n} : \sqrt{n} \leq c_1 \cdot \log_4(n)$$

$$1 \leq c_1 \cdot \log_4(n)$$

$$n_1 = 4 \\ c_1 = 1$$

$$\log_4(n) \cdot \sqrt{n} \leq c_2 \cdot \log_4(n) \cdot \sqrt{n}$$

$$(\log_4(n) \cdot \sqrt{n}) : (\log_4(n) \cdot \sqrt{n}) \leq c_2$$

$$1 \leq c_2$$

$$n_2 = 1 \\ c_2 = 1$$

$$\sqrt{n} + \log_4(n) \cdot \sqrt{n} \leq (c_1 + c_2) \cdot \log_4(n) \cdot \sqrt{n}$$

$$T(n) \leq (1+1) \cdot f(n)$$

$$T(n) \leq c \cdot f(n)$$

$$T(n) \leq O(\log_4(n) \cdot \sqrt{n}) \text{ para todo}$$

$$n \geq n_0, \text{ con } n_0 = 4; \text{ y con } c=2$$

$$T(n) = \begin{cases} 1 & n=1 \\ 4T(n/2) + n^2 & n \geq 2 \end{cases}$$

Paso 1: $4T(n/2) + n^2$

Paso 2: $4[4T(n/4) + (n/2)^2] + n^2$

$$T(n) = 16T(n/4) + 4(n^2/4) + n^2 = 16T(n/4) + 2n^2$$

Paso 3: $T(n) = 16[4T(n/8) + (n/4)^2] + 2n^2$

$$T(n) = 64T(n/8) + 16(n^2/16) + 2n^2 = 64T(n/8) + 3n^2$$

Paso i: $T(n) = 4^i T(n/2^i) + i \cdot n^2$

$n/2^i = 1$
 $i = \log_2 n$

Reemplazo i: $T(n) = 4^{\log_2 n} \cdot T(n/2^{\log_2 n}) + \log_2 n \cdot n^2$

$$T(n) = (2 \cdot 2)^{\log_2 n} \cdot T(1) + \log_2 n \cdot n^2$$

$$T(n) = n \cdot n \cdot 1 + \log_2 n \cdot n^2 \therefore \log_2(n) n^2$$

Existen constantes $c > 0$ y n_0 tales que:

$$n^2 + \log_2(n) \cdot n^2 \leq c \cdot \log_2(n) \cdot n^2 \text{ para todo } n > n_0, n \in \mathbb{N}$$

$$n^2 \leq c_1 \cdot \log_2(n) \cdot n^2 \quad \log_2(n) \cdot n^2 \leq c_2 \cdot \log_2(n) \cdot n^2$$

$$n^2/n^2 \leq c_1 \cdot \log_2(n) \quad (\log_2(n) \cdot n^2) / (\log_2(n) \cdot n^2) \leq c_2$$

$$1 \leq c_1 \cdot \log_2(n)$$

$$1 \leq c_2$$

$$n_1 = 2$$

$$c_1 = 1$$

$$n_2 = 0$$

$$c_2 = 1$$

$$\hookrightarrow n^2 + \log_2(n) \cdot n^2 \leq (c_1 + c_2) \cdot (\log_2(n) \cdot n^2)$$

$$T(n) \leq (1+1) \cdot (\log_2(n) \cdot n^2)$$

$$T(n) \leq c \cdot (\log_2(n) \cdot n^2)$$

$$T(n) \leq O(\log_2(n) \cdot n^2), \text{ con } c=2 \text{ para todo } n \geq n_0, \text{ con } n_0=2$$