

Sistemas de Informação e Bases de Dados 2020/2021

Aula 02: Introdução às Bases de Dados e ao SQL

Prof. Paulo Carreira

Class Outline

- How a database Organises Data
 - ☐ Tables, Columns, Rows, and Cells
 - Key, Primary Key, and Foreign Key
 - Table Schema and Database Schema
- Introduction to SQL
 - Creating database tables and populating them with information
 - Basic query structure: SELECT
 - Filters and Pattern Matching
 - Set operations and Joins



How a Database Organises Data



The Table



A two-dimensional grid

A **two-dimensional** grid that keeps data (facts) about (real-world) **Entities**

ID	ID Name		T-Shirt
76440	João Guilherme Silva da Cunha	12345678	М
76469	Tomás Pinto dos Santos	91234567	М
76500	David Miguel Redwanz Duque	89012345	L
64833	Pedro Daniel Diz Pinela	67890123	М
76464	Guilherme de Queiróz Rebelo Brum Gomes	22394856	XL
78230	Marta Isabel de Almeida Cardoso	34562732	S
78083	Filipe Emanuel Lourenço Ramalho Fernandes	82533235	L
78081	Gabriel Filipe Queirós Mesquita Delgado Freire	23134539	М
78001	João Gomes Vultos Freitas	22231233	L
76504	Ricardo Afonso Rodrigues da Silva Oliveira	56372848	L



Columns, lines, and cells

Data is organised into columns, lines, and cells

Columns are named and identified by Attributes

Lines are called 'Records' or 'Rows' ----

Data is recorded in Cells called ——
'Fields'

	V		7
ID	Name	TAX ID	T-Shirt
76440	João Guilherme Silva da Cunha	12345678	М
76469	Tomás Pinto dos Santos	91234567	М
76500	David Miguel Redwanz Duque	89012345	L
64833	Pedro Daniel Diz Pinela	67890123	М
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78001	João Gomes Vultos Freitas	22231233	L
76504	Ricardo Afonso Rodrigues da Silva Oliveira	56372848	L



Data Types, Domains, and **Formats**

For each column, the data in every cell has the same:

- Data Domain: Non-negative, min and max value range, minimum and maximum length, allowed characters, words, or codes from a set of values
- Data Type: Numeric, Text, Date, Binary
- Data Format: 5 digits, Lowercase vs Capitalised, License Plate, YYYY/MM/DD, etc

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Scnema	= Columns	→ Data	pomain

ID	Name	TAX ID	T-Shirt
76440	João Guilherme Silva da Cunha	12345678	M
76469	Tomás Pinto dos Santos	91234567	М

Every Column has an associated Domain, Type and Format



Header

Table Data

Atomic values

Text

```
'John Smith', 'R2D2', 'Red', ''
```

Numeric values

Floats and Doubles

Dates



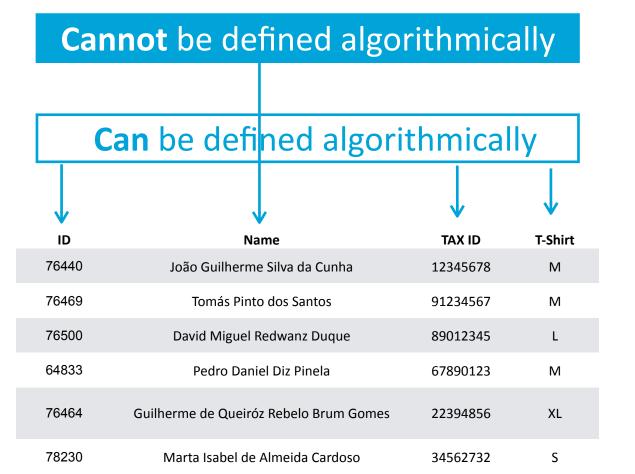
Table

Definition

An organisation of data into columns and lines where all <u>lines have a similar structure</u>, defined by <u>the schema</u> of the table



Algorithmically- vs. nonalgorithmically definable column domains





Ordering of columns

If the ordering of columns is changed, <u>information</u> remains the same

ID	Name	TAX ID	T-Shirt
76440	João Guilherme Silva da Cunha	12345678	M
76469	Tomás Pinto dos Santos	91234567	М
76500	David Miguel Redwanz Duque	89012345	L
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Ordering of columns

In terms of information contents, the order of the columns is irrelevant

Name	ID	TAX ID	T-Shirt
João Guilherme Silva da Cunha	76440	12345678	М
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David Miguel Redwanz Duque	76500	89012345	L
Pedro Daniel Diz Pinela	64833	67890123	М
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João Gomes Vultos Freitas	78001	22231233	L
Ricardo Afonso Rodrigues da Silva Oliveira	76504	56372848	L



Ordering of rows

When the order of rows is changed, <u>information</u> remains the same

ID	ID Name		T-Shirt	
76440	João Guilherme Silva da Cunha	12345678	М	
76469	Tomás Pinto dos Santos	91234567	М	
76500	David Miguel Redwanz Duque	89012345	L	
64833	Pedro Daniel Diz Pinela	67890123	М	
76464	Guilherme de Queiróz Rebelo Brum Gomes	22394856	XL	
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76504	Ricardo Afonso Rodrigues da Silva Oliveira	56372848	L	



Ordering of rows

In terms of information contents, the order of the rows is irrelevant

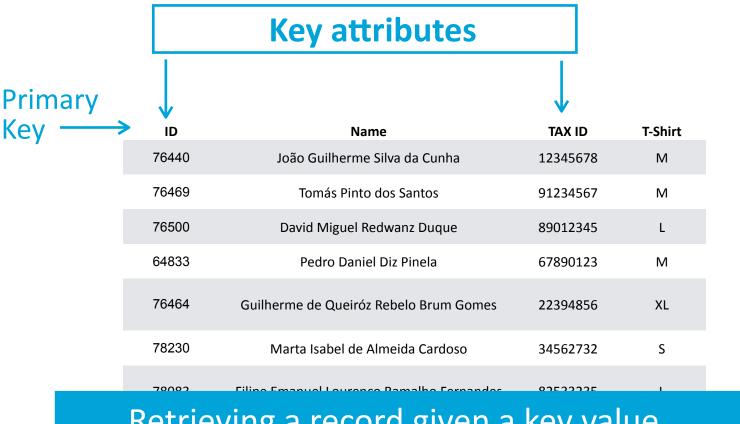
ID	Name	TAX ID	T-Shirt
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76469	Tomás Pinto dos Santos	91234567	М
76500	David Miguel Redwanz Duque	89012345	L
64833	Pedro Daniel Diz Pinela	67890123	М

The row number is irrelevant



Keys

Key values are the way to refer uniquely to a row



Retrieving a record given a key value is a fundamental feature/operation of database systems



15

Table Relationships

A foreign key connects two tables

Foreign Key

Employee



ID	Name	TAX ID	T-Shirt	DID				
76440	João Guilherme Silva da Cunha	12345678	М	EN				
76469	Tomás Pinto dos Santos	91234567	М	EN	_		Departmen	t
76500	David Miguel Redwanz Duque	89012345	L	MK		DID	Name	-
64833	Pedro Daniel Diz Pinela	67890123	M	HR		HR	Human Resources	
76464	Guilherme de Queiróz Rebelo Brum Gomes	22394856	XL	EN		EN	Software Engineering	
78230	Marta Isabel de Almeida Cardoso	34562732	S	HR		MK	Marketing	
78083	Filipe Emanuel Lourenço Ramalho Fernandes	82533235	L	EN		IVIIX	iviai ketilig	
78081	Gabriel Filipe Queirós Mesquita Delgado Freire	23134539	М	EN		T _{Pri} Ke	mary	
78001	João Gomes Vultos Freitas	22231233	L	EN		Ke	y	
76504	Ricardo Afonso Rodrigues da Silva Oliveira	56372848	L	MK				

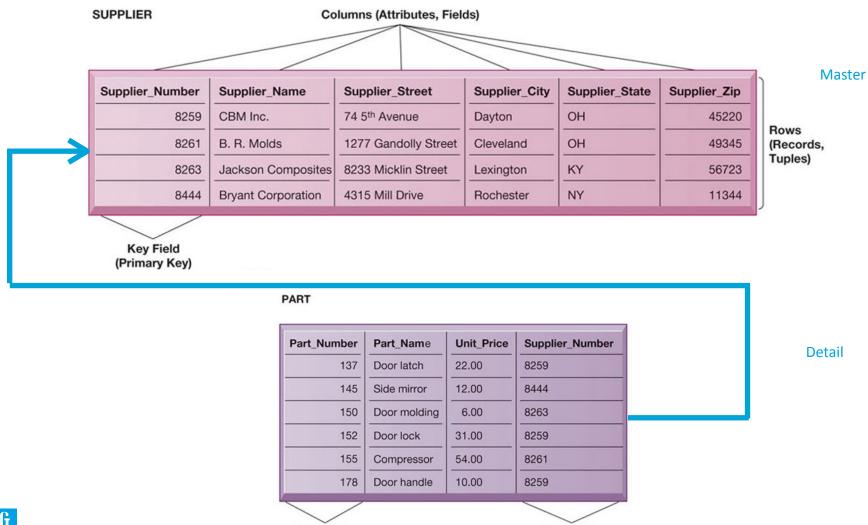
Budget

50 000

1 200 000

150 000

Master-detail related tables



Foreign Key

Primary Key



Tables Concepts Summary

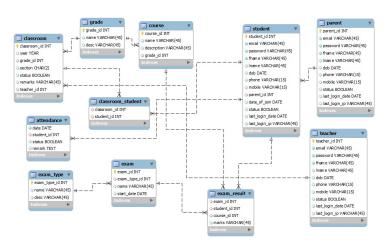
- Rows: Records data different entities
- Columns: Represents attributes for entity
- Fields: Represents attribute for entity
- Key field: Field used to uniquely identify each record
- Primary key: Field in table used for key fields
- Foreign key: Primary key used in second table as look-up field to identify records from original table



Databases and Database Schemas



Database Schema



- Individual schemas of tables: Define which data can be captured
- 2. Relationships between tables: Define how data can be related and navigated (the navigational structure)
- Integrity Constraints: Define which data are not valid in the database



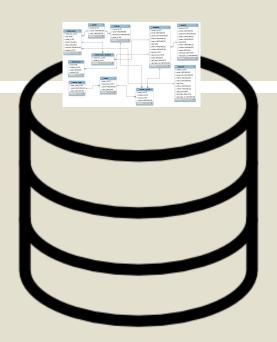
What is kept in the database

Schema

The description of data

Data

The actual data contents





Database

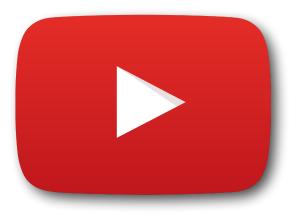
Definition

- 1. An **organised collection** of **data** (typically of large size) that is inter-relatable
- 2. A computerised system, that allows data to be easily accessed, manipulated and updated



See also

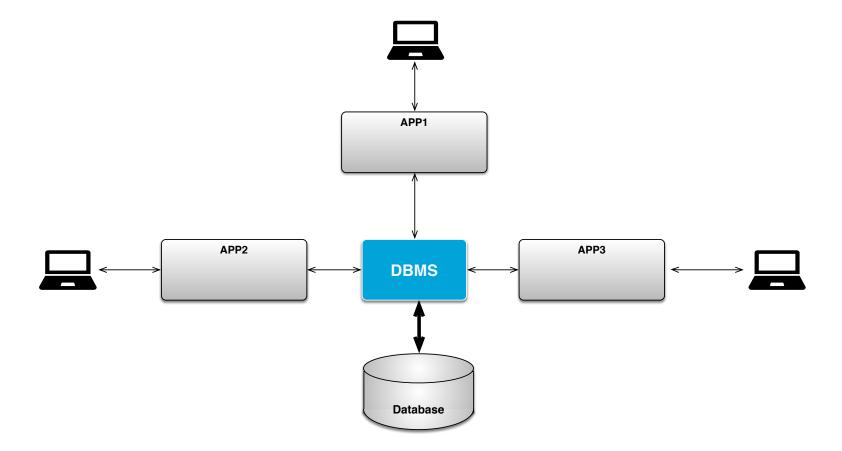
What is a Database?



https://www.youtube.com/watch?v=t8jgX1f8kc4



DBMS



Applications do not contact the database directly



Database Management System (DBMS)

Definition

A **software package** that manages a Database allows data to be easily accessed, manipulated, and updated



Introducing SQL



SQL Language

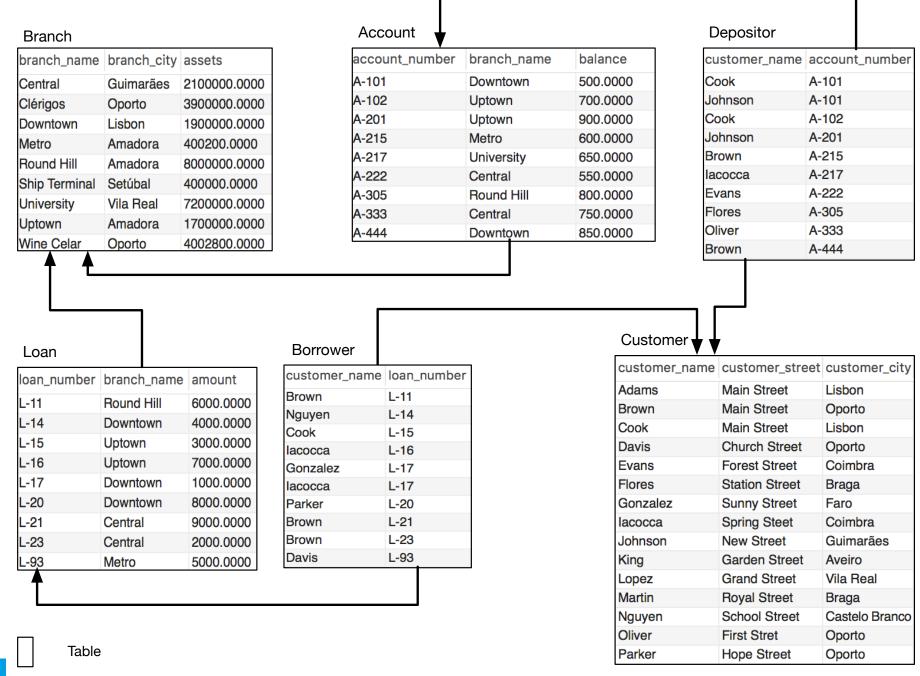
Organised into two sub-languages (parts):

- 1. Data Definition Language (SQL-DDL)
- 2. Data Query Language (SQL-DQL)



The Bank running example





→ Value dependency (field)

Creating Tables



The CREATE TABLE statement

A new table is created with the command:

```
CREATE TABLE <table-name> (
    column<sub>1</sub> type<sub>1</sub>,
    ...
    column<sub>n</sub> type<sub>n</sub>,
    <integrity-constraint<sub>1</sub>>,
    ...
    <integrity-constraint<sub>m</sub>>
); [storage-options]
```

The syntax of the storage options is dependent on the SGBD manufacturer



CREATE TABLE: Example

```
CREATE TABLE department(
did INTEGER,
name VARCHAR(80),
budget NUMERIC(12,4),
PRIMARY KEY(did));
```





Removing a table

Removing a table

DROP TABLE < table-name>

Removing the table department from the database

DROP TABLE department;



Specifying foreign key constraints

```
CREATE TABLE account
  (account_number CHAR(5),
  branch_name VARCHAR(80),
  balance NUMERIC(16,4),
  PRIMARY KEY(account_number),
  FOREIGN KEY(branch_name)
    REFRENCES branch(branch_name));
```



Creating, updating, and deleting rows



Inserting and querying rows

```
INSERT INTO <table-name>
    VALUES(<field-values>)
```

Inserting a record into a table

```
INSERT INTO department VALUES(1,
   'Finance', 1000.0);
```

Querying the records of a table

```
SELECT * FROM < table-name>
```



Schema Definition – Example 1

```
CREATE TABLE department(
    did INTEGER,
    name VARCHAR(80),
    budget NUMERIC(12,2));

INSERT INTO department VALUES(1, 'Finance', 1000.0);
INSERT INTO department VALUES(2, 'Marketing', 2000.0);
```

```
SELECT * FROM department;
```

```
did | name | budget
----+
1 | Finance | 1000.0000
2 | Marketing | 2000.0000
```



Schema Definition – Example 2

```
INSERT INTO department VALUES(1, 'Finance', 1000.0);
INSERT INTO department VALUES(2, 'Marketing', 2000.0);
INSERT INTO department VALUES(1, 'Engineering', 5000.0);
INSERT INTO department VALUES(3, 'Marketing', 3000.0);
SELECT * FROM department;
```

did	name	budget
1	Finance	1000.0000
2	Marketing	2000.0000
1	Engineering	5000.0000
3	Marketing	3000.0000



Schema Definition – Example 3

```
CREATE TABLE department(
    did INTEGER,
    name VARCHAR(80),
    budget NUMERIC(12,4),
    PRIMARY KEY(did),
    UNIQUE(name));
```

```
INSERT INTO department VALUES(1, 'Finance', 1000.0);
INSERT INTO department VALUES(2, 'Marketing', 2000.0);
```

```
INSERT INTO department VALUES(1, 'Engineering', 5000.0);
INSERT INTO department VALUES(3, 'Marketing', 3000.0);
```





Retrieving Data



What is the result of a query?

name	grade
Maria	10
Joseph	11
Chris	12
Ana	13
John	14
	Maria Joseph Chris Ana

Who are the students with grades above 11?

sidnamegrade53650Chris1253831Ana1353832John14

What is the percent of students that finished with a grade greater that 13?





Basic Query Structure

- SQL is based on set and relational operations with certain modifications and enhancements
- The fundamental idea is that of selecting the records that satisfy a given set of criteria; SQL allows for the definition of arbitrarily complex criteria.
- A typical SQL query has the form:

```
select C<sub>1</sub>, C<sub>2</sub>, ..., C<sub>n</sub>
from table
where condition
```



The result of an SQL query is also a table

The **SELECT** clause

- The select clause lists the attributes desired in the result of a query
- **Example**: find the names of all branches in the loan relation:

```
select branch_name
from loan;
```



NB: SQL names are case insensitive (i.e., you may use upper- or lower-case characters)



The **SELECT** clause (cont.)

An asterisk in the select clause denotes "all attributes"

```
select *
from loan;
```

loan_number	branch_name +	amount
L-17 L-23 L-15 L-14 L-93 L-11 L-16 L-20 L-21	Downtown Central Uptown Downtown Metro Round Hill Uptown Downtown Central	1000.0000 2000.0000 3000.0000 4000.0000 5000.0000 6000.0000 7000.0000 8000.0000



Expressions and derived fields



The **SELECT** clause (cont.)

The select clause can contain arithmetic expressions involving the operations +, −, *, and /, as well as functions operating on constants or attributes of tuples

```
select loan_number, branch_name, amount * 1.1
from loan
```

Returns a relation that is the same as the loan relation, except that the value of the attribute amount is multiplied by 1.1



Field and expression aliases

SQL allows renaming tables and columns using the AS clause:

old-name AS new-name

Example: Find the new balance of every account after subtracting 10€ in expenses

```
select account_number as acct, (balance-10) as new_bal
from account;
```

```
acct |
       new bal
A-101
                1 490.0000
A-215
                  590.0000
A-102
                  690.0000
A-305
                  790.0000
A-201
                 890.0000
A-222
                 540.0000
A-217
                  640.0000
A - 333
                  740.0000
                  840.0000
A - 444
```



Filtering rows



The WHERE clause

- The where clause specifies conditions that the table rows must satisfy (i.e., acts as a filter)
- **Example**: Find all loan numbers of loans made at the 'Central' branch with loan amounts greater than 5000€

```
select loan_number, amount
from loan
where branch_name = 'Central'
and amount > 5000
```

```
loan_number | amount
-----
L-21 | 9000.0000
```



The WHERE clause

- Comparison results can be combined using the logical connectives and, or, and not
- Comparisons can be applied to results of arithmetic expressions



The WHERE clause

- > SQL includes the **between** ... **and** operator
- Example: Find the loan number of those loans whose amounts are between 5 000€ and €10 000 (i.e., ≥ €5 000 and ≤ €10 000)

```
select loan_number
from loan
where amount between 5000 and 10000
```

```
loan_number
-----
L-93
L-11
L-16
L-20
L-21
```



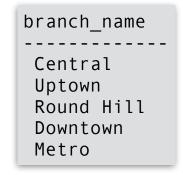
Handling Duplicates



The **SELECT DISTINCT** clause

- SQL allows duplicates in relations as well as in query results
- To force the elimination of duplicates, insert the keyword **distinct**
- **Example**: find the names of all branches in the loan relations, and remove the duplicates

```
SELECT DISTINCT branch_name
FROM loan;
```





Searching and Pattern Matching



String Pattern Matching

- The operator **LIKE** uses patterns that are described using two special characters:
 - Underscore '_': matches any (at least one) single character.
 - Percent '%': matches any sequence of variable length (including the empty sequence)

Pattern	Matches
Ana%	Ana, Anastasia, Anacleto, Anafora, Ananás,
%ana%	Banana, Cana, Panacota, Feliciana,
Ana_%s	Ana Santos, Ana Martins, Ananás
Ana %s	Ana Santos, Ana Martins,
	XT-09-10, 20-45-XY, 10-AD-90, AB-CD-EF,



String Pattern Matching

Example: Find the street of all customers whose name starts with "Ana"

```
select customer_name, customer_street
from customer
where customer_name like 'Ana%'
```

Ana, Anastasia, Anacleto, Anafora, Ananás, ...

Example: Find the street of all customers whose name has "ana"

```
select customer_name, customer_street
from customer
where customer_name like '%ana%'
```



Exercise



Exercise

- A. Get the data of all accounts
- B. Get the number of all accounts with a balance between 700€ and 900€
- C. Get the name of all customers whose street name starts in 'C'

```
FROM account

SELECT account_number
FROM account
WHERE balance BETWEEN 700 AND 900

SELECT customer_name
FROM customer
WHERE customer_street LIKE 'C%'
```



Sorting Rows



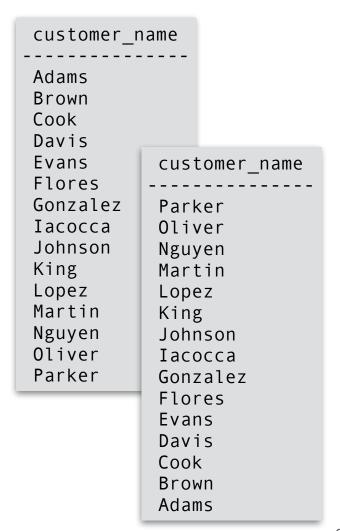
Changing the display order of rows

 List in alphabetic order the names of all customers

```
FROM customer
ORDER BY customer_name;
```

We may specify desc for descending order or asc for ascending order (the default)

```
FROM customer
ORDER BY customer_name DESC
```





Order in DBMS

Order-dependent Storage

- Each application may require data sorted in a different way; even distinct screens in the same application
- If records were stored ordered by a certain column, then sorting them by other order is costly

Order-independent Storage

- The order by which records are organised in the storage should be irrelevant (data are as a principle independent from the application)
- DBMSs provide efficient ways to sort records in whatever way the application requires



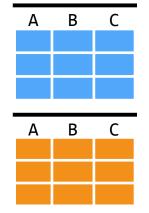
Set Operations

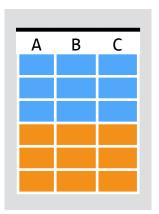


Set Operations

Union of tables

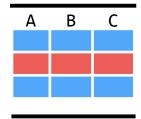
SELECT ...
UNION
SELECT ...



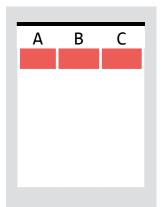


Intersection of tables

SELECT ...
intersect
SELECT ...

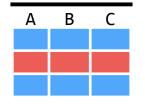


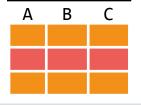


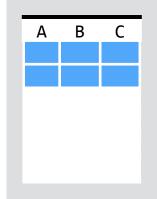


Difference of tables

SELECT ...
except
SELECT ...









Set Operations

Find all customers who have a loan, an account, or both

```
(SELECT customer_name FROM depositor)
UNION
(SELECT customer_name FROM borrower)
```

Find all customers who have both a loan and an account

```
(SELECT customer_name FROM depositor)
INTERSECT
(SELECT customer_name FROM borrower)
```

Find all customers who have an account but no loan

```
(SELECT customer_name FROM depositor)
EXCEPT
(SELECT customer_name FROM borrower)
```



Navigating by Joining Relations



JOIN Example

List the names of all employees along their corresponding departments

<u>Employee</u>			_	Department		
eid	name	did		did	name	loc
1	Alice	X	•••	X	Marketing	Damaia
2	Barbara	Υ	•	Y	Sales	Amadora
3	Carlos	Z	•••	Z	Production	Buraca
4	Duarte	W	•	W	Support	Cacém

FROM Employee e JOIN Department d
ON e.did = d.did

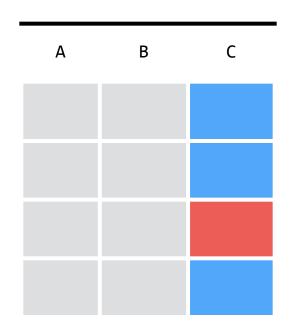
e.name	d.name		
Alice	Marketing		
Barbara	Sales		
Carlos	Production		
Duarte	Support		

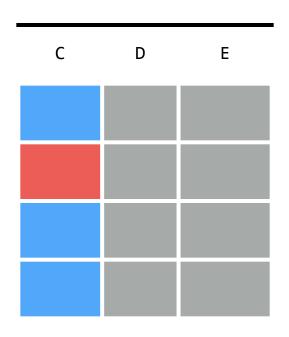
JOIN Operations



Join Operations

Joining two tables

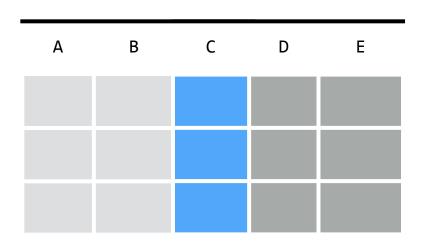






Join Operations

Results on a table that ties together the records that agree on (that match) the values of a column (or columns) subject to a condition



Like extending one table with another



Behaviour of INNER JOIN



Record Mapping

Employee			Department			
eid	name	did		did	name	loc
1	Alice	X	•••	X	Marketing	Damaia
2	Barbara	Υ		Υ	Sales	Amadora
3	Carlos	Z	•••	Z	Production	Buraca
4	Duarte	W	•••	W	Support	Cacém

SELECT *
FROM Employee e JOIN Department d
 ON e.did = d.did

e.eid	e.name	did	d.name	e.loc
1	Alice	X	Marketing	Damaia
2	Barbara	Υ	Sales	Amadora
3	Carlos	Z	Production	Buraca
4	Duarte	W	Support	Cacém

Values are all distinct

One-to-One = **Mapping**

Record Multiplication

Employee			Department			
eid	name	did		did	name	loc
1	Alice	X	•	X	Marketing	Damaia
2	Barbara	X		Υ	Sales	Amadora
3	Carlos	Z	•••	Z	Production	Buraca
4	Duarte	W	•••	W	Support	Cacém

SELECT *
FROM Employee e JOIN Department d
ON e.did = d.did

Some values repeat

e.ei	ename	did	dname	d.loc
1	Alice	X	Marketing	Damaia
2	Barbara	X	Marketing	Damaia
3	Carlos	Z	Production	Buraca
4	Duarte	W	Support	Cacém

Many-to-One = Multiplication

Record Filtering

Employee			_	Department		
eid	ename	did		did	dname	loc
1	Alice	X		X	Marketing	Damaia
2	Barbara	Υ		Υ	Sales	Amadora
3	Carlos	Т		W	Support	Cacém
4	Duarte	W				

```
FROM Employee e JOIN Department d
ON e.did = d.did
```

Some values don't have a match

e.eid	e.ename	did	d.dname	e.loc
1	Alice	X	Marketing	Damaia
2	Barbara	Υ	Sales	Amadora
4	Duarte	W	Production	Buraca

Zero-to-One = **Filter**

Behaviour of OUTTER JOIN



LEFT OUTER JOIN

Employee _					Departme	nt
eid	name	did		did	name	loc
1	Alice	X	•••	Χ	Marketing	Damaia
2	Barbara	Υ	•••	Υ	Sales	Amadora
3	Carlos	Т	•	Z	Production	Buraca
4	Duarte	W	•••	W	Support	Cacém

SELECT e.name, d.name
FROM Employee e LEFT OUTER JOIN Department d
ON d.did = d.did

e.ei	e.name	did	d.name	d.loc
1	Alice	X	Marketing	Damaia
2	Barbara	Υ	Sales	Amadora
3	Carlos	T	NULL	NULL
4	Duarte	W	Support	Buraca

Blank line with NULL values

One output row for every row of the table on the **LEFT**

RIGHT OUTER JOIN

Employee				Department		
eid	name	did		did	name	loc
1	Alice	X	•••	X	Marketing	Damaia
2	Barbara	Υ	•••	Υ	Sales	Amadora
3	Carlos	Т	•	Z	Production	Buraca
4	Duarte	W	•••	W	Support	Cacém

SELECT e.name, d.name
FROM Employee e RIGHT OUTER JOIN Department d
ON d.did = d.did

Blank line with NULL values

	e.ei	e.name	did	d.name	e.loc
	1	Alice	X	Marketing	Damaia
	2	Barbara	Υ	Sales	Amadora
-	NULL	NULL	Z	Production	Buraca
	4	Duarte	W	Support	Buraca

One output row for every row of the table on the **RIGHT**

FULL OUTER JOIN

Employee

Department

	1 3		_		<u> </u>	
eid	name	did		did	name	loc
1	Alice	X	•••	Χ	Marketing	Damaia
2	Barbara	Υ	•••	Υ	Sales	Amadora
3	Carlos	Т	• •	Z	Production	Buraca
4	Duarte	W	•••	W	Support	Cacém

FROM Employee e FULL OUTER JOIN Department d
ON d.did = d.did

e.eid	e.name	did	d.name	e.loc
1	Alice	X	Marketing	Damaia
2	Barbara	Υ	Sales	Amadora
3	Carlos	Т	NULL	NULL
	NULL	Z	Production	Buraca
4	Duarte	W	Support	Cacém

One output row for every row on EACH table

Querying the Bank Database with a Join



Find the names of the depositors with accounts with more than 750 € in balance

Step 1: Find all Accounts with balance > 750€

```
SELECT *
FROM account
WHERE balance > 750
ORDER BY balance;
```

account_number	· —	•
A-305 A-444	Round Hill Downtown Uptown	·

Step 2: Find the depositors of those accounts

```
FROM depositor
ORDER BY account_number;
```

```
account number
customer name |
Cook
                  A-101
Johnson
                  A-101
Cook
                  A - 102
Johnson
                  A - 201
Brown
                  A - 215
Lacocca
                  A - 217
Evans
                  A - 222
Flores
                  A - 305
                  A-333
01iver
                  A - 444
Brown
```



Find the **names of the customers** with accounts having more than 750 € in balance

Joining Account and Depositor on the attribute 'account_number'

```
SELECT *
FROM account a JOIN depositor d
ON a.account_number = d.account_number
WHERE balance > 750;
```

We are selecting all attributes

```
account number
                   branch name
                                   balance
                                               customer name
A - 305
                   Round Hill
                                  800.0000
                                               Flores
                   Uptown
A - 201
                                  900.0000
                                               Johnson
A - 444
                   Downtown
                                  850.0000
                                               Brown
```



Find the **names of the customers** with accounts having more than 750 € in balance

```
SELECT customer_name
FROM account a JOIN depositor d
ON a.account_number = d.account_number
WHERE balance > 750;
```

customer_name
----Flores
Johnson
Brown

Still joining tables but only selection the customer name



Cartesian Product



Cartesian Product

Employee

Department

	1 2		_		•	
eid	name	did	_	did	name	loc
1	Alice	X		X	Marketing	Damaia
2	Barbara	Υ		Υ	Sales	Amadora
3	Carlos	Z		Z	Production	Buraca

SELECT e.name, d.name
FROM Employee e, Department d

e.ei	e.name	e.did	d.di	d.name	e.loc	
1	Alice	X	X	Marketing	Damaia	
2	Barbara	Υ	Χ	Marketing	Damaia	
3	Carlos	Т	X	Marketing	Damaia	
1	Alice	X	Υ	Sales	Amadora	
2	Barbara	Υ	Υ	Sales	Amadora	
2	Canloc	Т	V	Sales	Amadona	
Employee X Department						
3	Carlos	Т	Z	Production	Buraca	

The FROM clause lists the tables involved in the query

Cartesian Product with a Filter

Employee			_	Department		
eid	name	did		did	name	loc
1	Alice	X		Χ	Marketing	Damaia
2	Barbara	Υ		Υ	Sales	Amadora
3	Carlos	Z		Z	Production	Buraca

SELECT e.name, d.name
FROM Employee e, Department d
WHERE e.did = d.did

e.eid	e.name	did	d.name	e.loc
1	Alice	X	Marketing	Damaia
2	Barbara	Υ	Sales	Amadora
3	Carlos	Z	Production	Buraca

Query has a filter

Rows that do not match the filter are dropped

```
SELECT *
FROM Employee e JOIN Department d
ON e.did = d.did
```

The Cartesian product

depositor × account

```
SELECT *
FROM depositor d, account a
```

customer_name	account_numbe	r account_number	branch_name	balance	
Johnson Johnson Johnson Johnson Johnson Johnson Johnson Johnson Brown Brown Brown Brown Brown Brown Brown	A-101 A-101 A-101 A-101 A-101 A-101 A-101 A-101 A-215 A-215 A-215	A-101 A-215 A-102 A-305 A-201 A-222 A-217 A-333 A-444 A-101 A-215 A-102	Downtown Metro Uptown Round Hill Uptown Central University Central Downtown Downtown Metro Uptown Round Hill	+	
Brown Brown Brown	d.account_number = a.account number?				
Brown	A-215	i A-444	Downtown	750.0000 850.0000	



Cartesian product with a filter

What happens when we filter a Cartesian product?

```
SELECT *
FROM depositor d, account a
WHERE d.account_number = a.account_number;
```

```
account number | account number
customer_name
                                                  branch name
                                                                balance
 Johnson
                A-101
                                 A-101
                                                   Downtown
                                                                 500.0000
                                 A-215
                A-215
                                                  Metro
                                                                 600.0000
 Brown
                A-102
                                 A-102
                                                                700.0000
Cook
                                                   Uptown
Cook
                A-101
                                 A-101
                                                   Downtown
                                                                 500.0000
 Flores
                A - 305
                                 A-305
                                                   Round Hill
                                                                 800.0000
 Johnson
                A-201
                                 A-201
                                                   Uptown
                                                                 900.0000
                A-217
                                 A-217
                                                   University
                                                                 650.0000
 Iacocca
                                                                 550.0000
                                                                 750.0000
     SELECT *
                                                                 850.0000
     FROM account a JOIN depositor d
(1(
       ON a.account_number = d.account_number
```



Find the names of the cities of the customers with accounts having more than 750 € in balance

```
SELECT customer_city
FROM account a, depositor d, customer c
WHERE a.account_number = d.account_number
AND c.customer_name = d.customer_name
AND balance > 750;
```

```
customer_city
Oporto
Braga
Cascais
```

