## Mathematical Logic (MATH6/70132;P65) Problem Sheet 5

[1] Prove (or at least, sketch a proof of) the following version of the Lindenbaum Lemma (2.5.2) which was used in the proof of 2.5.3:

Suppose  $\mathcal L$  is a countable first-order language and  $\Sigma$  is a consistent set of closed  $\mathcal L$ -formulas. Then there is a consistent set  $\Sigma^* \supseteq \Sigma$  of closed  $\mathcal L$ -formulas such that, for every closed  $\mathcal L$ -formula  $\phi$  either  $\Sigma^* \vdash \phi$  or  $\Sigma^* \vdash (\neg \phi)$ .

- [2] Describe a language with equality  $\mathcal{L}^{=}$  which is appropriate for groups (see the example after 2.2.7).
- (a) Write down a closed  $\mathcal{L}^=$ -formula  $\gamma$  whose normal models are precisely the groups. (You can use traditional mathematical notation if you wish.)
- (b) Write down a set  $\Sigma$  of closed  $\mathcal{L}^=$ -formulas whose normal models are precisely the infinite groups. [Hint: you may use the formulas  $\sigma_n$  in Question 2(a) for Problem class 2 (Lecture 15).]
- (c) Suppose that  $\phi$  is a closed  $\mathcal{L}^=$ -formula such that for every  $n \in \mathbb{N}$  there is a group with at least n elements which is a model of  $\phi$ . Show that there is an infinite group which is a model of  $\phi$ .
- (d) Show that there is no set  $\Delta$  of closed  $\mathcal{L}^{=}$ -formulas whose normal models are precisely the finite groups.
- (e) Show that there is no closed  $\mathcal{L}^{=}$  formula whose normal models are precisely the infinite groups.
- (f) (Harder) Is there a closed  $\mathcal{L}^=$ -formula  $\sigma$  which has a normal model and is such that any normal model of  $\sigma$  is an infinite group?
- [3] Suppose  $\mathcal{L}^{=}$  is a first-order language with equality and a single binary relation symbol R. A graph is a normal model of the closed formula  $\gamma$ :

$$(\forall x_1)(\forall x_2)((\neg R(x_1, x_1)) \land (R(x_1, x_2) \to R(x_2, x_1))).$$

- (i) Find a closed formula  $\tau$  with the property that there is a finite normal model of  $\gamma \wedge \tau$  whose domain has n elements iff n is divisible by 3.
- (ii) (Hard) Can you find a closed  $\mathcal{L}^=$ -formula which has no finite models and has some infinite graph as a normal model?
- [4] Suppose  $\mathcal{L}^=$  is a first-order language with equality (=) and a single binary relation symbol  $\leq$ . A linear order is a normal  $\mathcal{L}^=$ -structure  $\langle A; \leq_A \rangle$  such that the relation  $\leq_A$  is reflexive, transitive and such that for distinct  $a,b\in A$  exactly one of  $a\leq_A b$ ,  $b\leq_A a$  holds.

Let  $\Sigma$  be the set of all closed  $\mathcal{L}^{=}$ -formulas which are true in all *finite* linear orders.

- (i) Prove that any normal model of  $\Sigma$  is a linear order with a least element and a greatest element.
- (ii) Prove that any normal model of  $\Sigma$  (with at least 2 elements) is not dense.
- (iii) Prove that  $\Sigma$  has an infinite normal model.
- (iv) Find a closed  $\mathcal{L}^{=}$ -formula  $\phi$  such that neither  $\phi$  nor  $(\neg \phi)$  is a consequence of  $\Sigma$ .

- [5] Suppose  $\mathcal{L}^{=}$  is a first order language with equality ( = ) and a single binary relation symbol R. Write down what it means for two normal  $\mathcal{L}^{=}$ -structures to be isomorphic (see the problem class in week 6)?
- (i) Write down a set  $\Sigma$  of closed  $\mathcal{L}^=$ -formulas such that the normal models of  $\Sigma$  are the normal  $\mathcal{L}^=$ -structures in which R is interpreted as an equivalence relation in which all equivalence classes have size 2 or 3 and there are infinitely many equivalence classes of size 2 and infinitely many of size 3.
- (ii) Explain why any two countable normal models of  $\Sigma$  are isomorphic.
- [6] Suppose  $\mathcal{L}^{=}$  is a first-order language with equality having just a single 1-ary function symbol f (and no other relation, function or constant symbols apart from =).
- (i) What does it mean to say that two normal  $\mathcal{L}^{=}$ -structures  $\mathcal{A}$ ,  $\mathcal{B}$  are isomorphic?
- (ii) Write down a set  $\Sigma$  of closed  $\mathcal{L}^=$ -formulas such that  $\langle A; \bar{f} \rangle$  is a normal model of  $\Sigma$  if and only if:  $\bar{f}: A \to A$  is a bijection and for every  $n \in \mathbb{N}$ , the function  $\bar{f}^n: A \to A$  (obtained by applying  $\bar{f}$  n times) has no fixed points.
- (iii) Find countable normal models  $A_0, A_1, A_2, \ldots$  of  $\Sigma$  such that no two of these models are isomorphic and any countable model of  $\Sigma$  is isomorphic to one of these.

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